



Lecture 4

P2P with TomP2P

<http://tomp2p.net/doc>



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0. Lecture Overview

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2. Tracker / PEX with examples

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6. References

- **5.3.2015 - Newest version of uTorrent has Bitcoin mining offer during install**
 - ▶ Silently installs EpicScale
 - ▶ EpicScale is a bitcoin miner
 - ▶ Not easy to uninstall, many angry users
- **6.3.2015 - BTCPoint Creates 10,000 Bitcoin-Enabled ATMs Using Spanish Bank Network**
 - ▶ Bitcoin-to-cash withdrawals at 10,000 additional bank ATMs
 - ▶ 0% commission → they build ATMs
- **3.3.2015 - Evolution of the Internet: from Decentralized to Centralized**
 - ▶ Choosing convenience over decentralization
 - ▶ Try to make your CT prototype convenient to use!

- **3.3.2015 - Sync 2.0: Skip The Cloud, Share Direct**
 - ▶ New features compared to 1.x
 - ▶ Pro version for business workgroups
 - ▶ Available on NAS devices
- **18.2.2015 - MegaNET – New Decentralized, Non-IP Based and Encrypted Network**
 - ▶ Kim Dotcom wants to start own Internet using the blockchain
 - convicted of several crimes
 - suspended prison sentence in 1994 for computer fraud and data espionage / suspended prison sentence in 2003 for insider trading
 - 2012: Megaupload
 - ▶ Series of tweets...
 - ▶ MaidSafe – decentralized data storage

1. Introduction

What is TomP2P
History and project information

TomP2P

A P2P-based high performance key-value pair storage library

- **TomP2P is a P2P framework/library**
 - Implements DHT (structured), broadcasts (unstructured), direct messages (can implement super-peers)
 - NAT handling: UPNP, NATPMP, new addition: relays, hole punching
 - Direct / indirect (tracker / mesh) storage
 - Direct / indirect replication (churn prediction and ~rsync)
 - Modes: key,value / multi-key (versioned) value
 - Java 6, Maven, [Github](#), Netty, TCP/UDP, MapDB, (Android)

TomP2P

A P2P-based high performance key-value pair storage library

- **TomP2P extends DHT**
 - Distributed hash table concept → put(key,value) / get(key)
 - Extended DHT operations →
 - put(key1,key2,value)
 - put prepare / put confirm
 - add(key, value)
 - digest(key) / bloomfilters / versions
 - get(key) / bloomfilters / versions

- **TomP2P history**

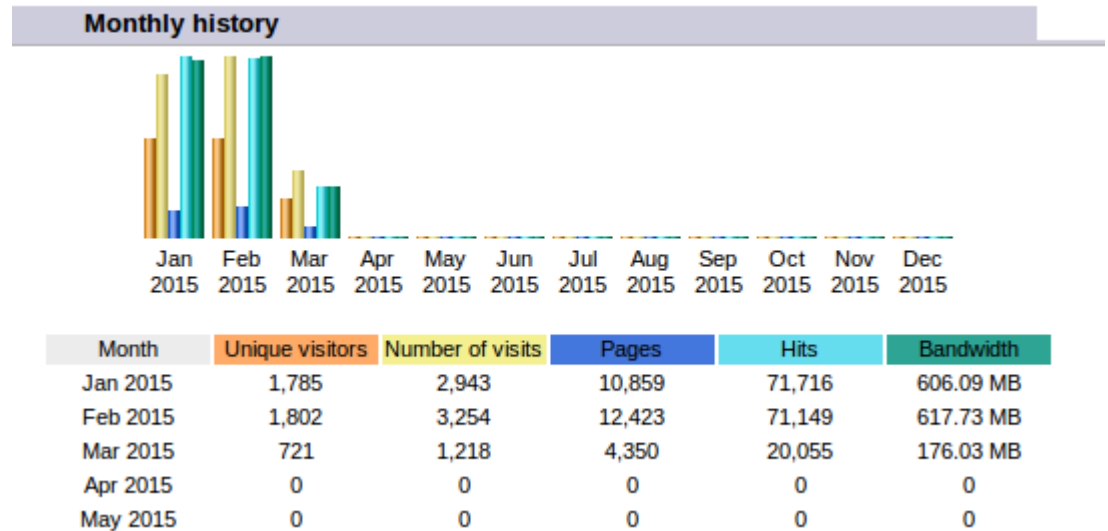
- ▶ TomP2P v1: Created in 2004 and used for a distributed DNS project
 - ▶ This version used blocking IO operations (1 thread / socket)
- ▶ TomP2P v2: Apache MINA (java.nio framework) / 6K LoC
 - ▶ Not well designed for non-blocking operations (event-driven)
- ▶ TomP2P v3: Redesigned for non-blocking operations
 - ▶ Switched to Netty / 14K LoC, 6K LoC JUnits
- ▶ TomP2P v4: API refinements, new features
 - ▶ Current release (preview) 4.1
 - ▶ Latest feature (work in progress) MapReduce
 - ▶ 22K LoC, 8K LoC Junits
- ▶ TomP2P v5 (core 18K LoC): modularization, relays, API refinements

Introduction

TomP2P

A P2P-based high performance key-value pair storage library

- **TomP2P started in 24.05.2004 and still going strong**
 - ▶ [Github](#): watch 30 / star 131 / fork 53 (not sure if good/bad)
 - ▶ TomP2P website (although documentation is outdated)
 - ▶ 72 users on the mailinglist
 - ▶ [9 contributors](#) on github
 - ▶ Don't buy the book!



