$\S1$ LADDERS INTRODUCTION 1

Important: Before reading LADDERS, please read or at least skim the programs for GB_WORDS and GB_DIJK.

1. Introduction. This demonstration program uses graphs constructed by the GB_WORDS module to produce an interactive program called ladders, which finds shortest paths between two given five-letter words of English.

The program assumes that UNIX conventions are being used. Some code in sections listed under 'UNIX dependencies' in the index might need to change if this program is ported to other operating systems.

To run the program under UNIX, say 'ladders (options)', where (options) consists of zero or more of the following specifications in any order:

- -v Verbosely print all words encountered during the shortest-path computation, showing also their distances from the goal word.
- -a Use alphabetic distance instead of considering adjacent words to be one unit apart; for example, the alphabetic distance from 'words' to 'woods' is 3, because 'r' is three places from 'o' in the alphabet.
- -f Use distance based on frequency (see below), instead of considering adjacent words to be one unit apart. This option is ignored if either -a or -r has been specified.
- -h Use a lower-bound heuristic to shorten the search (see below). This option is ignored if option -f has been selected.
- -e Echo the input to the output (useful if input comes from a file instead of from the terminal).
- $-n\langle \text{number}\rangle$ Limit the graph to the *n* most common English words, where *n* is the given $\langle \text{number}\rangle$.
- $-r\langle number \rangle$ Limit the graph to $\langle number \rangle$ randomly selected words. This option is incompatible with -n.
- $-s\langle \text{number}\rangle$ Use $\langle \text{number}\rangle$ instead of 0 as the seed for random numbers, to get different random samples or to explore words of equal frequency in a different order.

Option -f assigns a cost of 0 to the most common words and a cost of 16 to the least common words; a cost between 0 and 16 is assigned to words of intermediate frequency. The word ladders that are found will then have minimum total cost by this criterion. Experience shows that the -f option tends to give the "friendliest," most intuitively appealing ladders.

Option $\neg h$ attempts to focus the search by giving priority to words that are near the goal. (More precisely, it modifies distances between adjacent words by using a heuristic function hh(v), which would be the shortest possible distance between v and the goal if every five-letter combination happened to be an English word.) The GB_DIJK module explains more about such heuristics; this option is most interesting to watch when used in conjunction with $\neg v$.

2. The program will prompt you for a starting word. If you simply type $\langle \text{return} \rangle$, it exits; otherwise you should enter a five-letter word (with no uppercase letters) before typing $\langle \text{return} \rangle$.

Then the program will prompt you for a goal word. If you simply type (return) at this point, it will go back and ask for a new starting word; otherwise you should specify another five-letter word.

Then the program will find and display an optimal word ladder from the start to the goal, if there is a path from one to the other that changes only one letter at a time.

And then you have a chance to start all over again, with another starting word.

The start and goal words need not be present in the program's graph of "known" words. They are temporarily added to that graph, but removed again whenever new start and goal words are given. (Thus you can go from sturm to drang even though those words aren't English.) If the -f option is being used, the cost of the goal word will be 20 when it is not in the program's dictionary.

2 INTRODUCTION LADDERS §3

3. Here is the general layout of this program, as seen by the C compiler:

```
#include <ctype.h>
                            /* system file for character types */
#include "gb_graph.h"
                               /* the standard GraphBase data structures */
#include "gb_words.h"
                               /* routines for five-letter word graphs */
#include "gb_dijk.h"
                              /* routines for shortest paths */
  \langle Preprocessor definitions \rangle
  \langle Global variables 4\rangle
  (Subroutines 11)
  main(argc, argv)
                     /* the number of command-line arguments */
       int argc;
       \mathbf{char} * argv[];
                          /* an array of strings containing those arguments */
     \langle Scan \text{ the command-line options 5} \rangle;
    \langle Set up the graph of words 6\rangle;
    while (1) {
       (Prompt for starting word and goal word; break if none given 26);
       (Find a minimal ladder from start to goal, if one exists, and print it 13);
                    /* normal exit */
    return 0;
```

