Java Sorting Algorithms Comparison

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Abstract

I. MOTIVATION

In order to show how an algorithm might run on a given set of hardware, and how the algorithm will perform when given large amounts of data, algorithms are analysed. Sorting algorithms sort data into a natural order. By analysing sorting algorithms, the fastest algorithm for a given problem can be determined.

II. BACKGROUND

A sorting algorithm is used to sort data with a natural order. One such sorting algorithm is insertion sort, which sorts by iterating through a list of data, taking the current position, and repositioning it into a more appropriate place in the list. Heap sort is another sorting algorithm that greatly differs than insertion sort in that it uses a divide, conquer, and combine method; meaning that it breaks the set it is sorting into subsets until the subsets can no longer be broken up and then merge sort combines the subsets together rendering the correct answer.

III. Procedure

An insertion sort can be implemented in a multitude of languages using the pseudocode provided in Algorithm 1.

Insertion Sort Pre-Condition: A is a non-empty array of data with a natural order.

Insertion Sort Post-Condition: A' is a permutation of A (containing all the same elements) in strictly non-decreasing order.

Algorithm 1 Insertion-Sort(A)

```
1: procedure Insertion-Sort(A)
      if A.length < 2 then
2:
          return A
3:
      end if
4:
      i=2
5:
      while i upto A.length do
6:
          key = A[i]
7:
         j = i - 1
8:
          while j downto 1 and key < A[j] do
9:
             A[j+1] = A[j]
10:
             j = j - 1
11:
          end while
12:
          A[i+1] = key
13:
          i = i + 1
14:
      end while
15:
      return A
16:
17: end procedure
```

Insertion Sort Outer-Loop Invariant: The subarray A'[1 ... i - 1] contains all the same elements as the subarray A[1 ... i - 1].

Insertion Sort Outer-Loop Initialization: The outer-loop invariant holds because $A'[1 \dots i-1]$ and $A[1 \dots i-1]$ both contain the same one element.

Insertion Sort Outer-Loop Maintenance: The outer-loop invariant holds because A'[1 ... i - 1] and A[1 ... i - 1] both contain the same elements, although they may be in different orders.

Insertion Sort Outer-Loop Termination: When the outer-loop terminates, i = A.length, which implies that the entire array has been traversed and the guard has been negated. The negation of the guard implies that $A'[1 \dots i-1]$ contains all the elements in $A[1 \dots i-1]$.

Insertion Sort Inner-Loop Invariant: A'[1 ... j] is sorted in strictly non-decreasing order.

Insertion Sort Inner-Loop Initialization: Before the first iteration of the loop, j = 1, meaning the subarray A'[1 ... j] contains exactly one element, which is already sorted.

Insertion Sort Inner-Loop Maintenance: At the beginning of each iteration of the loop the inner-loop invariant holds because j counts down from i, and A'[j+1] is swapped with A'[j] only if A'[j+1] is less than A[j].

Insertion Sort Inner-Loop Termination: The negation of the guard implies that j = A.length and that $A'[1 \dots j]$ has been entirely traversed and sorted in strictly non-decreasing order, which maintains the inner-loop invariant.

Insertion Sort Conclusion: The termination of both the inner and outer loops implies that the entire array has been traversed, A' is a permutation of A containing all the same elements in strictly non-decreasing order. This satisfies the post condition.

A merge sort can be implemented in a variety of languages using the pseudocode provided in Algorithm 2.

Merge Sort Pre-Condition: A is a non-empty array of a comparable data with a natural order. Merge Sort Post-Condition: A' is a permutation of A (containing all the same elements) in strictly non-decreasing order.

Merge Pre-Condition: Left and right are both non-empty arrays of a comparable data type with a natural order in strictly non-decreasing order.

Merge Post-Condition: Combined has all the elements of both left and right in strictly non-decreasing order.

Merge Sort First While Loop Invariant: Left contains i elements, all of which can be found in A[0 ... i up to mid].

Merge Sort First While Loop Invariant Initialization: The invariant holds true vacuously because i is 0 and left contains no elements.

Merge Sort First While Loop Invariant Maintenance: The invariant holds true because i is incremented at the same rate that elements are added to left from the same i index value in A.

Merge Sort First While Loop Invariant Termination: The invariant holds true because i has been incremented at the same rate that elements have been added to left from the same i index value in A upto mid.

Merge Sort Second While Loop Invariant: Right contains i elements, all of which can be found in A[mid ... A.length].

Merge Sort Second While Loop Invariant Initialization: The invariant holds true vacuously because i is 0 and right contains no elements.

Merge Sort Second While Loop Invariant Maintenance: The invariant holds true because i is incremented at the same rate that elements are added to right from the same i index value in A.

