

# Scalable Tensor Algorithms for Modern Computing Systems

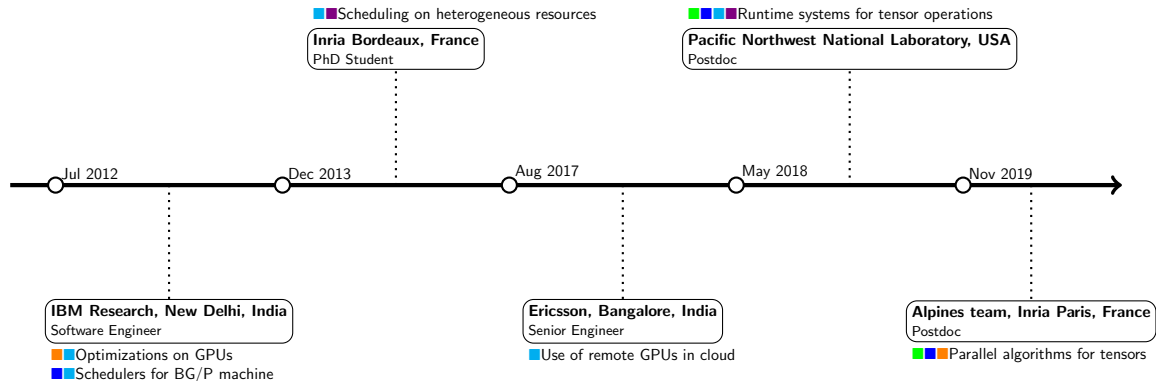
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June 4, 2021

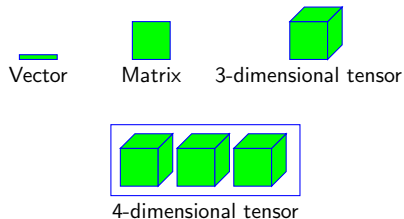
# Resume

Interests: **Tensors**, **Scalable Algorithms**, **Scheduling**, **Runtime Systems**, **Performance Optimizations**



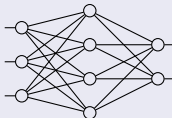
# Tensors are used in Several Domains

- **Neuroscience:** Neuron  $\times$  Time  $\times$  Trial
- **Transportation:** Pickup  $\times$  Dropoff  $\times$  Time
- **Media:** User  $\times$  Movie  $\times$  Time
- **Ecommerce:** User  $\times$  Product  $\times$  Time
- **Cyber-Traffic:** IP  $\times$  IP  $\times$  Port  $\times$  Time
- **Social-Network:** Person  $\times$  Person  $\times$  Time  $\times$  Interaction-Type



## High Dimensional Tensors

- **Neural Network:**



- **Molecular Simulation:** To represent wave functions
- **Quantum Computing:** To represent qubit states

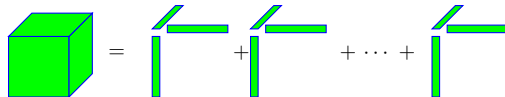
# Tensor Computations

- Memory and computation requirements are exponential in the number of dimensions
  - A simulation involving just 100 spatial orbitals manipulates a huge tensor with  $4^{100}$  elements
- People work with low dimensional structure (decomposition) of the tensors
  - A tensor is represented with smaller objects
  - Improves memory and computation requirements
- Most tensor decompositions rely on Singular Value Decomposition (SVD) of matrices
  - SVD represents a matrix as the sum of rank one matrices,  $A = U\Sigma V^T = \sum_i \Sigma(i; i) U_i V_i^T$

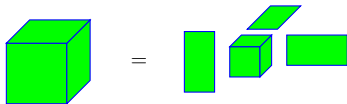
$$\text{Matrix} = \text{Matrix} \times \text{Matrix} = \text{Matrix} + \text{Matrix} + \dots + \text{Matrix}$$

# Popular Tensor Decompositions (Higher Order Generalization of SVD)

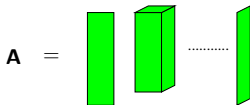
- Canonical decomposition (Also known as Canonical Polyadic or CANDECOMP/PARAFAC)


$$\text{Cube} = \text{rank-1 tensor} + \text{rank-1 tensor} + \dots + \text{rank-1 tensor}$$

- Tucker decomposition


$$\text{Cube} = \text{core cube} \times \text{rectangular slices}$$

- Tensor Train decomposition (equivalently known as Matrix Product States)


$$\mathbf{A} = \text{rectangular slice} \times \text{core box} \times \text{rectangular slice}$$

Tensor notation: bold letters denote tensors, i.e.,  $\mathbf{A} \in \mathbb{R}^{n_1 \times \dots \times n_d}$  is a  $d$ -dimensional tensor.

