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```
.bashrc
                                          bdbc
b() ( g++ -DLOC -O2 -std=c++20 -Wall -W
→ -Wfatal-errors -Wconversion -Wshadow
→ -Wlogical-op -Wfloat-equal -o $1.e $@ )
d() ( b $@ -OO -q -D GLIBCXX DEBUG

→ -fsanitize=address, undefined )

run()( $@ && echo start >&2 && time ./$2.e )
                                          68f6
.vimrc
se ai cin cul ic is nu scs sw=4 ts=4 so=7 ttm=9
vn _ :w !cpp -dD -P -fpreprocessed \|
\hookrightarrow sed -z sg\\sggg \| md5sum \| cut -c-4 <cr>
                                          1989
template.cpp
#include <bits/stdc++.h>
using namespace std;
using 11 = long long;
using vi = vector<int>;
using pii = pair<int,int>;
#define pb push back
#define x first
#define v second
#define rep(i,b,e) for (int i=(b); i<(e); i++)
#define each (a, x) for (auto \& a : (x))
#define all(x)
                   (x).begin(), (x).end()
#define sz(x)
                   (int)(x).size()
#define PP(x,y) auto operator<<(auto&o, auto a)</pre>
→ ->decltype(y,o) {o<<"("; x; return o<<")";}</pre>
PP(a.print(), a.print());
PP (o << a.x << ", " << a.y, a.y);
PP (for (auto i : a) o << i << ", ", all(a));
void DD(auto s, auto... k) {
  [&]
    while (cerr << *s++, 45 % ~*s);</pre>
    cerr << ": " << k;
  }(), ...);
} // 606c
#ifdef LOC
auto SS = signal(6, [](int) { *(int*)0=0; });
#define deb(x...)
→ DD(":, "#x, __LINE__, x), cerr << endl</pre>
#else
#define deb(...)
#endif
#define DBP(x...) void print() { DD(\#x, x); }
int main() {
 cin.tie(0)->sync_with_stdio(0);
  cout << fixed << setprecision(10);</pre>
} // 04a0
                                          9841
template.java
import java.io.*;
import java.util.*;
public class Task extends PrintWriter {
  BufferedReader reader = new BufferedReader(
     new InputStreamReader(System.in), 32768); | using namespace rel_ops;
```

```
StringTokenizer tok:
  public static void main(String[] a) {
   try (Task t = new Task()) { t.solve(); }
  } // ffa0
  Task() { super(System.out); }
  String scan() {
    while (tok == null || !tok.hasMoreTokens())
        tok = new StringTokenizer(
          reader.readLine());
      } catch (Exception e) {
       throw new RuntimeException(e):
     } // fdda
   return tok.nextToken();
 } // 969e
  int scanInt() {
   return Integer.parseInt(scan());
 ) // 4a91
 void solve() {
   int n = scanInt();
   printf("hello %d", n);
 } // 3fbe
} // 0966
various.bash
loo()( # loo b/d prog.cpp gen.cpp
 set -e; $1 $2; $1 $3
 for ((;;)) {
   ./$3.e > gen.in
   time ./$2.e < gen.in > gen.out
cmp()( # cmp b/d prog.cpp brute.cpp gen.cpp
 set -e; $1 $2; $1 $3; $1 $4
 for ((;;)) {
    ./$4.e > gen.in;
                             echo -n 0
    ./$2.e < gen.in > p1.out; echo -n 1
    ./$3.e < gen.in > p2.out; echo -n 2
   diff pl.out p2.out
# Other compilation flags:
# -Wformat=2 -Wshift-overflow=2 -Wcast-qual
# -Wcast-align -Wduplicated-cond
# -D_GLIBCXX_DEBUG_PEDANTIC -D_FORTIFY_SOURCE=2
# -fno-sanitize-recover -fstack-protector
# -fopt-info-all -fopt-info-missed
various.h
// If math constants like M PI are undefined:
#define USE MATH DEFINES
// Pragmas
#pragma GCC optimize("Ofast, unroll-loops")
#pragma GCC target("arch=???, tune=???")
#define GLIBCXX GTHREAD USE WEAK 0
// Exit without calling destructors
cout << flush; _Exit(0);</pre>
// Clock
while (clock() < duration *CLOCKS_PER_SEC)
// Automatically implement operators:
// 1. != if == is defined
// 2. >, <= and >= if < is defined
```

```
// Mersenne twister for randomization.
mt19937_64 rnd(chrono::steady_clock::now()
 .time_since_epoch().count());
// To shuffle randomly use:
shuffle(all(vec), rnd);
// To pick random integer from [A;B] use:
uniform_int_distribution <> dist(A, B);
int value = dist(rnd);
// To pick random real number from [A;B] use:
uniform_real_distribution  dist(A, B);
double value = dist(rnd);
// Floats can represent integers up to 19*10^6
// Doubles can represent integers up to 9*10^15
// lg(x) == floor(log2(x)), undefined for x=0
various.pv
input().split(' ') # Read and split into words
print('abc', end='') # Print without newline
>>> from fractions import *
>>> Fraction (16, -10)
Fraction (-8, 5)
>>> Fraction (123)
Fraction (123, 1)
>>> Fraction()
Fraction (0, 1)
>>> Fraction('3/7')
Fraction (3, 7)
>>> Fraction(' -3/7 ')
Fraction (-3, 7)
>>> Fraction('1.414213 \t\n')
>>> Fraction('7e-6')
Fraction (7, 1000000)
>>> Fraction (2.25)
Fraction (9, 4)
>>> Fraction (1.1)
Fraction (2476979795053773, 2251799813685248)
>>> Fraction('1/2') * Fraction('4/3')
Fraction(2, 3)
>>> Fraction (16, 5).numerator
>>> Fraction (16, 5).denominator
>>> from decimal import *
>>> getcontext().prec = 28
>>> Decimal (10)
Decimal('10')
>>> Decimal('3.14')
Decimal('3.14')
>>> Decimal (3.14)
Decimal('3.140000000000001243449787580175...')
>>> Decimal((0, (3, 1, 4), -2))
Decimal('3.14')
>>> Decimal(str(2.0 ** 0.5))
Decimal('1.4142135623730951')
>>> Decimal(2) ** Decimal('0.5')
Decimal('1.414213562373095048801688724')
>>> Decimal('NaN')
Decimal('NaN')
>>> Decimal('-Infinity')
Decimal ('-Infinity')
geo2d/circle.h
                                          3dc2
#include "vector.h"
#include "line.h" // For line intersections.
```

```
// 2D circle structure; UNIT-TESTED
struct circle {
           // Center
 vec p;
 sc r2 = 0; // Squared radius
 DBP (p, r2);
 // Returns -1 if point g lies outside circle,
 // 0 if on the edge, 1 if strictly inside.
 // Depends on vec: -, len2
 int side(vec a) {
   return sqn(r2 - (p-a).len2());
 } // f399
#if FLOATING_POINT_GEOMETRY
 // Intersect with another circle.
 // Returns number of intersection points
 // (3 means circles are identical).
 // Arc is CCW w.r.t to 'this', CW for 'a'.
 // Depends on vec: +, -, *, len2, perp
 int intersect(circle a, pair<vec, vec>& out) {
   vec d = a.p - p;
   sc d2 = d.len2():
   if (!sqn(d2)) return !sqn(r2-a.r2) * 3;
   sc pd = (d2+r2-a.r2)/2, h2 = r2-pd*pd/d2;
   vec h, t = p + d*(pd/d2);
   int s = sgn(h2)+1;
   if (s > 1) h = d.perp() * sqrt(h2/d2);
   out = \{t-h, t+h\};
   return s;
 1 // 8289
 // Intersect with line.
 // Returns number of intersection points.
 // Points are in order given by a.v.perp().
 int intersect(line a, pair<vec, vec>& out) {
   sc d = a.dist(p), h2 = r2 - d*d;
   vec h, t = a.proj(p);
   int s = sqn(h2)+1;
   if (s > 1)
     h = a.v.perp() * sqrt(h2 / a.v.len2());
   out = \{t-h, t+h\};
   return s;
 } // 1ee2
 // Find normal vectors of tangents.
 // Returns number of tangent points
 // (3 means circle is degenerated to `a`).
 // Covered arc is CCW between vectors.
 int tangents(vec a, pair<vec, vec>& out) {
   vec d = a - p;
   sc d2 = d.len2(), h2 = d2 - r2;
   if (!sqn(d2)) return !sqn(h2) * 3;
   vec h, t = d * sqrt(r2);
   int s = sgn(h2)+1;
   if (s > 1) h = d.perp() * sqrt(h2);
   out = \{(t-h)/d2, (t+h)/d2\};
   return s;
 } // 9bf4
 // Find normal vectors of tangents.
 // Returns number of tangent points
 // (3 means circles are identical).
 // Arc for 'this' is CCW between vectors.
 // For `a`, it is CW for outer, CCW for inner
 int tangents (circle a, bool inner,
              pair<vec, vec>& out) {
   vec d = a.p - p;
   sc d2 = d.len2();
   sc dr = sgrt(r2) + sgrt(a.r2) * (inner*2-1);
   sc h2 = d2 - dr*dr;
```