Algorithm 2 MERGE-SORT(A)

```
1: procedure Merge-Sort(A)
       if A.length < 2 then
3:
          return A
       end if
4:
      mid = A.length/2
5:
       for i = 1 upto mid do
6:
          left[left.length] = A[i]
7:
       end for
8:
       for i = mid upto A.length do
9:
          right[right.length] = A[i]
10:
       end for
11:
       left = Merge-Sort(left)
12:
       right = Merge-Sort(right)
13:
       A = Merge(left, right)
14:
       return A
15:
16: end procedure
17: procedure MERGE(left, right)
       var i = 0
18:
19:
       var y = 0
       var x = 0
20:
       while left.length! = i and right.length! = y do
21:
          if left[i] < right[y] then
22:
             combined[x] = left[i]
23:
             i = i + 1
24:
             x = x + 1
25:
          end if
26:
          if right[y] < left[i] then
27:
              combned[x] = right[y]
28:
             y = y + 1
29:
             x = x + 1
30:
          end if
31:
       end while
32:
       while left.length ! = i do
33:
          combined[x] = left[i]
34:
          i = i + 1
35:
          x = x + 1
36:
       end while
37:
       while right.length != y do
38:
          combined[x] = right[y]
39:
          y = y + 1
40:
          x = x + 1
41:
       end while
42:
       return combined
43:
44: end procedure
```

Merge Sort Second While Loop Invariant Termination: The invariant holds true because i has been incremented at the same rate that elements have been added to right from the same i index value in A upto to A.length.

Merge Together First While Loop Invariant: Combined contains x number of elements where x is the sum of i and y and those elements are contained in left[0 ... i] or right[0 ... y] in strictly non-decreasing order.

Merge Together First While Loop Invariant Initialization: Our invariant holds true vacuously before the first iteration of the loop because i = 0, y = 0, i + y = 0, x = 0 therefore x = i + y and combined contains no elements and is vacuously in order.

Merge Together First While Loop Invariant Maintenance: Our invariant holds true at the beginning of each iteration of the loop because x is incremented whenever i and y are incremented, and elements are added to combined from left only when left[i] < right[y] and added from right only when left[i] >= right[y].

Merge Together First While Loop Invariant Termination: Our invariant holds true when the loop terminates because x has been incremented whenever i and y are incremented, and elements have only been added to combined from left when left[i] < right[y] and from right only when left[i] >= right[y].

Merge Together Second While Loop Invariant: Combined contains x number of elements where x >= i and those elements are contained in left[0 ... i] in strictly non-decreasing order.

Merge Together Second While Loop Invariant Initialization: The invariant holds true before the first iteration of the loop because x has been incremented when i has been incremented and elements have been added to combined from left if and only if left[i] < right[y] and from right if and only if left[i] >= right[y].

Merge Together Second While Loop Invariant Maintenance: The invariant holds true at the beginning of each iteration of the loop because x has been incremented when i has been incremented and elements have been added to combined from left in order.

Merge Together Second While Loop Invariant Termination: The invariant holds true at the termination of the loop because x has been incremented when i has been incremented and elements have been added to combined from left in order.

Merge Together Third While Loop Invariant: Combined contains x number of elements where x is greather than or equal to y and those elements are contained in right[0 ... y] in strictly non-decreasing order.

Merge Together Third While Loop Invariant Initialization: The invariant holds true before the first initialization of the loop because x has been incremented when y has been incremented and elements have been added to combined from left if and only if left[i] < right[y] and from right if and only if left[i] >= right[y].

Merge Together Third While Loop Invariant Maintenance: The invariant holds true at the beginning of each iteration of the loop because x has been incremented when y has been incremented and elements have been added to combined from right in order.

Merge Together Third While Loop Invariant Termination: The invariant holds true at the termination of the loop because x has been incremented when y has been incremented and elements have been added to combined from right in order.

A heap sort can be implemented in a variety of languages using the pseudocode below in Algorithm 3.

Algorithm 3 Heap-Sort(A)

```
1: procedure Heapify(A, i, total)
      var\ left = i * 2
      var\ right = left + 1
3:
      var\ iPrime = i
4:
      if left \le total and A[left] > A[i] then
5:
          i = left
6:
      end if
7:
      if right <= total and A[right] > A[i] then
8:
9:
          i = right
      end if
10:
      if i! = iPrime then
11:
          var\ temp = A[iPrime]
12:
          A[iPrime] = A[i]
13:
          A[i] = tmp
14:
          A = heapify(A, i, total)
15:
      end if
16:
      return A
17:
18: end procedure
19: procedure HEAP-SORT(A)
      var\ size = A.length
20:
      for var\ i = size/2 downto 1 do
21:
          A = heapify(A, i, size)
22:
      end for
23:
      for var i = size downto 1 do
24:
          var\ tmp = A[1]
25:
          A[0] = A[i]
26:
          A[i] = tmp
27:
          size = size - 1
28:
          A = heapify(A, 1, size)
29:
      end for
30:
      return A
31:
32: end procedure
```

IV. Testing

A. Testing Plan and Results

All arrays used in testing are Java ArrayList<Integer> unless otherwise specified. All times are recorded in milliseconds using a stopwatch class borrowed from Robert Sedgewock and Kevin Wayne It is important to note that the soptwatch class used takes the elapsed real-time between the start of the sort algorithm and the end of the sort algorithm as opposed to taking the elapsed processor-time becasue these tests were run on a multi-core computer. In the table below, A denotes Array. Times in the table below are given as averages out of 10 trials.