§4 LADDERS

4. Parsing the options. Let's get the UNIX command-line junk out of the way first, so that we can concentrate on meatier stuff. Our job in this part of the program is to see if the default value zero of external variable *verbose* should change, and/or if the default values of any of the following internal variables should change:

```
\langle \text{Global variables 4} \rangle \equiv
                      /* nonzero if the alphabetic distance option is selected */
  char alph = 0;
                       /* nonzero if the frequency-based distance option is selected */
  char freq = 0;
  char heur = 0;
                       /* nonzero if the heuristic search option is selected */
  char echo = 0;
                       /* nonzero if the input-echo option is selected */
                              /* maximum number of words in the graph (0 means infinity) */
  unsigned long n = 0;
                         /* nonzero if we will ignore the weight of words */
  char randm = 0;
                       /* seed for random number generator */
  long seed = 0;
See also sections 7, 12, and 23.
This code is used in section 3.
5. \langle Scan \text{ the command-line options 5} \rangle \equiv
  while (--argc) {
    if (strcmp(argv[argc], "-v") \equiv 0) verbose = 1;
    else if (strcmp(argv[argc], "-a") \equiv 0) alph = 1;
    else if (strcmp(argv[argc], "-f") \equiv 0) freq = 1;
    else if (strcmp(argv[argc], "-h") \equiv 0) heur = 1;
    else if (strcmp(argv[argc], "-e") \equiv 0) echo = 1;
    else if (sscanf(argv[argc], "-n\%lu", \&n) \equiv 1) \ randm = 0;
    else if (sscanf(argv[argc], "-r\%lu", \&n) \equiv 1) \ randm = 1;
    else if (sscanf(argv[argc], "-s\%ld", \&seed) \equiv 1);
       fprintf(stderr, "Usage: \_%s_{\_}[-v][-a][-f][-h][-e][-nN][-rN][-sN] \n", argv[0]);
       return -2;
  if (alph \lor randm) freq = 0;
  if (freq) heur = 0;
```

This code is used in section 3.

4 CREATING THE GRAPH LADDERS §6

6. Creating the graph. The GraphBase words procedure will produce the five-letter words we want, organized in a graph structure.

```
#define quit_if(x,c)
          if (x) {
            fprintf(stderr, "Sorry, _|I_| couldn't_| build_| a_| dictionary_| (trouble_| code_| %|d)! \n", c);
\langle Set up the graph of words 6 \rangle \equiv
  g = words(n, (randm ? zero\_vector : \Lambda), 0_L, seed);
  quit_if(g \equiv \Lambda, panic_code);
  \langle Confirm the options selected 8 \rangle;
  \langle Modify the edge lengths, if the alph or freq option was selected 9\rangle;
  (Modify the priority queue algorithm, if unequal edge lengths are possible 10);
This code is used in section 3.
7. \langle \text{Global variables } 4 \rangle + \equiv
  Graph *g;
                    /* graph created by words */
  long zero\_vector[9];
                             /* weights to use when ignoring all frequency information */
   The actual number of words might be decreased to the size of the GraphBase dictionary, so we wait
until the graph is generated before confirming the user-selected options.
```

```
 \begin{split} &\langle \operatorname{Confirm} \ \operatorname{the} \ \operatorname{options} \ \operatorname{selected} \ 8 \rangle \equiv \\ & \quad \operatorname{if} \ (\mathit{verbose}) \ \{ \\ & \quad \operatorname{if} \ (\mathit{alph}) \ \mathit{printf} \left( \text{"(alphabetic\_distance\_selected)} \right); \\ & \quad \operatorname{if} \ (\mathit{freq}) \ \mathit{printf} \left( \text{"(frequency-based\_distances\_selected)} \right); \\ & \quad \operatorname{if} \ (\mathit{heur}) \ \mathit{printf} \left( \text{"(lowerbound\_heuristic\_will\_be\_used\_to\_focus\_the\_search)} \right); \\ & \quad \operatorname{if} \ (\mathit{randm}) \ \mathit{printf} \left( \text{"(random\_selection\_of\_%ld\_words\_with\_seed\_%ld)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\%ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\$ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\$ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\$ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\$ld\_words)} \right), \\ & \quad \operatorname{else} \ \mathit{printf} \left( \text{"(the\_graph\_has}\_\$ld\_words)}
```

This code is used in section 6.

This code is used in section 6.

