

# IDT: Intelligent Data Placement for Multi-tiered Main Memory with Reinforcement Learning

Juneseo Chang<sup>†</sup>, Wanju Doh<sup>†</sup>, Yaebin Moon<sup>‡</sup>,  
Eojin Lee<sup>§</sup>, and Jung Ho Ahn<sup>†</sup>

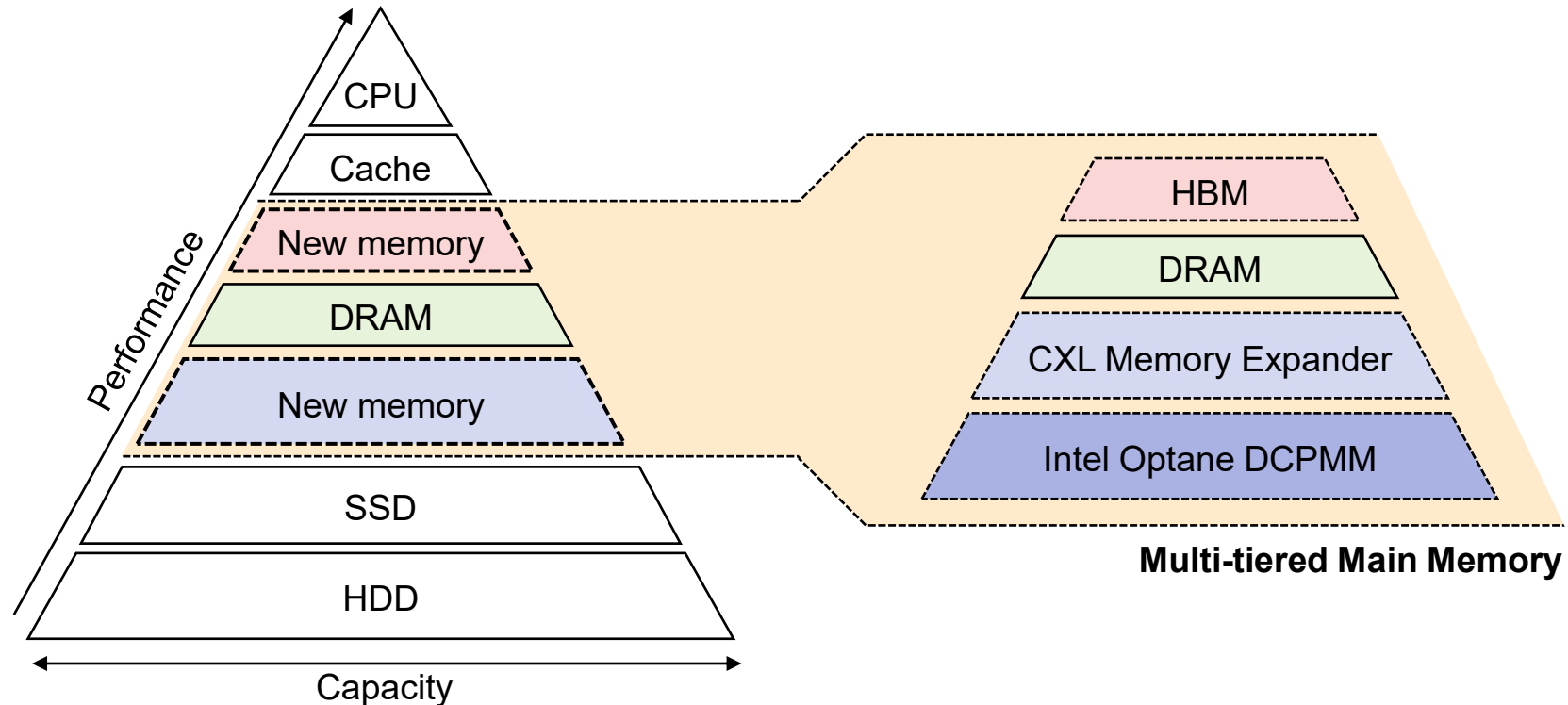
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<sup>‡</sup>This work was done while at Seoul National University

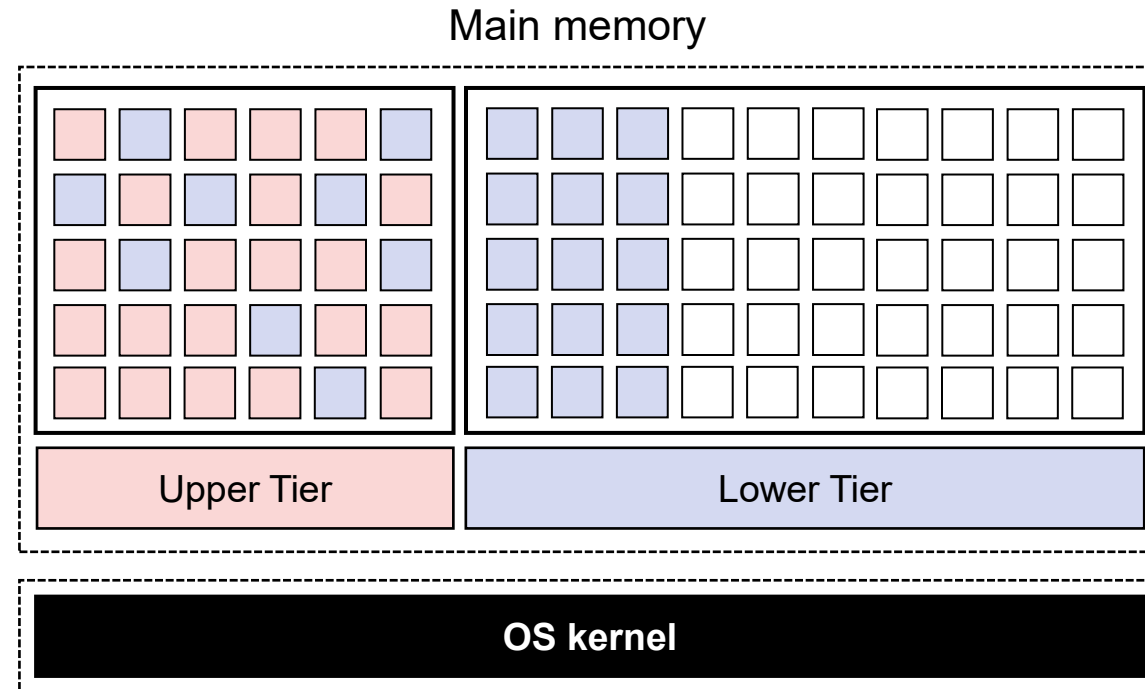
# Tiered Memory Systems

- Emerging memory technologies are introducing **multiple tiers** in the **main memory**
  - CXL Memory, HBM-enabled processors, Intel Optane DCPMM, ...



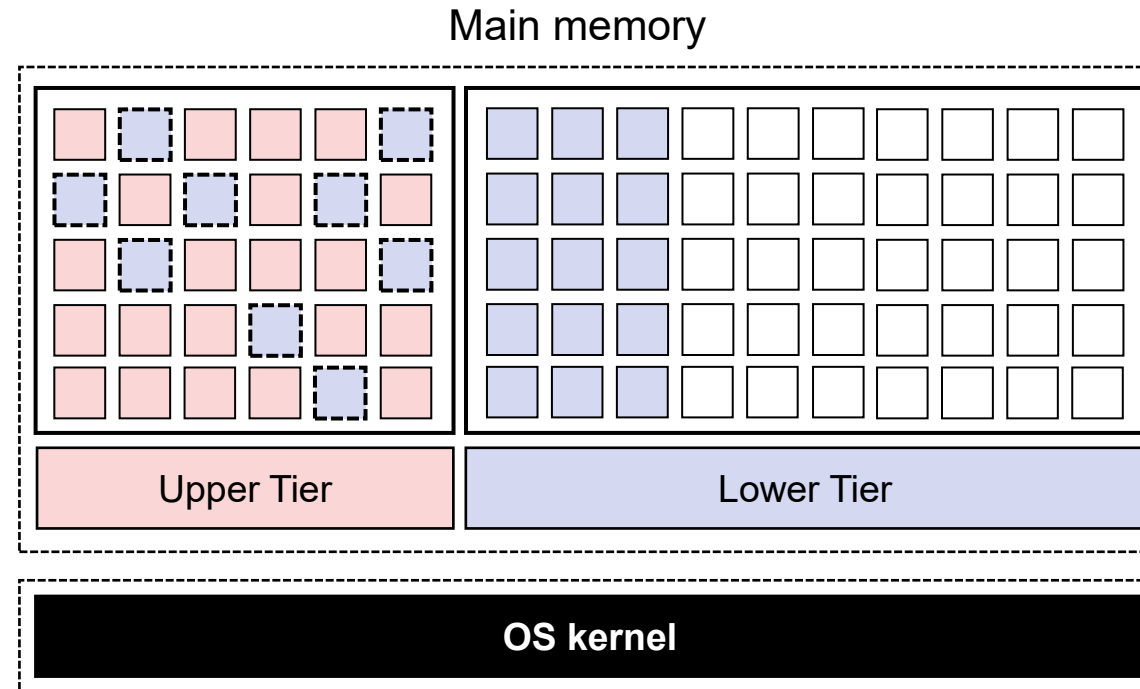
# OS-level Tiered Memory Management

- **OS** kernel manages **data placement** across tiers



# OS-level Tiered Memory Management

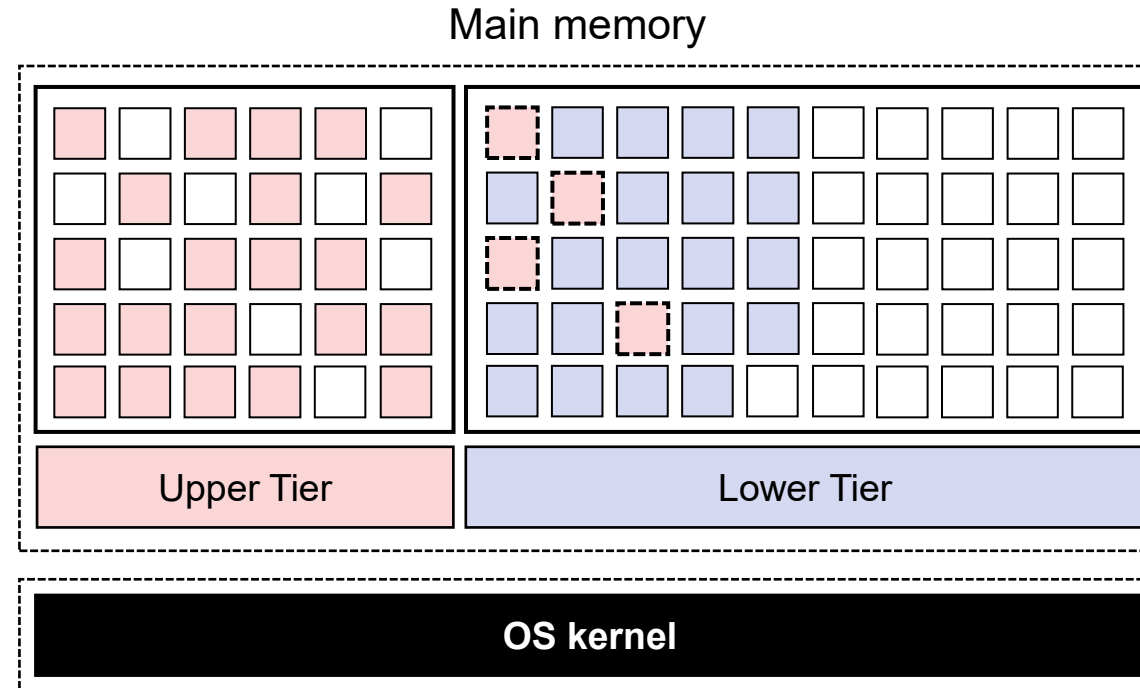
- OS kernel manages **data placement** across tiers
- OS kernel **demotes cold** pages to **lower-tier** memory



Accurately **identifying data hotness** and effective **demotion criteria** are necessary!

# OS-level Tiered Memory Management

- OS kernel manages **data placement** across tiers
- OS kernel **demotes cold** pages to **lower-tier** memory
- OS kernel **promotes hot** pages to **upper-tier** memory



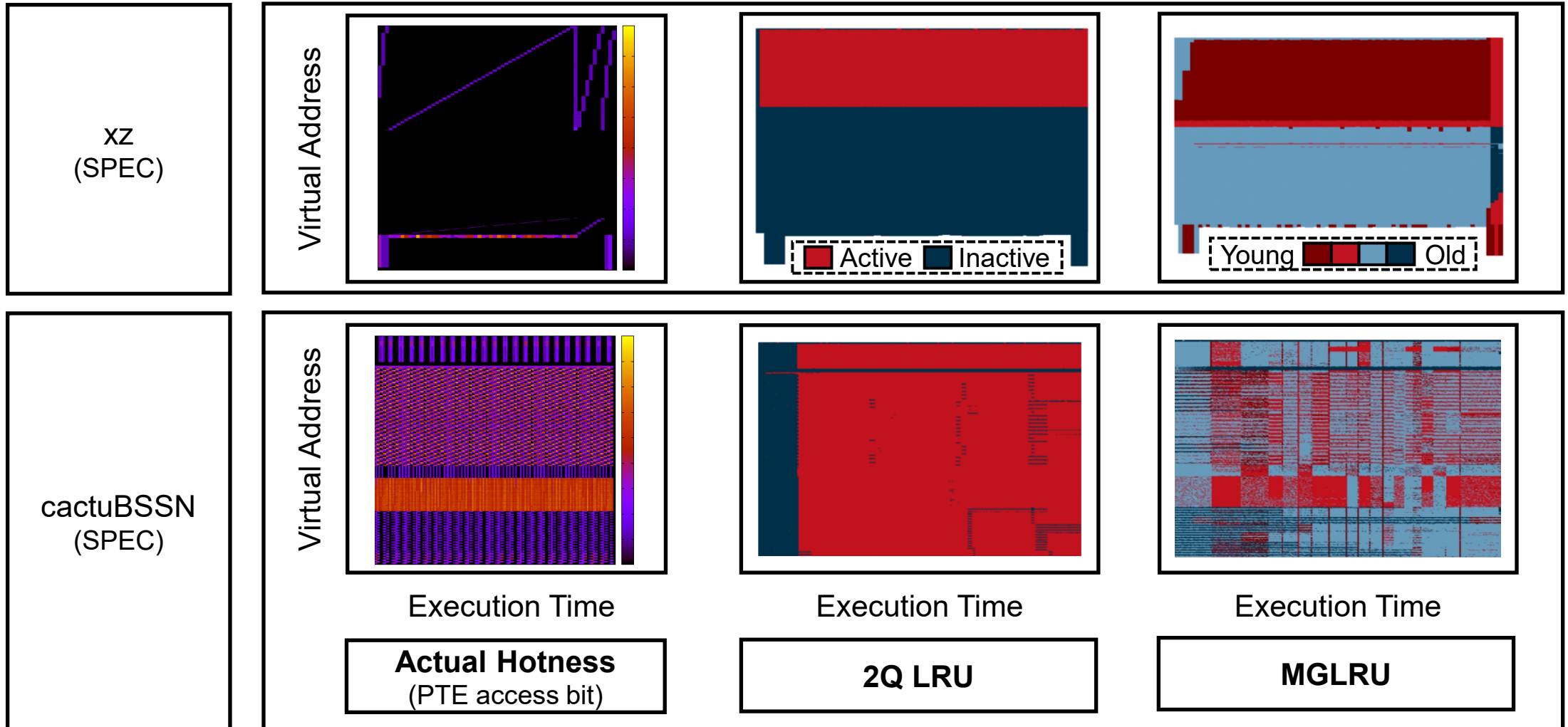
# Selecting Demotion Candidates: 2Q LRU and MGLRU

- Effective demotion candidate selection is crucial
  - Impacts **promotion**
  - Incorrectly identifying demotion targets causes **ping-pong** of demotion and promotion
- Prior works used Linux kernel's **active/inactive LRU lists (2Q LRU)**
  - Since 2022, **multi-generational LRU lists**<sup>[1]</sup> (**MGLRU**) for more fine-grained policy

[1] Yu Zhao. 2022. Multigenerational LRU Framework. <https://lwn.net/Articles/880393/>.

