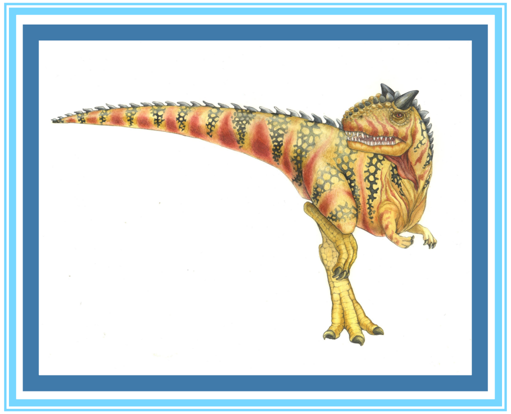
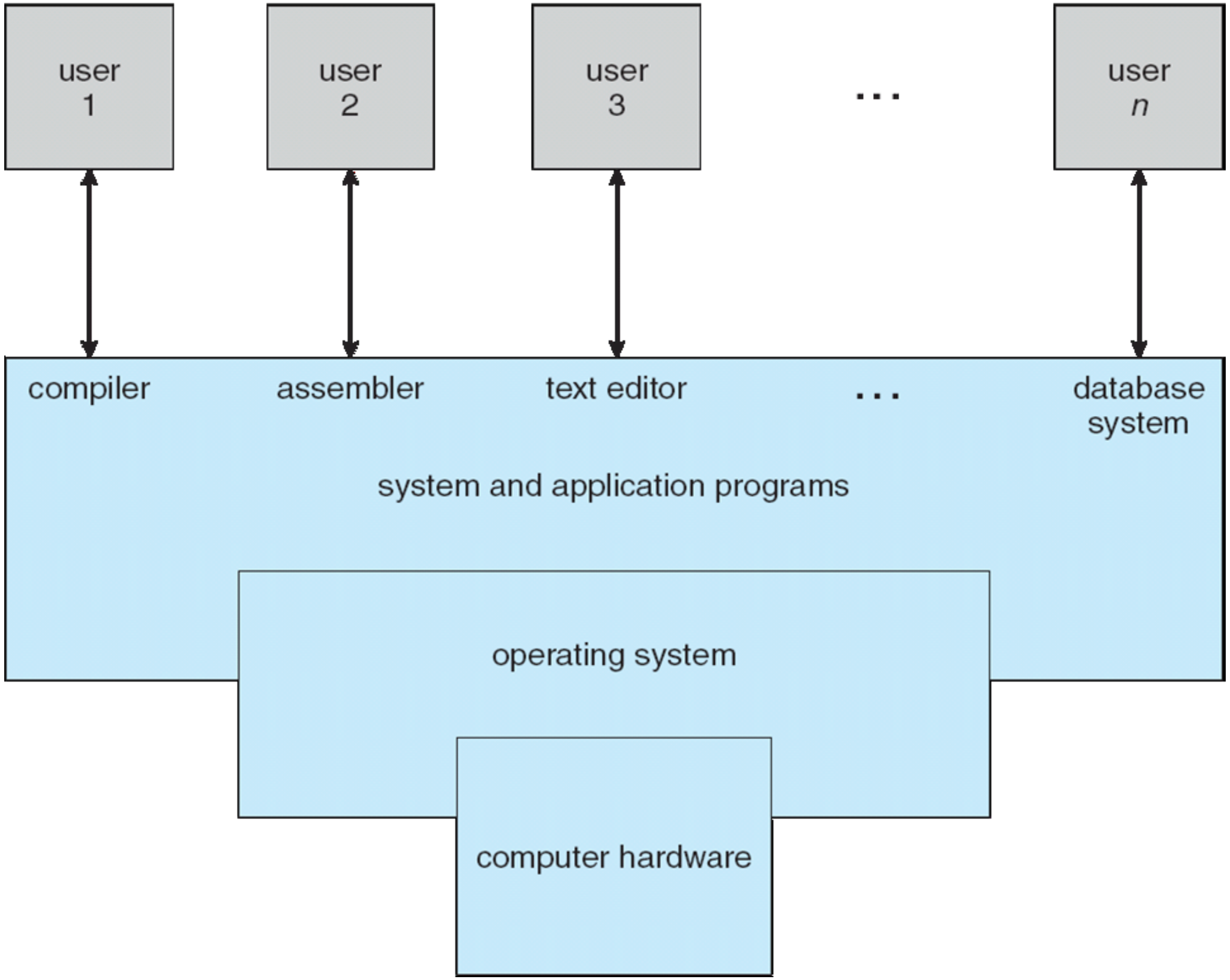
Some Operating Systems Concepts



What is an Operating System?

- A program that acts as an intermediary between a user of a computer and the computer hardware
- Operating system goals:
  - Execute user programs and make solving user problems easier
  - Make the computer system convenient to use
  - Use the computer hardware in an efficient manner

Four Components of a Computer System



Accessing Operating System Services

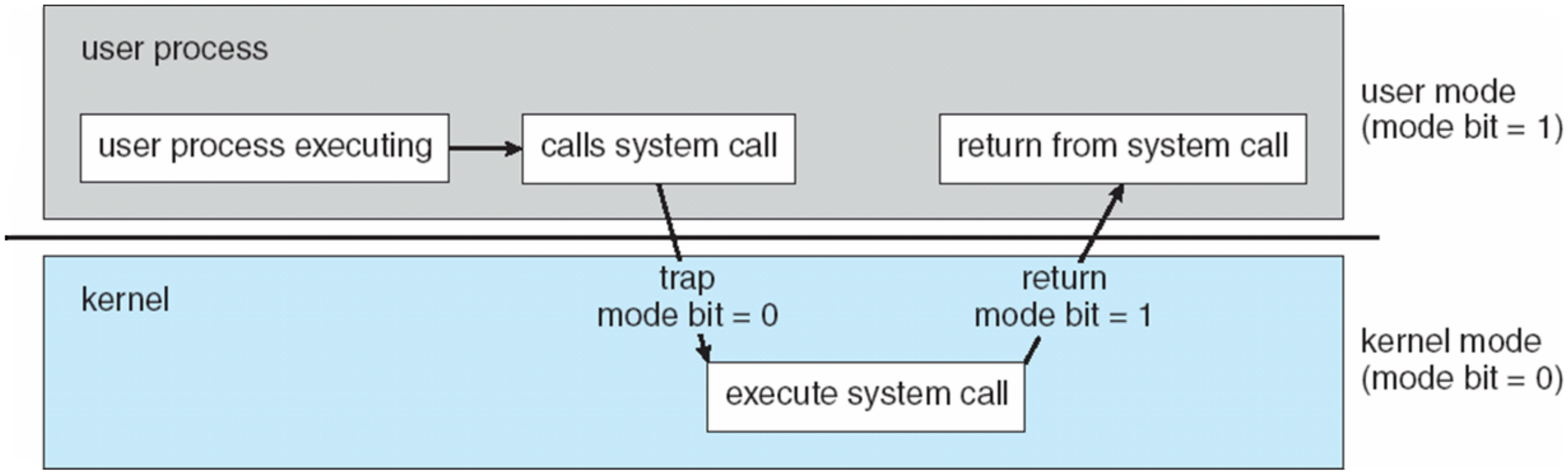
- **Interrupt driven** (hardware and software)
  - Hardware interrupt by one of the devices
  - Software interrupt (**exception** / **trap / syscall):**
    - Software error (e.g., division by zero)
    - Request for operating system service

