

Eine Woche, ein Beispiel

5.15 Category

Everybody knows a little about category theory, but nobody can conclude all the terms emerged in the category theory.

In this document I try to collect the notations and basic examples used in the course "Condensed Mathematics and Complex Geometry". I'm sure that it won't be better than the wikipedia, I just collect results I'm happy with.

I have to divide it into two parts which interact with each other, but you can always jump through examples which you're not familiar. You can also find a "complete" list of categories here: <http://katmat.math.uni-bremen.de/acc/acc.pdf>

For Chinese, the theory of category has been summed up in detail in [<https://wwli.asia/downloads/books/Al-jabr-1.pdf>], Chapter 2-3.

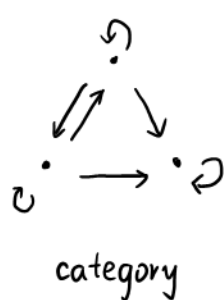
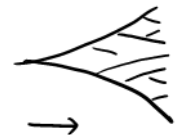
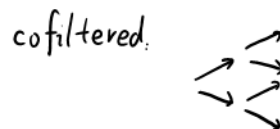
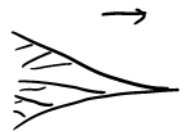
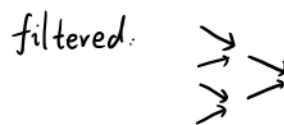
Process

0. Well-know concepts
1. Individual category
 - Complete/Cocomplete/Bicomplete category
 - Cartesian closed category / Closed category
 - Monoidal category = Tensor category
2. Functors between categories
 - Exactness
 - Adjoints
3. Examples of categories
 - Well-known examples
 - Cat
 - Hausdorff and compactness
 - Categories in condensed mathematics

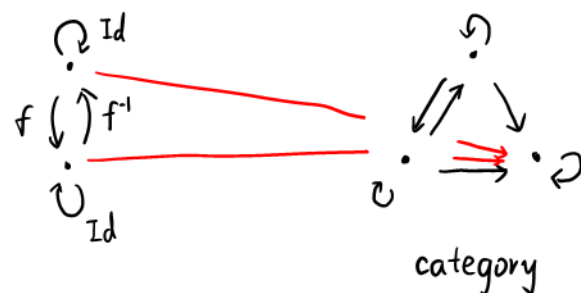
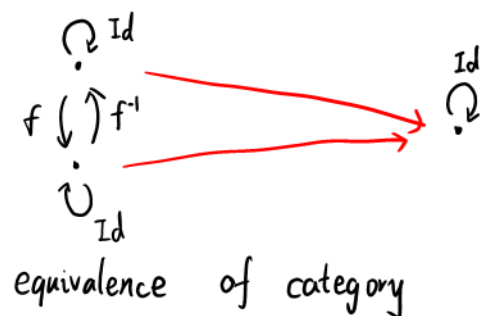
Appendix.

0. Well-known concepts
 \mathcal{C} is always a category.

	$Ob(\mathcal{C})$	$Mor(X, Y)$
small	Set	Set
loc. small	—	Set
large	not set	or not set



<https://math.stackexchange.com/questions/2147377/are-fully-faithful-functors-injective>



fully faithful

<https://blog.juliosong.com/linguistics/mathematics/category-theory-notes-8/>

1. Individual category

Complete/Cocomplete/Bicomplete category

Def. \mathcal{C} is **complete** if

$$\forall \text{ small category } \Delta, \forall \text{ fctor } F: \Delta \rightarrow \mathcal{C} \quad i \mapsto F_i, \\ \lim_{i \in \Delta} F_i \text{ exists} \quad \left(\lim_{i \in \Delta} F_i \text{ is called the small limit} \right)$$

\mathcal{C} is **cocomplete** if

$$\forall \text{ small category } \Delta, \forall \text{ fctor } F: \Delta \rightarrow \mathcal{C} \quad i \mapsto F_i, \\ \lim_{i \in \Delta} F_i \text{ exists} \quad \left(\lim_{i \in \Delta} F_i \text{ is called the small colimit} \right)$$

bicomplete = complete + cocomplete

\mathcal{C} is **finitely complete** if \forall finite limit exists

\mathcal{C} is **finitely cocomplete** if \forall finite colimit exists.

