§1 SAT-MUTEX INTRO 1

1. Intro. This is a quick-and-dirty way to go from a slightly symbolic description of proposed mutual-exclusion algorithms to a corresponding set of clauses, so that I can use the clauses for bounded model checking.

In other words, I want to see whether the given concurrent algorithms can violate the mutex property by permitting simultaneous execution of two critical sections, or whether they can lead to livelock or starvation, in a given number of steps. To test this, I'll see if certain extensions of the clauses are satisfiable.

First I have to describe the input language. Each step/state of an algorithm is given a name, which begins with an uppercase letter and has at most four characters. Every shared variable is also given a number, which begins with a lowercase letter and has at most two characters.

Only four elementary kinds of primitive operations are permitted at each step:

- 1) Compute non-critically, then optionally go to step l. (Here l is a step name.)
- 2) Compute critically, then go to step l. (Likewise.)
- 3) Set $V \leftarrow v$, then goto l. (Here V is a shared variable and v is a constant.)
- 4) If V = v, goto l, else goto l'. (Likewise.)

These steps specify state transitions in an fairly obvious way; precise semantics will be explained later. Here's a simple example of possible input:

~ separate locks
A0 maybe goto A1
A1 a=1 goto A2
A2 if b=1 goto A2 else A3
A3 critical goto A4
A4 a=0 goto A0
B0 maybe goto B1
B1 b=1 goto B2
B2 if a=1 goto B2 else B3
B3 critical goto B4
B4 b=0 goto B0

The first line, which begins with '~', is simply a comment that will be passed to the output file. It is followed by steps of types 1, 3, 4, 2, 3, 1, 3, 4, 2, 3, respectively. The shared variables are a and b. The concurrent occurrence of critical states should never occur.

(I do not claim that these programs solve the mutex problem; they simply provide an example.)

At present I assume that all step names begin with either A or B, and that all shared variables are Boolean. But those restrictions might well be lifted later, after I get some experience with this simpler scheme.

2 INTRO SAT-MUTEX §2

Here then is the basic outline of this program. #define massteps 100 /* at most this many steps */ #define bufsize 1024 /* must exceed the length of the longest input line */ #include <stdio.h> #include <stdlib.h> #include <string.h> \langle Type definitions $4\rangle$; **step** state[maxsteps]; /* internal representation of the programs */ char vars[maxsteps][2];/* the distinct shared-variable names */ int astep[maxsteps], bstep[maxsteps]; /* steps for processes A and B */ /* command-line parameter, the number of time steps to emulate */ **char** buf [bufsize]; /* input from stdin goes here */ $main(\mathbf{int} \ argc, \mathbf{char} * argv[])$ register int i, j, k, m, n, t, ma, mb; $\langle \text{ Process the command line } 3 \rangle$: $\langle \text{ Parse the input into the } state \text{ table } 5 \rangle;$ \langle Generate the initial clauses 16 \rangle ; for (t = 0; t < r; t++) (Generate the transitions from time t to time t + 1 17); \langle Generate clauses to force concurrent critical sections at time r 25 \rangle ; } 3. $\langle \text{Process the command line } 3 \rangle \equiv$ if $(argc \neq 2 \lor sscanf(argv[1], "%d", \&r) \neq 1)$ { exit(-1);**if** $(r \le 0)$ { $fprintf(stderr, "Parameter _r _must _be _positive! \n");$ $printf("~_{\sqcup}sat-mutex_{\sqcup}%d\n",r);$ This code is used in section 2. 4. Every non-comment line of input is recorded in an abbreviated form. $\langle \text{Type definitions 4} \rangle \equiv$ typedef struct state_struct { **char** name[4], lab[4], elab[4]; /* the name of this step and its successors */ char var[2]; /* the shared variable */ /* its value */ char val; char crit; /* is this a critical step? */ } step; This code is used in section 2.