TomP2P

A P2P-based high performance key-value pair storage library

- **Academic background (CSG - UZH):**
 - ▶ Used in EU projects: EC-GIN, Emanics, SmoothIT, SmartenIT, Flamingo
 - ▶ Used in research projects: [LiveShift](#), [DRFS](#), [Radiommender](#), [Box2Box](#), [Hive2Hive](#), [B-Tracker](#), [PiCsMu](#), [peerwasp](#), (and non-academic)
- **<http://tomp2p.net>**
 - ▶ For questions: mailinglist (<http://lists.tomp2p.net/cgi-bin/mailman/listinfo>)
 - ▶ Specific questions: bocek -at- ifi.uzh.ch or tom -at- tomp2p.net
 - ▶ Documentation: <http://tomp2p.net/doc/> (TomP2P v4.4)
Overview: <http://en.wikipedia.org/wiki/TomP2P>
 - If something is missing, ask! – **Documentation for v5 is missing!**
 - ▶ Development: <https://github.com/tomp2p>
 - Feature request possible if good reasons provided
- **(Demo: how to setup TomP2P with IntelliJ/git/maven)**

TomP2P

A P2P-based high performance key-value pair storage library

- [A Declarative Interface for Smart-phone Based Sensor Network Systems](#), Asanka Sayakkara and Kasun De Zoysa, [IWMS 2012](#), Beijing, China
- [Hybrid Peer-to-Peer DNS](#), Ricardo Sancho and Ricardo Lopes Pereira, Instituto Superior Tecnico, Porto, Portugal
- [A Semantic Publish-Subscribe Coordination Framework for IHE based Cross-Community Health Record Exchange](#), Visara Urovi, Alex C. Olivieri, Stefano Bromuri, Nicoletta Fornara, Michael Schumacher, [ACM SIGAPP](#) Applied Computing Review, 2013 ([slides](#))

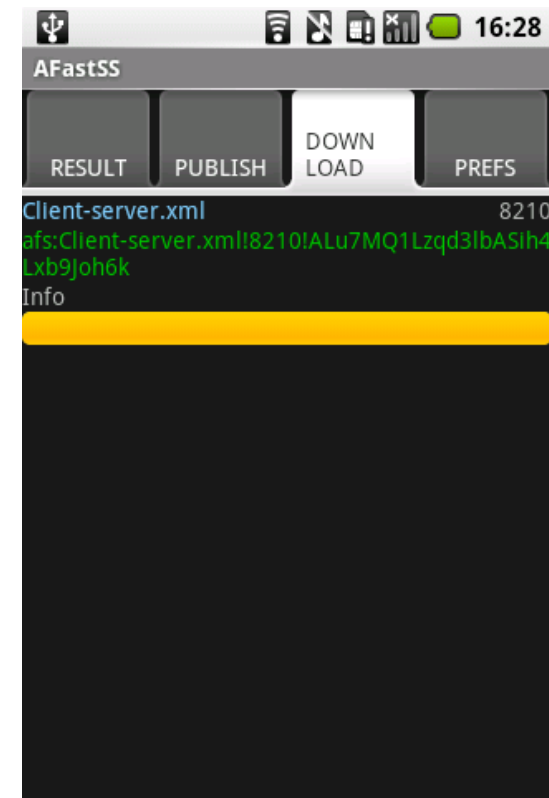
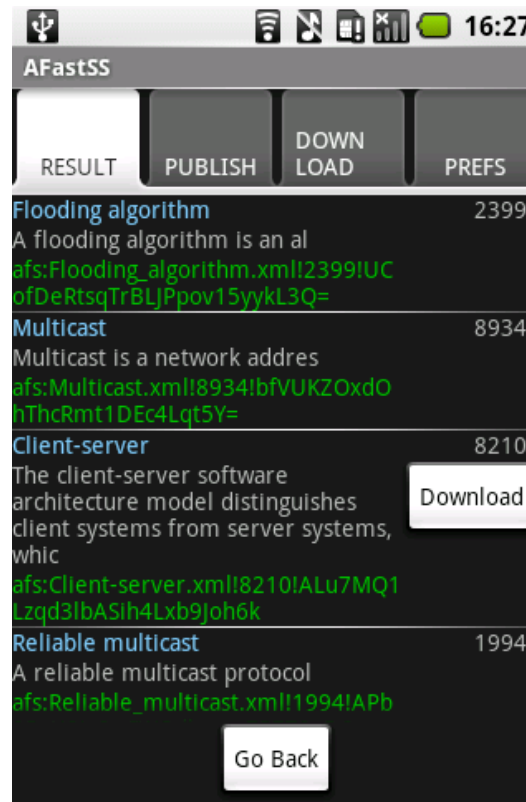
- [Adding Cryptographically Enforced Permissions to Fully Decentralized File Systems](#) – TUM, Bernhard Amann, Thomas Fuhrmann, April 2011
- [Optimis FP7 IP project](#): Optimized Infrastructure Services, D1.2.1.3, Architecture Design document, ended May 2013
- [Distributed Name-based Entity Search](#), Fausto Giunchiglia and Alethia Hume, [ISWC 2012](#), Boston
- [A Distributed Directory System](#), Fausto Giunchiglia and Alethia Hume, SSWS 2013, Sydney

