



Screencast: OMPI OpenFabrics Protocols (v1.2 series)

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May 2008



Short Messages

- For short messages
 - memcpy() into / out of pre-registered buffers
 - “short” = “memcpy cost does not matter”
- Once copied in, do one of two things:
 - Use “eager” RDMA
 - Use send / receive semantics

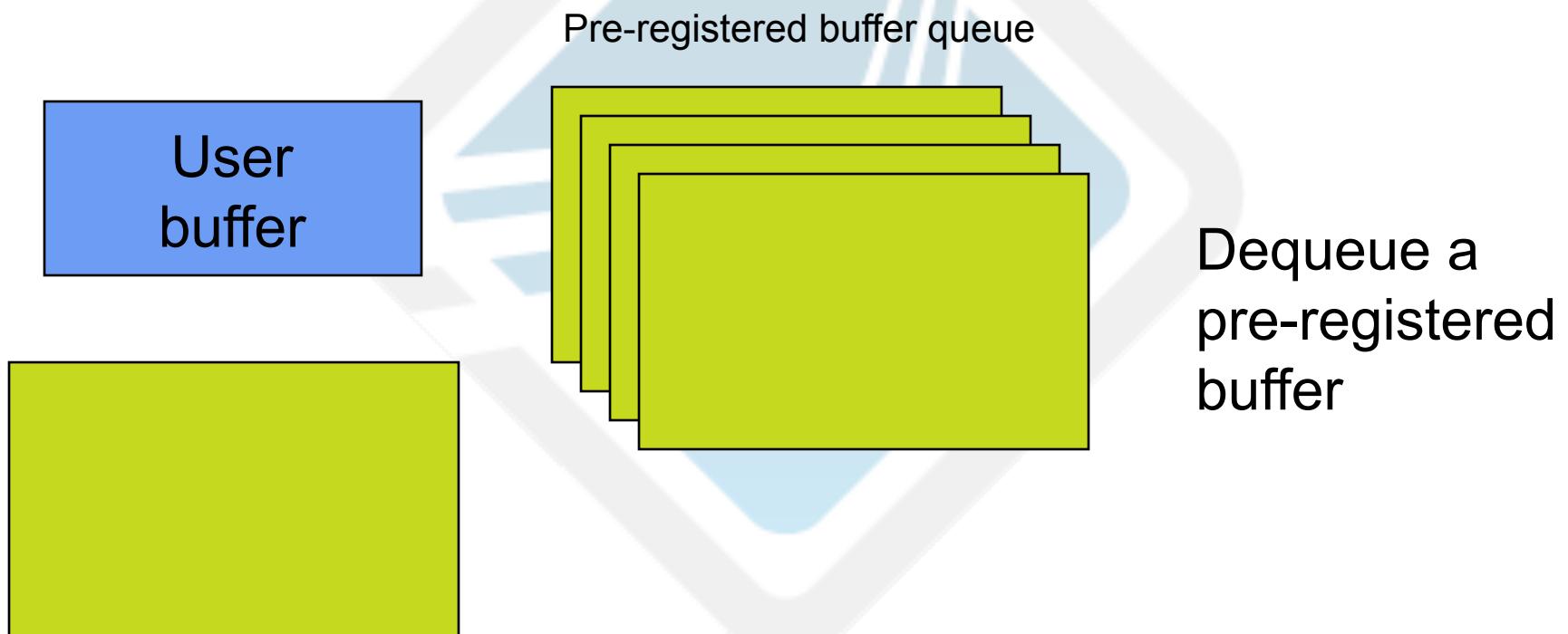
Short Message Protocol

