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```
.bashrc
build()(
 q++ $0 -o $1.e -DLOC -std=c++11
      -Wall -Wextra -Wfatal-errors -Wshadow \
      -Wlogical-op -Wconversion -Wfloat-equal
b() ( build $@ -02 )
d() ( build $@ -fsanitize=address, undefined \
              -D_GLIBCXX_DEBUG -q )
run()($1 $2 && echo start >&2 && time ./$2.e)
100()(
  set -e; $1 $2; $1 $3
  for ((;;)) {
    ./$3.e > gen.in
    time ./$2.e < gen.in > gen.out
cmp()(
 set -e; $1 $2; $1 $3; $1 $4
  for ((;;)) {
                             echo -n 0
    ./$4.e > gen.in;
    ./$2.e < gen.in > p1.out; echo -n 1
    ./$3.e < gen.in > p2.out; echo -n 2
    diff pl.out p2.out;
                             echo -n Y
# Other flags:
# -Wformat=2 -Wshift-overflow=2 -Wcast-qual
# -Wcast-align -Wduplicated-cond
# -D_GLIBCXX_DEBUG_PEDANTIC -D_FORTIFY_SOURCE=2
# -fno-sanitize-recover -fstack-protector
.vimrc
se ai aw cin cul ic is nocp nohls nu sc scs
se bg=dark sw=4 ts=4 so=7 ttm=9
sy on
template.cpp
#include <bits/stdc++.h>
using namespace std;
using 11 = long long;
using Vi = vector<int>;
using Pii = pair<int,int>;
#define mp make_pair
#define pb push back
#define x first
#define y second
#define rep(i,b,e) for(int i=(b); i<(e); i++)
#define each(a,x) for(auto& a : (x))
#define all(x)
                   (x).begin(),(x).end()
#define sz(x)
                   int((x).size())
int main() {
  cin.sync with stdio(0); cin.tie(0);
  cout << fixed << setprecision(18);
  return 0:
```

```
// > Debug printer
#define tem template < class t, class u, class...w>
#define pri(x,y,z)tem auto operator<<(t&o,u a) \
 ->decltype(x,o) { o << z; return o << y; }
pri(a.print(), '}', '{'; a.print())
pri(a.y, ')', '(' << a.x << ", " << a.y)
pri(all(a), ']', '['; auto d="";
 for (auto i : a) (o << d << i, d = ", "))
void DD(...) {}
tem void DD(t s, u a, w... k) {
  for (int b=1; *s && *s - b*44; cerr << *s++)</pre>
   b += 2 / (*s*2 - 81);
  cerr << ": " << a << *s++; DD(s, k...);
#ifdef LOC
#define deb(...) (DD("[,\b :] "#__VA_ARGS___, \
 __LINE__, __VA_ARGS__), cerr << endl)
#else
#define deb(...)
#endif
#define DBP(...) void print() { \
 DD (#__VA_ARGS__, __VA_ARGS__); }
// > Utils
// Return smallest k such that 2^k > n
// Undefined for n = 0!
int uplg(int n) { return 32-__builtin_clz(n); }
int uplg(ll n) { return 64-__builtin_clzll(n); }
// Compare with certain epsilon (branchless)
// Returns -1 if a < b; 1 if a > b; 0 if equal
// a and b are assumed equal if |a-b| <= eps</pre>
int cmp(double a, double b, double eps=1e-10) {
 return (a > b+eps) - (a+eps < b);
various.h
// If math constants like M_PI are not found
// add this at the beginning of file
#define _USE_MATH_DEFINES
// Pragmas
#pragma GCC optimize("Ofast, unroll-loops,
                      no-stack-protector")
#pragma GCC target("popcnt,avx,tune=native")
while (clock() < duration*CLOCKS PER SEC)
// Automatically implement operators:
// 1. != if == is defined
// 2. >, <= and >= if < is defined
using namespace rel ops;
// Mersenne twister for randomization.
mt19937 64 rnd(chrono::steady clock::now()
  .time since epoch().count());
```

```
#include "segment2.h"
// To shuffle randomly use:
shuffle(all(vec), rnd)
// To pick random integer from [A:B] use:
uniform int_distribution<> dist(A, B);
int value = dist(rnd);
// To pick random real number from [A;B] use:
uniform real distribution <> dist(A, B);
double value = dist(rnd);
geometry/convex hull.h
#include "vec2.h"
// Translate points such that lower-left point
// is (0, 0). Returns old point location; O(n)
vec2 normPos(vector<vec2>€ points) {
 auto q = points[0].yxPair();
  each(p, points) q = min(q, p.yxPair());
  vec2 ret{q.y, q.x};
  each (p, points) p = p-ret;
 return ret:
// Find convex hull of points; time: O(n lq n)
// Points are returned counter-clockwise.
vector<vec2> convexHull(vector<vec2> points) {
 vec2 pivot = normPos(points);
  sort(all(points));
  vector<vec2> hull:
  each(p, points) {
    while (sz(hull) >= 2) {
      vec2 a = hull.back() - hull[sz(hull)-2];
      vec2 b = p - hull.back();
      if (a.cross(b) > 0) break;
      hull.pop_back();
    hull.pb(p);
  // Translate back, optional
  each (p, hull) p = p+pivot;
 return hull:
geometry/convex hull dist.h
#include "vec2.h"
// Check if p is inside convex polygon. Hull
// must be given in counter-clockwise order.
// Returns \overset{\circ}{2} if inside, 1 if on border,
// 0 if outside; time: O(n)
int insideHull(vector<vec2>& hull, vec2 p) {
 int ret = 1:
  rep(i, 0, sz(hull)) {
    auto v = hull[(i+1) sz(hull)] - hull[i];
    auto t = v.cross(p-hull[i]);
```

ret = min(ret, cmp(t, 0)); // For doubles

//ret = min(ret, (t>0) - (t<0)); // Ints

return int(max(ret+1, 0));

```
// Get distance from point to hull; time: O(n)
double hullDist(vector<vec2>& hull, vec2 p) {
 if (insideHull(hull, p)) return 0;
 double ret = 1e30:
 rep(i, 0, sz(hull)) {
   seg2 seg{hull[(i+1)%sz(hull)], hull[i]};
   ret = min(ret, seg.distTo(p));
 return ret:
// Compare distance from point to hull
// with sqrt(d2); time: O(n)
// -1 if smaller, 0 if equal, 1 if greater
int cmpHullDist(vector<vec2>& hull,
                vec2 p, 11 d2) {
 if (insideHull(hull,p)) return (d2<0)-(d2>0);
 int ret = 1;
 rep(i, 0, sz(hull)) {
   seg2 seg{hull[(i+1)%sz(hull)], hull[i]};
   ret = min(ret, seg.cmpDistTo(p, d2));
 return ret:
geometry/convex hull sum.h
#include "vec2.h"
// Get edge sequence for given polygon
// starting from lower-left vertex; time: O(n)
// Returns start position.
vec2 edgeSeg(vector<vec2> points,
            vector<vec2>& edges) {
 int i = 0, n = sz(points);
 rep(j, 0, n) {
   if (points[i].yxPair()>points[j].yxPair())
      i = j;
 rep(j, 0, n) edges.pb(points[(i+j+1)%n] -
                        points[(i+j)%n]);
 return points[i];
// Minkowski sum of given convex polygons.
// Vertices are required to be in
// counter-clockwise order; time: O(n+m)
vector<vec2> hullSum(vector<vec2> A,
                     vector<vec2> B) {
 vector\langle vec2 \rangle sum, e1, e2, es(sz(A) + sz(B));
 vec2 pivot = edgeSeg(A, e1) + edgeSeg(B, e2);
 merge(all(e1), all(e2), es.begin());
 sum.pb(pivot);
 each (e, es) sum.pb(sum.back() + e);
 sum.pop back();
 return sum:
```

geometry/halfplane intersect.h

// Intersect given halfplanes and output

#include "vec2.h"

#include "line2.h"

```
// hull vertices to out.
// Returns 0 if intersection is empty,
// 1 if intersection is non-empty and bounded.
// 2 if intersection is unbounded.
// Output vertices are valid ONLY IF
// intersection is non-empty and bounded.
// Works only with floating point vec2/line2.
// CURRENTLY DOESN'T WORK FOR NON-EMPTY
// AND UNBOUNDED CASES!
int intersectHalfPlanes(vector<line2> lines.
                        vector<vec2>& out) {
  deque<line2> H;
  out.clear():
  if (sz(lines) <= 1) return 2:</pre>
  sort(all(lines), [](line2 a, line2 b) {
    int t = cmp(a.norm.angle(),b.norm.angle());
    return t ? t < 0 : cmp(a.off*b.norm.len(),</pre>
     b.off*a.norm.len()) < 0;
  });
  auto bad = [](line2 a, line2 b, line2 c) {
    if (cmp(a.norm.cross(c.norm), 0) <= 0)</pre>
      return false:
    vec2 p; assert(a.intersect(c, p));
    return b.side(p) <= 0;</pre>
  each(e, lines) {
    if (!H.empty() &&
      !cmp(H.back().norm.angle(),
      e.norm.angle())) continue;
    while (sz(H) > 1 \&\& bad(H[sz(H)-2]),
     H.back(), e)) H.pop_back();
    while (sz(H) > 1 \& bad(e, H[0], H[1]))
     H.pop_front();
   H.pb(e);
  while (sz(H) > 2 \&\& bad(H[sz(H)-2]),
   H.back(), H[0])) H.pop_back();
  while (sz(H) > 2 && bad(H.back(),
   H[0], H[1])) H.pop_front();
  out.resize(sz(H));
  rep(i, 0, sz(H)) {
    auto a = H[i], b = H[(i+1)%sz(H)];
    if (a.norm.cross(b.norm) <= 0)</pre>
      return cmp (a.off*b.norm.len(),
        -b.off*a.norm.len()) <= 0 ? 0 : 2;
    assert(a.intersect(b, out[i]));
  rep(i, 0, sz(H)) {
    auto a = out[i], b = out[(i+1)%sz(H)];
    if (H[i].norm.perp().cross(b-a) <= 0)</pre>
      return 0;
  return 1:
geometry/line2.h
#include "vec2.h"
// 2D line/halfplane structure
```

```
// Parallel line through point
 static S parallel(P a, S b) {
   return { b.norm, b.norm.dot(a) };
 // Perpendicular line through point
 static S perp(P a, S b) {
   return { b.norm.perp(), b.norm.cross(a) };
 // Distance from point to line
 double distTo(P a) {
   return fabs(norm.dot(a)-off) / norm.len();
};
// Version for integer coordinates (long long)
struct line2i : bline2<11, vec2i, line2i> {
 line2i() : bline2{{}, 0} {}
 line2i(vec2i n, 11 c) : bline2{n, c} {}
 // Returns 0 if point a lies on the line,
 // 1 if on side where normal vector points,
 // -1 if on the other side.
 int side(vec2i a) {
   11 d = norm.dot(a);
   return (d > off) - (d < off);</pre>
// Version for double coordinates
// Requires cmp() from template
struct line2d : bline2<double, vec2d, line2d> {
 line2d() : bline2{{}, 0} {}
 line2d(vec2d n, double c) : bline2{n, c} {}
 // Returns 0 if point a lies on the line,
 // 1 if on side where normal vector points,
 // -1 if on the other side.
 int side(vec2d a) {
   return cmp(norm.dot(a), off);
 // Intersect this line with line a, returns
 // true on success (false if parallel).
 // Intersection point is saved to `out`.
 bool intersect (line2d a, vec2d& out) {
   double d = norm.cross(a.norm);
```

// PARTIALLY TESTED

DBP (norm, off);

struct bline2 {

// Base class of versions for ints and doubles

T off; // Offset (C parameter of equation)

template<class T, class P, class S>

// For lines: norm*point == off

P norm; // Normal vector [A; B]

static S through(P a, P b) {

// For halfplanes: norm*point <= off

// (i.e. normal vector points outside)

// Line through 2 points; normal vector

return { (a-b).perp(), a.cross(b) };

// points to the right of ab vector

```
if (cmp(d, 0) == 0) return false;
    out = (norm*a.off-a.norm*off).perp() / d;
    return true:
 1
};
using line2 = line2d;
geometry/rmst.h
#include "../structures/find union.h"
// Rectilinear Minimum Spanning Tree
// (MST in Manhattan metric); time: O(n lq n)
// Returns MST weight. Outputs spanning tree
// to G, vertex indices match point indices.
// Edge in G is pair (target, weight).
ll rmst (vector<Pii>€ points,
        vector<vector<Pii>>& G) {
  int n = sz(points);
  vector<pair<int, Pii>> edges:
  vector<Pii> close;
  Vi ord(n), merged(n);
  iota(all(ord), 0);
  function<void(int.int)> octant =
      [&] (int begin, int end) {
    if (begin+1 >= end) return;
    int mid = (begin+end) / 2;
    octant (begin, mid);
    octant (mid, end);
    int j = mid;
    Pii best = \{INT_MAX, -1\};
    merged.clear();
    rep(i, begin, mid) {
      int v = ord[i];
      Pii p = points[v];
      while (j < end) {</pre>
        int e = ord[j];
        Pii q = points[e];
        if (q.x-q.y > p.x-p.y) break;
        best = min(best, make_pair(q.x+q.y, e));
        merged.pb(e);
        j++;
      if (best.y != -1) {
        int alt = best.x-p.x-p.y;
        if (alt < close[v].x)</pre>
          close[v] = {alt, best.y};
      merged.pb(v);
    while (j < end) merged.pb(ord[j++]);</pre>
    copv(all(merged), ord.begin()+begin);
  rep(i, 0, 4) {
    rep(j, 0, 2) {
      sort(all(ord), [&](int 1, int r) {
        return points[1] < points[r];</pre>
      1);
```

```
close.assign(n, {INT MAX, -1});
      octant(0, n);
      rep(k, 0, n) {
       Pii p = close[k];
        if (p.y != -1) edges.pb(\{p.x, \{k, p.y\}\}\);
        points[k].x \star = -1;
    each (p, points) p = \{p.y, -p.x\};
 11 \text{ sum} = 0;
 FAU fau(n):
 sort(all(edges));
 G.assign(n, {});
  each (e, edges) if (fau.join(e.y.x, e.y.y)) {
    sum += e.x;
    G[e.y.x].pb({e.y.y, e.x});
    G[e.y.y].pb({e.y.x, e.x});
 return sum;
geometry/segment2.h
#include "vec2.h"
// 2D segment structure; NOT HEAVILY TESTED
// Base class of versions for ints and doubles
template < class P, class S> struct bseq2 {
 P a, b; // Endpoints
  // Distance from segment to point
 double distTo(P p) const {
    if ((p-a).dot(b-a) < 0) return (p-a).len();</pre>
    if ((p-b).dot(a-b) < 0) return (p-b).len();</pre>
    return double(abs((p-a).cross(b-a)))
                  / (b-a).len();
};
// Version for integer coordinates (long long)
struct seq2i : bseq2<vec2i, seq2i> {
 seq2i() {}
 seg2i(vec2i c, vec2i d) : bseg2{c, d} {}
  // Check if segment contains point p
 bool contains(vec2i p) {
    return (a-p).dot(b-p) <= 0 &&
           (a-p).cross(b-p) == 0;
  // Compare distance to p with sqrt(d2)
  // -1 if smaller, 0 if equal, 1 if greater
 int cmpDistTo(vec2i p, 11 d2) const {
    if ((p-a).dot(b-a) < 0) {
      11 1 = (p-a).len2();
      return (1 > d2) - (1 < d2);
    if ((p-b).dot(a-b) < 0) {
      11 1 = (p-b).len2();
      return (1 > d2) - (1 < d2);
```

11 c = abs((p-a).cross(b-a));

d2 = (b-a).len2();

