

Formal Analysis

Hilary Term, Uncertainty and PBE

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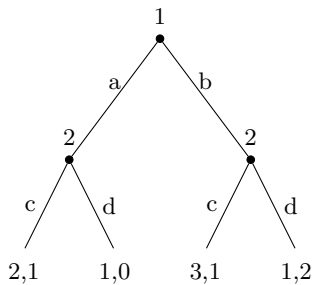
February 15, 2022

Agenda

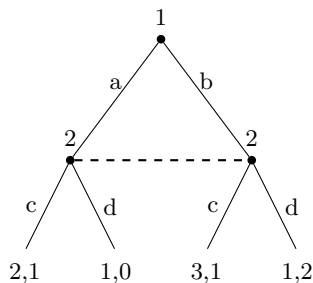
1. Information, uncertainty and Perfect Bayesian equilibria
2. Type 1: Games where uninformed player moves first
3. Type 2: Games where both players are uninformed
4. Type 3: Games where uninformed player moves second

Perfect information

An extensive game

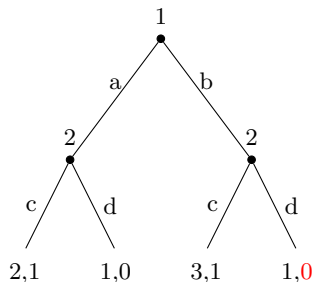
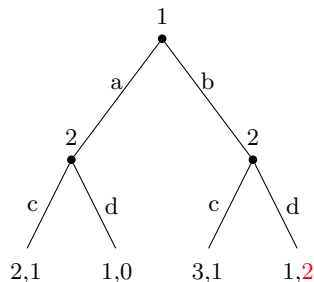


Uncertainty I



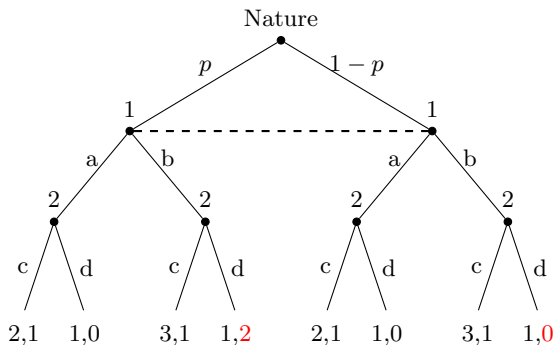
- ▶ 2 does not observe 1's move
- ▶ 2 does not know whether she is at left or right node
- ▶ a game with *imperfect information*

Uncertainty II



- ▶ 1 does not know 2's payoff (or, utility function of player 2, the *type* of player 2, the *state of the world*, ...)
- ▶ 1 does not know whether they are playing left or right game
- ▶ 'incomplete' description, not really a game ...

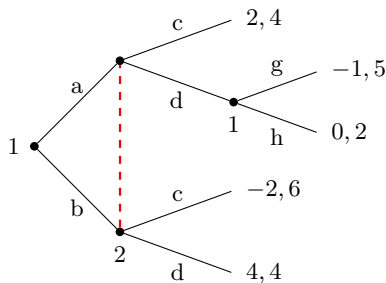
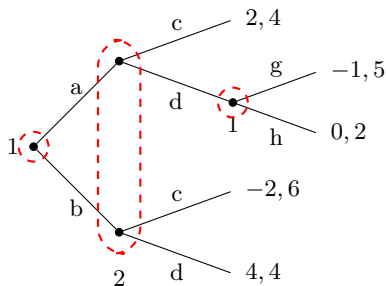
Uncertainty II



- ▶ add *Nature* as an additional, non-strategic player selecting *types* or *states* according to some (known) probability distribution
- ▶ a game with *incomplete information*

Information sets

Definition (Information set). An *information set* of a player is a set of nodes of the game tree among which she cannot distinguish. The player does not “know” where she is. A *singleton* is an information set that consists of only one node.



Perfect and complete information

Perfect information

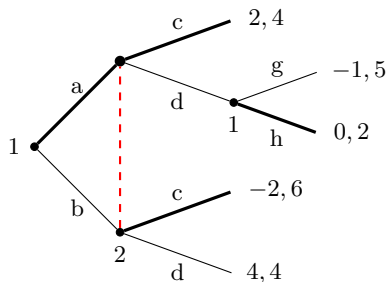
- ▶ in a game with perfect information every information set is a singleton
- ▶ otherwise it's a game with imperfect information

Complete information

- ▶ in a game with *incomplete information* “Nature” moves first and is not observed by at least one player
- ▶ otherwise the game is referred to as one of complete information

Beliefs

The previous example is instructive



- ▶ we kind of know that (ah, c) is the equilibrium
- ▶ so while 2 cannot be certain where she is, the only reasonable *belief* is: “Hey, I am at the upper node!”

