

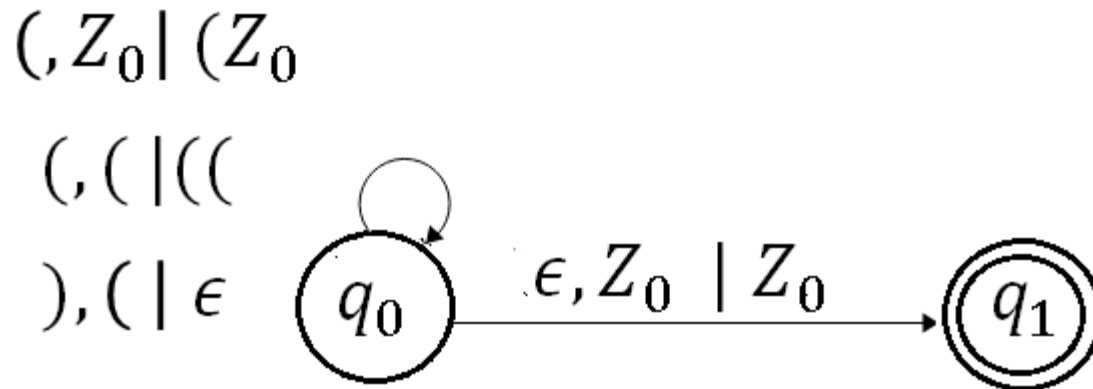
# Exercise Problems

- Problem where the product machine for PDA and DFA is constructed.
- If  $L$  is CFL but not regular, and  $R$  is regular, then is it possible for  $L \cap R$  to be regular?
  - Answer for this is given towards the end.

# Problem where the product machine for PDA and DFA is constructed.

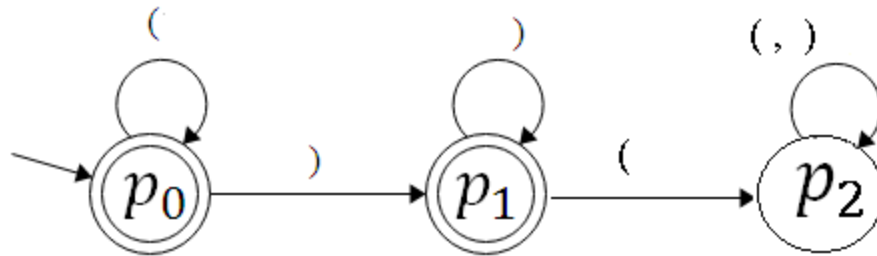
- Create PDA for well formed parentheses language (Dyck set).
- What is your idea ?

- Create PDA for well formed parentheses language (Dyck set).
- What is your idea ?

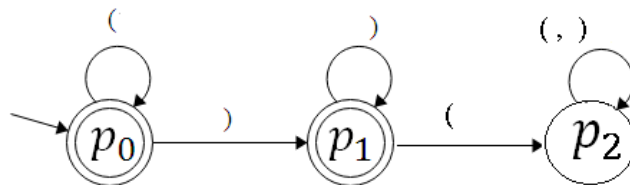
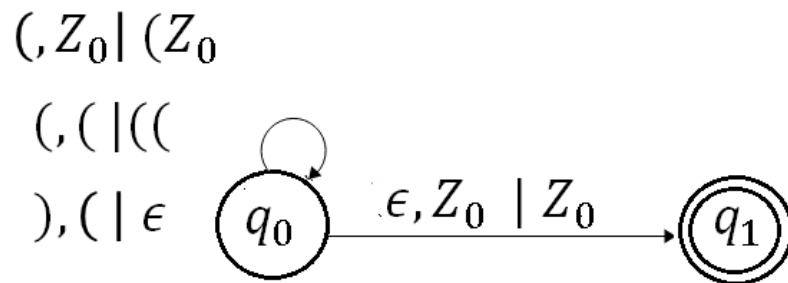


- Find DFA for  $(^* )^*$
- Give transition diagram for this.

- Find DFA for  $(^* )^*$
- Give transition diagram for this.



- Now create PDA for the intersection of the two languages.



# Verify your solution

- Your answer should have six states with all relevant arrows added.
  - In the class what is given is only a partial answer.
- Your constructed PDA should recognize the language  $\{ ({}^k )^k \mid k \geq 0 \}$ .
- Verify  $\epsilon$  is recognized by the product machine.
- Give ID seq. to verify  $()$  is in, but  $(( ))$  is not in the intersection.



# The second question -- answer

- If  $L$  is CFL but not regular, and  $R$  is regular, then is it possible for  $L \cap R$  to be regular?
- Yes. It is possible.
- Take for example  $L = \{a^k b^k \mid k \geq 1\}$  and  $R = \{a^l \mid l \geq 1\}$ .
- Their intersection is empty which is regular !!