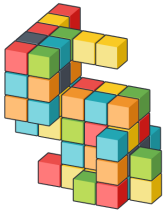

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Least squares for programmers

— with color plates —

Dmitry Sokolov

November 17, 2020

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Coin toss experiment

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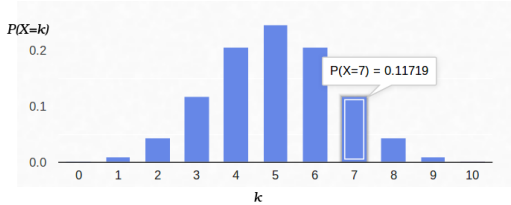

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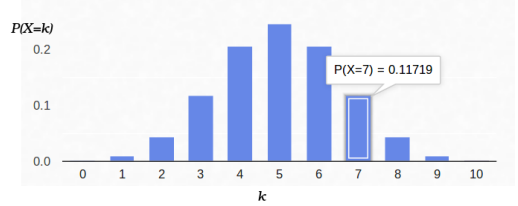

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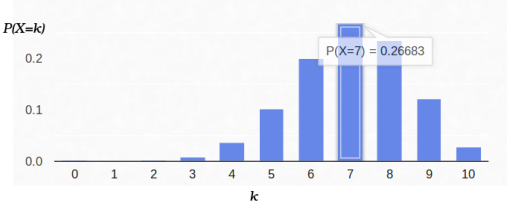

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Coin toss: the likelihood function

Suppose we have a real coin, but we do not know p. However, we can toss it ten times. For example, we have counted seven tails. Would it help us to evaluate p?

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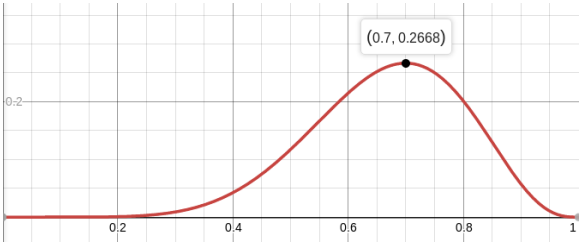

N.B. the function is continuous!

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Just in case, let us check the second derivative:

$$\frac{d^2 \log \mathcal{L}}{dp^2} = -\frac{7}{p^2} - \frac{3}{(1-p)^2}$$

At the point $p = 7/10$ it is negative, therefore this point is indeed a maximum of the function $\mathcal{L}$:

$$\frac{d^2 \log \mathcal{L}}{dp^2}(0.7) \approx -48 < 0$$

Least squares through maximum likelihood

Let us measure a constant value; all measurements are inherently noisy.

For example, if we measure the battery voltage N times, we get N different measurements:

$$\{U_j\}_{j=1}^{N}$$

Suppose that each measurement U_j is i.i.d. and subject to a Gaussian noise, e.g. it is equal to the real value plus the Gaussian noise. The probability density can be expressed as follows:

$$p(U_j) = \frac{1}{\sqrt{2\pi}\sigma} \exp\left(-\frac{(U_j - U)^2}{2\sigma^2}\right),$$

where U is the (unknown) value and σ is the noise amplitude (can be unknown).

Least squares through maximum likelihood

$$\log \mathcal{L}(U, \sigma) = \log \left(\prod_{j=1}^{N} \frac{1}{\sqrt{2\pi}\sigma} \exp \left(-\frac{(U_j - U)^2}{2\sigma^2} \right) \right)$$

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$$= \underbrace{-N \left(\log \sqrt{2\pi} + \log \sigma \right)}_{\text{does not depend on } \{U_j\}_{j=1}^{N}} - \frac{1}{2\sigma^2} \sum_{j=1}^{N} (U_j - U)^2$$

Under Gaussian noise

$$\arg\max_{U} \log \mathcal{L} = \arg\min_{U} \sum_{j=1}^{N} (U_j - U)^2$$

Least squares through maximum likelihood

$$\frac{\partial \log \mathcal{L}}{\partial U} = -\frac{1}{\sigma^2} \sum_{j=1}^{N} (U_j - U) = 0$$

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$$U = \frac{\sum_{j=1}^{N} U_j}{N}$$

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$$\frac{\partial \log \mathcal{L}}{\partial \sigma} = -\frac{N}{\sigma} + \frac{1}{\sigma^3} \sum_{j=1}^{N} (U_j - U)^2 = 0 \qquad \sigma = \sqrt{\frac{\sum_{j=1}^{N} (U_j - U)^2}{N}}$$

Such a convoluted way to obtain a simple average of all measurements...

Linear regression

It is much harder for less trivial examples. Suppose we have N measurements $\{x_j, y_j\}_{j=1}^{N}$, and we want to fit a straight line onto it.

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As before, $\arg\max_{a,b} \log \mathcal{L} = \arg\min_{a,b} S(a, b)$.

Linear regression

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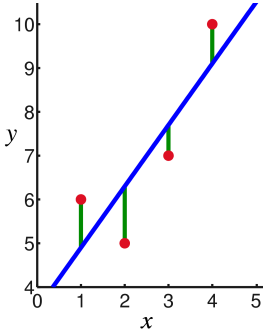

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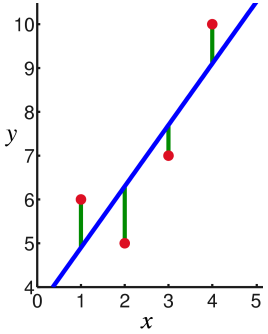
$$a = \frac{N\sum_{j=1}^{N} x_j y_j - \sum_{j=1}^{N} x_j \sum_{j=1}^{N} y_j}{N\sum_{j=1}^{N} x_j^2 - \left(\sum_{j=1}^{N} x_j\right)^2}$$



$$b = \frac{1}{N}\left(\sum_{j=1}^{N} y_j - a\sum_{j=1}^{N} x_j\right)$$

The takeaway message

The least squares method is a particular case of maximizing likelihood in cases where the probability density is Gaussian.

The more we parameters we have, the more cumbersome the analytical solutions are. Fortunately, we are not living in XVIII century anymore, we have computers!

Next we will try to build a geometric intuition on least squares, and see how can least squares problems be efficiently implemented.

