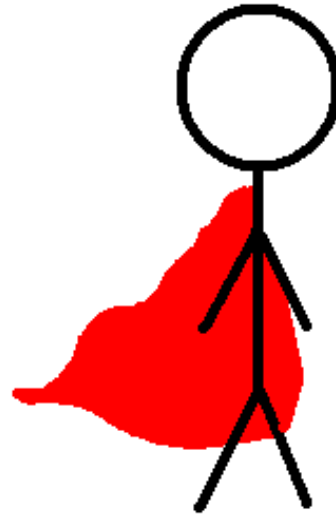


Faster Secure Multi-Party Computation of AES and DES Using Lookup Tables

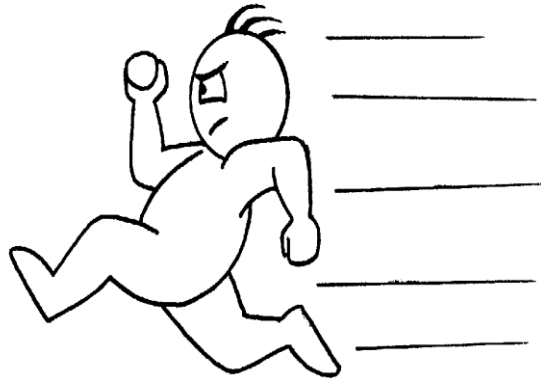
Marcel Keller, Emmanuela Orsini, **Dragos Rotaru**, Peter Scholl,
Eduardo Soria-Vazquez, and Srinivas Vivek

ACNS 2017

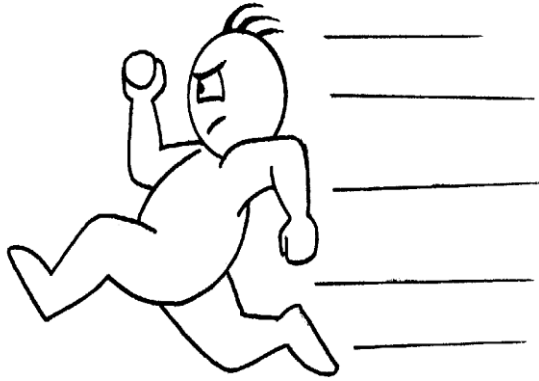
Meet Bob



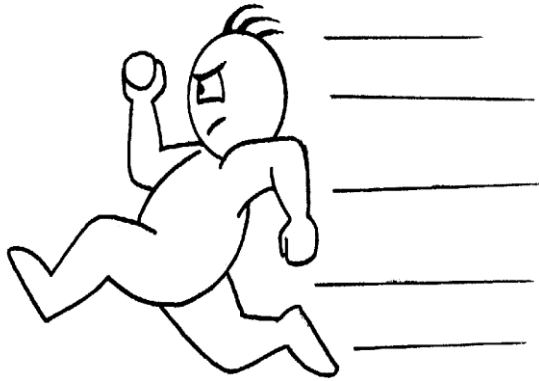
Bob approach to our title: Faster Secure Multi-Party Computation of AES and DES Using Lookup Tables



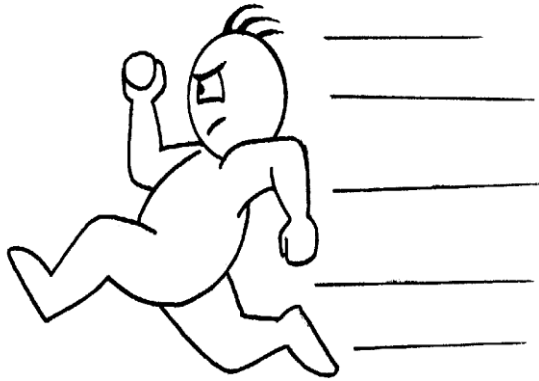
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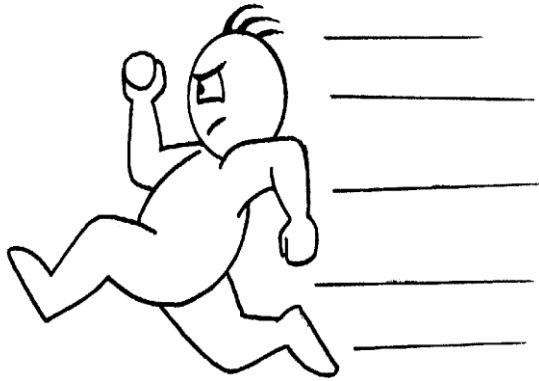
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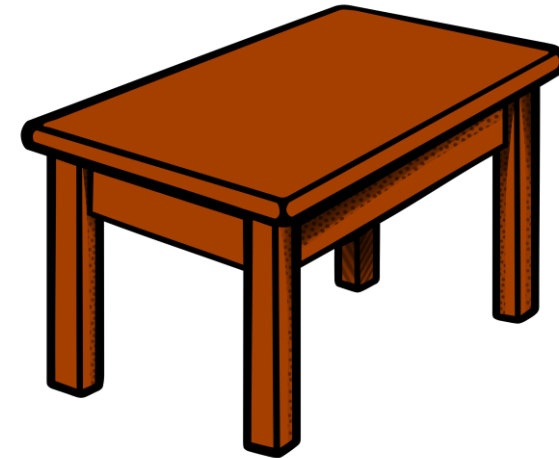
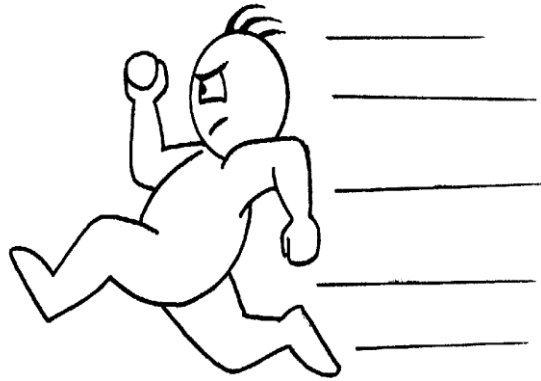
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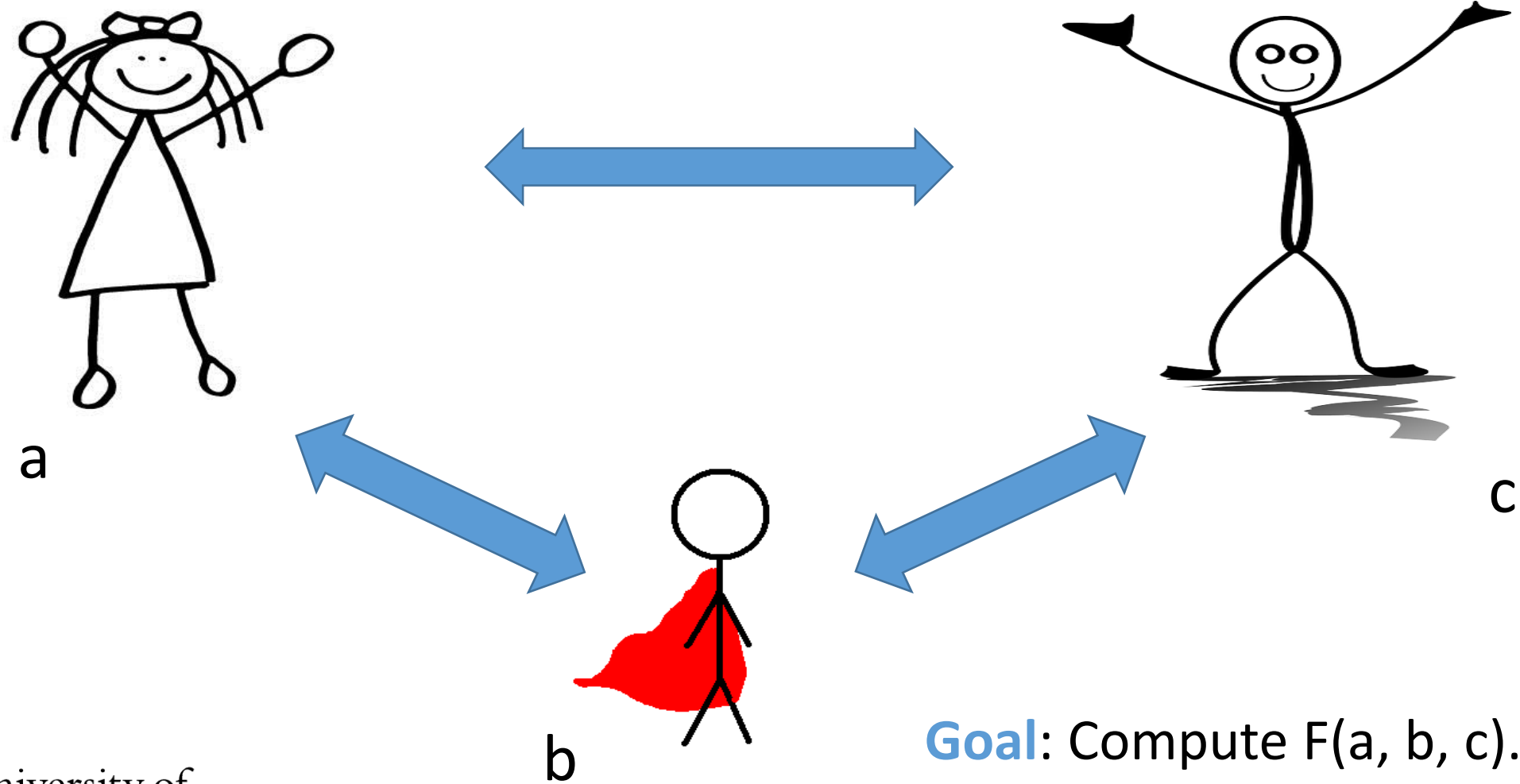
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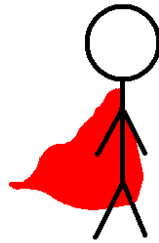


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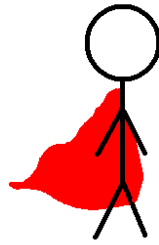


Multi-Party Computation



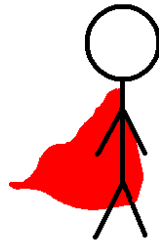


has problems.

