## Math 578B – Fall 2015 – Homework #1

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## due 1 September

1. Let X be a Markov chain on  $\mathcal{A} = \{0,1\}$  with transition probabilities

$$\mathbb{P}\{X_{t+1} = 1 \mid X_t = 0\} = \alpha \tag{1}$$

$$\mathbb{P}\{X_{t+1} = 0 \mid X_t = 1\} = \beta. \tag{2}$$

a. Show that if  $X_0 = 0$  then the distribution of

$$\tau := \min\{n > 0 : X_n = 1\} \tag{3}$$

is geometric.

b. Find the spectral decomposition of the chain and use it to find an expression for

$$\mathbb{P}\{X_n = 1 \mid X_0 = 0\}. \tag{4}$$

- c. What happens to the spectral expansion if  $\alpha + \beta = 1$ ? Explain.
- **2.** For the Markov chain in (1), suppose we choose a random starting state with  $\mathbb{P}\{X_1=0\}=\beta/(\alpha+\beta)$ , and run it for n steps, producing a string of n digits,  $X=X_1,\ldots,X_n$ . Let Y be the reversed string, i.e.,  $Y_k=X_{n+1-k}$  for  $1\leq k\leq n$ .
  - a. For a given string of digits  $a_1, \ldots, a_n \in \{0, 1\}^n$ , what is  $\mathbb{P}\{Y_1 = a_1, \ldots, Y_n = a_n\}$ ? (Don't use the spectral decomposition.)
  - b. Suppose we generate a string of length n as in (a) but it is reversed with probability 1/2: formally, let  $\theta = H$  with probability 1/2 and  $\theta = T$  otherwise, and define

$$Z = \begin{cases} X & \text{if } \theta = H \\ Y & \text{otherwise.} \end{cases}$$
 (5)

A colleague has developed a procedure that guesses whether the string was reversed, i.e., a function  $f:\{0,1\}^n \to \{H,T\}$  that is given Z and guesses the value of  $\theta$ . Show that

$$\mathbb{P}\{f(Z) = \theta\} = 1/2. \tag{6}$$

**3.** Let  $X_n$  be the simplified "polymerase complex assembly" Markov chain defined in class, with transition matrix (where "†" means transcription):

$$P = \begin{array}{c} \varnothing & \alpha & \beta & \alpha + \beta & \text{pol} & \dagger \\ \alpha & k_{\alpha} & k_{\beta} & 0 & 0 & 0 \\ k_{\alpha} & * & 0 & k_{\beta} & 0 & 0 \\ \alpha + \beta & k_{\beta} & 0 & * & k_{\alpha} & 0 & 0 \\ \text{pol} & 0 & k_{\beta} & k_{\alpha} & * & k_{\text{pol}} & 0 \\ \dagger & 0 & 0 & 0 & 0 & * & 1 \\ 0 & 0 & 0 & 1 & 0 & * \end{array} \right]$$
(7)

Here the "\*"s on the diagonal are set so that rows sum to 1. Let

$$\tau = \min\{n \ge 0 : X_n = \dagger\}. \tag{8}$$

a. Set  $k_{\alpha} = k_{\beta} = 0.2$  and  $k_{\text{pol}} = 0.5$ , and compute numerically  $\mathbb{E}[\tau]$  for all starting states. (And, explain how you do it!)

Now, since  $\mathbb{E}^a[\tau] = h$  solves  $(I - Q)^{-1}1$  where Q is

pander::pander(M[1:5,1:5])

0.6	0.2	0.2	0	0
0.2	0.6	0	0.2	0
0.2	0	0.6	0.2	0
0	0.2	0.2	0.1	0.5
0	0	0	0	0

the solution is

```
h <- solve( diag(5) - M[1:5,1:5], rep(1,5) )
pander::pander(h)

19, 16.5, 16.5, 9 and 1
```

b. Compute numerically the stationary distribution. What is the long-term average rate of transcription (i.e., mean number of visits to † per unit time)?

The stationary distribution  $\pi$  is the first left eigenvector of P:

```
pivec <- eigen(t(M))$vectors[,1]
pivec <- Re( zapsmall(pivec)/sum(pivec) )</pre>
```

and the mean time between visits to  $\dagger$  is 9.9999983.

c. Verify your computations in (a) and (b) using a simulation of the chain.