Table of Contents

Smooth an array

```
x = [0, .8, 1, .6, .1, .4, .2, .1, .6, .3, 1, .7, .4, 0, .6, 1]

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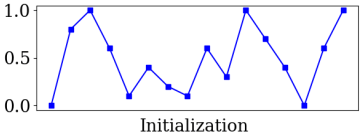

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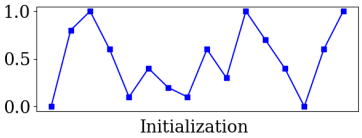

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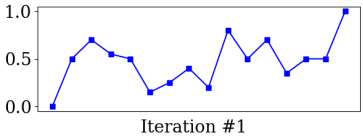

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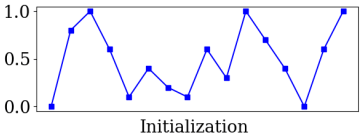

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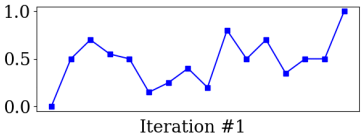

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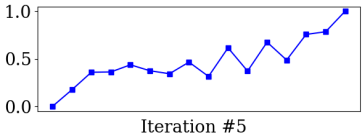

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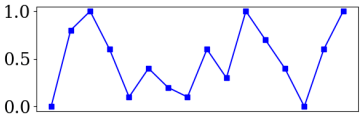

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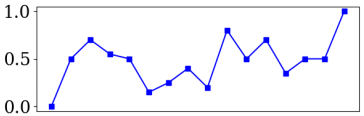

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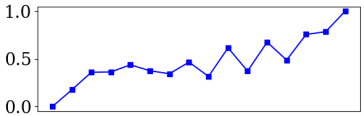
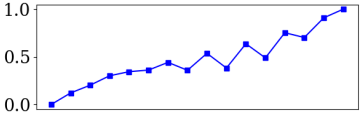


Iteration #5



Iteration #10

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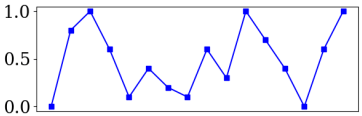

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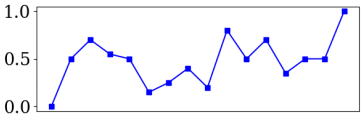

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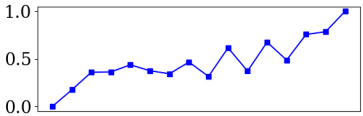
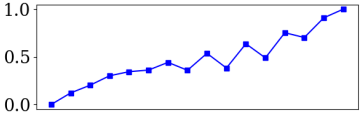
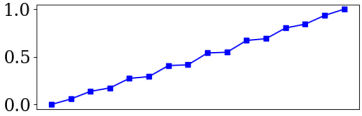


Iteration #5



Iteration #10



Iteration #50

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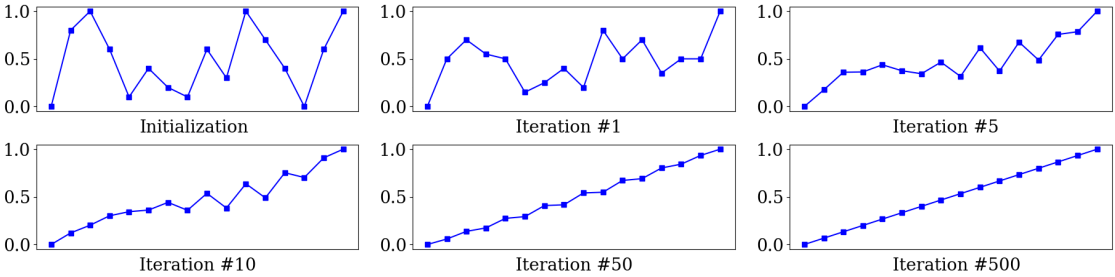

The Jacobi iterative method

Given an ordinary system of linear equations:

$$\begin{cases} a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2 \\ \vdots \\ a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n \end{cases}$$

The Jacobi iterative method

Given an ordinary system of linear equations:

$$\begin{cases} a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2 \\ \vdots \\ a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n \end{cases}$$

Let us rewrite it as follows:

$$\begin{aligned} x_1 &= \frac{1}{a_{11}}(b_1 - a_{12}x_2 - a_{13}x_3 - \cdots - a_{1n}x_n) \\ x_2 &= \frac{1}{a_{22}}(b_2 - a_{21}x_1 - a_{23}x_3 - \cdots - a_{2n}x_n) \\ &\vdots \\ x_n &= \frac{1}{a_{nn}}(b_n - a_{n1}x_1 - a_{n2}x_2 - \cdots - a_{n,n-1}x_{n-1}) \end{aligned}$$

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we can define $\vec{x}^{(1)}$ as follows:

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\begin{aligned}
x_1^{(1)} &= \frac{1}{a_{11}}(b_1 - a_{12}x_2^{(0)} - a_{13}x_3^{(0)} - \cdots - a_{1n}x_n^{(0)}) \\
x_2^{(1)} &= \frac{1}{a_{22}}(b_2 - a_{21}x_1^{(0)} - a_{23}x_3^{(0)} - \cdots - a_{2n}x_n^{(0)}) \\
&\vdots \\
x_n^{(1)} &= \frac{1}{a_{nn}}(b_n - a_{n1}x_1^{(0)} - a_{n2}x_2^{(0)} - \cdots - a_{n,n-1}x_{n-1}^{(0)})
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\end{aligned}
$$

Repeating the process k times, the solution can be approximated by the vector
$\vec{x}^{(k)} = \left(x_1^{(k)}, x_2^{(k)}, \ldots, x_n^{(k)}\right)$.

Back to the array smoothing

```
x = [0, .8, 1, .6, .1, .4, .2, .1, .6, .3, 1, .7, .4, 0, .6, 1]

for _ in range(512):
    x = [ x[0] ] + \
        [ (x[i-1]+x[i+1])/2. for i in range(1, len(x)-1) ] + \
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$$\begin{cases} x_0 = 0 \\ x_1 - x_0 = x_2 - x_1 \\ x_2 - x_1 = x_3 - x_1 \\ \vdots \\ x_{13} - x_{12} = x_{14} - x_{13} \\ x_{14} - x_{13} = x_{15} - x_{14} \\ x_{15} = 1 \end{cases}$$

The Gauß-Seidel iterative method

Jacobi:

$$x_i^{(k)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j=1, j \neq i}^{n} a_{ij} x_j^{(k-1)} \right), \quad \text{for } i = 1, 2, \ldots, n$$

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Smooth an array : Gauß-Seidel

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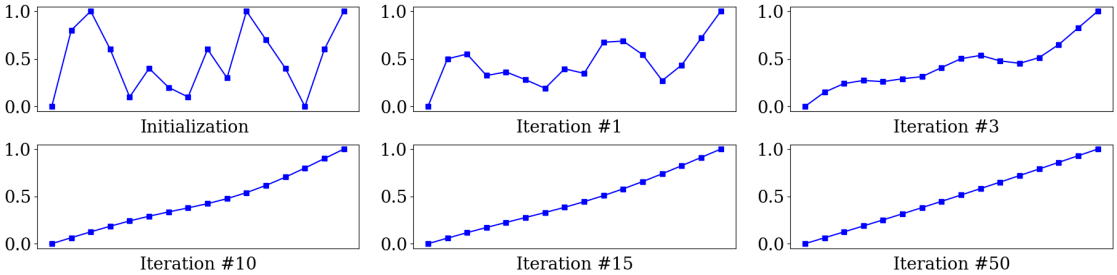

