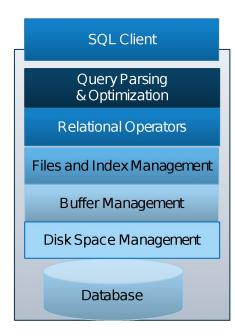
# Logical Database Design: Entity-Relation Models

R&G 2



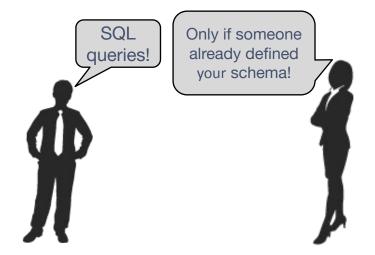
#### Architecture of a DBMS

- Gives us a good sense of how to build a DBMS
- How about using one?



#### Architecture of a DBMS, Pt 2

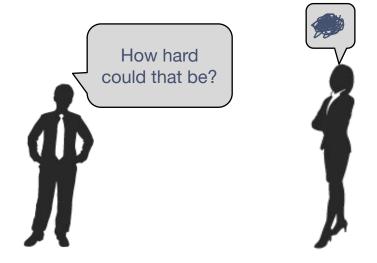
- Gives us a good sense of how to build a DBMS
- How about using one?





#### Architecture of a DBMS, Pt 3

- Gives us a good sense of how to build a DBMS
- How about using one?





### Design of a Database

- Gives us a good sense of how to build a DBMS
- How about using one?
- Today let's talk about how to design a database
  - Not a database system







## Steps in Database Design

#### Requirements Analysis

user needs; what must database do?

#### Conceptual Design

- high level description (often done w/ER model)
- Object-Relational Mappings (ORMs: Hibernate, Rails, Django, etc) encourage you to program here

You are here

#### Logical Design

- translate ER into DBMS data model
- ORMs often require you to help here too

#### Schema Refinement

- consistency, normalization
- Physical Design indexes, disk layout
- Security Design who accesses what, and how

#### Describing Data: Data Models

- <u>Data model</u>: collection of concepts for describing data.
- <u>Schema:</u> description of a particular collection of data, using a given data model.
- Relational model of data
  - Main concept: relation (table), rows and columns
  - Every relation has a schema
    - describes the columns
    - column names and domains

## Levels of Abstraction Users

Views describe how users see the data. View 1 View 2 View 3 Conceptual schema defines logical structure Conceptual Schema Physical Schema Physical schema describes the files and DB indexes used.

### Example: University Database

#### Conceptual schema:

- Students(sid text, name text, login text, age integer, gpa float)
- Courses (cid text, cname text, credits integer)
- Enrolled(sid text, cid text, grade text)

#### Physical schema:

- Relations stored as unordered files.
- Index on first column of Students.

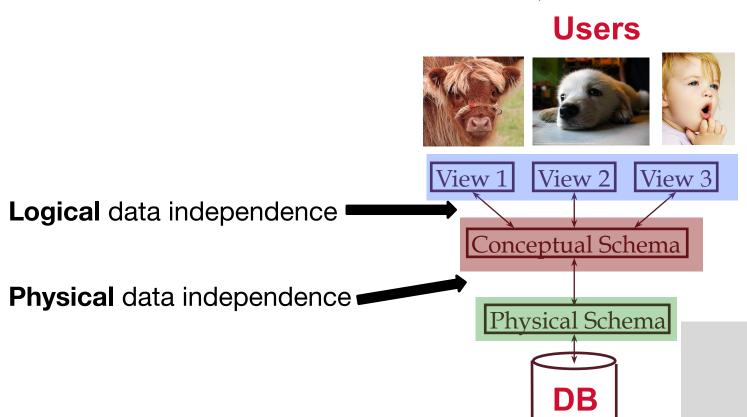
#### External Schema (View):

Course\_info(cid text, enrollment integer)

### Data Independence

- Insulate apps from structure of data
- Logical data independence:
  - Maintain views when logical structure changes
- Physical data independence:
  - Maintain logical structure when physical structure changes

#### Levels of Abstraction, cont



#### Data Independence, cont

- Insulate apps from structure of data
- Logical data independence:
  - Maintain views when logical structure changes
- Physical data independence:
  - Maintain logical structure when physical structure changes