Thm.

\mathcal{C} is complete $\Leftrightarrow \mathcal{C}$ has equalizers & products

$\Leftrightarrow \mathcal{C}$ has pullbacks & products

\mathcal{C} is cocomplete $\Leftrightarrow \mathcal{C}$ has coequalizers & coproducts

$\Leftrightarrow \mathcal{C}$ has pushouts & coproducts

\mathcal{C} is finitely complete $\Leftrightarrow \mathcal{C}$ has equalizers & finite products

$\Leftrightarrow \mathcal{C}$ has equalizers, binary products & terminal obj

$\Leftrightarrow \mathcal{C}$ has pullbacks & terminal obj

For small category \mathcal{C} ,

complete \Leftrightarrow cocomplete

\Downarrow

\Downarrow

thin $(\# \text{Mor}(X, Y) \leq 1)$

The "closedness" of the category is that taking the morphisms between two objects gives another morphism in the same category, rather than the category of sets or some other category.
 from: <https://math.stackexchange.com/questions/3486846/definition-of-cartesian-closed-category-why-do-we-need-exponential-objects>

Cartesian closed category / Closed category

Def. \mathcal{C} is **Cartesian closed** if

\mathcal{C} has terminal obj, binary product and exponential, where

$$- \times Y \vdash (-)^Y \quad \text{a bifunctor } F: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C} \text{ which is functorial in } Y$$

$$\text{ie. } \text{Mor}(X \times Y, Z) \cong \text{Mor}(X, Z^Y)$$

\mathcal{C} is **loc. Cartesian closed** if all its slice category is Cartesian closed.

<https://ncatlab.org/nlab/show/over+category>

Rmk. When \mathcal{C} is loc. Cartesian closed,

\mathcal{C} is Cartesian closed $\Leftrightarrow \mathcal{C}$ has a terminal object.

But \mathcal{C} is Cartesian closed $\nRightarrow \mathcal{C}$ is loc. Cartesian closed

For the closed category, we use the definition in <https://ncatlab.org/nlab/show/closed+category>.

Def. A **closed category** is a category \mathcal{C} together with the following data.

- bifunctor $[-, -]: \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$

called internal hom-functor

- $I \in \text{Ob}(\mathcal{C})$

called unit object

$$- i: \text{Id}_{\mathcal{C}} \xrightarrow{\cong} [I, -] \rightsquigarrow i_A: A \xrightarrow{\cong} [I, A]$$

$$- j_X: I \longrightarrow [X, X]$$

extranatural in X

$$- L_{Y,Z}^X: [Y, Z] \rightarrow [X, Y], [X, Z]$$

functorial in Y and Z

extranatural in X .

- Compatibilities

$$\begin{array}{ccccc} I & \xrightarrow{j_Y} & [Y, Y] & & [X, Y] & \xrightarrow{L_{XY}^X} & [X, X], [X, Y] & & [Y, Z] & \xrightarrow{L_{YZ}^I} & [I, Y], [I, Z] \\ & \searrow j_{[X,Y]} & \downarrow L_{Y,Y}^X & & \searrow i_{[X,Y]} & \downarrow [j_X, 1] & & \searrow [1, i_Z] & \downarrow [i_Y, 1] \\ & & [X, Y], [X, Y] & & & [I, [X, Y]] & & & [Y, [I, Z]] \end{array}$$

$$\begin{array}{ccc} & [U, V] & \\ L_{UV}^X \swarrow & & \searrow L_{UV}^Y \\ [X, U], [X, V] & & [Y, U], [Y, V] \\ \downarrow L_{[X,U],[X,V]}^{[X,Y]} & & \downarrow [1, L_{YV}^X] \\ [X, U], [X, V] & \xrightarrow{[L_{YU}^X, 1]} & [Y, U], [X, Y], [X, V] \end{array}$$

$$\gamma: \text{Mor}(X, Y) \longrightarrow \text{Mor}(I, [X, Y]) \quad \text{is an iso.}$$

$$f \longmapsto [1, f] \circ j_X$$

Monoidal category = Tensor category

么半范畴

Def A **monoidal category** is a category \mathcal{C} together with the following data.

