## Low rank approximations of tensors

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https://surakuma.github.io/courses/daamtc.html

## Properties of matrix Frobenius norm for real matrices

$$||A||_F^2 = \sum_{i,j} A^2(i,j) = \mathit{Trace}(AA^T) = \mathit{Trace}(A^TA)$$

$$||A + B||_F^2 = ||A||_F^2 + ||B||_F^2 + 2\langle A, B \rangle_F$$

Here  $\langle A, B \rangle_F$  is known as Frobenius inner product and defined as  $\langle A, B \rangle_F = Trace(A^T B) = Trace(B^T A)$ .

If Q is an orthonormal matrix then,

$$||A||_F^2 = ||QQ^T A||_F^2 + ||(I - QQ^T)A||_F^2,$$

$$||QC||_F = ||C||_F,$$

$$||Q^T A||_F = ||QQ^T A||_F \le ||A||_F,$$

$$\langle A - QQ^T A, QQ^T A \rangle_F = 0.$$

#### Tensor norm

• The norm of a tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times \cdots \times n_d}$  is analogous to the matrix Frobenius norm, and defined as

$$||\mathcal{A}||_F = \sqrt{\sum_{i_1=1}^{n_1} \sum_{i_2=1}^{n_2} \cdots \sum_{i_d=1}^{n_d} \mathcal{A}^2(i_1, i_2, \cdots, i_d)}$$

We will only focus on Frobenius norm in this course.

# Singular Value Decomposition (SVD)

- It decomposes a matrix  $A \in \mathbb{R}^{m \times n}$  to the form  $U \Sigma V^T$ 
  - U is an  $m \times m$  orthogonal matrix
  - V is an  $n \times n$  orthogonal matrix
  - $\Sigma$  is an  $m \times n$  rectangular diagonal matrix
- The diagonal entries  $\sigma_i = \Sigma_{ii}$  of  $\Sigma$  are called singular values
  - $\sigma_i \geq 0$  and  $\sigma_1 \geq \sigma_2 \geq \cdots \geq \sigma_{\min(m,n)}$
- The largest r such that  $\sigma_r \neq 0$  is called the rank of the matrix
- SVD represents a matrix as the sum of r rank one matrices



# Low rank approximations of matrices using SVD

SVD decomposition:  $A = U\Sigma V^T$ 

Let  $u_i$  and  $v_i$  be the column vectors of U and V, respectively.

### r'-rank approximation

If  $\tilde{A} = \sum_{i=1}^{r'} \sigma_i u_i v_i^T$ , then  $\tilde{A}$  is an r'-rank approximation of A.

$$||A - \tilde{A}||_F^2 = \sum_{i=r'+1}^{\min(m,n)} \sigma_i^2$$

SVD gives the best r'-rank approximation of any matrix.

### Approximation for $\epsilon$ accuracy

We select minimum r' such that  $\sum_{i=r'+1}^{\min(m,n)} \sigma_i^2 \le \epsilon^2$ . The approximation is

$$\tilde{A} = \sum_{i=1}^{r'} \sigma_i u_i v_i^T$$
.

$$||A - \tilde{A}||_F^2 = \sum_{i=r'+1}^{\min(m,n)} \sigma_i^2 \le \epsilon^2$$

## Properties of SVD

The SVD of  $A \in \mathbb{R}^{m \times n}$  can be written as  $A = U \Sigma V^T$ . Here  $U \in \mathbb{R}^{m \times m}$  and  $V \in \mathbb{R}^{n \times n}$  are orthogonal matrices and  $\Sigma \in \mathbb{R}^{m \times n}$  is a rectangular diagonal matrix.

- Columns of U are also eigen vectors of  $AA^T$
- Similarly, columns of V are eigen vectors of  $A^TA$
- If  $\sigma_i > 0$  is a singular value of A then  $\sigma_i^2$  is an eigen value of  $AA^T$  and  $A^TA$

 $\Sigma\Sigma^T$  and  $\Sigma^T\Sigma$  are diagonal matrices. Their diagonal entries are the eigen values of  $AA^T$  and  $A^TA$ , respectively.

We can also express SVD as

$$A = \begin{pmatrix} U_1 & U_2 \end{pmatrix} \begin{pmatrix} \Sigma_1 & 0 \\ 0 & \Sigma_2 \end{pmatrix} \begin{pmatrix} V_1 V_2 \end{pmatrix}^T = U_1 \Sigma_1 V_1^T + U_2 \Sigma_2 V_2^T.$$

This is equivalent to

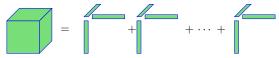
$$A = U_1 U_1^T A + U_2 U_2^T A = A V_1 V_1^T + A V_2 V_2^T.$$

### Table of Contents

- CP decomposition
- 2 Tucker decomposition
- Tensor Train decomposition
- Compact representations of tensor operations
- Miscellaneous

## CP decomposition of $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times \cdots \times n_d}$

It factorizes a tensor into a sum of rank one tensors.



CP decomposition of a 3-dimensional tensor.

$$\mathcal{A} = \sum_{\alpha=1}^{r} U_1(:,\alpha) \circ U_2(:,\alpha) \circ \cdots \circ U_d(:,\alpha)$$

It can be concisely expressed as  $\mathcal{A} = [\![U_1, U_2, \cdots, U_d]\!]$ . CP decomposition for a 3-dimensional tensor in matricized form can be written as:

$$A_{(1)} = U_1(U_3 \odot U_2)^T$$
,  $A_{(2)} = U_2(U_3 \odot U_1)^T$   $A_{(3)} = U_3(U_2 \odot U_1)^T$ .

