# fn make\_quantile\_score\_candidates

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This proof resides in "contrib" because it has not completed the vetting process.

Proves soundness of make\_quantile\_score\_candidates in mod.rs at commit f5bb719 (outdated<sup>1</sup>). make\_quantile\_score\_candidates returns a Transformation that takes a numeric vector database and a vector of numeric quantile candidates, and returns a vector of scores, where higher scores correspond to more accurate candidates.

## Vetting History

• Pull Request #456

## 1 Intuition

The quantile score function scores each c in a set of candidates C.

$$s_i = -|(1 - \alpha) \cdot \#(X < C_i) - \alpha \cdot \#(X > C_i)| \tag{1}$$

Where  $\#(X < C_i) = |\{x \in X | x < C_i\}|$  is the number of values in X less than  $C_i$ , and similarly for other variations of inequalities. The scalar score function can be equivalently stated:

$$s_i = -|(1 - \alpha) \cdot \#(X < c) - \alpha \cdot \#(X > c)| \tag{2}$$

$$= -|(1 - \alpha) \cdot \#(X < c) - \alpha \cdot (|X| - \#(X < c) - \#(X = c))| \tag{3}$$

$$= -|\#(X < c) - \alpha \cdot (|X| - \#(X = c))| \tag{4}$$

It has an intuitive interpretation as  $-|candidate\_rank - ideal\_rank|$ , where the absolute distance between the candidate and ideal rank penalizes the score. The ideal rank does not include values in the dataset equal to the candidate. This scoring function considers higher scores better, and the score is maximized at zero when the candidate rank is equivalent to the rank at the ideal  $\alpha$ -quantile.

The scalar scorer is almost equivalent to Smith's[1], but adjusts for a source of bias when there are values in the dataset equal to the candidate. For comparison, we can equivalently write the OpenDP scorer as if there were some  $\alpha$ -discount on dataset entries equal to the candidate.

$$\begin{array}{ll} OpenDP & -|\#(X < c) + \alpha \cdot \#(X = c) - \alpha \cdot |X|| \\ Smith & -|\#(X < c) + 1 \cdot \#(X = c) - \alpha \cdot |X|| \end{array}$$

Observing that  $\#(X \le c) = \#(X < c) + 1 \cdot \#(X = c)$ .

 $<sup>^1\</sup>mathrm{See}$  new changes with git diff f5bb719..31c80c6d rust/src/transformations/quantile\_score\_candidates/mod.rs

## 1.1 Examples

Let  $X = \{0, 1, 2, 3, 4\}$  and  $\alpha = 0.5$  (median):

$$score(X, 0, \alpha) = -|0 - .5 \cdot (5 - 1)| = -2$$

$$score(X, 1, \alpha) = -|1 - .5 \cdot (5 - 1)| = -1$$

$$score(X, 2, \alpha) = -|2 - .5 \cdot (5 - 1)| = -0$$

$$score(X, 3, \alpha) = -|3 - .5 \cdot (5 - 1)| = -1$$

$$score(X, 4, \alpha) = -|4 - .5 \cdot (5 - 1)| = -2$$

The score is maximized by the candidate at the true median. Let  $X = \{0, 1, 2, 3, 4, 5\}$  and  $\alpha = 0.5$  (median):

$$score(X, 0, \alpha) = -|0 - .5 \cdot (6 - 1)| = -2.5$$
  
 $score(X, 1, \alpha) = -|1 - .5 \cdot (6 - 1)| = -1.5$   
 $score(X, 2, \alpha) = -|2 - .5 \cdot (6 - 1)| = -0.5$   
 $score(X, 3, \alpha) = -|3 - .5 \cdot (6 - 1)| = -0.5$   
 $score(X, 4, \alpha) = -|4 - .5 \cdot (6 - 1)| = -1.5$   
 $score(X, 5, \alpha) = -|5 - .5 \cdot (6 - 1)| = -2.5$ 

