

② Differential entropy in nats

b) $f(x) = \frac{1}{\beta - \alpha} \quad \alpha \leq x \leq \beta$

sol:- Differential entropy = $-\int f \log f \, dx$

$$= -\int_{\alpha}^{\beta} \frac{1}{\beta - \alpha} \log \left(\frac{1}{\beta - \alpha} \right) dx$$

$$= \frac{1}{\beta - \alpha} \log(\beta - \alpha) (\beta - \alpha)$$

$$= \log(\beta - \alpha)$$

d) $f(x) = \frac{x}{b^2} e^{-\frac{x^2}{2b^2}}$

The given distribution is Rayleigh distribution

$$h(x) = -\int f \log f \, dx$$

$$\log f \Rightarrow \log \frac{x}{b^2} - \frac{x^2}{2b^2}$$

$$h(x) = - \underbrace{\int_0^{\infty} \frac{x}{b^2} e^{-\frac{x^2}{2b^2}} \log \frac{x}{b^2} dx}_{I_1} + \underbrace{\int_0^{\infty} \frac{x}{b^2} e^{-\frac{x^2}{2b^2}} \frac{x^2}{2b^2} dx}_{I_2}$$

Integral 1 (I_1):

$$I_1 = - \int_0^{\infty} x e^{-\frac{x^2}{2b^2}} [\log(x) - 2 \log b] dx$$

$$= - \int_0^{\infty} \frac{x}{b^2} e^{-\frac{x^2}{2b^2}} \log x \, dx + \int_0^{\infty} 2x e^{-\frac{x^2}{2b^2}} \log b \, dx$$

$$\left(\begin{array}{l} \text{let } x^2 = t \longrightarrow 2 \log x = \log t \\ \downarrow \\ 2x dx = dt \end{array} \right. \quad \log x = \frac{\log t}{2} \quad \Rightarrow \text{Substitution}$$

$$= - \int_0^{\infty} e^{-\frac{t}{2b^2}} \frac{\log t}{4} dt + \int_0^{\infty} e^{-\frac{t}{2b^2}} \log b \, dt$$

$$\left(\begin{array}{l} \text{let, } \underline{z} \quad \frac{t}{2b^2} = z \longrightarrow \log t = \log(2b^2) + \log z \\ \quad \quad \quad \downarrow \\ \quad \quad \quad t = 2b^2 z \\ \quad \quad \quad dt = 2b^2 dz \end{array} \right)$$

$$\begin{aligned} I_1 &= -\frac{1}{4} \int_0^{\infty} e^{-z} (\ln z + \ln 2b^2) 2b^2 dz + 2b^2 \log b \\ &= -\frac{b^2}{2} \underbrace{\int_0^{\infty} e^{-z} \ln z dz}_{\frac{1}{2}} - \frac{b^2}{2} \int_0^{\infty} \ln 2b^2 e^{-z} dz + 2b^2 \log b \\ &= -\frac{b^2}{2} \cdot \frac{1}{2} - \frac{b^2}{2} \ln 2b^2 + 2b^2 \log b \quad \text{--- (1)} \end{aligned}$$

Integral 2 (I₂):

$$I_2 = \int_0^{\infty} \frac{x}{b^2} e^{-\frac{x^2}{2b^2}} \frac{x^2}{2b^2} dx$$

$$\left(\text{let } x^2 = t \quad 2x dx = dt \right)$$

$$I_2 = \int_0^{\infty} \frac{1}{4b^4} e^{-\frac{t}{2b^2}} t dt$$

(By using the concept of By parts)

$$\begin{aligned} I_2 &= \frac{1}{4b^4} \left[t \int e^{-\frac{t}{2b^2}} dt \Big|_0^{\infty} - \int_0^{\infty} 1 \cdot \int e^{-\frac{t}{2b^2}} dt dt \right] \\ &= \frac{1}{4b^4} \left[-2b^2 t e^{-\frac{t}{2b^2}} \Big|_0^{\infty} - \int_0^{\infty} -2b^2 e^{-\frac{t}{2b^2}} dt \right] \\ &= \frac{1}{4b^4} \left[0 - (2b^2)^2 [0-1] \right] = \frac{4b^4}{4b^4} \end{aligned}$$