```
vec h, t = d * dr;
    int s = sgn(h2)+1;
    if (s > 1) h = d.perp() * sqrt(h2);
   out = \{(t-h)/d2, (t+h)/d2\};
   return s:
 } // 55a9
#endif
}; // 3f47
#if FLOATING_POINT_GEOMETRY
// Circumcircle. Points must be non-aligned.
// Depends on vec: +,-,*,/, cross, len2, perp
circle circum(vec a, vec b, vec c) {
 b = b-a; c = c-a;
  sc s = b.cross(c);
  assert (sqn(s));
  vec p = a+(b*c.len2()-c*b.len2()).perp()/s/2;
  return { p, (p-a).len2() };
} // fbaa
#endif
geo2d/circle min.h
                                          9444
#include "vector.h" // FLOATING POINT GEOMETRY
#include "circle.h"
mt19937 rnd(123);
// Minimum circle enclosing a set of points.
circle minDisk(vector<vec> p) { // time: O(n)
  shuffle(all(p), rnd);
  rep(i, 0, sz(p)) if (c.side(p[i]) < 0) {
   c = \{p[i], 0\};
   rep(j, 0, i) if (c.side(p[j]) < 0) {
     c = \{(p[i]+p[j])/2, (p[i]-p[j]).len2()/4\};
     rep(k, 0, j) if (c.side(p[k]) < 0)
       c = circum(p[i], p[j], p[k]);
   } // 6e1c
  } // 458a
 return c;
} // 38d5
geo2d/convex hull.h
                                          8159 | } // 6d03
#include "vector.h"
// Find convex hull of points; time: O(n lg n)
// Points are returned counter-clockwise,
// first point is the bottom-leftmost.
// Depends on vec: -, cross, cmpXY
vector<vec> convexHull(vector<vec> points) {
  if (sz(points) <= 1) return points;</pre>
  sort(all(points), [](vec 1, vec r) {
   return l.cmpYX(r) < 0;</pre>
  }); // 4b89
  vector<vec> h(sz(points)+1);
  int s = 0, t = 0;
  rep(i, 0, 2) {
   each (p, points) {
     for (; t >= s+2; t--)
       if ((p-h[t-2]).cross(p-h[t-1]) > eps)
         break:
     h[t++] = p;
    1 // 9306
    reverse(all(points));
   s = --t:
  1 // 3419
  h.resize(t - (t == 2 \& h[0] == h[1]);
  return h;
} // 349e
```

```
// Find point p that maximizes dot product p*q.
// Returns point index in hull; time: O(lg n)
// If multiple points have same dot product
// one with smallest index is returned.
// Points are expected to be in the same order
// as output from convexHull function.
// Depends on vec: -, cross, perp, upper, cmpAngle
int maxDot(const vector<vec>& h, vec q) {
 int b = 0, e = sz(h);
  while (b+1 < e) {
   int m = (b+e) / 2;
    vec s = h[m] - h[m-1];
    (q.perp().cmpAngle(s) > 0 ? b : e) = m;
 return q.dot(h[b]-h[0]) > eps ? b : 0;
} // 26f4
#include "segment.h"
// Get distance from point to a hull; O(lq n)
// Returns -1 if point is strictly inside.
// Points are expected to be in the same order
// as output from convexHull function.
// Depends on: maxDot
// Depends on vec: -, dot, cross, len, perp,
                  upper, cmpAngle
// Depends on seg: side, dist
double hullDist(const vector<vec> € h, vec q) {
 if (sz(h) == 1) return (q-h[0]).len();
  int b = (h[0]-g).upper() ? maxDot (h, \{0,1\}):0;
 int n = sz(h), e = b+n;
  vec p = h[b]:
  while (b+1 < e) {
   int m = (b+e) / 2;
    vec s = h[m%n], t = h[++m%n];
    sc x = ((s-t).cross(q-s) < -eps ?
      (q-p).cross(s-p) : (s-t).dot(q-s));
    (sgn(x) < !(m%n) ? b : e) = m-1;
 1 // eba8
  seg s{h[b%n], h[e%n]}, t{h[e%n], h[++e%n]};
  return s.side(q) + t.side(q) < 2 ?</pre>
   s.dist(q) : -1;
geo2d/convex hull sum.h
                                          a 935
#include "vector.h"
// Minkowski sum of given convex polygons.
// Points are expected to be in the same order
// as output from convexHull function; O(n+m)
// Depends on vec: +, -, cross, upper, cmpAngle
vector<vec> hullSum(const vector<vec>& A,
                    const vector<vec>& B) {
  int n = sz(A), m = sz(B), i = 0, j = 0;
 if (!n || !m) return {};
 vector<vec> C = \{A[0]+B[0]\};
  while (i+j < n+m) {
   vec a = A[(i+1)%n] - A[i%n];
    vec b = B[(j+1)%m] - B[j%m], v = C.back();
    int s = (i==n) - (j==m) ?: a.cmpAngle(b);
   if (s <= 0) v = v+a, i++;
   if (s >= 0) v = v+b, j++;
   C.pb(v);
 } // f93a
 C.pop_back();
 return C;
} // 00a3
geo2d/delaunay.h
#include "vector.h"
```

```
#include "../geo3d/convex hull.h"
// Delaunay triangulation using 3D convex hull.
// Faces are in CCW order. Doesn't work if
// all points are colinear or on a same circle!
// time and memory: O(n log n)
vector<Triple> delaunay(vector<vec>& p) {
 assert(sz(p) >= 3);
 if (sz(p) == 3) {
    int d = ((p[1]-p[0]).cross(p[2]-p[0]) < 0);
    return {{0, 1+d, 2-d}};
  } // 5c2e
  vector<vec3> p3;
  each(e, p) p3.pb({e.x, e.y, e.len2()});
  auto hull = convexHull(p3);
  erase_if(hull, [&](auto& t) {
   vec a = p[t[0]], b = p[t[1]], c = p[t[2]];
    swap(t[1], t[2]);
    return (b-a).cross(c-a) > -eps;
  }); // 794c
 return hull;
} // 294b
geo2d/halfplanes.h
                                         714f
#include "vector.h" // FLOATING POINT GEOMETRY
#include "line.h"
// Halfplane intersection; time: O(n lg n)
// Behaviour is undefined if intersection
// is unbounded, add bounding-box if necessary!
// Returns:
// - vertices of intersection area in CCW order
// starting from bottom-leftmost vertex;
// - empty vector if intersection is empty.
// Degenerate cases are supported.
// Works only with floating-point geometry.
// Depends on vec: -, *, /, dot, cross, len,
                  perp, ==, upper, cmpAngle
// Depends on line: side, intersect
vector<vec> intersectHalfs(vector<line> in) {
 sort(all(in), [](line a, line b) {
    return (a.v.perp().cmpAngle(b.v.perp()) ?:
           a.c*b.v.len() - b.c*a.v.len()) < 0;
  }); // 8e73
  int a = 1, b = 1, k = 0, n = sz(in);
  vector<line> dq(n+2);
  vector<vec> out (n+1);
  dq[1] = in[0];
  rep(i, 1, n+1) {
    line t = (i < n ? in[i] : dq[a]);
    while (a < b && t.side(out[b-1]) > 0) b--;
    while (a < b && t.side(out[a]) > 0) a++;
    if (t.intersect(dq[b], out[b])) dq[++b]=t;
 } // fecd
  out[0] = out[--b];
  rep(i, a, b)
   if (out[i] != out[0] && out[i] != out[k])
      out[++k] = out[i];
  out.resize(k+1);
  each(t, in) if (t.side(out[0]) > 0) return{};
 return out;
} // 33ee
geo2d/line.h
                                         03dd
#include "vector.h"
// 2D line/halfplane structure; UNIT-TESTED
```

```
struct line {
  // For lines: v * point == c
  // For halfplanes: v * point <= c
  // (i.e. normal vector points outside)
  vec v; // Normal vector
  sc c = 0; // Offset
  DBP(v, c);
  // Distance from point to line.
  // Depends on vec: dot, len
  double dist(vec a) {
    return fabs (v.dot(a) - c) / v.len();
  } // 79e6
  // Returns 0 if point a lies on the line,
  // 1 if on side where normal vector points.
  // -1 if on the other side.
  // Depends on vec: dot
  int side(vec a) { return sgn(v.dot(a)-c); }
#if FLOATING POINT GEOMETRY
  // Orthogonal projection of point on line.
  // Depends on vec: -, *, dot, len2
  vec proj(vec a) {
    return a - v * ((v.dot(a)-c) / v.len2());
  } // 406e
  // Intersect this line with line a, returns
  // true on success (false if parallel).
  // Intersection point is saved to 'out'.
  // Depends on vec: -, *, /, cross, perp
  bool intersect (line a, vec& out) {
    sc d = v.cross(a.v);
    if (!sqn(d)) return 0;
    out = (v*a.c - a.v*c).perp() / d;
    return 1;
 } // c152
#endif
}; // c1c3
// Line through 2 points with normal vector
// pointing to the right of ab vector.
// Depends on vec: -, cross, perp
line through (vec a, vec b) {
 return { (a-b).perp(), a.cross(b) };
} // 9ac7
// Parallel line through point.
// Depends on vec: dot
line parallel (vec a, line b) {
 return { b.v, b.v.dot(a) };
} // 8e1c
// Perpendicular line through point.
// Depends on vec: cross, perp
line perp(vec a, line b) {
 return { b.v.perp(), b.v.cross(a) };
} // 7b75
geo2d/rmst.h
                                          9cc3
#include "vector.h"
#include "../structures/find union.h"
// Rectilinear Minimum Spanning Tree
// (MST in Manhattan metric); time: O(n lg n)
// Returns MST weight. The spanning tree edges
// are saved in 'out' as triples (dist, (u,v)).
// Depends on vec: -
11 rmst (vector < vec> points,
        vector<pair<ll, pii>>& edges) {
  vector<pair<ll, pii>> span;
  vi id(sz(points));
```