- **`MPI_SEND(short_message, ...)`**



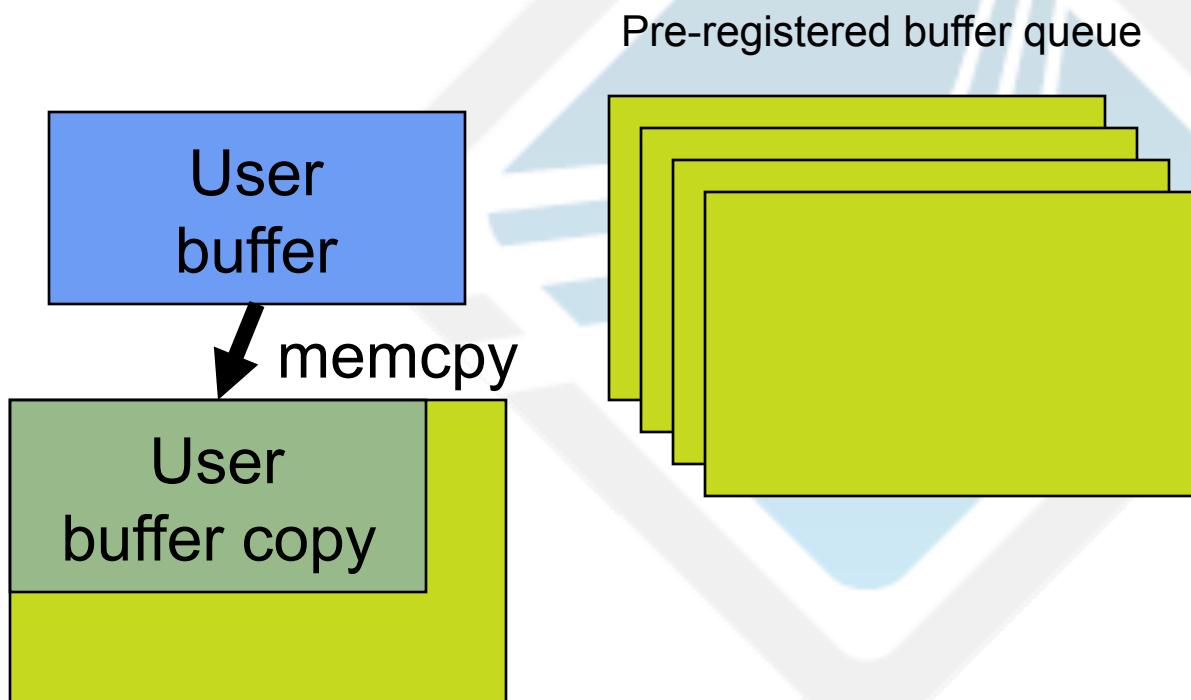
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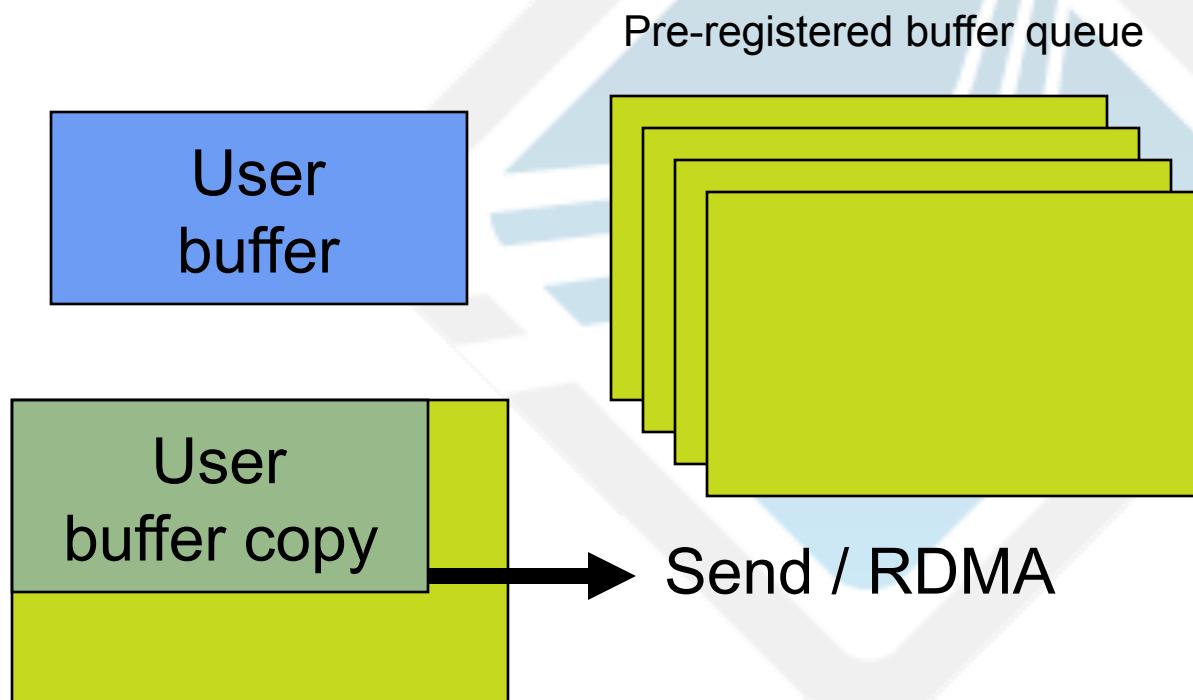
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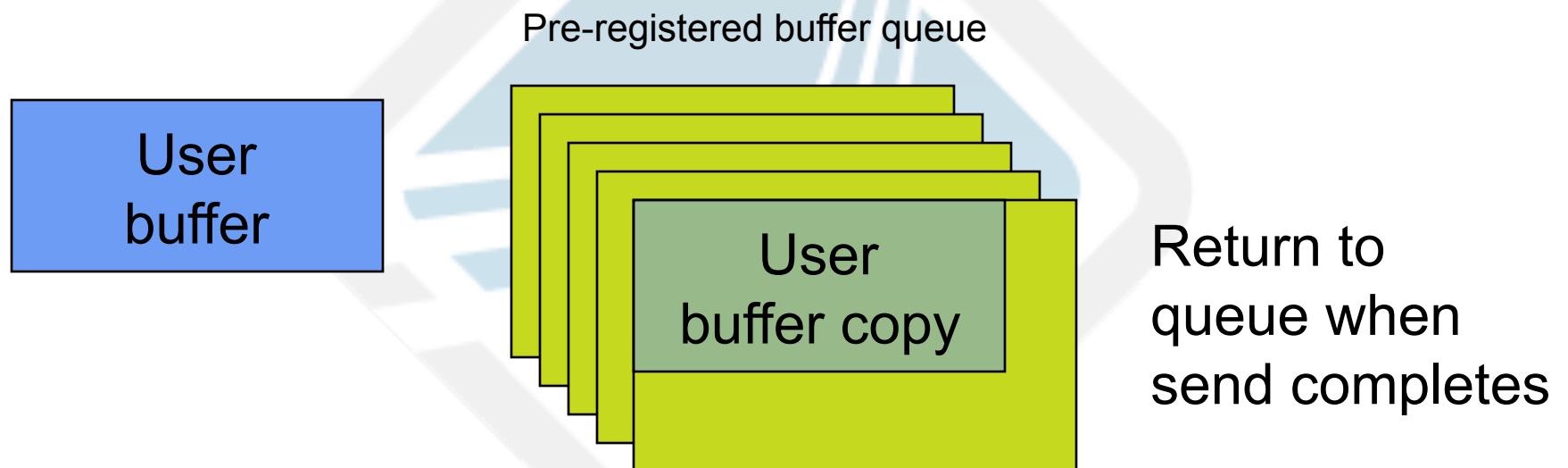
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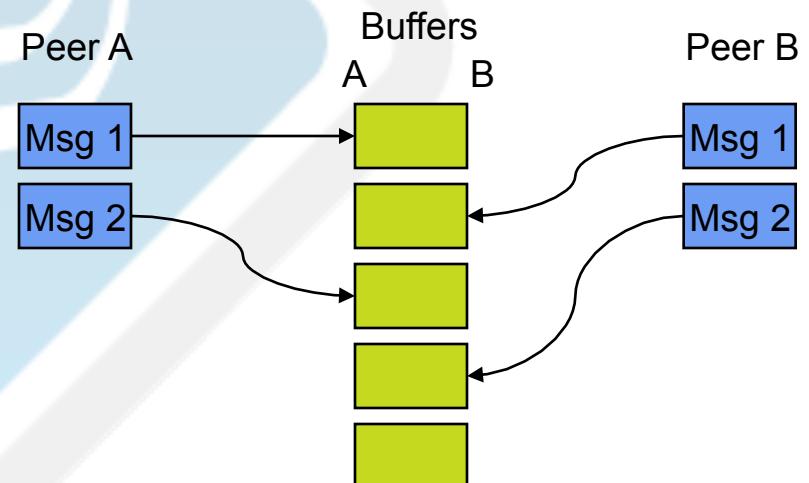
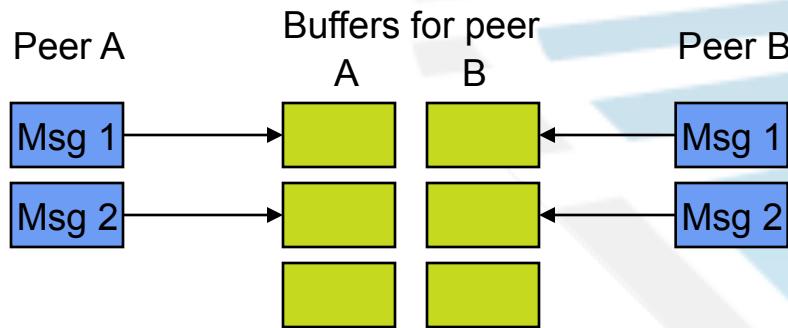


