

Topological data analysis

Lecture 5

Anton Ayzenberg

ATA Lab, FCS NRU HSE
Noeon Research

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Faculty of Computer Science / Yandex Data School

Persistent homology: practical algorithm

- Previously we had D_1, D_2, \dots matrices of simplicial differentials $\partial_i: C_i(K) \rightarrow C_{i-1}(K)$
- D_i is the incidence matrix between i -dim simplices and $(i-1)$ -dim simplices at least over \mathbb{Z}_2 .
- Now we combine them all in a huge matrix D of size $N \times N$,
- where N is the total number of simplices in a filtration.
- The order of rows and columns as in the array Simplices.

Gauss elimination on columns

- Reduce D by elementary operations **on columns**.
- It is allowed to any column, **to add another column multiplied by a scalar**.
- We move from left to right checking all $1 \leq j \leq N$.
- Let $\text{low}(j)$ denote **the index of the lowest nonzero element** in j -th column.
- For each j we loop through $r < j$, and look for r such that $\text{low}(r) = \text{low}(j)$.
- When we meet such r , we reduce j -th column by r -th column.
- We loop until $\text{low}: [N] \rightarrow [N]$ becomes injective. This is called **the reduced form of a matrix**.
- **Permutations are not allowed**.

You have the Matrix...

	1	2	3	4	23	34	12	24	234	13	14	123	124	134	1234
1					0	0	1	0		1	1				
2					1	0	1	1		0	0				
3					1	1		0		1	0				
4						1		1			1				
23									1			1			
34									1			0	0	1	
12									0			1	1	0	
24									1			0	1	0	
234												0	0	0	1
13												1	0	1	0
14													1	1	0
123	0			0			0			0					1
124	0	0	0				0			0	0				1
134	0	0	0				0			0	0				1
1234	0	0	0				0			0	0				

After Gauss elimination in columns

	1	2	3	4	23	34	12	24	234	13	14	123	124	134	1234
1							1								
2					1		(1)								
3					(1)	1									
4						(1)									
23								1							
34								1							
12												1	1		
24									(1)				1		
234															1
13												(1)			
14													(1)		
123															1
124															1
134															(1)
1234															

Look at the reduced matrix M .

Theorem

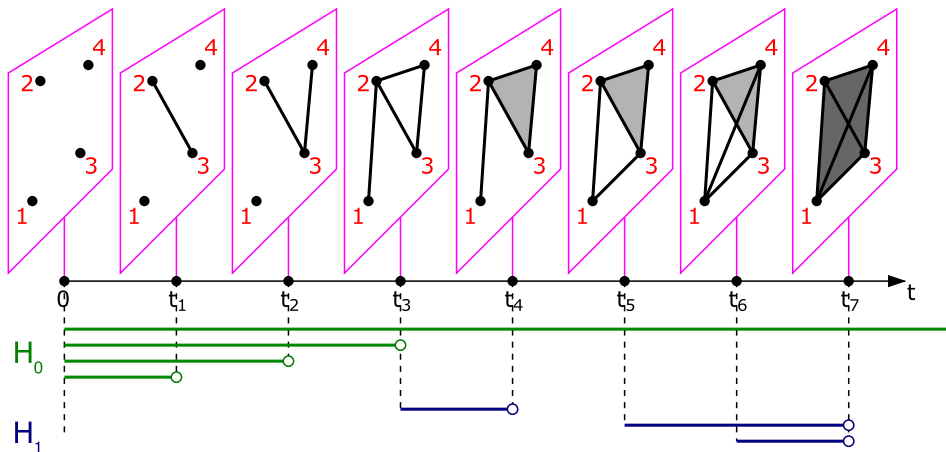
- Let $j = \text{low}(s) \neq 0$ in M . Then **Simplices[j]** gives birth to homology of rank $\dim \text{Simplices}[j]$ and **Simplices[s]** kills this homology. Each such pair (j, s) gives rise to **the interval module** $I_{[\text{BirthTimes}[j], \text{BirthTimes}[s])}$.
- If the whole r -th column vanish and moreover $r \notin \text{low}([N])$, then $\text{Simplices}[r]$ gives birth to homology of rank $\dim \text{Simplices}[r]$ which never dies. It gives **the interval module** $I_{[\text{BirthTimes}[r], +\infty)}$.

