$\S1$ SAT-TOMOGRAPHY-2ND INTRO 1

1.* Intro. Given row sum, column sums, and diagonal sums on stdin, this program outputs clauses by which a SAT solver can determine if they are compatible with the existence of an $m \times n$ matrix x_{ij} of zeros and ones.

The row sums are $r_i = \sum_{j=1}^n x_{ij}$, for $1 \le i \le m$. The column sums are $c_j = \sum_{i=1}^m x_{ij}$, for $1 \le j \le n$. And the diagonal sums are $a_d = \sum \{x_{ij} \mid i+j=d+1\}$ and $b_d = \sum \{x_{ij} \mid i-j=d-n\}$, for 0 < d < m+n. They should appear one per line in the input, in a format such as 'r3=20'. Zero sums need not be given. The program deduces m and n from the largest subscripts that appear, and it makes fairly careful syntax checks.

Well actually, the above should be amended: This version works not only with sums of ones, it also uses sums of '11's (that is, consecutive occurrences of 1s in the same line). The second-order sums are given after a comma; for example, 'r3=20,8'.

```
#define mmax 200
                           /* should be at most 255 unless I use bigger radix than hex */
#define nmax 100
                          /* should be at most 255 unless I use bigger radix than hex */
#include <stdio.h>
#include <stdlib.h>
  int r[mmax + 1], c[mmax + 1], a[mmax + nmax], b[mmax + nmax];
                                                                             /* the given data */
  int rr[mmax + 1], cc[mmax + 1], aa[mmax + nmax], bb[mmax + nmax]; /* and its extensions */
  int count[mmax + mmax + nmax + nmax];
                                                /* leaf counts for the BB method */
  char buf[80];
  char name[mmax + nmax][9];
  (Subroutines 10);
  main()
    register int d, i, j, k, l, ll, m, n, nn, t;
    register char *p;
    \langle \text{Input the data } 2^* \rangle;
     \langle \text{ Check the data } 7^* \rangle;
     \langle \text{Output the clauses } 8 \rangle;
  }
```

