

TPM Genie

Attacking the Hardware Root of Trust For Less Than \$50

Introduction

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- Focus on hardware and embedded systems security
- Previously: 10 years product security @ BlackBerry & Motorola Mobility

Agenda

1. Overview of Trusted Platform Module
2. Threat Model & Attack Surface
3. Bug Hunting
4. Interposer Design & Build
5. Demo
6. Conclusions

The Trusted Platform Module

What is a TPM?

- A crypto-processor
- Trusted Computing Group (TCG) consortium
- Multiple versions: TPM v1.2 and v2.0
- Used practically everywhere
 - Servers, laptops, desktops, IoT devices, ...
 - Over 2 billion computers use a TPM (allegedly)
 - The U.S. Army and DoD require that every new PC has one



Functions of a TPM

- **Command-Response protocol w/ 100+ ordinals**
- **Hardware random number generation**
- **Secure (aka on-chip) generation of cryptographic keys**
 - ... plus many other crypto primitives
- **Primary use in “Trusted Computing” applications:**
 - Measured Boot
 - Remote Attestation
 - Sealed Storage

TPM Functions – Measured Boot

- **Each boot stage is hashed (measured) by previous stage**
 - BIOS, MBR, UEFI, kernel command line ...
- **Hashes extended into 160-bit Platform Configuration Registers (PCRs)**
 - PCR is a shielded memory space within TPM chip.
 - `PCR[i].new := HASH(PCR[i].old || new_data)`
- **Chain of trust: Code shouldn't be executed until it's been measured.**
 - PCR set can be audited at any point to inspect measurements.
 - Intel Trusted eXecution Technology (TXT)

TPM Functions – Remote Attestation

- Prove to authorized remote party that platform is in a specific state.
- The basic idea:
 - Remote party sends nonce to TPM
 - TPM generates a Quote:
 - $\text{Quote} = \text{sign}(\text{PCR}[n..m] + \text{nonce})$
 - TPM returns Quote to remote party
 - Remote party verifies Quote and decides if it can trust host
 - Next steps are application specific
 - Ex: Hand over some kind of secret, such as DRM key

TPM Functions – Sealed Storage

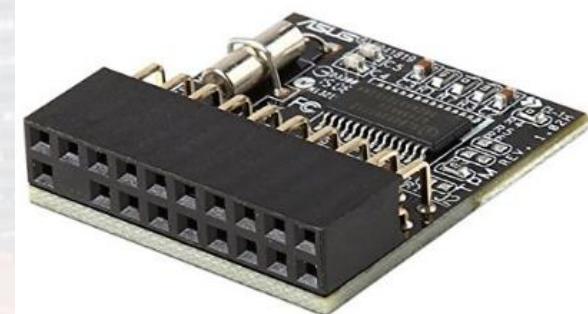
- **Protects a secret stored in the TPM's non-volatile memory**
 - Ex: Bitlocker or dm-crypt keys
- **Binds the secret to a specific PCR set**
 - `cipher_txt = tpm_seal(plain_text, PCRs, [password, locality])`
 - `plain_txt = tpm_unseal(cipher_text, PCRs, [password, locality])`
- **The secret can only be released when the system is in the correct state**

Attack Surface / Threat Model



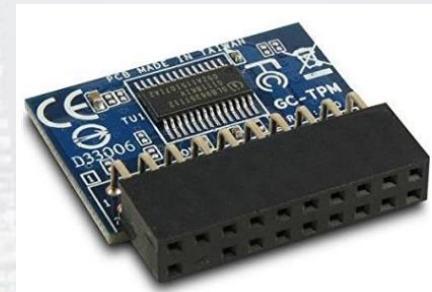
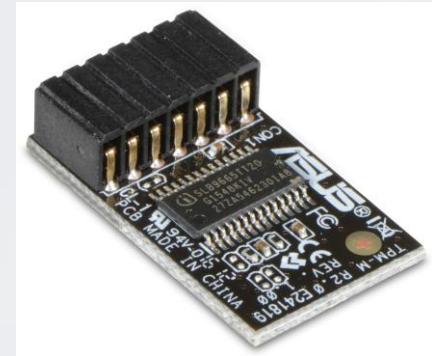
Discrete TPM

- **Manufacturers claim secure type of TPM**
 - Tamper “resistant” against invasive silicon, fault injection and side channel attacks
 - Enter Chris Tarnovsky to prove everyone wrong
- **But... invasive attacks cost \$\$\$**
 - I wanted an attack that works in 5 mins for < \$50
- **Threat Model – Those with physical access**
 - Rogue data center employee
 - Supply chain interdiction attacks (NSA ANT-style implant)
 - Evil Maid scenario