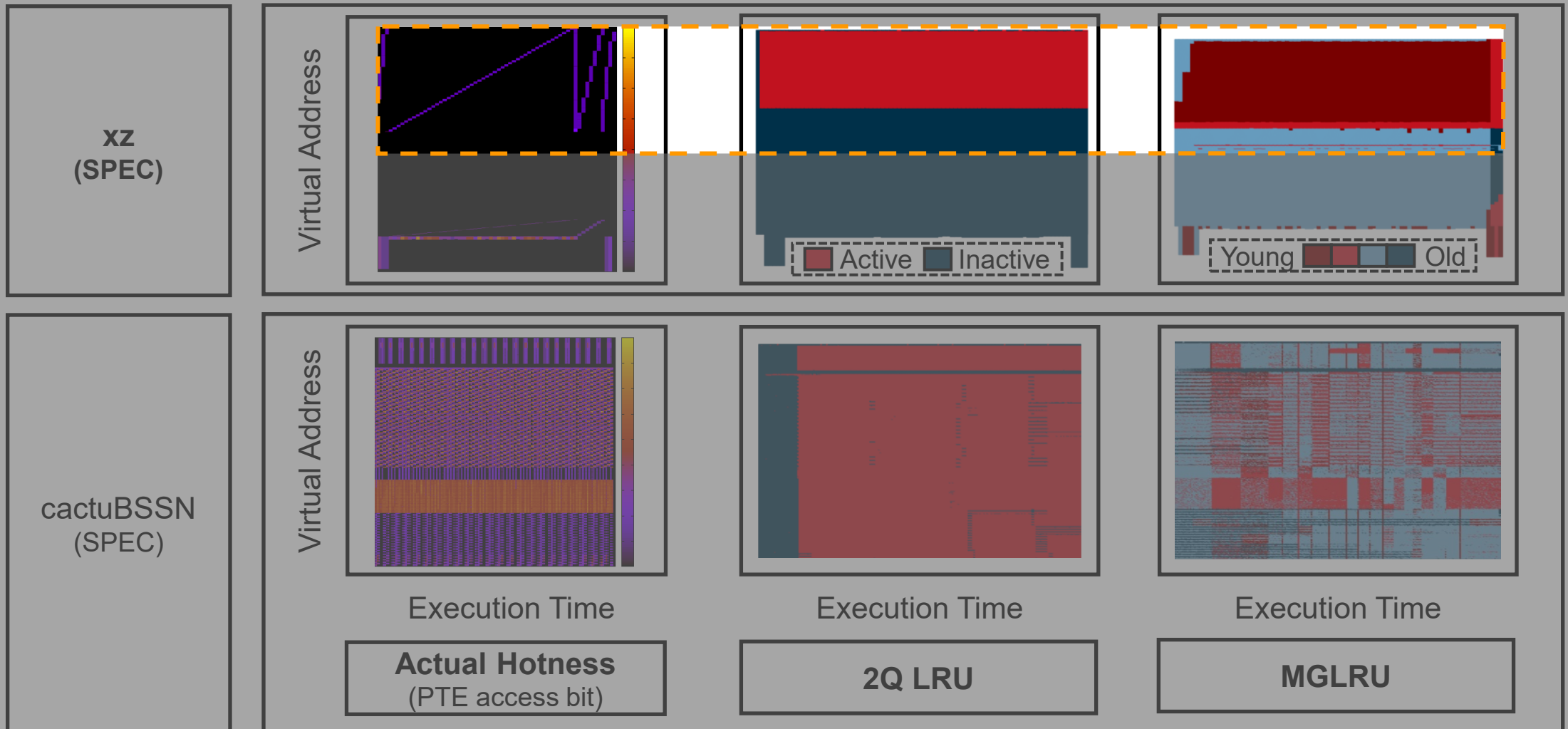
# Selecting Demotion Candidates: 2Q LRU and MGLRU

- However, **2Q LRU** and **MGLRU** often deviate from the **actual data hotness** (PTE access bit scanning)



# Selecting Demotion Candidates: 2Q LRU and MGLRU

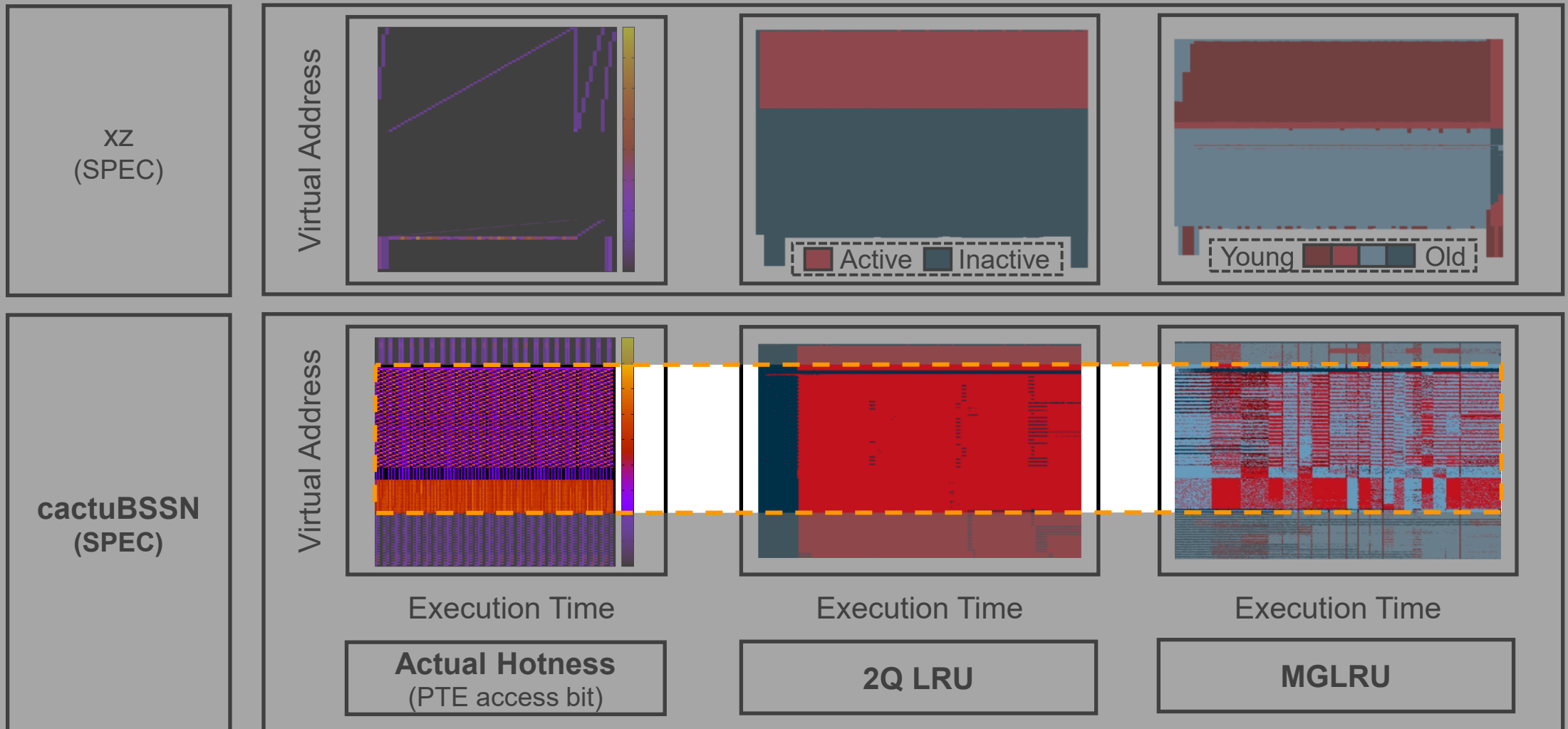
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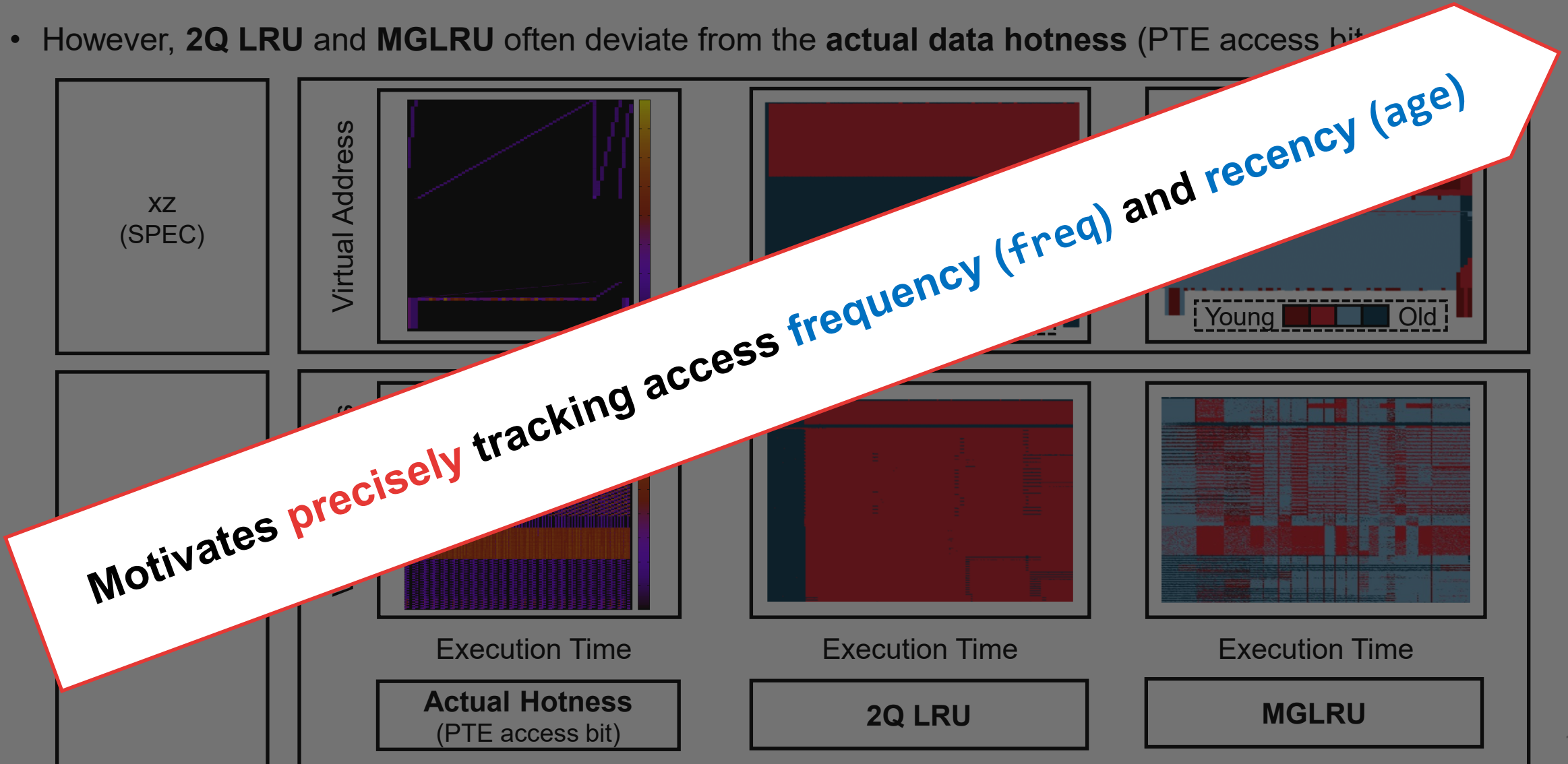
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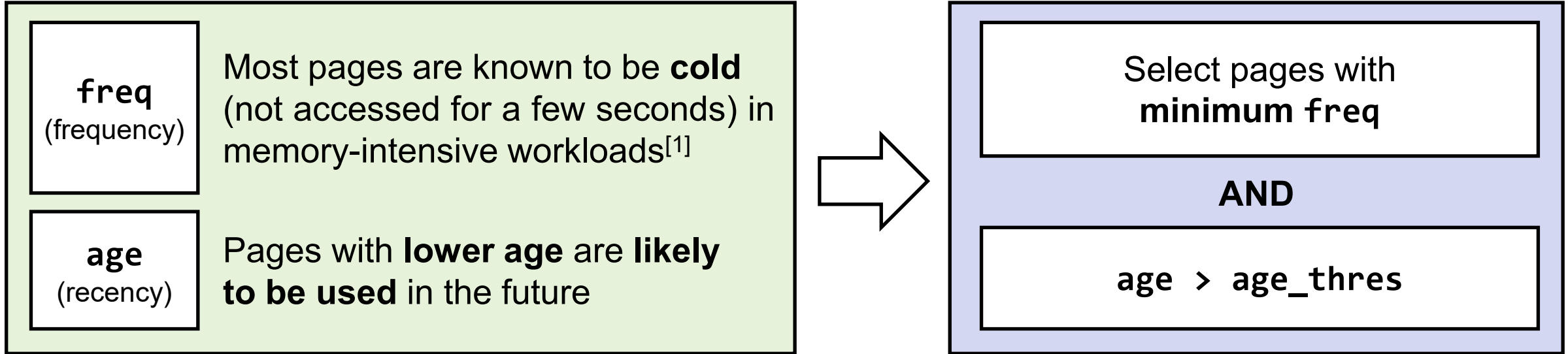


## Selecting Demotion Candidates: 2Q LRU and MGLRU

- However, **2Q LRU** and **MGLRU** often deviate from the **actual data hotness** (PTE access bit)