Short: RDMA vs. Send

- RDMA semantics
 - Sender specifies [remote] target buffer address
 - Requires N pre-registered buffers **for each peer**
 - Quickly becomes non-scalable
- Send / receive semantics
 - Receiver specifies target buffer address
 - Can use common pool of pre-registered buffers

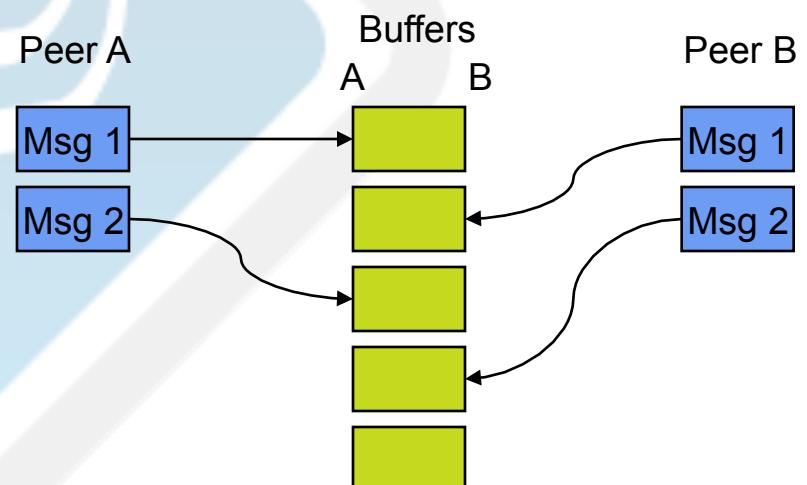
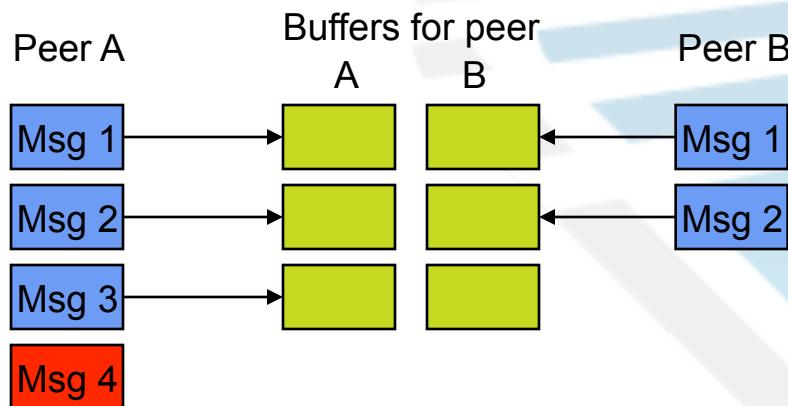
Short: RDMA vs. Send