$$\text{By (1), (2)} \quad = 1 \quad \text{--- (2)}$$

$$h(x) = -\frac{I_1}{b^2} + I_2$$

$$= \frac{1}{2} - \frac{\ln 2b^2}{2} + 2 \log b + 1$$

$$= 1 + \frac{1}{2} + \log \frac{b^2}{\sqrt{2b^2}}$$

$$h(x) = 1 + \frac{1}{2} + \log\left(\frac{b}{\sqrt{2}}\right)$$

$$\therefore h(x) = 1 + \frac{1}{2} + \log\left(\frac{b}{\sqrt{2}}\right)$$

a) $f(x) = \frac{c}{x} x^{c-1} e^{-\frac{x^c}{d}} ; x, c, d > 0$

sol: $\log f(x) = \log \frac{c}{x} + (c-1) \log x - \frac{x^c}{d}$

$$h(x) = - \underbrace{\int_0^{\infty} \frac{c}{x} x^{c-1} e^{-\frac{x^c}{d}} \log \frac{c}{x} dx}_{I_1} - \underbrace{\int_0^{\infty} \frac{c}{x} x^{c-1} e^{-\frac{x^c}{d}} (c-1) \log x dx}_{I_2} + \underbrace{\int_0^{\infty} \frac{c}{x} x^{c-1} e^{-\frac{x^c}{d}} \frac{x^c}{x} dx}_{I_3}$$

$$I_1 = \int_0^{\infty} \left(\frac{c}{x} x^{c-1} \right) e^{-\frac{x^c}{d}} \log\left(\frac{c}{x}\right) dx$$

let $x^c = t$

~~$c \log x = \log t$~~

$c x^{c-1} dx = dt$

$$= \int_0^{\infty} \frac{e^{-\frac{t}{d}}}{x} \log\left(\frac{c}{x}\right) dt$$

$$= \frac{\log\left(\frac{c}{x}\right)}{x} \int_0^{\infty} \frac{e^{-t/d}}{1} dt$$

$$I_1 \Rightarrow \frac{\log \frac{c}{x}}{x} \left[\frac{e^{-t/d}}{-1/d} \right]_0^{\infty}$$

$$= \frac{\log \frac{c}{x}}{\left(\frac{-1}{d}\right)x} [0 - 1]$$

$$= \log\left(\frac{c}{x}\right)$$

I_2

$$I_2 = \int_0^{\infty} \left(\frac{c}{\alpha} x^{c-1} \right) e^{-\frac{x^c}{\alpha}} (c-1) \log x \, dx$$

$$\text{let } x^c = t \implies c \log x = \log t$$

$$c x^{c-1} dx = dt$$

$$= \int_0^{\infty} \frac{e^{-\frac{t}{\alpha}}}{c \alpha} \ln t \, dt$$

$$= \int_0^{\infty} \frac{1}{c \alpha} e^{-\frac{t}{\alpha}} \ln t \, dt$$

$$\text{let } \frac{t}{\alpha} = z \quad \ln t = \ln(z \alpha)$$

$$dt = \alpha \, dz$$

$$I_2 = \int_0^{\infty} \frac{1}{c \alpha} e^{-z} (\ln z + \ln \alpha) \alpha \, dz$$

$$= \int_0^{\infty} \frac{1}{c} e^{-z} \ln \alpha \, dz + \int_0^{\infty} \frac{1}{c} e^{-z} \ln z \, dz$$

$$= -\frac{\ln \alpha}{c} [0 - 1] + \frac{1}{c} \int_0^{\infty} e^{-z} \ln z \, dz$$