## Selecting Demotion Candidates: Using more precise standards



## Selecting Demotion Candidates: Using more precise standards

freq



Minimum freq

AND

age

age > age\_thres

## Selecting Demotion Candidates: Using more precise standards

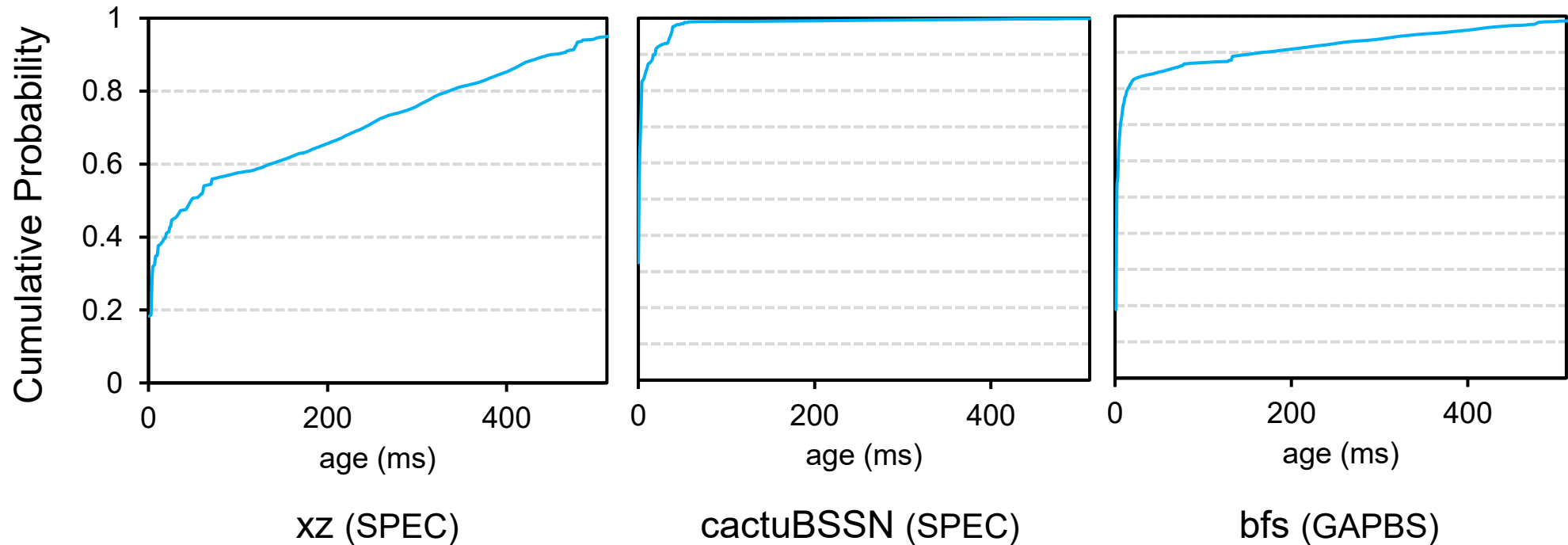
freq

Minimum freq

AND

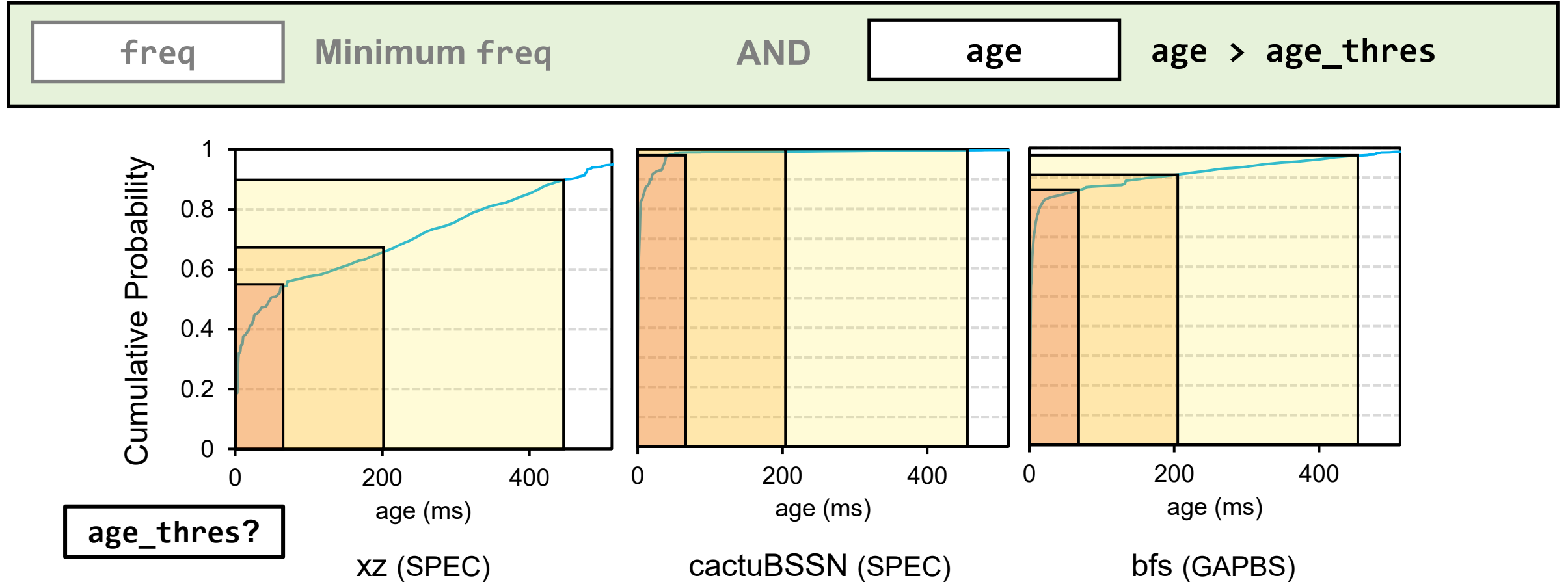
age

age > age\_thres



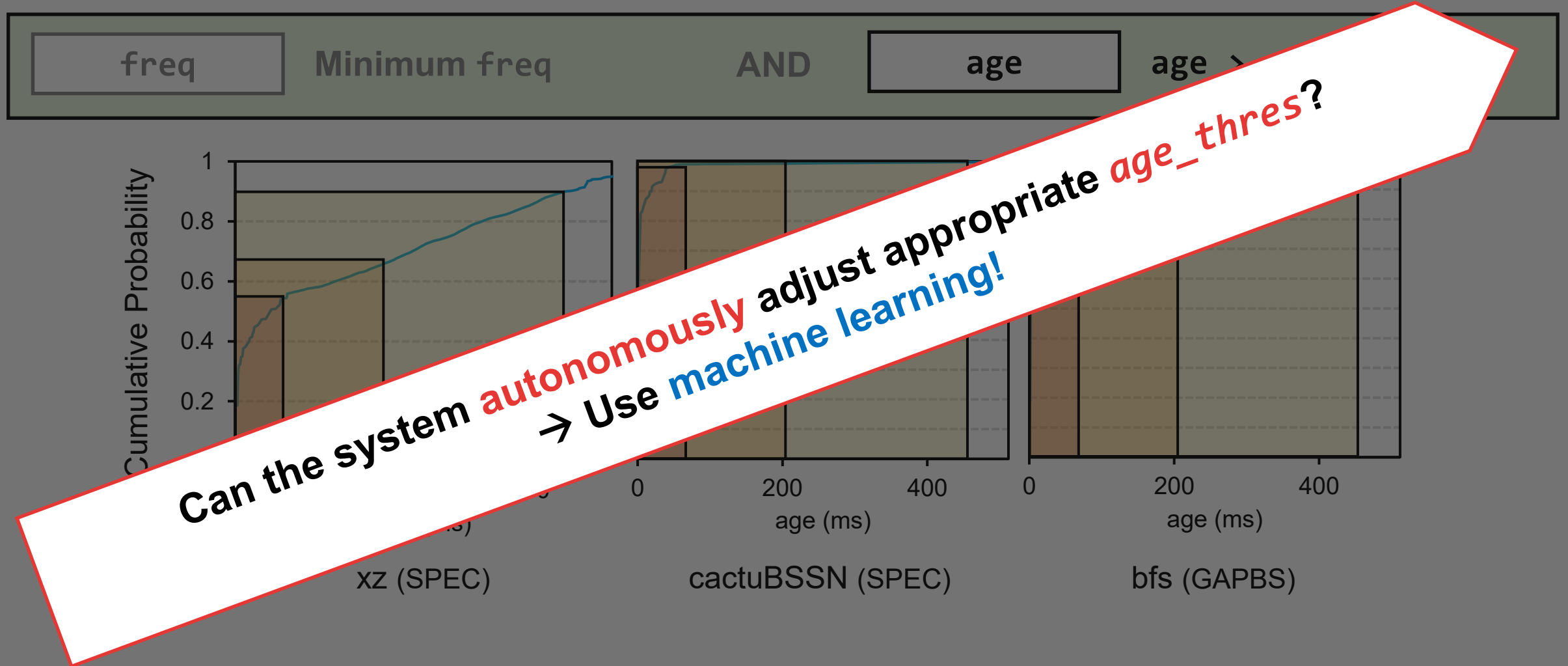
Cumulative probability distribution of accessed page's age varies across workloads

# Selecting Demotion Candidates: Using more precise standards



Cumulative probability distribution of accessed page's age varies across workloads

## Selecting Demotion Candidates: Using more precise standards



Cumulative probability distribution of accessed page's age varies across workloads

# ML for Demotion Policy

## Lightweight

Prior **supervised learning** approaches have high **execution time overhead** and **memory usage**

## Adaptability

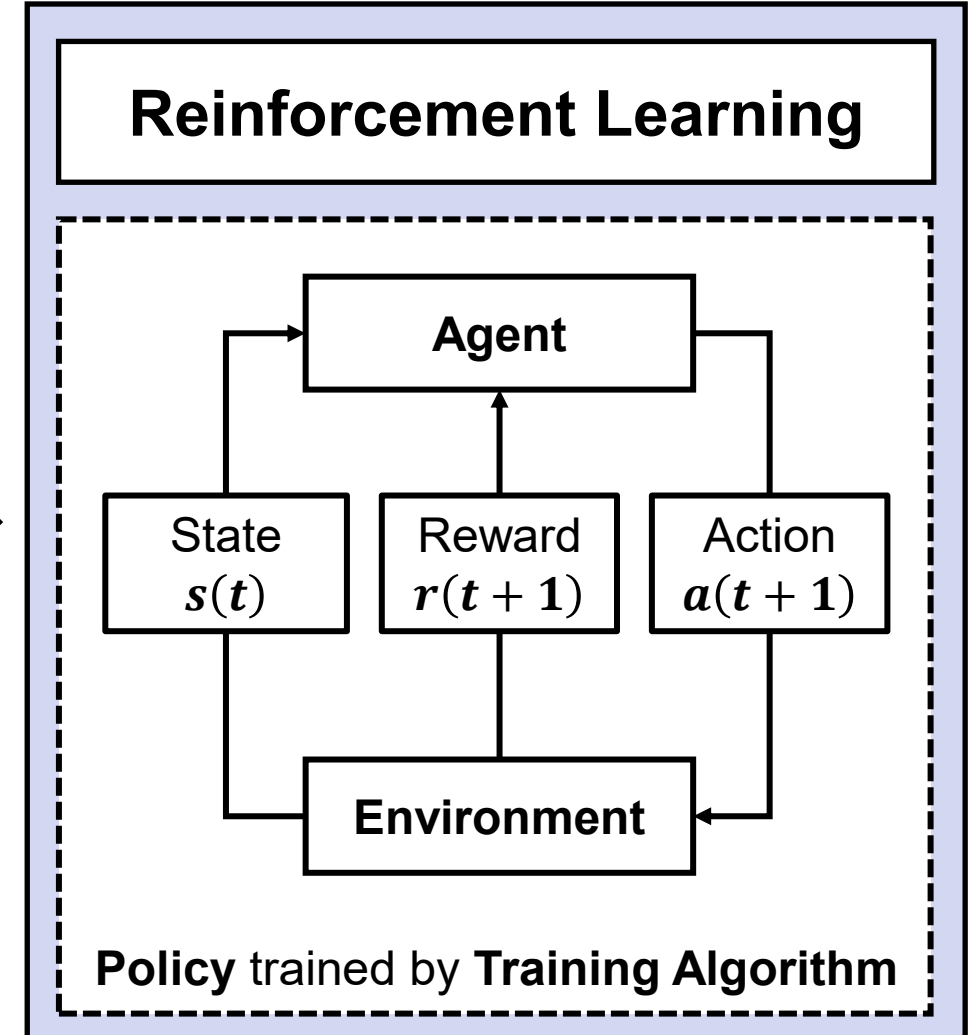
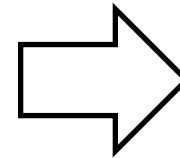
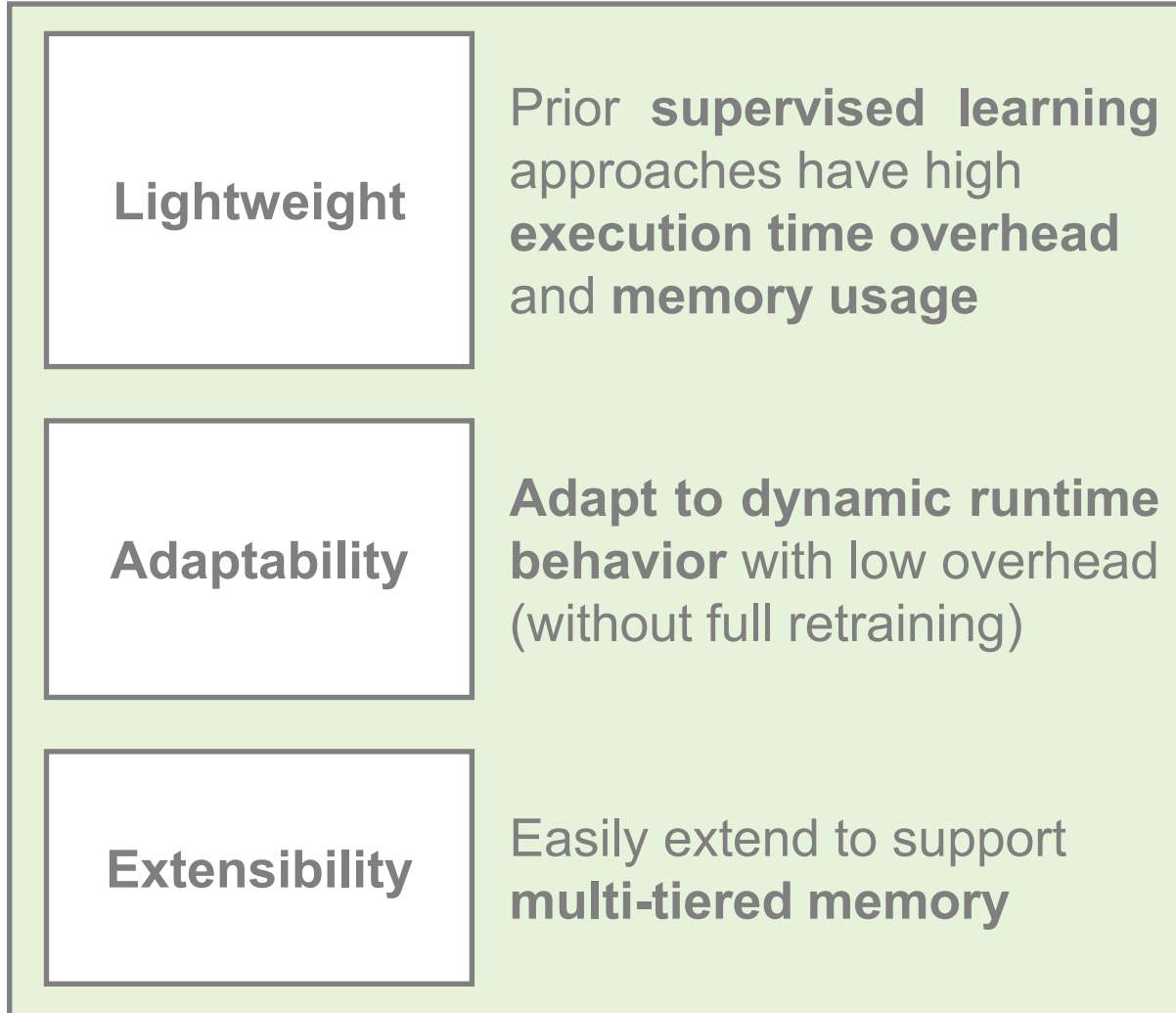
**Adapt to dynamic runtime behavior** with low overhead (without full retraining)

## Extensibility

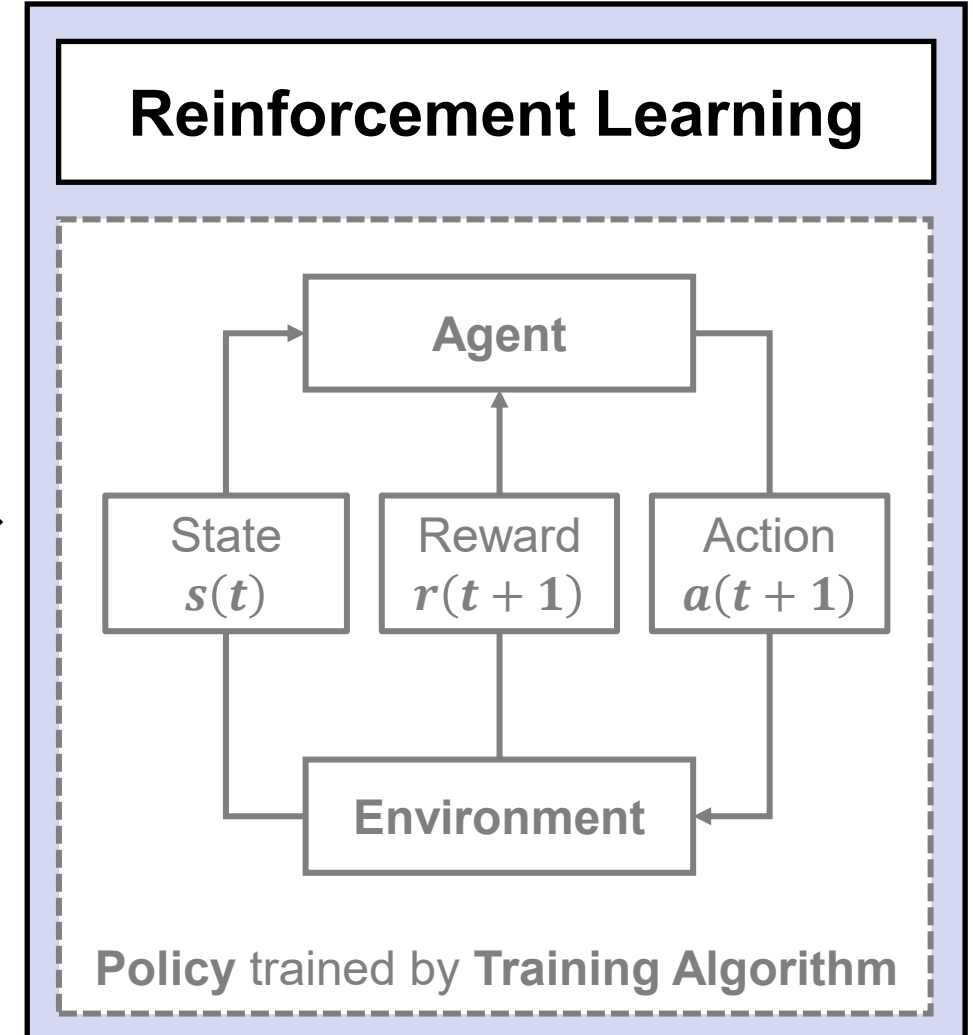
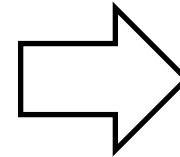
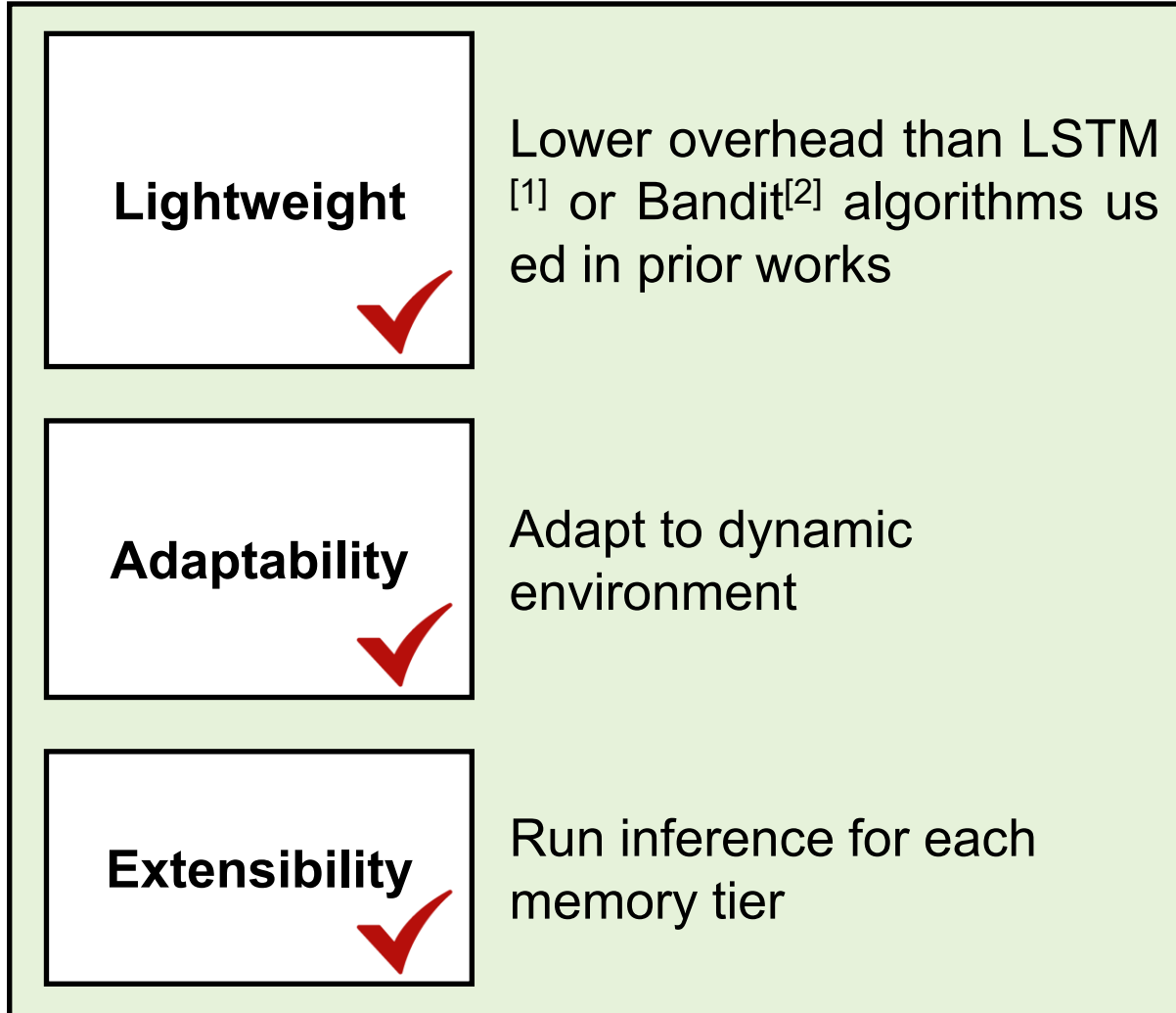
Easily extend to support **multi-tiered memory**



