

## Unit 5

→ Computational & complexity theory.

→ Decidable or Undecidable

1. Recursive → Tm Halt  $\begin{cases} \text{w} \in L \text{ (accept)} \\ \text{w} \notin L \text{ (reject)} \end{cases}$

2. Recursively enumerable → Tm Halt  $\begin{cases} \text{w} \in L \text{ (accept)} \\ \text{w} \notin L \text{ (reject)} \\ \text{loop} \end{cases}$

3. Decidable → Recursive

4. Partial decidable → Recursively enumerable

5. Undecidable → No Tm exists.

Halting Problem → Undecidable.

Tm → prgm — Halt  $\begin{cases} \text{Yes} \\ \text{No} \end{cases}$   
→ i/p

Check prgm with Halt / not with i/p

Given prgm.  $P$  and an i/p  $I$ , will

$P$  halt on  $I$ ?

proof by contradiction

Assume  $H \rightarrow TM$

$H$  is TM

$H(P, I)$   $C(x)$

if  $H(x, x) == \text{Halt}$ :

loop forever

else

if  $H(x, x) == \text{not Halt}$

return.

$C(C)$

if  $H(C, C) == \text{Halt}$ :

loop forever

else.

if ~~not~~ return.

$H(C, C) == \text{Halt}$

↓

Inf. loop

[not Halt]

proof of  
contradiction.

$H(C, C) \neq \text{Halt}$

↓

Return  
(Halt)

## Rice Theorem

A property of TM lang.  $P$  is a set of descriptions so that for any two TM's

$M_1$  &  $M_2$  with  $L(M_1) = L(M_2)$

(i)  $\langle M_1 \rangle, \langle M_2 \rangle \in P$

(ii)  $\langle M_1 \rangle, \langle M_2 \rangle \notin P$

A non trivial prop. is one statement.

Some  $\langle M_1 \rangle \in P$ , and some  $\langle M_2 \rangle \notin P$ .

Rice Theorem: Every non trivial prop is

Undecidable.

proof: Let  $P$  be arbitrary non trivial property

Suppose  $P$  is decidable

Build a decider for  $A_{TM}$

$A_{TM} \Rightarrow \langle M, w \rangle$ :  $M$  is TM  
 $w$  is string

on input  $\langle M, w \rangle$

1. Const. TM

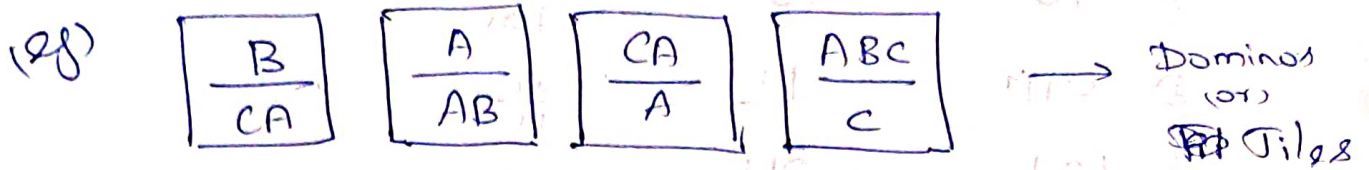
$M$  accepts  $w$  iff this  $mle \in P$ .

$M$  on input  $x$ :

(a) Simulate  $M$  on  $w$

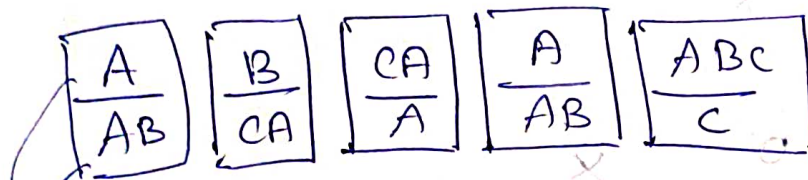
If  $M$  rejects  $w$  then reject  $x$

# Post Correspondence Problem - Undecidable Problem



we need to find a seq. of dominoes such that the top & bottom string are same

\* can use domino any no. of times.



$\rightarrow$  ABCAABCA  
 $\rightarrow$  ABCAABCA

$A = x_1, x_2, x_3, \dots$

$B = w_1, w_2, w_3, \dots$

(eg)

	A	B
①	1	111
②	10111	10
③	10	0
1	10111	10
111	10	0

A - top

B - Bottom

② ① ① ②  $\rightarrow$  10111110  
 $\rightarrow$  10111110



eg.      A                  B

①	10	101	$\frac{10}{101}$
②	011	11	$\frac{011}{11}$
③	101	011	$\frac{101}{011}$

1010	10	10	
101	101	101	X
10	011		X
101	11		
10	101	10	
101	011	101	X

### Undecidability of PCP

\* Take a problem that is already proven to be undecidable and try to convert it to PCP

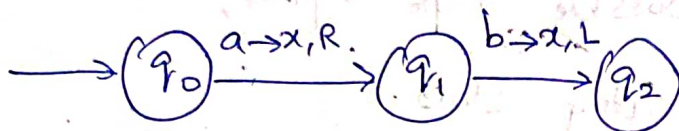
\* If Successfully convert it to an eq. PCP then we prove PCP is also undecidable.

