

Lecture Notes in Category Theory

Paul Ossientis

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Chapter 1

Category

1.1 Small Category

Before we define a category in full generality, we shall focus our attention on the notion of *small category*. This notion is interesting to us because while it essentially describes the notion of *category* itself, it remains simple enough to be compared with various other algebraic structures. For example, consider the case of a monoid: a monoid is essentially a set M together with a binary relation \circ defined on M which is associative, and an element e of M which acts as an identity element for \circ . In short a monoid is a tuple (M, \circ, e) containing some data, and which satisfy certain properties. The same is true of a *small category*: it is also a tuple containing some data, and which satisfy certain properties:

Definition 1 We call small category any tuple $(\text{Ob}, \text{Arr}, \text{dom}, \text{cod}, \text{id}, \circ)$ with:

- (1) Ob is a set
- (2) Arr is a set
- (3) $\text{dom} : \text{Arr} \rightarrow \text{Ob}$ is a function
- (4) $\text{cod} : \text{Arr} \rightarrow \text{Ob}$ is a function
- (5) $\text{id} : \text{Ob} \rightarrow \text{Arr}$ is a function
- (6) $\circ : \text{Arr} \times \text{Arr} \rightarrow \text{Arr}$ is a partial function
- (7) $g \circ f$ is defined $\Leftrightarrow \text{cod}(f) = \text{dom}(g)$
- (8) $\text{cod}(f) = \text{dom}(g) \Rightarrow \text{dom}(g \circ f) = \text{dom}(f)$
- (9) $\text{cod}(f) = \text{dom}(g) \Rightarrow \text{cod}(g \circ f) = \text{cod}(g)$
- (10) $\text{cod}(f) = \text{dom}(g) \wedge \text{cod}(g) = \text{dom}(h) \Rightarrow (h \circ g) \circ f = h \circ (g \circ f)$
- (11) $\text{dom}(\text{id}(a)) = a = \text{cod}(\text{id}(a))$
- (12) $\text{dom}(f) = a \Rightarrow f \circ \text{id}(a) = f$
- (13) $\text{cod}(f) = a \Rightarrow \text{id}(a) \circ f = f$

where (7) – (13) hold for all $f, g, h \in \text{Arr}$ and $a \in \text{Ob}$:

So if $\mathcal{C} = (\text{Ob}, \text{Arr}, \text{dom}, \text{cod}, \text{id}, \circ)$ is a small category, we have two sets Ob and Arr together with some structure defined on these sets. This feels very much like a monoid, except that we have two sets instead of one and it all looks more complicated. The set Ob is called the *set of objects* of the small category \mathcal{C} and is denoted $\text{Ob } \mathcal{C}$, while the set Arr is called the *set of arrows* of the small category \mathcal{C} and is denoted $\text{Arr } \mathcal{C}$. An element $x \in \text{Ob } \mathcal{C}$ is called an *object* of \mathcal{C} , while an element $f \in \text{Arr } \mathcal{C}$ is called an *arrow* of \mathcal{C} .

As part of the structure defined on the small category \mathcal{C} , we have two functions $\text{dom} : \text{Arr} \rightarrow \text{Ob}$ and $\text{cod} : \text{Arr} \rightarrow \text{Ob}$. Hence, given an arrow f of the small category \mathcal{C} , we have two objects $\text{dom}(f)$ and $\text{cod}(f)$ of the small category \mathcal{C} . The object $\text{dom}(f)$ is called the *domain* of f . The object $\text{cod}(f)$ is called the *codomain* of f . Note that an arrow f of the small category \mathcal{C} is simply an element of the set $\text{Arr } \mathcal{C}$. So it is itself a set but it may not be a function. The words *domain* and *codomain* are therefore overloaded as we are using them in relation to a set f which is possibly not a function. Whenever f is an arrow of the small category \mathcal{C} and a, b are objects, it is common to use the notation $f : a \rightarrow b$ as a notational shortcut for the equations $\text{dom}(f) = a$ and $\text{cod}(f) = b$. Once again, it is important to guard against the possible confusion induced by the notation $f : a \rightarrow b$ which does not mean that f is function. It simply means that f is an arrow with domain a and codomain b in the small category \mathcal{C} .

One of the main ingredients of the structure defining a small category \mathcal{C} is the partial function $\circ : \text{Arr} \times \text{Arr} \rightarrow \text{Arr}$, called the *composition operator* in the small category \mathcal{C} . The domain of this partial function is made of all ordered pairs (g, f) of arrows in \mathcal{C} for which $\text{cod}(f) = \text{dom}(g)$. As already indicated in definition (1), we use the infix notation $g \circ f$ rather than $\circ(g, f)$ and the arrow $g \circ f$ is called the *composition* of g and f . Once again, we should remember that the notation $g \circ f$ does not mean that g or f are functions. They are simply arrows in the small category \mathcal{C} . One key property of the composition operator \circ is the associativity postulated by (10) of definition (1). Note that if $f : a \rightarrow b$ and $g : b \rightarrow c$, then from properties (8) and (9) of definition (1) we obtain $g \circ f : a \rightarrow c$. Furthermore, if $h : c \rightarrow d$ we have $h \circ g : b \rightarrow d$ and the arrows $(h \circ g) \circ f$ and $h \circ (g \circ f)$ are therefore well-defined arrows with domain a and codomain d . This shows that the expression involved in the associativity condition (10) of definition (1) is always meaningful, involving terms which are well-defined provided $g \circ f$ and $h \circ g$ are themselves well-defined, i.e. provided $\text{cod}(f) = \text{dom}(g)$ and $\text{cod}(g) = \text{dom}(h)$.