- bifactor $- \otimes - : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$

- $I \in \text{Ob}(\mathcal{C})$

called unit object

- $\alpha_{A,B,C} : A \otimes (B \otimes C) \xrightarrow{\cong} (A \otimes B) \otimes C$

- $\lambda_A : I \otimes A \xrightarrow{\cong} A$

lambda: left

- $\rho_A : A \otimes I \xrightarrow{\cong} A$

rho: right

- Compatibilities

$$\begin{array}{ccc}
 & A \otimes (B \otimes (C \otimes D)) & \\
 1_A \otimes \alpha_{B,C,D} \swarrow & & \searrow \alpha_{A,B,C \otimes D} \\
 A \otimes ((B \otimes C) \otimes D) & \cong & (A \otimes B) \otimes (C \otimes D) \\
 \alpha_{A,B \otimes C,D} \downarrow & & \downarrow \alpha_{A \otimes B,C,D} \\
 (A \otimes (B \otimes C)) \otimes D & \xrightarrow{\alpha_{A,B,C} \otimes 1_A} & (A \otimes B) \otimes (C \otimes D) \\
 & & \\
 A \otimes (I \otimes B) & \xrightarrow{\alpha_{A,I,B}} & (A \otimes I) \otimes B \\
 1_A \otimes \lambda_B \searrow & & \swarrow \rho_A \otimes 1_B \\
 & A \otimes B &
 \end{array}$$

For **strict monoidal category**, we require in addition that $\alpha_{A,B,C}, \lambda_A, \rho_A$ are identities.

Eg. **Cartesian monoidal category** \mathcal{C} : category with finite products

$\otimes = \prod$

$I = \text{terminal object}$

e.g. Set, Cat.

Cocartesian monoidal category \mathcal{C} : category with finite coproducts

$\otimes = \coprod$

$I = \text{initial object}$

Abelian category is monoidal.

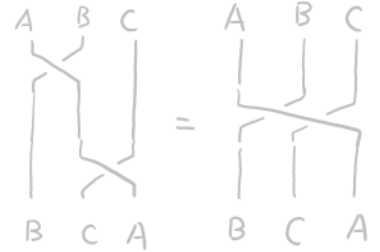
Def (Specializations)

Let \mathcal{C} be a monoidal category.

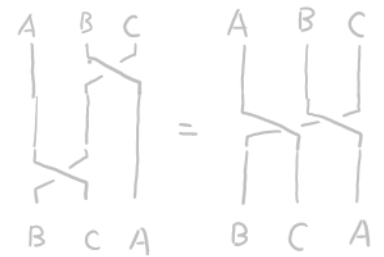
If in addition we have $\gamma_{A,B}: A \otimes B \rightarrow B \otimes A$,

then \mathcal{C} is **braided monoidal category** if

$$\begin{array}{ccc}
 & (A \otimes B) \otimes C & \\
 \gamma_{A,B} \otimes 1_C \swarrow & & \searrow \alpha_{A,B,C} \\
 (B \otimes A) \otimes C & & A \otimes (B \otimes C) \\
 \downarrow \alpha_{B,A,C} & & \downarrow \gamma_{A,B \otimes C} \\
 B \otimes (A \otimes C) & & (B \otimes C) \otimes A \\
 1_B \otimes \gamma_{A,C} \searrow & & \swarrow \alpha_{B,C,A} \\
 & B \otimes (C \otimes A) &
 \end{array}$$