9. The edges in a *words* graph normally have length 1, so we must change them if the user has selected *alph* or *freq*. The character position in which adjacent words differ is recorded in the *loc* field of each arc. The frequency of a word is stored in the *weight* field of its vertex.

```
#define a\_dist(k) (*(p+k) < *(q+k)? *(q+k) - *(p+k) : *(p+k) - *(q+k)) \langle Modify the edge lengths, if the alph or freq option was selected 9\rangle \equiv if (alph) \{ register Vertex *u; for (u=g\neg vertices+g\neg n-1;\ u\geq g\neg vertices;\ u--) \{ register Arc *a; register char *p=u\neg name; for (a=u\neg arcs;\ a;\ a=a\neg next) \{ register char *q=a\neg tip\neg name; a\neg len=a\_dist(a\neg loc); \} \} else if (freq) \{ register Vertex *u; for (u=g\neg vertices+g\neg n-1;\ u\geq g\neg vertices;\ u--) \{ register Arc *a; for (a=u\neg arcs;\ a;\ a=a\neg next) a\neg len=freq\_cost(a\neg tip); \} \}
```

 $\S10$ Ladders creating the graph 5

10. The default priority queue algorithm of *dijkstra* is quite efficient when all edge lengths are 1. Otherwise we change it to the alternative method that works best for edge lengths less than 128.

```
 \langle \mbox{ Modify the priority queue algorithm, if unequal edge lengths are possible 10} \rangle \equiv \mbox{if } (alph \lor freq \lor heur) \ \{ \\ init\_queue = init\_128 \, ; \\ del\_min = del\_128 \, ; \\ enqueue = enq\_128 \, ; \\ requeue = req\_128 \, ; \\ \}
```

This code is used in section 6.

11. The frequency has been computed with the default weights explained in the documentation of words; it is usually less than 2^{16} . A word whose frequency is 0 costs 16; a word whose frequency is 1 costs 15; a word whose frequency is 2 or 3 costs 14; and the costs keep decreasing by 1 as the frequency doubles, until we reach a cost of 0.

```
⟨Subroutines 11⟩ ≡ long freq\_cost(v) Vertex *v; { register long acc = v \neg weight; /* the frequency, to be shifted right */register long k = 16; while (acc) k - -, acc \gg = 1; return (k < 0? 0: k); } See also sections 17, 18, 20, 22, and 27. This code is used in section 3.
```

6 minimal ladders ladders 12

12. Minimal ladders. The guts of this program is a routine to compute shortest paths between two given words, *start* and *goal*.

The *dijkstra* procedure does this, in any graph with nonnegative arc lengths. The only complication we need to deal with here is that *start* and *goal* might not themselves be present in the graph. In that case we want to insert them, albeit temporarily.

The conventions of GB_GRAPH allow us to do the desired augmentation by creating a new graph gg whose vertices are borrowed from g. The graph g has space for two more vertices (actually for four), and any new memory blocks allocated for the additional arcs present in gg will be freed later by the operation $gb_recycle(gg)$ without confusion.

```
\langle \text{Global variables 4} \rangle + \equiv
                      /* clone of g with possible additional words */
  Graph *gg;
                                    /* words dear to the user's heart, plus '\0' */
  char start[6], goal[6];
                                 /* start and goal vertices in gg */
   Vertex *uu, *vv;
13. (Find a minimal ladder from start to goal, if one exists, and print it 13) \equiv
   \langle Build the amplified graph gg 14\rangle;
    Let dijkstra do the hard work 21 \rangle;
   \langle \text{ Print the answer 24} \rangle;
   \langle \text{ Remove all traces of } gg 25 \rangle;
This code is used in section 3.
14. \langle Build the amplified graph gg 14\rangle \equiv
  qq = qb_new_qraph(0_L);
   quit\_if(gg \equiv \Lambda, no\_room + 5);
                                            /* out of memory */
  gg \neg vertices = g \neg vertices;
   gg \rightarrow n = g \rightarrow n;
   \langle \text{ Put the } start \text{ word into } gg, \text{ and let } uu \text{ point to it } 15 \rangle;
   \langle \text{ Put the } goal \text{ word into } gg, \text{ and let } vv \text{ point to it } 16 \rangle;
  if (gg \neg n \equiv g \neg n + 2) (Check if start is adjacent to goal 19);
   quit\_if(gb\_trouble\_code, no\_room + 6);
                                                      /* out of memory */
This code is used in section 13.
```

