I wanted to try out making a programming language vaguely corresponding to the double category **Rel**.

You have a calculus of "values" corresponding to one edge of **Rel** and a sort of relational calculus of "predicates" corresponding to the other edge. Satisfies judgements stating a value satisfies a predicate ought to correspond to squares.

The language of "predicates" is based off the Simply Typed Lambda Calculus (in category theory terms is closed/has exponential objects) and the language of "values" has product types (in category theory terms is Cartesian.)

This is the core framework. I've been thinking about further extensions but I want to see how far I can get characterizing **Rel** with as few language features as possible. Later on I give a few possible extensions.

Core Calculus

Types $t ::= \top \mid t \times t$

Sorts $s ::= t? \mid t!$

Values $v ::= x \mid \mathbf{tt} \mid \mathbf{fst}(v) \mid \mathbf{snd}(v) \mid (v, v)$

Predicates $p ::= x \mid p\, p \mid \mu(x\colon t?).p$

Environment $\Gamma ::= \cdot \Gamma, \mid x\colon s$

Substitutions $\sigma ::= \cdot \mid \sigma,\, v \models x$

I need a better name for the abstraction for predicates. It's a little like the μ abstraction from the $\bar{\lambda}\mu\tilde{\mu}$ calculus but different.

Typing judgments.

$$\frac{}{\Gamma \vdash \mathbf{tt}\colon \top!}$$

$$\frac{\Gamma \vdash v\colon t_0 \times t_1!}{\Gamma \vdash \mathbf{fst}(v)\colon t_0!}$$

$$\frac{\Gamma \vdash v\colon t_0 \times t_1!}{\Gamma \vdash \mathbf{snd}(v)\colon t_1!}$$

$$\frac{\Gamma \vdash v_0\colon t_0! \quad \Gamma \vdash v_1\colon t_1!}{\Gamma \vdash (v_0, v_1)\colon t_0 \times t_1!}$$

$$\frac{\Gamma \vdash p_0\colon t_0 \times t_1? \quad \Gamma \vdash p_1\colon t_0?}{\Gamma \vdash p_0\, p_1\colon t_1?}$$

$$\frac{\Gamma,\, x\colon t_0? \vdash p\colon t_1?}{\Gamma \vdash \mu x\colon t_0?.p\colon t_0 \times t_1?}$$

When a value satisfies a predicate (I need a better symbol here.)

$$\frac{v_0 \models p \quad [\sigma]}{\mathbf{fst}(v_0, v_1) \models p \quad [\sigma]}$$

$$\frac{v_1 \models p \quad [\sigma]}{\mathbf{snd}(v_0, v_1) \models p \quad [\sigma]}$$

$$\frac{v \models [x\colon\ = p_1]p_0 \quad [\sigma]}{v \models (\mu(x\colon t?).p_0)\, p_1 \quad [\sigma]}$$

$$\frac{v_1 \models p \quad [\sigma,\, v_0 \models x]}{(v_0, v_1) \models \mu(x\colon t?).p \quad [\sigma]}$$

1 Examples

Pattern matching on equality

$$(v, v) \models \mu(x\colon t?).x$$

Transposition

$$\mu(p\colon t \times t \times \top)(x\colon t)(y\colon t).p\, y\, x$$

Disjoint Unions

I am fairly confident in a simple extension to disjoint unions of sets which are sum types in the value calculus and product types in the predicate calculus.

Types $t ::= \ldots \mid \bot \mid t + t$

Values $v ::= \ldots \mid \mathbf{inl}(t, v) \mid \mathbf{inr}(t, v) \mid$ $\mathbf{match}\ v \begin{cases} v & \leftarrow \mathbf{inl}(x) \\ v & \leftarrow \mathbf{inr}(x) \end{cases}$

Predicates $p ::= \ldots \mid \mathbf{false} \mid \mathbf{left}(p) \mid \mathbf{right}(p) \mid [p; p]$

Typing judgments.