# ML for Demotion Policy



# ML for Demotion Policy



[1] Thaleia Dimitra Doudali et al., "Kleio: A Hybrid Memory Page Scheduler with Machine Intelligence," HPDC, 2019

[2] Andres Lagar-Cavilla et al., "Software-Defined Far Memory in Warehouse-Scale Computers," ASPLOS, 2019

# ML for Demotion Policy

Lightweight



Lower overhead than LSTM<sup>[1]</sup> or Bandit<sup>[2]</sup> algorithms used in prior works

Adaptability

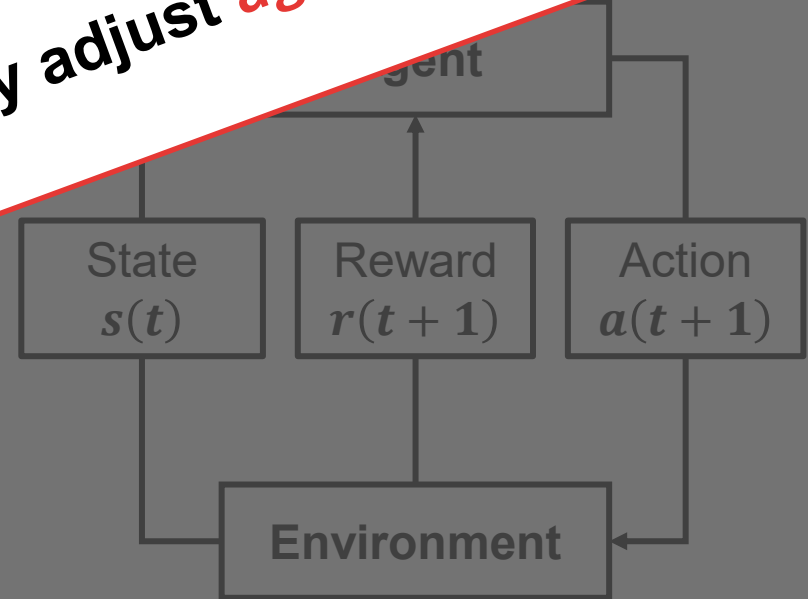
Adaptability

Reinforcement Learning (RL) can effectively adjust **age\_thres!**

Run inference for each memory tier



Reinforcement

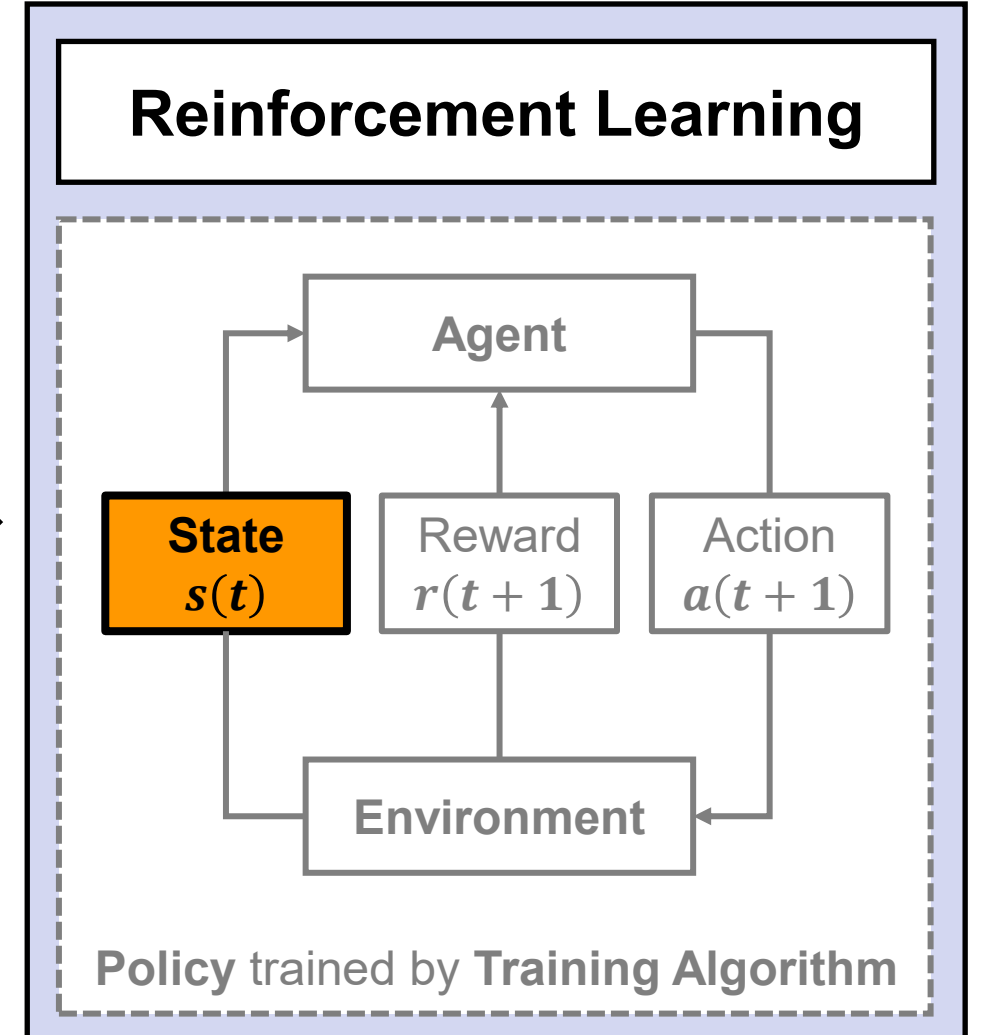
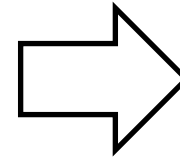
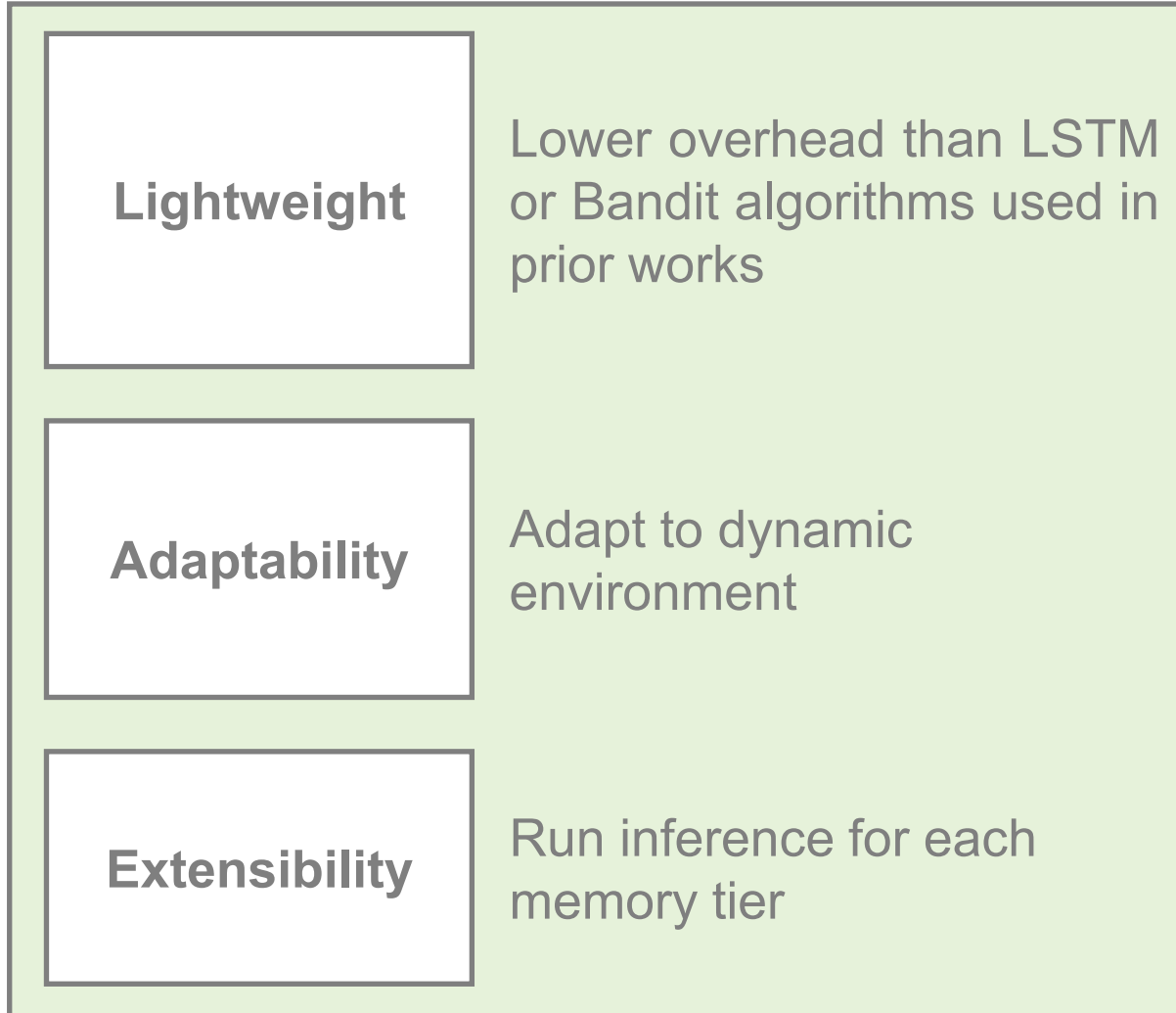


Policy trained by Training Algorithm

[1] Thaleia Dimitra Doudali et al., "Kleio: A Hybrid Memory Page Scheduler with Machine Intelligence," HPDC, 2019

[2] Andres Lagar-Cavilla et al., "Software-Defined Far Memory in Warehouse-Scale Computers," ASPLOS, 2019