- [A P2P Semantic Query Framework for the Internet of Things](#), Richard Mietz, Sven Groppe, Oliver Kleine, Daniel Bimschas, Stefan Fischer, Kay Römer and Dennis Pfisterer, PIK, Volume 36, Issue 2 (May 2013)
- [P2P Minecraft](#): “The mods described below are about adding peer-to-peer functionalities to Minecraft.”
- [Bitcoin Gateway - A Peer-to-peer Bitcoin Vault and Payment Network](#), Omar Syed & Aamir Syed, [July 2011](#)
- [Bitsquare.io](#) – The decentralized bitcoin exchange, ~ongoing

Introduction

- TomP2P with Android (early research)

- ▶ CSG: early adopter
- ▶ TomP2P 5 and Android: work in progress



2. Example

Example and Demo

Example

- Demo: a simple put / get example
- Package net.tomp2p.examples. ExamplePutGet

```
private static void examplePutGet(final PeerDHT[] peers, final Number160 nr)
    throws IOException, ClassNotFoundException {
    FuturePut futurePut = peers[PEER_NR_1].put(nr).data(new Data("hallo")).start();
    futurePut.awaitUninterruptibly();
    System.out.println("peer " + PEER_NR_1 + " stored [key: " + nr + ", value: \"hallo\"]");
    FutureGet futureGet = peers[PEER_NR_2].get(nr).start();
    futureGet.awaitUninterruptibly();
    System.out.println("peer " + PEER_NR_2 + " got: \"" + futureGet.data().object() + "\" for the key " + nr);
    // the output should look like this:
    // peer 30 stored [key: 0xba419d350dfe8af7aee7bbe10c45c0284f083ce4, value: "hallo"]
    // peer 77 got: "hallo" for the key 0xba419d350dfe8af7aee7bbe10c45c0284f083ce4
}
```

- Defaults
 - ▶ Replication factor 6, replication not enabled,
 - ▶ domain, content, version are zero if not specified

3. Fundamental Concepts

XOR-based iterative routing

Futures

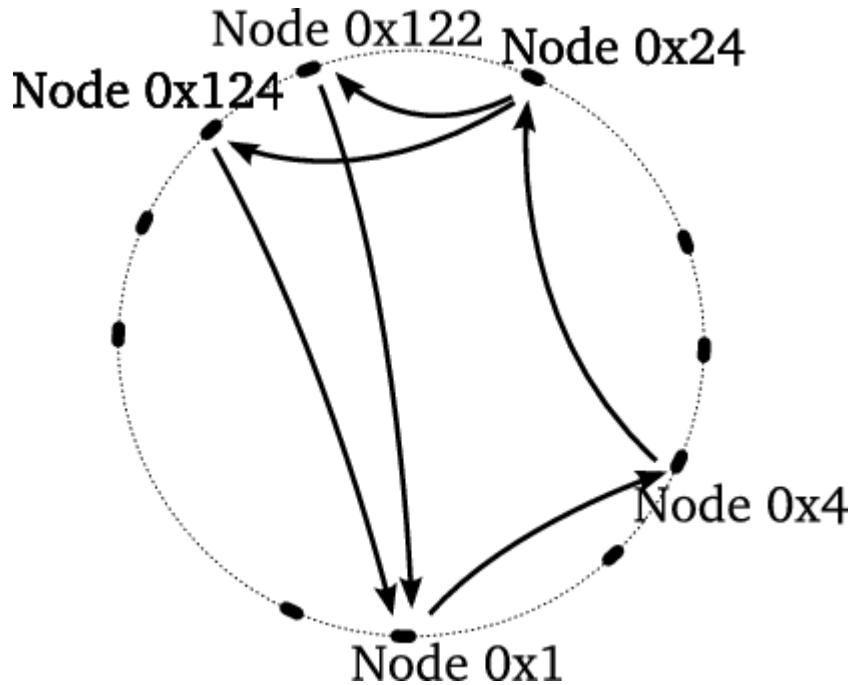
API Overview

Fundamental Concepts (repetition)

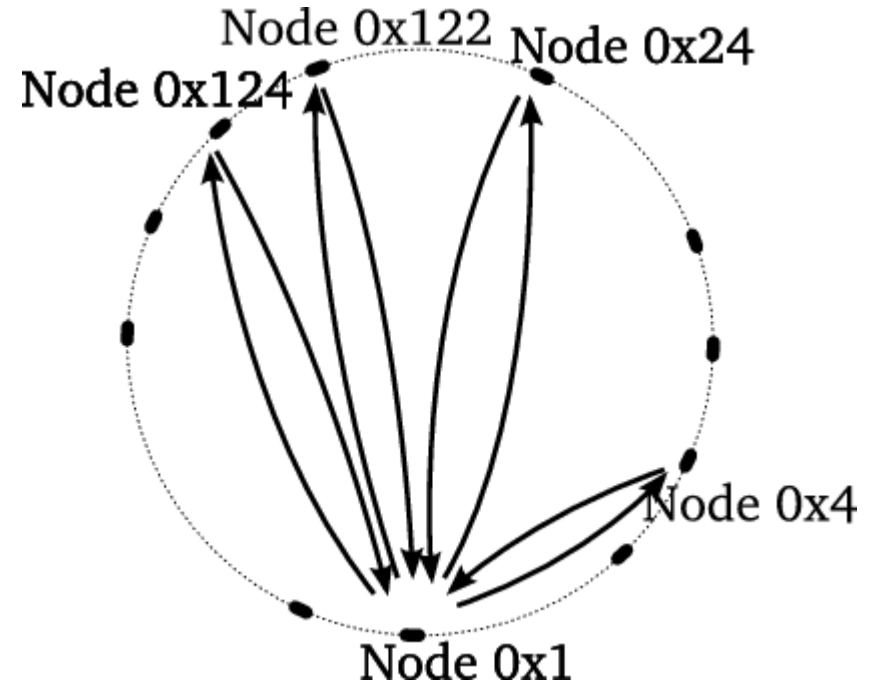
- Recursive routing

vs.

