

Riya Gaur

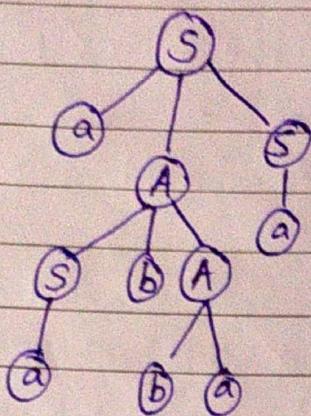
1. $S \rightarrow aAS1a$

$A \rightarrow SbA1ss1ba$

By LMD \rightarrow

$S \rightarrow aAS \rightarrow aSbAS \rightarrow aabAS \rightarrow aabbAS \rightarrow aabbbaa \dots$

Derivation Tree \rightarrow



2. $S \rightarrow 0B|1A$

$B \rightarrow 010S11AA$

$B \rightarrow 111S10BB$

String = 00110101

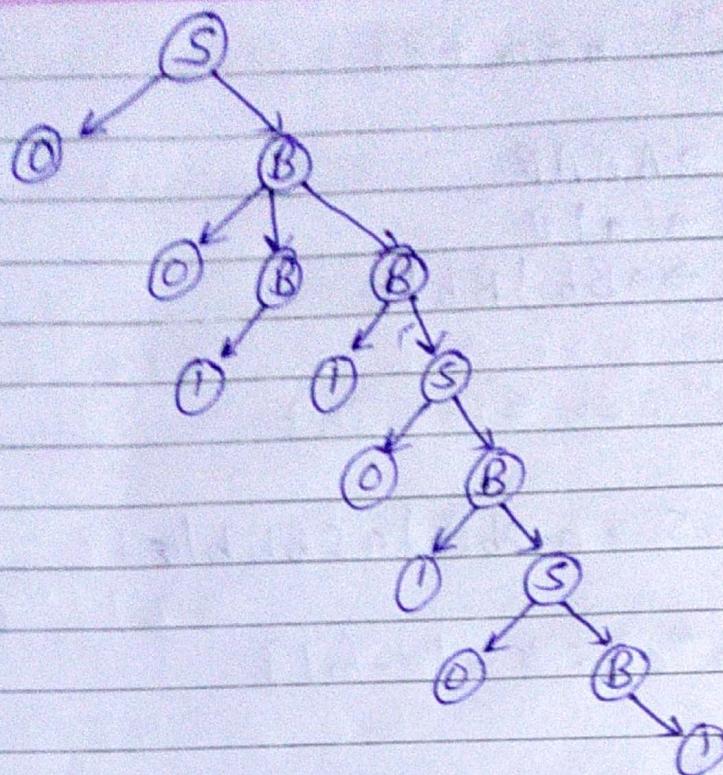
(a) LMD \rightarrow

$S \rightarrow 0B \rightarrow 00BB \rightarrow 001B \rightarrow 0011S \rightarrow 00110B \rightarrow 001101S \rightarrow 0011010B \rightarrow 00110101$

(b) RMD \rightarrow

$S \rightarrow 0B \rightarrow 00BB \rightarrow 00B1S \rightarrow 00B10B \rightarrow 00B101S \rightarrow 00B1010B \rightarrow 00B10101 \rightarrow 00110101$

(c) Derivation Tree \rightarrow



3.

$$(i) L = (aaa^* + b) \\ = \{b, aa, aaa, aaaa, \dots\}$$

$$G(L) = S \rightarrow aaA \mid b \\ A \rightarrow aA \mid \epsilon$$

$$(ii) L = \{a^n b^n : n \geq 1\}$$

$$G(L) = S \rightarrow aSb \mid ab$$

$$(iii) L = \{a^n b^{n+1} : n \geq 2\}$$

$$G(L) = S \rightarrow aSb \mid abb$$

$$(iv) L = w\omega^n : w \in \{a, b\}^*$$

$$G(L) = S \rightarrow aS_a \mid bS_b \mid aa \mid bb \mid \epsilon$$

(V) $L = \{a^n b^m : n \leq m + 3\}$

$$G(L) = S \rightarrow AAAAB$$

$$A \rightarrow a \mid \epsilon$$

$$B \rightarrow ab \mid Bb \mid B \mid \epsilon$$

(VI) $L = \{a^n b^m : 2n \leq m \leq 3n\}$

$$G(L) = S \rightarrow aSbbblasbbbbl\epsilon$$

(VII) $L = \{a^n b^m c^k : k = |n - m|\}$

$$K = |n - m|$$

$K = n - m$ for $n \geq m$ or $K = m - n$ for $m \geq n$

$$n = k + m \text{ or } m = k + n.$$

$$L = L_1 \cup L_2 \text{ where,}$$

$$L_1 = \{a^n b^m c^k : n = k + m\} \quad \& \quad L_2 = \{a^n b^m c^k : m = k + n\}$$