15. The $find_word$ procedure returns Λ if it can't find the given word in the graph just constructed by words. In that case it has applied its second argument to every adjacent word. Hence the program logic here does everything needed to add a new vertex to gg when necessary.

```
\langle \text{ Put the } start \text{ word into } gg, \text{ and let } uu \text{ point to it } 15 \rangle \equiv
                                                    /* a tentative new vertex */
  (gg \rightarrow vertices + gg \rightarrow n) \rightarrow name = start;
  uu = find\_word(start, plant\_new\_edge);
  if (\neg uu) uu = gg \neg vertices + gg \neg n ++;
                                                       /* recognize the new vertex and refer to it */
This code is used in section 14.
16. (Put the goal word into gg, and let vv point to it 16) \equiv
  if (strncmp(start, qoal, 5) \equiv 0) vv = uu;
                                                          /* avoid inserting a word twice */
  else {
                                                      /* a tentative new vertex */
     (gg \neg vertices + gg \neg n) \neg name = goal;
     vv = find\_word(goal, plant\_new\_edge);
     if (\neg vv) vv = gg \neg vertices + gg \neg n ++;
                                                         /* recognize the new vertex and refer to it */
  }
This code is used in section 14.
```

§17 LADDERS MINIMAL LADDERS 7

17. The $alph_dist$ subroutine calculates the alphabetic distance between arbitrary five-letter words, whether they are adjacent or not.

```
\langle Subroutines 11\rangle + \equiv
  long alph\_dist(p,q)
        register char *p, *q;
      return a_{-}dist(0) + a_{-}dist(1) + a_{-}dist(2) + a_{-}dist(3) + a_{-}dist(4);
  }
18. \langle \text{Subroutines } 11 \rangle + \equiv
  void plant_new_edge(v)
        Vertex *v;
   { Vertex *u = gg \neg vertices + gg \neg n;
                                                      /* the new edge runs from u to v */
      gb\_new\_edge(u, v, 1_L);
                                      /* we have u > v, hence v \rightarrow arcs = u \rightarrow arcs - 1 */
      if (alph) u \rightarrow arcs \rightarrow len = (u \rightarrow arcs - 1) \rightarrow len = alph\_dist(u \rightarrow name, v \rightarrow name);
      else if (freq) {
        u-arcs-len = freq\_cost(v); /* adjust the arc length from u to v */
        (u \rightarrow arcs - 1) \rightarrow len = 20; /* adjust the arc length from v to u */
  }
```

19. There's a bug in the above logic that could be embarrassing, although it will come up only when a user is trying to be clever: The $find_word$ routine knows only the words of g, so it will fail to make any direct connection between start and goal if they happen to be adjacent to each other yet not in the original graph. We had better fix this, otherwise the computer will look stupid.

```
 \begin{array}{l} \langle \, {\rm Check} \, \, {\rm if} \, \, start \, \, {\rm is} \, \, {\rm adjacent} \, \, {\rm to} \, \, goal \, \, \, 19 \, \rangle \equiv \\ {\rm if} \, \, \left( hamm\_dist(start,goal) \equiv 1 \right) \, \, \left\{ \\ gg \neg n - : \, \quad / * \, \, {\rm temporarily} \, \, {\rm pretend} \, \, vv \, \, {\rm hasn't} \, \, {\rm been} \, \, {\rm added} \, \, {\rm yet} \, \, \, * / \\ plant\_new\_edge(uu); \quad / * \, \, {\rm make} \, \, vv \, \, {\rm adjacent} \, \, {\rm to} \, \, uu \, \, * / \\ gg \neg n + + ; \quad / * \, \, {\rm and} \, \, {\rm recognize} \, \, {\rm it} \, \, {\rm again} \, \, * / \\ \end{array} \right\}
```

20. The Hamming distance between words is the number of character positions in which they differ.

```
#define h\_dist(k) (*(p+k) \equiv *(q+k)? 0:1)

\langle Subroutines 11 \rangle +\equiv

long hamm\_dist(p,q)

register char *p, *q;

\{

return h\_dist(0) + h\_dist(1) + h\_dist(2) + h\_dist(3) + h\_dist(4);

\}
```

21. OK, now we've got a graph in which *dijkstra* can operate.

```
\langle \text{Let } \textit{dijkstra} \text{ do the hard work } 21 \rangle \equiv  if (\neg \textit{heur}) \ \textit{min\_dist} = \textit{dijkstra}(uu, vv, gg, \Lambda); else if (\textit{alph}) \ \textit{min\_dist} = \textit{dijkstra}(uu, vv, gg, \textit{alph\_heur}); else \textit{min\_dist} = \textit{dijkstra}(uu, vv, gg, \textit{hamm\_heur}); This code is used in section 13.
```

This code is used in section 14.