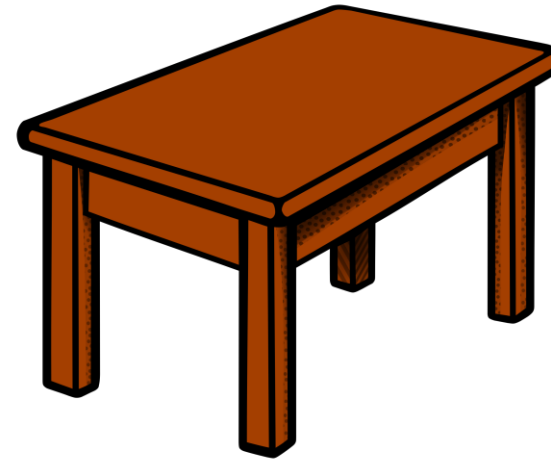


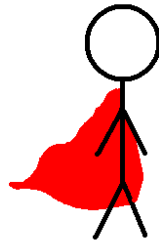
has problems.





has problems.





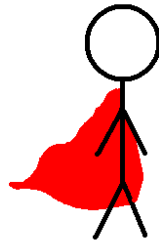
has problems?

Look-up tables are everywhere in MPC.

Floating
Point

Oblivious
RAM

Non-linear
functions



has problems?

Look-up tables are everywhere in MPC.

Floating
Point

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Non-linear? AES and 3-DES



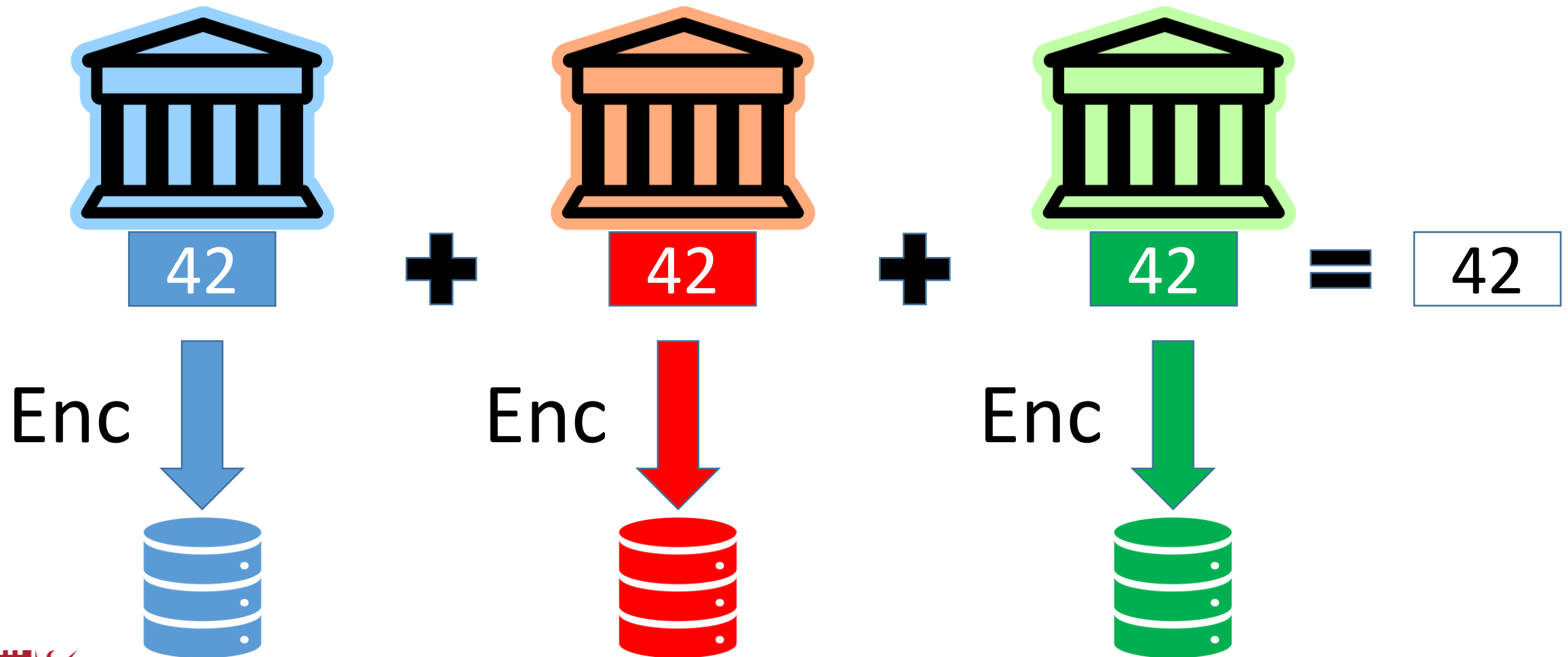
Non-linear? AES and 3-DES

The diagram illustrates the non-linearity of AES and 3-DES using a visual equation. It features three building icons, each with a colored outline and a corresponding colored box containing the number 42. The first building is blue, the second is orange, and the third is green. These are followed by plus signs and an equals sign, leading to a final box containing the number 42. This visualizes the concept that the output of a non-linear function (like AES or 3-DES) is not simply the sum of its inputs, but rather a complex transformation that results in a different value (in this case, 42).

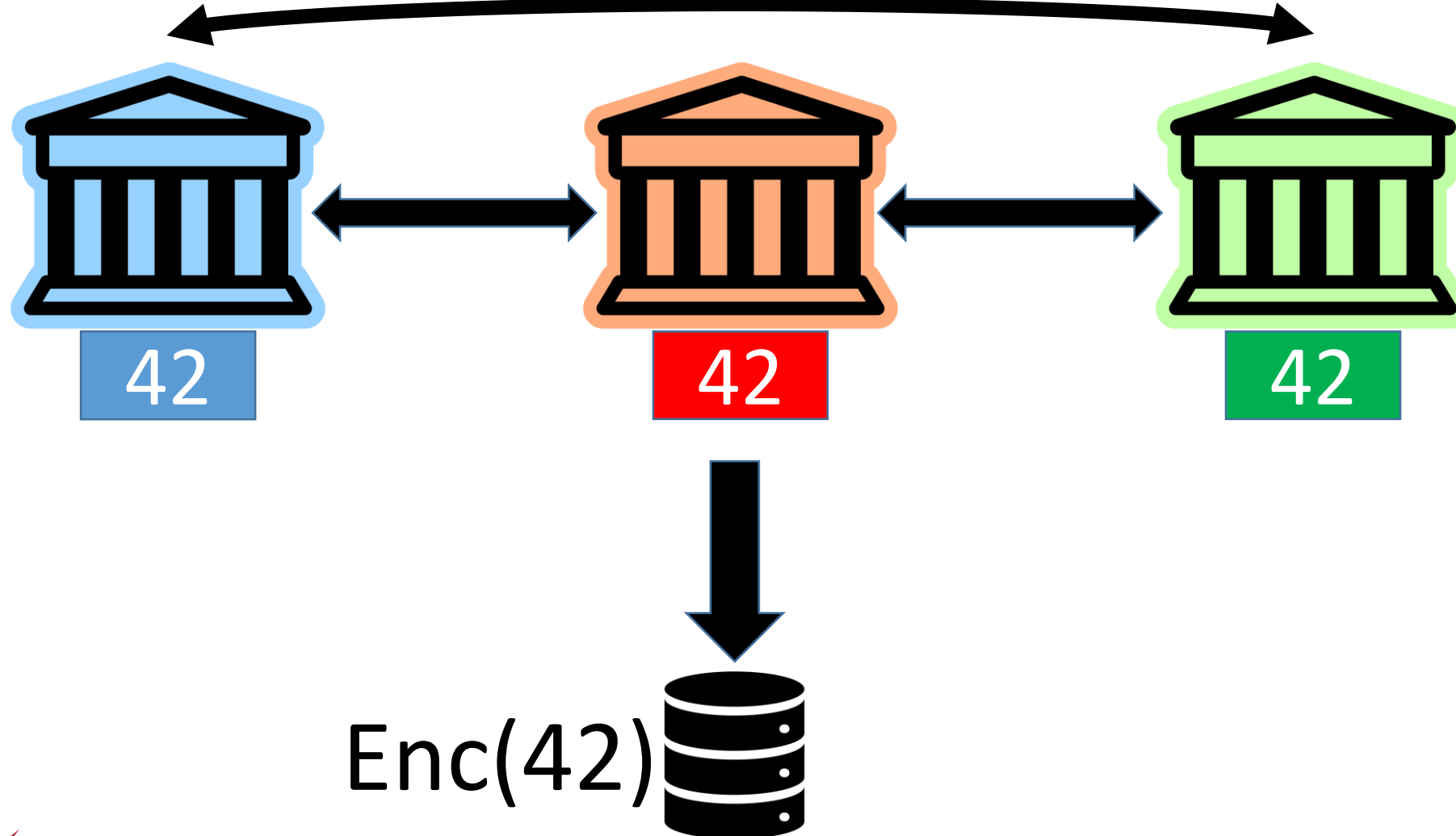
$$\text{Building 1} + \text{Building 2} + \text{Building 3} = \text{Result}$$

42 + 42 + 42 = 42

Non-linear? AES and 3-DES



Non-linear? AES and 3-DES



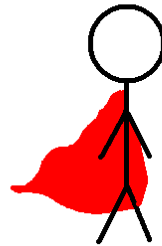
Fastest AES and 3-DES in MPC with malicious security

- Improve on previous AES TinyTable by at least 50 times.
- 3-DES has now 100 times faster online time.
- Apply side-channel countermeasures in the MPC land.

