# Cryptographic Secret Sharing

Girls Talk Math

#### Introduction

In this problem set, you will learn about ...

One last note about reading mathematical texts: it is very normal when reading math to read a passage or even a single sentence several times before understanding it properly. Also, never trust the author! Check every claim and calculation (time permitting). Take your time and never give up. Let's talk math!

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# 1 Probability and Randomness

#### 1.1 Introduction

#### 1.2 Randomness in Cryptography

(NG: Define information-theoretic security. Define what a security parameter is  $(\lambda)$ .)

#### 1.3 Sharing Secrets

Secret sharing is a way to "split" a secret value (call it s) into pieces. When we distribute those pieces among a large set of people, and some subset (or maybe the whole set) can recover the secret only if they pool their information (their shares).

(NG: Define the privacy property and the cryptographic privacy game? The goal can be to have a (bonus) exercise in which they write a real cryptographic proof of privacy of a secret sharing scheme (likely the XOR scheme).)

# 2 A simple secret sharing

## 2.1 Binary Arithmetic

The exclusive-OR (XOR) operation is denoted by the symbol  $\oplus$  and defined by the following truth table:

a	b	$a \oplus b$
0	0	0
0	1	1
1	0	1
1	1	0

(NG: explain binary representations of numbers and binary arithmetic.)

## 2.2 Sharing Secrets using XOR

$$\frac{\mathsf{Share}(s)}{s_1 \leftarrow_{\$} \{1, \dots, 2^{\lambda}\}}$$
$$\mathsf{return} \ (s_1, s \oplus s_2)$$

$$\frac{\mathsf{Reconstruct}(s_1,s_2)}{\mathsf{return}\ s_1 \oplus s_2}$$

### 2.2.1 Proving Security\*

(NG: Optional section. Introduce the privacy game, have them act as the adversary and try to beat the privacy and get some intuition about why it's hard, then work through a simple proof.)

## 3 Shamir's Secret Sharing

### 3.1 Polynomials

(NG: y-intercept, zeroes)

#### 3.1.1 Uniqueness

(NG: How many points uniquely define a polynomial)

#### 3.1.2 Lagrange Interpolation

**Exercise 3.1** (NG: Practice some manual Lagrange interpolation. Pick a polynomial, evaluate it at 3 points, then use those 3 points in Lagrance Interpolation and recover the polynomial.)

#### 3.2 Sharing Secrets Using Polynomials

Shamir secret sharing is a (t+1)-out-of-n secret sharing protocol, for some numbers t and n. This means that we split the secret s into n values and distribute them to n people. Then, at least t+1 of those people must work together to recover s.

(NG: Don't introduce finite fields but maybe make a note that this should be done over finite fields.)

(NG: Use the Jupyter Notebook to play around with this.)