 $\S5$ SAT-MUTEX INTRO 3

5. I don't attempt to provide much syntactic sugar for the user (since I expect to be the only user). If I need something fancier, I'll probably write a preprocessor to convert fancy output into the primitive form that is understood by this program.

```
\langle \text{ Parse the input into the } state \text{ table } 5 \rangle \equiv
  for (m = n = ma = mb = 0; ; ) {
     if (\neg fgets(buf, bufsize, stdin)) break;
     if (buf[0] \equiv "") printf("%s", buf);
     else {
        \mathbf{char} * curp = buf;
        if (m \ge maxsteps) {
           fprintf(stderr, "Recompile_lme_l--_lI_lonly_lhave_lroom_lfor_l%d_lsteps! \n", massteps);
        \langle Scan the name field 6 \rangle;
        if (strncmp(curp, "maybe_{\perp}", 6) \equiv 0) \langle Scan \ a \ maybe \ step \ 7 \rangle
        else if (strncmp(curp, "critical_{\sqcup}", 9) \equiv 0) \langle Scan \ a \ critical \ step \ 8 \rangle
        else if (strncmp(curp, "if_{\sqcup}", 3) \equiv 0) \langle Scan \text{ an if step } 10 \rangle
        else \langle Scan \text{ an assignment step } 14 \rangle;
        m++;
  }
   (Check for missing steps 15);
  if (state[astep[0]].crit + state[bstep[0]].crit > 1) {
     fprintf(stderr, "Both_processes_are_initially_in_critical_sections! \n");
     exit(-555);
  fprintf(stderr, "(\d+\d_steps_with_\d_shared_variables_successfully\_input)\n", ma, mb, n);
This code is used in section 2.
     #define abrt(m,t)
           { fprintf(stderr, "Oops, | \s! \n> | \s|, m, buf); exit(t); }
\langle \text{Scan the } name \text{ field } 6 \rangle \equiv
  for (j = 0; *curp \land *curp \neq ` \Box ` \land *curp \neq ` \" ; j++, curp ++)
     if (j < 4) state [m].name [j] = *curp;
  \textbf{if } (j>4) \ abrt("\texttt{the}\_\texttt{name}\_\texttt{is}\_\texttt{too}\_\texttt{long"}, -10);\\
  if (state[m].name[0] < `A` \lor state[m].name[0] > `B`)
     abrt("the_{\sqcup}step_{\sqcup}name_{\sqcup}must_{\sqcup}begin_{\sqcup}with_{\sqcup}A_{\sqcup}or_{\sqcup}B", -11);
  for (j = 0; j < m; j ++)
     if (strncmp(state[j].name, state[m].name, 4) \equiv 0) abrt("that_name_nhas_nalready_nbeen_nused", -12);
  if (state[m].name[0] \equiv 'A') astep[ma \leftrightarrow] = m;
  else bstep[mb++] = m;
  if (*curp \leftrightarrow \neq ' \cup ') abrt("step_{\sqcup}is_{\sqcup}incomplete", -13);
This code is used in section 5.
     \langle Scan \ a \ maybe \ step \ 7 \rangle \equiv
  {
      curp += 5;
      \langle \text{Scan the } lab \text{ field } 9 \rangle;
     if (*curp \neq '\n') abrt("maybe_lstep_lends_lbadly", -14);
This code is used in section 5.
```