8 MINIMAL LADDERS \$22

22. \langle Subroutines 11 $\rangle + \equiv$

 $gb_recycle(gg);$ /* th This code is used in section 13.

```
long alph\_heur(v)
        Vertex *v;
  { return \ alph\_dist(v \neg name, goal); }
  long hamm\_heur(v)
        Vertex *v;
  { return hamm\_dist(v \rightarrow name, goal); }
23. \langle Global variables 4 \rangle + \equiv
                         /* length of the shortest ladder */
  long min\_dist;
24. \langle \text{ Print the answer 24} \rangle \equiv
  if (min\_dist < 0) printf("Sorry, there's inoladder from % sito % s. \n", start, goal);
  else print\_dijkstra\_result(vv);
This code is used in section 13.
      Finally, we have to clean up our tracks. It's easy to remove all arcs from the new vertices of gg to the
old vertices of g; it's a bit trickier to remove the arcs from old to new. The loop here will also remove arcs
properly between start and goal vertices, if they both belong to gg not g.
\langle Remove all traces of gg 25 \rangle \equiv
  for (uu = g \neg vertices + gg \neg n - 1; uu \ge g \neg vertices + g \neg n; uu - -) { register Arc *a;
     for (a = uu \neg arcs; a; a = a \neg next) {
                        /* now vv \rightarrow arcs \equiv a - 1, since arcs for edges come in pairs */
        vv = a \rightarrow tip;
        vv \neg arcs = vv \neg arcs \neg next;
                        /* we needn't clear uu¬name */
     uu \rightarrow arcs = \Lambda;
```

/* the $gg \rightarrow data$ blocks disappear, but $g \rightarrow data$ remains */

§26 LADDERS

26. Terminal interaction. We've finished doing all the interesting things. Only one minor part of the program still remains to be written.

```
\langle Prompt for starting word and goal word; break if none given 26\rangle \equiv
  putchar('\n');
                     /* make a blank line for visual punctuation */
            /* if we try to avoid this label, the break command will be broken */
  if (prompt\_for\_five("Starting", start) \neq 0) break;
  if (prompt\_for\_five("uuuuGoal", goal) \neq 0) goto restart;
This code is used in section 3.
27. \langle Subroutines 11\rangle + \equiv
  long prompt\_for\_five(s, p)
                   /* string used in prompt message */
      char *s;
      register char *p; /* where to put a string typed by the user */
  { register char *q;
                           /* current position to store characters */
                         /* current character of input */
    register long c;
    while (1) {
       printf("%s\_word:_{\sqcup}", s);
       fflush(stdout); /* make sure the user sees the prompt */
       q = p;
       while (1) {
         c = getchar();
         if (c \equiv EOF) return -1; /* end-of-file */
         if (echo) putchar(c);
         if (c \equiv '\n') break;
         if (\neg islower(c)) q = p + 5;
         else if (q 
      if (q \equiv p + 5) return 0; /* got a good five-letter word */
                              /* got just \(\text{return}\) */
      if (q \equiv p) return 1;
       printf("(Please_type_five_lowercase_letters_and_RETURN.)\n");
  }
```

