### PhD Progress Report 1st year

### Tommaso Papini

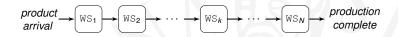
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#### 31st October 2017

- Analysis of assembly lines
- A hybrid technique for MRP transient analysis
- Other projects
  - the LINFA project
  - Activity Recognition for Ambient Assisted Living
- Research plan for the next year

Analysis of assembly lines

# Assembly line



*N* sequential workstations  $WS_1, \ldots, WS_N$ 

- with transfer blocking
- and no buffering capacity

Workstation WSk can be in one of three states

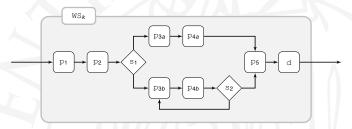
- ▶ producing: ₩Sk is working on a product
- ► done: ₩Sk is done working on a product
- ▶ idling: WSk is waiting for a new product



### Workstation

### Each workstation WSk

- has no internal parallelism
  - at most one item being processed in each workstation
- can implement complex workflows
  - sequential/alternative/cyclic phases with random choices
- and has GEN phases' durations



The last phase has no duration and encodes the done state

### Underlying stochastic process

The underlying stochastic process of each isolated workstation is a Semi Markov Process (SMP)

- due to GEN durations
- and the absence of internal parallelism

The whole assembly line finds a renewal in any case where

- every done station is in a queue before a bottleneck
- and everything else is idling



### Inspection with partial observability

The assembly line can be inspected by external observers

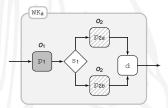
- the line can be considered at steady-state at inspection
- there can be ambiguity about the current phase

An observation is a tuple  $\omega = \langle \omega_0, \omega_1, \dots, \omega_N \rangle$ 

- lacksquare  $\omega_0$  indicates if a new product is ready to enter the line or not
- $\omega_k = \langle \sigma_k, \phi_k \rangle$  refers to  $WS_k$ 
  - $\sigma_k$  indicates if  $WS_k$  is idle/producing/done
  - $lackbox{}{} \phi_k$  identifies the set of possible current phases

#### Two kinds of uncertainty

- about the actual current phase
  - discrete
- about the remaining time in the current phase
  - continuous



### Performance measures

#### **Time To Done**

 The remaining time until workstation k, according to observation ω, reaches the *done* state

#### Time To Idle

The remaining time until workstation k, according to observation ω, reaches the idling state

#### **Time To Start Next**

The remaining time until workstation k, according to observation ω, starts the production of a new product



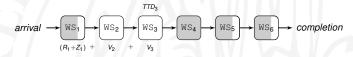




### Time To Done

$$\mathsf{TTD}(k,\omega) := \left\{ \begin{array}{ll} \displaystyle \sum_{\gamma \in \phi_k} P_{k,\gamma,\omega} \cdot (R(k,\gamma) + Z(k,\gamma)), & \text{if } \sigma_k = \textit{producing} \\ \\ \mathsf{TTD}(k-1,\omega) + V(k), & \text{if } \sigma_k = \textit{idling} \\ \\ 0, & \text{if } \sigma_k = \textit{done} \end{array} \right.$$

- $ightharpoonup P_{k,\gamma,\omega}$  probability weight that  $\mathbb{WS}_k$  is in phase  $\gamma$  according to  $\omega$
- ▶  $R(k, \gamma)$  remaining time in phase  $\gamma$  of  $WS_k$
- ▶  $Z(k, \gamma)$  execution time of phases of  $WS_k$  that follow  $\gamma$
- ▶ V(k) production time of  $WS_k$



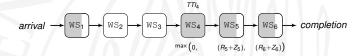
Backward recursive evaluation

Evaluation of performance measures

### Time To Idle

$$\mathsf{TTI}(k,\omega) := \left\{ \begin{array}{ll} \max\{\mathsf{TTD}(k,\omega),\mathsf{TTI}(k+1,\omega)\}, & \text{if } \sigma_k \in \{\textit{producing},\textit{done}\} \\ \\ 0, & \text{if } \sigma_k = \textit{idling} \end{array} \right.$$

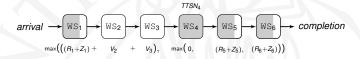
- ►  $\mathsf{TTI}(k, \omega) = \mathsf{max}\{\mathsf{TTD}(k, \omega), \dots, \mathsf{TTD}(k + n, \omega)\}$ 
  - ▶  $WS_i$  producing/done  $\forall j \in [k, k+n]$
  - either  $WS_{k+n}$  last workstation or  $WS_{k+n+1}$  idling
- ▶ ₩S<sub>K</sub> becomes idle when the bottleneck finishes its production



Forward recursive evaluation

### Time To Start Next

$$\mathsf{TTSN}(k,\omega) := \mathsf{max}\{\mathsf{TTI}(k,\omega),\mathsf{TTD}(k-1,\omega)\}$$



