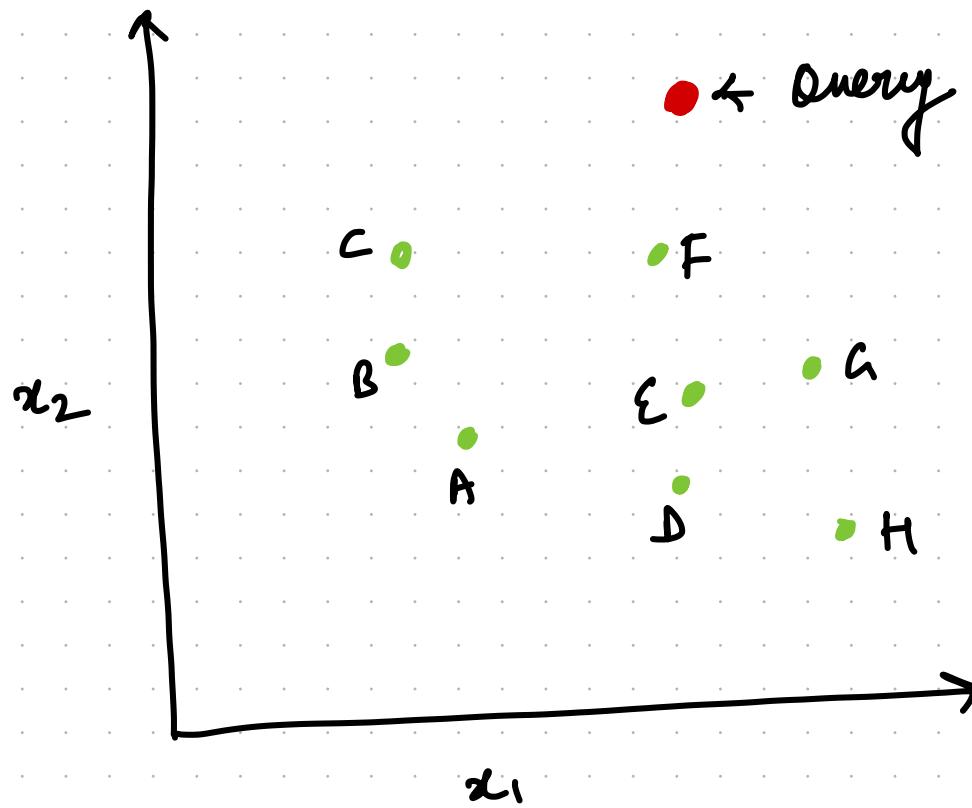
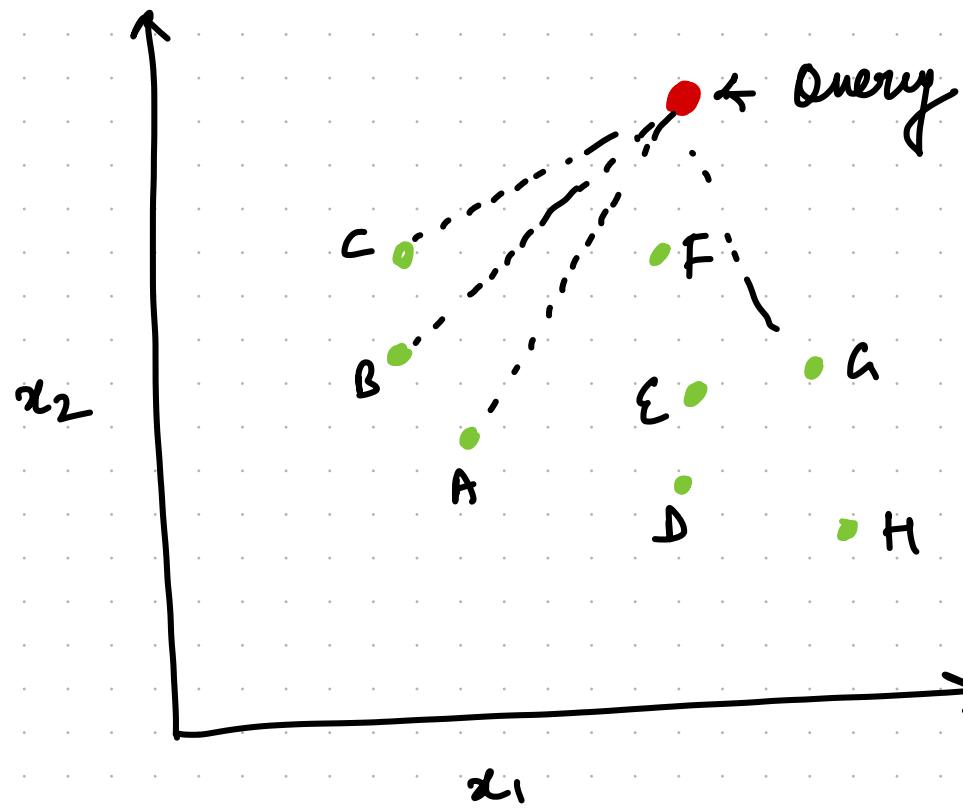


Find 1-NN for query point  $\vec{q} \in \mathbb{R}^D$

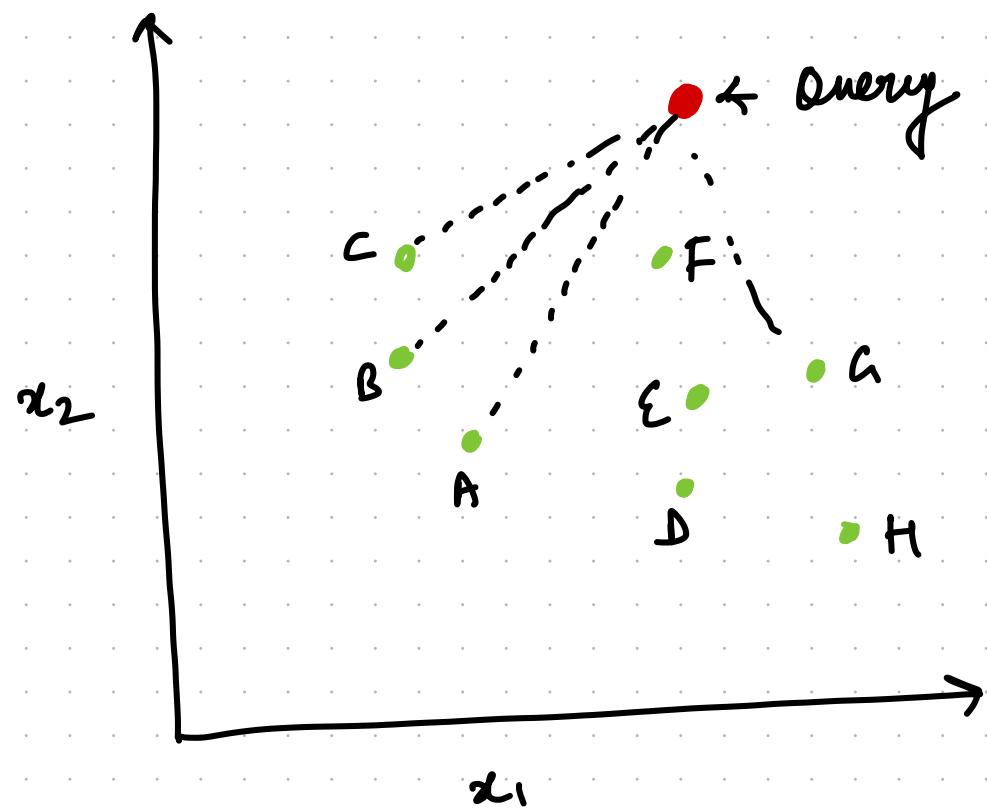


Find 1-NN for query point  $\vec{q} \in \mathbb{R}^D$   
 TRAIN SET is  $X \in \mathbb{R}^{N \times D}$



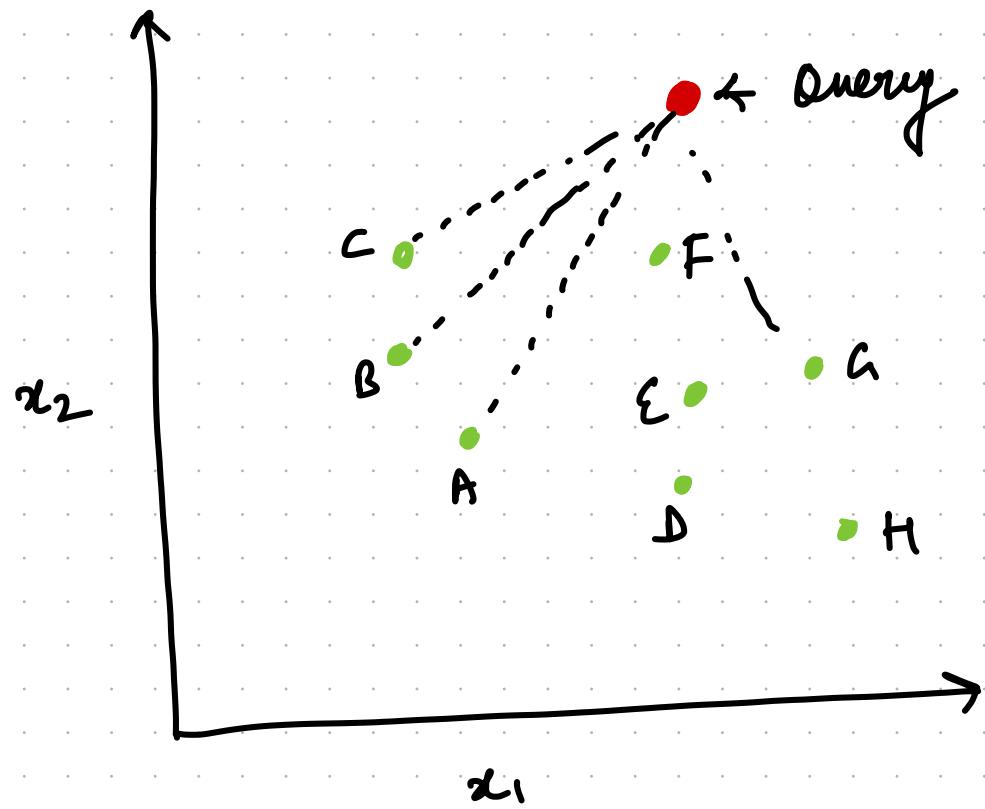
DISTANCE	..	..	..	..	..
$q - A$	..	..	..	..	..
$q - B$	..	..	..	..	..
..	..	..	..	..	..
..	..	..	..	..	..
..	..	..	..	..	..

Find 1-NN for query point  $\vec{q} \in \mathbb{R}^D$   
 TRAIN SET is  $X \in \mathbb{R}^{N \times D}$



DISTANCE	
$q_V - A$	..
$q_V - B$	..
	..
	..
	..
	..
	..
	..

STEPS FOR FINDING  $D(q_V, A)$



DISTANCE				
$q_V - A$	..	..	..	..
$q_V - B$	..	..	..	..
$q_V - C$	..	..	..	..
$q_V - D$	..	..	..	..
$q_V - E$	..	..	..	..
$q_V - F$	..	..	..	..
$q_V - G$	..	..	..	..
$q_V - H$	..	..	..	..

STEPS FOR FINDING  $D(q_V, A)$

$$D(q_V, A) = \sqrt{(q_{V1} - A_1)^2 + \dots + (q_{VD} - A_D)^2}$$

STEPS FOR FINDING  $D(q, A)$

$$D(q, A) = \sqrt{(q_1 - A_1)^2 + \dots + (q_D - A_D)^2}$$

# SUBTRACTIONS

# MULTIPLICATIONS

# ADDITIONS

# SQRT

STEPS FOR FINDING  $D(q, A)$

$$D(q, A) = \sqrt{(q_1 - A_1)^2 + \dots + (q_D - A_D)^2}$$

# SUBTRACTIONS D

# MULTIPLICATIONS D

# ADDITIONS D

# SQRT 1

STEPS FOR FINDING  $D(q, A)$

$$D(q, A) = \sqrt{(q_1 - A_1)^2 + \dots + (q_D - A_D)^2}$$

# SUBTRACTIONS D

# MULTIPLICATIONS D

# ADDITIONS D

# SQRT 1

Total time for  $D(q, A)$  or  
=  $O(D)$   
DISTANCE B/W 1 PAIR

Q) Time required for creating table

DISTANCE

Q) Time required for  
creating table

$O(N \cdot D)$

↑      ↗  
# samples      Dimensionality

DISTANCE	..	..	..	..	..	..
q-A	..	..	..	..	..	..
q-B	..	..	..	..	..	..
..	..	..	..	..	..	..
..	..	..	..	..	..	..
..	..	..	..	..	..	..
..	..	..	..	..	..	..

Q) Time required for  
finding 1-NN?

DISTANCE

q-A	20
q-B	30
.	2
.	8
.	9
.	..
.	-
	34

Q) Time required for  
finding 1-NN?

$O(n)$  : linear search

DISTANCE

	DISTANCE
q-A	20
q-B	30
.	2
.	8
9	.
.	..
.	-
	34

Q) Overall time complexity

$O(ND)$



Linear in # samples

(sometimes  
millions of  
samples)

DISTANCE

	DISTANCE
q-A	20
q-B	30
.	2
.	8
.	9
:	..
	-
	34

Goal :

Reduce  $O(N^D)$   
↑  
Target

How ?

