$\S1$  SAT-POSET-NOMAX INTRO 1

1. Intro. This quickie outputs clauses that are unsatisfiable because they state that there exists a partial ordering on m elements in which no element is maximal. (All backtrack proofs of this fact are known to require  $\Omega(2^m)$  steps.)

```
Variable j \cdot k means that j \prec k in the partial ordering.
#include <stdio.h>
#include <stdlib.h>
  int m;
  main(\mathbf{int} \ argc, \mathbf{char} * argv[])
     register i, j, k;
     \langle \text{Process the command line } 2 \rangle;
     (Generate the clauses for irreflexivity 3);
     ⟨ Generate the clauses for transitivite 4⟩;
     \langle Generate the clauses for nonmaximality 5\rangle;
  }
2. \langle \text{Process the command line } 2 \rangle \equiv
  if (argc \neq 2 \lor sscanf(argv[1], "%d", \&m) \neq 1) {
     fprintf(stderr, "Usage: \_\%s\_m\n", argv[0]);
     exit(-1);
  }
  printf("\"alpha",m);
This code is used in section 1.
3. \langle Generate the clauses for irreflexivity 3\rangle \equiv
  \textbf{for } (j=1; \ j \leq m; \ j +\!\!\!\!+\!\!\!\!+) \ \mathit{printf}(\texttt{"~\%d.\%d\n"}, j, j);
This code is used in section 1.
4. \langle Generate the clauses for transitivite 4\rangle \equiv
  for (i = 1; i \le m; i++)
     for (j = 1; j \le m; j++)
       if (i \neq j) {
          for (k = 1; k \le m; k++)
            if (j \neq k) {
               printf(\verb"""%d.%d."%d.%d.%d.",i,j,j,k,i,k);
       }
This code is used in section 1.
5. \langle Generate the clauses for nonmaximality 5\rangle \equiv
  for (j = 1; j \le m; j ++) {
     printf("\n");
  }
```

This code is used in section 1.

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## 6. Index.

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