```
return (c*c > d2) - (c*c < d2);
};
// Version for double coordinates
// Requires cmp() from template
struct seq2d : bseq2<vec2d, seq2d> {
  seq2d() {}
  seg2d(vec2d c, vec2d d) : bseg2{c, d} {}
  bool contains (vec2d p) {
    return cmp((a-p).dot(b-p), 0) <= 0 &&
           cmp((a-p).cross(b-p), 0) == 0;
};
using seg2 = seg2d;
geometry/vec2.h
// 2D point/vector structure; PARTIALLY TESTED
// Base class of versions for ints and doubles
template < class T, class S> struct bvec2 {
  T x, y;
  S operator+(S r) const {return{x+r.x,y+r.y};}
  S operator-(S r) const {return{x-r.x,y-r.y};}
  S operator*(T r) const { return {x*r, y*r}; }
  S operator/(T r) const { return {x/r, y/r}; }
  T dot(S r) const { return x*r.x+y*r.y; }
  T cross(S r) const { return x*r.y-y*r.x; }
               const { return x*x + y*y; }
  T len2()
  double len() const { return sqrt(len2()); }
  S perp()
               const { return {-y,x}; } //90deg
  pair<T, T> yxPair() const { return {y,x}; }
  double angle() const { //[0;2*PI] CCW from OX
   double a = atan2(y, x);
    return (a < 0 ? a+2*M_PI : a);</pre>
// Version for integer coordinates (long long)
struct vec2i : bvec2<11, vec2i> {
  vec2i() : bvec2{0, 0} {}
  vec2i(11 a, 11 b) : bvec2{a, b} {}
  bool operator==(vec2i r) const {
    return x == r.x && y == r.y;
  // Compare by angle, length if angles equal
  bool operator<(vec2i r) const {</pre>
    if (upper() != r.upper()) return upper();
    auto t = cross(r);
    return t > 0 || (!t && len2() < r.len2());</pre>
  bool upper() const {
    return y > 0 \mid \mid (y == 0 && x >= 0);
};
// Version for double coordinates
```

```
// Requires cmp() from template
struct vec2d : bvec2<double, vec2d> {
  vec2d() : bvec2{0, 0} {}
  vec2d(double a, double b) : bvec2{a, b} {}
  vec2d unit() const { return *this/len(); }
  vec2d rotate(double a) const { // CCW
    return {x*cos(a) - y*sin(a),
            x*sin(a) + y*cos(a);
  bool operator == (vec2d r) const {
    return !cmp(x, r.x) && !cmp(y, r.y);
  // Compare by angle, length if angles equal
  bool operator<(vec2d r) const {</pre>
    int t = cmp(angle(), r.angle());
    return t < 0 || (!t && len2() < r.len2());
};
using vec2 = vec2d;
graphs/2sat.h
// 2-SAT solver; time: O(n+m), space: O(n+m)
// Variables are indexed from 1 and
// negative indices represent negations!
// Usage: SAT2 sat(variable count);
// (add constraints...)
// bool solution_found = sat.solve();
// sat[i] = value of i-th variable, 0 or 1
            (also indexed from 1!)
// (internally: positive = i*2-1, neg. = i*2-2)
struct SAT2 : Vi {
  vector<Vi> G:
 Vi order, flags;
  // Init n variables, you can add more later
  SAT2 (int n = 0) : G(n*2) {}
  // Add new var and return its index
  int addVar() {
   G.resize(sz(G)+2); return sz(G)/2;
  // Add (i => j) constraint
  void imply(int i, int j) {
   i = max(i*2-1, -i*2-2);
    j = \max(j*2-1, -j*2-2);
    G[i].pb(j); G[j<sup>1</sup>].pb(i<sup>1</sup>);
  // Add (i v i) constraint
  void either(int i, int j) { imply(-i, j); }
  // Constraint at most one true variable
  void atMostOne(Vi& vars) {
    int x = addVar();
    each(i, vars) {
      int v = addVar();
      imply(x, y); imply(i, -x); imply(i, y);
      x = y;
    }
  }
```

```
// Solve and save assignments in `values`
 bool solve() { // O(n+m), Kosaraju is used
   assign(sz(G)/2+1, -1);
   flags.assign(sz(G), 0);
   rep(i, 0, sz(G)) dfs(i);
   while (!order.empty()) {
     if (!propag(order.back()^1, 1)) return 0;
     order.pop_back();
   return 1;
 void dfs(int i) {
   if (flags[i]) return:
   flags[i] = 1;
   each(e, G[i]) dfs(e);
   order.pb(i);
 bool propag(int i, bool first) {
   if (!flags[i]) return 1;
   flags[i] = 0;
   if (at(i/2+1) >= 0) return first;
   at (i/2+1) = i & 1;
   each(e, G[i]) if (!propag(e, 0)) return 0;
   return 1;
};
graphs/bellman inequalities.h
struct Ineq {
11 a, b, c; // a - b >= c
};
// Solve system of inequalities of form a-b>=c
// using Bellman-Ford; time: O(n*m)
bool solveIneq(vector<Ineq>& edges,
              vector<ll>& vars) {
 rep(i, 0, sz(vars)) each(e, edges)
   vars[e.b] = min(vars[e.b], vars[e.a]-e.c);
 each (e, edges)
   if (vars[e.a]-e.c < vars[e.b]) return 0;</pre>
 return 1;
graphs/bridges online.h
// Dynamic 2-edge connectivity queries
// Usage: Bridges bridges (vertex_count);
// - bridges.addEdge(u, v); - add edge (u, v)
// - bridges.cc[v] = connected component ID
// - bridges.bi(v) = 2-edge connected comp ID
struct Bridges {
 vector<Vi> G; // Spanning forest
 Vi cc, size, par, bp, seen;
 int cnt{0};
  // Initialize structure for n vertices; O(n)
 Bridges (int n = 0) : G(n), cc(n), size(n, 1),
                      par(n, -1), bp(n, -1),
                       seen(n) {
   iota(all(cc), 0);
 // Add edge (u, v); time: amortized O(lg n)
```

void addEdge(int u, int v) {

```
if (cc[u] == cc[v]) {
     int r = lca(u, v);
      while ((v = root(v)) != r)
       v = bp[bi(v)] = par[v];
      while ((u = root(u)) != r)
       u = bp[bi(u)] = par[u];
   } else {
     G[u].pb(v); G[v].pb(u);
     if (size[cc[u]] > size[cc[v]]) swap(u,v);
     size[cc[v]] += size[cc[u]];
     dfs(u, v);
   }
 }
 // Get 2-edge connected component ID
 int bi(int v) { // amortized time: < O(lq n)</pre>
   return bp[v] == -1 ? v : bp[v] = bi(bp[v]);
 int root(int v) {
   return par[v] == -1 || bi(par[v]) != bi(v)
     ? v : par[v] = root(par[v]);
 void dfs(int v, int p) {
   par[v] = p; cc[v] = cc[p];
   each(e, G[v]) if (e != p) dfs(e, v);
 int lca(int u, int v) { // Don't use this!
   for (cnt++;; swap(u, v)) if (u != -1) {
     if (seen[u = root(u)] == cnt) return u;
      seen[u] = cnt; u = par[u];
graphs/dense dfs.h
#include "../math/bit_matrix.h"
// DFS over adjacency matrix; time: O(n^2/64)
// G = graph, V = not visited vertices masks
// UNTESTED
struct DenseDFS {
 BitMatrix G, V; // space: O(n^2/64)
 DenseDFS(int n = 0) : G(n, n), V(1, n) {
   reset();
 void reset() { each(x, V.M) x = -1; }
 void setVisited(int i) { V.set(0, i, 0); }
 bool isVisited(int i) { return !V(0, i); }
 // DFS step: func is called on each unvisited
 // neighbour of i. You need to manually call
 // setVisited(child) to mark it visited.
 template < class T > // Single step: O(n/64)
 void step(int i, T func) {
   ull* E = G.row(i);
   for (int w = 0; w < G.stride;) {
     ull x = E[w] & V.row(0)[w];
     if (x) func((w<<6) | __builtin_ctzll(x));</pre>
     else w++;
```

};

graphs/edmonds_karp.h

```
using flow t = int:
constexpr flow t INF = 1e9+10;
// Edmonds-Karp algorithm for finding
// maximum flow in graph; time: O(V*E^2)
// NOT HEAVILY TESTED
struct MaxFlow {
  struct Edge {
   int dst. inv:
   flow t flow, cap;
  vector<vector<Edge>> G:
  vector<flow t> add:
  Vi prev;
  // Initialize for n vertices
  MaxFlow(int n = 0) : G(n) {}
  // Add new vertex
  int addVert() {
   G.emplace_back(); return sz(G)-1;
  // Add edge between u and v with capacity cap
  // and reverse capacity rcap
  void addEdge(int u, int v,
               flow_t cap, flow_t rcap = 0) {
   G[u].pb({ v, sz(G[v]), 0, cap });
   G[v].pb({ u, sz(G[u])-1, 0, rcap });
  // Compute maximum flow from src to dst.
  // Flow values can be found in edges,
  // vertices with `add` >= 0 belong to
  // cut component containing `s`.
  flow_t maxFlow(int src, int dst) {
   flow t f = 0;
    each (v, G) each (e, v) e.flow = 0;
    do {
     queue<int> 0;
     O.push (src);
     prev.assign(sz(G), -1);
      add.assign(sz(G), -1);
     add[src] = INF;
      while (!Q.empty()) {
       int i = Q.front();
        flow_t m = add[i];
       ()qoq.0;
       if (i == dst) {
          while (i != src) {
            auto& e = G[i][prev[i]];
           e.flow -= m;
           G[e.dst][e.inv].flow += m;
            i = e.dst:
          f += m;
          break;
```

```
each(e, G[i]) if (add[e.dst] < 0) {
          if (e.flow < e.cap) {</pre>
            O.push(e.dst);
            prev[e.dst] = e.inv;
           add[e.dst] = min(m, e.cap-e.flow);
       }
    } while (prev[dst] != -1);
   return f:
  // Get if v belongs to cut component with src
 bool cutSide(int v) {
   return add[v] >= 0;
};
graphs/gomory hu.h
#include "edmonds karp.h"
//#include "push_relabel.h" // if you need
struct Edge {
 int a, b; // vertices
 flow_t w; // weight
};
// Build Gomory-Hu tree; time: O(n*maxflow)
// Gomory-Hu tree encodes minimum cuts between
// all pairs of vertices: mincut for u and v
// is equal to minimum on path from u and v
// in Gomory-Hu tree. n is vertex count.
// Returns vector of Gomory-Hu tree edges.
vector<Edge> gomoryHu (vector<Edge>& edges,
                     int n) {
 MaxFlow flow(n):
 each (e, edges) flow.addEdge (e.a,e.b,e.w,e.w);
  vector<Edge> ret(n-1);
  rep(i, 1, n) ret[i-1] = \{i, 0, 0\};
 rep(i, 1, n) {
   ret[i-1].w = flow.maxFlow(i, ret[i-1].b);
   rep(i, i+1, n)
     if (ret[j-1].b == ret[i-1].b &&
          flow.cutSide(j)) ret[j-1].b = i;
 return ret:
graphs/matroids.h
// Find largest subset S of [n] such that
// S is independent in both matroid A and B.
// A and B are given by their oracles,
// see example implementations below.
// Returns vector V such that V[i] = 1 iff
// i-th element is included in found set;
// time: O(r^2*init + r^2*n*add).
// where r is max independent set,
// `init` is max time of oracles init
// and `add` is max time of oracles canAdd.
```

template < class T, class U>

vector<bool> intersectMatroids(T& A, U& B,

```
int n) {
 vector<bool> ans(n);
 bool ok = 1:
 A.init(ans);
 B.init(ans);
 rep(i, 0, n) if (A.canAdd(i) && B.canAdd(i))
   ans[i] = 1, A.init(ans), B.init(ans);
 while (ok) {
   vector<Vi> G(n);
   vector<bool> good(n);
   queue<int> que;
   Vi prev(n, -1):
   A.init(ans);
   B.init(ans);
   ok = 0;
   rep(i, 0, n) if (!ans[i]) {
     if (A.canAdd(i)) que.push(i), prev[i]=-2;
     good[i] = B.canAdd(i);
   rep(i, 0, n) if (ans[i]) {
     ans[i] = 0;
     A.init(ans);
     B.init(ans);
     rep(j, 0, n) if (i != j && !ans[j]) {
       if (A.canAdd(j)) G[i].pb(j);
        if (B.canAdd(j)) G[j].pb(i);
     ans[i] = 1;
   while (!que.empty()) {
     int i = que.front();
     que.pop();
     if (good[i]) {
        ans[i] = 1;
        while (prev[i] >= 0) {
         ans[i = prev[i]] = 0;
          ans[i = prev[i]] = 1;
        ok = 1:
       break;
     each(j, G[i]) if (prev[j] == -1)
        que.push(j), prev[j] = i;
 return ans;
// Matroid where each element has color
// and set is independent iff for each color c
// #{elements of color c} <= maxAllowed[c].</pre>
struct LimOracle {
 Vi color: // color[i] = color of i-th element
 Vi maxAllowed: // Limits for colors
 Vi tmp;
 // Init oracle for independent set S; O(n)
```

```
void init(vector<bool>& S) {
    tmp = maxAllowed;
    rep(i, 0, sz(S)) tmp[color[i]] -= S[i];
  // Check if S+{k} is independent; time: O(1)
 bool canAdd(int k) {
    return tmp[color[k]] > 0;
 }
};
// Graphic matroid - each element is edge,
// set is independent iff subgraph is acyclic.
struct GraphOracle {
 vector<Pii> elems; // Ground set: graph edges
 int n; // Number of vertices, indexed [0;n-1]
 Vi par:
 int find(int i) {
    return par[i] == -1 ? i
      : par[i] = find(par[i]);
  // Init oracle for independent set S; ~O(n)
 void init(vector<bool>& S) {
    par.assign(n, -1);
    rep(i, 0, sz(S)) if (S[i])
      par[find(elems[i].x)] = find(elems[i].y);
  // Check if S+{k} is independent; time: ~O(1)
 bool canAdd(int k) {
    return
      find(elems[k].x) != find(elems[k].y);
};
// Co-graphic matroid - each element is edge,
// set is independent iff after removing edges
// from graph number of connected components
// doesn't change.
struct CographOracle {
 vector<Pii> elems; // Ground set: graph edges
 int n: // Number of vertices, indexed [0:n-1]
 vector<Vi> G;
 Vi pre, low:
 int cnt;
 int dfs(int v, int p) {
    pre[v] = low[v] = ++cnt;
    each(e, G[v]) if (e != p)
      low[v] = min(low[v], pre[e] ?: dfs(e,v));
    return low[v];
  // Init oracle for independent set S; O(n)
 void init(vector<bool>& S) {
    G.assign(n, {});
    pre.assign(n, 0);
    low.resize(n);
    cnt = 0:
    rep(i, 0, sz(S)) if (!S[i]) {
     Pii e = elems[i]:
      G[e.x].pb(e.v);
      G[e.y].pb(e.x);
```

```
rep(v, 0, n) if (!pre[v]) dfs(v, -1);
  // Check if S+{k} is independent; time: O(1)
  bool canAdd(int k) {
    Pii e = elems[k]:
    return max(pre[e.x], pre[e.y])
      != max(low[e.x], low[e.y]);
};
// Matroid equivalent to linear space with XOR
struct XorOracle {
  vector<11> elems:
  vector<11> base:
  // Init for independent set S: O(n+r^2)
  void init(vector<bool>& S) {
   base.assign(63, 0);
    rep(i, 0, sz(S)) if (S[i]) {
     ll e = elems[i];
      rep(j, 0, sz(base)) if ((e >> j) & 1) {
       if (!base[j]) {
          base[j] = e;
          break:
        e ^= base[j];
  // Check if S+{k} is independent; time: O(r)
  bool canAdd(int k) {
   11 e = elems[k];
    rep(i, 0, sz(base)) if ((e >> i) & 1) {
     if (!base[i]) return 1;
      e ^= base[i];
    return 0;
};
graphs/push_relabel.h
using flow_t = int;
constexpr flow t INF = 1e9+10;
// Push-relabel algorithm with global relabel
// heuristic for finding maximum flow; O(V^3),
// but very fast in practice.
// Preflow is not converted to flow!
struct MaxFlow {
  struct Vert {
    int head{0}, cur{0}, label;
   flow t excess;
  struct Edge {
   int dst, nxt;
   flow t avail, cap;
  vector<Vert> V;
  vector<Edge> E;
  queue<int> que, bfs;
```

// Initialize for n vertices