- Q: Why particularly important for DBMS?
  - Because databases and their associated applications persist

#### An Anecdote



## Hellerstein's Inequality

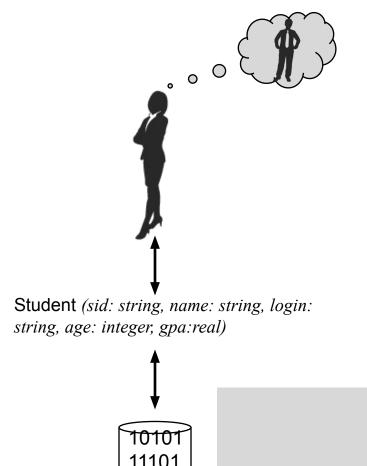
$$\frac{dapp}{dt} << \frac{denv}{dt}$$

Data independence is most important when the rate of change of your environment exceeds the rate of change of your applications.



#### **Data Models**

- Connect concepts to bits!
- Many models exist
- We will ground ourselves in the Relational model
  - clean and common
  - generalization of key/value
- Entity-Relationship model also handy for design
  - Translates down to Relational



## **Entity-Relationship Model**

- Relational model is a great formalism
  - But a bit detailed for design time
  - Too fussy for brainstorming
  - Hard to communicate to "customers"
- Entity-Relationship model: a graph-based model
  - can be viewed as a graph, or a veneer over relations
    - "feels" more flexible, less structured
  - corresponds well to "Object-Relational Mapping"
    - (ORM) SW packages
    - Ruby-on-Rails, Django, Hibernate, Sequelize, etc.

## Steps in Database Design, again

- Requirements Analysis
  - user needs; what must database do?
- Conceptual Design
  - high level description (often done w/ER model)

**ORM** encourages you to program here

You are here

- Logical Design
  - translate ER into DBMS data model
  - ORMs often require you to help here too
- Schema Refinement
  - consistency, normalization
- Physical Design indexes, disk layout
- Security Design who accesses what, and how

## Conceptual Design

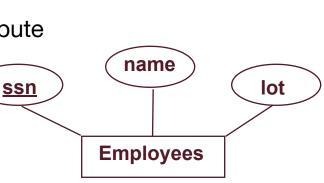
- What are the entities and relationships?
  - And what info about E's & R's should be in DB?
- What integrity constraints ("business rules") hold?
- ER diagram is the "schema"
- Can map an ER diagram into a relational schema.
- Conceptual design is where the data engineering begins
  - If you're familiar with the jargon, these are the "models" of the MVC pattern in ORMs

#### **ER Model Basics: Entities**

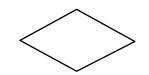
#### Entity:

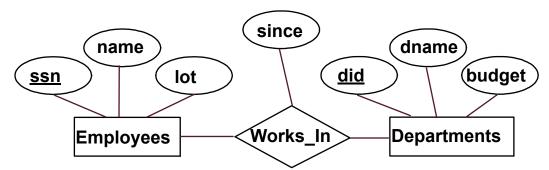
A real-world object described by a set of attribute values.

- Entity Set: A collection of similar entities.
  - E.g., all employees.
  - All entities in an entity set have the same attributes.
  - Each entity set has a key (underlined)
  - Each attribute has a domain



## ER Model Basics: Relationships



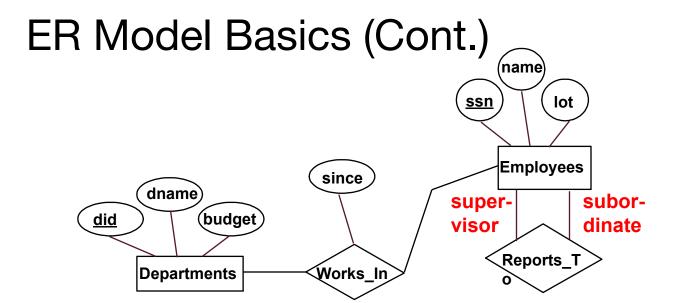


#### **Relationship:** Association among two or more entities.

- E.g., Attishoo works in Pharmacy department.
- Relationships can have their own attributes.