## 1 Parallel Tensor Train Algorithms

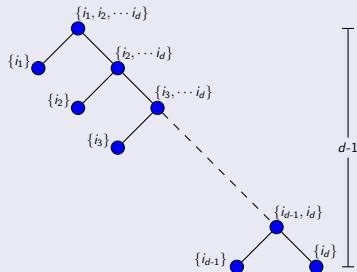
## 2 Scheduling on Heterogeneous Systems

- Scheduling of Dense Linear Algebra Kernels
- Communication Computation Overlap

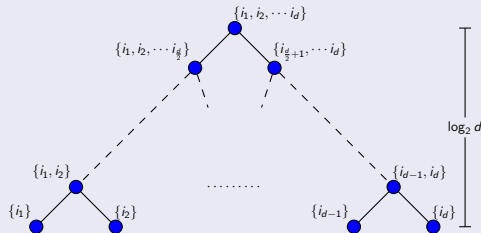
# Tensor Train algorithms and Separation of dimensions

- A sequential algorithm to compute Tensor Train decomposition exists [Oseledets, 2011]

## Sequential algorithm



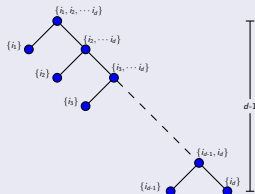
## For better parallelization



- Can obtain better parallelism by expressing the operation in a balanced binary tree shape
  - Proposed a parallel algorithm based on this idea (joint work with L. Grigori, Inria Paris)

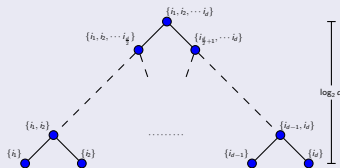
# Tensor Train approximation algorithms

## Sequential algorithm [Oseledets, 2011]



- Unfolding matrix: matricized representation of the tensor
- Perform truncated SVD of unfolding matrix  $A$ ,  $A = U\Sigma V^T + E_A$
- Work with  $\Sigma V^T$  on the right subtree

## Our parallel algorithm



- Perform truncated SVD of unfolding matrix  $A$ ,  $A = U\Sigma V^T + E_A$
- Find diagonal matrices  $X$ ,  $Y$ , and  $S$ , such that  $\Sigma = XSY$
- Call left (resp. right) subtree with  $UX$  (resp.  $YV^T$ )

Approach 1:  $X = I$ ,  $Y = \Sigma$ ,  $S = I$

Approach 2:  $X = Y = \Sigma^{\frac{1}{2}}$ ,  $S = I$

Approach 3:  $X = Y = \Sigma$ ,  $S = \Sigma^{-1}$



# Comparison of our approaches

- A 12-dimensional tensor with  $4^{12}$  elements (generated with a popular low rank function)
- prescribed accuracy =  $10^{-6}$
- Compr: compression ratio, NE: number of elements, AA: approximation accuracy

Metric	Sequential Algo	Parallel Algo		
		Approach 1	Approach 2	Approach 3
Compr	99.993	99.817	99.799	99.993
NE	1212	30632	33772	1212
AA	2.271e-07	3.629e-08	2.820e-08	2.265e-07

## SVD is expensive

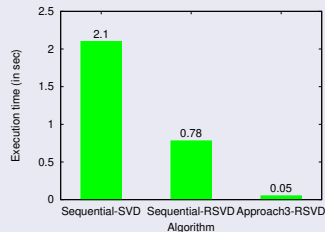
- Good alternatives to SVD: QR factorization with column pivoting (QRCP), randomized SVD (RSVD)

Approach	Rank	Compr	NE	Sequential-AA	Approach3-AA
SVD	5	99.994	992	6.079e-06	6.079e-06
QRCP+SVD				1.016e-05	1.384e-05
RSVD				6.079e-06	6.079e-06
SVD	6	99.992	1376	1.323e-07	1.340e-07
QRCP+SVD				3.555e-07	5.737e-07
RSVD				1.322e-07	1.322e-07

# Performance Comparison

## Single core performance

- Number of computations for both RSVD algorithms =  $\mathcal{O}(n^d)$
- Approach3-SVD is very slow
- Approach3-RSVD is much faster



