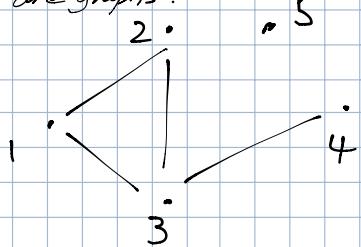


What are graphs?

2015-09-15



"Not a graph, but a representation of a graph"

$$V(G) = \{1, 2, 3, 4, 5\}$$

$$E(G) = \{\{1, 3\}, \{1, 2\}, \{2, 3\}, \{3, 4\}\}$$

Def: A graph G is an object consisting of

- a vertex set $V(G)$ which is finite and nonempty

- an edge set $E(G)$ which contains two element subsets of the vertex set

We denote $v = \#$ of vertices, $e = \#$ of edges

Def through example:

- the vertices 1, 2 are adjacent
2, 4 are not adjacent
- the edge $\{2, 3\}$ is incident to each 2 & 3
- 5 is an isolated vertex

Ex. Are these 2 graphs equal?



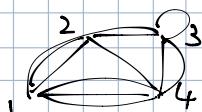
The only thing we cared about is $E(G)$ & $V(G)$

Def: Two graphs G, H are equal iff $V(G) = V(H)$ and $E(G) = E(H)$

Remark: there are 2 other more general notion of graphs:

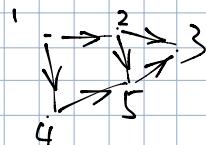
Multigraphs:

Directed graphs.



$$V = \{1, 2, 3, 4\}$$

$E = \{ \text{use tuples} \}$ b/c set cannot have repeated elements



Some important graphs

- The cyclic graph C_v $v \geq 3$

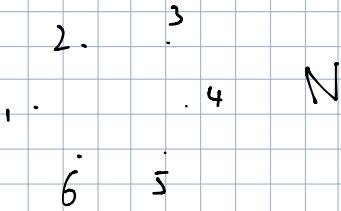
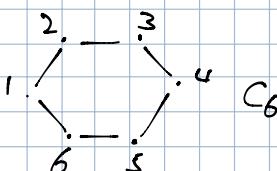
$$V(C_v) = \{1, 2, \dots, v\}$$

$$E(C_v) = \{\{1, 2\}, \dots, \{v-1, v\}, \{v, 1\}\}$$

- The null graph N_v

$$V(N_v) = \{1, \dots, v\}$$

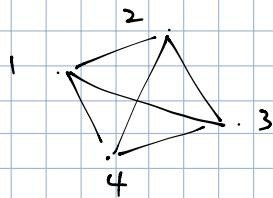
$$E(N_v) = \emptyset$$



The complete graph K_v

$$V(K_v) = \{1, \dots, v\}$$

$E(K_v)$ = {all possible edges}



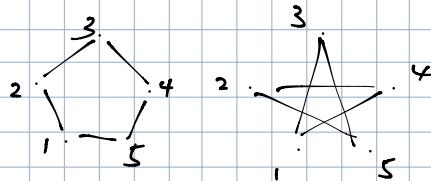
Thm: K_v has $\frac{v(v-1)}{2}$ edges

Proof: Each vertex is connected to $(v-1)$ vertices. Therefore there are $v(v-1)$ "edge-ends". Every edge has two ends. Therefore, divide by 2.

Def: The complement of a graph G , denoted by \overline{G} is the graph with $V(\overline{G}) = V(G)$

The edge set of \overline{G} consists of all two element subsets of $V(G)$ that are not in $E(G)$

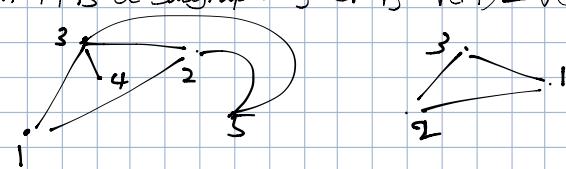
Ex: $\overline{K}_v = K_{v-1}$



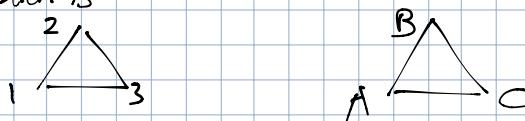
Corollary: $\overline{\overline{G}} = G$

Def: A graph H is a subgraph of G if $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$

Ex:



Question is



They are not equal, but are they "the same"

We'll use the definition of a bijective map

bijection

Def: Two graphs G & H are called isomorphic if there is a one-to-one map $V(G) \rightarrow V(H)$ such that two vertices are adjacent in G iff their corresponding vertices in H are adjacent

Ex: In the example above

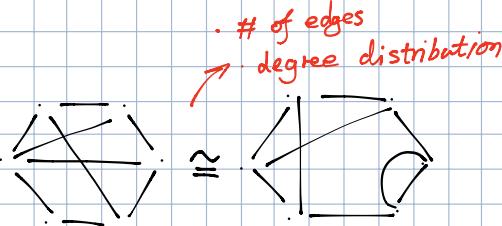
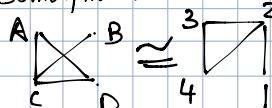
$$1 \rightarrow A$$

$$2 \rightarrow B$$

$$3 \rightarrow C$$

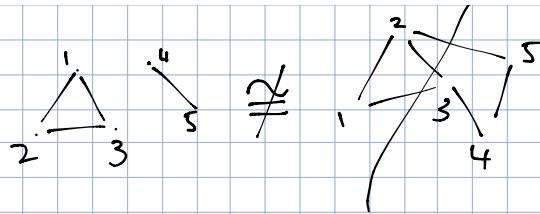
is such a map inducing an isomorphism

Are these graphs isomorphic?



