0/1 Knapsack

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Abstract

I. MOTIVATION

A 0/1 Knapsack problem is most often described as how does a theif steal the most valuable items with the amount of space he has available to carry them. However, the applications of a 0/1 Knapsack far beyond just this. When solved, a 0/1 Knapsack problem will render the optimal solution, whether it be how one should construct their schedule so they can get the most done in one day, or how to pack a truck such that the fewest amount of trips are taken.

II. Background

Unfortunately, solving a 0/1 Knapsack problem often times proves difficult because in order to find the best solution, all possible solutions must be attempted, a method known as brute force. It is not pratical to brute force a problem of this kind because it would take $\theta n!$ to solve, which is unacceptable for large sets of numbers. However, there are other methods beside brute force that allow for the correct or 'good-enought' solution to be found. One of these methods is known as dynamic programming, which trades off speed for space. With dynamic programming the correct solution can be found in $\theta n * m$ time, where n is the amount of items and m is the size of the knapsack. Another method for solving this type of problem is called a greedy algorithm, which works by making the best decision on which item to take locally. A greedy solution to the 0/1 Knapsack problem traditionally runs in θn time. Unfortunately a greedy solution is not always correct and therefore highly unreliable.

III. Procedure

A multidimentional array insertion sort can be implemented in a multitude of languages using the pseudocode provided in Algorithm 1.

Insertion Sort Pre-Condition: A is an unsorted non-empty array of non-empty arrays containing a comparable data type with a natural order such that v is an index value of the inner array.

Insertion Sort Post-Condition: A' is a permutation of A that is in strictly non-increasing order. **Insertion Sort Outer-Loop Invariant**: The subarray A'[1 ... i - 1] contains all the same elements as the subarray A[1 ... i - 1].

Insertion Sort Outer-Loop Initialization: The outer-loop invariant holds because $A'[1 \dots i-1]$ and $A[1 \dots i-1]$ both contain the same one element.

Insertion Sort Outer-Loop Maintenance: The outer-loop invariant holds because A'[1 ... i - 1] and A[1 ... i - 1] both contain the same elements, although they may be in different orders.

Insertion Sort Outer-Loop Termination: When the outer-loop terminates, i = A.length, which implies that the entire array has been traversed and the guard has been negated. The negation of the guard implies that $A'[1 \dots i-1]$ contains all the elements in $A[1 \dots i-1]$.

Insertion Sort Inner-Loop Invariant: A'[1 ... j] is sorted in strictly non-decreasing order.

Insertion Sort Inner-Loop Initialization: Before the first iteration of the loop, j = 1, meaning the subarray A'[1 ... j] contains exactly one element, which is already sorted.

Insertion Sort Inner-Loop Maintenance: At the beginning of each iteration of the loop the inner-loop invariant holds because j counts down from i, and A'[j+1] is swapped with A'[j] only if

Algorithm 1 Insertion-Sort(A, v)

```
1: procedure Insertion-Sort(A, v)
      if A.length < 2 then
2:
          return A
3:
4:
      end if
      i=2
5:
      while i upto A.length do
6:
          key = A[i][v]
7:
          a = A[i]
8:
          i = i - 2
9:
          while j downto 1 and key > A[j][v] do
10:
             A[j+1] = A[j]
11:
             j = j - 1
12:
          end while
13:
          A[j+1] = a
14:
          i = i + 1
15:
      end while
16:
      return A
17:
18: end procedure
```

A'[j+1] is less than A[j].

Insertion Sort Inner-Loop Termination: The negation of the guard implies that j = A.length and that $A'[1 \dots j]$ has been entirely traversed and sorted in strictly non-decreasing order, which maintains the inner-loop invariant.

Insertion Sort Conclusion: The termination of both the inner and outer loops implies that the entire array has been traversed, A' is a permutation of A containing all the same elements in strictly non-decreasing order. This satisfies the post condition.

A greedy solution to a 0/1 knapsack problem can be implemented in a variety of languages using the pseudocode in Algorithm 2.

Greedy Knapsack Pre-Condition: Weights and Prices both have in them n number of elements Greedy Knapsack Post-Condition: The returned value will be a reasonable solution for the largest value of price combinations such that the aggregate of the corresponding weights does not exceed knapsack capacity c.

A dynamic solution to 0/1 knapsack can be implemented in a variety of languages using the pseudocode provided in Algorithm 3.

Dynamic Knapsack Pre-Condition: Weights and Prices both have n number of elements Dynamic Knapsack Post-Condition: The returned value will be the correct solution for the largest value of price combinations such that the aggregate of the corresponding weights does not exceed knapsack capacity c.

Algorithm 2 Greedy-Knapsack(n, Weights, Prices, c)

```
1: procedure Greedy-Knapsack(n, Weights, Prices, c)
      if n == 1 then
2:
          return Prices[1]
3:
      end if
4:
      if n \le 0 then
5:
          return 0
6:
      end if
7:
      profit = 0
8:
      ratio = newArray
9:
      for v = 1 upto n do
10:
          a = Prices[v]/Weights[v]
11:
          ratio[v] = a
12:
      end for
13:
      ratio = Insertion - Sort(ratio, 1)
14:
      i = 1
15:
      while c > 0 and i < n do
16:
          if c - ratio[i][3] >= 0 then
17:
             profit = profit + ratio[i][2]
18:
              c = c - ratio[i][3]
19:
          end if
20:
          i = i + 1
21:
      end while
22:
      return profit
23:
24: end procedure
```

IV. Testing

A. Testing Plan and Results

All arrays used in testing are Java ArrayList<Integer> except for tab in the dynamic solution, which is a primitive Java nested array. All times are recorded in milliseconds using a stopwatch class borrowed from Robert Sedgwick and Kevin Wayne [?] It is important to note that the stopwatch class used takes the elapsed real-time between the start of the call to the knapsack solution and the end of that call as opposed to taking the elapse processor-time because these tests were run on a multi-core computer. When the number of elements was below 100, the size of the knapsack was 50; when values of 100 and greater were used the size of the knapsack was 2n where n is the number of elements. For example, if the number of elements was 100, the size of the knapsack was 200, in order so the knapsack was sure to be able to make reasonable comparisons.