Beliefs

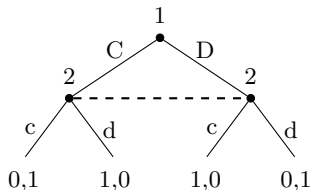
We shall define this term properly

Definition (Belief). A belief b at an information set I is a probability distribution over all elements of I .

In the previous example, we would write $b = (1, 0)$.

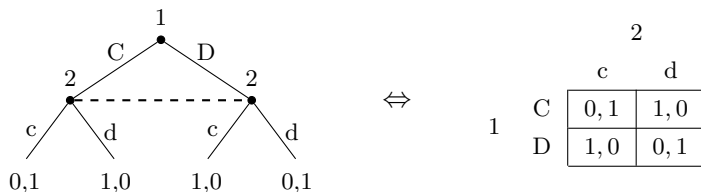
Another example (and more on normal form vs extensive games)

What is 2's belief?



- ▶ with *perfect* information, $c|C$ and $d|D$ would be best responses
- ▶ whereas with imperfect information ...
- ▶ players move simultaneously ... A normal form game!

Another example (and more on normal form vs extensive games)



- ▶ matching pennies has a unique NE (in mixed strategies):
 $s = ((\frac{1}{2}C, \frac{1}{2}D), (\frac{1}{2}c, \frac{1}{2}d))$
- ▶ hence, given s , 2's belief must be $b = (\frac{1}{2}, \frac{1}{2})$
- ▶ ... and given b , $s_2 = (\frac{1}{2}c, \frac{1}{2}d)$ indeed is a best response

Sequentially rational strategies and weakly consistent beliefs

Two more definitions

Definition (Sequential rational). Given a profile of beliefs b , a strategy profile s is *sequentially rational* if at any information set, s is a best response given b

Definition (Weakly consistent). A belief profile b is *weakly consistent* relative to strategy s if the beliefs are formed “according to Bayes’ rule whenever possible”.

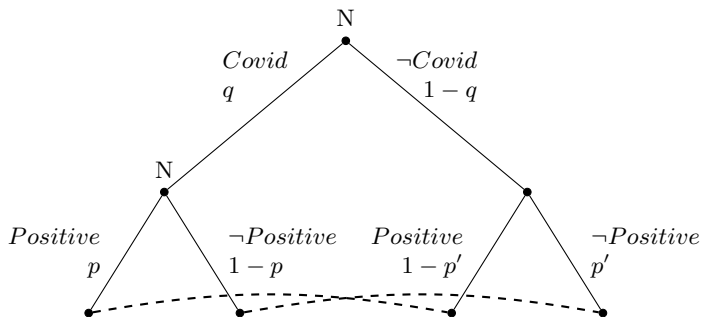
Bayes' rule

We still have to define what is meant by “according to Bayes' rule whenever possible”

- Recall Bayes' rule

$$\begin{aligned}\Pr(A|B) &= \frac{\Pr(B|A) \Pr(A)}{\Pr(B)} \\ &= \frac{\Pr(B|A) \Pr(A)}{\Pr(B|A) \Pr(A) + \Pr(B|\textit{not}A) \Pr(\textit{not}A)}\end{aligned}$$

Belief updating using Bayes' rule



What is a consistent belief for the first information set?

$$\begin{aligned}\Pr(C \mid P) &= \frac{\Pr(P \mid C) \Pr(C)}{\Pr(P \mid C) \Pr(C) + \Pr(P \mid \neg C) \Pr(\neg C)} \\ &= \frac{pq}{pq + (1 - p')(1 - q)}\end{aligned}$$

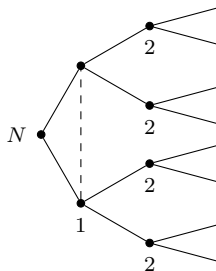
Perfect Bayesian equilibrium

Definition (Perfect Bayesian equilibrium). In a sequential game with imperfect information the pair (s, b) of strategy profile s and belief profile b is called a *perfect Bayesian equilibrium* if

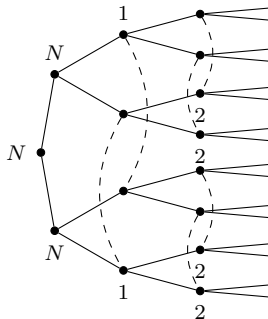
- ▶ s is sequentially rational with regard to b , and
- ▶ b is weakly consistent with regard to s

