

CS 202 : Data Structure and Algorithm  
Assignment 2

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## 0.1 Problem Statement

Implement the following sorting algorithms using C++ programming language and sort the input sequence in ascending order.

1. Insertion sort
2. Bubble sort
3. Selection sort
4. Rank sort
5. Merge sort
6. Quick sort

## 0.2 My Conclusions From Assignment :

1. Selection Sort, Insertion Sort and Bubble sort are all **\*\*in place\*\*** sorting algorithms so these can be used where input is small (also mostly sorted) and memory is costly.
2. One of the good things about selection sort is that it never makes more than  $n$  swaps so it can be used quite efficiently where memory write is a costly operation (as selection sort first iterates and then selects the only the required element to be swapped).
3. Also the algorithm like merge sort needs  $O(n)$  extra space so the program reported segmentation fault as our OS keeps bounds on allocation of memory to a certain program and that limit was violated in case of this large input for merge sort and hence program crashed.
4. Insertion sort should not be used where memory write is a costly operation as it makes too much swaps.
5. We also got to know about the stability of sorting algorithms which means that if we have some identical elements in our array then after sorting the indexing of those identical elements should not change. Insertion Sort, Merge Sort, Bubble Sort are some examples of stable sorting algorithms. Some sorting algorithms are not stable, like Heap Sort, Quick Sort, etc.
6. In computer graphics bubble sort is popular for its capability to detect a very small error (like swap of just two elements) in almost-sorted arrays and fix it with just linear complexity ( $2n$ )
7. For smaller input size insertion sort worked better.

## 0.3 Insertion Sort

### 0.3.1 Pseudo Code

```
1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3: for  $j = 2$  to  $n$  do
4:      $key = A[j]$                                  $\triangleright$  Insert  $A[j]$  into the sorted sequence  $A[1..j-1]$ 
5:      $i = j - 1$ 
6:     while  $i > 0$  and  $A[i] > key$  do
7:          $A[j + 1] = A[i]$ 
8:          $i = i + 1$ 
9:     end while
10:     $A[i + 1] = key$ 
11: end for
```

### 0.3.2 Running Time Analysis

Best case :  $O(n)$

Worst case :  $O(n^2)$

## 0.4 Bubble Sort

### 0.4.1 Pseudo Code

```

1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3:  $j = n$ 
4: while  $j \geq 2$  do
    ▷ Bubble up the smallest element to its correct position
5:   for  $i = 1$  to  $j - 1$  do
6:     if  $A[i] > A[i + 1]$  then
7:        $temp = A[i]$ 
8:        $A[i] = A[i + 1]$ 
9:        $A[i + 1] = temp$ 
10:    end if
11:    $j = j - 1$ 
12: end for
13: end while

```

### 0.4.2 Running Time Analysis

Best case :  $O(n)$

Worst case :  $O(n^2)$

## 0.5 Selection Sort

### 0.5.1 Pseudo Code

```

1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3:  $sorted = false$ 
4:  $j = n$ 
5: while  $j > 1$  and  $sorted = false$  do
6:    $pos = 1$ 
7:    $sorted = true$ 
    ▷ Find the position of the largest element
8:   for  $i = 2$  to  $j$  do
9:     if  $A[pos] < A[i]$  then
10:       $pos = i$ 
11:     else
12:        $sorted = false$ 
13:     end if

```

```
14:   end for
    ▷ Move  $A[j]$  to the position of largest element by swapping
15:    $temp = A[pos]$ 
16:    $A[pos] = A[j]$ 
17:    $A[j] = temp$ 
18:    $j = j - 1$ 
19: end while
```