Table I Greedy Solution Real-Time Run-Time

Number of Elements	Size of Knapsack	Time
0	50	0.0
1	50	0.0
10	50	0.0006
100	200	0.0043
1000	2000	0.0231
10000	20000	0.252
100000	200000	59.1176

Algorithm 3 DYNAMIC-KNAPSACK(n, Weights, Prices, c)

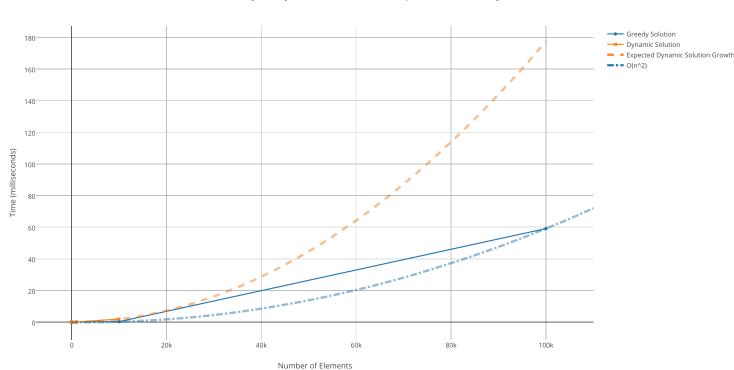
```
1: procedure DYNAMIC-KNAPSACK(n, Weights, Prices, c)
       if n == 1 then
          return Prices[1]
3:
       end if
4:
       if n \le 0 then
5:
          return 0
6:
       end if
7:
       tab[n][c] = newNestedArray
8:
       for x = 1 upto n do
9:
          for y = 1 upto c do
10:
             tab[x][y] = 0
11:
          end for
12:
       end for
13:
       for i = 1 upto n do
14:
          for j = 0 upto c do
15:
             if Weights[i] \le j and Prices[i] + tab[i][j - Weights[i]] > tab[i][j] then
16:
                 tab[i+1][j] = Price[i] + tab[i][j - weights[i]]
17:
             end if
18:
             if !(Weights[i] \le j) and !(Prices[i] + tab[i][j - Weights[i]] > tab[i][j]) then
19:
                 tab[i+1][j] = tab[i][j]
20:
             end if
21:
          end for
22:
       end for
23:
       return tab[n][c]
24:
25: end procedure
```

Table II
DYNAMIC SOLUTION REAL-TIME RUN-TIME

Number of Elements	Size of Knapsack	Time
0	50	0.0001
1	50	0.0
10	50	0.0002
100	200	0.0119
1000	2000	0.0356
10000	20000	1.9243

Table III
Dynamic vs Greedy Solutions Accuracy Comparison

Number of Elements	Size of Knapsack	Percent Error
0	50	0%
1	50	0%
10	50	5.3%
100	200	0.032887%
1000	2000	0.032887%
10000	20000	0.00555725%



Greedy and Dynamic Solutions to 0/1 Knapsack Run-Time Analysis

Figure 1. Java Implementation of Greedy and Dynamic Solutions to 0/1 Knapsack Run-Time Analysis

B. Problems Encountered

One major issue encountered during the development of the dynamic programming solution was that when the second pair of loops would make calls to any of ArrayLists (including tab, which at the time was also an ArrayList implementation) there would be an out of bounds error because of a negative index value. The out of bounds error was eminating from the conditional in the second nested loop. The cause for this error is still unknown. To aid the issue tab was changed from a nested ArrayList to a primitive nested Java array.

Another issue encountered during the development of this program was executing Java assert statements. When assertions were wrapped in parenthesis, they evaluated to a boolean, and the assert would throw an Assertion Error with the explanation 'cannot compare to boolean'. Removing the parenthesis wrapping the assertion fixed this issue.

When the dynamic solution was executed with values of anything above n=10000 and c=20000 Java gave a "Out Of Memory Error: Java heap space". This can be attributed the extreme number of elements being stored in tab. The computer's system monitor reported RAM usage at 95%. Unfortunately, this issue cannot be aided.

The last issue encountered during the development process of this program was generic programming. The program would not compile because not all Objects have a compare method. Adding the statement $\langle T | extends | Comparable \langle T \rangle >$ fixed this issue. [?]