10 INDEX LADDERS §28

28. Index. Finally, here's a list that shows where the identifiers of this program are defined and used.

 $a: \ \ \underline{9}, \ \underline{25}.$ $a_dist: \underline{9}, 17.$ $acc: \underline{11}.$ alph: $\underline{4}$, 5, 8, 9, 10, 18, 21. $alph_dist: 17, 18, 22.$ $alph_heur$: 21, 22. Arc: 9, 25. arcs: 9, 18, 25. $argc: \underline{3}, 5.$ $argv: \underline{3}, 5.$ c: 27. data: 25. $del_min: 10.$ $del_{-}128: 10.$ dijkstra: 10, 12, 21. $echo\colon \ \underline{4},\ 5,\ 27.$ $enq_{-}128: 10.$ enqueue: 10. EOF: 27. fflush: 27.find_word: 15, 16, 19. fprintf: 5, 6.freq: $\underline{4}$, 5, 8, 9, 10, 18. $freq_cost$: 9, $\underline{11}$, 18. $g: \underline{7}$. gb_new_edge : 18. gb_new_graph : 14. $gb_recycle$: 12, 25. $gb_trouble_code$: 14. getchar: 27.gg: 12, 14, 15, 16, 18, 19, 21, 25. $goal \colon \ \underline{12}, \ 16, \ 19, \ 22, \ 24, \ 26.$ **Graph**: 7, 12. $h_{-}dist$: 20. $hamm_{-}dist: 19, 20, 22.$ $hamm_heur$: 21, 22. Hamming, Richard Wesley, distance: heur: $\underline{4}$, 5, 8, 10, 21. $init_queue\colon \ \ 10.$ $init_{-}128: 10.$ is lower: 27.k: 11. len: 9, 18. loc: 9. $main: \underline{3}.$ $min_dist\colon \ \ 21,\ \underline{23},\ 24.$ $n: \underline{4}.$ name: 9, 15, 16, 18, 22, 25. next: 9, 25. no_room : 14. $p: \quad \underline{9}, \ \underline{17}, \ \underline{20}, \ \underline{27}.$

 $panic_code$: 6. plant_new_edge: 15, 16, <u>18</u>, 19. $print_dijkstra_result$: 24. printf: 8, 24, 27. $prompt_for_five: 26, \underline{27}.$ putchar: 26, 27. q: 9, 17, 20, 27. $\mathit{quit_if} : \quad \underline{6}, \ 14.$ $randm: \underline{4}, 5, 6, 8.$ $req_{-}128: 10.$ requeue: 10. restart: 26.s: $\underline{27}$. seed: $\underline{4}$, 5, 6, 8. sscanf: 5. $start: \ \underline{12}, \ 15, \ 16, \ 19, \ 24, \ 26.$ stderr: 5, 6. $stdout\colon 27.$ strcmp: 5. strncmp: 16. tip: 9, 25. $u: \ \underline{9}, \ \underline{18}.$ UNIX dependencies: 3, 5. uu: 12, 15, 16, 19, 21, 25.v: 11, 18, 22.verbose: 4, 5, 8.Vertex: 9, 11, 12, 18, 22. $vertices \colon \ \ 9, \ 14, \ 15, \ 16, \ 18, \ 25.$ vv: 12, 16, 19, 21, 24, 25.weight: 9, 11.words: 6, 7, 9, 11, 15. $zero_vector$: 6, $\underline{7}$.

LADDERS NAMES OF THE SECTIONS 11

```
\langle Build the amplified graph gg 14\rangle Used in section 13.
(Check if start is adjacent to goal 19) Used in section 14.
 Confirm the options selected 8 \rangle Used in section 6.
 Find a minimal ladder from start to goal, if one exists, and print it 13 \rangle Used in section 3.
 Global variables 4, 7, 12, 23 \ Used in section 3.
 Let dijkstra do the hard work 21 \rangle Used in section 13.
 Modify the edge lengths, if the alph or freq option was selected 9 \rangle Used in section 6.
 Modify the priority queue algorithm, if unequal edge lengths are possible 10 \( \) Used in section 6.
 Print the answer 24 \rangle Used in section 13.
 Prompt for starting word and goal word; break if none given 26 \> Used in section 3.
 Put the goal word into gg, and let vv point to it 16 \rangle Used in section 14.
 Put the start word into gg, and let uu point to it 15 \rangle Used in section 14.
 Remove all traces of gg 25 \ Used in section 13.
 Scan the command-line options 5 Used in section 3.
(Set up the graph of words 6) Used in section 3.
(Subroutines 11, 17, 18, 20, 22, 27) Used in section 3.
```

LADDERS

| | Section | ι Pag ϵ |
|----------------------|---------|------------------|
| Introduction | 1 | . 1 |
| Parsing the options | 4 | . 3 |
| Creating the graph | 6 | 4 |
| Minimal ladders | 12 | 6 |
| Terminal interaction | 26 | 9 |
| Index | 28 | 10 |

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