$$\frac{\Gamma \vdash v \colon t_0!}{\Gamma \vdash \mathbf{inl}(t_1, v) \colon t_0 + t_1!}$$

$$\frac{\Gamma \vdash v \colon t_1!}{\Gamma \vdash \mathbf{inr}(t_0, v) \colon t_0 + t_1!}$$

$$\frac{\Gamma \vdash v_0 \colon t_0 + t_1! \; \Gamma, \, x_0 \colon t_0! \vdash v_1 \colon t_2! \; \Gamma, \, x_1 \colon t_1! \vdash v_2 \colon t_2!}{\Gamma \vdash \mathbf{match} \; v_0 \begin{cases} v_1 & \leftarrow \mathbf{inl}(x_0) \\ v_2 & \leftarrow \mathbf{inr}(x_1) \end{cases} \colon t_2!}$$

$$\frac{}{\Gamma \vdash \mathbf{false} \colon \bot?}$$

$$\frac{\Gamma \vdash p \colon t_0 + t_1?}{\Gamma \vdash \mathbf{left}(p) \colon t_0?}$$

$$\frac{\Gamma \vdash p \colon t_0 + t_1?}{\Gamma \vdash \mathbf{right}(p) \colon t_1?}$$

$$\frac{\Gamma \vdash p_0 \colon t_0? \quad \Gamma \vdash p_1 \colon t_1?}{\Gamma \vdash [p_0; p_1] \colon t_0 + t_1?}$$

Satisfies

$$\frac{[x_0 := v_0]v_1 \models p \quad [\sigma]}{\mathbf{match} \; \mathbf{inl}(t, v_0) \begin{cases} v_1 & \leftarrow \mathbf{inl}(x_0) \\ v_2 & \leftarrow \mathbf{inr}(x_1) \end{cases} \models p \; [\sigma]}$$

$$\frac{[x_1 := v_0]v_2 \models p \quad [\sigma]}{\mathbf{match} \; \mathbf{inr}(t, v_0) \begin{cases} v_1 & \leftarrow \mathbf{inl}(x_0) \\ v_2 & \leftarrow \mathbf{inr}(x_1) \end{cases} \models p \; [\sigma]}$$

$$\frac{v \models p_0 \; [\sigma]}{v \models \mathbf{left}([p_0; p_1]) \quad [\sigma]}$$

$$\frac{v \models p_1 \; [\sigma]}{v \models \mathbf{right}([p_0; p_1]) \quad [\sigma]}$$

$$\frac{v \models p_0 \; [\sigma]}{\mathbf{inl}(t, v) \models [p_0; p_1] \quad [\sigma]}$$

$$\frac{v \models p_1 \; [\sigma]}{\mathbf{inr}(t, v) \models [p_0; p_1] \quad [\sigma]}$$

Dependent Sums

If product of sets becomes an exponential in the predicate calculus then dependent sums ought to become like Π types. So the predicate calculus effectively becomes like System-F.

Some things become awkward to interpret here though.

I also really can't figure out unpacking. It's messy if you don't want full dependent types.

Types $t ::= \ldots \mid x \mid \Sigma(x \colon \; *).t$

Sorts $s ::= \ldots \mid *$

Values $v ::= \ldots \mid \langle x \colon \; = t, v \rangle$

Predicates $p ::= \ldots \mid p \, t \mid \mathbf{M}(x \colon \; *).p$

Substitutions $\sigma ::= \ldots \mid \sigma, \, t \models x$

Not really good at the typing judgements for dependent sum types.

$$\frac{\Gamma \vdash t_0 \colon \; * \quad \Gamma, \, x \colon \; * \vdash v \colon t_1!}{\Gamma \vdash \langle x := t_0, v \rangle \colon \Sigma(x \colon \; *).t_0!}$$

$$\frac{\Gamma \vdash p \colon \Sigma(x \colon \; *).t_1? \quad \Gamma \vdash t_0 \colon \; *}{\Gamma \vdash p_0 \, t_0 \colon [x \colon \; = t_0]t_1?}$$

$$\frac{\Gamma, \, x \colon \; * \vdash p \colon t?}{\Gamma \vdash \mathbf{M}(x \colon \; *).p \colon \Sigma(x \colon \; *).t?}$$

I can't figure out satisfaction at all.

$$\frac{v \models [x := t]p \quad [\sigma]}{v \models (\mathbf{M}(x \colon \; *).p) \, t \quad [\sigma]}$$

$$\frac{v \models p \; [\sigma, \, t \models x]}{\langle x := t, v \rangle \models \mathbf{M}(x \colon \; *).p \quad [\sigma]}$$