

Enabling Event Triggered Monitoring of Traffic Clusters

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Network event	Management task
(Hierarchical) Heavy Hitters	accounting, traffic engineering
Changes in traffic patterns	anomaly detection, DoS detection
Superspreaders	worm, scan, DDoS detection

Let's first create a common ground

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(Hierarchical) Heavy Hitters

Heavy Hitters (HH): a prefix that contribute with a traffic volume larger than a given threshold T during a fixed time interval t .

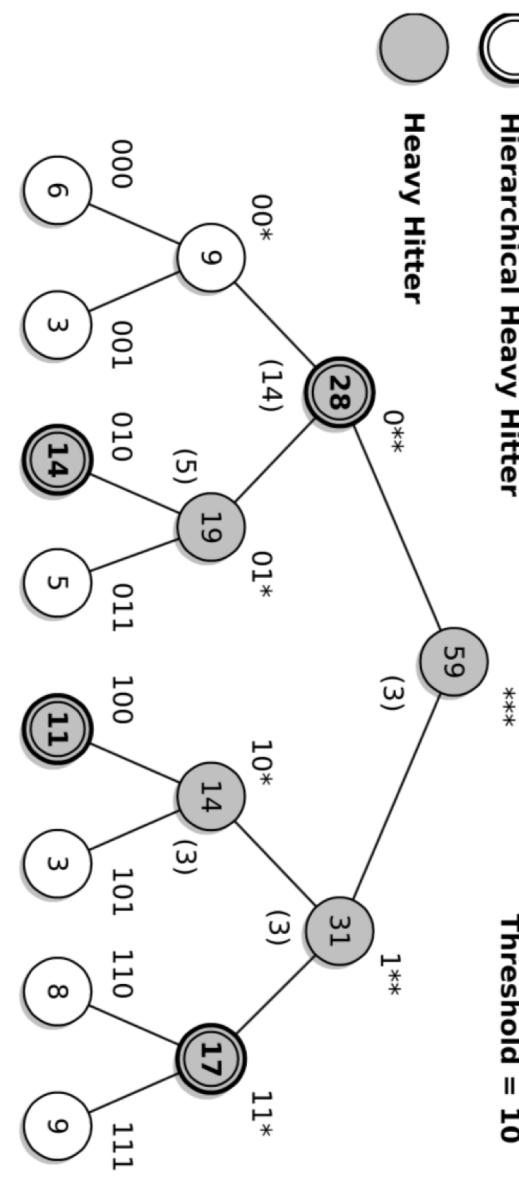


Hierarchical Heavy Hitter



Heavy Hitter

Threshold = 10



Hierarchical Heavy Hitter (HHH): a prefix that exceeds a threshold T after excluding the contribution of all its HHH descendants.

Changes in traffic patterns

Identifying the flows that contribute the most for the changes in the traffic patterns over two consecutive time intervals.

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Superspreaders

A host that contacts at least a given number of distinct destinations over a short time period.

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Superspreaders	worm, scan, DDoS detection

All those network events can be seen as a
traffic cluster detection problem

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HHH and change detection: packets or bytes per second.

Superspreaders: flows per second.

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the detection of those events?

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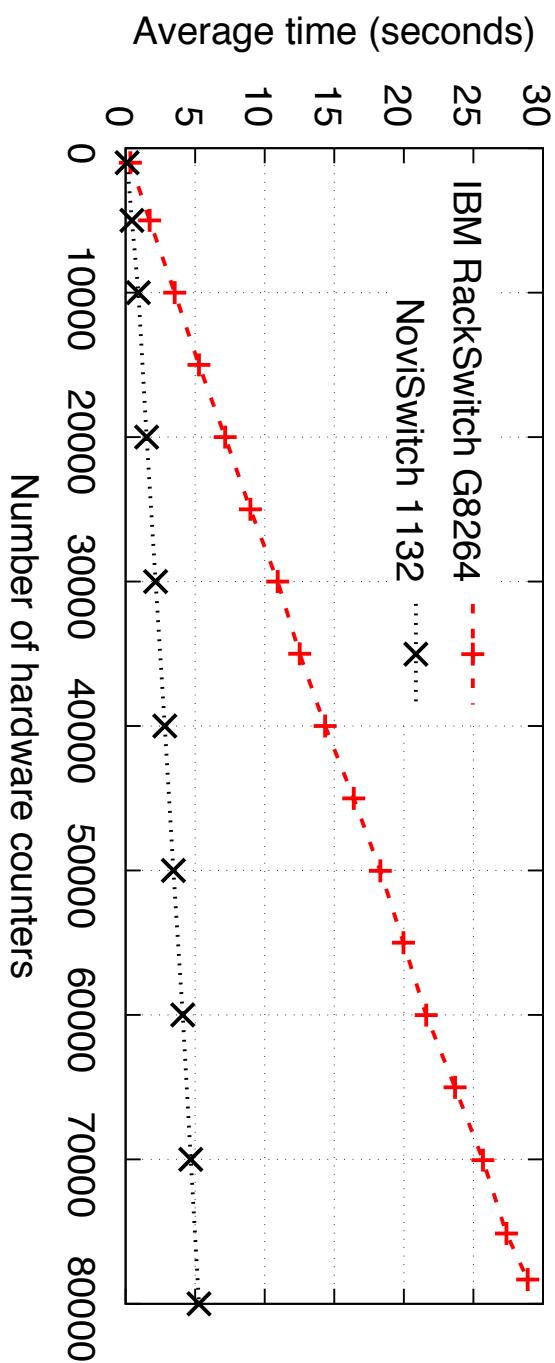
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Wait a minute. Is this a problem?