Equality of derivatives vs zero curvature

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\begin{cases}
x_0 & & & & & & & & & & & & = 0 \\
-x_0 & +2x_1 & -x_2 & & & & & & & & & & = 0 \\
 & & -x_2 & +2x_3 & -x_4 & & & & & & & & = 0 \\
 & & & & & \ddots & & & & & & & \vdots \\
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\end{cases}
$$

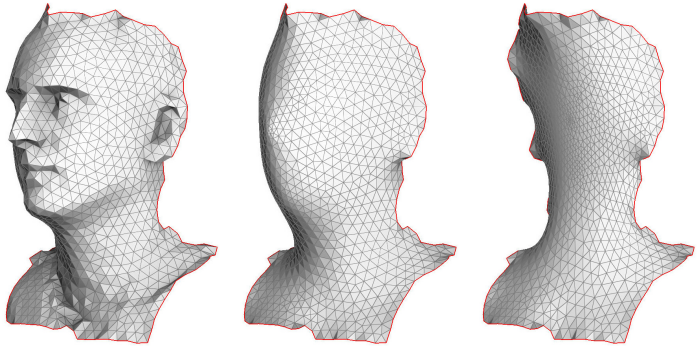
It also works for 3d surfaces

```
#include "model.h"
int main(void) {
    Model m("../input.obj"); // parse the input mesh

    // smooth the surface through Gauss-Seidel iterations
    for (int it=0; it<1000; it++)
        for (int v=0; v<m.nverts(); v++) // for all vertices
            if (!m.is_boundary_vert(v)) // fix the boundary
                m.point(v) = m.one_ring_barycenter(v);

    std::cout << m; // drop the result
    return 0;
}
```

It also works for 3d surfaces



input 10 iterations 1000 iterations

Prescribe the right hand side

```
x = [0, .8, 1, .6, .1, .4, .2, .1, .6, .3, 1, .7, .4, 0, .6, 1]

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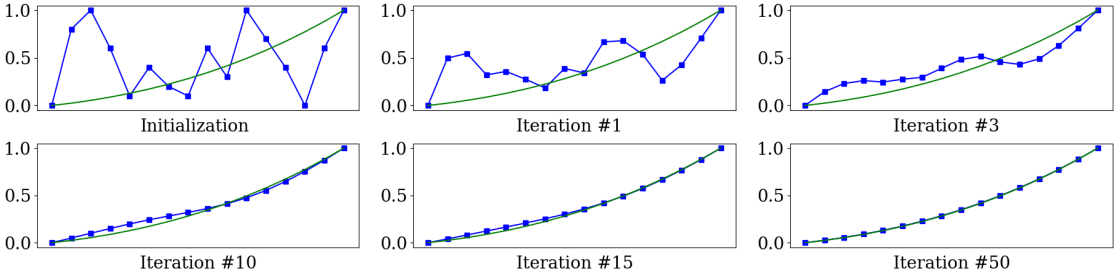

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Matrices and numbers

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Or $f(x) = x^\top Ax = \sum_i \sum_j a_{ij} x_i x_j : \mathbb{R}^2 \to \mathbb{R}$?

