

# Computer Networks

EDA387 / DIT663 SP1 2014/2015

#### **Internet Interconnections**

Ali Salehson

2

0

1

### **Course Literature**

Douglas Comer

"Internetworking with TCP/IP"

Principles, Protocols, and Architectures

Volume 1, 5th edition Prentice Hall 2006

➤ J. F. Kurose, K. W. Ross
"Computer Networking"

\*\*A Top-Down Approach\*\*

6<sup>th</sup> edition, Pearson Education 2012

# Internetworking with TCP/IP



**Ch 3:** Internetworking Concept (review)

**Ch 4:** IPv4 Addressing

**Ch 5:** ARP protocol & operation

Ch 6: IPv4 Protocol (self-study)

**Ch 7:** Forwarding IP Datagrams

Ch 8: ICMP and TCP/IP Utilities (self-study)

Ch 9: IPv4 CIDR



2

0

1

# NANCE

# Internetworking with TCP/IP

**Ch 23:** Domain Name System (DNS)

**Ch 31: IPv6 & ICMPv6** 

2

0

4

# NANCES

### **Internet Services**

## Practical Perspective (lab 4h)

- ➤ Domain Name System (DNS)
  - Using dig in Linux
  - DNS-client interacts with DNS-servers
- > IPv6 and ICMPv6
  - Using ping6 and traceroute6 in Linux
  - Capturing IPv6 packets with Wireshark
  - Neighbor Discovery and Autoconfig









# Internet Interconnections Underlying Network Technologies

# Ethernet and MAC frames Review

Ali Salehson







# Network Technologies (review)

### Outline

- > Review basic network concepts
- > Examine physical networks
- > Learn physical addressing
- > Focus on Ethernet technology
- > Study MAC frame









# The TCP/IP Concept

- - 2
  - 0
  - 1
  - 4

- > Use existing network hardware
- > Interconnect networks

The challenge is to accommodate all possible network hardware.

# **Basic Categories of Network**

- 2
- 0
- 1
- 4

- Connection-oriented
  - circuit-switching or virtual-circuit packet-switching
- Connectionless
  - datagram packet-switching frame broadcasting



### **Connection Oriented**

#### **Paradigm**

- establish a "connection" through the network
- send / receive data over the same connection
- terminate the connection
- > Can guarantee bandwidth and reserve resources
- ➤ Works well with real-time applications (e.g. voice)

#### **Examples**

- PSTN and ISDN (circuit-switching)
- Frame Relay and ATM (VC packet-switching)

9

0

1



#### **Connectionless Networks**

#### **Paradigm**

- form a "packet" of data to transmit
- send packet as soon as possible over network
- > Each packet travels as datagram independently
- > Packet includes identification of the destination
- **Each packet can be of a different size**
- Works well with file transfer applications

#### **Packet-switching Networks**

- Local Area Network (LAN)
- Wide Area Network (WAN)

2

0

1

#### **Local Area Networks**

- > Engineered for
  - Limited distance to a locale
  - Direct connection among computers
- > Low cost
- High speed
- > Need Medium Access Control (MAC)

2

0

1

#### Wide Area Networks

- Engineered for
  - Long distances
  - Indirect connection via special-purpose hardware (DTE/DCE)
- > Higher cost
- Lower capacity in shared links
- **Need Intermediate Network Devices**

2

0

1

#### **Packet-switched Networks**

- Wide Area Networks
  - ARPAnet, Internet
  - Common carrier services
- Leased Line services
  - Point-to-point connections
- Local Area Networks
  - Ethernet
  - Wi-Fi

2

0

1

## ARPAnet (1969-1989)

Original backbone of Internet

Wide area network around which TCP/IP was developed

Funding from Advanced Research Project Agency (ARPA)







# NANCES

#### **WANs from Carriers**

- Point-to-point digital circuits
  - T-lines (e.g. T1 = 1.544 Mbps)
  - E-lines (e.g. E1 = 2.048 Mbps)
  - OC-line (e.g. OC-3 = 155 Mbps)

- Packet-switching services also available by:
  - X.25, Frame Relay, ATM







# NANCE

#### **Point-to-Point Network**

- > Any direct connection between two nodes
  - Leased line
  - Connection between two routers
  - Dialup connection
  - DSL broadband connection
- Link-level protocol required for framing
- > TCP/IP views it as an independent network