# ML for Demotion Policy



## ML for Demotion Policy

**State:** Memory access information of the system

**Page granularity** memory access monitoring has a high **overhead**

→ Group **similar** pages with **region-granularity monitoring**

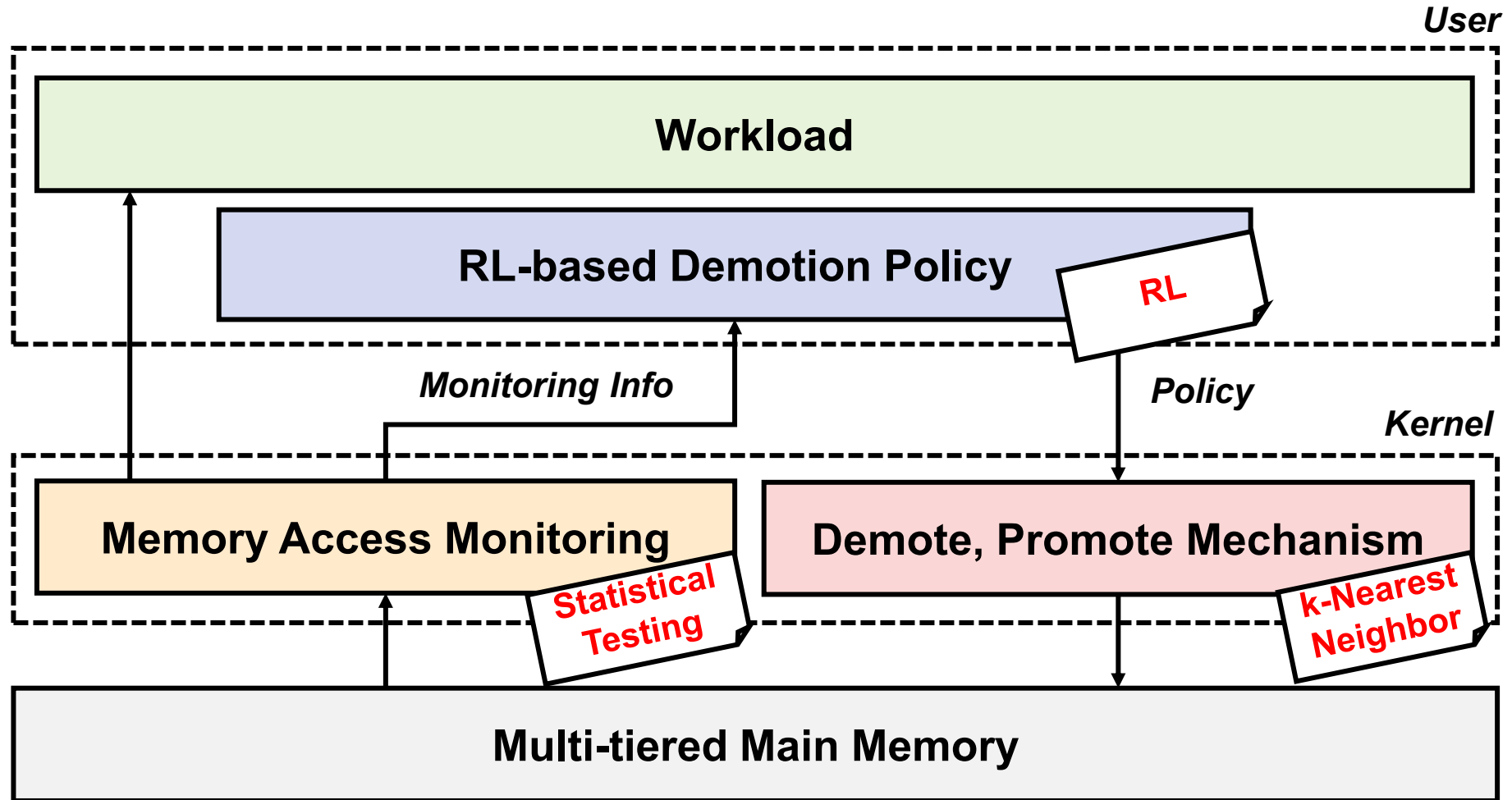
Extensibility

Run inference for each  
memory tier

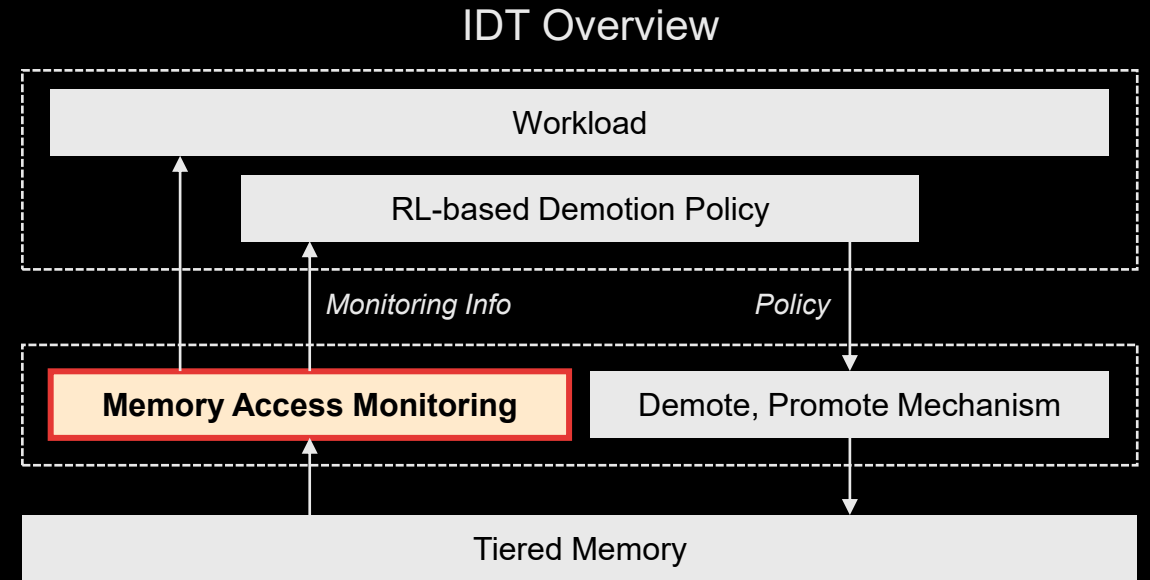
Policy trained by Training Algorithm

## **IDT:** Design and Implementation

## IDT: Overview



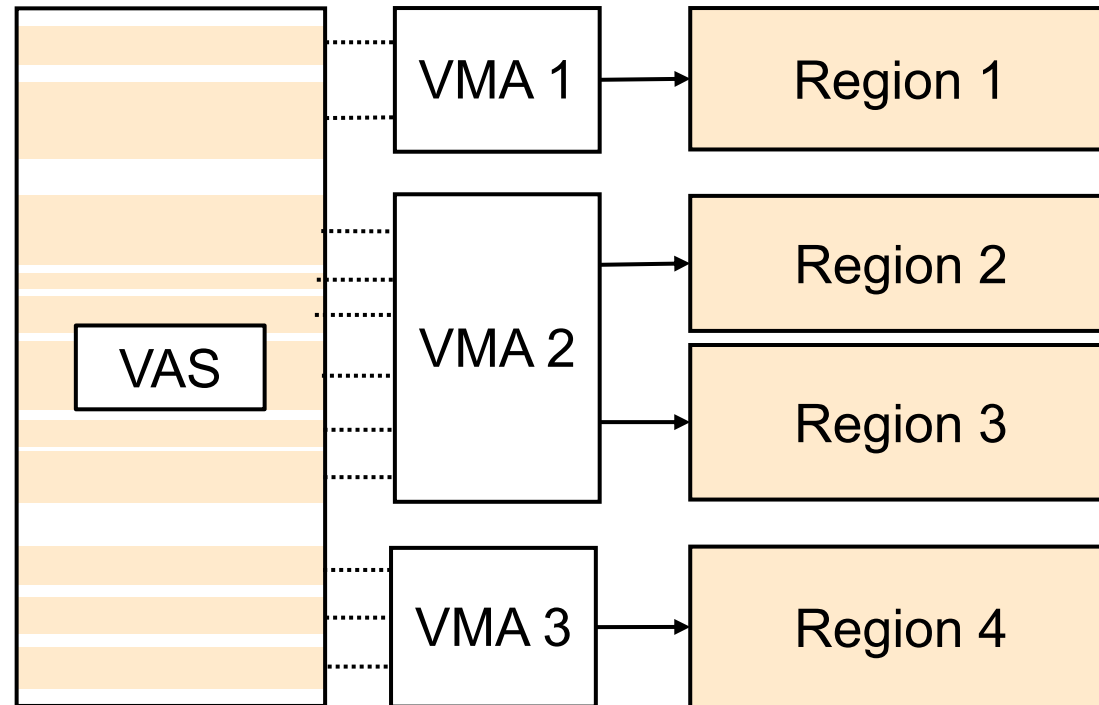
# Memory Access Monitoring





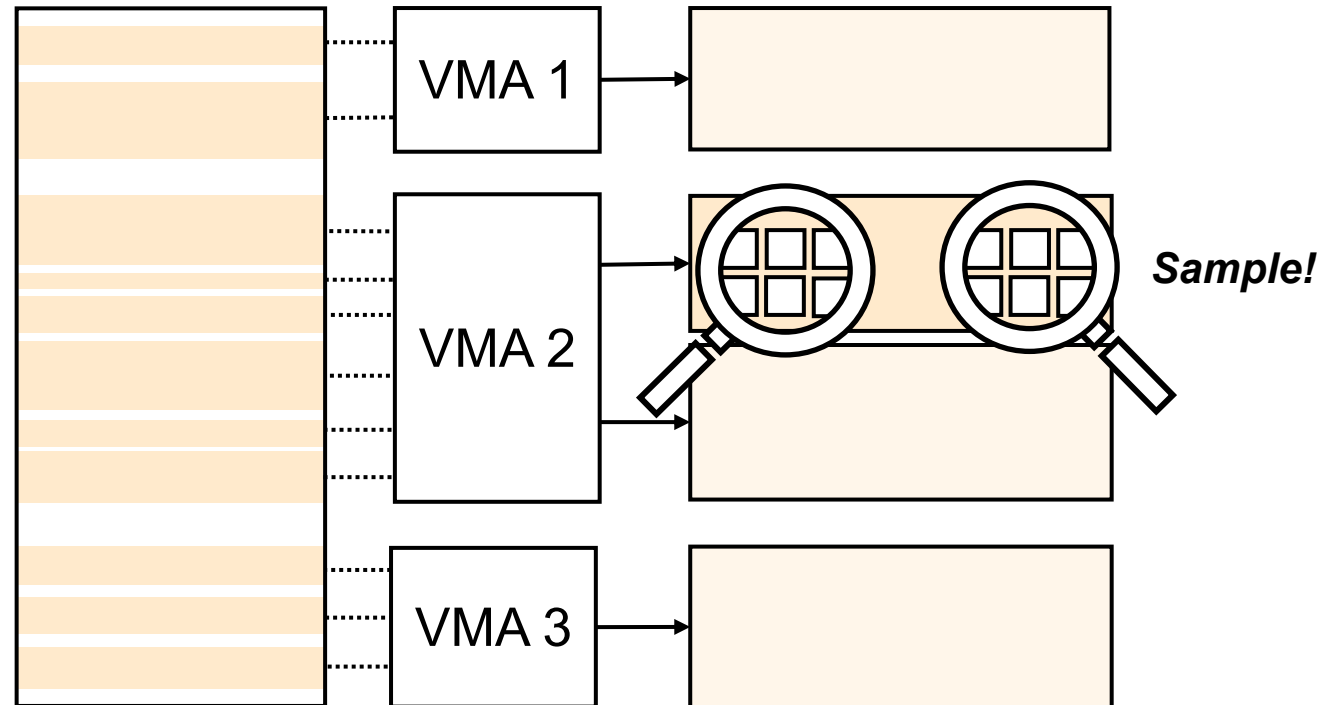
# Region-granularity Monitoring

- Monitor **group** of pages with **similar access patterns**
  - **Partition** Virtual Memory Area (VMA) into **regions**



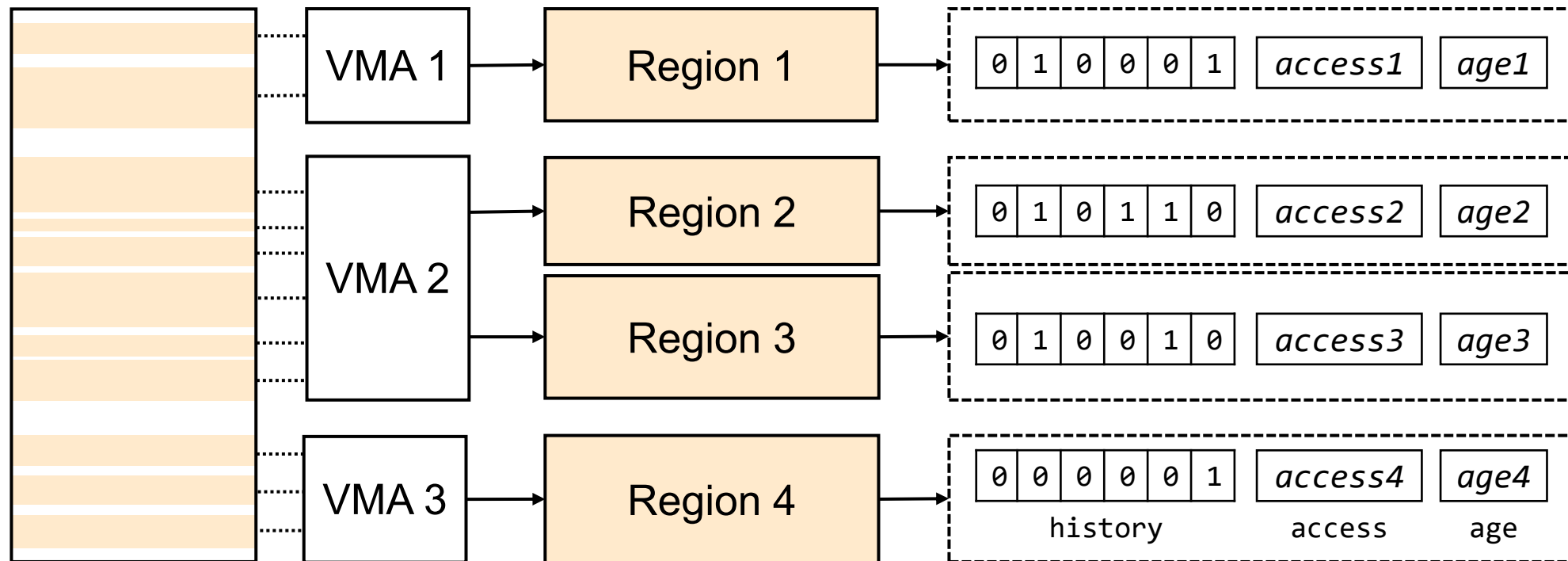
# Region-granularity Monitoring

- Monitor **group** of pages with **similar access patterns**
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- **Sample 2 pages** at each **sample\_interval**



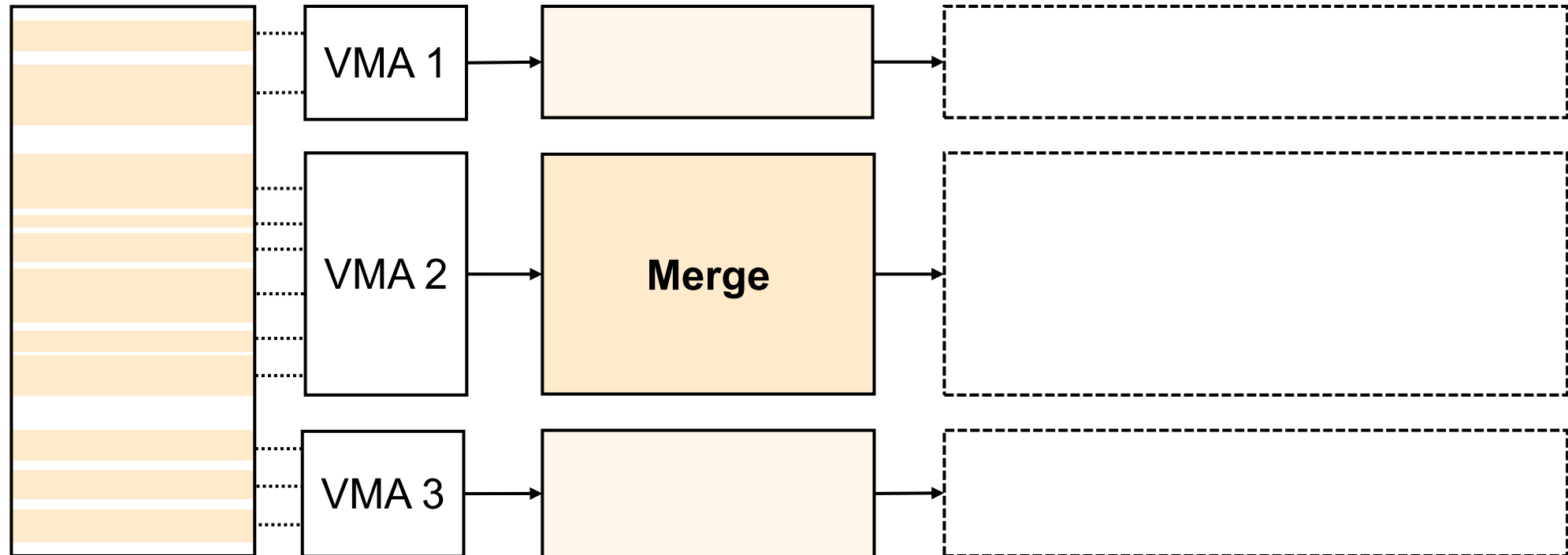
# Region-granularity Monitoring

- Monitor **group** of pages with **similar access patterns**
  - Partition Virtual Memory Area (VMA) into **regions**
- **Sample 2 pages** at each **sample\_interval**
  - Manage **history**, **access**, **age**<sup>[1]</sup>



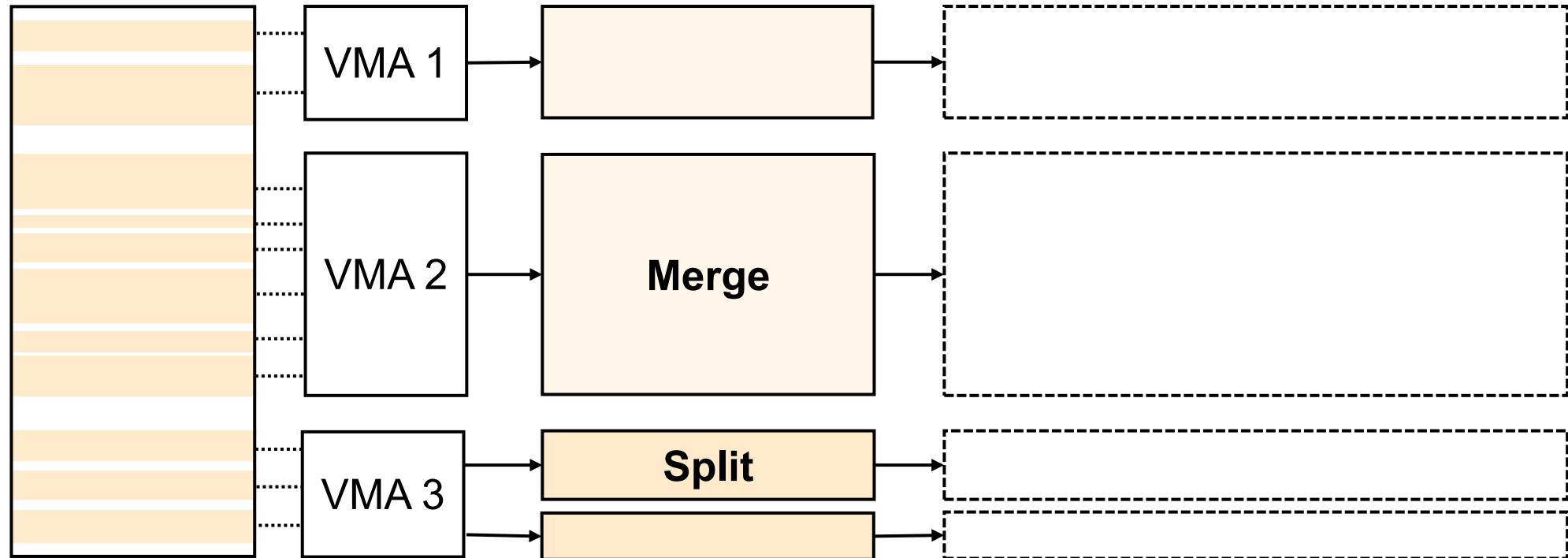
# Region Reconfiguration

- **Merge** or **split** adjacent regions for reconfiguration at each `aggregate_interval`
  - **Merge** regions with **similar access patterns** to **reduce** monitoring overhead



# Region Reconfiguration

- **Merge** or **split** adjacent regions for reconfiguration at each `aggregate_interval`
  - **Merge** regions with **similar access patterns** to **reduce** monitoring overhead
  - **Split** when pages in a region have **different access patterns**



# Region Reconfiguration

- Merge or split adjacent regions for reconfiguration at each `aggregate_interval`

Assume **similar** pages are **grouped** in the same region

**Sampling page's** information determines the **similarity** of regions

→ **Statistical testing** problem (Infer population similarity with samples)



## Region Reconfiguration: Merge

- Validate the similarity of region's history vector by **Fisher's exact test** with a 90% significance level

	Accessed	Not	Total
Region $i$	$a_i$	$n - a_i$	$n$
Region $(i + 1)$	$a_{i+1}$	$n - a_{i+1}$	$n$

window size =  $n$

$$P_{i,i+1} = \frac{\binom{n}{a_i} \times \binom{n}{a_{i+1}}}{\binom{2n}{a_i + a_{i+1}}}$$

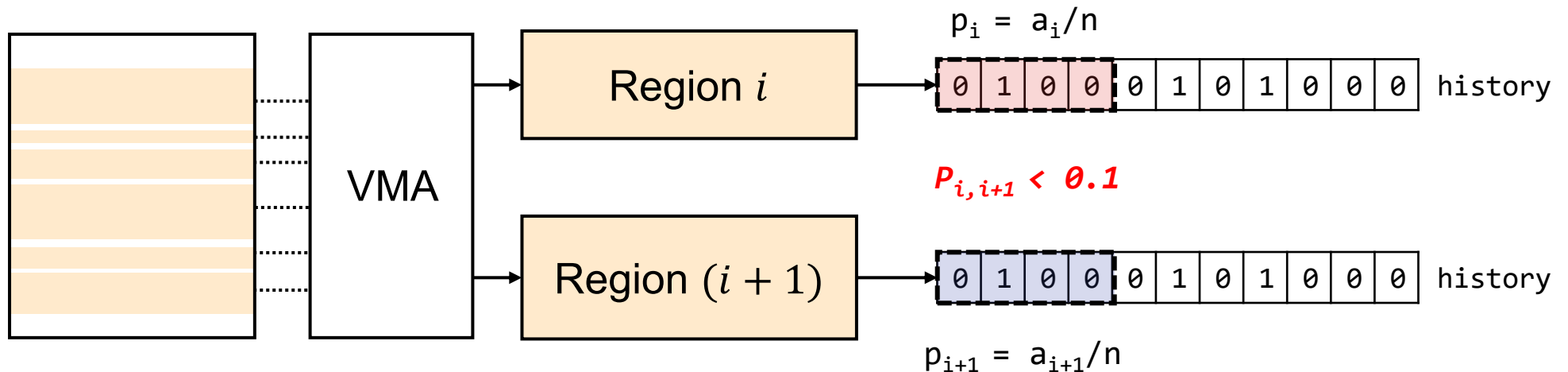
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- **Sliding window** → Compare the **access ratio** of each region's window