```
iota(all(id), 0);
  rep(k, 0, 4) {
   map<11, 11> S;
    sort(all(id), [&](int i, int j) {
     return (points[i]-points[j]).x <</pre>
             (points[j]-points[i]).y;
    }); // f699
    each(i, id) {
     auto it = S.lower_bound(-points[i].y);
      for (; it != S.end(); S.erase(it++)) {
       vec d = points[i] - points[it->y];
       if (d.y > d.x) break;
        span.push_back(\{d.x+d.y, \{i, it->y\}\});
      } // 490e
     S[-points[i].y] = i;
    } // Oadf
    each (p, points) {
     if (k % 2) p.x = -p.x;
      else swap(p.x, p.y);
   } // 9ec4
  } // 87be
  FAU fau(sz(id));
  11 sum = 0;
  sort (all (span));
  edges.clear();
  each (e, span) if (fau.join(e.y.x, e.y.y))
   edges.pb(e), sum += e.x;
  return sum;
} // b2f5
geo2d/segment.h
                                          08b3
#include "vector.h"
// 2D segment structure; UNIT-TESTED
struct seq {
  vec a, b; // Endpoints
  DBP (a, b);
  // Check if segment contains point p.
  // Depends on vec: -, dot, cross
  bool contains(vec p) {
    return (a-p).dot(b-p) <= eps &&
           !sgn((a-p).cross(b-p));
  // Returns 0 if point p lies on the line ab,
  // 1 if to the left of the vector ab,
  // -1 if on the right of the vector ab.
  // Depends on vec: cross
  int side(vec p) {
   return sqn((b-a).cross(p-a));
  } // 20a4
  // Distance from segment to point.
  // Depends on vec: -, dot, cross, len
  double dist (vec p) const {
    if ((p-a).dot(b-a) \le eps)
     return (p-a).len();
    if ((p-b).dot(a-b) <= eps)
      return (p-b).len();
    return double (abs ((p-a).cross(b-a))) /
      (b-a).len();
  } // a4c6
#if not FLOATING POINT GEOMETRY
  // Compare distance to p with sgrt(d2).
  // -1 if smaller, 0 if equal, 1 if greater
  // Depends on vec: -, dot, cross, len2
  int cmpDist(vec p, 11 d2) const {
   if ((p-a).dot(b-a) <= 0)</pre>
```

```
return sqn((p-a).len2()-d2);
    if ((p-b).dot(a-b) <= 0)</pre>
      return sqn((p-b).len2()-d2);
    11 c = (p-a).cross(b-a);
    return sqn(c*c - d2 * (b-a).len2());
 } // 7808
#endif
}; // b2ff
geo2d/vector.h
// Scalar type: float or integer.
#if FLOATING_POINT_GEOMETRY
  using sc = double;
  constexpr sc eps = 1e-9;
  using sc = 11;
  constexpr sc eps = 0;
#endif
// -1 if a < -eps, 1 if a > eps, 0 otherwise
int sqn(sc a) { return (a>eps) - (a < -eps); }</pre>
// 2D point/vector structure; UNIT-TESTED
struct vec {
 using P = vec;
  sc x = 0, y = 0;
  // The following methods are optional
  // and dependencies on them are noted
  // appropriately in library snippets.
  P operator+(P r) const {return{x+r.x,y+r.y};}
  P operator-(P r) const {return{x-r.x,y-r.y};}
  P operator*(sc r) const { return {x*r,y*r}; }
  P operator/(sc r) const { return {x/r,y/r}; }
  sc dot (P r) const { return x*r.x + y*r.y; }
  sc cross (P r) const { return x*r.y - y*r.x; }
  sc len2()
                const { return x*x + y*y; }
  double len() const { return hypot(x, y); }
                const { return {-v,x}; } // CCW
  P perp()
  double angle() const { //[0;2*PI] CCW from OX
    double a = atan2(y, x);
    return (a < 0 ? a+2*M PI : a);
  1 // 7095
  // Equality (with epsilon)
  bool operator == (vec r) const {
    return !sqn(x-r.x) && !sqn(y-r.y);
  // Lexicographic compare by (y,x) (with eps)
  int cmpYX(P r) const {
    return sqn(y-r.y) ?: sqn(x-r.x);
  } // 1f37
  // Is above OX or on its non-negative part?
  bool upper() const {
   return (sgn(y) ?: sgn(x)) >= 0;
  // Compare vectors by angles.
  // Depends on: cross, upper
  int cmpAngle(P r) const {
   return r.upper() -upper() ?: -sgn(cross(r));
  1 // 2ab8
#if FLOATING POINT GEOMETRY
 // Rotate counter-clockwise by given angle.
 P rotate (double a) const {
    return {x*cos(a) - y*sin(a),
            x*sin(a) + y*cos(a); // 1890
```

```
#endif
}; // c380
geo3d/convex hull.h
                                          5f7c
#include "vector.h"
using Triple = array<int, 3>;
mt19937 rnd(123);
// 3D convex hull; time and memory: O(n log n)
// Returns list of hull faces with vertices
// in CCW order when "looking from outside".
// Doesn't work if all points are coplanar!
// Depends on vec3: -, dot, cross, len2
vector<Triple> convexHull(vector<vec3>& in) {
 int n = sz(in), g = 1;
 vector<Triple> ret, fv, fe;
 vector<vi> fb, bad(n);
  vector<vec3> fq, p(n);
  vi dead, ord(n), link(n, -1);
  iota(all(ord), 0);
  shuffle (all (ord), rnd);
  rep(i, 0, n) p[i] = in[ord[i]];
  // Only needed if there are 4 coplanar points
  vec3 = p[0], b, c;
  rep(i, 1, n) if (q < 4) {
    swap(p[g], p[i]); swap(ord[g], ord[i]);
    if (q == 1)
     g += sgn((b = p[1]-a).len2());
    else if (q == 2)
      g \leftarrow sgn((c = b.cross(p[2]-a)).len2());
    else
      q += !!sqn(c.dot(p[3]-a));
  } // 633d
  assert(q == 4); // Not everything coplanar
  auto add = [&](int i, int j, int k) {
    fv.pb(\{i, j, k\});
    fe.pb(\{-1, -1, -1\});
    fq.pb((p[j]-p[i]).cross(p[k]-p[i]));
    fb.pb({});
    dead.pb (1e9);
    return sz(fv)-1;
  }; // 4652
  rep(i, 0, 2) {
    fe[add(0, i+1, 2-i)] = \{!i, !i, !i\};
    rep(i, 3, n) {
      sc t = fq[i].dot(p[j]-p[0]);
      if (t >= -eps) {
        fb[i].pb(j);
        if (t > eps) bad[j].pb(i);
      } // d64f
   } // a567
  } // e5be
  rep(i, 3, n) {
    int v = -1;
    each(f, bad[i]) dead[f] = min(dead[f], i);
    each(f, bad[i]) if (dead[f] == i) {
      rep(j, 0, 3) if (dead[fe[f][j]] > i) {
        int u = fv[f][(j+1)%3], e = fe[f][j];
        v = fv[f][j];
        fe[g = link[v] = add(v, u, i)][0] = e;
        set_union(all(fb[f]), all(fb[e]),
         back_inserter(fb[q]));
        erase_if(fb[g], [&](int k) {
          return k <= i ||
            fq[g].dot(p[k]-p[fv[g][0]]) \leftarrow eps;
        }); // 3119
```

```
each(k, fb[q]) bad[k].pb(q);
        rep(k, 0, 3) if (fv[e][k] == u) {
          fe[e][k] = q;
          break:
       } // c71e
      } // de51
      vi().swap(fb[f]);
    } // 9d4a
    while (v != -1 && fe[link[v]][1] == -1) {
      int u = fv[link[v]][1];
      fe[link[v]][1] = link[u];
      fe[link[u]][2] = link[v];
      v = u;
    } // 5cf7
    vi().swap(bad[i]);
 } // 343c
  rep(i, 0, sz(fv)) if (dead[i] >= n) {
    each(j, fv[i]) j = ord[j];
    ret.pb(fv[i]);
 } // 3c3b
 return ret;
} // 6324
geo3d/line.h
                                         dd10
#include "vector.h"
// 3D line structure: UNTESTED
struct line3 { // p + d*k == point
 vec3 p, d; // Point and direction
  // Distance from point to line.
  // Depends on vec3: dot, len
 double dist(vec3 a) {
    return d.cross(a-p).len() / d.len();
 } // eb4b
  // Distance between two lines.
  // Depends on vec3: -, dot, cross, len2
 double dist(line3 a) {
    vec3 n = d.cross(a.d);
    sc t = n.len2();
    if (!sqn(t)) return dist(a.p);
    return fabs(n.dot(a.p-p)) / sqrt(t);
 } // 22fa
#if FLOATING POINT GEOMETRY
 // Closest point to another line.
 // Assumes lines are not parallel!
 // Depends on vec3: -, dot, cross, len
 vec3 closest(line3 a) {
    vec3 n2 = a.d.cross(d.cross(a.d));
    return p + d * n2.dot(a.p-p) / d.dot(n2);
 } // fab4
  // Orthogonal projection of point on line.
  // Depends on vec3: -, *, dot, len2
 vec3 proj(vec3 a) {
   return p + d * (d.dot(a-p) / d.len2());
 } // 0187
#endif
}; // c870
// Line through 2 given points.
// Depends on vec: -
line3 through (vec3 a, vec3 b) {
 return {a, b-a};
} // 5b42
                                         09fe
geo3d/plane.h
#include "vector.h"
```

#include "line.h" // For intersections

