

# CS258: Information Theory

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## 1. Maximum entropy

The normal maximizes the entropy for a given variance

If  $\text{Var}(X) = a$  is fixed, then

$$h(x) \leq \frac{1}{2} \log 2\pi e a$$

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Proof.

$D(f||g) \geq 0$ , where  $g \sim \mathcal{N}(0, a)$

## 2. Entropy of a disjoint mixture

Let  $X_1$  and  $X_2$  be discrete random variables drawn according to probability mass functions  $p_1(\cdot)$  and  $p_2(\cdot)$  over the respective alphabets  $\mathcal{X}_1 = \{1, 2, \dots, m\}$  and  $\mathcal{X}_2 = \{m+1, \dots, n\}$ . Let

$$X = \begin{cases} X_1, & \text{with probability } \alpha, \\ X_2, & \text{with probability } 1 - \alpha. \end{cases}$$

- Find  $H(X)$  in terms of  $H(X_1)$ ,  $H(X_2)$ , and  $\alpha$ .

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By definition.

Question: what if  $\mathcal{X}_1 = \mathcal{X}_2$ ?

### 3. Entropy of a sum

Let  $X$  and  $Y$  be random variables that take on values  $x_1, x_2, \dots, x_r$  and  $y_1, y_2, \dots, y_s$ , respectively. Let  $Z = X + Y$ .

- (a) Show that  $H(Z|X) = H(Y|X)$ . Argue that if  $X, Y$  are independent, then  $H(Y) \leq H(Z)$  and  $H(X) \leq H(Z)$ . Thus, the addition of independent random variables adds uncertainty.
- (b) Give an example of (necessarily dependent) random variables in which  $H(X) > H(Z)$  and  $H(Y) > H(Z)$ .
- (c) Under what conditions does  $H(Z) = H(X) + H(Y)$ ?

## 4. Entropy and pairwise independence

Let  $X, Y, Z$  be three binary *Bernoulli* $(\frac{1}{2})$  random variables that are pairwise independent; that is,  $I(X; Y) = I(X; Z) = I(Y; Z) = 0$ .

- (a) Under this constraint, what is the minimum value for  $H(X, Y, Z)$ ?
- (b) Give an example achieving this minimum.

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$$Z = X + Y \pmod{2}$$



## 5. Subset inequality

Prove that

$$\frac{1}{2}[H(X_1, X_2) + H(X_2, X_3) + H(X_3, X_1)] \geq H(X_1, X_2, X_3)$$

## 5. Subset inequality

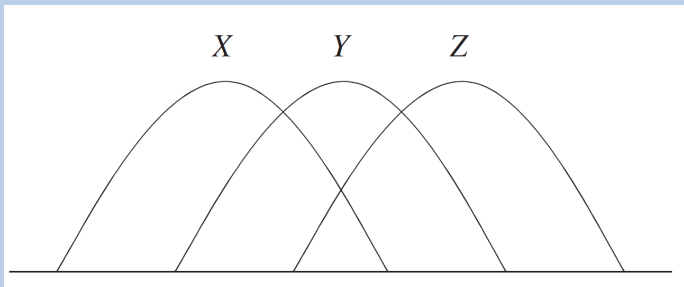
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Information diagram

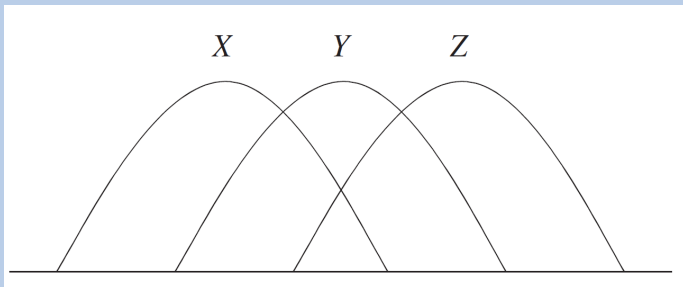
## 6. Information diagram for Markov chain

Verify that if  $X \rightarrow Y \rightarrow Z$ , then their information diagram can be simplified as



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The results can be extended to  $X_1 \rightarrow X_2 \rightarrow \dots \rightarrow X_n$

## 7. Implication and Markov chain

- (a) Prove that under the constraint that  $X \rightarrow Y \rightarrow Z$  forms a Markov chain,  $X \perp Y|Z$  and  $X \perp Z$  imply  $X \perp Y$ .
- (b) Prove that the implication in (a) continues to be valid without the Markov chain constraint.
- (c) Prove that  $Y \perp Z|T$  implies  $Y \perp Z|(X, T)$  conditioning on  $X \rightarrow Y \rightarrow Z \rightarrow T$ .

## 8. Markov chain with 4 random variables

Let  $X \rightarrow Y \rightarrow Z \rightarrow T$  form a Markov chain. Determine which of the following inequalities always hold:

- (i)  $I(X; T) + I(Y; Z) \geq I(X; Z) + I(Y; T)$
- (ii)  $I(X; T) + I(Y; Z) \geq I(X; Y) + I(Z; T)$
- (iii)  $I(X; Y) + I(Z; T) \geq I(X; Z) + I(Y; T)$

## 9. Imperfect secrecy theorem

Let  $X$  be the plain text,  $Y$  be the cipher text, and  $Z$  be the key in a secret key cryptosystem. Since  $X$  can be recovered from  $Y$  and  $Z$ , we have

$$H(X|Y, Z) = 0$$

We will show that this constraint implies

$$I(X; Y) \geq H(X) - H(Z)$$