$$G(L_1) = S_1 \rightarrow T_1 \mid T_2$$

$$T_1 \rightarrow aT_1 b \mid \epsilon$$

$$T_2 \rightarrow aT_2 c \mid T_1 \mid \epsilon$$

$$G(L_2) = S_2 \rightarrow T_3 \mid T_4$$

$$T_3 \rightarrow aT_3 b \mid \epsilon$$

$$T_4 \rightarrow bT_4 c \mid \epsilon$$

$$G(L) = S \rightarrow S_1 \mid S_2$$

$$S_1 \rightarrow T_1 \mid T_2$$

$$T_1 \rightarrow aT_1 b \mid \epsilon$$

$$T_2 \rightarrow aT_2 c \mid T_1 \mid \epsilon$$

$$S_2 \rightarrow T_3 \mid T_4$$

$$T_3 \rightarrow aT_3 b \mid \epsilon$$

$$T_4 \rightarrow bT_4 c \mid \epsilon$$

4. $S \rightarrow aSbs \mid bsas \mid \epsilon$

Using Left Most Derivative

$LMD_1 \rightarrow S \rightarrow aSbs \rightarrow abSasbs \rightarrow ababS \rightarrow abab$

$LMD_2 \rightarrow S \rightarrow aSbS \rightarrow abS \rightarrow abasbs \rightarrow ababS \rightarrow abab$

We have have two LMD for same string

Hence grammar is ambiguous

5. $S \rightarrow ABlaaB$,

$A \rightarrow a \mid Aa$,

$B \rightarrow b$

$LMD_1 \rightarrow S \rightarrow AB \rightarrow AaB \rightarrow aaB \rightarrow aab$

$LMD_2 \rightarrow S \rightarrow aaB \rightarrow aab$

Hence given grammar is ambiguous.

6. A regular grammar is also a context free grammar as regular grammar is a subset of context free grammar. So a regular grammar may or may not be ambiguous depending upon the language given in a grammar but if a grammar is having only regular languages it can never be ambiguous or inherently ambiguous.

$$7. E \rightarrow E + E$$

$$E \rightarrow E * E$$

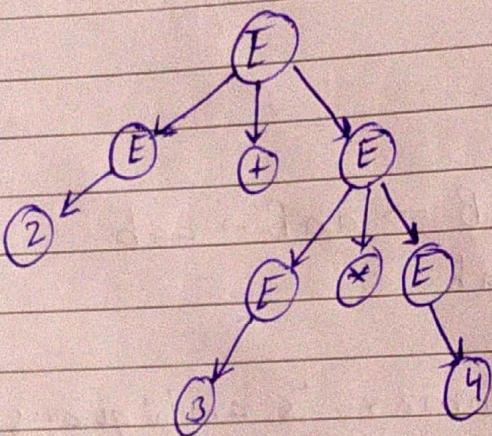
$$E \rightarrow id(2|3|4)$$

$$\text{Input } t \rightarrow 2 + 3 * 4$$

$$\text{LMD} \rightarrow E \rightarrow E + E \rightarrow 2 + E \rightarrow 2 + E * E \rightarrow 2 + 3 * E \rightarrow 2 + 3 * 4$$

$$\text{RMD} \rightarrow E \rightarrow E + E \rightarrow E + E * E \rightarrow E + E * 4 \rightarrow E + 3 * 4 \rightarrow 2 + 3 * 4$$

Derivation Tree \rightarrow



8.

$$(i) S \rightarrow AB,$$

$$A \rightarrow a,$$

$$B \rightarrow b,$$

$$B \rightarrow C,$$

$$E \rightarrow c$$

$$V_1' = \{A, B, E\}$$

$$V_2' = \{A, B, E, S\}$$

$$V_3' = \{A, B, E, S\}$$

$$= V_2'$$

$$A \rightarrow a, B \rightarrow b, E \rightarrow c$$

$$S \rightarrow AB$$

We get,

$$V' = \{S, A, B, E\}$$

$$P' = S \rightarrow AB, A \rightarrow a, B \rightarrow b \& E \rightarrow c$$

After phase 1 \Rightarrow

$$G' \rightarrow S \rightarrow AB, A \rightarrow a, B \rightarrow b \& E \rightarrow c$$