It is useful to assume that  $U_1, U_2 \cdots U_d$  are normalized to length one with the weights given in a vector  $\lambda \in \mathbb{R}^r$ .

$$\mathcal{A} = [\![\lambda; U_1, U_2, \cdots, U_d]\!] = \sum_{\alpha=1}^r \lambda_\alpha U_1(:, \alpha) \circ U_2(:, \alpha) \circ \cdots \circ U_d(:, \alpha)$$

### Tensor rank

$$\mathcal{A} = \sum_{\alpha=1}^{r} \lambda_{\alpha} U_{1}(:,\alpha) \circ U_{2}(:,\alpha) \circ \cdots \circ U_{d}(:,\alpha)$$

ullet The minimum r required to express  ${\mathcal A}$  is called the rank of  ${\mathcal A}$ 

The rank of a real-valued tensor may be different over  $\mathbb R$  and  $\mathbb C$ . For example, consider the frontal slices of  $\mathcal A\in\mathbb R^{2\times 2\times 2}$ 

$$\mathcal{A}(:,:,1) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \text{ and } \mathcal{A}(:,:,2) = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}.$$

This has rank three over  $\mathbb R$  and two over  $\mathbb C$ . The CP decomposition over  $\mathbb R$  has the following factor matrices:

$$U_1 = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & -1 \end{pmatrix}, U_2 = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \end{pmatrix}, \text{ and } U_3 = \begin{pmatrix} 1 & 1 & 0 \\ -1 & 1 & 1 \end{pmatrix}.$$

The CP decomposition over  $\mathbb C$  has the following factor matrices:

$$\label{eq:U1} \textit{U}_1 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ -\textit{i} & \textit{i} \end{pmatrix}, \, \textit{U}_2 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ \textit{i} & -\textit{i} \end{pmatrix}, \, \, \text{and} \, \, \textit{U}_3 = \begin{pmatrix} 1 & 1 \\ \textit{i} & -\textit{i} \end{pmatrix}.$$

# Rank and low-rank approximations

• Determining the rank of a tensor is an NP-complete problem

• If  $A = \sum_{\alpha=1}^{r} \lambda_{\alpha} U_1(:, \alpha) \circ U_2(:, \alpha) \circ \cdots \circ U_d(:, \alpha)$ , summing k < r terms may not yield a best rank-k approximation

Possible that the best rank-k approximation of a tensor may not exist

# CP decomposition: example

Let  $\mathcal{A} \in \mathbb{R}^{2 \times 4 \times 3}$  and  $A = [U_1, U_2, U_3]$ . The rank of  $\mathcal{A}$  is 2.

$$U_1 = \begin{pmatrix} 1 & 3 \\ 2 & 4 \end{pmatrix}, \quad U_2 = \begin{pmatrix} 1 & 4 \\ 2 & 5 \\ 4 & 6 \\ 3 & 7 \end{pmatrix}, \quad U_3 = \begin{pmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{pmatrix}$$

Computation of  $\mathcal{A}(2,3,1)$ ,

$$\mathcal{A}(2,3,1) = \sum_{\alpha=1}^{2} U_1(2,\alpha) U_2(3,\alpha) U_3(1,\alpha)$$
$$= 2 \cdot 4 \cdot 1 + 4 \cdot 6 \cdot 4 = 104$$

 ${\cal A}$  has total 24 elements, while the CP representation has 18 elements.

### Table of Contents

- CP decomposition
  - Computing CP with Alternating Least Squares
- 2 Tucker decomposition
  - Computing Tucker decomposition
- Tensor Train decomposition
  - Computing Tensor Train decomposition
- 4 Compact representations of tensor operations
- Miscellaneous
  - Research topics/articles for the project
  - Randomized SVD
  - Strassen's algorithm: application of CP-decomposition

# CP optimization problem for a 3-dimensional tensor



For fixed rank k, we want to solve

$$\min_{U_1,U_2U_3}||\mathcal{A}-\sum_{\alpha=1}^k\lambda_\alpha U_1(:,\alpha)\circ U_2(:,\alpha)\circ U_3(:,\alpha)||_F.$$

- It is a nonlinear, nonconvex optimization problem
- In the matrix case, the SVD provides us the optimal solution
- In the tensor case, convergence to optimum not guaranteed

# Alternating Least Squares (ALS) method

Fixing all but one factor matrix, we have a linear least squares problem:

$$\min_{\hat{U}_1} ||\mathcal{A} - \sum_{\alpha=1}^k \hat{U}_1(:,\alpha) \circ U_2(:,\alpha) \circ U_3(:,\alpha)||_F$$

or equivalently

$$\min_{\hat{U}_1} ||A_{(1)} - \hat{U}_1(U_3 \odot U_2)^T||_F$$

ALS works by alternating over factor matrices, updating one at a time.

# CP-ALS algorithm

Repeat until maximum iterations reached or no further improvement obtained

- **1** Solve  $U_1(U_3 \odot U_2)^T = A_{(1)}$  for  $U_1 \Rightarrow U_1 = A_{(1)}(U_3 \odot U_2)(U_3^T U_3 * U_2^T U_2)^{\dagger}$
- ② Normalize columns of  $U_1$
- **3** Solve  $U_2(U_3 \odot U_1)^T = A_{(2)}$  for  $U_2 \Rightarrow U_2 = A_{(2)}(U_3 \odot U_1)(U_3^T U_3 * U_1^T U_1)^{\dagger}$
- 4 Normalize columns of  $U_2$
- **5** Solve  $U_3(U_2 \odot U_1)^T = A_{(3)}$  for  $U_3 \Rightarrow U_3 = A_{(3)}(U_2 \odot U_1)(U_2^T U_2 * U_1^T U_1)^{\dagger}$
- **1** Normalize columns of  $U_3$

