# fn make\_bounded\_range\_to\_pureDP

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This proof resides in "contrib" because it has not completed the vetting process.

Proves soundness of fn make\_bounded\_range\_to\_pureDP. This proof is an adaptation of Lemma 2 here, which proves the conversion between bounded range [DR19] and pure DP.

## 1 Hoare Triple

### Preconditions

#### Compiler-verified

- Variable meas is a valid measurement of type Measurement<DI, TO, MI, RangeDivergence>
- Generic DI (input domain) is a type with trait Domain.
- Generic MI (input metric) is a type with trait Metric.
- MetricSpace is implemented for (DI, MI). Therefore MI is a valid metric on DI.

#### **User-verified**

None

### Pseudocode

#### Postcondition

Theorem 1.1 (Postcondition). For every setting of the input parameters (meas, DI, TO, MI) to make\_bounded\_range\_to\_pureDP such that the given preconditions hold, make\_bounded\_range\_to\_pureDP raises an exception (at compile time or run time) or returns a valid measurement. A valid measurement has the following property:

1. (Privacy guarantee). For every pair of elements x, x' in input\_domain and for every pair (d\_in, d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_measure, if x, x' are d\_in-close under input\_metric, privacy\_map(d\_in) does not raise an exception, and privacy\_map(d\_in)  $\leq$  d\_out, then function(x), function(x') are d\_out-close under output\_measure.

To prove the postcondition, we will need to establish the following definitions and theorem.

**Definition 1.2** (Range Divergence). For any two distributions Y, Y' and any non-negative d, Y, Y' are d-close under the bounded-range privacy measure whenever

$$D_{BR}(Y,Y') = \sup_{y_0,y_1 \in \text{Supp}(Y)} \mathcal{L}_{Y,Y'}(y_0) - \mathcal{L}_{Y,Y'}(y_1)$$
 (1)

**Definition 1.3** (Privacy Loss). The *privacy loss* of an outcome y with respect to random variables Y and Y' is defined as

 $\mathcal{L}_{Y,Y'}(y) = \ln\left(\frac{\mathbb{P}[Y=y]}{\mathbb{P}[Y'=y]}\right). \tag{2}$ 

If y is not in the support of Y' then we define the privacy loss as infinite. The privacy loss random variable Z is distributed according to  $\mathcal{L}_{Y,Y'}(y)$  where y is obtained by sampling  $y \sim Y$ .

**Definition 1.4** (Max Divergence Privacy Loss Random Variable). For a privacy loss random variable Z with respect to two distributions Y, Y' and any non-negative d, Y, Y' are d-close under the max divergence measure if  $|Z| \leq d$ .

**Theorem 1.5** (Range Divergence implies Max Divergence). If two random variables Y and Y' are  $\eta$ -close under range divergence, then they are also  $\eta$ -close under max divergence.

*Proof.* Claimed in Corollary 4.2 of [DR19]. A proof is provided here for completeness. Since the support must include zero, and the range divergence measures the width of the support, then  $|Z| \leq \eta$ . By Definition 1.4 this implies that Y, Y' are  $\eta$ -close under max divergence.

Proof of Theorem 1.1. By the postcondition of Measurement.with\_map, on line 2, make\_bounded\_range\_to\_pureDP returns a measurement with the same input metric, output metric MaxDivergence, and logically equivalent privacy map. The privacy guarantee holds by the precondition that meas is valid measurement, together with Theorem 1.5. Therefore the returned measurement is a valid measurement.

## References

[DR19] David Durfee and Ryan Rogers. Practical differentially private top-k selection with pay-what-you-get composition. *CoRR*, abs/1905.04273, 2019.