- iterative routing



- + online status update
- faulty peers cause delay



- + control
- neighbor maintenance

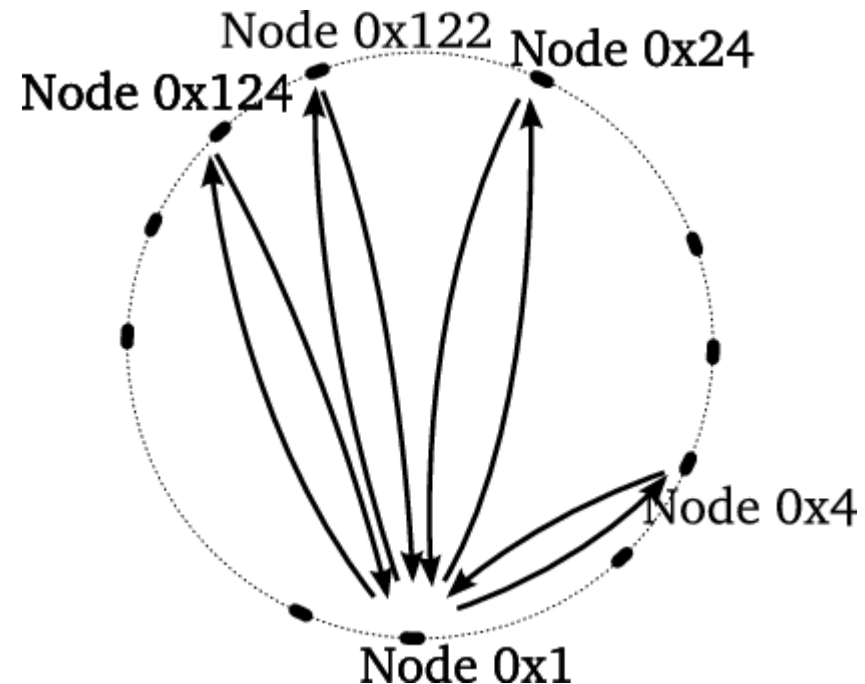
Fundamental Concepts (repetition)

- **TomP2P: iterative XOR-based routing**

- ▶ Node and data item unique 160bit identifier
- ▶ Keys are located on the nodes whose node ID is closest to the key
- ▶ Search for a key:
 - Lookup in neighbor table for closest peer (e.g. peers with ID: 0x1, 0x2, 0x3, 0x4)

My ID	Neighbor ID	Distance (XOR)
1	2	3
1	3	2
1	4	5

- Difference to Pastry: one metric, no leaf set / routing table



- **TomP2P iterative XOR-based routing**

- ▶ Neighbors stored in 159 “bags”, bag has capacity c (Kademlia, $c=20$)
- ▶ Routing takes $O(\log n) \rightarrow$ M03, slides 11
- ▶ By default UDP, message header 58 bytes
- ▶ Routing Mechanism variables, can be tuned
 - ▶ directHits, potentialHits – routing sends digest
 - ▶ forceTCP – use TCP instead of UDP
 - ▶ maxDirectHits, maxNoNewInfo, maxSuccess, maxFailure – stop conditions
 - ▶ parallel – number of parallel connections
- ▶ For the CT - don't worry, default settings are fine 😊

Fundamental Concepts

- All distributed operations use futures
- Future objects
 - ▶ Keeps track of future events, while the “normal” program flow continues → `addListener()` or `await()`
 - ▶ `await()`: Operations are executed in same thread
 - ▶ `addListener()`: Operations are executed in same or other thread
- Demo: blocking operation (`net.tomp2p.examples.ExamplePutGet`)

```
private static void exampleGetBlocking(final PeerDHT[] peers, final Number160 nr)
    throws ClassNotFoundException, IOException {
    FutureGet futureGet = peers[PEER_NR_2].get(nr).start();
    // blocking operation
    futureGet.awaitUninterruptibly();
    System.out.println("result blocking: " + futureGet.data().object());
    System.out.println("this may *not* happen before printing the result");
}
```

Fundamental Concepts

- Demo: non - blocking operation (net.tomp2p.examples.ExamplePutGet)
 - ▶ New utilities necessary (loops as recursions)
 - ▶ Advise: use `addListener(...)` as much as possible!
 - ▶ `operationComplete(...)` must be **always** called ([problem if not](#))

```
private static void exampleGetNonBlocking(final PeerDHT[] peers, final Number160 nr) {
    FutureGet futureGet = peers[PEER_NR_2].get(nr).start();
    // non-blocking operation
    futureGet.addListener(new BaseFutureAdapter<FutureGet>() {
        @Override
        public void operationComplete(FutureGet future) throws Exception {
            System.out.println("result non-blocking: " + future.data().object());
        }
    });
    System.out.println("this may happen before printing the result");
}
```

- **Future utilities**

- ▶ `FutureForkJoin(int nr, boolean cancel, K... Forks)`
 - Joins already “forked” futures. Waits until all or `nr` future finished. If `nr` reached, futures may be cancelled (e.g. abort download)
- ▶ `FutureLateJoin(int nrMaxFutures, int minSuccess)`
`FutureLaterJoin()`
 - No need to add the futures in the constructor, can be added later
- ▶ `FutureDone()`
 - A generic future used in many places, can be placeholder