$$\begin{array}{ccc}
 & A \otimes (B \otimes C) & \\
 1_A \otimes \gamma_{B,C} \swarrow & & \searrow \alpha_{A,B,C}^{-1} \\
 A \otimes (C \otimes B) & & (A \otimes B) \otimes C \\
 \downarrow \alpha_{A,C,B}^{-1} & & \downarrow \gamma_{A \otimes B, C} \\
 (A \otimes C) \otimes B & & C \otimes (A \otimes B) \\
 \gamma_{C,A} \otimes 1_B \searrow & & \swarrow \alpha_{C,A,B}^{-1} \\
 & (C \otimes A) \otimes B &
 \end{array}$$



$$\begin{array}{ccc}
 I \otimes A & \xrightarrow{\gamma_{I,A}} & A \otimes I \\
 \lambda_A \searrow & & \swarrow \rho_A \\
 & A &
 \end{array}$$

$$\begin{array}{ccc}
 A \otimes I & \xrightarrow{\gamma_{A,I}} & I \otimes A \\
 \rho_A \searrow & & \swarrow \lambda_A \\
 & A &
 \end{array}$$

\mathcal{C} is **symmetric monoidal category** if

$$\gamma_{B,A} \circ \gamma_{A,B} = 1_{A \otimes B}.$$

+ \mathcal{C} is braided.

closed monoidal category = closed category + monoidal category
+ compatabilite $- \otimes A \dashv [A, -]$

2. Functors between categories

Exactness

Ref: <https://ncatlab.org/nlab/show/exact+functor>

Prop/Def For a functor $F: \mathcal{C} \rightarrow \mathcal{D}$ between finitely complete categories, TFAE:

- F preserves finite limits,
- F preserves equalizers & finite products,
- F preserves equalizers, binary products & terminal objects,
- F preserves pullbacks & terminal objects,
- $\forall d \in \text{Ob}(\mathcal{D})$, the comma category F/d is filtered.

If so, we call F is a **left exact functor**.

↑ We require also that \mathcal{C} & \mathcal{D} are finitely complete categories

When \mathcal{C}, \mathcal{D} are abelian categories, this is equivalent to

- F preserves kernels, i.e.
- F sends left exact sequences to left exact sequences.

See [<https://stacks.math.columbia.edu/tag/010M>]. You may get the following results from the argument:

|| F is a left exact functor between abelian categories $\Rightarrow F$ preserve binary products $\Downarrow [\text{oDLP}]$
 F is additive

Def. A contravariant functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is called left exact, if
 $F: \mathcal{C}^{\text{op}} \rightarrow \mathcal{D}$ is left exact (In ptc. \mathcal{C} is finitely cocomplete, \mathcal{D} is finitely complete)

Similarly, we can define right exactness, and
 $\text{exact}_{\mathcal{C}, \mathcal{D} \text{ abelian}} = \text{left exact} + \text{right exact}$
 \equiv sends SES to SES.