```
// 3D plane/halfspace structure; UNTESTED
struct plane {
 // For planes: v * point == c
  // For halfspaces: v * point <= c
  // (i.e. normal vector points outside)
 vec3 v; // Normal vector
  sc c = 0; // Offset
  // Distance from point to plane.
  // Depends on vec3: dot, len
  double dist(vec3 a) {
   return fabs(v.dot(a) - c) / v.len();
  // Returns 0 if point a lies on the plane,
  // 1 if on side where normal vector points,
  // -1 if on the other side.
  // Depends on vec3: dot
  int side(vec3 a) { return sqn(v.dot(a)-c); }
#if FLOATING_POINT_GEOMETRY
  // Orthogonal projection of point on plane.
  // Depends on vec3: -, *, dot, len2
  vec3 proj(vec3 a) {
   return a - v * ((v.dot(a)-c) / v.len2());
  // Intersect this plane with line a, returns
  // true on success (false if parallel).
  // Intersection point is saved to 'out'.
  // Depends on vec3: -, *, dot
  bool intersect (line3 a, vec3& out) {
   sc t = v.dot(a.d);
   if (!sqn(t)) return 0;
   out = a.p - a.d * ((v.dot(a.p)-c) / t);
   return 1;
  } // a8f2
  // Intersect this plane with plane a, returns
  // true on success (false if parallel).
  // Depends on vec3: -, *, /, dot, cross, len2
  bool intersect (plane a, line3& out) {
   sc t = (out.d = v.cross(a.v)).len2();
   if (!sgn(t)) return 0;
   out.p = (a.v*c - v*a.c).cross(out.d) / t;
   return 1:
 } // Ofa6
#endif
}; // e26a
// Plane through 3 points with normal vector
// pointing upward when viewed CCW.
// Depends on vec3: -, dot, cross
plane span (vec3 a, vec3 b, vec3 c) {
 vec3 v = (b-a).cross(c-a);
 return {v, v.dot(a)};
} // 4fd9
geo3d/polyhedron_volume.h
                                         f3d3
#include "vector.h"
// Signed volume of a polyhedron; UNTESTED
// Faces orientation needs to be consistent.
// Depends on vec3: cross, dot
double volume (vector < vec3> € p, auto € faces) {
  double v = 0:
 for (auto [a, b, c] : faces)
   v += double(p[a].cross(p[b]).dot(p[c]));
```

```
return v / 6;
} // 423d
```

```
geo3d/sphere.h
                                         fcf1
#include "vector.h"
#include "line.h" // For line intersections.
// 3D sphere structure; UNTESTED
struct sphere {
 vec3 p; // Center
 sc r2 = 0; // Squared radius
 DBP (p, r2);
 // Returns -1 if point q lies outside sphere,
 // 0 if on the edge, 1 if strictly inside.
 // Depends on vec3: -, len2
 int side(vec3 a) {
   return sgn(r2 - (p-a).len2());
#if FLOATING POINT GEOMETRY
 // Intersect with line.
 // Returns number of intersection points.
 // Points are in order given by direction.
 int intersect(line3 a, pair<vec3, vec3>& out){
   sc d = a.dist(p), h2 = r2 - d*d;
   vec3 h, t = a.proj(p);
   int s = sqn(h2)+1;
   if (s > 1) h = a.d * sqrt(h2 / a.d.len2());
   out = \{t-h, t+h\};
   return s;
 } // 685b
#endif
}; // 2a59
geo3d/vector.h
                                         83e4
// Scalar type: float or integer.
#if FLOATING POINT GEOMETRY
 using sc = double;
 constexpr sc eps = 1e-9;
 using sc = 11;
 constexpr sc eps = 0;
#endif
// -1 if a < -eps, 1 if a > eps, 0 otherwise
int sqn(sc a) { return (a>eps) - (a < -eps); }</pre>
// 3D point/vector structure; UNTESTED
struct vec3 {
 using P = vec3;
 sc x = 0, y = 0, z = 0;
 // The following methods are optional
 // and dependencies on them are noted
 // appropriately in library snippets.
 P operator+(P r) const {
   return {x+r.x, y+r.y, z+r.z};
 } // 9aa2
 P operator-(P r) const {
   return {x-r.x, y-r.y, z-r.z};
 P operator*(sc r) const {
   return {x*r, y*r, z*r};
 } // f63f
 P operator/(sc r) const {
   return {x/r, y/r, z/r};
 } // c0d6
 sc dot (P r) const {
   return x*r.x + y*r.y + z*r.z;
 } // af4a
 P cross(P r) const {
   return {y*r.z - z*r.y, z*r.x - x*r.z,
```

```
x*r.y - y*r.x; // aa5e
  } // 28aa
               const { return x*x+y*y+z*z; }
  sc len2()
  double len() const { return hypot(x,y,z); }
  // Equality (with epsilon)
  bool operator==(vec3 r) const {
    return !sqn(x-r.x) && !sqn(y-r.y) &&
           !sqn(z-r.z);
  // Angle between vectors in [0,2*PI]
  double angle(vec3 r) const {
    return atan2(cross(r).len(), dot(r));
  } // 8349
#if FLOATING_POINT_GEOMETRY
 // Rotate counter-clockwise around axis.
 P rotate (double angle, vec3 axis) const {
    auto s = sin(angle), c = cos(angle);
    P u = axis / axis.len();
    return u*dot(u)*(1-c)+(*this)*c-cross(u)*s;
 1 // Ocd9
#endif
}; // 8cce
graphs/2sat.h
                                         4h33
// 2-SAT solver; time: O(n+m), space: O(n+m)
// Variables are indexed from 1 and
// negative indices represent negations!
// Usage: SAT2 sat(variable count);
// (add constraints...)
// bool solution_found = sat.solve();
// sat[i] = value of i-th variable, 0 or 1
            (also indexed from 1!)
// (internally: positive = i*2-1, neg. = i*2-2)
struct SAT2 : vi {
 vector<vi> G;
 vi order, flags;
  // Init n variables, you can add more later
  SAT2(int n = 0) : G(n*2) {}
  // Add new var and return its index
  int addVar() {
    G.resize(sz(G)+2); return sz(G)/2;
  1 // 98f3
  // Add (i => j) constraint
  void implv(int i, int i) {
   i = i*2 ^ i >> 31;
    j = j*2 ^ j >> 31;
    G[--i].pb(--j); G[j^1].pb(i^1);
  } // 8e25
  // Add (i v j) constraint
  void either(int i, int j) { imply(-i, j); }
  // Constraint at most one true variable
  void atMostOne(vi& vars) {
    int y, x = addVar();
    each(i, vars) {
      imply(x, y = addVar());
      imply(i, -x); imply(i, x = y);
   } // 24aa
  } // 3ed7
  // Solve and save assignments in 'values'
  bool solve() { // O(n+m), Kosaraju is used
    assign(sz(G)/2+1, -1);
    flags.assign(sz(G), 0);
    rep(i, 0, sz(G)) dfs(i);
    while (!order.empty()) {
```

```
if (!propag(order.back()^1, 1)) return 0;
      order.pop_back();
   } // 5594
    return 1;
 } // 1e58
  void dfs(int i) {
    if (flags[i]) return;
    flags[i] = 1;
    each(e, G[i]) dfs(e);
    order.pb(i);
 } // d076
 bool propag(int i, bool first) {
   if (!flags[i]) return 1;
    flags[i] = 0;
    if (at(i/2+1) >= 0) return first;
    at (i/2+1) = i&1;
    each(e, G[i]) if (!propag(e, 0)) return 0;
    return 1;
 } // 4c1b
}; // 7be4
graphs/bellman ineq.h
                                         cd51
struct Inea {
 11 a, b, c; // a - b >= c
}; // 663a
// Solve system of inequalities of form a-b>=c
// using Bellman-Ford; time: O(n*m)
bool solveIneg(vector<Ineg>& edges,
              vector<11>& vars) {
 rep(i, 0, sz(vars)) each(e, edges)
    vars[e.b] = min(vars[e.b], vars[e.a]-e.c);
 each (e, edges)
   if (vars[e.a]-e.c < vars[e.b]) return 0;</pre>
 return 1:
} // 241e
graphs/biconnected.h
                                         41fc
// Biconnected components; time: O(n+m)
// Usage: Biconnected bi(graph);
// bi[v] = indices of components containing v
// bi.verts[i] = vertices of i-th component
// bi.edges[i] = edges of i-th component
// Bridges <=> components with 2 vertices
// Articulation points <=> vertices that belong
                          to > 1 component
// Isolated vertex <=> empty component list
struct Biconnected : vector<vi> {
 vector<vi> verts;
 vector<vector<pii>>> edges;
  vector<pii> S:
  Biconnected() {}
  Biconnected(vector<vi>& G) : S(sz(G)) {
    resize(sz(G));
    rep(i, 0, sz(G)) S[i].x ?: dfs(G, i, -1);
   rep(c, 0, sz(verts)) each(v, verts[c])
      at (v) .pb(c);
 } // cfce
 int dfs(vector<vi>& G, int v, int p) {
    int low = S[v].x = sz(S)-1;
    S.pb(\{v, -1\});
    each(e, G[v]) if (e != p) {
      if (S[e].x < S[v].x) S.pb({v, e});</pre>
      low = min(low, S[e].x ?: dfs(G, e, v));
    ) // 446d
```

if (p+1 && low >= S[p].x) {

