

Task 1b:

pseudocode for Kruskal's algorithm:

sort edges by ascending order of weight values

$O(E \log E)$

initialize connected components

$O(V)$

while-loop that iterates through edges:

$O(V \log V) \leq O(E \log E)$

if adding the edge would not introduce a cycle:

add vertex to the tree

update connected components

Since $|E| \geq |V| - 1$, the theoretical runtime is $O(E \log E)$. The upper bound is caused by the sorting algorithm.