1

1.* Intro. This program generates clauses that enforce the constraint $x_1 + \cdots + x_n \le r$, using a method due to Olivier Bailleux and Yacine Boufkhad [Lecture Notes in Computer Science 2833 (2003), 108–122]. It introduces at most (n-2)r new variables Bi.j for $2 \le i < n$ and $1 \le j \le r$, and a number of clauses that I haven't yet tried to count carefully, but it is at most O(nr). All clauses have length 3 or less.

This version inputs a graph (specified as a third parameter) and and color number (specified fourth). The output clauses will limit the number of vertices of that color.

```
#define nmax 10000
#include <stdio.h>
#include <stdlib.h>
#include "gb_graph.h"
#include "gb_save.h"
                          /* the given parameters */
   int n, r, kk;
   int count[nmax + nmax];
                                           /* the number of leaves below each node */
   main(\mathbf{int} \ argc, \mathbf{char} *argv[])
      register int i, j, k, jl, jr, t, tl, tr;
      Graph * q;
      \langle \text{Process the command line } 2^* \rangle;
      if (r \equiv 0) (Handle the trivial case directly 6)
      else {
         \langle Build the complete binary tree with n leaves 3\rangle;
         for (i = n - 2; i; i--) (Generate the clauses for node i \ 4^*);
         ⟨Generate the clauses at the root 5⟩;
   }
2* \langle \text{Process the command line } 2^* \rangle \equiv
   \textbf{if} \; (\mathit{argc} \neq 5 \lor \mathit{sscanf} \, (\mathit{argv} \, [1], \, \texttt{"%d"}, \&n) \neq 1 \lor \mathit{sscanf} \, (\mathit{argv} \, [2], \, \texttt{"%d"}, \&r) \neq 1 \lor \mathit{sscanf} \, (\mathit{argv} \, [4], \, \texttt{"%d"}, \&kk) \neq 1)
      fprintf(stderr, "Usage: \_\%s \_n \_r \_foo.gb \_k n", argv[0]);
      exit(-1);
   }
   g = restore\_graph(argv[3]);
   if (\neg g) {
      fprintf(stderr, "I_{\square}can't_{\square}input_{\square}the_{\square}graph_{\square}'%s'! \n", argv[3]);
      exit(-2);
    \textbf{if} \ (g \neg n \neq n) \ \textit{fprintf} \ (\textit{stderr}, \texttt{"Warning:} \bot \texttt{The} \bot \texttt{graph} \bot \texttt{has} \bot \texttt{"Id} \bot \texttt{vertices}, \bot \texttt{not} \bot \texttt{"d!} \setminus \texttt{n"}, g \neg n, n); \\
   if (n > nmax) {
      fprintf(stderr, "Recompile\_me: \_I'd\_don't\_allow\_n>%d\n", nmax);
      exit(-2);
   if (r < 0 \lor r \ge n) {
      fprintf(stderr, "Eh? \_r \_should \_be \_between \_0 \_and \_n-1! \ ");
   This code is used in section 1*.
```

3. The tree has 2n-1 nodes, with 0 as the root; the leaves start at node n-1. Nonleaf node k has left child 2k+1 and right child 2k+2. Here we simply fill the *count* array.

```
 \langle \text{ Build the complete binary tree with } n \text{ leaves } 3 \rangle \equiv \\ \text{ for } (k=n+n-2; \ k \geq n-1; \ k--) \ count[k] = 1; \\ \text{ for } (\ ; \ k \geq 0; \ k--) \ count[k] = count[k+k+1] + count[k+k+2]; \\ \text{ if } (count[0] \neq n) \ fprintf(stderr, "I'm_totally_confused.\n"); \\ \text{ This code is used in section } 1^*.
```

4.* If there are t leaves below node i, we introduce $k = \min(r, t)$ variables Bi+1.j for $1 \le j \le k$. This variable is 1 if (but not only if) at least j of those leaf variables are true. If t > r, we also assert that no r+1 of those variables are true.

```
#define xbar(k) printf("~%s.%d", (g \rightarrow vertices + (k) - n + 1) \rightarrow name, kk)
\langle Generate the clauses for node i \ 4^* \rangle \equiv
     t = count[i], tl = count[i+i+1], tr = count[i+i+2];
     if (t > r + 1) t = r + 1;
     if (tl > r) tl = r;
     if (tr > r) tr = r;
     for (jl = 0; jl \le tl; jl ++)
       for (jr = 0; jr \le tr; jr ++)
          if ((jl + jr \le t) \land (jl + jr) > 0) {
            if (jl) {
               if (i+i+1 \ge n-1) xbar(i+i+1);
               else printf("~B%d.%d", i + i + 2, jl);
            if (jr) {
               printf("_{\sqcup}");
               if (i + i + 2 \ge n - 1) xbar(i + i + 2);
               else printf("~B%d.%d", i + i + 3, jr);
            if (jl + jr \le r) printf("\squareB%d.%d\n", i + 1, jl + jr);
            else printf("\n");
This code is used in section 1*.
```

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5. Finally, we assert that at most r of the x's are true, by implicitly asserting that the (nonexistent) variable B1.r+1 is false.

```
\langle Generate the clauses at the root 5\rangle \equiv
  tl = count[1], tr = count[2];
  if (tl > r) tl = r;
  for (jl = 1; jl \le tl; jl ++) {
     jr = r + 1 - jl;
     if (jr \leq tr) {
        if (1 \ge n - 1) \ xbar(1);
        else printf("~B2.%d", jl);
        printf("
_{\sqcup}");
        if (2 \ge n - 1) \ xbar(2);
       else printf("~B3.%d", jr);
        printf("\n");
  }
This code is used in section 1^*.
```

```
6. \langle Handle the trivial case directly 6 \rangle \equiv
     for (i = 1; i \le n; i++) {
       xbar(n-2+i);
       printf("\n");
  }
```

This code is used in section 1^* .

7* Index.

The following sections were changed by the change file: 1, 2, 4, 7.

 $argc: \underline{1}, 2.$ argv: $\underline{1}^*$, 2^* count: $\underline{1}^*$, 3, 4^* , 5. exit: 2.* $\begin{array}{ll} \textit{fprintf}: & 2 \\ \textit{Graph}: & 1 \\ \end{array} .$ i: 1*
j: 1*
j: 1*
ji: 1* 4*, 5.
jr: 1*, 4*, 5.
k: 1*
k: 1* 2* 4* kk: 1* 2* 4* main: <u>1</u>* n: $\underline{1}$ * name: 4.* $nmax \colon \ \underline{1},^{*} \ 2.^{*}$ printf: 2,* 4,* 5, 6. $r: \underline{1}^*$ $restore_graph$: 2* sscanf: 2* stderr: 2*, 3. t: 1*
tl: 1*, 4*, 5.
tr: 1*, 4*, 5. vertices: 4* $xbar: \underline{4}^*, 5, 6.$

```
 \begin{tabular}{ll} $\langle$ Build the complete binary tree with $n$ leaves $3$ \rangle Used in section $1^*$. \\ $\langle$ Generate the clauses at the root $5$ \rangle$ Used in section $1^*$. \\ $\langle$ Generate the clauses for node $i$ $4^*$ \rangle$ Used in section $1^*$. \\ $\langle$ Handle the trivial case directly $6$ \rangle$ Used in section $1^*$. \\ $\langle$ Process the command line $2^*$ \rangle$ Used in section $1^*$. } \end{tabular}
```

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