```
MaxFlow(int n = 0) {
 V.assign(n, {});
  E.resize(2);
// Add new vertex
int addVert() {
  V.emplace back():
  return sz(V)-1;
}
// Add edge between u and v with capacity cap
// and reverse capacity rcap
void addEdge(int u, int v,
             flow_t cap, flow_t rcap = 0) {
  E.pb({ v, V[u].head, 0, cap });
  E.pb({ u, V[v].head, 0, rcap });
  V[u].head = sz(E)-2;
  V[v].head = sz(E)-1;
}
void push(int v, int e) {
  flow_t f = min(V[v].excess, E[e].avail);
  E[e].avail -= f;
  E[e^1].avail += f;
  V[v].excess -= f;
  if ((V[E[e].dst].excess += f) == f)
    que.push(E[e].dst);
// Compute maximum flow from src to dst
flow_t maxFlow(int src, int dst) {
  each (v, V) v.excess = v.label = v.cur = 0;
  each (e, E) e.avail = max(e.cap, flow t(0));
  int cnt, n = cnt = V[src].label = sz(V);
  V[src].excess = INF:
  for (int e = V[src].head; e; e = E[e].nxt)
   push (src, e);
  for (; !que.empty(); que.pop()) {
    if (cnt >= n/2) {
      each (v, V) v.label = n;
      V[dst1.label = 0;
      bfs.push(dst);
      cnt = 0;
      for (; !bfs.empty(); bfs.pop()) {
        auto& v = V[bfs.front()];
        for (int e=v.head; e; e = E[e].nxt) {
          int x = E[e].dst;
          if (E[e^1].avail &&
              V[x].label > v.label+1) {
            V[x].label = v.label+1;
            bfs.push(x);
     }
    int v = que.front(), &l = V[v].label;
    if (v == dst) continue:
    while (V[v].excess && 1 < n) {
     if (!V[v].cur) {
       1 = n;
```

```
for (int e=V[v].head; e; e=E[e].nxt){
            if (E[e].avail)
             l = min(l, V[E[e].dst].label+1);
          V[v1.cur = V[v1.head;
         cnt++;
        int e = V[v].cur;
        V[v].cur = E[e].nxt;
        if (E[e].avail &&
         l == V[E[e].dst].label+1) push(v, e);
   }
   return V[dst].excess;
 // Get if v belongs to cut component with src
 bool cutSide(int v) {
   return V[v].label >= sz(V);
};
graphs/scc.h
// Tarjan's SCC algorithm; time: O(n+m)
// Usage: SCC scc(graph);
// scc[v] = index of SCC for vertex v
// scc.comps[i] = vertices of i-th SCC
struct SCC : Vi {
 vector<Vi> comps;
 Vi S:
 int cnt{0};
 SCC() {}
 SCC(vector\langle Vi \rangle \& G) : Vi(sz(G),-1), S(sz(G)) {
   rep(i, 0, sz(G)) if (!S[i]) dfs(G, i);
 int dfs(vector<Vi>& G, int v) {
   int low = S[v] = ++cnt, t = -1;
   S.pb(v);
   each(e, G[v]) if (at(e) < 0)
     low = min(low, S[e] ?: dfs(G, e));
   if (low == S[v]) {
     comps.emplace back();
     for (; t != v; S.pop_back()) {
        at (t = S.back()) = sz(comps) - 1;
        comps.back().pb(t);
   }
   return low;
graphs/turbo matching.h
// Find maximum bipartite matching; time: ?
// G must be bipartite graph!
// Returns matching size (edge count).
```

// match[v] = vert matched to v or -1

int matching (vector < Vi>& G, Vi& match) {

```
vector<bool> seen;
  int n = 0, k = 1;
 match.assign(sz(G), -1);
  function<int(int)> dfs = [&](int i) {
    if (seen[i]) return 0;
    seen[i] = 1;
    each(e, G[il) {
      if (match[e] < 0 || dfs(match[e])) {</pre>
        match[i] = e; match[e] = i;
        return 1:
   return 0:
 };
  while (k) {
    seen.assign(sz(G), 0);
    k = 0:
    rep(i, 0, sz(G)) if (match[i] < 0)
      k += dfs(i);
    n += k;
 return n:
// Convert maximum matching to vertex cover
// time: O(n+m)
Vi vertexCover(vector<Vi>& G, Vi& match) {
 Vi ret, col(sz(G)), seen(sz(G));
  function<void(int, int)> dfs =
      [&] (int i, int c) {
    if (col[i]) return;
    col[i] = c+1;
    each(e, G[i]) dfs(e, !c);
  function<void(int)> aug = [&](int i) {
    if (seen[i] || col[i] != 1) return;
    seen[i] = 1;
    each(e, G[i]) seen[e] = 1, aug(match[e]);
 };
  rep(i, 0, sz(G)) dfs(i, 0);
  rep(i, 0, sz(G)) if (match[i] < 0) aug(i);
  rep(i, 0, sz(G))
    if (seen[i] == col[i]-1) ret.pb(i);
 return ret:
math/bit gauss.h
```

```
Vi col;
  ans.assign(m, 0);
  rep(i, 0, sz(A)) {
   int c = int(A[i]._Find_first());
   if (c >= m) {
     if (c == m) return 0;
     continue;
    rep(k, i+1, sz(A)) if (A[k][c]) A[k]^=A[i];
   swap(A[i], A[sz(col)]);
   col.pb(c);
  for (int i = sz(col); i--;) if (A[i][m]) {
   ans[col[i]] = 1;
   rep(k,0,i) if(A[k][col[i]]) A[k][m].flip();
  return sz(col) < m ? 2 : 1;</pre>
math/bit matrix.h
using ull = uint64 t;
// Matrix over Z 2 (bits and xor)
// UNTESTED and UNFINISHED
struct BitMatrix {
  vector<ull> M;
  int rows, cols, stride;
  BitMatrix(int n = 0, int m = 0) {
   rows = n: cols = m:
   stride = (m+63)/64;
   M.resize (n*stride);
  ull* row(int i) { return &M[i*stride]; }
  bool operator()(int i, int j) {
```

math/crt.h

};

if (val) w |= m;

else w &= ~m;

```
using Pll = pair<ll, 11>;
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 11 d = egcd(b%a, a, y, x);
  x = b/a*v;
  return d:
// Chinese Remainder Theoerem; time: O(lg lcm)
// Solves x = a.x \pmod{a.y}, x = b.x \pmod{b.y}
// Returns pair (x mod lcm, lcm(a.v, b.v))
// or (-1, -1) if there's no solution.
```

return (row(i)[j/64] >> (j%64)) & 1;

ull &w = row(i)[j/64], m = 1 << (j%64);

void set(int i, int j, bool val) {

```
// WARNING: a.x and b.x are assumed to be
// in [0;a.v) and [0;b.v) respectively.
// Works properly if lcm(a.y, b.y) < 2^63.
Pll crt(Pll a, Pll b) {
 if (a.v < b.v) swap(a, b);
  ll x, y, q = \operatorname{egcd}(a.y, b.y, x, y);
  11 c = b.x-a.x, d = b.y/g, p = a.y*d;
  if (c % g) return {-1, -1};
 11 s = (a.x + c/q*x % d * a.v) % p;
  return {s < 0 ? s+p : s, p};
math/discrete logarithm.h
```

```
#include "../modular.h"
// Baby-step giant-step algorithm; O(sqrt(p))
// Finds smallest x such that a^x = b \pmod{p}
// or returns -1 if there's no solution.
11 dlog(ll a, ll b, ll p) {
 int m = int(min(ll(sqrt(p))+2, p-1));
 unordered map<11, int> small;
 11 t = 1;
 rep(i, 0, m) {
   int& k = small[t];
   if (!k) k = i+1;
   t = t*a % p;
 t = modInv(t, p);
 rep(i, 0, m) {
   int j = small[b];
   if (j) return i*m + j - 1;
   b = b*t % p;
 return -1;
```

math/fft complex.h

```
using dbl = double;
using cmpl = complex<dbl>;
// Default std::complex multiplication is slow.
// You can use this to achieve small speedup.
cmpl operator*(cmpl a, cmpl b) {
 dbl ax = real(a), ay = imag(a);
 dbl bx = real(b), by = imag(b);
 return {ax*bx-ay*by, ax*by+ay*bx};
cmpl operator*=(cmpl& a,cmpl b) {return a=a*b;}
// Compute DFT over complex numbers; O(n lg n)
// Input size must be power of 2!
void fft(vector<cmpl>& a) {
 static vector<cmpl> w(2, 1);
 int n = sz(a);
  for (int k = sz(w); k < n; k *= 2) {
   w.resize(n);
   rep(i,0,k) w[k+i] = \exp(\text{cmpl}(0, M PI*i/k));
```

```
Vi rev(n):
  rep(i,0,n) rev[i] = (rev[i/2] | i\%2*n) / 2;
  rep(i,0,n) if(i<rev[i]) swap(a[i],a[rev[i]]);
  for (int k = 1; k < n; k *= 2) {
    for (int i=0; i < n; i += k*2) rep(j,0,k) {</pre>
      auto d = a[i+j+k] * w[j+k];
      a[i+j+k] = a[i+j] - d;
      a[i+j] += d;
   1
 }
// Convolve complex-valued a and b.
// store result in a; time: O(n lg n), 3x FFT
void convolve(vector<cmpl>& a, vector<cmpl> b) {
 int len = sz(a) + sz(b) - 1;
  int n = 1 \ll (32 - builtin clz(len));
  a.resize(n); b.resize(n);
  fft(a); fft(b);
  rep(i, 0, n) a[i] *= b[i] / dbl(n);
  reverse (a.begin()+1, a.end());
 fft(a);
 a.resize(len);
// Convolve real-valued a and b, returns result
// time: O(n lg n), 2x FFT
// Rounding to integers is safe as long as
// (max\_coeff^2)*n*log_2(n) < 9*10^14
// (in practice 10^16 or higher).
vector<dbl> convolve(vector<dbl>& a,
                     vector<dbl>& b) {
  int len = sz(a) + sz(b) - 1;
 int n = 1 << (32 - __builtin_clz(len));</pre>
  vector<cmpl> in(n), out(n);
  rep(i, 0, sz(a)) in[i].real(a[i]);
  rep(i, 0, sz(b)) in[i].imag(b[i]);
  fft(in);
  each(x, in) x \star = x;
  rep(i,0,n) out[i] = in[-i&(n-1)]-conj(in[i]);
  fft (out):
  vector<dbl> ret(len);
  rep(i, 0, len) ret[i] = imag(out[i]) / (n*4);
  return ret;
constexpr 11 MOD = 1e9+7;
// High precision convolution of integer-valued
// a and b mod MOD; time: O(n lq n), 4x FFT
// Input is expected to be in range [0; MOD)!
// Rounding is safe if MOD*n*log 2(n) < 9*10^14
// (in practice 10^16 or higher).
vector<ll> convMod(vector<ll>& a,
                   vector<11>& b) {
 vector<11> ret(sz(a) + sz(b) - 1);
  int n = 1 \ll (32 - builtin clz(sz(ret)));
 11 \text{ cut} = 11 (\text{sgrt}(MOD)) + 1;
  vector<cmpl> c(n), d(n), g(n), f(n);
  rep(i, 0, sz(a))
```

```
c[i] = {dbl(a[i]/cut), dbl(a[i]%cut)};
rep(i, 0, sz(b))
  d[i] = {dbl(b[i]/cut), dbl(b[i]%cut)};
fft(c); fft(d);
rep(i, 0, n) {
  int i = -i & (n-1);
  f[j] = (c[i]+conj(c[j])) * d[i] / (n*2.0);
  a[i] =
    (c[i]-conj(c[j])) * d[i] / cmpl(0, n*2);
fft(f); fft(a);
rep(i, 0, sz(ret)) {
  11 t = llround(real(f[i])) % MOD * cut;
  t += llround(imag(f[i]));
  t = (t + llround(real(g[i]))) % MOD * cut;
  t = (t + llround(imag(g[i]))) % MOD;
  ret[i] = (t < 0 ? t+MOD : t);
return ret:
```

math/fft mod.h

```
// Number Theoretic Tranform (NTT)
// For functions below you can choose 2 params:
// 1. M - prime modulus that MUST BE of form
         a*2^k+1, computation is done in Z_M
// 2. R - generator of Z M
// Modulus often seen on Codeforces:
// M = (119<<23)+1, R = 62; M is 998244353
// Parameters for 11 computation with CRT:
// M = (479 << 21) + 1, R = 62; M is > 10^9
// M = (483 << 21) + 1, R = 62; M is > 10^9
11 modPow(11 a, 11 e, 11 m) {
 11 t = 1 % m;
 while (e) {
   if (e % 2) t = t*a % m;
    e /= 2; a = a*a % m;
 return t;
// Compute DFT over Z_M with generator R.
// Input size must be power of 2; O(n lq n)
// Input is expected to be in range [0; MOD)!
// dit == true <=> inverse transform * 2^n
                   (without normalization)
template<11 M, 11 R, bool dit>
void ntt(vector<ll>& a) {
 static vector<11> w(2, 1);
 int n = sz(a):
  for (int k = sz(w); k < n; k *= 2) {
    w.resize(n, 1);
    11 c = modPow(R, M/2/k, M);
    if (dit) c = modPow(c, M-2, M);
    rep(i, k+1, k*2) w[i] = w[i-1]*c % M;
```

```
for (int t = 1; t < n; t *= 2) {
    int k = (dit ? t : n/t/2);
    for (int i=0; i < n; i += k*2) rep(j,0,k) {</pre>
     ll &c = a[i+j], &d = a[i+j+k];
     ll e = w[i+k], f = d;
     d = (dit ? c - (f=f*e%M) : (c-f)*e % M);
     if (d < 0) d += M;
     if ((c += f) >= M) c -= M;
// Convolve a and b mod M (R is generator).
// store result in a; time: O(n lg n), 3x NTT
// Input is expected to be in range [0; MOD)!
template<11 M = (119<<23)+1, 11 R = 62>
void convolve(vector<11>& a, vector<11> b) {
  int len = sz(a) + sz(b) - 1;
  int n = 1 << (32 - __builtin_clz(len));</pre>
  ll t = modPow(n, M-2, M);
  a.resize(n); b.resize(n);
  ntt < M, R, 0 > (a); ntt < M, R, 0 > (b);
  rep(i, 0, n) a[i] = a[i] *b[i] % M * t % M;
  ntt<M,R,1>(a);
  a.resize(len);
ll egcd(ll a, ll b, ll& x, ll& y) {
  if (!a) return x=0, y=1, b;
  ll d = egcd(b%a, a, y, x);
  x = b/a*v;
  return d;
// Convolve a and b with 64-bit output,
// store result in a; time: O(n lg n), 6x NTT
// Input is expected to be non-negative!
void convLong(vector<11>& a, vector<11> b) {
  const 11 M1 = (479 << 21) +1, M2 = (483 << 21) +1;
  const 11 MOD = M1*M2, R = 62;
  vector<11> c = a, d = b;
  each(k, a) k %= M1; each(k, b) k %= M1;
  each(k, c) k %= M2; each(k, d) k %= M2;
  convolve<M1, R>(a, b);
  convolve<M2, R>(c, d);
  ll x, y; egcd(M1, M2, x, y);
  rep(i, 0, sz(a)) {
   a[i] += (c[i]-a[i]) *x % M2 * M1;
    if ((a[i] %= MOD) < 0) a[i] += MOD;</pre>
```

math/fwht.h

