# impl TopKMeasure for MaxDivergence

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This proof resides in "contrib" because it has not completed the vetting process.

This document proves soundness of permute\_and\_flip [2] in mod.rs at commit e62b0aa2 (outdated<sup>1</sup>). permute\_and\_flip noisily selects the index of the greatest score from a vector of input scores.

Permute and flip is equivalent to report noisy max with exponential noise [1]. Report noisy max exponential is implemented via permute and flip because of its discrete nature. Implementation-wise, we will follow permute-and-flip, yet prove the correctness of the algorithm via this equivalence.

## 1 Hoare Triple

### Precondition

#### Compiler-verified

- Method noisy\_top\_k Types consistent with pseudocode.
- Method privacy\_map Types consistent with pseudocode.

#### Caller-verified

- Method noisy\_top\_k
  - x elements are non-null.
  - scale is finite and non-negative.
- Method privacy\_map
  - d\_in is non-null and positive.
  - scale is non-null and positive.

#### Pseudocode

```
# MaxDivergence
def noisy_top_k(x: list[TIA], scale: f64, k: usize, negate: bool) -> list[usize]:
    return exponential_top_k(x, scale, k, negate)

def privacy_map(d_in: f64, scale: f64) -> f64:
    return d_in.inf_div(scale)
```

 $<sup>^{1}\</sup>mathrm{See}\ \mathrm{new}\ \mathrm{changes}\ \mathrm{with}\ \mathrm{git}\ \mathrm{diff}\ \mathrm{e}62\mathrm{b}0\mathrm{aa}2..\mathrm{b}7\mathrm{f}f303\ \mathrm{rust/src/measurements/noisy\_top\_k/mod.rs}$ 

#### Postcondition

Theorem 1.1. The implementation is consistent with all associated items in the TopKMeasure trait.

- 1. Method noisy\_top\_k:
  - Returns the index of the top element  $z_i$ , where each  $z_i \sim \text{DISTRIBUTION}(\text{shift} = y_i, \text{scale} = \text{scale})$ , and each  $y_i = -x_i$  if negate, else  $y_i = x_i$ , k times with removal.
  - Errors are data-independent, except for exhaustion of entropy.
- 2. Method privacy\_map: For any x, x' where  $d_{\text{in}} \geq d_{\text{Range}}(x, x')$ , return  $d_{\text{out}} \geq D_{\text{self}}(f(x), f(x'))$ , where f(x) = noisy top k(x = x, k = 1, scale = scale).

**Definition 1.2.** A random variable follows the Exponential distribution if it has density

$$f(x) = \frac{1}{\beta}e^{-z} \tag{1}$$

where  $z = \frac{x-\mu}{\beta}$ ,  $\mu$  is the shift (location) parameter and  $\beta$  is the scale parameter.

Proof of postcondition: noisy\_top\_k. The preconditions of exponential\_noisy\_max are met, therefore by the postcondition of exponential\_top\_k, the postcondition of noisy\_top\_k is satisfied.

Before proving the privacy guarantees, we state a few required definitions and lemmas:

**Definition 1.3.** Report noisy max with exponential noise computes the index of the maximum element from a set of candidates  $u \in d_i$ n, adds isotropic exponential noise  $Z_i \sim \text{Exp}(1/\lambda)$  to each element in the candidate set u and returns the maximum index as follows:

$$RNM-Exp(s) = \operatorname{argmax}_{i}(s_{i} + Z_{i}), Z_{i} \sim Exp(\lambda)$$
(2)

**Lemma 1.4.** The permute-and-flip mechanism is equivalent to the report-noisy-max with exponential noise mechanism.

See [1] for proof of Lemma 1.4.

**Lemma 1.5.** Let  $X_1, X_2 \sim \text{Exp}(\lambda), \Delta \geq 0$ , then

$$\Pr[X_1 - X_2 \ge \Delta] = e^{-\Delta \lambda} \Pr[X_1 - X_2 \ge 0]$$
 (3)

Proof of Lemma 1.5.

$$\Pr[X_1 - X_2 \ge \Delta] \tag{4}$$

$$=1-\Pr[X_1 \le \Delta + X_2] \tag{5}$$

$$=1-\int_0^\infty \Pr[X_1 \le \Delta + X_2 | X_2 = x] \Pr[X_2 = x] dx \quad \text{by Law of Total Probability}$$
 (6)

$$=1-\int_0^\infty \Pr[X_1 \le \Delta + x] \Pr[X_2 = x] dx \qquad \text{by the fact that } \Delta > 0$$
 (7)

$$=1-\int_{0}^{\infty}\lambda(1-e^{-(x+\Delta)\lambda})e^{-x\lambda}dx\tag{8}$$

$$=1-\lambda \int_0^\infty e^{-x\lambda} dx + \lambda e^{-\Delta\lambda} \int_0^\infty e^{-2x\lambda} dx \tag{9}$$