### 0.5.2 Running Time Analysis

Best case :  $O(n)$

Worst case :  $O(n^2)$

## 0.6 Rank Sort

### 0.6.1 Pseudo Code

```
1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3: for  $j = 1$  to  $n$  do
4:    $R[j] = 1$ 
5: end for
    $\triangleright$  Rank the  $n$  elements in  $A$  into  $R$ 
6: for  $j = 2$  to  $n$  do
7:   for  $i = 1$  to  $j - 1$  do
8:     if  $A[i] \leq A[j]$  then
9:        $R[j] = R[j] + 1$ 
10:    else
11:       $R[i] = R[i] + 1$ 
12:    end if
13:  end for
14: end for
    $\triangleright$  Move to correct place in  $U[1..n]$ 
15: for  $j = 1$  to  $n$  do
16:    $U[R[j]] = A[j]$ 
17: end for
    $\triangleright$  Move the sorted entries into  $A$ 
18: for  $j = 1$  to  $n$  do
19:    $A[j] = U[j]$ 
20: end for
```

### 0.6.2 Running Time Analysis

Best case :  $O(n^2)$

Worst case :  $O(n^2)$



## 0.7 Merge Sort

### 0.7.1 Pseudo Code

```
1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3: procedure MERGE( $A, p, q, r$ )
4:    $n_1 = q - p + 1$ 
5:    $n_2 = r - q$ 
6:    $A_1[n_1 + 1] = \infty$ 
7:    $A_2[n_2 + 1] = \infty$ 
8:    $i = 1$ 
9:    $j = 1$ 
10:  for  $k = p$  to  $r$  do
11:    if  $A_1[i] \leq A_2[j]$  then
12:       $A[k] = A_1[i]$ 
13:       $i = i + 1$ 
14:    else
15:       $A[k] = A_2[j]$ 
16:       $j = j + 1$ 
17:    end if
18:  end for
19: end procedure
20: procedure MERGE-SORT( $A, p, r$ )
21:  if  $p < r$  then
22:     $q = (p + r) / 2$ 
23:    MERGE-SORT ( $A, p, q$  )
24:    MERGE-SORT ( $A, q + 1, r$ )
25:    MERGE ( $A, p, q, r$ )
26:  end if
27: end procedure
```

### 0.7.2 Running Time Analysis

Best case :  $O(n \log n)$

Worst case :  $O(n \log n)$

## 0.8 Quick Sort

### 0.8.1 Pseudo Code

```

1: INPUT :  $A[1..n]$ , array of integers
2: OUTPUT : Rearrangement of  $A$  such that  $A[1] \leq A[2] \leq \dots \leq A[n]$ 
3: procedure PARTITION( $A, low, high$ )
4:    $pivot = A[high]$  ▷ Pivot
5:    $i = low - 1$ 
6:   for int  $j = low$  to  $high$  do
7:     if  $A[j] \leq pivot$  then
8:        $i++$ 
9:       exchange  $A[i]$  with  $A[j]$ 
10:    end if
11:    exchange  $A[i + 1]$  with  $A[high]$ 
12:    return  $i + 1$ 
13:  end for
14: end procedure
15: procedure QUICK-SORT( $A, low, high$ )
16:  if  $low < high$  then
17:     $pivot = \text{PARTITION}(A, low, high)$ 
18:    QUICK-SORT( $A, low, pivot$ )
19:    QUICK-SORT( $A, pivot + 1, high$ )
20:  end if
21: end procedure