Here  $A^{\dagger}$  denotes the Moore–Penrose pseudoinverse of A. We use the following identity to get expressions for  $U_1, U_2$  and  $U_3$ :

$$(A \odot B)^T (A \odot B) = A^T A * B^T B$$

# ALS for computing a CP decomposition

## **Algorithm 1** CP-ALS method to compute CP decomposition

**Require:** input tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ , desired rank k, initial factor matrices  $U_i \in \mathbb{R}^{n_j \times k}$  for 1 < j < d

**Ensure:**  $[\![\lambda; U_1, \cdots, U_d]\!]$  : a rank-k CP decomposition of  $\mathcal{A}$  repeat

$$\begin{aligned} & \textbf{for } i = 1 \text{ to } d \textbf{ do} \\ & V \leftarrow U_1^\mathsf{T} U_1 * \cdots * U_{i-1}^\mathsf{T} U_{i-1} * U_{i+1}^\mathsf{T} U_{i+1} * \cdots * U_d^\mathsf{T} U_d \\ & U_i \leftarrow A_{(i)} \big( U_d \odot \cdots \odot U_{i+1} \odot U_{i-1} \odot U_1 \big) \\ & U_i \leftarrow U_i V^\dagger \\ & \lambda \leftarrow \text{normalize colums of } U_i \end{aligned}$$

end for

until converge or the maximum number of iterations

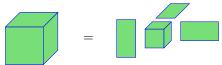
- The collective operation  $A_{(i)}(U_d \odot \cdots \odot U_{i+1} \odot U_{i-1} \odot U_1)$  is known as Matricized tensor times Khatri-Rao product (MTTKRP) computation
- $U_j$  can be chosen randomly or by setting k left singular vectors of  $A_{(j)}$  for 1 < i < d

### Table of Contents

- CP decomposition
- 2 Tucker decomposition
- Tensor Train decomposition
- 4 Compact representations of tensor operations
- Miscellaneous

## Tucker decomposition of $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times \cdots \times n_d}$

It represents a tensor with d matrices (usually orthonormal) and a small core tensor.



Tucker decomposition of a 3-dimensional tensor.

$$\mathcal{A} = \mathfrak{G} \times_1 U_1 \cdots \times_d U_d$$

$$\mathcal{A}(i_1, \cdots, i_d) = \sum_{i=1}^{r_1} \cdots \sum_{i=1}^{r_d} \mathfrak{G}(\alpha_1, \cdots, \alpha_d) U_1(i_1, \alpha_1) \cdots U_d(i_d, \alpha_d)$$

It can be concisely expressed as  $\mathcal{A} = \llbracket \mathcal{G}; U_1, \cdots, U_d 
rbracket$ .

Here  $r_j$  for  $1 \leq j \leq d$  denote a set of ranks. Matrices  $U_j \in \mathbb{R}^{n_j \times r_j}$  for  $1 \leq j \leq d$  are usually orthonormal and known as factor matrices. The tensor  $\mathfrak{G} \in \mathbb{R}^{r_1 \times r_2 \times \cdots \times r_d}$  is called the core tensor.

## Tucker decomposition: an example

Let 
$$\mathcal{A} \in \mathbb{R}^{3 \times 3 \times 3}$$
,  $\mathcal{G} \in \mathbb{R}^{2 \times 2 \times 2}$  and  $\mathcal{A} = \llbracket \mathcal{G}; \textit{U}_1, \textit{U}_2, \textit{U}_3 \rrbracket$ .

$$U_{1} = \frac{1}{3} \begin{pmatrix} 2 & -2 \\ 1 & 2 \\ 2 & 1 \end{pmatrix}, \quad U_{2} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 0 \end{pmatrix}, \quad U_{3} = \frac{1}{5} \begin{pmatrix} 0 & 4 \\ 3 & 3 \\ 4 & 0 \end{pmatrix}$$
$$S(:,:,1) = \begin{pmatrix} 1 & 4 \\ 2 & 5 \end{pmatrix}, \qquad S(:,:,2) = \begin{pmatrix} 7 & 10 \\ 8 & 11 \end{pmatrix}$$

$$\begin{split} \mathcal{A}(3,2,1) &= \sum_{\alpha_1=1}^2 \sum_{\alpha_2=1}^2 \sum_{\alpha_3=1}^2 \mathcal{G}(\alpha_1,\alpha_2,\alpha_3) U_1(3,\alpha_1) U_2(2,\alpha_2) U_3(1,\alpha_3) \\ &= \mathcal{G}(1,1,1) U_1(3,1) U_2(2,1) U_3(1,1) + \mathcal{G}(1,1,2) U_1(3,1) U_2(2,1) U_3(1,2) \\ &+ \mathcal{G}(1,2,1) U_1(3,1) U_2(2,2) U_3(1,1) + \mathcal{G}(1,2,2) U_1(3,1) U_2(2,2) U_3(1,2) \\ &+ \mathcal{G}(2,1,1) U_1(3,2) U_2(2,1) U_3(1,1) + \mathcal{G}(2,1,2) U_1(3,2) U_2(2,1) U_3(1,2) \\ &+ \mathcal{G}(2,2,1) U_1(3,2) U_2(2,2) U_3(1,1) + \mathcal{G}(2,2,2) U_1(3,2) U_2(2,2) U_3(1,2) \\ &= 1 \cdot \frac{2}{3} \cdot 0 \cdot 0 + 7 \cdot \frac{2}{3} \cdot 0 \cdot \frac{4}{5} + 4 \cdot \frac{2}{3} \cdot 1 \cdot 0 + 10 \cdot \frac{2}{3} \cdot 1 \cdot \frac{4}{5} \\ &+ 2 \cdot \frac{1}{3} \cdot 0 \cdot 0 + 8 \cdot \frac{1}{3} \cdot 0 \cdot \frac{4}{5} + 5 \cdot \frac{1}{3} \cdot 1 \cdot 0 + 11 \cdot \frac{1}{3} \cdot 1 \cdot \frac{4}{5} = \frac{124}{15}. \end{split}$$