Goal:

Reduce  $O(N^D)$   
↑  
Target

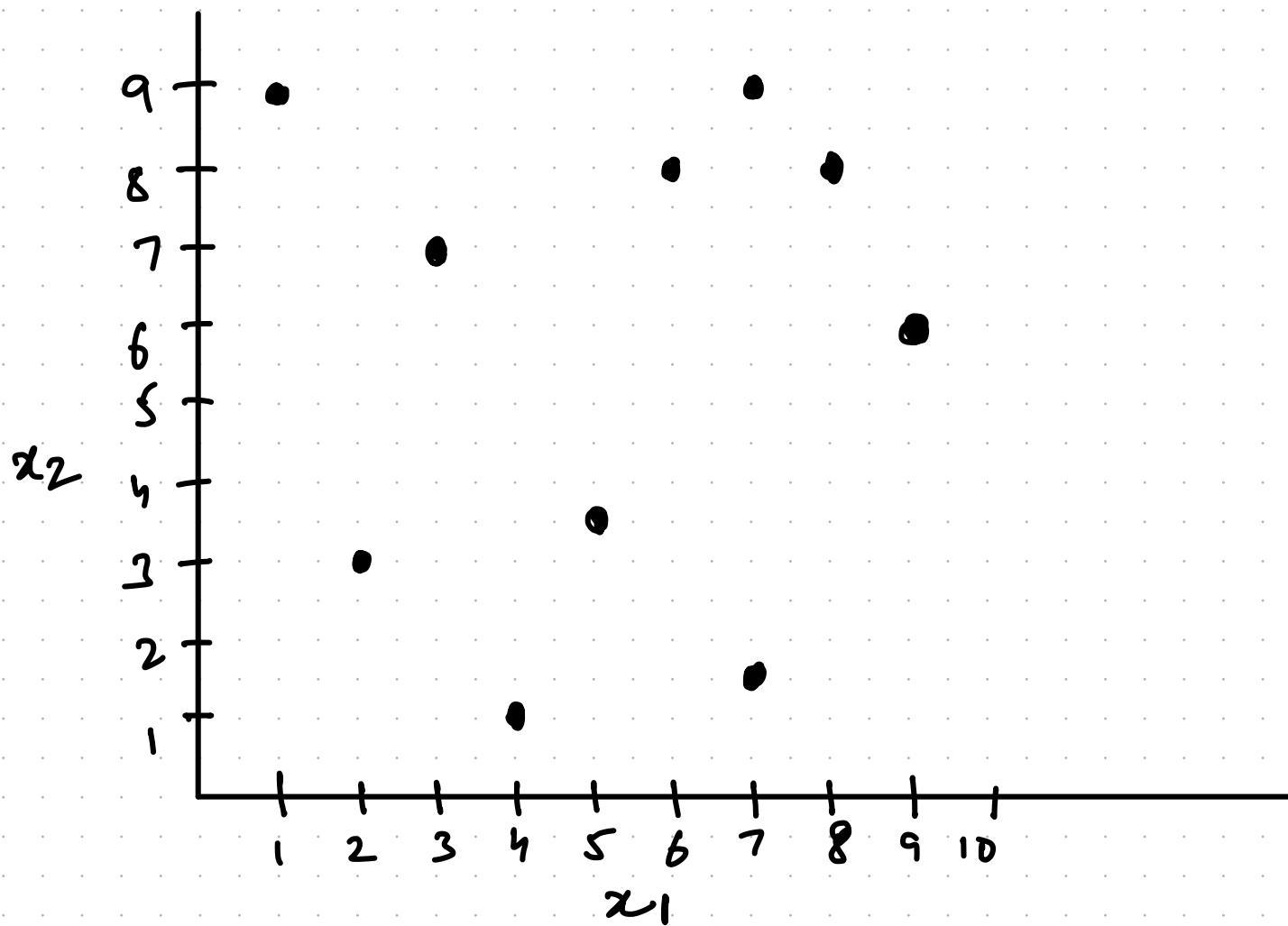
How? (Hints)

- Decision trees
- Search for subset of examples
- Current algorithm does nothing at training time.

## K-D trees

(Victor Lamarcq slides)

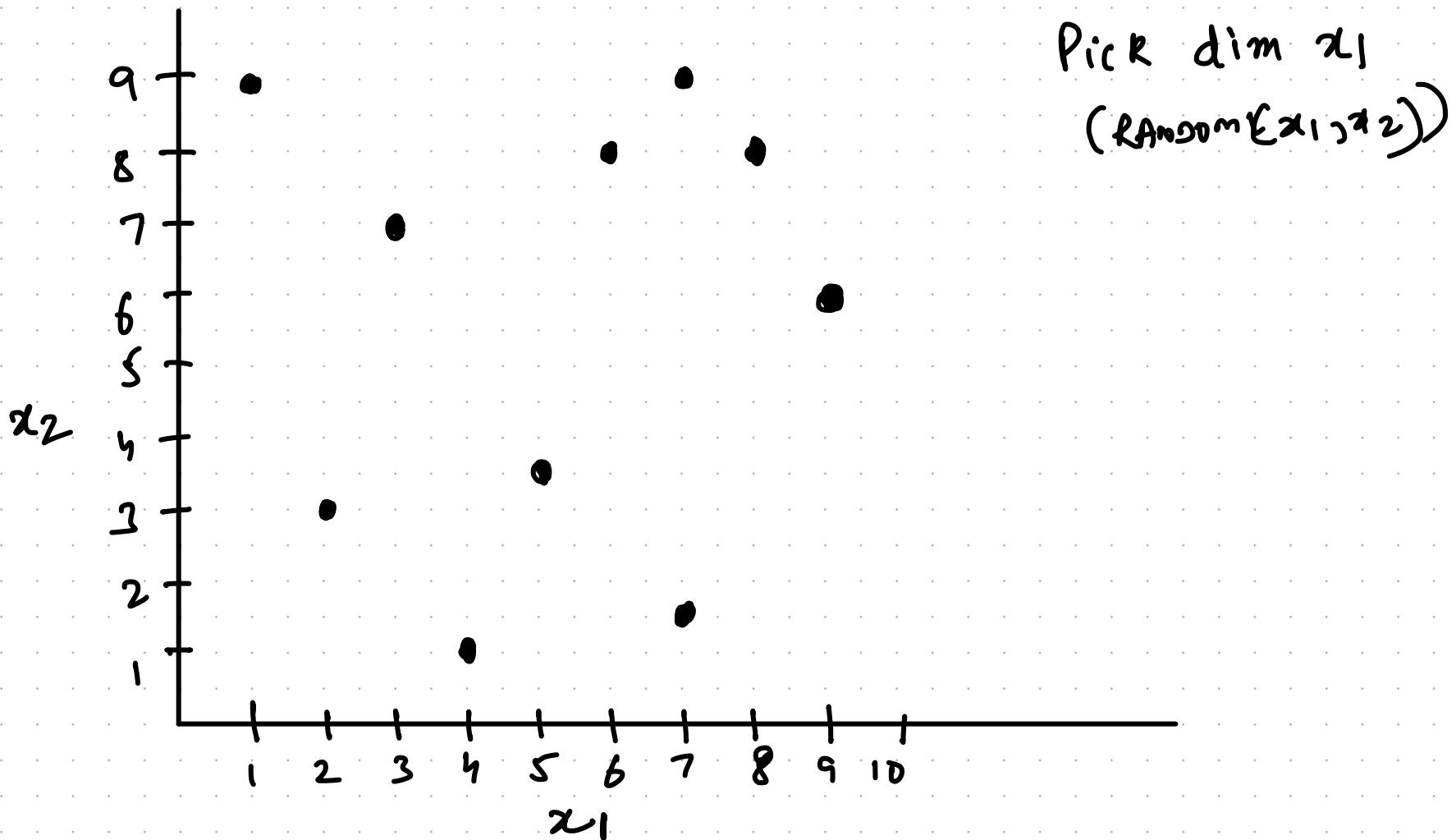
- $X = \{(1, 9), (2, 3), (4, 1), (3, 7), (5, 4), (6, 8), (7, 2), (8, 8), (7, 9), (9, 6)\}$



## K-D trees

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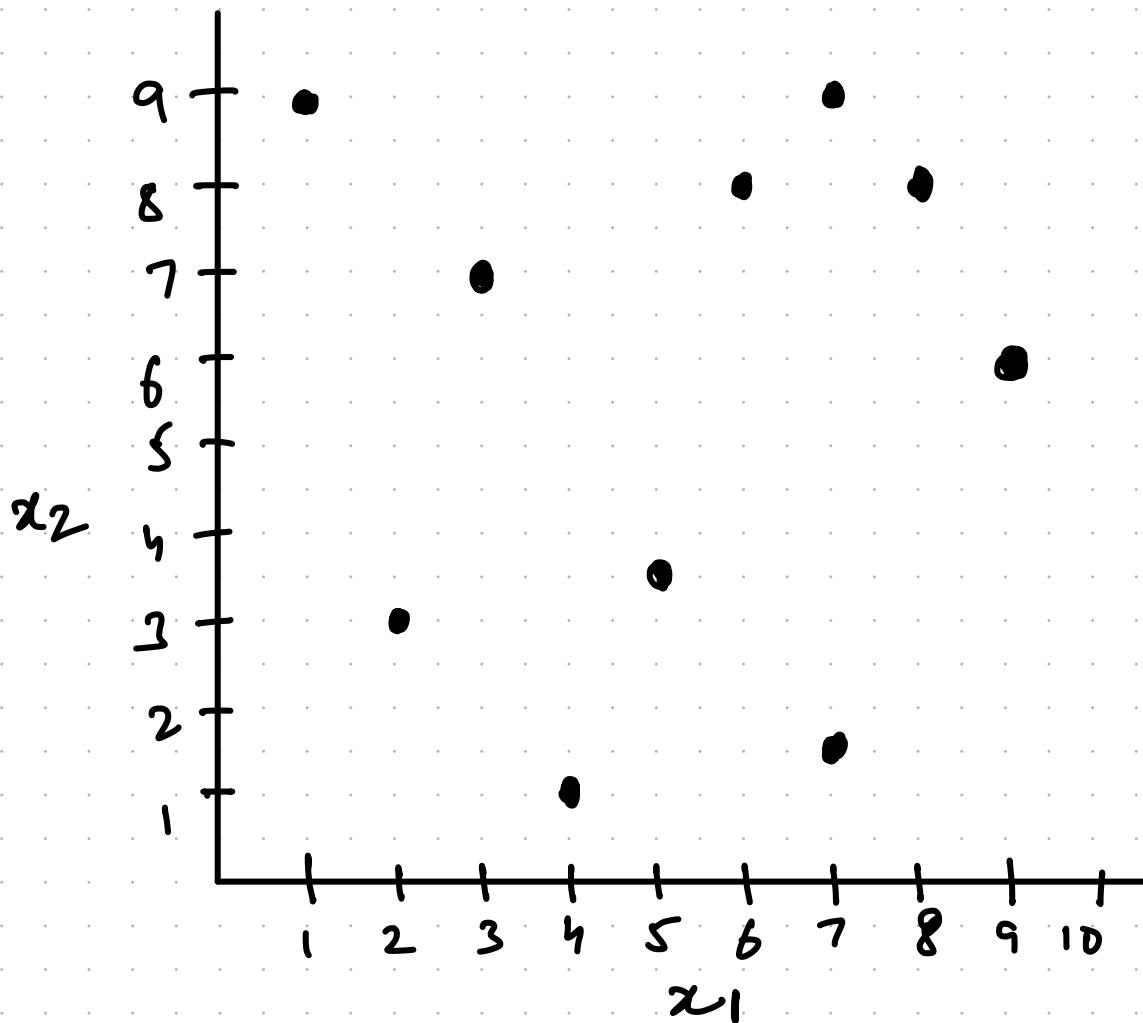
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## K-D trees

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- $X = \{(1, 9), (2, 3), (4, 1), (3, 7), (5, 5), (6, 8), (7, 2), (8, 8), (7, 9), (9, 6)\}$



Find median  
along  $x_1$   
and split  
data

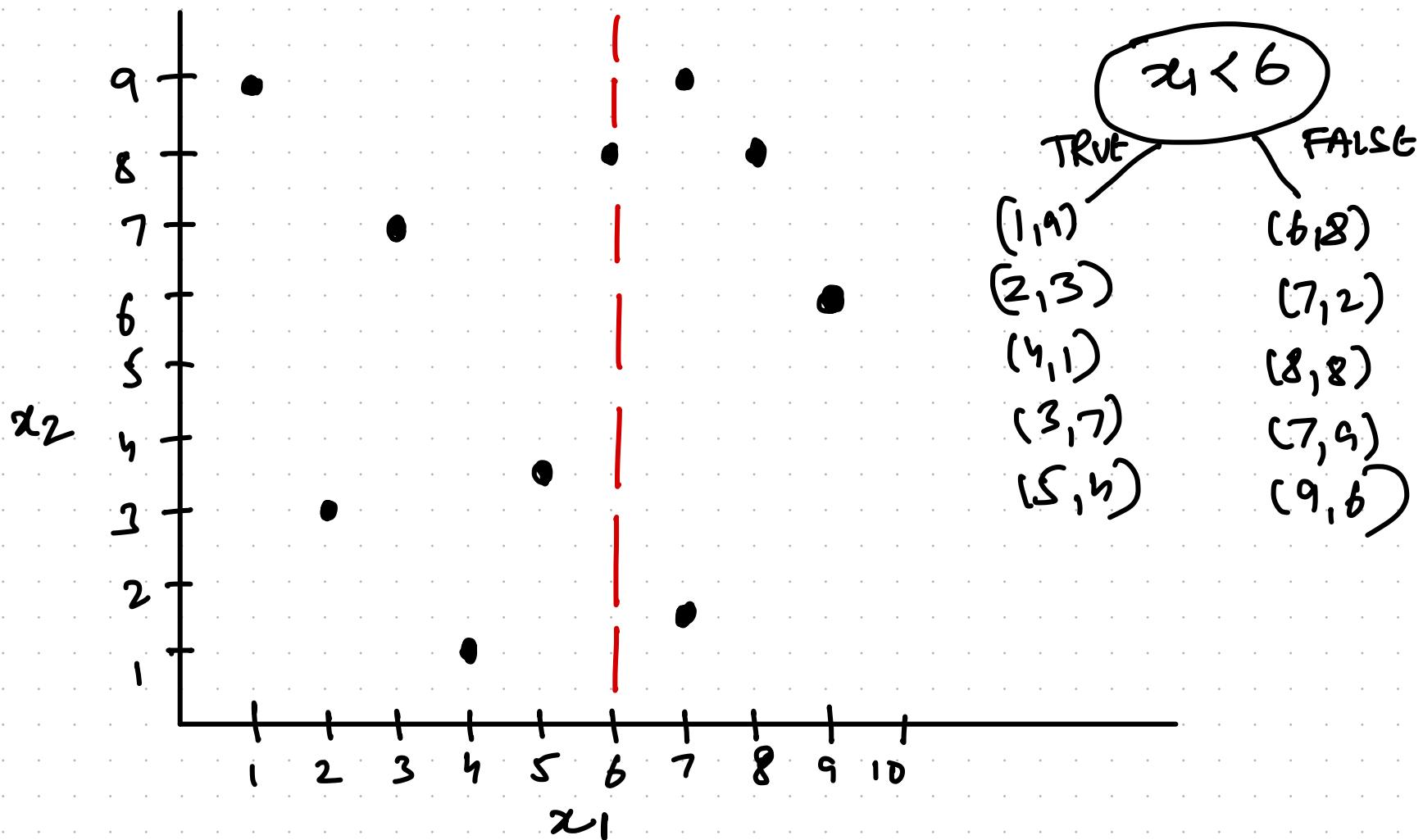
$$\text{MED}(1, 2, 4, 3, 5, 6, 7, 8, 7, 9)$$

$$= \\ \text{MED}(1, 2, 3, 4, 5, \underline{6}, 7, 7, 8, 9)$$

## K-D trees

(Victor Lamarcq slides)