Note: if an information set is not reached under the strategy then Bayes' rule does not work and **all** beliefs are weakly consistent (these are called “out-of-equilibrium-beliefs”)

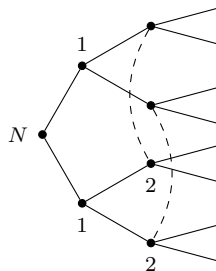
Typical information structures in incomplete information games



type 1

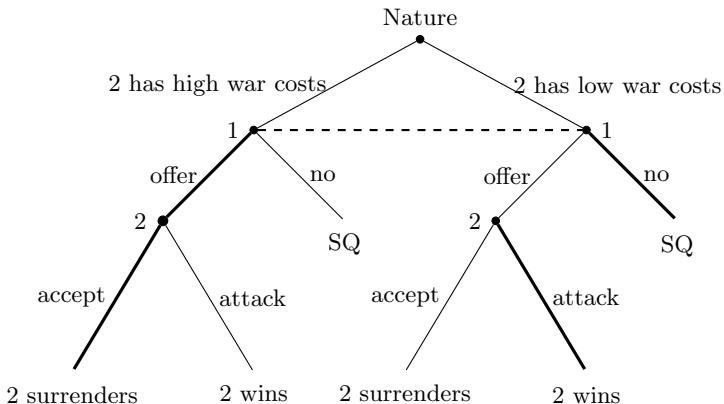


type 2

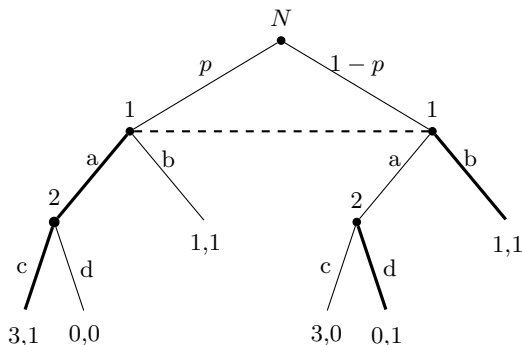


type 3

Type 1: Simplified bargaining with uncertainty over costs



Same story with parameters ...



- ▶ 1's belief simply is $b_1 = (p, 1 - p)$
- ▶ expected utility from a and b :

$$EU_1(a) = 3p + 0(1 - p) = 3p$$

$$EU_2(b) = 1p + 1(1 - p) = 1$$

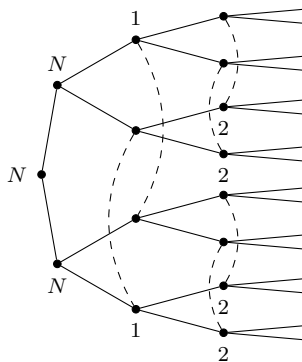
Type 1

$$EU_1(a) = 3p$$

$$EU_2(b) = 1$$

- ▶ a is best response iff $p \geq \frac{1}{3}$, b is best response iff $p \leq \frac{1}{3}$
- ▶ thus $s = (a, cd)$ is sequentially rational with regard to b_1 iff $p \geq \frac{1}{3}$, $s = (b, cd)$ is sequentially rational with regard to b_1 iff $p \leq \frac{1}{3}$
- ▶ b_1 is weakly consistent with regard to s
- ▶ $PBE = \{(a, cd; (p, 1-p)) \text{ if } p \geq \frac{1}{3}, (b, cd; (p, 1-p)) \text{ if } p \leq \frac{1}{3}\}$

Type 2: Both are uninformed about other player's type



Preventive war as an assurance game

		2	
		<i>Not</i>	<i>Attack</i>
1	<i>Not</i>	s_1, s_2 (1, 1)	$p_1^s - c_1, p_2^f - c_2$ (.1, .9)
	<i>Attack</i>	$p_1^f - c_1, p_2^s - c_2$ (.9, .1)	$p_1 - c_1, p_2 - c_2$ (.5, .5)

- ▶ p 's are winning probabilities, c 's are costs of war
- ▶ assumption: $p_i^{second} < p_i < p_i^{first}$
- ▶ if $s_i \geq p_i^f - c_i$: two NE, one is Pareto-superior