```
verts.pb({p}); edges.pb({});
      rep(i, S[v].x, sz(S)) {
       if (S[i].y == -1)
          verts.back().pb(S[i].x);
          edges.back().pb(S[i]);
     } // 4fab
     S.resize(S[v].x);
   } // 6d66
    return low;
  } // 7fcc
}; // 3d4a
graphs/bridges online.h
                                         ce5a
// Dynamic 2-edge connectivity queries
// Usage: Bridges bridges(vertex_count);
// - bridges.addEdge(u, v); - add edge (u, v)
// - bridges.cc[v] = connected component ID
// - bridges.bi(v) = 2-edge connected comp ID
struct Bridges {
  vector<vi> G; // Spanning forest
 vi cc, size, par, bp, seen;
  int cnt = 0;
  // Initialize structure for n vertices; O(n)
  Bridges(int n = 0) : G(n), cc(n), size(n, 1),
                       par(n, -1), bp(n, -1),
                       seen(n) {
   iota(all(cc), 0);
  } // ed70
  // Add edge (u, v); time: amortized O(lq n)
  void addEdge(int u, int v) {
   if (cc[u] == cc[v]) {
     int r = lca(u, v);
      for (int x : {u, v})
       while ((x = root(x)) != r)
          x = bp[bi(x)] = par[x];
      G[u].pb(v); G[v].pb(u);
     if (size[cc[u]] > size[cc[v]]) swap(u,v);
     size[cc[v]] += size[cc[u]];
     dfs(u, v);
   } // abc7
  } // a6fd
  // Get 2-edge connected component ID
  int bi(int v) { // amortized time: < O(lq n)</pre>
   return bp[v] + 1? bp[v] = bi(bp[v]) : v;
  1 // 3206
  int root(int v) {
    return par[v] == -1 || bi(par[v]) != bi(v)
     ? v : par[v] = root(par[v]);
  } // 2d27
  void dfs(int v, int p) {
   cc[v] = cc[par[v] = p];
   each(e, G[v]) if (e != p) dfs(e, v);
  int lca(int u, int v) { // Don't use this!
   for (cnt++;; swap(u, v)) if (u != -1) {
     if (seen[u = root(u)] == cnt) return u;
     seen[u] = cnt; u = par[u];
   } // afed
  } // 7f56
}; // 3685
```

```
graphs/chordal graph.h
                                         a894
vi perfectEliminationOrder(vector<vi>€ q) {
 int top = 0, n = sz(g);
 vi ord, vis(n), indeg(n);
 vector<vi> bucket(n);
 rep(i, 0, n) bucket[0].push back(i);
 for (int i = 0; i < n;) {</pre>
   while(bucket[top].empty()) --top;
                                                  } // f045
   int u = bucket[top].back();
                                                 }; // af18
   bucket[top].pop_back();
   if(vis[u]) continue;
   ord.push_back(u);
   vis[u] = 1:
   ++i;
   each(v, g[u]) {
     if (vis[v]) continue;
                                                  int a, b;
     bucket[++indeg[v]].push back(v);
                                                  11 w:
     top = max(top, indeg[v]);
   } // 8043
 } // 858b
                                                  } // 5e2c
 reverse (all (ord));
                                                 }; // 701d
 return ord;
} // 429f
bool isChordal (vector < vi>& g, vi ord) {
 int n = sz(q);
 set<pii> edg;
 rep(i, 0, n) each(v, g[i]) edg.insert({i,v});
 vi pos(n); rep(i, 0, n) pos[ord[i]] = i;
 rep(u, 0, n) {
   int mn = n;
   each (v, g[u]) if (pos[u] < pos[v])
     mn = min(mn, pos[v]);
   if (mn != n) {
     int p = ord[mn];
     each(v, q[u]) if (pos[v] > pos[u]
       && v != p && !edg.count({v, p}))
         return 0;
   } // 755f
 } // b0a2
 return 1:
} // 2dec
                                         5dcd
graphs/dense dfs.h
#include "../math/bit matrix.h"
// DFS over bit-packed adjacency matrix
// G = NxN adjacency matrix of graph
      G(i,i) \iff (i,i) \text{ is edge}
// V = 1xN matrix containing unvisited vertices
      V(0,i) <=> i-th vertex is not visited
// Total DFS time: O(n^2/64)
struct DenseDFS {
 BitMatrix G, V; // space: O(n^2/64)
 // Initialize structure for n vertices
 DenseDFS(int n = 0) : G(n, n), V(1, n) {
   reset():
 } // 79e4
 // Mark all vertices as unvisited
 void reset() { each(x, V.M) x = -1; }
 // Get/set visited flag for i-th vertex
 void setVisited(int i) { V.set(0, i, 0); }
 bool isVisited(int i) { return !V(0, i); }
 // DFS step: func is called on each unvisited
 // neighbour of i. You need to manually call
 // setVisited(child) to mark it visited
 // or this function will call the callback
```

```
// with the same vertex again.
  void step(int i, auto func) {
    ull \star E = G.row(i);
    for (int w = 0; w < G.stride;) {</pre>
     ull x = E[w] & V.row(0)[w];
      if (x) func((w<<6) | __builtin_ctzll(x));</pre>
      else w++;
    } // 4c0a
graphs/directed mst.h
                                          94d4
#include "../structures/find_union_undo.h"
#include <ext/pb_ds/priority_queue.hpp>
struct Edge {
 bool operator<(Edge r) const {</pre>
    return w > r.w;
// Find directed minimum spanning tree
// rooted at vertex `root`; O(m log n)
// Returns weight of found spanning tree.
// par[i] = parent of i-th vertex in the tree,
// par[root] = -1
11 dmst (vector<Edge>& edges,
        int n, int root, vi& par) {
  RollbackFAU dsu(n);
  vector<__gnu_pbds::priority_queue<Edge>>Q(n);
  vector<ll> delta(n);
  each (e, edges) Q[e.b].push (e);
  11 ans = 0:
  vi seen(n, -1), path(n);
  vector<Edge> ed(n), in(n, \{-1, -1, 0\});
  vector<tuple<int, int, vector<Edge>>> cycs;
  seen[root] = root;
  rep(s, 0, n)
    for (int u = s, pos = 0; seen[u] < 0;) {</pre>
      if (0[u].emptv()) return -1;
      auto e = Q[u].top();
      Q[u].pop();
      ans += e.w - delta[u];
      delta[u] = e.w;
      ed[pos] = in[u] = e;
      seen[path[pos++] = u] = s;
      if (seen[u = dsu.find(e.a)] == s) {
        int w, end = pos, t = dsu.time();
        while (dsu.join(u, w = path[--pos])) {
          if (sz(Q[w]) > sz(Q[u])) swap(u, w);
          for (auto f : O[w]) {
            f.w += delta[u] - delta[w];
            Q[u].push(f);
          } // e37e
        } // f27c
        Q[w = dsu.find(u)].swap(Q[u]);
        delta[w] = delta[u];
        seen[u=w] = -1;
        cycs.pb({u, t, {&ed[pos], &ed[end]}});
      1 // cc93
    ) // f264
  reverse (all (cycs));
  for (auto &[u, t, e] : cycs) {
    auto s = in[u];
    dsu.rollback(t);
    each(f, e) in[dsu.find(f.b)] = f;
```