Persistent homology of the filtration is the direct sum of all the listed interval modules.

Example

Simplices = [1, 2, 3, 4, 23, 34, 12, 24, 234, 13, 14, 123, 124, 134, 1234]

BirthTimes = [0, 0, 0, 0, t₁, t₂, t₃, t₃, t₄, t₅, t₆, t₇, t₇, t₇, t₇]

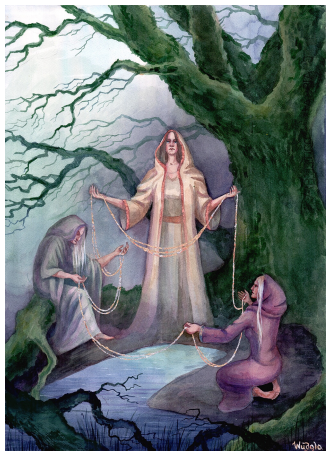


Our reduced matrix

	1	2	3	4	23	34	12	24	234	13	14	123	124	134	1234
1							1								
2					1		(1)								
3					(1)	1									
4						(1)									
23								1							
34								1							
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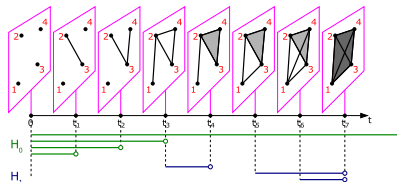
Circled elements have indices: $[3, 23]$, $[4, 34]$, $[2, 12]$, $[24, 234]$, $[13, 123]$, $[14, 124]$, $[134, 1234]$.

Reading births and deaths



Norns determine the fate of
homology

Simplices = [1, 2, 3, 4, 23, 34, 12, 24, 234, 13, 14, 123, 124, 134, 1234]
BirthTimes = [0, 0, 0, 0, t₁, t₂, t₃, t₄, t₅, t₆, t₇, t₇, t₇, t₇]



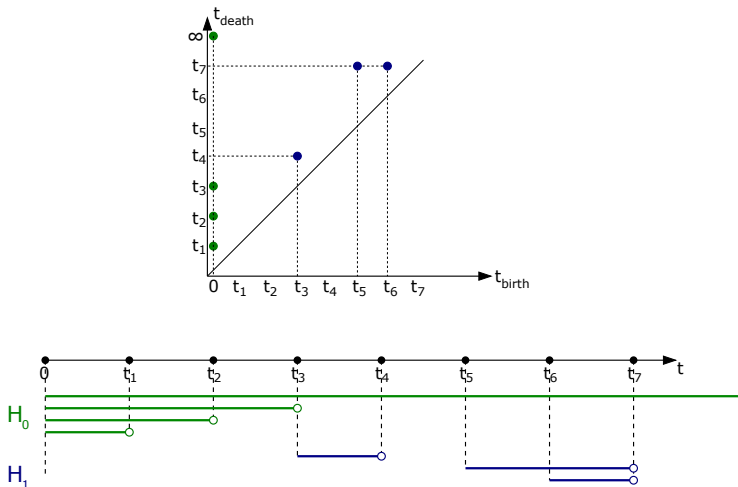
- We see pairs: [3, 23], [4, 34], [2, 12], [24, 234], [13, 123], [14, 124], [134, 1234].
- For example, vertex 3 gives rise to homology which dies when 23 appears.
- 3 appears at time 0, and 23 — at time t_1 .
- Hence we have interval module $I_{[0; t_1]}$ in 0-th homology.
- etc...

We have decomposition: $I_{[0; t_1]} \oplus I_{[0; t_2]} \oplus I_{[0; t_3]} \oplus I_{[t_3; t_4]} \oplus I_{[t_5; t_7]} \oplus I_{[t_6; t_7]}$.

- Let K_t , $t \in \mathbb{R}$ be a collection of simplicial complexes on vertex set $[m]$,
- such that $t_1 < t_2$ implies $K_{t_1} \subseteq K_{t_2}$.
- This is called a **filtration with real time**.
- Changes may occur only at discrete time moments.
- Therefore, the only difference is that BirthTimes stores real values.
- The algorithm above outputs the interval decomposition.

How to encode persistent homology?

Persistence diagram instead of barcode.



