

Inferences in Bayesian Networks

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Outline

- 1 Introduction
- 2 Exact inference by enumeration
- 3 Simple Bayesian Network Inferences
- 4 Exact inference in polytrees
- 5 Summary

Inference tasks

- Simple queries: compute posterior marginal $P(X_i|\mathbf{E} = \mathbf{e})$
e.g., $P(P = low|S = T, X = pos, D = T)$
- Conjunctive queries:
$$P(X_i, X_j|\mathbf{E} = \mathbf{e}) = P(X_i|\mathbf{E} = \mathbf{e}) \times P(X_j|X_i, \mathbf{E} = \mathbf{e})$$
- Optimal decisions: decision networks include utility information; probabilistic inference required for $P(outcome|action, evidence)$
- Value of information: which evidence to seek next?
- Sensitivity analysis: which probability values are most critical?
- Explanation: why do I need a new starter motor?

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Inference by enumeration

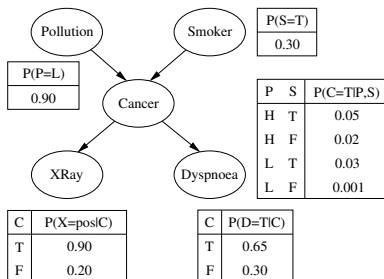
$$\begin{aligned}P(X|\mathbf{E} = \mathbf{e}) &= \sum_{\mathbf{y} \in \text{Domain}(\mathbf{Y})} P(X, \mathbf{Y} | \mathbf{E} = \mathbf{e}) \\&= \alpha \sum_{\mathbf{y} \in \text{Domain}(\mathbf{Y})} P(X, \mathbf{Y}, \mathbf{E} = \mathbf{e}) \\P(X_1, X_2 | \mathbf{E} = \mathbf{e}) &= \sum_{\mathbf{y} \in \text{Domain}(\mathbf{Y})} P(X_1, X_2, \mathbf{Y} | \mathbf{E} = \mathbf{e}) \\&= \alpha \sum_{\mathbf{y} \in \text{Domain}(\mathbf{Y})} P(X_1, X_2, \mathbf{Y}, \mathbf{E} = \mathbf{e})\end{aligned}$$

where $\alpha = \frac{1}{P(\mathbf{E} = \mathbf{e})}$.

- Recursive depth-first enumeration: $O(n)$ space, $O(|\text{Domain}(X_i)|^n)$ time

Inference by enumeration (example)

Simple query on the lung cancer example



$$\begin{aligned} P(P|X = pos, D = T) \\ &= P(P, X = pos, D = T) / P(X = pos, D = T) \\ &= \alpha P(P, X = pos, D = T) \\ &= \alpha \sum_{v_c \in \{T, F\}} \sum_{v_s \in \{T, F\}} P(P, S = v_s, C = v_c, X = pos, D = T) \end{aligned}$$

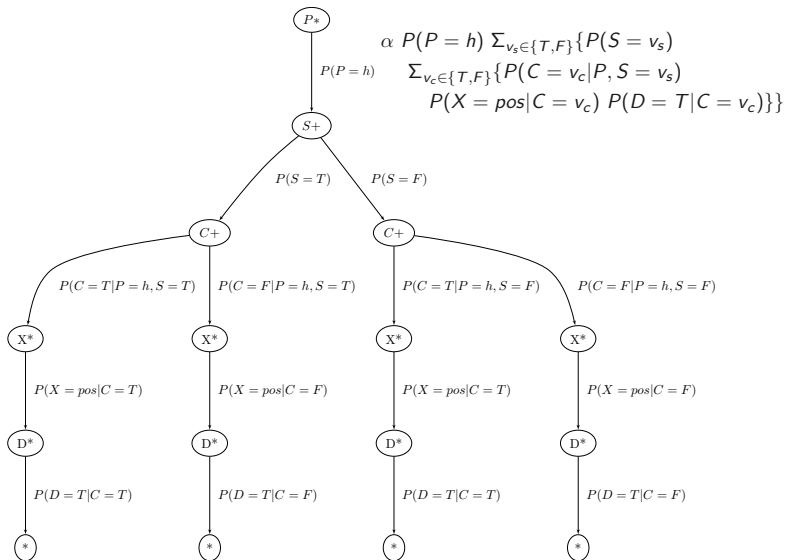
where $\alpha = \frac{1}{P(X=pos, D=T)}$ is the normalizer.

Inference by enumeration (example cont.)

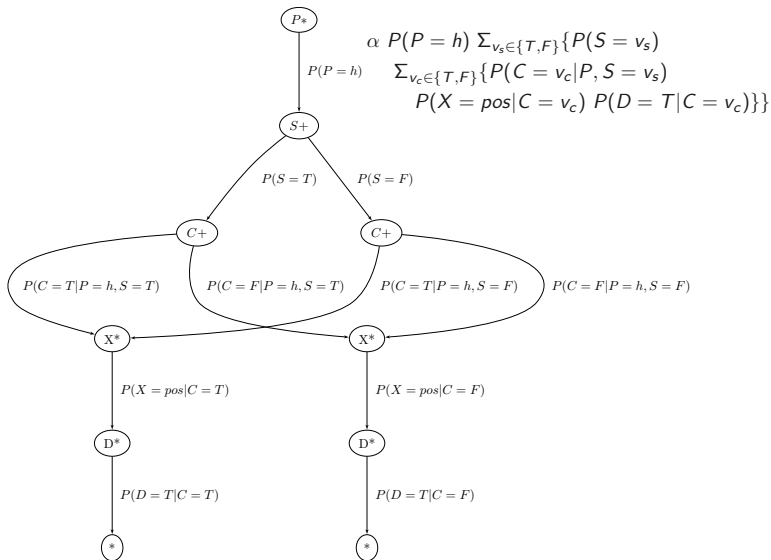
Rewrite full joint entries using product of CPT entries:

$$\begin{aligned} & P(P|X = pos, D = T) \\ &= \alpha \sum_{v_c \in \{T, F\}} \sum_{v_s \in \{T, F\}} \{ \\ &\quad P(P) \times P(S = v_s) \\ &\quad \times P(C = v_c|P, S = v_s) \\ &\quad \times P(X = pos|C = v_c) \times P(D = T|C = v_c)\} \\ &= \alpha \times P(P) \\ &\quad \times \sum_{v_s \in \{T, F\}} \{ \\ &\quad \quad P(S = v_s) \\ &\quad \quad \times \sum_{v_c \in \{T, F\}} \{P(C = v_c|P, S = v_s) \\ &\quad \quad \quad \times P(X = pos|C = v_c) \\ &\quad \quad \quad \times P(D = T|C = v_c)\} \\ &\quad \quad \} \end{aligned}$$