```
// Fast Walsh-Hadamard Transform; O(n lg n)
// Input must be power of 2!
// Uncommented version is for XOR.
// OR version is equivalent to sum-over-subsets
// (Zeta transform, inverse is Moebius).
// AND version is same as sum-over-supersets.
// TESTED ON RANDS
template < bool inv, class T>
void fwht(vector<T>& b) {
```

```
for (int s = 1; s < sz(b); s *= 2) {
    for (int i = 0; i < sz(b); i += s*2) {
      rep(j, i, i+s) {
        auto &x = b[j], &y = b[j+s];
        tie(x, y) =
          mp(x+y, x-y);
        // inv ? mp(x-y, y) : mp(x+y, y); //AND
        // inv ? mp(x, y-x) : mp(x, x+y); //OR
   }
 }
  // ONLY FOR XOR:
 if (inv) each (e, b) e /= sz(b);
// Compute convolution of a and b such that
// ans[i#i] += a[i]*b[i], where # is OR, AND
// or XOR, depending on FWHT version.
// Stores result in a; time: O(n lg n)
// Both arrays must be of same size = 2^n!
template<class T>
void bitConv(vector<T>& a, vector<T> b) {
 fwht<0>(a):
 fwht<0>(b):
 rep(i, 0, sz(a)) a[i] *= b[i];
 fwht<1>(a);
math/gauss.h
        columns are variables.
        m-th column is equation result
// - ans - output for variables values
// - m - variable count
// 2 if more than 1 solution exist.
int gauss (vector<vector<double>>& A,
          vector<double>& ans, int m) {
 Vi col:
 ans.assign(m, 0);
```

```
// Solve system of linear equations; O(n^2*m)
// - A - extended matrix, rows are equations,
        (A[i][j] - i-th row and j-th column)
// Returns 0 if no solutions found, 1 if one,
 rep(i, 0, sz(A)) {
   int c = 0;
   while (c <= m && !cmp(A[i][c], 0)) c++;</pre>
    // For Zp:
    //while (c <= m && !A[i][c].x) c++;
    if (c >= m) {
     if (c == m) return 0;
      continue;
    rep(k, i+1, sz(A)) {
      auto mult = A[k][c] / A[i][c];
     rep(j, 0, m+1) A[k][j] -= A[i][j]*mult;
    swap(A[i], A[sz(col)]);
   col.pb(c);
  for (int i = sz(col); i--;) {
    ans[col[i]] = A[i][m] / A[i][col[i]];
```

```
rep(k, 0, i)
     A[k][m] = ans[col[i]] * A[k][col[i]];
 return sz(col) < m ? 2 : 1;
math/miller rabin.h
```

```
#include "modular64.h"
// Miller-Rabin primality test
// time O(k*lg^2 n), where k = number of bases
// Deterministic for p <= 10^9
// constexpr 11 BASES[] = {
// 336781006125, 9639812373923155
// 1:
// Deterministic for p <= 2^64
constexpr 11 BASES[] = {
 2,325,9375,28178,450775,9780504,1795265022
bool isPrime(ll p) {
 if (p == 2) return true;
 if (p <= 1 || p%2 == 0) return false;</pre>
 11 d = p-1, times = 0;
 while (d%2 == 0) d /= 2, times++;
 each(a, BASES) if (a%p) {
   // 11 a = rand() % (p-1) + 1;
   ll b = modPow(a%p, d, p);
   if (b == 1 || b == p-1) continue;
   rep(i, 1, times) {
     b = modMul(b, b, p);
     if (b == p-1) break;
   if (b != p-1) return false;
 return true;
```

math/modinv precompute.h

```
constexpr 11 MOD = 234567899;
vector<11> modInv(MOD); // You can lower size
// Precompute modular inverses; time: O(MOD)
void initModInv() {
 modInv[1] = 1;
 rep(i, 2, sz(modInv)) modInv[i] =
   (MOD - (MOD/i) ★ modInv[MOD%i]) % MOD;
```

math/modular.h

```
// Big prime number, about 2*10^9
constexpr int MOD = 15*(1<<27)+1;
ll modInv(ll a, ll m) { // a^{(-1)} \mod m
 if (a == 1) return 1;
 return ((a - modInv(m%a, a)) *m + 1) / a;
```

```
11 modPow(11 a, 11 e, 11 m) { // a^e mod m
 11 t = 1 % m;
 while (e) {
    if (e % 2) t = t*a % m;
    e /= 2; a = a*a % m;
 return t;
// Wrapper for modular arithmetic
struct Zp {
 11 x; // Contained value, in range [0; MOD-1]
 Zp() : x(0) {}
 Zp(11 a) : x(a\%MOD) { if (x < 0) x += MOD; }
  #define OP(c,d) Zp& operator c##=(Zp r) { }
     x = x d; return *this; } 
    Zp operator c(Zp r) const { \
      Zp t = *this; return t c##= r; }
 OP(+, +r.x - MOD*(x+r.x >= MOD));
 OP(-, -r.x + MOD*(x-r.x < 0));
 OP(*, *r.x % MOD);
 OP(/, *r.inv().x % MOD);
  // For composite modulus use modInv, not pow
 Zp inv() const { return pow(MOD-2); }
 Zp pow(11 e) const{ return modPow(x,e,MOD); }
 void print() { cerr << x; } // For deb()</pre>
// Extended Euclidean Algorithm
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 ll d = egcd(b%a, a, y, x);
 x = b/a*y;
 return d;
```

math/modular64.h

```
// Modular arithmetic for modulus < 2<sup>62</sup>
11 modAdd(11 x, 11 y, 11 m) {
 x += v;
 return x < m ? x : x-m;</pre>
11 modSub(11 x, 11 y, 11 m) {
 return x >= 0 ? x : x+m;
// About 4x slower than normal modulo
11 modMul(11 a, 11 b, 11 m) {
 11 c = 11((long double) a * b / m);
 ll r = (a*b - c*m) % m;
 return r < 0 ? r+m : r;
11 modPow(11 x, 11 e, 11 m) {
 11 t = 1:
 while (e) {
   e >>= 1:
```

```
math/montgomery.h
#include "modular.h"
// Montgomery modular multiplication
// MOD < MG_MULT, gcd(MG_MULT, MOD) must be 1
// Don't use if modulo is constexpr; UNTESTED
constexpr 11 MG SHIFT = 32;
constexpr ll MG MULT = 1LL << MG SHIFT;</pre>
constexpr ll MG_MASK = MG_MULT - 1;
const 11 MG_INV = MG_MULT-modInv(MOD, MG_MULT);
// Convert to Montgomery form
11 MG(11 x) { return (x*MG_MULT) % MOD; }
// Montgomery reduction
// redc(mg * mg) = Montgomery-form product
ll redc(ll x) {
  11 q = (x * MG_INV) & MG_MASK;
  x = (x + q*MOD) >> MG_SHIFT;
  return (x \geq= MOD ? x-MOD : x);
math/nimber.h
// Nimbers are defined as sizes of Nim heaps.
// Operations on nimbers are defined as:
// a+b = mex({a'+b : a' < a} u {a+b' : b' < b})
// ab = mex(\{a'b+ab'+a'b' : a' < a, b' < b\})
// Nimbers smaller than M = 2^2k form a field.
// Addition is equivalent to xor, meanwhile
// multiplication can be evaluated
// in O(lg^2 M) after precomputing.
using ull = uint64_t;
ull nbuf[64][64]; // Nim-products for 2^i * 2^j
// Multiply nimbers; time: O(lq^2 M)
// WARNING: Call initNimMul() before using.
ull nimMul(ull a, ull b) {
  ull ret = 0;
  for (ull s = a; s; s &= (s-1))
    for (ull t = b; t; t &= (t-1))
      ret ^= nbuf[__builtin_ctzll(s)]
                 [__builtin_ctzll(t)];
  return ret;
// Initialize nim-products lookup table
void initNimMul() {
  rep(i, 0, 64)
    nbuf[i][0] = nbuf[0][i] = 1ull << i;
  rep(b, 1, 64) rep(a, 1, b+1) {
   int i = 1 << (63 - __builtin_clzll(a));
int j = 1 << (63 - __builtin_clzll(b));</pre>
    ull t = nbuf[a-i][b-j];
    if (i < j)
     t = nimMul(t, 1ull << i) << j;
     t = nimMul(t, 1ull << (i-1)) ^ (t << i);
    nbuf[a][b] = nbuf[b][a] = t;
```

x = modMul(x, x, m);

return t:

```
}
}
// Compute a^e under nim arithmetic; O(1g^3 M)
// WARNING: Call initNimMul() before using.
ull nimPow(ull a, ull e) {
 ull t = 1;
 while (e) {
   if (e % 2) t = nimMul(t, a);
   e \neq 2; a = nimMul(a, a);
 return t;
// Compute inverse of a in 2^64 nim-field;
// time: 0(1a^3 M)
// WARNING: Call initNimMul() before using.
ull nimInv(ull a) {
 return nimPow(a, ull(-2));
// If you need to multiply many nimbers by
// the same value you can use this to speedup.
// WARNING: Call initNimMul() before using.
struct NimMult {
 ull M[64] = \{0\};
  // Initialize lookup; time: O(lg^2 M)
 NimMult(ull a) {
   for (ull t=a; t; t &= (t-1)) rep(i, 0, 64)
     M[i] ^= nbuf[__builtin_ctzll(t)][i];
  // Multiply by b; time: O(lq M)
 ull operator()(ull b) {
   ull ret = 0;
    for (ull t = b; t; t \&= (t-1))
     ret ^= M[__builtin_ctzll(t)];
    return ret;
math/phi large.h
#include "pollard rho.h"
// Compute Euler's totient of large numbers
// time: O(n^(1/4)) \leftarrow factorization
11 phi(ll n) {
 each (p, factorize (n)) n = n / p.x * (p.x-1);
 return n;
math/phi precompute.h
Vi phi(1e7+1);
// Precompute Euler's totients: time: O(n lg n)
void calcPhi() {
 iota(all(phi), 0);
 rep(i, 2, sz(phi)) if (phi[i] == i)
   for (int j = i; j < sz(phi); j += i)</pre>
     phi[j] = phi[j] / i * (i-1);
math/phi prefix sum.h
```

```
#include "phi precompute.h"
vector<11> phiSum: // [k] = sum \ from \ 0 \ to \ k-1
// Precompute Euler's totient prefix sums
// for small values; time: O(n lg n)
void calcPhiSum() {
 calcPhi():
 phiSum.resize(sz(phi)+1);
 rep(i, 0, sz(phi))
   phiSum[i+1] = phiSum[i] + phi[i];
// Get prefix sum of phi(0) + ... + phi(n-1).
// WARNING: Call calcPhiSum first!
// For n > 4*10^9, answer will overflow.
// If you wish to get answer mod M use
// commented lines.
ll getPhiSum(ll n) { // time: O(n^{(2/3)})
 static unordered_map<11, 11> big;
  if (n < sz(phiSum)) return phiSum[n];</pre>
  if (big.count(--n)) return big[n];
  ll ret = n*(n+1)/2;
  // 11 ret = (n \cdot 2 \cdot 2 \cdot n \cdot M) * ((n+1)/2 \cdot 8 \cdot M)
                   : n/2%M * (n%M+1)) % M;
  for (11 s, i = 2; i \le n; i = s+1) {
   s = n / (n/i);
   ret -= (s-i+1) * getPhiSum(n/i+1);
    // ret -= (s-i+1) %M * getPhiSum(n/i+1) % M;
  // ret = (ret%M + M) % M;
 return big[n] = ret;
math/pi large precomp.h
#include "sieve.h"
// Count primes in given interval
// using precomputed table.
// Set MAX P to sgrt (MAX N) and run sieve()!
// Precomputed table will contain N_BUCKETS
// elements - check source size limit.
constexpr ll MAX_N = 1e11+1;
constexpr ll N BUCKETS = 10000;
constexpr 11 BUCKET_SIZE = (MAX_N/N_BUCKETS)+1;
constexpr ll precomputed[] = {/* ... */};
11 sieveRange(11 from, 11 to) {
 bitset < BUCKET SIZE > elems:
  from = max(from, 2LL);
  to = max(from, to);
  each (p, primesList) {
    ll c = max((from+p-1) / p, 2LL);
    for (11 i = c*p; i < to; i += p)
      elems.set(i-from);
 return to-from-elems.count();
// Run once on local computer to precompute
// table. Takes about 10 minutes for n = 1e11.
// Sanity check (for default params):
```

```
// 664579, 606028, 587253, 575795, ...
void localPrecompute() {
 for (11 i = 0; i < MAX N; i += BUCKET SIZE) {
    11 to = min(i+BUCKET SIZE, MAX N);
    cout << sieveRange(i, to) << ',' << flush;</pre>
 cout << endl;
// Count primes in [from; to) using table.
// O(N BUCKETS + BUCKET_SIZE*lq lq n + sqrt(n))
11 countPrimes(11 from, 11 to) {
 ll bFrom = from/BUCKET SIZE+1,
     bTo = to/BUCKET SIZE:
 if (bFrom > bTo) return sieveRange(from, to);
 11 ret = accumulate (precomputed+bFrom,
                      precomputed+bTo, 0);
  ret += sieveRange(from, bFrom*BUCKET SIZE);
 ret += sieveRange(bTo*BUCKET_SIZE, to);
 return ret:
math/pollard rho.h
#include "modular64.h"
#include "miller rabin.h"
using Factor = pair<11, int>;
void rho(vector<11>& out, 11 n) {
 if (n <= 1) return;
 if (isPrime(n)) out.pb(n);
  else if (n^2 == 0) rho(out,2), rho(out,n/2);
  else for (ll a = 2;; a++) {
   11 x = 2, y = 2, d = 1;
    while (d == 1) {
     x = modAdd(modMul(x, x, n), a, n);
      y = modAdd (modMul(y, y, n), a, n);
      y = modAdd(modMul(y, y, n), a, n);
      d = gcd(abs(x-y), n);
    if (d != n) {
      rho(out, d);
      rho(out, n/d);
      return;
 }
// Pollard's rho factorization algorithm
// Las Vegas version; time: n^(1/4)
// Returns pairs (prime, power), sorted
vector<Factor> factorize(ll n) {
 vector<Factor> ret;
 vector<ll> raw:
 rho(raw, n);
  sort(all(raw));
 each(f, raw) {
   if (ret.emptv() || ret.back().x != f)
      ret.pb({ f, 1 });
      ret.back().v++;
 return ret;
math/polynomial interp.h
```

8

```
// Interpolates set of points (i, vec[i])
// and returns it evaluated at x; time: O(n^2)
// TODO: Improve to linear time
template<typename T>
T polvExtend(vector<T>& vec, T x) {
 T ret = 0:
  rep(i, 0, sz(vec)) {
   T = vec[i], b = 1;
    rep(j, 0, sz(vec)) if (i != j) {
     a *= x-j; b *= i-j;
   ret += a/b;
  return ret:
math/sieve.h
constexpr int MAX P = 1e6:
bitset<MAX_P+1> primes;
Vi primesList:
// Erathostenes sieve; time: O(n lg lg n)
void sieve() {
  primes.set():
  primes.reset(0):
  primes.reset(1);
  for (int i = 2; i*i <= MAX_P; i++)</pre>
    if (primes[i])
      for (int j = i*i; j <= MAX_P; j += i)</pre>
       primes.reset(i);
  rep(i, 0, MAX_P+1) if (primes[i])
    primesList.pb(i);
math/sieve factors.h
constexpr int MAX P = 1e6;
Vi factor (MAX_P+1);
// Erathostenes sieve with saving smallest
// factor for each number; time: O(n lg lg n)
void sieve() {
  for (int i = 2; i*i <= MAX_P; i++)</pre>
    if (!factor[i])
      for (int j = i*i; j <= MAX_P; j += i)</pre>
       if (!factor[j])
          factor[i] = i;
  rep(i, 0, MAX_P+1) if (!factor[i]) factor[i]=i;
// Factorize n <= MAX_P; time: O(lq n)</pre>
// Returns pairs (prime, power), sorted
vector<Pii> factorize(ll n) {
  vector<Pii> ret;
  while (n > 1) {
    int f = factor[n];
    if (ret.emptv() | ret.back().x != f)
      ret.pb({ f, 1 });
      ret.back().v++;
   n /= f;
  return ret;
```