The two candidates nearest the median are scored equally and highest. Let  $X = \{0, 1, 2, 3, 4\}$  and  $\alpha = 0.25$  (first quartile):

$$score(X, 0, \alpha) = -|0 - .25 \cdot (5 - 1)| = -1$$

$$score(X, 1, \alpha) = -|1 - .25 \cdot (5 - 1)| = -0$$

$$score(X, 2, \alpha) = -|2 - .25 \cdot (5 - 1)| = -1$$

$$score(X, 3, \alpha) = -|3 - .25 \cdot (5 - 1)| = -2$$

$$score(X, 4, \alpha) = -|4 - .25 \cdot (5 - 1)| = -3$$

As expected, the score is maximized when c=1. Let  $X=\{0,1,2,3,4,5\}$  and  $\alpha=0.25$  (first quartile):

$$score(X, 0, \alpha) = -|0 - .25 \cdot (6 - 1)| = -1.25$$

$$score(X, 1, \alpha) = -|1 - .25 \cdot (6 - 1)| = -0.25$$

$$score(X, 2, \alpha) = -|2 - .25 \cdot (6 - 1)| = -0.75$$

$$score(X, 3, \alpha) = -|3 - .25 \cdot (6 - 1)| = -1.75$$

$$score(X, 4, \alpha) = -|4 - .25 \cdot (6 - 1)| = -2.75$$

$$score(X, 5, \alpha) = -|5 - .25 \cdot (6 - 1)| = -3.75$$

The ideal rank is 1.25. The nearest candidate, 1, has the greatest score, followed by 2, and then 0.

# 2 Finite Data Types

The previous equation assumes the existence of real numbers to represent  $\alpha$ . We instead assume  $\alpha$  is rational, such that  $\alpha = \frac{\alpha_{num}}{\alpha_{den}}$ . Multiply the equation through by  $\alpha_{den}$  to get the following, which only uses integers:

$$score(X, c, \alpha_{num}, \alpha_{den}) = -|\alpha_{den} \cdot \#(X < c) - \alpha_{num} \cdot (|X| - \#(X = c))|$$

$$\tag{5}$$

This adjustment also increases the sensitivity by a factor  $\alpha_{den}$ , but does not affect the utility. We now make the scoring strictly non-negative.

- Drop the negation and instead configure the exponential mechanism to minimize the score.
- Compute the absolute difference in a function that swaps the order of arguments to keep the sign positive.

$$score(X, c, \alpha_{num}, \alpha_{den}) = abs \quad diff(\alpha_{den} \cdot \#(X < c), \alpha_{num} \cdot (|X| - \#(X = c)))$$
(6)

To prevent a numerical overflow when computing the arguments to abs\_diff, first choose a data type that the scores are to be represented in. If the number of records is greater than can be represented in this data type, then sample the dataset down to at most this number of records. Notice that when any given record is added or removed, the counts differ by no more than they would have without this sampling down. In the OpenDP implementation, the dataset size may be no greater than the max value of a Rust usize, because each index into the dataset maps to a distinct computer memory address.

Now allocate some of the bits of the data type for the alpha denominator, and use the remaining bits for counts of up to l, where l is the effective dataset size. From this set-up, we choose an  $\alpha_{den}$  such that  $\alpha_{den} \cdot l$  is representable. Since  $\alpha_{num} \leq \alpha_{den}$ ,  $\alpha_{num} \cdot l$  is representable. Since the dataset size fits in the choice of data type, then |X| is representable. Therefore, no quantity in the following equation is not representable.

$$score(X, c, \alpha_{num}, \alpha_{den}, l) = abs \quad diff(\alpha_{den} \cdot \min(\#(X < c), l), \alpha_{num} \cdot \min(|X| - \#(X = c), l))$$
 (7)

Should we compute counts with a 64-bit integer, we might choose  $\alpha_{den}$  to be 10,000. This would allow for a fine fractional approximation of alpha, while still leaving enough room for datasets on the order of  $10^{15}$  elements.

## 3 Hoare Triple

### Precondition

- TIA (input atom type) is a type with trait Number.
- A (alpha type) is a type with trait Float.
- MI is a type with trait ARDatasetMetric.

#### **Function**

```
def make_quantile_score_candidates(
   input_domain: VectorDomain[AtomDomain[TIA]],
   input_metric: MI,
   candidates: List[TIA],
   alpha: A
   ) -> Transformation:
   input_domain.element_domain.assert_non_null()

for i in range(len(candidates) - 1):
   assert candidates[i] < candidates[i + 1]</pre>
```

```
alpha_numer, alpha_denom = alpha.into_frac(size=None)
13
       if alpha_numer > alpha_denom or alpha_denom == 0:
14
           raise ValueError("alpha must be within [0, 1]")
16
17
      if input_domain.size is not None:
          # to ensure that the function will not overflow
18
          input_domain.size.inf_mul(alpha_denom)
19
          size_limit = input_domain.size
20
21
          size_limit = (usize.MAX).neg_inf_div(alpha_den)
22
23
      def function(arg: List[TIA]) -> List[usize]:
24
          return compute_score(arg, candidates, alpha_numer, alpha_denom, size_limit)
25
26
27
      if input domain.size is not None:
28
          def stability_map(d_in: u32) -> usize:
               return TOA.inf_cast(d_in // 2).inf_mul(4).inf_mul(alpha_denom)
29
      else:
30
           abs_dist_const: usize = max(alpha_numer, alpha_denom.inf_sub(alpha_numer))
31
          sup_dist_const: usize = abs_dist_const.inf_mul(2)
32
33
           stability_map = new_stability_map_from_constant(sup_dist_const, Q0=usize)
34
      return Transformation (
35
          input_domain=input_domain,
36
          output_domain=VectorDomain(
37
               element_domain=AtomDomain(T=usize),
38
               size=len(candidates)),
39
          function=function,
40
41
          input_metric=input_metric,
           output_metric=LInfDiffDistance(Q=usize),
42
           stability_map=stability_map,
43
```

