Assignment Project Exam Help Add WeChat powcoder

21. Cache Problems

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EECS 370 – Introduction to Computer Organization – Fall 2020

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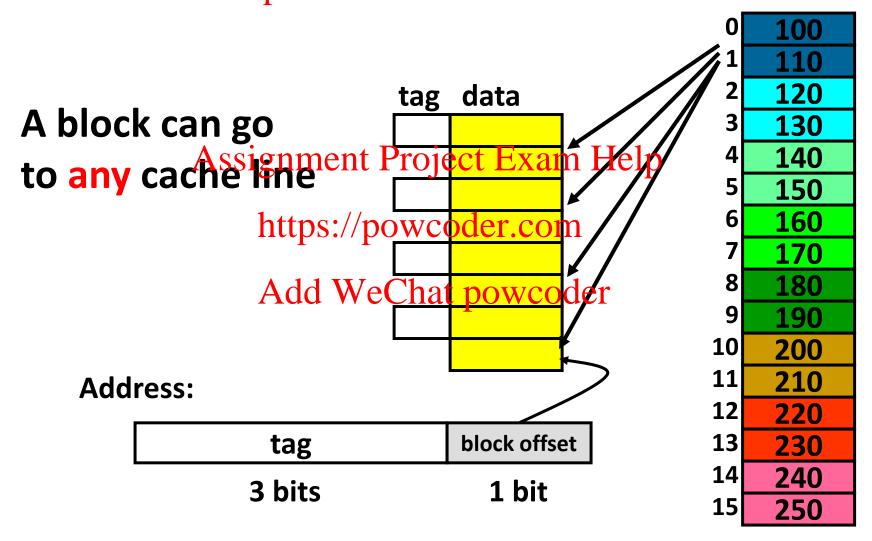
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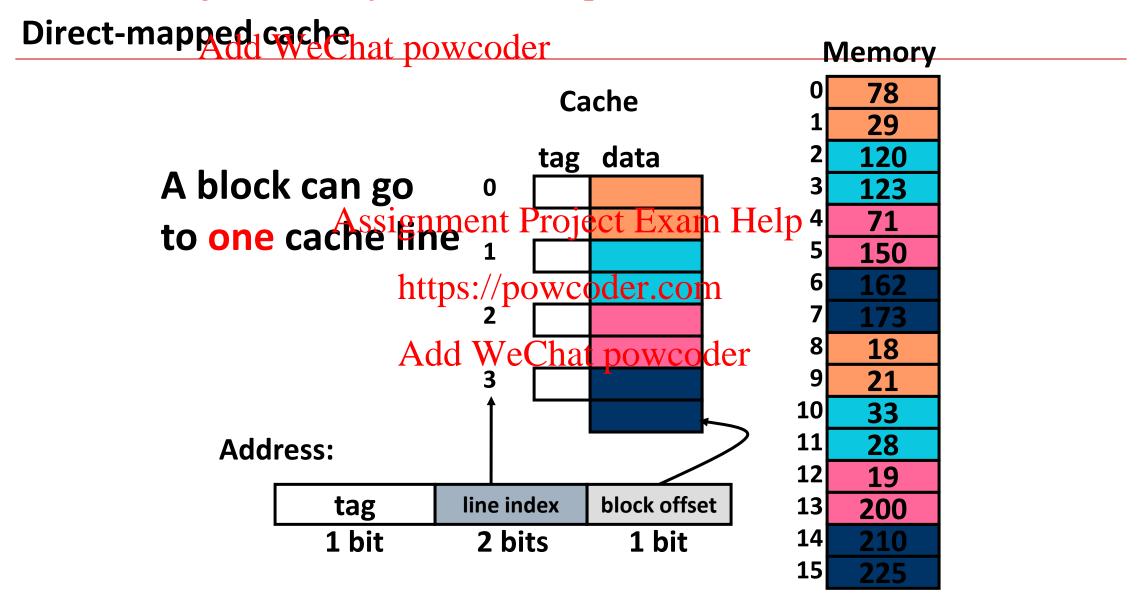
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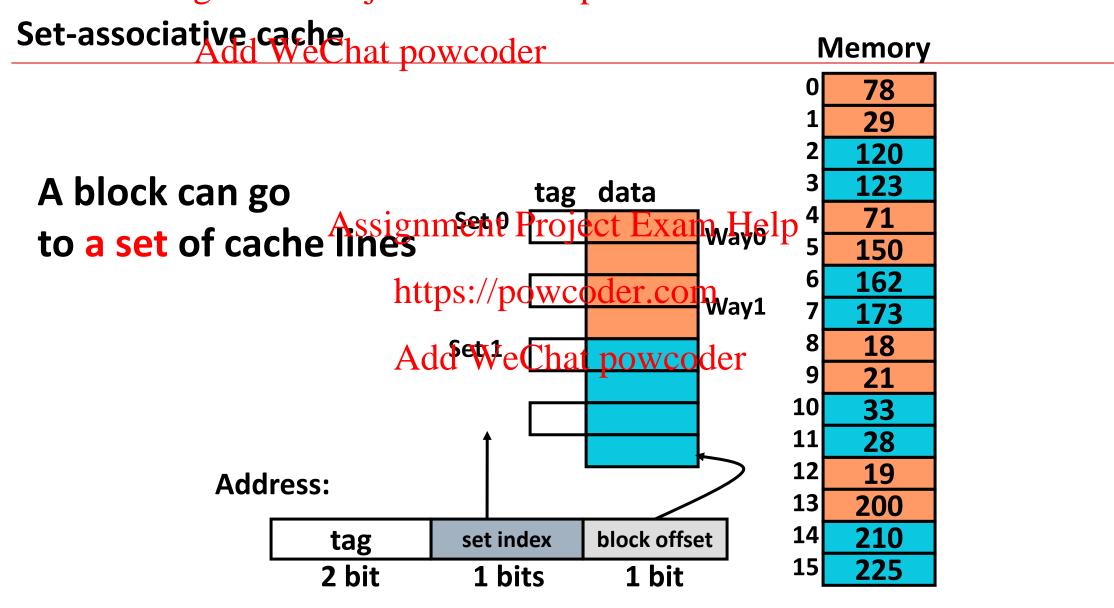
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Fully-associative cache powcoder









