

Object-Oriented Software Development

Advanced Principles

Faktorisierung: Sachverhalt nur an einer Stelle im SW System vorkommen

Statisch vs. dynamischer Typ: dynamischer erst zu Laufzeit bekannt, statischer Typ der Variablen ergibt sich durch Deklaration

Overview

Chapter 1: Principles

Chapter 2: Advanced Principles

Chapter 3: Class Libraries

Chapter 4: Design Patterns

Chapter 5: Design and Implementation

Chapter 6: Testing

Chapter 7: Refactoring

Chapter 8: Frameworks

Contents

- Type checking
-

- **Classes**
 - Nouns, first letter of each word is capitalized
 - Example: `ShadowedBox`
- **Interfaces**
 - Start with `I`, first letter of each internal word is capitalized
 - Example: `IMouseListener`
- **Methods**
 - First letter lowercase, first letter of each internal word is capitalized
 - Example: `setCoordinate`
- **Instance variables**
 - Starts with `f`, first letter of each internal word is capitalized
 - Example: `fFillPattern`

- Class variables
 - Starts with `s`, first letter of each internal word is capitalized
 - Example: `sMinSize`
- Constants
 - All uppercase, words are separated by “_”
 - Example: `MAX_ITEMS`

- Value
 - Describe abstract values
 - No lifecycle
 - No side effects (when shared)
 - Immutable state (conceptually),
can only be interpreted (accessible by different representations)
 - Examples: integer, float, character, monetary value, account number
- Object
 - Describe a phenomenon in a problem domain
 - Can be distinctively referenced
 - Can be shared → side effects
 - Lifecycle (creation, mutation, deletion)
 - Examples: account, payment, person

- Implementation has to follow object or value semantics!
- Implementation of value semantics:
 - Immutability
 - Copy on write
- Advantages: simple, easy sharing and threadsafeness
- Java
 - Class “Integer” is implemented with value semantics
 - Instances of Integer are immutable
 - Same holds true for BigInteger, BigDecimal etc
 - Class “Date” is implemented with object semantics
 - Mutable → side effects
 - Date is deprecated and (partially) replaced by Calendar

Types and Classes

- Type: Abstract description of properties of a set of objects
- Class: Implementation of a type
Object → Class → Interface → Type).
- Static type checking
 - Class/interface is a type (e.g. in Java, C++,...)
 - Assignment compatibility only along the class/interface hierarchy
- Dynamic type checking
 - Protocol of an object = type
 - Protocol is the set of messages understood by an object
 - More flexible than static type checking
 - Assignment compatibility is not restricted by class hierarchy

Prototypes

- Some languages get by without classes, e.g. JavaScript
- Object creation in JavaScript

```
var obj = {};  
obj.prop = 'value';  
obj.method = function() {  
    // ...  
};
```

- Constructor definition in JavaScript

```
function Person(name, age) {  
    this.name = name;  
    this.age = age;  
    this.greet = function() {  
        // ...  
    };  
}
```

- Static type checking
 - Type errors are usually detected at compile time
 - Allows for generating efficient target code (e.g. byte code)
 - Type errors can still occur at run-time (e.g. by type casting)

```
    =    ();  
    .    ();
```

← Compiler detects invalid method call

- Dynamic type checking
 - Flexibility at the expense of reliability
 - Developer is responsible for invoking valid methods

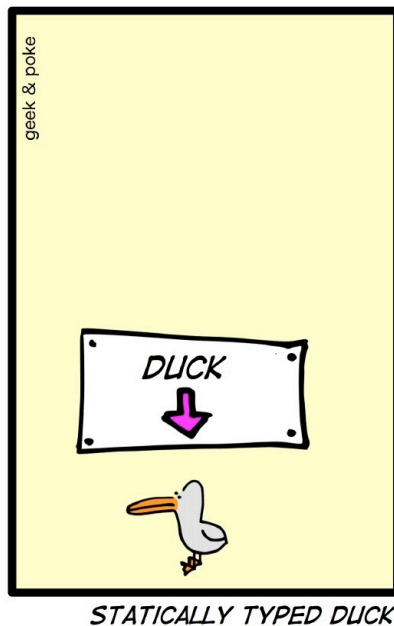
```
    =    A    <    >();  
    .    ();
```

← „MissingMethodException“ at run-time

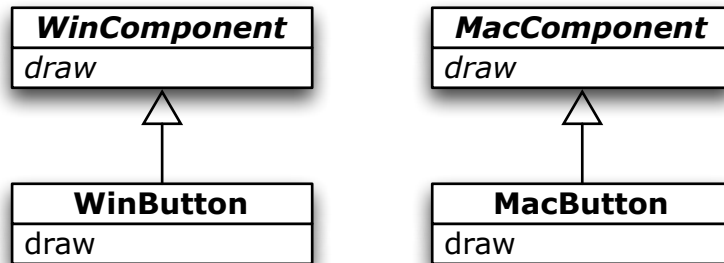
- **Dynamic type checking \neq weak type checking**

