

FPOP with labels

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We have a sequence of d data points to segment, $\mathbf{z} \in \mathbb{R}^d$. We also have l labeled regions $R = \{(\underline{p}_1, \bar{p}_1, c_1), \dots, (\underline{p}_l, \bar{p}_l, c_l)\}$:

- \underline{p}_j is the start of a labeled region,
- \bar{p}_j is the end of a labeled region,
- $c_j \in \{0, 1\}$ is the number of changes in the region $[\underline{p}_j, \bar{p}_j]$.

For example $R = \{(\underline{p}_1 = 1, \bar{p}_1 = 2, c_1 = 0), (\underline{p}_2 = 2, \bar{p}_2 = 5, c_2 = 1)\}$ means that there is no change after the first data point, and there must be exactly one change somewhere between data points 2 and 5 (three possibilities). Assume the labeled regions are ordered:

$$1 \leq \underline{p}_1 < \bar{p}_1 \leq \underline{p}_2 < \bar{p}_2 \leq \dots \leq \underline{p}_l < \bar{p}_l \leq d \quad (1)$$

Thus the number of possible labels is $l \in \{0, 1, \dots, d-1\}$.

1 Optimization problem

We would like to find the best mean vector $\mathbf{m} \in \mathbb{R}^d$. Best is defined using the Optimal Partitioning objective function – a loss ℓ term for each data point, and a per-change penalty of $\lambda \geq 0$ [Maidstone et al., 2016]. We also add the constraints that the labels must be obeyed.

$$\begin{aligned} \hat{C}_t = \min_{\mathbf{m} \in \mathbb{R}^t} & \lambda \sum_{i=1}^{t-1} I(m_i \neq m_{i+1}) + \sum_{i=1}^d \ell(z_i, m_i) \\ \text{subject to} & \sum_{i=\underline{p}_j}^{\bar{p}_j-1} I(m_i \neq m_{i+1}) = c_j \text{ for all } j \in \{1, \dots, l\}. \end{aligned} \quad (2)$$

Note that I is the indicator function (1 if true, 0 otherwise).

SegAnnot is one algorithm that can be used in this context, but it is limited in that it can not detect any changes that are not labeled [Hocking and Rigai, 2012]. In fact SegAnnot computes a model with exactly $N_1(R) = \sum_{(\underline{p}_j, \bar{p}_j, c_j) \in R} I(c_j = 1)$ changes.

The labels cover $S(R) = \sum_{(\underline{p}_j, \bar{p}_j, c_j) \in R} \bar{p}_j - \underline{p}_j$ possible changes, in which there are exactly $N_1(R)$ changes. There are thus $d - S(R)$ positions which are not labeled, which may or may not have changes. Thus the number of possible changes is in $[N_1(R), d - S(R) + N_1(R)]$.

2 Algorithm

We propose a new algorithm, LabeledFPOP, which solves (2). We define $C_t(\mu)$ as the optimal cost up to data point t if there is a mean of μ on the last data point t . The initialization is $C_1(\mu) = \ell(z_1, \mu)$.

2.1 Usual FPOP

The usual FPOP algorithm does not have label constraints [Maidstone et al., 2016]. The dynamic programming update rule for $t > 1$ is

$$C_t(\mu) = \ell(z_t, \mu) + \min \begin{cases} \hat{C}_{t-1} + \lambda & \text{if there is a change between } t-1 \text{ and } t, \\ C_{t-1}(\mu) & \text{if there is no change.} \end{cases} \quad (3)$$

Note that $\hat{C}_{t-1} = \min_{\mu} C_{t-1}(\mu)$ is the optimal cost up to data point $t-1$.

For LabeledFPOP the update rule (3) is used if there is no label between data points $t-1$ and t . Indeed, if $R = \{\}$ then there are no labels at all, so the LabeledFPOP problem is the same as the usual FPOP.

2.2 Negative label

If there is a $c_j = 0$ label such that $t \in \{\underline{p}_j + 1, \dots, \bar{p}_j\}$, then the update rule is simpler: the only possibility is no change.

$$C_t(\mu) = \ell(z_t, \mu) + C_{t-1}(\mu). \quad (4)$$

2.3 Positive label at the start

For a $c_j = 1$ label the update rule is a bit more complicated. Take a simple example to start. For a $(\underline{p}_j = 1, \bar{p}_j = 3, c_j = 1)$ label, the cost up to the end of the label is

$$\min_{\mu_1, \mu_2} \begin{cases} \lambda + \ell(z_1, \mu_1) + \ell(z_2, \mu_2) + \ell(z_3, \mu_2) & \text{for a change after 1} \\ \lambda + \ell(z_1, \mu_1) + \ell(z_2, \mu_1) + \ell(z_3, \mu_2) & \text{for a change after 2} \end{cases} \quad (5)$$