- RDMA
 - Requires initial setup (exchange addresses)
 - Completely hardware driven
- Send / receive
 - Less initial setup
 - Driver picks buffer
 - Involves remote software



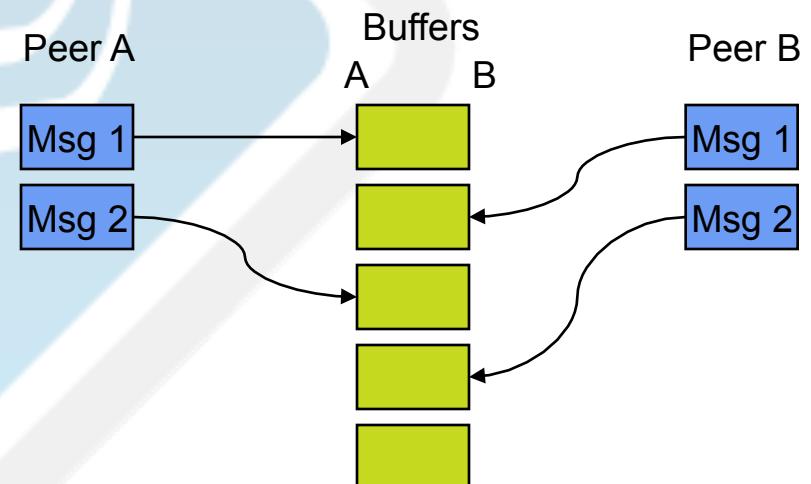
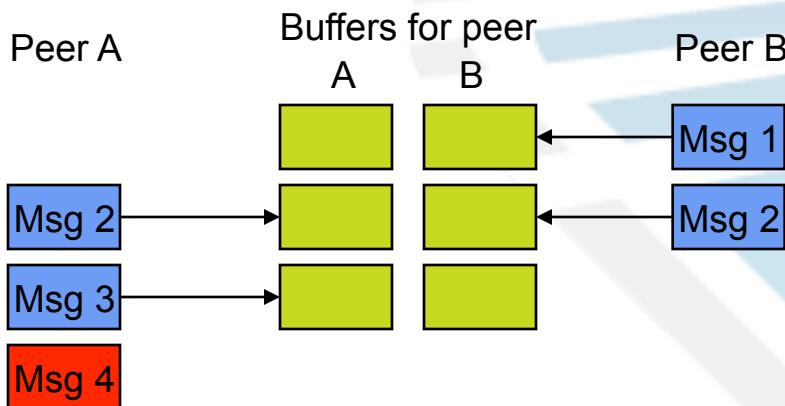
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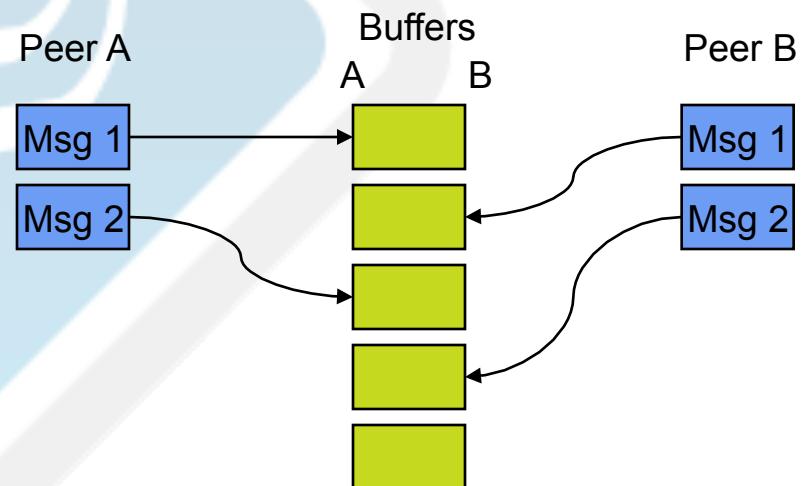
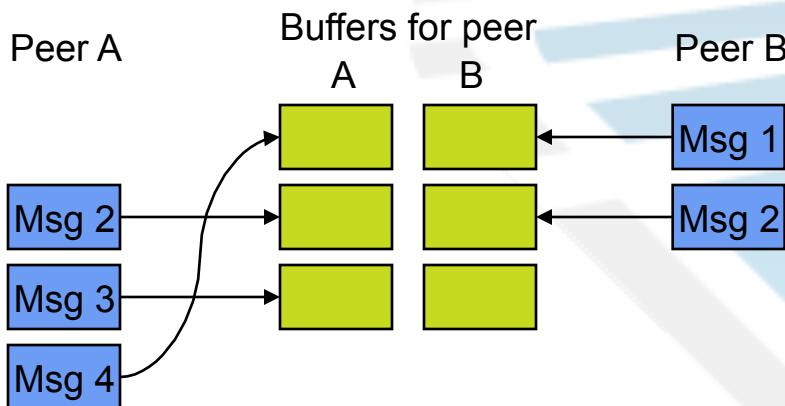