V. Experimental Analysis VI. Conclusion

References

APPENDIX

Listing 1 Driver

```
/*
 @Author: Preston Stosur-Bassett
* @Date: 3, 3, 2015
* @Class: Driver
* @Description: This class will serve as a driver function for our
  Knapsack class
*/
import java.util.ArrayList;
public class Driver {
  public static void main(String args[]) {
    Knapsack theif = new Knapsack();
   DummyData testData = new DummyData();
    Stopwatch watchman = new Stopwatch();
    ArrayList<Integer > prices = new ArrayList<Integer >();
    ArrayList < Integer > weights = new ArrayList < Integer > ();
    if (args [0]. equals ("greedy") == true && args [1] != null && args [2] !=
       null) {
      System.out.println("Running greedy algorithm...");
      int numberOfElements = Integer.parseInt(args[1]);
      int knapsackSize = Integer.parseInt(args[2]);
      System.out.println("Max Knapsack Capacity: "+knapsackSize);
      prices = testData.runArrayList(numberOfElements, 1, 1000, prices);
      weights = testData.runArrayList(numberOfElements, 1, 50, weights);
      System.out.println("Set P:"+prices);
      System.out.println("Set W:"+weights);
      System.out.println("");
      watchman.startTime();
      int totalProfit = theif.greedyKnapsack(numberOfElements, weights,
         prices, knapsackSize);
      double elapsedTime = watchman.elapsedTime();
      Integer totalProfitObject = new Integer(totalProfit);
      System.out.println("The total profit according to this greedy
         algorithm is: "+totalProfitObject);
      System.out.println("Time to Complete: "+elapsedTime);
    else if (args [0]. equals ("dynamic") == true && args [1] != null && args
```

```
[2] != null) {
 System.out.println("Running dynamic algorithm...");
 int numberOfElements = Integer.parseInt(args[1]);
 int knapsackSize = Integer.parseInt(args[2]);
 System.out.println("Max Knapsack Capacity: "+knapsackSize);
 prices = testData.runArrayList(numberOfElements, 1, 1000, prices);
 weights = testData.runArrayList(numberOfElements, 1, 50, weights);
 System.out.println("Set P:"+prices);
 System.out.println("Set W:"+weights);
 System.out.println("");
 watchman.startTime();
 int totalProfit = theif.dynamicKnapsack(numberOfElements, weights,
    prices, knapsackSize);
 double elapsedTime = watchman.elapsedTime();
 Integer totalProfitObject = new Integer(totalProfit);
 System.out.println("The total profit according to this dynamic
    algorithm is: "+totalProfitObject);
 System.out.println("Time to Complete: "+elapsedTime);
else if (args [0]. equals ("compare") == true && args [1] != null && args
   [2] != null) {
 System.out.println("Comparing greedy solution to dynamic solution
 Stopwatch dynamicWatch = new Stopwatch();
 Stopwatch greedyWatch = new Stopwatch();
 int numberOfElements = Integer.parseInt(args[1]);
 int knapsackSize = Integer.parseInt(args[2]);
  if (numberOfElements <= 0) {
    System.out.println("Cannot have 0 or negative number for number
       of elements. Changing to 1.");
   numberOfElements = 1;
 System.out.println("Max Knapsack Capacity: "+knapsackSize);
 prices = testData.runArrayList(numberOfElements, 1, 1000, prices);
 weights = testData.runArrayList(numberOfElements, 1, 50, weights);
 System.out.println("Set P:"+prices);
 System.out.println("Set W:"+weights);
 System.out.println("");
```

```
greedyWatch.startTime();
      int greedyProfit = theif.greedyKnapsack(numberOfElements, weights,
         prices, knapsackSize);
      double greedyTime = greedyWatch.elapsedTime();
      System.out.println("Greedy Solution: $"+greedyProfit+". Completed
         in "+greedyTime+" milliseconds.");
      dynamicWatch.startTime();
      int dynamicProfit = theif.dynamicKnapsack(numberOfElements, weights
         , prices , knapsackSize);
      double dynamicTime = dynamicWatch.elapsedTime();
      System.out.println("Dynamic Solution: $"+dynamicProfit+". Completed
          in "+dynamicTime+" milliseconds.");
      System.out.println("");
      int difference = greedyProfit - dynamicProfit;
      int top = Math.abs(difference);
      int bottom = Math.abs(dynamicProfit);
      double error = new Integer(top).doubleValue() / new Integer(bottom)
         . doubleValue();
      double percentError = error * 100;
      System.out.println("Percent Error: "+percentError+"%");
    }
    else {
      System.out.println("Invalid argument error!!");
      System.out.println("The correct format is: Driver [method] [number
         of elements [size of knapsack]");
 }
                                  Listing 2
                                  Knapsack
* @Auhtor: Preston Stosur-Bassett
* @Description: This class implements a greedy and dynamic solution for
  the 0/1 Knapsack problem. These two solutions will return the max
  value that can be obtained only, and not how to obtain them.
* @Class: Knapsack
* @Date: 3, 3, 2015
*/
import java.util.ArrayList;
import java.lang.Math;
```

```
public class Knapsack<T extends Comparable<T>>> {
  /*
  * @Pre-Condition: <code>weights</code> and <code>prices</code> both
    have in them <code>elems</code> amount of elements
  * @Post-Condition: The returned value will be a reasonable solution for
      the largest value in price where the aggregate of the corresponding
      weight does not excede the <code>backpackSize</code>
  * @Description: greedyKnapsack implements a greedy algorithm to find a
     reasonable solution for a 0/1 Knapsack problem
  * @param int elems is the amount of elements in both <code>ArrayList<
    Integer > weights </code > and <code > ArrayList < Integer > prices </code >
  * @param ArrayList<Integer> weights is a non-empty ArrayList of Integer
      objects with exactly <code>elems</code> amount of elements in it
    and contains absolutely no zeros
  * @param ArrayList < Integer > prices is a non-empty ArrayList of Integer
     object with exactly <code>elems</code> amount of elements in it
  * @param int backpackSize is the maximum value of weights that the
    knapsack can hold
   @return int profit is a reasonable solution to the given 0/1 Knapsack
      problem
  //INVARIANT (First Loop): TODO: write this
  //INVARIANT (Second Loop): TODO: write this
  public int greedyKnapsack(int elems, ArrayList<Integer> weights,
     ArrayList < Integer > prices, int backpackSize) {
    int returnValue;
    if (elems = 1)  {
      returnValue = prices.get(0).intValue();
    else if (elems < 1) {
      returnValue = 0;
    else {
      int profit = 0;
      ArrayList < ArrayList < Integer >> ratioListings = new ArrayList <
         ArrayList<Integer>>(elems);
      int v = 0;
      /*INITIALIZATION (First Loop): */
      while (v < elems) {
        /*MAINTANENCE (First-Loop): */
        int ratio = prices.get(v).intValue() / weights.get(v).intValue();
        ArrayList<Integer> innerRatioListing = new ArrayList<Integer>(3);
        innerRatioListing.add(new Integer(ratio));
        innerRatioListing.add(new Integer(prices.get(v).intValue()));
        innerRatioListing.add(new Integer(weights.get(v).intValue()));
        ratioListings.add(innerRatioListing);
```

```
v++;
    /*TERMINATION (First Loop): */
    Sort sorter = new Sort();
    ratioListings = sorter.insertionSortNestedArray(ratioListings, 0);
    int i = 0;
    /*INITIALIZATION (Second Loop): */
    while (backpackSize > 0 && i < elems) {
      /*MAINTENANCE (Second Loop): */
      if (backpackSize - ratioListings.get(i).get(2).intValue() >= 0) {
        profit = profit + ratioListings.get(i).get(1).intValue();
        backpackSize = backpackSize - ratioListings.get(i).get(2).
           int Value ();
      i++;
    /*TERMINATION (Second Loop): */
    returnValue = profit;
  return returnValue;
* @Pre-Condition: <code>weights</code> and <code>prices</code> both
  have in them <code>elems</code> amount of elements
* @Post-Condition: The returned value will be the correct solution for
  the largest value in price where the aggregate of the corresponding
  weights do not excede the <code>backpackSize</code>
* @Description: dynamicKnapsack implements a dynamic programming
  solution to find the correct answer for a 0/1 Knapsack problem
* @param int elems is the amount of elemnts in both <code>ArrayList<
  Integer > weights </code > and <code > ArrayList < Integer > prices </code >
* @param ArrayList<Integer> weights is a non-empty ArrayList of Integer
    objects with esactly <code>elems</code> amount of elements in it
  and contains absolutely no zeros
* @param ArrayList < Integer > prices is a non-empty ArrayList of Integer
  objects with exactly <code>elemes</code> amount of elemnts in it
* @param int backpackSize is the maximum value of the weights that the
  knapsack can hold
* @return int return Value is the correct value of the maximum amount of
    values you can get from prices where their correcsponding weights
  do not excede the backpackSize
//INVARIANT (First Outer-Loop): TODO: write this
```

```
//INVARIANT (First Inner-Loop): TODO: write this
//INVARIANT (Second Outer-Loop): TODO: write this
//INVARIANT (Second Inner-Loop): TODO: write this
public int dynamicKnapsack(int elems, ArrayList<Integer> weights,
  ArrayList < Integer > prices, int backpackSize) {
 int return Value = 0;
 backpackSize = backpackSize + 1;
  if (elems = 1)  {
    returnValue = prices.get(0).intValue();
 else if (elems \leq 0) {
    returnValue = 0;
  else {
    int [][] tab = new int [elems] [backpackSize];
    int x = 0:
    int y = 0;
    /*INITIALIZATION (First Outer-Loop): TODO: write this */
    while (x < elems) {
      /*MAINTENANCE (First Outer-Loop): TODO: write this*/
      /*INITIALIZATION (First Inner-Loop): TODO: write this */
      while (y < backpackSize) {
        /*MAINTENANCE (First Inner-Loop): TODO: write this */
        tab[x][y] = 0;
        y++;
      /*TERMINATION (First Inner-Loop): TODO: write this */
     x++;
    /*TERMINATION (First Outer-Loop): TODO: write this */
    // Second Outer For Loop Initialization: TODO: write this
    for (int i = 0; i < elems - 1; i++) {
      // Second Outer For Loop Maintenance: TODO: write this
      // Second Inner For Loop Initialization: TODO: write this
      for (int j = 0; j < backpackSize; j++) {
        // Second Inner For Loop Maintenance: TODO: write this
        if (weights.get(i).intValue() <= j && prices.get(i).intValue() +
            tab[i][j - weights.get(i).intValue()] > tab[i][j]) {
          tab[i+1][j] = prices.get(i).intValue() + tab[i][j - weights.
             get(i).intValue();
        else {
          tab[i+1][j] = tab[i][j];
      // Second Inner For Loop Termination: TODO: write this
```

```
// Second Outer For Loop Termination: TODO: write this
      returnValue = tab [elems - 1][backpackSize - 1];
    return returnValue;
}
                                   Listing 3
                                    Sort
/*
        @Author: Preston Stosur-Bassett
        @Date: Jan 24, 2015
        @Class: Sort
        @Description: This class will contain many methods that will sort
    generic data types using common sorting algorithms.
*/
import java.util.ArrayList;
import java.util.List;
public class Sort<T extends Comparable<T>>> {
        /*
                @Pre-Condition: ArrayList<T> is a non-empty set of data
           where T is a comparable data type with a natural order
                @Post-Condition: Each parent node is more extreme than
           its child node.
                @Description: heapify is a helper method for heapSort
           that keeps the heap in order so that the root node is the most
            extreme element in the heap.
                @param ArrayList<T> unsorted is a non-empty set of data
           where T is a comparable data type with a natural order
                @param int i
                @param int total
        *
                @return ArrayList<T> unsorted
        private ArrayList<T> heapify(ArrayList<T> unsorted, int i, int
           total) {
                int left = i * 2;
                int right = left + 1;
                int original I = i;
                if (left <= total && unsorted.get(left).compareTo(unsorted
                   . get(i)) > 0)  {
                         i = left;
                if (right <= total && unsorted.get(right).compareTo(
                   unsorted.get(i) > 0) {
```

```
i = right;
        if (i != originalI) {
                T tmp = unsorted.get(originalI);
                unsorted.set(originalI, unsorted.get(i));
                unsorted.set(i, tmp);
                unsorted = heapify (unsorted, i, total);
        }
        return unsorted;
}
        @Pre-Condition: ArrayList<T> unsorted is a non-empty
  ArrayList <T> where T is a comparable data type with a natural
  order.
        @Post-Condition: ArrayList<T> sorted is a permutation of
  unsorted (it contains all the same elements) in stricly non-
  decreasing order
        @Description: heapSort will sort a given set of data in
  an ArrayList<T> in strictly non-decreasing order using the
  heap sort method.
        @param ArrayList<T> unsorted is a non-empty ArrayList<T>
  where T is a comparable data type with a natural order
        @return ArrayList<T> sorted is a permutation of unsorted
  in strictly non-decreasing order
//Invariant for First While Loop: unsorted[i] is the parent
  element in a heap
//Invariant for Second While Loop: All elements in unsorted
  greater than the index value of y are in stricly non-
  decreasing order
public ArrayList<T> heapSort(ArrayList<T> unsorted) {
        //Debug
        Debug debugger = new Debug();
        int arrSize = unsorted.size() - 1;
        int i = arrSize / 2;
        //Initialization: Our invariant holds true before the
           first iteration of the loop because unsorted[i] must
           have child elements
        debugger.assertChildren(unsorted, i);
        while (i >= 0)
                //Maintanance: Our invariant holds true at the
                   beginning of each iteration of the loop
                   because unsorted[i] must have children
                   elements
                debugger.assertChildren(unsorted, i);
                unsorted = heapify (unsorted, i, arrSize);
```