Remark: A graph isomorphism preserves
(Thm)

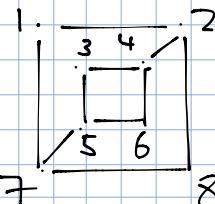
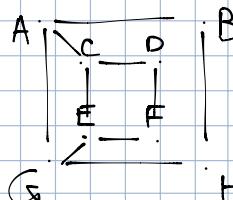
- number of vertices
- number of edges
- number of connected components
- "degree distribution"



Def. The degree of a vertex is the number of edges incident to it

Remark: Even if all conditions as above are met, the graphs might not be isomorphic

Ex:



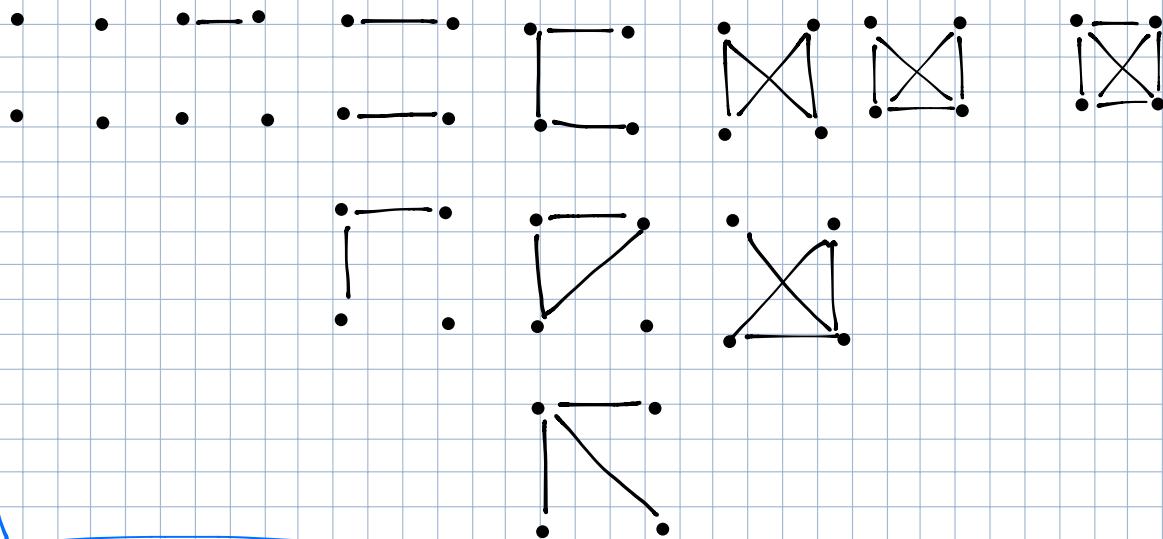
d	#
0	0
1	8
2	4
3	4

No vertices of dog 2 that are adjacent here

Remark: from now on, we'll distinguish graphs only up to isomorphism
This is why we will omit labels in most cases.

How many graphs are there with a given # of vertices? Starting from $v=4$

$e=0 \quad e=1 \quad e=2 \quad e=3 \quad e=4 \quad e=5 \quad e=6$



Remark: The # of graphs with a given # of vertices grows fast

of vertices | # of graphs

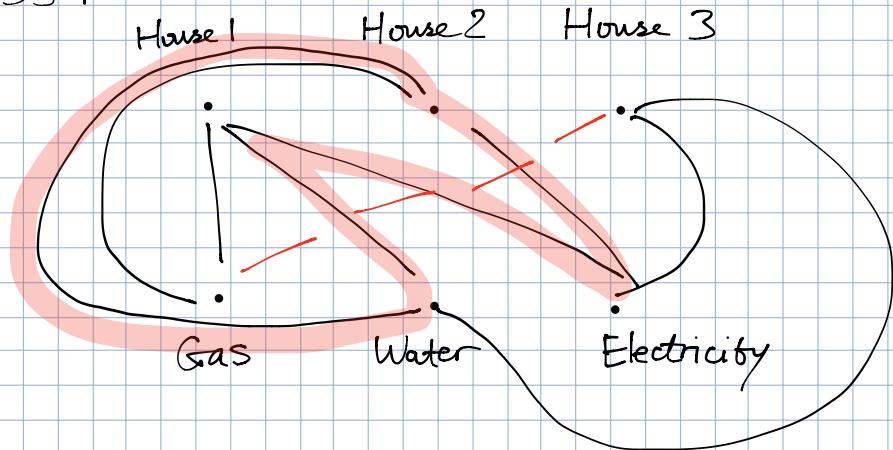
1	1
2	2
3	4
4	11
5	34

6	156
7	1044
8	12346
9	308708

Outlook: Planar graphs

Def: A graph is planar if it is isomorphic to a graph that can be drawn in the plane without edge crossings. Otherwise it is nonplanar.

Ex: The "utility graph"



The study of "planar" graphs involves an important seemingly trivial theorem

Thm (Jordan Curve Thm): If C is a continuous simple closed curve in a plane, then the curve C divides the rest of the plane into two regions.
↓(not self-intersecting)

If a point P in one of these regions is joined to a point Q in the other region by a cts. curve L in this plane, then L intersects C .

