

# The matrix 2-norm

If  $A$  is sym ( $A^T = A$ ),  $\|A\|_2 \geq |\lambda_{\max}(A)|$ .

◁ Example:

The 2-norm is the natural norm for LSQ problems (minimizing  $\|r\|_2$ )  $\Rightarrow$  can no longer avoid the matrix 2-norm :  
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for a square matrix  $A$  (see MatrixNorms for proof)

$$\|A\|_2 = \sqrt{\lambda_{\max}(A^T A)}$$

$\lambda_{\max}(A^T A)$  is the largest eigenvalue of  $A^T A$   
(all eigenvalues are positive since  $A^T A$  is positive definite).

# Singular value decomposition SVD

It is easier to characterize the condition number in the 2-norm in terms of the **singular values** of  $\mathbf{A}$ . To do that we need the

## Definition

A  $m \times n$  real matrix  $\mathbf{A}$  has the **singular value decomposition**

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where

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- 1  $\mathbf{U}$  is  $m \times m$  orthogonal matrix
- 2  $\Sigma$  is a diagonal  $m \times n$  real matrix
- 3  $\mathbf{V}$  is  $n \times n$  orthogonal matrix

The non-negative diagonal entries  $\{\sigma_k \geq 0\}$  in  $\Sigma$  are called the **singular values** of  $\mathbf{A}$ .

In our case, where  $m > n$ , there are  $n$  positive singular values  $\sigma_1 \geq \sigma_2 \dots \geq \sigma_n > 0$ , if  $\mathbf{A}$  is of full rank.

$$m \times n \quad m \times m \quad m \times n \quad n \times n$$

$$\left[ \quad \right] = \left[ \quad \right] \left[ \begin{matrix} b_1 & b_2 & \dots & b_n \\ 0 & & & \end{matrix} \right] \left[ \quad \right]$$

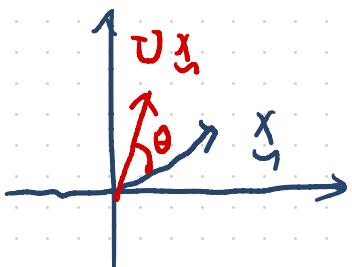
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$$A = U \Sigma V^T$$

$U$  &  $V$  are orthogonal, i.e.  $U^T U = V V^T = I_{m \times m}$ ,  $\Leftrightarrow U^{-1} = U^T$   
 $V^T V = V V^T = I_{n \times n}$   $\Leftrightarrow V^{-1} = V^T$

Example:

$$U = \begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix}$$



orthogonal matrix  
 $\Leftrightarrow$  rotation transformation

$m > n$

$$b_1 \geq b_2 \geq b_3 \geq \dots \geq b_n \geq 0$$



Singular Value of matrix A

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If A is full rank <https://powcoder.com>

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**Lemma:** If  $U$  is an orthogonal matrix, then  $\|Ux\|_2 = \|x\|_2$ .

**Proof:**  $\|Ux\|_2^2 = (Ux)^T Ux = x^T \underline{U^T} Ux = x^T x = \|x\|_2^2$

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**Lemma:** If  $U \in \mathbb{R}^{m \times m}$  and  $V \in \mathbb{R}^{n \times n}$  are orthogonal, then  $\|UAV\|_2 = \|AV\|_2 = \|A\|_2$ .

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**Proof:**  $\|UAV\|_2 = \max_{x \neq 0} \frac{\|UAVx\|_2}{\|x\|_2} = \max_{x \neq 0} \frac{\|AVx\|_2}{\|x\|_2} = \|A\|_2$

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$$\|AV\|_2 = \max_{x \neq 0} \frac{\|AVx\|_2}{\|x\|_2} \stackrel{y=Vx}{=} \max_{y \neq 0} \frac{\|AVy\|_2}{\|y\|_2} = \max_{y \neq 0} \frac{\|Ay\|_2}{\|y\|_2} = \|A\|_2$$

$$\|y\|_2 = \|Vx\|_2 = \|x\|_2$$

# The matrix 2-norm

Then from the definition above, we get

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Proof:

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$$A = U \Sigma V^T$$

$$\|A\|_2 = \left\| U \Sigma V^T \right\|_2 = \|\Sigma\|_2. \leq 6_1 \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}$$

$$\Sigma = \begin{bmatrix} \sigma_1 & & \\ & \sigma_2 & \\ & \ddots & \ddots & \sigma_n \\ & 0 & & \\ & & & \in \mathbb{R}^{m \times n} \end{bmatrix}$$

$$\|\Sigma\|_2 = \max_{x \neq 0} \frac{\|\Sigma x\|_2}{\|x\|_2} = \max_{x \neq 0} \frac{\sqrt{\sigma_1^2 x_1^2 + \dots + \sigma_n^2 x_n^2}}{\sqrt{x_1^2 + \dots + x_n^2}} \leq \sigma_1$$

$$\|\Sigma\|_2 = \max_x \frac{\|\Sigma x\|_2}{\|x\|_2} \geq \frac{\|\Sigma e_1\|_2}{\|e_1\|_2} = \frac{6}{1} = 6,$$

$\sum e_i =$  Assignment Project Exam Help  
 $\begin{bmatrix} 6 \\ 6 \\ 6 \\ 6 \\ 6 \\ 6 \\ 6 \\ 6 \\ 6 \\ 6 \end{bmatrix}$   
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# The pseudoinverse

The pseudoinverse  $\mathbf{A}^\dagger = (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T$  can be expressed in terms of the SVD of  $\mathbf{A}$ :

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$$\mathbf{A}^\dagger = \mathbf{V} \Sigma^\dagger \mathbf{U}^T$$

where  $\Sigma^\dagger$  is the  $n \times m$  diagonal matrix, with entries  $\{1/\sigma_k\}$ .  
Proof:

$$A^+ = (A^T A)^{-1} A^T$$

$$A = V \Sigma V^T$$

$$\Sigma = \begin{bmatrix} 6_1 & & \\ & \ddots & \\ & & 6_n \\ & & 0 \end{bmatrix}$$

$$A^T A = V \Sigma^T \underbrace{U^T}_{\text{---}} U \Sigma V^T$$

$$= V \Sigma^T \Sigma V^T$$

$$= V \Lambda V^T$$

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$$\Lambda = \Sigma^T \Sigma = \text{diag} \left( \begin{bmatrix} 6_1^2 & & \\ & \ddots & \\ & & 6_n^2 \end{bmatrix} \right) = \begin{bmatrix} 6_1^2 & & \\ & \ddots & \\ & & 6_n^2 \end{bmatrix}$$

https://powcoder.com value decomposition  
Add WeChat powcoder of  $A^T A$

$$(A^T A)^{-1} = (V \Lambda V^T)^{-1} = (V^T)^{-1} \Lambda^{-1} V^{-1} = V \Lambda^{-1} V^T$$

$$\Lambda^{-1} = \begin{bmatrix} 1/6_1^2 & & \\ & \ddots & \\ & & 1/6_n^2 \end{bmatrix}$$

the largest singular value of  $A^T A = 6_1^2$   
-- smallest - .

$$(A^T A)^{-1} A^T = V \Delta^{-1} V^T (U \Sigma V)^{-1}$$

$$= V \Delta^{-1} \underbrace{V^T V}_{\sim} \Sigma^{-1} U^T$$

$$\Sigma = \begin{bmatrix} \sigma_1 & & & \\ & \ddots & & \\ & & \sigma_n & \\ & 0 & & \end{bmatrix}$$

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$\Sigma^+$   
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$$\Sigma^+ = \Delta^{-1} \Sigma^{-1} = \left[ \begin{smallmatrix} 1/\sigma_1^2 & & & \\ & \ddots & & \\ & & 1/\sigma_n^2 & \\ & & & \end{smallmatrix} \right] \left[ \begin{smallmatrix} \sigma_1 & & & \\ & \ddots & & \\ & & \sigma_n & \\ & 0 & & \end{smallmatrix} \right]$$

