



OPERATIONAL RESEARCH

Section 7



Section 7: Transportation Model

Objective: To minimize the transportation cost.

The values in the table cells represents the cost of transportation between different sources and different destinations.

Decision Variables: x_{ij}

↳ No. of items transferred from source "I" to destination "J"

Objective Function: $Z = \sum C_{ij} \cdot x_{ij}$

The problem should be balanced to start solving the problem.

To make it balanced

→ Add Dummy Destination "If Demand < Supply"

→ Add Dummy Source "If Supply < Demand"

To Get IBFS, there is 3 ways to solve:

① North west Corner method

② Least Cost method

③ Vogel method

Problem 1

	D ₁	D ₂	D ₃	
S ₁	5	1	0	20
S ₂	3	2	4	10
S ₃	7	5	9	15
S ₄	9	6	0	15
Demand	5	10	15	

Balanced Condition

- Total Supply = 60
- Total Demand = 30

Supply > Demand

Add "Dummy Destination" with demand value = $60 - 30 = 30$

	D ₁	D ₂	D ₃	D ₄	Dummy Destination
S ₁	5	1	0	0	20
S ₂	3	2	4	0	10
S ₃	7	5	9	0	15
S ₄	9	6	0	0	15
Demand	5	10	15	30	

Using North west method:

Start

	D ₁	D ₂	D ₃	D ₄	
S ₁	5			15	
S ₂				10	
S ₃				15	
S ₄				15	
	5	10	15	30	
	0				

	D ₁	D ₂	D ₃	D ₄	
S ₁	5 → 10				15 15
S ₂					10
S ₃					15
S ₄					15
	0	10	15	30	
					10

	D ₁	D ₂	D ₃	D ₄	
S ₁	5 → 10 → 5				5 10
S ₂					10
S ₃					15
S ₄					15
	0	0	18	30	
					10

	D ₁	D ₂	D ₃	D ₄	
S ₁	5 → 10 → 5				0
S ₂			10		10 10
S ₃					15
S ₄					15
	0	0	10	30	
					10

هذا هو المنهج بالـ Col Lines ← Farcate ← Row Lines

	D ₁	D ₂	D ₃	D ₄	
S ₁	5 → 10 → 5				0
S ₂			10		0
S ₃			0		15
S ₄					15
	0	0	0	30	

	D1	D2	D3	D4	
S1	5 → 10	5	1	0	0
S2		10	1	0	0
S3		0	15	15	15
S4		0	0	30	15
	0	0	0	30	15

	D1	D2	D3	D4	
S1	5 → 10	5	1	0	0
S2		10	1	0	0
S3		0	15	15	15
S4		0	0	15	15
	0	0	0	15	15

$$\begin{aligned}
 Z &= (5 \times 5) + (10 \times 1) + (5 \times 0) + (10 \times 1) + (0 \times 2) + (15 \times 0) + (15 \times 0) \\
 &= 75
 \end{aligned}$$

Note → North west method is the worst as we choose the cell whatever its cost which is wrong

Using "Vogel" Method:

Steps of Solution:

1. Diff. between 2 least Costs in Row/Col
2. Select max. diff
3. Allocate in least Cost cell.
4. If 3 only one source or only one destination
allocate in least cost cell.

	D ₁	D ₂	D ₃	D ₄		
S ₁	5	11	9	9	20	0
S ₂	3	2	4	9	10	16 10 → Max
S ₃	7	5	2	9	15	9
S ₄	9	6	9	9	15	0
	5	10	15	36		
				20		
	2	1	0	0		

	D ₁	D ₂	D ₃	D ₄		
S ₁	5	11	9	9	26 10	0
S ₂	3	2	4	9	10	0 X
S ₃	7	5	2	9	15	9
S ₄	9	6	9	9	15	0
	5	10	15	36		
		10	0			
	2	4	0	0		