Table I Insertion Sort Test Results

Tested Input	Expected Results	Actual Results	Time
Empty A	Empty A	Empty A	0.0003
A of 1000 Strings	Sorted A 1000 Strings	Sorted A 1000 Strings	0.021
A 1 Element	Original A	Original A	0.0003
A 10 Elements	Sorted A 10 Elements	Sorted A 10 Elements	0.0005
A 100 Elements	Sorted A 100 Elements	Sorted A 100 Elements	0.0021
A 1000 Elements	Sorted A 1000 Elements	Sorted A 1000 Elements	0.019
A 10000 Elements	Sorted A 10000 Elements	Sorted A 10000 Elements	0.129
A 100000 Elements	Sorted A 100000 Elements	Sorted A 100000 Elements	6.4923
A 1000000 Elements	Sorted A 1000000 Elements	Sorted A 1000000 Elements	2135.5007
A 10000000 Elements	Sorted A 10000000 Elements	OS Crash	N/A
A 1000 Identical Elements	Original Array	Original Array	0.0052

Table II
MERGE SORT TEST RESULTS

Tested Input	Expected Results	Actual Results	Time
Empty A	Empty A	Empty A	0.0002
A of 1000 Strings	Sorted A 1000 Strings	Sorted A 1000 Strings	0.0297
A 1 Element	Original A	Original A	0.0001
A 10 Elements	Sorted A 10 Elements	Sorted A 10 Elements	0.0004
A 100 Elements	Sorted A 100 Elements	Sorted A 100 Elements	0.0028
A 1000 Elements	Sorted A 1000 Elements	Sorted A 1000 Elements	0.0294
A 10000 Elements	Sorted A 10000 Elements	Sorted A 10000 Elements	0.2098
A 100000 Elements	Sorted A 100000 Elements	Sorted A 100000 Elements	2.1904
A 1000000 Elements	Sorted A 1000000 Elements	Sorted A 1000000 Elements	22.9314
A 10000000 Elements	Sorted A 10000000 Elements	Sorted A 10000000 Elements	241.0322
A 1000 Identical Elements	Original A	Original A	0.0298

Table III HEAP SORT TEST RESULTS

Tested Input	Expected Results	Actual Results	Time
Empty A	Empty A	Empty A	0.0002
A of 1000 Strings	Sorted A 1000 Strings	Sorted A 1000 Strings	.0188
A 1 Element	Original A	Original A	0.0003
A 10 Elements	Sorted A 10 Elements	Sorted A 10 Elements	0.0006
A 100 Elements	Sorted A 100 Elements	Sorted A 100 Elements	0.002
A 1000 Elements	Sorted A 1000 Elements	Sorted A 1000 Elements	0.0162
A 10000 Elements	Sorted A 10000 Elements	Sorted A 10000 Elements	0.0468
A 100000 Elements	Sorted A 100000 Elements	Sorted A 100000 Elements	0.889
A 1000000 Elements	Sorted A 1000000 Elements	Sorted A 1000000 Elements	23.0635
A 10000000 Elements	Sorted A 10000000 Elements	Sorted A 10000000 Elements	244.1952
A 1000 Identical Elements	Original A	Original A	0.0156

B. Problems Encountered

One major issue encountered during the development of this insertion sort was that after completing the sort, A' was sorted properly except for the first element in the array. No matter what value the

first element of A had, it did not change position in A'. For example, if A[5, 6, 3, 4, 7] was passed to the insertion sort algorithm, the returned array would look like A'[5, 3, 4, 6, 7]. Changing the guard for the inner for loop (see Algorithm 1 line 6) from key < A[j] and j downto 1 to j downto 1 and key < A[j] corrected this issue.

Another issue encountered during the development of this program was executing Java assert statements. When assertions were wrapped in parenthesis, they evaluated to a boolean, and the assert would throw an Assertion Error with the explanation 'cannot compare to boolean'. Removing the parenthesis wrapping the assertion fixed this issue.

The last issue encountered during the development process of this program was generic programming. The program would not compile because not all Objects have a compare To method. Adding the statement $\langle T | extends | Comparable \langle T \rangle >$ fixed this issue. [?]

V. Experimental Analysis
VI. Conclusions

APPENDIX

Listing 1 Driver

```
/*
* @Author: Preston Stosur-Bassett
* @Date: Feb, 23, 2015
* @Class: Driver
* @Description: This class will test the functionality of the overall
   program by serving as a driver that runs through and calls all other
   required classes
*/
import java.util.ArrayList;
public class Driver {
  public static void main(String args[]) {
   DummyData testData = new DummyData();
    ArrayList < Integer > testList = new ArrayList < Integer > ();
    testList = testData.runArrayList(10, 0, 100, testList);
    //testList = testData.runArrayList(1000, testList);
    System.out.println("Unsorted list: ");
    System.out.println(testList);
    Sort sorter = new Sort();
    Stopwatch watchman = new Stopwatch();
    //testList = sorter.mergeSort(testList);
    testList = sorter.heapSort(testList);
    System.out.println("Sorted List: ");
    System.out.println(testList);
    System.out.println("Time To Complete: "+watchman.elapsedTime());
}
                                   Listing 2
                                    Debug
/*
        @Author Preston Stosur-Bassett
        @Date Jan 21, 2015
        @Class Debug
        @Description This class will help debugging by being able to turn
   on and turn off debug messages easily
*/
import java.util.List;
import java.util.ArrayList;
```

```
public class Debug<T> {
        boolean debugOn; //Variable to keep track of whether or not debug
            is on
        /*
                @Description constructor method that sets the default
           value of debugOn to false so that debug statements will not
           automatically print
        */
        public void Debug() {
                debugOn = false;
        }
        /*
                @Description turn on debugging print statements
        */
        public void turnOn() {
                debugOn = true;
        }
        /*
                @Description turn off debugging print statements
        */
        public void turnOff() {
                debugOn = false;
        }
        /*
                @Description will print messages only when debugOn
           boolean is set to true
                @param String message the string to print when debugging
           is turned on
        public void print(T message) {
                if (debugOn == true) {
                        System.out.println(message);
                }
        }
        /*
                @Pre-Condition <code>T expected </code> and <code>T actual
           </code> are both of the same type T
                @Post-Condition If <code>T expected </code> and <code>T
           actual </code> are found to be equal, the program moves on,
           otherwise the program halts with <code>AssertionError</code>
           is thrown
                @Description runs an assert statement against an expected
            value and the actual value that are passed as parameters only
```