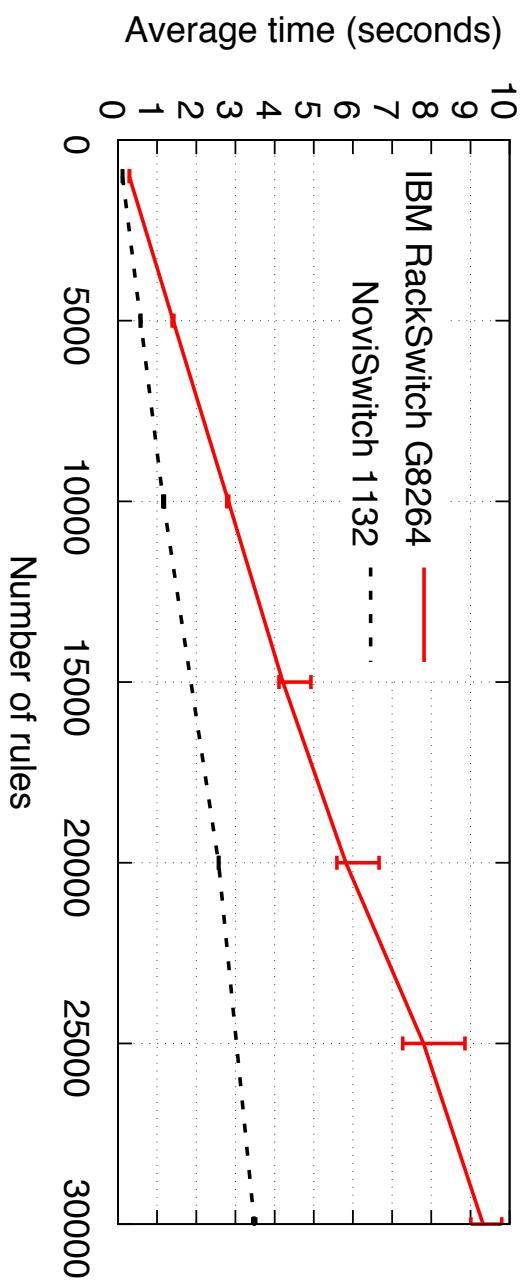


Retrieving a large number of counters from hardware is
time consuming!!!



Note: probabilistic data structures (i.e., sketches) require large amount of counters to lower false positive ratio.

Updating forwarding state and statistic retrieval are two **competing operations** that are commonly run sequentially



Note: having large chunk of forwarding updates is a pretty common case during blackholing.

Can we leverage **dataplane programmability** to assist in
the detection of those events?

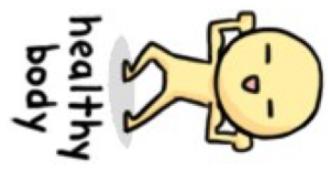
Can we leverage dataplane programmability to assist in
the detection of those events? ~~assist~~

Can we leverage dataplane programmability to **enable in-**
network detection of those events?

Because if you have in-network detection..

As soon as you detect you can take pre-defined actions.

Good for network reactivity.



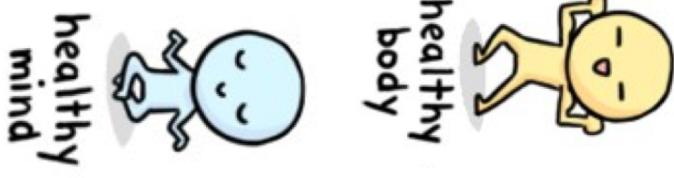
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You can directly export the detection result to the control plane.

Control plane does not have to receive lot of data and understand what is going on.



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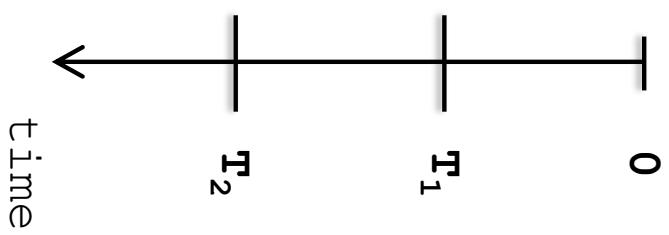


Elastic Trie in a nutshell

- Prefix tree that **grows** or **collapses**: focus on who account for a *large* share of the traffic.
- **Starting condition:** a single node corresponding with zero-length prefix *.
- Each node consists of **three elements**: (1) left child counter, (2) right child counter, (3) node timestamp.
- Use **timeouts** to detect heavy prefixes and to grow or collapse the trie
(i.e., if in the time interval t , prefix A exceeds a threshold, then refine the prefix)