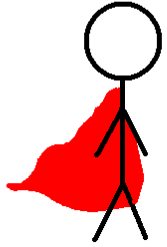
Concurrent Work

- [DNNR16] – TinyTable. Improved version now at Crypto17.
- [DKS+17] – Dessouky et al. in NDSS17. Semi-Honest setting based on 1-out-of-N OT. Also built a compiler which can be used with our protocol.

MPC with Secret Sharing 101



MPC with Secret Sharing 101



$$x = x_1 + \dots + x_n$$

Each P_i has $[x] \leftarrow x_i$



$$[x] \leftarrow x_1$$

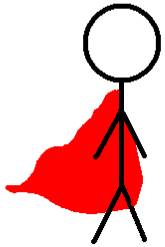


$$[x] \leftarrow x_2$$



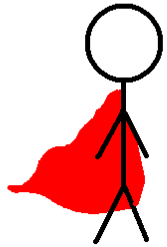
$$[x] \leftarrow x_3$$

MPC Preprocessing Phase



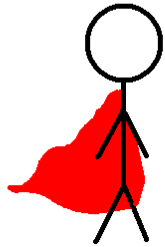
Generate Triples.
 $[c] = [a][b]$

MPC Preprocessing Phase

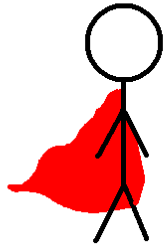


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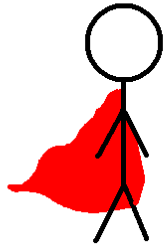


MPC Online Phase

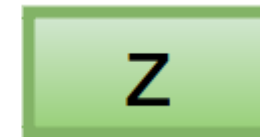
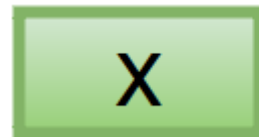


Use Triples.