```
float f(vector<float> x) {
    return x[0]*a11*x[0] + x[0]*a12*x[1] +
           x[1]*a21*x[0] + x[1]*a22*x[1];
}
```

What is a positive number?

We have a great tool called the predicate “greater than” $>$.

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Definition

The real number a is positive if and only if for all non-zero real $x \in \mathbb{R},\ x \neq 0$ the condition $ax^2 > 0$ is satisfied.

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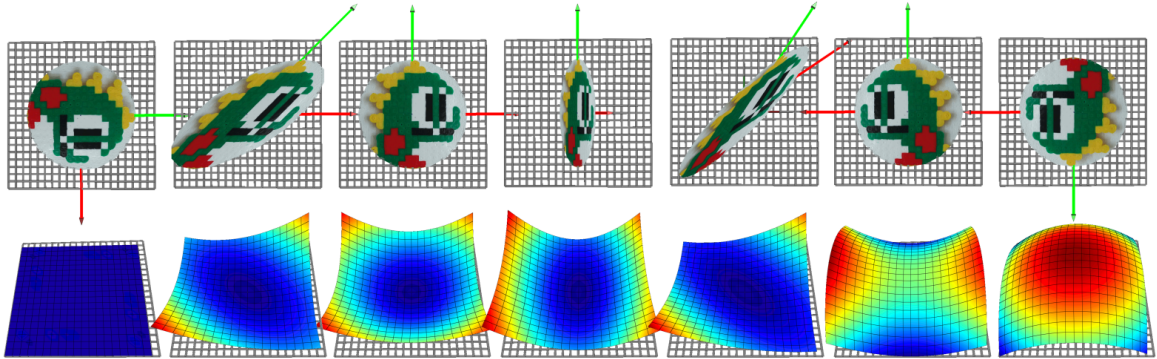
This definition looks pretty awkward, but it applies perfectly to matrices:

Definition

The square matrix A is called positive definite if for any non-zero x the condition $x^\top Ax > 0$ is met, i.e. the corresponding quadratic form is strictly positive everywhere except at the origin.

What is a positive number?



Minimizing a 1d quadratic function

Let us find the minimum of the function $f(x) = ax^2 - 2bx$ (with a positive).

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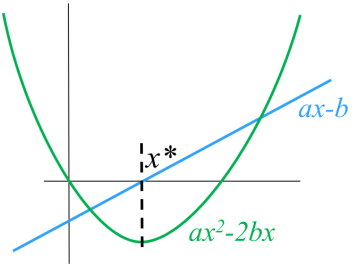

In 1d, the solution x^* of the equation $ax - b = 0$ solves the minimization problem $\arg\min_x (ax^2 - 2bx)$ as well.

Differentiating matrix expressions

The first theorem states that 1×1 matrices are invariant w.r.t the transposition:

Theorem

$x \in \mathbb{R} \Rightarrow x^\top = x$

The proof is left as an exercise.

Differentiating matrix expressions

For a 1d function bx we know that $\frac{d}{dx}(bx) = b$, but what happens in the case of a real function of n variables?

Theorem

$\nabla b^\top x = \nabla x^\top b = b$

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Theorem

$\nabla b^\top x = \nabla x^\top b = b$

$$\nabla(b^\top x) = \begin{bmatrix} \frac{\partial(b^\top x)}{\partial x_1} \\ \vdots \\ \frac{\partial(b^\top x)}{\partial x_n} \end{bmatrix} = \begin{bmatrix} \frac{\partial(b_1 x_1 + \cdots + b_n x_n)}{\partial x_1} \\ \vdots \\ \frac{\partial(b_1 x_1 + \cdots + b_n x_n)}{\partial x_n} \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix} = b$$

Differentiating matrix expressions

For a 1d function ax^2 we know that $\frac{d}{dx}(ax^2) = 2ax$, but what about quadratic forms?

Theorem

$\nabla(x^\top Ax) = (A + A^\top)x$

Note that if A is symmetric, then $\nabla(x^\top Ax) = 2Ax$.

The proof is straightforward, let us express the quadratic form as a double sum:

$$x^\top Ax = \sum_i \sum_j a_{ij} x_i x_j$$

Differentiating matrix expressions

$$\frac{\partial (x^\top A x)}{\partial x_i} = \frac{\partial}{\partial x_i} \left(\sum_{k_1} \sum_{k_2} a_{k_1 k_2} x_{k_1} x_{k_2} \right) =$$

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$$= \sum_j (a_{ij} + a_{ji}) x_j \qquad \Rightarrow \nabla (x^\top A x) = (A + A^\top) x$$

Minimum of a quadratic form and the linear system

Recall that for $a > 0$ solving the equation $ax = b$ is equivalent to the quadratic function $\arg\min_x(ax^2 - 2bx)$ minimization.

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To minimize a quadratic form $\arg\min_{x \in \mathbb{R}^n}(x^\top Ax - 2b^\top x)$ with a symmetric positive definite matrix A, equate the derivative to zero:

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Recall that A is symmetric: $Ax = b$.

Back to the linear regression

Given two points (x_1, y_1) and (x_2, y_2), find the line that passes through: $y = \alpha x + \beta$.

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$$\begin{cases} \alpha x_1 + \beta = y_1 \\ \alpha x_2 + \beta = y_2 \end{cases} \qquad \underbrace{\begin{bmatrix} x_1 & 1 \\ x_2 & 1 \end{bmatrix}}_{:=A} \underbrace{\begin{bmatrix} \alpha \\ \beta \end{bmatrix}}_{:=x} = \underbrace{\begin{bmatrix} y_1 \\ y_2 \end{bmatrix}}_{:=b} \qquad \Rightarrow x^* = A^{-1} b$$

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Now add a **third** point (x_3, y_3):

$$\begin{cases} \alpha x_1 + \beta = y_1 \\ \alpha x_2 + \beta = y_2 \\ \alpha x_3 + \beta = y_3 \end{cases}$$

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Given two points (x_1, y_1) and (x_2, y_2), find the line that passes through: $y = \alpha x + \beta$.

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$$\underbrace{\begin{bmatrix} x_1 & 1 \\ x_2 & 1 \\ x_3 & 1 \end{bmatrix}}_{:=A\,(3\times 2)} \underbrace{\begin{bmatrix} \alpha \\ \beta \end{bmatrix}}_{:=x\,(2\times 1)} = \underbrace{\begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}}_{:=b\,(3\times 1)}$$

A is rectangular, and thus it is not invertible. Oops!

Back to the linear regression

No biggie, let us rewrite the system:

$$\alpha \underbrace{\begin{bmatrix} x_1 & x_2 & x_3 \end{bmatrix}^\top}_{:=\vec{i}} + \beta \underbrace{\begin{bmatrix} 1 & 1 & 1 \end{bmatrix}^\top}_{:=\vec{j}} = \begin{bmatrix} y_1 & y_2 & y_3 \end{bmatrix}^\top$$

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$$\alpha\vec{i} + \beta\vec{j} = \vec{b}$$

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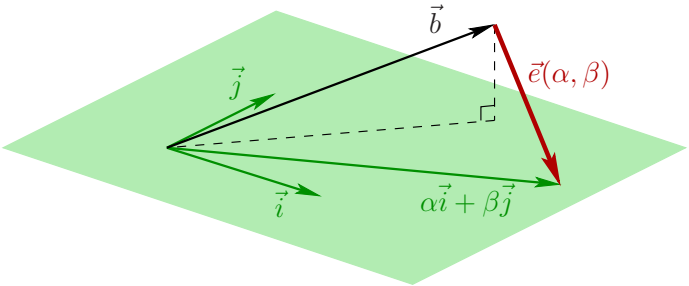
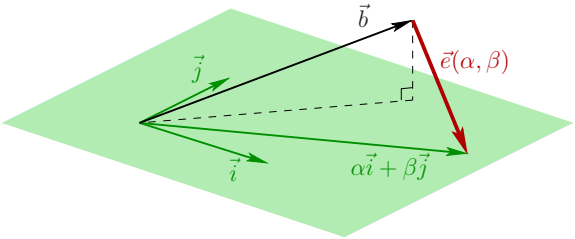
Solve for $\arg\min_{\alpha,\beta} \|\vec{e}(\alpha,\beta)\|$, where $\vec{e}(\alpha,\beta) := \alpha\vec{i} + \beta\vec{j} - \vec{b}$:



Back to the linear regression



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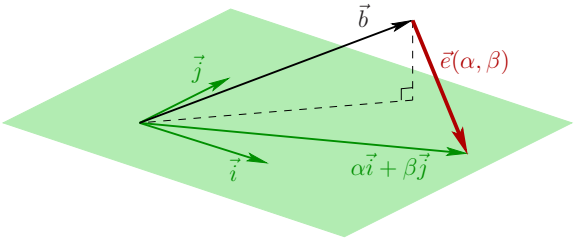

The $\|\vec{e}(\alpha, \beta)\|$ is minimized when $\vec{e}(\alpha, \beta) \perp \text{span}\{\vec{i}, \vec{j}\}$:

$$\begin{cases} \vec{i}^{\top} \vec{e}(\alpha, \beta) = 0 \\ \vec{j}^{\top} \vec{e}(\alpha, \beta) = 0 \end{cases}$$

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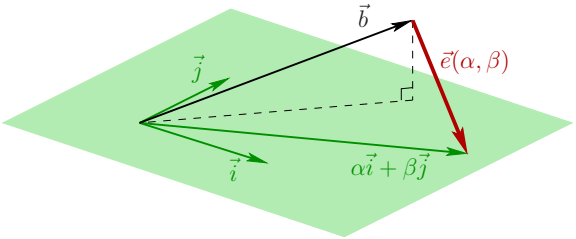

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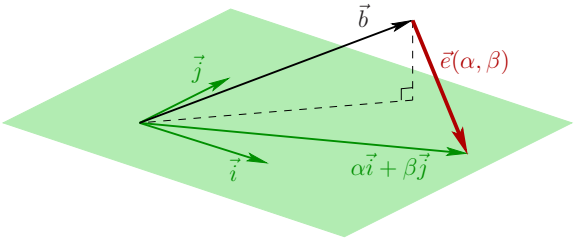
$$\begin{bmatrix} x_1 & x_2 & x_3 \\ 1 & 1 & 1 \end{bmatrix} \left(\alpha \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} + \beta \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} - \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} \right) = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

Back to the linear regression



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$$A^{\top}(Ax - b) = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

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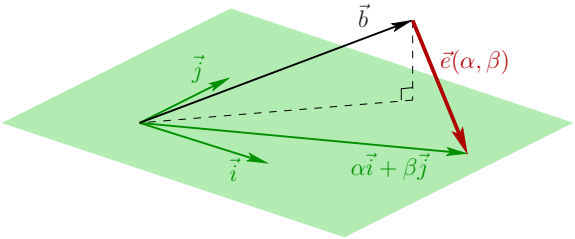

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$$A^{\top}(Ax - b) = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

In a general case the matrix $A^{\top}A$ can be invertible!

$$A^{\top}Ax = A^{\top}b.$$

Some nice properties of $A^\top A$

Theorem

$A^\top A$ is symmetric.

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It follows from the fact that $x^\top A^\top A x = (Ax)^\top Ax > 0$. Moreover, $A^\top A$ is positive definite in the case where A has linearly independent columns (rank A is equal to the number of the variables in the system).

Least squares in more than two dimensions

The same reasoning applies, here is an algebraic way to show it:

$$\arg\min \|Ax - b\|^2$$

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The takeaway message

The least squares problem $\arg\min \|Ax-b\|^2$ is equivalent to minimizing the quadratic function $\arg\min \left(x^\top A' x - 2b'^\top x\right)$ with (in general) a symmetric positive definite matrix A'. This can be done by solving a linear system $A'x = b'$.

Table of Contents

Linear-quadratic regulator

Imagine a car going at $v_0 = 0.5$ m/s. The goal is to accelerate to $v_n = 2.3$ m/s in $n = 30$ s maximum. We can control the acceleration u_i via the gas pedal:

$$v_{i+1} = v_i + u_i$$

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So, we need to find $\{u_i\}_{i=0}^{n-1}$ that optimizes some quality criterion $J(\vec{v}, \vec{u})$:

$$\min J(\vec{v}, \vec{u}) \quad s.t. \quad v_{i+1} = v_i + u_i = v_0 + \sum_{j=0}^{i-1} u_j \quad \forall i \in 0..n-1$$

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What happens if we ask for the car to reach the final speed as quickly as possible?
It can be written as follows:

$$J(\vec{v}, \vec{u}) := \sum_{i=1}^{n} (v_i - v_n)^2 = \sum_{i=1}^{n} \left(\sum_{j=0}^{i-1} u_j - v_n + v_0 \right)^2$$

Linear-quadratic regulator

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Solve in the least squares sense:

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```
import numpy as np
n,v0,vn = 30,0.5,2.3
A = np.matrix(np.tril(np.ones((n,n)))) # lower triangular matrix
b = np.matrix(np.full((n, 1), vn-v0))
u = np.linalg.inv(A.T*A)*A.T*b
v = [v0 + np.sum(u[:i]) for i in range(0,n+1)]
```