- **ForkJoin in Java7**

- ▶ Fork and join framework – future utilities in TomP2P focus on join, forking is done “manually”

- **Needs face-lifting, add `Futures.waitAny / waitAll` etc.**

- Fun with futures: loops

```
Future loop() {
    Future future = new Future();
    recLoop(future);
    return future;
}

void recLoop(Future future) {
    int active = 0;
    for (int i = 0; i < parallel; i++) {
        //if future finished, it will be set to null
        if (futureResponses[i] == null) {
            active++;
            futureResponses[i] = doSomething();
        }
        else if (futureResponses[i] != null) active++;
    }
    if (active == 0) future.weAreDone();
    FutureForkJoin<FutureResponse> fp = new FutureForkJoin<FutureResponse>(1, futureResponses);
    fp.addListener(new BaseFutureAdapter<FutureForkJoin<FutureResponse>>() {
        @Override
        public void operationComplete(FutureForkJoin<FutureResponse> future)
            throws Exception {
            boolean finished = evaluate(future);
            if(finished) future.weAreDone();
            else recLoop(future);
        }
    });
}
```


Fundamental Concepts

- **Java 8 lambda expressions not used**
- **.NET and other languages have better support for async**
 - ▶ Example (tbd)

Fundamental Concepts

- **API Overview: Peer.java**

- ▶ Core methods, network related

- `sendDirect()`
- `bootstrap()`
- `announceShutdown()`
- `ping()`
- `discover()`
- `broadcast()`

- ▶ Methods for DHTs (PeerDHT.java)

- `put(key, value),`
- `get(key)`
- `add()`
- `digest()`
- `remove()`
- `send()`
- `parallelRequest()` // mostly used internally



```
public class Number160 {
    // ...
    // Create an instance with an array of 160 integers
    // copied into the backing array
    // @param val the value to copy to the backing
    // array (160 integers)
    //
    public Number160(final int... val) {
        if (val.length > ARRAY_SIZE)
            throw new IllegalArgumentException(
                "Array size " + ARRAY_SIZE + " exceeded");
        this.val = new int[ARRAY_SIZE];
        for (int i = 0; i < val.length; i++)
            this.val[i] = val[i];
    }
}
```

• Extensions

▶ TomP2P can store multiple values for a key

- `put() (location_key, content_key, value) → content_key`
specified in Builder
- `get().setAll()`
→ returns a map with `[content_key, value]`
- `add() (location_key, value) → is translated to`
`put() (location_key, hash(value), value)`

▶ TomP2P support domains

- Avoid collision for same keys
- Domains are used for protection (more details later)
- Domains specified in Builder
- `put() (key, domain, value) → get() (key, domain)`

- **Configurations Example**

- ▶ Configuration with builder pattern

```
RoutingConfiguration rc = new RoutingConfiguration(0, 0, 1);
RequestP2PConfiguration pc = new RequestP2PConfiguration(1, 0, 0);
FuturePut fdht = peers[444].put(peers[30].peerID()).data(new Number160(5), data)
    .domainKey(Number160.createHash("test")).routingConfiguration(rc)
    .requestP2PConfiguration(pc).start();
```

- ▶ System-wide configuration when creating Peer

```
Number160 peerId = new Number160(rnd);
PeerMap peerMap = new PeerMap(new PeerMapConfiguration(peerId));
master = new PeerBuilder(peerId)
    .ports(port).enableMaintenance(maintenance)
    .bindings(bindings).peerMap(peerMap).start().addAutomaticFuture(automaticFuture);
peers[0] = new PeerBuilderDHT(master).start();
```

- ▶ ... and per module

5. Components with Examples

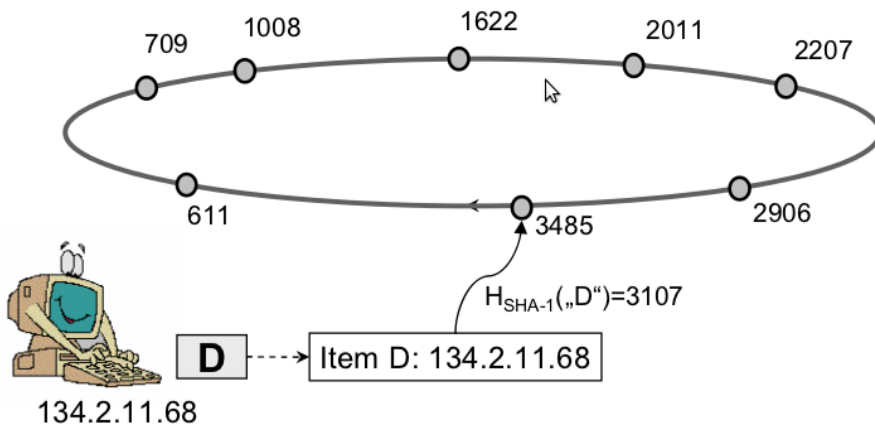
DHT
Tracker

Components with Examples (repetition)