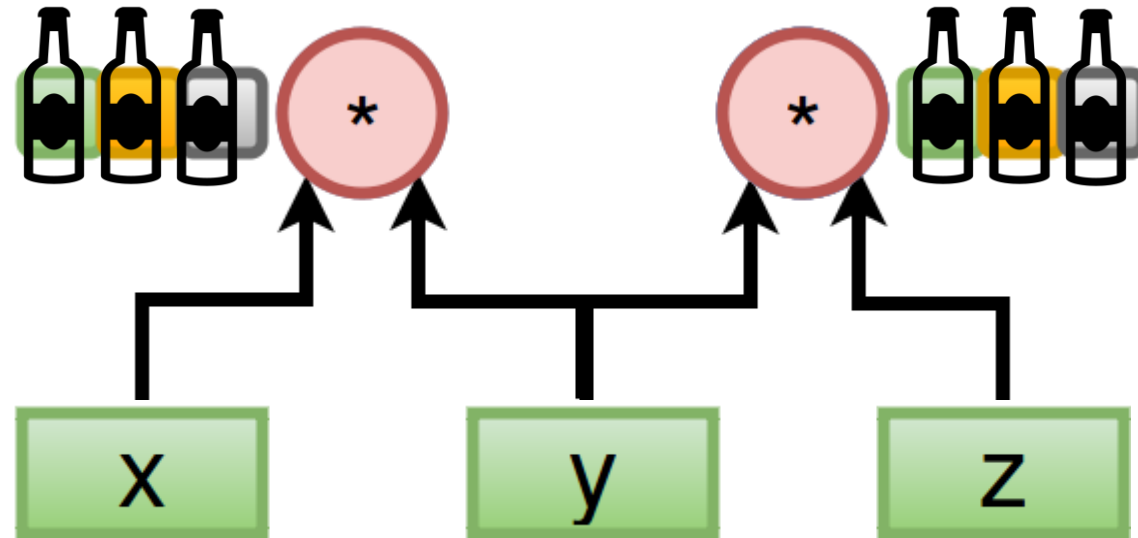
MPC Online Phase



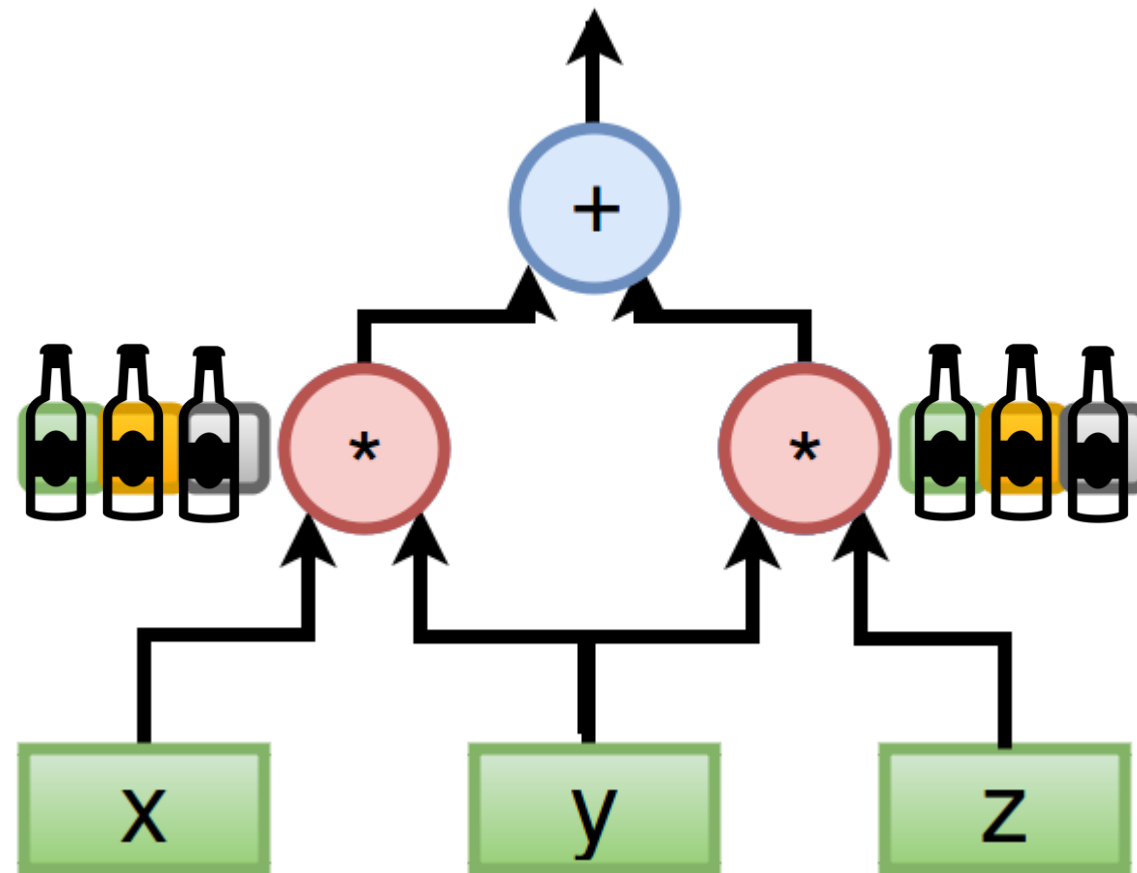
MPC Circuit Evaluation



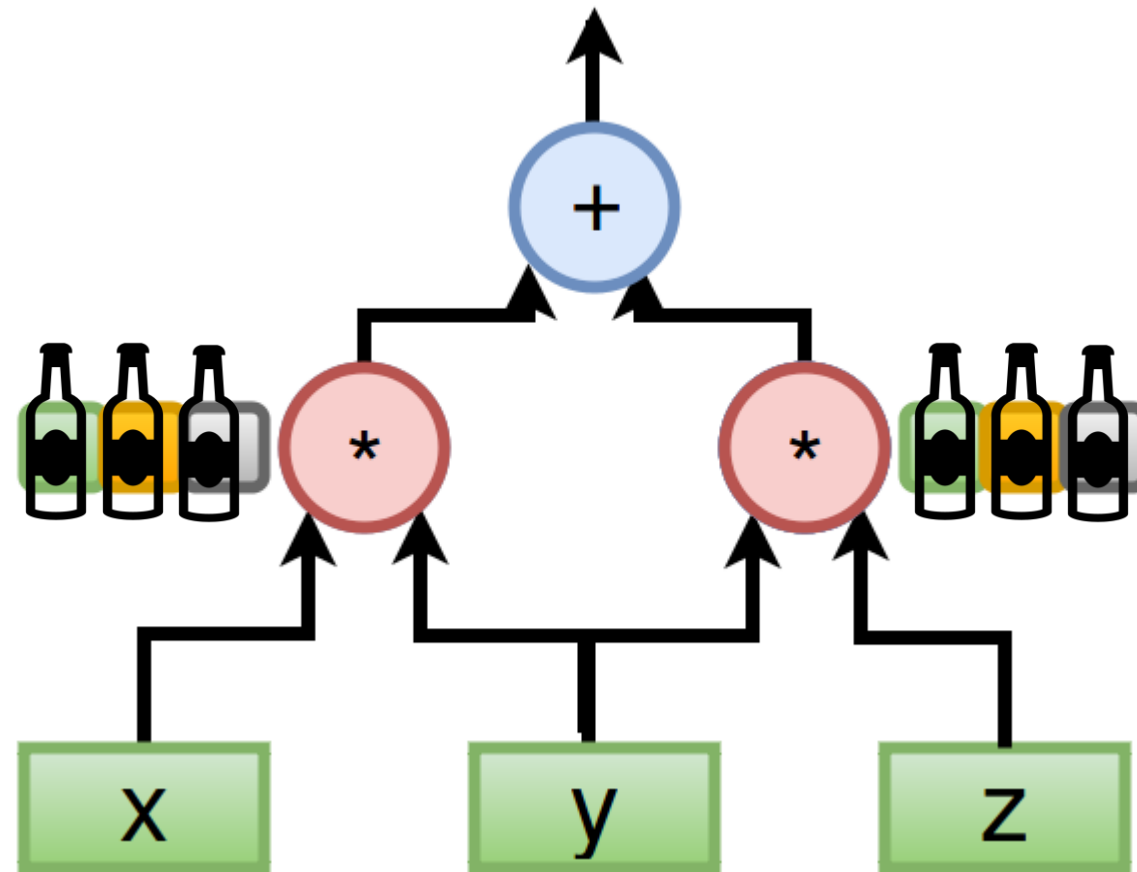
MPC Circuit Evaluation



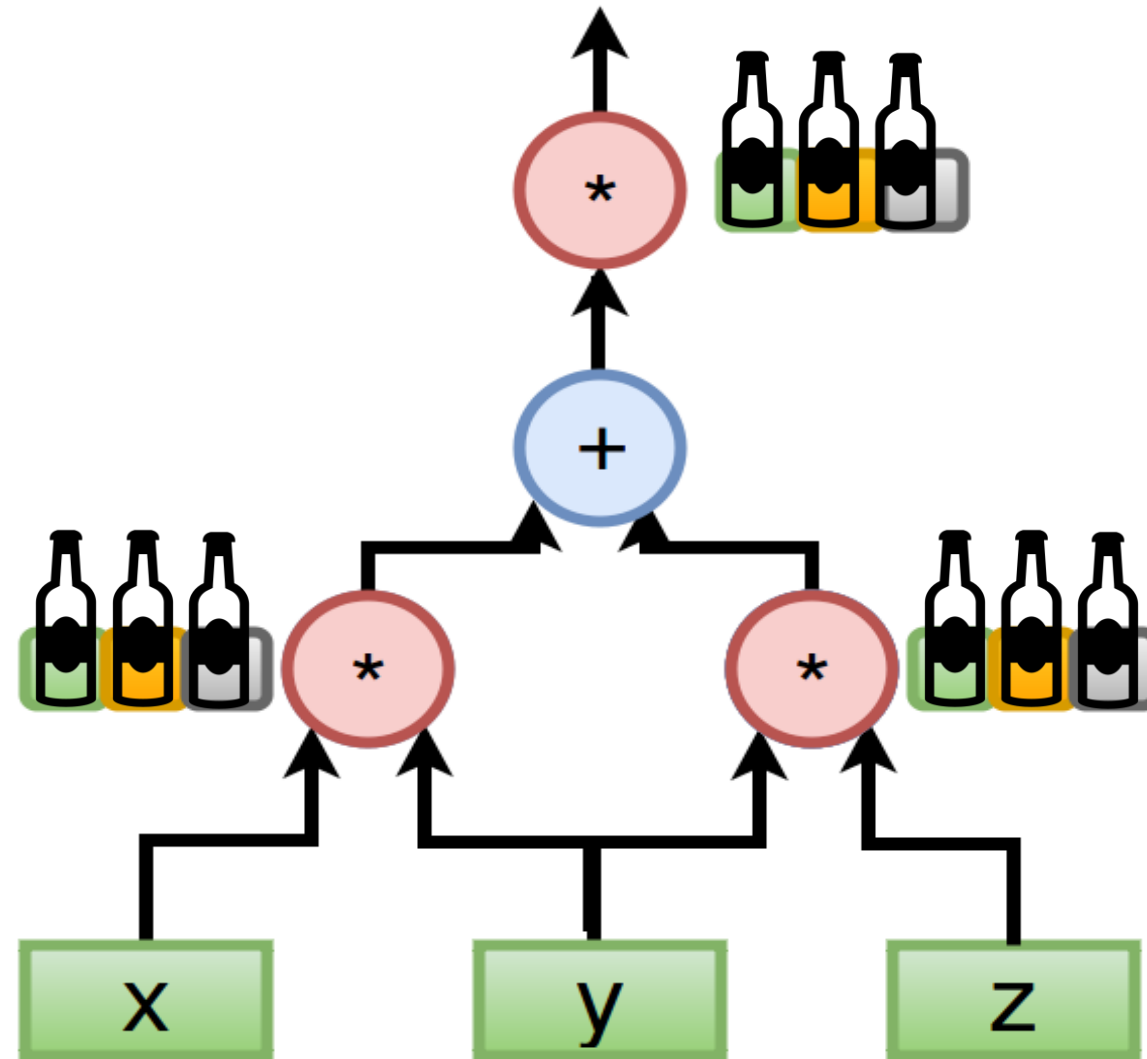
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MPC Circuit Evaluation



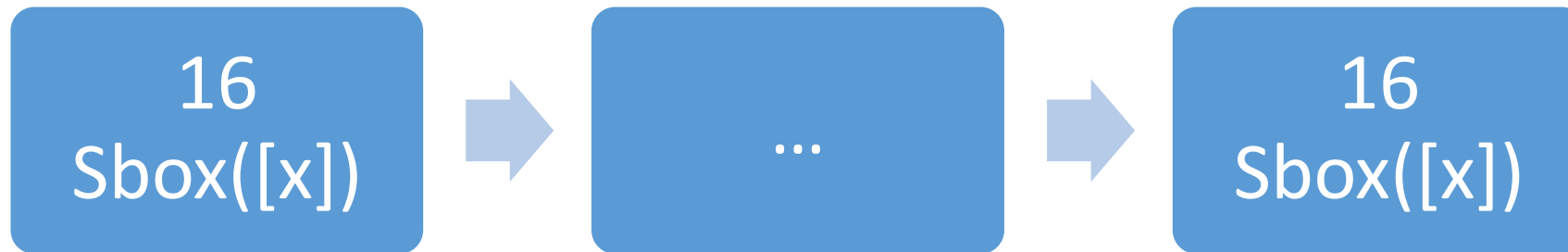
MPC Circuit Evaluation



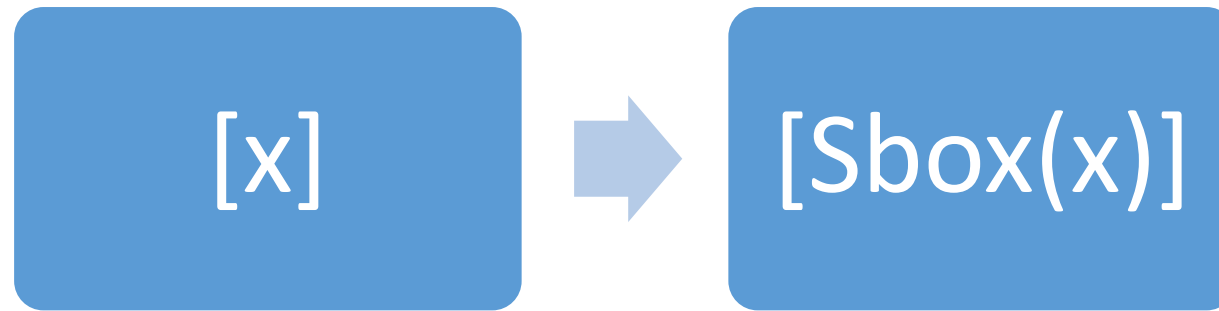
3 triples.
2 rounds.

AES-128

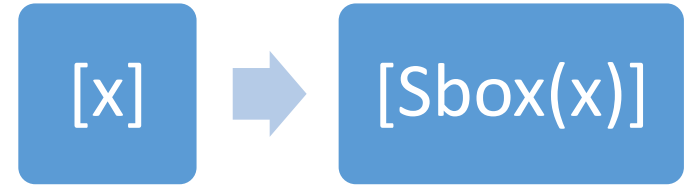
10 rounds



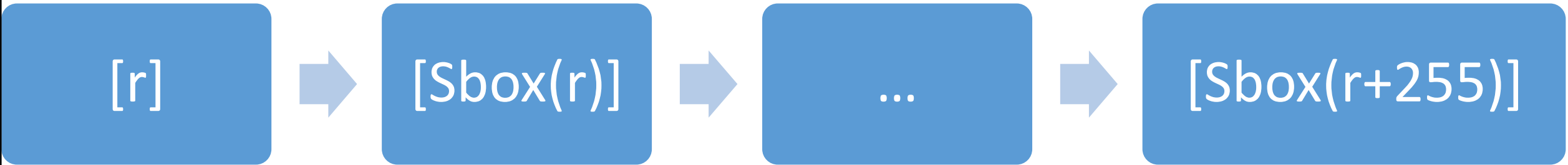
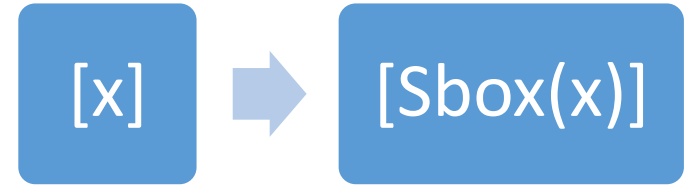
How to Sbox