- Static type checking \Leftrightarrow dynamic binding
 - Static type checking ensures that a message is understood by an object
 - Dynamic binding selects the appropriate method based on the dynamic type of the receiver
- Static type checking \Leftrightarrow inheritance
 - An object of type T' can be assigned to a variable of static type T where T' is derived from T
 - Assignment compatibility is restricted by type hierarchy
 - Dynamic binding is only possible along the type hierarchy

- Dynamic type checking (dynamic typing)
 - Allows for dynamic binding outside of the type hierarchy
 - Criterion: does an object understand a message?
 - Smalltalk developers: „True Dynamic Binding“
 - Groovy developers: “Duck Typing”
 - Looks like a duck, walks like a duck, quacks like a duck → D ck



- Static type checking

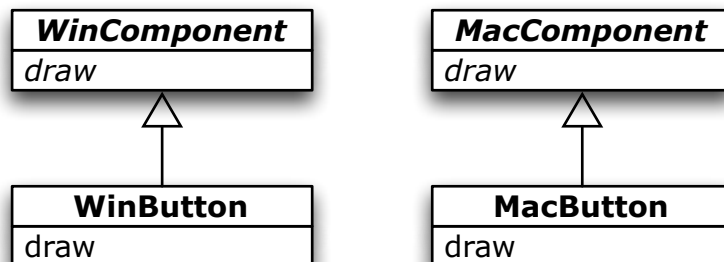


```

C      ;
C      ;
B      =      B      ();

=      ; //
=      ; //C      !
    
```

- Dynamic type checking

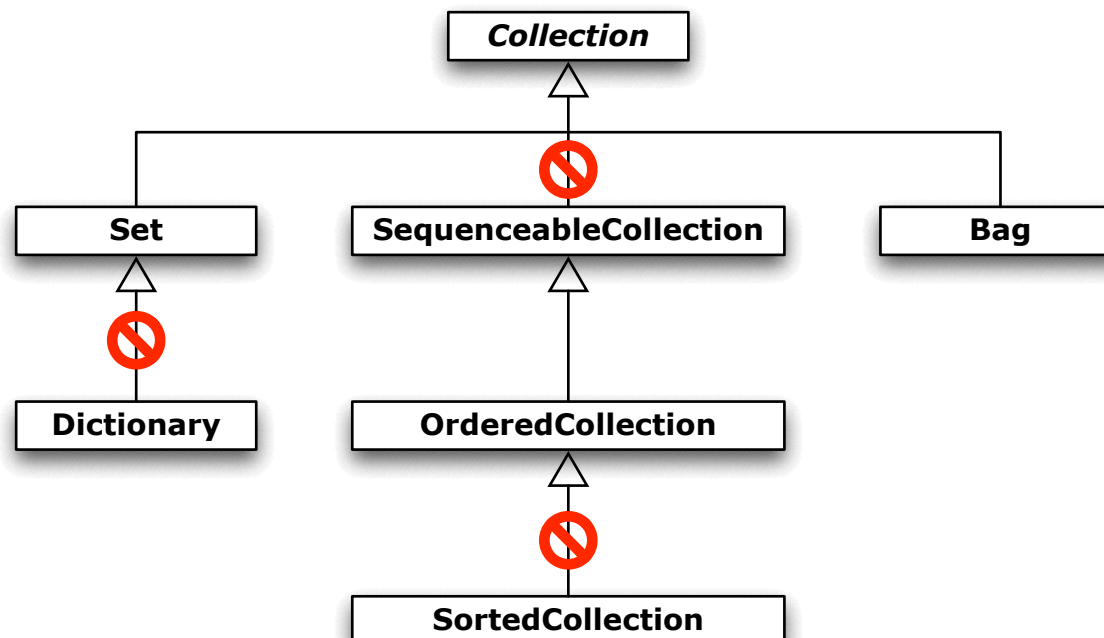


```

      =      B      ()

=      //
=      //
.      () // B      .      ()
    
```

- Subtype: T' is a subtype of T if T -objects can be replaced by T' -objects
„Principle of Substitutability“, Is-A-Relationship
- Subclass: T' is a subclass of T if T' is derived from T but T -objects cannot be replaced by T' -objects
- Example from the Smalltalk class library:



Principle of Substitutability – Requirements

- Subtype must provide at least the interface of its base type
 - Restrictions are not permitted!
- Pre- and postconditions
 - Preconditions must not be narrowed by a subtype
 - Postconditions must not be widened by a subtype
- Exceptions
 - Subtypes must not throw new exceptions
 - *Unless*: exception is a subtype of the exception thrown by the base type

- Interface inheritance:
 - Enables polymorphism
 - Modeling conceptual relationship among types
 - Implementation inheritance:
 - Reusing the functionality of the base class
 - Code sharing
 - Interface inheritance → subtype
Implementation inheritance → „only“ subclass
- Discrimination between subtype and subclass is essential in software design!