Dual-mode of Operation

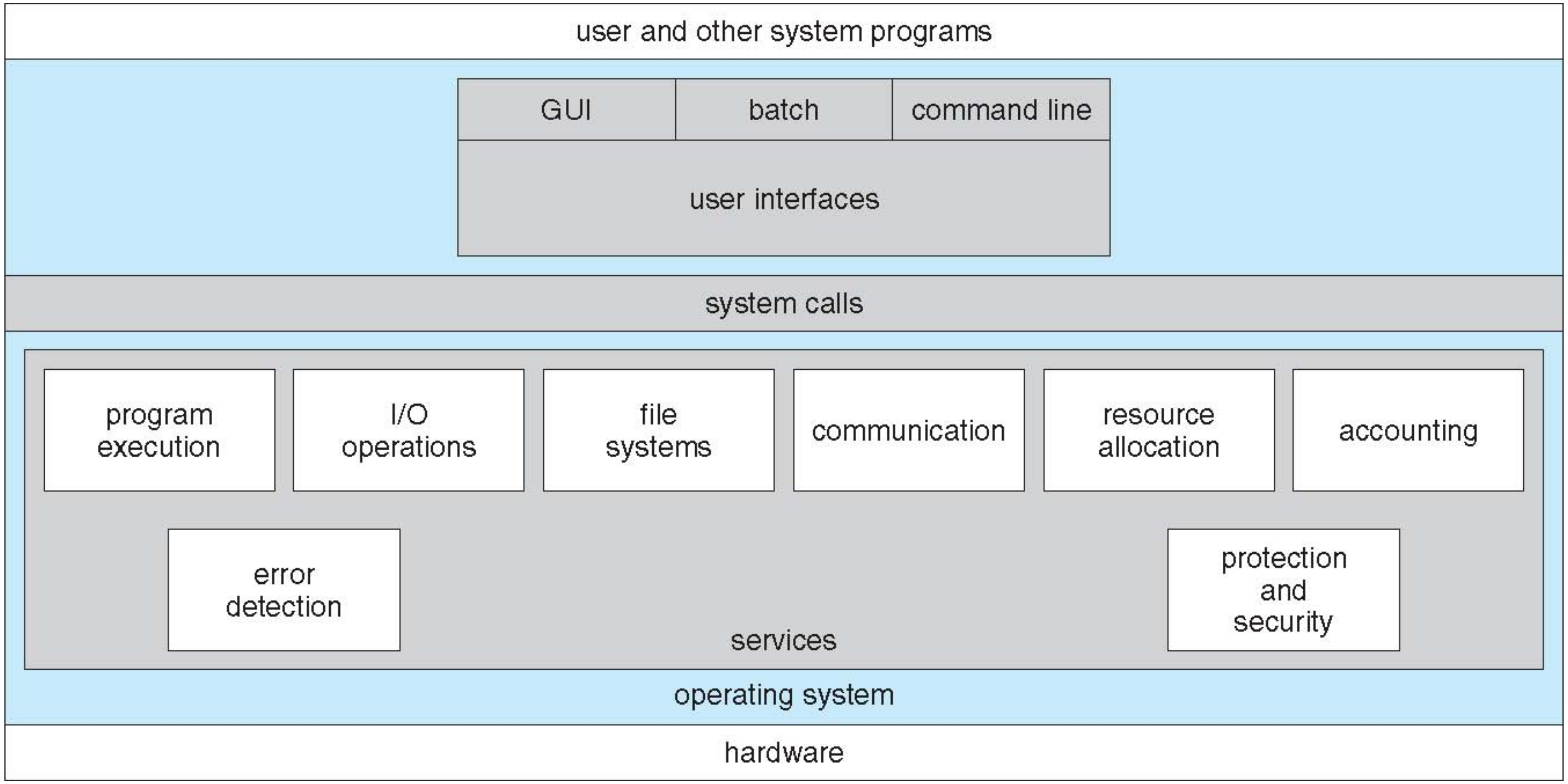
- **Dual-mode** operation allows OS to protect itself and other system components
  - **User mode** and **kernel mode**
  - **Mode bit** provided by hardware
    - Provides ability to distinguish when system is running user code or kernel code
    - Some instructions designated as **privileged**, only executable in kernel mode
    - System call changes mode to kernel, return from call resets it to user

Transition from User to Kernel Mode

- Timer to prevent infinite loop / process hogging resources
  - Timer is set to interrupt the computer after some time period
  - Keep a counter that is decremented by the physical clock.
  - Operating system set the counter (privileged instruction)
  - When counter zero generate an interrupt
  - Set up before scheduling process to regain control or terminate program that exceeds allotted time



A View of Operating System Services



System Calls

- Programming interface to the services provided by the OS
- Typically written in a high-level language (C or C++)
- Mostly accessed by programs via a high-level **Application Programming Interface** **(API)** rather than direct system call use
- Three most common APIs:
  - Win32 API for Windows
  - POSIX API for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X)
  - Java API for the Java virtual machine (JVM)

Example of Standard API

EXAMPLE OF STANDARD API

As an example of a standard API, consider the `read()` function that is available in UNIX and Linux systems. The API for this function is obtained from the `man` page by invoking the command

```
man read
```

on the command line. A description of this API appears below:

```
#include <unistd.h>

ssize_t    read(int fd, void *buf, size_t count)
```

return value	function name	parameters
`ssize_t`	`read`	`(int fd, void *buf, size_t count)`

A program that uses the `read()` function must include the `unistd.h` header file, as this file defines the `ssize_t` and `size_t` data types (among other things). The parameters passed to `read()` are as follows:

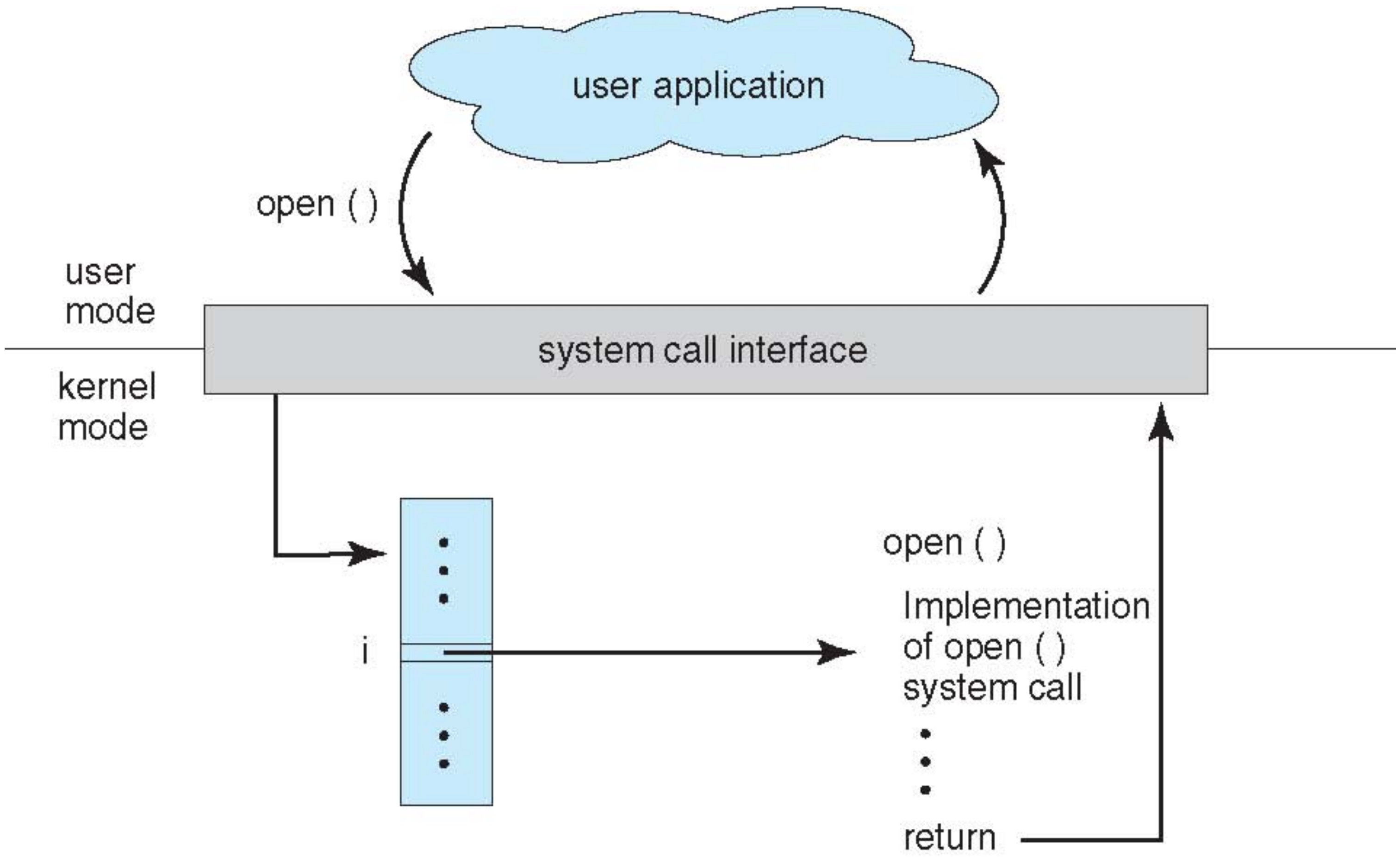
- `int fd`—the file descriptor to be read
- `void *buf`—a buffer where the data will be read into
- `size_t count`—the maximum number of bytes to be read into the buffer

On a successful read, the number of bytes read is returned. A return value of 0 indicates end of file. If an error occurs, `read()` returns –1.