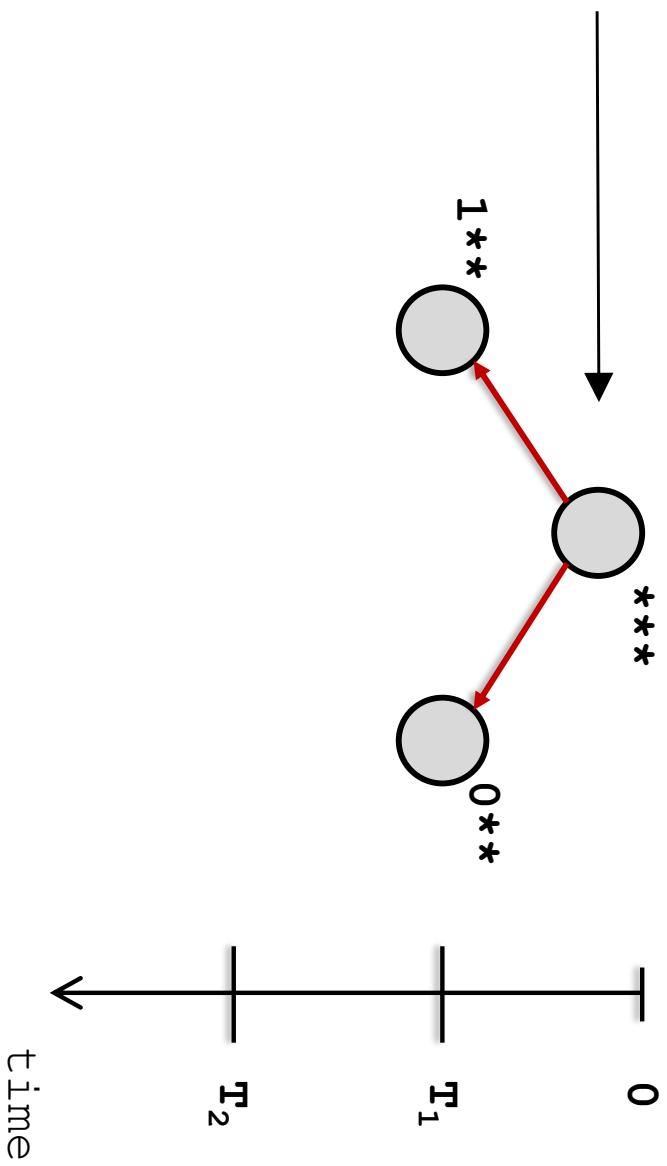
Elastic Trie in action

- counter-Left
- timestamp
- counter-Right

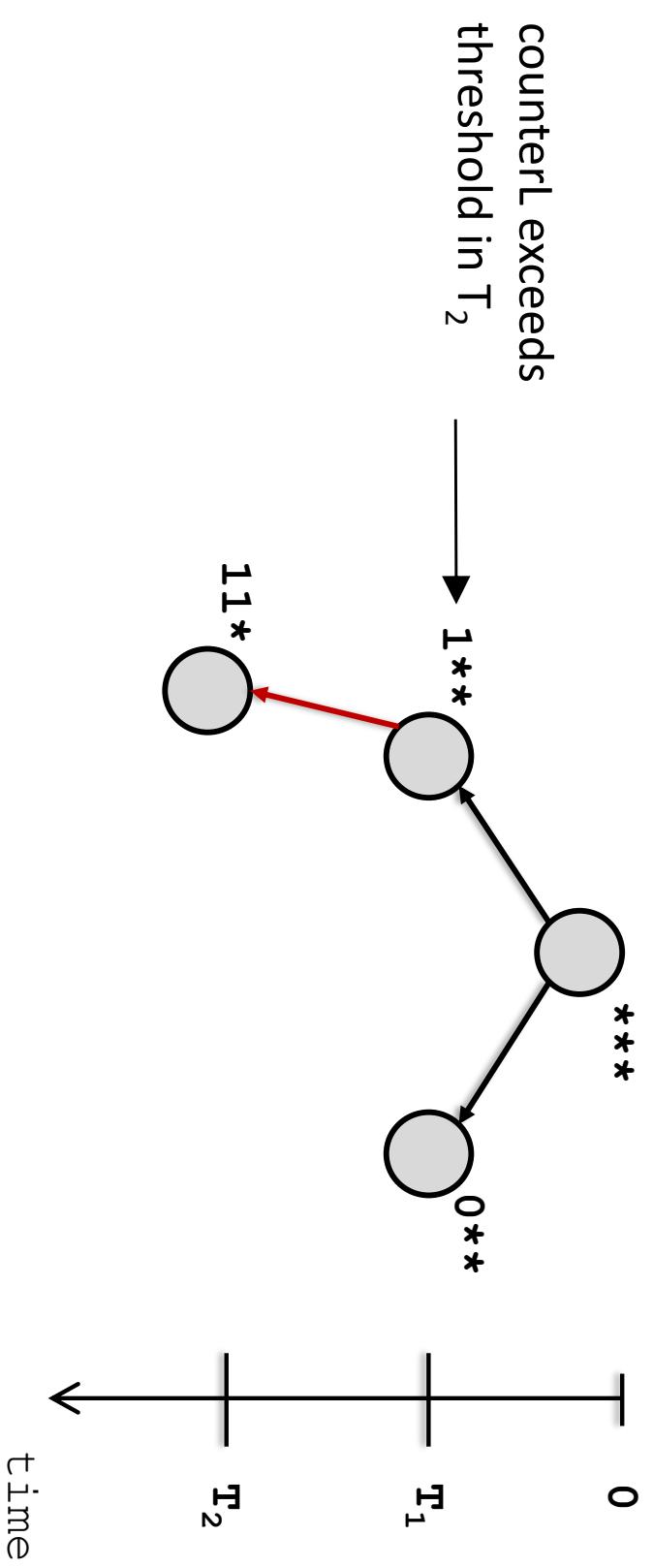


Elastic Trie in action

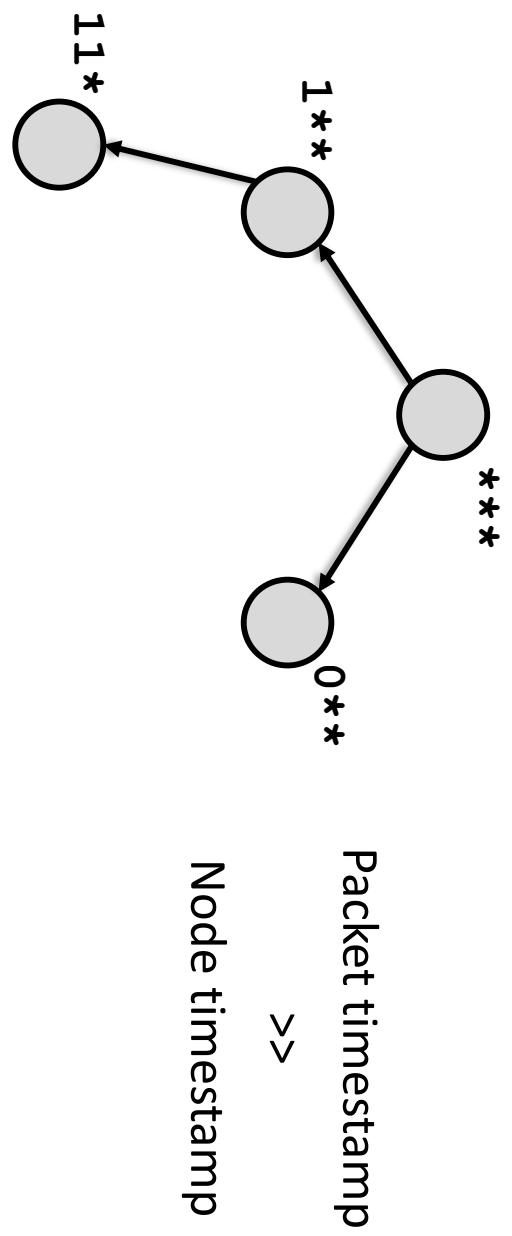
Both counterL and
counterR exceed