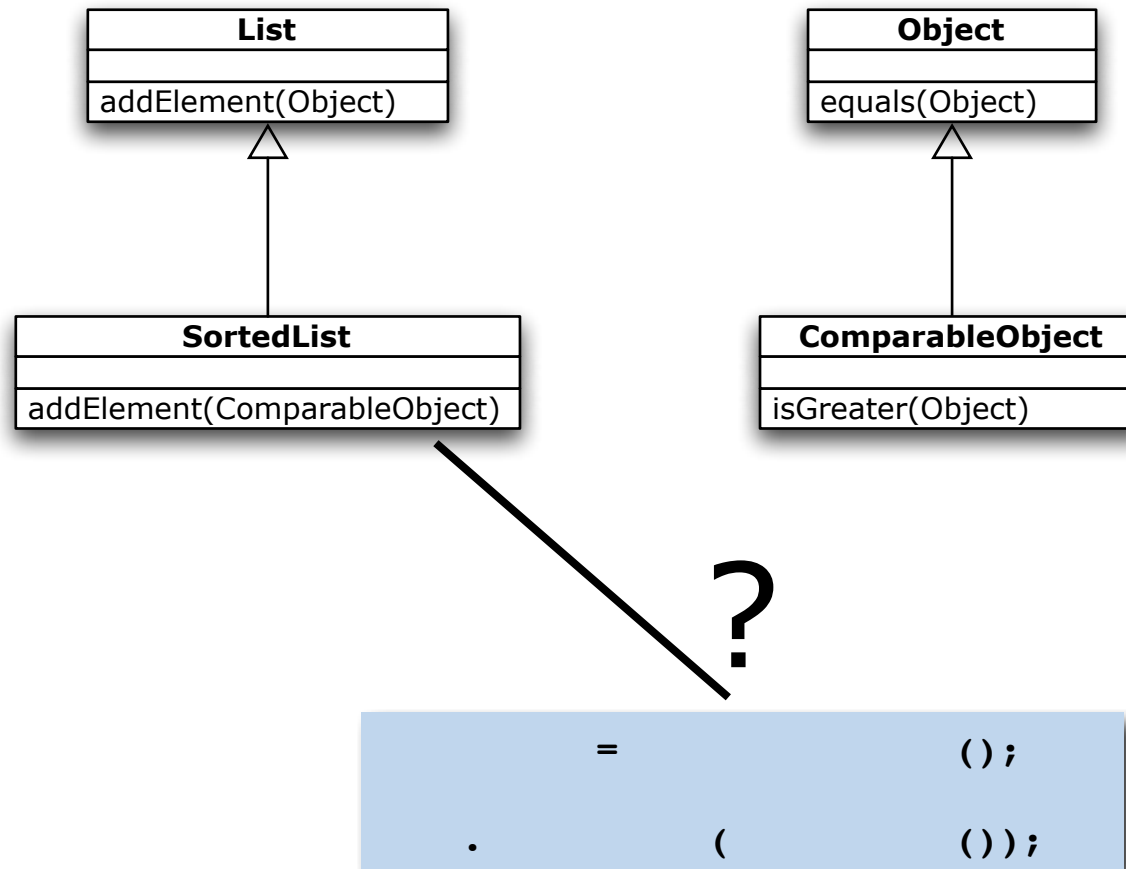
- Design goal:

Minimize implementation inheritance

→ promotes comprehensiveness/maintainability of software

- Most object-oriented programming languages provide language features to differentiate between interface and implementation inheritance
- Examples: C++, Java, Eiffel

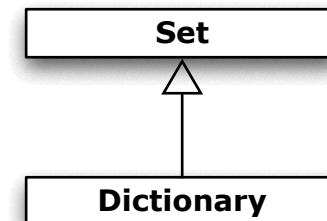
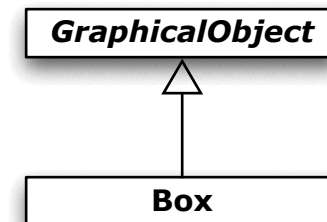
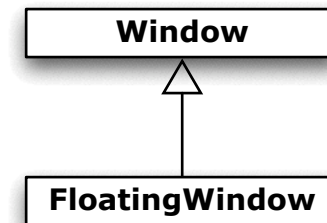
Covariance and Contravariance