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```
\langle Scan \ a \ critical \ step \ 8 \rangle \equiv
      curp += 8;
      state[m].crit = 1;
      \langle Scan the lab field 9 \rangle;
      if (*curp \neq '\n') abrt("critical_step_ends_badly", -15);
This code is used in section 5.
9. \langle Scan the lab field 9 \rangle \equiv
   if (strncmp(curp, "\_goto\_", 6) \neq 0) abrt("missing\_goto", -16);
   curp += 6;
   for (j = 0; *curp \land *curp \neq ` \Box ` \land *curp \neq ` \" ; j++, curp ++)
       \textbf{if} \ (j<4) \ state[m].lab[j] = *curp; \\
   if (j > 4) abrt("the_{\sqcup}label_{\sqcup}is_{\sqcup}too_{\sqcup}long", -17);
This code is used in sections 7, 8, 10, and 14.
      \langle Scan an if step 10 \rangle \equiv
      curp += 3;
      \langle Scan the var field 11 \rangle;
      if (*curp ++ \neq '=') abrt("missing_{\sqcup}'='_{\sqcup}in_{\sqcup}an_{\sqcup}if_{\sqcup}step", -18);
      \langle Scan the val field 12 \rangle;
      \langle \text{ Scan the } lab \text{ field } 9 \rangle;
      \langle Scan the elab field 13 \rangle;
      if (*curp \neq '\n') abrt("that_if_step_ends_badly", -19);
   }
This code is used in section 5.
11. \langle \text{Scan the } var \text{ field } 11 \rangle \equiv
   for (j = 0; *curp \land *curp \neq `=` \land *curp \neq `\n'; j++, curp++)
      if (j < 2) state [m]. var[j] = vars[n][j] = *curp;
   \textbf{if } (j>2) \ \ abrt(\texttt{"the} \verb"variable" \verb"name" \verb"is" \verb"too" \verb"long", -20);\\
   if (state[m].var[0] < 'a' \lor state[m].var[0] > 'z')
      abrt("a_{\sqcup}variable_{\sqcup}name_{\sqcup}must_{\sqcup}begin_{\sqcup}with_{\sqcup}a_{\sqcup}lowercase_{\sqcup}letter", -21);
   for (j = 0; j < n; j++)
      \mathbf{if} \ (strncmp(vars[j], state[m].var, 2) \equiv 0) \ \mathbf{break};
   if (j \equiv n) n++;
   else vars[n][1] = 0;
This code is used in sections 10 and 14.
12. \langle \text{Scan the } val \text{ field } 12 \rangle \equiv
   if (*curp < 0, \forall *curp > 1, abrt("the value must be 0 or 1, -22);
   state[m].val = *curp ++ - '0';
This code is used in sections 10 and 14.
```

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```
13. \langle \text{Scan the } elab \text{ field } 13 \rangle \equiv
  if (strncmp(curp, "uelseu", 6) \neq 0) abrt("missinguelse", -23);
  curp += 6;
  for (j = 0; *curp \land *curp \neq `\n'; j++, curp ++)
     if (j < 4) state [m]. elab [j] = *curp;
  \textbf{if } (j>4) \ \ abrt(\texttt{"the\_else\_label\_is\_too\_long"}, -24);\\
This code is used in section 10.
     \langle Scan \text{ an assignment step } 14 \rangle \equiv
14.
     \langle Scan the var field 11 \rangle;
     if (*curp \leftrightarrow \neq '=') abrt("missing_{\square}'='_{\square}in_{\square}an_{\square}assignment_{\square}step", -25);
     \langle Scan the val field 12 \rangle;
     \langle Scan the lab field 9 \rangle;
     if (*curp \neq '\n') abrt("assignment_step_ends_badly", -26);
This code is used in section 5.
15. \langle Check for missing steps 15\rangle \equiv
  if (ma \equiv 0) {
     fprintf(stderr, "There = are = no = steps = for = process = A! \n");
     exit(-99);
  if (mb \equiv 0) {
     fprintf(stderr, "There \_are \_no \_steps \_for \_process \_B! \n");
     exit(-98);
  for (k = t = 0; k < m; k++) {
     if (state[k].lab[0]) {
        for (j = 0; j < m; j ++)
          if (strncmp(state[j].name, state[k].lab, 4) \equiv 0) break;
        if (j \equiv m) {
           fprintf(stderr, "Missing_step_\%.4s!\n", state[k].lab);
          t++;
        }
     if (state[k].elab[0]) {
        for (j = 0; j < m; j ++)
          if (strncmp(state[j].name, state[k].elab, 4) \equiv 0) break;
        if (j \equiv m) {
          fprintf(stderr, "Missing_step_\%.4s!\n", state[k].elab);
          t++;
        }
     }
  if (t) exit(-30);
This code is used in section 5.
```

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16. The generated clauses involve variables like '2A1', meaning that process A is in state A1 at time 2; also variables like '3b', meaning that shared variable b is 1 (true) at time 3; also variables like '10', meaning that process A took a turn at time 1. (The negations of these variables, namely ~2A1, ~3b, ~10, mean respectively that A is not in state A1 at time 2, b is 0 (false) at time 3, and process B took a turn at time 1.) At time 0, all shared variables are 0 and each process is in its first-mentioned state.

```
 \left\{ \begin{array}{ll} & \text{ for } (j=0;\; j< n;\; j++)\;\; printf\,(\verb"""000%.2s\n", vars[j]);\\ & printf\,(\verb"000%.4s\n", state[astep[0]].name);\\ & \text{ for } (j=1;\; j< ma;\; j++)\;\; printf\,(\verb"""000%.4s\n", state[astep[j]].name);\\ & printf\,(\verb"000%.4s\n", state[bstep[0]].name);\\ & \text{ for } (j=1;\; j< mb;\; j++)\;\; printf\,(\verb"""000%.4s\n", state[bstep[j]].name);\\ & \right\} \\ & \text{ This code is used in section 2.}
```

17. Speaking of turns reminds me that I promised to define precise semantics.

At each time t one of the processes, chosen nondeterministically, is granted permission to take a turn, which means intuitively that it performs the step corresponding to its current state. We say that the selected process is "bumped."