2

0

1

### LANs: Focus on Ethernet

- extremely cheap and popular
- can run over
  - copper (twisted pair)
  - optical fiber
- many physical standards
  - 10 Base T (Legacy Ethernet 10 Mbps)
  - 100 Base T (Fast Ethernet 100 Mbps)
  - 1000 Base T (Gigabit Ethernet 1 Gbps)
  - 10 Gigabit Ethernet over fiber
- one IEEE MAC standard 802.3

2

0

1

# NANCES

#### **Hardware Address**

- Unique number assigned to each network interface (NIC) on a machine
- > Used to identify destination for a frame
- **Known as (all terms are almost equal!)** 
  - Physical address
  - Link-layer address (layer-2 address)
  - Hardware address
  - Medium Access Control (MAC) address

2

0

1



#### Use of Hardware Address

- > Sender supplies
  - Destination address
  - Source address (in most technologies)
- > Network Hardware
  - Uses destination address to forward frame
  - Delivers frame to proper machine interface
- > Each network technology defines its own scheme of addressing
  - Static: assigned by hardware vendor
  - Configurable: assigned by customer
  - Dynamic: assigned by software at startup

2

0

1



#### LAN MAC Address

- > MAC (link-layer or physical) address:
  - used to get frame from one interface to another physically-connected interface (same network)
  - 48-bit MAC address (for most LANs)
  - normally burned in the adapter ROM
- ➤ MAC address allocation administered by IEEE:
  Institute of Electrical and Electronics Engineers
- ➤ Manufacturer gets portion of MAC address space (to assure uniqueness)
- > MAC addressing is flat, no hierarchy
  - LAN card may be moved from one LAN to another



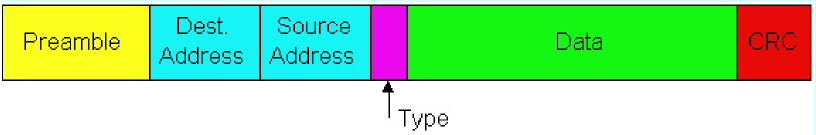




### **Ethernet Frame Format**

 Sending adapter encapsulates IP datagram (or other higher layer protocol data) in a frame

Sync | Header | Payload | Trailer

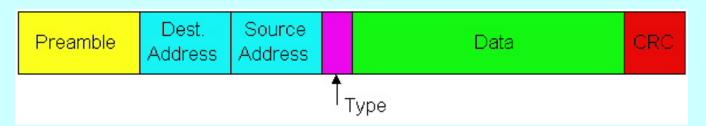


#### Preamble: (7 + 1) bytes

7 bytes with pattern 10101010 followed by one byte as: Start of Frame Delimiter (SFD) with pattern 10101011 both used to synchronize receiver and sender clock rates and indicating the start of the frame.



#### **Ethernet Frame Format**



Addresses: 2 x 6 bytes

if adapter receives frame with matching destination address, or with broadcast address (e.g. ARP packet), it passes data in frame to network layer protocol

Type/Length: 2 bytes (0x0800 for IP) / (0x002E → 0x05DC) in Ethernet type II, indicates higher layer protocol (mostly IP) in IEEE 802.3 indicates length of data field (LLC frame)

**Data:** 46 – 1500 bytes (variable)

Min. length makes sure collisions are always detected

Max. length indicates Maximum Transmission Unit (MTU) on Ethernet

CRC: 4 bytes

at sender inserted FCS bits, to check for errors at receiver

2

0

1

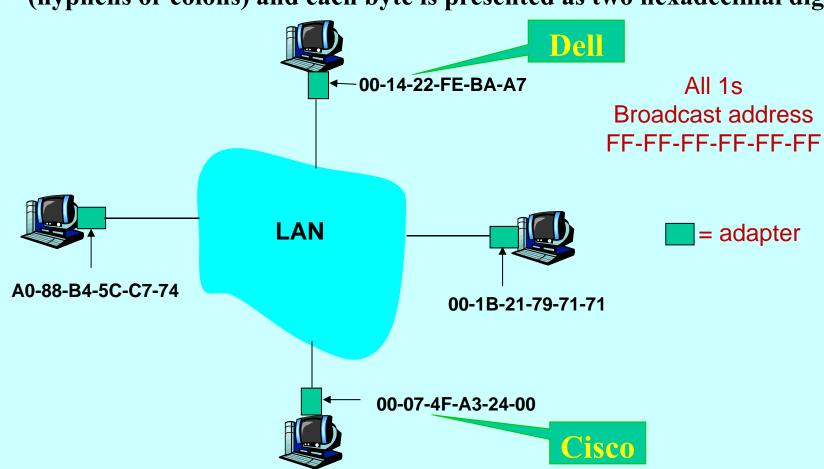
# NANCES

#### **MAC Addresses**

 The 48-bit MAC address is given as 6 bytes, separated by periods (hyphens or colons) and each byte is presented as two hexadecimal digits 2