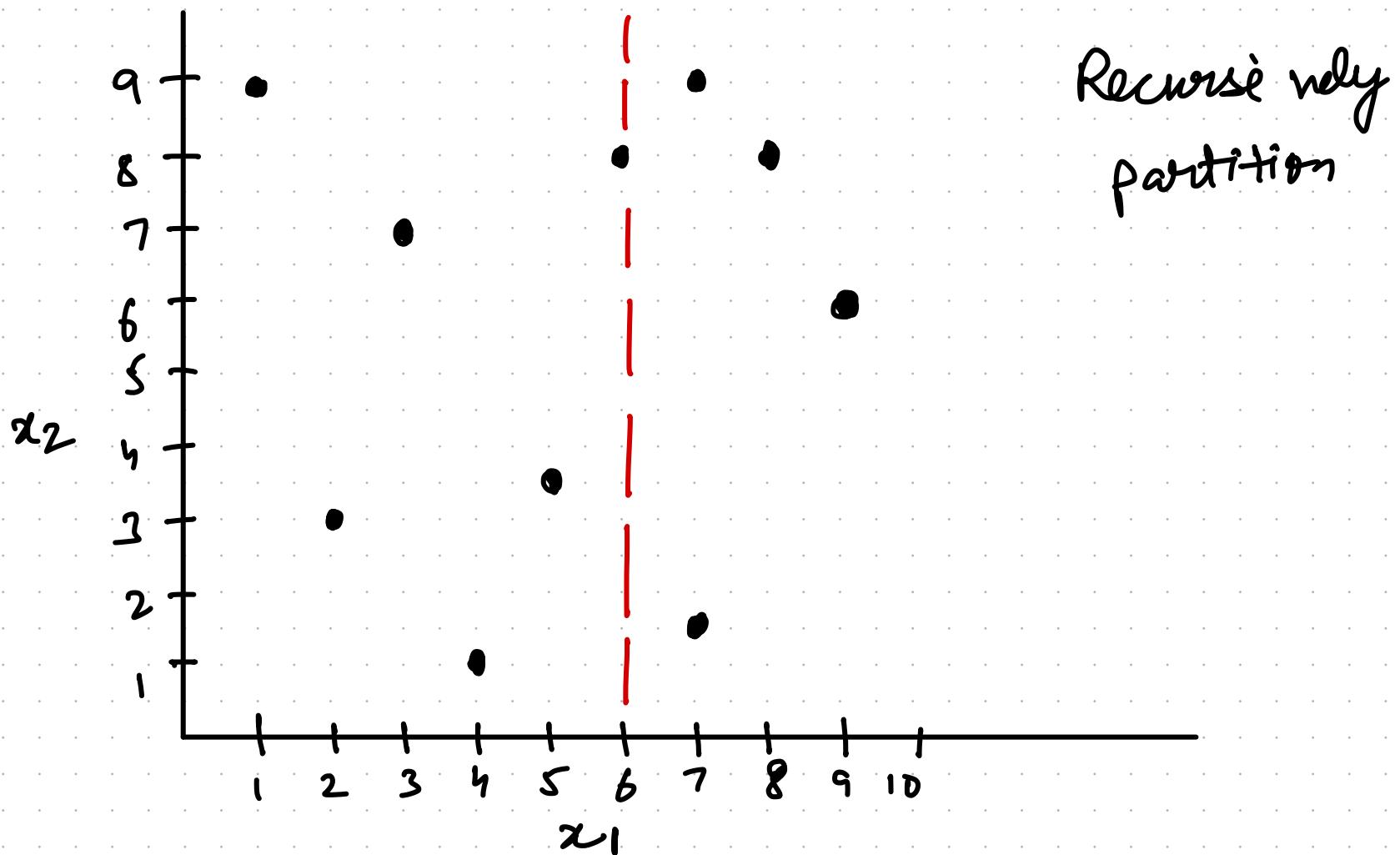
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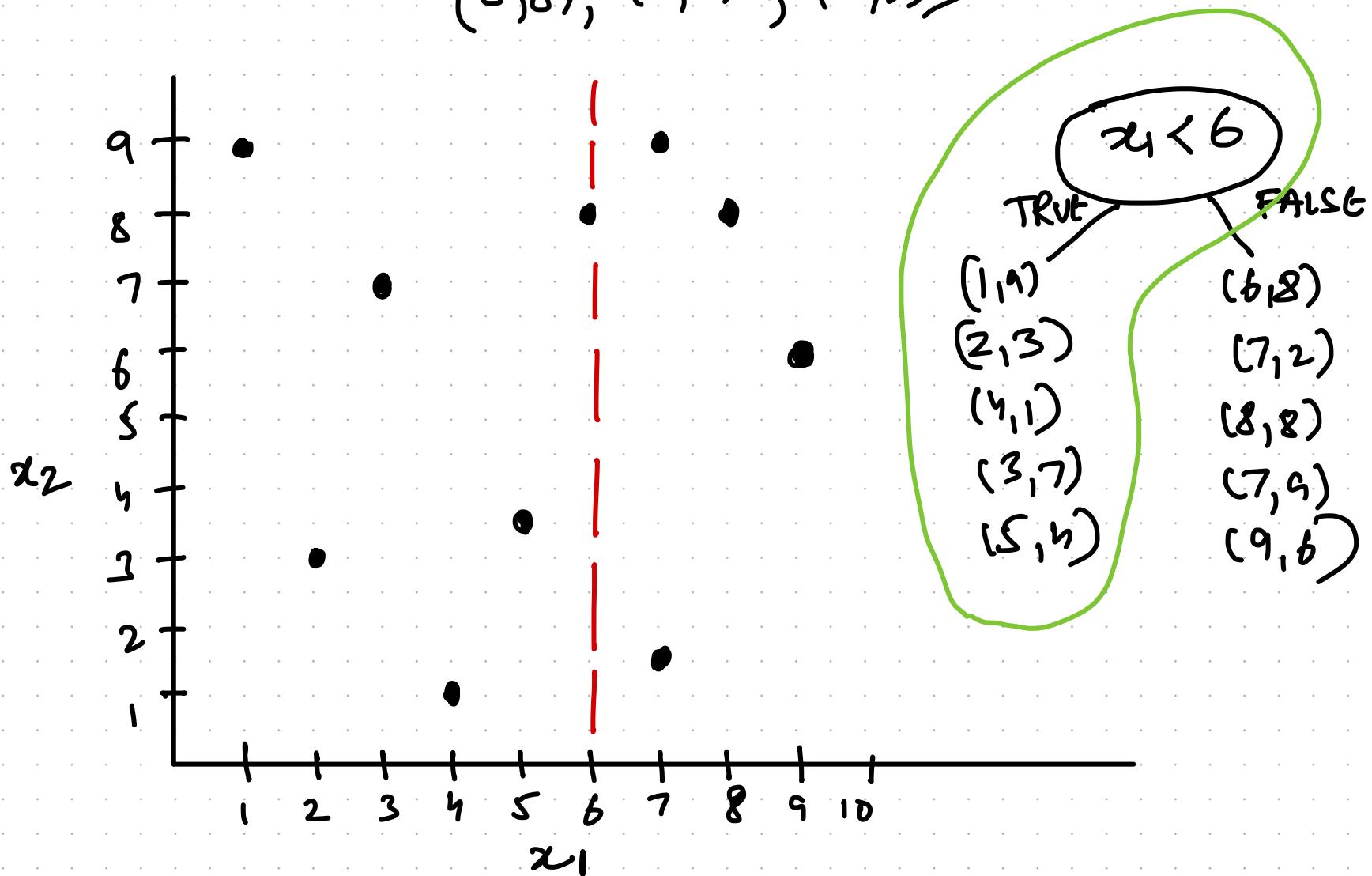
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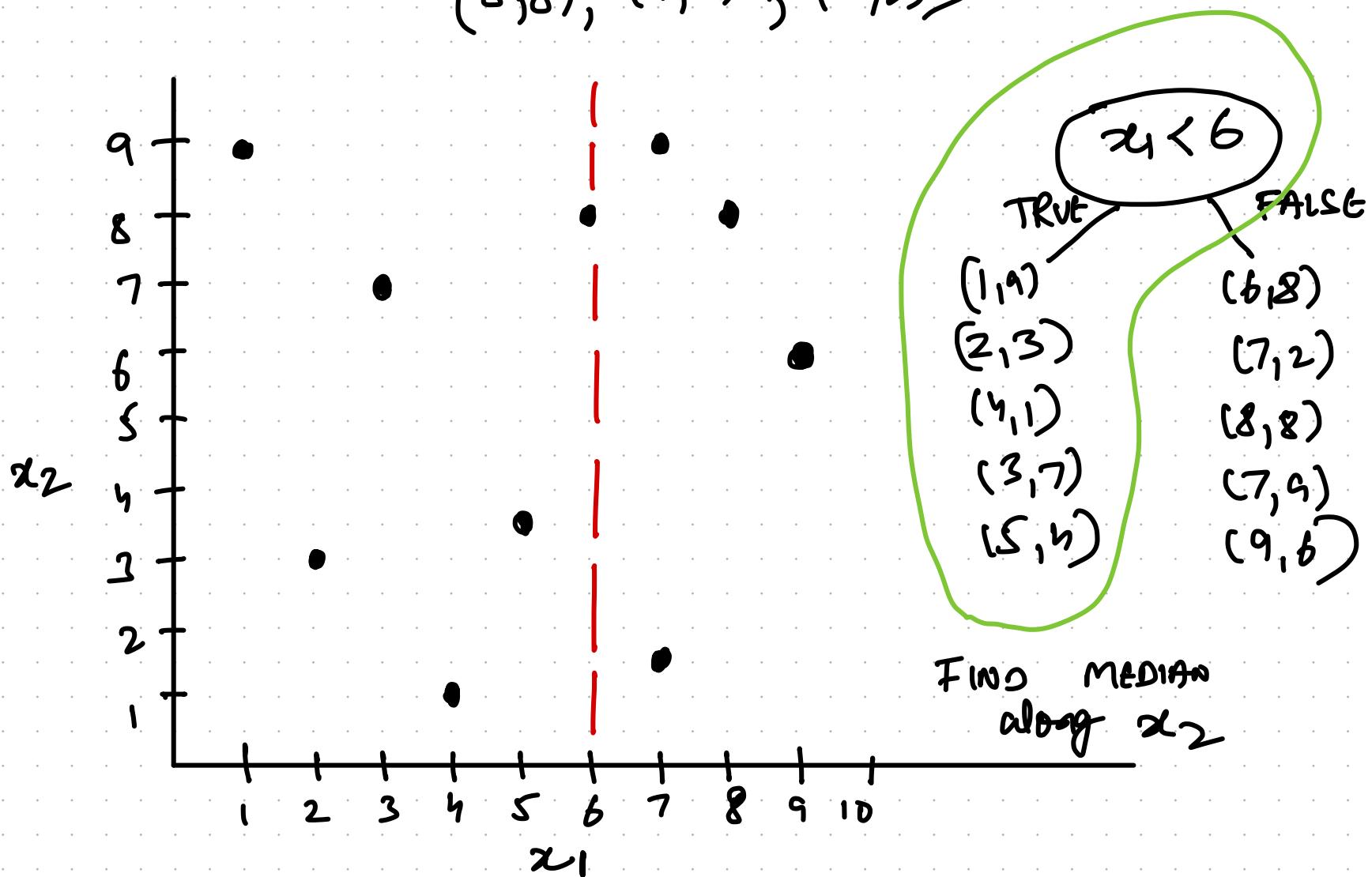
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(Victor Lamarcqo slides)