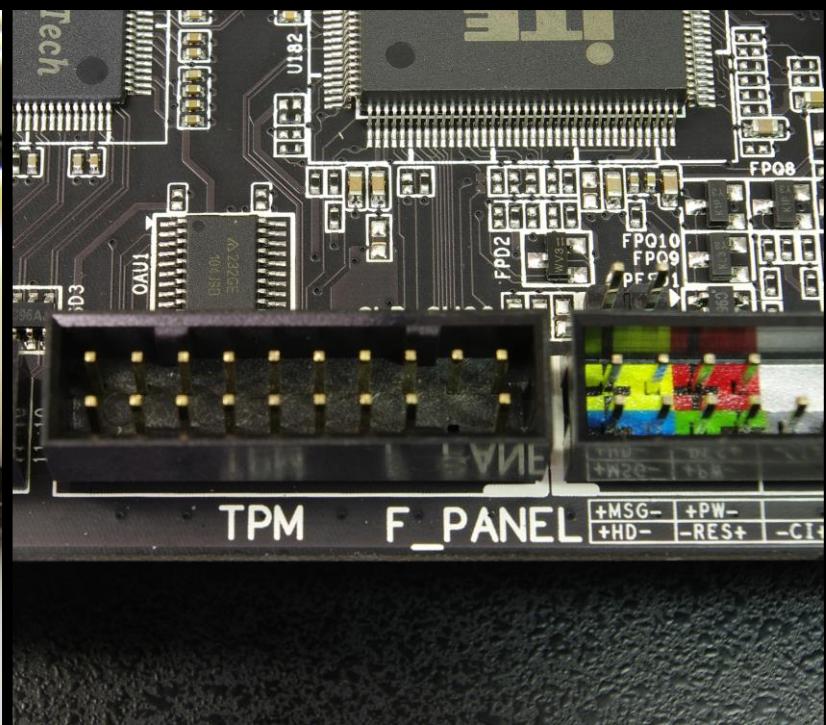
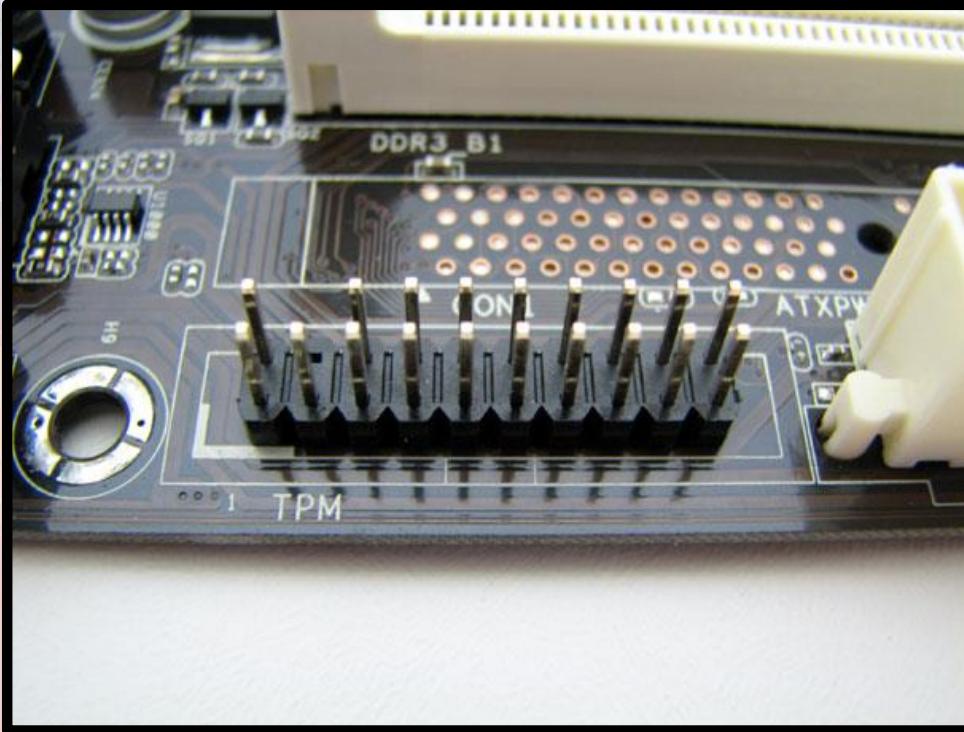


Discrete TPM Risks

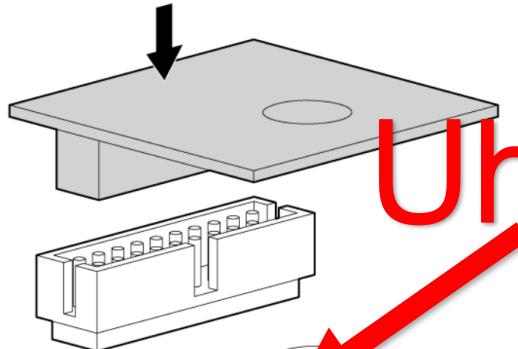
- Often on a daughter card
- Connected to main board via a header
- Communicate with host via serial bus: I2C, SPI, LPC
- Exposes serial bus to tampering
 - A MITM on the bus can sniff/modify traffic
 - Non-invasive attacks via an “interposer” device
 - No need to cut traces or desolder TPM



