

Feb 6th.

Recall: Euler's Theorem  
if  $\gcd(a, m) = 1$      $a, m > 1$   
then  $a^{\varphi(m)} \equiv 1 \pmod{m}$

Let  $b_1, b_2, \dots, b_{\varphi(m)}$  be all #'s relatively prime to  $m$ .

look at  $ab_1, ab_2, ab_3, \dots, ab_{\varphi(m)}$

①  $\gcd(a, m) = 1$   
 $\gcd(b_i, m) = 1$   
 $\Rightarrow a \cdot b_1, a \cdot b_2, \dots, a \cdot b_{\varphi(m)}$  are relatively prime to  $m$ .

② Claim:  
 $a \cdot b_i \not\equiv a \cdot b_j \pmod{m}$  if  $i \neq j$   
if  $a \cdot b_i \equiv a \cdot b_j \pmod{m} \Rightarrow m | a \cdot b_i - a \cdot b_j = a(b_i - b_j)$

last  $\gcd(m, a) = 1$

$\Rightarrow m | b_i - b_j$  but  $|b_i - b_j| < m$

$\Rightarrow b_i - b_j = 0 \Rightarrow b_i = b_j$

$\Rightarrow a \cdot b_1, a \cdot b_2, \dots, a \cdot b_{\varphi(m)}$  are distinct  $\pmod{m}$

③ Let  $a \cdot b_i \equiv m_i \pmod{m}$  divide  $a \cdot b_i$  by  $m$  with remainder  $0 \leq m_i < m$

$$\begin{array}{ccccccc} ab_1 & a \cdot b_2 & \cdots & ab_{\varphi(m)} \\ \parallel & \parallel & & \parallel \\ m_1 \pmod{m} & m_2 \pmod{m} & & m_{\varphi(m)} \pmod{m} \end{array}$$

④ Claim  $\gcd(m_i, m) = 1$   
In general if  $x \equiv y \pmod{m} \Rightarrow x = y + km$   
Then  $\gcd(x, m) = \gcd(y, m)$        $d | x \& d | m \Leftrightarrow d | y \& d | m$

in particular if  $\gcd(x, m) = 1 \Leftrightarrow \gcd(y, m)$   
 $\Rightarrow$  have  $m_1, \dots, m_{\varphi(m)}$  are all  $\leq m$ , distinct & relatively prime to  $m$

$\Rightarrow m_1, \dots, m_{\varphi(m)}$  are the same as  
 $b_1, \dots, b_{\varphi(m)}$  in the same order

$$\begin{array}{c} m_1, \dots, m_{\varphi(m)} \equiv b_1, \dots, b_{\varphi(m)} \pmod{m} \\ \parallel \\ (ab_1)(ab_2) \cdots (ab_{\varphi(m)}) \end{array}$$

$$\begin{aligned} a^{\varphi(m)} \cdot b_1 \cdots b_{\varphi(m)} &\equiv b_1 \cdots b_{\varphi(m)} \pmod{m} \\ \Rightarrow m | a^{\varphi(m)} b_1 \cdots b_{\varphi(m)} - b_1 \cdots b_{\varphi(m)} &= (a^{\varphi(m)} - 1) b_1 \cdots b_{\varphi(m)} \\ \text{but } \gcd(m_1, b_1, \dots, b_{\varphi(m)}) &= 1 \quad \Rightarrow m | a^{\varphi(m)} - 1 \Rightarrow a^{\varphi(m)} \equiv 1 \pmod{m} \end{aligned}$$

$$\text{Ex: } m=10 \quad \varphi(m) = \varphi(10) = \varphi(5 \cdot 2) = (5^1 - 5^0)(2^1 - 2^0) = (5-1)(2-1) = 4$$

If  $(a, m) = 1 \Rightarrow a^4 \equiv 1 \pmod{10}$

Ex:  $a=3 \quad 3^4 = 81 \equiv 1 \pmod{10}$   
 $a=7 \quad 7 \equiv -3 \pmod{10}$   
 $7^2 \equiv (-3)^2 \equiv 9 \pmod{10}$   
 $7^4 \equiv (-3)^4 \equiv 81 \equiv 1 \pmod{10}$

$$\text{Ex: Find } 5^{999} \pmod{9} = ?$$

$$\gcd(5, 9) = 1$$

$$\Rightarrow 5^{\varphi(9)} \equiv 1 \pmod{9}$$

$$*\varphi(9) = \varphi(3 \cdot 3) = 3^2 - 3^1 = 6$$

$$5^6 \equiv 1 \pmod{9}$$

$$*999 \equiv 996 + 3$$

$\overbrace{\text{L}}^{\text{divisible by 2 \& 3}}$

$$999 = 6k + 3 \text{ for some } k$$

$$5^{999} \equiv 5^{\frac{6k+3}{6k}} \equiv 5^{\frac{3}{6k}} \equiv 5^3 \pmod{9} = -1 \pmod{9}$$