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$$= \left[ \begin{smallmatrix} 1/\sigma_1 & & & \\ & \ddots & & \\ & & 1/\sigma_n & \\ & & & 0 \end{smallmatrix} \right]$$

# The condition number of a rectangular matrix

$$\mathbf{A}^\dagger = \mathbf{V} \Sigma^\dagger \mathbf{U}^\top \quad \|\mathbf{A}^\dagger\|_2 = \|\Sigma^\dagger\|_2$$

By the same argument, we get

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 $\|\mathbf{A}^\dagger\|_2 = 1/\sigma_n(\mathbf{A})$

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By the same argument as before, the condition number of a rectangular matrix is given by  $\kappa_2(\mathbf{A}) = \|\mathbf{A}\|_2 \|\mathbf{A}^\dagger\|_2 = \sigma_1(\mathbf{A})/\sigma_n(\mathbf{A})$

Proof:



the ratio of the largest singular value

to the smallest singular value.

# Sensitivity of the normal equations

The normal equations

$$\mathbf{A}^T \mathbf{A} \mathbf{x} = \mathbf{A}^T \mathbf{b}$$

are a linear system, so the sensitivity is given by the condition number of  $\mathbf{A}^T \mathbf{A}$ . But

$$\kappa_2(\mathbf{A}^T \mathbf{A}) = \kappa_2(\mathbf{A})^2$$

Proof:

$$\begin{aligned} \kappa_2(\mathbf{A}^T \mathbf{A}) &= \frac{\text{the largest singular value of } \mathbf{A}^T \mathbf{A}}{\text{smallest singular value of } \mathbf{A}^T \mathbf{A}} \\ &= \frac{6_1^2}{6_n^2} = \left( \frac{6_1}{6_n} \right)^2 \\ &= \kappa_2(\mathbf{A})^2. \end{aligned}$$

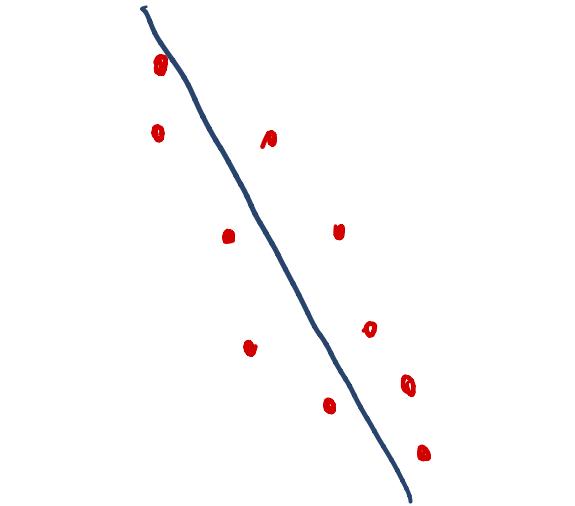
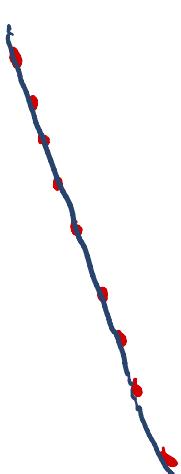
# Sensitivity of the LSQ problem

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It turns out : condition number of LSQ problem is

- $\approx \kappa_2(\mathbf{A})$  if the fit to the data is good (not much scatter)
- $\approx \kappa_2(\mathbf{A})^2$  if the fit to the data is poor (a lot of scatter)

⇒ using normal equations **worsens conditioning of problem** (if the fit is good)



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# Better ways to solve the LSQ problem

$A \in \mathbb{R}^{m \times n}, \quad m > n$

if rank(A) = n,  $\Rightarrow$  full rank A.

rank(A) < n  $\Rightarrow$  rank-deficient

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- If **A** is of full rank, use another matrix factorization — the **QR factorization** <https://powcoder.com>
- For rank-deficient matrices, use the **QR factorization with column pivoting** (MATLAB) aka Rank-revealing QR or the **singular value decomposition SVD**

We'll assume **A** is of full rank.