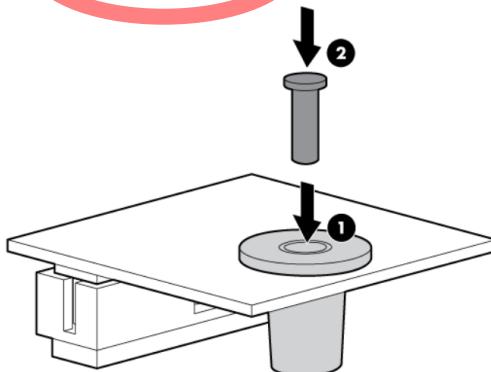
TPM Headers



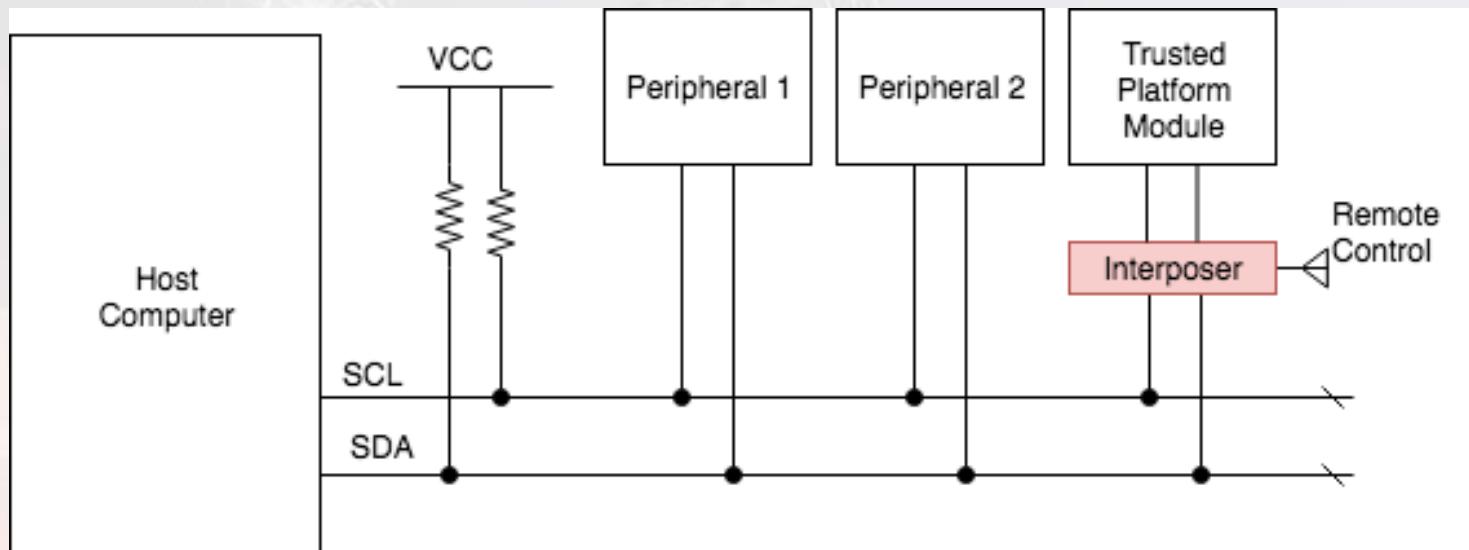
5. Install the TPM board. Press down on the connector to seat the board ("System board components (A-side)" on page 9).