i --;

```
//Termination: Our invariant holds true at the
           termination of the loop because i will be the smallest
            index value of the loop and must have children
           elements
        debugger.assertChildren(unsorted, i);
        int y = arrSize;
        //Initialization: Our invariant holds vacuously true
           before the first iteration of the loop because there
           are no elements in unsorted that are at an index value
            greater than y.
        debugger.assertStrictLess(arrSize, y+1);
        while (y > 0) {
                //Maintanance: Our invariant holds true at the
                   beginning of each iteration of the loop
                   because all elements greater than y are in
                   strictly non-decreasing order
                T \text{ tmp} = unsorted.get(0);
                unsorted.set(0, unsorted.get(y));
                unsorted.set(y, tmp);
                arrSize --;
                unsorted = heapify (unsorted, 0, arrSize);
                y--;
        //Termination: Our invariant holds true at the
           termination of the loop because y decreases as each
           largest element is moved to the end of the list until
           the entire array has been traversed, so that all
           elements greater than y are in stricly non-decreasing
           order
        debugger.assertOrder(unsorted);
        ArrayList<T> sorted = unsorted;
        return sorted;
}
/*
        @Pre-Condition: ArrayList<T> left is a non-empty sorted
  array in stricly non-decreasing order where T is a comparable
  data type with a natural order
        @Post-Condition: ArrayList<T> right is a non-empty sorted
   array in strictly non-decreasing order where T is a
  comparable data type with a natural order
        @Description: mergeTogether is used by the mergeSort
```

```
method to recombine the left and right sections of the
  ArrayList<T> that is being sorted by merge sort. Note that
  this is a helper method for the mergeSort method, and should
  not be called externally of this class.
        @param ArrayList<T> left a non-empty ArrayList<T> where T
   is a comparable data type with a natural order.
        @param ArrayList<T> right a non-empty ArrayList<T> where
  T is a comparable data type with a natural order.
        @return ArrayList<T> combined should contain all the
  elements of left and right in stricly non-decreasing order
//Invariant for First While Loop: combined contains x number of
  elements where x is the sum of i and y and those elements are
  contained in left [0 ... i] or right [0 ... y] in stricly non-
  decreasing order
//Invaraint for Second While Loop: combined contains x number of
  elements where x is greater than or equal to i and those
  elements are contined in left [0 ... i] in stricly non-
  decreasing order
//Invariant for Third While Loop: combined contains x number of
  elements where x is greater than or equal to y and those
  elements are contained in right [0 ... y] in stricly non-
  decreasing order
private ArrayList<T> mergeTogether(ArrayList<T> left, ArrayList<T
  > right) {
        ArrayList<T> combined = new ArrayList<T>();
        int i = 0;
        int y = 0;
        int x = 0:
        //Debug
        Debug debugger = new Debug();
        //Initialization: Our invariant holds true vacuously
           before the first execution of the loop because x, i,
           and y are all equal to zero, combined is empty and
           therefore in order
        debugger.assertEquals(0, i);
        debugger.assertEquals(0, y);
        debugger.assertEquals(0, x);
        debugger.assertEquals(i, combined.size());
        while (left.size() != i && right.size() != y) {
                //Maintanance: Our invariant holds true at the
                   beginning of each iteration of the loop
                   because x is incremented whenever i or y is
                   incremented and elements are added to combined
                    from left and right in order
                debugger.assertEquals(i+y, x);
                debugger.assertEquals(x, combined.size());
```

```
debugger.assertOrder(combined);
        debugger.assertContains(right, left, combined);
        if(left.get(i).compareTo(right.get(y)) < 0) {
                combined.add(x, left.get(i));
                i++;
                x++;
        else {
                combined.add(x, right.get(y));
                y++;
                x++;
        }
//Termination: Our invariant holds tur at the termination
   of the loop because x has been incremented whenever i
   or y has been incremented, and elements are added to
  combined from left and right in order
debugger.assertEquals(i+y, x);
debugger.assertEquals(x, combined.size());
debugger.assertOrder(combined);
debugger.assertContains(right, left, combined);
//Initialization: Our invariant holds true before the
   first execution of the loop because x has been
  incremented whenever i has been incremented and
  elements have been added to combined from left in
  order
debugger.assertGreatEquals(x, i);
debugger.assertEquals(x, combined.size());
debugger.assertOrder(combined);
debugger.assertContains(left, combined);
while (left.size() != i) {
        //Maintanance: Our invariant holds true at the
           beginning of each iteration of the loop
           because x has been incremented whenever i is
           incremented and elements have been added to
           combined from left in order
        debugger.assertGreatEquals(x, i);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(left, combined);
        combined.add(x, left.get(i));
        i++;
        x++;
//Termination: Our invariant holds true at the
   termination of the loop because x has been incremented
```

```
added to combined from left in order
        debugger.assertGreatEquals(x, i);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(left, combined);
        //Initialization: Our invariant holds true before the
           first execution of the loop because x has been
           incremented whenever y has been incremented and
           elements have been added to combined from right in
           order
        debugger.assertGreatEquals(x, y);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(right, combined);
        while (right.size() != y) {
                //Maintanance: Our invariant holds tur at the
                   beinning of each iteration of the loop because
                    x has been incremented whenever y is
                   incremented and elements have been added to
                   combined from right in order
                debugger.assertGreatEquals(x, y);
                debugger.assertEquals(x, combined.size());
                debugger.assertOrder(combined);
                debugger.assertContains(right, combined);
                combined.add(x, right.get(y));
                y++;
                x++;
        //Termination: Our invariant holds true at the
           terminatino of the loop because x has been incremented
            whenever y was incremented and elements have been
           added to combined from right in order.
        debugger.assertGreatEquals(x, y);
        debugger.assertEquals(x, combined.size());
        debugger.assertOrder(combined);
        debugger.assertContains(right, combined);
        return combined;
}
        @Pre-Condition: ArrayList<T> unsorted is a set of data
  type T, where T is a Comparable data type with a natural order
        @Post-Condition: ArrayList<T> returnValue is a
  permutation of unsorted in strictly non-decreasing order.
```