when <code>debugOn == true</code>

```
Operam T expected the expected value to assert against
  the actual value
        Oparam T actualt he actual value to assert against the
   expected value
public void assertEquals (T expected, T actual) {
        if (debugOn == true) {
                assert actual.equals(expected);
        }
}
/*
* @Pre-Condition: <code>List<Integer> actual</code> is a iterable
    list of Integer objects
        @Post-Conditions: If the List of Integer objects is in
   stricly non-decreasing order, the program moves on normally,
   if not, the program halts with an <code>AssertionError</code>
* @Description: runs an assertion statement against a list of
  Integer objects to ensure that for \langle code \rangle k = actual.size(); A[
  |k - 2| \le A[k - 1]; </code>
        @param List<Integer> actual the list to assert is in
   stricly non-decreasing order
public void assertOrder(List<Integer> actual) {
        if (debugOn == true) {
                int i = actual.size();
                while (i > 1)
                         assert actual.get(i - 1).compareTo(actual
                            . get(i - 2)) >= 0;
                         i --;
                }
        }
}
/*
        @Pre-Condition: <code>ArrayList<Integer> actual</code> is
   an ArrayList of Integer Objects
        @Post-Condition: If the ArrayList of Integer Objects is
   in stricly non-decreasing order, the program moves on normally
   , if not, the pgoram halts with an <code>AssertionError</code>
        @Description: runs an assertion statement against an
   ArrayList of Integer Objects to ensure that for \langle code \rangle k =
   actual.size(); A[k-2] \le A[k-1]; </code>
        @param ArrayList < Integer > actual the ArrayList to assert
   is in stricly non-decreasing order
public void assertOrder(ArrayList<Integer> actual) {
        if (debugOn == true) {
                int i = actual.size();
```

```
while (i > 1)
                        assert actual.get(i - 1).compareTo(actual
                           . get(i - 2)) >= 0;
                        i --;
                }
        }
}
/*
        @Pre-Condition: actual and expected both contain Integer
  Objects
        @Post-Condition: If all the elements inside of the actual
    arraylist are also contained in the expected arraylist, then
  the assertion holds true
        @Description: Tests to ensure a given ArrayList of
  Integer Objects contains all the elements of another given
  ArrayList of Integer Objects
        @param ArrayList<Integer> expected the list to check
  contains against
        @param ArrayList<Integer> actual the list to check to
  make sure all its elements are contained in the other
  arravlist
*/
public void assertContains(ArrayList<Integer> expected, ArrayList
  <Integer> actual) {
        if (debugOn = true) {
                for (int i = 0; i < actual.size(); i++) {
                        assert expected.contains(actual.get(i));
                }
        }
}
/*
        @Pre-Condition: expectedOne, expectedTwo, and actual all
  contain Integer Objects
        @Post-Condition: If all the elemts inside of the actual
  ArrayList are also contined in either the expectedOne
  ArrayList or the expectedTwo ArrayList, then the assertion
  holds true
        @Description: Tests to ensure a given ArrayList of
  Integer Objects contains all the elements of another given
  ArrayList of Integer Objects
        @param: ArrayList<Integer> expectedOne one of the lists
  to check to see if the given ArrayList actual's elements are
  contained in
        @param: ArrayList < Integer > expected Two one of the lists
  to check to see if the given ArrayList actual's elements are
* @param: ArrayList<Integer> actual the list to check of make
```

```
sure all its elements are contained in either expectedOne or
  expectedTwo
*/
public void assertContains(ArrayList < Integer > expectedOne,
   ArrayList < Integer > expectedTwo, ArrayList < Integer > actual) {
        if (debugOn == true) {
                for (int i = 0; i < actual.size(); i++) {
                         assert expectedOne.contains(actual.get(i)
                            ) | expected Two.contains(actual.get(i
                }
        }
}
/*
* @Description: asserts that the first arguement is stricly
  greator than the second arguement
        @param int large an integer primative value to assert is
   strictly greator than the second arguement
        Oparam int small an integer primative value to assert the
    first arguement is strictly greator than.
*/
public void assertStrictGreat(int large, int small) {
        if (debugOn == true) {
                assert large > small;
        }
}
/*
        @Description: asserts that the first arguement is
   strictly less than the second arguement
        Oparam int small an integer primative value to assert is
   stricly less than the second arguement
        Oparam int large an integer primative value to assert the
    first arguement is strictly less than.
*/
public void assertStrictLess(int small, int large) {
        if (debugOn == true) {
                assert small < large;
        }
}
/*
        @Description: asserts that the first arguement is greator
    than or equal to the second arguement
        Oparam int large an integer primative value to assert is
   greator than or equal to the second arguement
        @param int small an integer primative value to assert the
    first arguement is greator than or equal to.
*/
```

```
public void assertGreatEquals(int large, int small) {
                if (debugOn == true) {
                         assert large >= small;
                }
        }
        /*
                @Description: asserts that the first arguement is less
           than or equal to the second arguement
                Oparam int small an integer primative value to assert is
           less than or equal to the second arguement
                Operam int large an integer primative value to assert the
            first arguement is less than or equal to.