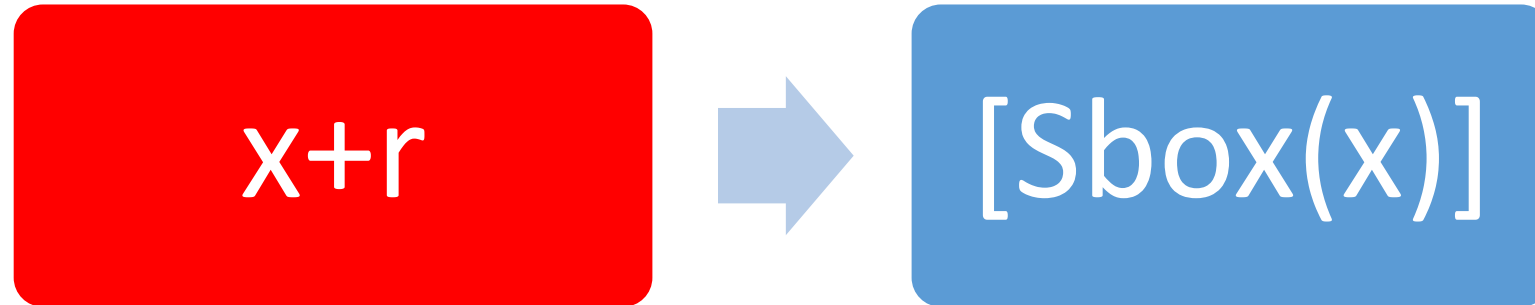
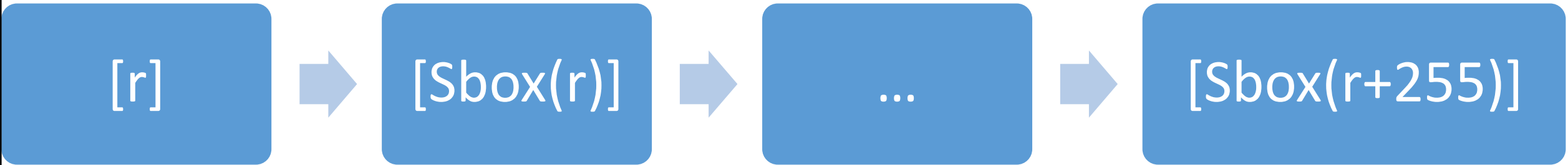
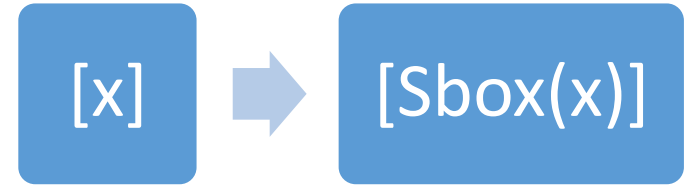
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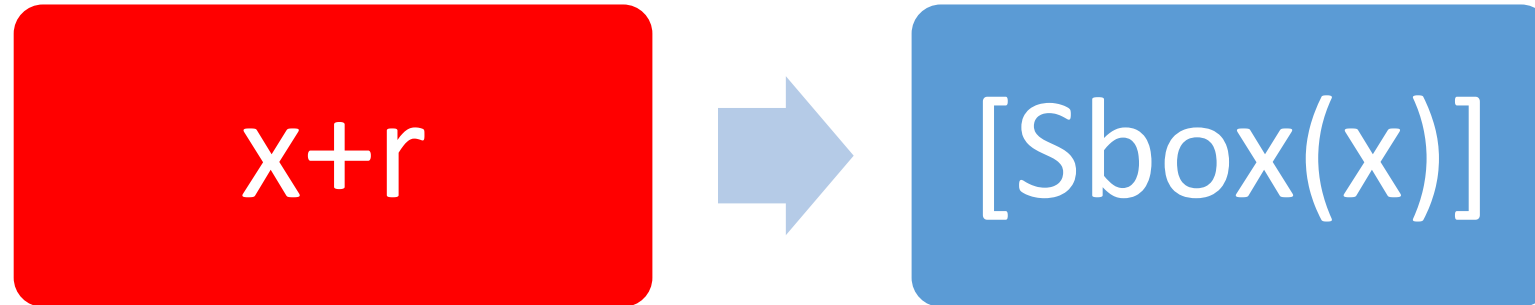
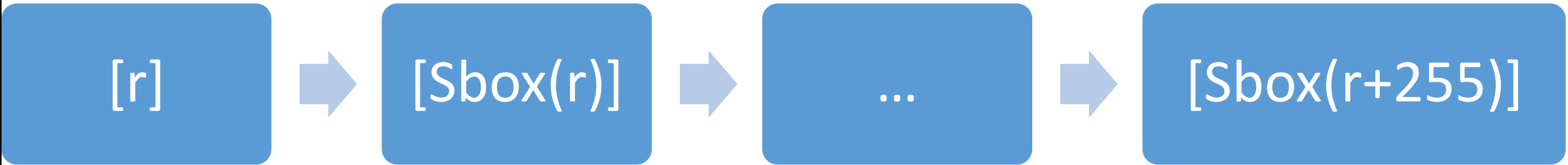
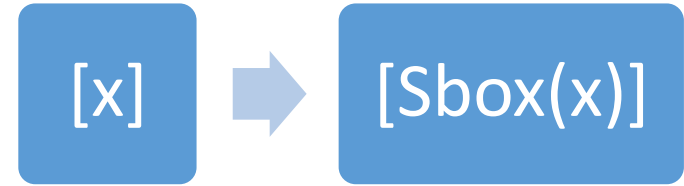
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How to Sbox

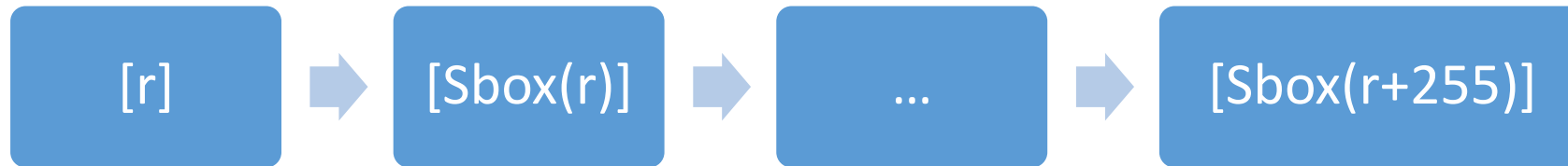


How to Sbox



At pos $(x+r) \Rightarrow Sbox(r + x + r)$

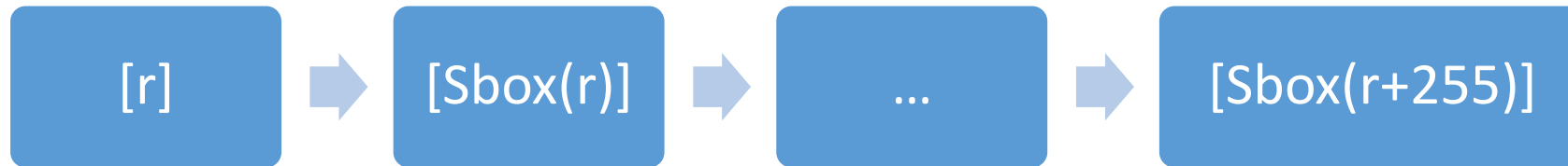
How to pre-Sbox



Take random $[r]$.

Compute $[Sbox(r)]$, ... $[Sbox(r+255)]$

How to pre-Sbox

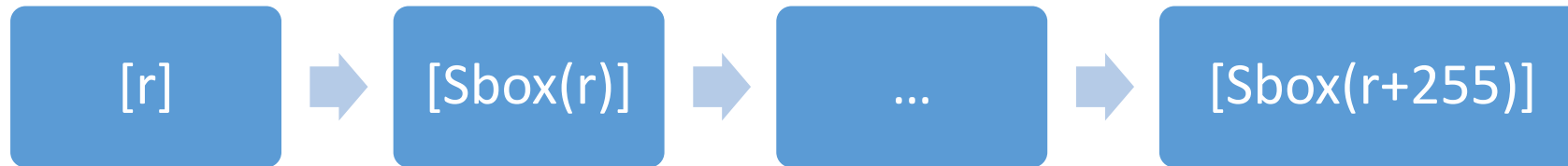


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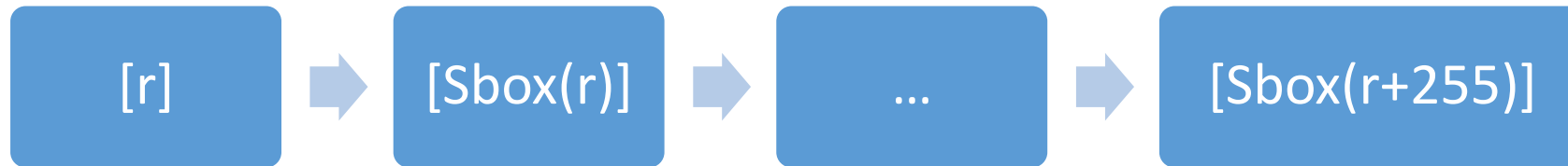


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7 mults.



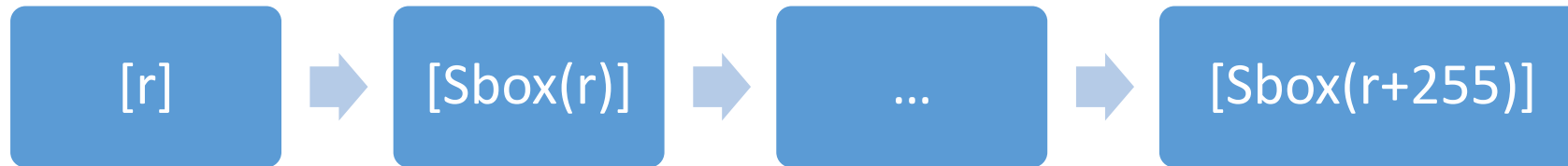
7 mults.



1792 mults.