2 INTRO SAT-TOMOGRAPHY-2ND §2

```
2* \langle \text{Input the data } 2^* \rangle \equiv
  m = n = 0;
  while (1) {
     if (\neg fgets(buf, 80, stdin)) break;
     for (d = 0, p = buf + 1; *p \ge '0' \land *p \le '9'; p++) d = 10 *d + *p - '0';
     if (*p++ \neq '=') {
       fprintf(stderr, "Missing, '=', sign!\nBad, line:, %s", buf);
       exit(-1);
     for (l = 0; *p \ge 0, \wedge p \le 9; p++) l = 10 * l + p - 0;
     if (*p++ \neq ', ') {
       fprintf(stderr, "Missing\_comma!\nBad\_line\_%s", buf);
       exit(-12);
     for (ll = 0; *p \ge 0, \land *p \le 9; p++) ll = 10 * ll + p - 0;
     if (*p \neq '\n') {
       fprintf(stderr, "Missing | \n_character! \nBad_line_%s", buf);
       exit(-2);
     switch (buf[0]) {
       (Cases for row, column, and diagonal sums 3*)
     \mathbf{default}: fprintf(stderr, "\mathtt{Data}_{\mathtt{must}}_{\mathtt{begin}}_{\mathtt{with}}_{\mathtt{lr}}, \mathtt{uc}, \mathtt{ua}, \mathtt{uor}_{\mathtt{b}}! \mathtt{NBad}_{\mathtt{line}}_{\mathtt{ls}}", buf);
       exit(-3);
This code is used in section 1*.
3.* (Cases for row, column, and diagonal sums 3^*) \equiv
case 'r':
  if (d < 1 \lor d > mmax) {
     fprintf(stderr, "Row_index_out_of_range!\nBad_line_%s", buf);
     exit(-4);
  if (l < 0 \lor l > nmax) {
     fprintf(stderr, "Row_data_out_of_range! \nBad_line_%s", buf);
     exit(-5);
  if (d > m) m = d;
  if (r[d]) {
     fprintf(stderr, "The \ value \ of \ r\%d \ has \ already \ been \ given! \ had \ line \ %s", d, buf);
     exit(-6);
  r[d] = l, rr[d] = ll;
  break;
See also sections 4^*, 5^*, and 6^*.
This code is used in section 2*.
```

```
\S 4
```

```
4* (Cases for row, column, and diagonal sums 3^*) +\equiv
case 'c':
 if (d < 1 \lor d > nmax) {
    fprintf(stderr, "Column_index_out_of_range! \nBad_line_%s", buf);
  if (l < 0 \lor l > mmax) {
    fprintf(stderr, "Column_data_out_of_range! \nBad_line_%s", buf);
    exit(-15);
  if (d > n) n = d;
  if (c[d]) {
    fprintf(stderr, \verb"The_uvalue_of_uc%d_has_ualready_been_given! \verb"\nBad_line_ks", d, buf);
    exit(-16);
  c[d] = l, cc[d] = ll;
  break:
5.* (Cases for row, column, and diagonal sums 3*) +\equiv
  if (d < 1 \lor d \ge mmax + nmax) {
    fprintf(stderr, "Diagonal\_index\_out\_of\_range! \nBad\_line\_%s", buf);
    exit(-24);
  if (l < 0 \lor l > mmax \lor l > nmax) {
    fprintf(stderr, "Diagonal_data_out_of_range!\nBad_line_%s", buf);
    exit(-25);
  if (a[d]) {
    fprintf(stderr, "The \ value \ of \ a \ d \ has \ already \ been \ given! \ had \ line \ %s", d, buf);
    exit(-26);
  a[d] = l, aa[d] = ll;
  break;
6* (Cases for row, column, and diagonal sums 3^*) +\equiv
  if (d < 1 \lor d \ge mmax + nmax) {
    fprintf(stderr, "Diagonal_index_out_of_range! \nBad_line_%s", buf);
    exit(-34);
  if (l < 0 \lor l > mmax \lor l > nmax) {
    fprintf(stderr, "Diagonal_data_out_of_range! \nBad_line_%s", buf);
    exit(-35);
  if (b[d]) {
    exit(-36);
  b[d] = l, bb[d] = ll;
  break;
```

```
7* \langle Check the data 7^* \rangle \equiv
  for (i = 1, l = 0; i \le m; i++) l += r[i];
  nn = l;
   {\bf for} \ (j=1,l=0; \ j \le n; \ j +\!\!\!+) \ l \ +\!\!\!+ c[j]; 
  if (l \neq nn) {
     fprintf(stderr, "The_{l}total_{l}of_{l}the_{l}r's_{l}is_{l}%d,_{l}but_{l}the_{l}total_{l}of_{l}the_{l}c's_{l}is_{l}%d! \n", nn, l);
     exit(-40);
  for (d = 1, l = 0; d < m + n; d++) l += a[d];
  if (l \neq nn) {
     fprintf(stderr, "The\_total\_of\_the\_r's\_is\_%d,\_but\_the\_total\_of\_the\_a's\_is\_%d! \n", nn, l);
  for (d = 1, l = 0; d < m + n; d++) l += b[d];
  if (l \neq nn) {
     fprintf(stderr, "The_{\sqcup}total_{\sqcup}of_{\sqcup}the_{\sqcup}r's_{\sqcup}is_{\sqcup}\%d,_{\sqcup}but_{\sqcup}the_{\sqcup}total_{\sqcup}of_{\sqcup}the_{\sqcup}b's_{\sqcup}is_{\sqcup}\%d! \setminus n", nn, l);
     exit(-41);
  fprintf(stderr, "Input_{\square}for_{\square}%d_{\square}rows_{\square}and_{\square}%d_{\square}columns_{\square}successfully_{\square}read", m, n);
  fprintf(stderr, " (total %d) n", nn);
  This code is used in section 1*.
```

8. The variables x_{ij} of the unknown Boolean matrix are denoted by 'ixj'. Auxiliary variables by which we check lower and upper bounds for row sum r_i are denoted by 'iRl'. And similar conventions hold for the column sums and the diagonal sums.

