1 Tableaux for classical propositional logic

Initial list for $A_1,...A_n \models B$ is with $\neg B$. Closure condition: A and $\neg A$ for some formula A occur on a branch.

2 Semantic tableaux for FDE

Initial list for $A_1, ... A_n \models B$ is with all $A_n, +$ and B, -. Closure condition: A, + and A, - for some formula A occur on a branch.

formula
$$A$$
 occur on a branch. $A \wedge B, +$ $A \wedge B, A \vee B, +$ $A \vee B, A \vee B, A \vee B, +$ $A \vee B, A \vee B, -$

3 Semantic tableaux for K_3 , LP, L_3 , RM_3

3.1 K_3

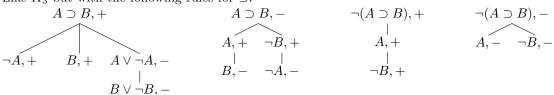
Same as for FDE with extra closure condition: A branch also closes if it contains A, + and $\neg A$, + for some formula A.

3.2 *LP*

Same as for FDE with extra closure condition: A branch also closes if it contains A, — and $\neg A$, — for some formula A.

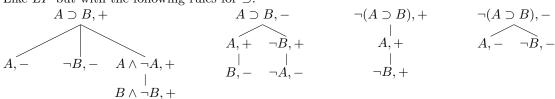
3.3 L_3

Like K_3 but with the following rules for \supset :



3.4 RM_3

Like LP but with the following rules for \supset :



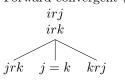
4 Tableaux for modal logic

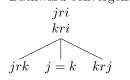
Initial list for $A_1,...A_n \models B$ is with all $A_n,0$ and $\neg B,0$. Closure condition: A,i and $\neg A,i$ for some formula A occur on a branch with the same number i.

4.1 Restrictions

Reflexive (ρ) : Symmetric (σ) : Extendable (η) : irjirijriirjTransitive (τ) : Universal (v): Dense (δ) : $\Box A, i$ irjirj $\Diamond A, i$ jrkirkA, jirkkrjExtra rules for =: Forward convergent (φ) : Backward convergent (β) : $\alpha(i)$ $\alpha(i)$ irjjri

$\begin{array}{ccc} \alpha(i) & \alpha(i) \\ i = j & j = i \\ \beta(j) & \alpha(j) \end{array}$



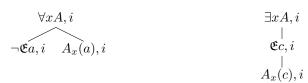


4.2 Tense logic

5 First-order modal logic

5.1 *CK*

5.2 VK



a is old, if possible; otherwise introduce a new a, c is a new variable