threshold in T_1



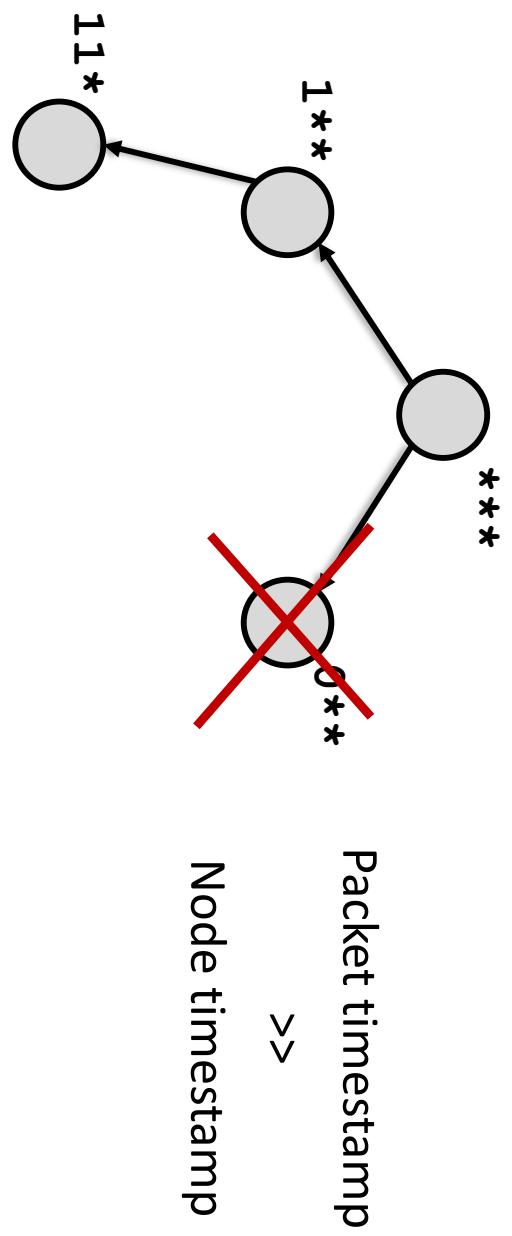
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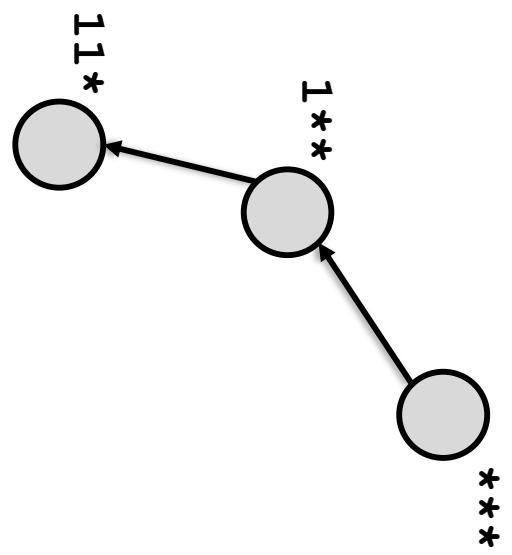
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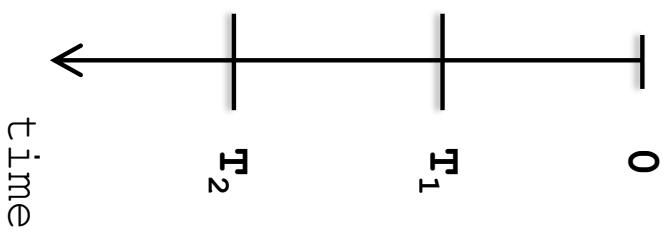