        public void assertLessEquals(int small, int large) {
                if (debugOn == true) {
                         assert small <= large;
                }
        }
}
                                   Listing 3
                                  DummyData
/*
        @Author Preston Stosur-Bassett
        @Date Jan 25, 2015
        @Class DummyData
        @Description This class contains methods to generate dummy data
   given a set of parameters.
*/
import java.util.ArrayList;
import java.util.Random;
public class DummyData {
        /*
                @Description runArrayList < Integer > will take an ArrayList
            of Integer Objects and add a given amount of values to it
                @param int end the ending value to denote when to stop
           adding to the array list
                Oparam int min the minimum value of the randomly
        *
           generated data.
                Operam int max the maximum value of the randomly
           generated data.
                @param ArrayList < Integer > list the list to add value to
           and return
                @return ArrayList<Integer> the list after it has been
           updated with the randomly generated data
```

```
public static ArrayList<Integer> runArrayList(int end, int min,
  int max, ArrayList < Integer > list) {
        Random random = new Random();
        Debug debugger = new Debug();
        int start = 0;
        // INVARIANT: A.length >= start
        // INITIALIZATION: start = 0, A.length can be longer than
            0 when initially passed, but not smaller, so our
           invariant holds
        debugger.assertGreatEquals(list.size(), start);
        while (start < end) {
                // MAINTANANCE: At the beginning of each
                   iteration, one element was added to A and
                   start was increased by one, therefore, our
                   invariant holds true.
                debugger.assertGreatEquals(list.size(), start);
                Integer intToAdd = new Integer (random.nextInt((
                   \max - \min + 1) + \min);
                list.add(intToAdd);
                //Count up on the iterator
                start++;
        /*TERMINATION: The negation of the guard implies that (
           end - start) number of elements have been added to A,
           since start is initialized as 0 at the beginning of
           the method and is
                incremented by 1 each iteration of the loop,
                   which means that start amount of elements have
                    been added to A, and so our invariant holds
                   true.
        debugger.assertGreatEquals(list.size(), start);
        return list;
}
/*
        @Description: runArrayList<String> will take an ArrayList
   of String Objects and add a given amount of String numerical
  values to it
        @param int end the ending value to denote when to stop
  adding to the array list
        @param ArrayList < String > list the list to add String
  values to and return
        @return ArrayList<String> the list after it has been
  updated with the randomly generated numerical String values
public static ArrayList<String> runArrayList(int end, ArrayList<
  String> list) {
        Random random = new Random();
```

```
Debug debugger = new Debug();
        int start = 0;
        // INVARIANT: A.length >= start
        // INITIALIZATION: Before the first iteration of the loop
           , start = 0 and A.length cannot be less than 0, so our
            invariant holds true
        debugger.assertGreatEquals(list.size(), start);
        while (start < end) {
                // MAINTENANCE: At the beginning of each
                   iteration of the loop our invariant holds
                   because for each iteration of the loop one
                   element is added to A and start is incremented
                    by 1
                debugger.assertGreatEquals(list.size(), start);
                Integer intToString = new Integer (random.nextInt
                   ((1000000 - 1) + 1));
                String intString = String.valueOf(intToString);
                list.add(intString);
                //Count up on the iterator
                start++;
        /* TERMINATION: The negation of the guard implies that (
           end - start) number of elements have been added to A,
           since start is initialized as 0 at the beginning of
           the method and is
                        incremented by 1 each iteration of the
                           loop, which means that start amount of
                            elements have been added to A, and so
                            our invariant holds true. */
        debugger.assertGreatEquals(list.size(), start);
        return list;
}
/*
        @Description: identicalElement will take an element and
  add it to the ArrayList<Integer> for a given amount of times
        @param int end the ending value to denote when to stop
  adding elements to the array
        Oparam int element the element to add over and over again
   to the array
        @param ArrayList<Integer> list the list to add elements
  to
        @return ArrayList<Integer> the list after it has been
  updated with the given data
public static ArrayList < Integer > identicalElement (int end, int
  element, ArrayList<Integer> list) {
        // INVARIANT: A.length >= start
```

```
int start = 0;
              Debug debugger = new Debug();
              //The element to add over and over again
              Integer iden = new Integer (element);
              // INITIALIZATION: Before the first iteration of the eloop
                  , start = 0 and A.length cannot equal anything less
                  than 0, so our invariant holds true
              debugger.assertGreatEquals(list.size(), start);
               while (start < end) {
                       // MAINTENANCE: At the beginning of each
                          iteration of the loop our invariant holds