Evaluation tree



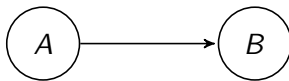
Reduced evaluation tree



Outline

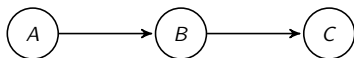
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Inference for two node network



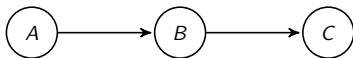
- $P(B|A)$ can be read from the CPTs
- $P(A|B) = \frac{P(B|A)P(A)}{P(B)}$ by Bayes' rule
 - ▶ $P(A), P(B|A)$ can be read from the CPTs
 - ▶ $\alpha = \frac{1}{P(B)}$ can be computed by $\sum_{a \in \text{Domain}(A)} P(A = a|B) = 1$.

Inference for three node chain network I



- $P(B|A)$ and $P(C|B)$ can be read from the tables
- $P(A|B)$ and $P(C|B)$ can be computed with the Bayes' rule
- $P(C|A)$ can be computed by
$$\sum_{b \in \text{Domain}(B)} P(B|A)P(C|B, A) = \sum_{b \in \text{Domain}(B)} P(B|A)P(C|B)$$
- $P(A|C)$ can be computed by $P(A|C) = \alpha P(C|A) P(A)$
- $P(C|A, B) = P(C|B)$ since C is independent of A given B
- $P(B|A, C) = \alpha P(C|A, B)P(B|A) = \alpha P(C|B)P(B|A)$ – first apply Bayes' rule conditional on A , and then simplify $P(C|A, B)$ by the conditional independence in the network
- $P(A|B, C) = \alpha P(C|A, B)P(A|B) = \alpha P(C|B)P(A|B)$ – first apply Bayes' rule conditional on B , and then simplify $P(C|A, B)$ into $P(C|B)$; $P(A|B)$ are computed above.

Inference for three node chain network II



- $P(A, B|C) = \alpha P(C|A, B)P(A, B) = \alpha P(C|B)P(B|A)P(A)$ — first apply Bayes' rule by taking A, B and C as two events, and then simplify.
- $P(A, C|B) = P(A|B)P(C|B) = \alpha P(B|A)P(A)P(C|B)$ — first apply conditional independence of A and C given B , and apply Bayes' rule on $P(A|B)$
- $P(B, C|A) = P(B|A)P(C|A, B) = P(B|A)P(C|B)$ — apply chain rule conditional on A , and the conditional independence

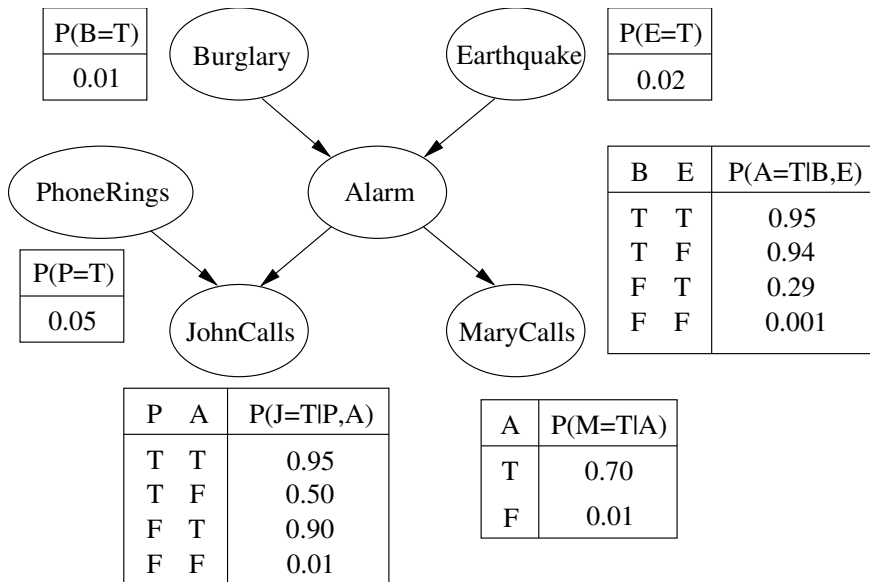
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Polytrees (“forest”)

- Polytrees have at most one path between any pair of nodes
- Also called **singly-connected** networks

Earthquake Example



The query

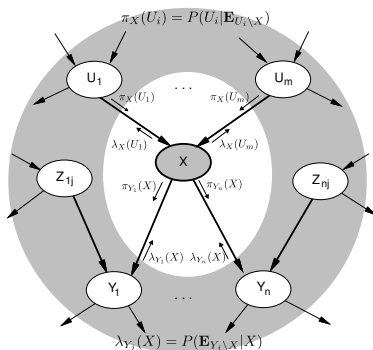
$$\begin{aligned} P(X|\mathbf{E}) &= \frac{P(X, \mathbf{E})}{P(\mathbf{E})} \\ &= \alpha P(X, \mathbf{E}) \end{aligned}$$

Look at the neighbors, marginalize, and re-organize the variables

$$P(X, \mathbf{E})$$

$$= \sum_{U_1, \dots, U_m} [P(X, U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X})]$$

$$= \sum_{U_1, \dots, U_m} [P(U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, X, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X})]$$