## Parallel performance counts along the critical path on $P$ processors

Algorithm	# Computations	Communications	# Messages
Sequential-RSVD	$\mathcal{O}(\frac{n^d}{P})$	$\mathcal{O}(\frac{n^{d-1}}{\sqrt{P}} \log P)$	$\mathcal{O}(d \log P)$
Parallel-RSVD	$\mathcal{O}(\frac{n^d}{P})$	$\mathcal{O}(\frac{n^{\frac{d}{2}}}{\sqrt{P}} \log P)$	$\mathcal{O}(\log d \log P)$

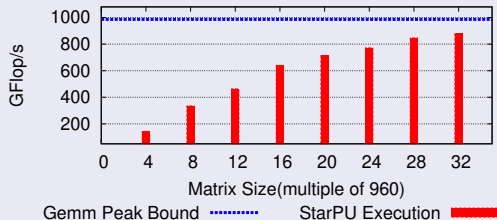
# Scheduling on Heterogeneous Systems

- Heterogeneous systems are common in High Performance Computing (HPC) (147 out of 500 in TOP500 list)
- Task based runtimes are a popular approach to exploit these systems



- Task based runtimes: StarPU, OmpSS, Legion, PaRSEC
- Application is represented as a graph of tasks (computations)
- E.g., Cholesky graph for  $4 \times 4$  tile matrix

## StarPU scheduler performance

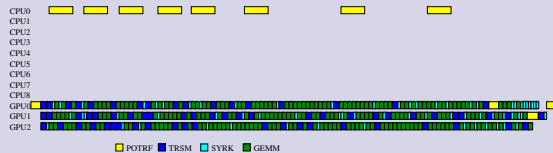


- A platform with 9 CPUs and 3 GPUs
- Scheduler is based on popular heft strategy
- Goal: Enhance performance bounds and propose better scheduling strategies

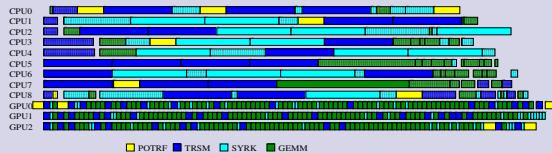
Joint work with E. Agullo, O. Beaumont, L. Eyraud-Dubois, and S. Thibault during my PhD at Inria Bordeaux

# Our strategy and Performance comparison

## Trace for 12 X 12 tile matrix of Cholesky factorization



StarPU scheduling strategy, performance = 686 GFlop/s

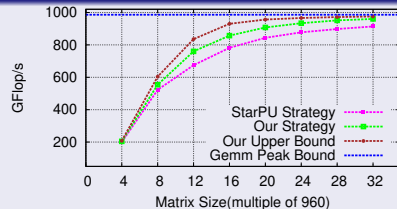


Our strategy, performance = 760 GFlop/s

## Theoretical guarantees of our strategy and performance comparison

- Each resource selects the best suited task
- A fast resource restarts the blocking task

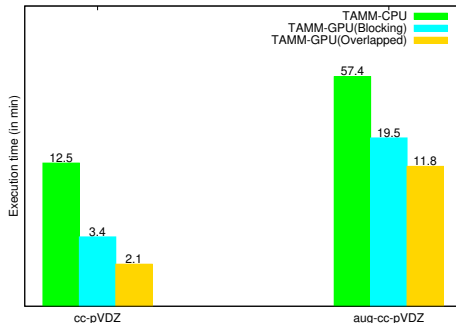
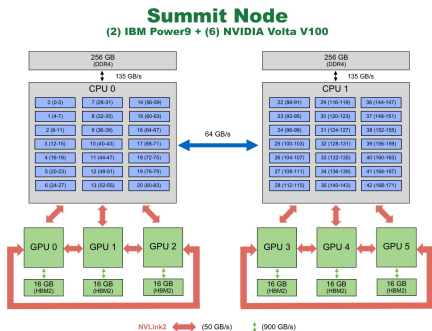
(#CPUs, #GPUs)	For a set of independent tasks	
	Approximation ratio	Worst case ex.
(1,1)	$\frac{1+\sqrt{5}}{2}$	$\frac{1+\sqrt{5}}{2}$
(m,n)	$2 + \sqrt{2} \approx 3.41$	$2 + \frac{2}{\sqrt{3}} \approx 3.15$



- Our upper bound is obtained by a linear program

# Minimizing impact of communications on Summit supercomputer

- Maximizing the overlap of communications and computations
- Implemented proposed approaches in Tensor Algebra for Manybody Methods (TAMM) library
- Molecular chemistry application (CCSD), Ubiquitin molecule, cc-pVDZ (737 basis functions, 220 nodes), aug-cc-pVDZ (1243 basis functions, 256 nodes)