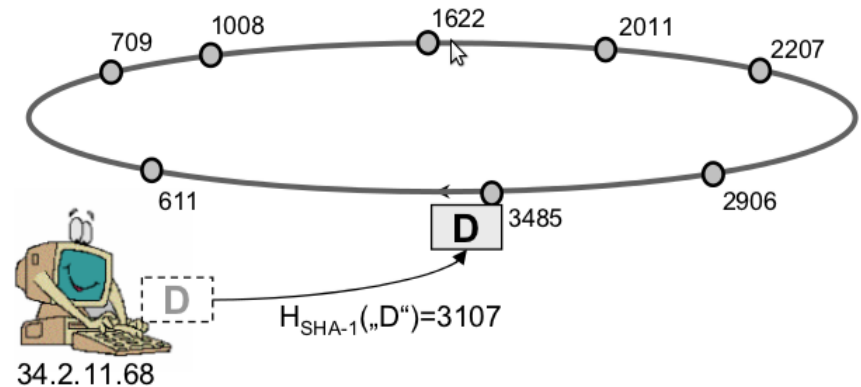
- **DHT vs. Tracker**

- ▶ M03, slide 23: DHT “stored by value” – direct storage
- ▶ M03, slide 24: Tracker “stored by reference” – indirect storage

indirect (Tracker)



direct (DHT)



- **B-Tracker**

- ▶ Centralized tracker – one peer gets traffic
- ▶ DHT: store reference on 20 peers – 20 peers gets traffic
- ▶ PEX: exchange information every minute (push)
- ▶ B-Tracker, every downloading peer becomes a tracker → forms mesh
 - Better balance of load
 - To avoid duplicates send compressed list of known peers
- ▶ B-Tracker in TomP2P enabled by default

Components with Examples

- **Demo: Tracker (net.tomp2p.examples.ExampleTracker)**
 - ▶ Create 100 peers,
 - ▶ Add to tracker, get from tracker
 - ▶ Stored on 3 peers: TrackerBuilder.java (can be configured)
 - ▶ Attachment of data is possible (attachement(Data))

```
private static void example(final PeerTracker[] peers) throws IOException, ClassNotFoundException {  
    FutureTracker futureTracker = peers[12].addTracker(Number160.createHash("song1")).start().awaitUninterruptibly();  
    System.out.println("added myself to the tracker with location [song1]: "+futureTracker.isSuccess()+" I'm: "+peers[12].peerAddress());  
  
    FutureTracker futureTracker2 = peers[24].getTracker(Number160.createHash("song1")).start().awaitUninterruptibly();  
  
    System.out.println("peer24 got this: "+futureTracker2.trackers());  
    System.out.println("currently stored on: "+futureTracker2.trackerPeers());  
}
```


- **Demo: Tracker**

- ▶ Although demo uses `await()`, try not to use it
- ▶ Tracker vs. DHT what is better for the CT? You decide!

- **Further interesting aspects for the challenge task:**

- ▶ To be discussed on Thursday
- ▶ Reminder: Thursday starts the challenge task
- ▶ Task presentation and Scrum introduction

