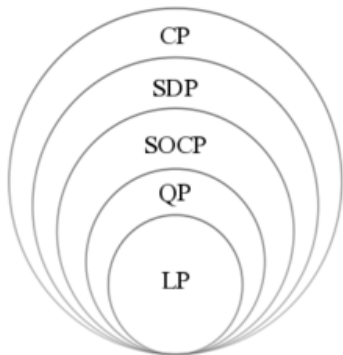


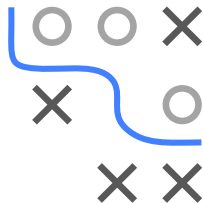
Optimization in Machine Learning

Linear Programming



Learning goals

- Definition and different forms of an LP
- Geometric intuition of LPs
- Characteristics of vertices
- Simplex algorithms



LINEAR PROGRAMMING

Linear program (LP):

optimization problem with **linear** objective function + **linear** constraints

General form

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^\top \mathbf{x} \\ \text{s.t.} & \mathbf{A}_1 \mathbf{x} \leq \mathbf{b}_1 \\ & \mathbf{A}_2 \mathbf{x} \geq \mathbf{b}_2 \\ & \mathbf{A}_3 \mathbf{x} = \mathbf{b}_3\end{array}$$

Examples:

$$\min_{\mathbf{x}} \quad \|\mathbf{Ax} - \mathbf{b}\|_1 \Leftrightarrow$$

$$\min_{\mathbf{x}, \mathbf{s}} \quad \mathbf{1}^\top \mathbf{s}$$

$$\text{s.t.} \quad \mathbf{Ax} - \mathbf{b} \leq \mathbf{s}$$

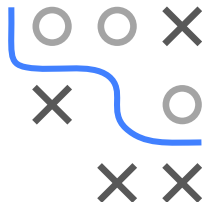
$$\mathbf{Ax} - \mathbf{b} \geq -\mathbf{s}$$

$$\min_{\mathbf{x}} \quad \|\mathbf{Ax} - \mathbf{b}\|_\infty \Leftrightarrow$$

$$\min_{\mathbf{x}, t} \quad t$$

$$\text{s.t.} \quad \mathbf{Ax} - \mathbf{b} \leq t\mathbf{1}$$

$$\mathbf{Ax} - \mathbf{b} \geq -t\mathbf{1}$$



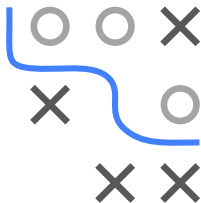
VERTICES

We assume that the rows of $\mathbf{A} \in \mathbb{R}^{m \times n}$ are linearly independent and $m \leq n$ to form a bounded unempty feasible set.

$\mathbf{Ax} = \mathbf{b}$ imposes m equality constraints:

- Each equality constraint reduces the dimension of the feasible set by 1.
- Starting with n -dimensional space, applying m independent equality constraints leaves a solution space of dimension $n - m$.

$\mathbf{x} \geq \mathbf{0}$ imposes n non-negativity.



SIMPLEX ALGORITHMS

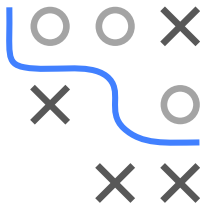
The **simplex algorithm** solves linear programs by moving from vertex to vertex of the feasible set, and produces an optimal vertex.

It operates on equality-form linear programs $\mathbf{Ax} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}$. We still assume that the rows of $\mathbf{A} \in \mathbb{R}^{m \times n}$ are linearly independent and $m \leq n$.

The method is guaranteed to arrive at an optimal solution so long as the linear program is feasible and bounded.

The simplex algorithm operates in two phases:

- **Initialization** phase: identifies a vertex partition.
- **Optimization** phase: transitions between vertex partitions toward a partition corresponding to an optimal vertex.



EXAMPLES

$$\mathbf{A}_{\mathcal{V}} = \begin{pmatrix} 2 & 1 & 1 & 0 \\ -4 & 2 & 0 & 1 \end{pmatrix}, \mathbf{b} = (9, 2)^{\top}, \mathbf{c} = (3, -1, 0, 0)^{\top}$$

Solution: $\mathcal{V} = \{1, 2\}, \mathcal{B} = \{3, 4\}$

$$\mathbf{x}_{\mathcal{B}} = \mathbf{A}_{\mathcal{B}}^{-1} \mathbf{b} = (9, 2)^{\top}$$

$$\boldsymbol{\lambda} = \mathbf{A}_{\mathcal{B}}^{-1} \mathbf{c}_{\mathcal{B}} = \mathbf{0}$$

$$\boldsymbol{\mu}_{\mathcal{V}} = \mathbf{c}_{\mathcal{V}} - (\mathbf{A}_{\mathcal{B}}^{-1} \mathbf{A}_{\mathcal{V}})^{\top} \mathbf{c}_{\mathcal{B}} = (3, -1)^{\top}$$

$\boldsymbol{\mu}_{\mathcal{V}}$ contains negative elements, so our current \mathcal{B} is suboptimal.

We will pivot on the negative one with $q = 2$, $-\mathbf{A}_{\mathcal{B}}^{-1} \mathbf{A}_{\{q\}} = (1, 2)^{\top}$.

This causes $x_4 = 0$, so updated $\mathcal{B} = \{2, 3\}$.

In the second iteration, we find

$$\mathbf{x}_{\mathcal{B}} = (1, 8)^{\top}, \boldsymbol{\lambda} = (0, -\frac{1}{2})^{\top}, \boldsymbol{\mu}_{\mathcal{V}} = (1, \frac{1}{2})^{\top}.$$

This is optimal with no negative entry, thus we have $\mathbf{x}^* = (0, 1, 8, 0)^{\top}$