6. Install the TPM security rivet by pressing the rivet firmly into the system board.



Serial Bus Interposer

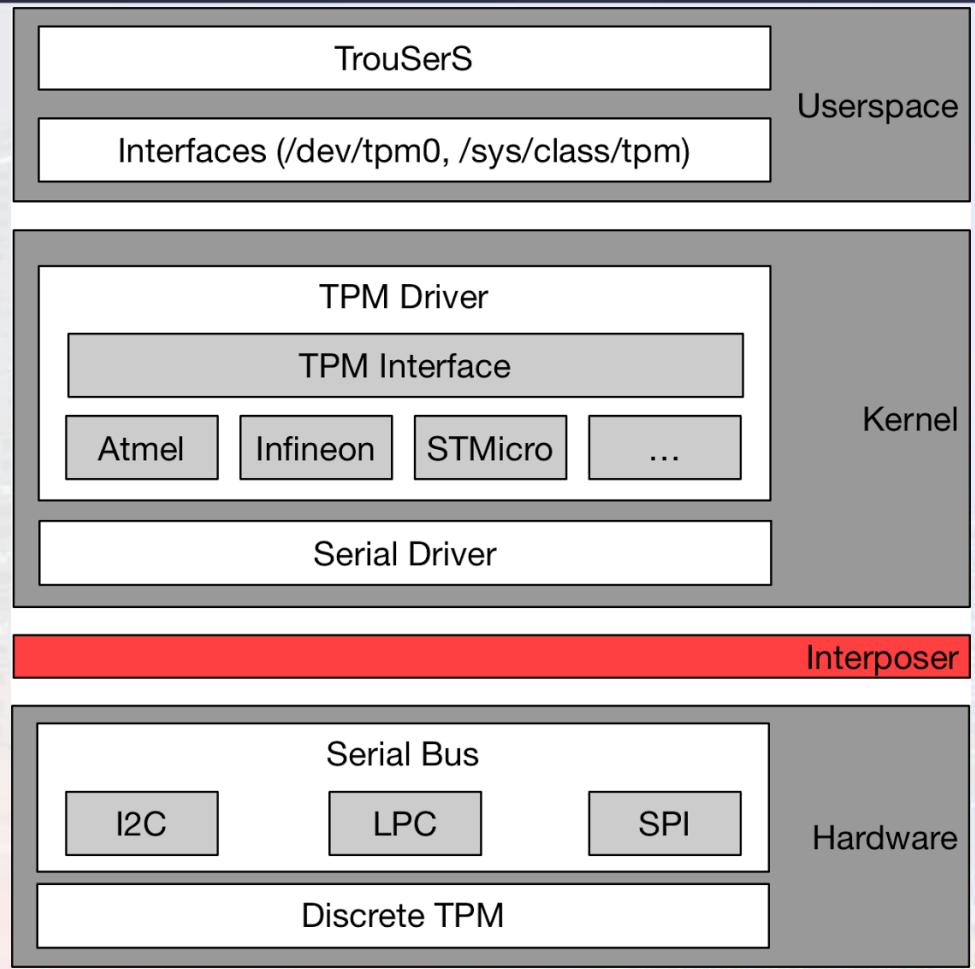


Hunting For Bugs

Target Enumeration & Selection

- **Operating Systems:**
 - Linux kernel: Hardware RNG, integrity subsystem
 - Plus other kernels that implement TPM drivers
- **Pre-kernel environments (Bootloader, Legacy BIOS, UEFI):**
 - tboot, coreboot, GRUB, Tianocore EDK2, +more
- **Userspace stuff:**
 - TrouSerS, OpenSSL TPM Engine, +more

Linux Kernel TPM Driver Architecture



Kernel Code Review

```
int tpm_get_random(u32 chip_num, u8 *out, size_t max) {
    struct tpm_chip *chip;
    struct tpm_cmd_t tpm_cmd;
    u32 recd, num_bytes = min_t(u32, max, TPM_MAX RNG DATA);
    ...
    tpm_cmd.header.in = tpm_getrandom_header;
    tpm_cmd.params.getrandom_in.num_bytes = cpu_to_be32(num_bytes);
    err = tpm_transmit_cmd( chip, &tpm_cmd,
                           TPM_GETRANDOM_RESULT_SIZE + num_bytes );
    ...
    recd = be32_to_cpu(tpm_cmd.params.getrandom_out.rng_data_len);
    memcpy(out, tpm_cmd.params.getrandom_out.rng_data, recd);
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```

The Results

- **Many memory corruption issues:**
 - Linux Kernel: 6
 - U-Boot: 2
 - Coreboot: 1
 - tboot: 13
 - Tianocore EDK2: 10
- **Root cause: Fragile response payload parsing**
- **Some bugs are specific to a TPM chipset vendor (lowest driver layer)**
- **Other bugs affect the generic TPM command/response interface**
 - Both TPM v1.2 and v2.0
 - PCR Read, Get Random, Seal, Unseal, NV Read, Get Capability

At the Packet Level: Request

00	C1	00	00	00	0E	00	00	00	46	00	00	00	10
+	-	-	-	-	+	-	-	-	+	-	-	-	+
tag	length		ordinal			size_req							
+-----+	+-----+		+-----+		+-----+		+-----+		+-----+		+-----+		+
	header								body				

```
tpm_cmd.header.in = tpm_getrandom_header;  
tpm_cmd.params.getrandom_in.num_bytes = cpu_to_be32(num_bytes);  
tpm_transmit_cmd( chip, &tpm_cmd, TPM_GETRANDOM_RESULT_SIZE + num_bytes );
```

At the Packet Level: Response

00	C4	00	00	00	1E	00	00	00	00	00	00	00	00	00	10	EF	F8	2C	6A	...	
+---+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	...	+	
tag	length				ret	code															
+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	...	+
		header																		body	

```
tpm_transmit_cmd(chip, &tpm_cmd, 0x1A);  
recd = be32_to_cpu(tpm_cmd.params.getrandom_out.rng_data_len);  
memcpy(dest, tpm_cmd.params.getrandom_out.rng_data, recd);
```

At the Packet Level: Response

00	C4	00	00	00	1E	00	00	00	00	00	00	00	00	FF	FF	AA	AA	AA	AA	...
+---+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	...	
tag	length			ret	code			data_len							rng_data					
+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	+-----+	...	
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memcpy(dest, tpm_cmd.params.getrandom_out.rng_data, recd);
```

Didn't The TCG Anticipate This?

Authorization Sessions

- **HMAC**
 - Appended to command and response packets.
 - Defense against payload tampering.
 - Guarantees Integrity: Packet hasn't been tampered with.
 - Guarantees Authenticity: By knowledge of *shared secret*.
- **Parameter Encryption**
 - Guarantees Confidentiality: Can transmit secrets on the bus w/o exposure.
- **Rolling Nonces**
 - Prevent replay attacks.

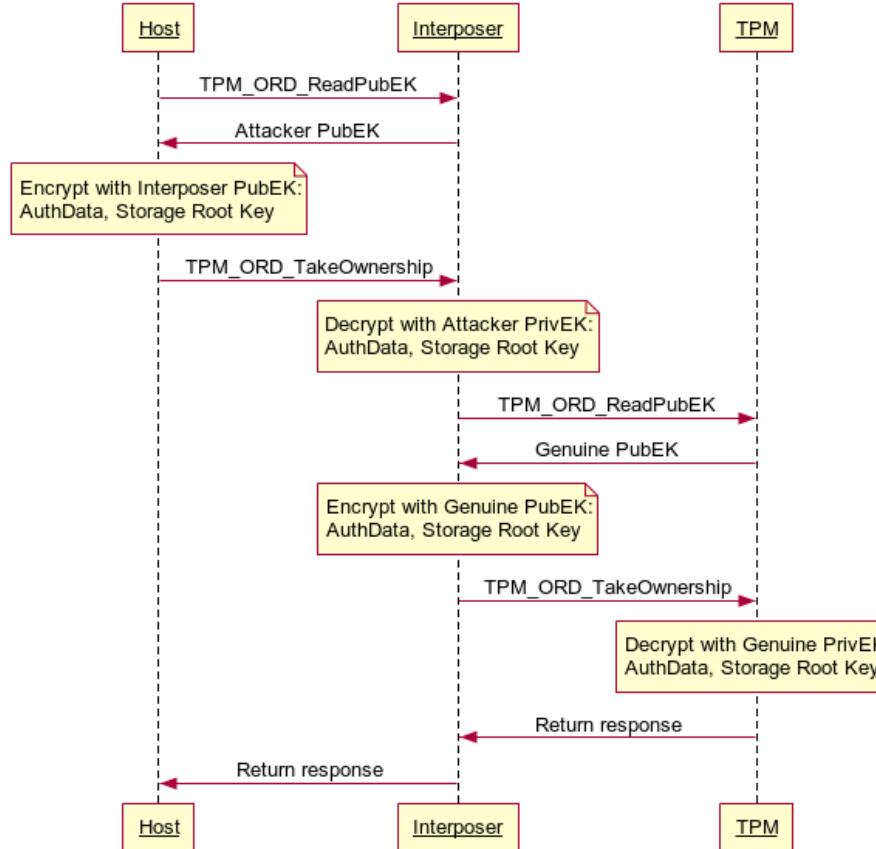
A Few Problems With That...

