

Lecture 2

- Matrix form: $P^{(n)} = (P_{ij}^{(n)}) = P^n = P^2$
- By induction, $P^{(n)} = P^n$, i.e. to compute probabilities of n -step jumps, you can take n th powers of the transition matrix P .
- Convention: $P^{(0)} = I$
- Observe: $P_{ij}^{(m+n)} = P(X_{m+n}=j | X_0=i) = \sum_{k \in S} P(X_{m+n}=j, X_m=k | X_0=i) = \sum_{k \in S} P_{ik}^{(m)} P_{kj}^{(n)}$
matrix form: $P^{(m+n)} = P^{(m)} P^{(n)}$

Classification of States:

- Shorthand: write $P_i(\dots)$ for $P(\dots | X_0=i)$ And write $E_i(\dots)$ for $E(\dots | X_0=i)$
- Def'n: a state i of a Markov chain is recurrent (or, persistent) if $P_i(X_n=i \text{ for some } n \geq 1) = 1$. Otherwise, i is transient.
- Let $N(i) = \#\{n \geq 1 : X_n=i\}$ = total # times the chain hits i . (R.v.; could be ∞)
 - Recurrence theorem
 - i recurrent iff $\sum_{n=1}^{\infty} P_{ii}^{(n)} = \infty$ iff $P_i(N(i)=\infty) = 1$
 - And, i transient iff $\sum_{n=1}^{\infty} P_{ii}^{(n)} < \infty$ iff $P_i(N(i)=\infty) = 0$
 - To prove this let $f_{ij} = P_i(X_n=j \text{ for some } n \geq 1)$ prob of starting from i , and eventually hitting j
 - Then i recurrent iff $f_{ii} = 1$
 - And, i transient iff $f_{ii} < 1$
- Also $P_i(N(i) \geq 1) = f_{ii}$ and $P_i(N(i) \geq 2) = (f_{ii})^2$, etc.
 $P_i(N(i) \geq k) = (f_{ii})^k$

Also, recall that if Z is any non-negative integer-valued r.v. then $\sum_{k=1}^{\infty} P(Z \geq k) = E(Z)$

PROOF OF RECURRENCE THM: First, by continuity of probabilities

$$P_i(N(i)=\infty) = \lim_{k \rightarrow \infty} P_i(N(i) \geq k) = \lim_{k \rightarrow \infty} (f_{ii})^k = \begin{cases} 1 & : f_{ii} = 1 \\ 0 & : f_{ii} < 1 \end{cases}$$

Second, using countable linearity,

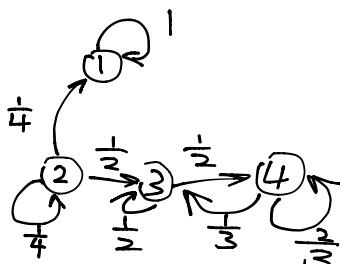
$$\begin{aligned} \sum_{n=1}^{\infty} P_{ii}^{(n)} &= \sum_{n=1}^{\infty} P_i(X_n=i) = \sum_{n=1}^{\infty} E_i(1_{X_n=i}) \\ &= E_i(\sum_{n=1}^{\infty} 1_{X_n=i}) = E_i(N(i)) = \sum_{k=1}^{\infty} P_i(N(i) \geq k) \\ &= \sum_{k=1}^{\infty} (f_{ii})^k = \begin{cases} \infty, & f_{ii} = 1 \\ \frac{f_{ii}}{1-f_{ii}} < \infty, & f_{ii} < 1 \end{cases} \end{aligned} \quad \text{in 347}$$

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Example:

$$S = \{1, 2, 3, 4\}$$

$$P = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1/4 & 1/4 & 1/2 & 0 \\ 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 1/3 & 2/3 \end{pmatrix}$$



$$\begin{aligned} f_{11} &= 1 \\ f_{22} &= \frac{1}{4} \quad \text{not } P_{22}! \\ f_{33} &= 1 \\ f_{44} &= 1 \quad f_{43} = 1 \\ f_{21} &= \frac{1}{3} \quad \text{conditional on leaving 2} \end{aligned}$$

$$\begin{aligned} f_{21} &= P_{21} + P_{22}f_{11} + \\ P_{22}f_{21} + P_{24}f_{41} &= \frac{1}{4} + \frac{1}{4}f_{11} + 0 + 0 \Rightarrow f_{21} = \frac{1}{3} \quad \text{so } f_{23} = \frac{2}{3} \\ f_{24} &= \frac{2}{3} \end{aligned}$$

• Example: Simple Random Walk (s.r.w.) ($S = \mathbb{Z}$, and $P_{i,i+1} = p$, and $P_{i,i-1} = 1-p$)

• Is the state 0 recurrent?

$$\text{If } n \text{ odd, } P_{00}^{(n)} = 0$$

$$\sum_{n=1}^{\infty} P_{00}^{(n)} \quad n \text{ even}$$

$$P_{00}^{(2)} = p(1-p) + (1-p)p = 2p(1-p)$$

$$P_{00}^{(4)} = pp(1-p)(1-p) + \dots = \binom{4}{2} p^2 (1-p)^2$$

$$P_{00}^{(n)} = P(\text{n/2 heads \& n/2 tails on first } n \text{ tosses}) = \binom{n}{n/2} p^{n/2} (1-p)^{n/2} = \frac{n!}{\left(\frac{n}{2}\right)!^2} p^{\frac{n}{2}} (1-p)^{\frac{n}{2}}$$

$$n! \approx \left(\frac{n}{e}\right)^n \sqrt{2\pi n} \quad \text{Stirling Approx}$$