$\overbrace{\text{L}}^{\text{1 mod 9}}$

\*  $6^{3^{101}} \pmod{22} = ?$

$$\varphi(22) = \varphi(2 \cdot 11) = (2-1)(11-1) = 10$$

$$\cancel{6^{10} \equiv 1 \pmod{22}}$$

b/c 6 & 22 aren't relatively prime  
 $\gcd(6, 22) \neq 1$

Let's first find  $6^{3^{101}} \pmod{11}$

$$\gcd(6, 11) = 1$$

$$\Rightarrow 6^{\varphi(11)} \equiv 1 \pmod{11}$$

$$6^{10} \equiv 1 \pmod{11} \quad * \quad \varphi(11) = 10$$

Need to find  $3^{101} \pmod{10}$ ,  $\gcd(3, 10) = 1$

$$3^{\varphi(10)} \equiv 1 \pmod{10}$$

$$\varphi(10) = \varphi(2 \cdot 5) = (2-1)(5-1) = 4$$

$$3^4 \equiv 1 \pmod{10}$$

$$101 \equiv 1 \pmod{4}$$

$$* 101 = 4 \cdot 25 + 1$$

$$3^{101} = 3^{100} \cdot 3 = (3^4)^{25} \cdot 3 \equiv 1 \cdot 3 \equiv 3 \pmod{10}$$

$$\Rightarrow 3^{101} = 10k + 3 \text{ for some } k$$

$$6^{3^{101}} = 6^{10k+3} = (6^{10})^k \cdot 6^3 \equiv 6^3 \pmod{11}$$

$$6^2 = 36 \equiv 3 \pmod{11}$$

$$\overbrace{36}^{33+3}$$

$$\Rightarrow 6^3 \equiv 36 \cdot 6 \equiv 3 \cdot 6 \equiv 18 \equiv 7 \pmod{11}$$

What about  $6^{3^{101}} \pmod{22}$ ?

$$(6^3)^{101} \equiv 7 \pmod{11} \Rightarrow 7 + 11 \equiv 18 \pmod{11}$$

$$\Rightarrow 6^{3^{101}} \equiv 18 \pmod{22}$$

$$* 11 \mid 6^{3^{101}} - 18 \text{ even}$$

$$2 \mid 6^{3^{101}} - 18 \text{ also even}$$

$$\Rightarrow 22 \mid 6^{3^{101}} - 18$$

## RSA Encryption

!!!

Lemma: Let  $N=p \cdot q$  where  $p$  and  $q$  are prime and  $p \neq q$  then  $a^{\frac{kp(a)}{N}+1} \equiv a \pmod{N}$

for any  $a$

Proof:  $\gcd(a, N) = 1$ ,  $p, q$  or  $p|q$

Case 1:  $\gcd(a, N) = 1 \Rightarrow a^{\varphi(N)} \equiv 1 \pmod{N}$  By Euler's Thm  
 $\Rightarrow a^{\frac{kp(a)}{N}} = a^{\frac{kp(a)}{p-1}(q-1)} \equiv 1 \pmod{N}$   
 $a^{\frac{kp(a)}{N}+1} \equiv a \cdot 1 \equiv a \pmod{N}$

Case 2:  $\gcd(a, N) = p \cdot q = N \therefore N/a$   
 $\Rightarrow a \equiv 0 \pmod{N} \quad a^{\frac{kp(a)}{N}+1} \equiv 0 \pmod{N}$   
 $a^{\frac{kp(a)}{N}+1} \equiv 0 \equiv a \pmod{N}$

Case 3:  $\gcd(a, N) = p$   
 $\Rightarrow \tilde{a}^{\varphi(q)} \equiv 1 \pmod{q} \quad \begin{cases} \gcd(\tilde{a}, q) = 1 \\ \tilde{a}^{q-1} \equiv 1 \pmod{q} \end{cases} \quad \begin{cases} \varphi(q) = q-1 \\ N = pq \\ \varphi(p) = p-1 \end{cases} \quad \tilde{a} \text{ is not divisible by } q$   
 $\Rightarrow (\tilde{a}^{q-1})^{p-1} \equiv 1 \pmod{q}$   
 $\tilde{a}^{\varphi(N)} \equiv 1 \pmod{q}$

$$a = \tilde{a} \cdot p$$

$$\underline{a^{\varphi(N)} = \tilde{a}^{\varphi(N)} p^{\varphi(N)} =}$$

$\gcd(a, N) = p$   $a$  is divisible by  $p$  but not by  $q$   
 $\gcd(a, q) = 1$

$$\begin{aligned} a^{\varphi(q)} &\equiv 1 \pmod{q} \\ a^{q-1} &\equiv 1 \pmod{q} \end{aligned} \Rightarrow (a^{q-1})^{p-1} \equiv 1^{p-1} \pmod{q}$$