Chain rule

$$\begin{aligned} P(U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, X, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X}) = \\ P(U_1, \mathbf{E}_{U_1 \setminus X}) \\ P(U_2, \mathbf{E}_{U_2 \setminus X} | U_1, \mathbf{E}_{U_1 \setminus X}) \\ P(U_3, \mathbf{E}_{U_3 \setminus X} | U_1, \mathbf{E}_{U_1 \setminus X}, U_2, \mathbf{E}_{U_2 \setminus X}) \\ \dots \\ P(U_m, \mathbf{E}_{U_m \setminus X} | U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_{m-1}, \mathbf{E}_{U_{m-1} \setminus X}) \\ P(X | U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}) \\ P(\mathbf{E}_{Y_1 \setminus X} | X, U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}) \\ P(\mathbf{E}_{Y_2 \setminus X} | X, U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, \mathbf{E}_{Y_1 \setminus X}) \\ \dots \\ P(\mathbf{E}_{Y_n \setminus X} | X, U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_{n-1} \setminus X}) \end{aligned}$$

Independence/conditional independence

$$P(U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, X, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X}) =$$

$$P(U_1, \mathbf{E}_{U_1 \setminus X})$$

// $U_i, \mathbf{E}_{U_i \setminus X}$ is independent of $U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_{i-1}, \mathbf{E}_{U_{i-1} \setminus X}$:

$$P(U_2, \mathbf{E}_{U_2 \setminus X})$$

...

$$P(U_m, \mathbf{E}_{U_m \setminus X})$$

// X is independent of $\mathbf{E}_{U_i \setminus X}$ given U_i ($i = 1, \dots, m$):

$$P(X | U_1, U_2, \dots, U_m)$$

// $\mathbf{E}_{Y_k \setminus X}$ is independent of U_i s, $\mathbf{E}_{U_i \setminus X}$ s and $\mathbf{E}_{Y_{k'} \setminus X}$ ($k' \neq k$) given X :

$$P(\mathbf{E}_{Y_1 \setminus X} | X)$$

$$P(\mathbf{E}_{Y_2 \setminus X} | X)$$

...

$$P(\mathbf{E}_{Y_n \setminus X} | X)$$

One more re-write with the definition of conditionals

$$\begin{aligned} & P(U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, X, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X}) \\ &= \\ & \quad P(U_1 | \mathbf{E}_{U_1 \setminus X}) P(\mathbf{E}_{U_1 \setminus X}) \\ & \quad P(U_2 | \mathbf{E}_{U_2 \setminus X}) P(\mathbf{E}_{U_2 \setminus X}) \\ & \quad \dots \\ & \quad P(U_m | \mathbf{E}_{U_m \setminus X}) P(\mathbf{E}_{U_m \setminus X}) \\ & \quad P(X | U_1, U_2, \dots, U_m) \\ & \quad P(\mathbf{E}_{Y_1 \setminus X} | X) \\ & \quad P(\mathbf{E}_{Y_2 \setminus X} | X) \\ & \quad \dots \\ & \quad P(\mathbf{E}_{Y_n \setminus X} | X) \\ &= P(X | U_1, U_2, \dots, U_m) \\ & \quad \cdot \prod_{i=1, \dots, m} P(\mathbf{E}_{U_i \setminus X}) \cdot \prod_{i=1, \dots, m} P(U_i | \mathbf{E}_{U_i \setminus X}) \cdot \prod_{k=1, \dots, n} P(\mathbf{E}_{Y_k \setminus X} | X) \end{aligned}$$

Back to the query

$$\begin{aligned} P(X|\mathbf{E}) &= \alpha \sum_{U_1, \dots, U_m} \{ \\ &\quad P(U_1, \mathbf{E}_{U_1 \setminus X}, \dots, U_m, \mathbf{E}_{U_m \setminus X}, X, \mathbf{E}_{Y_1 \setminus X}, \dots, \mathbf{E}_{Y_n \setminus X}) \} \\ &= \alpha \sum_{U_1, \dots, U_m} \{ P(X|U_1, U_2, \dots, U_m) \\ &\quad \prod_{i=1, \dots, m} P(\mathbf{E}_{U_i \setminus X}) \cdot \prod_{i=1, \dots, m} P(U_i|\mathbf{E}_{U_i \setminus X}) \cdot \prod_{k=1, \dots, n} P(\mathbf{E}_{Y_k \setminus X}|X) \} \\ &= \alpha \cdot \prod_{i=1, \dots, m} P(\mathbf{E}_{U_i \setminus X}) \cdot \sum_{U_1, \dots, U_m} \{ P(X|U_1, U_2, \dots, U_m) \\ &\quad \prod_{i=1, \dots, m} P(U_i|\mathbf{E}_{U_i \setminus X}) \cdot \prod_{k=1, \dots, n} P(\mathbf{E}_{Y_k \setminus X}|X) \} \end{aligned}$$

As $\prod_{i=1, \dots, m} P(\mathbf{E}_{U_i \setminus X})$ doesn't change when X takes different values, we can put it into the normalizing **constant** α :

$$\begin{aligned} P(X|\mathbf{E}) &= \alpha \sum_{U_1, \dots, U_m} \{ P(X|U_1, U_2, \dots, U_m) \\ &\quad \prod_{i=1, \dots, m} P(U_i|\mathbf{E}_{U_i \setminus X}) \cdot \prod_{k=1, \dots, n} P(\mathbf{E}_{Y_k \setminus X}|X) \} \end{aligned}$$

Recursively compute $P(U_i|\mathbf{E}_{U_i \setminus X})$ and $P(\mathbf{E}_{Y_k}|X)$

$$P(X|\mathbf{E}) = \alpha \sum_{U_1, \dots, U_m} \{ \\ P(X|U_1, U_2, \dots, U_m) \\ \prod_{i=1, \dots, m} P(U_i|\mathbf{E}_{U_i \setminus X}) \\ \prod_{k=1, \dots, n} P(\mathbf{E}_{Y_k \setminus X}|X) \}$$