$$= \frac{\ln \alpha}{c} + \frac{1}{c} \int_0^{\infty} e^{-z} \ln z \, dz$$

$$\int e^{-z} \ln z \, dz = \ln z \int e^{-z} dz - \int \frac{1}{z} \int e^{-z} dz \, dz$$

$$= -e^{-z} \ln z - \int \frac{1}{z} e^{-z} dz$$

$$-\int \frac{e^{-z}}{z} dz = \Gamma(z)$$

$$I_2 = \frac{\ln \alpha}{c} + \frac{1}{c} \left[-e^{-z} \ln z - \Gamma(z) \right]_0^{\infty}$$

$$= \frac{\ln \alpha}{c} + \frac{\gamma}{c}$$

$$\text{let } I_3 = \int_0^{\infty} c \frac{x^{c-1}}{d} e^{-\frac{x^c}{d}} \frac{x^c}{d} dx$$

$$x^c = t$$

$$c x^{c-1} dx = dt$$

$$I_3 = \int_0^{\infty} \frac{t e^{-t/d}}{d^2} dt$$

$$\int_0^{\infty} t e^{-\frac{t}{d}} dt = \left[t \int e^{-t/d} dt \right]_0^{\infty} - \int_0^{\infty} 1 \cdot \int e^{-t/d} dt \cdot dt$$

$$= \left[-td e^{-t/d} \right]_0^{\infty} + \int_0^{\infty} d e^{-t/d} dt$$

$$= \left[-d^2 e^{-t/d} \right]_0^{\infty}$$

$$= -d^2 [0 - 1]$$

$$= d^2$$

$$I_3 = \frac{\int_0^{\infty} t e^{-t/d} dt}{d^2} = \frac{d^2}{d^2} = 1$$

Finally,

$$h(x) = -I_1 - I_2 + I_3$$

$$= -\log\left(\frac{c}{d}\right) - \frac{\ln d}{c} + \frac{x^c}{c} + 1$$

$$= 1 + \log\left(\frac{d}{c}\right) - \frac{\log d}{c} + \frac{x^c}{c}$$

Rearranging terms,

$$h(x) = \frac{(c-1)d}{c} + \log\left(d^{1/c}/c\right) + 1$$

4) a) $f(x) = \frac{a c^a}{x^{a+1}}$ Generalized Pareto

let x_1, x_2, \dots, x_n be random sample of pdf $f_X(x) = \frac{a \left(\frac{x-\mu}{\theta} \right)^{\frac{1}{\theta}-1}}{\theta^{\frac{1}{\theta}} \left(1 + \left(\frac{x-\mu}{\theta} \right)^{\frac{1}{\theta}} \right)^{a+1}}; x > \mu$

$y_{1:n} \leq y_{2:n} \leq \dots \leq y_{n-1:n} \leq y_{n:n}$

$g_{i:n}(y) = n \binom{n-1}{i-1} [F_{i:n}(y)]^{i-1} [1 - F_{i:n}(y)]^{n-i} f_X(y)$

Entropy

$H_{i:n}(y) = -\ln \left[\frac{n a}{\theta^{\frac{1}{\theta}}} \binom{n-1}{i-1} \right] + (i-1) E \left[\ln \left(\frac{y-\mu}{\theta} \right) \right] + [a(n-i+1)+1] E \left[\ln \left(1 + \frac{y-\mu}{\theta} \right)^{\frac{1}{\theta}} \right] + (1-i) E \left[\ln \left(1 - \left(1 + \frac{y-\mu}{\theta} \right)^{\frac{1}{\theta}} \right)^{-a} \right]$

We need to calculate,

$E \left[\left(\frac{y-\mu}{\theta} \right)^r \right] = \int_{\mu}^{\infty} \left(\frac{y-\mu}{\theta} \right)^r g_{i:n}(y) dy$
 $= n a \binom{n-1}{i-1} \sum_{l=0}^{i-1} (-1)^l \binom{i-1}{l} \frac{\Gamma(r \frac{1}{\theta} + 1) \Gamma(-r \frac{1}{\theta} + a(n-i+1+l))}{\Gamma(1 + a(n-i+1+l))}$