In general for a $(\underline{p}_j = 1, \bar{p}_j, c_j = 1)$ label, the cost up to the end of the label is

$$\min_{\mu_1, \mu_2} \min_{t \in [1, \bar{p}_j - 1]} \lambda + \underbrace{\sum_{j=1}^t \ell(z_j, \mu_1)}_{\mathcal{L}_{1,t}(\mu_1)} + \underbrace{\sum_{j=t+1}^{\bar{p}_j} \ell(z_j, \mu_2)}_{\mathcal{L}_{t+1, \bar{p}_j}(\mu_2)}. \quad (6)$$

We can use the recursive update rule for $t \in \{2, \dots, \bar{p}_j - 1\}$:

$$\mathcal{L}_{1,t}(\mu) = \mathcal{L}_{1,t-1}(\mu) + \ell(z_t, \mu). \quad (7)$$

And then for the optimal cost we initialize $t = 2$ to the cost of a change between 1 and 2:

$$C_2(\mu) = \hat{\mathcal{L}}_{1,1} + \lambda + \ell(z_2, \mu). \quad (8)$$

And then the update rule for $t \in \{3, \dots, \bar{p}_j\}$ is:

$$C_t(\mu) = \ell(z_t, \mu) + \min \begin{cases} \hat{\mathcal{L}}_{1,t-1} + \lambda & \text{if there is a change between } t-1 \text{ and } t, \\ C_{t-1}(\mu) & \text{if the labeled change was before that.} \end{cases} \quad (9)$$

2.4 Positive label general case

In general for a $(\underline{p}_j, \bar{p}_j, c_j = 1)$ label, we need to compute the following quantities. First at $t = \underline{p}_j + 1$ we only need to consider one possibility (a change between $t-1 = \underline{p}_j$ and $t = \underline{p}_j + 1$) to compute the optimal cost function,

$$C_t(\mu) = \ell(z_t, \mu) + \hat{C}_{t-1} + \lambda. \quad (10)$$

If the label spans more than one possible changepoint ($\underline{p}_j + 1 < \bar{p}_j$) then for all $t \in \{\underline{p}_j + 1, \dots, \bar{p}_j - 1\}$ we also need to compute the cost of no change since the start of the label, and store it in a separate function,

$$V_t(\mu) = \ell(z_t, \mu) + \begin{cases} C_{t-1}(\mu) & \text{if } t = \underline{p}_j + 1, \\ V_{t-1}(\mu) & \text{if } t \in \{\underline{p}_j + 2, \dots, \bar{p}_j - 1\}. \end{cases} \quad (11)$$

And for all $t \in \{\underline{p}_j + 2, \dots, \bar{p}_j\}$ the optimal cost is computed by taking the minimum of two functions,

$$C_t(\mu) = \ell(z_t, \mu) + \min \begin{cases} \hat{V}_{t-1} + \lambda & \text{if there is a change between } t-1 \text{ and } t, \\ C_{t-1}(\mu) & \text{if the change was before that.} \end{cases} \quad (12)$$

2.5 Pseudocode

Algorithm 1 Labeled Functional Pruning Optimal Partitioning Algorithm.

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1: Input: data  $\mathbf{z} \in \mathbb{R}^d$ , non-negative penalty  $\lambda \in \mathbb{R}_+$ , sorted labels  $(\underline{p}_j, \bar{p}_j, c_j) \in R$ .
2: Output: optimal cost  $\hat{\mathbf{C}} \in \mathbb{R}^d$ , mean  $\mathbf{m} \in \mathbb{R}^d$ , previous segment end  $\mathbf{T} \in \{1, \dots, d\}^d$ .
3:  $j \leftarrow 1$  // label counter
4: labelType  $\leftarrow$  unlabeled
5: for  $t$  from 1 to  $d$ :
6:   if  $t == 1$ :
7:     cost  $\leftarrow 0$ 
8:   else if labelType == 0:
9:     cost  $\leftarrow$  prev_cost
10:  else if labelType == 1:
11:    change_cost  $\leftarrow$  MinPiece( $V$ ) +  $\lambda$ 
12:    if  $t == \underline{p}_j + 1$ :
13:      cost  $\leftarrow$  change_cost
14:    else:
15:      cost  $\leftarrow$  MinOfTwo(change_cost, prev_cost)
16:     $V \mathrel{+}=$  CostFun( $z_t$ )
17:  else if labelType == unlabeled:
18:    change_cost  $\leftarrow$  MinPiece(cost) +  $\lambda$ 
19:    cost  $\leftarrow$  MinOfTwo(change_cost, prev_cost)
20:  cost  $\mathrel{+}=$  CostFun( $z_t$ )
21:  if  $t == \bar{p}_j$ : // what update rule should we use next?
22:    labelType  $\leftarrow$  unlabeled
23:     $j \mathrel{++}$ 
24:  if  $t == \underline{p}_j$ :
25:    labelType  $\leftarrow c_j$ 
26:    if labelType == 1:
27:       $V \leftarrow$  cost
28:   $\hat{C}_t, m_t, T_t \leftarrow$  Minimize(cost)
29:  prev_cost  $\leftarrow$  cost

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References

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