```
math/sieve segmented.h
constexpr int MAX P = 1e9;
bitset<MAX_P/2+1> primes; // Only odd numbers
// Cache-friendly Erathostenes sieve
// ~1.5s on Intel Core i5 for MAX_P = 10^9
// Memory usage: MAX P/16 bytes
void sieve() {
  constexpr int SEG SIZE = 1<<18;</pre>
  int pSqrt = int(sqrt(MAX_P)+0.5);
  vector<Pii> dels;
  primes.set():
  primes.reset(0);
  for (int i = 3; i <= pSqrt; i += 2) {</pre>
   if (primes[i/2]) {
      for (j = i*i; j <= pSqrt; j += i*2)</pre>
        primes.reset(j/2);
      dels.pb(\{i, i/2\});
   }
 }
  for (int seg = pSgrt/2;
       seg <= sz(primes); seg += SEG SIZE) {
    int lim = min(seg+SEG SIZE, sz(primes));
    each(d, dels) for (;d.y < lim; d.y += d.x)
      primes.reset(d.v);
}
bool isPrime(int x) {
  return x == 2 \mid \mid (x \cdot 2 \cdot \xi \cdot primes[x/2]);
structures/bitset plus.h
// Undocumented std::bitset features:
// - Find first() - returns first bit = 1 or N
// - _Find_next(i) - returns first bit = 1
                     after i-th bit
                     or N if not found
// Bitwise operations for vector<bool>
// UNTESTED
#define OP(x) vector<bool>& operator x##=(
    vector<bool>& 1, const vector<bool>& r) { \
  assert(sz(1) == sz(r));
  auto a = 1.begin(); auto b = r.begin();
  while (a<1.end()) *a._M_p++ x##= *b._M_p++; \
  return 1; }
OP (&) OP (|) OP (^)
structures/fenwick tree.h
// Fenwick tree (BIT tree); space: O(n)
// Default version: prefix sums
struct Fenwick {
 using T = int;
  static const T ID = 0;
 T f(T a, T b) { return a+b; }
  vector<T> s;
  Fenwick(int n = 0) : s(n, ID) {}
```

//A[i] = f(A[i], v); time: O(lq n)

```
void modify(int i, T v) {
   for (; i < sz(s); i = i+1) s[i]=f(s[i],v);
 // Get f(A[0], ..., A[i-1]); time: O(lq n)
 T query(int i) {
   T v = ID;
   for (; i > 0; i \in [i-1]) v = f(v, s[i-1]);
   return v;
 1
 // Find smallest i such that
 // f(A[0],...,A[i-1]) >= val; time: O(lq n)
  // Prefixes must have non-descreasing values.
 int lowerBound(T val) {
   if (val <= ID) return 0;</pre>
   int i = -1, mask = 1;
   while (mask \leq sz(s)) mask \star= 2;
   T off = ID:
   while (mask /= 2) {
     int k = mask+i:
     if (k < sz(s)) {
       T x = f(off, s[k]);
        if (val > x) i=k, off=x;
   return i+2;
};
structures/fenwick tree 2d.h
// Fenwick tree 2D (BIT tree 2D); space: O(n*m)
// Default version: prefix sums 2D
// Change s to hashmap for O(g lg^2 n) memory
struct Fenwick2D {
 using T = int;
 static constexpr T ID = 0;
 T f(T a, T b) { return a+b; }
 vector<T> s:
 int w, h;
 Fenwick2D(int n = 0, int m = 0)
   : s(n*m, ID), w(n), h(m) {}
  // A[i,j] = f(A[i,j], v); time: O(1g^2 n)
 void modify(int i, int j, T v) {
   for (; i < w; i |= i+1)
     for (int k = j; k < h; k = k+1)
        s[i*h+k] = f(s[i*h+k], v);
  // Query prefix; time: O(lg^2 n)
 T query(int i, int j) {
   T v = ID;
   for (; i>0; i&=i-1)
     for (int k = i; k > 0; k \le k-1)
       v = f(v, s[i*h+k-h-1]);
   return v;
structures/find union.h
// Disjoint set data structure; space: O(n)
// Operations work in amortized O(alfa(n))
```

```
struct FAU {
 Vi G;
 FAU(int n = 0) : G(n, -1) {}
  // Get size of set containing i
  int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
  int find(int i) {
   return G[i] < 0 ? i : G[i] = find(G[i]);</pre>
  // Union sets containing i and i
 bool join(int i, int j) {
    i = find(i); j = find(j);
    if (i == j) return 0;
    if (G[i] > G[j]) swap(i, j);
    G[i] += G[j]; G[j] = i;
    return 1:
 }
};
structures/hull offline.h
constexpr 11 INF = 2e18;
// constexpr double INF = 1e30;
// constexpr double EPS = 1e-9;
// MAX of linear functions; space: O(n)
// Use if you add lines in increasing `a` order
// Default uncommented version is for int64
struct Hull {
 using T = 11; // Or change to double
 struct Line {
   T a, b, end;
   T intersect (const Line& r) const {
      // Version for double:
      //if (r.a-a < EPS) return b>r.b?INF:-INF;
      //return (b-r.b) / (r.a-a);
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b, d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u^d));
 };
  vector<Line> S;
 Hull() { S.pb({ 0, -INF, INF }); }
  // Insert f(x) = ax+b; time: amortized O(1)
 void push(T a, T b) {
    Line 1{a, b, INF};
    while (true) {
     T e = S.back().end=S.back().intersect(1);
      if (sz(S) < 2 | | S[sz(S)-2].end < e)
       break:
      S.pop_back();
    S.pb(1);
  // Ouerv max(f(x) for each f): time: O(lg n)
 T query (T x) {
    auto t = *upper_bound(all(S), x,
      [](int 1, const Line& r) {
        return 1 < r.end;</pre>
```

1);

```
return t.a*x + t.b;
};
structures/hull online.h
constexpr 11 INF = 2e18;
// MAX of linear functions online; space: O(n)
struct Hull {
  static bool modeQ; // Toggles operator< mode</pre>
  struct Line {
   mutable 11 a, b, end;
    ll intersect (const Line& r) const {
     if (a==r.a) return b > r.b ? INF : -INF;
     11 u = b-r.b. d = r.a-a;
     return u/d + ((u^d) >= 0 || !(u%d));
   }
   bool operator<(const Line& r) const {</pre>
      return modeO ? end < r.end : a < r.a;</pre>
  };
  multiset<Line> S:
  Hull() { S.insert({ 0, -INF, INF }); }
  // Updates segment end
  bool update(multiset<Line>::iterator it) {
   auto cur = it++; cur->end = INF;
   if (it == S.end()) return false;
   cur->end = cur->intersect(*it);
   return cur->end >= it->end;
  // Insert f(x) = ax+b; time: O(lq n)
  void insert(ll a, ll b) {
    auto it = S.insert({ a, b, INF });
    while (update(it)) it = --S.erase(++it);
   rep(i, 0, 2)
     while (it != S.begin() && update(--it))
       update(it = --S.erase(++it));
  // Query max(f(x) for each f): time: O(lg n)
  ll query(ll x) {
   mode0 = 1:
   auto 1 = \starS.upper bound({ 0, 0, x });
   mode0 = 0;
   return l.a*x + l.b;
};
bool Hull::mode0 = false;
structures/max queue.h
// Queue with max query on contained elements
struct MaxQueue {
  using T = int:
  deque<T> 0, M;
```

```
// Add v to the back; time: amortized O(1)
```

while (!M.emptv() && M.back() < v)</pre>

void push(T v) {

```
M.pop back();
   M.pb(v); O.pb(v);
 // Pop from the front; time: O(1)
 } () gog biov
   if (M.front() == 0.front()) M.pop front();
   Q.pop_front();
 // Get max element value; time: O(1)
 T max() const { return M.front(); }
structures/pairing heap.h
// Pairing heap implementation; space O(n)
// Elements are stored in vector for faster
// allocation. It's MINIMUM queue.
// Allows to merge heaps in O(1)
template < class T, class Cmp = less < T>>
struct PHeap {
 struct Node {
   T val:
   int child{-1}, next{-1}, prev{-1};
   Node (T x = T()) : val(x) \{\}
 };
 using Vnode = vector<Node>;
 Vnode& M:
 int root{-1};
 int unlink(int& i) {
   if (i >= 0) M[i].prev = -1;
   int x = i; i = -1;
   return x:
 void link(int host, int& i, int val) {
   if (i >= 0) M[i].prev = -1;
   i = val:
   if (i >= 0) M[i].prev = host;
 int merge(int 1, int r) {
   if (1 < 0) return r;
   if (r < 0) return 1;
   if (Cmp()(M[1].val, M[r].val)) swap(l, r);
   link(1, M[1].next, unlink(M[r].child));
   link(r, M[r].child, 1);
   return r;
 int mergePairs(int v) {
   if (v < 0 || M[v].next < 0) return v;</pre>
   int v2 = unlink(M[v].next);
   int v3 = unlink(M[v2].next);
   return merge(merge(v, v2), mergePairs(v3));
 // Initialize heap with given node storage
 // Just declare 1 Vnode and pass it to heaps
 PHeap(Vnode& mem) : M(mem) {}
```

```
// Add given key to heap, returns index; O(1)
  int push (const T& x) {
    int index = sz(M);
   M.emplace back(x):
    root = merge (root, index);
    return index;
  // Change key of i to smaller value; O(1)
  void decrease(int i, T val) {
    assert(!Cmp()(M[i].val, val));
    M[i].val = val;
    int prev = M[i].prev;
    if (prev < 0) return;</pre>
    auto& p = M[prev];
    link (prev, (p.child == i ? p.child
        : p.next), unlink(M[i].next));
    root = merge(root, i);
  bool empty() { return root < 0; }</pre>
  const T& top() { return M[root].val; }
  // Merge with other heap. Must use same vec.
  void merge(PHeap& r) { // time: O(1)
   assert (&M == &r.M);
   root = merge(root, r.root); r.root = -1;
  // Remove min element; time: O(lq n)
 } () gog biov
   root = mergePairs (unlink (M[root].child));
};
structures/rmq.h
// Range Minimum Query; space: O(n lg n)
struct RMO {
 using T = int;
  static constexpr T ID = INT_MAX;
 T f(T a, T b) { return min(a, b); }
 vector<vector<T>>> s;
  // Initialize RMQ structure; time: O(n lq n)
  RMO(const vector<T>& vec = {}) {
   s = \{vec\};
   for (int h = 1; h \le sz(vec); h *= 2) {
      s.emplace_back();
      auto& prev = s[sz(s)-2];
      rep(i, 0, sz(vec)-h*2+1)
        s.back().pb(f(prev[i], prev[i+h]));
  // Query f(s[b], ..., s[e-1]); time: O(1)
 T guerv(int b, int e) {
    if (b >= e) return ID;
    int k = 31 - __builtin_clz(e-b);
    return f(s[k][b], s[k][e - (1<<k)]);
```

```
structures/segment tree.h
```

```
// Optionally dynamic segment tree with lazy
// propagation. Configure by modifying:
// - T - data type for updates (stored type)
// - ID - neutral element for extra
// - Node - details in comments
struct SeamentTree {
 using T = int;
 static constexpr T ID = 0;
 // static constexpr T ID = INT MIN; // max/=
 struct Node {
   T extra{ID}; // Lazy propagated value
   // Aggregates: sum, max, count of max
   T sum{0}, great{INT_MIN}, nGreat{0};
   // Initialize node with default value x
   void init(T x, int size) {
     sum = x*size; great = x; nGreat = size;
    // Merge with node R on the right
   void merge(const Node& R) {
      if (great < R.great) nGreat =R.nGreat;</pre>
     else if(great==R.great) nGreat+=R.nGreat;
     sum += R.sum;
     great = max(great, R.great);
   // + version
   // Apply modification to node, return
    // value to be applied to node on right
   T apply (T x, int size) {
     extra += x;
     sum += x*size:
     great += x;
     return x;
   // MAX
    // T apply(T x, int size) {
   // if (great <= x) nGreat = size;</pre>
   // extra = max(extra, x);
   // great = max(great, x);
       // sum doesn't work here
   // return x;
   1/ }
    // T apply(T x, int size) {
   // extra = x;
       sum = x*size;
        great = x;
       nGreat = size;
   // return x;
   1/ }
 };
 vector<Node> V:
 int len:
 // vector<array<int, 3>> links; // [DYNAMIC]
 // T defVal:
 SegmentTree(int n=0, T def=ID) {init(n,def);}
```

```
void init(int n, T def) {
  for (len = 1; len < n; len *= 2);
  // [STATIC] version
 V.assign(len*2, {});
  rep(i, len, len+n) V[i].init(def, 1);
  for (int i = len-1; i > 0; i--) update(i);
  // [DYNAMIC] version
  // defVal = def:
  // links.assign(2, {-1, -1, len});
  // V.assign(2, {});
 // V[1].init(def, len);
// [STATIC] version
int getChild(int i, int j) { return i*2+j; }
// [DYNAMIC] version
// int getChild(int i, int j) {
// if (links[i][j] < 0) {
      int size = links[i][2] / 2;
      links[i][i] = sz(V);
      links.push_back({ -1, -1, size });
      V.emplace_back();
      V.back().init(defVal, size);
// return links[i][j];
1/ 7
int L(int i) { return getChild(i, 0); }
int R(int i) { return getChild(i, 1); }
void update(int i) {
 int a = L(i), b = R(i);
 V[i] = {};
 V[i].merge(V[a]);
 V[i].merge(V[b]);
void push(int i, int size) {
 T e = V[i].extra;
 if (e != ID) {
   e = V[L(i)].apply(e, size/2);
   V[R(i)].apply(e, size/2);
   V[i].extra = ID;
 }
// Modify [vBegin; end) with x; time: O(lq n)
T modify(int vBegin, int vEnd, T x,
        int i = 1.
        int begin = 0, int end = -1) {
  if (end < 0) end = len;</pre>
 if (vEnd <= begin || end <= vBegin)
   return x;
  if (vBegin <= begin && end <= vEnd) {
   return V[i].apply(x, end-begin);
  int mid = (begin + end) / 2;
 push (i, end-begin):
  x = modify(vBegin, vEnd, x, L(i), begin, mid);
  x = modify(vBegin, vEnd, x, R(i), mid, end);
  update(i);
```

```
return x;
  // Ouery [vBegin; vEnd); time: O(lq n)
  // Returns base nodes merged together
 Node guery (int vBegin, int vEnd, int i = 1,
             int begin = 0, int end = -1) {
    if (end < 0) end = len;</pre>
    if (vEnd <= begin || end <= vBegin)</pre>
      return {};
    if (vBegin <= begin && end <= vEnd)
      return V[i];
    int mid = (begin + end) / 2;
   push (i, end-begin);
   Node x = query(vBegin, vEnd, L(i), begin, mid);
   x.merge(query(vBegin, vEnd, R(i), mid, end));
   return x;
  // TODO: generalize?
  // Find longest suffix of given interval
  // such that max value is smaller than val.
  // Returns suffix begin index; time: O(lq n)
 T search (int vBegin, int vEnd, int val,
           int i=1, int begin=0, int end=-1) {
    if (end < 0) end = len;</pre>
    if (vEnd <= begin || end <= vBegin)</pre>
      return begin;
    if (vBegin <= begin && end <= vEnd) {</pre>
     if (V[i].great < val) return begin;</pre>
     if (begin+1 == end) return end;
   }
    int mid = (begin+end) / 2;
   push(i, end-begin);
   int ind = search(vBegin, vEnd, val,
                     R(i), mid, end);
   if (ind > mid) return ind;
   return search (vBegin, vEnd, val,
                  L(i), begin, mid);
structures/segment tree beats.h
constexpr 11 INF = 1e18;
// Segment tree with min/+ update and
// sum/max query; time: amortized O(n lg^2 n)
// or O(n lg n) if not using + operation
struct SegmentTree {
 using T = 11;
 vector<T> sum, plus, max1, max2, cnt1;
 int len;
  SegmentTree(int n = 0) {
   for (len = 1; len < n; len *= 2);
    sum.resize(len*2);
    plus.resize(len*2);
    max1.resize(len*2);
    max2.assign(len*2, -INF);
    cnt1.assign(len*2, 1);
    for (int i = len-1; i > 0; i--) update(i);
```