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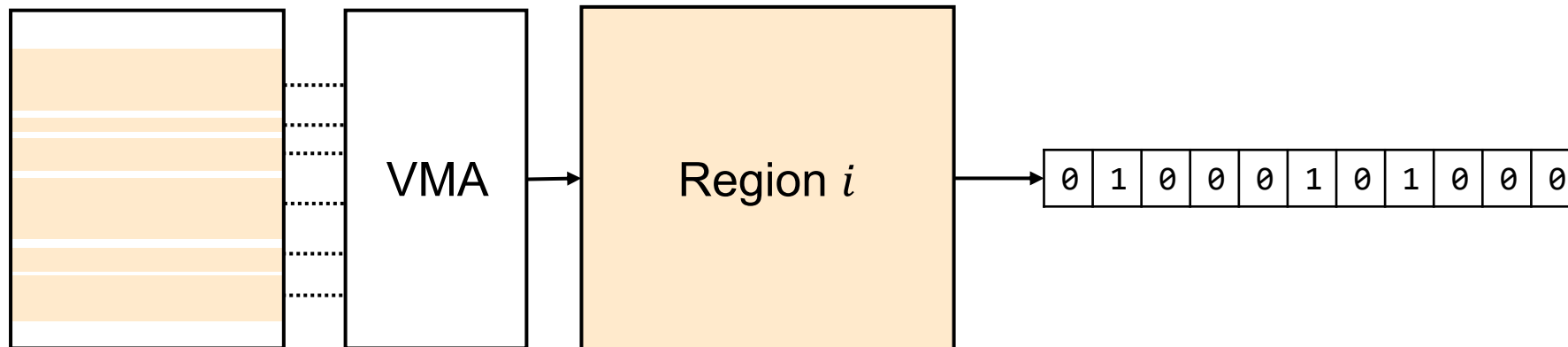
# Region Reconfiguration: Merge

- Validate the similarity of region's history vector by **Fisher's exact test** with a 90% significance level
- **Sliding window** → Compare the **access ratio** of each region's window
  - If every window yields a *similar access ratio* → **Merge**

	Accessed	Not	Total
Region $i$	$a_i$	$n - a_i$	$n$
Region $(i + 1)$	$a_{i+1}$	$n - a_{i+1}$	$n$

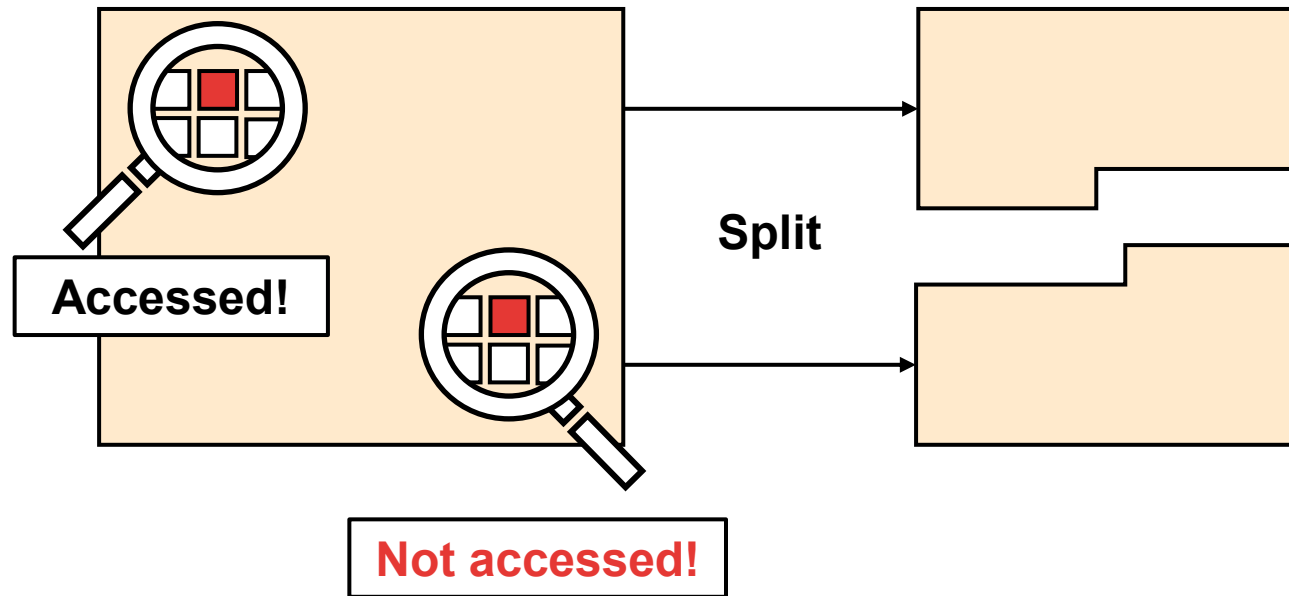
window size =  $n$

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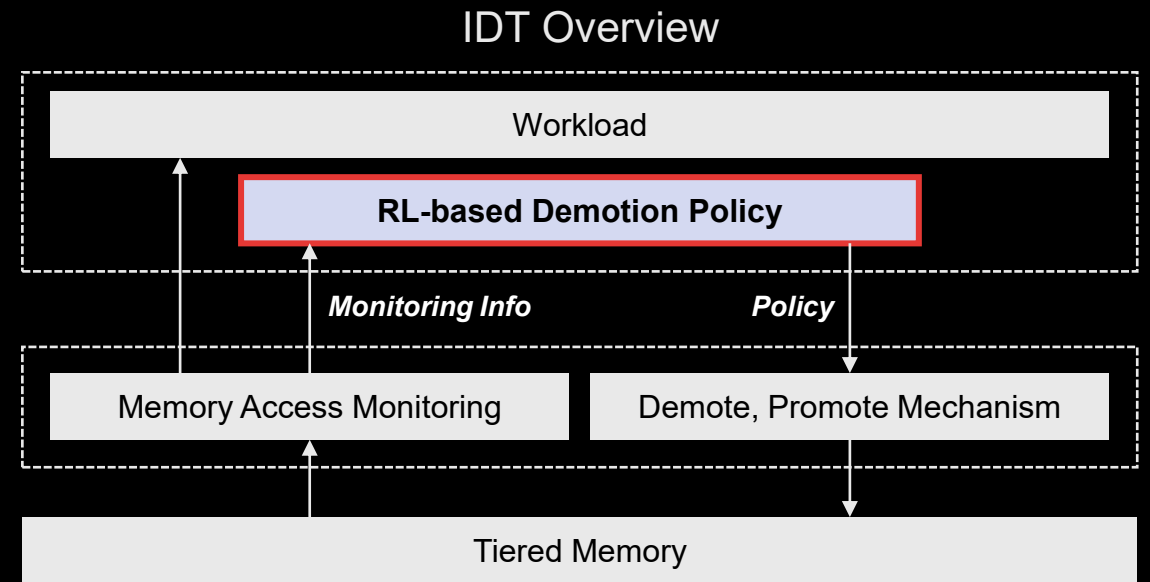


# Region Reconfiguration: Split

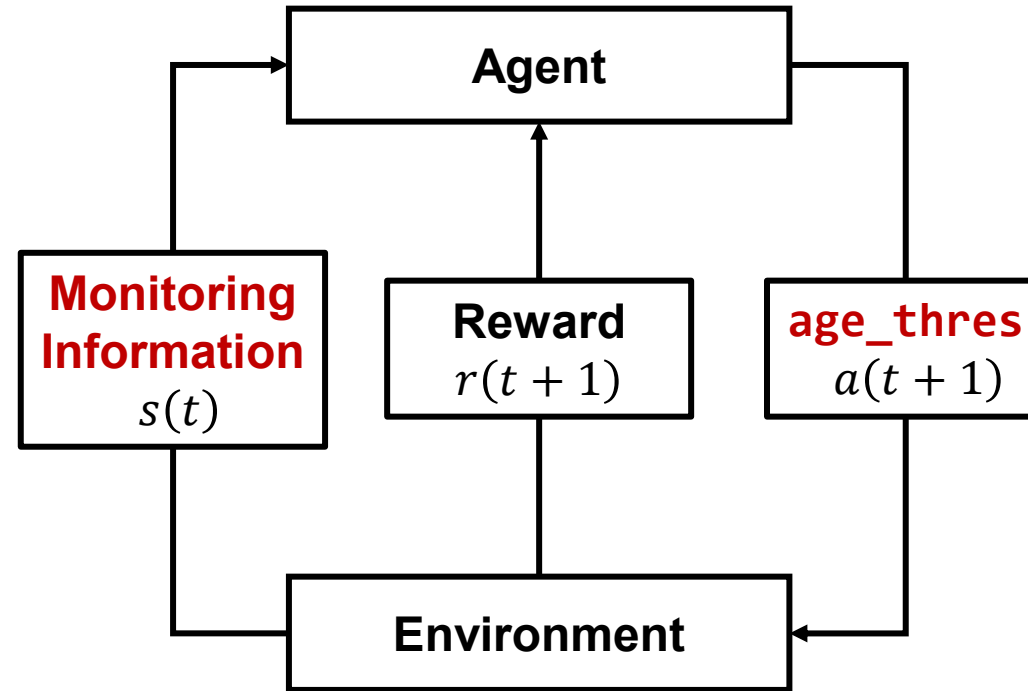
- **Split** region when the **access status of the sampling pages differs** at `sample_interval`



## RL-based Demotion Policy

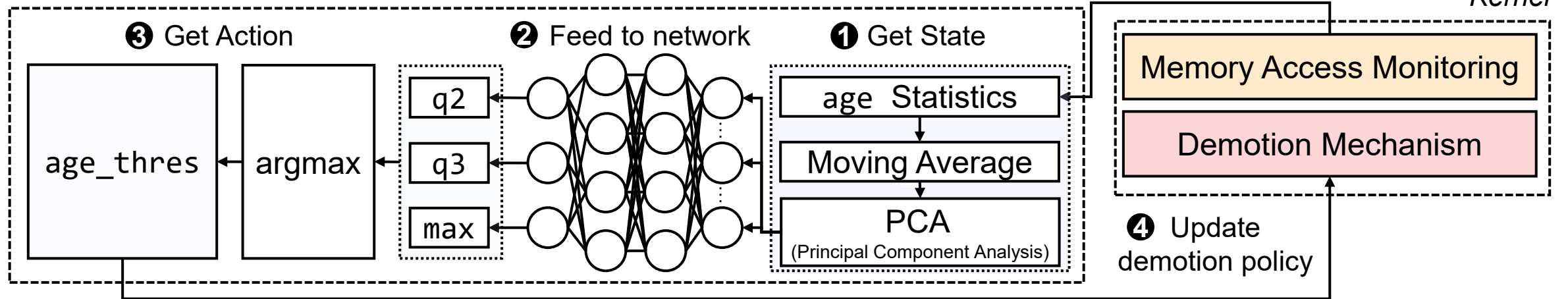


## RL: Recall



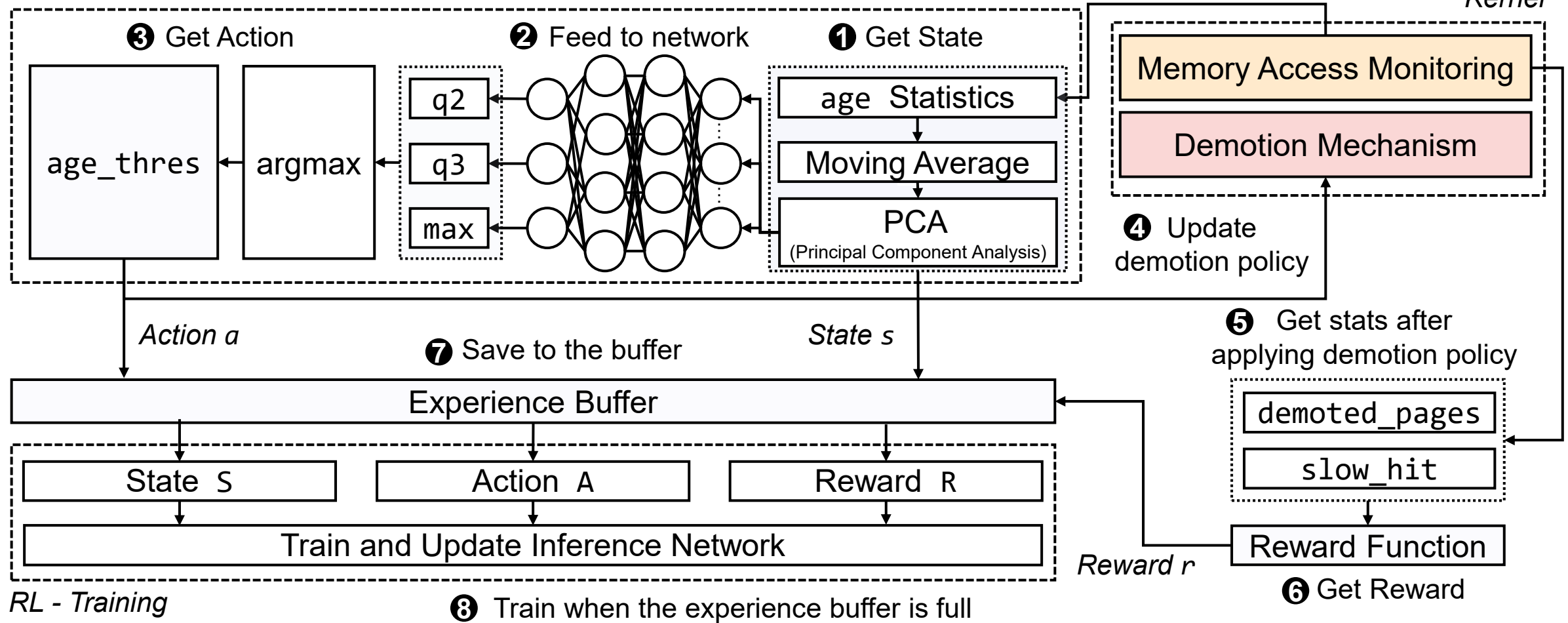
# RL: Design

RL - Inference



# RL: Design

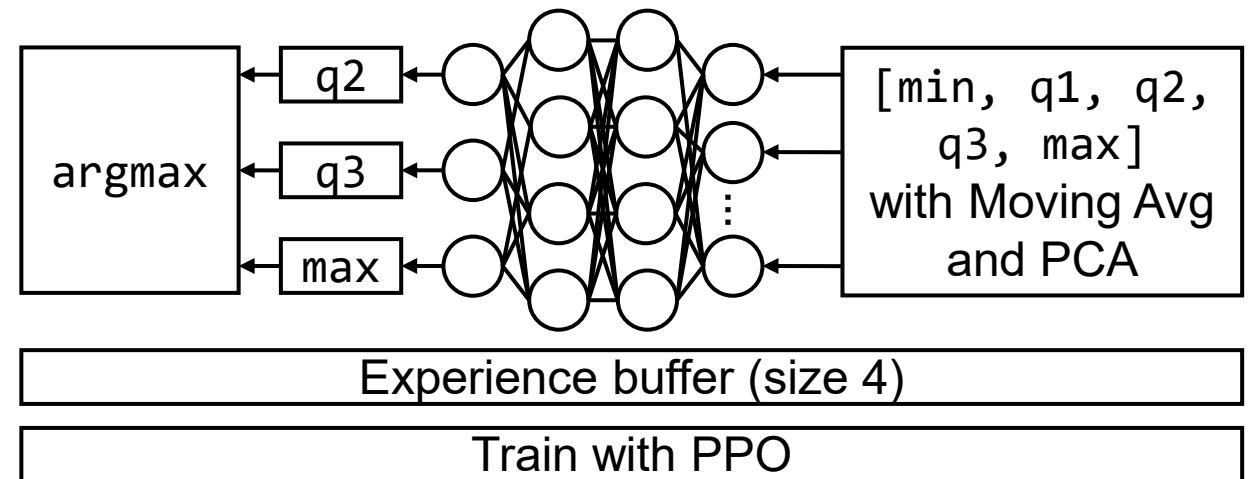
RL - Inference



$$r(t) = \log(\text{demoted\_pages}(t) / \text{slow\_hit}(t))$$

# RL: Detail

- Input Layer
  - min, q1 (25 percentile), q2 (50 percentile), q3 (75 percentile), max age distribution
  - 1x5 state vector
- **2 Hidden Layers**
  - 16, 32 nodes
- **Proximal Policy Optimization<sup>[1]</sup> (PPO)**  
Training Algorithm
- Experience buffer size: 4
  - Trained every 4 inferences
- **Pre-train** with GUPS microbenchmark
  - 3 memory access patterns used in HeMem<sup>[2]</sup>
- Implemented with PyTorch-based Rllib