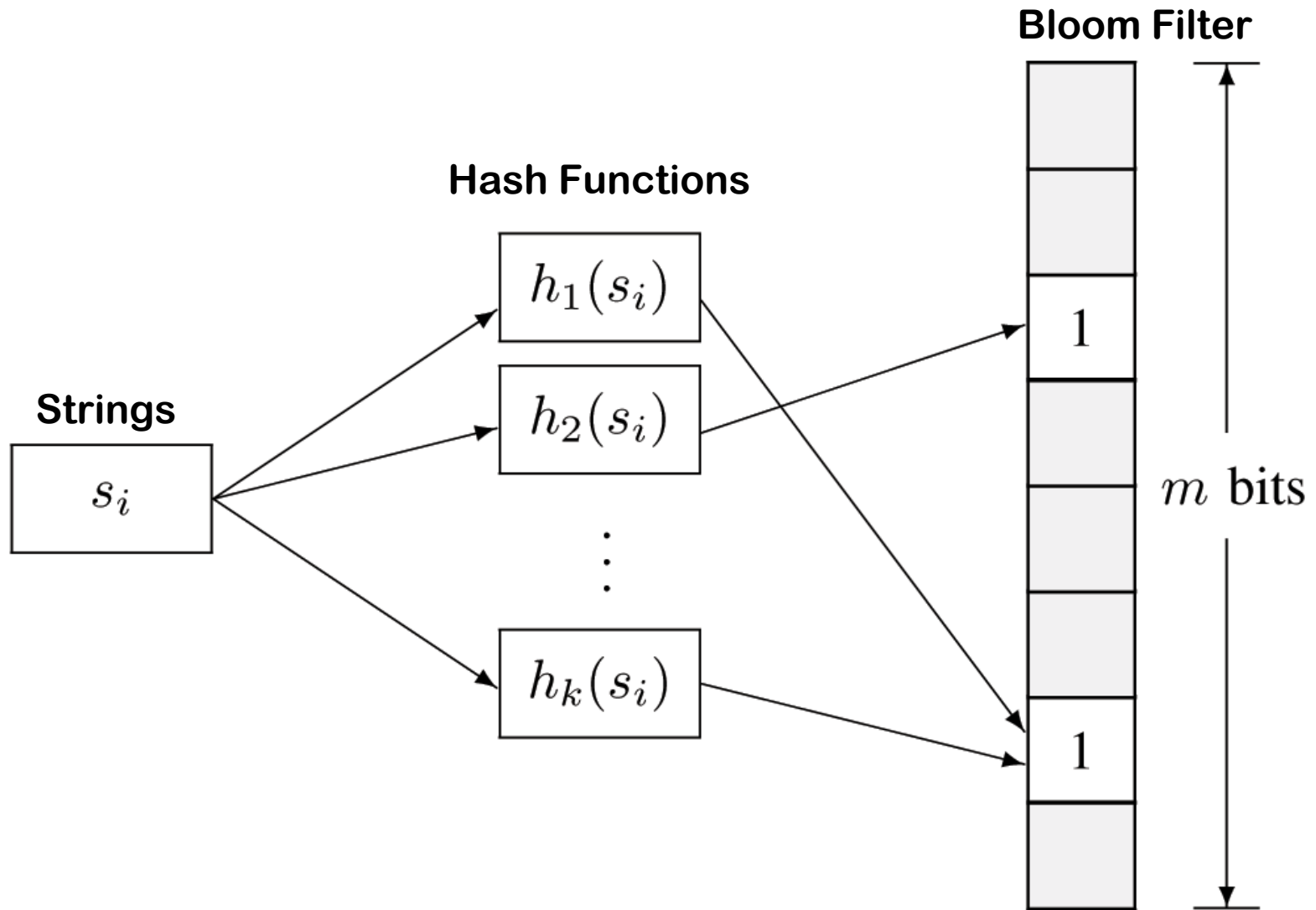
5. Bloom Filters

Traditional Bloom Filter, Attenuated Bloom Filter

Traditional Bloom Filter

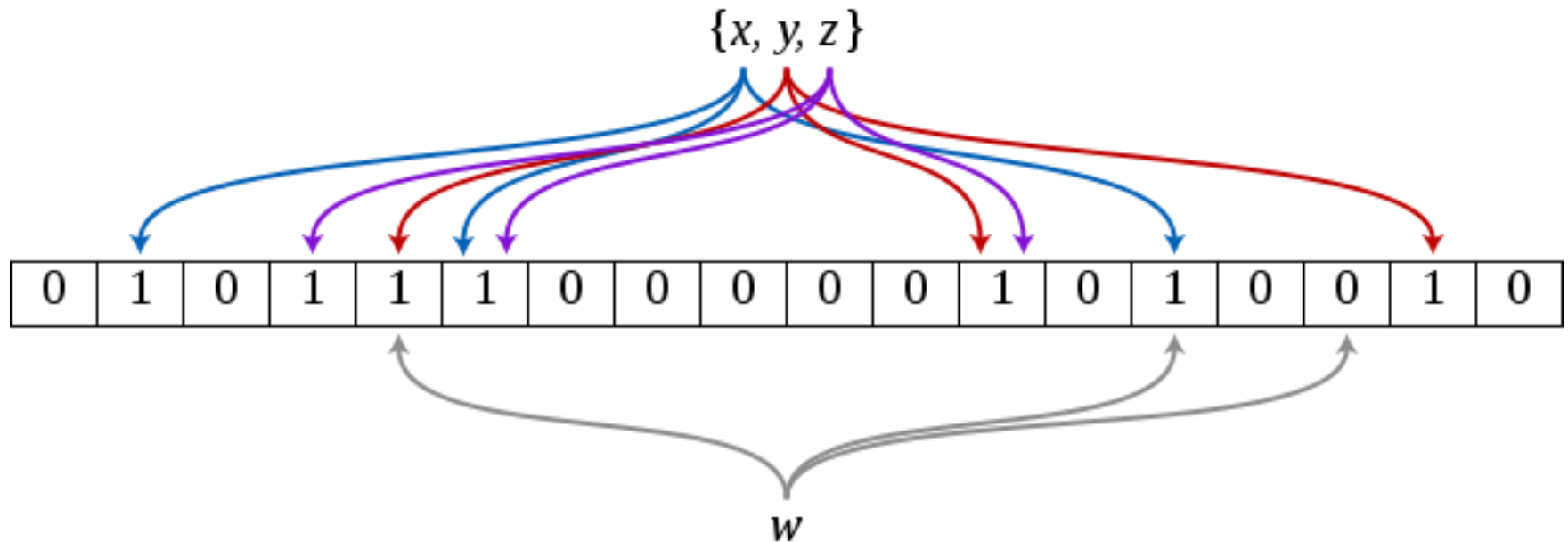
- An array of m bits, initially all bits set to 0
- A bloom filter uses k independent hash functions
 - ▶ h_1, h_2, \dots, h_k with range $\{1, \dots, m\}$
- Each key is hashed with every hash function
 - ▶ Set the corresponding bits in the vector
- Operations
 - ▶ Insertion
 - The bit $A[h_i(x)]$ for $1 < i < k$ are set to 1
 - ▶ Query
 - Yes if all of the bits $A[h_i(x)]$ are 1, no otherwise
 - ▶ Deletion
 - Removing an element from this simple Bloom filter is impossible

Insertion of an Element



Query of an Element, $m=18$, $k=3$

- Insert x, y, z
- Query w



http://en.wikipedia.org/wiki/Bloom_filter

- **Space Efficiency**

- ▶ Any Bloom filter can represent the entire universe of elements
 - In this case, all bits are 1

- **No Space Constraints**

- ▶ Add never fails
- ▶ But false positive rate increases steadily as elements are added