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How to solve problems on Caches Exam Help

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Programmer's view Chat powcoder

Address size: defined by ISA

e.g., 64-bit ISA has 64-bit addresses

Cache design defines how address is split into the following (hidden from the programmers/compiler):



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Byte vs word addressable powcoder

Byte addressable ISA: each address refers to a **byte** in memory

Word addressable ISA: each address refers to a **word** in memory.

(Annipardress in Pividuet by tes and Whell)

Common practice:

In 64-bit ISA, a "word" is 64-bits (8 bytes) oder.com
In 32-bit ISA, a "word" is 32-bits (4 bytes)

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This difference impacts how we determine block offset size:

Example: Block-offset for a 256-byte block:

In 64-bit **byte** addressable ISA: a 256-byte block has **256 bytes**. So block_offset_size is 8 bits In 64-bit **word** addressable ISA: a 256-byte block has **32 words**. So block_offset_size is 5 bits

tag	line index	block offset

Cache Organization: Fquations der

```
Block
```

```
#blocks = cache size / block_size #cache lines = #blocks
block_offset_size = log2(block_size)
```

Set

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#sets = ##ines/: #waysvcodtlines of thines per set)

Direct-mapped: #sets = #lines/1

2-way associative: #sets = Aides/WeChat powcoder

n-way associative: #sets = #lines/n

fully-associative: #sets = 1 (all lines are in 1 set)

set_index_size = log2(#sets)

Tag size = address size - set_index_size - block_offset_size

Calculating Average Memory Access Time (AMAT)

If we assume memory latency includes time to cache access and determine hit/miss:

AMAT = cache latency Antique ment employed tency Histrate

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If we do not assume memory latency includes time to cache access and determine hit/miss: Add WeChat powcoder

AMAT = cache latency + memory latency × miss rate

3C problem_{Add WeChat powcoder}

```
Simulate three different caches:
```

```
(same block size)
Infinite cache
```

(same block size and cache size) Fully associative cache

Assignment Project Exam Help Given cache

https://powcoder.com
How to classify a miss in the given cache?

Is it also a miss in the iAddt Wacher powcoder

Yes, then Compulsory miss

No. Is it a miss in the fully associative cache?

Yes, then capacity miss

No, then conflict miss

Simulating caches WeChat powcoder

Figure out fields in the address:

tag, set index, block-offset

Determine (tag, set-index) of given steel p

https://powcoder.com Sequence of (tag, set-index) is sufficient to simulate a cache. Data is not necessary.

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Keep track of LRU rank of cache lines within each cache set for set-associative caches. (replacement policies other than LRU exist)

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Sample problems ignmente Project Exam Help

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Problem 0: Cache Block Size powcoder

Assume 32-bit word-addressable ISA, and a fully-associative cache of size 64 KB. If the block size 32 bytes. Determine the size of block offset and tag.

