

CS258: Information Theory

Fan Cheng



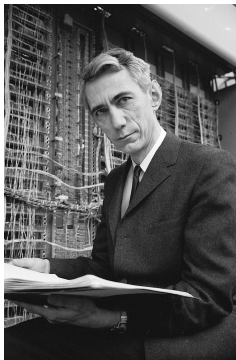
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Lecture 1: Introduction

- Instructor: Fan Cheng, Rm 3-513, SEIEE
(<http://www.cs.sjtu.edu.cn/~chengfan/>)
Office hour: By appointment
TA: TBA
- Textbook: David J.C. MacKay, “Information Theory, Inference, and Learning Algorithms,” Cambridge Press, 2005
(<http://www.inference.org.uk/itprnn/book.html>)
- 16 Weeks := 14 lectures + 1 in-class midterm + 1 Q&A
- Grade policy := 50% final + 30% midterm + 10% attendance + 10% homework

Birth of Information Theory

“A Mathematical Theory of Communication,” Bell System Technical Journal, 27 (3): 379-423, July, 1948.



Claude. E. Shannon
(1916-2001)

[https://en.wikipedia.org/
wiki/Claude_Shannon](https://en.wikipedia.org/wiki/Claude_Shannon)

The fundamental problem of communication is that of reproducing at one point either exactly or approximately a message selected at another point.

C. E. Shannon, 1948

IEEE Information Theory Society

<http://www.itsoc.org>

IEEE Transactions on Information Theory

[http://ieeexplore.ieee.org/xpl/
RecentIssue.jsp?punumber=18](http://ieeexplore.ieee.org/xpl/RecentIssue.jsp?punumber=18)

Shannon: father of information theory

Mathematician

Ph.D. in Mathematics from MIT. Worked at AT&T Bell Labs and RLE in MIT

Electrical engineer

Mater's Thesis: electrical applications of Boolean algebra could construct any logical, numerical relationship

Cryptographer

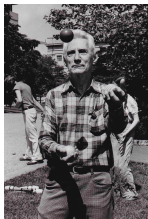
"A Mathematical Theory of Cryptography," 1949.

Friend of Turing

For two months early in 1943, Shannon came into contact with the leading British mathematician Alan Turing. Shannon and Turing met at teatime in the cafeteria. Turing showed Shannon his 1936 paper that defined what is now known as the "Universal Turing machine"



Magnetic mouse



Juggling



Unicycling

Topics in IT

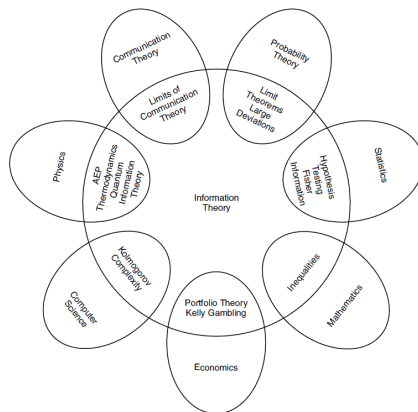
Big Data Analytics
Coding for Communication and Storage
Coding Theory
Combinatorics and Information Theory
Communication Theory
Complexity and Computation Theory
Compressed Sensing and Sparsity
Cryptography and Security

Detection and Estimation
Distributed Storage
Emerging Applications of Information Theory
Information Theory and Statistics
Information Theory in Biology
Information Theory in Computer Science
Statistical/Machine Learning
Network Coding and Applications

Network Data Analysis
Network Information Theory
Optical Communication
Quantum Information and Coding Theory
Shannon Theory
Signal Processing
Source Coding and Data Compression
Wireless Communication and Networks

<https://www.isit2018.org/authors/call-for-papers/>

Information theory to other fields



1

- Information Theory and Reliable Communication, 1st Edition, Robert G. Gallager
- Elements of Information Theory, 2nd Edition (Wiley Series in Telecommunications and Signal Processing), Thomas M. Cover, Joy A. Thomas
- Information Theory: Coding Theorems for Discrete Memoryless Systems, 2nd Edition, Imre Csiszar, Janos Korner
- Information Theory, Inference and Learning Algorithms, David J. C. MacKay
- A First Course in Information Theory (Information Technology: Transmission, Processing and Storage), 1st Edition, Raymond W. Yeung