Forward and backward recursive evaluation

# Disambiguation of observed phases

Resolve observed (producing) phases' ambiguity

• steady-state probability that  $\mathtt{WS}_k$  is in phase  $\gamma$  according to  $\omega$ 

Given observation  $\phi_k$  for workstation  $WS_k$ 

- we compute probability  $P_{k,\gamma,\omega}$
- that it was actually  $\gamma$  that produced  $\phi_k$

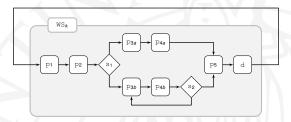
$$extstyle extstyle P_{m{k},\gamma,\omega} = rac{ ilde{\pi}(\gamma)}{\displaystyle\sum_{\gamma' \in \phi_{m{k}}} ilde{\pi}(\gamma')}$$

•  $\tilde{\pi}(\gamma)$  steady-state probability of phase  $\gamma$  in an isolated model of  $\mathtt{WS}_k$ 

### Isolated workstation model

The isolated workstation model represents a workstation repeatedly processing a product

- one product being processed
- after its production, it's moved back to the entry point of the workstation



It can be used for two reasons

- steady-state probabilities of producing phases are independent
- the inspection is at steady-state
  - arrivals and productions can be considerer in equilibrium

# Remaining time

Evaluation of  $F_{R(k,\gamma)}(t) = \text{CDF of } R(k,\gamma)$ 

▶  $R(k, \gamma)$  remaining time in phase  $\gamma$  of  $WS_k$ 

#### Problem!

- remaining times of enabled GEN transitions are dependent
- joint probabilities don't allow for a compositional approach

### 1/3 Immediate approximation

- ightharpoonup assume that phase  $\gamma$  is inspected at its ending
  - $\tilde{F}_{R(k,\gamma)}(t) = 1 \quad \forall t$
- represents an upper bound

### <sup>2</sup>/<sub>3</sub> Newly enabled approximation

- lacktriangleright assume that phase  $\gamma$  is inspected at its beginning
  - $\tilde{F}_{R(k,\gamma)}(t) = F_{\gamma}(t)$
  - $F_{\gamma}(t)$  original CDF of the duration of  $\gamma$
- represents a lower bound

### Remaining time

<sup>3</sup>/<sub>3</sub> *Independent remaining times* approximation

- consider the remaining times of ongoing phases as independent
- represents a (better) lower bound

### Theorem: positive correlation & stochastic order

If  $\hat{R}$  is an independent version of vector R of positively correlated remaining times of ongoing phases, then  $\hat{R} \ge_{st} R$ 

Steady-state distribution of  $\hat{R}(k, \gamma)$  computed according to the Key Renewal Theorem<sup>1</sup>

$$\tilde{F}_{R(k,\gamma)}(t) = \frac{1}{\mu} \int_0^t [1 - F_{\gamma}(s)] ds$$

μ expected value of F<sub>γ</sub>(t)

<sup>&</sup>lt;sup>1</sup>Serfozo, R., 2009. Basics of applied stochastic processes. Springer Science & Business Media.

# Execution and production time

### Evaluation of $F_{Z(k,\gamma)}(t)$ and $F_{V(k)}$

- ▶  $Z(k, \gamma)$  execution time of phases of  $WS_k$  that follow  $\gamma$
- $\triangleright$  V(k) production time of  $WS_k$

### CDFs of $Z(k, \gamma)$ and V(k) are computed as transient probabilities

- ▶  $F_{Z(k,\gamma)}$  transient probability from phase after  $\gamma$  to final phase of  $WS_k$
- $ightharpoonup F_{V(k)}$  transient probability from first phase to final phase of WS<sub>k</sub>

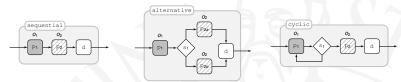
### Upper/lower bounds for TTD, TTI and TTSN can be evaluated

convolution and max operations maintain stochastic order

Experimentation

### Case study assembly lines

### Sequential, alternative and cyclic workstations



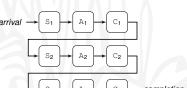
arrival

### Simple assembly line

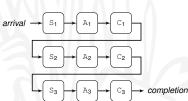
- two sequential workstations
- ▶ both in phase p₁ at inspection

### Complex assembly line

- three repetitions
- of sequential/alternative/cyclic ws
- all observed in producing

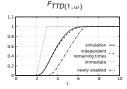


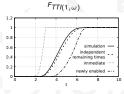
completion

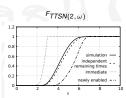


### Simple assembly line

TTDone, TTIdle, TTStartNext







### TTD, TTI and TTSN computed in

- ▶ 41/45/42 min for simulation
- ▶ 0.15/0.18/0.10 s for bounds

### Very good approximation results

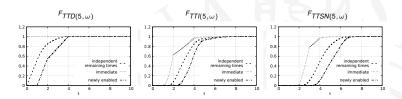
especially for independent remaining times