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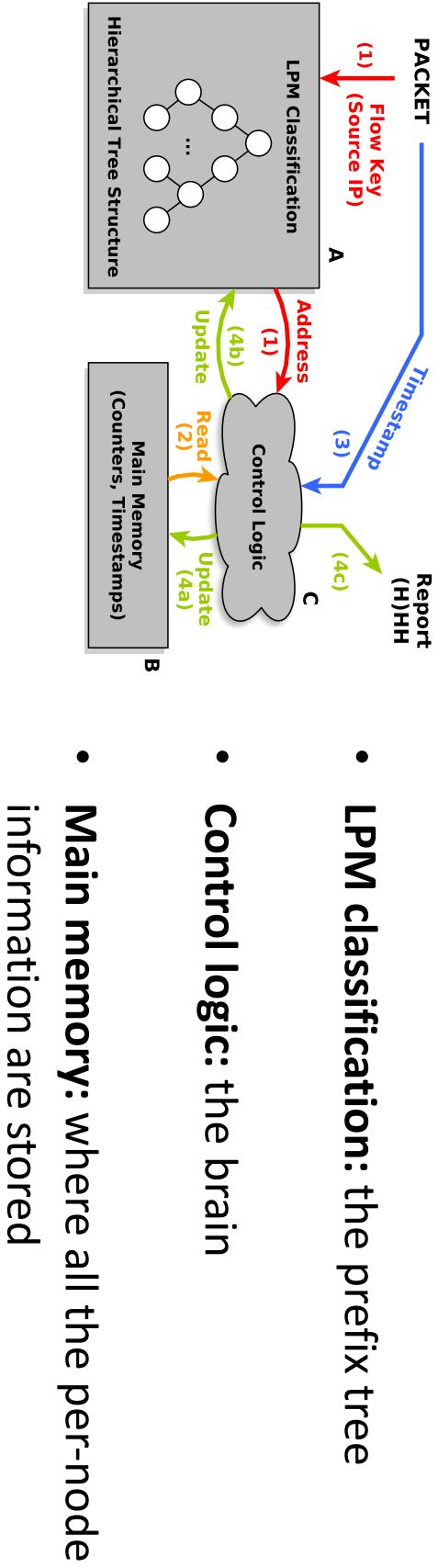


Elastic Trie implications

- The dataplane iteratively refine the responsible IP prefixes: the controller can receive a **flexible granularity information**.
- Each prefix tree layer can have a **different timeout**:
trade-off between trie building process and memory consumption.
- By looking at the **growing rate** of the trie it is possible to: identify changes in the traffic patterns.