Solving,

$H_{i:n}(y) = -\ln \left(\frac{n a}{\theta^{\frac{1}{\theta}}} \right) + (i-1) [\psi(1) - \psi(nd)] = H_X(nd)$

By setting $a=1 \Rightarrow F_X(x) = 1 - \left(1 + \left(\frac{x-\mu}{\theta} \right)^{\frac{1}{\theta}} \right)^{-1} x > \mu$
 $f_X(x) = \frac{1}{\theta^{\frac{1}{\theta}}} \left(\frac{x-\mu}{\theta} \right)^{\frac{1}{\theta}-1} \left(1 + \left(\frac{x-\mu}{\theta} \right)^{\frac{1}{\theta}} \right)^{-2} x > \mu$

substituting

$$i=1 \text{ and } i=n$$

$$H_{1:n}(Y) = -\ln\left(\frac{n}{\theta}\right) + (n-1) [\psi(1) - \psi(n)] \\ = H_X(X=n)$$

Entropy of order statistics,

$$H_{n:n} = -\ln\left(\frac{n\theta}{\theta}\right) + \left[\frac{\psi(1) - \psi(n+1)}{1+1} \right] \\ + \left(\frac{n-1}{n}\right) + \left(\frac{n+1}{n}\right) [\psi(n+1) - \psi(1)]$$

$$H_{1:n} = \ln\left(\frac{\theta}{n\theta}\right) + \frac{n\theta+1}{n\theta} = H_X(nd)$$

let $\theta=1$ $\mu=0$ for Pareto Type I

$n=1$

$H_X(x) \Rightarrow \ln\left(\frac{\theta}{x}\right) + \frac{x+1}{x}$

(our approach:-
We solved first generalized
Pareto and substituted
for given Pareto (I))

For given question $\theta=k$
 $\alpha=a$

substituting,

$$H_X(a) = \ln\left(\frac{k}{a}\right) + \frac{a+1}{a}$$

$H_X(a) = \ln\left(\frac{k}{a}\right) + 1 + \frac{1}{a}$

$H_X(a) = \log\left(\frac{k}{a}\right) + 1 + \frac{1}{a}$

 { "answer" }

"Some results of integration taken from"

IEEE
Pages: "PARETO-TYPES AND ITS ORDER STATISTICS DISTRIBUTIONS"
By G.H. Yan and G.R.M. Bozadayan.

② Let $\{X_i\}$ be stationary stochastic process with entropy rate $H(X)$. Relation b/w $H(X)$ and $H(X_i)$

Sol:

We know entropy rate is defined as,

$$H(X) = H(X_1 / X_0, X_{-1}, \dots)$$

We also know that, conditioning reduces entropy, hence.

$$H(X_1 / X_0, X_{-1}, \dots) \leq H(X_1)$$

Let's see, when equality holds,

If X_i is independent of past X_0, X_{-1}, \dots i.e. if and only if X_i is an iid process.

↳ Then equality holds.

③ Let $Y_i = (X_{i-1}, X_i)$ X_i is a stationary process.
 $Z_i = (X_{3i-1}, X_{3i})$ Relation b/w entropy rates
 $W_i = X_{3i}$ $H(Y)$, $H(X)$, $H(Z)$ and $H(W)$

Sol:- $H(Y_1, Y_2, \dots, Y_n) = H(X_0, X_1, \dots, X_{n-1}, X_n)$

(Y_i part) $H(Y_1, \dots, Y_n) = H(X_1, \dots, X_{n+1})$ (By stationarity)

Dividing by n and taking limit, we get equality.

