

# fn make\_permute\_and\_flip

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This proof resides in “**contrib**” because it has not completed the vetting process.

This document proves soundness of `make_permute_and_flip` [3] in `mod.rs` at commit `e62b0aa2` (outdated<sup>1</sup>). `make_permute_and_flip` returns a `Measurement` that noisily selects the index of the greatest score from a vector of input scores. This released index can later be used to index into a public candidate set (postprocessing).

## Vetting history

- Pull Request #1678

It is appealing to implement `permute` and `flip` due to its discrete nature. The algorithm requires two primitives that need to be carefully implemented to work out using floating point arithmetics:

1. A shuffling algorithm to shuffle the order in which the elements are being analysed. This is implemented via the Fisher-Yates shuffling algorithm, and it requires uniformly selecting an integer from a set, implemented using `SampleIntBelow`.
2. A exact bernoulli sampler for distributions of the form  $Bern(\exp(-\gamma))$ . The naive technique would require computing the binary expansions of  $\exp(-\gamma)$ . We bypass this issue by using the exact sampler presented in section 5.2 in [1], essentially reducing sampling  $Bern(\exp(-\gamma))$  to sampling  $Bern(\frac{\gamma}{k})$ ,  $k \in N_+$ , for which Bernoulli factories are known.

`Permute` and `flip` is equivalent to report noisy max with exponential noise [2]. Implementation-wise, we will follow `permute-and-flip`, yet proving the correctness of the algorithm will be done via this equivalence.

## Metric

`Permute` and `flip` reuses the same metric as `RNM-gumbel`.

## 1 Hoare Triple

### Preconditions

- TIA (input atom type) is a type with traits `Number` and `CastInternalRational`
- Q0 (output distance type) is a type with traits `Float`, `CastInternalRational` and `DistanceConstant` from type TIA

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<sup>1</sup>See new changes with `git diff e62b0aa2..595c70e1 rust/src/measurements/permute_and_flip/mod.rs`

## Pseudocode

```
1 def make_permute_and_flip(  
2     input_domain: VectorDomain[AtomDomain[TIA]],  
3     input_metric: RangeDistance[TIA],  
4     scale: Q0,  
5     optimize: Union[Literal["max"], Literal["min"]]  
6 ) -> Measurement:  
7     if input_domain.element_domain.nullable:  
8         raise ValueError("input domain must be non-nullable")  
9  
10    if scale < 0:  
11        raise ValueError("scale must be non-negative")  
12  
13    if optimize == "max":  
14        sign = +1  
15    elif optimize == "min":  
16        sign = -1  
17    else:  
18        raise ValueError("must specify optimization")  
19  
20    scale_frac = Fraction(scale)  
21  
22    def function(scores: list[TIA]):  
23        scores = sign * scores  
24  
25        best_score = max(scores)  
26        indexes = range(len(scores))  
27        shuffled_indexes = shuffle(indexes)  
28  
29        for current_index in shuffled_indexes:  
30            coin_bias = (best_score - scores[current_index])/scale  
31            if Bern(coin_bias):  
32                return current_index  
33  
34    def privacy_map(d_in: TIA):  
35        # convert to range distance  
36        # will multiply by 2 if not monotonic  
37        d_in = input_metric.range_distance(d_in)  
38  
39        d_in = Q0.inf_cast(d_in)  
40        if d_in < 0:  
41            raise ValueError("input distance must be non-negative")  
42  
43        if d_in == 0:  
44            return 0  
45  
46        return d_in.inf_div(scale_frac)  
47  
48    return Measurement(  
49        input_domain=input_domain,  
50        function=function,  
51        input_metric=input_metric,  
52        output_metric=MaxDivergence(Q0),  
53        privacy_map=privacy_map,  
54    )
```

## Postcondition

For every setting of the input parameters `input_domain`, `input_metric`, `scale`, `optimize`, `TIA`, `Q0` to `make_permute_and_flip` such that the given preconditions hold, `make_permute_and_flip` raises an exception (at compile time or run time) or returns a valid measurement. A valid measurement has the following property:

1. (Privacy guarantee). For every pair of elements  $u, v$  in `input_domain` and for every pair  $(\mathbf{d\_in}, \mathbf{d\_out})$ , where  $\mathbf{d\_in}$  has the associated type for `input_metric` and  $\mathbf{d\_out}$  has the associated type for `output_measure`, if  $u, v$  are  $\mathbf{d\_in}$ -close under `input_metric`, `privacy_map(d_in)` does not raise an exception, and `privacy_map(d_in) ≤ d_out`, then `function(u), function(v)` are  $\mathbf{d\_out}$ -close under `output_measure`.