#### **Relationship Set:** Collection of similar relationships.

An n-ary relationship set R relates n entity sets E1 ... En;
 each relationship in R involves entities e1 ∈ E1, ..., en ∈
 En

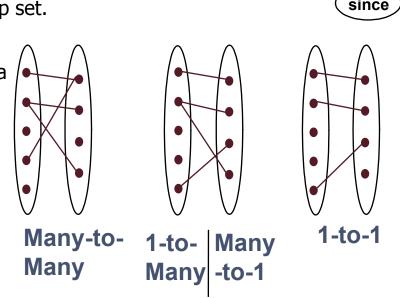


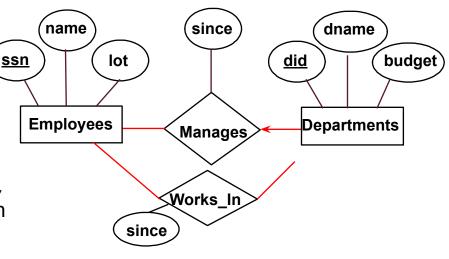
Same entity set can participate in different relationship sets, or in different "roles" in the same relationship set.



## **Key Constraints**

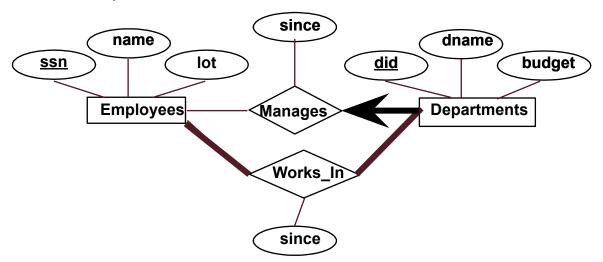
- An employee can work in many departments;
   a dept can have many employees.
- In contrast, each dept has at most one manager, according to the <u>key constraint</u> on **Department** in the **Manages** relationship set.
- A <u>key constraint</u> gives a 1-to-many relationship.





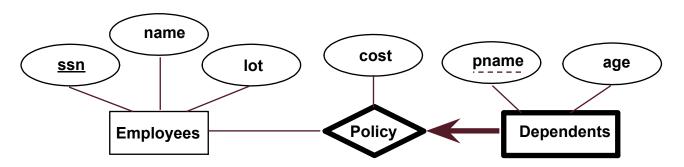
## Participation Constraints

- Does every employee work in a department?
- If so: a participation constraint
  - participation of Employees in Works\_In is total (vs. partial)
  - What if every department has an employee working in it?
- Basically means at least one.



#### Weak Entities

- A weak entity can be identified uniquely only by considering the primary key of another (owner) entity.
  - Owner entity set and weak entity set must participate in a one-to-many relationship set (one owner, many weak entities).
  - Weak entity set must have total participation in this identifying relationship set.



Weak entities have only a "partial key" (dashed underline)