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Problem 0: Cache Block Size powcoder

Assume 32-bit word-addressable ISA, and a fully-associative cache of size 64 KB. If the block size 32 bytes. Determine the size of block offset and tag.

A 32-byte block contains 32/ws/psiznemperofeetal That is 18216 = 8 words.

 \Rightarrow block_offset is log(8) = 3 bits. https://powcoder.com

Tag = 32 (address size) – 3 bits = 29 dd WeChat powcoder

No set-index, because the cache is fully-associative (1 set).

Practice Problem 1: CPI with caches

The *grinder* application is run on the LC2k with full data forwarding. All branches are predicted not-taken. Assume the following.

45% R-type

20% Branches (of these, 40% of branches are taken.)

15% Loads A (Sthese, 10% of branches are taken.)

20% Stores

https://powcoder.com Cache performance:

Assume cache hit is 1 cycle

Add WeChat powcoder I-cache miss rate: 3%

D-cache miss rate: 6%

(assume misses do not overlap)

Main memory latency: 100 ns

Processor clock frequency: 500 MHz

What is the CPI of *grinder* on the LC2k?

Practice Problem 1: Solution powcoder

What is the CPI of *grinder* on the LC2k?

```
Cycle time: 2 ns per cycle (ASSIGNMENT Project Exam Help Stalls per cache miss = 100 ns / 2ns = 50 cycles

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CPI = 1 + data hazard stalls + control hazard stalls + icache stalls + dcache stalls

CPI = 1 + 0.15*0.50*1 + 0.20*0.00*50
```

Practice Problem 2: CPI with caches

The 370 application is run on the LC2k. Assume the following:

30% of instructions are loads/stores

Cache:

I-cache and D-cache: Project - way associative: 1-cycle access latency I-cache hit rate: 95%

D-cache hit rate: https://powcoder.com

(assume misses do not overlap)

Main memory latency: Add We Chatspowcoder

Processor clock frequency:

Assume CPI is 1, if cache has 100% hit rate and 1 cycle latency.

continued....

Practice Problem 2: CPI w/ Caches 2

Suppose you have 2 options for the next generation processor, which do you pick?

Option 1: Double the clock frequency—aspume this will increase your memory latency to 40 cycles. Also assume a base CPI of 1 can still be achieved after this change.

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Option 2: Double the size of your caches, this will increase the instruction cache hit rate to 98% and the data cache hit rate to 95% (Aswere that hip attention) at the cache hit rate to 95% (Aswere that hip attention) are the cache hit rate to 98%

Practice Problem 2: Solution Add WeChat powcoder

200 MHz \rightarrow clock cycle = 5 ns.

Option 1: (double clock freq to 400 MHz, so new cycle time is 2.5 ns)

```
CPI = baseCPI + IcacheSta Assignmenta Project Exam Help
CPI = 1.0 + 0.05*40 + 0.3*0.1*40 = 4.2
```

Execution time = 4.2 * NumInstrps: 2.BRYCOSERSON umInstrs

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Option 2 (icache/dcache miss rates lowered to 2% and 5%)

```
CPI = baseCPI + IcacheStallCPI + DcacheStallCPI

CPI = 1.0 + 0.02*20 + 0.3*0.05*20 = 1.7
```

Execution time = 1.7 * NumInstrs * 5ns = 8.5ns * NumInstrs

Therefore, Option 2 is the better choice

Practice Problem 3: Write-back vs Write-through

Say you have the following:

A program that generates 2 Billion loads and 1 Billion stores, each 4 bytes in size.

A cache with a 32-byte block which gets a 95% hit rate on that program.

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How many bytes of memory would be read and written if:

- 1. We had no cache? https://powcoder.com
- 2. We had a write-though cache with a no-write allocate policy?
- 3. We had a write-back cache with a write-allocate policy? (Assume 25% of all misses result in a dirty eviction)

Practice Problem 3: Solution (1/3)

No-cache

```
All stores go to memory and are 4 bytes each
     Writes: 1 billion stores* 4 bytes = 4 billion bytes
All loads go to memory and are 4 bytes each oject Exam Help Reads: 2 billion loads* 4 bytes = 8 billion bytes
```

https://powcoder.com

Cache: Write-through, no write-allocate

All stores still go to memory. Cache block not read from memory, due to no write-allocate

```
Writes: 1 billion stores* 4 bytes = 4 billion bytes
```

Only loads that miss in the cache go to memory. But they read the full cache block.