### Feasible approach

- very fast bounds evaluation
- compared to simulation

# Complex assembly line

TTDone, TTIdle, TTStartNext



TTD, TTI and TTSN computed in

▶ 0.126/0.123/0.75 s for bounds

#### Scalable solution

- in a complex scenario
- simulation would be infeasible

A hybrid technique for MRP transient analysis

# A hybrid technique for MRP transient analysis

Transient analysis of Markov Regenerative Processes (MRP) employing different techniques for different epochs

#### The basics:

- Exact techniques require specific conditions to be met
  - different techniques require different conditions
- Kernel rows of different epochs can be evaluated independently

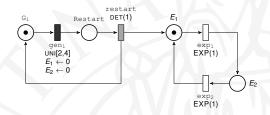
#### The idea:

- Evaluate each kernel row with a different technique
  - corresponding to the condition enabled in that epoch
  - eventually with an approximate technique, if no conditions are met
- Compute transient probabilities with Markov Renewal Equations

# Analysis under enabling restriction<sup>2</sup>

### **Enabling restriction**

at most one GEN enabled in each state



### Kernel rows computed by

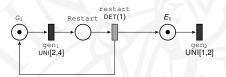
- analysing the CTMC
- subordinated to the activity of the GEN

<sup>&</sup>lt;sup>2</sup>German, R., Logothetis, D., & Trivedi, K. S. (1995, October). Transient analysis of Markov regenerative stochastic Petri nets: A comparison of approaches. In Petri Nets and Performance Models, 1995., Proceedings of the Sixth International Workshop on (pp. 103-112). IEEE.

# Analysis with stochastic state classes<sup>3</sup>

#### **Bounded regeneration restriction**

- a regeneration is always reached within a bounded number of events
- i.e. there are no cycles without regenerations



### Kernel rows computed by

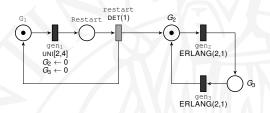
- enumerating stochastic transient trees
- of reached classes
- between two regenerations

<sup>&</sup>lt;sup>3</sup>Horváth, A., Paolieri, M., Ridi, L., & Vicario, E. (2012). Transient analysis of non-Markovian models using stochastic state classes. Performance Evaluation, 69(7), 315-335.

# Approximate analysis with stochastic state classes

#### No restriction

usable when no conditions are met



### Based on analysis with stochastic state classes

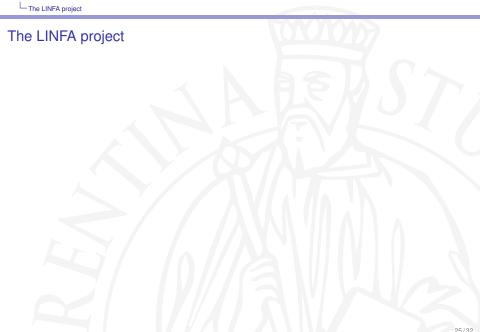
truncated after enough precision is met



PhD Progress Report: 1st year

Other projects

The LINFA project



PhD Progress Report: 1st year

Other projects

Activity Recognition for Ambient Assisted Living

# Activity Recognition for Ambient Assisted Living

PhD Progress Report: 1st year

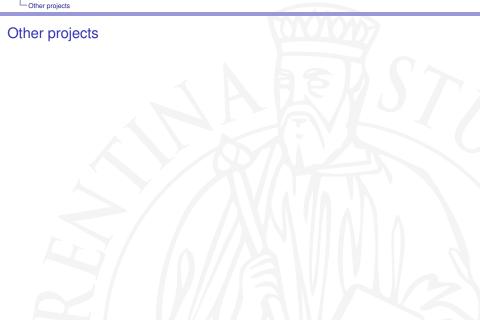
Other projects

LActivity Recognition for Ambient Assisted Living

Activity Recognition for Ambient Assisted Living

PhD Progress Report: 1st year

Cother projects
Other projects



Research plan for the next year

### Research plan for the next year

Analysis of assembly lines

### Introduction of buffering capacity

- with fixed/variable capacity
- so to model more realistic scenarios

#### Derivation of additional performance measures

- ▶ in the same compositional fashion
- e.g. production time of a certain product in the line
- or of the next N products

Derivation of a more educated upper bound

# Research plan for the next year The LINFA project

Model more aspects to refine the ward model

- introduce personalised healthcare protocols
- employ process mining techniques

### State-space optimisation

- avoid state-space explosion
- investigate other tools
  - Storm

Activity Recognition for Ambient Assisted Living

### Research plan for the next year

Activity Recognition for Ambient Assisted Living

#### Refine model based AR

- exploit fuzzy logic to include support for continuous sensors
  - accelerometer/thermometer/...

### Identify good AR datasets for AAL

- ▶ investigate the literature
- generate new datasets

In order to apply process mining techniques