```
in[dsu.find(s.b)] = s;
  } // fla6
  par.resize(n);
  rep(i, 0, n) par[i] = in[i].a;
  return ans;
} // 428e
graphs/dominators.h
                                          8db3
// Tarian's algorithm for finding dominators
// in directed graph; time: O(m log n)
// Returns array of immediate dominators idom.
// idom[root] = root
// idom[v] = -1 if v is unreachable from root
vi dominators (const vector < vi>& G, int root) {
  int n = sz(G);
  vector<vi> in(n), bucket(n);
  vi pre(n, -1), anc(n, -1), par(n), best(n);
  vi ord, idom(n, -1), sdom(n, n), rdom(n);
  auto dfs = [&] (auto f, int v, int p) -> void {
    if (pre[v] == -1) {
      par[v] = p;
      pre[v] = sz(ord);
      ord.pb(v);
      each(e, G[v]) in[e].pb(v), f(f, e, v);
    } // 9c70
  1: // 495a
  auto find = [&](auto f, int v)->pii {
    if (anc[v] == -1) return {best[v], v};
    int b; tie(b, anc[v]) = f(f, anc[v]);
    if (sdom[b] < sdom[best[v]]) best[v] = b;</pre>
    return {best[v], anc[v]};
  1: // cf13
  rdom[root] = idom[root] = root;
  iota(all(best), 0);
  dfs(dfs, root, -1);
  rep(i, 0, sz(ord)) {
    int v = ord[sz(ord)-i-1], b = pre[v];
    each(e, in[v])
      b = min(b, pre[e] < pre[v] ? pre[e] :</pre>
                 sdom[find(find, e).x]);
    each (u, bucket[v]) rdom[u]=find(find,u).x;
    sdom[v] = b;
    anc[v] = par[v];
    bucket[ord[sdom[v]]].pb(v);
  1 // 3663
  each(v, ord) idom[v] = (rdom[v] == v ?
    ord[sdom[v]] : idom[rdom[v]]);
  return idom:
} // b856
graphs/edge_color_bipart.h
                                          dbb2
// Bipartite edge coloring; time: O(nm)
// 'edges' is list of (left vert, right vert),
// where vertices on both sides are indexed
// from 0 to n-1. Returns number of used colors
// (which is equal to max degree).
// col[i] = color of i-th edge [0..max deg-1]
int colorEdges(vector<pii>& edges,
               int n, vi& col) {
  int m = sz (edges), c[2] = {}, ans = 0;
  vi deg[2];
  vector<vector<pii>>> has[2];
  col.assign(m, 0);
  rep(i, 0, 2) {
    deg[i].resize(n+1);
    has[i].resize(n+1, vector<pii>(n+1));
```

```
1 // 693b
  auto dfs = [&] (auto f, int x, int p)->void {
    pii i = has[p][x][c[!p]];
    if (has[!p][i.x][c[p]].y) f(f, i.x, !p);
    else has[!p][i.x][c[!p]] = {};
    has[p][x][c[p]] = i;
    has[!p][i.x][c[p]] = {x, i.y};
    if (i.y) col[i.y-1] = c[p]-1;
  }; // 08b0
  rep(i, 0, m) {
   int x[2] = \{edges[i].x+1, edges[i].y+1\};
    rep(d, 0, 2) {
      deg[d][x[d]]++;
      ans = max(ans, deg[d][x[d]]);
      for (c[d] = 1; has[d][x[d]][c[d]].y;)
        c[d]++;
    } // 9454
    if (c[0]-c[1]) dfs(dfs, x[1], 1);
    rep(d, 0, 2)
     has[d][x[d]][c[0]] = \{x[!d], i+1\};
    col[i] = c[0]-1;
  } // 46c6
  return ans;
} // 5ab4
```

```
graphs/edge color vizing.h
                                        a53f
// General graph edge coloring; time: O(nm)
// Finds (D+1)-edge-coloring of given graph,
// where D is max vertex degree.
// Returns vector of edge colors 'col'.
// col[i] = color of i-th edge [0..D]
vi vizing(vector<pii>€ edges, int n) {
  vi cc(n+1), ret(sz(edges)),
    fan(n), fre(n), loc;
  each(e, edges) cc[e.x]++, cc[e.y]++;
  int u, v, cnt = *max element(all(cc)) + 1;
  vector<vi> adj(n, vi(cnt, -1));
  each (e, edges) {
   tie(u, v) = e;
    fan[0] = v;
   loc.assign(cnt, 0);
   int at = u, end = u, d, c = fre[u],
        ind = 0, i = 0;
    while (d = fre[v],
           !loc[d] && (v = adj[u][d]) != -1)
     loc[d] = ++ind, cc[ind] = d, fan[ind] = v;
    cc[loc[d]] = c;
    for (int cd = d; at+1; cd ^= c ^ d,
        at = adj[at][cd])
      swap(adj[at][cd], adj[end=at][cd^c^d]);
    while (adj[fan[i]][d] + 1) {
     int x = fan[i], y = fan[++i], f = cc[i];
     adj[u][f] = x; adj[x][f] = u;
     adj[y][f] = -1; fre[y] = f;
    } // 0024
    adj[u][d] = fan[i];
    adi[fan[i]][d] = u;
    for (int y : {fan[0], u, end})
      for (int& z = fre[y] = 0; adj[y][z]+1;)
       z++;
  } // 4240
 rep(i, 0, sz(edges))
   for (tie(u, v) = edges[i];
        adj[u][ret[i]] != v;) ++ret[i];
 return ret;
} // 028a
```

```
graphs/flow edmonds karp.h
                                          4cbc
using flow_t = int;
constexpr flow_t INF = 1e9+10;
// Edmonds-Karp algorithm for finding
// maximum flow in graph; time: O(V*E^2)
struct MaxFlow {
 struct Edge {
   int dst, inv;
   flow t flow, cap;
 }; // a53c
 vector<vector<Edge>> G;
 vector<flow t> add:
 vi prev;
 // Initialize for n vertices
 MaxFlow(int n = 0) : G(n) {}
 // Add new vertex
 int addVert() { G.pb({}); return sz(G)-1; }
 // Add edge from u to v with capacity cap
 // and reverse capacity rcap.
 // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
              flow_t cap, flow_t rcap = 0) {
   G[u].pb({ v, sz(G[v]), 0, cap });
   G[v].pb({u, sz(G[u])-1, 0, rcap });
   return sz(G[u])-1;
 } // c96a
  // Compute maximum flow from src to dst.
 flow_t maxFlow(int src, int dst) {
   flow t i, m, f = 0;
   each (v, G) each (e, v) e.flow = 0;
   queue<int> 0:
   O.push (src);
   prev.assign(sz(G), -1);
   add.assign(sz(G), -1);
   add[src] = INF;
   while (!O.empty()) {
     m = add[i = Q.front()];
     Q.pop();
     if (i == dst) {
       while (i != src) {
         auto& e = G[i][prev[i]];
         e.flow -= m;
         G[i = e.dst][e.inv].flow += m;
       } // 1f86
       f += m;
       goto nxt;
     } // 43a2
     each (e, G[i])
       if (add[e.dst] < 0 && e.flow < e.cap) {</pre>
         Q.push (e.dst);
         prev[e.dst] = e.inv;
         add[e.dst] = min(m, e.cap-e.flow);
       } // 4cdb
   } // 887e
   return f;
 } // cec0
 // Get flow through e-th edge of vertex v
 flow_t getFlow(int v, int e) {
   return G[v][e].flow;
 } // Ofaf
```

```
// Get if v belongs to cut component with src
 bool cutSide(int v) { return add[v] >= 0; }
}; // d858
graphs/flow min cost.h
                                         7182
using flow t = 11:
constexpr flow_t INF = 1e18;
// Min cost max flow using cheapest paths;
// time: O(nm + |f|*(m log n))
// or O(|f|*(m log n)) if costs are nonnegative
struct MCMF {
 struct Edge {
   int dst, inv;
   flow_t flow, cap, cost;
 }; // 20f7
 vector<vector<Edge>> G;
 vector<flow_t> add;
  // Initialize for n vertices
 MCMF(int n = 0) : G(n) {}
  // Add new vertex
 int addVert() { G.pb({}); return sz(G)-1; }
 // Add edge from u to v.
  // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
              flow_t cap, flow_t cost) {
   G[u].pb({ v, sz(G[v]), 0, cap, cost });
   G[v].pb({u, sz(G[u])-1, 0, 0, -cost });
   return sz(G[u])-1;
 } // 1095
  // Compute minimum cost maximum flow
 // from src to dst. `f` is set to flow value,
 // 'c' is set to total cost value.
 // Returns false iff negative cycle
  // is reachable from from source.
 bool maxFlow(int src, int dst,
              flow_t& f, flow_t& c) {
   flow t m;
   f = c = 0;
   each (v, G) each (e, v) e.flow = 0;
  #if FLOW NONNEGATIVE COSTS
   vector<flow_t> pot(sz(G));
   // Bellman-Ford O(n*m)
   vector<flow_t> pot(sz(G), INF);
   pot[src] = 0;
   int it = sz(G), ch = 1;
   while (ch-- && it--)
     rep(s, 0, sz(G)) if (pot[s] != INF)
        each(e, G[s]) if (e.cap)
          if ((m = pot[s]+e.cost) < pot[e.dst])</pre>
           pot[e.dst] = m, ch = 1;
   if (it < 0) return 0;</pre>
  #endif
   vi prev(sz(G), -1);
   vector<flow_t> dist(sz(G), INF);
   priority_queue<pair<flow_t, int>> Q;
   add.assign(sz(G), -1);
   Q.push({0, src});
   add[src] = INF;
   dist[src] = 0;
   while (!Q.empty()) {
     auto [d, i] = Q.top();
     Q.pop();
```