### Postcondition

For every setting of the input parameters (input\_domain, input\_metric, candidates, alpha, TIA, A, MI) to make\_quantile\_ score\_candidates such that the given preconditions hold, make\_quantile\_ score\_candidates raises an exception (at compile time or run time) or returns a valid transformation. A valid transformation has the following properties:

- 1. (Appropriate output domain). For every element v in input\_domain, function(v) is in output\_domain or raises a data-independent runtime exception.
- 2. (Stability guarantee). For every pair of elements u, v in input\_domain and for every pair (d\_in, d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_metric, if u, v are d\_in-close under input\_metric and stability\_map(d\_in)  $\leq$  d\_out, then function(u), function(v) are d\_out-close under output\_metric.

## 4 Proof

### 4.1 Appropriate Output Domain

The raw type and domain are equivalent, save for potential nullity in the atomic type. The scalar scorer structurally cannot emit null. Therefore the output of the function is a member of the output domain.

## 4.2 Stability Guarantee

The constructor first performs checks to ensure that the preconditions on compute\_score are met. It checks that vectors in the input domain do not contain null values, that the candidates are strictly increasing, that

alpha is fractional and in the range [0,1], and computes a size\_limit for which size\_limit · alpha\_den does not overflow a usize. Thus by the definition of compute\_score, for each candidate, the response from the function is:

compute 
$$\operatorname{score}(X, c, \alpha_{num}, \alpha_{den}, l) = |\alpha_{den} \cdot \min(\#(X < c), l), \alpha_{num} \cdot \min(|X| - \#(X = c), l)|$$
 (8)

The sensitivity of this function differs depending on if the size of the input vector is known.

#### 4.2.1 Unknown Size Stability

First, consider the case where the size is unknown.

**Lemma 4.1.** If  $d_{Sym}(X, X') = 1$ , then  $d_{\infty}(\text{function}(X), \text{function}(X')) \leq \max(\alpha_{num}, \alpha_{den} - \alpha_{num})$ . Proof. Assume  $d_{Sym}(X, X') = 1$ .

$$\begin{split} &d_{\infty}(\operatorname{function}(X)_{i},\operatorname{function}(X')_{i}) \\ &= \max_{i} |\operatorname{function}(X)_{i} - \operatorname{function}(X')_{i}| \qquad \qquad \text{by definition of } d_{\infty} \\ &= \max_{i} |\operatorname{abs\_diff}(\alpha_{den} \cdot \min(\#(X < C_{i}), l), \alpha_{num} \cdot \min(|X| - \#(X = C_{i}), l)) \qquad \text{by definition of function} \\ &= \operatorname{abs\_diff}(\alpha_{den} \cdot \min(\#(X' < C_{i}), l), \alpha_{num} \cdot \min(|X'| - \#(X' = C_{i}), l))| \\ &= \alpha_{den} \cdot \max_{i} ||\min(\#(X < C_{i}), l) - \alpha \cdot \min(|X| - \#(X = C_{i}), l)|| \\ &= |\min(\#(X' < C_{i}), l) - \alpha \cdot \min(|X'| - \#(X' = C_{i}), l)|| \\ &\leq \alpha_{den} \cdot \max_{i} ||\#(X < C_{i}) - \alpha \cdot (|X| - \#(X = C_{i}))|| \\ &= |\#(X' < C_{i}) - \alpha \cdot (|X| - \#(X' = C_{i}))|| \end{split}$$

Consider each of the three cases of adding or removing an element in X.