Adjoint

left adjoint \dashv right adjoint

$$f_! \dashv f^* \dashv f_* \dashv f^!$$

free

forget

$$-\otimes_A N$$

$$\text{Hom}_A(N, -)$$

$$\Delta$$

$$\varprojlim$$

$$(\)^\sim$$

$$\Gamma^*$$

$$f_p$$

$$f^*$$

$$c\text{-Ind}$$

$$\text{Res}$$

$$\text{Res}$$

$$\text{Ind}$$

$$-\otimes_k A^e$$

$$\text{Res}$$

$$\text{sh}(-)$$

$$\text{Res}_{\text{sh} \rightarrow \text{Psh}}$$

$$\pi_0$$

$$T \mapsto \coprod_{t \in T} \text{Spec } \mathbb{Z}$$

$$G^{[p]}$$

$$\text{Lie}$$

$$|-|$$

$$S$$

$$\tau_1 = H_0(-)$$

$$N : \text{nerve}$$

$$\beta : \text{Stone-}\check{\text{Cech}}$$

$$U : \text{forget}$$

ad \Rightarrow

preserve colimits

preserve limits

\Rightarrow

right exacts

left exacts

in (co)complete category

$$\text{coker } f$$

$$\ker f$$

$$\coprod A \oplus A$$

$$\prod A \quad A \times_c B$$

$$\text{pushforward}$$

$$\text{pullback}$$

$$\text{coequalizer}$$

$$\text{equalizer}$$

$$\bar{K} = \varinjlim_{L/K} L$$

$$\text{Spec } \bar{K} = \varprojlim_{L/K} \text{Spec } L \quad \text{Gal}(\bar{K}/K) = \varprojlim_{\substack{L \text{ finite} \\ \text{Galois}}} \text{Gal}(L/K)$$

$$\text{Spec } \mathbb{Z}_p = \varprojlim_n \text{Spec } \mathbb{Z}/p^n \mathbb{Z}$$

$$\mathbb{Z}_p = \varprojlim_n \mathbb{Z}/p^n \mathbb{Z}$$

$$\mathcal{F}_p = \varinjlim_{\text{stalk}} \mathcal{F}(U)$$

co limit

limit

$$\left. \begin{array}{l} \text{direct} \\ \text{inductive} \\ \text{injective} \end{array} \right\} \text{limit.}$$

$$\left. \begin{array}{l} \text{inverse} \\ \text{projective} \end{array} \right\} \text{limit.}$$

How to memorize:

limit

colimit

$$0 \rightarrow \ker \rightarrow M \rightarrow N \rightarrow \text{coker} \rightarrow 0$$

3. Examples of categories

Well-known examples

Set Top Grp Ab Vect(k) Mod(R)

Ring: identity + preserve identity

CRing Rng

Field: full subcategory of CRing

$$0: Ob(0) = \emptyset$$

$$1: Ob(1) = \{*\} \quad Mor(*, *) = \{1_{*}\}$$

$$K(2): Ob(K(2)) = \{V, E\} \quad Mor(V, V) = \{1_V\} \quad Mor(E, E) = \{1_E\} \\ Mor(V, E) = \emptyset \quad Mor(E, V) = \{s, t\}$$

$$\begin{array}{ccc} 1_E & & \\ \downarrow & s & \\ 0 & E & \xrightarrow{t} V & 1_V \end{array}$$

$$\Delta: Ob(\Delta) = \{[n] := \{0, 1, 2, \dots, n\} \mid n \geq 0\} \\ Mor([m], [n]) = \{\text{weakly monotone maps}\}$$

$$sSet: Ob(sSet) = \left\{ X: \Delta^{op} \rightarrow Set \mid [n] \mapsto X_n \right\} \quad Mor(X, Y) = \left\{ \alpha: \Delta^{op} \begin{array}{c} \xrightarrow{X} \\ \Downarrow \alpha \\ Y \end{array} Set \right\}$$

$$CHaus: Ob(CHaus) = \left\{ \underbrace{cpt \text{ Hausdorff space}}_{cptum/cpta} X \right\}$$

<https://ncatlab.org/nlab/show/compactum>

$$Mor(X, Y) = \{f: X \rightarrow Y \mid f \text{ cont}\}$$

Met: full subcategory of CHaus whose objects are metric spaces.

! For the category of Graph, there're different realizations.

$$Quiv(e): Ob(Quiv(e)) = \{fctor \Gamma: K(2) \rightarrow e\} \\ Mor(\Gamma_1, \Gamma_2) = \left\{ \alpha: K(2) \begin{array}{c} \xrightarrow{\Gamma_1} \\ \Downarrow \alpha \\ \Gamma_2 \end{array} e \right\}$$

$$Quiv = Quiv(Set) \\ = \text{Category of presheaves on } Q^{op}.$$

Cat

$\mathbf{Cat} = \{\text{the category of small categories}\}$ is a 2-category.