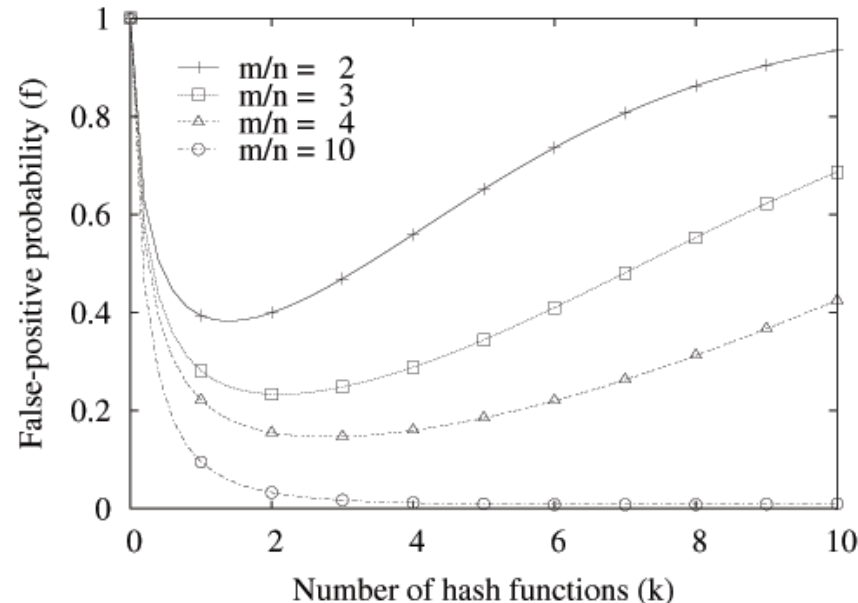
- **Simple Operations**

- ▶ Union of Bloom filters: bitwise OR
- ▶ Intersection of Bloom filters: bitwise AND

False-Positive Probability

- No false negative, but false positive
- False-positive probability:
 - ▶ n number of strings; k hash functions; m -bit vector

$$f = \left(1 - e^{-\frac{nk}{m}}\right)^k$$



=> Given m/n , there is an optimal number of hash functions (opt. $k = m/n \ln 2$) (when 50% of the bits are set)

- **Example for False-positives**

- ▶ Insertions

- Hash („color printer“) => (1,4,6)
 - Hash („digital camera“) => (3,4,5)
 - Bloom filter (1,3,4,5,6)

- ▶ Query

- Hash („heat sensor“) => (3,4,6)
 - Matches since bits 3,4,6 are all set to 1

- ▶ [Online](#)

- **False-negative**

- ▶ Query

- Hash (“color printer”) => (1,4,6) , matches (1,3,4,5,6) → no false-negative

Bloom Filter Variants (1)

- Compressed Bloom Filters

- ▶ When the filter is intended to be passed as a message
- ▶ False-positive rate is optimized for the compressed bloom filter (uncompressed bit vector m will be larger but sparser)
- ▶ However, compression/decompression, more memory

- Generalized Bloom Filter

- ▶ Two type of hash functions g_i (reset bits to 0) and h_j (set bits to 1)
- ▶ Start with an arbitrary vector (bits can be either 0 or 1)
- ▶ In case of collisions between g_i and h_j , bit is reset to 0
- ▶ Store more bits with low false positive
- ▶ Produces either false positives or false negatives

Bloom Filter Variants (2)

- **Counting Bloom Filters**

- ▶ Entry in the filter not be a single bit but a counter
- ▶ Delete operation possible (decrementing counter)
- ▶ [Variable-Increment Counting Bloom Filter](#)

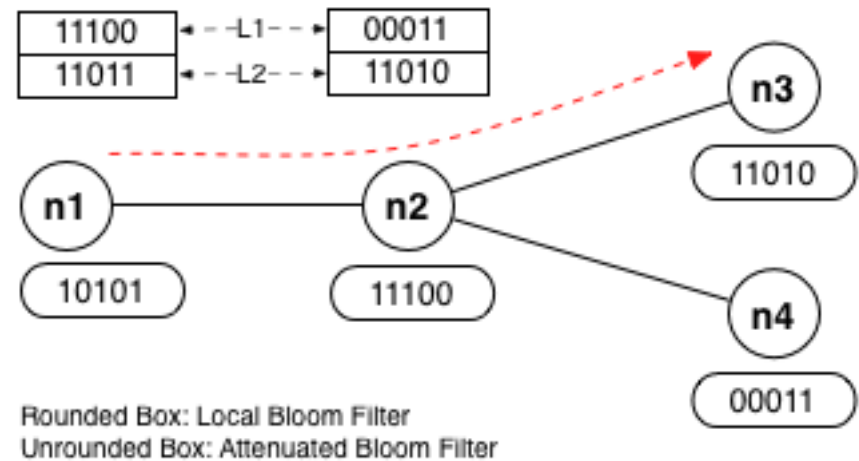
- **Scalable Bloom Filter**

- ▶ Adapt dynamically to number of elements, consist of regular Bloom filters
- ▶ “A SBF is made up of a series of one or more (plain) Bloom Filters; when filters get full due to the limit on the fill ratio, a new one is added; querying is made by testing for the presence in each filter”

Bloom Filter Variants (3)

- Attenuated Bloom Filter

- ▶ Use arrays of Bloom filters to store shortest path distance information
- ▶ Each neighbor link is associated with an attenuated Bloom filter
- ▶ An attenuated Bloom filter consists of d normal Bloom filters
→ d th filter keeps track of resources reachable via d hops.
- ▶ Finds fast only resources within d hops, false positive



http://en.wikipedia.org/wiki/Bloom_filter

Example and Applications

- **Demo**

- ▶ Setup: Bloom Filter of size 128 bits, 20 Number160 objects

- **Applications: Distributed Caching, Spell checking, Routing, (distributed) Databases**

- **B-Tracker uses Bloom Filters**

- ▶ “To avoid duplicates send compressed list of known peers”
- ▶ Idea: store peers in Bloom Filter and send it. Other peers only send us peer not in the Bloom Filter
 - Less traffic (request is larger, reply may be smaller)
 - False positive are possible

- **Demo: Bloom Filter for get() in TomP2P**