# NoiseThresholdPrivacyMap<L02I<AbsoluteDistance<UBig», Approximate<ZeroConcentratedDivergence» for ZExpFamily<2>

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This proof resides in "contrib" because it has not completed the vetting process.

Proves soundness of the implementation of NoiseThresholdPrivacyMap for ZExpFamily<2> in mod.rs at commit f5bb719 (outdated¹).

# 1 Hoare Triple

### Precondition

## Compiler-Verified

NoiseThresholdPrivacyMap is parameterized as follows:

- MI, the input metric, is of type LO2I<AbsoluteDistance<UBig>
- MO, the output measure, is of type Approximate<ZeroConcentratedDivergence>

### **User-Verified**

None

# Pseudocode

```
_{
m 1} # analogous to impl NoiseThresholdPrivacyMap < LO2I < AbsoluteDistance < RBig >> , Approximate <
      ZeroConcentratedDivergence>> for ZExpFamily<2> in Rust
  class ZExpFamily2:
      def noise_threshold_privacy_map(
           _input_metric: L02InfDistance[AbsoluteDistance[RBig]],
          output_measure: Approximate[ZeroConcentratedDivergence],
          threshold: UBig,
      ) -> PrivacyMap[L02InfDistance[AbsoluteDistance[RBig]], Approximate[
      ZeroConcentratedDivergence]]:
          noise_privacy_map = self.noise_privacy_map(L2Distance.default(), output_measure[0])
10
11
          scale = self.scale
12
          def privacy_map(d_in: tuple[u32, RBig, RBig]):
13
               10, 12, 1i = d_in
```

 $<sup>^{1}\</sup>mathrm{See}\ \mathrm{new}\ \mathrm{changes}\ \mathrm{with}\ \mathrm{git}\ \mathrm{diff}\ \mathrm{f5bb719..7d35dbd}\ \mathrm{rust/src/measurements/noise\_threshold/distribution/gaussian/mod.}$ 

```
li_sign, li = li.floor().into_parts() #
16
               if li_sign != Sign.Positive: #
17
18
                   raise f"1-infinity sensitivity ({li}) must be non-negative"
19
               12 = 12.min(li * RBig.try_from(f64.from_(10).inf_sqrt())) #
20
               li = li.min(12.floor()) #
21
22
               if 12.is_zero(): #
                   return 0.0, 0.0
24
25
               if scale.is_zero(): #
26
                   return f64.INFINITY, 1.0
27
28
29
               rho = noise_privacy_map.eval(12) #
30
               if li > threshold: #
31
                   raise f"threshold must not be smaller than {li}"
32
33
               d_instability = threshold - li #
34
35
36
               delta_single = conservative_discrete_gaussian_tail_to_alpha(
                   scale.
37
                   d_instability,
38
39
40
               delta_joint: f64 = (1.0).inf_sub(
41
                    (1.0).neg_inf_sub(delta_single).neg_inf_powi(IBig.from_(10)),
42
43
44
               # delta is only sensibly at most 1
45
               return rho, delta_joint.min(1.0)
46
           return PrivacyMap.new_fallible(privacy_map)
```

# Postcondition

Theorem 1.1. Given a distribution self, returns Err(e) if self is not a valid distribution. Otherwise the output is Ok(privacy\_map) where privacy\_map observes the following:

Define function(x) as a function that updates each pair  $(k_i, v_i + Z_i)$ , where  $Z_i$  are iid samples from self, and discards pairs where  $v_i + Z_i$  has smaller magnitude than |threshold|. The ordering of returned pairs is independent from the input ordering.

For every pair of elements x, x' in VectorDomain<AtomDomain<IBig>>, and for every pair (d\_in, d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_measure, if x, x' are d\_in-close under input\_metric, privacy\_map(d\_in) does not raise an exception, and privacy\_map(d\_in)  $\leq$  d\_out, then function(x), function(x') are d\_out-close under output\_measure.

*Proof.* Line 9 rejects self if self does not represent a valid distribution, satisfying the error conditions of the postcondition.