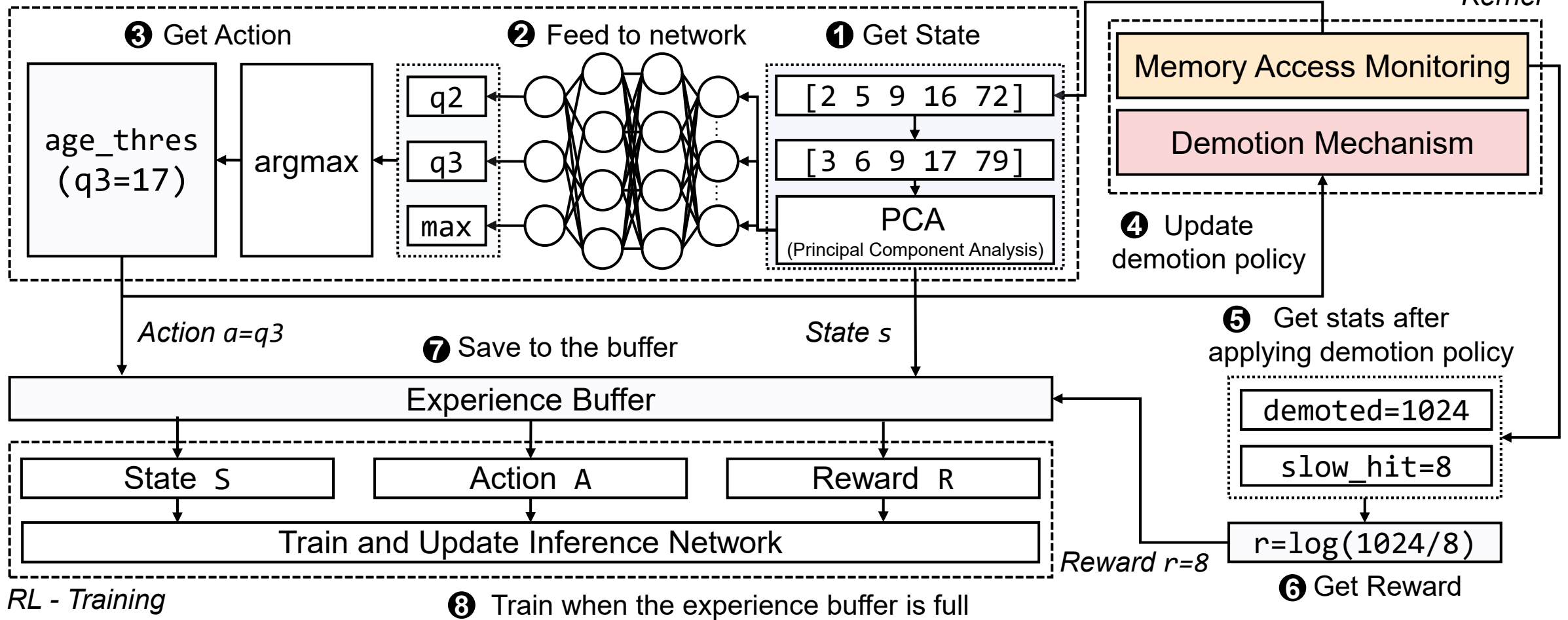


[1] John Schulman et al., "Proximal Policy Optimization Algorithms.", arXiv 2017

[2] Amanda Raybuck et al., "HeMem: Scalable Tiered Memory Management for Big Data Applications and Real NVM.", SOSP 2021

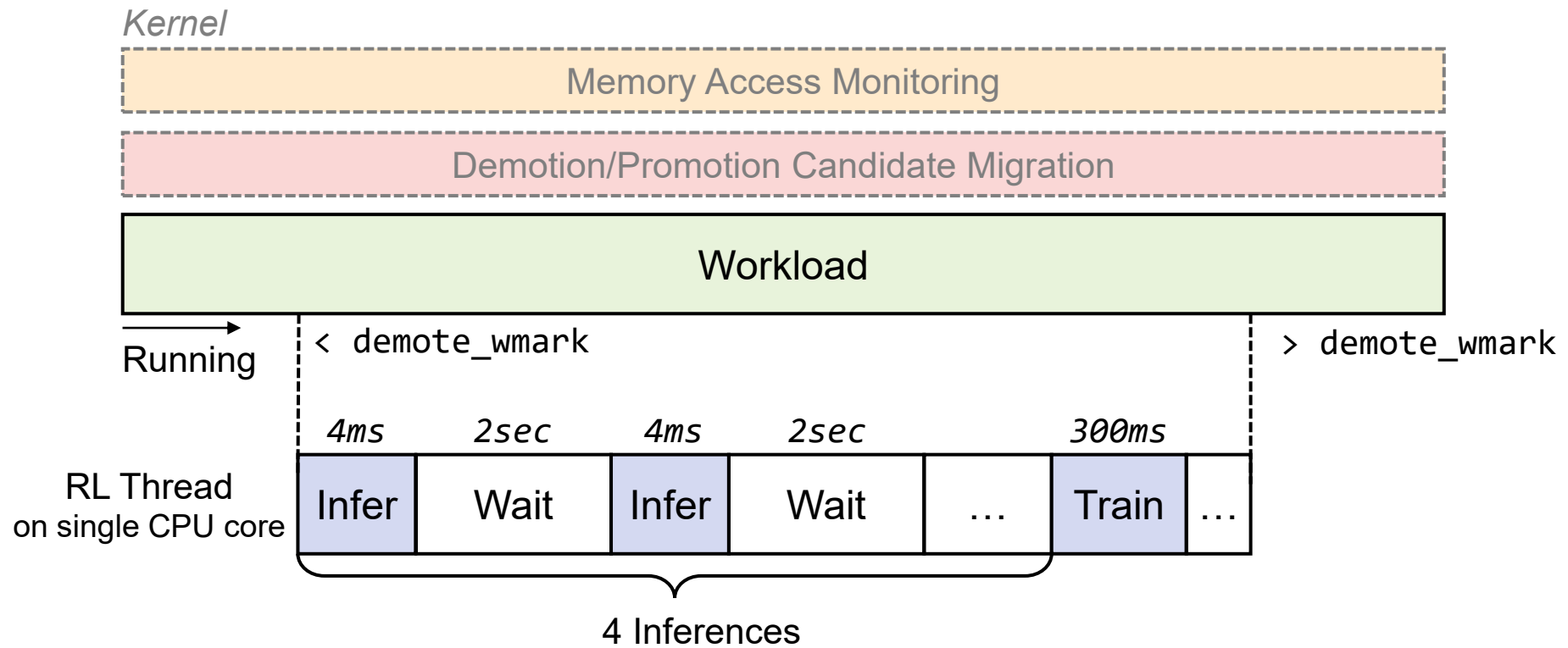
# RL: Example

RL - Inference





# RL: Execution Phases



## RL Execution Phases

Kernel

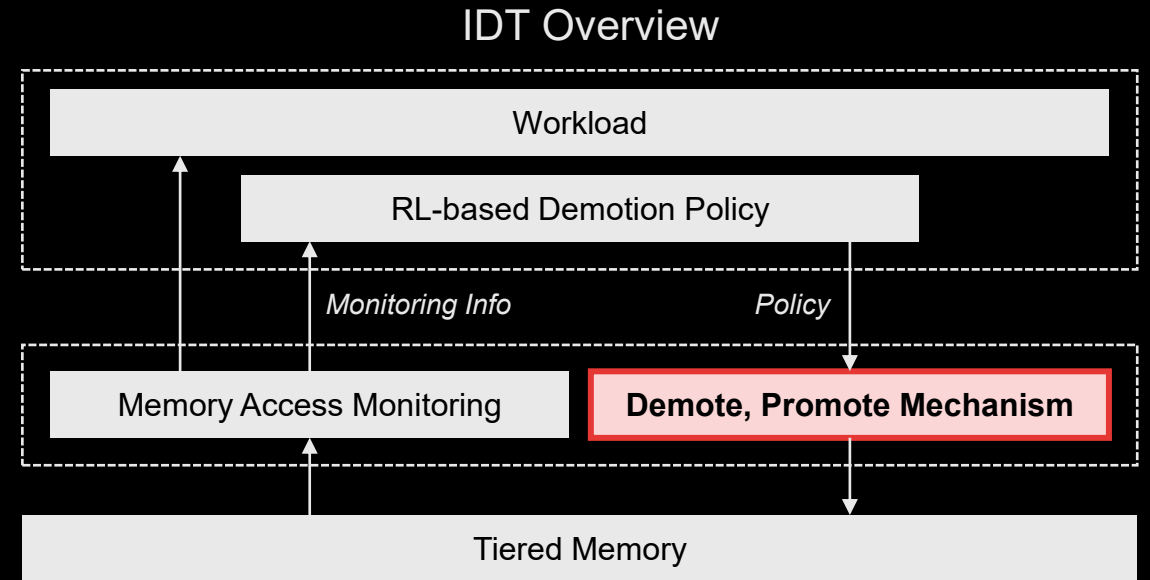
Memory Access Monitoring

**Theoretical Overhead:**  $(4\text{ms} \times 4 + 300\text{ms}) / (2\text{s} \times 4) = 3.95\%$  of a **single core**

**Actual overhead:** **Average 1.35%**, peak 3.75% of a single core

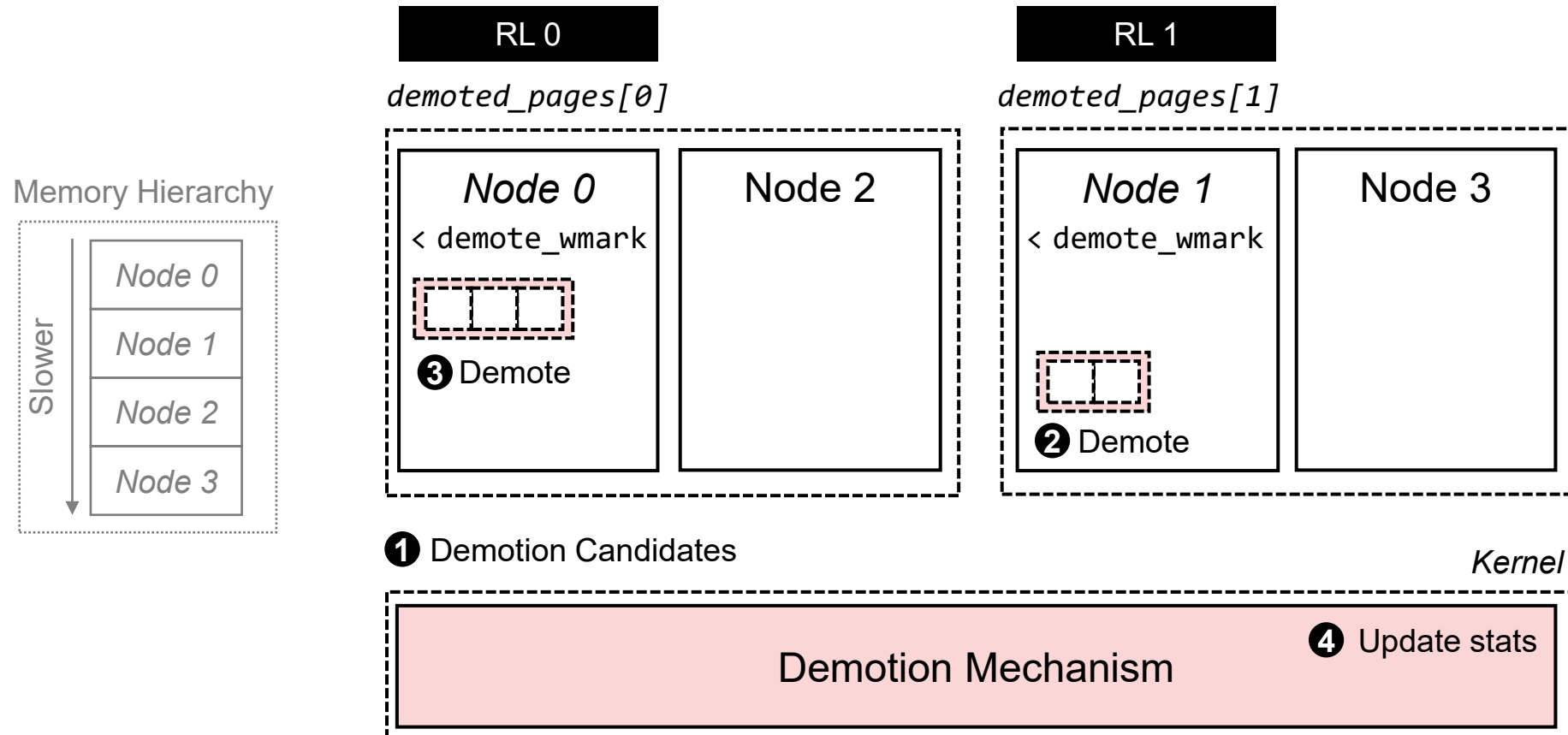
4 Inferences

## Demotion, Promotion Mechanism



# Demotion

- When a memory node's available space < demote\_wmark (Set to 10%)
- Demote regions with age > age\_thres and minimum access

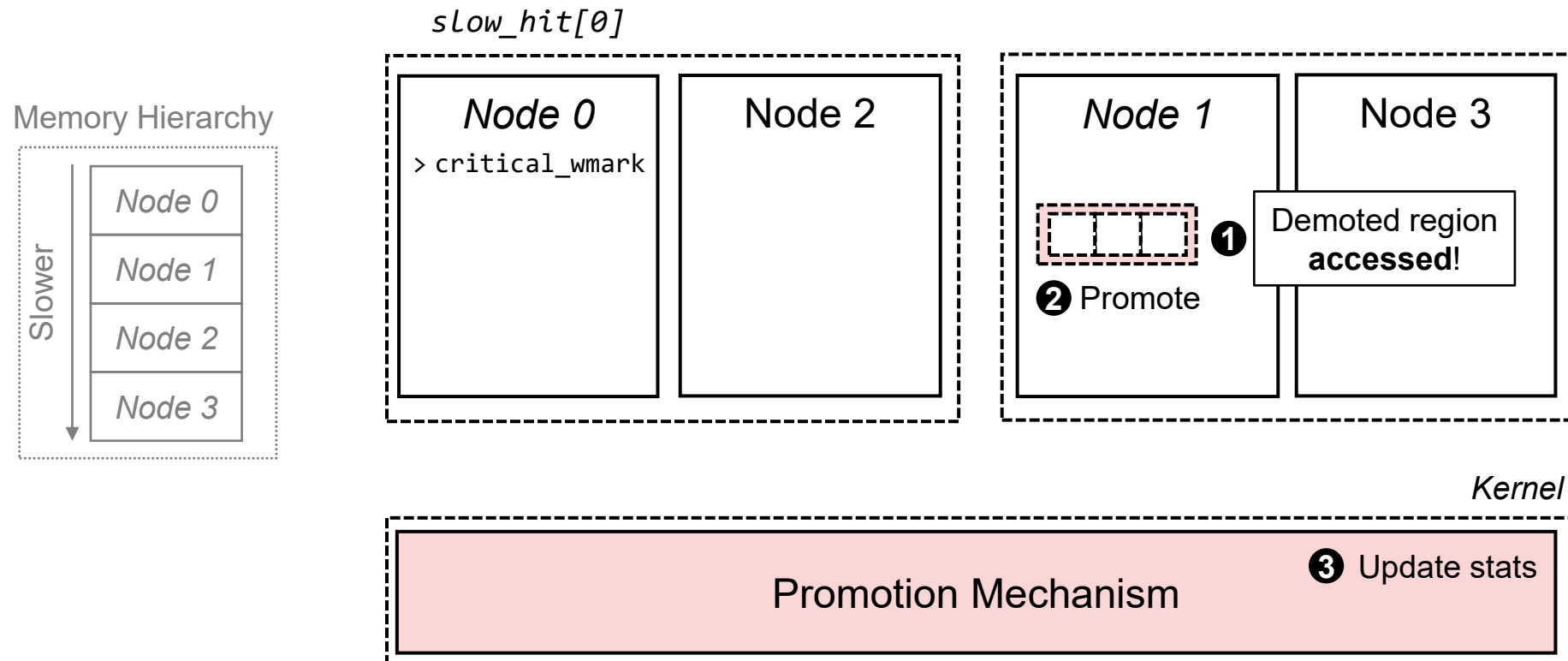


# Promotion

- **ARP** (**Accessed Region Promotion**)
- **PRP** (**Predictive Region Promotion**)

# Promotion: ARP (Accessed Region Promotion)

- Promote when demoted region is **accessed**
  - Destination node should have available space > `critical_wmark` (Set to 1%)

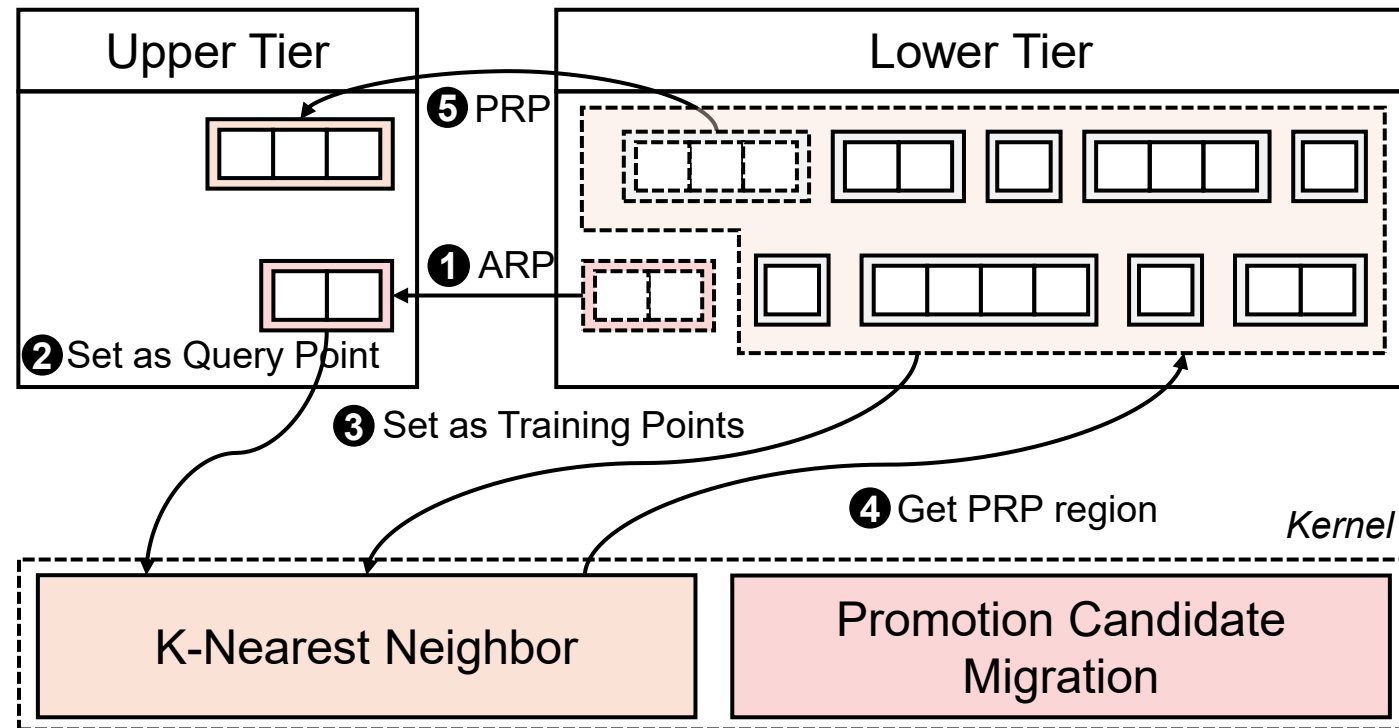


## Promotion: PRP (Predictive Region Promotion)

- ARP does not promote until access to the region's sampling pages is observed
  - Preemptively promoting regions similar to ARP region may be beneficial
- Identify a **similar** region with **k-Nearest Neighbor** and promote