- $P(X|U_1, U_2, \dots, U_m)$ can be looked from the CPTs
- Let $BN_X^\pi(U_i)$ be a sub-BN composed of U_i and all the nodes connecting to X through U_i .
 - ▶ Within $BN_X^\pi(U_i)$, $P(U_i|\mathbf{E}_{U_i \setminus X})$ can be computed recursively in the same manner
- Let $BN_{Y_k}^\lambda(X)$ be a sub-BN composed of Y_k and all the nodes connecting to X through Y_k .
 - ▶ Within $BN_{Y_k}^\lambda(X)$, $P(\mathbf{E}_{Y_k \setminus X}|X) = \sum_{Y_k} P(Y_k, \mathbf{E}_{Y_k \setminus X}|X)$ can be computed recursively in the same manner

Recursively compute $P(U_i | \mathbf{E}_{U_i \setminus X})$

Within $BN_X^\pi(U_i)$,

$$P(U_i | \mathbf{E}_{U_i \setminus X}) = \alpha \sum_{Parents^{BN_X^\pi(U_i)}(U_i)} \{ \\ P(U_i | Parents^{BN_X^\pi(U_i)}(U_i)) \cdot \\ \prod_{Z_j \in Parents^{BN_X^\pi(U_i)}(U_i)} P(Z_j | \mathbf{E}_{Z_j \setminus U_i}) \cdot \\ \prod_{Y_k \in Children^{BN_X^\pi(U_i)}(U_i)} P(\mathbf{E}_{Y_k \setminus X} | U_i) \}$$

- $Parents^{BN_X(U_i)}(U_i)$ are the parents of U_i in $BN_X(U_i)$
- $Children^{BN_X(U_i)}(U_i)$ are the children of U_i in $BN_X(U_i)$
- When U_i is one of the evidence input
 - ▶ $P(U_i = u_{i,e} | \mathbf{E}_{U_i \setminus X}) = 1$ if $u_{i,e}$ is the evidence value entered
 - ▶ $P(U_i = u_{i,e} | \mathbf{E}_{U_i \setminus X}) = 0$ if $u_{i,e}$ is not the evidence value entered

Recursively compute $P(\mathbf{E}_{Y_k}|X)$ I

Take $P(\mathbf{E}_{Y_k}|X)$ as a joint probability on the evidences connecting to X through Y_k and all the probabilities in $BN_{Y_k}^\lambda(X)$ as conditional on X

$$\begin{aligned}P(\mathbf{E}_{Y_k \setminus X}|X) &= \sum_{Y_k} P(Y_k, \mathbf{E}_{Y_k \setminus X}|X) \\P(Y_k, \mathbf{E}_{Y_k \setminus X}|X) &= \alpha \sum_{Parents^{BN_{Y_k}^\lambda(X)}(Y_k), Children^{BN_{Y_k}^\lambda(X)}(Y_k)} \{ \\&\quad P(Y_k | Parents^{BN_{Y_k}^\lambda(X)}(Y_k), X) \cdot \\&\quad \prod_{Z_{k,j} \in Parents^{BN_{Y_k}^\lambda(X)}(Y_k)} P(Z_{k,j} | \mathbf{E}_{Z_{k,j} \setminus Y_k}, X) \cdot \\&\quad \prod_{W_{k,l} \in Children^{BN_{Y_k}^\lambda(X)}(Y_k)} P(\mathbf{E}_{W_{k,l} \setminus Y_k} | Y_k, X) \} \end{aligned}$$

- Observations

- ▶ $P(Y_k | Parents^{BN_{Y_k}^\lambda(X)}(Y_k), X)$ is in the CPTs

Recursively compute $P(\mathbf{E}_{Y_k}|X)$ II

- ▶ $P(Z_{k,j}|\mathbf{E}_{Z_{k,j}\setminus Y_k}, X) = P(Z_{k,j}|\mathbf{E}_{Z_{k,j}\setminus Y_k})$ by conditional independence ($Z_{k,j}$ and X are marginal independence and there is no path between $Z_{k,j}$ and X through $\mathbf{E}_{Z_{k,j}}$)
- ▶ $P(\mathbf{E}_{W_{k,l}\setminus Y_k}|Y_k, X) = P(\mathbf{E}_{W_{k,l}\setminus Y_k}|Y_k)$ by conditional independence (Given Y_k , $\mathbf{E}_{W_{k,l}}$ and X are independence)

$$\begin{aligned} P(\mathbf{E}_{Y_k\setminus X}|X) &= \alpha \sum_{Y_k} \prod_{Parents^{BN_{Y_k}^\lambda(X)}(Y_k)} P(Y_k | Parents^{BN_{Y_k}^\lambda(X)}(Y_k), X) \cdot \\ &\quad \prod_{Z_{k,j} \in Parents^{BN_{Y_k}^\lambda(X)}(Y_k)} P(Z_{k,j} | \mathbf{E}_{Z_{k,j}\setminus Y_k}) \cdot \\ &\quad \prod_{W_{k,l} \in Children^{BN_{Y_k}^\lambda(X)}(Y_k)} P(\mathbf{E}_{W_{k,l}\setminus Y_k} | Y_k) \Big\} \end{aligned}$$

Recursively compute $P(\mathbf{E}_{Y_k}|X)$ III

- When Y_k is one of the evidence input, let $y_{k,e}$ be the evidence value entered, the marginalization over Y_k is replaced by setting $Y_k = y_{k,e}$ in the calculation.

$$\begin{aligned} P(\mathbf{E}_{Y_k \setminus X}|X) = & \alpha \sum_{Parents^{BN_{Y_k}^\lambda(X)}(Y_k), Children^{BN_{Y_k}^\lambda(X)}(Y_k)} \{ \\ & P(Y_k = y_{k,e} | Parents^{BN_{Y_k}^\lambda(X)}(Y_k), X) \cdot \\ & \prod_{Z_{k,j} \in Parents^{BN_{Y_k}^\lambda(X)}(Y_k)} P(Z_{k,j} | \mathbf{E}_{Z_{k,j} \setminus Y_k}) \cdot \\ & \prod_{W_{k,l} \in Children^{BN_{Y_k}^\lambda(X)}(Y_k)} P(\mathbf{E}_{W_{k,l} \setminus Y_k} | Y_k = y_{k,e}) \} \end{aligned}$$

