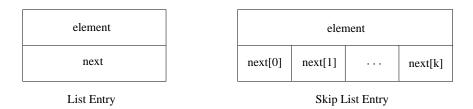
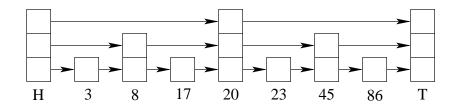
Skip Lists

Generalization of sorted linked lists for implementing Dictionary ADT (insert, delete, find, min, succ) in $O(\log n)$ expected time per operation. Skip lists compete with balanced search trees like AVL, Red-Black, and B-Trees.



The elements are stored in sorted order, in a linked list of nodes. Each skip list entry has an array of next pointers, where next[i] points to an element that is roughly 2^i nodes away from it. The next array at each entry has random size between 1 and maxLevel, the maximum number of levels in the current skip list. Ideally, $maxLevel \approx \log n$. Each skip list has dummy head and tail nodes, both of maxLevel height, storing sentinels $-\infty$ and $+\infty$, respectively. Iterating through the list using next[0] will go through the nodes in sorted order. A reference to the previous element can also be stored by adding a prev field to Skip List Entry.



Search starts at the top level, goes as far as possible at each level, without going past target, descending one level at a time, until reaching the target node. Addition/Removal of nodes makes it difficult to maintain an ideal skip list, in which next[i] of a node points to a node that is exactly 2^i away from it. Skip lists solve this problem by selecting the number of levels (size of $next[\]$) of a new node probabilistically.

Skip List implementation:

```
Entry class:
        T element
        Entry[] next
       Entry prev // prev is optional
       int[] span // for indexing
Entry(x, lev): // constructor
       element \leftarrow x
        next ← new Entry[lev]
       span ← new int[lev]
SkipList class:
        Entry head, tail // dummy nodes
       int size, maxLevel
        Entry[] last // used by find()
        Random random
SkipList(): // Constructor
        head \leftarrow new Entry(null, 33) // sentinel -\infty, with maximum number of levels
       tail \leftarrow new Entry(null, 33) // sentinel +\infty, with maximum number of levels
       size \leftarrow 0
       maxLevel ← 1
       last ← new Entry[33]
       random ← new Random()
find(x): // helper method to search for x. Sets last[i] = node at which search came down from level i to i-1
       p ← head
       for i \leftarrow maxLevel-1 down to 0 do
               while p.next[i].element < x do // watch out for NPE because of null element in tail
                       p \leftarrow p.next[i]
               last[i] ← p
contains(s): // is x there in list?
       find(x)
       return last[0].next[0].element == x
remove(x): // delete x
       if not contains(x) then return null
       ent \leftarrow last[0].next[0]
       for i \leftarrow 0 to ent.next.length-1 do
               last[i].next[i] ← ent.next[i] // bypass ent at level i
       size ← size - 1
        return ent.element
```

```
add(x): // insert x
       if contains(x) then return false // reject duplicate
       lev ← chooseLevel() // length of next[] for x's entry
       ent \leftarrow new Entry(x, lev)
       for i \leftarrow 0 to lev-1 do
               ent.next[i] ← last[i].next[i]
               last[i].next[i] ← ent
       ent.next[0].prev ← ent; ent.prev ← last[0] // if prev link is defined in Entry
       size ← size + 1
       return true
chooseLevel(): // Prob(choosing level i) = 1/2 Prob(choosing level i-1)
       // Slow method:
       lev ← 1
       while random.nextBoolean() do lev++ // should limit to 33
       if lev > maxLevel then maxLevel ← lev
       return lev
       // fast method:
       lev ← 1 + Integer.numberOfTrailingZeros(random.nextInt())
       // Optionally, lev ← min(lev, maxLevel+1), to allow maxLevel to grow gradually
       if lev > maxLevel then maxLevel ← lev
       return lev
Indexing in skip lists:
get(index): // return element at index (first element is at index 0)
       if index < 0 or index > size - 1 then throw NoSuchElementException
       p ← head
       for i \leftarrow 0 to index do
               p \leftarrow p.next[0]
       return p.element
```

Running time of get(index) is O(1+index), which is O(n) in the worst case.

Improving the RT of get(): Each Entry stores the span of next[i] in span[i] (number of elements between the two nodes). Code for find() has to be updated to calculate the indexes of the nodes in last[]. Code for add(), and remove() need to be modified to update the appropriate span[] values to reflect the changes to the list made by the operations.