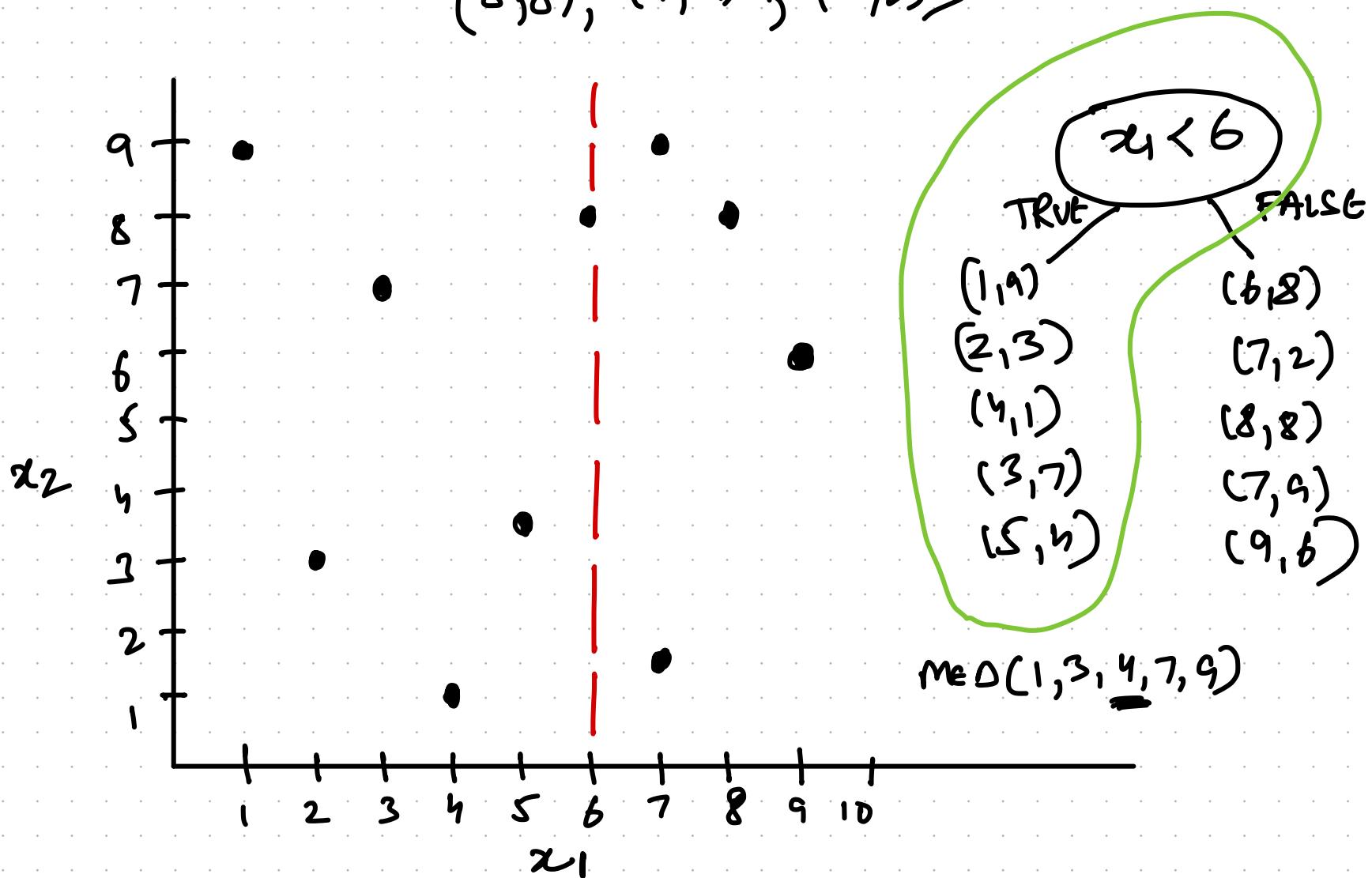
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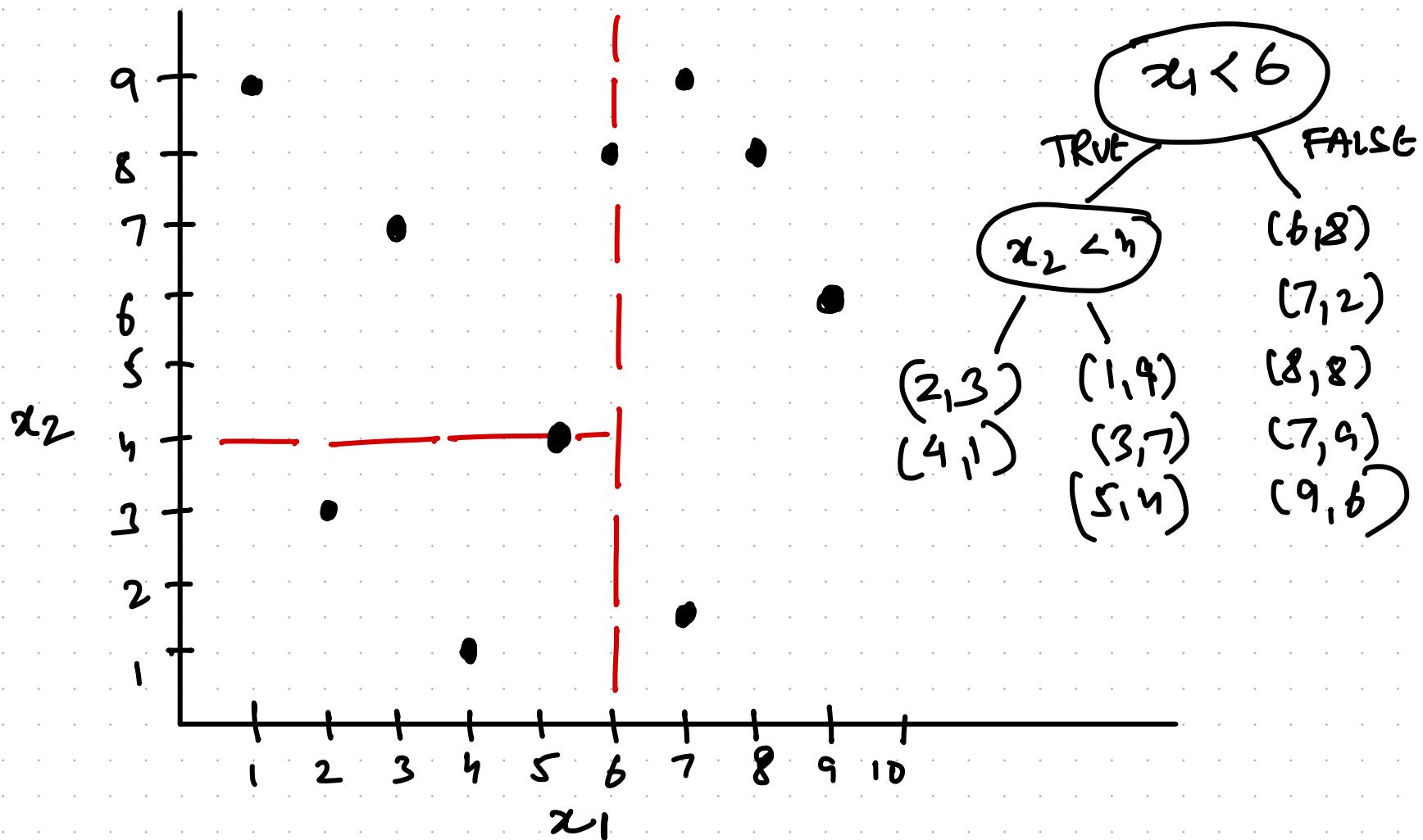
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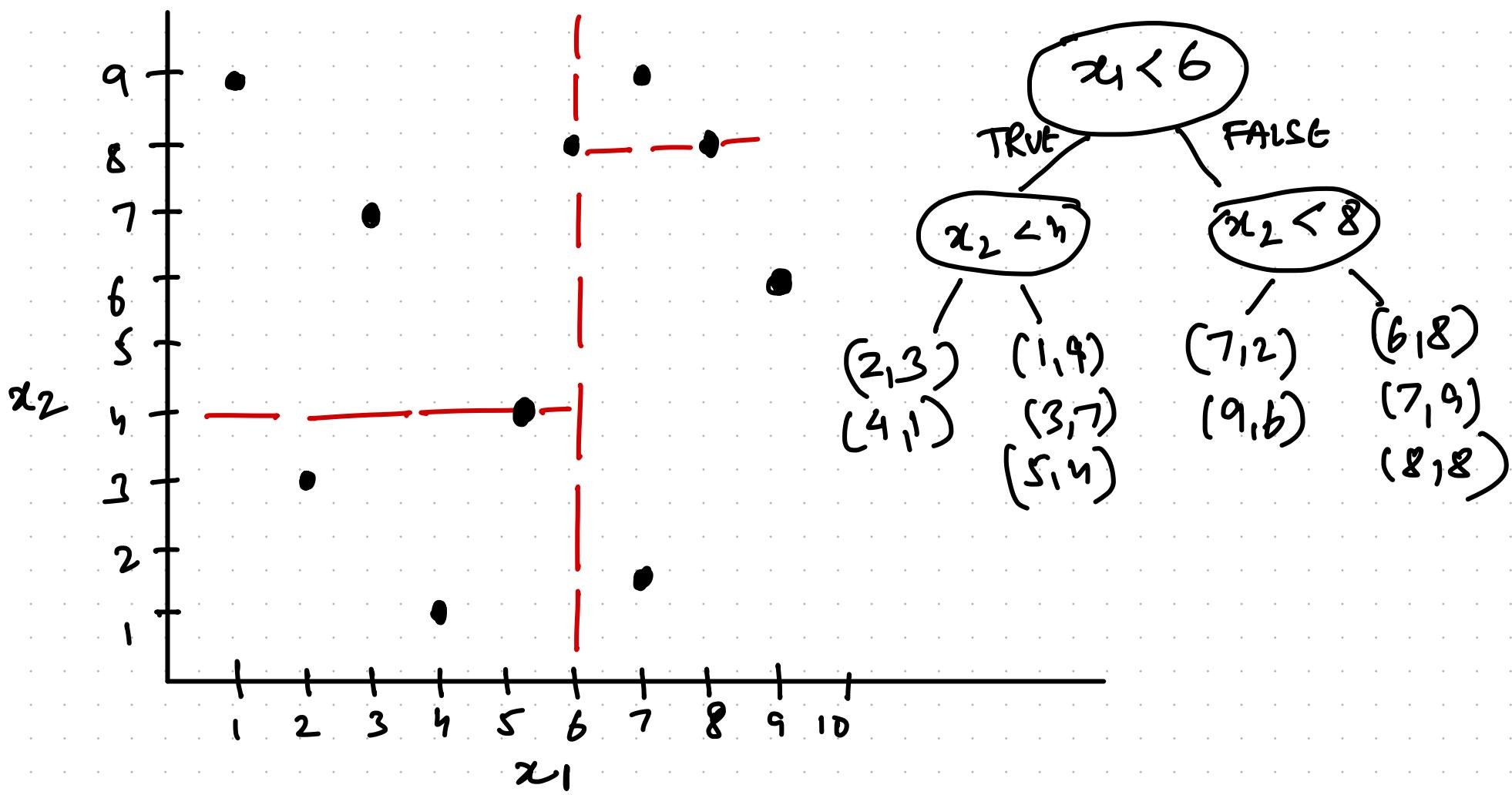
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## K-D trees

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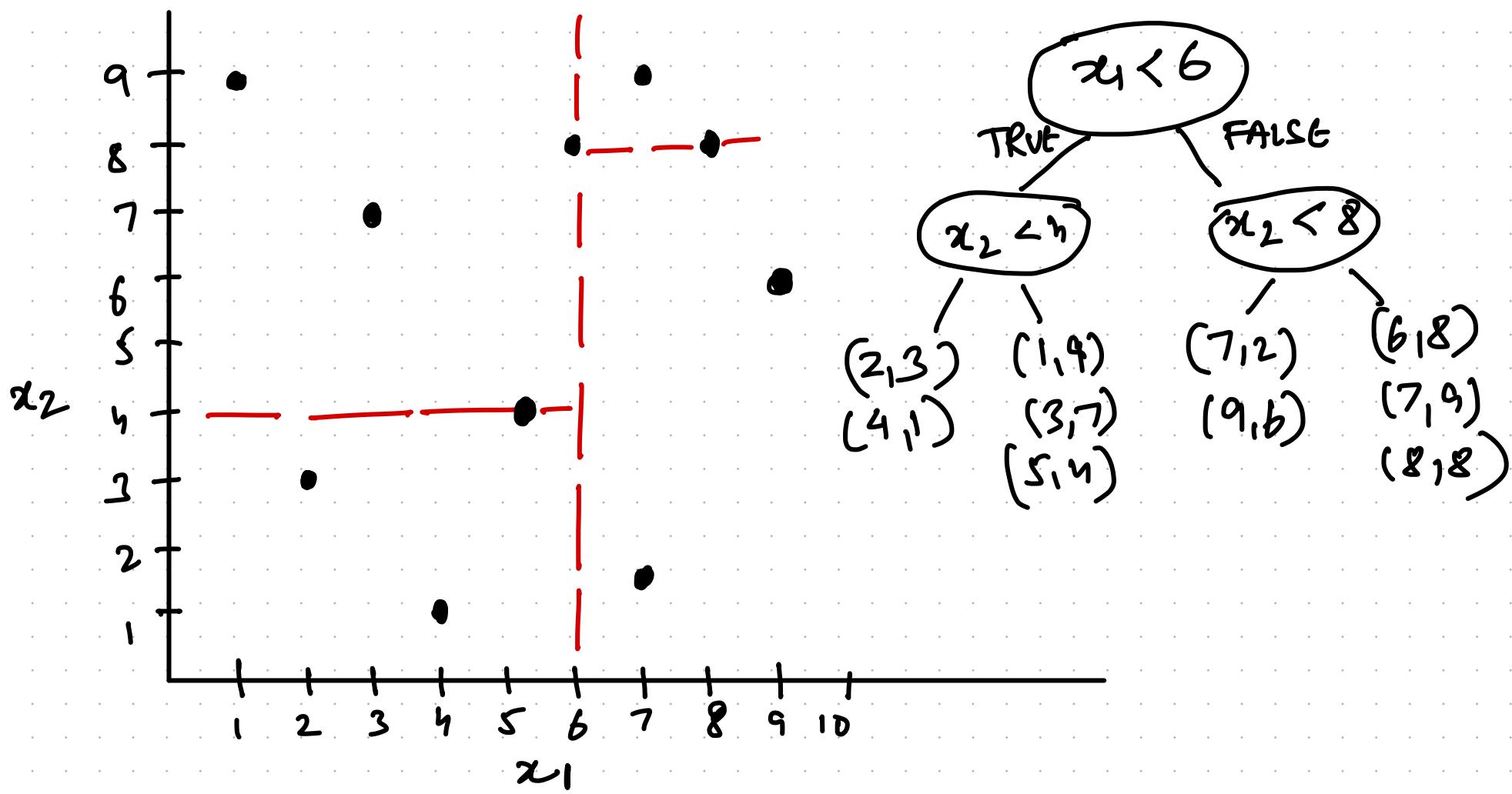
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(Victor Lamarcq slides)