Shorten the notation for $\pi_X(U_i)$

Let

$$\pi_X(U_i) = P(U_i | \mathbf{E}_{U_i \setminus X})$$

$$\lambda_{Y_k}(X) = P(\mathbf{E}_{Y_k \setminus X} | X)$$

$$\begin{aligned} \pi_X(U_i) &= P(U_i | \mathbf{E}_{U_i \setminus X}) \\ &= \alpha \sum_{Parents^{BN_X^\pi(U_i)}(U_i)} \{ \\ &\quad P(U_i | Parents^{BN_X^\pi(U_i)}(U_i)) \cdot \\ &\quad \prod_{Z_j \in Parents^{BN_X^\pi(U_i)}(U_i)} \pi_{U_i}(Z_j) \cdot \\ &\quad \prod_{Y_k \in Children^{BN_X^\pi(U_i)}(U_i)} \lambda_{Y_k}(U_i) \} \end{aligned}$$

When U_i is one of the evidence input

- $\pi_X(U_i = u_{i,e}) = 1$ if $u_{i,e}$ is the evidence value entered
- $\pi_X(U_i = u_{i,e}) = 0$ if $u_{i,e}$ is not the evidence value entered

Shorten the notation for $\pi_X(U_i)$ (cont.)

Let

$$\begin{aligned}\pi_X(U_i) &= P(U_i | \mathbf{E}_{U_i \setminus X}) \\ \lambda_{Y_k}(X) &= P(\mathbf{E}_{Y_k \setminus X} | X)\end{aligned}$$

$$\begin{aligned}\pi_X(U_i) &= P(U_i | \mathbf{E}_{U_i \setminus X}) \\ &= \alpha \sum_{Parents(U_i)} \{ \\ &\quad P(U_i | Parents(U_i)) \cdot \\ &\quad \prod_{Z_j \in Parents(U_i)} \pi_{U_i}(Z_j) \cdot \\ &\quad \prod_{Y_k \in Children(U_i) \setminus \{X\}} \lambda_{Y_k}(U_i) \}\end{aligned}$$

When U_i is one of the evidence input

- $\pi_X(U_i = u_{i,e}) = 1$ if $u_{i,e}$ is the evidence value entered
- $\pi_X(U_i = u_{i,e}) = 0$ if $u_{i,e}$ is not the evidence value entered

Shorten the notation for $\lambda_{Y_k}(X)$

Let

$$\pi_X(U_i) = P(U_i | \mathbf{E}_{U_i \setminus X})$$

$$\lambda_{Y_k}(X) = P(\mathbf{E}_{Y_k \setminus X} | X)$$

$$\begin{aligned} \lambda_{Y_k}(X) &= \sum_{Y_k} P(Y_k, \mathbf{E}_{Y_k \setminus X} | X) \\ &= \sum_{Y_k} \sum_{Parents^{BN^{\lambda_{Y_k}(X)}(Y_k)}} \{ \\ &\quad P(Y_k | Parents^{BN^{\lambda_{Y_k}(X)}(Y_k)}(Y_k), X) \cdot \\ &\quad \prod_{Z_{k,j} \in Parents^{BN^{\lambda_{Y_k}(X)}(Y_k)}} \pi_{Y_k}(Z_{k,j}) \cdot \\ &\quad \prod_{W_{k,l} \in Children^{BN^{\lambda_{Y_k}(X)}(Y_k)}} \lambda_{W_{k,l}}(Y_k) \} \end{aligned}$$

When Y_k is one of the evidence input, let $y_{k,e}$ be the evidence value entered, the marginalization over Y_k is replaced by setting $Y_k = y_{k,e}$ in the calculation.

Shorten the notation for $\lambda_{Y_k}(X)$ (cont.) I

Let

$$\begin{aligned}\pi_X(U_i) &= P(U_i | \mathbf{E}_{U_i \setminus X}) \\ \lambda_{Y_k}(X) &= P(\mathbf{E}_{Y_k \setminus X} | X)\end{aligned}$$

$$\begin{aligned}\lambda_{Y_k}(X) &= \sum_{Y_k} P(Y_k, \mathbf{E}_{Y_k \setminus X} | X) \\ &= \sum_{Y_k} \sum_{Parents(Y_k) \setminus \{X\}} \{ \\ &\quad P(Y_k | Parent(Y_k)) \cdot \\ &\quad \prod_{Z_{k,j} \in Parents(Y_k) \setminus \{X\}} \pi_{Y_k}(Z_{k,j}) \cdot \\ &\quad \prod_{W_{k,l} \in Children(Y_k)} \lambda_{W_{k,l}}(Y_k) \} \end{aligned}$$

Shorten the notation for $\lambda_{Y_k}(X)$ (cont.) II

When Y_k is one of the evidence input, let $y_{k,e}$ be the evidence value entered, the marginalization over Y_k is replaced by setting $Y_k = y_{k,e}$ in the calculation.

$$\begin{aligned}\lambda_{Y_k}(X) &= P(Y_k = y_{k,e}, \mathbf{E}_{Y_k \setminus X} | X) \\ &= \sum_{Parents(Y_k) \setminus \{X\}} \{ \\ &\quad P(Y_k = y_{k,e} | Parent(Y_k)) \cdot \\ &\quad \prod_{Z_{k,j} \in Parents(Y_k) \setminus \{X\}} \pi_{Y_k}(Z_{k,j}) \cdot \\ &\quad \prod_{W_{k,l} \in Children(Y_k)} \lambda_{W_{k,l}}(Y_k) \} \end{aligned}$$

Message passing (bottom-up instead of recursion) I

- Query decomposition

$$P(X|\mathbf{E}) = \alpha \sum_{U_1, \dots, U_m} \{ \\ P(X|U_1, U_2, \dots, U_m) \cdot \\ \prod_{i=1, \dots, m} \pi_{U_i}(X) \cdot \\ \prod_{k=1, \dots, n} \lambda_{Y_k}(X) \}$$