- Is the principle of substitutability applicable to method parameters in overriding methods?
- *Covariance*: Type of a method parameter is narrowed in a subclass
 - Principle of substitutability applied to method arguments
 - Useful, but expensive to be checked statically!
- *Contravariance*: Type of a method parameter is widened in a subclass (or stays the same).
- Most object-oriented languages only support contravariance
 - Exception: Eiffel

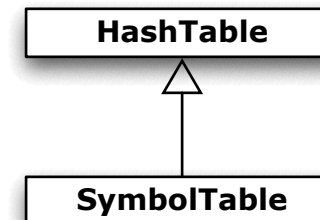
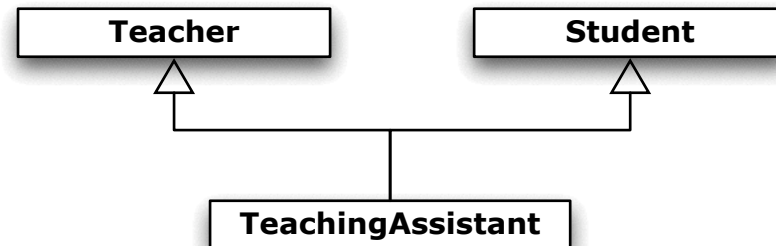
Types of Inheritance

- Specialization
 - Derive concrete subclass from concrete base class
- Specification
 - Concrete subclass implements abstract base class
- Restriction
 - Subclass restricts the base class interface



Types of Inheritance

- Combination
 - Multiple inheritance
- Construction
 - Implementation inheritance

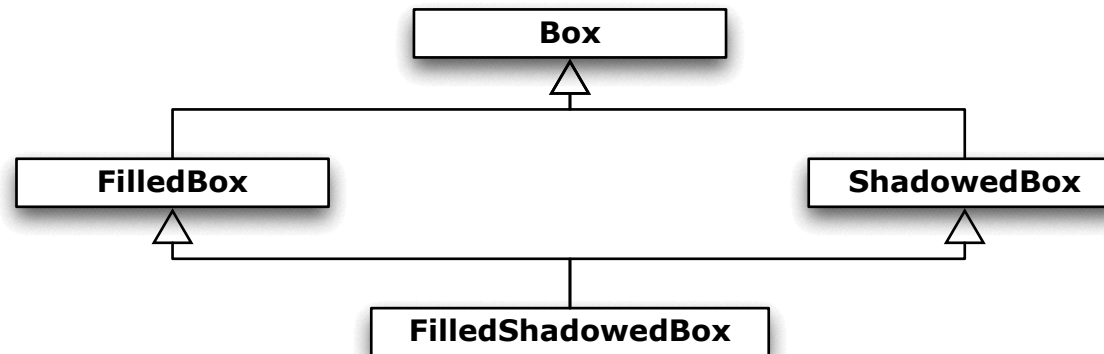




Multiple Inheritance



- Multiple inheritance from the same class? (Fork-Join)



- C++: FilledShadowedBox has two Box objects
- Virtual inheritance from Box:
 - A single Box object will be shared by all derived classes

```
B ;
  B :
    B :
      B :
        B :
          B ;
            B ;
              B ;
                B ;
```

- Virtual inheritance is conceptually „expensive“
 - → Requires assumptions on future derived classes