Every process is in a unique state at time t. The state of a process remains the same at time t+1 if it's not bumped. But if it's bumped, the next state is (1) either the same or lab, nondeterministically, after a maybe step; (2) lab after a critical step or an assignment step; (2) either lab or elab after an if step, depending on whether or not the shared variable has the specified value.

The value of a shared variable at time t + 1 is the same as the value that it had at time t, unless the bumped process assigned another value to it. In particular, if two processes are trying to change the same shared variable, the bumped process changes it first.

When the bumped process executes an if statement at the same time as another process is trying to write the same variable, the other process does not influence the result of the if; the change it wants to make will have to wait. [This rule means that weaker algorithms can get by, but they need stronger (and presumably more expensive and/or slower) hardware support. I'm using this rule in all the early examples of mutex in TAOCP, because it is easier to explain; the harder rule can be considered later, after algorithms pass this simpler criterion.]

```
 \left \langle \text{ Generate the transitions from time } t \text{ to time } t+1 \text{ 17} \right \rangle \equiv \left \{ \right. \\ \left. \left. \left\langle \text{ Generate clauses to forbid nonunique states for A at time } t+1 \text{ 18} \right \rangle; \right. \\ \left. \left\langle \text{ Generate clauses to forbid nonunique states for B at time } t+1 \text{ 19} \right \rangle; \right. \\ \left. \left\langle \text{ Generate the state transition clauses when A is bumped 20} \right \rangle; \right. \\ \left. \left\langle \text{ Generate the state transition clauses when B is bumped 22} \right \rangle; \right. \\ \left. \left\langle \text{ Generate the variable transition clauses 24} \right \rangle; \right. \\ \left. \left. \right. \right \}
```

This code is used in section 2.

§18 SAT-MUTEX INTRO 7

18. I introduce auxiliary variables here, using Heule's exclusion clauses, so that we don't have quadratic blowup when the programs are large.

```
#define printprevA()
        if (j) printf("\%03d_A\%d", t + 1, i - 1);
        else printf("~\%03d\%.4s", t+1, state[astep[k-1]].name);
\langle Generate clauses to forbid nonunique states for A at time t+1 18\rangle \equiv
  k = ma;
  if (k > 1) {
    i = j = 0;
    while (k > 4) {
      printprevA();
      printf("_{\perp}^{\sim}\%03d\%.4s\n", t+1, state[astep[k-2]].name);
      printprevA();
      printf(\verb"\"-"%03d%.4s\n",t+1,state[astep[k-3]].name);
      printprevA();
      printf("_{\square}^{\sim}\%03d_A\%d\n", t+1, i);
      printf("~\%03d\%.4s_{\sim}\%03d\%.4s_{\sim},t+1,state[astep[k-2]].name,t+1,state[astep[k-3]].name);\\
      printf("~\%03d\%.4s_~\%03d_A\%d\n", t+1, state[astep[k-2]].name, t+1, i);
      printf("~\%03d\%.4s_{\sim}\%03d_A\%d\n", t+1, state[astep[k-3]].name, t+1, i);
      i++, j=1, k-=2;
    printprevA();
    printf("_{\perp}^{\sim}\%03d\%.4s\n", t+1, state[astep[k-2]].name);
    printprevA();
    printf("_{\square} \%03d\%.4s\n", t+1, state[astep[k-3]].name);
    printf("^\%03d\%.4s_"\%03d\%.4s_", t+1, state[astep[k-2]].name, t+1, state[astep[k-3]].name);
    if (k > 3) {
      printprevA();
      printf("_{\square}^{\sim}\%03d\%.4s\n", t+1, state[astep[k-4]].name);
      printf("~\%03d\%.4s_~\%03d\%.4s_n", t+1, state[astep[k-2]].name, t+1, state[astep[k-4]].name);
      printf("^\%03d\%.4s_{^\%}03d\%.4s_{^\%}) + 1, state[astep[k-3]].name, t+1, state[astep[k-4]].name);
  }
```

This code is used in section 17.

