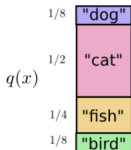
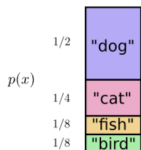


# Introduction to Machine Learning

## Cross-Entropy, KL and Source Coding

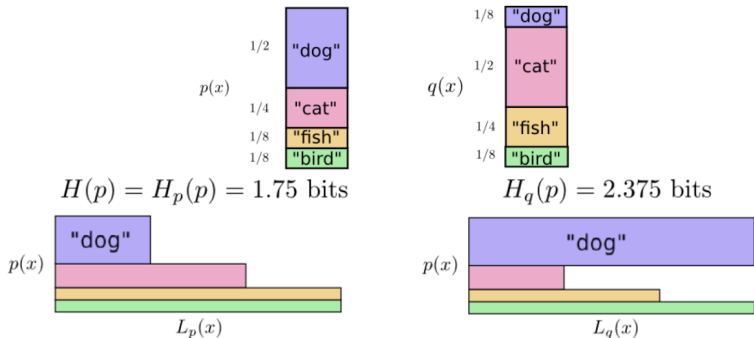


### Learning goals

- Know the cross-entropy
- Understand the connection between entropy, cross-entropy, and KL divergence

# CROSS-ENTROPY - DISCRETE CASE

- For a random source / distribution  $p$ , the minimal number of bits to optimally encode messages from is the entropy  $H(p)$ .
- If the optimal code for a different distribution  $q(x)$  is instead used to encode messages from  $p(x)$ , expected code length will grow.

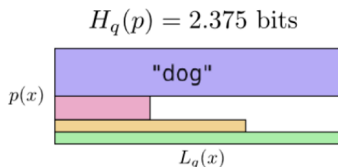
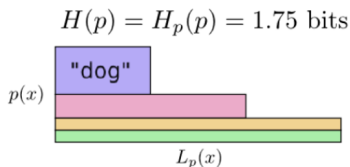


**Figure:**  $L_p(x)$ ,  $L_q(x)$  are the optimal code lengths for  $p(x)$  and  $q(x)$

# CROSS-ENTROPY - DISCRETE CASE

**Cross-entropy** is the average length of communicating an event from one distribution with the optimal code for another distribution (assume they have the same domain  $\mathcal{X}$  as in KL).

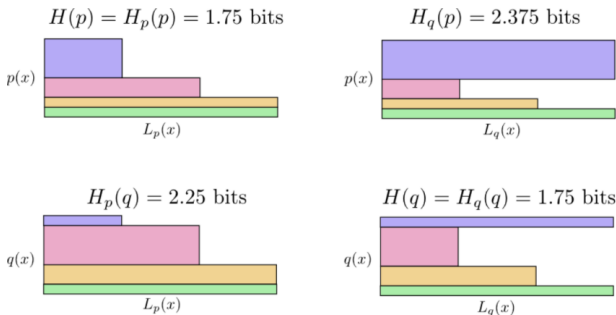
$$H_q(p) = \sum_{x \in \mathcal{X}} p(x) \log \left( \frac{1}{q(x)} \right) = - \sum_{x \in \mathcal{X}} p(x) \log (q(x))$$



**Figure:**  $L_p(x)$ ,  $L_q(x)$  are the optimal code lengths for  $p(x)$  and  $q(x)$

We directly see: cross-entropy of  $p$  with itself is entropy:  $H_p(p) = H(p)$ .

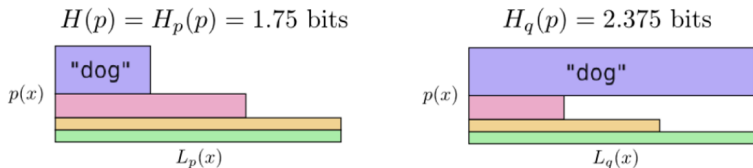
# CROSS-ENTROPY - DISCRETE CASE



Credit: Chris Olah

- In top,  $H_q(p)$  is greater than  $H(p)$  primarily because the blue event that is very likely under  $p$  has a very long codeword in  $q$ .
- Same, in bottom, for pink when we go from  $q$  to  $p$ .
- Note that  $H_q(p) \neq H_p(q)$ .

