







Rank









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Maximal AND Subsequences





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Basic knowledge:

- Every number can be represented as a binary number. Largesity of a number depends on its most significant bits. Therefore to make the end result maximum we have to make its most significant bits one.
- And operation on a bit only gives one if all of the numbers in the operation has one on that particular bit.

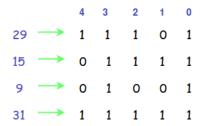
I will explain the problem with an example.

Lets the input is:

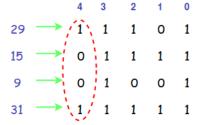
43

29 15 9 31

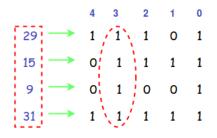
The binary representation of the four numbers are as follows:



We start iterating from the most significant bit which in this case is **4**. On standing bit-**4** we check all the numbers and count how many of them has **1** in this bit.



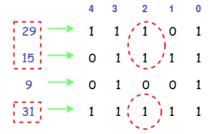
We found two, which is less than k that means and result always has 0 in bit-4 for all possible combination of k numbers. Now move to the next bit.



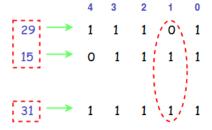
Statistics

Difficulty: Medium Time Complexity: $O(n \cdot log A_i)$ Required Knowledge: Bitwise operation, loop, vector

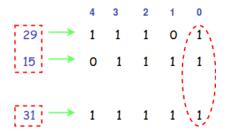
operation, loop, vector Publish Date: Mar 02 2017 Here all the four number has $\bf 1$ in bit-3. Therefore **and** result always has $\bf 1$ in bit- $\bf 3$ for all possible combination of $\bf k$ numbers. Now move to bit- $\bf 2$.



Here 29, 15 and 31 has 1 in bit-2 and 9 has 0. Now if we consider all the combinations we can observe that only one combination (29, 15, 31) has 1 in bit-2 and all other combinations $\{(29, 15, 9), (29, 9, 31), (15, 9, 31)\}$ has 0 in this bit. We can also observe that the combinations have 0 due to the presence of 9 because 9 has 0 in bit-2. So we delete 9, and move to bit 1.



Here we get ${\bf 1}$ only in two numbers which is less than ${\bf k}$ that means ${\bf and}$ result always has ${\bf 0}$ in bit- ${\bf 1}$ for all possible combination of ${\bf k}$ numbers. Now move to the next bit.



In bit-0 we find all the three number has one.

And value of all the remaining numbers will be our final result.

Subsequence:

The question basically wants nCk. That is calculate all the possible ways to choose k numbers from n numbers. Here n is the size of our remaining numbers.



```
Problem Setter's code:

C++

    #include <bits/stdc++.h>
    #include<assert.h>

#define vlong long long
#define mod 1000000007

using namespace std;

inline vlong bigmod ( vlong a, vlong p, vlong m ) {
    vlong res = 1 % m, x = a % m;
    while ( p ) {
        if ( p & 1 ) res = ( res * x ) % m;
            x = ( x * x ) % m; p >>= 1;
    }
}
```

```
return res;
vector<vlong>v;
void solution() {
    int n, k;
    //Taking input
    cin >> n >> k;
    assert(n>1 && n<=1000000);
    assert(k>1 && k<=n);
    for (int i = 0; i < n; i++) {
        vlong j;
        v.push_back(j);
    //Checking bits
    for (int b = 62, len; b >= 0; b--) {
        vector <vlong> nv;
        len = v.size();
        for (int i=0; i<len; i++) {</pre>
            //if i bit is 1
            if (v[i]&(1LL<<b)) {</pre>
                nv.push_back(v[i]);
        //Decreasing number of possible values
        if (nv.size() >= k) {
            v = nv;
    //Calculating and value
    vlong ans = v[0];
    for(int i=1; i<k; i++)</pre>
       ans = (ans & v[i]);
    cout<<ans<<end1;</pre>
    //Claculating nCk
    int len = v.size();
    long long a = 1;
    for(int i=1; i<=len; i++)</pre>
       a = (a*i)\%(mod);
    long long b = 1;
    for(int i=1; i<=k; i++)</pre>
       b = (b*i)\%(mod);
    long long c = 1;
    for(int i=1; i<=(len-k); i++)</pre>
       c = (c*i)\%(mod);
    b = bigmod(b, mod-2, mod);
    c = bigmod(c, mod-2, mod);
    a = (a*b)%mod;
    a = (a*c)%mod;
    cout<<a<<endl;</pre>
    return;
}
int main () {
        solution();
```

```
return 0;
```



Tested by allllekssssa

```
Problem Tester's code:
 #include<bits/stdc++.h>
 using namespace std;
 const int sz=64;
 const int maxi=1e6;
 const long long big=1e18;
 const long long mo=1e9+7;
 int n,k;
 long long st[sz+1],a[maxi];
 long long f[maxi],fi[maxi];
 vector<long long> v[sz+1];
 long long step(long long x, long long y, long long mo)
     if (y==0) return 1;
       long long g=step(x,y/2,mo);
         if (y\%2) return (((g*g)\% mo)*x)\%mo; else return (g*g)\%mo;
 }
 int main()
     scanf("%d%d",&n,&k);
     assert(n>=2 && n<=100000);
     assert(k>=2 && k<=n);
     f[0]=1;
     for (long long i=1;i<=n;i++)
         f[i]=(f[i-1]*i)%mo;
     for (long long i=0;i<=n;i++)</pre>
    fi[i]=step(f[i],mo-2,mo);
     for (int i=0;i<n;i++)
         scanf("%lld",&a[i]);
         assert(a[i]>=0 && a[i]<=big);
         v[sz].push_back(a[i]);
      st[0]=1;
     for (int i=1;i<sz;i++)
         st[i]=st[i-1]*2;
   for (int i=sz-1;i>=0;i--)
       for (int j=0;j<v[i+1].size();j++)</pre>
         if (st[i]&v[i+1][j]) v[i].push_back(v[i+1][j]);
         if (v[i].size()<k)</pre>
         {
             v[i].clear();
             for (int j=0;j<v[i+1].size();j++)</pre>
                 v[i].push_back(v[i+1][j]);
  long long ans=v[0][0];
  int sz=v[0].size();
  long long cnt=(((f[sz]*fi[k])%mo)*fi[sz-k])%mo;
  for (int i=1;i<sz;i++)
    ans&=v[0][i];
    printf("%lld\n",ans);
    printf("%lld\n",cnt);
     return 0;
 }
```

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