0

1



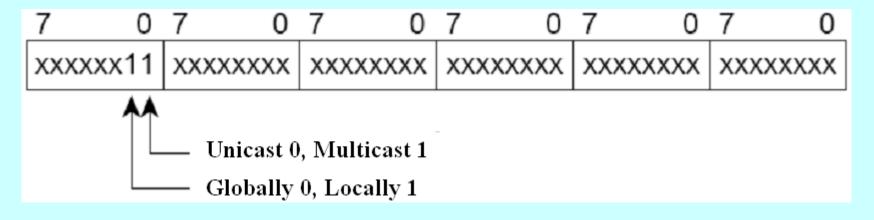
# NANCE OF THE PARTY OF THE PARTY

### **MAC** Unicast Address

- Ethernet hardware addresses are 48 bits, expressed as 12 hexadecimal digits (0-9, A-F often capitalized).
- ➤ The first (left) 6 digits should match the vendor of the Ethernet interface. These high-order 3 octets (6 hex digits) are also known as the Organizationally Unique Identifier (OUI).
- The last (right) 6 digits specify the serial number given for the interface by that vendor.
- These addresses are physical unicast addresses, and not *multicast* nor *broadcast*, so the second hex digit (reading from the left) will be even, not odd.

### **MAC** Multicast Address

- ➤ An Ethernet multicast address consists of the multicast bit, the 23-bit vendor component, and the 24-bit group identifier assigned by the vendor.
- ➤ For example, DEC is assigned the vendor component (first 3 bytes) 08-00-2B, so multicast addresses assigned by DEC will have the first 24-bits 09-00-2B.
- > The multicast bit is the low-order bit of the first byte, which is "the first bit on the wire".











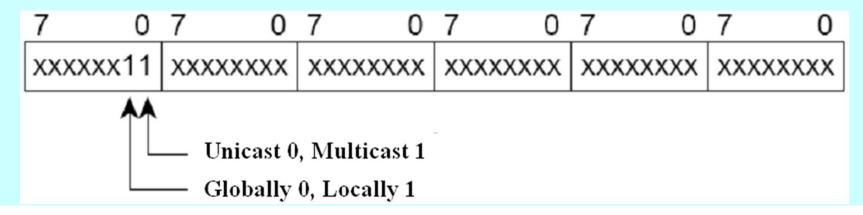
#### **MAC Multicast Address**

#### **ONLY** destination address

#### **Examples:**

01-80-C2-00-00 used by Spanning Tree Protocol 01-00-5E-00-00 is the Internet Multicast

Octets transmitted left-to-right, and within octets: on the wire, bits transmitted from right to left (LSB) in memory, bits placed from left to right (MSB)









### **CSMA/Collision Detection**

- 1. Adapter receives datagram from network layer and creates frame
- 2. Adapter waits until it senses channel idle, plus 96 bit times, then transmits frame.
- 3. If adapter transmits entire frame without detecting another transmission, the adapter is done with frame.

- 4. If adapter detects another transmission while it is self transmitting, aborts and sends jam signal (48 bits) to warn other nodes about collision
- 5. After aborting, adapter enters **exponential backoff**: after *m*th collision, adapter chooses *K* at random from {0,1,2,...,2<sup>m</sup>-1} until K is max 1023 (m = 10).