# QR factorization

The idea of ('economy-size') QR factorization is:

- form a factorization

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$$\text{https://powcoder.com}$$

where  $\mathbf{Q}$  is orthogonal  $m \times n$  matrix,  $\mathbf{R}$  is upper triangular  $n \times n$  matrix i.e.  $\mathbf{Q}^T \mathbf{Q} = \mathbf{I}_n$

- to solve  $\mathbf{Ax} = \mathbf{b}$ ,

$$\mathbf{QRx} = \mathbf{b} \Rightarrow \mathbf{Q}^T \mathbf{QRx} = \mathbf{Q}^T \mathbf{b} \Rightarrow \mathbf{Rx} = \mathbf{Q}^T \mathbf{b}$$

so solve a triangular system for  $\mathbf{x}$  in  $O(n^2)$  ops!

upper triangular matrix

# Gram-Schmidt process

You've seen a QR factorization before (in disguise) in **Gram-Schmidt orthogonalization**:

given a set of linearly independent vectors  $\{a_1, a_2 \dots a_n\}$  forming an  $n$ -D subspace of  $\mathbb{R}^m$ , Gram-Schmidt orthogonalization produces a set of orthonormal vectors  $\{q_1, q_2 \dots q_n\}$ , an orthonormal basis of the same subspace.

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$$\mathbf{A} = \mathbf{Q}\mathbf{R}$$

i.e. use triangular transformations to produce an orthogonal matrix. We don't do it this way because it's numerically unstable; instead we use orthogonal transformations to turn  $\mathbf{A}$  into  $\mathbf{R}$ .

Given's transform, House hold transform

## Gram-Schmidt process

$\{q_1, q_2, \dots, q_n\}$  linearly independent in  $\mathbb{R}^M$

$$q_1 = b_{11} q_1 \quad q_1 \cdot q_1 = 1 \Rightarrow b_{11}$$

$$q_2 = b_{12} q_1 + b_{22} q_2 \quad q_2 \cdot q_1 = 0, \quad q_2 \cdot q_2 = 1 \Rightarrow b_{12}, b_{22}$$

$$\vdots \quad q_1, \dots, q_{k-1}, \quad q_k = b_{1k} q_1 + b_{2k} q_2 + \dots + b_{k-1,k} q_{k-1} + b_{kk} q_k$$

$$q_n = b_{1n} q_1 + \dots + b_{n-1,n} q_{n-1} + b_{nn} q_n$$

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$a_1 = r_{11} q_1$  <https://powcoder.com>

$$a_2 = r_{12} q_1 + r_{22} q_2$$

$\vdots$  Add WeChat powcoder

$$a_n = r_{1n} q_1 + r_{2n} q_2 + \dots + r_{n-1,n} q_{n-1} + r_{nn} q_n$$

$$[a_1 \ a_2 \ \dots \ a_n] = [q_1 \ q_2 \ \dots \ q_n] \begin{bmatrix} r_{11} & r_{12} & \dots & r_{1n} \\ & r_{22} & & \\ & & \ddots & \\ & & & r_{nn} \end{bmatrix}$$

A

$\mathbb{R}^{M \times n}$

Q

$\mathbb{R}^{M \times n}$

R

$\mathbb{R}^{n \times n}$

# Orthogonal transformations

$$\mathbf{x}_{LS} = \arg \min_{\mathbf{x}} \|\mathbf{r}\|_2 = \|\mathbf{b} - \mathbf{Ax}\|$$

$$\mathbf{Ax} = \mathbf{b} \rightarrow Q\mathbf{Ax} = Q\mathbf{b}$$

$$\mathbf{r} \rightarrow Q\mathbf{r}$$

Orthogonal transformations are good because:

- they involve perfectly-conditioned matrices
- they don't change the conditioning of the problem
- they don't change the solution of the LSQ problem

Proofs:

$$k_2(U) = 1.$$

$$\|U\|_2 = 1 \Rightarrow \|U\|_2 = \max_{x \neq 0} \frac{\|Ux\|_2}{\|x\|_2} \geq \max_{x \neq 0} \frac{\|x\|_2}{\|x\|_2} = 1$$

$$k_2(QA) = \|QA\|_2 \|Q(A^+)^T\|_2$$

$$= \|A\|_2 \|A^+ Q^T\|_2$$

$$= \|A\|_2 \|A^+\|_2 = k_2(A)$$

$$(QA)^+ = A^+ Q^+ = A^+ Q^T$$

# Complexity of QR

The QR factorization takes  $n^2(m - n/3)$  ops  
i.e. for  $m \gg n$ ,  $\approx$  twice as expensive as Cholesky factorization of normal equations  
<https://powcoder.com>  
but allows us to handle a larger class of matrices.