whenever i was incremented and elements have been

```
@Description: mergeSort will sort a given set of data in
  ArrayList<T> using the merge sort method
        @param a non-empty ArrayList<T> unsorted where T is a
  Comparable data type with a natural order
        @return ArrayList<T> returnValue which is a permutation
   of unsorted, in strictly non-decreasing order,
//Invariant for First While Loop: left contains i elements, all
   of which can be found in sorted
//Invariant for Second While Loop: right contains y elements, all
    of which can be found in sorted
public ArrayList<T> mergeSort(ArrayList<T> unsorted) {
        ArrayList<T> sorted = unsorted;
        ArrayList<T> left = new ArrayList<T>();
        ArrayList < T > right = new ArrayList < T > ();
        ArrayList <T> returnValue;
        //Debug
        Debug debugger = new Debug();
        debugger.turnOn();
        if(sorted.size() \ll 1) {
                returnValue = sorted;
        }
        else {
                int mid = (sorted.size() / 2);
                int i = 0:
                //Initialization: Our invariant holds true because
                    i is zero and left contains 0 elements before
                    the first iteration of the loop.
                debugger.assertEquals(i, left.size());
                while (i < mid)
                        //Maintanance: Our invariant holds true
                            because i is increased at the same
                           rate elements are added to left from
                            the same i index in sorted
                         debugger.assertEquals(i, left.size());
                         debugger.assertContains(sorted, left);
                        T \text{ temp} = \text{sorted.get(i)};
                         left.add(temp);
                         i++;
                //Termination: Our invariant holds true because i
                    has been incremented at the same rate
                   elements are added to left from the same index
                    i in sorted
                debugger.assertEquals(i, left.size());
```

```
debugger.assertContains(sorted, left);
                int y = mid;
                //Initialization: Our invariant holds true
                   because i is zero and right contains 0
                   elements before the first iteration of the
                   loop.
                debugger.assertEquals(y, right.size());
                while (y < sorted.size()) {
                        //Maintanance: Our invariant holds true
                           because y is increased at the same
                           rate elements are added to right from
                           the same y index in sorted.
                        debugger.assertEquals(y, right.size());
                         debugger.assertContains(sorted, right);
                        T \text{ temp} = \text{sorted.get}(y);
                        right.add(temp);
                        y++;
                //Termination: Our invariant holds true because y
                    has been incremented at the same rate
                   elements are added to left from the same index
                    y in sorted
                debugger.assertEquals(y, right.size());
                debugger.assertContains(sorted, right);
                left = mergeSort(left);
                right = mergeSort(right);
                returnValue = mergeTogether(left, right);
        return return Value;
}
/*
        @Pre-Condition: ArrayList<T> unsorted is an unsorted
   ArrayList of a comparable data type that is non-empty
        @Post-Condition: ArrayList<T> will return a permutation
   of <code>unsorted</code> that will be in increasing order
        @Description: insertionSort will sort an ArrayList of
*
   generic type T in increasing order using an insertion sort
        @param ArrayList<T> unsorted is a non-empty unsorted
  array list of T, where T is a comparable type
        @return sorted is a permutation of <code>unsorted</code>
  where all the elements are sorted in increasing order
// INVARIANT (Outer-Loop): The pre condition implies that sorted
   [0 \ldots i-1] will contain all the same data as unsorted [0 \ldots
```

```
i - 1].
// INVARIANT (Inner-Loop): sorted[0 ... j] is sorted in stricly
  non-decreasing order.
public ArrayList<T> insertionSort(ArrayList<T> unsorted) {
        Debug debugger = new Debug<List<T>>();
        debugger.turnOn();
        ArrayList<T> sorted = unsorted;
        if (sorted.size() > 1) {
                int i = 1;
                /* INITIALIZATION (Outer-Loop): The invariant
                   holds because i = 1, and there is one element
                   in the subarray of sorted [0 \dots i-1] and
                   unsorted [0 \ldots i-1], */
                List < T > subSortedOI = sorted.subList(0, i - 1);
                List <T> subUnsortedOI = unsorted.subList(0, i -
                debugger.assertEquals(subUnsortedOI, subSortedOI)
                while (i < sorted.size()) {
                        /* MAINTENANCE (Outer-Loop): At the
                           beginning of each iteration of the
                           loop, the loop invariant is maintained
                            because the subarray of sorted [0 ...
                           i-1 contains all the same elements
                                 unsorted [0 \ldots i-1] */
                        List <T> subSortedOM = sorted.subList(0, i
                            -1):
                        List <T> subUnsortedOM = unsorted.subList
                           (0, i-1);
                        debugger.assertEquals(subUnsortedOM,
                           subSortedOM);
                        T value = sorted.get(i);
                        int i = i - 1;
                        // INITIALIZATION (Inner-Loop): Before
                           the first iteration of the loop, j =
                           0, the subarray of sorted [0 \dots 0]
                           contains one elements and therefore
                           the invariants holds vacuously.
                         List subSortedII = sorted.subList(0, j);
                        debugger.assertOrder(subSortedII);
                         while (j \ge 0 \&\& (value.compareTo(sorted.
                           get(j) < 0)
                                 // MAINTENANCE: (Inner-Loop): At
                                    the beginning of each
                                    iteration sorted [0 \dots j] is
                                    sorted in stricly non-
```

```
decreasing order
                                List subSortedIM = sorted.subList
                                    (0, j);
                                debugger.assertOrder(subSortedIM)
                                sorted.set(j+1, sorted.get(j));
                        }
                        sorted set (j+1, value);
                        // TERMINATION (Inner-Loop): The negation
                            of the guard implies that the sorted
                           [0 ... j] has been traversed and is
                           stricly non-decreasing order.
                        List subSortedIT = sorted.subList(0, j+1)
                        debugger.assertOrder(subSortedIT);
                        //Count up on the iterator
                        i++;
                /* TERMINATION (Outer-Loop): When the loop
                   terminates, i is equal to sorted.size()
                   meaning the entire array has been traversed
                   and that the guard has been negated.
                        The negation of the guard implies that
                           sorted [0 \dots i-1] contains all the
                           elements of unsorted [0 \dots i-1] */
                List subSortedOT = sorted.subList(0, i - 1);
                debugger.assertOrder(subSortedOT);
                Integer integer I = new Integer (i);
                Integer sortedSizeO = new Integer(sorted.size());
                debugger.assertEquals(sortedSizeO, integerI);
                debugger.assertEquals(unsorted, sorted);
        return sorted;
}
/*
        @Pre-Condition: ArrayList<ArrayList<T>> unsorted is an
  unsorted a nested non-empty ArrayList of a non-empty ArrayList
   (in tabular format) of a comparable data type with a natural
  order where sortingIndex is an index value of the nest
  ArrayList.
        @Post-Condition: ArrayList<ArrayList<T>> will return a
  permutation of <code>unsorted</code> that will be in stricly
  non-increasing order.
        @Description: insertionSortNestedArray will sort a nested
   ArrayList in tabular format of a comparable data type and
```