So n large & even

$$P_{00}^{(n)} \approx \frac{\left(\frac{n}{e}\right)^{\frac{n}{2}} \sqrt{2\pi n}}{\left(\frac{n}{2}e\right)^{\frac{n}{2}} \sqrt{2\pi \frac{n}{2}}} p^{\frac{n}{2}} (1-p)^{\frac{n}{2}} = [4p(1-p)]^{\frac{n}{2}} \sqrt{\frac{n}{\pi}}$$

Now, if $p = \frac{1}{2}$, then $4p(1-p) = 1$, so $\sum_{n=1}^{\infty} P_{00}^{(n)} \approx \sum_{n=2,4,6,\dots} \sqrt{2/\pi n} = \infty$, so state 0 is recurrent.

But if $p \neq \frac{1}{2}$, then $4p(1-p) < 1$, so $\sum_{n=1}^{\infty} P_{00}^{(n)} \approx \sum_{n=2,4,6,\dots} [4p(1-p)]^{n/2} \sqrt{2/\pi n} < \infty$ so state 0 is transient

- Similarly true for all other states besides 0, too.

Communicating States:

• Say that state i communicates with state j , $i \rightarrow j$ if $f_{ij} > 0$, i.e. if it is possible to get from i to j i.e. $\exists m \geq 1$ with $P_{ij}^{(m)} > 0$.

• $i \leftrightarrow j$

• Say a MC is irreducible if $i \rightarrow j$ for all $i, j \in S$.

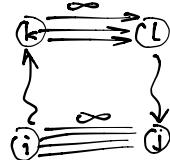
CASES THEOREM: For an irreducible MC, either

(a) $\sum_{n=1}^{\infty} P_{ij}^{(n)} = \infty$ for all $i, j \in S$. so all states are recurrent. "Recurrent MC"

(b) $\sum_{n=1}^{\infty} P_{ij}^{(n)} < \infty$ for all $i, j \in S$. so all states are transient. "transient MC"

• It follows from

SUM LEMMA: if $i \rightarrow k$, and $k \rightarrow l$, and $\sum_{n=1}^{\infty} P_{kl}^{(n)} = \infty$, then $\sum_{n=1}^{\infty} P_{il}^{(n)} = \infty$.



Proof of Cases Thm from Sum Lemma

Either (a) $\exists k, l \in S$ with $\sum_n P_{kl}^{(n)} = \infty$, then $\sum_n P_{il}^{(n)} = \infty, \forall i, j$

(b) or $\nexists k, l \in S$ with $\sum_n P_{kl}^{(n)} = \infty$, so $\sum_n P_{il}^{(n)} < \infty \forall i, j$

Proof of SUM LEMMA

Find $m, r \geq 1$ with $P_{ik}^{(m)} > 0, P_{kj}^{(r)} > 0$

Note that $P_{ij}^{(m+r)} \geq P_{ik}^{(m)} P_{kl}^{(s)} P_{lj}^{(r)}$. Hence

$$\sum_{n=1}^{\infty} P_{ij}^{(n)} \geq \sum_{n=m+r+1}^{\infty} P_{ij}^{(n)} = \sum_{s=1}^{\infty} P_{ij}^{(m+s+r)} \geq \sum_{s=1}^{\infty} P_{ik}^{(m)} P_{kl}^{(s)} P_{lj}^{(r)} = P_{ik}^{(m)} P_{kl}^{(r)} \sum_{s=1}^{\infty} P_{kl}^{(s)} = (+)(+)(\infty) = \infty$$

Simple random walk. Irreducible.

$p = \frac{1}{2}$ case (a)

$p \neq \frac{1}{2}$, case (b)

Finite SPACE THM: an irreducible MC on a finite state space always falls into case (a), i.e.

$\sum_{n=1}^{\infty} P_{ij}^{(n)} = \infty$ for all $i, j \in S$, and all states are recurrent

PROOF OF FINITE SPACE THM:

• Choose any state $i \in S$, then $\sum_{j \in S} \sum_{n=1}^{\infty} P_{ij}^{(n)} = \sum_{n=1}^{\infty} \sum_{j \in S} P_{ij}^{(n)} = \sum_{n=1}^{\infty} 1 = \infty$

Since S is finite, \exists at least one $j \in S$ with $\sum_{n=1}^{\infty} P_{ij}^{(n)} = \infty$, so must be case (a)

F-LEMMA: if $j \rightarrow i$ & $f_{jj} = 1$, then $f_{ij} = 1$

$\begin{array}{c} \uparrow \\ \text{possibly } j \text{ to } i \\ \downarrow \\ \text{eventually} \\ \text{in } j \end{array}$ so $i \xrightarrow{\text{eventually}} j$

PROOF OF F-LEMMA.

- Assume $i \neq j$ (otherwise trivial)

Let $T_i = \min \{n \geq 1 : X_n = i\}$ ($T_i = \infty$ if never hit i)

Since $j \rightarrow i$, $P_j(T_i < \infty) > 0$

In fact, must have $P_j(T_i < T_j) > 0$ (o.w. would always return to j before hitting i , so would never reach i from j)

But $1 - f_{jj} = P_j(T_j = \infty) \geq P_j(T_i < T_j)P_i(T_j = \infty) = P_j(T_i < T_j)(1 - f_{ij})$

If $f_{jj} = 1$, then $1 - f_{jj} = 0$, so must have $1 - f_{ij} = 0$, i.e. $f_{ij} = 1$.