## 2 Proof

Before proving the privacy guarantees, we state a few required definitions and lemmas:

**Definition 2.1.** Report noisy max with exponential noise computes the index of the maximum element from a set of candidates  $u \in \mathbf{d\_in}$ , add isotropic exponential noise  $Z \sim \text{Exp}(1/\lambda)$  to each element in the candidate set  $u$  and returns the maximum index as follows:

$$\text{RNM-Exp}(u) = \text{argmax}(u + Z), Z \sim \text{Exp}(1/\lambda_d) \quad (1)$$

**Lemma 2.2.** The permute-and-flip mechanism is equivalent to the report-noisy-max with exponential noise mechanism. Providing proof for this is out of the scope of this document; the details on this equivalence are presented in [2]. The pseudocode for report noisy max with exponential noise is implemented in Algorithm 2 for completeness.

**Lemma 2.3.** Let  $X_1, X_2 \sim \text{Exp}(\lambda)$ ,  $\Delta \geq 0$ , then

$$\Pr[X_1 - X_2 \geq \Delta] = e^{-\Delta\lambda} \Pr[X_1 - X_2 \geq 0] \quad (2)$$

*Proof.*

$$\Pr[X_1 - X_2 \geq \Delta] \quad (3)$$

$$= 1 - \Pr[X_1 \geq \Delta + X_2] \quad \text{by Law of Total Probability} \quad (4)$$

$$= 1 - \int_0^\infty \Pr[X_1 \geq \Delta + X_2 | X_2 = x] \Pr[X_2 = x] dx \quad (5)$$

$$= 1 - \int_0^\infty \Pr[X_1 \geq \Delta + x] \Pr[X_2 = x] dx \quad \text{by the fact that } \Delta > 0 \quad (6)$$

$$= 1 - \int_0^\infty \lambda(1 - e^{-(x+\Delta)\lambda})e^{-x\lambda} dx \quad (7)$$

$$= 1 - \lambda \int_0^\infty e^{-x\lambda} dx + \lambda e^{-\Delta\lambda} \int_0^\infty e^{-2x\lambda} dx \quad (8)$$

$$= 1 - 1 + e^{-\Delta\lambda}/2 \quad \text{Observe } \Pr[X_1 - X_2 \geq 0] = 1/2 \quad (9)$$

$$= e^{-\Delta\lambda} \Pr[X_1 - X_2 \geq 0] \quad (10)$$

□

**Theorem 2.4.** The permute-and-flip mechanism (a.k.a `function`) satisfies  $(\epsilon, 0)$ -DP

*Proof.* Let  $u, v \in \text{input\_domain}$  be two candidate sets associated with input neighbouring datasets  $D$ , respectively  $D'$ . Assume  $u, v$  in `input_domain` are `LInfDistance` and `privacy_map(d_in) ≤ d_out`. We will proceed by proving both directions of the differential privacy inequality:

$$\Pr[\text{function}(u) = i] \geq e^{-\epsilon} \Pr[\text{function}(v) = i], \forall i \in [m] \iff \quad \text{by Lemma 2.2} \quad (11)$$

$$\Pr[\text{RNM-Exp}(u) = i] \geq e^{-\epsilon} \Pr[\text{RNM-Exp}(v) = i], \forall i \in [m] \iff \quad \text{by Definition 2.1} \quad (12)$$

$$\Pr[\text{argmax}_k(u_k + Z_k) = i] \geq e^{-\epsilon} \Pr[\text{argmax}_k(v_k + Z_k) = i], \forall i \in [m] \quad (13)$$

