fn sample_uniform_ubig_below

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This document proves that the implementation of sample_uniform_ubig_below in mod.rs at commit f5bb719 (outdated¹) satisfies its definition. This algorithm uses the same algorithm and argument as used for unsigned native integers, but this time the bit depth is dynamically chosen to fill the last byte of a series of bytes long enough to hold upper.

1 Hoare Triple

Preconditions

Compiler-verified

Argument upper is of type UBig, a bignum

Caller-verified

 $upper \neq 0$

Pseudocode

```
def sample_uniform_ubig_below(upper: UBig) -> UBig:
    byte_len = upper.bit_len().div_ceil(8)
    max = Ubig.from_be_bytes([u8.MAX] * byte_len)
    threshold = max - max % upper

buffer = [0] * byte_len
    while True:
        fill_bytes(buffer)

sample = UBig.from_be_bytes(buffer)
    if sample < threshold:
        return sample % upper</pre>
```

Postcondition

sample_uniform_ubig_below either

- raises an exception if there is a lack of system entropy,
- returns out where out is uniformly distributed between [0, upper).

¹See new changes with git diff f5bb719..a8642172 rust/src/traits/samplers/uniform/mod.rs

2 Proof

Proof. byte_len is the fewest number of bytes necessary to represent upper, which works out to $ceil(ceil(log_2(upper))/8)$. Let max denote the largest integer representable in this many bytes $(2^{(byte_len\cdot 8)}-1)$. Let sample be a uniform sample between [0, max] by flipping byte_len $\cdot 8$ even coins.

You could then (naively) update the range of sample to [0, upper) by rejecting any sample greater than or equal to upper. To reduce the probability of rejection (and improve computational performance), partition the numbers into two sets:

- the leading upper $\cdot k =$ threshold numbers that wrap evenly modulo upper
- the remaining trailing (max mod upper) numbers

It is equivalent to only reject trailing numbers, and return the sample modulo upper. Since $max = threshold + (max \mod upper)$, then $threshold = max - (max \mod upper)$.

Therefore, for any value of upper, the function satisfies the postcondition.

For an intuitive understanding of the sampling approach, consider when upper is two. The naive approach would sample one byte, and then reject if the integer is not zero or one, with probability $\frac{2^8-2}{2^8} \approx .99$. Alternatively, two valid outcomes wrap 2^7 times into 2^8 , with no trailing numbers. Never reject sample because there are no trailing numbers, and return sample mod 2.

Now consider when upper is three. Three valid outcomes wrap 85 times into 2^8 , with one trailing number (255). Reject if sample is 255, otherwise return sample mod 3.