Linear-quadratic regulator

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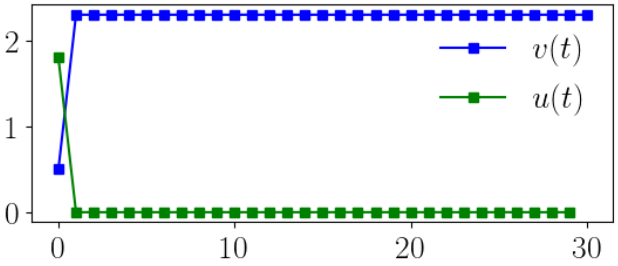

Ouch... Quite brutal accelerations!

Linear-quadratic regulator

Ok, no problem, let us penalize large accelerations:

$$J(\vec{v}, \vec{u}) := \sum_{i=0}^{n-1} {u_i}^2$$

Solve in the least squares sense:

$$\begin{cases} u_0 &= 0 \\ u_1 &= 0 \\ &\vdots \\ u_{n-1} &= 0 \end{cases}$$

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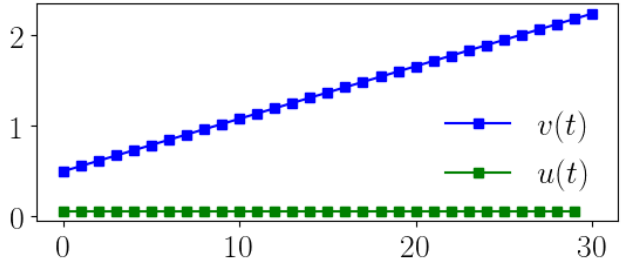

Low acceleration, however the transient time becomes unacceptable.

Linear-quadratic regulator

Optimize for competing goals:

$$J(\vec{v}, \vec{u}) := \sum_{i=1}^{n} (v_i - v_n)^2 + 4 \sum_{i=0}^{n-1} {u_i}^2$$

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$$\sum_{i=1}^{n} \left(\sum_{j=0}^{i-1} u_j - v_n + v_0 \right)^2 + 4 \sum_{i=0}^{n-1} {u_i}^2$$

N.B. Note the tradeoff coefficients !

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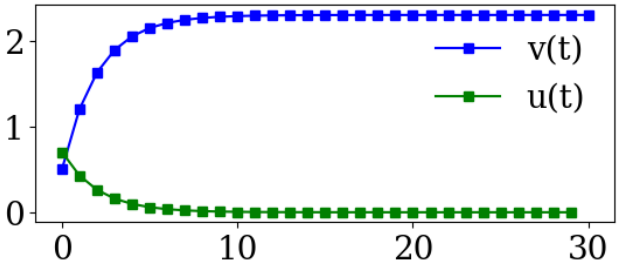
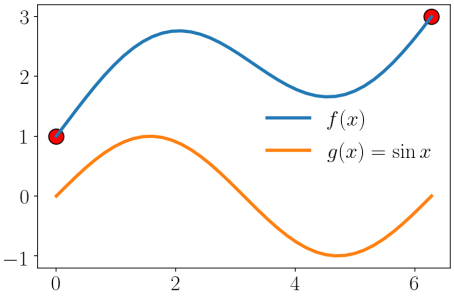


Poisson's equation

Problem: find $f(x)$ defined on $x \in [0, 2\pi]$ as close as possible to $g(x) := \sin x$, constrained to $f(0) = 1$ and $f(2\pi) = 3$.