$\text{Ob}(\mathbf{Cat}) = \{\text{small category } \mathcal{C}\}$

$\text{Mor}(\mathcal{C}, \mathcal{D})$ is a category by

$\text{Ob}(\text{Mor}(\mathcal{C}, \mathcal{D})) = \{F: \mathcal{C} \rightarrow \mathcal{D}\}$

$\text{Mor}(F, G) = \left\{ \alpha: \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{D} \right\}$

Basic properties of \mathbf{Cat} :

1. Initial object 0 , Terminal object 1 .
2. \mathbf{Cat} is loc. small but not small
3. \mathbf{Cat} is bicomplete
4. \mathbf{Cat} is Cartesian closed but not loc. Cartesian closed
5. \mathbf{Cat} is **loc. finitely presentable** <https://ncatlab.org/nlab/show/locally+finitely+presentable+category>

6. $\mathbf{Cat} \begin{array}{c} \xleftarrow{\text{free}} \\ \text{forget} \end{array} \mathbf{Quiv}$

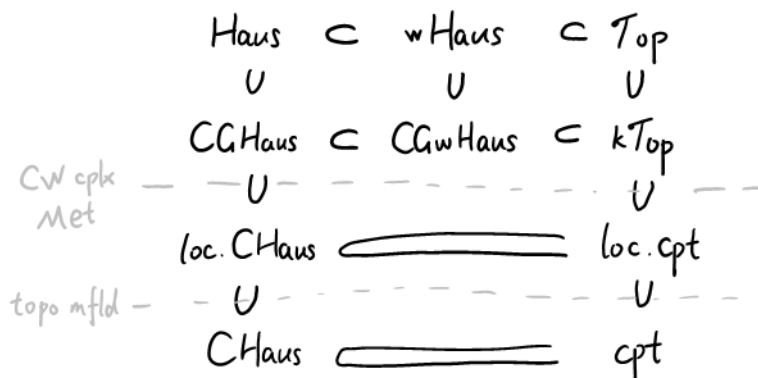
e.g of "free"

$$f \circ \begin{array}{c} \circlearrowleft \\ \circlearrowright \end{array} \circ 1 \quad \Leftarrow \quad \circlearrowleft f$$

$$1_a \circ \begin{array}{c} \circlearrowleft \\ \circlearrowright \end{array} \circ 1_b \quad \Leftarrow \quad a \begin{array}{c} \xrightarrow{e} \\ \xleftarrow{f} \end{array} b$$

$$1_a \circ \begin{array}{c} \circlearrowleft \\ \circlearrowright \end{array} \circ 1_c \quad \Leftarrow \quad a \xrightarrow{f} b \xrightarrow{g} c$$

Hausdorff and compactness $\leftarrow \text{cpt} \approx (\text{quasi})\text{cpt}$



Def. $X \in \text{Top}$ is a **weak Hausdorff space** (in wHaus) if
 $\forall K \in \text{CHaus}, \forall g: K \rightarrow X$ cont, $g(K) \subset X$ is closed.

Def. $X \in \text{Top}$ is **locally compact** (in loc.cpt) if
 $\forall p \in X, \exists \text{ cpt nbhd } V$ (i.e. $p \in U \subseteq V \subseteq X$ $U \subseteq X$ open, V cpt)
 $\text{loc. CHaus} = \text{loc.cpt} \cap \text{Haus}$

see https://en.wikipedia.org/wiki/Locally_compact_space for other common definitions which are not equivalent in general.

Def. $X \in \text{Top}$ is a **compactly generated / a k-space** (in kTop) if

cpt gen in Condensed Math

Hausdorff-cpt gen/k-space in wiki

cpt gen/k-space in nlab

k-space in ATII

$$\begin{aligned}
 \text{CGwHaus} &= \text{kTop} \cap \text{wHaus} \\
 \text{CGHaus} &= \text{kTop} \cap \text{Haus}
 \end{aligned}$$

