

# fn find\_nearest\_multiple\_of\_2k

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This proof resides in “**contrib**” because it has not completed the vetting process.

Proves soundness of the implementation of `find_nearest_multiple_of_2k` in `mod.rs` at commit `f5bb719` (outdated<sup>1</sup>).

## 1 Hoare Triple

### Precondition

#### Compiler-Verified

None

#### User-Verified

$k \neq \text{i32.MIN}$

### Pseudocode

```
1 def find_nearest_multiple_of_2k(x: RBig, k: i32) -> IBig:
2   # exactly compute x/2^k and break into fractional parts
3   num, den = x_mul_2k(x, -k).into_parts() #
4
5   # argmin_i |i * 2^k - x|, the index of nearest multiple of 2^k
6   return (floor_div(num << 1, den) + 1) >> 1 #
```

### Postcondition

**Theorem 1.1.** Return  $\max \argmin_i |i2^k - x|$ .

*Proof.* We first rewrite the expression into an efficiently computable integer form.

$$\max \argmin_i |i2^k - x| \tag{1}$$

$$= \max \argmin_i |i - x2^{-k}| \quad \text{argmin doesn't change when rescaled} \tag{2}$$

$$= \lfloor x2^{-k} \rfloor \quad \text{where } \lfloor \cdot \rfloor \text{ rounds to nearest int with ties towards infinity} \tag{3}$$

$$= \lfloor n/d \rfloor \quad \text{where } x2^{-k} = n/d, n \in \mathbb{Z}, n \in \mathbb{Z}^+ \tag{4}$$

$$= ((n \cdot 2) // d + 1) // 2 \quad \text{where } // \text{ denotes integer floor division} \tag{5}$$

$$= ((n << 1) // d + 1) >> 1 \quad \text{where } \ll \text{ and } \gg \text{ denote left and right shift} \tag{6}$$

Since  $x$  is a rational, line 3 splits  $x2^{-k}$  into its numerator  $n$  and denominator  $d$ . Then line 6 directly computes the quantity of interest. The formula works by dividing twice the numerator, which retains whether

<sup>1</sup>See new changes with `git diff f5bb719..f10ff690 rust/src/measurements/noise/nature/float/utilities/mod.rs`

the original number was in the lower or upper half of the interval. Addition by one moves the upper half of each interval to the next interval. Finally, floor division by two maps the data into the expected space.

Several implementation considerations are worth noting:

- Negation of  $k$  is well-defined for all values of `i32`, except for `i32.MIN`, which is not allowed by the precondition.
- The left shift operator `<<` is equivalent to integer multiplication by 2.
- The right shift operator `>>` is equivalent to floor integer division by 2, as the right-most bit is discarded, which always floors in two's complement arithmetic.
- Integer floor division is implemented via `floor_div`, because `IBig` integer division rounds towards zero, giving the wrong result on negative numerators.

Therefore the implementation follows the formula, which is equivalent to the postcondition. □