4

9.* We use the methods of Bailleux and Boufkhad (see SAT-THRESHOLD-BB-EQUAL). Indeed, Bailleux and Boufkhad introduced those methods because they wanted to study digital tomography problems.

```
 \begin{split} \langle \text{ Output clauses to check } r_i \text{ 9*} \rangle &\equiv \\ \{ & sprintf(buf, \text{"%dR"}, i); \\ & \textbf{for } (j=1; \ j \leq n; \ j++) \ sprintf(name[j], \text{"%dx%d"}, i, j); \\ & gen\_clauses(n, r[i]); \\ & gen\_clauses1(n-1, rr[i]); \\ \} \end{split}
```

This code is used in section 8.

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```
10. \langle \text{Subroutines } 10 \rangle \equiv
  void gen\_clauses(\mathbf{int} \ n, \mathbf{int} \ r)
     register int i, j, k, jl, jr, t, tl, tr, swap = 0;
     if (r > n - r) swap = 1, r = n - r;
     if (r < 0) {
        \textit{fprintf} \, (\textit{stderr}\,, \texttt{"Negative} \, \texttt{\_parameter} \, \texttt{\_for} \, \texttt{\_case} \, \texttt{\_\%s!} \, \texttt{`n"} \,, \, buf \,);
     if (r \equiv 0) (Handle the trivial case directly 16)
     else {
        \langle Build the complete binary tree with n leaves 11\rangle;
        for (i = n - 2; i; i--) {
           \langle Generate the clauses for node i 12\rangle;
           \langle Generate additional clauses for node i 13\rangle;
        (Generate the clauses at the root 14);
        (Generate additional clauses at the root 15);
This code is used in section 1^*.
11. The tree has 2n-1 nodes, with 0 as the root; the leaves start at node n-1. Nonleaf node k has left
child 2k + 1 and right child 2k + 2. Here we simply fill the count array.
\langle Build the complete binary tree with n leaves 11 \rangle \equiv
  for (k = n + n - 2; k \ge n - 1; k - -) count[k] = 1;
  for (; k \ge 0; k--) count[k] = count[k+k+1] + count[k+k+2];
  if (count[0] \neq n) {
     fprintf(stderr, "I'm_{\sqcup}totally_{\sqcup}confused.\n");
     exit(-666);
```

This code is used in section 10.

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12. If there are t leaves below node i, we introduce $k = \min(r, t)$ auxiliary variables, beginning with the symbolic name in buf and ending with two hex digits of i + 1 and two hex digits of j, for $1 \le j \le k$. This variable will be 1 if and only if at least j of those leaf variables are true. If t > r, we also assert that no r + 1 of those variables are true.

```
#define x(k) printf ("%s%s", swap? "~": "", name[(k) - n + 2])
#define xbar(k) printf ("%s%s", swap? "": "~", name[(k) - n + 2])
\langle Generate the clauses for node i 12\rangle \equiv
    t = count[i], tl = count[i+i+1], tr = count[i+i+2];
    if (t > r + 1) t = r + 1;
    if (tl > r) tl = r;
    if (tr > r) tr = r;
    for (jl = 0; jl \le tl; jl ++)
       for (jr = 0; jr \le tr; jr ++)
         if ((jl + jr \le t) \land (jl + jr) > 0) {
            if (jl) {
              if (i+i+1 \ge n-1) xbar(i+i+1);
              else printf("~\%s\%02x\%02x", buf, i + i + 2, jl);
            if (jr) {
              printf (", ");
              if (i+i+2 \ge n-1) xbar(i+i+2);
              else printf("~\%s\%02x\%02x", buf, i + i + 3, jr);
            if (jl + jr \le r) printf("\\\sum_\%s\%02x\\02x\\n\\, buf, i + 1, jl + jr);
            else printf("\n");
  }
```

This code is used in section 10.

13. So far we've only propagated the effects of the known 1s among the x's. Now we propagate the effects of the 0s: If there are fewer than tl 1s in the leaves of the left subtree and fewer than tr 1s in those of the right subtree, there are fewer than tl + tr - 1 in the leaves of below node i.