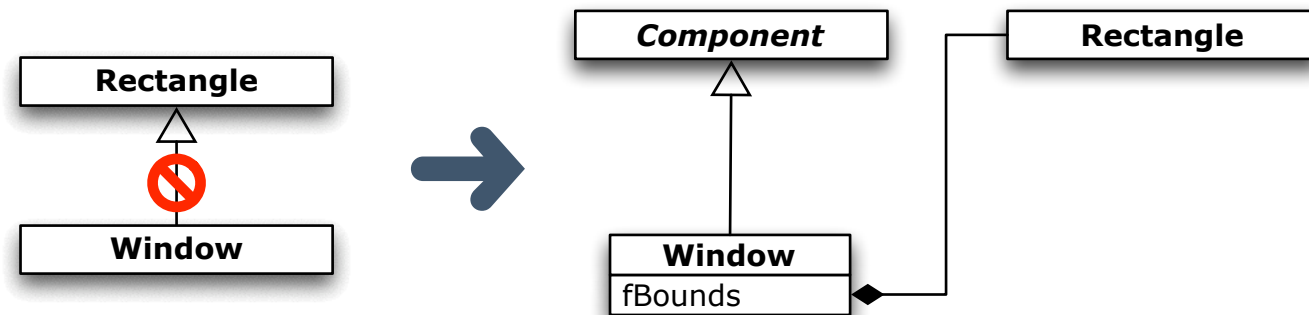


Restricting the Inheritance Interface

- Prevent method overriding
 - Java: keyword “final”
 - Turns off dynamic binding
 - Enables optimizations by the compiler or run-time system
 - Excludes alternative implementations
 - For example: export restrictions in cryptographic libraries
- Prevent derived classes
 - Java: keyword “final”
 - All methods are final in a final class
 - Excludes alternative implementations
 - Prohibit mutable derivations of immutable classes

Inheritance vs. Composition

- Inheritance
 - Protocol is extended (restriction contradicts subtyping)
 - Subtyping (is-a relationship)
 - Conforms to principle of substitutability!
 - Protocols of base class and subclass are very similar
- Composition
 - Little compliance between base class and subclass
 - Base class is just a minor detail of the subclass
 - Base class is a means of implementation (has-a relationship)
 - Implementation is substitutable



Design for Inheritance

- Inheritance violates encapsulation
 - Subclass needs to know implementation details
 - Judiciously grant access to subclasses by using protected access
 - By default members should be private
- Classes have to rely on the sanity of base classes and subclasses
- Constructors must not invoke overridable methods
- Test inheritability by writing subclasses!
- Prohibit subclassing if class cannot be safely subclassed
 - → final
 - Immutable class members cannot be subclassed

- „Class“ which declares only abstract methods
 - Interface defines a type, most abstract type specification
 - No state allowed (except constants)
- Class can inherit from 0-1 classes and implement 0-n interfaces

```
( , );
```

```
B ( B , ) ;
```

```
( ) ;
```

- Use interfaces as static type declarations whenever possible!

- Tagging (or marker) interfaces
 - Empty interface
 - Purpose:
 - Marking of a class to be eligible for being processed by some other class
 - Default implementation
 - Examples
 - `java.lang.Cloneable`
 - Instances (of classes which are marked as `Cloneable`) may be copied
 - `java.io.Serializable`
 - Instances may be serialized
- Constant interface
 - JDK example: `java.io.ObjectStreamConstants`
 - Implementation detail becomes part of the class's exported API
 - Since Java 5: use static import instead

Abstract Classes

- Purpose:
 - Define common structural and functional properties
 - Do not allow for instance creation
- Factorization
 - Common properties of subclasses are defined in an abstract class
 - Inheritance allows to adapt resp. implement these properties
 - Implementation of a system can be based on abstract classes
- Ease the implementation of concrete classes
 - Several concrete methods can be based on a few abstract methods
 - Example: Comparison of objects
 - All comparison methods can be based on "=" and ">"
 - Small inheritance interface, broad method interface
- Facilitates comprehensibility of a class hierarchy

- Abstract method
 - Only the signature is defined
 - Implementation is left to derived classes
 - Abstract method implies an abstract class
- Abstract class
 - Explicitly defined as an abstract class
 - Can contain 0-n abstract methods

$$C \left(\frac{C}{C} \right) = i$$

Abstract Classes and Static Type Checking

- Smalltalk: Abstract classes are based on conventions (self subclassResponsibility)
- Java: Abstract classes are supported by the language

```
      B      ;  
      (      ) ;  
B      (      ) B = ;  
      B      () B ;  
  
=      ();
```



Compile time error

- Principle of Substitutability
 - *Static type checking*: Common (abstract) base class resp. interface
→ documentation of design
 - *Dynamic type checking*: compatible method signature

- Interfaces should be as small as possible
 - All methods of an interface are abstract
 - Implementing class must implement all methods of its interfaces
 - Small interfaces ease implementability and substitutability
- Combining abstract classes and interfaces
 - Abstract class “implements” interface
 - Methods are not abstract but implementation is empty
 - Client may:
 - Implement interface → forced to implement all methods
 - Inherit from abstract class → override only required methods

Abstract Classes and Interfaces



- Example from the Java class library

```
C    (    );  
      (    );  
      (    );  
      (    );  
      (    );
```

```
C    A  
      (    ) ;  
      (    ) ;  
      (    ) ;  
      (    ) ;  
      (    ) ;
```