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```
19. #define printprevB()
         if (j) printf("\%03d_B\%d", t + 1, i - 1);
         else printf("~\%03d\%.4s", t+1, state[bstep[k-1]].name);
\langle Generate clauses to forbid nonunique states for B at time t+1 19\rangle \equiv
  k = mb;
  if (k > 1) {
    i = j = 0;
    if (k \equiv 2) \ printf("~\%03d\%.4s_"~\%03d\%.4s_", t+1, state[bstep[0]].name, t+1, state[bstep[1]].name);
    while (k > 4) {
       printprevB();
       printf("_{\sqcup}~%03d%.4s\n", t+1, state[bstep[k-2]].name);
       printf("_{\square} \%03d\%.4s\n", t+1, state[bstep[k-3]].name);
       printprevB();
       printf("_{\square}^{\sim}\%03d_B\%d\n", t+1, i);
       printf("^{\infty}03d\%.4s_{1}^{\infty}03d\%.4s_{1}^{\infty}), t+1, state[bstep[k-2]].name, t+1, state[bstep[k-3]].name);
       printf("~\%03d\%.4s_{}~\%03d_B\%d\n", t+1, state[bstep[k-2]].name, t+1, i);
       printf("~\%03d\%.4s_{\sim}\%03d_B\%d\n", t+1, state[bstep[k-3]].name, t+1, i);
       i++, j=1, k-=2;
    printprevB();
    printf("_{\sqcup} \%03d\%.4s\n", t+1, state[bstep[k-2]].name);
    printprevB();
    printf("_1~%03d%.4s\n", t+1, state[bstep[k-3]].name);
    printf("^{\infty},03d\%.4s_1^{\infty},03d\%.4s_n",t+1,state[bstep[k-2]].name,t+1,state[bstep[k-3]].name);
    if (k > 3) {
       printprevB();
       printf("_{\perp}^{\sim}\%03d\%.4s\n", t+1, state[bstep[k-4]].name);
       printf("^\%03d\%.4s_{^\%}03d\%.4s_{^\%}) + 1, state[bstep[k-2]].name, t+1, state[bstep[k-4]].name);
       printf("^\%03d\%.4s_{^\%}03d\%.4s_{^\%}) + 1, state[bstep[k-3]].name, t+1, state[bstep[k-4]].name);
  }
This code is used in section 17.
20. #define tprime (t+1)
\langle Generate the state transition clauses when A is bumped 20\rangle \equiv
  for (k = 0; k < ma; k++) {
    printf("\%03d@_{\square}^{\infty}\%03d\%.4s_{\square}\%03d\%.4s_{\square}",t,t,state[astep[k]].name,tprime,state[astep[k]].name);
    if (state[astep[k]].var[0] \equiv 0) {
       if (state[astep[k]].crit \equiv 0)
                                         /* a maybe step */
         printf("~\%03d@_~\%03d\%.4s_\%03d\%.4s_\%03d\%.4s_n", t, t, state[astep[k]].name, tprime,
              state[astep[k]].name, tprime, state[astep[k]].lab);
       else printf("~\%03d@_~\%03d\%.4s_\m",t,t,state[astep[k]].name,tprime,state[astep[k]].lab);
             /* a critical step */
    } else if (state[astep[k]].elab[0] \equiv 0) {
                                                   /* an assignment step */
       printf("~\%03d@_1~\%03d\%.4s_1\%03d\%.4s_n",t,t,state[astep[k]].name,tprime,state[astep[k]].lab);
    } else (Generate clauses for when A is bumped in an if step 21);
This code is used in section 17.
```

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```
21.
             \langle Generate clauses for when A is bumped in an if step 21 \rangle \equiv
          printf("~\%03d@_{\bot}~\%03d\%.4s", t, t, state[astep[k]].name);
          printf("_{\square}\%s\%03d\%.2s_{\square}\%03d\%.4s\n", state[astep[k]].val?"^{"}:"",t, state[astep[k]].var, tprime,
                     state[astep[k]].lab);
          printf("~\%03d@_1~\%03d\%.4s", t, t, state[astep[k]].name);
          printf("\verb|||%s%03d%.2s||%03d%.4s\\n", state[astep[k]].val?"": "~", t, state[astep[k]].var, tprime, the printer of the printer
                     state[astep[k]].elab);
     }
This code is used in section 20.
22. \langle Generate the state transition clauses when B is bumped 22 \rangle \equiv
     for (k = 0; k < mb; k++) {
          printf("\"\%03d@_\"\%03d\%.4s\), t, t, state[bstep[k]].name, tprime, state[bstep[k]].name);
          if (state[bstep[k]].var[0] \equiv 0) {
               if (state[bstep[k]].crit \equiv 0)
                                                                                        /* a maybe step */
                    printf("\%03d@_{\sim}\%03d\%.4s_{\sim}\%03d\%.4s_{\sim}\%03d\%.4s_{\sim}", t, t, state[bstep[k]].name, tprime,
                               state[bstep[k]].name, tprime, state[bstep[k]].lab);
               else printf("\%03d@_~\%03d\%.4s_\%03d\%.4s_m",t,t,state[bstep[k]].name,tprime,state[bstep[k]].lab);
                          /* a critical step */
          } else if (state[bstep[k]].elab[0] \equiv 0) { /* an assignment step */
               printf("\%03d@_{\sim}\%03d\%.4s_{\sim}\%03d\%.4s_{\sim},t,t,state[bstep[k]].name,tprime,state[bstep[k]].lab);
          } else (Generate clauses for when B is bumped in an if step 23);
This code is used in section 17.
             \langle Generate clauses for when B is bumped in an if step 23\rangle \equiv
          printf("\%03d@_{\_}^{\sim}\%03d\%.4s", t, t, state[bstep[k]].name);
          printf("_{\square}\%s\%03d\%.2s_{\square}\%03d\%.4s\n", state[bstep[k]].val?"^{"}:"",t,state[bstep[k]].var,tprime,
                    state[bstep[k]].lab);
          printf("\%03d@_{\square}^{\sim}\%03d\%.4s", t, t, state[bstep[k]].name);
          printf("u\%s\%03d\%.2s_{\%}03d\%.4s\n", state[bstep[k]].val?"":"~", t, state[bstep[k]].var, tprime,
                     state[bstep[k]].elab);
     }
This code is used in section 22.
```