```
 \begin{array}{l} \text{ if } (t > r) \ t = r; \\ \text{ for } (jl = 1; \ jl \le tl + 1; \ jl + ) \\ \text{ for } (jr = 1; \ jr \le tr + 1; \ jr + +) \\ \text{ if } (jl + jr \le t + 1) \ \{ \\ \text{ if } (jl \le tl) \ \{ \\ \text{ if } (i+i+1 \ge n-1) \ x(i+i+1); \\ \text{ else } printf("\%s\%02x\%02x", buf, i+i+2, jl); \\ printf("\_"); \\ \} \\ \text{ if } (jr \le tr) \ \{ \\ \text{ /* note that we can't have both } jl > tl \ \text{and } jr > tr \ */ \\ \text{ if } (i+i+2 \ge n-1) \ x(i+i+2); \\ \text{ else } printf("\%s\%02x\%02x", buf, i+i+3, jr); \\ printf("\_"); \\ \} \\ printf("-\%s\%02x\%02x\n", buf, i+1, jl+jr-1); \\ \} \end{array}
```

This code is used in section 10.

Finally, we assert that at most r of the x's are true, by implicitly asserting that the (nonexistent) variable for i = 0 and j = r + 1 is false.

```
\langle Generate the clauses at the root 14\rangle \equiv
  tl = count[1], tr = count[2];
  for (jl = 1; jl \le tl; jl ++) {
     jr = r + 1 - jl;
      \text{if } (jr>0 \land jr \leq tr) \ \{ \\
        if (1 \ge n - 1) \ xbar(1);
        else printf ("~%s02%02x", buf , jl);
        printf("
_{\sqcup}");
        if (2 \ge n - 1) \ xbar(2);
        else printf("~\%s03\%02x", buf, jr);
        printf("\n");
  }
```

This code is used in section 10.

15. To make exactly r of the x's true, we also assert that the (nonexistent) variable for i=1 and j=r is true.

```
\langle Generate additional clauses at the root 15\rangle \equiv
  for (jl = 1; jl \le r; jl ++) {
     jr = r + 1 - jl;
     if (jr > 0) {
       if (jl \leq tl) {
          if (1 \ge n - 1) \ x(1);
          else printf ("%s02%02x", buf, jl);
          printf("_{\sqcup}");
       if (jr \leq tr) {
          if (2 \ge n-1) \ x(2);
          else printf("\%s03\%02x", buf, jr);
       printf("\n");
  }
```

This code is used in section 10.

```
\langle Handle the trivial case directly 16\rangle \equiv
for (i = 1; i \le n; i++) {
  xbar(n-2+i);
  printf("\n");
```

This code is used in section 10.

```
(Output clauses to check c_i 17) \equiv
    sprintf(buf, "%dC", j);
    for (i = 1; i \leq m; i \leftrightarrow) sprintf(name[i], "%dx%d", i, j);
    gen\_clauses(m, c[j]);
This code is used in section 8.
18. (Output clauses to check a_d 18) \equiv
    sprintf(buf, "%dA", d);
    if (m \le n) {
      if (d \leq m) {
         for (i = 1; i \le d; i++) sprintf (name[i], "%dx%d", i, d+1-i);
         gen\_clauses(d, a[d]);
       } else if (d \le n) {
         for (i = 1; i \le m; i++) sprintf (name[i], "%dx%d", i, d+1-i);
         gen\_clauses(m, a[d]);
       } else {
         for (t = 1; t \le m + n - d; t ++) sprintf (name[t], "%dx%d", d + t - n, n + 1 - t);
         gen\_clauses(m+n-d, a[d]);
    \} else \{
       if (d \le n) {
         for (i = 1; i \le d; i++) sprintf (name[i], "%dx%d", i, d+1-i);
         gen\_clauses(d, a[d]);
       } else if (d \le m) {
         for (j = 1; j \le n; j ++) sprintf(name[j], "%dx%d", d + 1 - j, j);
         gen\_clauses(n, a[d]);
       } else {
         for (t = 1; t \le m + n - d; t ++) sprintf (name[t], "%dx%d", d + t - n, n + 1 - t);
         gen\_clauses(m+n-d, a[d]);
    }
  }
```

This code is used in section 8.

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```
19.
      \langle \text{ Output clauses to check } b_d | 19 \rangle \equiv
     sprintf(buf, "%dB", d);
    if (m \le n) {
       if (d \leq m) {
         for (i = 1; i \le d; i ++) sprintf (name[i], "%dx%d", i, n + i - d);
         gen\_clauses(d, b[d]);
       } else if (d \le n) {
         for (i = 1; i \le m; i++) sprintf (name[i], "%dx%d", i, n+i-d);
         gen\_clauses(m,b[d]);
       } else {
         for (j = 1; j \le m + n - d; j ++) sprintf (name[j], "%dx%d", j + d - n, j);
         gen\_clauses(m+n-d, b[d]);
       }
    } else {
       if (d \le n) {
         for (i = 1; i \le d; i++) sprintf (name[i], "%dx%d", i, n+i-d);
         gen\_clauses(d, b[d]);
       } else if (d \le m) {
         for (j = 1; j \le n; j++) sprintf(name[j], "%dx%d", j + d - n, j);
         gen\_clauses(n, b[d]);
       } else {
         for (j = 1; j \le m + n - d; j ++) sprintf(name[j], "%dx%d", j + d - n, j);
         gen\_clauses(m+n-d,b[d]);
    }
  }
```