### Table of Contents

- CP decomposition
  - Computing CP with Alternating Least Squares
- 2 Tucker decomposition
  - Computing Tucker decomposition
- Tensor Train decomposition
  - Computing Tensor Train decomposition
- 4 Compact representations of tensor operations
- Miscellaneous
  - Research topics/articles for the project
  - Randomized SVD
  - Strassen's algorithm: application of CP-decomposition

## High Order SVD (HOSVD) for computing a Tucker decomposition

### Algorithm 2 HOSVD method to compute a Tucker decomposition

**Require:** input tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ , desired rank  $(r_1, \cdots, r_d)$ 

**Ensure:** 
$$A = 9 \times_1 U_1 \times_2 U_2 \cdots \times_d U_d$$

for 
$$k = 1$$
 to  $d$  do

 $U_k \leftarrow r_k$  leading left singular vectors of  $A_{(k)}$ 

#### end for

$$\mathfrak{G} = \mathcal{A} \times_1 U_1^\mathsf{T} \times_2 U_2^\mathsf{T} \cdots \times_d U_d^\mathsf{T}$$

- When r<sub>i</sub> < rank(A<sub>(i)</sub>) for one or more i, the decomposition is called the truncated-HOSVD (T-HOSVD)
- Output of T-HOSVD can be used as a starting point for an ALS algorithm
- The collective operation  $\mathcal{A} \times_1 U_1^\mathsf{T} \times_2 U_2^\mathsf{T} \cdots \times_d U_d^\mathsf{T}$  is known as Multiple Tensor-Times-Matrix (Multi-TTM) computation

21/50

## Quasi-optimality of T-HOSVD

Let  $\mathcal{A} = \mathcal{G} \times_1 U_1 \times_2 U_2 \cdots \times_d U_d$  be the tensor obtained from T-HOSVD.

$$\begin{split} ||\mathcal{A} - \tilde{\mathcal{A}}||_F^2 = &||\mathcal{A} - \mathcal{G} \times_1 U_1 \times_2 U_2 \cdots \times_d U_d||_F^2 = ||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \cdots \times_d U_d U_d^\mathsf{T}||_F^2 \\ = &||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} + \mathcal{A} \times_1 U_1 U_1^\mathsf{T} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \cdots \times_d U_d U_d^\mathsf{T}||_F^2 \\ = &||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T}||_F^2 + ||\mathcal{A} \times_1 U_1 U_1^\mathsf{T} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \cdots \times_d U_d U_d^\mathsf{T}||_F^2 \\ = &||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T}||_F^2 + ||\mathcal{A} \times_1 U_1 U_1^\mathsf{T} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \times_2 U_2 U_2^\mathsf{T}||_F^2 + \cdots \\ & \cdots + ||\mathcal{A} \times_1 U_1 U_1^\mathsf{T} \cdots \times_{d-1} U_{d-1} U_{d-1}^\mathsf{T} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \cdots \times_d U_d U_d^\mathsf{T}||_F^2 \\ \leq &||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T}||_F^2 + ||\mathcal{A} - \mathcal{A} \times_2 U_2 U_2^\mathsf{T}||_F^2 + \cdots + ||\mathcal{A} - \mathcal{A} \times_d U_d U_d^\mathsf{T}||_F^2 \end{split}$$

#### Theorem

Tensor A obtained from T-HOSVD satisfies quasi-optimality condition

$$||A-\tilde{\mathcal{A}}||_F \leq \sqrt{d}||\mathcal{A}-\mathcal{A}_{\textit{best}}||_F$$
 ,

where  $\mathcal{A}_{best}$  is the best approximation of  $\mathcal{A}$  with ranks  $(r_1, \dots, r_d)$ .

Proof:  $||\mathcal{A} - \mathcal{A} \times_i U_i U_i^{\mathsf{T}}||_F \le ||\mathcal{A} - \mathcal{A}_{best}||_F$  for  $1 \le i \le d$ . Substituting these in the previous result yields the specified inequality.

## Sequentially T-HOSVD (ST-HOSVD) for Tucker decomposition

- This method is more work efficient than T-HOSVD
- In each step, it reduces the size of one dimension of the tensor

## Algorithm 3 ST-HOSVD method to compute a Tucker decomposition

```
Require: input tensor \mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}, desired rank (r_1, \cdots, r_d)

Ensure: [\![\mathcal{G}; U_1, \cdots, U_d]\!]: a (r_1, \cdots, r_d)-rank Tucker decomposition of \mathcal{A}

\mathcal{B} \leftarrow \mathcal{A}

for k = 1 to d do

U_k \leftarrow r_k leading singular vectors of B_{(k)}

\mathcal{B} \leftarrow \mathcal{B} \times_k U_k^T

end for

\mathcal{G} = \mathcal{B}
```

# Quasi-optimality of ST-HOSVD

Let  $\tilde{A} = \mathfrak{G} \times_1 U_1 \times_2 U_2 \cdots \times_d U_d$  be the tensor obtained from ST-HOSVD.

$$\begin{aligned} ||\mathcal{A} - \tilde{\mathcal{A}}||_{F}^{2} &= ||\mathcal{A} - \mathcal{G} \times_{1} U_{1} \times_{2} U_{2} \cdots \times_{d} U_{d}||_{F}^{2} &= ||\mathcal{A} - \mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}} \cdots \times_{d} U_{d} U_{d}^{\mathsf{T}}||_{F}^{2} \\ &= ||\mathcal{A} - \mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}}||_{F}^{2} + ||\mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}} - \mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}} \times_{2} U_{2} U_{2}^{\mathsf{T}}||_{F}^{2} + \cdots \\ &\cdots + ||\mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}} \cdots \times_{d-1} U_{d-1} U_{d-1}^{\mathsf{T}} - \mathcal{A} \times_{1} U_{1} U_{1}^{\mathsf{T}} \cdots \times_{d} U_{d} U_{d}^{\mathsf{T}}||_{F}^{2} \end{aligned}$$

#### Theorem

Tensor  $\widehat{\mathcal{A}}$  obtained from ST-HOSVD satisfies quasi-optimality condition

$$||A-\tilde{\mathcal{A}}||_F \leq \sqrt{d}||\mathcal{A}-\mathcal{A}_{\textit{best}}||_F$$
 ,

where  $A_{best}$  is the best approximation of A with ranks  $(r_1, \dots, r_d)$ .