Example

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For square systems, can use QR (normwise backward stable)  $\rightarrow$  takes  $2n^3/3$  ops       $M \geq n$   
twice as expensive as GEPP but no issues re growth factor etc.

# QR in MATLAB

I have described what MATLAB calls ‘economy-size QR’ factorization.

$$\mathbf{A} = \mathbf{Q}\mathbf{R}$$

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where  $\mathbf{Q}$  is orthogonal  $m \times n$  matrix,  $\mathbf{R}$  is upper triangular  $n \times n$  matrix.

This is all we need for the LSQ problem.

MATLAB by default produces the ‘full QR’ factorization

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$$\mathbf{A} = \bar{\mathbf{Q}}\bar{\mathbf{R}}$$

where  $\bar{\mathbf{Q}}$  is orthogonal  $m \times m$  matrix,  $\bar{\mathbf{R}}$  is upper triangular  $m \times n$  matrix

$$\bar{\mathbf{Q}} = [\mathbf{Q} | \text{extra orthog. cols}] ; \quad \bar{\mathbf{R}} = \begin{bmatrix} \mathbf{R} \\ \mathbf{0} \end{bmatrix}$$

so that  $\bar{\mathbf{Q}}^T \bar{\mathbf{Q}} = \mathbf{I}_m$ . The extra columns are never used in the LSQ problem.

# Using QR

Since  $\mathbf{A}^T \mathbf{A} = \mathbf{R}^T \mathbf{Q}^T \mathbf{Q} \mathbf{R} = \mathbf{R}^T \mathbf{R}$ ,  $\mathbf{R}$  is the Cholesky factor of  $\mathbf{A}^T \mathbf{A}$ .

Hence to solve LSQ problem, we could:

1  $[q, r] = qr(A, 0); x = r \backslash (q' * b)$   
easiest to understand

2  $r = triu(qr(A)); x = r \backslash (r' \backslash (A' * b))$   
better since never need to form  $\mathbf{Q}$

→ 3  $x = A \backslash b;$   
\\ acting on overdetermined system does the same as 2 (unless  $\mathbf{A}^T \mathbf{A}$  is rank-deficient)

A has full rank

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<https://powcoder.com>

$$\begin{aligned} Ax &\approx b & q' A x &\approx q' b \\ q' q r x &\approx q' b \\ \approx rx &\approx q' b \end{aligned}$$

# The rank-deficient case

Suppose the matrix  $\mathbf{A}$  has rank  $k < n$  i.e. is rank-deficient or not of full rank. This means the columns of  $\mathbf{A}$  are not linearly independent. In this case, there is no unique solution, and we usually use the solution with minimum 2-norm. This solution can be found by either of:

- 1 QR factorization with column pivoting aka Rank revealing QR (RR-QR), based on

<https://powcoder.com>

$$\mathbf{A}\mathbf{E} = \mathbf{Q}\mathbf{R}$$

obtained in MATLAB by a 3-output call to `qr`

- 2 the SVD

In the latter case, let

$$\mathbf{A} = \mathbf{U}\Sigma\mathbf{V}^T = [\mathbf{U}_1 \ \mathbf{U}_2] \begin{bmatrix} \Sigma_1 & 0 \\ 0 & 0 \end{bmatrix} [\mathbf{V}_1 \ \mathbf{V}_2]$$

then

$$\mathbf{x}_{LS} = \mathbf{A}^\dagger \mathbf{b} = \mathbf{V}_1 \Sigma_1^{-1} \mathbf{U}_1^T \mathbf{b}$$

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End of Week 8

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