fn make_bounded_range_to_zCDP

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This proof resides in "contrib" because it has not completed the vetting process.

Proves soundness of bounded_range_to_zCDP in mod.rs at commit 0b8f4222 (outdated¹). The conversion between bounded range [DR19] and zCDP comes from Lemma 3.2 in [CR20]. The proof in this document is an adaptation of Theorem 5 here.

1 Hoare Triple

Preconditions

Compiler-verified

- Variable meas is a valid measurement of type Measurement<DI, TO, MI, RangeDivergence>
- Generic DI (input domain) is a type with trait Domain.
- Generic MI (input metric) is a type with trait Metric.
- MetricSpace is implemented for (DI, MI). Therefore MI is a valid metric on DI.

Human-verified

None

Pseudocode

```
def make_bounded_range_to_zCDP(meas: Measurement) -> Measurement:
    def privacy_map(d_in: f64) -> f64:
        return meas.map(d_in).inf_powi(ibig(2)).inf_div(8.0)

return meas.with_map( #
        meas.input_metric,
        ZeroConcentratedDivergence,
        PrivacyMap.new_fallible(privacy_map),
)
```

Postcondition

Theorem 1.1 (Postcondition). For every setting of the input parameters (meas, DI, TO, MI) to make_bounded_range_to_zCDP such that the given preconditions hold, make_bounded_range_to_zCDP raises an exception (at compile time or run time) or returns a valid measurement. A valid measurement has the following property:

 $^{^{1}\}mathrm{See}$ new changes with git diff 0b8f4222..35b51954 rust/src/combinators/measure_cast/bounded_range_to_zCDP/mod.rs

1. (Privacy guarantee). For every pair of elements x, x' in input_domain and for every pair (d_in, d_out) , where d_in has the associated type for input_metric and d_out has the associated type for output_measure, if x, x' are d_in-close under input_metric, privacy_map(d_in) does not raise an exception, and privacy_map(d_in) \leq d_out, then function(x), function(x') are d_out-close under output_measure.

By the precondition, meas is a valid measurement with RangeDivergence privacy measure.

Definition 1.2 (Range Divergence). For any two distributions Y, Y' and any non-negative d, Y, Y' are d-close under the bounded-range privacy measure whenever

$$D_{BR}(Y,Y') = \sup_{y_0,y_1 \in Supp(Y)} \mathcal{L}_{Y,Y'}(y_0) - \mathcal{L}_{Y,Y'}(y_1)$$
(1)

Definition 1.3 (Privacy Loss). The *privacy loss* of an outcome y with respect to random variables Y and Y' is defined as

$$\mathcal{L}_{Y,Y'}(y) = \ln\left(\frac{\mathbb{P}[Y=y]}{\mathbb{P}[Y'=y]}\right). \tag{2}$$

If y is not in the support of Y' then we define the privacy loss as infinite. The privacy loss random variable Z is distributed according to $\mathcal{L}_{Y,Y'}(y)$ where y is obtained by sampling $y \sim Y$.

Lemma 1.4. Given a privacy loss random variable Z with respect to random variables Y and Y', if supp(Y) = supp(Y'), then

$$\underset{y \sim Y}{\mathbb{E}} \left[\exp(-Z) \right] = 1. \tag{3}$$

Proof.

$$\underset{v \sim Y}{\mathbb{E}} \left[\exp(-Z) \right] \tag{4}$$

$$= \int_{\text{SUDD}(Y)} \exp(-Z)dY \tag{5}$$

$$= \int_{\text{supp}(Y)} \mathbb{P}[Y = y] \cdot \exp(-\mathcal{L}_{Y,Y'}(y)) dy \qquad \text{by Definition 1.3}$$
 (6)

$$= \int_{\text{supp}(Y)} \mathbb{P}[Y = y] \cdot \frac{\mathbb{P}[Y' = y]}{\mathbb{P}[Y = y]} dy \tag{7}$$

$$= \int_{\text{supp}(Y')} \mathbb{P}[Y' = y] dy \qquad \text{since } Y \text{ and } Y' \text{ have the same support}$$
 (8)

$$=1$$

Definition 1.5 (Hoeffding's Lemma). Let X be a random variable supported on [a, b]. Then for any $\lambda \in \mathbb{R}$,

$$\mathbb{E}\left[\exp(\lambda X)\right] \le \exp\left(\mathbb{E}[X] \cdot \lambda + \frac{(b-a)^2}{8} \cdot \lambda^2\right) \tag{10}$$

Lemma 1.6. Given two distributions Y and Y' that are η -close under range divergence, then the associated privacy loss random variable satisfies

$$\mathbb{E}\left[Z\right] \le \frac{1}{8}\eta^2\tag{11}$$

Proof.

$$\mathbb{E}\left[\exp(-Z)\right] \le \exp\left(-\mathbb{E}[Z] + \frac{\eta^2}{8}\right) \qquad \text{by 1.5 since } Z \in [-t, \eta - t], \text{ let } \lambda = -1$$
 (12)

$$\implies \mathbb{E}[Z] \le \frac{\eta^2}{8} - \log \mathbb{E}[\exp(-Z)]$$
 rearrange terms (13)

$$= \frac{\eta^2}{8}$$
 by Lemma 1.4 (14)

Lemma 1.4 can only be applied when Y and Y' have the same support. This requirement is satisfied via a proof by contradiction: if Y and Y' have different supports, then the privacy loss would be infinite, meaning that Y and Y' are not η -close under range divergence.

Definition 1.7 (zero-Concentrated Divergence Privacy Loss Random Variable). For a privacy loss random variable Z with respect to two distributions Y, Y' and any non-negative d, Y, Y' are d-close under the zero-concentrated divergence measure if, for every possible choice of $\alpha \in (1, \infty)$,

$$\mathbb{E}[\exp(\alpha Z)] \le \exp(\alpha(\alpha + 1)d) \tag{15}$$

Theorem 1.8 (Range Divergence implies zero-Concentrated Divergence). If two random variables Y and Y' are η -close under range divergence, then they are also $\frac{1}{8}\eta^2$ -close under zero-concentrated divergence.

Proof.

$$\mathbb{E}\left[\exp(\alpha Z)\right] \qquad \text{starting from Definition 1.7} \tag{16}$$

$$\leq \exp\left(\mathbb{E}[Z]\alpha + \frac{\eta^2}{8}\alpha^2\right)$$
 by 1.5 since $Z \in [-t, \eta - t]$, let $\lambda = \alpha$ (17)

$$\leq \exp\left(\frac{\eta^2}{8}\alpha + \frac{\eta^2}{8}\alpha^2\right)$$
 by Lemma 1.6 (18)

$$=\exp\left(\frac{\eta^2}{8}\alpha\left(\alpha+1\right)\right) \tag{19}$$

$$= \exp(\alpha (\alpha + 1) \rho) \qquad \text{where } \rho = \frac{\eta^2}{8}$$
 (20)

Proof of Theorem 1.1. By the postcondition of Measurement.with_map, on line 5, make_bounded_range_to_zCDP returns a measurement with the same input metric, output metric ZeroConcentratedDivergence, and a privacy map that computes $\eta^2/8$ with conservative rounding. The privacy guarantee holds by the precondition that meas is valid measurement, together with Theorem 1.8. Therefore the returned measurement is a valid measurement.

References

[CR20] Mark Cesar and Ryan Rogers. Unifying privacy loss composition for data analytics. CoRR, abs/2004.07223, 2020.

[DR19] David Durfee and Ryan Rogers. Practical differentially private top-k selection with pay-what-you-get composition. *CoRR*, abs/1905.04273, 2019.