Case 1. Assume X' is equal to X, but with some  $X_j < C_i$  added or removed.

$$\begin{split} &=\alpha_{den}\cdot\max_{i}||(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|\\ &-|(1-\alpha)\cdot(\#(X< C_{i})\pm 1)-\alpha\cdot\#(X> C_{i})||\\ &\leq\alpha_{den}\cdot\max_{i}||(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|\\ &-(|(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|+|\pm(1-\alpha)|)| & \text{by triangle inequality}\\ &=\alpha_{den}\cdot\max_{i}|1-\alpha| & \text{scores cancel}\\ &=\alpha_{den}-\alpha_{num} & \text{since }\alpha\leq 1 \end{split}$$

Case 2. Assume X' is equal to X, but with some  $X_j > C_i$  added or removed.

$$\begin{split} &=\alpha_{den}\cdot\max_{i}||(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|\\ &-|(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot(\#(X> C_{i})\pm 1)||\\ &\leq\alpha_{den}\cdot\max_{i}||(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|\\ &-(|(1-\alpha)\cdot\#(X< C_{i})-\alpha\cdot\#(X> C_{i})|+|\pm\alpha|)| & \text{by triangle inequality}\\ &=\alpha_{den}\cdot\max_{i}|\alpha| & \text{scores cancel}\\ &=\alpha_{num} & \text{since }\alpha\geq0 \end{split}$$

Case 3. Assume X' is equal to X, but with some  $X_j = C_i$  added or removed.

$$= \alpha_{den} \cdot \max_i ||(1-\alpha) \cdot \#(X < C_i) - \alpha \cdot \#(X > C_i)|$$

$$- |(1-\alpha) \cdot \#(X < C_i) - \alpha \cdot \#(X > C_i)||$$

$$= 0$$
no change in score

Take the union bound over all cases.

$$\leq \max(\alpha_{num}, \alpha_{den} - \alpha_{num})$$

Take any two elements X, X' in the input\_domain and any pair (d\_in,d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_metric. Assume X, X' are d\_in-close under input\_metric and that stability\_map(d\_in)  $\leq$  d\_out.

$$\begin{split} \operatorname{d\_out} &= \max_{X \sim X'} d_{IDD}(s,s') & \text{where } s = \operatorname{function}(X) \\ &= \max_{X \sim X'} \max_{ij} |(s_i - s_i') - (s_j - s_j')| & \text{by definition of LInfDiffDistance} \\ &\leq \max_{X \sim X'} \max_{ij} |(s_i - s_i')| + |-(s_j - s_j')| & \text{by triangle inequality} \\ &\leq 2 \max_{X \sim X'} \max_{i} |s_i - s_i'| & \text{since } s \text{ is not monotonic, we used the looser triangle inequality} \\ &\leq 2 \sum_{j} \max_{Z_j \sim Z_{j+1}} \max_{i} |s_{i,j} - s_{i,j+1}| & \text{by path property where } d_{Sym}(Z_i, Z_{i+1}) = 1, X = Z_0 \text{ and } X' = Z_{\operatorname{d\_in}} \\ &\leq 2 \sum_{j} \max(\alpha_{num}, \alpha_{den} - \alpha_{num}) & \text{by } 4.1 \\ &\leq 2 \cdot \operatorname{d\_in} \cdot \max(\alpha_{num}, \alpha_{den} - \alpha_{num}) \end{aligned}$$

This formula matches the stability map in the case where the dataset size is unknown.

#### 4.2.2 Known Size Stability

Now consider the case where the dataset size is known.

**Lemma 4.2.** If  $d_{CO}(X, X') \leq 1$ , then  $d_{\infty}(\text{function}(X), \text{function}(X')) \leq 2 \cdot \alpha_{den}$ . *Proof.* Assume  $d_{CO}(X, X') \leq 1$ .

$$\begin{split} &d_{\infty}(\operatorname{function}(X), \operatorname{function}(X')) \\ &= \max_{i} |\operatorname{function}(X)_{i} - \operatorname{function}(X')_{i}| & \text{by definition of } d_{\infty} \\ &= \max_{i} |\operatorname{abs\_diff}(\alpha_{den} \cdot \min(\#(X < C_{i}), l), \alpha_{num} \cdot \min(|X| - \#(X = C_{i}), l)) & \text{by definition of function} \\ &- \operatorname{abs\_diff}(\alpha_{den} \cdot \min(\#(X' < C_{i}), l), \alpha_{num} \cdot \min(|X'| - \#(X' = C_{i}), l))| \\ &= \alpha_{den} \cdot \max_{i} ||\min(\#(X < C_{i}), l) - \alpha \cdot \min(|X'| - \#(X = C_{i}), l)|| \\ &- |\min(\#(X' < C_{i}), l) - \alpha \cdot (|X| - \#(X = C_{i}))|| \\ &- |\#(X' < C_{i}) - \alpha \cdot (|X| - \#(X' = C_{i}))|| \end{split}$$

Consider each of the four cases of changing a row in X.