```
if (d != -dist[i]) continue;
     m = add[i];
     if (i == dst) {
       f += m;
       c += m * (dist[i]-pot[src]+pot[i]);
       while (i != src) {
         auto& e = G[i][prev[i]];
         e.flow -= m;
         G[i = e.dst][e.inv].flow += m;
       } // 1f86
       rep(j, 0, sz(G))
         pot[j] = min(pot[j]+dist[j], INF);
     } // 36d4
      each(e, G[i]) if (e.flow < e.cap) {
       d = dist[i]+e.cost+pot[i]-pot[e.dst];
       if (d < dist[e.dst]) {
         0.push({-d, e.dst});
         prev[e.dst] = e.inv;
         add[e.dst] = min(m, e.cap-e.flow);
         dist[e.dst] = d;
       } // 5ee6
     } // b6b2
   } // 9eba
   return 1;
 } // a34d
 // Get flow through e-th edge of vertex v
 flow t getFlow(int v, int e) {
   return G[v][e].flow;
 } // Ofaf
 // Get if v belongs to cut component with src
 bool cutSide(int v) { return add[v] >= 0; }
graphs/flow_push_relabel.h
                                         5c4b
using flow t = int;
// Push-relabel algorithm for maximum flow;
// O(V^2*sqrt(E)), but very fast in practice.
struct MaxFlow {
 struct Edge {
   int to, inv;
   flow t rem, cap;
 }; // bc77
 vector<basic_string<Edge>> G;
 vector<flow t> extra;
 vi hei, arc, prv, nxt, act, bot;
 queue<int> 0:
 int n, high, cut, work;
 // Initialize for k vertices
 MaxFlow(int k = 0) : G(k) {}
  // Add new vertex
 int addVert() { G.pb({}); return sz(G)-1; }
 // Add edge from u to v with capacity cap
 // and reverse capacity rcap.
 // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
             flow_t cap, flow_t rcap = 0) {
   G[u].pb({ v, sz(G[v]), 0, cap });
   G[v].pb({u, sz(G[u])-1, 0, rcap });
   return sz(G[u])-1;
 } // c96a
 void raise(int v, int h) {
   prv[nxt[prv[v]] = nxt[v]] = prv[v];
```

```
hei[v] = h;
 if (extra[v] > 0) {
   bot[v] = act[h]; act[h] = v;
   high = max(high, h);
 } // d7ee
 if (h < n) cut = max(cut, h+1);
 nxt[v] = nxt[prv[v] = h += n];
 prv[nxt[nxt[h] = v]] = v;
} // 5274
void global(int s, int t) {
 hei.assign(n, n*2);
 act.assign(n*2, -1);
 iota(all(prv), 0);
 iota(all(nxt), 0);
 hei[t] = high = cut = work = 0;
 hei[s] = n;
 for (int x : {t, s})
   for (Q.push(x); !Q.empty(); Q.pop()) {
     int v = Q.front();
     each(e, G[v])
       if (hei[e.to] == n*2 &&
           G[e.to][e.inv].rem)
         Q.push(e.to), raise(e.to,hei[v]+1);
   } // 1901
} // 3181
void push(int v, Edge& e, bool z) {
 auto f = min(extra[v], e.rem);
 if (f > 0) {
   if (z && !extra[e.to]) {
     bot[e.to] = act[hei[e.to]];
     act[hei[e.to]] = e.to;
   } // 9d90
   e.rem -= f; G[e.to][e.inv].rem += f;
   extra[v] -= f; extra[e.to] += f;
 } // Offb
} // da44
void discharge(int v) {
 int h = n*2, k = hei[v];
 rep(j, 0, sz(G[v])) {
   auto& e = G[v][arc[v]];
   if (e.rem) {
     if (k == hei[e.to]+1) {
       push (v, e, 1);
       if (extra[v] <= 0) return;</pre>
     } else h = min(h, hei[e.to]+1);
   if (++arc[v] >= sz(G[v])) arc[v] = 0;
 1 // 9741
 if (k < n \& \& nxt[k+n] == prv[k+n]) {
   rep(j, k, cut) while (nxt[j+n] < n)
     raise(nxt[j+n], n);
   cut = k;
 } else raise(v, h), work++;
} // b64f
// Compute maximum flow from src to dst
flow t maxFlow(int src, int dst) {
 extra.assign(n = sz(G), 0);
 arc.assign(n, 0);
 prv.resize(n*3);
 nxt.resize(n*3);
 bot.resize(n);
 each(v, G) each(e, v) e.rem = e.cap;
 each (e, G[src])
   extra[src] = e.cap, push(src, e, 0);
 global(src, dst);
```

```
for (; high; high--)
      while (act[high] != -1) {
        int v = act[high];
        act[high] = bot[v];
        if (v != src && hei[v] == high) {
         discharge (v);
          if (work > 4*n) global(src, dst);
       } // 7dcc
     } // 26d4
   return extra[dst];
 } // aa5e
  // Get flow through e-th edge of vertex v
 flow t getFlow(int v, int e) {
   return G[v][e].cap - G[v][e].rem;
 } // 812c
  // Get if v belongs to cut component with src
 bool cutSide(int v) { return hei[v] >= n; }
}; // b6f4
graphs/flow with demands.h
                                         e1c0
#include "flow_edmonds_karp.h"
//#include "flow_push_relabel.h" // if you need
// Flow with demands; time: O(maxflow)
struct FlowDemands {
 MaxFlow net;
 vector<vector<flow t>> demands;
 flow_t total = 0;
  // Initialize for k vertices
 FlowDemands(int k = 0): net(2) {
   while (k--) addVert();
  } // 7bdf
  // Add new vertex
  int addVert() {
   int v = net.addVert();
   demands.pb({});
   net.addEdge(0, v, 0);
   net.addEdge(v, 1, 0);
   return v-2;
  } // 48b6
  // Add edge from u to v with demand dem
  // and capacity cap (dem <= flow <= cap).
  // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
             flow t dem, flow t cap) {
    demands[u].pb(dem);
    demands[v1.pb(0);
    total += dem;
    net.G[0][v].cap += dem;
    net.G[u+2][1].cap += dem;
   return net.addEdge(u+2, v+2, cap-dem) - 2;
  // Check if there exists a flow with value f
  // for source src and destination dst.
  // For circulation, you can set args to 0.
 bool canFlow(int src, int dst, flow_t f) {
   net.addEdge(dst += 2, src += 2, f);
    f = net.maxFlow(0, 1);
   net.G[src].pop_back();
   net.G[dst].pop_back();
   return f == total;
 } // 6285
  // Get flow through e-th edge of vertex v
 flow_t getFlow(int v, int e) {
```

```
return net.getFlow(v+2,e+2)+demands[v][e];
 } // 6cf6
}; // f735
graphs/global min cut.h
                                         c9e3
// Find a minimum cut in an undirected graph
// with non-negative edge weights
// given its adjacency matrix M; time: O(n^3)
// 'out' contains vertices on one side.
ll minCut (vector<vector<ll>>> M, vi& out) {
 int n = sz(M);
 11 ans = INT64_MAX;
  vector<vi> co(n);
  rep(i, 0, n) co[i].pb(i);
  out.clear():
  rep(ph, 1, n) {
   auto w = M[0];
    size_t s = 0, t = 0;
    // O(V^2) -> O(E log V) with priority queue
    rep(it, 0, n-ph) {
      w[t] = INT64_MIN; s = t;
      t = max_element(all(w)) - w.begin();
      rep(i, 0, n) w[i] += M[t][i];
    } // 0831
    ll alt = w[t] - M[t][t];
    if (alt < ans) ans = alt, out = co[t];</pre>
    co[s].insert(co[s].end(), all(co[t]));
    rep(i, 0, n) M[s][i] += M[t][i];
    rep(i, 0, n) M[i][s] = M[s][i];
   M[0][t] = INT64 MIN;
  1 // df69
 return ans;
} // 6664
graphs/gomory hu.h
                                         a520
#include "flow_edmonds_karp.h"
//#include "flow push relabel.h" // if you need
struct Edge {
 int a, b; // vertices
 flow t w; // weight
}; // c331
// Build Gomory-Hu tree; time: O(n*maxflow)
// Gomory-Hu tree encodes minimum cuts between
// all pairs of vertices: mincut for u and v
// is equal to minimum on path from u and v
// in Gomory-Hu tree. n is vertex count.
// Returns vector of Gomory-Hu tree edges.
vector<Edge> gomoryHu (vector<Edge> edges,
                      int n) {
  MaxFlow flow(n);
  each(e, edges) flow.addEdge(e.a,e.b,e.w,e.w);
  vector<Edge> ret(n-1);
  rep(i, 1, n) ret[i-1] = {i, 0, 0};
  rep(i, 1, n) {
    ret[i-1].w = flow.maxFlow(i, ret[i-1].b);
    rep(j, i+1, n)
      if (ret[j-1].b == ret[i-1].b &&
         flow.cutSide(j)) ret[j-1].b = i;
 } // 5ae4
 return ret;
} // afdb
graphs/kth_shortest.h
                                         b346
constexpr ll INF = 1e18;
```

// Eppstein's k-th shortest path algorithm;