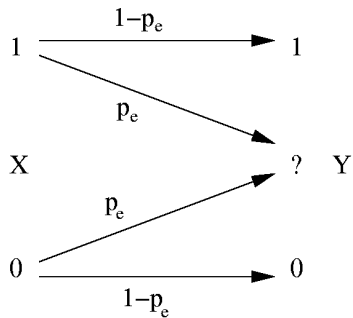
Course plan

- Elements: entropy, mutual information, information divergence, etc.
- Data compression
- Noisy-channel coding
- Probability and inference
- Neural networks
- Low-density parity-check codes

Prerequisite: Probability theory, mathematical analysis, matrix theory

In class: pen and paper

Information theory: An example



Binary Erasure Channel

For random variable X defined on alphabet \mathcal{X} , its mean and variance is defined as

$$\mathcal{E}(X) := \sum_{x \in \mathcal{X}} xp(x)$$

$$\text{Var}(x) := \mathcal{E}(X^2) - (\mathcal{E}(X))^2$$

- $\mathcal{E}(X_1 + X_2) = \mathcal{E}(X_1) + \mathcal{E}(X_2)$
- If X_1 and X_2 are independent, then
 $\text{Var}(X_1 + X_2) = \text{Var}(X_1) + \text{Var}(X_2)$

Some probability distributions: Bernoulli, Binomial, Poisson, Gauss, etc.

A function f on (a, b) is called convex iff for any x_1, x_2 in (a, b)

$$f\left(\frac{x_1 + x_2}{2}\right) \leq \frac{f(x_1) + f(x_2)}{2}$$

- If $f(x)$ is twice differentiable, then $f(x)$ is convex iff $f''(x) \geq 0$
- If $f(x)$ is twice differentiable, then $f(x)$ is minimized iff $f'(x) = 0$
- (Jensen's inequality) $f(x)$ is convex in (a, b) ,

$$f(\mathcal{E}(X)) \leq \mathcal{E}f(X)$$

Take $f := e^x, \sin(x), \cos(x), x^2, x^3$ for example

Binomial distribution

A bent coin has probability f of coming up heads. The coin is tossed N times. What is the probability distribution of the number of heads, r ? What are the mean and variance of r ?

$$p(r|f, N) = \binom{N}{r} f^r (1-f)^{N-r}$$

$$\mathcal{E}(r) = Nf$$

$$\text{Var}(r) = Nf(1-f)$$

Approximating $x!$ and $\binom{N}{r}$

Stirling's approximation

$$x! \simeq x^x e^{-x} \sqrt{2\pi x} \iff \ln x! = x \ln x - x + \frac{1}{2} \ln 2\pi x$$

- Poisson distribution: $P(r|\lambda) = e^{-\lambda} \frac{\lambda^r}{r!}$
- When λ is large and $r \rightarrow \lambda$, $P(r|\lambda) \rightarrow \frac{1}{\sqrt{2\pi\lambda}} e^{-\frac{(r-\lambda)^2}{2\lambda}}$
- Plug $r = \lambda$

$$\ln \binom{N}{r} \simeq (N-r) \ln \frac{N}{N-r} + r \ln \frac{N}{r}$$

Binary Entropy Function

$$H_2(x) = -x \log x - (1-x) \log(1-x)$$

$$\binom{N}{r} \simeq 2^{NH_2(r/N)}$$

Exercise

- Plot $H_2(x)$ in python
- $H_2(x)$ is symmetric at $x = \frac{1}{2}$
- $H_2(x)$ is maximized at $x = \frac{1}{2}$
- Let $H_2^{-1}(x)$ be the inverse of $H_2(x)$ and $H_2^{-1}(x) \in [0, 1/2]$. For $p \in [0, 1]$, define $p * x := (1 - p)x + p(1 - x)$. Prove that $H_2(p * H_2^{-1}(x))$ is convex in x