- If some homology lives long, then it is a **meaningful homology**.
- The lifetime = death time - birth time.
- Lifetime is long, if the point of persistence diagram is far from the diagonal $y = x$.
- Points close to the diagonal are considered noise.
- Long-live-homology are important topological features of a filtration.

Filtrations: Čech

Demonstration: press to play in browser

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- Let $X = \{x_1, \dots, x_m\} \subset \mathbb{R}^d$ be a point-cloud.
- Let $X_t = \bigcup B_{t/2}(x_i)$.
- X_t is a filtration, but not of simplicial complexes.

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- Let $X_t = \bigcup B_{t/2}(x_i)$.
- X_t is a filtration, but not of simplicial complexes.
- Replace X_t by its **nerve**!

Nerve complex

Let $K_t^{\check{C}}$ be a simplicial complex on $[m]$, such that $\{i_1, \dots, i_s\} \in K_t^{\check{C}}$ iff

$$B_{t/2}(x_{i_1}) \cap \dots \cap B_{t/2}(x_{i_s}) \neq \emptyset$$

Nerve complex and Nerve theorem

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- Therefore X_t and $K_t^{\check{C}}$ have the same homology for each t .
- Hence they have the same persistent homology.
- We can work with simplicial filtration $\{K_t^{\check{C}}\}$ called **\check{C} ech filtration**.

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- Simplex I is born at the time = minimal t for which balls of radii $t/2$ around x_i , $i \in I$, intersect.
- It may be **difficult to find this number** in practice.

Vietoris–Rips filtration

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Vietoris–Rips filtration

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- K_t^{VR} be a simplicial complex on $[m]$ such that
- $I \in K_t^{VR}$ iff all pairwise distances between x_i and x_j , for $i \neq j$, $i, j \in I$, are less than t .
- Simplex I is born at time moment $\max_{i \neq j \in I} \text{dist}(x_i, x_j)$.

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- $\{K_t^{VR}\}$ is called **Vietoris–Rips filtration**.
- Very easy to compute birth times!
- Can be adapted to any finite metric space, e.g. metric graph.

How to compare filtrations and diagrams?

- We have two filtrations $\{K_t\}$ and $\{L_t\}$ on the same set $[m]$.
- Set $\text{dist}(\{K_t\}, \{L_t\}) = \max_{I \subset [m]} |\text{BirthTime}_K[I] - \text{BirthTime}_L[I]|$.

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- If $X = \{x_1, \dots, x_m\}$ and $Y = \{y_1, \dots, y_m\}$ are two data clouds and
- $\text{dist}(X, Y) = \max_i \text{dist}(x_i, y_i)$, and
- K_t^{VR} and L_t^{VR} are their Vietoris–Rips filtrations

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- **Exercise:** $\text{dist}(\{K_t^{VR}\}, \{L_t^{VR}\}) \leq 2 \text{dist}(X, Y)$.

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- **Exercise:** $\text{dist}(\{K_t^{VR}\}, \{L_t^{VR}\}) \leq 2 \text{dist}(X, Y)$.
- What about their persistent diagrams?

Metric on diagrams

Bottleneck distance (or Wasserstein, or Kantorovich metric)

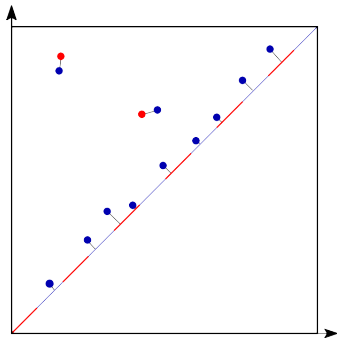


Fig.from GUDHI Library

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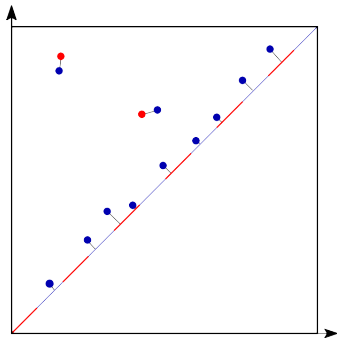







Fig. from GUDHI Library

Stability theorem

$$\text{dist}(\text{PD}(\{K_t\}), \text{PD}(\{L_t\})) \leq \text{dist}(\{K_t\}, \{L_t\}).$$

Generalizations of persistent modules

We can hardly see this slide. However if we do, then it is time to switch to whiteboard

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