Programmer's Study Group

Week 05

Last homework

- https://leetcode.com/problems/remove-duplicates-from-sorted-list/
- https://leetcode.com/problems/swap-nodes-in-pairs/
- https://leetcode.com/problems/rotate-list/
- https://leetcode.com/problems/partition-list/
- https://leetcode.com/problems/reverse-linked-list-ii/
- https://leetcode.com/problems/odd-even-linked-list/
- https://leetcode.com/problems/intersection-of-two-linked-lists/
- https://leetcode.com/problems/linked-list-cycle/
- https://leetcode.com/problems/linked-list-cycle-ii/
- https://leetcode.com/problems/clone-graph/
- https://leetcode.com/problems/course-schedule/
- https://leetcode.com/problems/course-schedule-ii/

Prove of https://leetcode.com/problems/linked-list-cycle-ii/

Using two pointers, one of them one step at a time. another pointer each take two steps. If there is a cycle, the slow and fast pointer will meet somewhere in the cycle.

Suppose the distance between the start node of list and the start node of cycle is s, the distance between start of cycle and the meet point is k, the length of the Cycle is r. When they meet, slow pointer walked n rounds of the circle, totally walked t steps from very beginning, faster pointer walked m rounds of the circle, totally walked m rounds of the circle,

So
$$t = s + n * r + k$$
, $2t = s + m * r + k => s = (m - 2n) * r - k$

Since m and n are integer, we denote it as p.

So
$$s = p * r - k$$

That means, using one pointer start from the start node of list, another pointer start from the first meet point, all of them walk one step at a time, the first time they meet each other is the start of the cycle.