Formulate it as the Poisson's equation with Dirichlet boundary conditions:
$\frac{d^2}{dx^2} f = \frac{d^2}{dx^2} g$ s.t. $f(0) = 1,\ f(2\pi) = 3$



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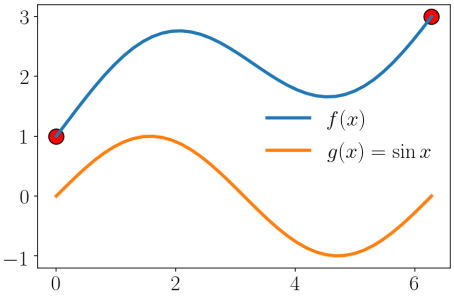

Least squares formulation:

$$\min_f \int_0^{2\pi} \|f' - g'\|^2$$

with $f(0) = 1,\ f(2\pi) = 3$

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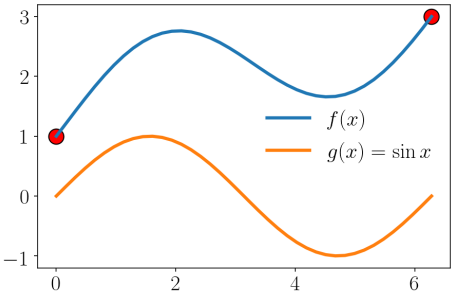

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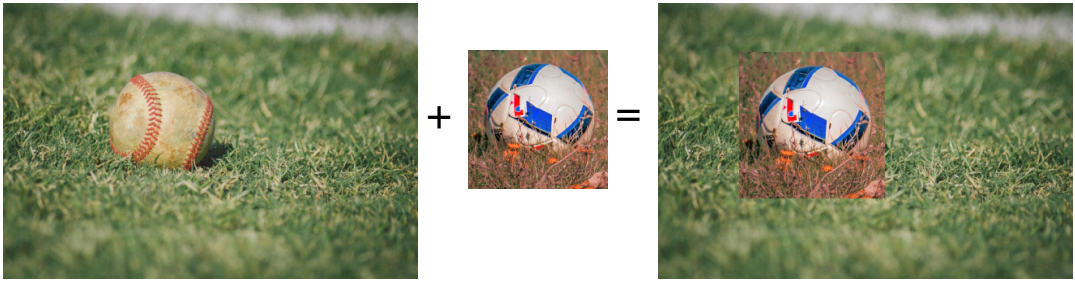
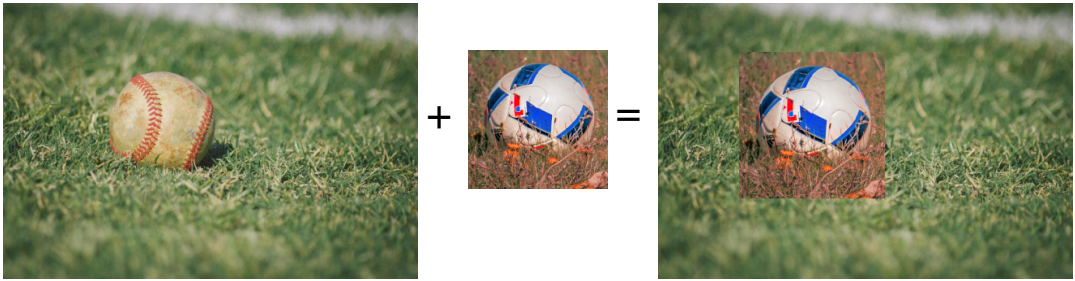
Discretization:

$$\begin{cases} f_1 & & & = g_1 - g_0 + f_0 \\ -f_1 & +f_2 & & = g_2 - g_1 \\ & & \vdots & \\ & -f_{n-3} & +f_{n-2} & = g_{n-2} - g_{n-3} \\ & & -f_{n-2} & = g_{n-1} - g_{n-2} - f_{n-1} \end{cases}$$

Poisson image editing

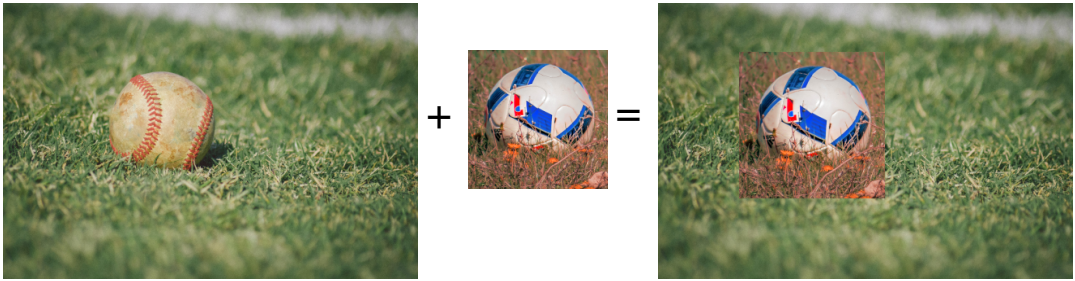


Poisson image editing



We can do better: solve a linear system par color channel.

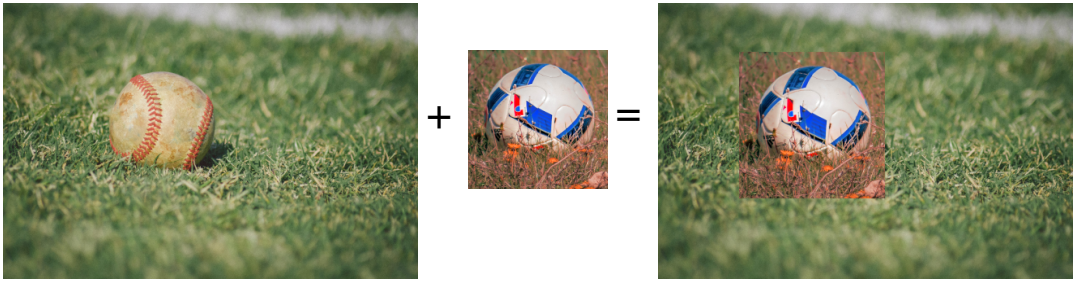
Poisson image editing



We can do better: solve a linear system par color channel.

Let a be:

Poisson image editing

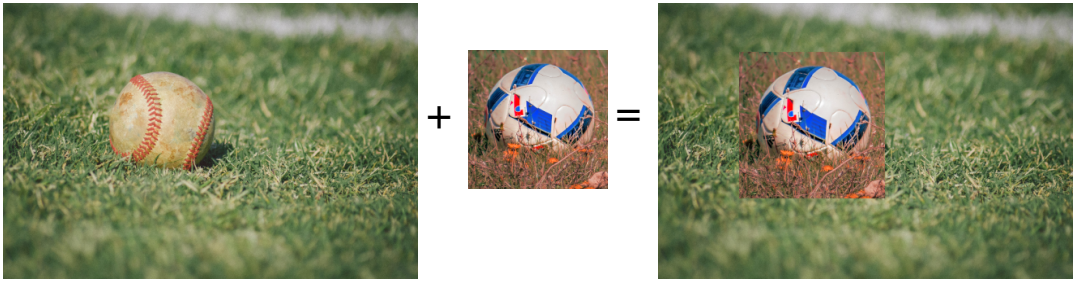


We can do better: solve a linear system par color channel.