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```
\langle Generate the variable transition clauses 24\rangle \equiv
  for (k = 0; k < n; k++) {
                                 /* first consider all cases where the value changes */
    for (j = 0; j < m; j ++)
      if (strncmp(state[j].var, vars[k], 2) \equiv 0 \land state[j].elab[0] \equiv 0)
         state[j].name, state[j].val \equiv 0? "~": "", tprime, state[j].var);
           /* now consider all cases where the value doesn't change */
    printf("~\%03d@_{\sqcup}\%03d\%.2s", t, t, vars[k]);
                                                 /* A bumped and val is 0 */
    for (j = 0; j < m; j ++)
      if (strncmp(state[j].var, vars[k], 2) \equiv 0 \land state[j].elab[0] \equiv 0 \land state[j].name[0] \equiv A
         /* not changed by A */
    /* it stays 0 */
    printf("\%03d@_{\sqcup}\%03d\%.2s", t, t, vars[k]);
                                                /* B bumped and val is 0 */
    for (j = 0; j < m; j ++)
      if (strncmp(state[j].var, vars[k], 2) \equiv 0 \land state[j].elab[0] \equiv 0 \land state[j].name[0] \equiv 'B')
         printf(" " \%03d\%.4s", t, state[j].name);
                                                   /* not changed by B */
                                                  /* it stays 0 */
    printf("_{\square}^{\sim}\%03d\%.2s\n", tprime, vars[k]);
    printf("~\%03d@_{\_}~\%03d\%.2s", t, t, vars[k]);
                                                  /* A bumped and val is 1 */
    for (j = 0; j < m; j ++)
      if (strncmp(state[j].var, vars[k], 2) \equiv 0 \land state[j].elab[0] \equiv 0 \land state[j].name[0] \equiv A
         /* not changed by A */
    /* it stays 1 */
    printf("\%03d@_{\square}^{\sim}\%03d\%.2s", t, t, vars[k]);
                                                 /* B bumped and val is 1 */
    for (j = 0; j < m; j ++)
      if (strncmp(state[j].var, vars[k], 2) \equiv 0 \land state[j].elab[0] \equiv 0 \land state[j].name[0] \equiv 'B')
         /* not changed by B */
    /* it stays 1 */
This code is used in section 17.
     The different ways of going jointly critical are CO, C1, etc.
\langle Generate clauses to force concurrent critical sections at time r 25 \rangle \equiv
  for (i = j = 0; j < ma; j ++)
    if (state[astep[j]].crit) {
      for (k = 0; k < mb; k++)
         if (state[bstep[k]].crit) {
           printf("\ C\d_{\sqcup}\%03d\%.4s\n", i, r, state[astep[j]].name);
           printf("\ C\ d_\ 03d\ .4s\ n", i, r, state[bstep[k]].name);
           i++;
         }
   \mathbf{for} \ (j = 0; \ j < i; \ j +\!\!\!+\!\!\!) \ \mathit{printf} ( \verb""$ $\square$ C\d", j ); 
  printf("\n");
This code is used in section 2.
```

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abrt: 6, 7, 8, 9, 10, 11, 12, 13, 14. $argc: \underline{2}, 3.$ $\mathit{argv}\colon \ \underline{2},\ 3.$ $astep \colon \ \underline{2}, \ 5, \ 6, \ 16, \ 18, \ 20, \ 21, \ 25.$ $bstep: \ \underline{2}, 5, 6, 16, 19, 22, 23, 25.$ $\mathit{buf}\colon \ \underline{2},\ 5,\ 6.$ bufsize: $\underline{2}$, 5. $crit: \underline{4}, 5, 8, 20, 22, 25.$ $curp\colon \ \underline{5},\ 6,\ 7,\ 8,\ 9,\ 10,\ 11,\ 12,\ 13,\ 14.$ elab: $\underline{4}$, 13, 15, 17, 20, 21, 22, 23, 24. exit: 3, 5, 6, 15.fgets: 5. fprintf: 3, 5, 6, 15. *i*: 2. j: $\underline{2}$. k: $\underline{2}$. $lab: \underline{4}, 9, 15, 17, 20, 21, 22, 23.$ $m: \underline{2}.$ $ma: \underline{2}, 5, 6, 15, 16, 18, 20, 25.$ $main\colon \ \underline{2}.$ $maxsteps: \underline{2}, 5.$ $mb: \ \underline{2}, 5, 6, 15, 16, 19, 22, 25.$ $name: \underline{4}, 6, 15, 16, 18, 19, 20, 21, 22, 23, 24, 25.$ printf: 3, 5, 16, 18, 19, 20, 21, 22, 23, 24, 25. $printprevA\colon \ \underline{18}.$ $printprevB: \underline{19}.$ r: $\underline{2}$. sscanf: 3. $state: \ \underline{2}, 5, 6, 8, 9, 11, 12, 13, 15, 16, 18, 19,$ $20,\ 21,\ 22,\ 23,\ 24,\ 25.$ $state_struct: \underline{4}.$ stderr: 3, 5, 6, 15. stdin: 2, 5. step: $2, \underline{4}$. strncmp: 5, 6, 9, 11, 13, 15, 24. t: 2. tprime: 20, 21, 22, 23, 24. $val\colon \ \ \underline{4},\ 12,\ 21,\ 23,\ 24.$ $var \colon \quad \underline{4}, \ 11, \ 20, \ 21, \ 22, \ 23, \ 24.$ vars: 2, 11, 16, 24.

12 NAMES OF THE SECTIONS SAT-MUTEX