Case 1. Assume X' is equal to X, but with some  $X_j < C_i$  replaced with  $X'_j > C_i$ .

$$= 2 \cdot \alpha_{den} \cdot \max_{i} ||(1 - \alpha) \cdot \#(X < C_{i}) - \alpha \cdot \#(X > C_{i})|$$

$$- (1 - \alpha) \cdot (\#(X < C_{i}) - 1) - \alpha \cdot (\#(X > C_{i}) + 1)||$$

$$\leq 2 \cdot \alpha_{den} \cdot \max_{i} ||(1 - \alpha) \cdot \#(X < C_{i}) - \alpha \cdot \#(X > C_{i})|$$

$$- (|(1 - \alpha) \cdot \#(X < C_{i}) - \alpha \cdot \#(X > C_{i})| + |1|)|$$

$$= 2 \cdot \alpha_{den} \cdot \max_{i} |1|$$

$$= 2 \cdot \alpha_{den}$$

by definition of function

by triangle inequality scores cancel

Case 2. Assume X' is equal to X, but with some  $X_j > C_i$  replaced with  $X'_j < C_i$ .

$$= 2 \cdot \alpha_{den}$$

by symmetry, follows from Case 1.

Case 3. Assume X' is equal to X, but with some  $X_j \neq C_i$  replaced with  $C_i$ .

$$\leq 2 \cdot \max(\alpha_{num}, \alpha_{den} - \alpha_{num})$$

equivalent to one removal (see make\_quantile\_score\_candidates)

Case 4. Assume X' is equal to X, but with some  $X_j = C_i$  replaced with  $X'_j \neq C_i$ .

$$\leq 2 \cdot \max(\alpha_{num}, \alpha_{den} - \alpha_{num})$$

equivalent to one addition (see make\_quantile\_score\_candidates)

Take the union bound over all cases.

$$d_{\infty}(s_i, s_i') \le \max(2 \cdot \alpha_{den}, 2 \cdot \max(\alpha_{num}, \alpha_{den} - \alpha_{num})) = 2 \cdot \alpha_{den}$$
  
since  $\max(\alpha, 1 - \alpha) \le 1$ 

Take any two elements X, X' in the input\_domain and any pair (d\_in,d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_metric. Assume X, X' are d\_in-close under input\_metric and that stability\_map(d\_in)  $\leq$  d\_out.

$$\begin{aligned} \operatorname{d\_out} &= \max_{X \sim X'} d_{IDD}(s,s') \\ &= \max_{X \sim X'} \max_{ij} |(s_i - s_i') - (s_j - s_j')| & \text{by definition of LInfDiffDistance} \\ &\leq \max_{X \sim X'} \max_{ij} |(s_i - s_i')| + |-(s_j - s_j')| & \text{by triangle inequality} \\ &\leq 2 \max_{X \sim X'} \max_{i} |s_i - s_i'| & \text{since $s$ is not monotonic, we used the looser triangle inequality} \\ &\leq 2 \sum_{j} \max_{Z_j \sim Z_{j+1}} \max_{i} |s_{i,j} - s_{i,j+1}| & \text{by path property where $d_{CO}(Z_i, Z_{i+1}) = 1, X = Z_0$ and $Z_{\mathtt{d\_in}} = X'$} \\ &\leq 2 \sum_{j} 2 \cdot \alpha_{den} & \text{by 4.2} \\ &\leq 4(\mathtt{d\_in}//2) \cdot \alpha_{den} \end{aligned}$$

This formula matches the stability map in the case where the dataset size is known.

#### 4.2.3 Conclusion

Take any two elements X, X' in the input\_domain and any pair (d\_in,d\_out), where d\_in has the associated type for input\_metric and d\_out has the associated type for output\_metric. Assume X, X' are d\_in-close under input\_metric and that stability\_map(d\_in)  $\leq$  d\_out.

By 4.2.1 and 4.2.2 it is shown that function(X), function(X') are  $d_out$ -close under output\_metric for any choice of input arguments.

## References

[1] Adam Smith. Privacy-preserving statistical estimation with optimal convergence rates. In *Proceedings* of the Forty-Third Annual ACM Symposium on Theory of Computing, STOC '11, page 813–822, New York, NY, USA, 2011. Association for Computing Machinery.