Joint work with S. Krishnamoorthy and M. Zalewski during my postdoc at PNNL, USA

Fig source: <https://www.olcf.ornl.gov>

# Project: Scalable Tensor Algorithms for Modern Computing Systems

- 1 Design of Scalable Communication Optimal Algorithms for Tensors (Main Focus)
- 2 Extension of Existing Approaches/Algorithms (Short/Mid Term Research Plans)
- 3 Exploratory Topics (Mid/Long Term Research Plans)

## Scalable communication optimal algorithms for tensors

- Analyze existing algorithms
- Determine communication lower bounds
- Propose communication optimal algorithms
- Implement the proposed algorithms

**Main focus**

## Extension of existing approaches

- Strassen's concepts to tensors
- Concepts of hierarchical matrices to tensors
- Separation order of dimensions in tensor train

**Short/Mid term plans**

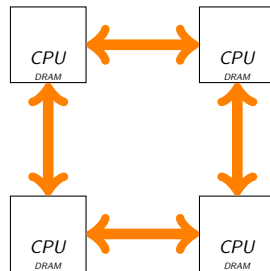
## Exploratory topics

- New tensor representations
- Architecture aware algorithms
- Randomization in tensors
- Factorizations of tensors

**Mid/Long term plans**

# Communication and its importance in HPC

- Running time of an algorithm depends on
  - Computations
    - Number of operations \* time-per-operation
  - Data movement
    - Volume of communication / Network-bandwidth
    - Number of messages \* Network-latency



- Gaps growing exponentially with time (Source: Getting up to speed: The future of supercomputing)

	time-per-operation	Network-bandwidth	Network-latency
Annual improvements	59 %	26 %	15 %

- Avoid communication to save time (and energy)

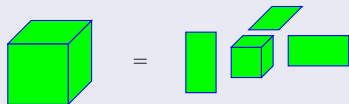
# Scalable algorithms for popular tensor operations

- Determine the communication lower bounds for tensor decompositions
- Analyse the popular decomposition algorithms and communications performed by them
- Propose new scalable communication optimal algorithms
  - If possible design tiles/tasks based algorithms
- Implement the proposed algorithms
  - Handle performance issues for homogeneous systems
    - Load balancing
    - Memory aware approaches
    - scheduling strategies
- Same for manipulation operations of popular tensor representations
- Extend implementation for heterogeneous systems (start with Nvidia GPUs based heterogeneous systems)
- Create a tensor library



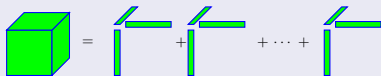
# Popular tensor decompositions

## Tucker decomposition



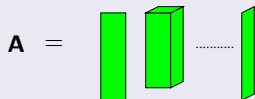
- Determine communication lower bounds for this operation
- Analyse communications performed by state of the art algorithms
- Propose and implement new scalable communication algorithms

## Canonical decomposition



- No deterministic algorithm to find the decomposition
- Analyse one iteration of the popular existing algorithms
- Propose and implement scalable algorithms for one iteration

## Tensor Train decomposition



- Determine communication lower bounds for this operation
- Analyse communication performed by popular algorithm
- Propose and implement new scalable communication algorithms

# Proving communication lower bounds for parallel computations

## How people did it for linear algebra operations?

- People obtain results for matrix multiplication operations
- Same lower bound apply to almost all direct linear algebra operations using reduction [Ballard et. al., 09] , for instance, bound for LU factorization

$$\begin{pmatrix} I & & -B \\ A & I & \\ & & I \end{pmatrix} = \begin{pmatrix} I & & \\ A & I & \\ & & I \end{pmatrix} \begin{pmatrix} I & -B \\ & I & AB \\ & & I \end{pmatrix}$$

## Approach to compute lower bounds for tensor computations

Notation: Tensors are denoted by solid shapes and number of lines denote the dimensions of the tensors. Connecting two lines implies summation (or contraction) over the connected dimensions.