```
void apply(int i, T m, T p, int size) {
  plus[i] += p; sum[i] += p*size;
  \max 1[i] += p; \max 2[i] += p;
  if (m < max1[i]) {</pre>
    sum[i] -= (max1[i]-m) *cnt1[i];
    max1[i] = m;
}
void update(int i) {
  int a = i*2, b = i*2+1;
  sum[i] = sum[a] + sum[b];
  plus[i] = 0;
  max1[i] = max1[a];
  max2[i] = max2[a];
  cnt1[i] = cnt1[a];
  if (max1[b] > max1[i]) {
    max2[i] = max1[i];
    max1[i] = max1[b];
    cnt1[i] = cnt1[b];
  } else if (max1[b] == max1[i]) {
    cnt1[i] += cnt1[b];
  } else if (max1[b] > max2[i]) {
    max2[i] = max1[b];
  max2[i] = max(max2[i], max2[b]);
void push(int i, int s) {
  rep(j, 0, 2)
    apply(i*2+j, max1[i], plus[i], s/2);
  plus[i] = 0;
// Apply min with x on [vBegin; vEnd)
// time: amortized O(lq n) or O(lq^2 n)
void setMin(int vBegin, int vEnd, T x,
            int i = 1,
            int begin = 0, int end = -1) {
  if (end < 0) end = len;</pre>
  if (vEnd <= begin || end <= vBegin ||</pre>
      max1[i] < x) return;</pre>
  if (begin >= vBegin && end <= vEnd &&
      max2[i] < x)
    return apply(i, x, 0, end-begin);
  int mid = (begin+end) / 2;
  push (i, end-begin);
  setMin(vBegin, vEnd, x, i*2, begin, mid);
  setMin(vBegin, vEnd, x, i*2+1, mid, end);
  update(i);
// Add x on [vBegin; vEnd); time: O(lg n)
void add(int vBegin, int vEnd, T x,
         int i = 1.
         int begin = 0, int end = -1) {
  if (end < 0) end = len;</pre>
  if (vEnd <= begin || end <= vBegin) return;
  if (begin >= vBegin && end <= vEnd)
    return apply(i, INF, x, end-begin);
  int mid = (begin+end) / 2;
```

```
add (vBegin, vEnd, x, i*2+1, mid, end);
    update(i);
 // Ouery sum of [vBegin; vEnd); time: O(lq n)
 T getSum(int vBegin, int vEnd, int i = 1,
           int begin = 0, int end = -1) {
    if (end < 0) end = len;</pre>
    if (vEnd<=begin || end<=vBegin) return 0;</pre>
    if (vBegin <= begin && end <= vEnd)
      return sum[i]:
    int mid = (begin+end) / 2;
    push (i, end-begin);
    return getSum(vBegin, vEnd, i*2, begin, mid) +
          getSum(vBegin, vEnd, i*2+1, mid, end);
  // Query max of [vBegin; vEnd); time: O(lg n)
 T getMax(int vBegin, int vEnd, int i = 1,
           int begin = 0, int end = -1) {
    if (end < 0) end = len;</pre>
    if (vEnd <= begin || end <= vBegin)</pre>
      return -INF:
    if (vBegin <= begin && end <= vEnd)
      return max1[i];
    int mid = (begin+end) / 2;
    push (i, end-begin);
    return max (
      getMax(vBegin, vEnd, i*2, begin, mid),
      getMax(vBegin, vEnd, i*2+1, mid, end)
 }
};
structures/segment tree point.h
// Segment tree (point, interval)
// Configure by modifying:
// - T - stored data type
// - ID - neutral element for query operation
// - merge(a, b) - combine results
struct SegmentTree {
 using T = int;
 static constexpr T ID = INT MIN;
  static T merge(T a, T b) { return max(a,b); }
  vector<T> V;
 int len:
  SegmentTree(int n = 0, T def = ID) {
    for (len = 1; len < n; len *= 2);
    V.resize(len*2, ID);
    rep(i, 0, n) V[len+i] = def;
    for (int i = len-1; i > 0; i--)
      V[i] = merge(V[i*2], V[i*2+1]);
  void set(int i, T val) {
    V[i+=len] = val;
    while ((i/=2) > 0)
      V[i] = merge(V[i*2], V[i*2+1]);
 T query (int begin, int end) {
    begin += len; end += len-1;
```

push (i, end-begin);

add(vBegin, vEnd, x, i*2, begin, mid);

```
if (begin > end) return ID;
    if (begin == end) return V[begin];
    T \times = merge(V[begin], V[end]);
    while (begin/2 < end/2) {
     if (~begin&1) x = merge(x, V[begin^1]);
     if (end&1) x = merge(x, V[end^1]);
     begin /= 2; end /= 2;
   return x;
};
constexpr SeamentTree::T SeamentTree::ID;
structures/treap.h
// "Set" of implicit keyed treaps; space: O(n)
// Nodes are keyed by their indices in array
// of all nodes. Treap key is key of its root.
// "Node x" means "node with kev x".
// "Treap x" means "treap with key x".
// Key -1 is "null".
// Put anv additional data in Node struct.
struct Treap {
  struct Node {
    // E[0] = left child, <math>E[1] = right child
    // weight = node random weight (for treap)
    // size = subtree size, par = parent node
    int E[2] = \{-1, -1\}, weight{rand()};
   int size{1}, par{-1};
   bool flip{0}; // Is interval reversed?
  vector<Node> G; // Array of all nodes
  // Initialize structure for n nodes
  // with keys 0, ..., n-1; time: O(n)
  // Each node is separate treap,
  // use join() to make sequence.
  Treap(int n = 0) : G(n) {}
  // Create new treap (a single node),
  // returns its key; time: O(1)
  int make() {
   G.emplace_back(); return sz(G)-1;
  // Get size of node x subtree. x can be -1.
  int size(int x) { // time: O(1)
   return (x \ge 0 ? G[x].size : 0);
  // Propagate down data (flip flag etc).
  // x can be -1; time: O(1)
  void push(int x) {
   if (x >= 0 && G[x].flip) {
     G[x].flip = 0;
      swap(G[x].E[0], G[x].E[1]);
      each (e, G[x].E) if (e>=0) G[e].flip ^= 1;
   } // + any other lazy operations
  // Update aggregates of node x.
  // x can be -1; time: O(1)
  void update(int x) {
   if (x >= 0) {
```

```
int & s = G[x].size = 1;
    G[x].par = -1;
    each (e, G[x].E) if (e >= 0) {
      s += G[e].size;
      G[e].par = x;
 } // + any other aggregates
// Split treap x into treaps 1 and r
// such that 1 contains first i elements
// and r the remaining ones.
// x, 1, r can be -1; time: ^{\circ}O(lg n)
void split(int x, int& l, int& r, int i) {
  push(x); l = r = -1;
  if (x < 0) return;
  int key = size(G[x].E[0]);
  if (i <= kev) {
    split(G[x].E[0], 1, G[x].E[0], i);
    r = x;
  } else {
    split(G[x].E[1], G[x].E[1], r, i-key-1);
    1 = x;
  update(x);
// Join treaps 1 and r into one treap
// such that elements of 1 are before
// elements of r. Returns new treap.
// 1, r and returned value can be -1.
int join(int 1, int r) { // time: ~O(lg n)
  push(1); push(r);
  if (1 < 0 || r < 0) return max(1, r);</pre>
  if (G[1].weight < G[r].weight) {</pre>
    G[1].E[1] = join(G[1].E[1], r);
    update(1);
    return 1;
  G[r].E[0] = join(l, G[r].E[0]);
  update(r);
  return r;
// Find i-th node in treap x.
// Returns its key or -1 if not found.
// x can be -1; time: ~O(lq n)
int find(int x, int i) {
  while (x \ge 0) {
    push(x);
    int key = size(G[x].E[0]);
    if (key == i) return x;
    x = G[x].E[kev < i];
    if (key < i) i -= key+1;</pre>
  return -1;
// Get key of treap containing node x
// (kev of treap root). x can be -1.
int root(int x) { // time: ~O(lg n)
  while (G[x].par \ge 0) x = G[x].par;
  return x:
```

```
// Get position of node x in its treap.
  // x is assumed to NOT be -1; time: ~O(lg n)
 int index(int x) {
   int p, i = size(G[x].E[G[x].flip]);
   while ((p = G[x].par) >= 0) {
     if (G[p].E[1] == x) i+=size(G[p].E[0])+1;
     if (G[p].flip) i = G[p].size-i-1;
     x = p;
   return i:
 // Reverse interval [1;r) in treap x.
  // Returns new key of treap; time: ~O(lg n)
 int reverse(int x, int 1, int r) {
   int a, b, c;
   split(x, b, c, r);
   split (b, a, b, 1);
   if (b >= 0) G[b].flip ^= 1;
   return join(join(a, b), c);
};
structures/ext/hash table.h
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
// gp_hash_table<K, V> = faster unordered_set
// Anti-anti-hash
const size_t HXOR = mt19937_64(time(0))();
template < class T > struct SafeHash {
 size t operator()(const T& x) const {
   return hash<T>()(x ^ T(HXOR));
};
structures/ext/rope.h
#include <ext/rope>
using namespace __gnu_cxx;
// rope<T> = implicit cartesian tree
structures/ext/tree.h
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template < class T, class Cmp = less < T>>
using ordered set = tree<
T, null_type, Cmp, rb_tree_tag,
 tree_order_statistics_node_update
// Standard set functions and:
// t.order of key(key) - index of first >= key
// t.find_by_order(i) - find i-th element
// t1.join(t2) - assuming t1<>t2 merge t2 to t1
structures/ext/trie.h
```

#include <ext/pb_ds/assoc_container.hpp>

#include <ext/pb ds/trie policy.hpp>

using namespace gnu pbds;

```
using pref trie = trie<
 string, null type,
 trie_string_access_traits , pat_trie_tag,
 trie prefix search node update
text/aho corasick.h
constexpr char AMIN = 'a'; // Smallest letter
constexpr int ALPHA = 26: // Alphabet size
// Aho-Corasick algorithm for linear-time
// multiple pattern matching.
// Add patterns using add(), then call build().
struct Aho {
 vector<array<int, ALPHA>> nxt{1};
 Vi suf = \{-1\}, accLink = \{-1\};
 vector<Vi> accept{1};
  // Add string with given ID to structure
  // Returns index of accepting node
 int add(const string& str, int id) {
    int i = 0:
    each(c, str) {
      if (!nxt[i][c-AMIN]) {
        nxt[i][c-AMIN] = sz(nxt);
        nxt.pb({}); suf.pb(-1);
        accLink.pb(1); accept.pb({});
      i = nxt[i][c-AMIN];
    accept[i].pb(id);
    return i;
  // Build automata; time: O(V*ALPHA)
 void build() {
    queue<int> que;
    each(e, nxt[0]) if (e) {
      suf[e] = 0; que.push(e);
    while (!que.empty()) {
      int i = que.front(), s = suf[i], j = 0;
      que.pop();
      each(e, nxt[i]) {
       if (e) que.push(e);
        (e ? suf[e] : e) = nxt[s][j++];
      accLink[i] = (accept[s].empty() ?
          accLink[s] : s);
  // Append `c` to state `i`
 int next(int i, char c) {
    return nxt[i][c-AMIN];
  // Call `f` for each pattern accepted
  // when in state `i` with its ID as argument.
  // Return true from `f` to terminate early.
  // Calls are in descreasing length order.
 template < class F > void accepted (int i, F f) {
    while (i !=-1) {
      each(a, accept[i]) if (f(a)) return;
      i = accLink[i];
```

};

```
text/kmp.h
// Computes prefsuf array; time: O(n)
// ps[i] = max prefsuf of [0;i); ps[0] := -1
template < class T > Vi kmp (const T& str) {
 Vi ps; ps.pb(-1);
  each(x, str) {
    int k = ps.back();
    while (k \ge 0 \&\& str[k] != x) k = ps[k];
   ps.pb(k+1);
  return ps;
// Finds occurences of pat in vec; time: O(n)
// Returns starting indices of matches.
template<class T>
Vi match (const T& str, T pat) {
  int n = sz(pat);
  pat.pb(-1); // SET TO SOME UNUSED CHARACTER
  pat.insert(pat.end(), all(str));
  Vi ret, ps = kmp(pat);
  rep(i, 0, sz(ps)) {
   if (ps[i] == n) ret.pb(i-2*n-1);
```

return ret;

```
text/kmr.h
// KMR algorithm for O(1) lexicographical
// comparison of substrings.
struct KMR {
 vector<Vi> ids:
  KMR() {}
  // Initialize structure; time: O(n lg^2 n)
  // You can change str type to Vi freely.
  KMR (const string& str) {
    ids.clear();
    ids.pb(Vi(all(str)));
    for (int h = 1; h \le sz(str); h *= 2) {
     vector<pair<Pii, int>> tmp;
      rep(j, 0, sz(str)) {
       int a = ids.back()[j], b = -1;
       if (j+h < sz(str)) b = ids.back()[j+h];
       tmp.pb({ {a, b}, j });
      sort(all(tmp));
     ids.emplace_back(sz(tmp));
      rep(j, 0, sz(tmp)) {
       if (j > 0 && tmp[j-1].x != tmp[j].x)
        ids.back()[tmp[j].y] = n;
```

```
// Get representative of [begin; end); 0(1)
 Pii get (int begin, int end) {
   if (begin >= end) return {0, 0};
   int k = 31 - builtin clz(end-begin);
   return {ids[k][begin], ids[k][end-(1<<k)]};</pre>
  // Compare [b1;e1) with [b2;e2); O(1)
  // Returns -1 if <, 0 if ==, 1 if >
 int cmp(int b1, int e1, int b2, int e2) {
   int 11 = e1-b1, 12 = e2-b2;
    int 1 = min(11, 12);
    Pii x = get(b1, b1+1), y = get(b2, b2+1);
    if (x == y) return (11 > 12) - (11 < 12);
   return (x > y) - (x < y);
  // Compute suffix array of string; O(n)
 Vi sufArray() {
   Vi sufs(sz(ids.back()));
    rep(i, 0, sz(ids.back()))
     sufs[ids.back()[i]] = i;
    return sufs:
};
```

text/lcp.h

```
// Compute Longest Common Prefix array for
// given string and it's suffix array; O(n)
// In order to compute suffix array use kmr.h
// or suffix_array_linear.h
template<class T>
Vi lcpArray(const T& str, const Vi& sufs) {
  int n = sz(str), k = 0;
  Vi pos(n), lcp(n-1);
  rep(i, 0, n) pos[sufs[i]] = i;
  rep(i, 0, n) {
    if (pos[i] < n-1) {</pre>
      int j = sufs[pos[i]+1];
      while (i+k < n && j+k < n &&
          str[i+k] == str[j+k]) k++;
      lcp[pos[i]] = k;
    if (k > 0) k--;
  return lcp;
```

text/lyndon factorization.h

```
// Compute Lyndon factorization for s; O(n)
// Word is simple iff it's stricly smaller
// than any of it's nontrivial suffixes.
// Lyndon factorization is division of string
// into non-increasing simple words.
// It is unique.
vector<string> duval(const string& s) {
 int n = sz(s), i = 0;
 vector<string> ret:
  while (i < n) {
   int j = i+1, k = i;
    while (j < n \&\& s[k] \leftarrow s[j])
     k = (s[k] < s[j] ? i : k+1), j++;
    while (i \le k)
      ret.pb(s.substr(i, j-k)), i += j-k;
```