- **Authorization Sessions:**
 - TCG specification says auth sessions are optional for many commands.
 - Across the board HMAC not applied to critical commands:
 - Ex: TPM_ORD_PcrExtend and TPM_ORD_GetRandom
- **HMAC Verification:**
 - tboot TPM_ORD_Unseal cmd uses HMAC, but resp HMAC is not verified.
- **Weak Nonces:**
 - tboot uses uninitialized stack memory for nonce.
 - Kernel generates nonce with TPM_ORD_GetRandom command.

Endorsement Key

- **How do I know if I'm speaking to a real TPM or an interposer?**
- **Endorsement Key (EK)**
 - Embedded in TPM.
 - Host can request EK using `TPM_ORD_ReadPubEk` cmd.
- **Used by critical operations**
 - How the Auth Session HMAC shared secret is provisioned.
 - Ex: Taking ownership of TPM (`TPM_ORD_TakeOwnership`)
 - Owner encrypts “AuthData” using PubEK.
 - TPM decrypts “AuthData” with PrivEK.
- **Critically Important:**
 - Must verify the EK Certificate to prove identity of TPM.

Interposer Falsifying an Endorsement Key



It Gets Worse...

- We could not find any host-side drivers that actually verify the EK.
- **Other problems:**
 - Not all manufacturers publish their EK Certs.
 - Some Nuvoton TPMs crash when requesting EK Cert from NVRAM.
 - Some TPMs don't ship with EK Cert, or allow buyers to provision their own.
- **Impact:**
 - Interposer provides owner with its PubEK
 - Owner sends AuthData to Interposer, encrypted with attacker PubEK
 - Authorization Sessions are defeated.

