

Exercise 2.1-3

```
for i = 1 to n
    if A[i] == v
        return i
return NIL
```

Invariant: Each element is accessed only once.

Proof: The variable i starts at 1. After the body of the for loop, i is incremented by 1. i is never incremented by more than one, and is never decremented.

Invariant: When v is found, i is at the correct index.

Proof: We proved that each element is accessed only once. This means that i must start at the beginning of the array and solely increment by 1 until it reaches the end of the array. Therefore, if $A[i]$ has the value v , i is the correct index.

Invariant: If v is not in A , NIL is returned.

Proof: We proved that each element is accessed only once. We also proved that when v is found, i is at the correct index. Therefore, if v is not found, the for loop will terminate, and NIL will be returned.

Exercise 3.1-1

We are trying to prove the following:

$$c_1 * (f(n) + g(n)) \leq \max(f(n), g(n)) \leq c_2 * (f(n) + g(n))$$

First, prove that $\max(f(n), g(n)) = O(f(n) + g(n))$:

1. $f(n) > 0$
2. $g(n) > 0$
3. $f(n) + g(n) \geq f(n)$
4. $f(n) + g(n) \geq g(n)$
5. $f(n) + g(n) \geq \max(f(n), g(n))$
6. $c_2 * f(n) + g(n) \geq \max(f(n), g(n))$ with $c_2 = 1$

Next, prove that $\max(f(n), g(n)) = \Omega(f(n) + g(n))$:

1. $f(n) > 0$
2. $g(n) > 0$
3. $\max(f(n) + g(n)) \geq f(n)$
4. $\max(f(n) + g(n)) \geq g(n)$
5. $2 * \max(f(n) + g(n)) \geq f(n) + g(n)$
6. $\max(f(n) + g(n)) \geq c_1 * f(n) + g(n)$ with $c_1 = \frac{1}{2}$

Therefore, $\max(f(n), g(n)) = \Theta(f(n) + g(n))$.