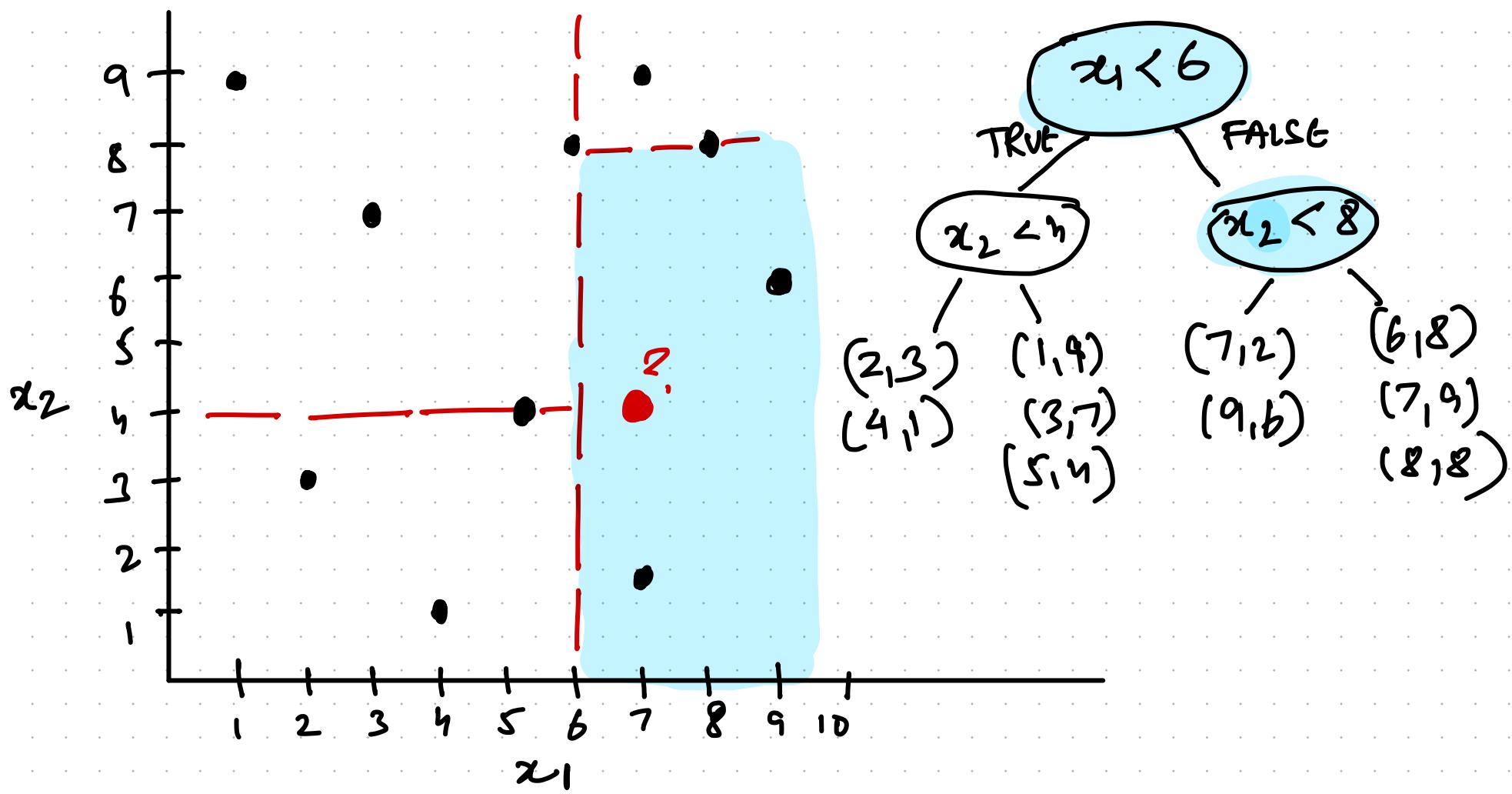
- query pt  $(7, 4)$



## K-D trees

(Victor Lamarcq slides)

- query pt  $(7, 5)$

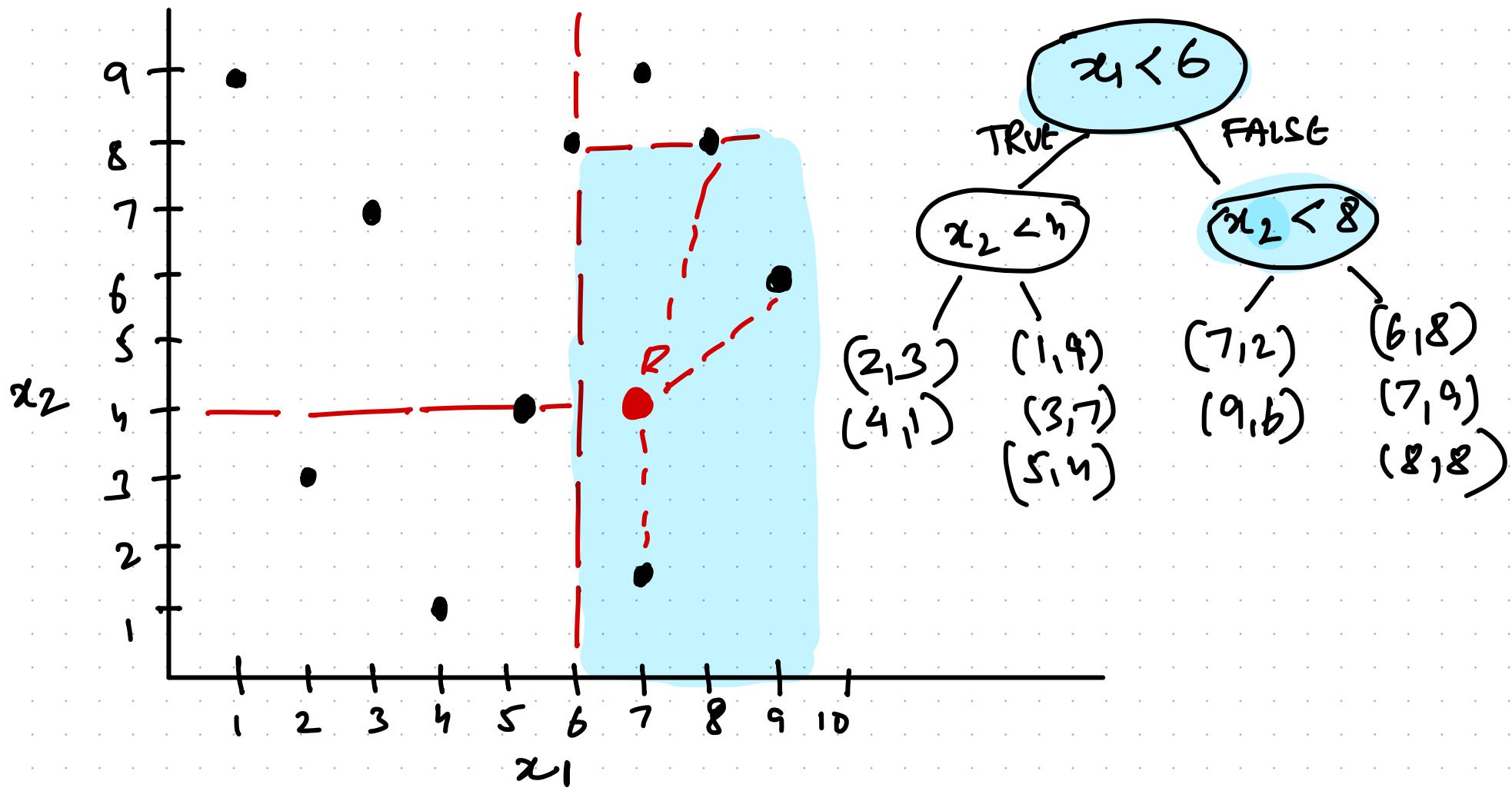


## K-D trees

(Victor Lamarcq slides)

- query pt  $(7, 5)$

FOR finding NBS, look  
in subspace



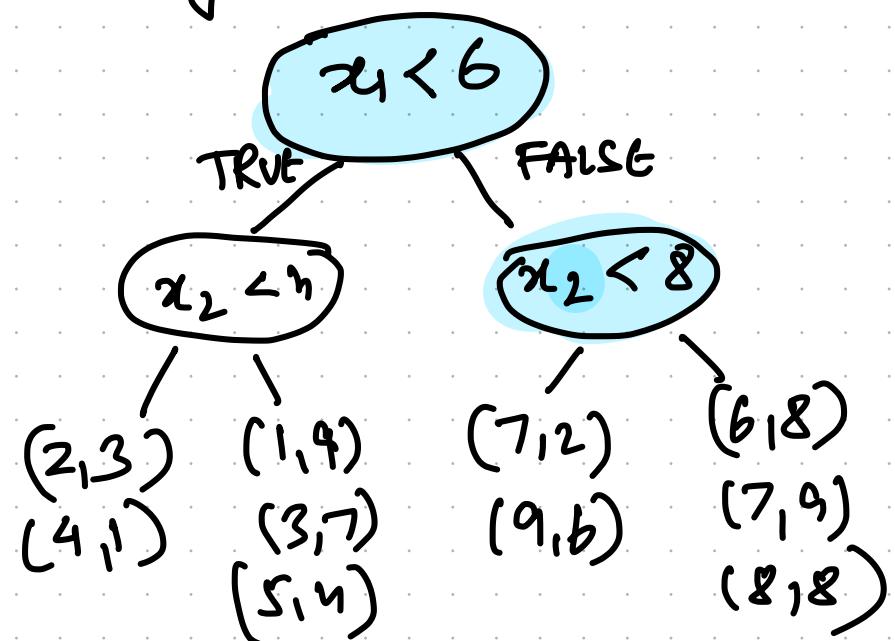
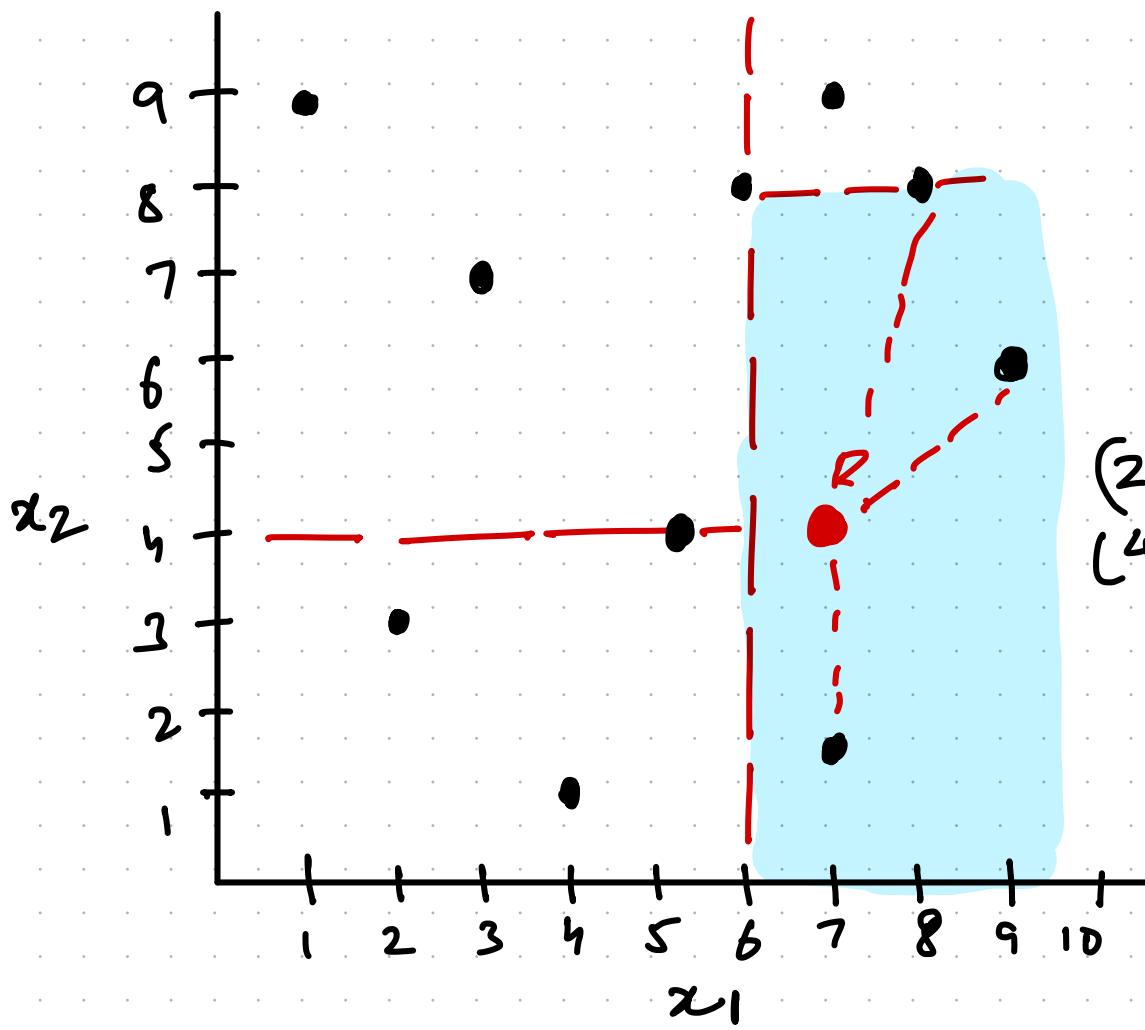
## K-D trees

(Victor Lamarcq slides)

