

Applied Algorithms Assignment 4

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Exercise 1 (10 pts)

$$1048579 = 2(2^{19} + 1) + 1 = 2 \cdot 524289 + 1$$

Using Miller-Rabin test with $r = 1$ and $d = 524289$, we choose a random $a = 3$. Using fast modulo exponentiation where:

$$\begin{aligned} a &= c \pmod{N} \\ a^2 &= (nN + c)^2 = (n^2N^2 + 2dnN + c^2) = c^2 \pmod{N} \\ 3^{524289} &\pmod{1048579} = 27320 \end{aligned}$$

And since $r = 1$, we find that this number is a composite.

Exercise 2 (10 pts)

- a) The probability of a hash collision on the k probe for item i is $\frac{i-1}{m}$. For $i < n < m/2$, $\frac{i-1}{m} < \frac{1}{2}$.
- b) The probability of a hash collision on $2\log(n)$ probe is $\frac{i-1}{m}^{2\log(n)}$. For $i < n < m/2$, $\frac{i-1}{m}^{2\log(n)} \leq 2^{-2\log(n)} = \frac{1}{n^2}$.
- c)

$$\begin{aligned} \Pr\{X > 2\log(n)\} &= \Pr\{\max_{1 \leq i < n} X_i > 2\log(n)\} \\ &= \Pr\{X_1 > 2\log(n)\} + \Pr\{X_2 > 2\log(n)\} + \dots + \Pr\{X_n > 2\log(n)\} \\ &\leq n \cdot \frac{1}{n^2} = \frac{1}{n} \end{aligned}$$

d) Assume $t = \log(n)$

$$\begin{aligned}
E[X] &= \sum_{j=1}^{\infty} jPr(X = j) \\
&= \sum_{j=1}^t jPr(X = j) + \sum_{j=t+1}^{\infty} jPr(X = j) \\
&\leq t \sum_{j=1}^t Pr(X = j) + \sum_{j=t+1}^{\infty} (j - t)Pr(X = j) \\
&= t + \sum_{j=t+1}^{\infty} (j - t)Pr(X = j) \\
&= \log(n) + \sum_{j=t+1}^{\infty} (j - t)Pr(X = j)
\end{aligned}$$

We use a summation range change to re-express the second term above:

$$E[X] \leq \log(n) + \sum_{j=t+1}^{\infty} \sum_{k=t}^{j-1} Pr(X = j)$$

Notice that the last summation can be re-ordered and once re-ordered is $1 - CDF(X)$:

$$\begin{aligned}
E[X] &\leq \log(n) + \sum_{k=t}^{\infty} \sum_{j=k+1}^{\infty} Pr(X = j) \\
&= \log(n) + \sum_{k=t}^{\infty} Pr(X > k)
\end{aligned}$$

We use result from a) and c) where $Pr(X > k) \leq \frac{n}{2^k}$:

$$\begin{aligned}
E[X] &\leq \log(n) + \sum_{k=t}^{\infty} \frac{n}{2^k} \\
&= \log(n) + \frac{n}{2^t} \sum_{j=0}^{\infty} 2^{-j} \\
&= \log(n) + n2^{-\log(n)} \times 2 \\
&= \log(n) + 2 \\
&= O(\log(n))
\end{aligned}$$

NOTE: I found a lot of help with this solution online.

Exercise 3 (10 pts)

You can logical OR the two halves of the bloom filter together. Using the same k hashes, the bit location can then be applied with modulo $n/2$. This has the effect of increasing the false positive rate from $f_1 = (1 - e^{-\frac{kn}{m}})^k$ to $f_2 = ((1 - e^{-\frac{2kn}{m}}))^k$.

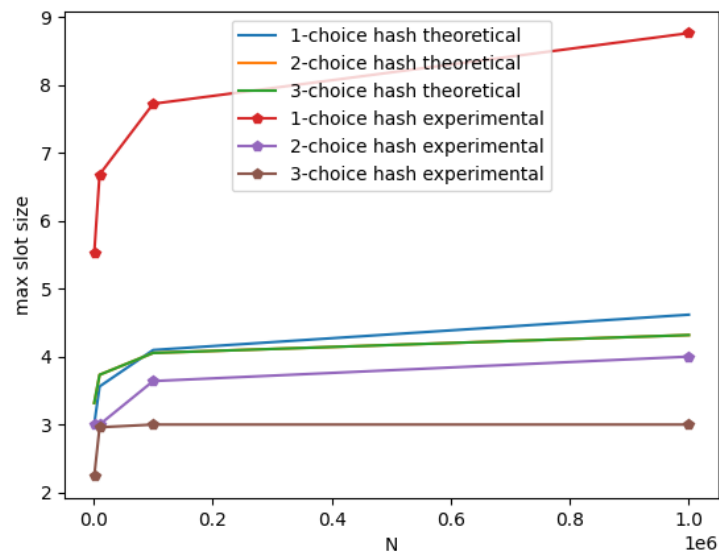


Figure 1: [1,2,3]-choice hashing and max items per slot

Exercise 4 (10 pts)

The results of 2-choice and 3 choice hashing seem to be quite a bit better than 1-choice hashing. Although even for upwards of a million items, 1-choice hashing seems to be not too bad.

Code

```
#!/usr/bin/python
import math
import random
import matplotlib.pyplot as plt

class HashSimulator():
    def __init__(self, N, hashes):
        self._N = N
        self._hashes = hashes

    def simulate(self, runs=1):
        maxSlotSizeAvg = 0
        for _ in range(runs):
            maxSlotSize = 0
            hash_table = [ 0 for i in range(self._N) ]
            for _ in range(self._N):
                indices = [ int(random.uniform(0, self._N))
                           for i in range(self._hashes) ]
                index_sizes = [ (i, hash_table[i]) for i in indices ]
                index = min( index_sizes, key=lambda x: x[1] )[0]
                hash_table[index] += 1
                if hash_table[index] > maxSlotSize:
                    maxSlotSize = hash_table[index]
            maxSlotSizeAvg += maxSlotSize
        return maxSlotSizeAvg / runs

if __name__ == "__main__":
    random.seed()
    N = [1000, 10000, 100000, 1000000]
    hashes = [1, 2, 3]
    runs = 25
    experimental_results = []
    theoretical_results = []
    for h in hashes:
        experimental_results.append([])
        theoretical_results.append([])
        for n in N:
            if h == 1:
                theoretical_results[-1].append(
                    math.log2(n)/math.log2(math.log2(n)))
            else:
                theoretical_results[-1].append(math.log2(math.log2(n)))
            hs = HashSimulator(n,h)
            res = hs.simulate(runs)
            print(f"hashes: {h} size: {n} avg: {res}")
            experimental_results[-1].append(res)
```

```

print("all experimental results:", experimental_results)
print("all theoretical results:", theoretical_results)
plt.figure()
plt.plot(N, theoretical_results[0], "-",
        label="1-choice hash theoretical")
plt.plot(N, theoretical_results[1], "-",
        label="2-choice hash theoretical")
plt.plot(N, theoretical_results[2], "-",
        label="3-choice hash theoretical")
plt.plot(N, experimental_results[0], "-p",
        label="1-choice hash experimental")
plt.plot(N, experimental_results[1], "-p",
        label="2-choice hash experimental")
plt.plot(N, experimental_results[2], "-p",
        label="3-choice hash experimental")
plt.ylabel("max slot size")
plt.xlabel("N")
plt.legend()
plt.show()

```