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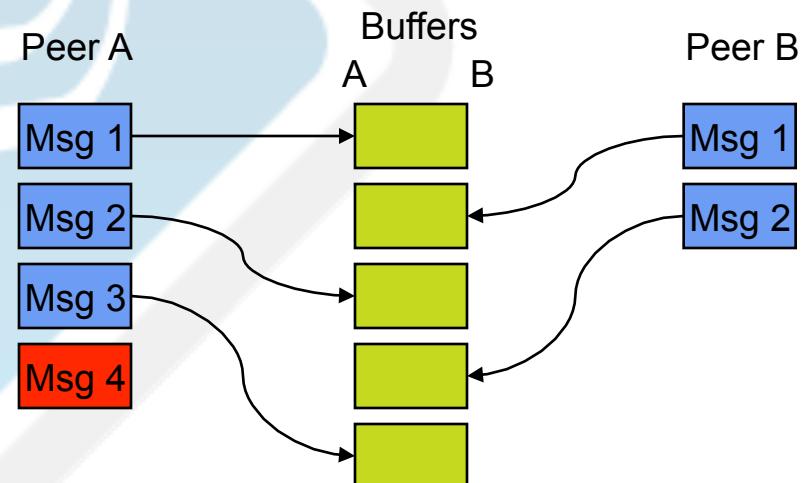
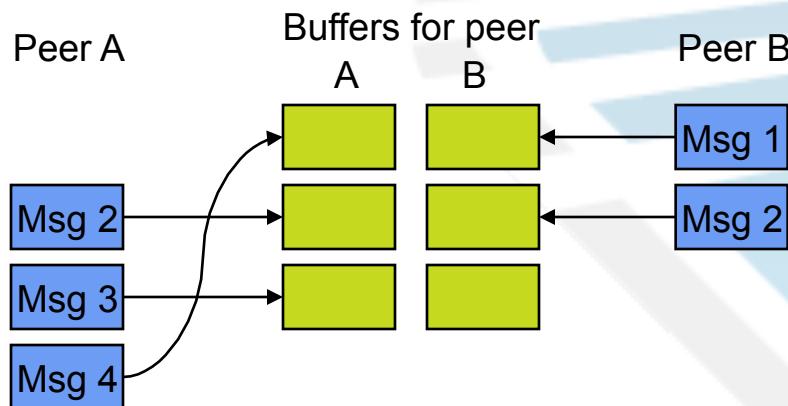
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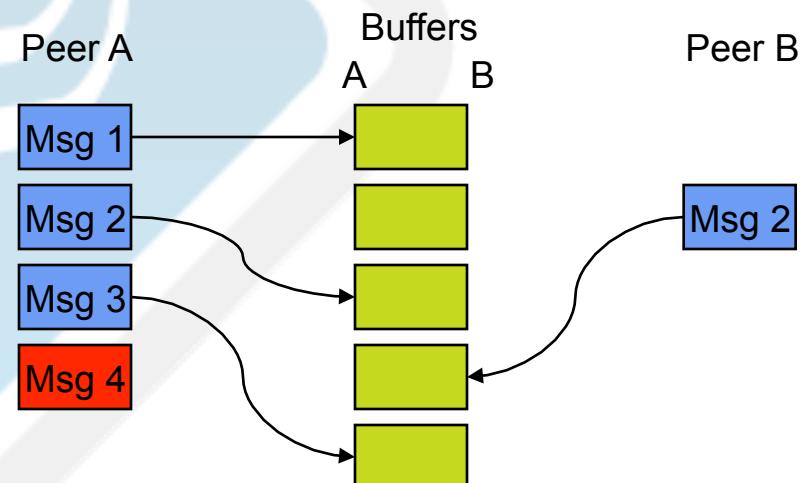
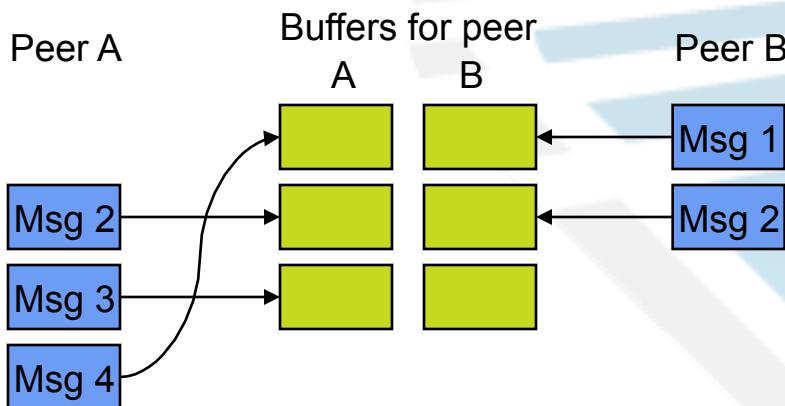
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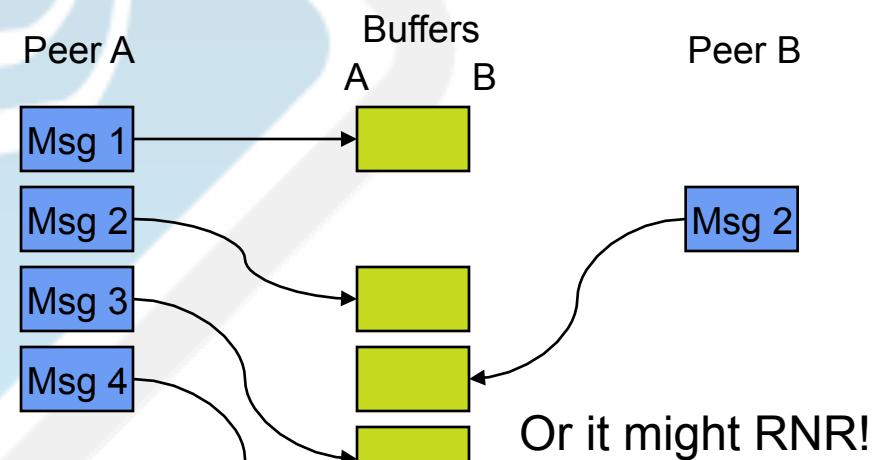
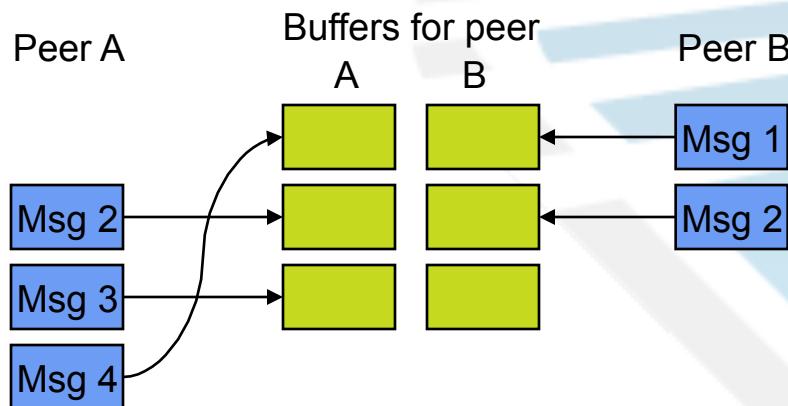
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RDMA vs. Send