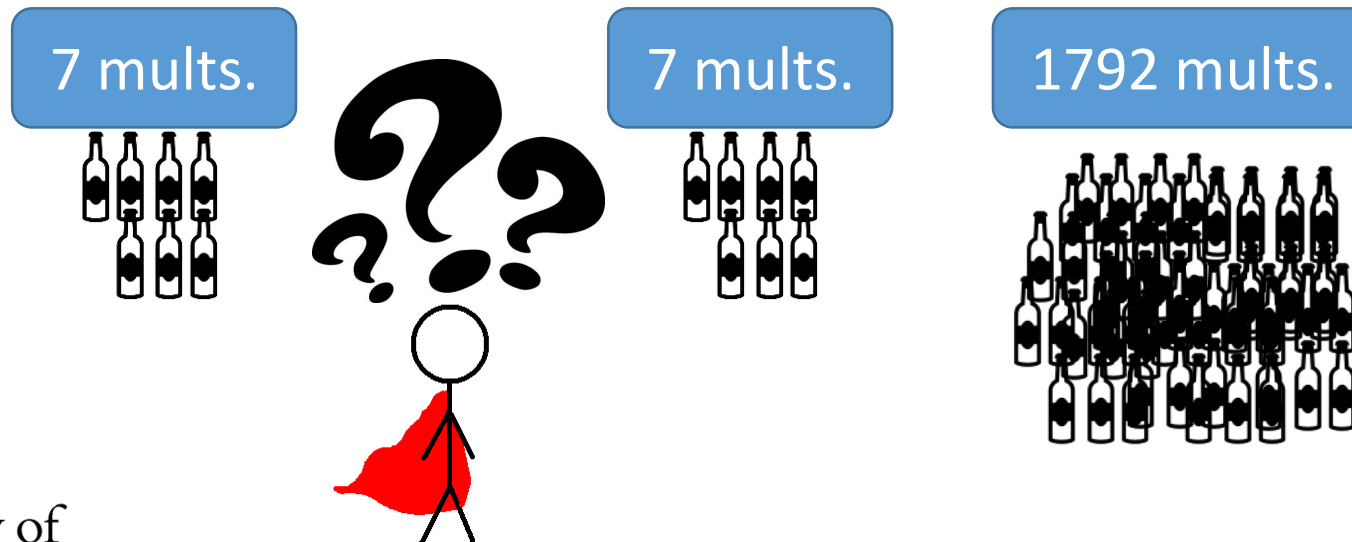


How to pre-Sbox

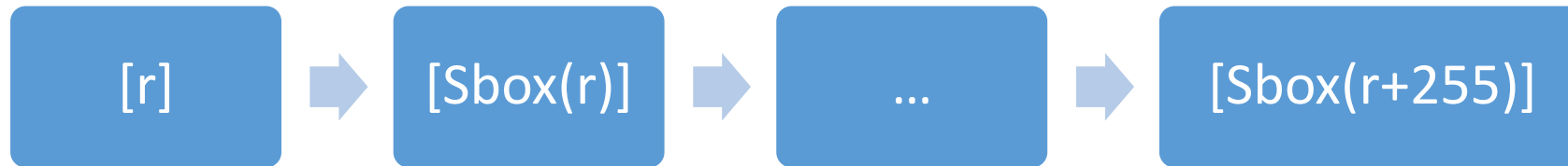


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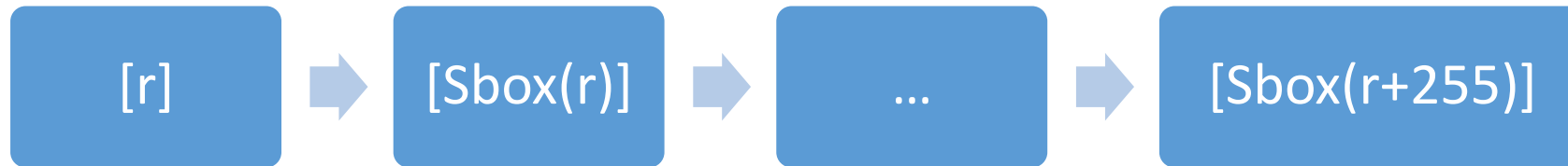


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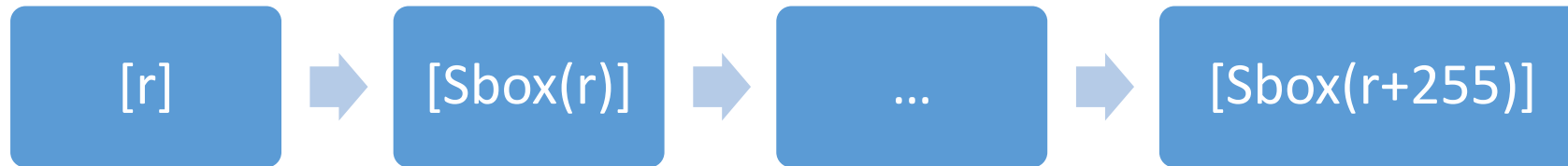
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How to pre-Sbox

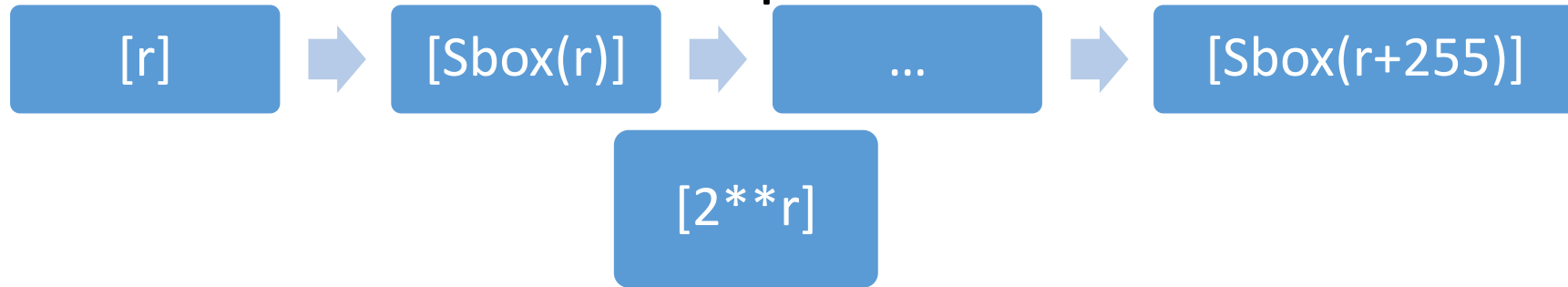


How to pre-Sbox

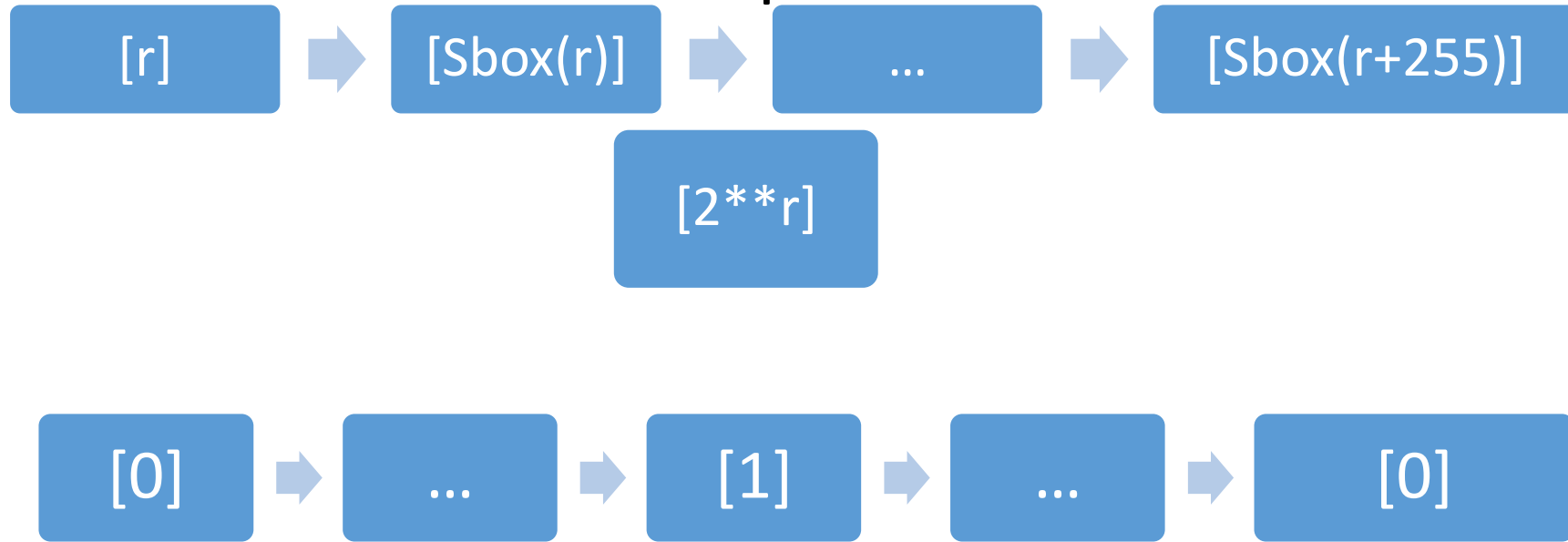


- Take random $[r]$.
- Compute $v = \text{Bits}(\text{pow}(2, [r]))$.
- Take every cyclic rotation of the Sbox row and compute $\langle v, \text{Sbox} \rangle$