Preventive war as an assurance game when players are either greedy (G) or status quo (S) types

		q_2		$1 - q_2$	
		N	A	N	A
q_1	N	s_1^S, s_2^S	$p_1^s - c_1,$ $p_2^f - c_2$	s_1^S, s_2^G	$p_1^s - c_1,$ $p_2^f - c_2$
	A	$p_1^f - c_1,$ $p_2^s - c_2$	$p_1 - c_1,$ $p_2 - c_2$	$p_1^f - c_1,$ $p_2^s - c_2$	$p_1 - c_1,$ $p_2 - c_2$
		N	A	N	A
$1 - q_1$	N	s_1^G, s_2^S	$p_1^s - c_1,$ $p_2^f - c_2$	s_1^G, s_2^G	$p_1^s - c_1,$ $p_2^f - c_2$
	A	$p_1^f - c_1,$ $p_2^s - c_2$	$p_1 - c_1,$ $p_2 - c_2$	$p_1^f - c_1,$ $p_2^s - c_2$	$p_1 - c_1,$ $p_2 - c_2$

► assume $s_i^G < p_i^f - c_i < s_i^S = 1$

Preventive war as an assurance game when players are either greedy (G) or status quo (S) types

- ▶ for greedy types: dominant strategy is to attack (check!)
- ▶ for status quo types:

$$EU_i(N) = q_{-i}s_i^S + (1 - q_{-i})(p_i^s - c_i)$$

$$EU_i(A) = q_{-i}(p_i^f - c_i) + (1 - q_{-i})(p_i - c_i)$$

Solving $EU_i(N) \geq EU_i(A)$ leads to

$$q_{-i} \geq \frac{1}{\frac{1-(p_i^f-c_i)}{p_i-p_i^s} + 1}$$

in plain words: S -types will not attack if the prob. that the opponent is an S -type is large enough (though they will be attacked with prob. q_{-i})

Preventive war as an assurance game when players are either greedy (G) or status quo (S) types

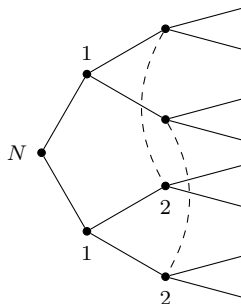
Let

$$T = \frac{1}{\frac{1-(p_i^f-c_i)}{p_i-p_i^s} + 1}.$$

As $\frac{\partial T}{\partial p_i^f} > 0$, $\frac{\partial T}{\partial c_i} < 0$, we have

- ▶ if the level of “trust” (that the opponent is an S -type) falls below T , conflict can occur *even if both* are S -types (the problem of mistrust)
- ▶ if first strike advantage (winning prob.) is large, or costs of war are small, more trust is needed

Type 3: Uninformed player moves second



- ▶ (probably) most interesting case: can 2 *learn* something about state from 1's move?
- ▶ cp. PA framework: (informed) agent moves first, (uninformed) principal second

Interest group lobbying

- ▶ sequence of moves
 - ▶ N chooses state of the world A or B (e.g. ‘policy a is effective in reducing unemployment’ vs ‘policy b is effective in reducing unemployment’)
 - ▶ 1 (lobbyist) sends a message α or β (α = ‘ a is effective in reducing unemployment’, β = ‘ b is effective in reducing unemployment’)
 - ▶ 2 (policymaker) chooses policy a or b
- ▶ payoffs
 - ▶ policymaker wants to match policy to state of the world

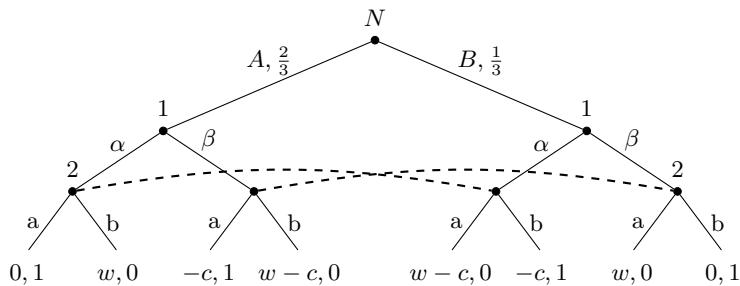
$$u_2(a, A), u_2(b, B) \geq u_2(a, B), u_2(b, A)$$

- ▶ lobbyist has opposed preferences

$$w = u_1(a, B), u_1(b, A) \geq u_1(a, A), u_1(b, B) = 0$$

but incurs costs $c > 0$ if caught lying

Interest group lobbying



Finding equilibria

Before we start

- ▶ focus on pure strategy equilibria (no mixed strategies for now)
- ▶ 1 has four pure strategies

$$(\alpha|A, \beta|B), (\alpha|A, \alpha|B), (\beta|A, \beta|B), (\beta|A, \alpha|B)$$

- ▶ 2 has four pure strategies

$$(a|\alpha, b|\beta), (a|\alpha, a|\beta), (b|\alpha, b|\beta), (b|\alpha, a|\beta)$$

- ▶ to find equilibria we simply posit one and then check whether someone has an incentive to defect

Truthful reporting?