- Obtain bounds for basic tensor operations: Tensor times matrix (TTM), Multiple tensor times matrix (Multi-TTM), Tensor contraction

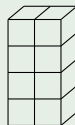


- Express decompositions and manipulations in terms of these basis operations

# Communication lower bounds on $P$ processors

- Recently started to work with L. Grigori (Inria Paris, France) and G. Ballard (Wake Forest University, USA)
- Revisited communication lower bounds for matrix multiplications
  - Expressed existing approaches in suitable forms for tensors
  - Lower bounds also instruct arrangement of processors in optimal algorithms
  - Improved the constants in the existing ranges of  $P$  (Demmel et.al [IPDPS 2013])
- Plan to continue this collaboration to compute lower bounds for tensor computations

## Arrangements of 8 processors



# Research Project

- 1 Design of Scalable Communication Optimal Algorithms for Tensors (Main Focus)
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# Strassen's concepts to tensors

## Matrix multiplication of $n \times n$ square matrices

- Complexity of traditional matrix multiplication is  $\mathcal{O}(n^3)$
- Strassen's matrix multiplication
  - Expressed matrix multiplication operation as a tensor operation
  - Canonical rank of the tensor determines the complexity of the operation
  - Complexity is  $\mathcal{O}(n^{2.81})$
- Plan to extend Strassen's concepts to tensor contractions

## Contraction of a 3-dimensional tensor with a matrix

```
for  $i_1 = 1 : n$  do
  for  $i_2 = 1 : n$  do
    for  $i_3 = 1 : n$  do
      for  $j_2 = 1 : n$  do
         $G(i_1, i_2, j_2) = G(i_1, i_2, j_2) + A(i_1, i_2, i_3) * B(i_3, j_2)$ 
      end for
    end for
  end for
end for
```

- Total  $\mathcal{O}(n^4)$  operations
- Apply Strassen's algorithm for each  $i_1$ , total  $\mathcal{O}(n^{3.81})$  operations
- Expressing as a canonical decomposition of  $8 \times 8 \times 4$  tensor can further reduce the number of operations

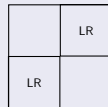
# Hierarchical Matrix concepts to Tensors

## Hierarchical Matrices

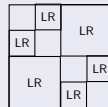
- Data sparse approximation of non-sparse matrices



Original Matrix



Step 1

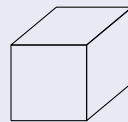


Step 2

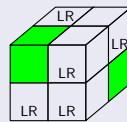
*LR*: low rank block

## Tensors

- $f(i, j, k) = \frac{1}{|i-j|+|j-k|+|k-i|}$
- Value is small if difference of any pair is large
- Formalize and evaluate this approach for tensors



Original Tensor



Step 1

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# Tensor Representations for High Dimensional Tensors

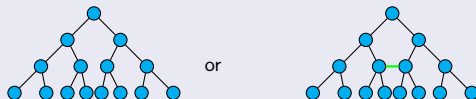
- Tensor Train is the popular representation to work with high dimensional tensors
- Adding tensors and applying an operator in this representation



- Requires a truncation process which iterates over cores one by one
- This representation is not much suited to work in parallel

## New Tensor Representations

- Look at new representations in tree format – suitable for parallelization

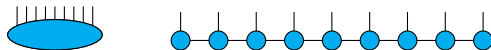


- Data will be stored at the leaf nodes
- Internal nodes will help to manipulate tensors in parallel

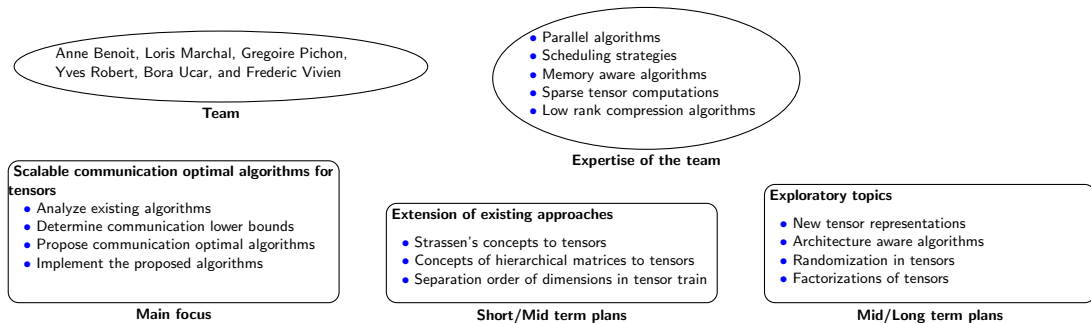


# Randomization in Tensor Computations

- Randomized SVD and UTV factorization are now well established
- Apply randomization to tensors
- Perform factorizations of tensors
  - For example: QR like factorization of a tensor



# Integration in the ROMA team



## Bringing additional skills in the team

- High dimensional dense tensor computations
- Communication lower bounds for linear algebra computations
- Scalable approaches for large HPC systems
- Use of tensors in quantum and molecular simulations