                          because for each iteration of the loop one
                          element is added to A and start is incremented
                       debugger.assertGreatEquals(list.size(), start);
                       list.add(iden);
                       //Count up on the iterator
                       start++;
               /* TERMINATION: The negation of the gaurd implies that (
                  end - start) number of elements hav ebeen added to A,
                  since start is initialied as 0 at the beginning of the
                  method and is
                               incremented by 1 each iteration of the
                                  loop, which means that start amount of
                                   elements have been added to A, and so
                                   our invariant holds true
              debugger.assertGreatEquals(list.size(), start);
              return list;
      }
The class Stopwatch has not been altered from its original form.
```

```
Listing 4
Stopwatch
```

```
See {@link StopwatchCPU} for a version that measures CPU time.
 *
    @author Robert Sedgewick
 *
    @author Kevin Wayne
 *
 * /
public class Stopwatch {
    private final long start;
     * Initialize a stopwatch object.
     */
    public Stopwatch() {
        start = System.currentTimeMillis();
    /**
     * Returns the elapsed time (in seconds) since this object was
        created.
     */
    public double elapsedTime() {
        long now = System.currentTimeMillis();
        return (now - start) / 1000.0;
    }
}
                                   Listing 5
                                    Sort
/*
        @Author: Preston Stosur-Bassett
        @Date: Jan 24, 2015
        @Class: Sort
        @Description: This class will contain many methods that will sort
    generic data types using common sorting algorithms.
*/
import java.util.ArrayList;
import java.util.List;
public class Sort<T extends Comparable<T>>> {
        *
                @Pre-Condition: ArrayList<T> is a non-empty set of data
           where T is a comparable data type with a natural order
                @Post-Condition:
                @Description: heapify is a helper method for heapSort
```

```
that keeps the heap in order so that the root node is the
  largest element in the heap.
        @param ArrayList<T> unsorted is a non-empty set of data
  where T is a comparable data type with a natural order
        @param int i
        @param int total
*
        @return ArrayList<T> unsorted
*
private ArrayList<T> heapify(ArrayList<T> unsorted, int i, int
  total) {
        int left = i * 2;
        int right = left + 1;
        int original I = i;
        if (left <= total && unsorted.get(left).compareTo(unsorted
           . get(i)) > 0)  {
                i = left:
        if (right <= total && unsorted.get(right).compareTo(
           unsorted.get(i) > 0 {
                i = right;
        if (i != originalI) {
                T tmp = unsorted.get(originalI);
                unsorted.set(originalI, unsorted.get(i));
                unsorted.set(i, tmp);
                unsorted = heapify (unsorted, i, total);
        }
        return unsorted;
}
/*
        @Pre-Condition: ArrayList<T> unsorted is a non-empty
  ArrayList <T> where T is a comparable data type with a natural
  order.
        @Post-Condition: ArrayList<T> sorted is a permutation of
  unsorted (it contains all the same elements) in stricly non-
  decreasing order
        @Description: heapSort will sort a given set of data in
  an ArrayList<T> in strictly non-decreasing order using the
  heap sort method.
        @param ArrayList<T> unsorted is a non-empty ArrayList<T>
  where T is a comparable data type with a natural order
        @return ArrayList<T> sorted is a permutation of unsorted
  in strictly non-decreasing order
//Invariant for First For Loop: TODO: write this
//Invariant for Second For Loop: TODO: write this
public ArrayList<T> heapSort(ArrayList<T> unsorted) {
```

```
int arrSize = unsorted.size() - 1;
        //Initialization: TODO: write this
        for (int i = arrSize / 2; i >= 0; i--) {
                //Maintanance: TODO: write this
                unsorted = heapify (unsorted, i, arrSize);
                //Termination: TODO: write this
        for (int i = arrSize; i > 0; i---) { //Initialization: TODO
           : write this
                //Maintanance: TODO: write this
                T \text{ tmp} = unsorted.get(0);
                unsorted.set(0, unsorted.get(i));
                unsorted.set(i, tmp);
                arrSize --:
                unsorted = heapify (unsorted, 0, arrSize);
                //Termination: TODO: write this
        ArrayList<T> sorted = unsorted;
        return sorted;
}
/*
        @Pre-Condition: ArrayList<T> left is a non-empty sorted
  array in stricly non-decreasing order where T is a comparable
  data type with a natural order
        @Post-Condition: ArrayList<T> right is a non-empty sorted
   array in strictly non-decreasing order where T is a
  comparable data type with a natural order
        @Description: mergeTogether is used by the mergeSort
  method to recombine the left and right sections of the
  ArrayList<T> that is being sorted by merge sort. Note that
  this is a helper method for the mergeSort method, and should
  not be called externally of this class.
        @param ArrayList<T> left a non-empty ArrayList<T> where T
   is a comparable data type with a natural order.
        @param ArrayList<T> right a non-empty ArrayList<T> where
  T is a comparable data type with a natural order.
        @return ArrayList<T> combined should contain all the
  elements of left and right in stricly non-decreasing order
//Invariant for First While Loop: combined contains x number of
  elements where x is the sum of i and y and those elements are
  contained in left [0 ... i] or right [0 ... y] in stricly non-
  decreasing order
//Invaraint for Second While Loop: combined contains x number of