- RDMA
 - Buffers for each peer: must be (MxN)
 - Exchange addresses
 - MPI maintains accounting/flow control
 - MPI must notice new received messages
 - Unordered
 - “One-sided”
- More work for MPI
- Send / receive
 - Pool of buffers -- can be less than (MxN)
 - Network maintains accounting / [some] flow control
 - Network notifies of new received messages
 - Ordered
 - Two-sided (ACK’ed)
- Less work for MPI



Limiting Short RDMA

- Open MPI allows N “short” RDMA peers
 - (N+1)th peer will use send/receive for short
 - Receiver’s choice
- If M buffers posted for each RDMA peer
 - Total registered memory for short RDMA

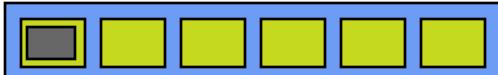
$$M \times N \times \text{short_size}$$

[Shared] Receive Queue

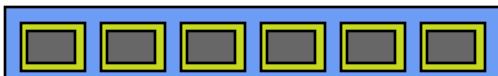
- How to receive messages?
- Per-peer resources
 - Flow control
 - Never error
- Pooled resources
 - No flow control
 - Better utilization
 - Can play stats game
 - Potential for retransmits

Per-peer receive queues

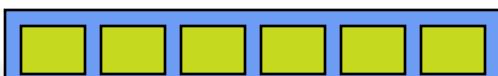
For peer 1:



For peer 2:



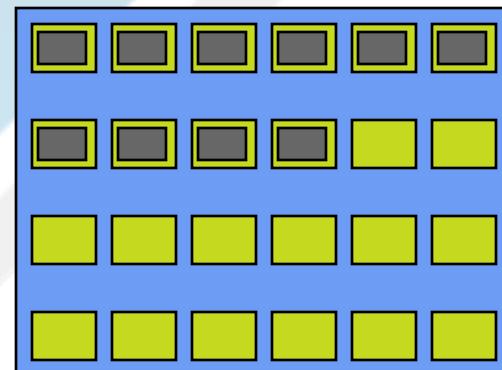
For peer 3:



For peer 4:



Shared receive queue

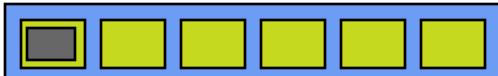


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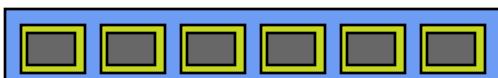
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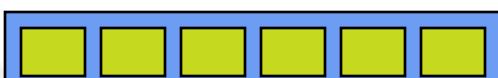
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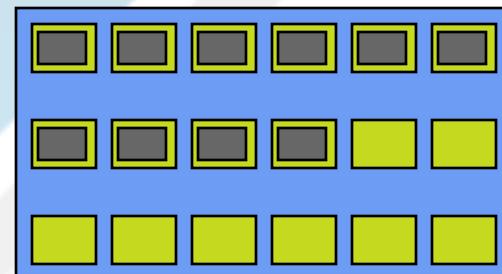
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Shared receive queue



Less than NxM buffers

Long Messages

- For long messages
 - Pipelined protocol
 - Fragment the message, event-driven queue
 - Register
 - Send / receive
 - Unregister
 - (...skipping many details...)
- More complicated if message is not contiguous (not described here)

Long Message Protocol

- **MPI_SEND (long_message, ...)**

Long message

- Sent in 3 phases:
 - Eager / match data
 - Send / receive data
 - RDMA data

Long Message Protocol

- **MPI_SEND (long_message, ...)**

1

- 1st phase: Find receiver match
 - Use (copy + send/receive) for first fragment
 - Only send enough to confirm receiver match; do not overwhelm receiver resources
 - Typical rendezvous protocol