System Call Implementation

- Typically, a number associated with each system call
  - **System-call interface** maintains a table indexed according to these numbers
- The system call interface invokes the intended system call in OS kernel and returns status of the system call and any return values
- The caller need know nothing about how the system call is implemented

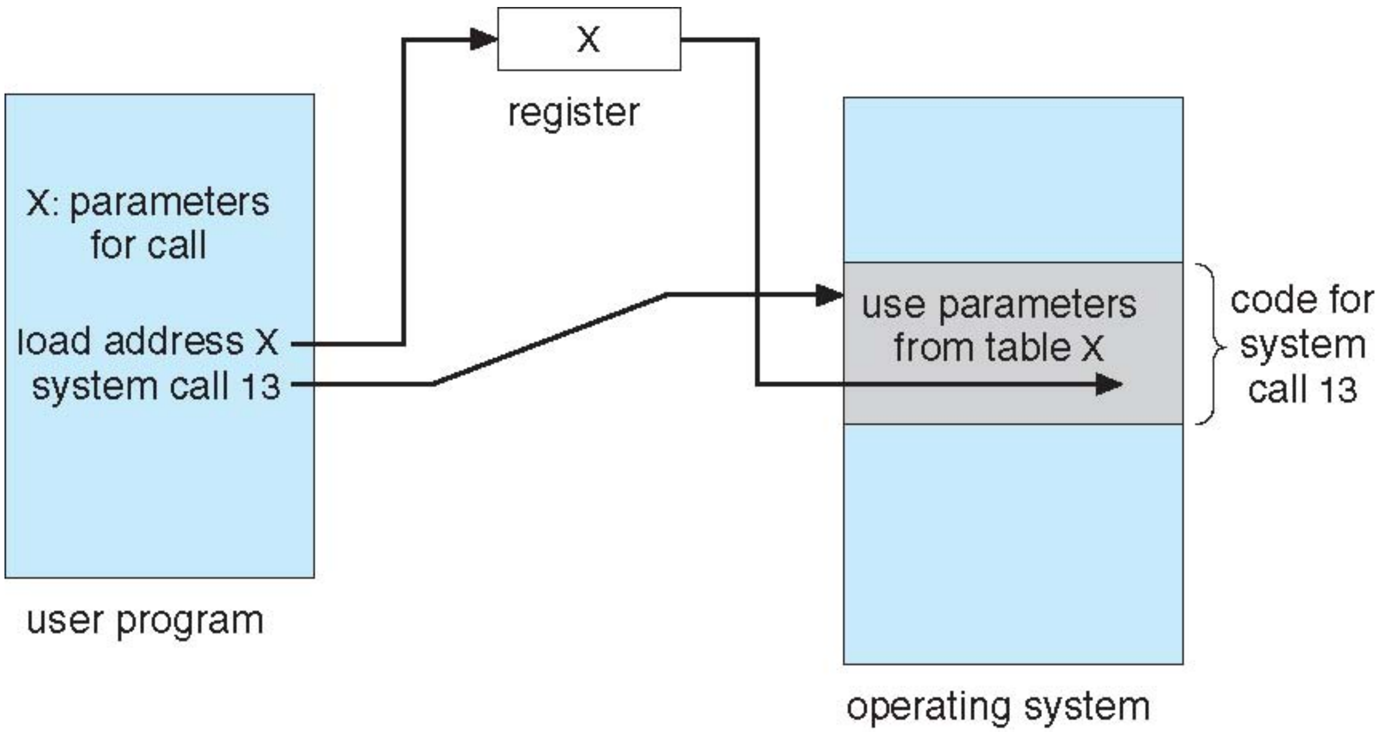
API – System Call – OS Relationship



System Call Parameter Passing

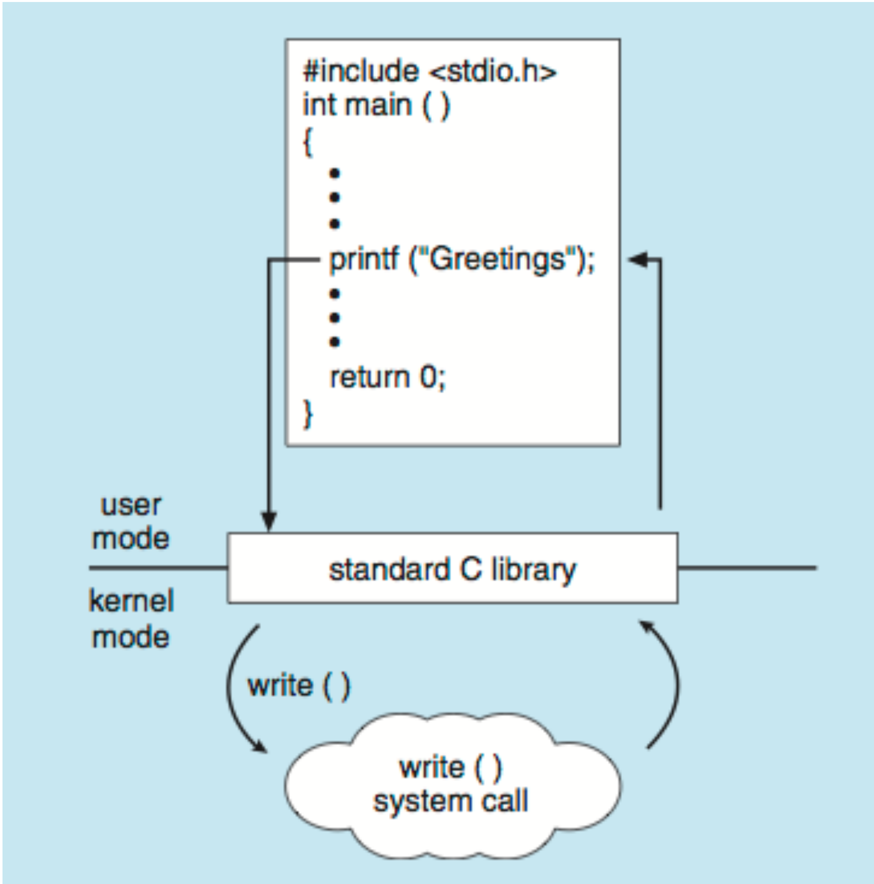
- Often, more information is required than simply identity of desired system call
- Three general methods used to pass parameters to the OS
  - Simplest: pass the parameters in registers
  - Parameters stored in a block, or table, in memory, and address of block passed as a parameter in a register
  - Parameters placed, or **pushed**, onto the **stack** by the program and **popped** off the stack by the operating system