Let *a* be:

Let *b* be:

Poisson image editing



We can do better: solve a linear system par color channel.

Let a be:

Let b be:

Solve for f who takes its boundary conditions from a and the gradients from b:

$$\min_f \int_\Omega \|\nabla f - \nabla b\|^2 \quad \text{with } f|_{\partial\Omega} = a|_{\partial\Omega}$$

Poisson image editing

Discretize the problem: having $w \times h$ pixels grayscale images a and b, we compute a $w \times h$ pixels image f, solve in the least squares sense:

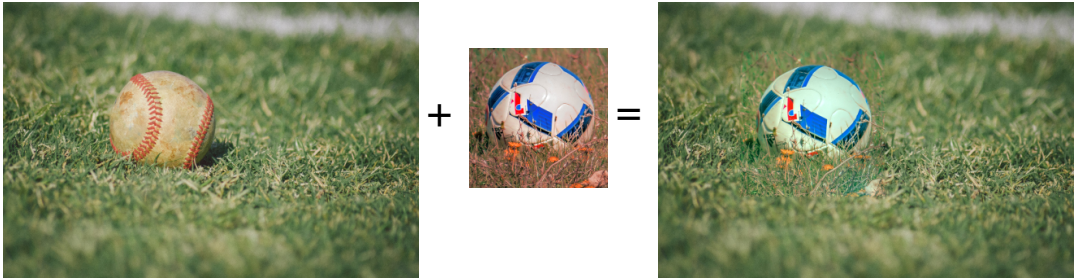
$$\begin{cases} f_{i+1,j} - f_{i,j} &= b_{i+1,j} - b_{i,j} \quad \forall (i,j) \in [0 \dots w-2] \times [0 \dots h-2] \\ f_{i,j+1} - f_{i,j} &= b_{i,j+1} - b_{i,j} \quad \forall (i,j) \in [0 \dots w-2] \times [0 \dots h-2] \\ f_{i,j} &= a_{i,j} \quad \forall (i,j) \text{ s.t. } i = 0 \text{ or } i = w-1 \text{ or } j = 0 \text{ or } j = h-1 \end{cases}$$

N.B: sparse system solver!

Poisson image editing

Discretize the problem: having $w \times h$ pixels grayscale images a and b, we compute a $w \times h$ pixels image f, solve in the least squares sense:

$$\begin{cases} f_{i+1,j} - f_{i,j} = b_{i+1,j} - b_{i,j} & \forall (i,j) \in [0 \dots w-2] \times [0 \dots h-2] \\ f_{i,j+1} - f_{i,j} = b_{i,j+1} - b_{i,j} & \forall (i,j) \in [0 \dots w-2] \times [0 \dots h-2] \\ f_{i,j} = a_{i,j} & \forall (i,j) \text{ s.t. } i = 0 \text{ or } i = w-1 \text{ or } j = 0 \text{ or } j = h-1 \end{cases}$$



Caricature

```
x = [100,100,97,93 ...  23,21,19] # 2d closed polyline
y = [0,25,27,28,30 ...  11,6,4,1]
n = len(x)                        # number of points
cx = [x[i] - (x[(i-1+n)%n]+x[(i+1)%n])/2 for i in range(n)] #precompute the
cy = [y[i] - (y[(i-1+n)%n]+y[(i+1)%n])/2 for i in range(n)] #discrete curvature
for _ in range(1000): # Gauss-Seidel iterations
    for i in range(n):
        x[i] = (x[(i-1+n)%n]+x[(i+1)%n])/2 + cx[i]*1.9
        y[i] = (y[(i-1+n)%n]+y[(i+1)%n])/2 + cy[i]*1.9
```