# Promotion: PRP (Predictive Region Promotion)

- ARP does not promote until access to the region's sampling pages is observed
  - Preemptively promoting regions similar to ARP region may be beneficial
- Identify a **similar** region with **k-Nearest Neighbor** and promote



$$distance = \text{Normalized}(\text{vaddr distance}) + \text{Normalized}(\text{access\_history distance})$$

**Spatial Locality** **Temporal Locality**



# More Details in the Paper

- Aggressive demotion
  - Tighten demotion criteria when scarce fast memory
- Misplaced region promotion
  - Handle promotion of regions demoted by kswapd
- RL formulation
  - Problem formulation
  - Approximation for feasible implementation
- Sensitivity study

## IDT: Intelligent Data Placement for Multi-tiered Main Memory with Reinforcement Learning

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### ABSTRACT

To address the limitation of a DRAM-based single-tier in satisfying the comprehensive demands of main memory, multi-tiered memory systems are gaining widespread adoption. To support these systems, operating-system-level solutions that analyze the application's memory access patterns and ensure data placement in the appropriate memory tier have been vastly explored. In this paper, we identify reinforcement learning (RL) as an effective solution for tiered memory management, and its policy can be formulated in a solvable form using RL. We also demonstrate that an effective region-granularity memory access monitoring method is necessary to provide an accurate environment state to the RL model. Thus, we propose **IDT**, an intelligent data placement for multi-tiered main memory. IDT incorporates an RL-based demotion policy autotuning and a mechanism that efficiently demotes cold pages to lower-tier memory. IDT also promotes hot pages to upper-tier memory to minimize access on slow memory, featuring a lightweight machine learning algorithm. IDT employs region-granularity memory access monitoring with statistical-testing-based adjacent region merge and split to improve precision and mitigate ambiguity observed in prior works. Experiments on an actual four-tiered memory system show that IDT achieves an average 2.08× speedup over the default Linux kernel and 11.2% performance improvement compared to the state-of-the-art solution.

### CCS CONCEPTS

• Software and its engineering → Memory management; • Computer systems organization → Heterogeneous (hybrid) systems; • Computing methodologies → Reinforcement learning.

### KEYWORDS

Memory Tiering, Emerging Memory Technologies, Memory Management, Reinforcement Learning

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### 1 INTRODUCTION

The growing demand for memory-intensive workloads, such as high-performance computing, graph analytics, and in-memory databases, is highlighting the scaling limitations of a DRAM-based single-tier main memory [39]. To tackle this issue, a variety of memory types with diverse performance characteristics have been adopted to compose tiered memory systems. Recently, the rising interest in memories attached to Compute Express Link (CXL-Memory [9]) underscores that the future lies in multi-tiered memory systems by integrating various heterogeneous memories with a main-memory-like appearance to a system [36]. Cloud vendors, such as Amazon and Google, already serve large memory cloud instances based on multi-tiered memory systems [20, 33].

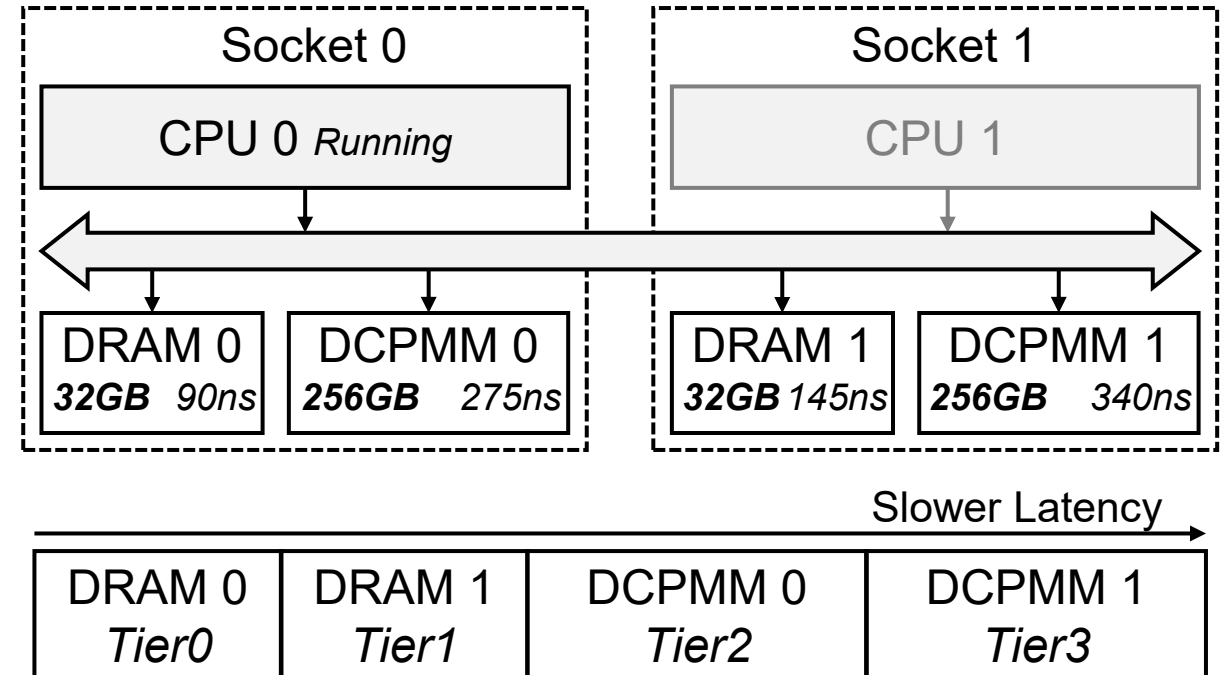
Tiered memory management requires a keen insight into an application's memory usage and placing the data in the proper memory tier according to its hotness. Thus, a number of prior studies have proposed operating-system-(OS)-level solutions to improve application performance by attentively exploiting the tiered memory system [2, 12, 15, 19, 23, 27, 36, 38, 48, 51, 55]. These OS-level tiered memory solutions typically consist of *data placement* to fully leverage diverse memory types and *memory access monitoring* to gather information for guiding data placement.

**Data placement.** Infrequently accessed pages in tiered memory should be *demoted* to lower-tier slow memory for efficient utilization of upper-tier fast memory. Moreover, to complement demotion, hot pages trapped in slow memory should be identified and *promoted* to upper-tier memory. Several prior studies have utilized the Linux kernel's 2Q LRU [19, 21, 35, 36, 56, 57] or multi-generational LRU (MGLRU) [58] to determine demotion candidates. However, the data hotness identified by 2Q LRU or MGLRU often fails to reflect the actual data hotness of the application. Therefore, precisely tracking both access frequency and recency for each page, and establishing a demotion policy with solid criteria would be more effective. Yet, implementing this method presents the challenge of

# Evaluation

# Experimental Setup

- Based on Linux kernel v6.0.19
  - Memory access monitoring developed with DAMON
- **Multi-tiered memory** setup
  - 2 socket machine with **DRAM** (fast memory) and Intel Optane **DCPMM** (slow memory)
- 4 State-of-the-art solutions for comparison
  - **Intel Tiering 0.8**<sup>[1]</sup>, **TPP**<sup>[2]</sup>, **AutoTiering**<sup>[3]</sup>, **AutoNUMA Tiering (MGRLU)**<sup>[4]</sup>
- Workloads: SPEC CPU2017, graph500, GAPBS
  - RSS set **96GB~110GB** to facilitate using 3 tiers
- Evaluation metric: Speedup (execution time) normalized against AutoNUMA Balancing



[1] Intel. 2022. Tiering-0.8. <https://git.kernel.org/pub/scm/linux/kernel/git/vishal/tiering.git/>.

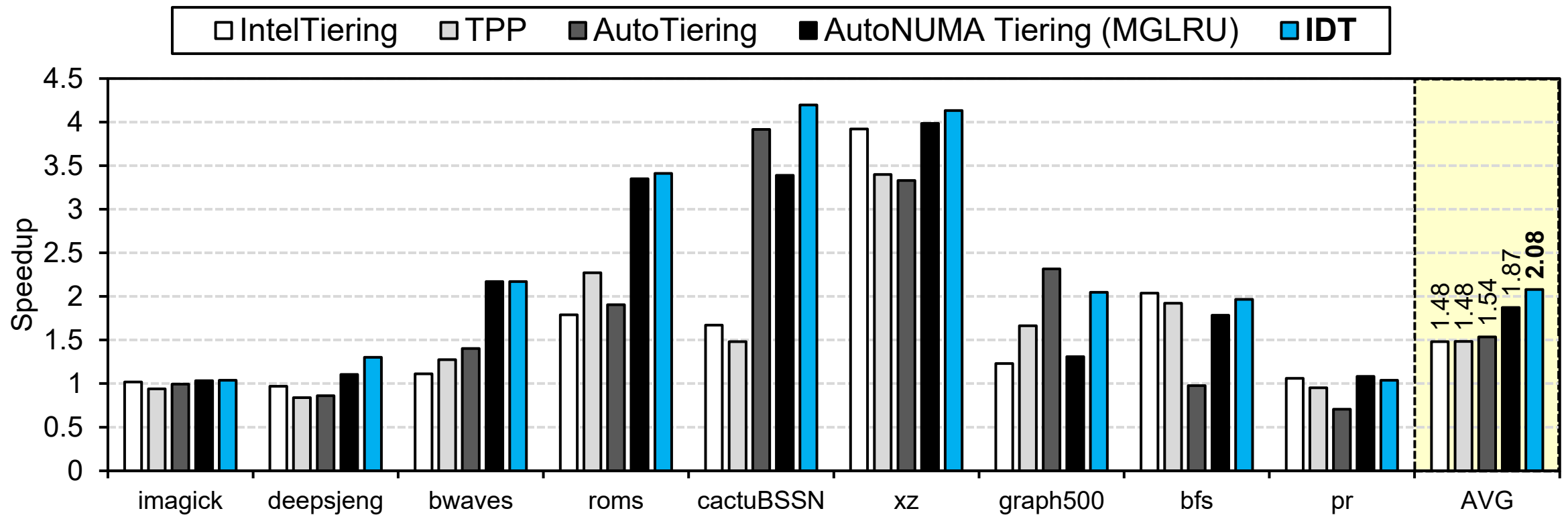
[2] Hasan Al Maruf et al., "TPP: Transparent Page Placement for CXL-Enabled Tiered-Memory," ASPLOS, 2023

[3] Jonghyeon Kim et al., "Exploring the Design Space of Page Management for Multi-Tiered Memory Systems," USENIC ATC (Virtual Event), 2021

[4] Yu Zhao. 2022. Multigenerational LRU Framework. <https://lwn.net/Articles/880393/>.

# Performance

- Outperforms the best-performing state-of-the-art solution AutoNUMA Tiering (MGLRU)) by **11.2%**
  - Average **2.08x** speedup against AutoNUMA Balancing



# Limitations

- Other parameters (e.g. 10% and 1% watermarks, sliding window size) are not determined by RL (or ML)
  - Our **goal was to advance the state-of-the-art** solution by **appropriately utilizing RL** (or ML)
  - Future works may apply ML to optimize other parameters
- **Blackbox**: Difficult to explain clear reasons for performance improvement by using RL

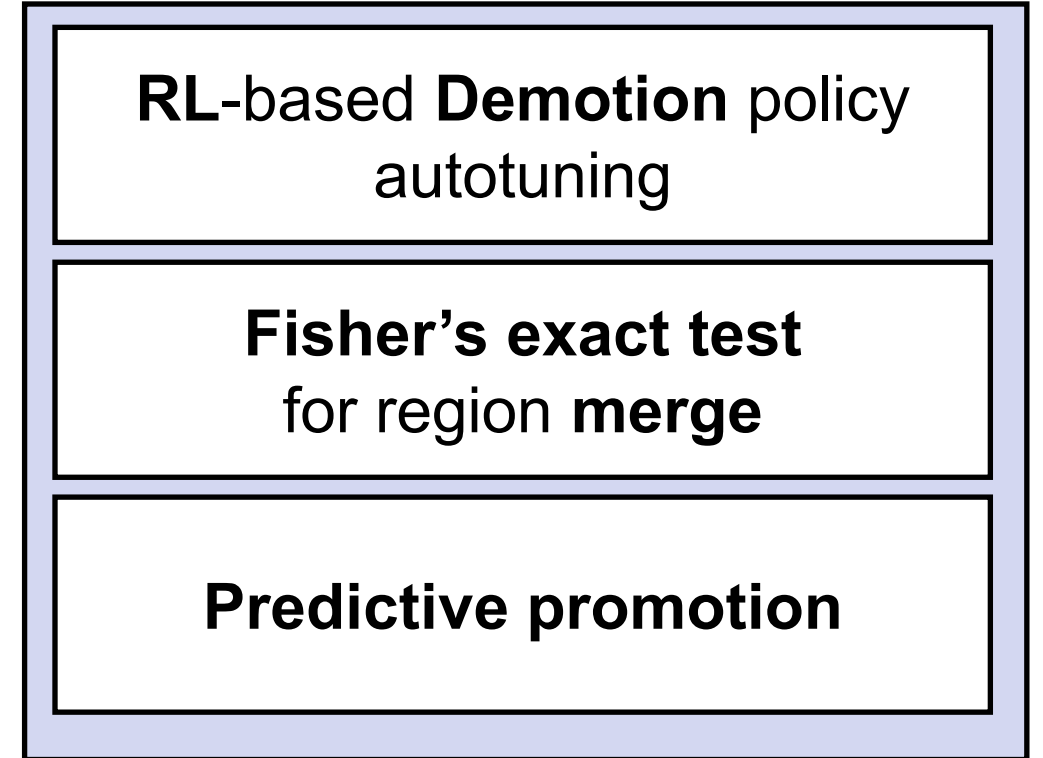
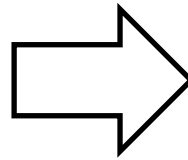
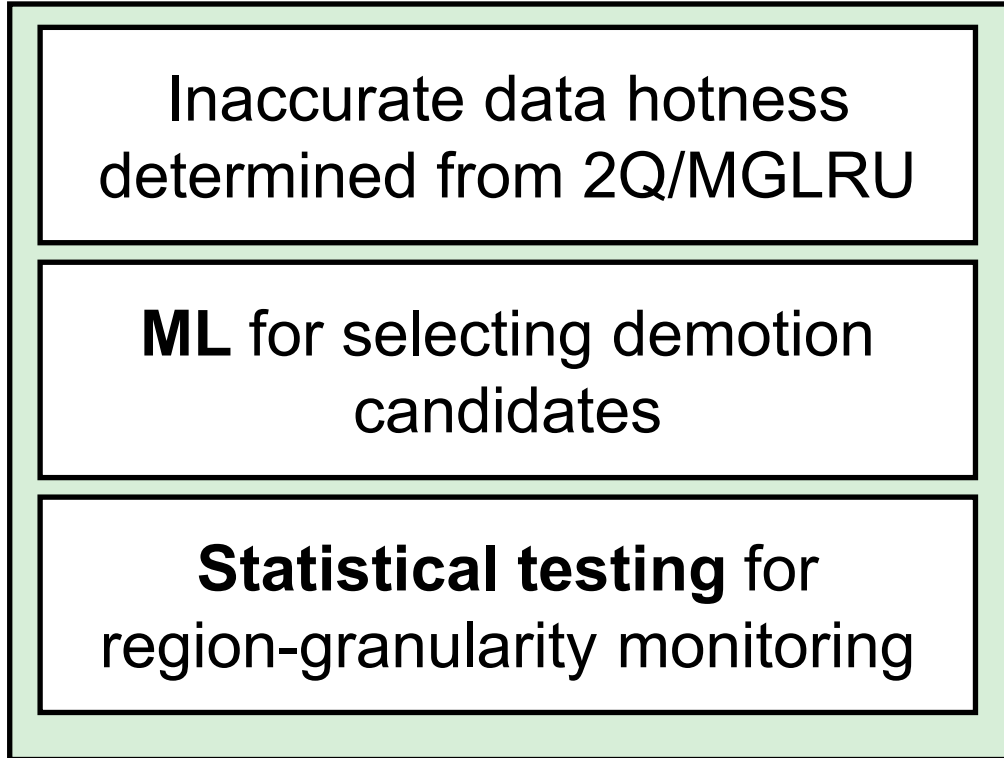
## Summary

Inaccurate data hotness  
determined from 2Q/MGLRU

**ML** for selecting demotion  
candidates

**Statistical testing** for  
region-granularity monitoring

## Summary



Outperforms the default Linux kernel by **2.08×**, state-of-the-art solution by **11.2%**

# Thank you!

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