



Selected Topics of Embedded Software Development 2

WS-2021/22

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**Testing and generating Prime Numbers and Safe Primes
using CryptoCore**

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Code enhancements

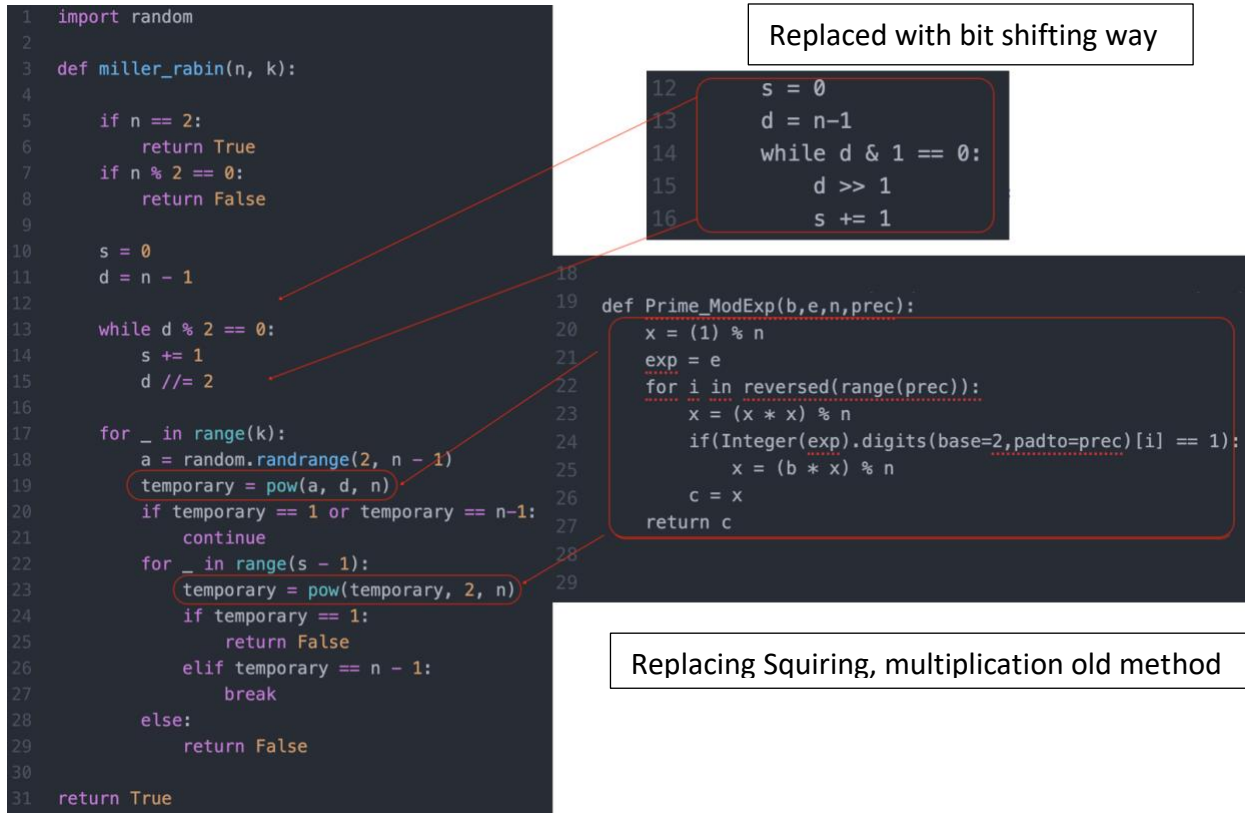


Figure 1: Replacing default pow() functionality

Enhanced Version of MRT

```

10 def miller_rabin(n, k = 20):
11
12     s = 0
13     d = n-1
14     while d & 1 == 0:
15         d = d >> 1
16         s += 1
17
18     for _ in range(k):
19         a = random.randrange(2, n-2)
20         prec = n.nbits()
21         temporary = Prime_ModExp(a, d, n, prec)
22         if temporary == 1 or temporary == n-1:
23             continue
24         for _ in range(s - 1):
25             temporary = (temporary * temporary) % n
26             if temporary == n - 1:
27                 break
28         else:
29             return False
30     return True
31
32 # Taken from >> CryptoCore_User_Story_2-20211123 (Replacment of Pow(a, d, n))
33 def Prime_ModExp(b,e,n,prec):
34     x = (1) % n
35     exp = e
36     for i in reversed(range(prec)):
37         x = (x * x) % n
38         if(Integer(exp).digits(base=2,padto=prec)[i] == 1):
39             x = (b * x) % n
40     c = x
41     return c

```

Figure 2: MRT Algorithm

Primes Generator

```
43 # ----- Prime Generator -----
44
45 def generate_prime(bits):
46     while True:
47         random_number = Integer((random.randrange(1 << bits - 1, 1 << bits) << 1) + 1)
48         if miller_rabin(random_number):
49             return random_number
50
51 # ----- Safe Prime Generator -----
52
53 def generate_safe_prime(bits):
54     while True:
55         random_number = Integer((random.randrange(1 << bits - 1, 1 << bits) << 1) + 1)
56         if safe_prime(random_number):
57             return random_number
58
59 def safe_prime(n):
60     if miller_rabin(n):
61         d = (n-1) >> 1
62         if miller_rabin(d):
63             return True
64     return False
```