Reads: 2 billion loads* 0.05* 32 bytes = 3.2 billion bytes

Practice Problem 3: Solution (2/3) Add WeChat powcoder

Cache: Write-back, write-allocate

Total memory accesses: 1 billion stores + 2 billion loads = 3 billion memory accesses

Data read from memory Assignment Project Exam Help

Store (due to write-allocate policy) and load misses result in a cache block being read.

Data read from memory: 3 billion memory accesses * 0.05* 32 bytes = 4.8 billion bytes

Data written to memory

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Store and load cache misses result in a cache block getting evicted.

Given: 25% of cache misses result in dirty eviction (write 32-byte block to memory).

Data written to memory = 3 billion memory accesses * 0.05* 0.25 * 32 bytes = 1.2 billion bytes

Practice Problem 3: Solution (3/3)

Write-back, write-allocate (data writes)

When do we write to memory?

Only on dirth soit gament Brojec by Exampled Stores.

Recall we are assuming 25% of all evictions are of dirty data. https://powcoder.com

Total memory accesses: 1 hillion stores + 2 billion loads = 3 billion memory accesses 25% of cache misses result in dirty eviction (write 32-byte block to memory).

3 billion memory accesses * 0.05* 0.25 * 32 bytes = 1.2 billion bytes

Practice Problem 4: Reverse engineer the cache!

Here is the series of address references (in hex) to a cache of size 512 bytes. You are asked to **determine the configuration of the cache**. Assume 12-bit addresses.

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0x30f – Miss https://powcoder.com

0x510 – Miss

0x31f – Hit Add WeChat powcoder

0x72d – Miss

0x72f - Hit

0x320 – Miss

0x520 - Miss

0x720 - Miss

Block size: ?

Associativity: ?

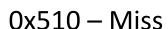
Number of sets: ?

Practice Problem 4: Solution coder

Here is the series of address references (in hex) to a cache of size 512 bytes. You are asked to **determine the configuration of the cache**. Assume 12-bit addresses

0x310 – Miss

0x30f – Miss **4**



0x31f – Hit 🛑

0x72d - Miss

0x72f - Hit

0x320 - Miss

0x520 - Miss

0x720 - Miss

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Determine block size

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Practice Problem 4: Solution coder

Here is the series of address references (in hex) to a cache of size 512 bytes. You are asked to **determine the configuration of the cache**. Assume 12-bit addresses

0x310 – Miss 🛑

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0x510 - Miss

0x31f – Hit 🛑

0x72d - Miss

0x72f - Hit

0x320 – Miss

0x520 - Miss

0x720 - Miss

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Determine block size

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First hit must be brought in by another miss

Take desest address: 0x310, 80 know block size must be at least 16 bytes so 0x31f brought in when 0x310 miss occurs

Now, is the block size larger? Know that 0x30f was a miss, thus 0x310 and 0x30f not in the same block. Thus, block size must be <= 16 bytes

Thus Block Size = 16 bytes

Practice Problem 4: Solution (2) er

Here is the series of address references (in hex) to a cache of size 512 bytes. You are asked to **determine the configuration of the**cache. Assume 12-bit addresses

Determine associativity

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0x310 – Miss

0x30f - Miss

0x510 – Miss 📥

0x31f – Hit

0x72d - Miss

0x72f – Hit 🛑

0x320 – Miss

0x520 – Miss **(**

0x720 - Miss <

https://powcoder.com

Practice Problem 4: Solution (2) er

Here is the series of address references (in hex) to a cache of size

512 bytes. You are asked to determine the configuration of the

cache. Assume 12-bit addresses

Determine associativity

0x310 – Miss

0x30f - Miss

0x510 – Miss 🛑

0x31f – Hit

0x72d - Miss

0x72f – Hit

0x320 - Miss

0x520 - Miss

0x720 - Miss

Assignment Project Exam Help Assume direct mapped: 3-bit tag, 5-bit index, 4-bit offset.

If direct mapped 0x310 and 0x510 would both map to index 9,

https://powgodercollom be a hit. So, not direct mapped.

Add We Charles the green accesses, and allows 0x31f to be a hit.

What about > 2-way associative?

Now we also know that 0x720 is a miss even though 3 accesses earlier 0x72f was a hit, and thus it is in the cache. The intervening 2 accesses must kick it out, 0x320 and 0x520. Both go to set 2. If the associativity was > 2, then 0x720 would be a hit. So, must conclude that cache is 2-way associative.

Lastly, number of sets = 512 / (2 * 16) = 16