- Goal: abstracting from the implementation
 - Minimal coupling between classes
 - Implementation can be replaced
- Restricted access rights:
 - Discrimination between clients and inheritors
 - Java: public vs. protected
 - Classes within a subsystem
 - Java: public vs. package private
 - C++: friends
- Accessors hide state implementations (representation of state)
 - Advantages:
 - Minimized coupling, implementation can be replaced
 - Disadvantages
 - More verbose and inefficient (optimizations as in Java are feasible)

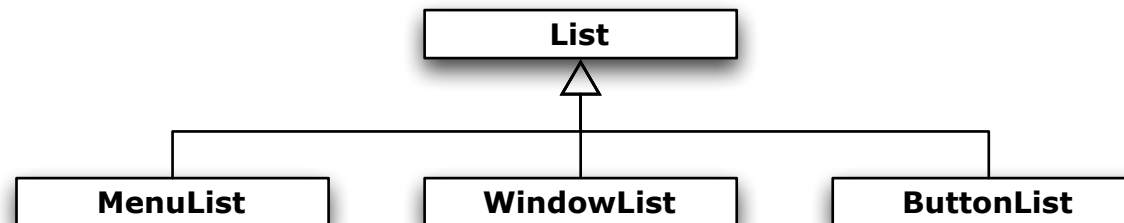
Inheritance and Information Hiding

- Two kinds of interfaces: clients and inheritors
 - What access rights are provided?
- Open-closed principle (B. Meyer)
 - Open: Class can be extended / extended in derived class
→ inheritors
 - Closed: Class features a stable interface and is provided in a library.
Class is reused by → clients
- Two flavors:
 - ① Inheritors have unrestricted access to class implementation
→ Inheritors are affected by changes to the implementation
 - ② Inheritors have no access to class implementation
→ Inheritors are only affected by changes in the interface
- Smalltalk: flavor ①
- Java: fine granularity with access rights

Immutability

- How to define an immutable class (in Java)?
 - No mutators or copy on write
 - Copy on write: all mutating operations work on a copy and return it
 - Class must not be extensible
 - Use public final class ...
 - All fields must be final
 - Prohibit changes after initialization
 - All fields must be private
 - Prohibit changes to final fields, e.g. non-zero arrays
 - Exclusive access to mutable components
 - References to mutable objects are not available to other clients

- What if implementations of a class only differ by type?



- Violates principle of substitutability!
- Implementation based on type “Object” is questionable
 - Frequent type conversions necessary
 - Error prone

```
//          =    A    ();
//          .    (    B    ());
//          .    (    ());
//
B    1 = (B    )    .    (0);
B    2 = (B    )    .    (1); //->
```


- Class specifies formal type parameters
 - Actual type parameters are provided upon instantiation / inheritance
 - Programming language feature
- Programming languages with generic classes:
 - Eiffel, C++
 - In Java since Java 5

```

A      < >      < >
      (      )      ;
      (      )      ;
    
```

```

// <B      >      =      A      <B      >();
//      .      (      B      ());
//      .      (      ()); //->      -
// B      1 =      .      (0);
//      2 =      .      (1); //->      -
    
```

- Java: generic classes are compatible with non-generic classes
 - *Example:* Collection classes can be instantiated without actual type parameter
 - “Object” is assumed upon missing type parameter
- Base types are not allowed for actual type parameters
 - Hardly restricting with since autoboxing / unboxing
- Java reflection (meta information) was extended as well
- Java also provides generic methods
- Generic class have no impact on the JVM
 - Compiler generates compatible byte code
 - Backward compatibility leads to restrictions

- Advantages of generic classes
 - Reliability (avoiding down-casts)
 - Readability

```
//      =      A      ();  
<B      >      =      A      <B      >();
```

- Autoboxing / -unboxing complements nicely

```
//      <      >      /-      =      A      <      >();  
      .      (      (2));  
      =      .      (0).      ();  
  
//      .      (2);      /-  
      =      .      (1);
```

- Disadvantage of generic classes
 - Language complexity (syntax and above all semantics)

- Inheritance is more powerful than genericity
- Combining inheritance and generics is desirable
 - Flexibility and reliability
 - Even more complex syntax and semantics
 - Is `List<String>` a subtype of `List<Object>`? (→ covariance)
- Restricting genericity
 - Actual type parameter must conform to a specific type

```
        <          >  
      ( (      ) )  
    (      ) ; ;
```

```
<    >      =      <    >  
<    >      =      <    > //      !!
```

- Design of a class API
 - Flexibility and type safety
 - Readability
 - Expressive and self documenting API
 - Client should not care about typing challenges

```
class A (< >) {  
    :  
}
```

hardly usable

```
class A (<?>) {  
    :  
}
```

flexible

Inheritance and Generic Classes



- Case study: API design of a generic class "Stack"

```

        < >
        ();
        ( );
        ();
        ();
    
```

```

//A
        (
        A (
        < > )
        ( : ) ( );
    
```