$$\therefore H(X) = H(Y) \text{ --- (1)}$$

$$H(Z_1, Z_2, \dots, Z_n) = H(X_2, X_3, \dots, X_{3n-1}, X_{3n})$$

$$H(Z_1 / Z_0, \dots) = H(X_2, X_3 / (X_{-1}, X_0) (X_{-3}, X_{-4}))$$

$$= H(X_2 / X_{-1}, X_0, X_{-4}, X_{-3}, \dots) + H(X_3 / X_{-1}, X_0, X_{-2})$$

①
②

$$\textcircled{1} H(x_2/x_{-1}, x_0, x_{-4}, x_{-3}) \leq H(x_2/x_1, x_0, x_{-1} \dots) = H(x)$$

$$\textcircled{2} H(x_3/x_2, x_{-1} \dots) \leq H(x_3/x_2, x_1, x_0, x_{-1} \dots) = H(x)$$

$$H(w_1, w_2 \dots w_n) = H(x_3, x_0 \dots x_{2n}) \quad (w_i \text{ part})$$

$$H(w_1/w_0 \dots) = H(x_3/x_0, x_{-3} \dots)$$

$$H(x_3/x_0, x_{-3} \dots) \leq \lim_{n \rightarrow \infty} H(x_3/x_2, x_1, x_0, x_{-1} \dots)$$

$$H(w) \leq H(x)$$

$$\therefore H(x) > H(w) \longrightarrow \textcircled{3}$$

Following z_i and y_i

we can notice that

$$\begin{array}{ccc} z_i & = & y_{3i} \\ \downarrow & & \downarrow \\ (x_{3i-1}, x_{3i}) & & (x_{3i-1}, x_{3i}) \end{array}$$

$$H(z_1/z_0, z_{-1} \dots) = H(y_3/y_0, y_{-3} \dots)$$

$$H(y_3/y_0, y_{-3} \dots) \leq \lim_{n \rightarrow \infty} H(y_3/y_2, y_1, y_{-1} \dots)$$

From $\textcircled{1}$

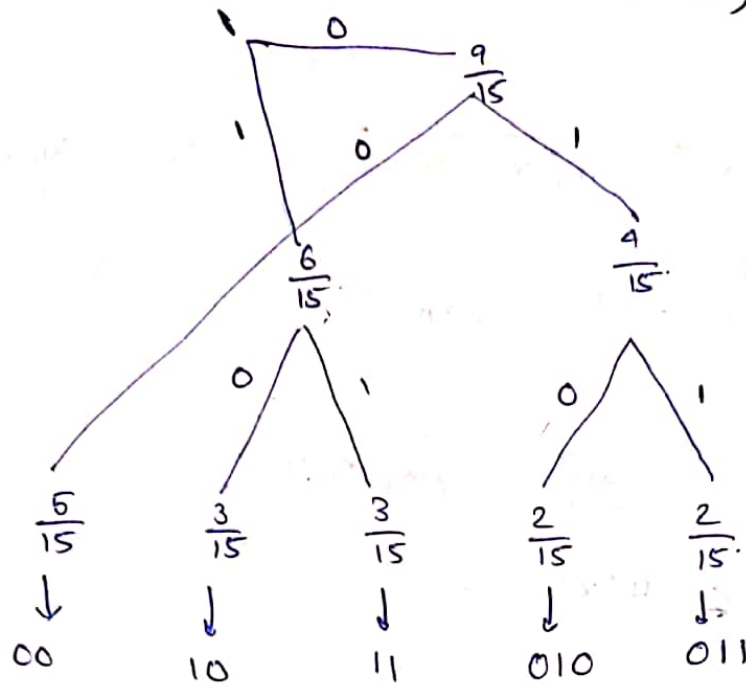
$$H(z) \leq H(y) = H(x)$$

$$\therefore H(z) < H(y) = H(x) \longrightarrow \textcircled{2}$$

- \Rightarrow Answers:
- $\textcircled{1} H(x) = H(y) \longrightarrow \textcircled{1}$
 - $\textcircled{2} H(x) > H(z) \longrightarrow \textcircled{2}$
 - $\textcircled{3} H(x) > H(w) \longrightarrow \textcircled{3}$

6) a) Coding scheme for,

$$\left(\frac{1}{3}, \frac{1}{7}, \frac{1}{5}, \frac{2}{15}, \frac{2}{15} \right)$$