	D1	D2	D3	D4	
S1	5	10	0	0	15 10 0
S2	3	2	4	0 10	0 X
S3	7	5	2	0	15 9
S4	9	6	0	0	15 0
	8	0	15	20	
	2	X	0	0	

	D1	D2	D3	D4	
S1	5	10	0	0	5 0
S2	3	2	4	0 10	0 X
S3	7	5	2	0	15 10 12
S4	9	6	0	0	15 0
	0	0	15	20	
	X	X	0	0	

	D1	D2	D3	D4	
S1	5	10	0	0	8 0 0
S2	3	2	4	0 10	0 X
S3	7	5	2	0	0 X
S4	9	6	0	0	15 0
	0	0	15	10	5 0
	X	X	0	0	

	D1	D2	D3	D4		
S1	5	10	5	0	0	x
S2	3	4	4	10	0	x
S3	7	5	2	15	0	x
S4	9	6	10	0	15	10
	0	0	10	5		
	x	x	0	0		

	D1	D2	D3	D4		
S1	5	10	5	0	0	x
S2	3	4	4	10	0	x
S3	7	5	2	15	0	x
S4	9	6	10	0	8	0
	0	0	0	8		
	x	x	0	0		

$$\begin{aligned}
 Z &= (S_1 * 5) + (1 * 10) + (0 * 5) + (0 * 10) + (0 * 15) + (0 * 10) + (0 * 5) \\
 &= 25 + 10 = 35
 \end{aligned}$$

Basic Variables = # Sources + # Destinations - 1

- Steps of Solution: 'Optimal Solution'

1. For B.V $\rightarrow C_{ij} = U_i + V_j$

2. For Non-B.V $\rightarrow C_{ij} = -U_i - V_j$

3. If all " $C_{ij} - U_i - V_j \geq 0$ " \rightarrow Optimal

else, Select most Negative value "EV"

لإيجادLoop || يكتب دوبيكوت

4. Loop.

Conditions of loop:

→ Horizontal and vertical lines.

→ Start and end EV

→ least no. of cells = 4

→ No 3 cells consecutive in Row/Col.

→ # cells should be even.

EV is Non-Basic to Basic

↳ So, we add +ve Sign.

Example 1

	D ₁	D ₂	D ₃	D ₄	Supply	
S ₁	0	2	15	20	11	15
S ₂	12	7	9	15	20	10
S ₃	4	5	14	16	18	5
	5	15	15	15		

	D ₁	D ₂	D ₃	D ₄	U _i
S ₁	0 $(0 - 11 + 14)$ = (3)	2 $(20 - 11 + 11)$ = (20)	15	20 $(20 - 11 + 11)$ = (20)	11 0
S ₂	12 $(12 - 20 + 14)$ = (6)	7 $(7 - 20 + 9)$ = (-4)	9 15	20 10	20 - 0 = 20
S ₃	4 5	14 $(14 - 18 + 9)$ = (5)	16 $(16 - 18 + 11)$ = (9)	18 5	18 - 0 = 18

$$U_j \quad -14 \quad 2-11=-9 \quad -11 \quad 0$$

	D ₁	D ₂	D ₃	D ₄	U _i
S ₁	0 (3)	2 15	0 (20)	20 0	11
S ₂	12 (6)	7 $\overset{0}{(-4)}$ E.V.	9 15	20 10	20
S ₃	4 5	14 (5)	16 (9)	18 5	18
	-14	-9	-11	0	

Donors = 15, 10 $\rightarrow \min(15, 10) = 10$
 So, 10 is D₀

	D1	D2	D3	D4	vi
S1	②	15-10 = 5	20	11 0+10 = 10	11
S2	12	10	15	20	20
S3	4 5	14	16	18 5	18

vj -14 -9 -11 0

Another example: "One Optimal Solution"

	D1	D2	D3	D4	U1	
S1	5 10	10	5 10	6 10	0,0=0	
S2	3 (-2)	2 (1)	3 (1)	6 10	0,0=0	buffer
S3	11 (2)	5 (4)	5 (2)	6 15	0,0=0	0
S4	5 (4)	6 (3)	6 10	5 10	0,0=0	
Uf	5,0,5	1,0,1	0,0,0	0		

Entry Variable = ?

Decors "ve" = (5, 10, 10)

L. Minimum = 5 → D4

depart

	D1	D2	D3	D4	
S1	5 0	10	5 10	6 10	(0)
S2	3 5	2 (1)	3 (4)	6 5	
S3	11 (2)	5 (4)	5 (2)	6 15	
S4	5 (4)	6 (3)	6 5	6 10	