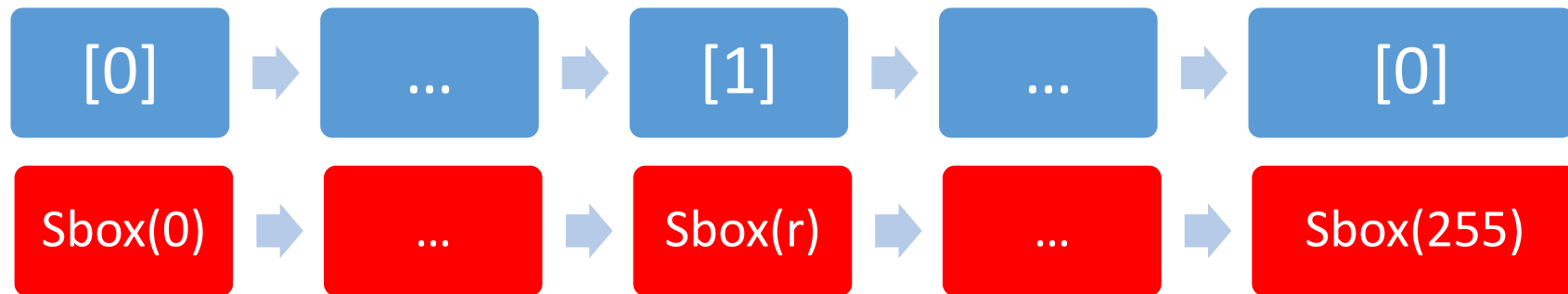
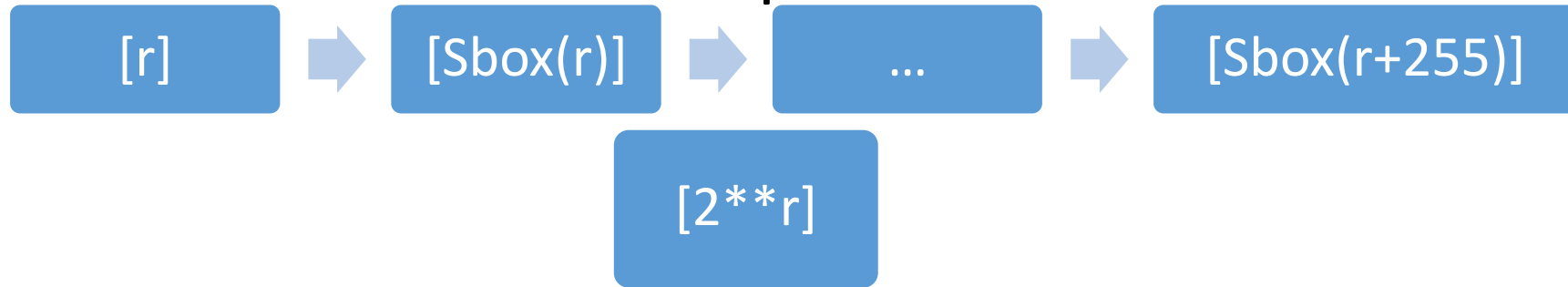
How to pre-Sbox



How to pre-Sbox



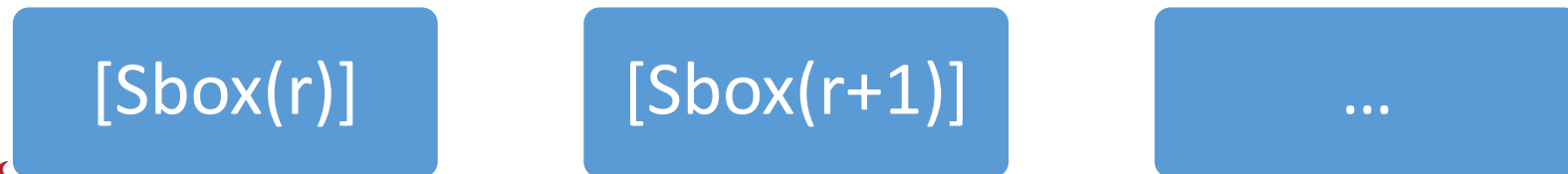
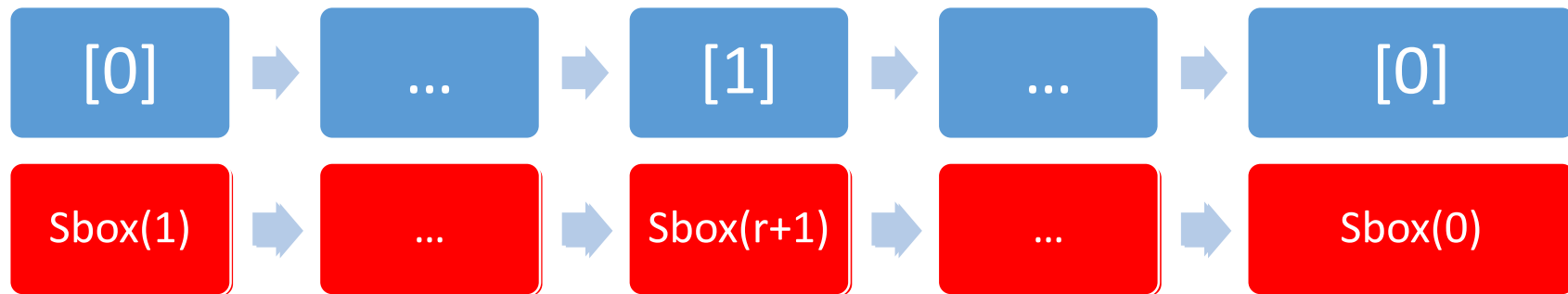
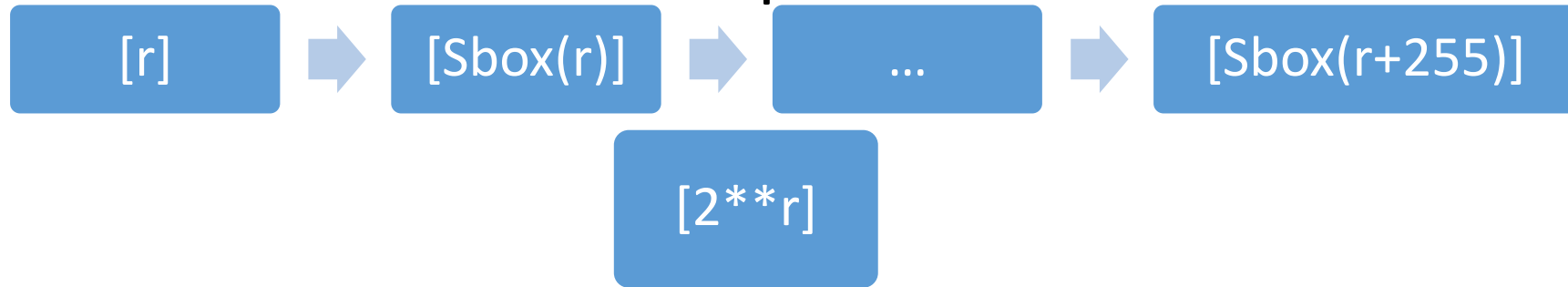
How to pre-Sbox