Code word	Probability
00	$\frac{1}{3}$
10	$\frac{1}{5}$
11	$\frac{1}{5}$
010	$\frac{2}{15}$
011	$\frac{2}{15}$

(b) x_1, x_2, \dots, x_n be iid accⁿ to $P(x)$

Find $\lim_{n \rightarrow \infty} [P(x_1, \dots, x_n)]^{1/n}$

Sol:- As x_1, x_2, \dots, x_n are iid,
 $\log P(x_i)$ are also iid's.

$$\begin{aligned} \lim_{n \rightarrow \infty} (P(x_1, \dots, x_n))^{1/n} &= \lim_{n \rightarrow \infty} 2^{\frac{1}{n} \log P(x_1, \dots, x_n)} \\ &= \lim_{n \rightarrow \infty} 2^{\frac{1}{n} \log P(x_1) + \dots + \log P(x_n)} \end{aligned}$$

$$\text{As } P(x_1, \dots, x_n) = P(x_1) P(x_2) \dots P(x_n)$$

$$= \lim_{n \rightarrow \infty} \left(2^{\frac{1}{n} \log (\prod_{i=1}^n P(x_i))} \right)$$

$$= \lim_{n \rightarrow \infty} 2^{\frac{1}{n} \sum_{i=1}^n \log P(x_i)}$$

From law of large numbers

$\log P(x_i)$ are also iids.

$$\frac{\sum_{i=1}^n \log P(x_i)}{n} \rightarrow \mathbb{E}(\log P(x))$$

$$= \frac{\mathbb{E}[\log P(x)]}{2}$$

$$= \frac{-H(x)}{2}$$

$$\therefore \lim_{n \rightarrow \infty} (P(x_1, \dots, x_n))^{1/n} = 2^{-H(x)}$$

①

$$f: X \rightarrow [0, 1]$$

$x \in X$, $f_A(x)$ degree of inclusion of x in set A .

(a) Design entropy to quantify uncertainty associated with inclusion of x in A . Here $\boxed{\mu_A(x) \Rightarrow f_A(x)}$

sol:- Let P be a probability measure over R^n .

$$P(A) = \int_A dP$$

(or) equivalently

$$P(A) = \int_{R^n} \mu_A(x) dP \quad \xrightarrow{\text{contin}} \quad = \sum_{x \in R^n} \mu_A(x) \quad \xrightarrow{\text{discrete}}$$

$$= E(\mu_A) \quad \text{--- ①}$$

$$\mu_A: R^n \rightarrow [0, 1]$$

Probability of fuzzy set Event A is defined by the Lebesgue-Stieltjes integral given in eq (1)

Some properties of fuzzy set

$$A \subset B \Rightarrow \mu_A(x) \leq \mu_B(x) \quad \forall x$$

$$A \cup B \quad \mu_{A \cup B}(x) = \max[\mu_A(x), \mu_B(x)] \quad \forall x$$

Turning to notion of entropy, we note that its usual defⁿ in information theory is as follows:

let x be a random variable takes x_1, \dots, x_n
 P_1, \dots, P_n

$$\text{Entropy given by } H(x) = - \sum_{i=1}^n P_i \log P_i$$

Here for fuzzy sets, we quantify the entropy as a weighted Shannon entropy.

For a random variable z an event occurs with probability P , here uncertainty defined as $\mu_A(x_i)$
 Event x .

$$\begin{aligned} & \rightarrow - \underbrace{\mu_A(x)}_{\text{weighted}} P \log P \Rightarrow \frac{-\mu_A(x) P \log P}{P(x)} \end{aligned}$$

For a binary $\begin{cases} p \rightarrow x_1 \\ 1-p \rightarrow x_2 \end{cases}$ events

$$H(A) = -\mu_A(x_1) P \log P - \mu_A(x_2) (1-P) \log(1-P)$$

$$= -\mu_A(x_1) P \log P - \mu_A(x_2) (1-P) \log(1-P)$$

(b) let $X = \{x_1, \dots, x_n\}$ with probabilities

$\{p_1, \dots, p_n\}$

$H(X)$ relation with (a).

