

fn get_rounding_distance

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This proof resides in “**contrib**” because it has not completed the vetting process.

Proves soundness of the implementation of `get_rounding_distance` in `mod.rs` at commit `f5bb719` (outdated¹).

1 Hoare Triple

Precondition

Compiler-Verified

- Generic T implements trait `Float`
- Type `i32` implements the trait `ExactIntCast<T.Bits>`, where `T.Bits` is the type of the native bit representation of T.

User-Verified

None

Pseudocode

```
1 def get_rounding_distance() -> RBig:
2     k_min = get_min_k(T) #
3     if k < k_min: #
4         raise f"k ({k}) must not be smaller than {k_min}"
5
6     # input has granularity 2^{k_min} (subnormal float precision)
7     input_gran = x_mul_2k(RBig.ONE, k_min) #
8
9     # discretization rounds to the nearest 2^k
10    output_gran = x_mul_2k(RBig.ONE, k) #
11
12    # the worst-case increase in sensitivity due to discretization is
13    # the range, minus the smallest step in the range
14    distance = output_gran - input_gran #
15
16    # rounding may occur on all vector elements
17    if not distance.is_zero(): #
18        if size is None: #
19            raise "domain size must be known if discretization is not exact"
20
21    match P:
22        case 1:
23            distance *= RBig.from_(size)
```

¹See new changes with `git diff f5bb719..be04bf3f rust/src/measurements/noise/nature/float/utilities/mod.rs`

```

24         case 2:
25             distance *= RBig.try_from(f64.inf_cast(P).inf_sqrt())
26         case _:
27             raise f"norm ({P}) must be one or two"
28
29     return distance

```

Postcondition

Theorem 1.1. Let D denote the space of **size**-dimensional vectors whose elements are in $\mathbb{Z}2^{k_{min}}$, where $2^{k_{min}}$ is the smallest distance between adjacent non-equal values in **T**. Let round_k be a function that rounds each element to the nearest multiple of 2^k , with ties rounding down. Return $\max_{x, x' \in D} \|\text{round}_k(x) - \text{round}_k(x')\|_P - \|x - x'\|_P$, the increase in the sensitivity due to rounding.

Proof. We first consider the increase in sensitivity due to rounding one element. The greatest increase in sensitivity occurs when one value x is rounded down, and x' rounded up. The greatest round down occurs when x is offset from the output grid by 2^{k-1} , and the greatest round up occurs when x' is offset from the output grid by $2^{k-1} + 2^{k_{min}}$.

$$\max_{x, x'} |\text{round}_k(x) - \text{round}_k(x')| \quad (1)$$

$$\leq \max_{x, x'} |(x - 2^{k-1}) - (x' + 2^{k-1} - 2^{k_{min}})| \quad (2)$$

$$= \max_{x, x'} |(x - x') - 2 \cdot 2^{k-1} + 2^{k_{min}}| \quad (3)$$

$$= \max_{x, x'} |x - x'| + 2^k - 2^{k_{min}} \quad (4)$$

$$(5)$$

Now apply this same logic to a vector of rounded values.

$$\max_{x, x'} \|\text{round}_k(x) - \text{round}_k(x')\|_P \quad (6)$$

$$\leq \max_{x, x'} \|x - x' + r\|_P \quad \text{where } r \text{ is a vector of } 2^k - 2^{k_{min}} \quad (7)$$

$$\leq \max_{x, x'} \|x, x'\|_P + \|r\|_P \quad \text{triangle inequality} \quad (8)$$

$$= \max_{x, x'} \|x, x'\|_P + n^{1/P} \cdot (2^k - 2^{k_{min}}) \quad (9)$$

$$(10)$$

Substituting into the return criteria of the postcondition, the return value is:

$$\max_{x, x'} \|\text{round}_k(x) - \text{round}_k(x')\|_P - \|x - x'\|_P \quad (11)$$

$$\leq \max_{x, x'} \|x, x'\|_P + n^{1/P} \cdot (2^k - 2^{k_{min}}) + \max_{x, x'} \|x, x'\|_P \quad (12)$$

$$= n^{1/P} \cdot (2^k - 2^{k_{min}}) \quad (13)$$

We now focus on showing correctness of the implementation. By the postcondition of `get_min_k` on line 2, `min_k` is the k such that the smallest distance between adjacent non-equal values of type **T** is 2^k (the distance between subnormals).

Line 3 ensures that k is not too small, as any smaller k would result in unused precision in the noise sample due to the output being rounded to the nearest **T**. This check is not necessary for privacy; it prevents wasted performance.

By the postcondition of `x_mul_2k` on line 7, `input_gran` is the distance between subnormals, 2^k . Similarly, `output_gran` on line 10 is the distance between adjacent values in the rounded space.

The greatest possible increase in distances between rounded values is thus $2^k - 2^{k_{min}}$, as defined on line 14.

When $k = k_{min}$, the rounding is a no-op and $2^k - 2^{k_{min}}$ is zero, so line 17 skips the vector calculations. Otherwise line 18 ensures the vector size n is known, and the following lines increase the distance by a factor of $n^{1/P}$, resulting in a conservative upper estimate of the expected bound. \square