  elements where x is greater than or equal to i and those
  elements are contined in left [0 ... i] in stricly non-
  decreasing order
//Invariant for Third While Loop: combined contains x number of
  elements where x is greater than or equal to y and those
```

```
elements are contained in right [0 ... y] in stricly non-
  decreasing order
private ArrayList<T> mergeTogether(ArrayList<T> left , ArrayList<T</pre>
  > right) {
        ArrayList < T > combined = new ArrayList < T > ();
        int i = 0;
        int y = 0;
        int x = 0;
        //Debug
        Debug debugger = new Debug();
        //Initialization: Our invariant holds true vacuously
           before the first execution of the loop because x, i,
           and y are all equal to zero, combined is empty and
           therefore in order
        debugger.assertEquals(0, i);
        debugger.assertEquals(0, y);
        debugger.assertEquals(0, x);
        debugger.assertEquals(i, combined.size());
        while (left.size() != i \&\& right.size() != y) {
                //Maintanance: Our invariant holds true at the
                   beginning of each iteration of the loop
                   because x is incremented whenever i or v is
                   incremented and elements are added to combined
                    from left and right in order
                debugger.assertEquals(i+y, x);
                debugger.assertEquals(x, combined.size());
                debugger.assertOrder(combined);
                debugger.assertContains(right, left, combined);
                if(left.get(i).compareTo(right.get(y)) < 0) {
                        combined.add(x, left.get(i));
                        i++;
                        x++;
                else {
                        combined.add(x, right.get(y));
                        y++;
                        x++;
                }
        //Termination: Our invariant holds tur at the termination
            of the loop because x has been incremented whenever i
            or y has been incremented, and elements are added to
           combined from left and right in order
        debugger.assertEquals(i+y, x);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
```

```
debugger.assertContains(right, left, combined);
//Initialization: Our invariant holds true before the
   first execution of the loop because x has been
  incremented whenever i has been incremented and
  elements have been added to combined from left in
  order
debugger.assertGreatEquals(x, i);
debugger.assertEquals(x, combined.size());
debugger.assertOrder(combined);
debugger.assertContains(left, combined);
while (left.size() != i) {
        //Maintanance: Our invariant holds true at the
           beginning of each iteration of the loop
           because x has been incremented whenever i is
           incremented and elements have been added to
           combined from left in order
        debugger.assertGreatEquals(x, i);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(left, combined);
        combined.add(x, left.get(i));
        i++;
        x++;
//Termination: Our invariant holds true at the
   termination of the loop because x has been incremented
   whenever i was incremented and elements have been
  added to combined from left in order
debugger.assertGreatEquals(x, i);
debugger.assertEquals(x, combined.size());
debugger.assertOrder(combined);
debugger.assertContains(left, combined);
//Initialization: Our invariant holds true before the
   first execution of the loop because x has been
  incremented whenever y has been incremented and
  elements have been added to combined from right in
  order
debugger.assertGreatEquals(x, y);
debugger.assertEquals(x, combined.size());
debugger.assertOrder(combined);
debugger.assertContains(right, combined);
while (right.size() != y) {
        //Maintanance: Our invariant holds tur at the
           beinning of each iteration of the loop because
            x has been incremented whenever y is
```

```
incremented and elements have been added to
                   combined from right in order
                debugger.assertGreatEquals(x, y);
                debugger.assertEquals(x, combined.size());
                debugger.assertOrder(combined);
                debugger.assertContains(right, combined);
                combined.add(x, right.get(y));
                v++;
                x++;
        //Termination: Our invariant holds true at the
           terminating of the loop because x has been incremented
            whenever y was incremented and elements have been
           added to combined from right in order.
        debugger.assertGreatEquals(x, y);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(right, combined);
        return combined;
}
/*
        @Pre-Condition: ArrayList<T> unsorted is a set of data
  type T, where T is a Comparable data type with a natural order
        @Post-Condition: ArrayList<T> returnValue is a
  permutation of unsorted in strictly non-decreasing order.
        @Description: mergeSort will sort a given set of data in
  ArrayList<T> using the merge sort method
        @param a non-empty ArrayList<T> unsorted where T is a
  Comparable data type with a natural order
        @return ArrayList<T> returnValue which is a permutation
  of unsorted, in strictly non-decreasing order,
//Invariant for First While Loop: left contains i elements, all
  of which can be found in sorted
//Invariant for Second While Loop: right contains y elements, all
   of which can be found in sorted
public ArrayList<T> mergeSort(ArrayList<T> unsorted) {
        ArrayList<T> sorted = unsorted;
        ArrayList<T> left = new ArrayList<T>();
        ArrayList < T > right = new ArrayList < T > ();
        ArrayList<T> returnValue;
        //Debug
        Debug debugger = new Debug();
        debugger.turnOn();
```

```
if(sorted.size() \le 1)
        returnValue = sorted;
}
else {
        int mid = (sorted.size() / 2);
        int i = 0;
        //Initialization: Our invariant holds true because