Parameter Passing via Table

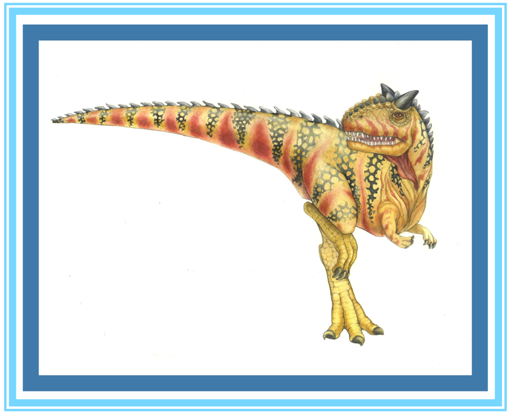


Standard C Library Example

- C program invoking printf() library call, which calls write() system call



Memory Management

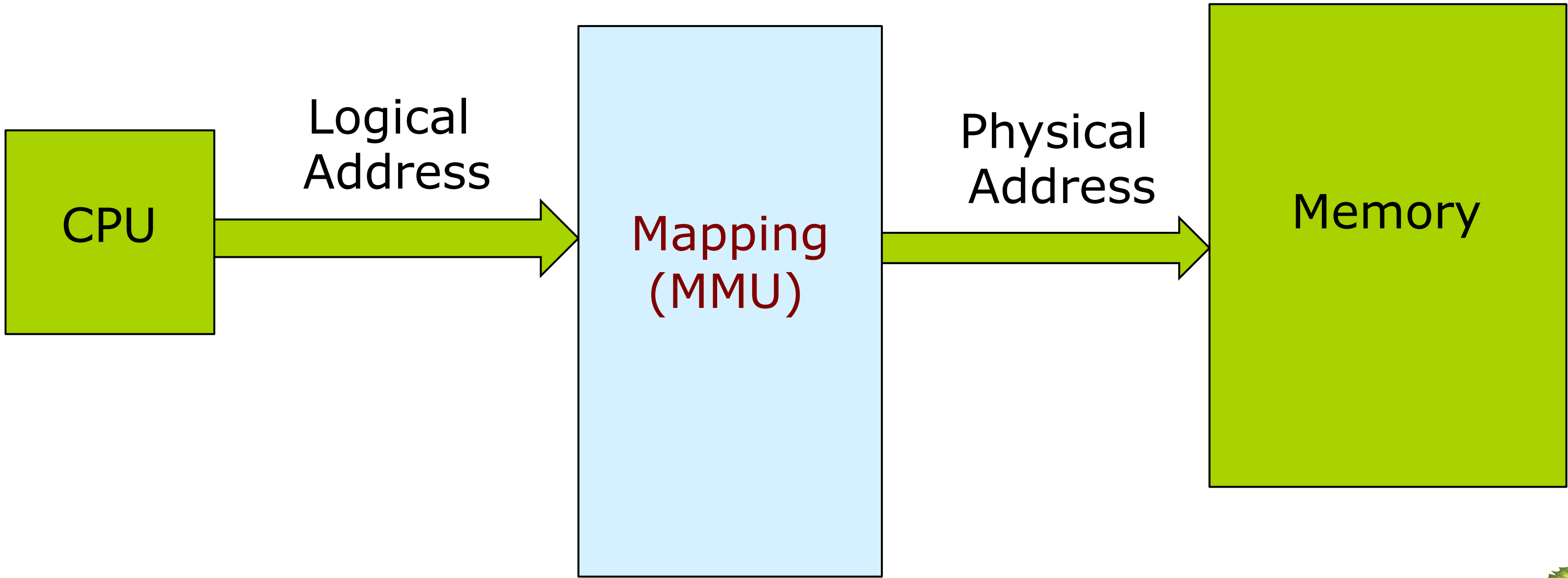


An Important Responsibility of OS

- How to allocate memory space to user programs?
  - A program in execution is called a process
- How to ensure memory protection at run time?
  - A program should not be able to overwrite other programs.
- Virtual memory is an important concept that is universally used in computer systems
  - Based on the principle of locality of reference

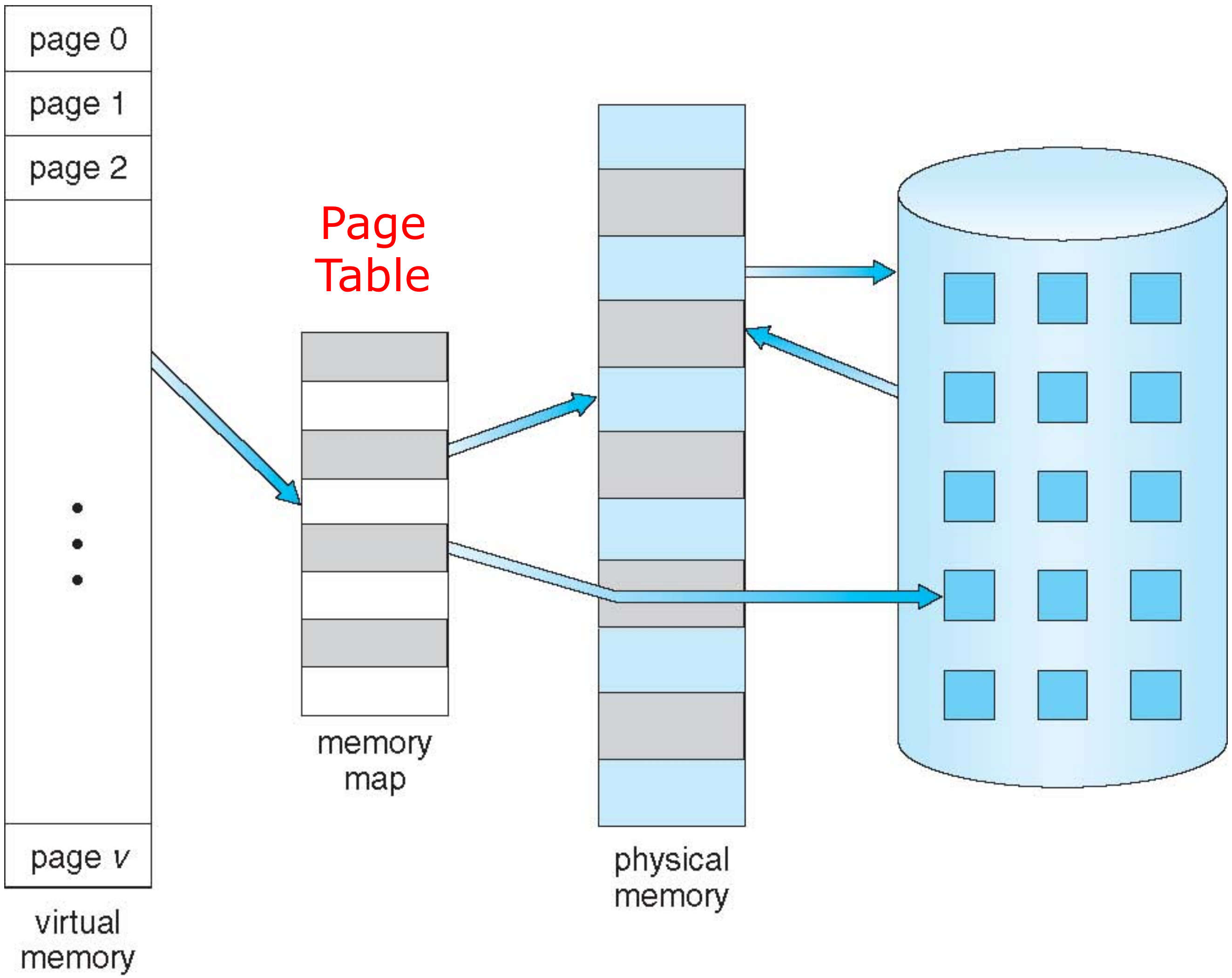
Virtual Memory: Background

- Code needs to be in memory to execute, but entire program rarely used
  - Error code, unusual routines, large data structures
- Entire program code not needed at same time
- Consider ability to execute partially-loaded program



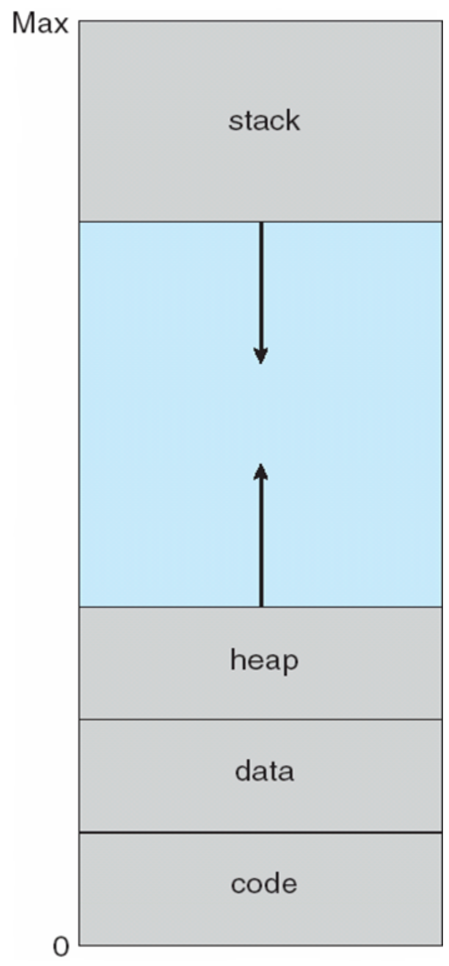
- **Virtual memory** – separation of user logical memory from physical memory
  - Only part of the program needs to be in memory for execution
  - Logical address space can be much larger than physical address space
- **Virtual address space** – logical view of how process is stored in memory
  - Usually start at address 0, contiguous addresses until end of space
  - Meanwhile, physical memory organized in page frames
  - MMU must map logical to physical