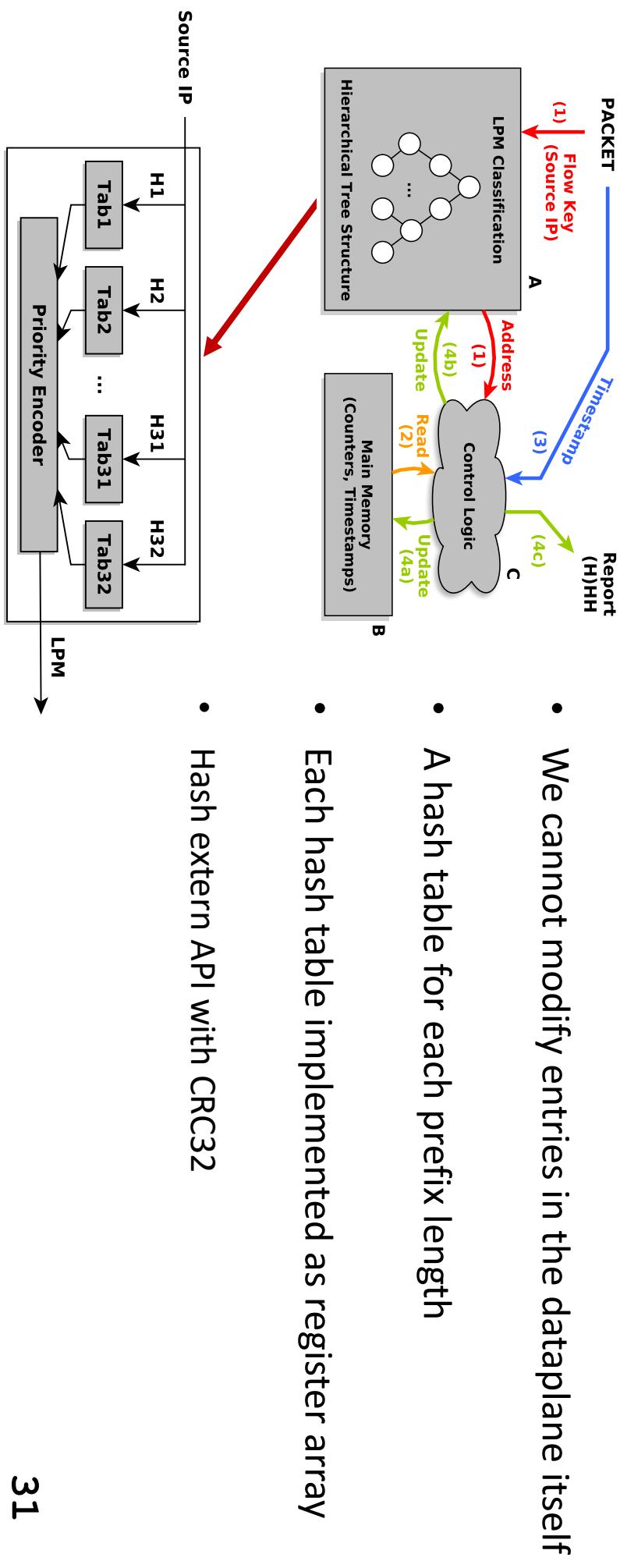


Elastic Trie in P4



- **LPM classification:** the prefix tree
- **Control logic:** the brain
- **Main memory:** where all the per-node information are stored

Elastic Trie in P4: LPM classification

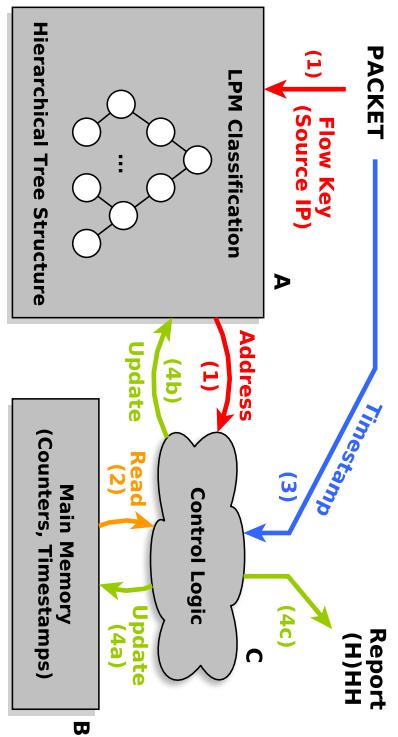


- We cannot modify entries in the dataplane itself

- A hash table for each prefix length
- Each hash table implemented as register array