Figure 3: Primes Generator Functionality

Test Result

```
62 # ----- TEST -----  
63  
64 start_prime = time.time()  
65 print(generate_prime(512))  
66 finish_prime = time.time()  
67 print('Prime takes', (finish_prime - start_prime), 'seconds')  
68 start_safe_prime = time.time()  
69 print(generate_safe_prime(512))  
70 finish_safe_prime = time.time()  
71 print('Safe Prime takes', (finish_safe_prime - start_safe_prime), 'seconds')  
72
```

Will give us

```
267260057033759207087520624889170005507325565809986811242061305136360374778  
297685741486602532015127800478557133170101563644662779952715459726100326644  
70373
```

```
Prime takes 0.08292651176452637 seconds
```

```
248781546171890584600550712540752545216684211642419810170196165277792994639  
259676752398792589844970615162452268578295009483577955283442953017618455154  
34523
```

```
Safe Prime takes 0.16614723205566406 seconds
```

Montgomery multiplication

Is a method for performing fast modular multiplication. It was introduced in 1985 by the American mathematician Peter L. Montgomery

Montgomery modular multiplication relies on a special representation of numbers called Montgomery form. The algorithm uses the Montgomery forms of a and b to efficiently compute the Montgomery form of $ab \bmod N$. The efficiency comes from avoiding expensive division operations. Classical modular multiplication reduces the double-width product ab using division by N and keeping only the remainder. This division requires quotient digit estimation and correction. The Montgomery form, in contrast, depends on a constant $R > N$ which is coprime to N , and the only division necessary in Montgomery multiplication is division by R . The constant R can be chosen so that division by R is easy, significantly improving the speed of the algorithm. In practice, R is always a power of two, since division by powers of two can be implemented by bit shifting.

For example, suppose that $N = 17$ and that $R = 100$.

The Montgomery forms of 3, 5, 7, and 15 are $300 \bmod 17 = 11$, $500 \bmod 17 = 7$, $700 \bmod 17 = 3$, and $1500 \bmod 17 = 4$.

Addition and subtraction in Montgomery form are the same as ordinary modular addition and subtraction because of the distributive law:

$$\begin{aligned} aR + bR &= (a + b) R, \\ aR - bR &= (a - b) R. \end{aligned}$$

This is a consequence of the fact that, because $\gcd(R, N) = 1$, multiplication by R is an isomorphism on the additive group $\mathbb{Z}/N\mathbb{Z}$.

For example, $(7 + 15) \bmod 17 = 5$, which in Montgomery form becomes $(3 + 4) \bmod 17 = 7$.

MRT in Montgomery Domain

```
7  def miller_rabin(n, k = 20, isMontgomery = False):
8
9      s = 0
10     d = n-1
11     prec = n.nbits()
12     temporary = 0
13
14     while d & 1 == 0:
15         d = d >> 1
16         s += 1
17
18     for _ in range(k):
19         a = random.randrange(2, n-1)
20
21         if isMontgomery:
22             temporary = Prime_MontExp(a, d, n, prec)
23         else:
24             temporary = Prime_ModExp(a, d, n, prec)
25
26         if temporary == 1 or temporary == n-1:
27             continue
28         for _ in range(s - 1):
29             temporary = (temporary * temporary) % n
30             if temporary == n - 1:
31                 break
32         else:
33             return False
34     return True
```

```

36 # Taken from >> CryptoCore_User_Story_2-20211123 (Replacment of Pow(a, d, n))
37 def Prime_ModExp(b,e,n,prec):
38     x = (1) % n
39     exp = e
40     for i in reversed(range(prec)):
41         x = (x * x) % n
42         if(Integer(exp).digits(base=2, padto=prec)[i] == 1):
43             x = (b * x) % n
44         c = x
45     return c
46
47 # Taken from >> CryptoCore_User_Story_2-20211123 (Replacment of Pow(a, d, n) and integrating the montgomery method)
48 def Prime_MontExp(b,e,n,prec):
49     r = 2^prec % n; r2 = (r*r) % n; rinv = inverse_mod(r,n)
50     x = (1 * r2 * rinv) % n
51     exp = e
52     for i in reversed(xrange(prec)):
53         x = (x * x * rinv) % n
54         if(Integer(exp).digits(base=2, padto=prec)[i] == 1):
55             x = (b * x * rinv) % n
56     c = x
57     return(c)
58

```


Test Result

```
62 # ----- TEST -----
63
64 start_prime = time.time()
65 print(generate_prime(512))
66 finish_prime = time.time()
67 print('Prime takes', (finish_prime - start_prime), 'seconds')
68 start_safe_prime = time.time()
69 print(generate_safe_prime(512))
70 finish_safe_prime = time.time()
71 print('Safe Prime takes', (finish_safe_prime - start_safe_prime), 'seconds')
72
```

Will give us

```
204466626762739354657685647960825630598057253125579812380780798972913290454
123028564312144708296967270995504201361321192429536808009692794068949897909
00497
```

```
Prime takes 3.300482988357544 seconds
```

```
Safe prime, does not work on bits > 256
```