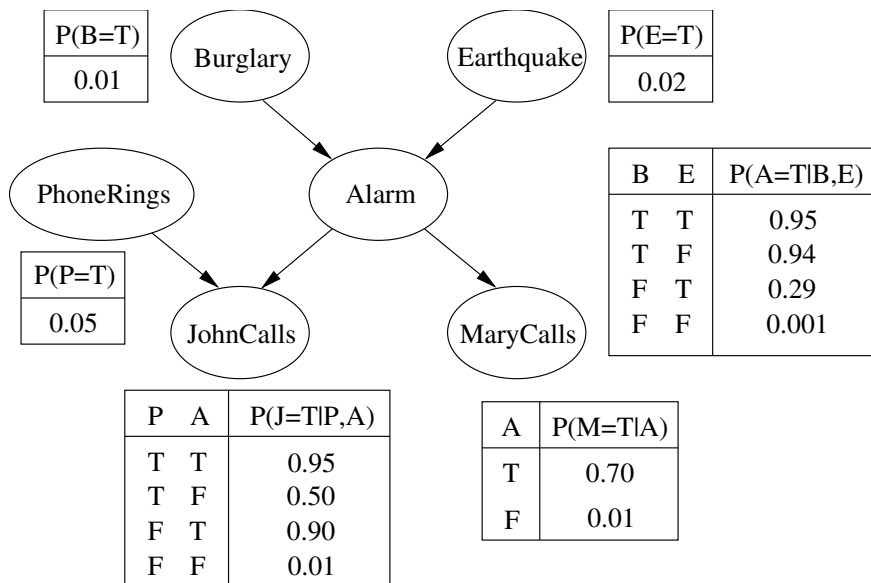
- Message passing:

- ▶ Start with the nodes that don't need any messages to compute π
- ▶ Start with the nodes that don't need any messages to compute λ
- ▶ Once a node gets the π and λ messages required to compute its own messages sent to its children or parents, compute the messages and send them out. Notice that
 - ★ U_i send to its child X : $\pi_X(U_i)$ — computing $\pi_X(U_i)$ only requires messages from nodes in $BN_X^\pi(U_i)$
 - ★ Y_k send to its parent X : $\lambda_{Y_k}(X)$ — computing $\lambda_{Y_k}(X)$ only requires messages from nodes in $BN_{Y_k}^\lambda(X)$

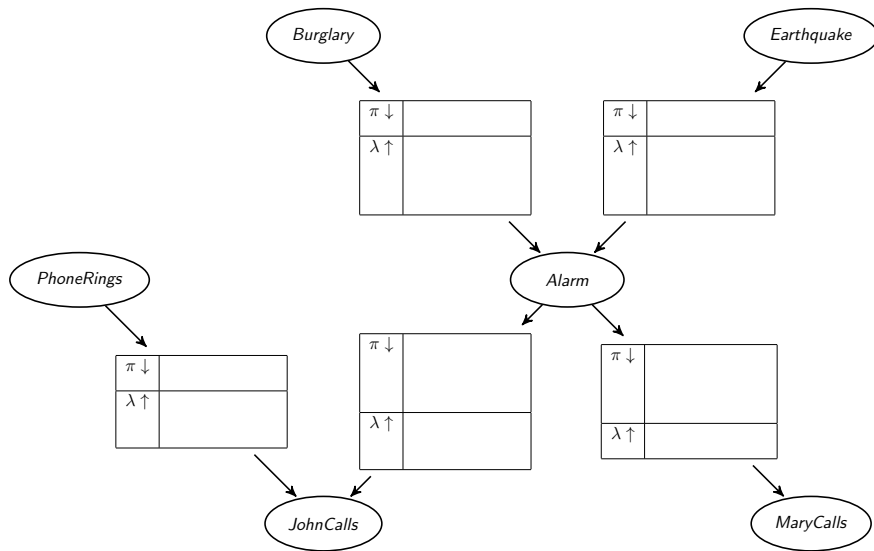
Message passing (bottom-up instead of recursion) II

- X computes query $P(X|\mathbf{E})$ using its conditional table $P(X|Parents(X))$, and the message $\pi_{U_i}(X)$ received from its parents, the message $\lambda_{Y_k}(X)$ received from its children
- For queries other than $P(X|\mathbf{E})$, X send messages
 - ▶ X send message $\pi_{Y_k}(X)$ to its child Y_k
 - ▶ X send message $\lambda_X(U_i)$ to its parent U_i