```
return ret;
```

text/main lorentz.h

#include "z function.h"

```
struct Sqr {
 int begin, end, len;
// Main-Lorentz algorithm for finding
// all squares in given word; time: O(n lg n)
// Results are in compressed form:
// (b, e, 1) means that for each b <= i < e
// there is square at position i of size 21.
// Each square is present in only one interval.
vector<Sqr> lorentz(const string  s) {
 int n = sz(s);
 if (n <= 1) return {};</pre>
  auto a = s.substr(0, n/2), b = s.substr(n/2);
  auto ans = lorentz(a):
  each (p, lorentz (b))
   ans.pb(\{p.begin+n/2, p.end+n/2, p.len\});
  string ra(a.rbegin(), a.rend());
  string rb(b.rbegin(), b.rend());
  rep(j, 0, 2) {
    Vi z1 = prefPref(ra), z2 = prefPref(b+a);
    z1.pb(0); z2.pb(0);
    rep(c, 0, sz(a)) {
      int 1 = sz(a) - c;
      int x = c - \min(1-1, z1[1]);
      int y = c - max(1-z2[sz(b)+c], j);
      if (x > y) continue;
      if (j)
        ans.pb(\{n-y-1*2, n-x-1*2+1, 1\});
      else
        ans.pb(\{x, y+1, 1\});
    a.swap(rb);
   b.swap(ra);
  return ans:
```

text/manacher.h

```
// Manacher algorithm; time: O(n)
// Finds largest radiuses for palindromes:
// r[2*i] = for center at i (single letter = 1)
// r[2*i+1] = for center between i and i+1
template < class T > Vi manacher (const T& str) {
 int n = sz(str) \star 2, c = 0, e = 1;
 Vi r(n, 1);
 auto get = [&] (int i) { return i%2 ? 0 :
   (i \ge 0 \&\& i < n ? str[i/2] : i); };
 rep(i, 0, n) {
   if (i < e) r[i] = min(r[c*2-i], e-i);
```

```
while (get(i-r[i]) == get(i+r[i])) r[i]++;
  if (i+r[i] > e) c = i, e = i+r[i]-1;
rep(i, 0, n) r[i] /= 2;
return r:
```

text/min rotation.h

```
// Find lexicographically smallest
// rotation of s; time: O(n)
// Returns index where shifted word starts.
// You can use std::rotate to get the word:
// rotate(s.begin(), v.begin()+minRotation(v),
         v.end());
int minRotation(string s) {
 int a = 0, n = sz(s); s += s;
 rep(b, 0, n) rep(i, 0, n) {
    if (a+i == b || s[a+i] < s[b+i]) {</pre>
      b += max(0, i-1); break;
    if (s[a+i] > s[b+i]) {
      a = b; break;
 return a;
```

text/palindromic tree.h

```
constexpr int ALPHA = 26; // Set alphabet size
// Tree of all palindromes in string,
// constructed online by appending letters.
// space: O(n*ALPHA); time: O(n)
// Code marked with [EXT] is extension for
// calculating minimal palindrome partition
// in O(n lq n). Can also be modified for
// similar dynamic programmings.
struct PalTree {
 Vi txt: // Text for which tree is built
 // Node 0 = empty palindrome (root of even)
 // Node 1 = "-1" palindrome (root of odd)
 Vi len{0, -1}; // Lengths of palindromes
 Vi link{1, 0}; // Suffix palindrome links
 // Edges to next palindromes
 vector<array<int, ALPHA>> to{ {}, {} };
 int last{0}; // Current node (max suffix pal)
 Vi diff{0, 0};  // len[i]-len[link[i]] [EXT]
 Vi slink{0, 0}; // Serial links
                                          [EXT]
 Vi series {0, 0}; // Series DP answer
                                          [EXT]
 Vi ans{0};
                  // DP answer for prefix[EXT]
 int ext(int i) {
   while (len[i]+2 > sz(txt) ||
           txt[sz(txt)-len[i]-2] != txt.back())
      i = link[i];
   return i:
 // Append letter from [0; ALPHA); time: O(1)
 // (or O(lq n) if [EXT] is enabled)
 void add(int x) {
   txt.pb(x);
```

```
last = ext(last);
    if (!to[last][x]) {
     len.pb(len[last]+2);
     link.pb(to[ext(link[last])][x]);
     to[last][x] = sz(to);
     to.emplace back();
     diff.pb(len.back() - len[link.back()]);
     slink.pb(diff.back() == diff[link.back()]
       ? slink[link.back()] : link.back());
     series.pb(0);
     // [/EXT]
    last = to[last][x];
    // [EXT]
    ans.pb(INT MAX);
    for (int i=last; len[i] > 0; i=slink[i]) {
     series[i] = ans[sz(ans) - len[slink[i]]
                    - diff[i] - 1];
     if (diff[i] == diff[link[i]])
       series[i] = min(series[i],
                        series[link[i]]);
     // If you want only even palindromes
      // set ans only for sz(txt) %2 == 0
     ans.back() = min(ans.back(), series[i]+1);
    // [/EXT]
};
```

text/suffix array linear.h

```
#include "../util/radix sort.h"
// KS algorithm for suffix array; time: O(n)
// Input values are assumed to be in [1;k]
Vi sufArray(Vi str, int k) {
  int n = sz(str);
 Vi suf(n):
  str.resize(n+15);
  if (n < 15) {
   iota(all(suf), 0);
    rep(j, 0, n) countSort(suf,
     [&](int i) { return str[i+n-j-1]; }, k);
    return suf;
  // Compute triples codes
  Vi tmp, code(n+2);
  rep(i, 0, n) if (i % 3) tmp.pb(i);
  rep(j, 0, 3) countSort(tmp,
   [&] (int i) { return str[i-j+2]; }, k);
  int mc = 0, j = -1;
  each(i, tmp) {
   code[i] = mc += (i == -1)
       str[i] != str[j] ||
       str[i+1] != str[j+1] ||
       str[i+2] != str[j+2]);
   j = i;
```

```
// Compute suffix array of 2/3
  tmp.clear():
  for (int i=1; i < n; i += 3) tmp.pb(code[i]);</pre>
  for (int i=2; i < n; i += 3) tmp.pb(code[i]);</pre>
  tmp = sufArray(move(tmp), mc);
  // Compute partial suffix arrays
 Vi third:
  int th = (n+4) / 3;
 if (n%3 == 1) third.pb(n-1);
 rep(i, 1, sz(tmp)) {
   int e = tmp[i]:
   tmp[i-1] = (e 
   code[tmp[i-1]] = i;
   if (e < th) third.pb(e*3);
  tmp.pop_back();
 countSort (third.
   [&](int i) { return str[i]; }, k);
  // Merge suffix arrays
 merge(all(third), all(tmp), suf.begin(),
    [&] (int 1, int r) {
      while (1%3 == 0 || r%3 == 0) {
        if (str[l] != str[r])
          return str[1] < str[r];</pre>
        1++; r++;
     return code[1] < code[r];</pre>
    1);
  return suf;
// KS algorithm for suffix array; time: O(n)
Vi sufArray(const string& str) {
 return sufArray(Vi(all(str)), 255);
text/suffix automaton.h
```

```
constexpr char AMIN = 'a'; // Smallest letter
constexpr int ALPHA = 26; // Set alphabet size
// Suffix automaton - minimal DFA that
// recognizes all suffixes of given string
// (and encodes all substrings);
// space: O(n*ALPHA); time: O(n)
// Paths from root are equivalent to substrings
// Extensions:
// - [OCC] - count occurences of substrings
// - [PATHS] - count paths from node
struct SufDFA {
 // State v represents endpos-equivalence
 // class that contains words of all lengths
  // between link[len[v]]+1 and len[v].
  // len[v] = longest word of equivalence class
 // link[v] = link to state of longest suffix
              in other equivalence class
  // to[v][c] = automaton edge c from v
 Vi len{0}, link{-1};
 vector<array<int, ALPHA>> to{ {} };
 int last{0}: // Current node (whole word)
```

```
vector<Vi> inSufs; // [OCC] Suffix-link tree
                  // [OCC] Occurence count
vector<ll> paths; // [PATHS] Out-path count
SufDFA() {}
// Build suffix automaton for given string
// and compute extended stuff; time: O(n)
SufDFA(const string& s) {
 each(c, s) add(c);
 finish();
// Append letter to the back
void add(char c) {
 int v = last, x = c-AMIN;
 last = sz(len);
 len.pb(len[v]+1);
 link.pb(0);
 to.pb({});
 cnt.pb(1); // [OCC]
 while (v != -1 && !to[v][x]) {
   to[v][x] = last;
   v = link[v];
 if (v != -1) {
   int q = to[v][x];
   if (len[v]+1 == len[q]) {
     link[last] = q;
   } else {
      len.pb(len[v]+1);
      link.pb(link[q]);
      to.pb(to[q]);
      cnt.pb(0); // [OCC]
      link[last] = link[q] = sz(len)-1;
      while (v != -1 \&\& to[v][x] == g) {
       to[v][x] = link[q];
       v = link[v];
// Compute some additional stuff (offline)
void finish() {
 inSufs.resize(sz(len));
 rep(i, 1, sz(link)) inSufs[link[i]].pb(i);
 dfsSufs(0);
 // [PATHS]
 paths.assign(sz(len), 0);
 dfs(0);
 // [/PATHS]
// Only for [OCC]
void dfsSufs(int v) {
 each(e, inSufs[v]) {
   dfsSufs(e);
   cnt[v1 += cnt[e1;
```

```
// Only for [PATHS]
  void dfs(int v) {
    if (paths[v]) return;
    paths[v] = 1:
    each(e, to[v]) if (e) {
      dfs(e):
      paths[v] += paths[e];
 }
 // Go using edge `c` from state `i`.
 // Returns 0 if edge doesn't exist.
 int next(int i, char c) {
    return to[i][c-AMIN];
 // Get lexicographically k-th substring
 // of represented string; time: O(|substr|)
 // Empty string has index 0.
 // Requires [PATHS] extension.
 string lex(ll k) {
    string s:
    int v = 0:
    while (k--) rep(i, 0, ALPHA) {
      int e = to[v][i];
      if (e) {
       if (k < paths[e]) {</pre>
         s.pb(char(AMIN+i));
         v = e;
         break;
        k -= paths[e];
    return s;
};
```

text/suffix tree.h

```
// Ukkonen's algorithm for online suffix tree
// construction; space: O(n*ALPHA); time: O(n)
// Real tree nodes are called dedicated nodes.
// "Nodes" lying on compressed edges are called
// implicit nodes and are represented
// as pairs (lower node, label index).
// Labels are represented as intervals [L;R)
// which refer to substrings [L;R) of txt.
// Leaves have labels of form [L:infinity].
// use getR to get current right endpoint.
// Suffix links are valid only for internal
// nodes (non-leaves).
struct SufTree {
 Vi txt: // Text for which tree is built
 // to [v][c] = edge with label starting with c
               from node v
 vector<array<int, ALPHA>> to{ {} };
 Vi L{0}, R{0}; // Parent edge label endpoints
 Vi par{0};
              // Parent link
 Vi link{O};
               // Suffix link
 Pii cur{0, 0}; // Current state
 // Get current right end of node label
 int getR(int i) { return min(R[i],sz(txt)); }
```

```
// Follow edge `e` of implicit node `s`.
  // Returns (-1, -1) if there is no edge.
  Pii next (Pii s, int e) {
   if (s.v < qetR(s.x))
     return txt[s.v] == e ? mp(s.x, s.v+1)
                           : mp(-1, -1);
    e = to[s.x][e];
   return e ? mp(e, L[e]+1) : mp(-1, -1);
  // Create dedicated node for implicit node
  // and all its suffixes
  int split(Pii s) {
   if (s.y == R[s.x]) return s.x;
   int t = sz(to); to.pb({});
   to[t][txt[s.y]] = s.x;
   L.pb(L[s.x]);
   R.pb(L[s.x] = s.y);
   par.pb(par[s.x]);
   par[s.x] = to[par[t]][txt[L[t]]] = t;
   link.pb(-1);
    int v = link[par[t]], l = L[t] + !par[t];
    while (1 < R[t]) {
     v = to[v][txt[1]];
     1 += getR(v) - L[v];
   v = split(\{v, getR(v)-l+R[t]\});
   link[t] = v;
   return t;
  // Append letter from [0;ALPHA) to the back
  void add(int x) { // amoritzed time: 0(1)
   Pii t; txt.pb(x);
    while ((t = next(cur, x)).x == -1) {
     int m = split(cur);
     to[m][x] = sz(to);
     to.pb({});
     par.pb(m);
     L.pb(sz(txt)-1);
     R.pb(INT MAX);
     link.pb(-1);
     cur = {link[m], getR(link[m])};
     if (!m) return;
    cur = t:
} ;
```

text/z function.h

```
// Computes Z function array; time: O(n)
// zf[i] = max common prefix of str and str[i:]
template < class T > Vi prefPref (const T& str) {
 int n = sz(str), b = 0, e = 1;
  Vi zf(n);
  rep(i, 1, n) {
   if (i < e) zf[i] = min(zf[i-b], e-i);</pre>
    while (i+zf[i] < n &&
     str[zf[i]] == str[i+zf[i]]) zf[i]++;
    if (i+zf[i] > e) b = i, e = i+zf[i];
  zf[0] = n;
  return zf:
```

trees/centroid decomp.h

```
// Centroid decomposition; space: O(n lg n)
// UNTESTED
struct CentroidTree {
 // child[v] = children of v in centroid tree
 // par[v] = parent of v in centroid tree
              (-1 for root)
 // depth[v] = depth of v in centroid tree
               (0 for root)
 // ind[v][i] = index of vertex v in i-th
                centroid subtree from root
 // size[v] = size of centroid subtree of v
 // subtree[v] = list of vertices
                 in centroid subtree of v
 // dists[v] = distances from v to vertices
               in its centroid subtree
                (in the order of subtree[v])
 // neigh[v] = neighbours of v
               in its centroid subtree
 // dir[v][i] = index of centroid neighbour
                that is first vertex on path
                from centroid v to i-th vertex
                of centroid subtree
                 (-1 for centroid)
 vector<Vi> child, ind, dists, subtree,
            neigh, dir:
 Vi par, depth, size;
 int root; // Root centroid
 CentroidTree() {}
 CentroidTree (vector<Vi>& G)
     : child(sz(G)), ind(sz(G)), dists(sz(G)),
       subtree(sz(G)), neigh(sz(G)),
       dir(sz(G)), par(sz(G), -2),
       depth(sz(G)), size(sz(G)) {
   root = decomp(G, 0, 0);
 int dfs(vector<Vi>& G, int v, int p) {
   size[v] = 1;
   each (e, G[v]) if (e != p && par[e] == -2)
     size[v] += dfs(G, e, v);
   return size[v];
 void layer(vector<Vi>& G, int v,
            int p, int c, int d) {
   ind[v].pb(sz(subtree[c]));
   subtree[c].pb(v);
   dists[c].pb(d);
   dir[c].pb(sz(neigh[c])-1);
   each(e, G[v]) if (e != p && par[e] == -2) {
     if (v == c) neigh[c].pb(e);
     laver(G, e, v, c, d+1);
   }
 }
 int decomp(vector<Vi>& G, int v, int d) {
   int p = -1, s = dfs(G, v, -1);
   bool ok = 1:
   while (ok) {
     ok = 0;
     each (e, G[v]) {
```

```
if (e != p && par[e] == -2 &&
           size[e] > s/2) {
         p = v; v = e; ok = 1;
         break;
     }
   par[v] = -1;
   size[v] = s;
   depth[v] = d;
   layer(G, v, -1, v, 0);
   each(e, G[v]) if (par[e] == -2) {
     int j = decomp(G, e, d+1);
     child[v].pb(j);
     par[j] = v;
   return v:
 }
};
```