#### **FYI: Crow's Foot Notation**

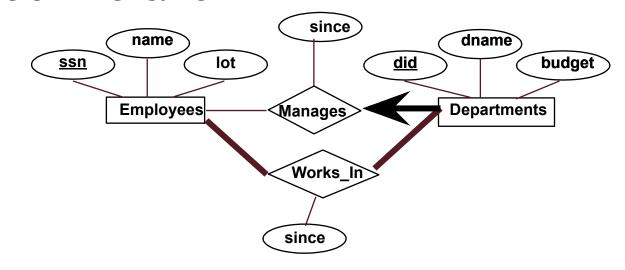
: 0 or more

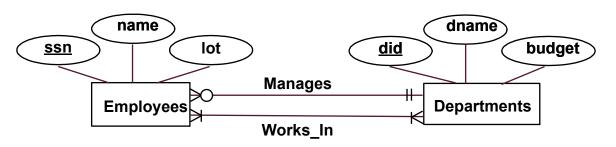
1:1 or more

 $| \cap : 1 \text{ or } 0$ 

: exactly one

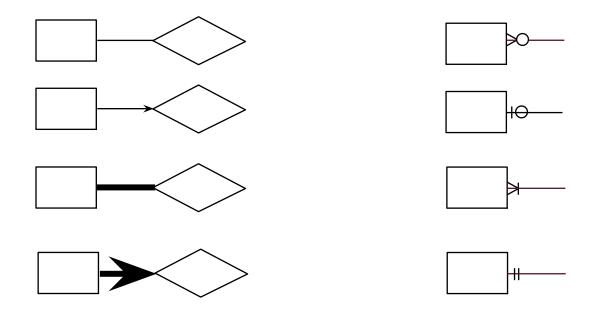
← : many





No relationship attributes

### Translating constraints across notations

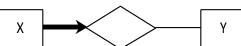


#### Translation to Math Terminology on Relations

• Relation R(X, Y) is a (partial) function



• Relation R(X, Y) is a total function



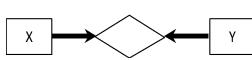
Relation R(X, Y) is surjective (onto)



Relation R(X, Y) is injective (1-1)



Relation R(X, Y) is a bijection



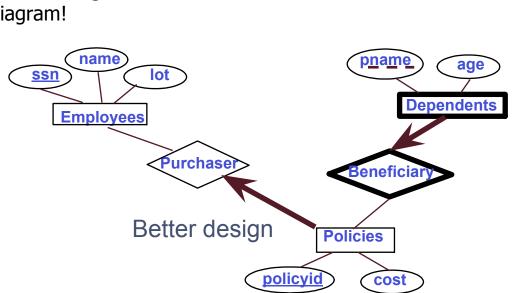


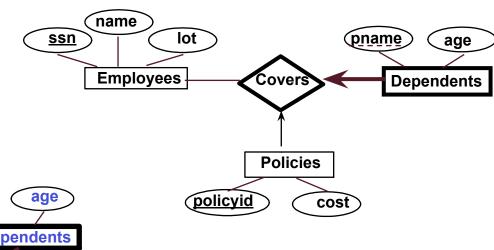
#### Binary vs. Ternary Relationships

If each policy is owned by just 1 employee:



Think through **all** the constraints in the 2nd diagram!

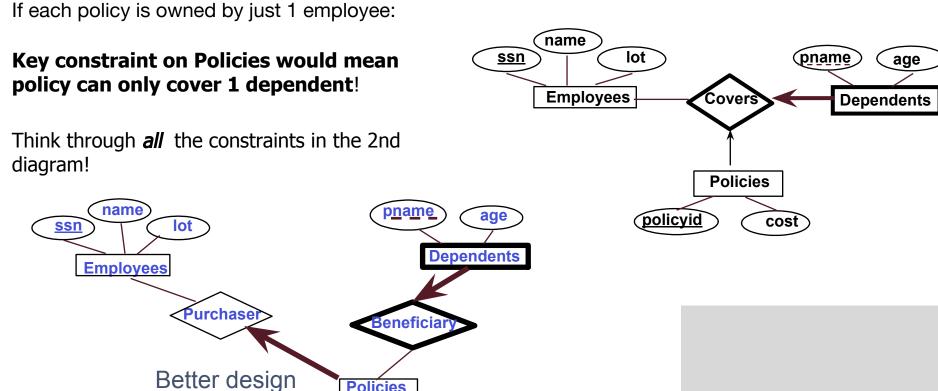




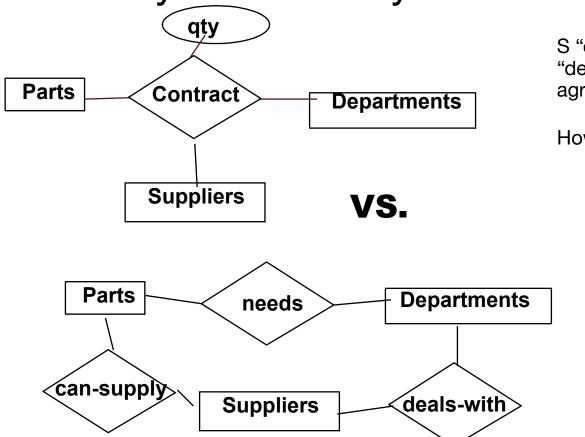
#### Binary vs. Ternary Relationships, cont

**Policies** 

cost



## Binary and Ternary Relationship (cont)

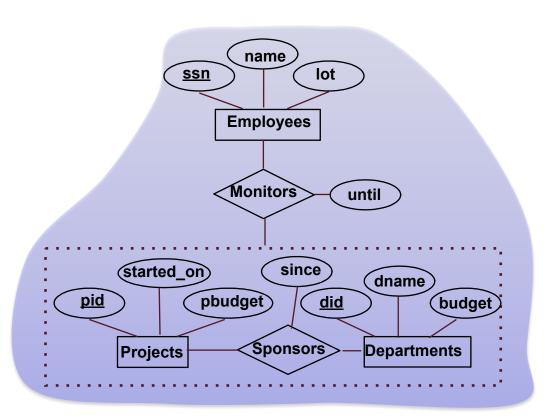


S "can-supply" P, D "needs" P, and D "deals-with" S does not imply that D has agreed to buy P from S.