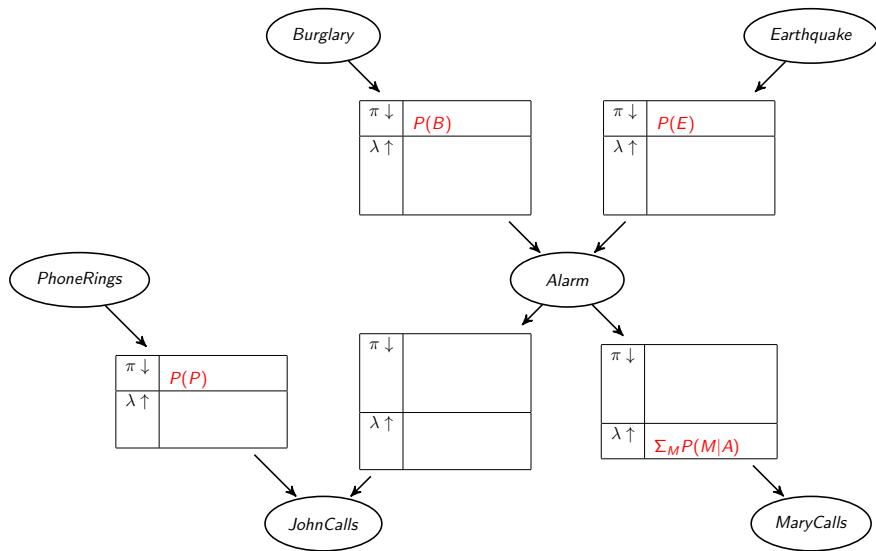
Earthquake Example



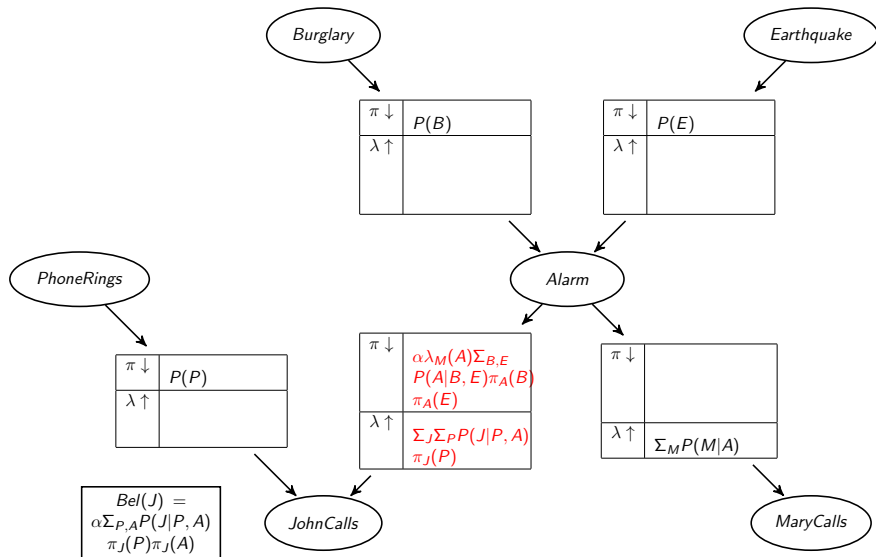
Message passing – No Evidence



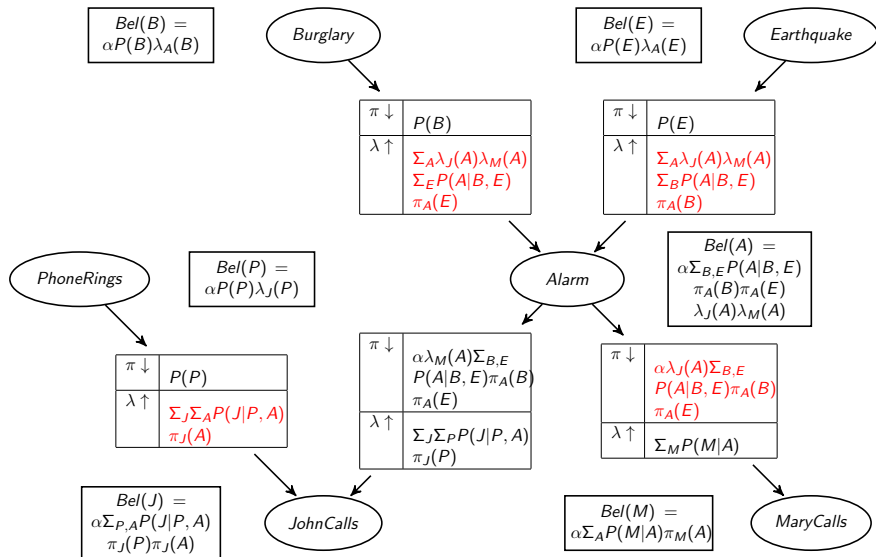
Message passing – No Evidence



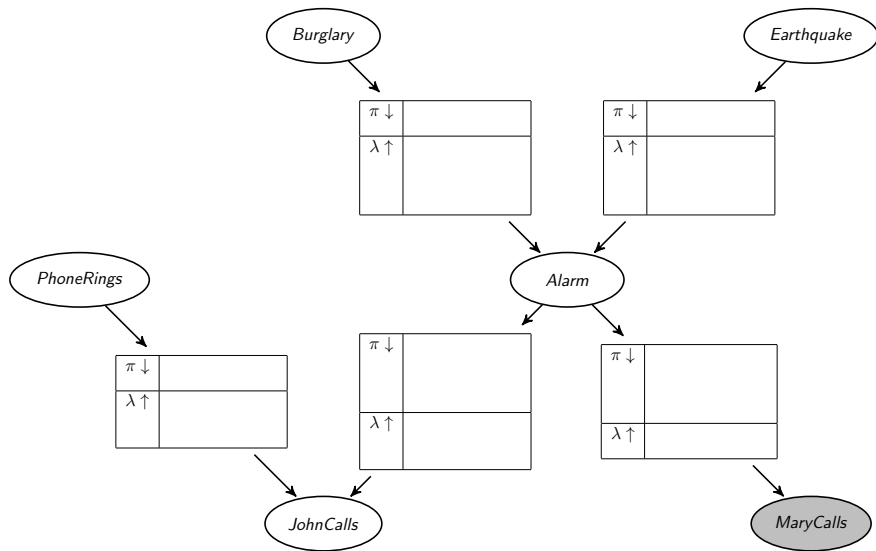
Message passing – No Evidence



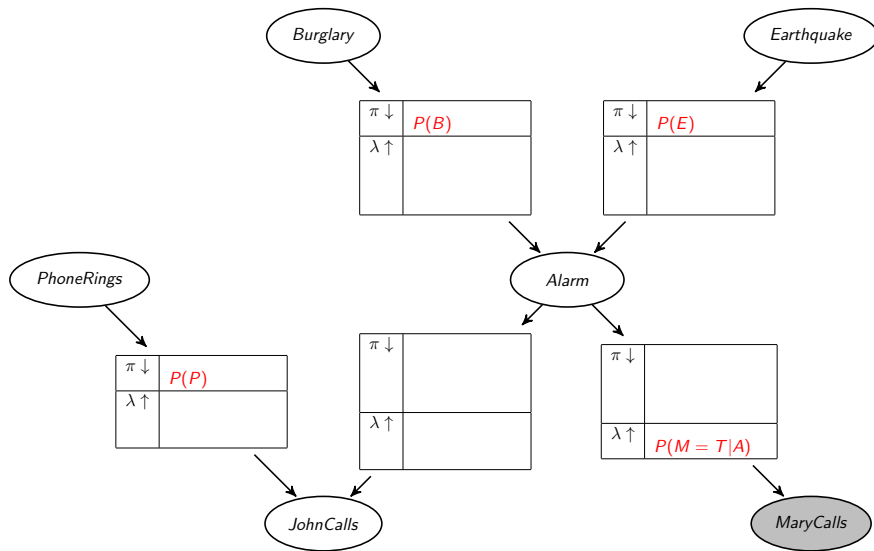
Message passing – No Evidence



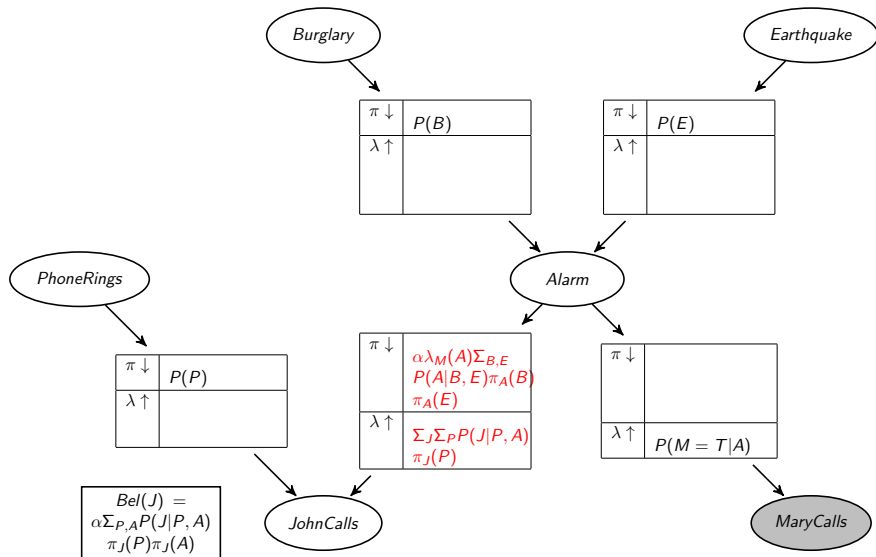
Message passing – Evidence Example



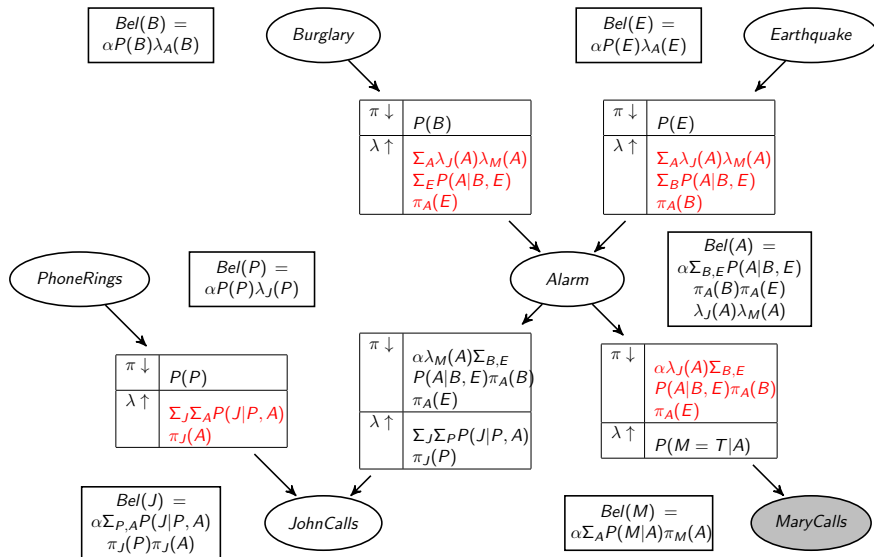
Message passing – Evidence Example



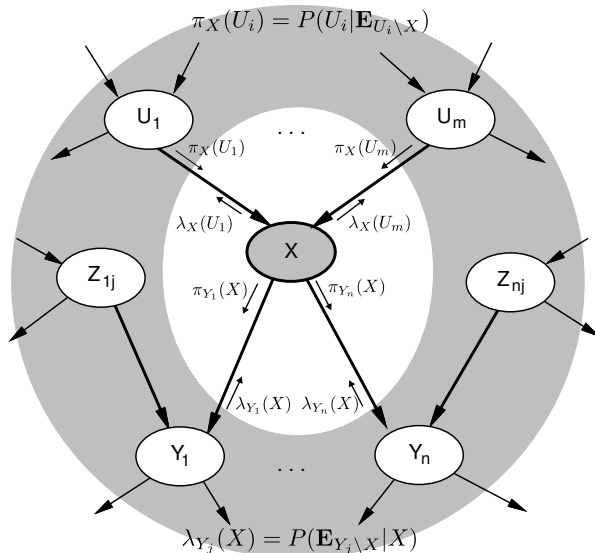
Message passing – Evidence Example



Message passing – Evidence Example



Overview of message passing



Message passing complexity

- Rounds of message passing: The maximum length of paths from X to roots (bounded by the number variables in the BN)
- Number of summations: It can be exponential in the number of parents and children of each node

Outline

- 1 Introduction
- 2 Exact inference by enumeration
- 3 Simple Bayesian Network Inferences
- 4 Exact inference in polytrees
- 5 Summary

Summary

- Probabilistic inference: compute the probability distribution for query variables, given evidence variables

$$P(X|\mathbf{E})$$

- BN Inference is very flexible: can enter evidence about any node and update beliefs in any other nodes.
- The speed of inference in practice depends on the structure of the network: how many loops; numbers of parents; location of evidence and query nodes.
- Inference methods
 - ▶ Exact inference by enumeration
 - ▶ Polytree message passing
 - ▶ Junction tree algorithms for general BNs (next lecture)
 - ▶ Approximation inference with stochastic simulation (next lecture)

Acknowledgments

Lecture 5 is composed the instructor's own understanding of the subject, and materials from [Korb and Nicholson, 2003, Chapter 3] and [Russell and Norvig, 2009, Chapter 16] with the instructor's own interpretations. The instructor takes full responsibility of any mistakes in the slides.

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