Finally, as part of the structure defining the small category \mathcal{C} , we have a function $\text{id} : \text{Ob} \rightarrow \text{Arr}$ called the *identity operator* on the small category \mathcal{C} . Hence, for every object a of \mathcal{C} we have an arrow $\text{id}(a)$ called the *identity at a* . Looking at property (11) of definition (1) we see that $\text{id}(a) : a \rightarrow a$. In other words, the arrow $\text{id}(a)$ has domain a and codomain a . Furthermore, looking at properties (12) and (13) of definition (1), for every arrow $f : a \rightarrow b$, the composition arrows $\text{id}(b) \circ f$ and $f \circ \text{id}(a)$ are well-defined and both equal to f .

1.2 Category

The notion of *small category* defined in definition (1) is similar to that of any other algebraic structure the reader may be familiar with. It can safely be encoded in set theory as a tuple (which is a set) containing data (which are other sets) which satisfies certain properties. In set theory, everything is a set. A small category \mathcal{C} is a set, its collection of objects $\text{Ob } \mathcal{C}$ is a set, its arrows $\text{Arr } \mathcal{C}$ form a set, the functions dom , cod , id and the composition operator \circ are all sets (functions are typically encoded as sets of ordered pairs).

Category theory falls outside of set theory. While the definition of a *category* we provide below is formally identical to that of a small category, the object we are defining can no longer be encoded in general as an object of set theory. For example, say we want to speak about the *universe of all sets* or the *universe of all monoids*. These *universes* which are known as *classes* cannot be represented as sets. They are not objects of set theory. Or say we are working within the formal framework of a proof assistant such as Coq, Agda or Lean. These tools are based on type theory and do not fall within the scope of set theory. When defining a *category*, we assume some form of meta-theoretic context, some form of logic, some way of reasoning about objects which may not be sets, where some meaning is attached to the words *tuple*, *collection*, *equality* and *map*. This may sound all very fuzzy, yet we cannot be more formal at this stage.

Definition 2 We call category any tuple $(\text{Ob}, \text{Arr}, \text{dom}, \text{cod}, \text{id}, \circ)$ such that:

- (1) Ob is a collection with equality
- (2) Arr is a collection with equality
- (3) $\text{dom} : \text{Arr} \rightarrow \text{Ob}$ is a map
- (4) $\text{cod} : \text{Arr} \rightarrow \text{Ob}$ is a map
- (5) $\text{id} : \text{Ob} \rightarrow \text{Arr}$ is a map
- (6) $\circ : \text{Arr} \times \text{Arr} \rightarrow \text{Arr}$ is a partial map
- (7) $g \circ f$ is defined $\Leftrightarrow \text{cod}(f) = \text{dom}(g)$
- (8) $\text{cod}(f) = \text{dom}(g) \Rightarrow \text{dom}(g \circ f) = \text{dom}(f)$
- (9) $\text{cod}(f) = \text{dom}(g) \Rightarrow \text{cod}(g \circ f) = \text{cod}(g)$
- (10) $\text{cod}(f) = \text{dom}(g) \wedge \text{cod}(g) = \text{dom}(h) \Rightarrow (h \circ g) \circ f = h \circ (g \circ f)$
- (11) $\text{dom}(\text{id}(a)) = a = \text{cod}(\text{id}(a))$
- (12) $\text{dom}(f) = a \Rightarrow f \circ \text{id}(a) = f$
- (13) $\text{cod}(f) = a \Rightarrow \text{id}(a) \circ f = f$

where (7) – (13) hold for all $f, g, h \in \text{Arr}$ and $a \in \text{Ob}$:

So let $\mathcal{C} = (\text{Ob}, \text{Arr}, \text{dom}, \text{cod}, \text{id}, \circ)$ be a category: then \mathcal{C} is a *tuple* but it is no longer a tuple in a set theoretic sense. We assume given some logical framework where the notion of *tuple* is clear, even if not formally defined. Furthermore, We are no longer imposing that Ob should be a set, but are instead using the phrase *collection with equality*, whatever this may mean in our given logical context. So we shall still make use of the notation $\text{Ob } \mathcal{C}$ but this will now refer to the *collection* of all *objects* of the category \mathcal{C} . In fact, if a is an object of the category \mathcal{C} , we shall abuse notations somewhat by writing $a \in \text{Ob } \mathcal{C}$ or even simply $a \in \mathcal{C}$ to express the fact that a is an object of \mathcal{C} , being understood that this use of the set membership symbol ' \in ' does not mean anything is a set. Since we are stepping out of set theory, the objects of the category \mathcal{C} may not be sets themselves. They are simply members of the *collection* $\text{Ob } \mathcal{C}$. However, properties (7) – (13) of definition (2) are all referring to equalities between objects such that $\text{cod}(f) = \text{dom}(g)$. So it must be the case that the notion of *equality* be meaningful on the collection $\text{Ob } \mathcal{C}$. This explains our use of the phrase *collection with equality*: given $a, b \in \mathcal{C}$, the statement $a = b$ while not a set-theoretic equality is nonetheless assumed to be defined.

Chapter 2

Functor

Chapter 3

Natural Transformation

Chapter 4

Adjunction

4.1 Definition

Definition 3 We call adjunction an ordered pair (F, G) where F is a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ and G is a functor $G : \mathcal{D} \rightarrow \mathcal{C}$ while \mathcal{C} and \mathcal{D} are two locally-small categories for which there exists a natural isomorphism:

$$\alpha : \mathcal{D} \circ (F \times I_{\mathcal{D}}) \Rightarrow \mathcal{C} \circ (I_{\mathcal{C}^{op}} \times G)$$

in the functor category $[\mathcal{C}^{op} \times \mathcal{D}, \mathbf{Set}]$, where F also denotes $F : \mathcal{C}^{op} \rightarrow \mathcal{D}^{op}$.