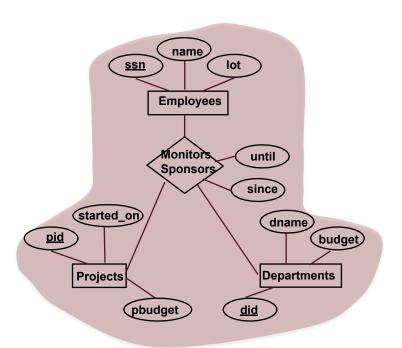
How do we record qty?

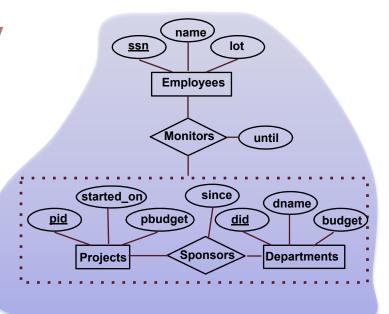
#### Aggregation

Allows relationships to have relationships.



# Aggregation vs. Ternary







#### Conceptual Design Using the ER Model

- ER modeling can get tricky!
- Design choices:
  - Entity or attribute?
  - Entity or relationship?
  - Relationships: Binary or ternary? Aggregation?
- ER Model goals and limitations:
  - Lots of semantics can (and should) be captured.
  - Some constraints cannot be captured in ER.
    - We'll refine things in our logical (relational) design

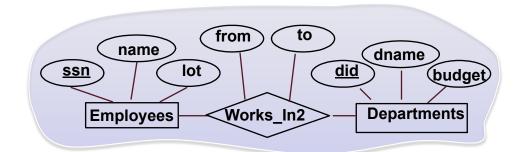
#### Entity vs. Attribute

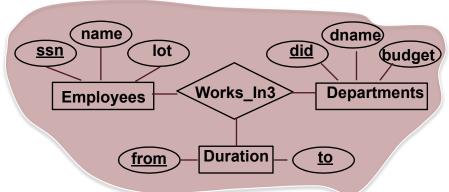
- "Address":
  - attribute of Employees?
  - Entity of its own?
- It depends! Semantics and usage.
  - Several addresses per employee?
    - must be an entity
    - atomic attribute types (no set-valued attributes!)
  - Care about structure? (city, street, etc.)
    - must be an entity!
    - atomic attribute types (no tuple-valued attributes!)

# Entity vs. Attribute (Cont.)

 Works\_In2: employee cannot work in a department for >1 period.

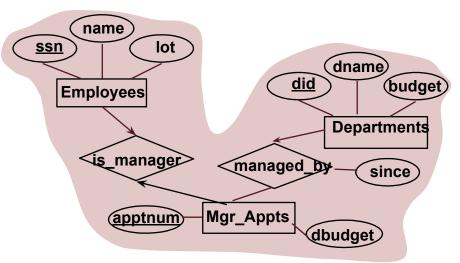
 Like multiple addresses per employee!

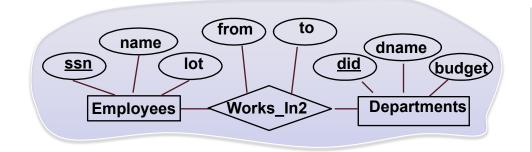




#### Entity vs. Relationship

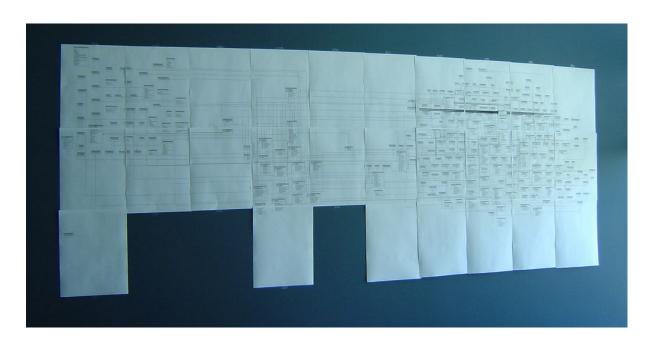
- Separate discretionary budget (dbudget) for each dept.
- What if manager's dbudget covers all managed depts
  - Could repeat value
  - But redundancy = problems
- Better design:





# E-R Diagram as Wallpaper

Very common for them to be wall-sized

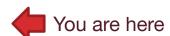




## Steps in Database Design, Part 4

- Requirements Analysis
  - user needs; what must database do?
- Conceptual Design
  - high level description (often done w/ER model)
  - ORM encourages you to program here
- Logical Design
  - translate ER into DBMS data model
  - ORMs often require you to help here too
- Schema Refinement
  - consistency, normalization
- Physical Design indexes, disk layout
- Security Design who accesses what, and how





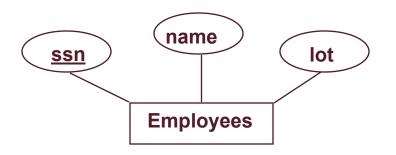
#### Converting ER to Relational

- Fairly analogous structure
- But many simple concepts in ER are subtle to specify in relations

# Logical DB Design: ER to Relational

Entity sets to tables.

Easy.



ssn	name	lot
123-22-3666	Attishoo	48
231-31-5368	Smiley	22
131-24-3650	Smethurst	35

```
CREATE TABLE Employees
  (ssn CHAR(11),
   name CHAR(20),
   lot INTEGER,
   PRIMARY KEY (ssn))
```

#### Relationship Sets to Tables

In translating a **many-to-many** relationship set to a relation, attributes of the relation must include:

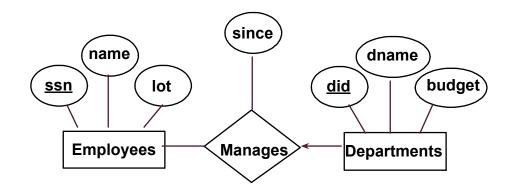
- 1) Keys for each participating entity set (as foreign keys). This set of attributes forms a *superkey* for the relation.
- 2) All descriptive attributes.

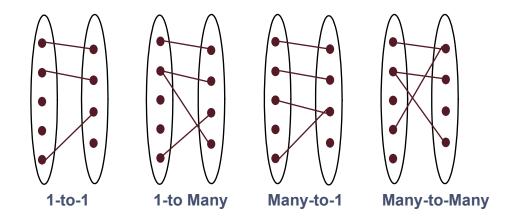
```
CREATE TABLE Works_In(
   ssn CHAR(1),
   did INTEGER,
   since DATE,
   PRIMARY KEY (ssn, did),
   FOREIGN KEY (ssn)
   REFERENCES Employees,
   FOREIGN KEY (did)
   REFERENCES Departments)
```

ssn	did	since
123-22-3666	51	1/1/91
123-22-3666	56	3/3/93
231-31-5368	51	2/2/92

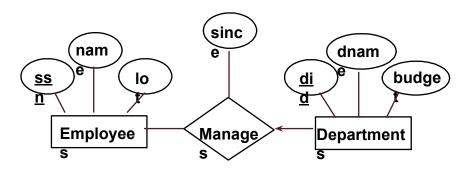
#### Review: Key Constraints

Each dept has at most one manager, according to the **key constraint** on Manages.



