Initial Exploration

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Objectives

- 1. model distributed semantics
- 2. verify a given semantic satisfies its specifications
- 3. check pair-wise compatibility
- 4. composition of two or more systems

Definitions: Operation

Operation are tuples: (proc, type, obj, ival, oval, stime, rtime) [1, pp.3]. class Operation(NamedTuple): proc: int # process id type: str # operation type obj: int # object id ival: Any # input value stime: int # start time rtime: int = None # return time oval: Any = None # output value

symbol: str = None # readable representation

Definitions: History

A set of operations, contains all operations invoked in a given execution, describes the **observable outcomes of executions** [1, pp.4].

It's possible that a set of symbols and their relations:

- does not form a history (invalid)
- forms multiple histories (ambiguous)

Types:

- $H|_{\mathrm{rd}} = \{ \mathrm{op} \in H : \mathrm{op.type} = \mathrm{rd} \}$
- $H|_{wr} = \{ op \in H : op.type = wr \}$

Definitions: History

Relations:

```
• returns-before: rb \triangleq \{(a,b) : a,b \in H \land a.rtime < b.stime\}
• same-session: ss \triangleq \{(a,b) : a,b \in H \land a.proc = b.proc\}
• session-order: so \triangleq rb \cap ss
class History:
     def init (self: Self,
       ops: set[Operation],
       **kwargs: set[Relation]
     ) -> None: ...
```

Definitions: Anstract Execution

An abstract execution is

- built on top of a history
- · captures the non-determinism, and constraints
- an event graph A = (H, vis, ar, hb) [2, pp.25-27,34-35]
- can be visualized as a directed graph

Definitions: Anstract Execution

```
Relations:
• vis (visibility): a \xrightarrow{\text{vis}} b
class AbstractExecution:
     def init (self: Self,
       hist: History,
       **kwargs: set[Relation]
     ) -> None: ...
```

Example: Monotonic Reads

```
1: op a: set("somekey", "someval")
                op b: get("somekey") -> "someval"
2:
                                op c: get("somekey") -> "someval"
• a \stackrel{\text{vis}}{\longrightarrow} b
• b \xrightarrow{\text{so}} c
  • b \xrightarrow{\mathrm{rb}} c
   • b \stackrel{\text{ss}}{\longrightarrow} c
```

References

- [1] P. Viotti and M. Vukolić, "Consistency in Non-Transactional Distributed Storage Systems". 2016.
- [2] S. Burckhardt, *Principles of Eventual Consistency*, Principles of Eventual Consistency., vol. 1. in Foundations and Trends® in Programming Languages, vol. 1. Now Publishers, 2014, pp. 1–150. [Online]. Available: https://www.microsoft.com/en-us/research/publication/principles-of-eventual-consistency/