sol: By using the defn of above entropy as a weighted Shannon entropy, for all events $1 - n$

This definition, suggests that entropy of a fuzzy set subset A , of finite set $\{x_1, \dots, x_n\}$ with respect to a probability distribution $P = \{p_1, \dots, p_n\}$ be defined as follows

$$H^P(A) = - \sum_{i=1}^n \mu_A(x_i) p_i \log p_i$$

$(\mu_A \rightarrow f_A)$ \rightarrow Entropy of fuzzy set A w.r.t probabilistic distribution P .

If x and y are independent R.V's with P, Q (probabilities (p_1, \dots, p_n) (q_1, \dots, q_m))

$$H(x, y) = H(x) + H(y)$$

$$H^{P \otimes Q}(A, B) = P(A) H^P(A) + P(B) H^Q(B)$$

5) Consider a channel $(x^n, H(y^n/x^n), y^n)$ DMC

↳ Entropy transition matrix.

channel capacity, Max. probab of error, avg. probab of error and rate.

sol: ⇒ we know Rate = $\frac{1}{n} \log M$. For a (n, M) code.

By property of chain rule,

$$\begin{aligned} H(x^n, y^n) &= H(x^n) + H(y^n/x^n) \quad \text{--- ①} \\ &= H(y^n) + H(x^n/y^n) \quad \text{--- ②} \end{aligned}$$

By property of DMC,

$$P(y^n/x^n) = \prod_{i=1}^n P(y_i/x_i)$$

Applying E and log on B.S.

$$-E \log P(y^n/x^n) = \sum_{i=1}^n -E[\log P(y_i/x_i)]$$

$$\boxed{H(y^n/x^n) = \sum_{i=1}^n H(y_i/x_i)} \quad \text{--- ②}$$

From the achievability of channel coding theorem, for every i/p distribution $P(x^n)$ $I(x^n; y^n)$ is achievable as rate for large n^{th} extension such that

① Rate R arbitrarily close to $I(x^n; y^n)$

② λ_{max} (max error probability) ~~close~~ arbitrarily small.

without loss of generality,

choose $P(x^n)$ to be one that achieves channel capacity i.e. $I(x^n; y^n) = C$

$$\begin{aligned} I(x^n; y^n) &= H(x^n) - H(x^n/y^n) \quad \text{--- ③} \\ &= H(y^n) - H(y^n/x^n) \end{aligned}$$

From ②,

$$H(y^n/x^n) = \sum_{i=1}^n H(y_i/x_i)$$

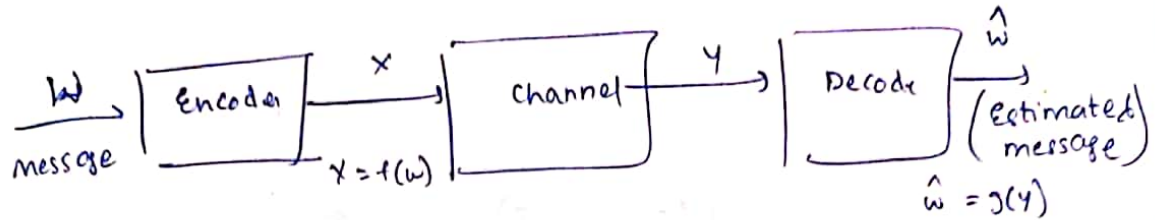
↳ This can be obtained from entropy transition matrix

We know $H(X^n)$ distribution at transmitter, $H(Y^n)$ at receiver

From ①,

$H(Y^n)$ can be known. [from $H(X^n/Y^n)$.]