- ▶ *suppose*

$$s_1 = (\alpha|A, \beta|B)$$

is an equilibrium strategy

- ▶ then $\Pr(A|\alpha) = \Pr(B|\beta) = 1$, $\Pr(B|\alpha) = \Pr(A|\beta) = 0$ i.e.
2's weakly consistent belief is

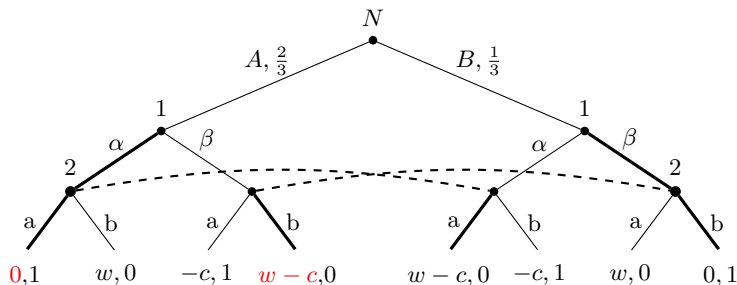
$$b_2 = ((1, 0), (0, 1))$$

- ▶ hence 2's sequentially rational strategy is

$$s_2 = (a|\alpha, b|\beta)$$

- ▶ now, is s_1 really a best response?

Truthful reporting?



- ▶ if $0 \geq w - c$ then 1 has no incentive to switch to $(\beta|A, \beta|B)$
- ▶ same holds for $(\alpha|A, \alpha|B)$ and $(\beta|A, \alpha|B)$
- ▶ thus $((\alpha|A, \beta|B), (a|\alpha, b|\beta); ((1, 0), (0, 1)))$ is a PBE if $w \leq c$
- ▶ this is a *separating* equilibrium

Same message?

- ▶ *assume*

$$s_1 = (\alpha|A, \alpha|B)$$

- ▶ then

$$\Pr(A|\alpha) = \frac{2}{3} \quad \checkmark$$

$$\Pr(B|\alpha) = \frac{1}{3} \quad \checkmark$$

$$\Pr(A|\beta) = \Pr(\beta|A) \Pr(A) / \Pr(\beta) = \frac{0p}{0} \quad \nlessgtr$$

$$\Pr(B|\beta) = \Pr(\beta|B) \Pr(B) / \Pr(\beta) = \frac{0(1-p)}{0} \quad \nlessgtr$$

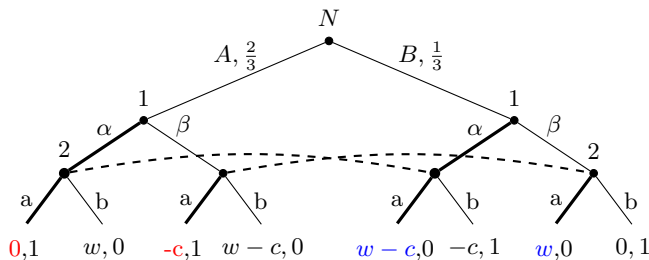
so *one* (of many) weakly consistent belief is

$$b_2 = ((\frac{2}{3}, \frac{1}{3}), (\frac{2}{3}, \frac{1}{3}))$$

- ▶ then 2's sequentially rational strategy is

$$s_2 = (a|\alpha, a|\beta)$$

Same message?



- ▶ as $0 \geq -c$, 1 has no incentive to switch to $(\beta|A, \alpha|B)$ ✓
- ▶ if $w - c \geq w$, 1 has no incentive to switch to $(\alpha|A, \beta|B)$ or $(\beta|A, \beta|B)$
- ▶ as $w - c \geq w \iff c \leq 0$
 $((\alpha|A, \alpha|B), (a|\alpha, a|\beta); ((\frac{2}{3}, \frac{1}{3}), (\frac{2}{3}, \frac{1}{3})))$ is a PBE if $c = 0$
- ▶ this is a *pooling* equilibrium

Note on readings

- ▶ Kydd, chap. 6.2 is slightly more complex version (using continuous strategy spaces) of type-1 games
- ▶ Kydd, chap. 6.3 and 6.5 are versions of type-2 games
- ▶ none of them actually requires PBE, using subgame perfection is perfectly fine
- ▶ type-3 games *do* require PBE; our example foreshadows Kydd, chap. 9 (see week 7 & 8)