Proof: We know that  $||\mathcal{A} - \mathcal{A} \times_i U_i U_i^\mathsf{T}||_F \le ||\mathcal{A} - \mathcal{A}_{best}||_F$  for  $1 \le i \le d$ .

$$||\mathcal{A} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T}||_F \le ||\mathcal{A} - \mathcal{A}_{\textit{best}}||_F$$

$$||\mathcal{A} \times_1 U_1 U_1^\mathsf{T} - \mathcal{A} \times_1 U_1 U_1^\mathsf{T} \times_2 U_2 U_2^\mathsf{T}||_F \leq ||\mathcal{A} - \mathcal{A} \times_2 U_2 U_2^\mathsf{T}||_F \leq ||\mathcal{A} - \mathcal{A}_{best}||_F$$

 $||\mathcal{A}\times_1 U_1 U_1^\mathsf{T} \cdots \times_{d-1} U_{d-1} U_{d-1}^\mathsf{T} - \mathcal{A}\times_1 U_1 U_1^\mathsf{T} \cdots \times_d U_d U_d^\mathsf{T}||_F \leq ||\mathcal{A}-\mathcal{A}\times_d U_d U_d^\mathsf{T}||_F \leq ||\mathcal{A}-\mathcal{A}_{best}||_F$  Summing the above terms yields the specified inequality.

## Tucker decomposition optimization problem for a 3-dimensional tensor



For fixed ranks orthonormal matrices  $U_1, U_2, U_3$ , we want to solve

$$\min_{U_1,U_2,U_3} ||\mathcal{A} - \mathcal{G} \times_1 U_1 \times_2 U_2 \times_3 U_3||_F \text{, where } \mathcal{G} = \mathcal{A} \times_1 U_1^T \times_2 U_2^T \times_3 U_3^T.$$

This is equivalent to

$$\max_{U_1,U_2,U_3} ||\mathcal{A} \times_1 U_1^T \times_2 U_2^T \times_3 U_3^T||_F.$$

It is a nonlinear, nonconvex optimization problem.

$$\begin{split} ||\mathcal{A} - \mathcal{G} \times_{1} U_{1} \times_{2} U_{2} \times_{3} U_{3}||_{F}^{2} = &||\mathcal{A}||_{F}^{2} + ||\mathcal{G} \times_{1} U_{1} \times_{2} U_{2} \times_{3} U_{3}||_{F}^{2} \\ &- 2\langle \mathcal{A} - \mathcal{G} \times_{1} U_{1} \times_{2} U_{2} \times_{3} U_{3}\rangle \\ = &||\mathcal{A}||_{F}^{2} + ||\mathcal{G}||_{F}^{2} - 2\langle \mathcal{A} \times_{1} U_{1}^{T} \times_{2} U_{2}^{T} \times_{3} U_{3}^{T}, \mathcal{G}\rangle \\ = &||\mathcal{A}||_{F}^{2} - ||\mathcal{G}||_{F}^{2} \end{split}$$

# Higher-order orthogonal iteration (HOOI) method

Fixing all but one factor matrix, we have a matrix problem:

$$\max_{\hat{U_1}} ||\mathcal{A} \times_1 \hat{U_1}^T \times_2 U_2^T \times_3 U_3^T||_F$$

HOOI works by alternating over factor matrices, updating one by computing left singular vectors

## HOOI method for computing a Tucker decomposition

## Algorithm 4 HOOI method to compute Tucker decomposition

**Require:** input tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ , desired ranks  $(r_1, \cdots, r_d)$ , initial factor matrices  $U_j \in \mathbb{R}^{n_j \times r_j}$  for  $1 \leq j \leq d$ 

**Ensure:**  $[\![\mathcal{G};U_1,\cdots,U_d]\!]$ : a  $(r_1,\cdots,r_d)$ -rank Tucker decomposition of  $\mathcal{A}$  repeat

for 
$$i = 1$$
 to  $d$  do
$$\mathcal{B} \leftarrow \mathcal{A} \times_1 U_1^T \cdots \times_{i-1} U_{i-1}^T \times_{i+1} U_{i+1}^T \cdots \times_d U_d^T$$

$$U_i \leftarrow r_i \text{ left singular vectors of } \mathcal{B}_{(i)}$$
end for

until converge or the maximum number of iterations

$$\mathfrak{G} \leftarrow \mathcal{A} \times_1 U_1^T \times_2 U_2^T \cdots \times_d U_d^T$$

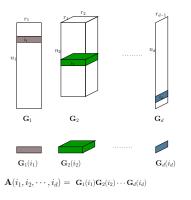
• Outputs of HOSVD ( $U_j$  for  $1 \le j \le d$ ) can be used as a starting point for this method

### Table of Contents

- CP decomposition
- 2 Tucker decomposition
- Tensor Train decomposition
- 4 Compact representations of tensor operations
- Miscellaneous

### Tensor Train (TT) decomposition: Product of matrices view

 A d-dimensional tensor is represented with 2 matrices and d-2 3-dimensional tensors.



An entry of  $A \in \mathbb{R}^{n_1 \times \cdots \times n_d}$  is computed by multiplying corresponding matrix (or row/column) of each matrix/tensor.