Compiler ✓

```

//
        < > = < > ();
        < > = ...;
        . A ( );
    
```

Compiler ✗

```

//A
        (
        A (
        <? > )
        ( : ) ( );
    
```



Inheritance and Generic Classes



```
//A
    A (C
    (!    () )
    .    (    () );
```

Compiler ✓

```
//
    <    >    =    <    >();
c    <    >    = ...;
    .    A    (    );
```

Compiler ✗

```
//A
    A (C
    (!    () )
    .    (    () );
```



- When to use *extends* resp. *super*?
 - **PECS**: **p**roducer **e**xtends, **c**onsumer **s**uper
 - Stack example: src = producer, dst = consumer

		Parameter produces T instances	
Parameter consumes T instances		Yes	No
	Yes	Type<T> (<i>Invariant</i> in T)	Type<? super T> (<i>Contravariant</i> in T)
	No	Type<? extends T> (<i>Covariant</i> in T)	Type<?> (<i>Independent</i> of T)

Meta-Information

- Information on properties of objects, classes, and the inheritance hierarchy
- Used for:
 - Basic functionality, such as persisting objects
 - Tools (e.g. debugger)
 - Type checking at run-time
 - Java provides meta-information at run-time (reflection)

: #

//

. c . (,)

- Meta classes
 - Usually provided by pure oo languages
 - A class is an object and hence instance of a class (= meta class)
 - Meta classes provide all meta-information
 - Meta-information can be modified

- Dynamic class loading
- Purpose:
 - Reducing dependencies
 - Multiple implementations of a type are available *but*
 - Only one is needed at run-time *or*
 - Only one is executable in a specific run-time environment

```
(  
=  
= c . (  
...  
);  
)
```

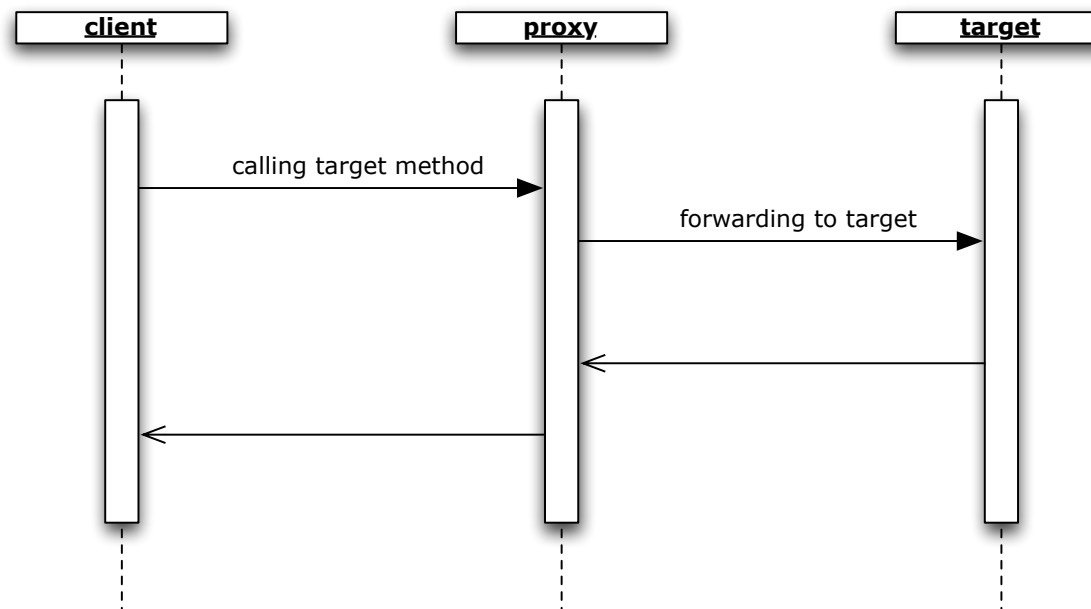
- Class "Class" provides access to meta-information at run-time
- Example: calling a private method from outside its declaring class

```
        (                )                ...  
        =    A    (    ) ;  
C    .    (    (    ) ) ;  
        =    .    C    (    ) ;  
        =    .    (    ,    .    ) ;  
        .    A    (    ) ;  
        .    (    ,    0 ) ;
```

- Limitations of object-oriented programming
 - Factorization is only possible:
 - along the class hierarchy
 - by delegation
 - Retrospective changes are hardly possible
- Cross-cutting concerns
 - Identical functionality spread across a software system
 - Example: database access
 - Functionality: caching, authorization, transaction, logging
 - Ramifications:
 - Duplicated code
 - Mixing domain logic and technical aspects
 - Decreased maintainability: complex and replicated code
 - Reusability questionable: technical aspects are platform dependent
- Aspect-Oriented Programming is an add-on to rather than a replacement of OO!