Where  $Z_k \sim \text{Exp}(1/\lambda)$ . Let  $Z^* = \min_{Z_i} \{u_i + Z_i \geq u_j + Z_j\}, \forall i \neq j$ . Observe that for a fixed  $i$ , report noisy max outputs  $i$  if:

$$u_i + Z^* \geq u_j + Z_j, \forall i \neq j \iff (14)$$

$$u_i + (v_i - v_i) + Z^* \geq u_j + (v_j - v_j) + Z_j \iff (15)$$

$$v_i + (u_i - v_i) + Z^* \geq v_j + (u_j - v_j) + Z_j \iff (16)$$

$$v_i + ((u_i - v_i) - (u_j - v_j) + Z^*) \geq v_j + Z_j \iff (17)$$

$$v_i + (\Delta + Z^*) \geq v_j + Z_j (18)$$

In other words, if  $Z_i \geq (\Delta + Z^*)$ , then  $\text{function}(u) = \text{function}(v) = i$ . Switching to a probability bound:

$$\Pr[\text{function}(v) = i] = \Pr[Z_i \geq \Delta + Z^*] \quad \text{by Lemma 2.3} \quad (19)$$

$$\geq e^{-\Delta/\lambda} \Pr[Z_i \geq Z^*] \quad \text{by the definition of } Z^* \quad (20)$$

$$= e^{-\Delta/\lambda} \Pr[\text{function}(u) = i] \quad \text{Set } \lambda = \Delta/\epsilon \quad (21)$$

$$= e^{-\epsilon} \Pr[\text{function}(u)] \iff (22)$$

$$\Pr[\text{function}(v) = i] \geq e^{-\epsilon} \Pr[\text{function}(u)] \iff (23)$$

$$\log \frac{\Pr[\text{function}(u) = i]}{\Pr[\text{function}(u) = i]} \leq \epsilon \quad (24)$$

$$= \text{d\_out} \quad \text{by RangeDistance} \quad (25)$$

The inverse case follows by symmetry, completing the proof. It has been shown that  $\text{function}(u)$  and  $\text{function}(v)$  are  $\text{d\_out}$ -close under  $\text{output\_measure}$  under the definitions of  $\text{function}$  and  $\text{privacy\_map}$ , and the conditions on the input distance and privacy map.

(26)

□

```

1 def report_noisy_max_exponential(
2     input_domain: VectorDomain[AtomDomain[TIA]],
3     input_metric: RangeDistance[TIA],
4     scale: Q0,
5     optimize: Union[Literal["max"], Literal["min"]])
6 ) -> Measurement:
7     if input_domain.element_domain.nullable:
8         raise ValueError("input domain must be non-nullable")
9
10    if scale < 0:
11        raise ValueError("scale must be non-negative")
12
13    if optimize == "max":
14        sign = +1
15    elif optimize == "min":
16        sign = -1
17    else:
18        raise ValueError("must specify optimization")
19
20    scale_frac = Fraction(scale)
21
22    def function(scores: list[TIA]):
23        scores = sign * scores
24        noised_scores = []
25
```

```

26         for score in scores:
27             Z = ExponentialNoise(scale_frac)
28             noised_scores.append(score + Z)
29
30         return argmax(noised_scores)
31
32     def privacy_map(d_in: TIA):
33         # convert to range distance
34         # will multiply by 2 if not monotonic
35         d_in = input_metric.range_distance(d_in)
36
37         d_in = Q0.inf_cast(d_in)
38         if d_in < 0:
39             raise ValueError("input distance must be non-negative")
40
41         if d_in == 0:
42             return 0
43
44         return d_in.inf_div(scale_frac)
45
46     return Measurement(
47         input_domain=input_domain,
48         function=function,
49         input_metric=input_metric,
50         output_metric=MaxDivergence(Q0),
51         privacy_map=privacy_map,
52     )

```

## References

- [1] Clément L Canonne, Gautam Kamath, and Thomas Steinke. The discrete gaussian for differential privacy. *Advances in Neural Information Processing Systems*, 33:15676–15688, 2020.
- [2] Zeyu Ding, Daniel Kifer, Thomas Steinke, Yuxin Wang, Yingtai Xiao, Danfeng Zhang, et al. The permute-and-flip mechanism is identical to report-noisy-max with exponential noise. *arXiv preprint arXiv:2105.07260*, 2021.
- [3] Ryan McKenna and Daniel R Sheldon. Permute-and-flip: A new mechanism for differentially private selection. *Advances in Neural Information Processing Systems*, 33:193–203, 2020.