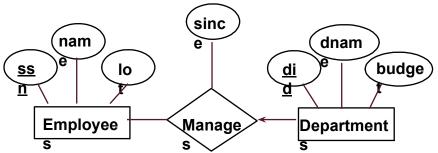


#### Translating ER with Key Constraints



```
CREATE TABLE Manages(
ssn CHAR(11),
did INTEGER,
since DATE,
PRIMARY KEY (did),
FOREIGN KEY (ssn)
REFERENCES Employees,
FOREIGN KEY (did)
REFERENCES Departments)
```

#### Translating ER with Key Constraints, cont



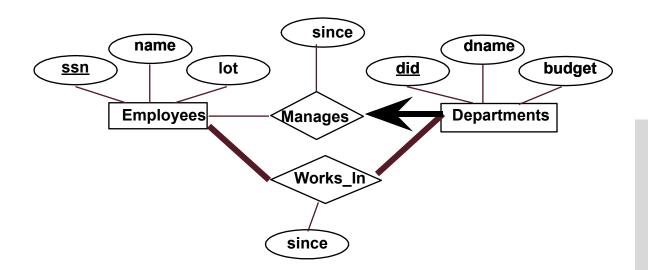
Since each department has a unique manager, we could instead combine Manages and Departments.

```
CREATE TABLE Manages(
                            CREATE TABLE
                                          Dept_Mgr(
 ssn CHAR(11),
                             did
                                  INTEGER,
     INTEGER,
                             dname CHAR(20),
did
 since DATE.
                             budget REAL,
PRIMARY KEY (did),
                             ssn CHAR(11),
 FOREIGN KEY (ssn)
                             since DATE,
   REFERENCES Employees,
                             PRIMARY KEY (did),
 FOREIGN KEY (did)
                             FOREIGN KEY (ssn)
                               REFERENCES Employees)
   REFERENCES Departments)
```



## Review: Key+Participation Constraints

- Every department has one manager.
  - Every did value in Departments table must appear in a row of the Manages table (with a non-null ssn value!)



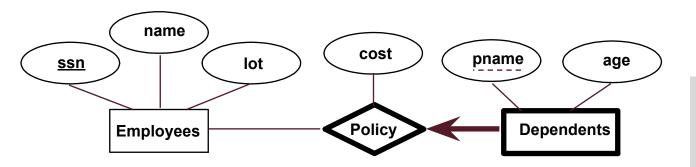
# Participation Constraints in SQL

 We can capture participation constraints involving one entity set in a binary relationship, but little else (without resorting to CHECK constraints which we'll learn later).

```
CREATE TABLE Dept_Mgr(
    did INTEGER,
    dname CHAR(20),
    budget REAL,
    ssn CHAR(11) NOT NULL, -- total participation!
    since DATE,
    PRIMARY KEY (did),
    FOREIGN KEY (ssn) REFERENCES Employees
        ON DELETE NO ACTION)
```

#### Review: Weak Entities

- A weak entity can be identified uniquely only by considering the primary key of another (owner) entity.
  - Owner entity set and weak entity set must participate in a one-to-many relationship set (1 owner, many weak entities).
  - Weak entity set must have total participation in this identifying relationship set.



## Translating Weak Entity Sets

- Weak entity set and identifying relationship set are translated into a single table.
  - When the owner entity is deleted, all owned weak entities must also be deleted.

```
CREATE TABLE Dep_Policy (
   pname CHAR(20),
   age INTEGER,
   cost REAL,
   ssn CHAR(11) NOT NULL,
   PRIMARY KEY (pname, ssn),
   FOREIGN KEY (ssn) REFERENCES Employees
   ON DELETE CASCADE)
```



#### Summary of Conceptual Design

- Conceptual design follows requirements analysis
  - Yields a high-level description of data to be stored
- ER model popular for conceptual design
  - Constructs are expressive, close to the way we think about applications.
  - Note: There are many variations on ER model
    - Both graphically and conceptually
- Basic constructs: entities, relationships, and attributes (of entities and relationships).
- Some additional constructs: weak entities, ISA hierarchies (see text if you're curious), and aggregation.

## Summary of ER (Cont.)

- Basic integrity constraints
  - key constraints
  - participation constraints
- Some foreign key constraints are also implicit in the definition of a relationship set.
- Many other constraints (notably, functional dependencies) cannot be expressed.
- Constraints play an important role in determining the best database design for an enterprise.

#### Summary of ER (Cont....)

- ER design is subjective. Many ways to model a given scenario!
- Analyzing alternatives can be tricky! Common choices include:
  - Entity vs. attribute, entity vs. relationship, binary or n-ary relationship, whether or not to use aggregation
- For good DB design: resulting relational schema should be analyzed and refined further.
  - Functional Dependency information
    - + normalization coming in subsequent lecture.