Caricature

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```

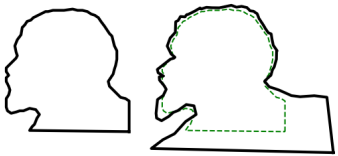

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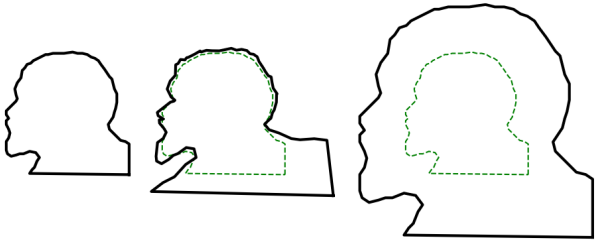
```



Almost, but no :(

Caricature

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```
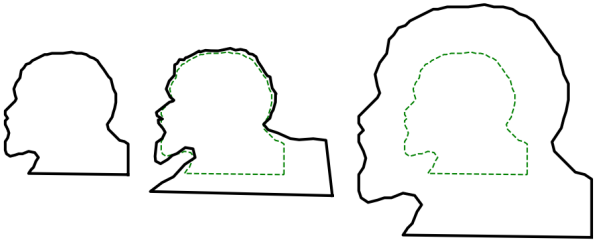



Almost, but no :(

Least squares equivalent:

$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c \cdot (x_j - x_i) \right)^2$$

x_i are the input coordinates and x'_i are the unknowns (separable in x and y)

Caricature

A quick fix:

$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c_0 \cdot (x_j - x_i) \right)^2$$

Caricature

A quick fix:

$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c_0 \cdot (x_j - x_i) \right)^2 + \sum_{\forall \text{ vertex } i} c_1^2 \left(x'_i - x_i \right)^2$$

Caricature

A quick fix:

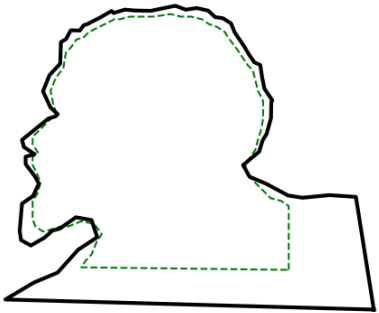
$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c_0 \cdot (x_j - x_i) \right)^2 + \sum_{\forall \text{ vertex } i} c_1^2 \left(x'_i - x_i \right)^2$$

$$\begin{cases}
-x'_0 + x'_1 = c_0 \cdot (x_1 - x_0) \\
-x'_1 + x'_2 = c_0 \cdot (x_2 - x_1) \\
\vdots \\
x'_{n-2} + x'_{n-1} = c_0 \cdot (x_{n-2} - x_{n-1}) \\
x'_0 - x'_{n-1} = c_0 \cdot (x_{n-1} - x_0) \\
c_1 \cdot x'_0 = c_1 \cdot x_0 \\
c_1 \cdot x'_1 = c_1 \cdot x_1 \\
\vdots \\
c_1 \cdot x'_{n-1} = c_1 \cdot x_{n-1}
\end{cases}$$

Caricature

A quick fix:

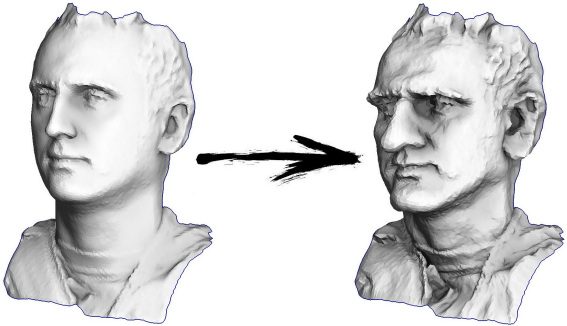
$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c_0 \cdot (x_j - x_i) \right)^2 + \sum_{\forall \text{ vertex } i} c_1^2 \left(x'_i - x_i \right)^2$$



Caricature

A quick fix:

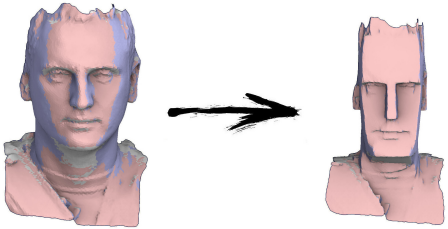
$$\min \sum_{\forall \text{ edge } (i,j)} \left(x'_j - x'_i - c_0 \cdot (x_j - x_i) \right)^2 + \sum_{\forall \text{ vertex } i} c_1^2 \left(x'_i - x_i \right)^2$$



And it works out of the box for 3d surfaces as well!

Cubify it!

$$\vec{a}_{ijk} := \underset{\vec{a} \in \{(1,0,0),(0,1,0),(0,0,1)\}}{\arg\max} \left| \vec{a} \cdot \vec{N}_{ijk} \right|$$



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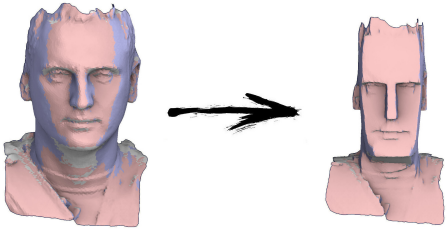


Let $\vec{e}_{ij}$ be the input geometry,
and $\vec{e'}_{ij}$ the unknowns.

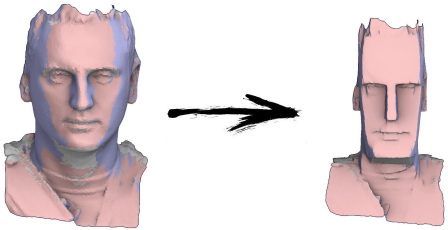
Quick test

What would be the result?

$$\min \sum_{\forall \text{ edge } ij} \left\| \vec{e'}_{ij} - \vec{e}_{ij} \right\|^2$$

Cubify it!

$$\vec{a}_{ijk} := \underset{\vec{a} \in \{(1,0,0),(0,1,0),(0,0,1)\}}{\arg\max} \left| \vec{a} \cdot \vec{N}_{ijk} \right|$$



$$\mathrm{proj}_{\vec{a}}\vec{v} := \vec{v} - \frac{\vec{v}\cdot\vec{v}}{\vec{a}\cdot\vec{a}}\vec{a}$$

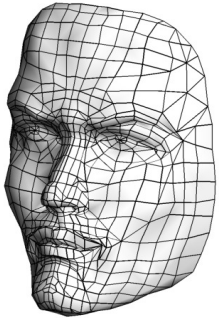
Let $\vec{e}_{ij}$ be the input geometry, and $\vec{e'}_{ij}$ the unknowns.

$$\min \quad c_0 \cdot \sum_{\forall \text{ edge } ij} \left\| \vec{e'}_{ij} - \vec{e}_{ij} \right\|^2 +$$

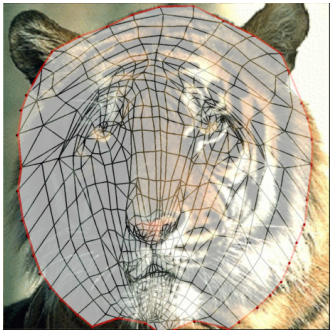
$$c_1 \cdot \sum_{\forall \text{ triangle } ijk} \left(\left\| \vec{e'}_{ij} - \mathrm{proj}_{\vec{a}_{ijk}} \vec{e}_{ij} \right\|^2 + \left\| \vec{e'}_{jk} - \mathrm{proj}_{\vec{a}_{ijk}} \vec{e}_{jk} \right\|^2 + \left\| \vec{e'}_{ki} - \mathrm{proj}_{\vec{a}_{ijk}} \vec{e}_{ki} \right\|^2 \right)$$

N.B: still a separable problem

Mix the coordinates: least squares conformal maps

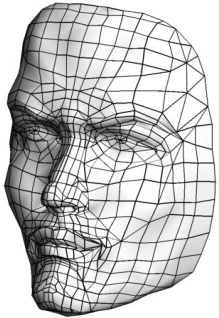


$$U(x, y, z) = (u, v)$$



$$U^{-1}(u, v) = (x, y, z)$$

Mix the coordinates: least squares conformal maps



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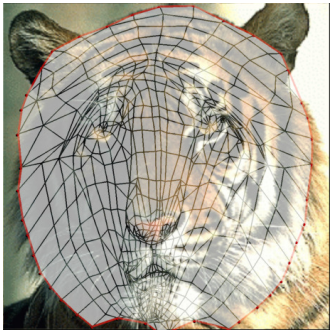
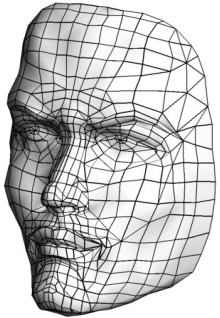


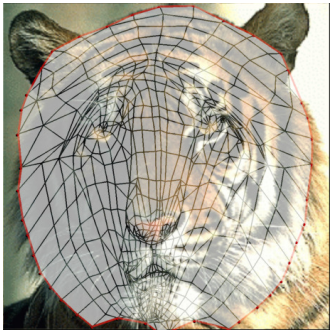
$$U^{-1}(u, v) = (x, y, z)$$

How to create such a map?

Mix the coordinates: least squares conformal maps



$$U(x, y, z) = (u, v)$$

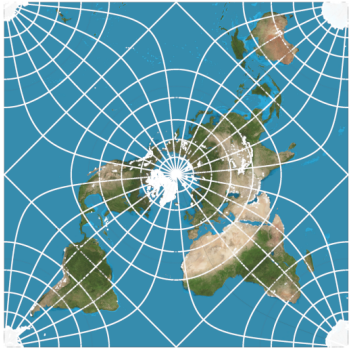


$$U^{-1}(u, v) = (x, y, z)$$

How to create such a map?

Let us try to compute a conformal map!

Mix the coordinates: least squares conformal maps



Dmitry Sokolov

Least squares for programmers

— with color plates —