given win
$$c(x,u)$$
 s.t. $x(s+i) = F(s,x(s),u(s))$
(DT-OCP) $c(x,u) = l(t,x(t)) + \sum_{s=0}^{t-1} z(s,x(s),u(s))$

there are 2 main interpretations/strategies

Show that (DT-OCP) can be rewritten as
$$\min_{x \in \mathbb{R}^m} \overline{c(x,u)} \cdot s.t. \ \overline{f(x,u)} = 0$$
 $\overline{u} \in \mathbb{R}^m$

(i.e. determine $\overline{u}, \overline{x}, \overline{c}, \overline{f}$)

- recalling that $(u: [o,t] \to \mathbb{R}^m) \in (\mathbb{R}^m)$,

let $\overline{u} \in \mathbb{R}^m$ be $[\overline{u}]_s = u(s)$,

and similarly $(\overline{x})_s = x^u(s)$

where x^u is the trig generated by u

- then $\overline{c}(\overline{x}, \overline{u}) = c(x, u)$
 $\overline{f(x, u)} = 0 \in x(s+i) = \overline{f(s, x(s), u(s))}$
 $\overline{c(x, u)} = c(x, u)$
 $\overline{c(x, u)} = c(x, u)$

there is a "natural" correspondence between discrete-time signals 3: [0,t) -> 1Rd and "stacked"/block vectors

$$3 = \begin{bmatrix} 3(0) \\ 3(1) \\ 3(2) \\ \vdots \\ 3(t-1) \end{bmatrix}$$

 $\vec{\beta} = \left[3(0), 3(1), \dots, 3(t-1)\right] \in \mathbb{R}^{d \times t}$ $\{3: [0,t) \rightarrow \mathbb{R}^d\} = 2 \xrightarrow{\psi} Z = \{3 \in \mathbb{R}^{t \cdot d}\}$

1°. dynamic programming problems/ 2°. trajectory optimization interp: output will be an optimal control policy $u^*: [0,t) \times X \rightarrow U$ that tells us the optimal action u*(t,x(t)) for every time t and state X(t)

· output will be an optimal central signal ext [o,t) -> U that tells us the optimal action ux(t) brevery time - from initial state x(o)=Xo

time to and state X(t)

stat: Solve a segunce of parametaged

NIP backnown in time

time to from initial state x(0)=X0

o solve one NLP "forward"

in time

* since these are working from some underlying problem, thegre consistent.

 $u^*(\tau, x^*(\tau)) = u^*_{x_0}(\tau)$ assuming $x^*(s) = x_0$ and $u^*_{x_0}$ is applied while τ