$\forall \text{ map } f: X \rightarrow Y,$

f is cont $\Leftrightarrow K \xrightarrow{g} X \xrightarrow{f} Y$ is cont

$\forall K \in \text{CHaus}, g: K \rightarrow X$ cont

equivalently,

$\forall A \subseteq X$ subspace,

$A \subseteq X$ is closed $\Leftrightarrow g^{-1}(A) \subseteq K$ is closed

$\forall K \in \text{CHaus}, g: K \rightarrow X$ cont

When X is Hausdorff, this is equivalent to

$\forall A \subseteq X$ subspace,

$A \subseteq X$ is closed $\Leftrightarrow A \cap K \subseteq K$ is closed

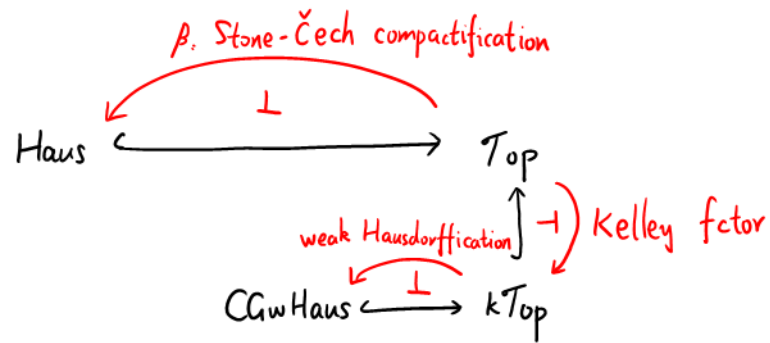
$\forall K \in \text{CHaus}$

Prop. $X \in \text{Top}$, then

X is a k-space $\Leftrightarrow X \cong \coprod_{i \in I} S_i / \sim$ $S_i \in \text{CHaus}$

Rmk. In the def/prop of kTop , CHaus can be replaced by Prof.

Adjoints



Kelley fctor $(-)^{cg} : \mathcal{T}_{op} \longrightarrow k\mathcal{T}_{op}$
 $X \longmapsto X^{cg}$ compactly generated

Set: $X^{cg} = X$

Topo: $A \subseteq X^{cg}$ is closed if $g^{-1}(A) \subseteq K$ is closed
 $\forall K \in \mathcal{CHaus}, g: K \rightarrow X$ cont

Categories in condensed mathematics

$$\begin{array}{ccccccc}
 & & \text{cpt} & \subset & \text{qc} & & \\
 & \swarrow & & & \searrow & & \\
 \text{qcProj} & \subset & \text{Prof} & \subset & \text{qcqs} & \subset & \text{qs} & \subset & \text{CondSet} \\
 \parallel & & \parallel & & \parallel & & \parallel & & \parallel \\
 \text{Cp/BoolAlg}^{\text{op}} & & \text{BoolAlg}^{\text{op}} & & \text{CHaus} & & \text{Ind(CHaus)} & & \bigcup \\
 & & & & & & & & \\
 & & & & & & \text{Liq}_p & \subset & \text{CondAb} \\
 & & & & & & \cup & & \cup \\
 & & & & & & \text{Liq}_p(R) & \subset & \text{Cond}(R)
 \end{array}$$

Appendix

I'm just too lazy to fill in this table. If you know more, tell me and I will fill in, thanks!

ReCRing

Category	cpl	fin cpl	cocpl	fin cocpl	Cartesian closed	closed	monoidal
Set		✓		✓	✓		✓
Top		✓		✓	×		✓
Grp		✓		✓	×		✓
Ab		✓		✓	×		✓
Vect(K)		✓		✓	×		✓
Mod(R)		✓		✓	×		✓
Ring		✓		✓			
CRing		✓		✓			
Rng		✓		✓			
Field	×	×	×	×			
0							
1							
K(z)							
Δ	×	×	×	×			
sSet		✓		✓			
CHaus		✓		✓			
Met	×	✓	×	×			
Quiv(e)							
Quiv							
Cat		✓		✓	✓		✓
kTop							✓
CGHaus					✓		
CGwHaus		✓		✓	✓		
Prof		✓					