Virtual Memory Larger Than Physical Memory

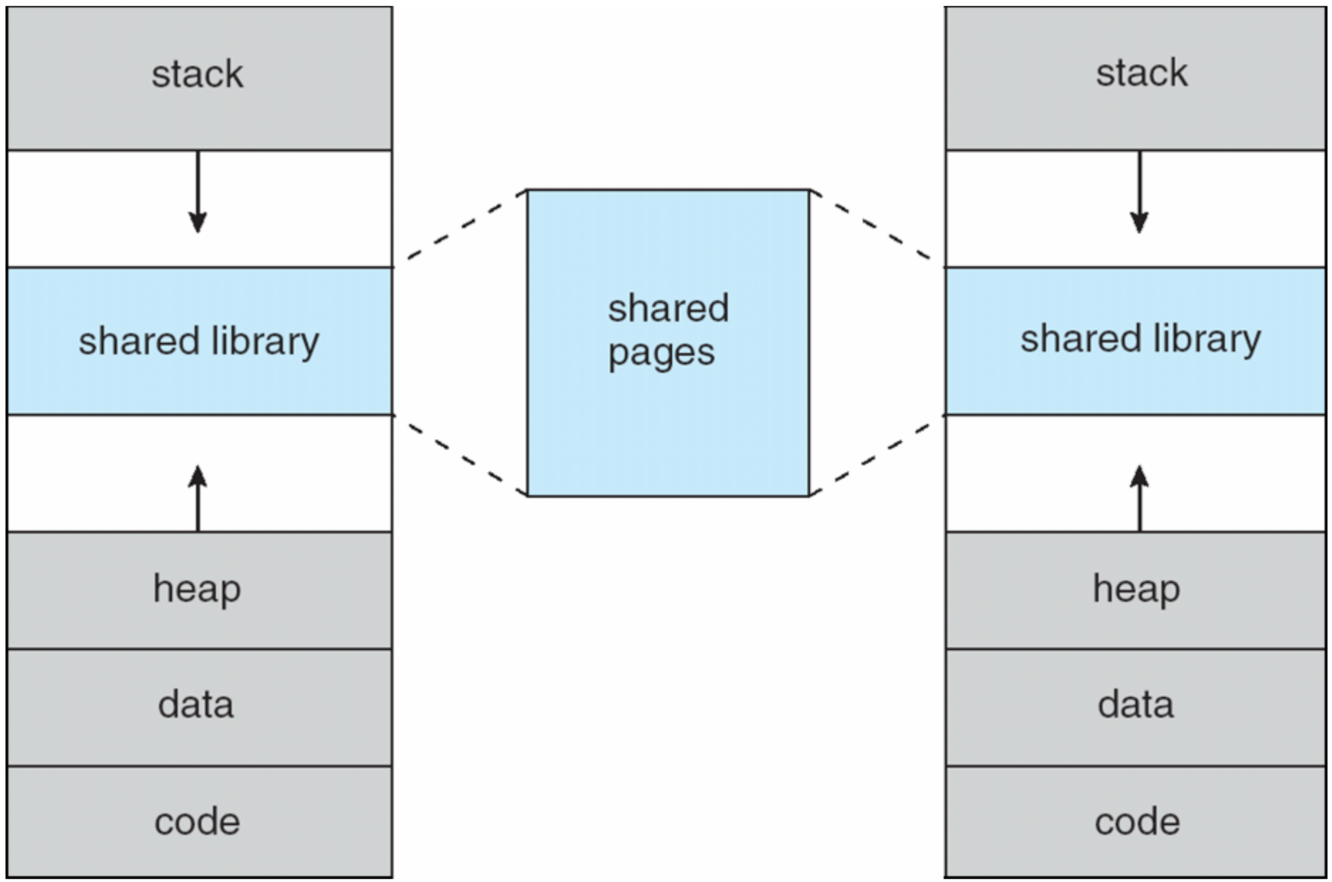


Virtual-address Space

- Usually design logical address space for stack to start at Max logical address and grow "down" while heap grows "up"
  - Maximizes address space use
  - Unused address space between the two is hole
    - No physical memory needed until heap or stack grows beyond a page
- System libraries shared via mapping into virtual address space



Shared Library Using Virtual Memory



Paging: Basic Concepts

- Divide physical memory into fixed-sized blocks called **frames**
  - Size is power of 2, between 512 bytes and 16 Mbytes
- Divide logical memory into blocks of same size called **pages**
- Keep track of all free frames
- To run a program of size ***N*** pages, need to find ***N*** free frames and load program
- Set up a **page table** to translate logical to physical addresses

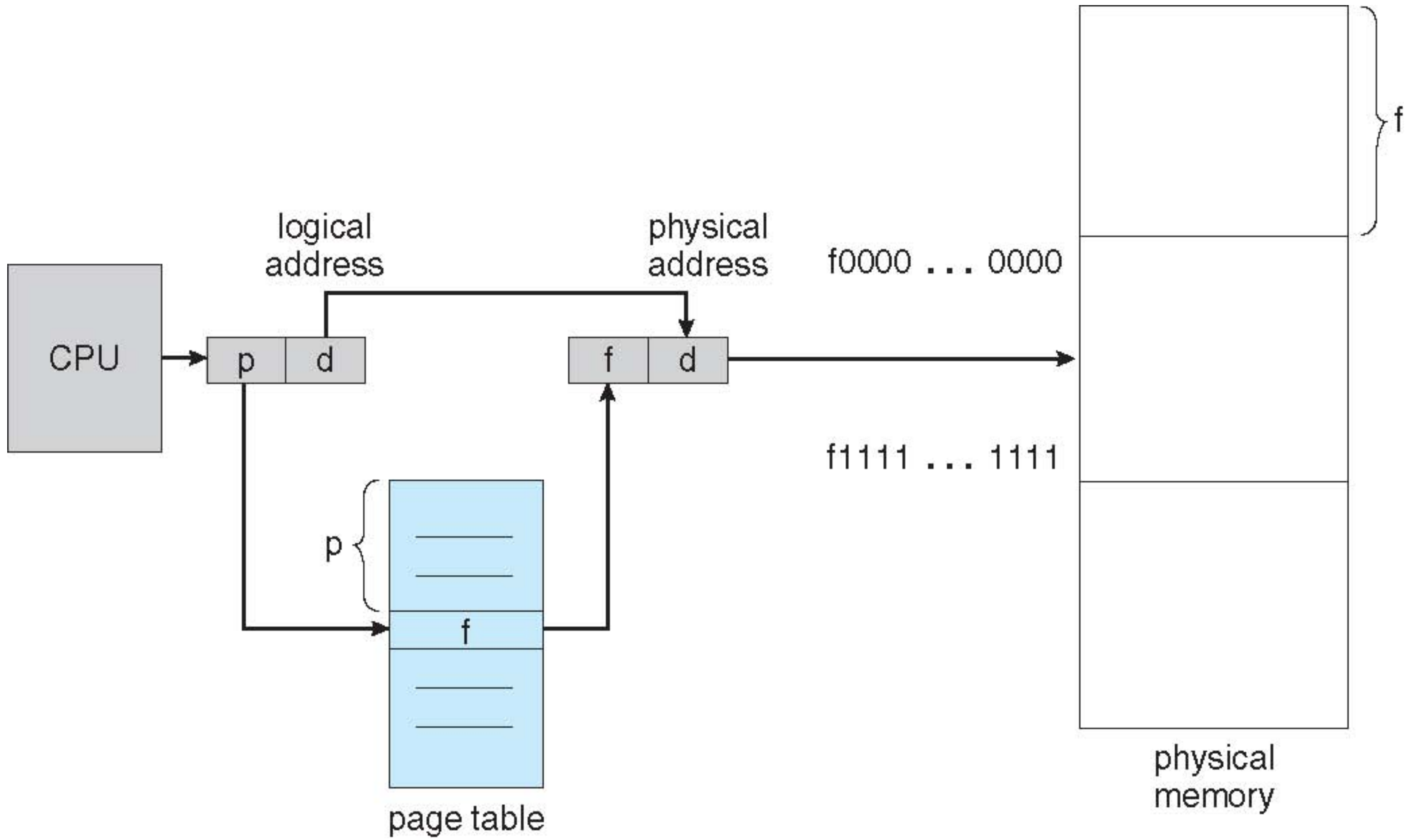
Address Translation Scheme

- Address generated by CPU is divided into:
  - **Page number** (*p*) – used as an index into a **page table** which contains base address of each page in physical memory
  - **Page offset** (*d*) – combined with base address to define the physical memory address that is sent to the memory unit