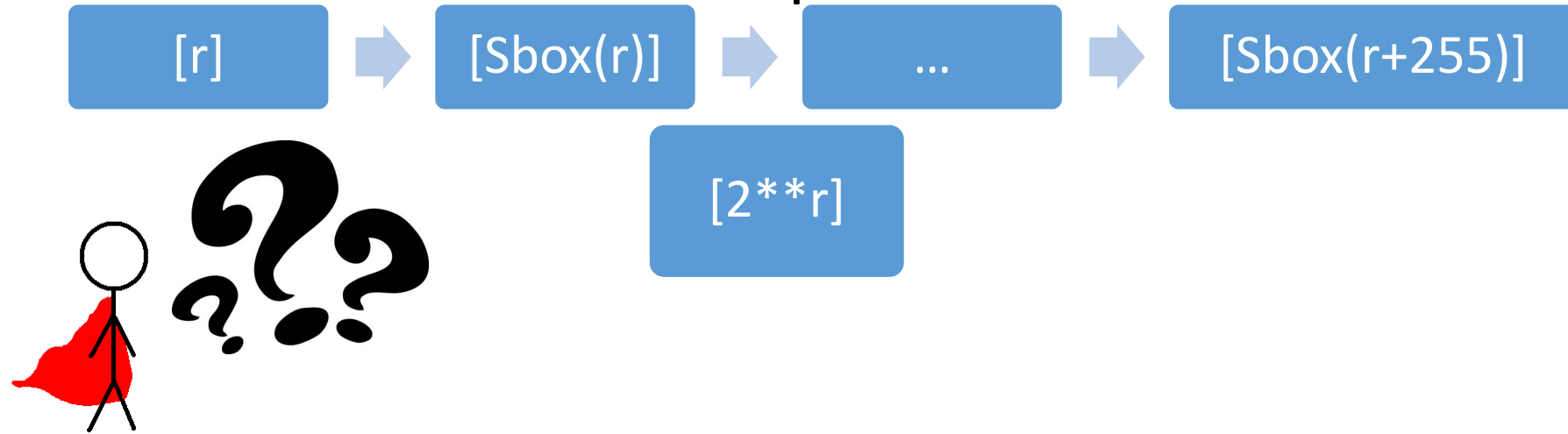
$[Sbox(r)]$



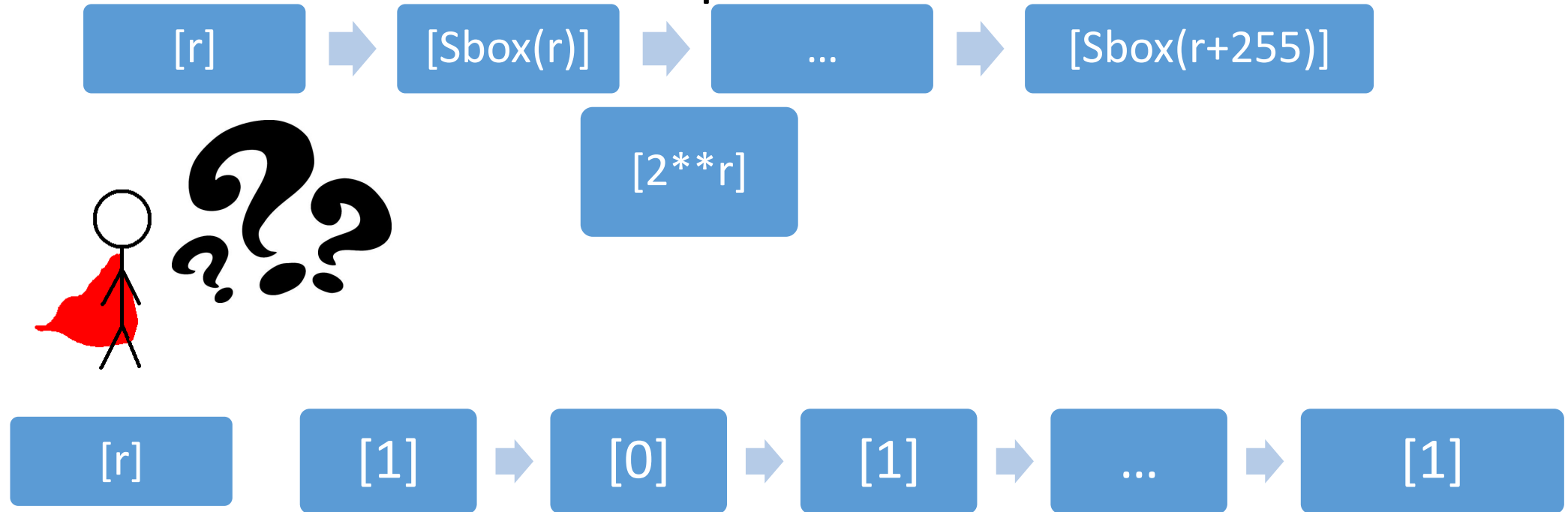
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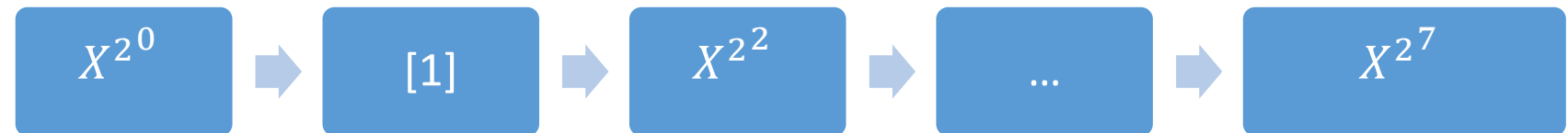
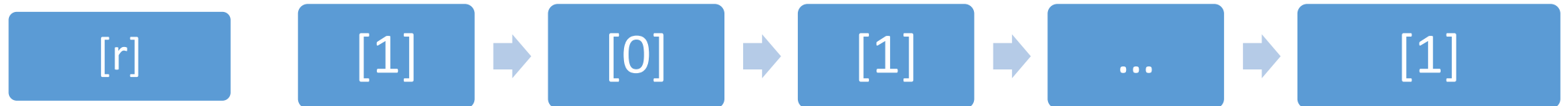
How to pre-Sbox



How to pre-Sbox



How to pre-Sbox



$$[X^r] = [2^r] \in GF(2^n)$$

How to pre-Sbox



$[2^{**}r]$



$[A] = [Z] \in GF(2^{256})$

How to pre-Sbox



[KOS16]

7 mults. in
 $GF(2^{256})$



How to pre-Sbox

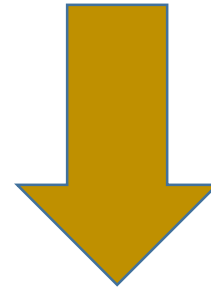


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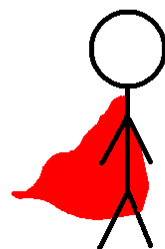


View ops. as polys in $GF(2^k)$



11 mults. in
 $GF(2^{40})$





TLDR;

N	$k = 1$	8	40	64	128
64	62	9	5	5	5
128	126	17	7	6	6
256	254	33	11	8	7
512	510	65	18	12	9
1024	1022	129	31	20	13

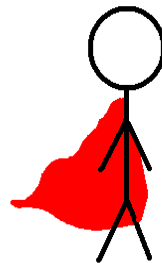
Table 1. Number of $\mathbb{F}_2 \times \mathbb{F}_{2^k}$ multiplications for creating a masked lookup table of size N , for varying k .

Side-Channel inspired techniques.

- Write $S_{\text{box}}(x)$ as a poly with minimal non-linear multiplications, ie squares are for free.
- AES Sbox requires 4 non-linear mults.
- DES Sbox requires 3 non-linear mults.

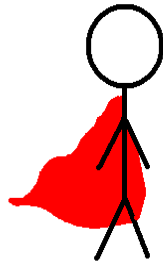
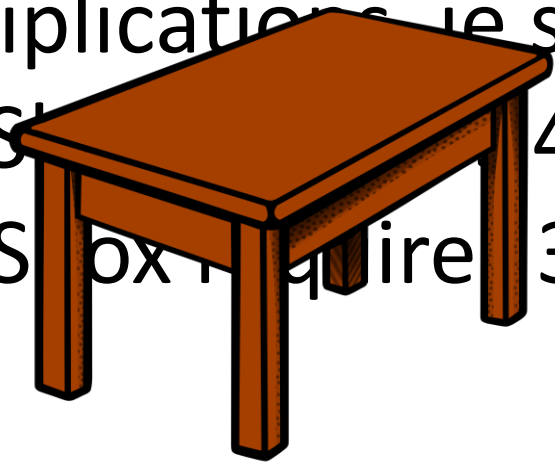
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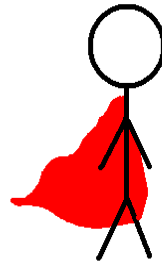
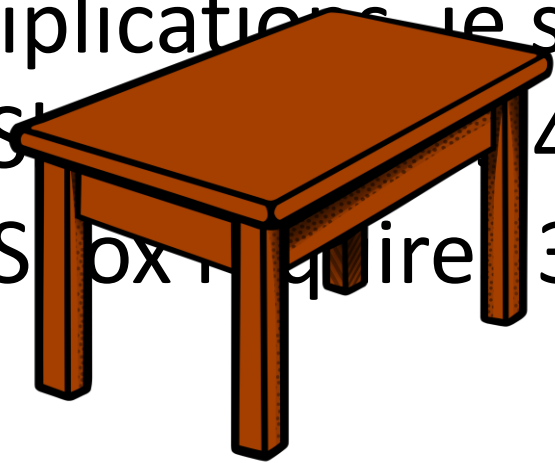
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Faster is...faster.

Protocol	Online		Comms. (total)	Notes
	Latency (ms)	Throughput (/s)		
TinyTable (binary) [DNNR16]	4.18	24500	3.07 MB	
TinyTable (optim.) [DNNR16]	1.02	339000	786.4 MB	
Wang et al. [WRK17]	0.93	1075	2.57 MB	10 Gbps
Rindal-Rosulek [RR16]	1.0	1000	1.6 MB	10 Gbps
OP-LUT [DKS ⁺ 17]	5	41670	0.103 MB	passive
SP-LUT [DKS ⁺ 17]	6	2208	0.044 MB	passive
AES-LT	0.93	236200	8.4 MB	
AES-RP	7.19	940	2.9 MB	

Table 6. Performance comparison with other 2-PC protocols for evaluating AES in a LAN setting.

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TinyTable (binary) [DNNR16]	4.18	24500	400kB	
TinyTable (optim.) [DNNR16]	1.02	339000	786.4 MB	
Wang et al. [WRK17]	0.93	1075	2.57 MB	10 Gbps
Rindal-Rosulek [RR16]	1.0	1000	1.6 MB	10 Gbps
OP-LUT [DKS ⁺ 17]	5	41670	0.103 MB	passive
SP-LUT [DKS ⁺ 17]	6	2208	0.044 MB	passive
AES-LT	0.93	236200	8.4 MB	
AES-RP	7.19	940	2.9 MB	

Table 6. Performance comparison with other 2-PC protocols for evaluating AES in a LAN setting.

Faster is...faster.

Protocol	Online		Comms. (total)	Notes
	Latency (ms)	Throughput (/s)		
TinyTable (binary) [DNNR16]	4.18	24500	3.07 MB	
TinyTable (optim.) [DNNR16]	1.02	339000	786.4 MB	
Wang et al. [WRK17]	0.93	1075	2.57 MB	10 Gbps
Rindal-Rosulek [RR16]	1.0	1000	1.6 MB	10 Gbps
OP-LUT [DKS ⁺ 17]	5	41670	0.103 MB	passive
SP-LUT [DKS ⁺ 17]	6	2208	0.044 MB	passive
AES-LT	0.93	236200	8.4 MB	
AES-RP	7.19	940	2.9 MB	

Table 6. Performance comparison with other 2-PC protocols for evaluating AES in a LAN setting.

Faster is...faster.

Protocol	Online		Comms. (total)	Notes
	Latency (ms)	Throughput (/s)		
TinyTable (binary) [DNNR16]	4.18	24500	3.07 MB	
TinyTable (optim.) [DNNR16]	1.02	339000	50MB	
Wang et al. [WRK17]	0.93	1075	2.57 MB	10 Gbps
Rindal-Rosulek [RR16]	1.0	1000	1.6 MB	10 Gbps
OP-LUT [DKS ⁺ 17]	5	41670	0.103 MB	passive
SP-LUT [DKS ⁺ 17]	6	2208	0.044 MB	passive
AES-LT	0.93	236200	8.4 MB	
AES-RP	7.19	940	2.9 MB	

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OP-LUT [DKS ⁺ 17]	5	41670	0.103 MB	passive
SP-LUT [DKS ⁺ 17]	6	2208	0.044 MB	passive
AES-LT	0.93	236200	8.4 MB	
AES-RP	7.19	940	2.9 MB	

Table 6. Performance comparison with other 2-PC protocols for evaluating AES in a LAN setting.

Thank you! #triples

Thank you! #triples



LAN results.

Cipher	Online (single-thread)			Online (multi-thread)			Preprocessing ^a
	Latency (ms)	Batch size	ops/s	Batch size	ops/s	Threads	ops/s
AES-BD	5.20	64	758	1024	3164	16	30.7
AES-RP	7.19	1024	940	64	3872	16	46.1
AES-LT	0.928	1024	51654	512	236191	32	16.79
3DES-Raw	270	512	130	-	-	-	1.24
3DES-PV	36.98	512	86	512	366	32	25.6
3DES-LT	4.254	1024	10883	512	45869	16	15.3

Table 3. 1 Gbps LAN timings for evaluating AES and 3DES in MPC.

#party #party #party