## Tensor Train decomposition

 $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$  is represented with cores  $g_k \in \mathbb{R}^{r_{k-1} \times n_k \times r_k}$ ,  $k=1,2,\cdots d$ ,  $r_0=r_d=1$  and its elements satisfy the following expression:

$$\mathcal{A}(i_1, \dots, i_d) = \sum_{\alpha_0=1}^{r_0} \dots \sum_{\alpha_d=1}^{r_d} \mathcal{G}_1(\alpha_0, i_1, \alpha_1) \dots \mathcal{G}_d(\alpha_{d-1}, i_d, \alpha_d)$$

$$= \sum_{\alpha_1=1}^{r_1} \dots \sum_{\alpha_{d-1}=1}^{r_{d-1}} \mathcal{G}_1(1, i_1, \alpha_1) \dots \mathcal{G}_d(\alpha_{d-1}, i_d, 1)$$

$$i_1\alpha_1 \dots \alpha_1 \dots \alpha_{d-1} \dots \alpha$$

The ranks  $r_k$  are called TT-ranks.

• The number of entries in this decomposition =  $\mathcal{O}(n_1r_1 + n_2r_1r_2 + n_3r_2r_3 + \cdots + n_{d-1}r_{d-2}r_{d-1} + n_dr_{d-1})$ 

## TT-decomposition: an example

Let  $\mathcal{A} \in \mathbb{R}^{3 \times 4 \times 5}$ .  $\mathcal{G}_1 \in \mathbb{R}^{3 \times 2}, \mathcal{G}_2 \in \mathbb{R}^{2 \times 4 \times 2}$ , and  $\mathcal{G}_3 \in \mathbb{R}^{2 \times 5}$  are the cores of a TT-decomposition.

$$\mathfrak{G}_1 = \begin{pmatrix} 1 & 1 \\ 2 & 1 \\ 3 & 1 \end{pmatrix}, \quad \mathfrak{G}_3 = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix},$$

$$\mathcal{G}_{2}(:,1,:) = \begin{pmatrix} 1 & 1 \\ 2 & 1 \end{pmatrix}, \mathcal{G}_{2}(:,2,:) = \begin{pmatrix} 1 & 1 \\ 3 & 1 \end{pmatrix}, \mathcal{G}_{2}(:,3,:) = \begin{pmatrix} 1 & 1 \\ 4 & 1 \end{pmatrix}, \mathcal{G}_{2}(:,4,:) = \begin{pmatrix} 1 & 1 \\ 5 & 1 \end{pmatrix}$$

Computation of  $\mathcal{A}(2,3,4)$ ,

$$\begin{split} \mathcal{A}(2,3,4) = & \mathcal{G}_1(2,:) \mathcal{G}_2(:,3,:) \mathcal{G}_3(:,4) \\ = & \begin{pmatrix} 2 & 1 \end{pmatrix} \begin{pmatrix} 1 & 1 \\ 4 & 1 \end{pmatrix} \begin{pmatrix} 4 \\ 1 \end{pmatrix} = 27 \end{split}$$

## Another representation of unfolding matrices of a tensor

 $A_k$  denotes k-th unfolding matrix of tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$ .

$$A_k = [A_k(i_1, i_2, \cdots, i_k; i_{k+1}, \cdots, i_d)]$$

• Size of  $A_k$  is  $(\prod_{\ell=1}^k n_\ell) \times (\prod_{\ell=k+1}^d n_\ell)$ 

### Table of Contents

- CP decomposition
  - Computing CP with Alternating Least Squares
- 2 Tucker decomposition
  - Computing Tucker decomposition
- Tensor Train decomposition
  - Computing Tensor Train decomposition
- 4 Compact representations of tensor operations
- 6 Miscellaneous
  - Research topics/articles for the project
  - Randomized SVD
  - Strassen's algorithm: application of CP-decomposition

# TT-SVD algorithm for TT approximation [Oseledets, 2011]

### **Algorithm 5** TT-SVD algorithm

**Require:** d-dimensional tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times \cdots \times n_d}$  and desired ranks  $(r_0 = 1, r_1, r_2, \cdots r_{d-1}, r_d = 1)$ 

**Ensure:** Cores  $\mathcal{G}_k \in \mathbb{R}^{r_{k-1} \times n_k \times r_k}$  for  $1 \leq k \leq d$  of a TT representation

- 1: Temporary tensor:  $\mathfrak{C} = \mathcal{A}$
- 2: **for** k = 1 : d 1 **do**
- 3:  $A_k = reshape(\mathfrak{C}, r_{k-1}n_k, \frac{numel(\mathfrak{C})}{r_{k-1}n_k})$
- 4: Compute SVD:  $A_k = U \Sigma V^T$
- 5: New core:  $g_k := reshape(U(; 1 : r_k), r_{k-1}, n_k, r_k)$
- 6:  $C = \Sigma(1:r_k;1:r_k)V^T(1:r_k;)$
- 7: end for
- 8:  $9_d = 0$
- 9: return  $\mathcal{G}_1, \cdots, \mathcal{G}_d$
- $reshape(A, m_1, \dots, m_\ell)$ : rearranges array A into a  $m_1 \times \dots \times m_\ell$  array
- numel(A): number of elements of array A

## Error with TT-SVD approximation

Suppose the unfolding matrices of  ${\mathcal A}$  satisfy the following:

 $A_k = R_k + E_k$ ,  $R_k$  is the best  $r_k$ - rank approximation of  $A_k$ , for  $1 \le k \le d-1$ .

The accuracy analysis of TT-SVD is similar to that of ST-HOSVD method (see [Oseledets, 2011]).