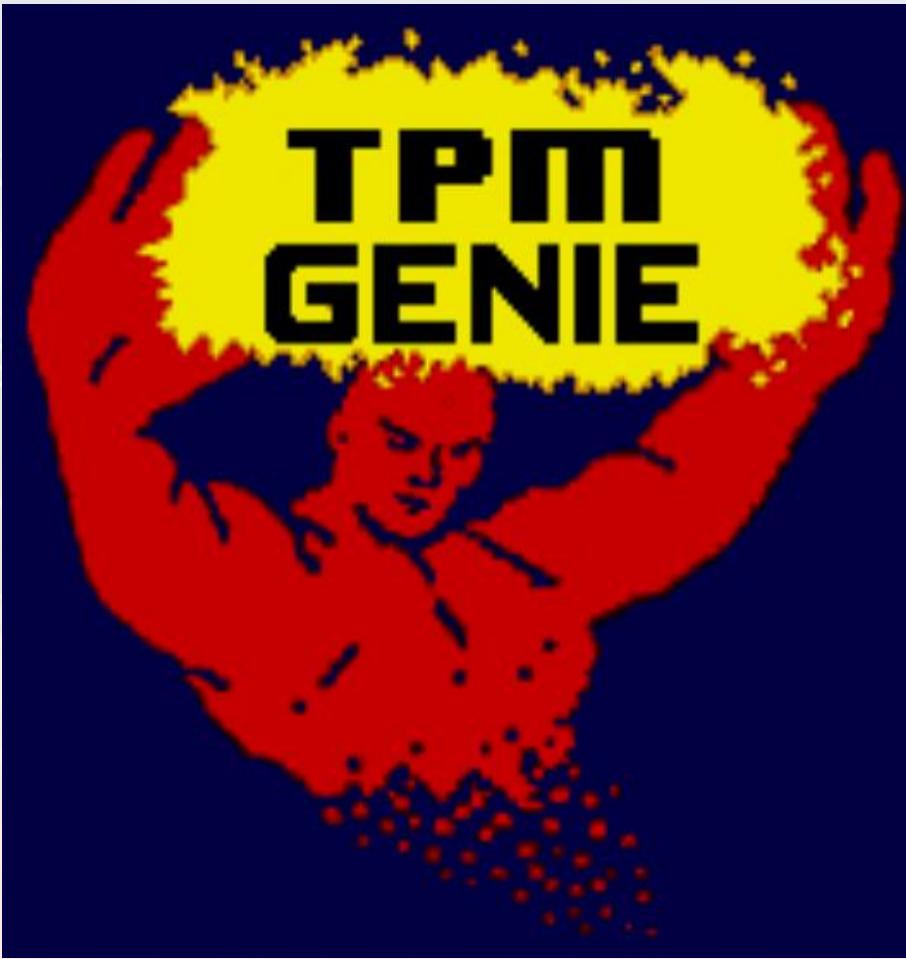
Designing an Interposer



Interposer Design

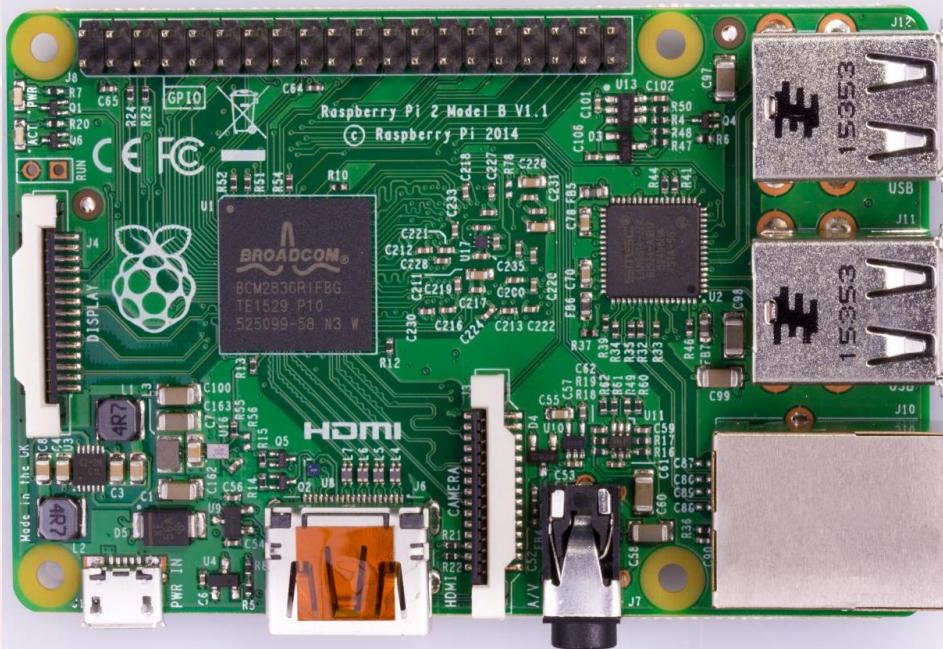
- **Implant in discrete TPM socket**
- **Sit between main board and TPM daughter card**
 - Fool host and TPM into thinking they're talking to each other
- **Allows most traffic to pass through without modification**
- **Modify traffic at opportune time to...**
 - Trigger memory corruption in host response parser.
 - Control PCR Extend operations
 - Control HW RNG output





Components – Victim Device

- Raspberry Pi 2 Model B
 - Why?
 - Runs Linux
 - Easy to set up
 - Inexpensive
 - Has GPIO header simulating discrete socket



Components – TPM

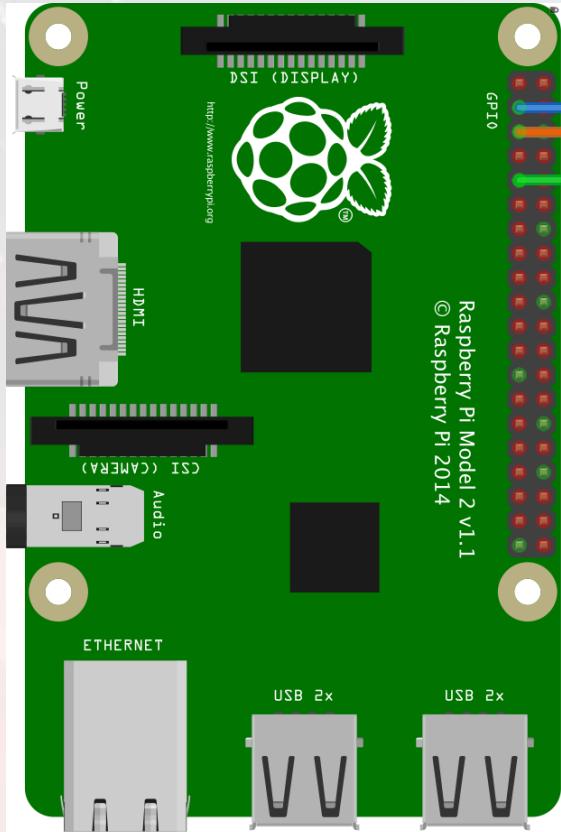
- Infineon TPM45 Iridium Board
 - Based on SLB9645
 - TPM v1.2 specification
- Why?
 - Inexpensive devkit
 - Designed to work with Raspberry Pi
 - I2C protocol simplifies PoC