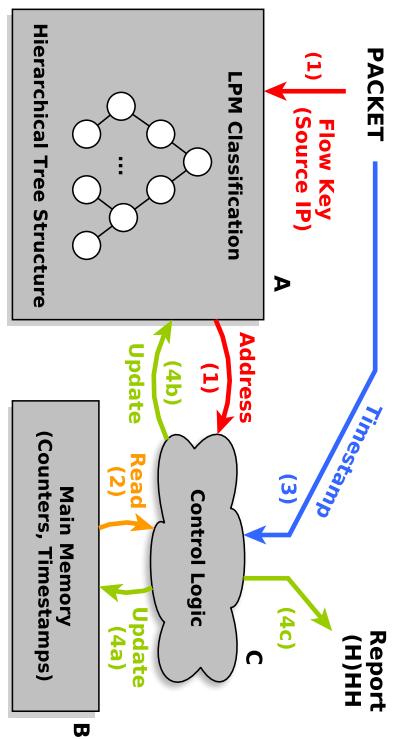
- Hash extern API with CRC32

Elastic Trie in P4: main memory



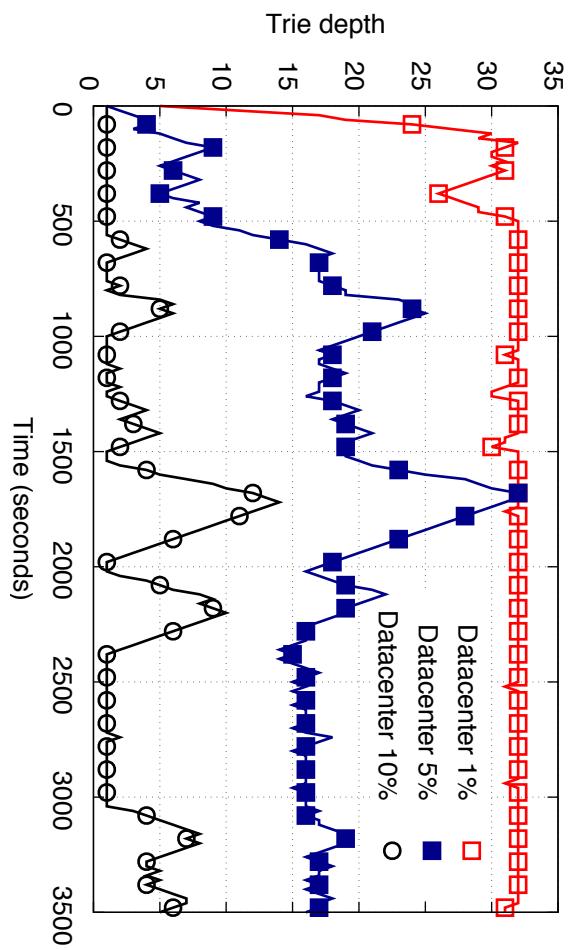
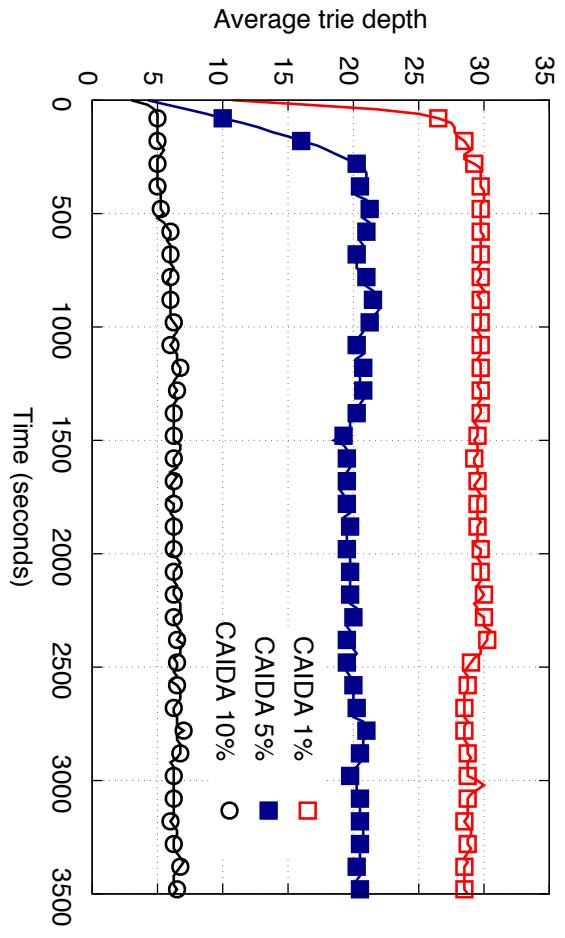
- The hash value of the LPM is the address to access a register that stores the node information

Elastic Trie in P4: control logic



- We compare node timestamp and packet timestamp
- It implements the node update logic, and the push-based mechanic with a *digest message*