```
(Check for missing steps 15) Used in section 5.
 Generate clauses for when A is bumped in an if step 21 \ Used in section 20.
 Generate clauses for when B is bumped in an if step 23 \ Used in section 22.
 Generate clauses to forbid nonunique states for A at time t+1 18 \rangle Used in section 17.
 Generate clauses to forbid nonunique states for B at time t+1 19 \) Used in section 17.
 Generate clauses to force concurrent critical sections at time r 25 \rangle Used in section 2.
 Generate the initial clauses 16 \rangle Used in section 2.
 Generate the state transition clauses when A is bumped 20 \rangle Used in section 17.
 Generate the state transition clauses when B is bumped 22 \ Used in section 17.
 Generate the transitions from time t to time t+1 17 \rangle Used in section 2.
 Generate the variable transition clauses 24 \rangle Used in section 17.
 Parse the input into the state table 5 \ Used in section 2.
 Process the command line 3 \ Used in section 2.
 Scan a critical step 8 \rangle Used in section 5.
 Scan a maybe step 7 Used in section 5.
 Scan an if step 10 \ Used in section 5.
 Scan an assignment step 14 \rangle Used in section 5.
 Scan the elab field 13 Vsed in section 10.
 Scan the lab field 9 \rangle Used in sections 7, 8, 10, and 14.
 Scan the name field 6 \ Used in section 5.
 Scan the val field 12 \rangle Used in sections 10 and 14.
 Scan the var field 11 \rangle Used in sections 10 and 14.
⟨ Type definitions 4⟩ Used in section 2.
```

SAT-MUTEX

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