  Adapter waits K\*512 bit times, and returns to Step 2 **Bit time is 10 ns for 100 Mbps**

max. wait time is about 5 ms

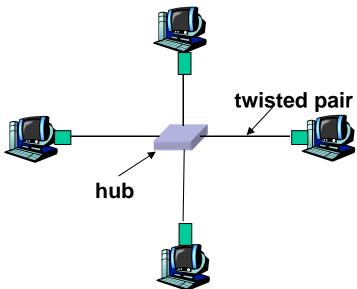
for K = 1023,

# NAME OF THE PARTY OF THE PARTY

### **Shared Ethernet: Hubs**

Hubs are essentially physical-layer multiport repeaters:

- bits coming in one link repeated out all other links
- at same rate
- no frame buffering
- adapters detect collisions and do CSMA/CD



2

0

1

## LAN Bridge

- ➤ Hardware device that connects two LAN segments and makes them appear to be a single LAN
- Inspects frames to learn which computers are on which side of the bridge
- > Uses hardware addresses to filter
- > Repeats frames from one LAN segment to another
- ➤ Introduces delay of 1 frame-time (store-and-forward)
- ➤ Does not forward collisions or noise (checks CRC and performs CSMA/CD)
- ➤ Called Layer 2 device











### **Switched Ethernet**

#### link-layer device

- multi-port LAN bridge
- store and forward Ethernet frames
- examine frame header and selectively forward frame based on MAC destination address

#### Auto-negotiation

detection of devices and their communication capabilities

#### • Transparent

hosts are unaware of presence of switches

#### Plug-and-play, self-learning

The switch forwarding table does not need to be configured











## **Self Learning**

- Switch learns which MAC addresses of hosts can be reached through which interface
  - When frame received, switch "learns" location of sender i.e. incoming LAN segment
  - Records sender/location pair in switch table
- Entry in switch table
   [MAC Address, Interface, Time Stamp]
- Stale entries in table dropped (aging time may be configured or set to a default)



2

0

1

## Summary

- > TCP/IP is designed to use all types of networks
  - Connection-oriented
  - Connectionless
  - Local Area Network (LAN)
  - Wide Area Network (WAN)
  - Point-to-point link
  - Set of bridged networks
- > Recall: each technology can define its own addressing scheme
- ➤ Heterogeneous networks imply potential for heterogeneous addressing
- Conclusion: cannot rely on hardware addressing











# **Questions?**



2

0

1



### Thank You!



### **Internet Interconnections**

# Internetworking Concept Architectural Model Review

Ali Salehson

2

0

1

# The TCP/IP Internet Concept

- > Use available physical networks
- > Interconnect underlying networks
  - Network of networks
  - Providing universal communication
- > Devise abstractions that hide
  - Underlying architecture
  - Hardware addresses
  - Routes and alternative paths



2

0

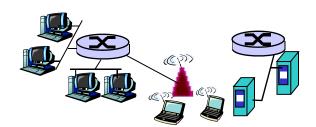
1

#### **Internet Interconnection**

- > Network uses active intermediate system
- Each network sees an additional computer system attached
- > IS is an IP router (originally called IP gateway)







In a TCP/IP internet, special computers called IP routers or IP gateways provide interconnections among physical networks.

2

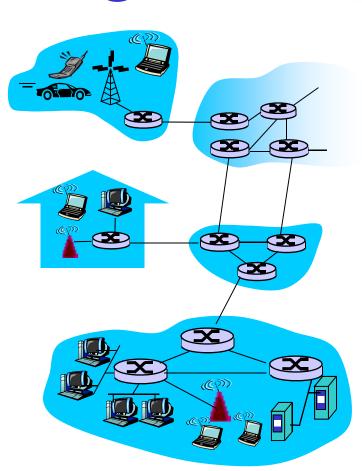
0

4

## Internetworking

Uses multiple IP routers

- > Ensures that each network is reachable
- Do not need router between *each* pair of networks (packets may be passed through other networks)



# Packet Transmission Paradigm

- > Source computer
  - Generates a packet
  - Sends across one network to a router
- > Intermediate router
  - Forwards packet to "next" router
- > Final router
  - Delivers packet to destination

Routers use the destination network, not the destination computer, when forwarding packets.







## **Architectural Terminology**

- > End-user System (ES) is called *host* 
  - Connects to physical network
  - Normally many hosts per network
  - Possibly more than one network connection per host (multi-homed)
- ➤ Dedicated systems called *IP gateways* or *IP routers* interconnect networks
  - Router connects two or more networks









## Summary

- > TCP/IP Internet is set of interconnected (possibly heterogeneous) networks
- > Routers provide interconnection
- > End-user systems are called host computers
- ➤ Internetworking introduces abstractions that hide details of underlying networks

#### **Needs:**

- Addressing model and relationship to hardware addresses
- Format of packet as it travels through Internet



2

0

1

# **Questions?**



2

0

1



### Thank You!