Acceptance  $\emptyset$  Tm (Undecidable problem)

↓

Convert to PCP (called modified PCP - MPCP)

ex)



$\Sigma = \{a, b\}$  :  $\Gamma = \{a, b, x, B\}$  tape symbols

i/p: w = aba



domino  $\frac{\#}{\#q_0aba\#}$   $\frac{\#}{\#q_0w\#}$

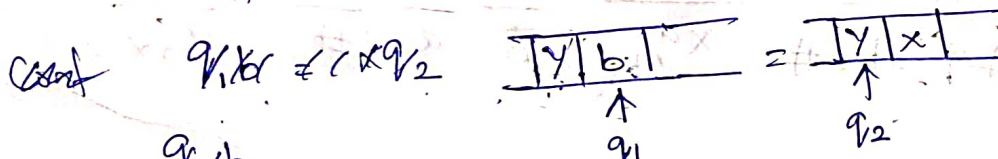
Step 2:  $\delta(q_0, a) = (q_1, x, R)$

configuration  $q_0a = xq_1$

$$\frac{q_0a}{xq_1}$$

domino  $\frac{q_0a}{xq_1}$

Step 3:  $\delta(q_1, b) = (q_2, x, L)$



domino  $\frac{q_1b}{xq_2}$

$\frac{yq_1b}{q_2yx}$

$\forall \gamma \in \Gamma$

$\Gamma = \{a, b, x, B\}$

$$\gamma q_1b = q_2\gamma x$$

$$\left[ \frac{a q_1b}{q_2ax} \right] \left[ \frac{b q_1b}{q_2bx} \right] \left[ \frac{x q_1b}{q_2xx} \right] \left[ \frac{B q_1b}{q_2Bx} \right]$$

Step 4: For all possible tape symbols

$$\Gamma = \{a, b, x, B\}$$

$$\left[ \frac{a}{a} \right] \left[ \frac{b}{b} \right] \left[ \frac{x}{x} \right] \left[ \frac{B}{B} \right]$$

Step 5: For all possible tape symbol, after reaching the accepting state ( $q_1$ )

$$\left[ \frac{a q_2}{q_2} \right] \left[ \frac{q_2 a}{q_2} \right], \left[ \frac{b q_2}{q_2} \right] \left[ \frac{q_2 b}{q_2} \right]$$

$$\left[ \frac{x q_2}{q_2} \right] \left[ \frac{q_2 x}{q_2} \right], \left[ \frac{B q_2}{q_2} \right] \left[ \frac{q_2 B}{q_2} \right]$$

Step 6:  $\left[ \frac{\#}{\#} \right] \left[ \frac{\#}{B\#} \right]$  (Blank & Hash symbol)

↓  
blank

Step 7:  $\left[ \frac{q_2 \#\#}{\#} \right]$  domino for final st.

$$\frac{\#}{\# q_0 a b a \#} \quad \frac{q_0 a}{x q_1} \quad \frac{b}{b} \quad \frac{a}{a} \quad \frac{\#}{\#} \quad \frac{x q_1 b}{q_2 x x} \quad \frac{a}{a} \quad \frac{\#}{\#}$$

$$\rightarrow \frac{\#}{\# q_2 x x a \#} \quad \frac{q_2 x}{q_2} \quad \frac{x}{x} \quad \frac{a}{a} \quad \frac{\#}{\#} \quad \frac{q_2 x}{q_2} \quad \frac{a}{a}$$

$$\frac{q_2 a}{q_2} \quad \frac{\#}{\#} \quad \frac{q_2 \#\#}{\#} \quad \frac{\#}{\#} \quad \checkmark$$

PCP is undecid



## MCP to PCP

→ Insert a symbol  $*$  after all symbol in list A.

→ For list B insert a symbol  $*$  before any symbol in list B.

→ These inserted symbol should be a new symbol.

→ The symbol  $A0 = B0$ , so we insert  $0*$  symbol in front of A.

→ For end of B, insert  $\$$  in A and  $*\$$  in B list.

op	A	B
11	10110	
111	000	
001	0101	
010	0	

I	A	B
1.	$*1*1$	$*1*0*1*1*0$
2.	$1*1*1$	$*0*0*0$
3.	$0*0*1*$	$*0*1*0*1$
4.	$0*1*0*$	$*0$
5.	$\$$	$*\$$



## MCP to PCP

→ Insert a symbol  $*$  after all symbol in list A.

→ For list B insert a symbol  $*$  before any symbol in list B

→ These inserted symbol should be a new symbol

→ The symbol  $A0 = B0$ , so we insert a symbol in front of A1

→ For end of str. insert  $\$$  in A and  $*\$$  in B list.

eq,	A	B	I	A	B
11	10110		1.	$(*)1*$	$*1*0*1*1*0$
111	000		2.	$1*1*1$	$*0*0*0$
001	0101		3.	$0*0*1*$	$*0*1*0*1$
010	0		4.	$0*1*0*$	$*0$
			5	$\$$	$*\$$

X