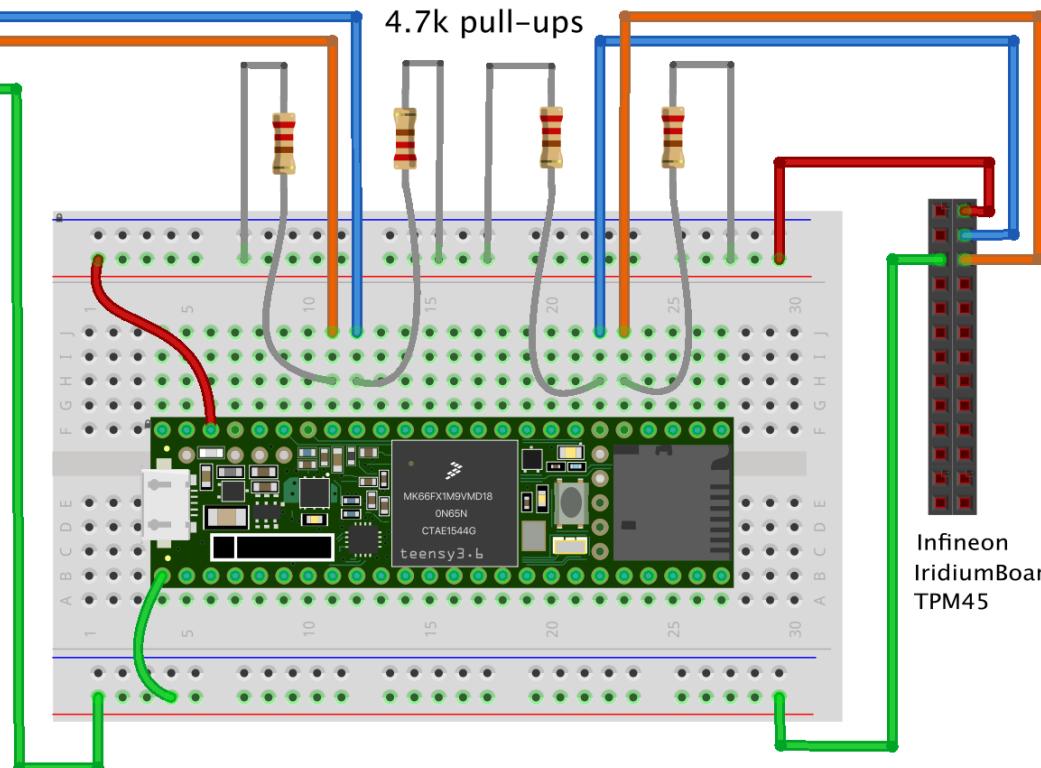
Components – Interposer Device

- **Teensy 3.6 Microcontroller**
- **Why?**
 - Inexpensive MCU
 - Powerful (180 MHz)
 - Arduino simplifies PoC
 - Has multiple I2C busses