given a specific index value of the inner ArrayList will sort the inner ArrayLists into stricly non-increasing order within the outer ArrayList * @param ArrayList < ArrayList < T>> list is a non-empty unsorted nested ArrayList of ArrayList of data type T, where T is a comparable data type with a natural order. @param int sortingIndex is an index value of the inner ArrayList to use for sorting comparisons @return ArrayList<ArrayList<T>>> list is a permutation of <code>unsorted</code> where all the elements in the outer ArrayList are sorted in stricly non-increasing order by inner ArrayLists index value of sortingIndex //INVARIANT (Outer-Loop): The Pre-Condition implies that A[0 ... i-1] will contain all the same data as A'[0 ... i-1] //INVARIANT (Inner-Loop): A[0 ... j] is sorted in stricly nonincreasing order public ArrayList<ArrayList<T>>> insertionSortNestedArray(ArrayList <ArrayList<T>> unsorted, int sortingIndex) { Debug debugger = new Debug<List<T>>();ArrayList<ArrayList<T>> list = unsorted; if(list.size() > 1) { int i = 1; /*INITIALIZATION (Outer-Loop): Before the first iteration of the loop the invariant holds beause i = 1, and there is one elment in the subarray of $A[0 \ldots i-1]$ and $A'[0 \ldots i-1]$ */ List < ArrayList < T>> subSortedOI = list.subList(0, i - 1);List < ArrayList < T>> subUnsortedOI = unsorted. subList(0, i-1);debugger.assertEquals(subUnsortedOI, subSortedOI) while (i < list.size()) { /*MAINTENANCE (Outer-Loop): At the beginning of each iteration of the loop, the loop invariant is maintained because the subarray of A'[0 ... i -1] contains all the same elements as A $[0 \dots i - 1] */$ List < ArrayList < T>> subSortedOM = list. subList(0, i-1);List < ArrayList < T>> subUnsortedOM = unsorted.subList(0, i - 1); debugger.assertEquals(subUnsortedOM,

subSortedOM);

```
ArrayList<T> currentElement = list.get(i)
       T value = list.get(i).get(sortingIndex);
        int j = i - 1;
        /*INITIALIZATION (Inner-Loop): Before the
            first iteration of the loop, j = 0,
           the subarray of sorted [0 ... 0]
           contains one element and therefore the
            invariants hold vacuously. */
        List subSortedII = list.subList(0, j);
        debugger.assertOrder(subSortedII,
           sortingIndex);
        while (j >= 0 && (value.compareTo(list.get
           (j).get(sortingIndex)) > 0)
                /*MAINTENANCE (Inner-Loop): At
                   the beginning of each
                   iteration A[0 \dots j] is sorted
                    in stricly non-increasing
                   order */
                List subSortedIM = list.subList
                   (0, j);
                debugger.assertOrder(subSortedIM,
                    sortingIndex);
                list.set(j+1, list.get(j));
                j --;
        /*TERMINATION (Inner-Loop): The negation
           of the gaurd implies that A'[0 ... j]
           has been entirely traversed and is in
           stricly non-increasing order. */
        List subSortedIT = list.subList(0, j+1);
        debugger.assertOrder(subSortedIT,
           sortingIndex);
        list.set(j+1, currentElement);
        i++;
/*TERMINATION (Outer-Loop): When the loop
  terminates, i is equal to A'. length meaning
  the entire array has been traversed and that
  the guard has been negated.
```

```
List subSortedOT = list.subList(0, i - 1);
debugger.assertOrder(subSortedOT, sortingIndex);
debugger.assertEquals(new Integer(list.size()),
   new Integer(i));
debugger.assertEquals(unsorted, list);
```

}

```
return list;
}
```

The class Stopwatch has been altered from its original form.