This code is used in section 8.

20* Index.

```
The following sections were changed by the change file: 1, 2, 3, 4, 5, 6, 7, 9, 20.
```

```
aa: 1^*, 5^*
b: <u>1</u>*
bb: <u>1</u>*, 6*
buf: \underline{1}, 2, 3, 4, 5, 6, 9, 10, 12, 13, 14, 15, 17, 18, 19.
c: \underline{1}*
cc: \underline{1}, 4.
count: \underline{1}^*, 11, 12, 14.
d: 1*
exit: 2,* 3,* 4,* 5,* 6,* 7,* 10, 11.
fgets: 2*
fprintf: 2,* 3,* 4,* 5,* 6,* 7,* 10, 11.
gen_clauses: 9, 10, 17, 18, 19.
gen\_clauses1: 9.*
i: \underline{1}^*, \underline{10}.
j: \ \underline{1}^*, \ \underline{10}.
jl: 10, 12, 13, 14, 15.
jr: 10, 12, 13, 14, 15.
k: \quad \underline{1}, \quad \underline{10}.
l: 1* ll: 1, 2, 3, 4, 5, 6.
m: \underline{1}^*
main: \underline{1}^*
mmax: 1, 3, 4, 5, 6.
n: \ \underline{1}^*, \ \underline{10}.
name: \underline{1}, 9, 12, 17, 18, 19.
nmax: 1, 3, 4, 5, 6.
nn: \underline{1}^*, 7^*
p: <u>1</u>*
printf: 7, 12, 13, 14, 15, 16.
r: \ \underline{1}^*, \ \underline{10}.
rr: 1*, 3*, 9*
sprintf: 9*, 17, 18, 19.
stderr: 2, 3, 4, 5, 6, 7, 10, 11.
stdin: 1,* 2.*
swap: 10, 12.
t: \underline{1}^*, \underline{10}.
tl: 10, 12, 13, 14, 15.
tr: 10, 12, 13, 14, 15.
x: 12.
```

xbar: 12, 14, 16.

SAT-TOMOGRAPHY-2ND NAMES OF THE SECTIONS 11

```
\langle Build the complete binary tree with n leaves 11 \rangle Used in section 10.
(Cases for row, column, and diagonal sums 3^*, 4^*, 5^*, 6^*) Used in section 2^*.
 Check the data 7^* Used in section 1^*.
 Generate additional clauses at the root 15 \rangle Used in section 10.
 Generate additional clauses for node i 13 \rangle Used in section 10.
 Generate the clauses at the root 14 \rangle Used in section 10.
 Generate the clauses for node i 12 \rangle Used in section 10.
 Handle the trivial case directly 16 \ Used in section 10.
 Input the data 2* \rangle Used in section 1*.
 Output clauses to check a_d 18 \rangle Used in section 8.
 Output clauses to check b_d 19 \rangle
                                       Used in section 8.
 Output clauses to check c_j 17 \rangle Used in section 8.
 Output clauses to check r_i 9* Used in section 8.
 Output the clauses 8 \rangle Used in section 1^*.
\langle Subroutines 10\rangle Used in section 1*.
```

SAT-TOMOGRAPHY-2ND

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