- query pt  $(7, 5)$

FOR finding NNS, look  
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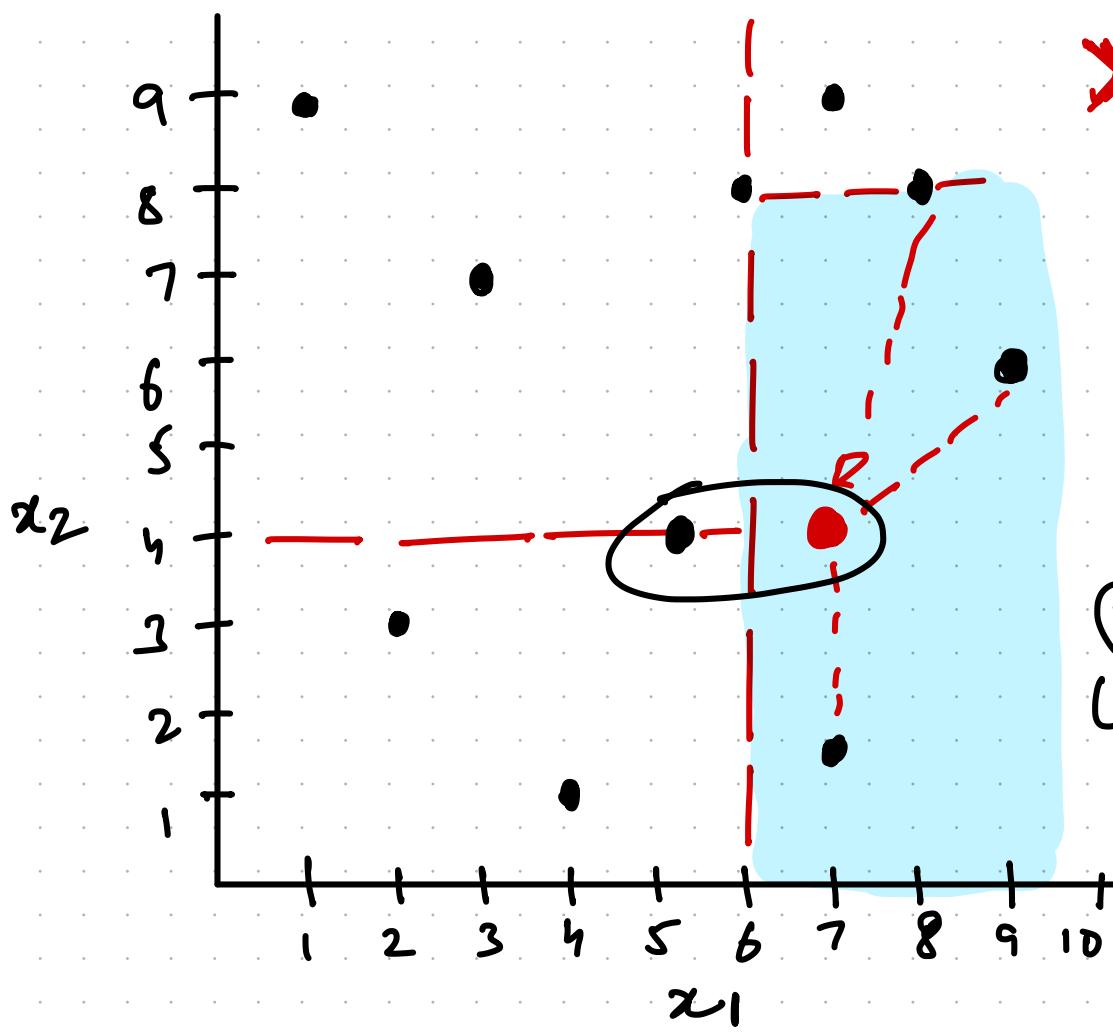
✓ way lesser comparisons



## K-D trees

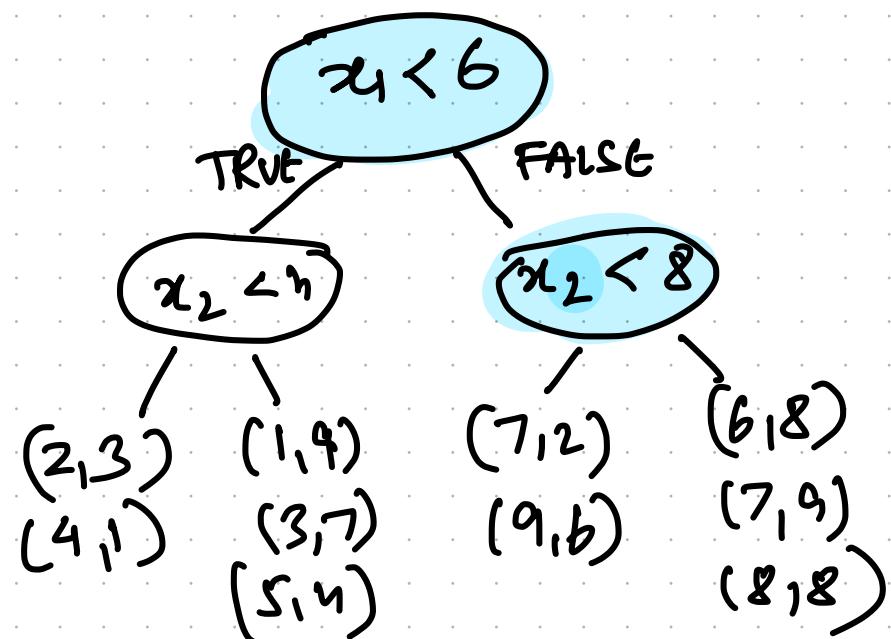
(Victor Lamarcq slides)

- query pt  $(7, 5)$



FOR finding NNS, look  
in subspace

- ✓ Way lesser comparisons
- ✗ May miss closest neighbors



KD trees (time complexity)

TRAINING TIME

(Ref. Wikipedia KD trees)

KD trees (time complexity)

TRAINING TIME

# samples in subtrees

# KD trees (time complexity)

TRAINING TIME

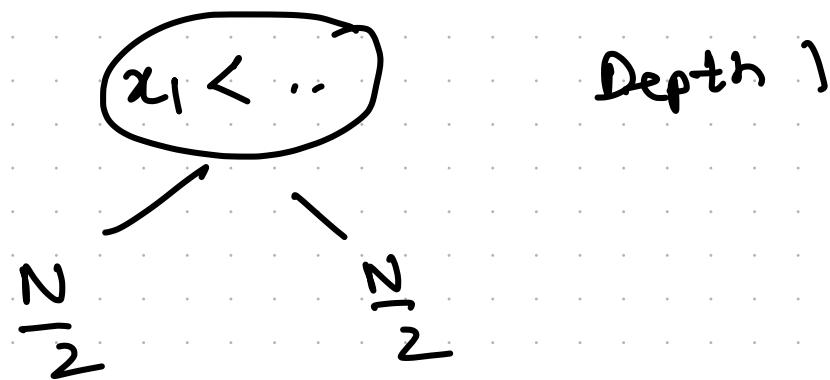
# samples in subtrees

Depth 0  
|  
N

# KD trees (time complexity)

TRAINING TIME

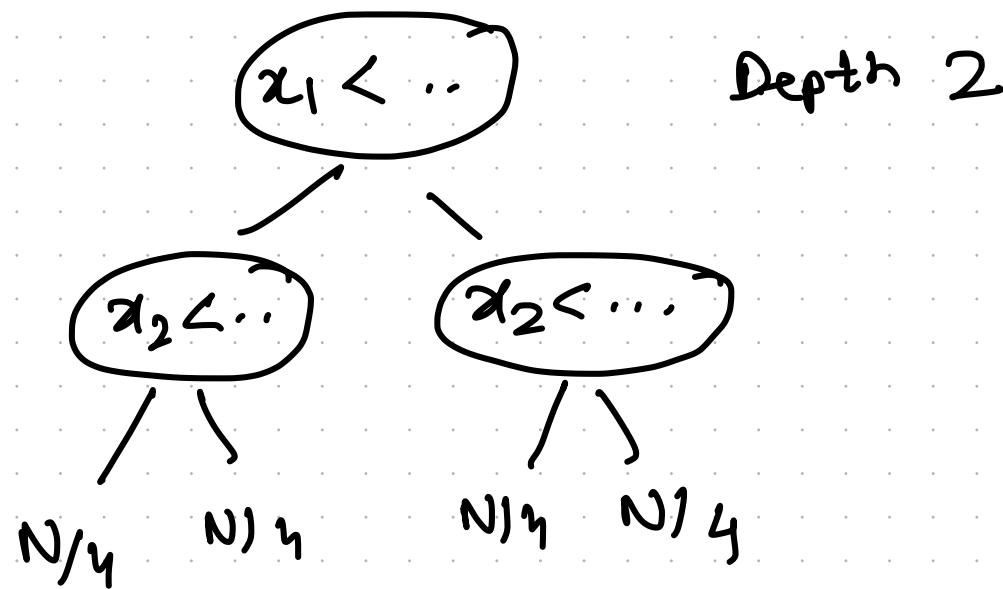
# samples in subtrees



# KD trees (time complexity)

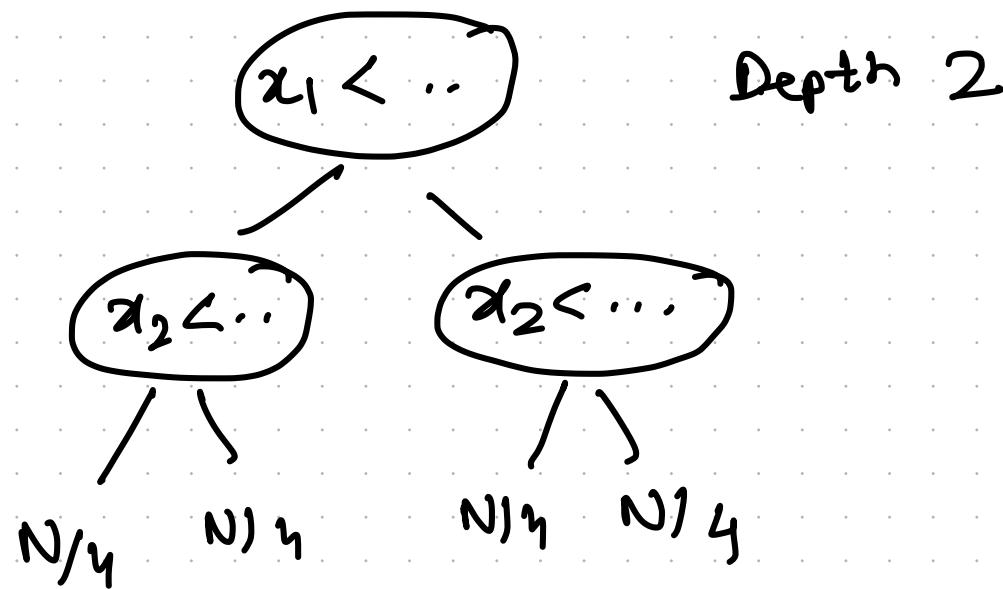
TRAINING TIME

# samples in subtrees



# KD trees (time complexity)

TRAINING TIME



Depth  $O(\log_2 N)$

# KD trees (time complexity)

## Training Time

- we have  $O(\log_2 N)$  levels.
- For each level,
  - sort "N" examples to find median
  - $O(N \log_2 N)$  time
- Total time to build datastructure  $O(N \log^2 N)$

## KD trees (time complexity)

- \* let's assume 'S' ( $\approx 40$ ) is leaf size  
(# samples at which we stop partitioning)
- \* what is time complexity of query?

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    - $O(D * S)$

## KD trees (time complexity)

\* What is time complexity of query?

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  - $O(\log N)$
- How many computat's per level?
  - One  $\rightarrow$  comparison (e.g.,  $x_1 < \delta$ )
- How many computat's at leaf?
  - $O(D \times S)$