$$V_1'' = \{S\}$$

$$V_2'' = \{S, A, B\} \quad S \rightarrow AB$$

$$V_3'' = \{S, A, B\} \quad A \rightarrow a, B \rightarrow b \\ = V_2''$$

We get,

$$V''' = \{S, A, B\}$$

$$P''' = S \rightarrow AB, A \rightarrow a, B \rightarrow b$$

$$\text{iii) } S \rightarrow AB \mid CA$$

$$B \rightarrow BC \mid AB$$

$$A \rightarrow a$$

$$C \rightarrow aB \mid b$$

$$\text{Phase 1} \quad V_1' = \{A, C\}$$

$$V_2' = \{A, C, S\}$$

$$V_3' = \{A, C, S\}$$

$$= V_2'$$

$$S' = \{ \{A, C, S\}, \{a, b\}, P', \{S\} \}$$

$$\text{when } P' = S \rightarrow CA, A \rightarrow a, C \rightarrow b$$

$$\text{Phase 2} \quad V_1'' = \{S\}$$

$$V_2'' = \{S, A, C\}$$

$$V_3'' = \{S, A, C\} = V''$$

$$S'' = \{\{S, A, C\}, \{a, b\}, P'', \{S\}\}$$

$$\text{when } P'' = S \rightarrow CA, A \rightarrow a, C \rightarrow b$$

(iii) $S \rightarrow aAa,$
 $A \rightarrow sblbcc | DaA,$
 $C \rightarrow abb | DD,$
 $E \rightarrow aC,$
 $D \rightarrow aDA$

Phase 1 $V_1' = \{C\}$

$V_2' = \{C, A, E\}$

$V_3' = \{S, C, A, E\}$

$V_4' = \{S, C, A, E\} = V_3'$

$G' = \{S, C, A, E\}, \{a, b\}, P', \{S\}$

where $P' = S \rightarrow aAa, A \rightarrow sblbcc, C \rightarrow abb, E \rightarrow ac$

Phase 2 $V_1'' = \{S\}$

$V_2'' = \{S, A\}$

$V_3'' = \{S, A, C\}$

$V_4'' = \{S, A, C\} = V_3''$

$G'' = \{S, A, C\}, \{a, b\}, P'', \{S\}$

where $P'' = S \rightarrow aAa, A \rightarrow sblbcc, C \rightarrow abb$

g.

(i) $S \rightarrow ASB | \epsilon$

$A \rightarrow aASla$

$B = bb | A | sbc$

Remove ϵ production \rightarrow

$S \rightarrow ASB$

$A \rightarrow aASla$

$B \rightarrow bb | A | sbc$

Remove Unit production \Rightarrow

$$\begin{aligned} S &\rightarrow ASB \\ A &\rightarrow aAS \mid a \\ B &\rightarrow b \mid bSB \end{aligned}$$

Remove Useless symbols

There is no useless symbols

$$\begin{aligned} S &\rightarrow ASB & \text{Let } P \rightarrow AS \\ \Rightarrow S &\rightarrow PB \end{aligned}$$

$$\begin{aligned} A &\rightarrow aAS & \text{Let } Q \rightarrow aA \\ \Rightarrow A &\rightarrow QS & \text{Let } X \rightarrow a \\ && Q \rightarrow XA \\ B &\rightarrow bb \\ \Rightarrow B &\rightarrow YY & \text{Let } Y \rightarrow b \end{aligned}$$

$$\begin{aligned} B &\rightarrow SBS \\ B &\rightarrow SYS \\ \Rightarrow B &\rightarrow RS \end{aligned} \quad \text{Let } R \rightarrow SY$$

So given grammar into CNF is

$$\begin{aligned} S &\rightarrow PB, P \rightarrow AS, A \rightarrow QS, Q \rightarrow aA, X \rightarrow a, Q \rightarrow XA, \\ B &\rightarrow YY, Y \rightarrow b, B \rightarrow RS, R \rightarrow SY \end{aligned}$$