```

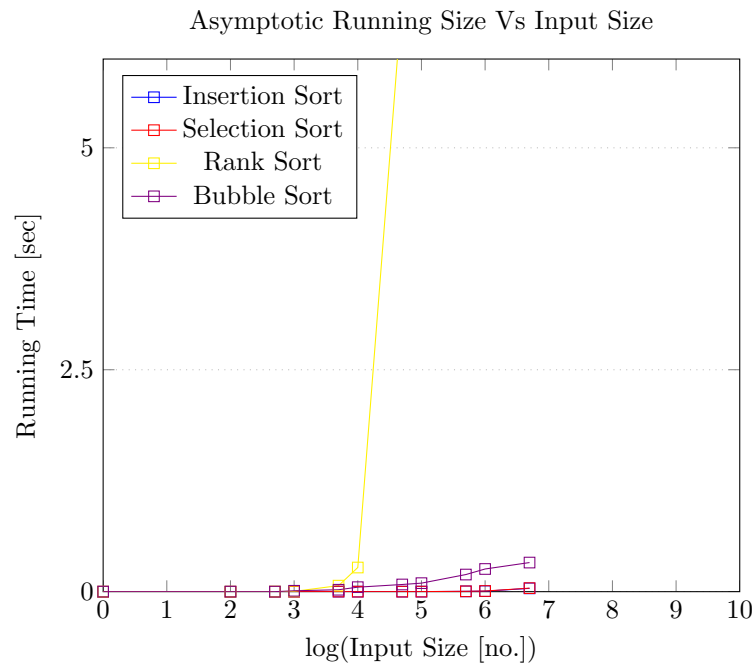
### 0.8.2 Running Time Analysis

Best case :  $O(n \log n)$

Worst case :  $O(n^2)$

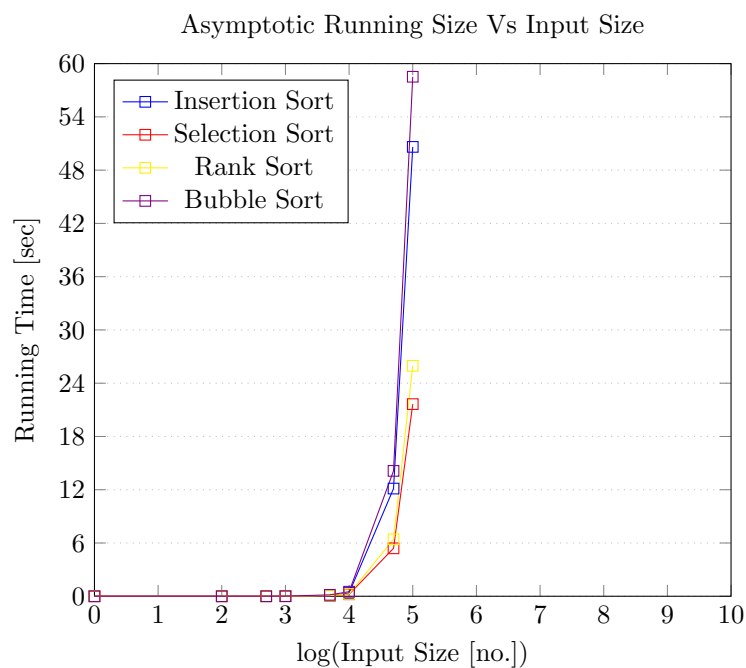
## 0.9 Plots of "Asymptotic Running Size Vs Input Size" for different algorithms

### 0.9.1 Comparison Algorithms (Sorted Input)



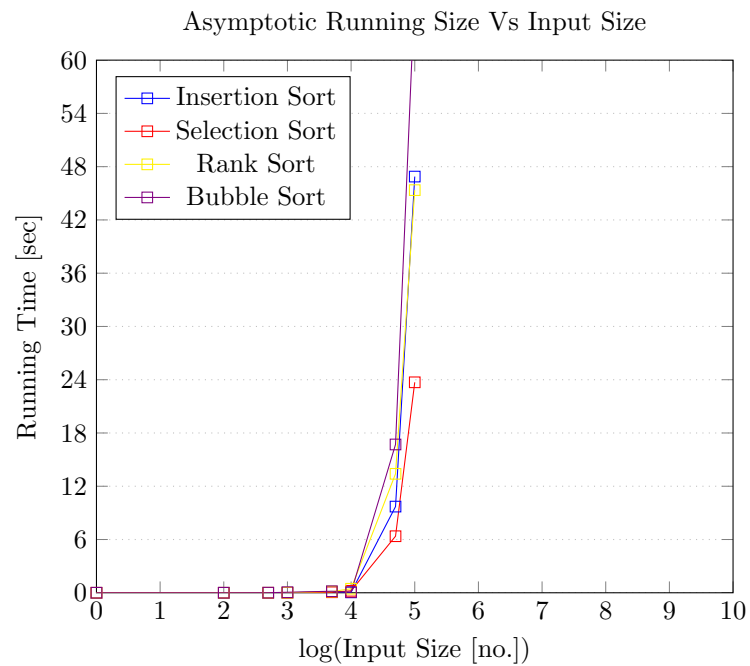
**Time Complexity :**  $Rank > Selection > Bubble > Insertion$