Long Message Protocol

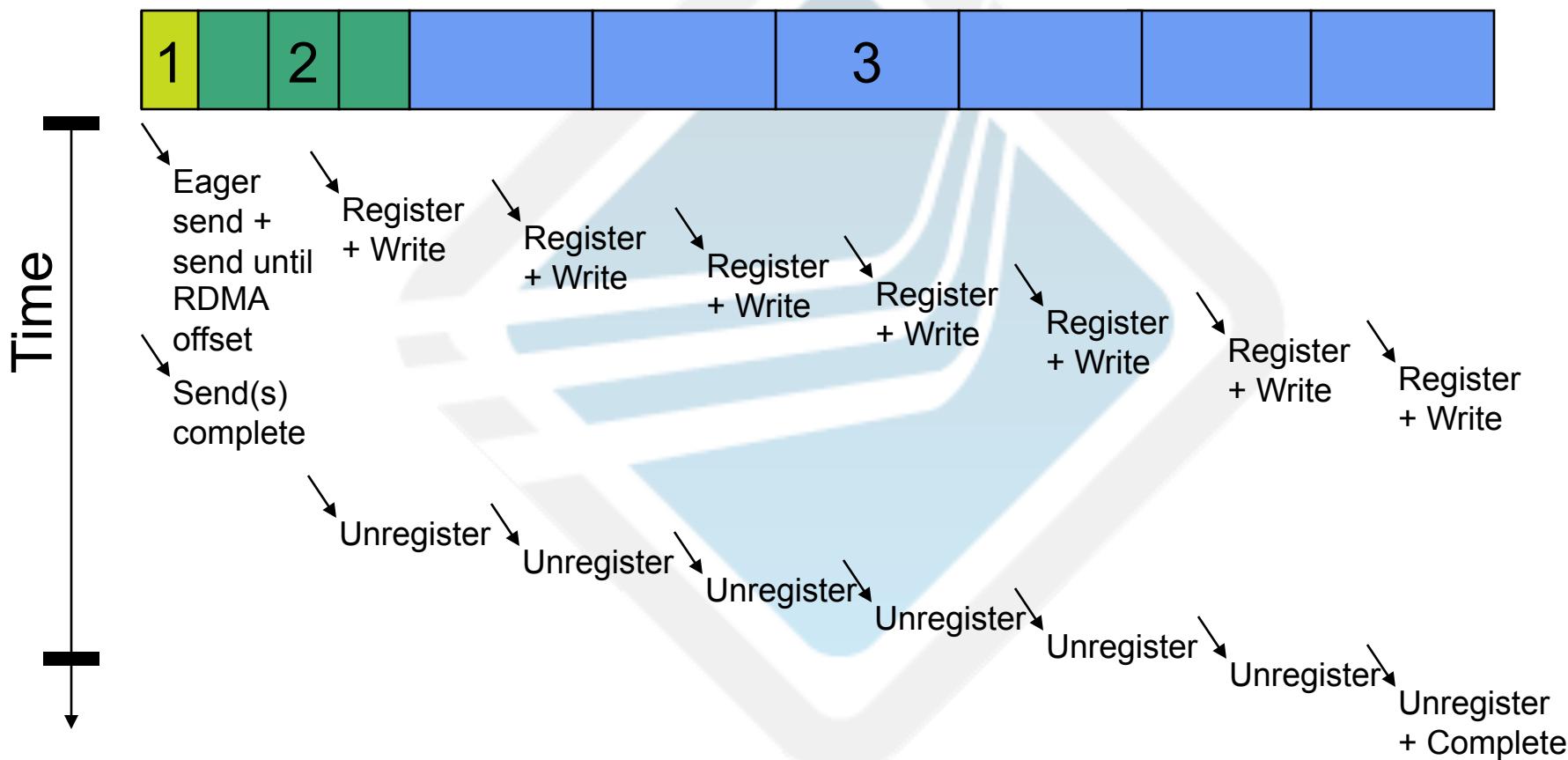
- **MPI_SEND (long_message, ...)**



- 2nd phase: Hide receiver register latency
 - Use (copy + send/receive) for next few frags
 - Allows overlap of memory registration
- Pipeline subject to max depth setting
 - Conserve resources

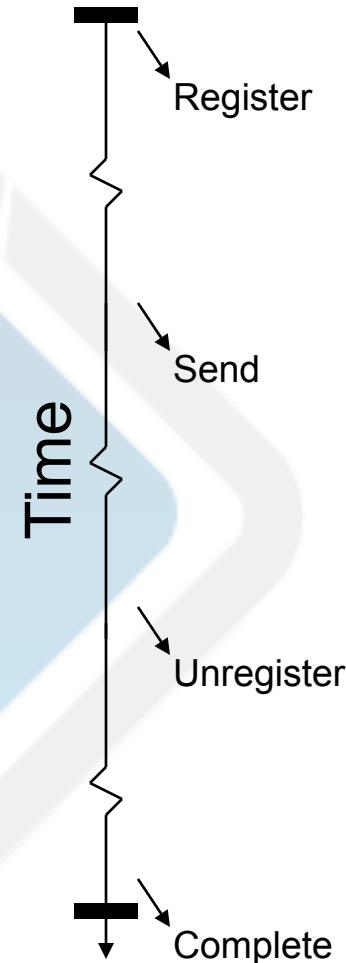
Long Message Protocol

- **MPI_SEND (long_message , ...)**



Why So Complex?