$$(ii) S \rightarrow 0A01IB1IBB$$

$$A \rightarrow C$$

$$B \rightarrow S1A$$

$$C \rightarrow S1\varepsilon$$

Remove ε -productions \rightarrow

$$S \rightarrow 0010A0111IB1IBB$$

$$A \rightarrow C$$

$$B \rightarrow S1A$$

$$C \rightarrow S$$

Remove unit productions \rightarrow

$$S \rightarrow 0010A01IB1IBB$$

$$A \rightarrow 0010A01IB1IBB$$

$$B \rightarrow 0010A01IB1IBB$$

$$C \rightarrow 0010A01IB1IBB$$

Remove Useless Symbols \rightarrow

Remove, A, B, C \rightarrow

$$S \rightarrow 0010\cancel{S}01IB1IBB$$

$$S \rightarrow 00\cancel{R}$$

Let $R \rightarrow 0$

$$\Rightarrow S \rightarrow PR$$

$$S \rightarrow 0\cancel{P}0$$

$$S \rightarrow \cancel{P}PS$$

Let $P \rightarrow PS$

$$\Rightarrow S \rightarrow PQ$$

$$S \rightarrow ISI$$

Let $R \rightarrow I$

$$S \rightarrow RSR$$

Let $T \rightarrow SR$

$$\Rightarrow S \rightarrow RT$$

CNF grammar is →

$$S \rightarrow A A A | A$$

$$S \rightarrow P P | P Q | R T | S S$$

$$P \rightarrow O$$

$$Q \rightarrow P S$$

$$R \rightarrow I$$

$$T \rightarrow S R$$

$$(iii) \quad S \rightarrow A A A | B,$$

$$A \rightarrow a A | B,$$

$$B \rightarrow \epsilon$$

Remove ϵ -productions →

$$S \rightarrow A A A | B | \epsilon$$

$$A \rightarrow a A | B | \epsilon$$

↓

$$S \rightarrow A A I A | A A A | B | \epsilon$$

$$A \rightarrow a A | a | B$$

↓

$$S \rightarrow A A I A | A A A | B$$

$$A \rightarrow a A | a | B$$

Remove unit productions →

$$S \rightarrow A A A | A A I a | a$$

$$A \rightarrow a A | a$$

CNF →

$$S \rightarrow S_1 A | A A | A, A | a$$

$$S_1 \rightarrow A A$$

$$A_1 \rightarrow a$$

$$A \rightarrow A, A | a$$

(iv) $S \rightarrow aAa|bBb|\epsilon$

$A \rightarrow c|a$

$B \rightarrow c|b$

$C \rightarrow CDE|\epsilon$

$D \rightarrow B|A|ab$

Removing ϵ -production \rightarrow

$S \rightarrow aAa|bBb$

$A \rightarrow c|a|\epsilon$

$B \rightarrow c|b|\epsilon$

$C \rightarrow CDE|DE$

$D \rightarrow B|A|ab$

\downarrow

$S \rightarrow aAa|aa|bBb|bb$

$A \rightarrow c|a$

$B \rightarrow c|b$

$C \rightarrow CDE|DE$

$D \rightarrow A|B|ab|\epsilon$

\downarrow

$S \rightarrow aAa|bBb|aa|bb$

$A \rightarrow c|a$

$B \rightarrow c|b$

$C \rightarrow CDE|DE|CE|E$

$D \rightarrow A|B|ab$

Removing Unit Production \rightarrow

$S \rightarrow aAa|bBb|aa|bb$

$A \rightarrow CDE|DE|CE|a$

$B \rightarrow CDE|DE|CE|b$

$C \rightarrow CDE|DE|CE$

$D \rightarrow CDE | DE | CE | aabbab$

Removing useless production :-

Since E is no variable with any production so all productions with E are useless.

$S \rightarrow aAa | bBb | aab | bb$

$A \rightarrow a$

$B \rightarrow b$

$D \rightarrow aabbab$

D cannot be reached from S hence it is also useless

$S \rightarrow aAa | bBb | aab | bb$

$A \rightarrow a$

$B \rightarrow b$

Converting to CNF -

$S \rightarrow S_1 A_1, S_2 B_1 | A_1 A_1 | B_1 B_1,$

$S_1 \rightarrow A_1 A_1$,

$S_2 \rightarrow B_1 B_1$

$A_1 \rightarrow a$

$B_1 \rightarrow b$

10. (i) $S \rightarrow aSb | bSA | aab$

GNF Conversion
 $S \rightarrow aSB | bSA | aab$

$B \rightarrow b$

$A \rightarrow a$

(ii) $S \rightarrow aSbab$

GNF \rightarrow

$S \rightarrow aSB1aB$

$B \rightarrow a$

(iii) $S \rightarrow ab1aslaas$

GNF \rightarrow

$S \rightarrow aB1asl1aAS$

$B \rightarrow b$

$A \rightarrow a$

(iv) $S \rightarrow ABb1a$

$A \rightarrow aaA1B$

$B \rightarrow bAB$

Removing unit production:-

$S \rightarrow ABb1a$

$A \rightarrow aaA1bAB$

$B \rightarrow bAB$

GNF \rightarrow

$S \rightarrow aA, ABB, 1bABABB, 1a$

$A \rightarrow aA, A1bAB$

$B \rightarrow bAB$

$A_1 \rightarrow a$

$B_1 \rightarrow b$