            i is zero and left contains 0 elements before
            the first iteration of the loop.
        debugger.assertEquals(i, left.size());
        while (i < mid)
                //Maintanance: Our invariant holds true
                   because i is increased at the same
                   rate elements are added to left from
                   the same i index in sorted
                 debugger.assertEquals(i, left.size());
                debugger.assertContains(sorted, left);
                T \text{ temp} = \text{sorted.get(i)};
                left.add(temp);
                i++;
        //Termination: Our invariant holds true because i
            has been incremented at the same rate
           elements are added to left from the same index
            i in sorted
        debugger.assertEquals(i, left.size());
        debugger.assertContains(sorted, left);
        int y = mid;
        //Initialization: Our invariant holds true
           because i is zero and right contains 0
           elements before the first iteration of the
           loop.
        debugger.assertEquals(y, right.size());
        while (y < sorted.size())
                //Maintanance: Our invariant holds true
                   because y is increased at the same
                   rate elements are added to right from
                   the same y index in sorted.
                debugger.assertEquals(v, right.size());
                debugger.assertContains(sorted, right);
                T \text{ temp} = \text{sorted.get}(y);
                right.add(temp);
                y++;
```

```
//Termination: Our invariant holds true because y
                    has been incremented at the same rate
                   elements are added to left from the same index
                    y in sorted
                debugger.assertEquals(y, right.size());
                debugger.assertContains(sorted, right);
                left = mergeSort(left);
                right = mergeSort(right);
                returnValue = mergeTogether(left, right);
        return return Value;
}
/*
        @Pre-Condition: ArrayList<T> unsorted is an unsorted
  ArrayList of a comparable data type that is non-empty
        @Post-Condition: ArrayList <T> will return a permutation
  of <code>unsorted</code> that will be in increasing order
        @Description: insertionSort will sort an ArrayList of
  generic type T in increasing order using an insertion sort
        @param ArrayList<T> unsorted is a non-empty unsorted
  array list of T, where T is a comparable type
        @return sorted is a permutation of <code>unsorted </code>
  where all the elements are sorted in increasing order
*/
// INVARIANT (Outer-Loop): The pre condition implies that sorted
  [0 \ldots i-1] will contain all the same data as unsorted [0 \ldots
   i - 1.
// INVARIANT (Inner-Loop): sorted[0 ... j] is sorted in stricly
  non-decreasing order.
public ArrayList<T> insertionSort(ArrayList<T> unsorted) {
        Debug debugger = new Debug<List<T>>();
        debugger.turnOn();
        ArrayList<T> sorted = unsorted;
        if(sorted.size() > 1) {
                int i = 1;
                /* INITIALIZATION (Outer-Loop): The invariant
                   holds because i = 1, and there is one element
                   in the subarray of sorted [0 \dots i-1] and
                   unsorted [0 \ldots i-1], */
                List < T > subSortedOI = sorted.subList(0, i - 1);
                List <T> subUnsortedOI = unsorted.subList(0, i -
                   1);
                debugger.assertEquals(subUnsortedOI, subSortedOI)
                while (i < sorted.size()) {
                        /* MAINTENANCE (Outer-Loop): At the
```

```
beginning of each iteration of the
   loop, the loop invariant is maintained
    because the subarray of sorted [0 ...
   i-1 contains all the same elements
        unsorted [0 \ldots i-1] */
List < T > subSortedOM = sorted.subList(0, i)
   -1);
List <T> subUnsortedOM = unsorted.subList
   (0, i-1);
debugger.assertEquals(subUnsortedOM,
   subSortedOM);
T value = sorted.get(i);
int j = i - 1;
// INITIALIZATION (Inner-Loop): Before
   the first iteration of the loop, j =
   0, the subarray of sorted [0 \dots 0]
   contains one elements and therefore
   the invariants holds vacuously.
List subSortedII = sorted.subList(0, j);
debugger.assertOrder(subSortedII);
while (j \ge 0 \&\& \text{ (value.compareTo(sorted.)})
   get(j) < 0)
        // MAINTENANCE: (Inner-Loop): At
           the beginning of each
           iteration sorted [0 \dots j] is
           sorted in stricly non-
           decreasing order
        List subSortedIM = sorted.subList
           (0, j);
        debugger.assertOrder(subSortedIM)
        sorted.set(j+1, sorted.get(j));
        j --;
sorted.set(j+1, value);
// TERMINATION (Inner-Loop): The negation
    of the guard implies that the sorted
   [0 ... j] has been traversed and is
   stricly non-decreasing order.
List subSortedIT = sorted.subList(0, j+1)
debugger.assertOrder(subSortedIT);
//Count up on the iterator
i++;
```

}

```
/* TERMINATION (Outer-Loop): When the loop
                           terminates, i is equal to sorted.size()
                           meaning the entire array has been traversed
                           and that the guard has been negated.
                                 The negation of the guard implies that
                                    sorted [0 \dots i-1] contains all the
                                    elements of unsorted [0 \dots i-1] */
                         List subSortedOT = sorted.subList(0, i - 1);
                         debugger.assertOrder(subSortedOT);
                         Integer integerI = new Integer(i);
                         Integer sortedSizeO = new Integer(sorted.size());
                         debugger.assertEquals(sortedSizeO, integerI);
                         debugger.assertEquals(unsorted, sorted);
                return sorted;
        }
}
```