Listing 4 Stopwatch

```
Compilation: javac Stopwatch.java
*
/**
   The <tt>Stopwatch</tt> data type is for measuring
   the time that elapses between the start and end of a
   programming task (wall-clock time).
*
   See {@link StopwatchCPU} for a version that measures CPU time.
*
   @author Robert Sedgewick
*
   @author Kevin Wayne
   @update @ 5.3.15 by Preston Stosur-Bassett, added start method.
*
*/
public class Stopwatch {
   private long start;
   * Starts the stopwatch timer
   */
   public void startTime() {
     start = System.currentTimeMillis();
   /**
    * Returns the elapsed time (in seconds) since this object was
      created.
   public double elapsedTime() {
      long now = System.currentTimeMillis();
      return (now - start) / 1000.0;
   }
}
```

Listing 5 Debug

```
/*
        @Author Preston Stosur-Bassett
        @Date Jan 21, 2015
        @Class Debug
        @Description This class will help debugging by being able to turn
   on and turn off debug messages easily
*/
import java.util.List;
import java.util.ArrayList;
public class Debug<T> {
        boolean debugOn; //Variable to keep track of whether or not debug
            is on
        /*
                @Description constructor method that sets the default
           value of debugOn to false so that debug statements will not
           automatically print
        */
        public void Debug() {
                debugOn = false;
        }
        /*
                @Description turn on debugging print statements
        public void turnOn() {
                debugOn = true;
        }
        /*
                @Description turn off debugging print statements
        */
        public void turnOff() {
                debugOn = false;
        }
        /*
                @Description will print messages only when debugOn
           boolean is set to true
                @param String message the string to print when debugging
           is turned on
        */
        public void print(T message) {
                if (debugOn == true) {
                         System.out.println(message);
```

```
}
}
/*
        @Pre-Condition < code>T expected </code> and <code>T actual
  </code> are both of the same type T
        @Post-Condition If <code>T expected </code> and <code>T
  actual </code> are found to be equal, the program moves on,
  otherwise the program halts with <code>AssertionError</code>
  is thrown
        @Description runs an assert statement against an expected
   value and the actual value that are passed as parameters only
   when <code>debugOn == true </code>
        Oparam T expected the expected value to assert against
  the actual value
        Oparam T actualt he actual value to assert against the
  expected value
*/
public void assertEquals(T expected, T actual) {
        if (debugOn == true) {
                assert actual.equals(expected);
        }
}
* @Pre-Condition: <code>List<Integer> actual</code> is a iterable
   list of Integer objects
        @Post-Conditions: If the List of Integer objects is in
  stricly non-decreasing order, the program moves on normally,
  if not, the program halts with an <code>AssertionError</code>
* @Description: runs an assertion statement against a list of
  Integer objects to ensure that for <code>k = actual.size(); A
  |k - 2| \le A[k - 1]; </code>
        @param List < Integer > actual the list to assert is in
  stricly non-decreasing order
public void assertOrder(List<Integer> actual) {
        if (debugOn == true) {
                int i = actual.size();
                while (i > 1)
                        assert actual.get(i - 1).compareTo(actual
                           . get(i - 2)) >= 0;
                        i --;
                }
}
/*
```

```
@Pre-Condition: <code>List<ArrayList<Integer>> actual</
  code> is an iterable list of ArrayList of Integer objects
  where sortingIndex is an index value of the ArrayList
        @Post-Condition: If the LIst of ArrayList of Integer
   objects is in stricly non-increasing order, the program moves
  on normally, if not, the program halts with an <code>
   AssertionError</code>
        @Description: runs an assertion statement against a list
   of ArrayList of Integer objects to ensure that for <code>k =
   actual.size(); A[k-2] >= A[k-1]; </code>
        @param List < Array List < Integer >> actual the list to assert
    is in stricly non-decreasing order
        @param int sortingIndex the index value to make sorting
   comparisons from
*/
public void assertOrder(List<ArrayList<Integer>> actual, int
   sortingIndex) {
        if (debugOn == true) {
                int i = actual.size();
                while (i > 1)
                         assert actual.get (i - 1).get (sortingIndex)
                            ).compareTo(actual.get(i - 2).get(
                            sortingIndex) <= 0;
                         i --:
                }
        }
}
/*
        @Pre-Condition: <code>ArrayList<Integer> actual</code> is
   an ArrayList of Integer Objects
        @Post-Condition: If the ArrayList of Integer Objects is
   in stricly non-decreasing order, the program moves on normally
   , if not, the pgoram halts with an <code>AssertionError</code>
        @Description: runs an assertion statement against an
   ArrayList of Integer Objects to ensure that for \langle code \rangle k =
   actual.size(); A[k-2] \le A[k-1]; </code>
        @param ArrayList < Integer > actual the ArrayList to assert
   is in stricly non-decreasing order
public void assertOrder(ArrayList<Integer> actual) {
        if (debugOn == true) {
                int i = actual.size();
                while (i > 1)
                         assert actual.get(i - 1).compareTo(actual
                            . get(i - 2)) >= 0;
                         i --;
                }
```

```
}
}
/*
        @Pre-Condition: ArrayList is an ArrayList of Integers and
   i is less than or equal to half of the size of actual
        @Post-Condition: If elements exist past i the assertion
  holds
        @Description: runs an assertion statement against an
  ArrayList of Integer Objects to ensure that there are children
   nodes of actual[i].
        @param ArrayList<Integer> actual the array to test
  against
        Operam int i the index value to check has children nodes.
*/
public void assertChildren(ArrayList<Integer> actual, int i) {
        if (debugOn == true) {
                assert actual.size() > i;
        }
}
/*
        @Pre-Condition: actual and expected both contain Integer
  Objects
        @Post-Condition: If all the elements inside of the actual
   arraylist are also contained in the expected arraylist, then
  the assertion holds true
        @Description: Tests to ensure a given ArrayList of
  Integer Objects contains all the elements of another given
  ArrayList of Integer Objects
        @param ArrayList<Integer> expected the list to check
  contains against
        @param ArrayList<Integer> actual the list to check to
  make sure all its elements are contained in the other
  arravlist
*/
public void assertContains(ArrayList<Integer> expected, ArrayList
  <Integer> actual) {
        if (debugOn == true) {
                for (int i = 0; i < actual.size(); i++) {
                        assert expected.contains(actual.get(i));
                }
        }
}
/*
        @Pre-Condition: expectedOne, expectedTwo, and actual all
  contain Integer Objects
        @Post-Condition: If all the elemts inside of the actual
  ArrayList are also contined in either the expectedOne
```

```
ArrayList or the expectedTwo ArrayList, then the assertion
  holds true
        @Description: Tests to ensure a given ArrayList of
  Integer Objects contains all the elements of another given
  ArrayList of Integer Objects
        @param: ArrayList<Integer> expectedOne one of the lists
  to check to see if the given ArrayList actual's elements are
  contained in
        @param: ArrayList < Integer > expected Two one of the lists
  to check to see if the given ArrayList actual's elements are
  contained in
* @param: ArrayList<Integer> actual the list to check of make
  sure all its elements are contained in either expectedOne or
  expectedTwo
*/
public void assertContains(ArrayList<Integer> expectedOne,
  ArrayList<Integer> expectedTwo, ArrayList<Integer> actual) {
        if (debugOn == true) {
                for (int i = 0; i < actual.size(); i++) {
                        assert expectedOne.contains(actual.get(i)
                           ) | expectedTwo.contains(actual.get(i
                           ));
                }
        }
}
* @Description: asserts that the first arguement is stricly
  greator than the second arguement
        Oparam int large an integer primative value to assert is
  strictly greator than the second arguement
        Oparam int small an integer primative value to assert the
    first arguement is strictly greator than.
public void assertStrictGreat(int large, int small) {
        if (debugOn == true) {
                assert large > small;
        }
}
/*
        @Description: asserts that the first arguement is
  strictly less than the second arguement
        @param int small an integer primative value to assert is
  stricly less than the second arguement
        Oparam int large an integer primative value to assert the
    first arguement is strictly less than.
*/
public void assertStrictLess(int small, int large) {
        if (debugOn == true) {
```

```
assert small < large;
                }
        }
        /*
                @Description: asserts that the first arguement is greator
            than or equal to the second arguement
                Oparam int large an integer primative value to assert is
           greator than or equal to the second arguement
                Oparam int small an integer primative value to assert the
            first arguement is greator than or equal to.
        */
        public void assertGreatEquals(int large, int small) {
                if (debugOn == true) {
                         assert large >= small;
                }
        }
        /*
                @Description: asserts that the first argument is less
           than or equal to the second arguement
                Oparam int small an integer primative value to assert is
           less than or equal to the second arguement
                Operam int large an integer primative value to assert the
            first arguement is less than or equal to.
        public void assertLessEquals(int small, int large) {
                if (debugOn == true) {
                         assert small <= large;
                }
        }
}
                                   Listing 6
                                  DummyData
/*
        @Author Preston Stosur-Bassett
        @Date Jan 25, 2015
        @Class DummyData
        @Description This class contains methods to generate dummy data
   given a set of parameters.
*/
import java.util.ArrayList;
import java.util.Random;
public class DummyData {
        /*
                @Description runArrayList<Integer> will take an ArrayList
```

```
of Integer Objects and add a given amount of values to it
        @param int end the ending value to denote when to stop
  adding to the array list
        Oparam int min the minimum value of the randomly
  generated data.
        Operam int max the maximum value of the randomly
  generated data.
        @param ArrayList < Integer > list the list to add value to
  and return
        @return ArrayList<Integer> the list after it has been
  updated with the randomly generated data
*/
public static ArrayList<Integer> runArrayList(int end, int min,
  int max, ArrayList<Integer> list) {
        Random random = new Random();
        Debug debugger = new Debug();
        int start = 0;
        // INVARIANT: A.length >= start
        // INITIALIZATION: start = 0, A.length can be longer than
            0 when initially passed, but not smaller, so our
           invariant holds
        debugger.assertGreatEquals(list.size(), start);
        while (start < end) {
                // MAINTANANCE: At the beginning of each
                   iteration, one element was added to A and
                   start was increased by one, therefore, our
                   invariant holds true.
                debugger.assertGreatEquals(list.size(), start);
                Integer intToAdd = new Integer (random.nextInt((
                   \max - \min + 1) + \min);
                if(intToAdd != 0) {
                        list.add(intToAdd);
                        start++;
                }
        /*TERMINATION: The negation of the guard implies that (
           end - start) number of elements have been added to A,
           since start is initialized as 0 at the beginning of
           the method and is
                incremented by 1 each iteration of the loop,
                   which means that start amount of elements have
                    been added to A, and so our invariant holds
                   true.
                           */
        debugger.assertGreatEquals(list.size(), start);
        return list;
}
/*
```

```
@Description: runArrayList<String> will take an ArrayList
   of String Objects and add a given amount of String numerical
  values to it
        @param int end the ending value to denote when to stop
  adding to the array list
        @param ArrayList < String > list the list to add String
  values to and return
        @return ArrayList<String> the list after it has been
  updated with the randomly generated numerical String values
*/
public static ArrayList<String> runArrayList(int end, ArrayList<
  String > list) {
        Random random = new Random();
        Debug debugger = new Debug();
        int start = 0;
        // INVARIANT: A.length >= start
        // INITIALIZATION: Before the first iteration of the loop
           , start = 0 and A.length cannot be less than 0, so our
            invariant holds true
        debugger.assertGreatEquals(list.size(), start);
        while (start < end) {
                // MAINTENANCE: At the beginning of each
                   iteration of the loop our invariant holds
                   because for each iteration of the loop one
                   element is added to A and start is incremented
                debugger.assertGreatEquals(list.size(), start);
                Integer intToString = new Integer (random.nextInt
                   ((1000000 - 1) + 1));
                String intString = String.valueOf(intToString);
                list.add(intString);
                //Count up on the iterator
                start++;
        /* TERMINATION: The negation of the guard implies that (
           end - start) number of elements have been added to A,
           since start is initialized as 0 at the beginning of
           the method and is
                        incremented by 1 each iteration of the
                           loop, which means that start amount of
                            elements have been added to A, and so
                            our invariant holds true. */
        debugger.assertGreatEquals(list.size(), start);
        return list;
}
/*
        @Description: identicalElement will take an element and
```