# CROSS-ENTROPY - DISCRETE CASE



**Figure:**  $L_p(x)$ ,  $L_q(x)$  are the optimal code lengths for  $p(x)$  and  $q(x)$

- Let  $x'$  denote the symbol "dog". The difference in code lengths is:

$$\log \left( \frac{1}{q(x')} \right) - \log \left( \frac{1}{p(x')} \right) = \log \frac{p(x')}{q(x')}$$

- If  $p(x') > q(x')$ , this is positive, if  $p(x') < q(x')$ , it is negative.
- The expected difference is KL, if we encode symbols from  $p$ :

$$D_{KL}(p||q) = \sum_{x \in \mathcal{X}} p(x) \cdot \log \frac{p(x)}{q(x)}$$

# CROSS-ENTROPY - DISCRETE CASE

- Entropy = Avg. nr. of bits if we optimally encode  $p$
- Cross-Entropy = Avg. nr. of bits if we suboptimally encode  $p$  with  $q$
- $DL_{KL}(p||q)$ : Difference in bits between the two

We can summarize this also through this identity:

$$H_q(p) = H(p) + D_{KL}(p||q)$$

This is because:

$$\begin{aligned} H(p) + D_{KL}(p||q) &= - \sum_{x \in \mathcal{X}} p(x) \log p(x) + \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)} \\ &= \sum_{x \in \mathcal{X}} p(x) (-\log p(x) + \log p(x) - \log q(x)) \\ &= - \sum_{x \in \mathcal{X}} p(x) \log q(x) = H_q(p) \end{aligned}$$

# CROSS-ENTROPY - CONTINUOUS CASE

For continuous density functions  $p(x)$  and  $q(x)$ :

$$H_p(q) = \int q(x) \ln \left( \frac{1}{p(x)} \right) dx = - \int q(x) \ln (p(x)) dx$$

- It is not symmetric.
- As for the discrete case,  $H_p(q) = h(q) + D_{KL}(q||p)$  holds.
- Can now become negative, as the  $h(q)$  can be negative!

# PROOF: MAXIMUM OF DIFFERENTIAL ENTROPY

**Claim:** For a given variance, the distribution that maximizes differential entropy is the Gaussian.

**Proof:** Let  $g(x)$  be a Gaussian with mean  $\mu$  and variance  $\sigma^2$  and  $f(x)$  an arbitrary density function with the same variance. Since differential entropy is translation invariant, we can assume  $f(x)$  and  $g(x)$  have the same mean.

The KL divergence (which is non-negative) between  $f(x)$  and  $g(x)$  is:

$$\begin{aligned} 0 \leq D_{KL}(f\|g) &= -h(f) + H_g(f) \\ &= -h(f) - \int_{-\infty}^{\infty} f(x) \ln(g(x)) dx \end{aligned} \tag{1}$$



# PROOF: MAXIMUM OF DIFFERENTIAL ENTROPY

The second term in (1) is,

$$\begin{aligned}\int_{-\infty}^{\infty} f(x) \log(g(x)) dx &= \int_{-\infty}^{\infty} f(x) \log \left( \frac{1}{\sqrt{2\pi}\sigma^2} e^{-\frac{(x-\mu)^2}{2\sigma^2}} \right) dx \\&= \int_{-\infty}^{\infty} f(x) \log \left( \frac{1}{\sqrt{2\pi}\sigma^2} \right) dx + \log(e) \int_{-\infty}^{\infty} f(x) \left( -\frac{(x-\mu)^2}{2\sigma^2} \right) dx \\&= -\frac{1}{2} \log(2\pi\sigma^2) - \log(e) \frac{\sigma^2}{2\sigma^2} = -\frac{1}{2} (\log(2\pi\sigma^2) + \log(e)) \\&= -\frac{1}{2} \log(2\pi e\sigma^2) = -h(g),\end{aligned}\tag{2}$$

where the last equality follows from the normal distribution example of the entropy chapter. Combining (1) and (2) results in

$$h(g) - h(f) \geq 0$$

with equality when  $f(x) = g(x)$  (following from the properties of Kullback-Leibler divergence).