- Why not a single register / send / deregister?
 - Each time is directly proportional to buffer length
- Total time is longer
- Many details skipped
 - See paper on
www.open-mpi.org



fork() Support

- fork() and registered memory problematic
 - Must differentiate between parent / child registered memory
- Not properly supported until:
 - OFED v1.2
 - Open MPI v1.2.1
 - Kernel 2.6.16 or later (but some Linux distros have backported, e.g., RHEL4U6)
- Query MCA param:
 - `btl_openib_have_fork_support`

Striping

- Automatic:
 - Short messages round robin across lowest latency BTL module
 - Long message fragments round robin across all available BTL modules
 - Size of long message fragments proportional to network capacity
- Developers investigating more complex scheduling scenarios

How to Activate Striping?

- Do nothing (!)

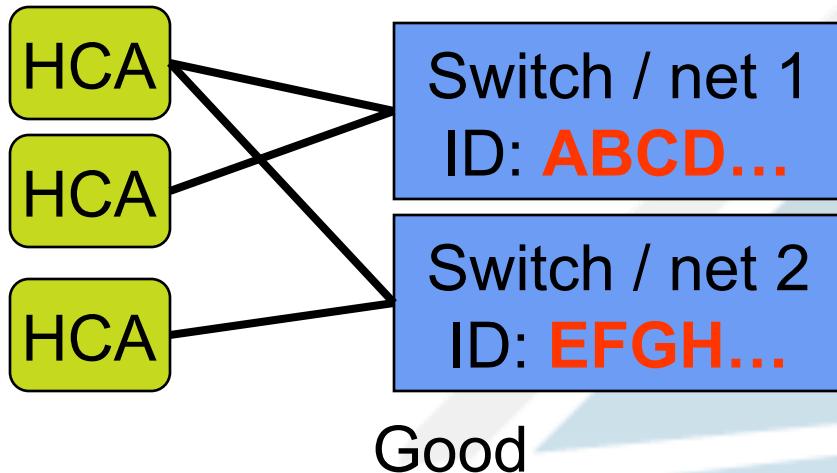
```
mpirun -np 4 a.out
```

- If Open MPI detects multiple active OF ports, it will use them all
 - Assuming each peer port is “reachable”
 - Determined by subnet ID

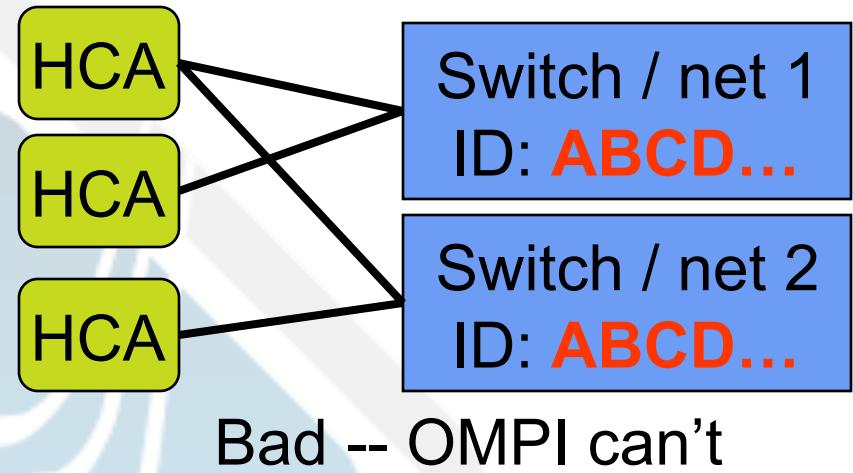
Warning: Default Subnet ID

- Subnet ID uniquely identifies an IB network
 - All Cisco IB switches ship with the default ID
 - FE:80:00:00...
- To compute peer process reachability, Open MPI *must* have different subnet IDs
 - But still, ambiguities exist

Warning: Default Subnet ID

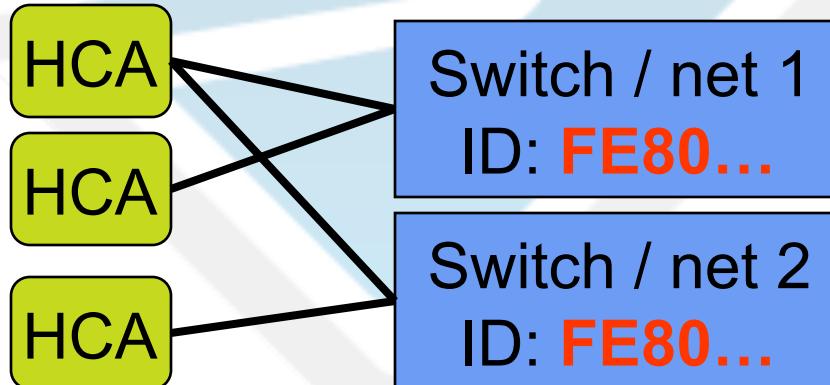


Good



Bad -- OMPI can't
compute reachability

Bad -- OMPI
can't compute
reachability,
but will warn



More Information

- Open MPI FAQ
 - General tuning
<http://www.open-mpi.org/faq/?category=tuning>
 - InfiniBand / OpenFabrics tuning
<http://www.open-mpi.org/faq/?category=openfabrics>