}

```
add it to the ArrayList<Integer> for a given amount of times
        @param int end the ending value to denote when to stop
  adding elements to the array
        Oparam int element the element to add over and over again
   to the array
        @param ArrayList<Integer> list the list to add elements
  to
        @return ArrayList<Integer> the list after it has been
  updated with the given data
*/
public static ArrayList < Integer > identicalElement (int end, int
  element, ArrayList < Integer > list) {
        // INVARIANT: A.length >= start
        int start = 0;
        Debug debugger = new Debug();
        //The element to add over and over again
        Integer iden = new Integer(element);
        // INITIALIZATION: Before the first iteration of the loop
           , start = 0 and A. length cannot equal anything less
           than 0, so our invariant holds true
        debugger.assertGreatEquals(list.size(), start);
        while (start < end) {
                // MAINTENANCE: At the beginning of each
                   iteration of the loop our invariant holds
                   because for each iteration of the loop one
                   element is added to A and start is incremented
                    by 1
                debugger.assertGreatEquals(list.size(), start);
                list.add(iden);
                //Count up on the iterator
                start++;
        /* TERMINATION: The negation of the gaurd implies that (
           end - start) number of elements hav ebeen added to A,
           since start is initialied as 0 at the beginning of the
            method and is
                        incremented by 1 each iteration of the
                           loop, which means that start amount of
                            elements have been added to A, and so
                            our invariant holds true
        debugger.assertGreatEquals(list.size(), start);
        return list;
}
```