$\therefore I(X^n; Y^n)$ can be calculated. — ④



λ_w : Error probability given message w sent,

$$\lambda_w = \Pr\{\hat{w} \neq w / w = w\}$$

$$\lambda_w = \Pr\{g(Y^n) \neq w / x^n = f^n(w)\}$$

$$\lambda_{\max} = \max_w \lambda_w$$

Avg. probability of error

$$\Pr\{\text{Err}\} = P_e = \frac{1}{M} \sum_w \lambda_w$$

$$\Pr\{\text{Err}\} = \sum_{w=1}^M \Pr\{\text{Err} / w = w\} \Pr\{w = w\}$$

$$= \Pr\{\text{Err} / w = 1\} \sum_{w=1}^M \Pr\{w = w\}$$

$$= \Pr\{\text{Err} / w = 1\}$$

$$\therefore P_e = \lambda_1 = \lambda_2 = \dots = \lambda_{\max}$$

$$E_w : \text{Event} = \left\{ \left(\tilde{x}^n(w), Y \right) \in T_{[X, Y]}^n \right\}$$

Jointly typical.

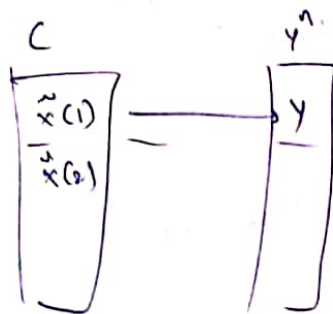
$$\Pr\{\text{Err}\} \leq \Pr\{E^c / w = 1\} + \sum_{w=2}^M \Pr\{E_w / w = 1\}$$

— ④a

By strong JAFEP,

$$\Pr\{E^c/w=1\} = \Pr\{(\tilde{x}(1), y) \notin T_{[x,y]}^n / w=1\} < v$$

↳ Not jointly typical



$$\Pr\{E^c/w=1\} \leq 2^{-n(I(x;y)-\epsilon)} \quad - (5)$$

From eq-4 getting $I(x;y)$ we can calculate

and

$$\frac{1}{n} \log M < I(x;y) - \frac{\epsilon}{4} \Rightarrow M < 2^{n(I(x;y) - \frac{\epsilon}{4})}$$

$$R < I(x;y) - \frac{\epsilon}{4} \Rightarrow \text{(Rate can be calculated from } I(x;y) \text{ calculated)}$$

Substituting (6) in (5)

$$\Pr\{E^c\} < v + 2^{\underbrace{n(I(x;y) - \frac{\epsilon}{4})}_{\substack{\downarrow \\ \text{m terms}}}} \underbrace{2^{-n(I(x;y) - \epsilon)}}_{\substack{\downarrow \\ \text{Each one.} \\ \Pr\{E^c/w\}}}$$

$$P_e \leq v + 2^{-n(\frac{\epsilon}{4} - \epsilon)} < \frac{\epsilon}{3} < \frac{\epsilon}{6}$$

(Pe < ε/2) ⇒ small.

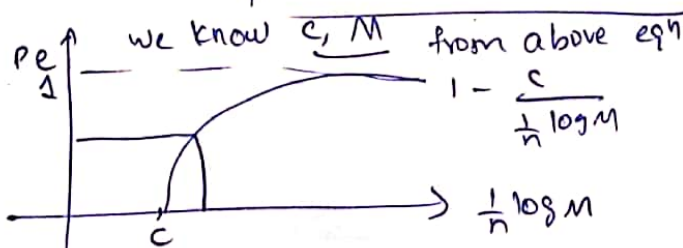
$\begin{matrix} \tau \rightarrow 0 & \delta \rightarrow 0 \\ \frac{\epsilon}{4} - \tau > 0 \end{matrix}$

Here (5), calculate P_e ,

By our assumption $\lambda_1, \lambda_2 - \lambda_n = \lambda_{\max} = \epsilon$

P_e can also be calculated by
Other calculation,

$$P_e \geq 1 - \frac{c}{\frac{1}{n} \log M} \quad - (7)$$



Hence from eq-① to ⑦,

Capacity, P_e , A_{\max} , R can be calculated
from the entropy transition matrix