

Practical 7: Floyd Warshall – Dynamic Programming

Aim

To Implement the Floyd-Warshall Algorithm for All Pair Shortest Path Problem using Dynamic Programming.

Algorithm

1. Initialize distance and pred matrices with infinity and null values respectively, except for the diagonal elements which are initialized to 0.
2. For each vertex k from 1 to n, do the following:
 - a. For each pair of vertices i and j from 1 to n, check if the path from i to k and then from k to j is shorter than the current path from i to j. If it is, update the distance and pred matrices accordingly.
3. Return the distance and pred matrices.

Program

```
public class FloydWarshall {
    static int INF = 9999;
    public static void main(String[] args) {
        int graph[][] = {
            {0, 5, INF, 10},
            {INF, 0, 3, INF},
            {INF, INF, 0, 1},
            {INF, INF, INF, 0}
        };
        floydWarshall(graph);
    }
    static void floydWarshall(int[][] graph) {
        int V = graph.length;
        int[][] matrix = new int[V][V];
        for(int i = 0; i < V; i++) {
            for(int j = 0; j < V; j++) {
                matrix[i][j] = graph[i][j];
            }
        }
        for(int i = 0; i < V; i++) {
            for(int j = 0; j < V; j++) {
                for(int k = 0; k < V; k++) {
                    if(matrix[j][k] > matrix[i][k] + matrix[j][i]) {
                        matrix[j][k] = matrix[i][k] + matrix[j][i];
                    }
                }
            }
        }
        printMatrix(matrix);
    }
    static void printMatrix(int[][] matrix) {
        System.out.println("Resultant Matrix using Floyd Warshall is: ");
        int V = matrix.length;
        for(int i = 0; i < V; i++) {
            System.out.print("[");
            for(int j = 0; j < V; j++) {
                if(matrix[i][j] == INF) {
                    System.out.print("INF ");
                }
                else {
                    System.out.print(matrix[i][j] + " ");
                }
            }
        }
    }
}
```

```
        }  
    }  
    System.out.print("]");  
    System.out.println();  
}  
}
```

Output:

```
Resultant Matrix using Floyd Warshall is:  
[0 5 8 9 ]  
[INF 0 3 4 ]  
[INF INF 0 1 ]  
[INF INF INF 0 ]
```

Analysis of Algorithm**Time Complexity:**

The time complexity of the Floyd-Warshall algorithm is $O(V^3)$, where V is the number of vertices in the graph. The algorithm iteratively considers all possible intermediate vertices in the shortest path calculation, resulting in a **nested triple loop** that performs V^3 operations. This time complexity is independent of the specific graph structure or edge weights, making it a suitable algorithm for a wide range of graph problems.

Space Complexity:

The space complexity of the Floyd-Warshall algorithm is $O(V^2)$, where V is the number of vertices in the graph. This is because the algorithm requires a matrix of size $V \times V$ to store the shortest path distances between all pairs of vertices.