### 0.9.2 Comparison Algorithms (Reverse Sorted Input)



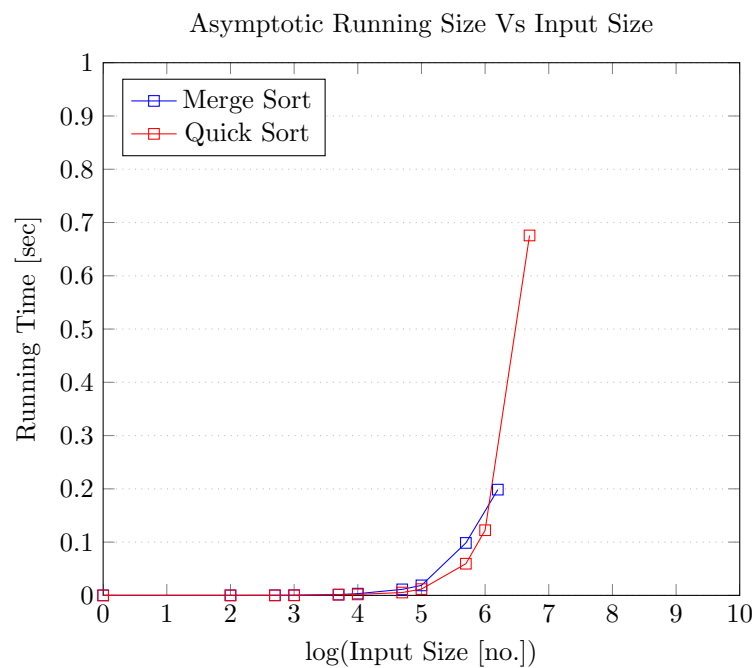
**Time Complexity :**  $Bubble > Insertion > Rank > Selection$

### 0.9.3 Comparison Algorithms (Random Input)

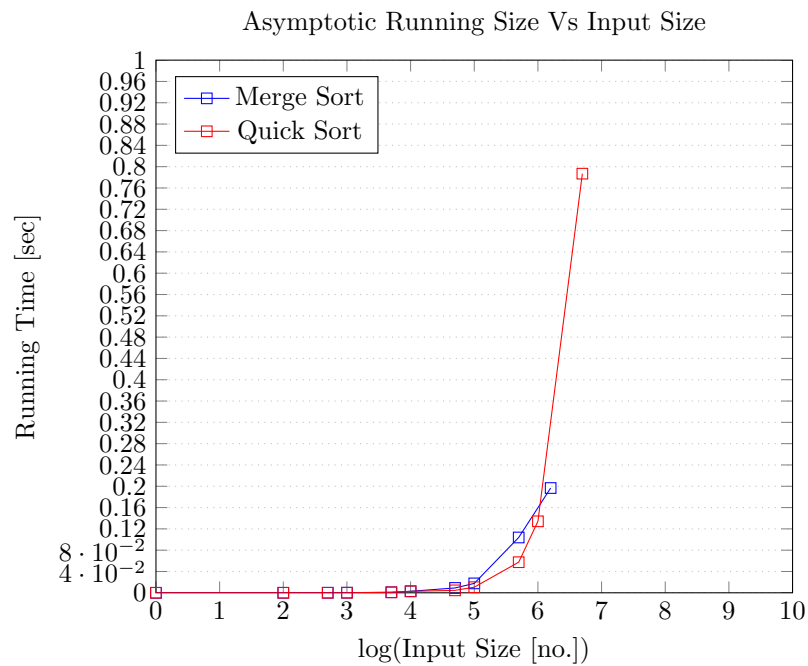


**Time Complexity :** *Bubble > Insertion > Rank > Selection*

### 0.9.4 Divide and Conquer Algorithms (Sorted Input)



### 0.9.5 Divide and Conquer Algorithms (Reverse Sorted Input)



### 0.9.6 Divide and Conquer Algorithms (Random Input)