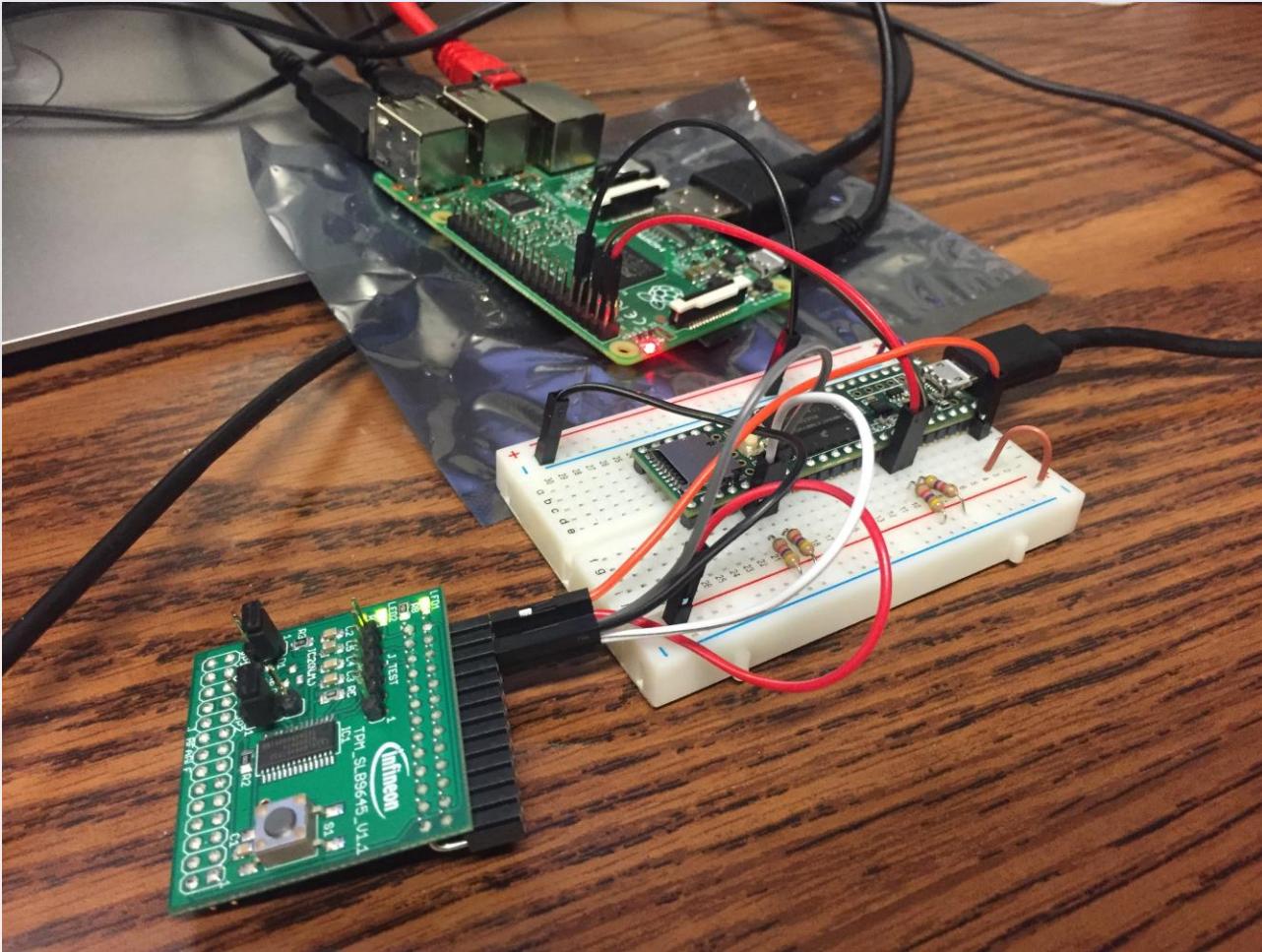




Raspberry Pi Model 2 v1.1
© Raspberry Pi 2014



fritzing



TPM Genie Firmware

- **Arduino based:**
 - About 1200 lines of C.
 - Most code is for pretty printing TPM packets on UART console
 - Use `i2c_t3` library for dual I2C bus support
- **Effort:**
 - 10 days to develop initial PoC
 - 15 days to polish up and make not suck
- **Main Challenges:**
 - Delicate timing constraints
 - Raspberry Pi I2C clock stretching bug
 - SDA/SCL cross-talk on messy bread board

Demos



Impair Hardware RNG

- Interpose TPM_ORD_GetRandom command
- **Impact:**
 - Undermine the Linux kernel hardware RNG
 - Weaken system entropy
 - Impair cryptographic operations on the host

```
# dd if=/dev/hwrng count=1 bs=16 | hexdump -C  
...  
00000000  43 61 6e 53 65 63 57 65  73 74 20 32 30 31 38
```

| CanSecWest 2018 |

Kernel Memory Corruption

- Kernel's TPM_ORD_GetRandom parser is fragile
- **Impact:**
 - This ordinal is used by hw_rng and integrity subsystems.
 - Compromise kernel by returning malformed response packet

```
# dd if=/dev/hwrng count=1 bs=16
<crash>
[65.650713] Unable to handle kernel paging request at virtual address c80255aa
<crash>
```

PCR Spoofing

- Interpose TPM_ORD_PcrExtend command ordinal
- **Impact:**
 - Remote Attestation: Fool trusted remote party into believing that a tampered OS has actually booted into an expected state
 - Sealed Storage: Allow secrets to be unsealed even if the firmware has been tampered with.

```
# python3 pcr_extend.py -i 0 -f data_good
# python3 pcr_extend.py -i 0 -f data_bad
# python3 pcr_read.py -i 0,1

PCR 00: d6 eb c4 e0 4e 16 12 a1 ae 46 5c 51 c0 90 60 8b c5 e6 e1 74
PCR 01: d6 eb c4 e0 4e 16 12 a1 ae 46 5c 51 c0 90 60 8b c5 e6 e1 74
```

Conclusions

Conclusions

- **Across the board TPM drivers are buggy**
 - Poor validation of variable-length response structures
- **Design does not defend against local attack vectors**
 - Interposer controls everything on the bus:
 - Register reads/writes, locality, all cmd/resp packet data
 - TCG specification doesn't require Session HMAC for many critical commands.
 - Cannot easily verify the identity of your TPM hardware.
- **Using a TPM can increase your system's attack surface**

Patch Availability

- **Linux Kernel:**
 - Very quick response!
 - Patches for memory corruption bugs now available in v4.16.
 - Plan to backport patches to stable branches: 3.18, 4.4, 4.9, 4.14, and 4.15
 - Unfortunately EK verification and mandatory Auth Session usage is not ready.
- **Coreboot**
 - Unfixed. Report can be found in public bug trackers.
- **U-Boot:**
 - Fixed in master branch.
- **Tianocore EDK2 / tboot**
 - Intel PSIRT has stated that they do not plan on patching the reported vulnerabilities.

Advice For TCG and TPM Manufacturers

- **Trusted Computing Group**
 - Make Authorization Sessions mandatory for all commands.
 - Publish clear guidance on the need to verify the EK Certificate.
 - Work closely with those that implement TPM drivers to ensure they follow best practices.
- **TPM Chip Manufacturers:**
 - Every TPM must have an Endorsement Certificate.
 - ... that is published online so the drivers can verify it.

Advice for Platform Engineers

- **For now, avoid using discrete TPMs on daughter cards.**
- **Slightly better: Make serial bus access difficult.**
 - Soldered onto mainboard.
 - Cover with a RF shield, under-fill with epoxy resin.
 - Bury serial bus in inner layer of PCB. Difficult because of pull up resistors.
- **Even better: Make bus access much more difficult.**
 - Integrate TPM within the processor.
 - Serial bus is never exposed externally.

Next Steps / Future Work

- **Improve PoC**
 - Add remote access capability for secret exfiltration
- **Improve hardware**
 - Add support for TPM v2.0, SPI and LPC serial buses
 - Miniaturize, or disguise as legitimate TPM daughter card.
- **Audit more drivers**
 - We only scratched the surface.

Questions?

<https://github.com/nccgroup/TPMGenie>

Thanks: Rob Wood, Sultan Qasim Khan,
Ian Robertson, Jason Meltzer, Jon Szymaniak