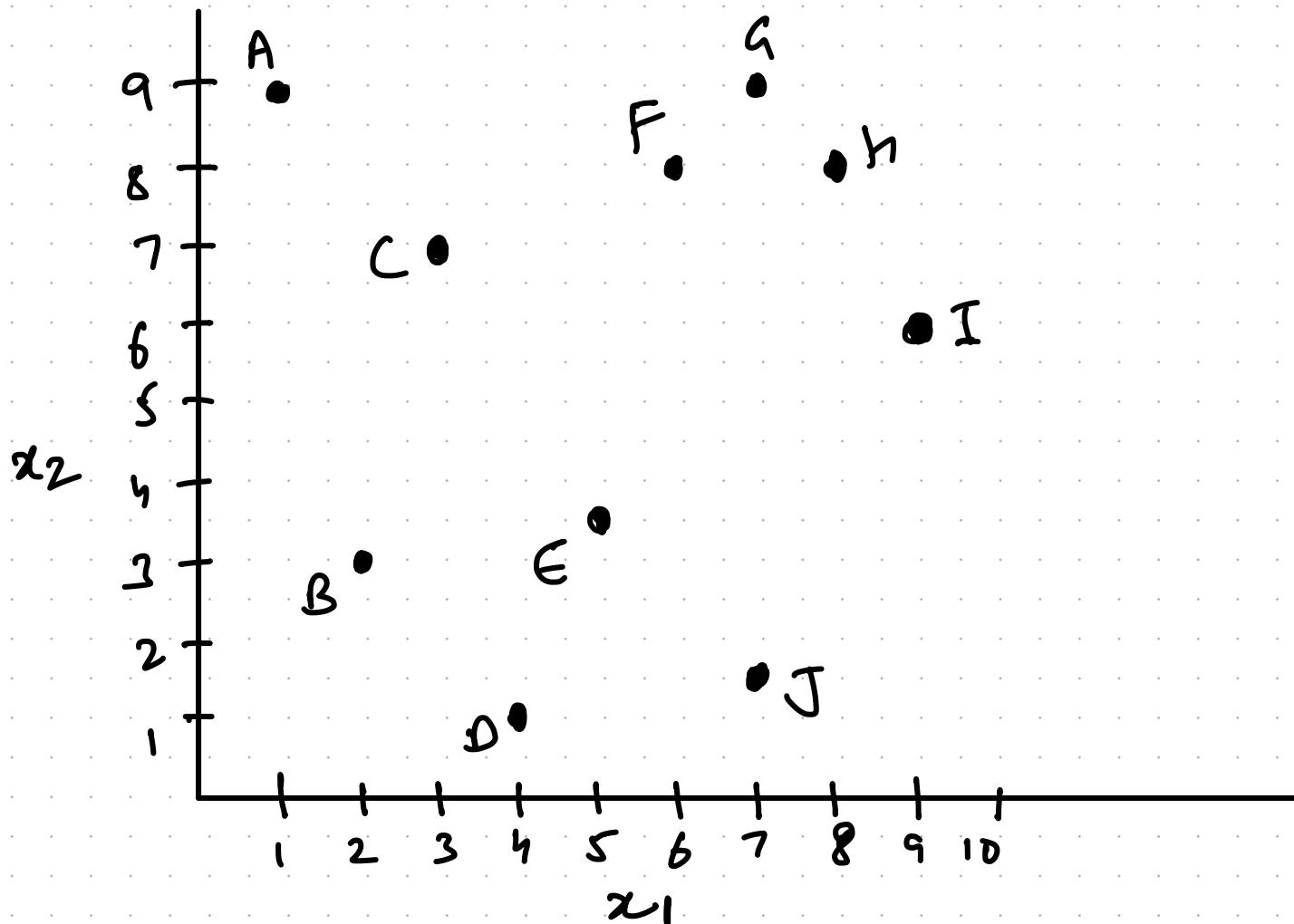
TOTAL QUERY  $O(\log N + D \times S)$

if  $D \ll N$  and  $S$  is small

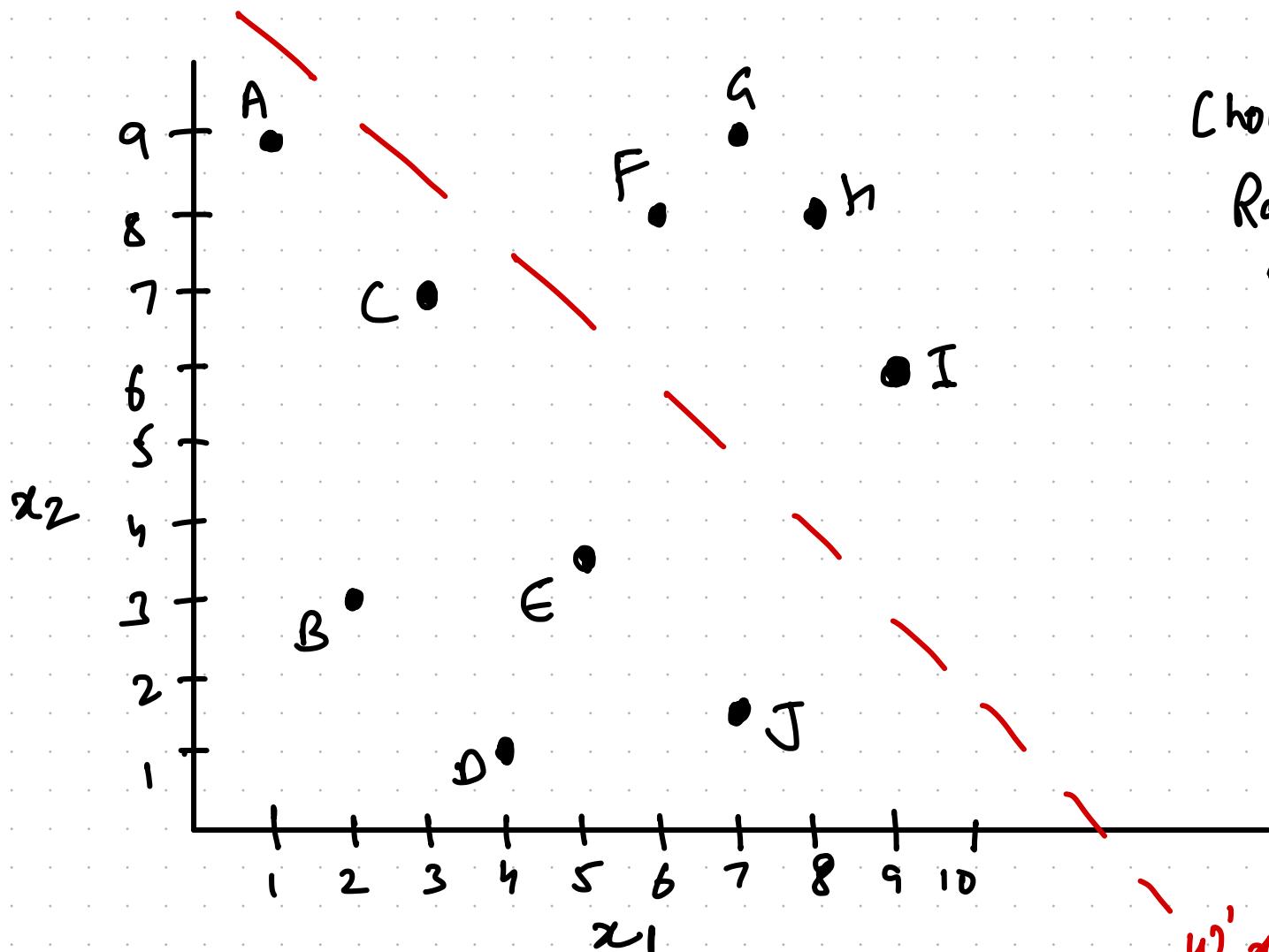
$$O(\log N)$$

# Locality Sensitive Hashing (LSH) w/ Random Projection

(MachineLearningInterview.com)