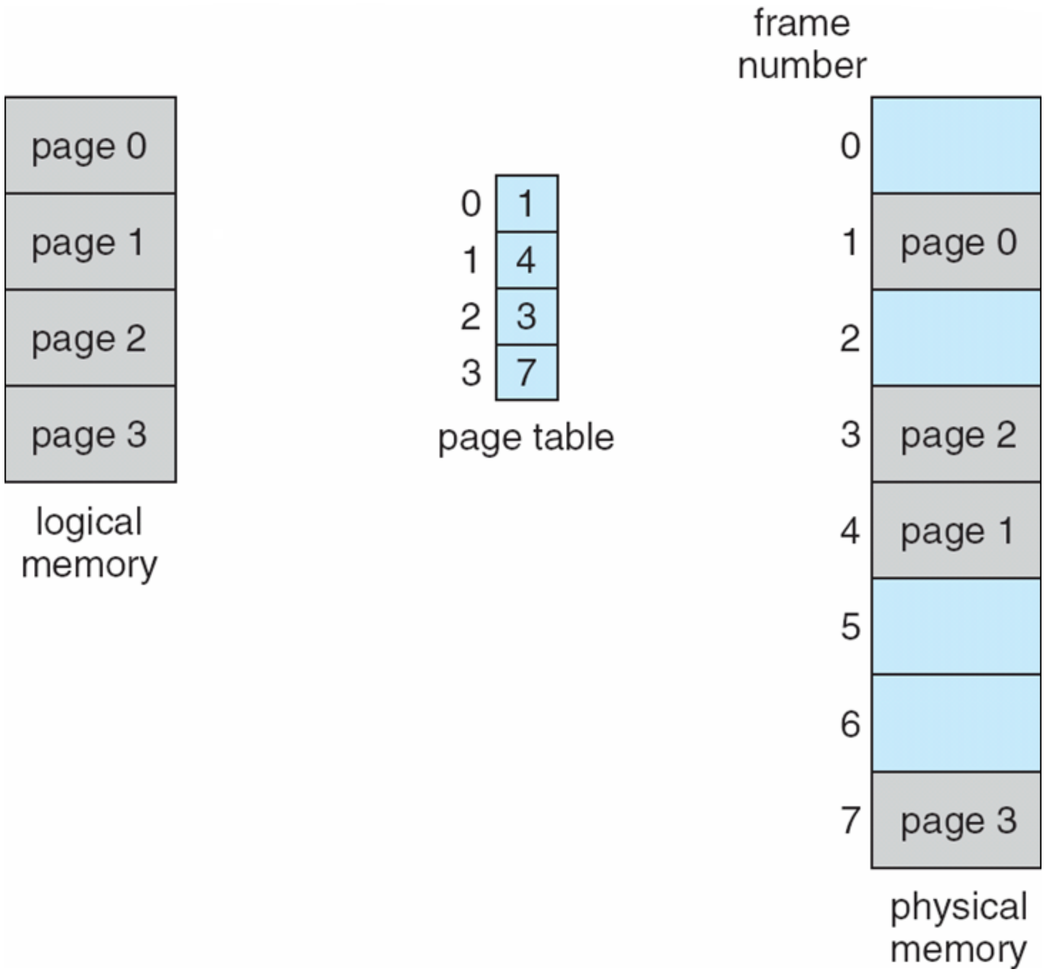
page number	page offset
p	d
m -n	n

  - For given logical address space 2^m and page size 2^n

Paging Hardware



Paging Model of Logical and Physical Memory

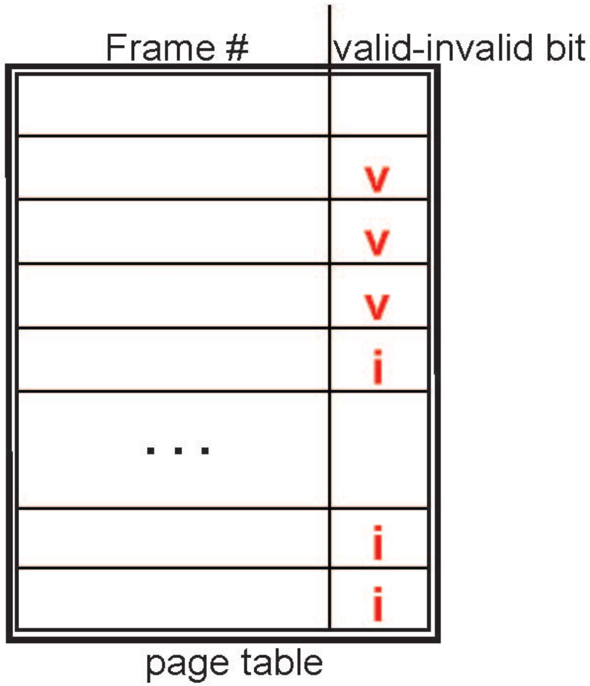


Demand Paging

- Could bring entire process into memory at load time
- Or bring a page into memory only when it is needed
  - No unnecessary I/O
  - Less memory needed
  - Faster response
- Page is needed $\Rightarrow$ reference to it
  - invalid reference $\Rightarrow$ abort
  - not-in-memory $\Rightarrow$ bring to memory

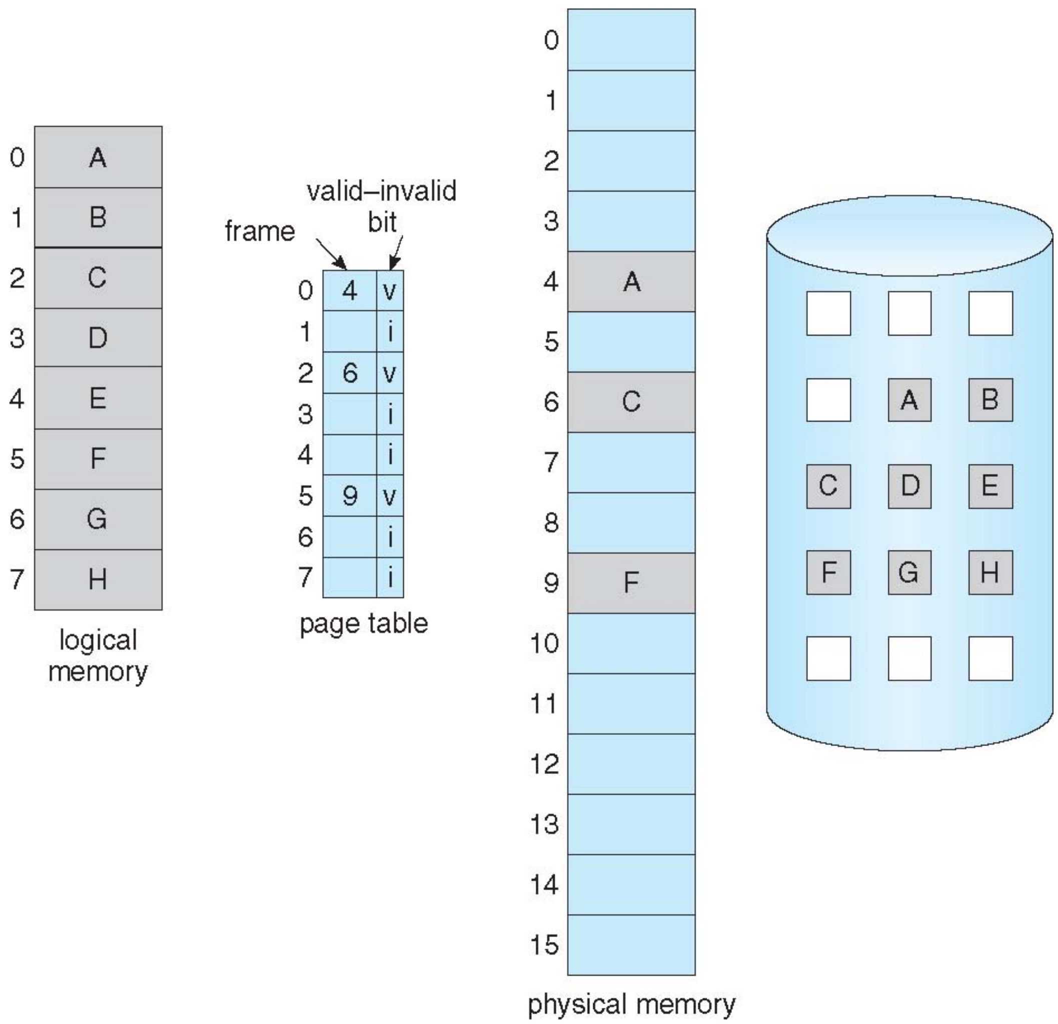
Valid-Invalid Bit

- With each page table entry a valid–invalid bit is associated
(**v** ⇒ in-memory – **memory resident**, **i** ⇒ not-in-memory)
- Initially valid–invalid bit is set to **i** on all entries
- Example of a page table snapshot:

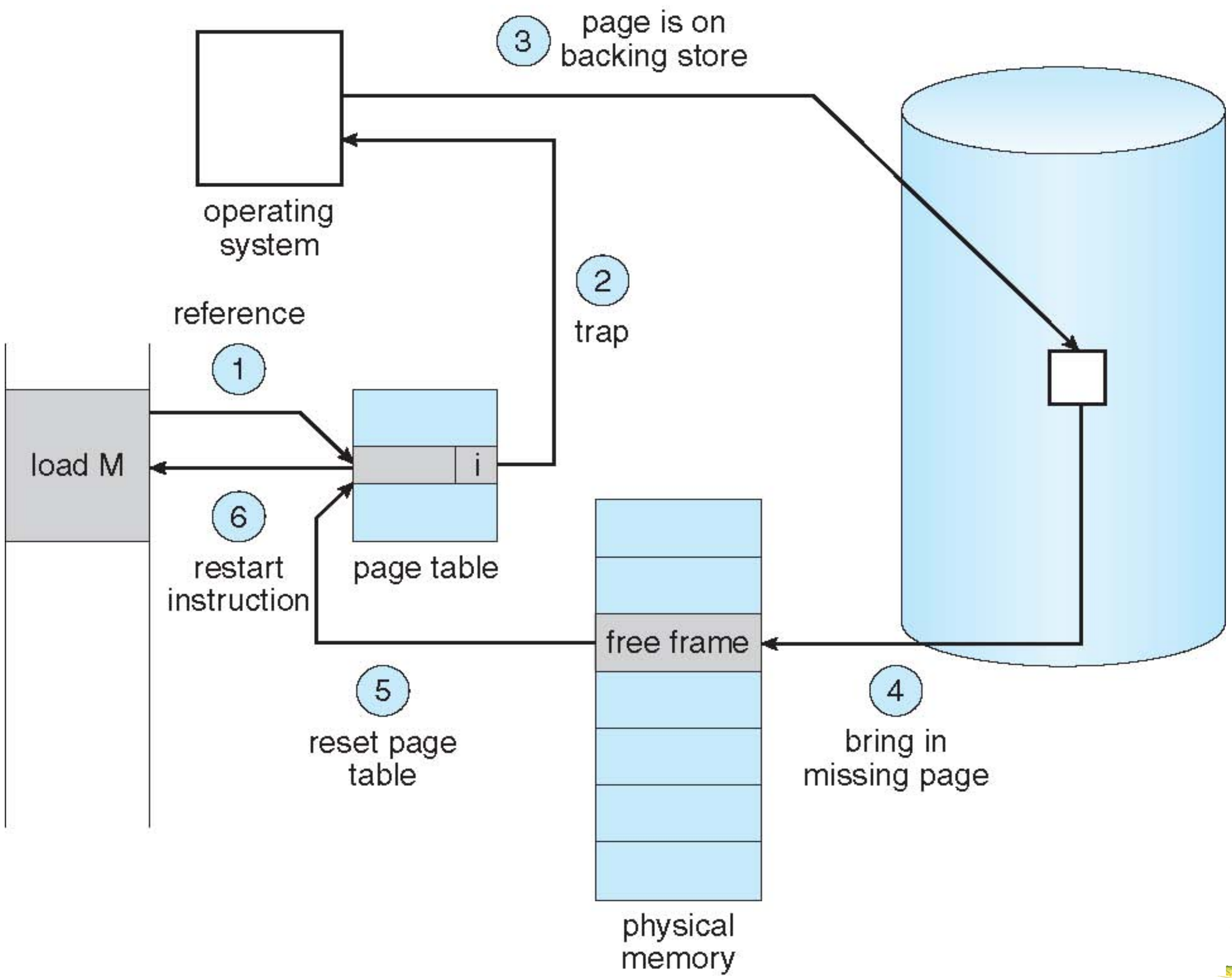


- During MMU address translation, if valid–invalid bit in page table entry is **i** ⇒ page fault

A Typical Snapshot



Steps in Handling a Page Fault



Segmentation

- Memory-management scheme that supports user view of memory
- A program is a collection of segments:
  - A segment is a logical unit such as:

    main program

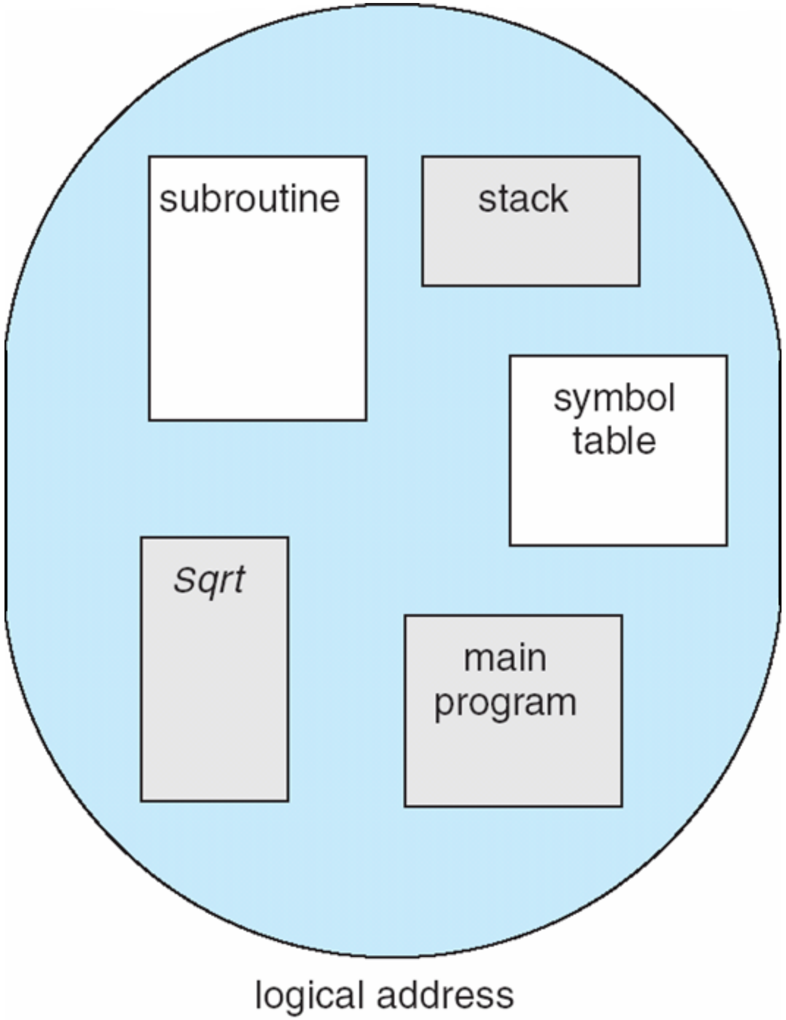
    procedure

    function / method

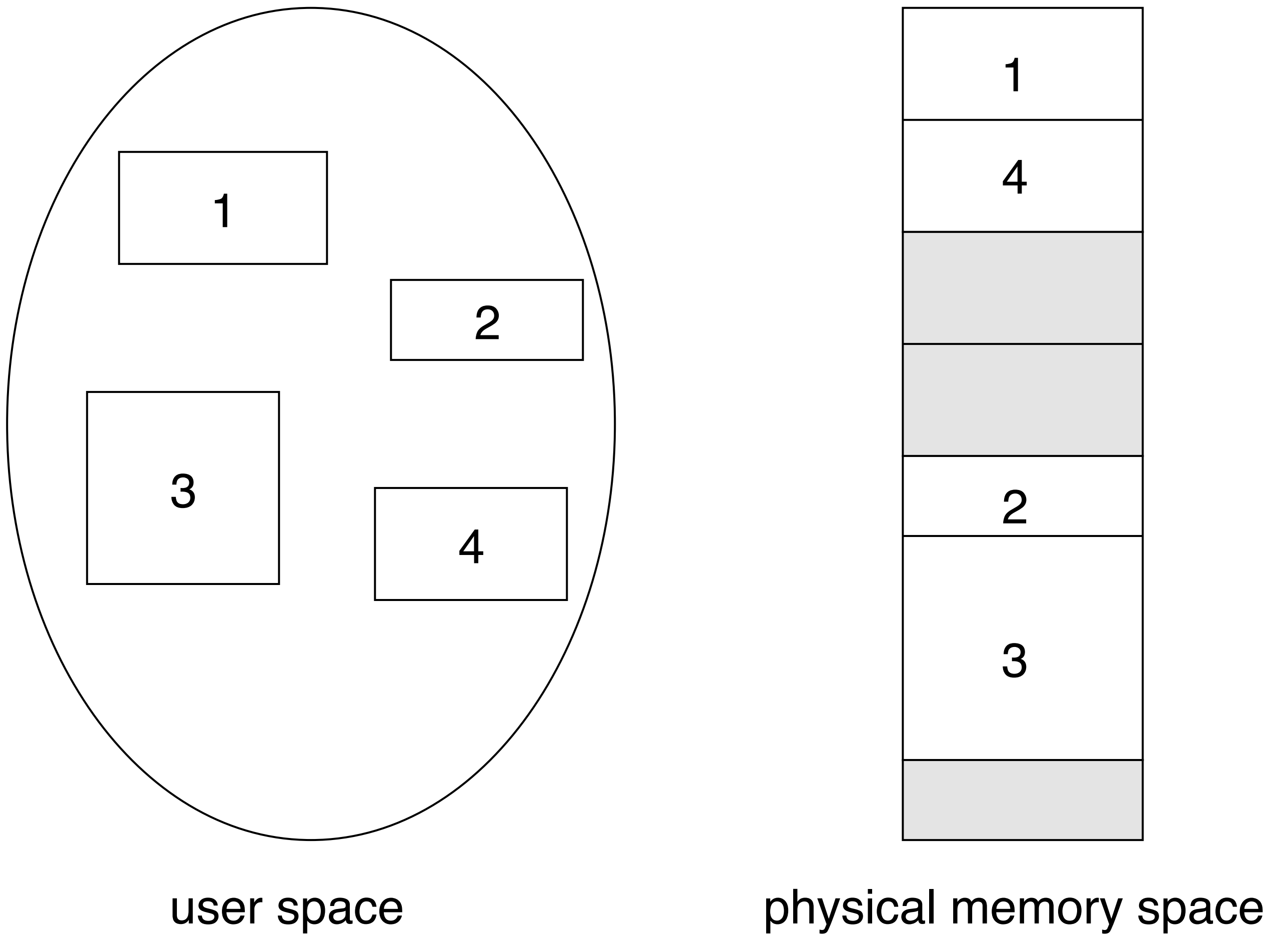
    stack

    arrays

User's View of a Program



Logical View of Segmentation



Segmentation Architecture

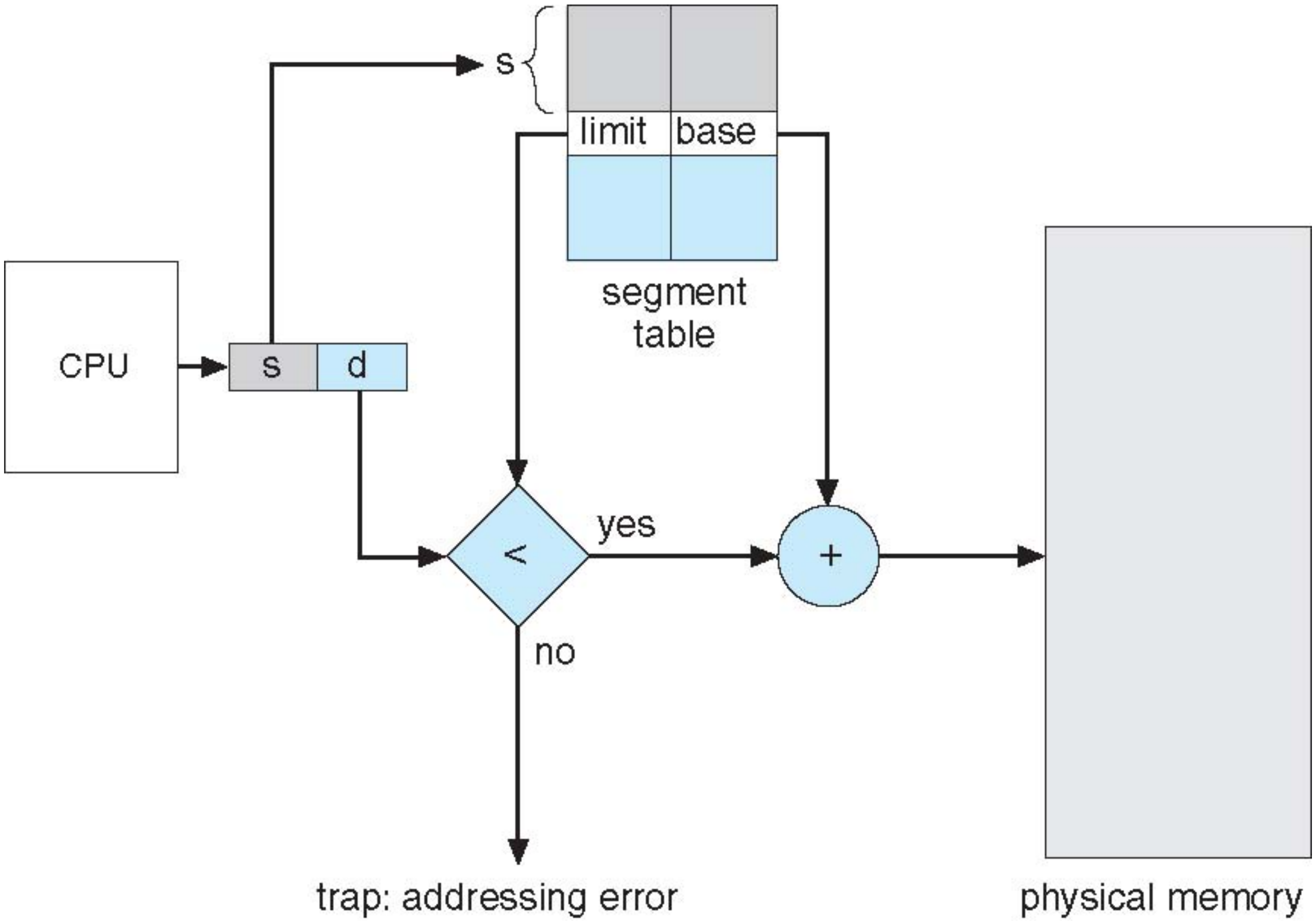
- Logical address consists of a two tuple:

  <segment-number, offset>,

- **Segment table** – maps two-dimensional physical addresses; each table entry has:
  - **base** – contains the starting physical address where the segments reside in memory
  - **limit** – specifies the length of the segment
- **Segment-table base register (STBR)** points to the segment table's location in memory
- **Segment-table length register (STLR)** indicates number of segments used by a program;

  segment number ***s*** is legal if ***s*** < **STLR**

Segmentation Hardware



To Summarize

- Segmentation provides memory protection.
  - While accessing a segment, we cannot access any memory location outside it.
- Paging allows efficient utilization of memory space.
  - A block of program or data need not be contiguous in memory.
- Often segmentation and paging are combined.
  - Program divided into logical segment.
  - Each segment is divided into pages.

Thank You