Tensor  ${\mathfrak B}$  obtained from the TT-SVD algorithm satisfies

$$||\mathcal{A} - \mathcal{B}||_F^2 \le \sum_{k=1}^{d-1} ||E_k||_F^2.$$

#### Theorem

Tensor B obtained from TT-SVD satisfies quasi-optimality condition

$$||A - \mathcal{B}||_F \leq \sqrt{d-1}||\mathcal{A} - \mathcal{A}_{best}||_F$$
,

where  $A_{best}$  is the best  $(r_1, \dots, r_{d-1})$ -ranks approximation of A in TT-format.

Proof: As SVD gives the best  $r_k$  rank approximation for  $A_k$ , we have

$$||E_k||_F \leq ||\mathcal{A} - \mathcal{A}_{best}||_F$$
 for  $1 \leq k \leq d$ .

Putting the values of  $||E_k||_F$  in the error expression of TT-SVD algorithm completes the proof.

# Why TT representation is good for high dimension tensors?

This representation allows one to perform various basic linear algebra operations in its own structure.

Addition: The addition of two tensors in the TT-representation ,

$$\mathcal{A} = \mathcal{A}_1(i_1)\cdots\mathcal{A}_d(i_d), \quad \mathcal{B} = \mathcal{B}_1(i_1)\cdots\mathcal{B}_d(i_d),$$

requires to merge cores for each mode. Auxiliary dimensions are added. The cores  $\mathcal{C}_k(i_k)$  of  $\mathcal{C} = \mathcal{A} + \mathcal{B}$  are defined as

$$\mathfrak{C}_k(i_k) = egin{pmatrix} \mathcal{A}_k(i_k) & 0 \ 0 & \mathcal{B}_k(i_k) \end{pmatrix}$$
 , for  $2 \leq k \leq d-1$ , and

$$\mathcal{C}_1(i_1) = \begin{pmatrix} \mathcal{A}_1(i_1) & \mathcal{B}_1(i_1) \end{pmatrix}, \quad \mathcal{C}_d(i_d) = \begin{pmatrix} \mathcal{A}_d(i_d) \\ \mathcal{B}_d(i_d) \end{pmatrix}.$$

- Multiplication by a number: requires to scale one of the cores
- Multidimensional contraction, Hadamard product and scalar product can also be performed
- Further approximation (or compression) can also be obtained

- CP decomposition
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## Tensor network representations

Notation: Tensors are denoted by solid shapes and number of lines denote the dimensions of the tensors. Connecting two lines implies summation (or contraction) over the connected dimensions.

Vector: Matrix: 3-dimensional tensor: Tucker decomposition of a 3-dimensional tensor:

TT decomposition of of a 4-dimensional tensor

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  - Randomized SVD
  - Strassen's algorithm: application of CP-decomposition

# Course project

- A list of topics/articles is given
- Each student or a group of two students will prepare a 5-6 pages report for the chosen topic/article
- Deadline for submitting the report: Nov 5
- Presentation would be after Nov 5
- Email me your or your group topic/article choices (atleast two) in preference order by Oct 15

If you want to propose another topic or article, your are more than welcome to discuss it with me.

# Research topics

- Communication costs of a specific matrix factorization
- Extending a specific matrix factorization for tensors
- Use of tensors in a particular domain
  - Neuroscience, data analysis, molecular simulations, quantum computing, face recognition

## What do I expect from you in the report?

- State-of-the-art of the field
- Bottleneck part of the operation
- Your idea of improvement and preliminary work on why it will be effective

#### Research articles

- Obtain lower bounds on data transfers for various computations on a sequential machine:
   Automated Derivation of Parametric Data Movement Lower Bounds for Affine Programs
- Performance optimizations for TSQR algorithm: Reconstructing Householder Vectors from Tall-Skinny QR
- Low rank approximation for stencil computations: LoRAStencil: Low-Rank Adaptation of Stencil Computation on Tensor Cores
- Sequential lower bounds and optimal algorithms for symmetric computations:
   I/O-Optimal Algorithms for Symmetric Linear Algebra Kernels
- Hypergraph partitioning-based methods to improve MTTKRP performance: Scalable Sparse Tensor Decompositions in Distributed Memory Systems
- A parallel method to perform MTTKRP on a parallel shared memory machine: SPLATT: Efficient and Parallel Sparse Tensor-Matrix Multiplication
- Randomization based parallel HOSVD and ST-HOSVD methods: Parallel Randomized Tucker Decomposition Algorithms
- Tucker decomposition to improve performance of convolution kernels: Stable Low-rank Tensor Decomposition for Compression of Convolutional Neural Network
- Tensor train representation for the weight matrices of the fully connected layers:
   Tensorizing Neural Networks
- Use of tensor train representation in quantum systems: The density-matrix renormalization group: a short introduction

# Contents of the report for a research article

- The general idea of the work
- A detailed analysis of some parts
- Overview of the state of the art
- Mention why the work of this paper is important
- Your feedback on the work (possible extensions, limitations of the work, ...)
- What challenges you faced while reading the paper (which parts are not clear, explanation is not appropriate, missing information, ...)

Each group (or person) will do a presentation of the selected topic/article for 20-35 minutes, followed by 5-10 minutes of questions/comments.

- CP decomposition
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## Main idea of randomized SVD

We want to find r-rank approximation of  $A \in \mathbb{R}^{m \times n}$ . We select a matrix Q with  $\ell$   $(r \le \ell \le n)$  orthonormal columns that well approximates the action of A,  $A \approx QQ^TA$ .

- **1** Construct  $B = Q^T A$
- **2** Perform SVD of B,  $B = \tilde{U}\Sigma V^T$
- $\odot$  Set  $U = Q\tilde{U}$
- Return  $U, \Sigma, V$

## A simple way to find Q

- **①** Construct a Gaussian random matrix  $\Omega$  of  $n \times \ell$  size
- **2** Form  $X = A\Omega$
- **3** Obtain an orthonormal matrix using QR factorization, X = QR

Usually  $\ell - r$  is a small constant, such as 5 or 10.