# Locality Sensitive Hashing (LSH) w/ Random Projection

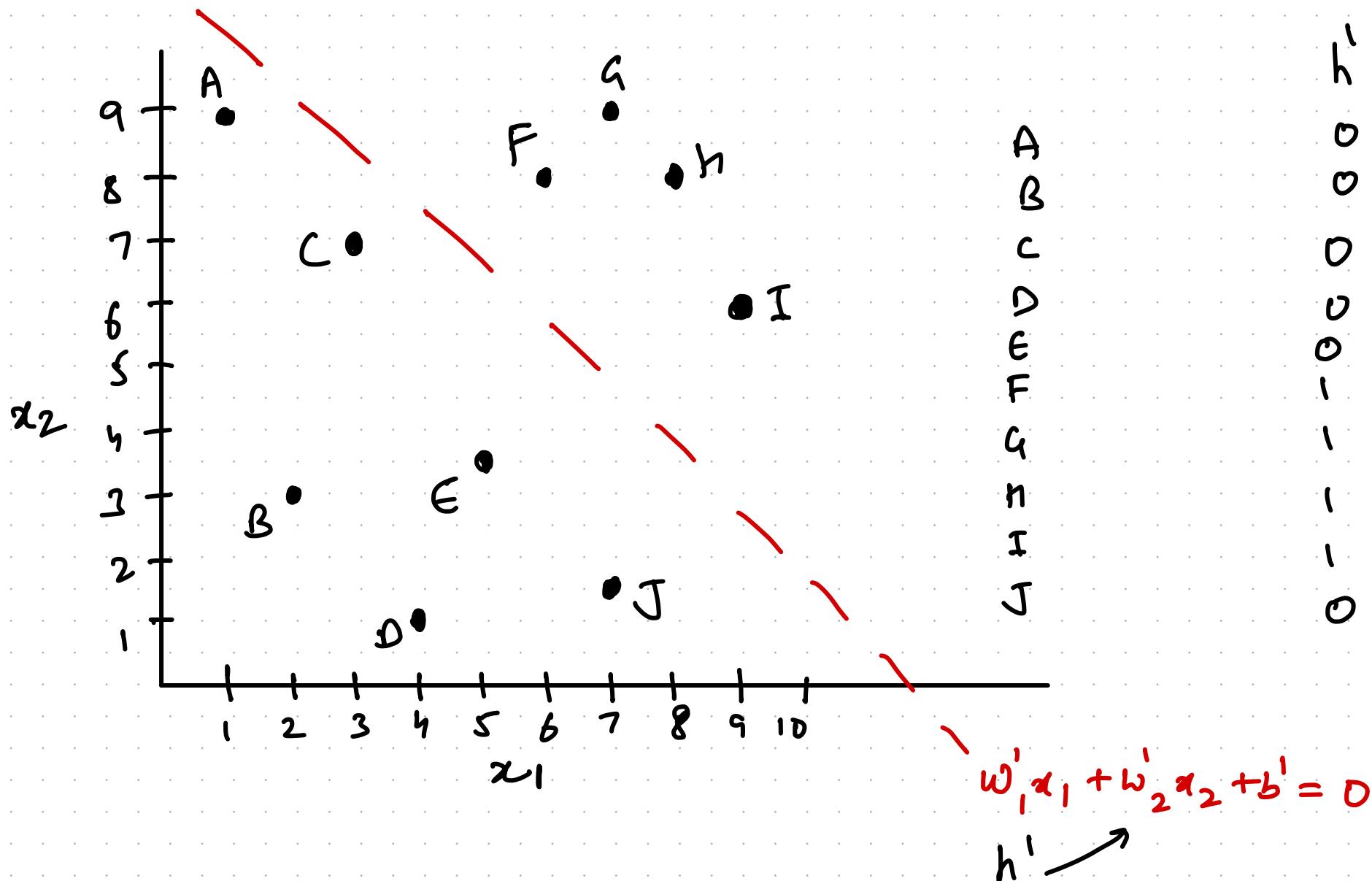


Choose ' $h'$  :  
Random hyperplane  
in 'D' dim

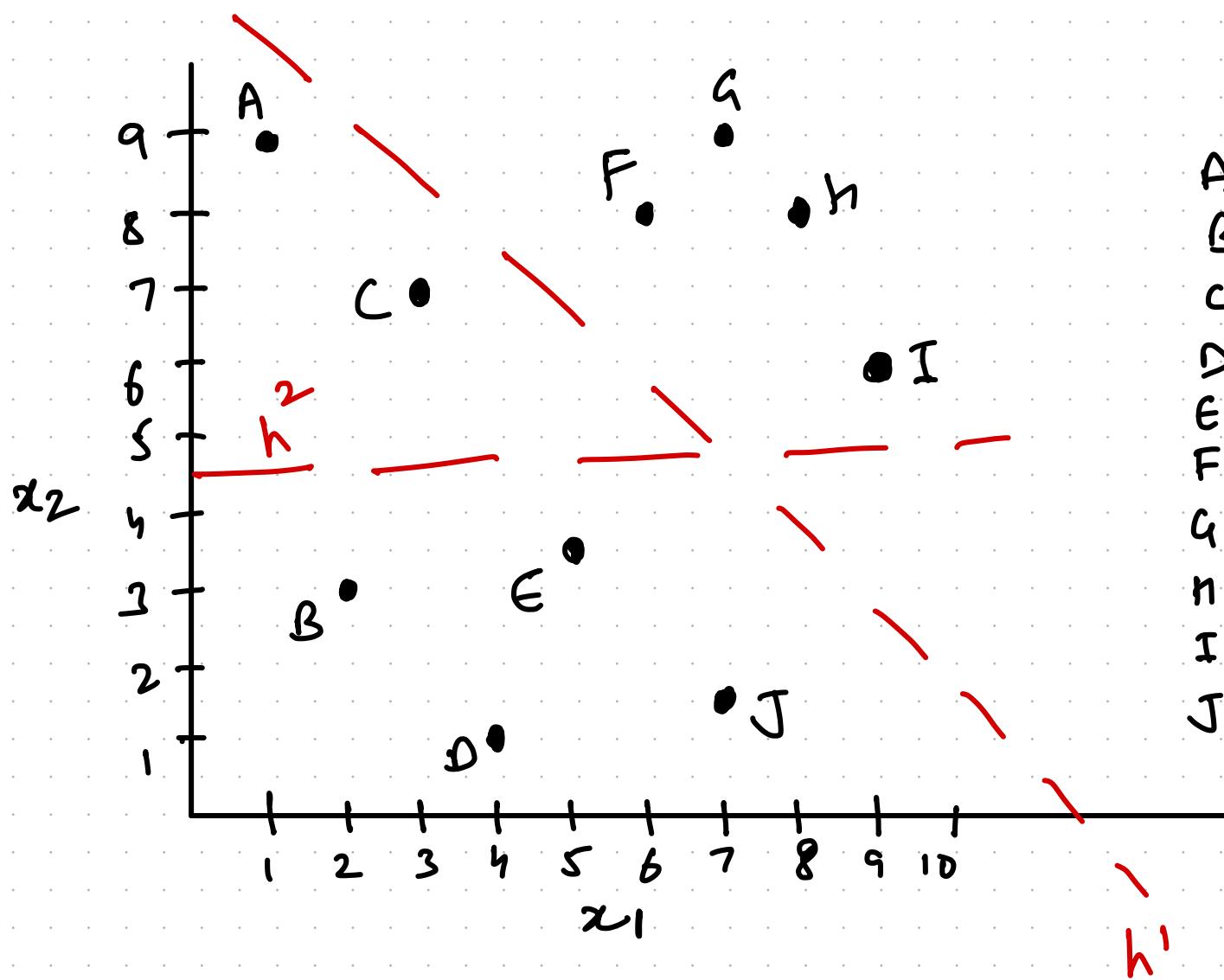
$$w'_1 x_1 + w'_2 x_2 + b' = 0$$

$h' \rightarrow$

# Locality Sensitive Hashing (LSH) w/ Random Projection

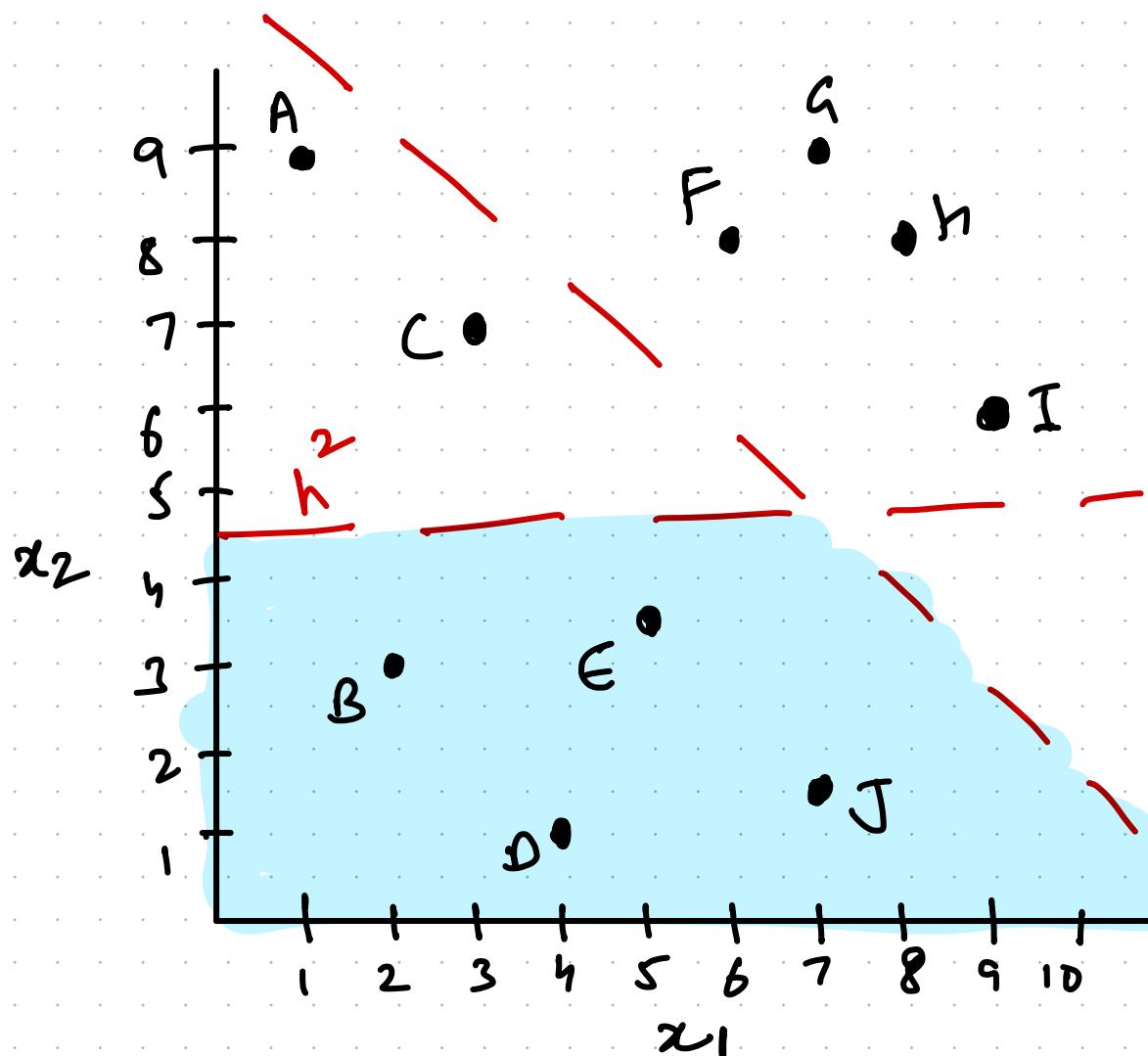


# Locality Sensitive Hashing (LSH) w/ Random Projection



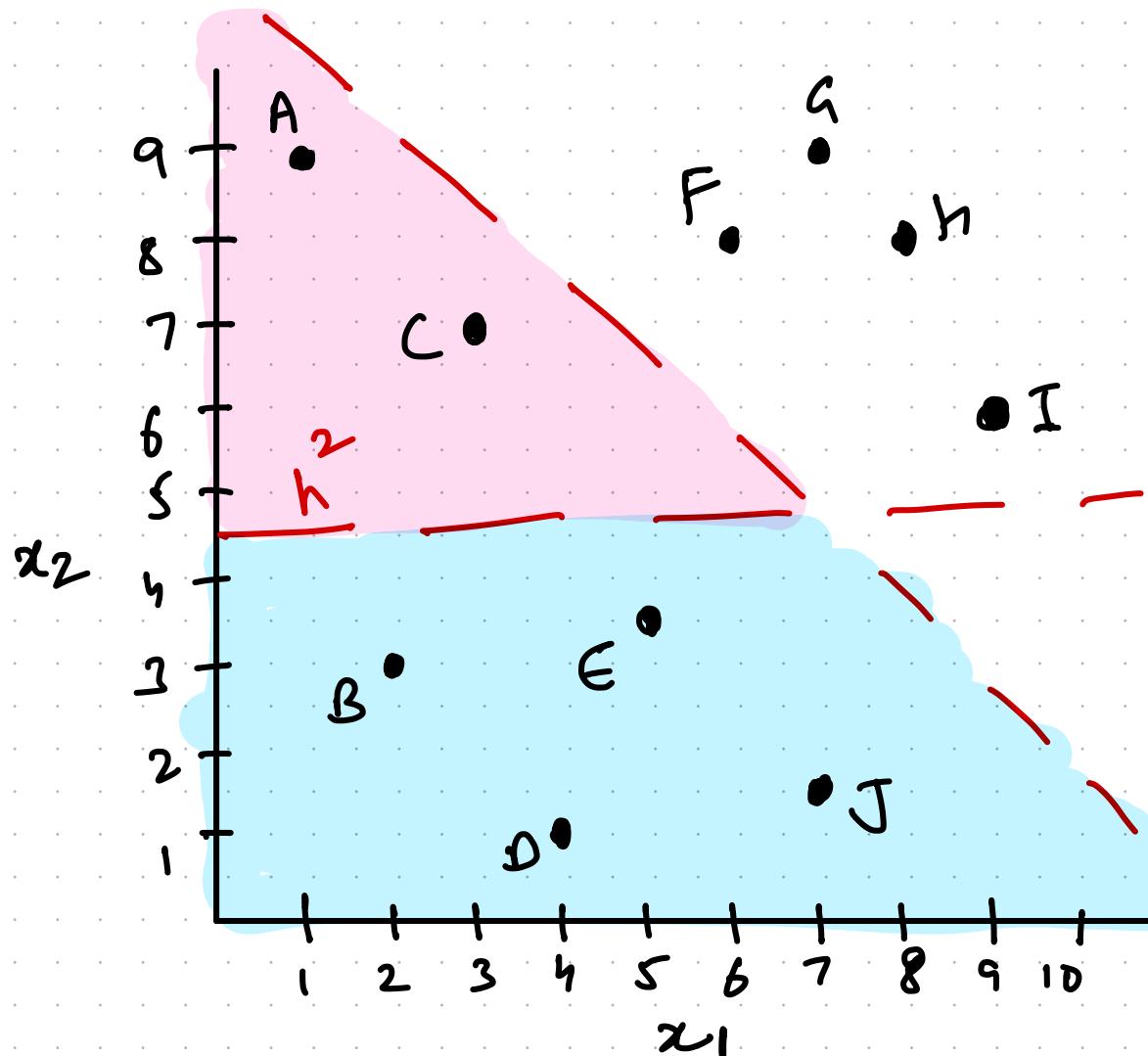
$\begin{array}{r} h^2 \\ - 1 \\ \hline 0 \end{array}$        $\begin{array}{r} h^1 \\ - 0 \\ \hline 0 \end{array}$

# Locality Sensitive Hashing (LSH) w/ Random Projection



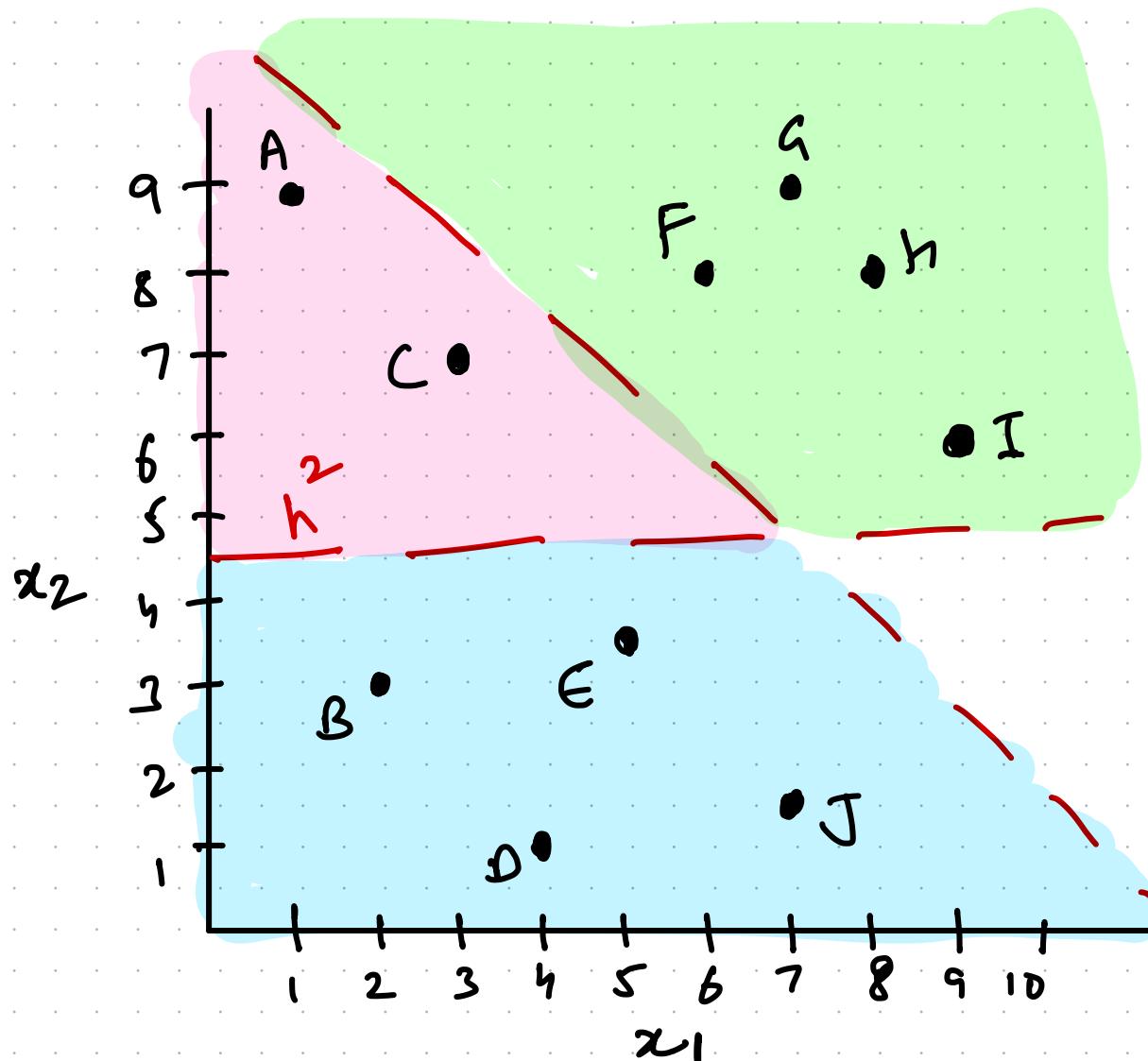
	$h^2$	$h^1$
A	1	0
B	0	0
C	1	0
D	0	0
E	0	0
F	-1	-1
G	-1	-1
H	1	1
I	-1	-1
J	0	0

# Locality Sensitive Hashing (LSH) w/ Random Projection



	$h^2$	$h'$
A	1	0
B	0	0
C	1	0
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# Locality Sensitive Hashing (LSH) w/ Random Projection



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## Practical implementation

\* How to sample hyperplane?

$i = 1 \dots D :$  sampled from

$$w_i \sim N(\dots, \dots)$$

$$b \sim N(\dots, \dots)$$

## Practical implementation

\* How to get  $h^i(x) = 0 \text{ or } 1$

$$h^i(x) = \text{SIGN}(b + \hat{w}_1^i x_1 + w_2^i x_2 + \dots)$$

## Practical implementation

- \* How to vectorize finding Hash Table

$$X = \begin{bmatrix} & \\ & \end{bmatrix}$$

## Practical implementation

\* How to vectorize finding Hash Table

$$X = \begin{bmatrix} & \\ & \\ & \\ & \\ & \end{bmatrix} \quad N \times D$$

$$x^l = \begin{bmatrix} | \\ | \\ . \\ | \end{bmatrix} \quad N \times D+1$$

## Practical implementation

\* How to vectorize finding Hash Table

$$x^i = \begin{bmatrix} 1 \\ - \\ \vdots \\ i \end{bmatrix}_{N \times D+1}$$

let # planes be 'P'

$$w = \begin{bmatrix} \dots & \dots & \dots \\ n(\dots, \dots) \\ \dots & \dots & \dots \end{bmatrix}_{D+1 \times P}$$

## Practical implementation

\* How to vectorize finding Hash Table

$$x' = \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}_{N \times D+1}$$

$$H = \text{Sgn}(x' w)_{N \times P}$$

let # planes be 'P'

$$w = \begin{bmatrix} \dots & \dots & \dots \\ N(\dots, \dots) \\ \dots & \dots & \dots \end{bmatrix}_{D+1 \times P}$$

## Time complexity

Assuming one (1) set of hash functions

$$\rightarrow h_{N \times P} = \text{SGN}(x_{N \times D} \cdot r_{D \times P})$$

$$\text{Time} = O(N \times D \times P)$$

Time to generate  $r \approx O(D \times P)$

Thus, at training time :  $O(N \times D \times P)$

At testing time,

- Computing hash on  $q_{1 \times D}$  takes  $O(D^P)$
- Assuming ' $T$ ' points in bucket:  
 $O(T * D)$  to find  
closest  
neighbour

At testing time,

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- What is T in terms of N and P (on average)

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$$T \approx \frac{N}{2^P}$$

At testing time,

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closest  
neighbour
- What is T in terms of N and P (on average)

$$T \approx \frac{N}{2^P}$$

$$\rightarrow \text{Test Time} = O(DP + \frac{DN}{2^P})$$