Elastic Trie in action

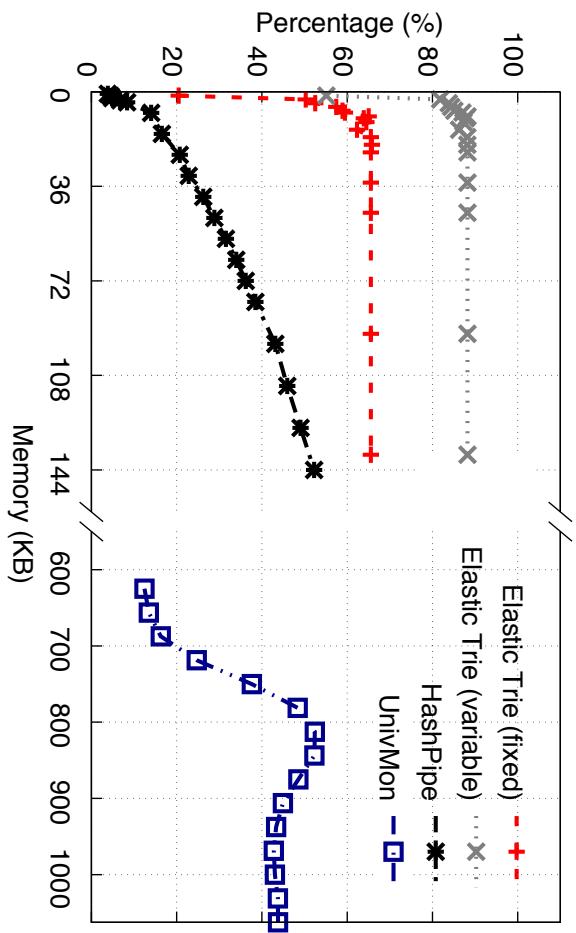


ISP

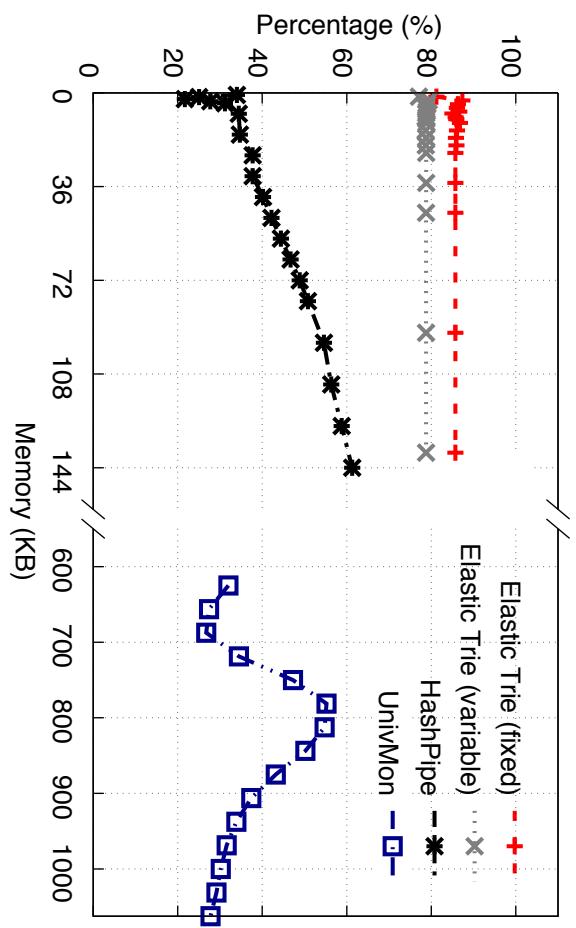
Datacenter

Elastic Trie in action

Recall



Precision

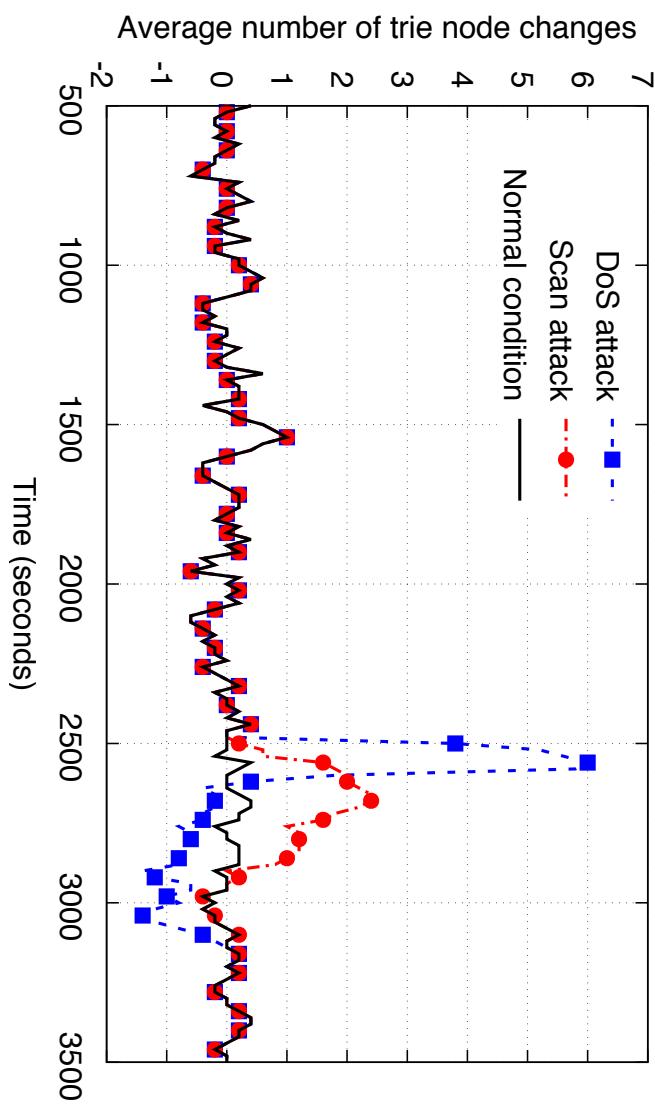


$\frac{\text{true positives}}{\text{true positives} + \text{false positives}}$

$\frac{\text{true positives}}{\text{true positives} + \text{false negatives}}$

$\text{true positives} + \text{false positives}$

Elastic Trie in action



Changes can be spotted!!!

Conclusions

- Elastic Trie enables **in-network detection** of traffic aggregates
- **Push-based monitoring** approach
- Suitable for **HH**, **HHH**, **Superspreaders** and **Change detection**.
- Low memory footprint!