\*  $q \mid a^{\frac{kp(a)}{N}+1} - a$   $\pmod{q}$  since  $\gcd(a, N) = a$

$$p/a^{\frac{kp(a)}{N}+1} \Rightarrow p/a^{\frac{kp(a)}{N}+1} - a \Rightarrow pq/a^{\frac{kp(a)}{N}+1} - a$$

$$q/a^{\frac{kp(a)}{N}+1} - a$$

$$\Rightarrow k^{\frac{kp(a)}{N}+1} \equiv a \pmod{pq}$$

$$\begin{aligned} a^{\frac{(p-1)(q-1)-kp(a)}{N}} &\equiv 1 \pmod{q} \\ a^{\varphi(N)} &\equiv 1 \pmod{q} \end{aligned}$$

$$\begin{aligned} a^{\frac{kp(a)}{N}+1} &\equiv 1 \pmod{q} \\ a^{\frac{kp(a)}{N}+1} &\equiv 1 \cdot a \equiv a \pmod{q} \end{aligned}$$

## How RSA Encryption Works?

Recevier picks 2 very large prime #'s  $p$  and  $q$

Compute  $N = p \cdot q$  and  $\varphi(N) = (p-1)(q-1)$

Receiver also picks an encoder which is a number  $E$  relatively prime to  $\varphi(N)$

$$\gcd(E, \varphi(N)) = 1$$

$\Rightarrow$  Receiver broadcasts the pair  $(N, E)$   
Sender waits to send secret message

$$0 < M < N$$

Sender takes  $M$  and computes  $M^E \bmod N$   
 $0 \leq R < N$

The Receiver gets  $R$  and the sender wants to recover  $M$   
 $E$  is relatively prime to  $\varphi(N) \Rightarrow$  there exists  $D$ -decoder such that  $ED \equiv 1 \pmod{\varphi(N)}$

$$\begin{aligned} \gcd(a, m) = 1 \Rightarrow a^{(p-1)m} &\equiv 1 \pmod{m} \Rightarrow a \cdot a^{(p-1)m-1} \equiv 1 \pmod{m} \\ \text{if } \gcd(a, m) = 1 \Rightarrow a^{(p-1)m} &\equiv 1 \pmod{m} \Rightarrow a \cdot a^{(p-1)m-1} \equiv 1 \pmod{m} \end{aligned}$$

Receiver computes

$$R^D \bmod N$$

Claim:  $R^D \equiv M \pmod{N} \rightarrow$  Why is this so?  
 $R = M^E \pmod{N}$

$$R^D = (M^E)^D = M^{ED} = M^{1+k\varphi(N)} \equiv M \pmod{N}$$

$\hookrightarrow$  by Lemma.

The spy hears  $(N, E)$  that the receiver broadcast and hears  $R$ .  
Why can't the spy recover  $M$ ?

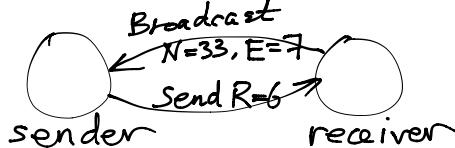
To recover  $M$ , one needs  $D$  s.t.  $DE \equiv 1 \pmod{\varphi(N)}$

if you know  $D \Rightarrow$  take  $R^D \bmod N \equiv M$

The problem is that the spy can't compute  $\varphi(N) = (p-1)(q-1)$   
The spy can't factor  $N$  or  $P, Q$

Ex)  $p=3, q=11, N=p \cdot q = 3 \cdot 11 = 33$   
 $\varphi(N) = (3-1)(11-1) = 2 \cdot 10 = 20$

Take  $E=7, \gcd(7, 20)=1$



What was the secret message  $M$  ( $M < 33$ )

Receiver needs  $D$  such that

$$D \cdot E \equiv 1 \pmod{\varphi(N)}$$

$$\begin{array}{c|c} \parallel & \parallel \\ 7 & 20 \end{array}$$

$$7D \equiv 1 \pmod{20}$$

D=3 works ✓

$$R^D = 6^3 \bmod 33 \equiv 18 \bmod 33$$

$$6^2 = 36 \equiv 3 \bmod 33$$

$$6^3 \equiv 3 \cdot 6 \equiv 18 \bmod 33$$

$N=18$  - Original Secret Message

Check  $M=18$  take  $\frac{M^E}{R} \bmod N$

$$18^7 \bmod 33 = 6 \bmod 33$$

Check