We now construct the privacy map. The li sensitivity can be floored, as neighboring integer datasets always differ in whole integer increments. This doesn't affect the check on line 17 to ensure li is non-negative.

There are two ways to bound the l2 sensitivity given the three bounds:

- 1. The 12 bound directly.
- 2. Define  $x \sim x'$  as  $||x x'||_0 \le 10$  and  $||x x'||_{\infty} \le 1i$ . Then

$$\max_{x \sim x'} ||x - x'||_2 = \max_{x \sim x'} \sqrt{\sum_{i=1}^{n} (x_i - x_i')^2} \le \sqrt{10 \cdot 1i^2} = \sqrt{10} \cdot 1i$$
 (1)

Line 20 updates 12 to the tighter of these bounds.

li similarly has multiple bounds:

- 1. The li bound directly.
- 2. Define  $x \sim x'$  as  $||x x'||_2 \le 12$ . Then

$$\max_{x \sim x'} ||x - x'||_{\infty} = \max_{x \sim x'} \max_{i} |x_i - x_i'| \le \max_{i} [0, \cdots, \lfloor 12 \rfloor, \cdots, 0] = \lfloor 12 \rfloor$$
 (2)

Line 21 updates 1i to the tighter of these bounds.

Now the 12 sensitivity is zero if any of the three bounds are zero, due to the tightening of the 12 bound on line 20. Line 23 then checks for zero sensitivity to return a privacy loss of zero rho, zero delta, because all neighboring datasets have identical pairs, and the ordering is randomized.

Otherwise sensitivity is nonzero, so if the scale is zero, the privacy loss is unbounded, shown on line 26. Now that the edge cases are handled, all sensitivities and scales are finite and strictly positive.

The mechanism that adds noise to the values in the hashmap, as described in the conditions of the post-condition of NoiseThresholdPrivacyMap, matches the mechanism that adds noise to a vector, as described in the conditions of the postcondition of NoisePrivacyMap. Therefore releasing the values incurs the privacy loss computed by noise\_privacy\_map as defined on line 29, The privacy map also ensures that 12 is non-negative.

We now focus on computing the remaining privacy loss parameter delta. To ensure that the distance to instability remains positive, the threshold must be at least as large as the sensitivity, as checked on line 31.

Adapting the proof from [Rog23] (Theorem 7). Consider S to be the set of labels that are common between x and x'. Define event E to be any potential outcome of the mechanism for which all labels are in S (where only stable partitions are released). We then lower bound the probability of the mechanism returning an event E. In the following,  $c_j$  denotes the exact count for partition j, and  $Z_j$  is a random variable distributed according to self.

$$\Pr[E] = \prod_{j \in x \setminus x'} \Pr[c_j + Z_j \le T]$$

$$\ge \prod_{j \in x \setminus x'} \Pr[\Delta_\infty + Z_j \le T]$$

$$\ge \Pr[\Delta_\infty + Z_j \le T]^{\Delta_0}$$

The probability of returning a set of stable partitions  $(\Pr[E])$  is the probability of not returning any of the unstable partitions. We now solve for the choice of threshold T such that  $\Pr[E] \geq 1 - \delta$ .

$$\Pr[\Delta_{\infty} + Z_j \le T]^{\Delta_0} = \Pr[Z_j \le T - \Delta_{\infty}]^{\Delta_0}$$
$$= (1 - \Pr[Z_j > T - \Delta_{\infty}])^{\Delta_0}$$

Let d\_instability denote the distance to instability of  $T - \Delta_{\infty}$ . By the postcondition of conservative\_discrete\_gaussian\_tail\_to\_alpha, the probability that a random noise sample exceeds d\_instability is at most delta\_single. Therefore  $\delta = 1 - (1 - \text{delta_single})^{\Delta_0}$ .

The privacy map returns (rho, delta.min(1)), as  $\delta$  is bounded by 1. It is shown that function(x), function(x') are d\_out-close under output\_measure.

# References

[Rog23] Ryan Rogers. A unifying privacy analysis framework for unknown domain algorithms in differential privacy, 2023.