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# Strassen's algorithm for matrix multiplication (C = AB)

Matrix is divided into 2×2 blocks

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$$

$$M_2 = (A_{21} + A_{22})B_{11}$$

$$M_3 = A_{11}(B_{12} - B_{22})$$

$$M_4 = A_{22}(B_{21} - B_{11})$$

$$M_5 = (A_{11} + A_{12})B_{22}$$

$$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$$

$$M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$$

$$C_{11} = M_1 + M_4 - M_5 + M_7$$
  
 $C_{12} = M_3 + M_5$   
 $C_{21} = M_2 + M_4$   
 $C_{22} = M_1 - M_2 + M_3 + M_6$ 

## $2 \times 2$ Matrix multiplication as a tensor operation

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

We can write this multiplication as a tensor operation,

$$\mathfrak{T} \times_1 \begin{pmatrix} A_{11} \\ A_{12} \\ A_{21} \\ A_{22} \end{pmatrix}^T \times_2 \begin{pmatrix} B_{11} \\ B_{12} \\ B_{21} \\ B_{22} \end{pmatrix}^T = \begin{pmatrix} C_{11} \\ C_{12} \\ C_{21} \\ C_{22} \end{pmatrix}$$

Where  $\mathfrak{T}$  is a  $4 \times 4 \times 4$  tensor with the following frontal slices:

$$T_1 = \left(\begin{smallmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{smallmatrix}\right) T_2 = \left(\begin{smallmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{smallmatrix}\right) \quad T_3 = \left(\begin{smallmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{smallmatrix}\right) T_4 = \left(\begin{smallmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{smallmatrix}\right)$$

# $2 \times 2$ Matrix multiplication as a tensor operation

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

We can write this multiplication as a tensor operation,

$$\mathfrak{T} \times_1 \begin{pmatrix} A_{11} \\ A_{12} \\ A_{21} \\ A_{22} \end{pmatrix}^T \times_2 \begin{pmatrix} B_{11} \\ B_{12} \\ B_{21} \\ B_{22} \end{pmatrix}^T = \begin{pmatrix} C_{11} \\ C_{12} \\ C_{21} \\ C_{22} \end{pmatrix}$$

For example,

$$T_{2} \times_{1} \begin{pmatrix} A_{11} \\ A_{12} \\ A_{21} \\ A_{22} \end{pmatrix}^{T} \times_{2} \begin{pmatrix} B_{11} \\ B_{12} \\ B_{21} \\ B_{22} \end{pmatrix}^{T} = (A_{11} \ A_{12} \ A_{21} \ A_{22}) \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} B_{11} \\ B_{12} \\ B_{21} \\ B_{22} \end{pmatrix} = A_{11} B_{12} + A_{12} B_{22} = C_{12}$$

# Matrix multiplication with CP decomposition

CP decomposition of  $\mathfrak{T}$ ,  $\mathfrak{T} = \llbracket U, V, W \rrbracket$  can be written as,

$$\mathfrak{T} = \sum_{r=1}^{R} u_r \circ v_r \circ w_r$$

Here  $u_r$ ,  $v_r$  and  $w_r$  are the columns of U, V and W, respectively. R is the rank of  $\mathfrak{T}$ . We can write matrix multiplication as,

$$\mathfrak{I} \times_{1} \begin{pmatrix} A_{11} \\ A_{12} \\ A_{22} \end{pmatrix}^{T} \times_{2} \begin{pmatrix} B_{11} \\ B_{12} \\ B_{21} \\ B_{22} \end{pmatrix}^{T} = \sum_{r=1}^{R} (u_{r} \circ v_{r} \circ w_{r}) \times_{1} \begin{pmatrix} A_{11} \\ A_{12} \\ A_{21} \\ A_{22} \end{pmatrix}^{T} \times_{2} \begin{pmatrix} B_{11} \\ B_{12} \\ B_{22} \end{pmatrix}^{T} \\
= \sum_{r=1}^{R} \left[ (A_{11} A_{12} A_{21} A_{22}) u_{r} (B_{11} B_{12} B_{21} B_{22}) v_{r} \right] w_{r} = \begin{pmatrix} C_{11} \\ C_{12} \\ C_{21} \\ C_{22} \end{pmatrix}$$

# Factor matrices and Strassen's algorithm

Factor matrices.

$$V = egin{pmatrix} 1 & 1 & 0 & -1 & 0 & 1 & 0 \ 0 & 0 & 1 & 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 1 & 0 & 0 & 1 \ 1 & 0 & -1 & 0 & 1 & 0 & 1 \end{pmatrix}$$

$$W = \begin{pmatrix} 1 & 0 & 0 & 1 & -1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & -1 & 1 & 0 & 0 & 1 & 0 \end{pmatrix} \qquad C_{11} = M_1 + M_4 - M_5 + M_7$$

$$C_{12} = M_3 + M_5$$

Strassen's algorithm,

$$U = \begin{pmatrix} 1 & 0 & 1 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & -1 \end{pmatrix}$$

$$M_1 = (A_{11} + A_{22})(B_{11} + B_{22})$$

$$M_2 = (A_{21} + A_{22})B_{11}$$

$$M_3 = A_{11}(B_{12} - B_{22})$$

$$M_4 = A_{22}(B_{21} - B_{11})$$

$$M_5 = (A_{11} + A_{12})B_{22}$$

$$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$$

$$C_{11} = M_1 + M_4 - M_5 + M_7$$

$$C_{12} = M_3 + M_5$$

$$C_{21} = M_2 + M_4$$

$$C_{22} = M_1 - M_2 + M_3 + M_6$$

 $M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$ 

Factor matrices U, V and W construct the algorithm.