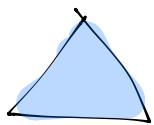


Today - §1.3 (end) and §1.4

Def'n: A subset $S \subseteq \mathbb{R}^n$ is convex provided, for any $x_1 \in S$ and $x_2 \in S$, the line segment joining x_1 and x_2 is a subset of S .

Remark: This means, for any $x_1 \in S$ and $x_2 \in S$ and $\lambda \in [0, 1]$, we have $x_1 + \lambda(x_2 - x_1) \in S$

Eg. In \mathbb{R}^2



is convex



is not convex

(because line isn't completely inside the slope.)

See page 79.

Theorem 1.1 A closed half-space is convex.

Proof: Let S be a closed half-space in \mathbb{R}^n .

There exist $c \in \mathbb{R}^n$, $c \neq 0 \in \mathbb{R}^n$ and a real number k such that $x \in S$ iff $c^T x \leq k$.

Now fix $x_1 \in S$ and $x_2 \in S$, and λ in $[0, 1]$ to show $x_1 + \lambda(x_2 - x_1) = (1-\lambda)x_1 + \lambda x_2 \in S$, that is, $c^T((1-\lambda)x_1 + \lambda x_2) \leq k$

$$\text{But } c^T((1-\lambda)x_1 + \lambda x_2) = \underbrace{(1-\lambda)c^T x_1}_{\geq 0} + \underbrace{\lambda c^T x_2}_{\geq 0} \leq (1-\lambda)k + \lambda k = k$$

Theorems 1.3 The intersection of finitely many convex sets is convex.

Proof: Let S_1, \dots, S_m be convex subsets of \mathbb{R}^n , to show $S_1 \cap \dots \cap S_m$ is convex. Pick x_1 and x_2 in $S_1 \cap \dots \cap S_m$ and pick $\lambda \in [0, 1]$. For any $i=1, \dots, m$, $x_1 + \lambda(x_2 - x_1) \in S_i$ (S_i is convex) Thus $x_1 + \lambda(x_2 - x_1) \in S_1 \cap \dots \cap S_m$.

Theorem 1.2 A hyperplane is convex.

Proof: Any hyperplane is the intersection of 2 half-spaces.

(The hyperplane given by $c^T x = k$ is the intersection of the half-spaces given by $c^T x \leq k$ and $c^T x \geq k$.)

Eg. The empty set is convex.

Theorem 1.4 (generalized)

The feasible region of a linear programming problem is convex.

Proof: The feasible region is the intersection of finitely many closed half-spaces.

Definition: A set S in \mathbb{R}^n is bounded provided there exists M so large that for any $x = [x_1 \ x_2 \ \dots \ x_n]^T$ we have $-M \leq x_i \leq M$ for any $i=1, \dots, n$.

Remark: If a problem is unbounded (has no optimal solution, although the feasible region is non-empty), its feasible region is unbounded. The converse of this is false: certain problems are unbounded (have an optional solution) even though they have unbounded feasible regions.

§ 1.4

Definition: Let S be a convex subset of \mathbb{R}^n and let $u \in S$. u is an extreme point of S provided for any $x_1 \in S$ and $x_2 \in S$ and any $\lambda \in [0, 1]$ with $(1-\lambda)x_1 + \lambda x_2 = u$ we have either $u = x_1$ or $u = x_2$.

Remark: This means : if u lies on a line segment which is entirely in S , then u is an endpoint of the line segment.

Eg: In \mathbb{R}^2 , let S be the solution set of $x \geq 0, y \geq 0$

Suppose $\begin{bmatrix} x_1 \\ y_1 \end{bmatrix}$ and $\begin{bmatrix} x_2 \\ y_2 \end{bmatrix} \in S$ and suppose $\lambda \in [0, 1]$ so that

$$(1-\lambda) \begin{bmatrix} x_1 \\ y_1 \end{bmatrix} + \lambda \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

We may assume $0 < \lambda < 1$

$$\begin{matrix} >0 & >0 \\ (1-\lambda)x_1 + \lambda x_2 = 0 \end{matrix}$$

$$(1-\lambda)y_1 + \lambda y_2 = 0$$

$$\begin{matrix} >0 & >0 \end{matrix}$$

$$So \quad x_1 = 0, x_2 = 0, y_1 = 0, y_2 = 0.$$