- AOP: dividing a problem into aspects rather than objects/classes
- Join-point model
 - Join Point
 - Point in a running program where behaviour can be added
 - Examples: method call, access to instance variable etc.
 - Advice:
 - Defines the behaviour (before, after, and around)
 - Point Cut:
 - Quantify join points for executing an advice
 - Point cuts can be defined via pattern matching
- AspectJ: aspects as an extension to Java for AOP
 - Defines point cuts and features advice implementations
 - Can add new instance and class variables / methods
 - Corresponds to a class in object-oriented programming
 - Since Java 5 aspects can be specified using annotations

- Application of reflection
- Enables simple aspect-oriented programming
- Task:
 - Adding aspects to classes at run-time
 - No source code changes to these classes
 - No inheritance (base class cannot be changed in retrospect)



- Class
 - Facilitates creating classes at run-time
 - Classes may implement several interfaces

```

        c
        c
        (c
        , c
        )
        , c
        )
        c
        (c
        )
        (
        )
    
```

```

        (
        ,
        ,
        )
    
```

- Example: logging aspect
 - InvocationHandler logs before and after a method call

```
        .getInvocationHandler()
        .invoke(this, method, args);

        // Log the method call
        System.out.println("Method " + method.getName() +
            " called with arguments " + args);

        // Log the return value
        System.out.println("Return value: " + result);

        return result;
    }

    // Log the method call
    System.out.println("Method " + method.getName() +
        " called with arguments " + args);

    // Log the return value
    System.out.println("Return value: " + result);

    return result;
}
```



```
(
    ,
    )

    =
    ;

    ( c
    +
    .
    ());

    =
    .
    (
    ,
    );

    (
    .
    (
    .
    () +
    +
    .
    c
    ());

    .
    c
    ();

    .
    (
    +
    .
    ());

    ;
```

- Example of usage

```
= ( )  
    ();  
    . ( ),
```

AOP – Google Guice

- Google Guice (pronounced “juice”)
 - Lightweight dependency injection framework
- Dependency Injection
 - Decoupling and configuration of dependencies
- Google Guice provides some infrastructure for AOP

AOP – Google Guice

- Approach
 - Module
 - Definition of advices
 - Objects must be created by injector
 - Interceptor specifies advice
 - Annotations can be used to define join point

```
@      (      .      ) @      (      .      )  
@
```

```
      =      ();  
  
      (  
      .      ("B      " +      .      ().      ());  
      =      .      ();  
      .      ("      " +      .      ().      ());  
      ;
```

```

A
    ()
(      .   (), //           ?
        (      .   ), //       ?
          ()); //
=
(      )
=      .   (      ());
      =      (      .   );
      .   ();

```

AOP – Google Guice

- Internals
 - Google Guice dynamically creates subclasses
 - Subclass overrides method and invokes interceptor
- Appraisal
 - Google Guice offers simple AOP for common use cases
 - Lightweight solution
 - Simpler but not as powerful as comprehensive AOP approaches, e.g. Spring
 - Way more usable than dynamic proxies (and more abstract, too)
 - Exemplary API
 - Flexibility and expressivity
 - Creating objects via injector can be cumbersome
 - Viral impact
 - Adding AOP retrospectively is hardly possible
 - Even if using patterns rather than annotations for point cuts
 - Debugging
 - Noisy run-time stack



Polymorphism - Java Examples

- Parametric polymorphism → generic classes
- Inheritance → chapter “Advanced Principles” and “Class Libraries”
- Coercion → e.g. autoboxing/-unboxing
- Overloading

```
        (      ) ...  
        (      ) ...  
  
//...  
  
        =        ();  
  
    .    (256);  
    .    (      );
```


Literature

- Concepts
 - H. Züllighoven: Das objektorientierte Konstruktionshandbuch
 - B. Meyer: Object Oriented Software Construction
- Reflection in Java
 - I. Forman, N. Forman: Java Reflection in Action
- Java
 - Java Generics: <http://java.sun.com/j2se/1.5/pdf/generics-tutorial.pdf>
 - J. Bloch: Effective Java
- Values
 - D. Bäumer et al: Values in Object Systems
- Aspect-oriented programming
 - G. Kiczales et al: Aspect-Oriented Programming
 - AspectJ: <http://www.eclipse.org/aspectj/>
 - Google Guice: <http://code.google.com/p/google-guice/>