$$= 1 - 1 + e^{-\Delta \lambda}/2 \qquad \qquad \Pr[X_1 - X_2 \le 0] = \Pr[X_1 - X_2 \ge 0] = 1/2 \quad (10)$$

$$=e^{-\Delta\lambda}\Pr[X_1 - X_2 \ge 0] \tag{11}$$

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**Lemma 1.6.** Let  $u, v \in \text{input\_domain}$  be two vectors of scores. Assume u, v in input\\_domain are d\_inclose under LInfDistance and privacy\_map(d\_in)  $\leq$  d\_out. Let  $Z^* = min_{Z_i} \{u_i + Z_i \geq u_j + Z_j\}, \forall i \neq j$ . Then

$$\ln\left(\frac{\Pr[\mathsf{function}(u)=i]}{\Pr[\mathsf{function}(v)=i]}\right) = \ln\left(\frac{\Pr[Z_i \ge Z^*]}{\Pr[Z_i \ge Z^* + \mathsf{d_in}]}\right) \tag{12}$$

Proof.

$$\ln\left(\frac{\Pr[\mathtt{function}(u)=i]}{\Pr[\mathtt{function}(v)=i]}\right) \tag{13}$$

$$\left( \Pr[\text{function}(v) = i] \right) \\
= \ln \left( \frac{\Pr[\text{RNM-Exp}(\mathbf{u}) = i]}{\Pr[\text{RNM-Exp}(v) = i]} \right) \quad \text{by Lemma 1.4}$$
(14)

$$= \ln \left( \frac{\Pr[\operatorname{argmax}_{k}(u_{k} + Z_{k}) = i]}{\Pr[\operatorname{argmax}_{k}(v_{k} + Z_{k}) = i]} \right)$$
 by Definition 1.3 (15)

Observe that for a fixed i, report noisy max outputs i if:

$$u_i + Z^* \ge u_j + Z_j, \forall i \ne j \tag{16}$$

$$u_i + (v_i - v_i) + Z^* \ge u_i + (v_i - v_i) + Z_i$$
  $\iff$  (17)

$$v_i + (u_i - v_i) + Z^* \ge v_i + (u_i - v_i) + Z_i \qquad \Longleftrightarrow \qquad (18)$$

$$v_i + ((u_i - v_i) - (u_j - v_j) + Z^*) \ge v_j + Z_j \iff (19)$$

$$v_i + (\Delta + Z^*) \ge v_j + Z_j \tag{20}$$

In other words, if  $Z_i \geq (\Delta + Z^*)$ , then function(u) = function(v) = i. This yields us:

$$\ln\left(\frac{\Pr[\operatorname{argmax}_k(u_k + Z_k) = i]}{\Pr[\operatorname{argmax}_k(v_k + Z_k) = i]}\right) = \ln\left(\frac{\Pr[Z_i \ge Z^*]}{\Pr[Z_i \ge \Delta + Z^*]}\right) \tag{21}$$

Proof of postcondition: privacy\_map.

$$\max_{u \sim v} D_{\infty}(M(u)|M(v)) \tag{22}$$

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[function(u) = i]}{\Pr[function(v) = i]} \right)$$

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[RNM-Exp(u) = i]}{\Pr[RNM-Exp(v) = i]} \right)$$

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[argmax_{k}(u_{k} + Z_{k}) = i]}{\Pr[argmax_{k}(v_{k} + Z_{k}) = i]} \right)$$

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[Z_{i} \geq Z^{*}]}{\Pr[Z_{i} \geq Z^{*} + d_{in}]} \right)$$
by Lemma 1.4 (24)
$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[Z_{i} \geq Z^{*}]}{\Pr[Z_{i} \geq Z^{*} + d_{in}]} \right)$$
by Lemma 1.6 (26)

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[\mathtt{RNM-Exp}(\mathtt{u}) = i]}{\Pr[\mathtt{RNM-Exp}(v) = i]} \right)$$
 by Lemma 1.4 (24)

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[\operatorname{argmax}_{k}(u_{k} + Z_{k}) = i]}{\Pr[\operatorname{argmax}_{k}(v_{k} + Z_{k}) = i]} \right)$$
 by Definition 1.3 (25)

$$= \max_{u \sim v} \max_{i} \ln \left( \frac{\Pr[Z_i \ge Z^*]}{\Pr[Z_i \ge Z^* + d_{in}]} \right)$$
 by Lemma 1.6 (26)

$$\leq \frac{d_{in}}{scale}$$
 by Lemma 1.5 (27)

References

[1] Zevu Ding, Daniel Kifer, Thomas Steinke, Yuxin Wang, Yingtai Xiao, Danfeng Zhang, et al. The permute-and-flip mechanism is identical to report-noisy-max with exponential noise. arXiv preprint arXiv:2105.07260, 2021.

[2]	Ryan McKenna selection, 2020.	and	Daniel	Sheldon.	Permute-and-flip:	A	new	mechanism	for	differentially	private