```
trees/heavylight decomp.h
#include "../structures/segment_tree_point.h"
// Heavy-Light Decomposition of tree
// with subtree query support; space: O(n)
struct HLD (
 // Subtree of v = [pos[v]; pos[v]+size[v])
 // Chain with v = [chBegin[v]; chEnd[v])
             // Vertex parent
 Vi par:
 Vi size:
             // Vertex subtree size
 Vi depth: // Vertex distance to root
 Vi pos:
             // Vertex position in "HLD" order
 Vi chBegin; // Begin of chain with vertex
 Vi chEnd: // End of chain with vertex
 Vi order; // "HLD" preorder of vertices
 SegmentTree tree; // Verts are in HLD order
 HLD() {}
 // Initialize structure for tree G
 // and given root; time: O(n lg n)
 // MODIFIES ORDER OF EDGES IN G!
 HLD(vector<Vi>& G, int root)
     : par(sz(G)), size(sz(G)),
       depth(sz(G)), pos(sz(G)),
       chBegin(sz(G)), chEnd(sz(G)) {
   dfs(G, root, -1);
   decomp(G, root, -1, 0);
   tree = {sz(order)};
 void dfs(vector<Vi>& G, int v, int p) {
   par[v] = p;
   size[v] = 1;
   depth[v] = p < 0 ? 0 : depth[p]+1;
   int& fs = G[v][0];
   if (fs == p) swap(fs, G[v].back());
   each(e, G[v]) if (e != p) {
     dfs(G, e, v);
     size[v] += size[e];
     if (size[e] > size[fs]) swap(e, fs);
```

```
void decomp(vector<Vi>& G.
            int v, int p, int chb) {
  pos[v] = sz(order);
  chBegin[v] = chb;
  chEnd[v] = pos[v]+1;
  order.pb(v);
  each(e, G[v]) if (e != p) {
    if (e == G[v][0]) {
      decomp(G, e, v, chb);
      chEnd[v] = chEnd[e];
    } else {
      decomp(G, e, v, sz(order));
}
// Get root of chain containing v
int chRoot(int v) {return order[chBegin[v]];}
// Level Ancestor Query; time: O(lq n)
int lag(int v, int level) {
  while (true) {
    int k = pos[v] - depth[v] + level;
    if (k >= chBegin[v]) return order[k];
    v = par[chRoot(v)];
// Lowest Common Ancestor; time: O(lg n)
int lca(int a, int b) {
  while (chBegin[a] != chBegin[b]) {
    int ha = chRoot(a), hb = chRoot(b);
    if (depth[ha] > depth[hb]) a = par[ha];
    else b = par[hb];
  return depth[a] < depth[b] ? a : b;</pre>
// Call func(chBegin, chEnd) on each path
// segment; time: O(lq n * time of func)
template<class T>
void iterPath(int a, int b, T func) {
  while (chBegin[a] != chBegin[b]) {
    int ha = chRoot(a), hb = chRoot(b);
    if (depth[ha] > depth[hb]) {
      func (chBegin[a], pos[a]+1);
      a = par[ha];
    } else {
      func(chBegin[b], pos[b]+1);
      b = par[hb];
  if (pos[a] > pos[b]) swap(a, b);
  // Remove +1 from pos[a]+1 for vertices
  // queries (with +1 -> edges).
  func (pos[a]+1, pos[b]+1);
// Ouerv path between a and b: O(1g^2 n)
SegmentTree::T guervPath(int a, int b) {
  auto ret = SegmentTree::ID;
  iterPath(a, b, [&](int i, int j) {
```

```
ret = SegmentTree::merge(ret,
          tree.querv(i, j));
    1);
    return ret;
  // Ouery subtree of v; time: O(lq n)
  SeamentTree::T quervSubtree(int v) {
    return tree.query(pos[v], pos[v]+size[v]);
};
trees/lca.h
// LAO and LCA using jump pointers
// space: O(n lg n)
struct LCA {
  vector<Vi> jumps;
  Vi level, pre, post;
  int cnt{0}, depth;
  LCA() {}
  // Initialize structure for tree G
  // and root r; time: O(n lg n)
  LCA (vector < Vi>& G, int root)
      : jumps(sz(G)), level(sz(G)),
        pre(sz(G)), post(sz(G)) {
    dfs(G, root, root);
    depth = int(log2(sz(G))) + 2;
    rep(j, 0, depth) each(v, jumps)
     v.pb(jumps[v[j]][j]);
  void dfs(vector<Vi>& G, int v, int p) {
    level[v] = p == v ? 0 : level[p]+1;
    jumps[v].pb(p);
   pre[v] = ++cnt;
   each(e, G[v]) if (e != p) dfs(G, e, v);
   post[v] = ++cnt;
  // Check if a is ancestor of b; time: O(1)
  bool isAncestor(int a, int b) {
    return pre[a] <= pre[b] &&
           post[b] <= post[a];</pre>
  // Lowest Common Ancestor; time: O(lq n)
  int operator() (int a, int b) {
    for (int j = depth; j--;)
     if (!isAncestor(jumps[a][j], b))
        a = jumps[a][j];
    return isAncestor(a, b) ? a : jumps[a][0];
  // Level Ancestor Ouery; time: O(lq n)
  int lag(int a, int lvl) {
    for (int j = depth; j--;)
     if (lvl <= level[jumps[a][j]])</pre>
       a = jumps[a][j];
    return a;
```

// Get distance from a to b; time: O(lq n)

int distance(int a, int b) {

```
return level[a] + level[b] -
           level[operator()(a, b)]*2;
 1
};
trees/link cut tree.h
// Link/cut tree; space: O(n)
// Represents forest of (un)rooted trees.
struct LinkCutTree {
 vector<array<int, 2>> child;
 Vi par, prev, flip;
  // Initialize structure for n vertices; O(n)
  // At first there's no edges.
 LinkCutTree(int n = 0)
     : child(n, \{-1, -1\}), par(n, -1),
        prev(n, -1), flip(n, -1) {}
 void auxLink(int p, int i, int ch) {
    child[p][i] = ch;
   if (ch >= 0) par[ch] = p;
 void push(int x) {
   if (x >= 0 && flip[x]) {
      flip[x] = 0;
     swap(child[x][0], child[x][1]);
     each(e, child[x]) if (e>=0) flip[e] ^= 1;
 }
 void rot(int p, int i) {
   int x = \text{child}[p][i], q = \text{par}[x] = \text{par}[p];
    if (q \ge 0) child[q][child[q][1] == p] = x;
    auxLink(p, i, child[x][!i]);
    auxLink(x, !i, p);
    swap(prev[x], prev[p]);
 void splay(int x) {
    while (par[x] >= 0) {
      int p = par[x], g = par[p];
      push(g); push(p); push(x);
     bool f = (child[p][1] == x);
      if (q >= 0) {
        if (\text{child}[q][f] == p) { <math>// ziq - ziq
          rot(g, f); rot(p, f);
        } else { // zig-zag
          rot(p, f); rot(q, !f);
      } else { // zig
        rot(p, f);
   push(x);
  void access(int x) {
   while (true) {
      splav(x):
      int p = prev[x];
     if (p < 0) break;
      prev[x] = -1;
      splav(p):
```

```
auxLink(p, 1, x);
 1
 void makeRoot(int x) {
   access(x);
   int& l = child[x][0];
   if (1 >= 0) {
     swap(par[1], prev[1]);
     flip[1] ^= 1;
     1 = -1:
   }
 }
 // Find representative of tree containing x
 int find(int x) { // time: amortized O(lg n)
   access(x):
   while (child[x][0] >= 0)
     push(x = child[x][0]);
   splay(x);
   return x:
 // Add edge x-y; time: amortized O(lg n)
 void link(int x, int y) {
   makeRoot(x); prev[x] = y;
 // Remove edge x-y; time: amortized O(lq n)
 void cut(int x, int y) {
   makeRoot(x); access(v);
   par[x] = child[y][0] = -1;
};
util/arc interval cover.h
using dbl = double;
// Find size of smallest set of points
// such that each arc contains at least one
// of them; time: O(n lq n)
int arcCover(vector<pair<dbl, dbl>>& inters,
            dbl wrap) {
 int n = sz(inters);
 rep(i, 0, n) {
   auto& e = inters[i]:
   e.x = fmod(e.x, wrap);
   e.y = fmod(e.y, wrap);
   if (e.x < 0) e.x += wrap, e.y += wrap;</pre>
   if (e.x > e.v) e.x += wrap;
   inters.pb({e.x+wrap, e.y+wrap});
 Vi nxt(n);
 deque<dbl> que;
 dbl r = wrap*4;
 sort(all(inters)):
 for (int i = n*2-1; i--;) {
   r = min(r, inters[i].v);
   que.push front(inters[i].x);
   while (!que.emptv() && que.back() > r)
```

int r = child[p][1];

if $(r \ge 0)$ swap(par[r], prev[r]);

```
que.pop back();
    if (i < n) nxt[i] = i+sz(que);
  int a = 0, b = 0;
  do {
    a = nxt[a] % n;
    b = nxt[nxt[b]%n] % n;
 } while (a != b);
  int ans = 0:
  while (b < a+n) {
    b += nxt[b%n] - b%n;
    ans++;
 return ans:
util/bit hacks.h
// builtin popcount - count number of 1 bits
// builtin clz - count most significant 0s
// builtin ctz - count least significant Os
// __builtin_ffs - like ctz, but indexed from 1
                  returns 0 for 0
// For 11 version add 11 to name
using ull = uint64 t;
#define T64(s,up)
 for (ull i=0; i<64; i+=s*2)
    for (ull j = i; j < i+s; j++) {
      ull \ a = (M[j] >> s) \& up;
      ull \ b = (M[j+s] \& up) << s;
      M[j] = (M[j] \& up) / b;
      M[j+s] = (M[j+s] & (up << s)) | a; 
// Transpose 64x64 bit matrix
void transpose64 (arrav<ull, 64>& M) {
 T64(1, 0x55555555555555);
 T64(2, 0x33333333333333333);
  T64 (4, 0xF0F0F0F0F0F0F0F);
  T64(8, 0xFF00FF00FF00FF);
  T64 (16, 0xFFFF0000FFFF);
 T64 (32. OxFFFFFFFLL):
// Lexicographically next mask with same
// amount of ones.
int nextSubset(int v) {
 int t = v | (v - 1);
  return (t + 1) | (((~t & -~t) - 1) >>
      (\underline{\phantom{a}}builtin_ctz(v) + 1));
util/bump alloc.h
// Allocator, which doesn't free memory.
char mem[400<<20]; // Set memory limit</pre>
size t nMem;
void* operator new(size_t n) {
 nMem += n; return &mem[nMem-n];
void operator delete(void*) {}
```

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util/compress vec.h

```
// Compress integers to range [0;n) while
// preserving their order; time: O(n lg n)
// Returns mapping: compressed -> original
Vi compressVec(vector<int*>& vec) {
    sort(all(vec),
      [](int* l, int* r) { return *l < *r; });
    Vi old;
    each(e, vec) {
      if (old.empty() || old.back() != *e)
         old.pb(*e);
    *e = sz(old)-1;
    }
    return old;
}</pre>
```

util/inversion_vector.h

```
// Get inversion vector for sequence of
// numbers in [0;n); ret[i] = count of numbers
// smaller than perm[i] to the left; O(n lg n)
Vi encodeInversions (Vi perm) {
 Vi odd, ret(sz(perm));
 int cont = 1;
  while (cont) {
   odd.assign(sz(perm)+1, 0);
    cont = 0;
    rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret[i] += odd[perm[i]+1];
     cont += perm[i] /= 2;
   }
  return ret;
// Count inversions in sequence of numbers
// in [0;n); time: O(n lg n)
11 countInversions(Vi perm) {
 ll ret = 0, cont = 1;
 Vi odd;
  while (cont) {
   odd.assign(sz(perm)+1, 0);
    cont = 0;
    rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret += odd[perm[i]+1];
     cont += perm[i] /= 2;
 return ret:
```

util/longest_increasing_sub.h

util/max rects.h

```
struct MaxRect {
 // begin = first column of rectangle
 // end = first column after rectangle
 // hei = height of rectangle
 // touch = columns of height hei inside
 int begin, end, hei;
 Vi touch; // sorted increasing
};
// Given consecutive column heights find
// all inclusion-wise maximal rectangles
// contained in "drawing" of columns; time O(n)
vector < MaxRect > getMaxRects (Vi hei) {
 hei.insert(hei.begin(), -1);
 hei.pb(-1);
 Vi reach (sz (hei), sz (hei) -1);
  vector < MaxRect > ans;
  for (int i = sz(hei)-1; --i;) {
   int j = i+1, k = i;
    while (hei[j] > hei[i]) j = reach[j];
    reach[i] = j;
    while (hei[k] > hei[i-1]) {
      ans.pb(\{ i-1, 0, hei[k], \{\} \});
      auto& rect = ans.back();
      while (hei[k] == rect.hei) {
       rect.touch.pb(k-1);
        k = reach[k];
     rect.end = k-1;
   }
 }
 return ans;
```

util/mo.h

```
// Modified MO's queries sorting algorithm,
// slightly better results than standard.
// Allows to process q queries in O(n*sqrt(q))
struct Query {
  int begin, end;
};
// Get point index on Hilbert curve
11 hilbert(int x, int y, int s, ll c = 0) {
 if (s <= 1) return c:
  s /= 2; c *= 4;
  if (y < s)
    return hilbert (x \in (s-1), y, s, c+(x>=s)+1);
    return hilbert (2*s-y-1, s-x-1, s, c);
  return hilbert(y-s, x-s, s, c+3);
// Get good order of gueries; time: O(n lg n)
Vi moOrder (vector < Query>& queries, int maxN) {
  int s = 1:
  while (s < maxN) s \star= 2;
  vector<ll> ord;
  each (q, queries)
    ord.pb(hilbert(q.begin, q.end, s));
```

```
Vi ret(sz(ord));
iota(all(ret), 0);
sort(all(ret), [&](int l, int r) {
    return ord[1] < ord[r];
});
return ret;
}</pre>
```

util/parallel_binsearch.h

```
// Run `count` binary searches on [begin; end),
// `cmp` arguments:
// 1) vector<Pii>& - pairs (value, index)
      which are queries if value of index is
     greater or equal to value,
     sorted by value
// 2) vector<bool>& - true at index i means
     value of i-th guery is >= gueried value
// Returns vector of found values.
// Time: O((n+c) lg n), where c is cmp time.
template<class T>
Vi multiBS(int begin, int end, int count, T cmp) {
 vector<Pii> ranges(count, {begin, end});
 vector<Pii> queries(count);
 vector<bool> answers(count);
  rep(i, 0, count) queries[i]={ (begin+end) /2, i};
  for (int k = uplg(end-begin); k > 0; k--) {
    int last = 0, j = 0;
    cmp (queries, answers);
    rep(i, 0, sz(queries)) {
      Pii &g = queries[i], &r = ranges[q.y];
      if (q.x != last) last = q.x, j = i;
      (answers[i] ? r.x : r.y) = q.x;
      q.x = (r.x+r.v) / 2;
      if (!answers[i])
        swap(queries[i], queries[j++]);
  each (p, ranges) ret.pb(p.x);
  return ret;
```

util/radix sort.h

```
// Stable countingsort; time: O(k+sz(vec))
// See example usage in radixSort for pairs.
template<class F>
void countSort(Vi& vec, F key, int k) {
    static Vi buf, cnt;
    vec.swap(buf);
    vec.resize(sz(buf));
    cnt.assign(k+1, 0);
    each(e, buf) cnt[key(e)]++;
    rep(i, 1, k+1) cnt[i] += cnt[i-1];
    for (int i = sz(vec)-1; i >= 0; i--)
        vec[--cnt[key(buf[i])]] = buf[i];
}

// Compute order of elems, k is max key; O(n)
```

```
Vi radixSort(const vector<Pii>& elems, int k) {
  Vi order(sz(elems));
  iota(all(order), 0);
  countSort(order,
    [&](int i) { return elems[i].y; }, k);
  countSort(order,
    [&](int i) { return elems[i].x; }, k);
  return order;
}
```