```
// time and space: O((m+k) log (m+k))
struct Eppstein {
 using T = 11; // Type for edge weights
 using Edge = pair<int, T>;
 struct Node {
   int E[2] = \{\}, s = 0;
   Edge x;
 }; // fc26
 T shortest; // Shortest path length
 priority_queue<pair<T, int>> Q;
 vector<Node> P{1};
 vi h;
 // Initialize shortest path structure for
 // weighted graph G, source s and target t;
 // time: O(m log m)
 Eppstein(vector<vector<Edge>>& G,
          int s, int t) {
   int n = sz(G);
   vector<vector<Edge>> H(n);
   rep(i,0,n) each(e,G[i]) H[e.x].pb({i,e.y});
   vi ord, par(n, -1);
   vector<T> d(n, -INF);
   Q.push(\{d[t] = 0, t\});
   while (!Q.empty()) {
     auto v = 0.top();
     () qoq.0
     if (d[v,v] == v,x) {
       ord.pb(v.v);
       each(e, H[v.y]) if (v.x-e.y > d[e.x]) {
         Q.push(\{d[e.x] = v.x-e.y, e.x\});
         par[e.x] = v.v;
       1 // 5895
     } // 1b62
   } // 1a6d
   if ((shortest = -d[s]) >= INF) return;
   h.resize(n):
   each (v, ord) {
     int p = par[v];
     if (p+1) h[v] = h[p];
     each(e, G[v]) if (d[e.x] > -INF) {
       T k = e.y - d[e.x] + d[v];
       if (k || e.x != p)
         h[v] = push(h[v], \{e.x, k\});
       else
         p = -1;
     } // 5e05
   1 // 31b9
   P[0].x.x = s;
   Q.push({0, 0});
 } // f546
 int push(int t, Edge x) {
   P.pb(P[t]);
   if (!P[t = sz(P)-1].s || P[t].x.y >= x.y)
     swap(x, P[t].x);
   if (P[t].s) {
     int i = P[t].E[0], j = P[t].E[1];
     int d = P[i].s > P[j].s;
     int k = push(d ? j : i, x);
     P[t].E[d] = k; // Don't inline k!
   } // 10e1
   P[t].s++;
   return t;
 } // a2dc
```

```
// Get next shortest path length,
  // the first call returns shortest path.
  // Returns -1 if there's no more paths;
  // time: O(log k), where k is total count
  // of nextPath calls.
  11 nextPath() {
    if (Q.empty()) return -1;
    auto v = Q.top();
   Q.pop();
    for (int i : P[v.y].E) if (i)
     Q.push({ v.x-P[i].x.y+P[v.y].x.y, i });
    int t = h[P[v.y].x.x];
    if (t) Q.push({ v.x - P[t].x.y, t });
    return shortest - v.x;
 } // 08af
}; // 9a8d
                                          4650
graphs/matching blossom.h
// Edmond's Blossom algorithm for maximum
// matching in general graphs; time: O(nm)
// Returns matching size (edge count).
// match[v] = vert matched to v or -1
int blossom(vector<vi>& G, vi& match) {
  int n = sz(G), cnt = -1, ans = 0;
  match.assign(n, -1);
  vi lab(n), par(n), orig(n), aux(n, -1), q;
  auto blos = [&](int v, int w, int a) {
    while (orig[v] != a) {
     par[v] = w; w = match[v];
     if (lab[w] == 1) lab[w] = 0, q.pb(w);
     orig[v] = orig[w] = a; v = par[w];
   } // 319e
  }; // ab9e
  rep(i, 0, n) if (match[i] == -1)
   each(e, G[i]) if (match[e] == -1) {
     match[match[e] = i] = e; ans++; break;
  rep(root, 0, n) if (match[root] == -1) {
    fill(all(lab), -1);
    iota(all(orig), 0);
    lab[root] = 0;
    q = \{root\};
    rep(i, 0, sz(q)) {
     int v = q[i];
      each(x, G[v]) if (lab[x] == -1) {
       lab[x] = 1; par[x] = v;
       if (match[x] == -1) {
          for (int y = x; y+1;) {
            int p = par[y], w = match[p];
           match[match[p] = y] = p; y = w;
          } // 30c1
          ans++;
          goto nxt;
        } // 6fd0
        lab[match[x]] = 0; q.pb(match[x]);
      } else if (lab[x] == 0 &&
                 orig[v] != orig[x]) {
        int a = orig[v], b = orig[x];
        for (cnt++;; swap(a, b)) if (a+1) {
         if (aux[a] == cnt) break;
          aux[a] = cnt;
          a = (match[a]+1 ?
           orig[par[match[a]]] : -1);
       } // 2776
       blos(x, v, a); blos(v, x, a);
     } // 45a1
   } // d488
```

```
nxt::
 } // 8d8a
  return ans;
} // f17b
graphs/matching_blossom_w.h 536a
// Edmond's Blossom algorithm for weighted
// maximum matching in general graphs; O(n^3)?
// Weights must be positive (I believe).
struct WeightedBlossom {
  struct edge { int u, v, w; };
  int n, s, nx;
  vector<vector<edge>> g;
  vi lab, match, slack, st, pa, S, vis;
  vector<vi> flo, floFrom;
  queue<int> q;
  // Initialize for k vertices
  WeightedBlossom(int k)
      : n(k), s(n*2+1),
        g(s, vector<edge>(s)),
        lab(s), match(s), slack(s), st(s),
        pa(s), S(s), vis(s), flo(s),
        floFrom(s, vi(n+1)) {
    rep(u, 1, n+1) rep(v, 1, n+1)
      g[u][v] = \{u, v, 0\};
  } // 5e51
  // Add edge between u and v with weight w
  void addEdge(int u, int v, int w) {
    q[u][v].w = q[v][u].w = max(q[u][v].w, w);
  } // d296
  // Compute max weight matching.
  // 'count' is set to matching size,
  // `weight` is set to matching weight.
  // Returns vector `match` such that:
  // match[v] = vert matched to v or -1
  vi solve(int& count, ll& weight) {
    fill(all(match), 0);
    nx = n;
    weight = count = 0;
    rep(u, 0, n+1) flo[st[u] = u].clear();
    int tmp = 0;
    rep(u, 1, n+1) rep(v, 1, n+1) {
      floFrom[u][v] = (u-v ? 0 : v);
      tmp = max(tmp, g[u][v].w);
    } // a881
    rep(u, 1, n+1) lab[u] = tmp;
    while (matching()) count++;
    rep(u, 1, n+1)
      if (match[u] && match[u] < u)</pre>
        weight += g[u][match[u]].w;
    vi ans(n):
    rep(i, 0, n) ans[i] = match[i+1]-1;
    return ans;
  } // 9ca0
  int delta(edge& e) {
    return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
  } // 7b58
  void updateSlack(int u, int x) {
    if (!slack[x] || delta(g[u][x]) <</pre>
      delta(g[slack[x]][x])) slack[x] = u;
  } // 1f7f
  void setSlack(int x) {
    slack[x] = 0;
    rep(u, 1, n+1) if (g[u][x].w > 0 &&
      st[u] != x && !S[st[u]])
```

```
updateSlack(u, x);
} // ee9c
void push(int x) {
  if (x \le n) q.push(x);
  else rep(i, 0, sz(flo[x])) push(flo[x][i]);
void setSt(int x, int b) {
  st[x] = b;
  if (x > n) rep(i, 0, sz(flo[x]))
    setSt(flo[x][i],b);
} // c5c8
int getPr(int b, int xr) {
  int pr = int(find(all(flo[b]), xr) -
    flo[b].begin());
  if (pr % 2) {
    reverse(flo[b].begin()+1, flo[b].end());
    return sz(flo[b]) - pr;
  } else return pr;
} // 399f
void setMatch(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = floFrom[u][e.u], pr = getPr(u,xr);
  rep(i, 0, pr)
    setMatch(flo[u][i], flo[u][i^1]);
  setMatch(xr, v);
  rotate(flo[u].begin(), flo[u].begin()+pr,
    flo[u].end());
} // f19d
void augment(int u, int v) {
  while (1) {
    int xnv = st[match[u]];
    setMatch(u, v);
    if (!xnv) return;
    setMatch(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 } // bd23
} // e61a
int getLca(int u, int v) {
  static int t = 0;
  for (++t; u||v; swap(u, v)) {
    if (!u) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  } // aa78
  return 0;
} // 9f28
void blossom(int u, int lca, int v) {
  int b = n+1;
  while (b <= nx && st[b]) ++b;
  if (b > nx) ++nx;
  lab[b] = S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].pb(lca);
  for (int x=u, y; x != lca; x = st[pa[y]]) {
    flo[b].pb(x);
    flo[b].pb(y = st[match[x]]);
    push (y);
  } // 63e0
  reverse (flo[b].begin()+1, flo[b].end());
  for (int x=v, y; x != lca; x = st[pa[y]]) {
    flo[b].pb(x);
    flo[b].pb(y = st[match[x]]);
    push(y);
```

```
} // 63e0
  setSt(b, b);
  rep(x, 1, nx+1) g[b][x].w = g[x][b].w = 0;
  rep(x, 1, n+1) floFrom[b][x] = 0;
  rep(i, 0, sz(flo[b])) {
    int xs = flo[b][i];
    rep(x, 1, nx+1) if (!g[b][x].w ||
      delta(g[xs][x]) < delta(g[b][x]))
        q[b][x]=q[xs][x], q[x][b]=q[x][xs];
    rep(x, 1, n+1) if (floFrom[xs][x])
      floFrom[b][x] = xs;
  } // 5833
  setSlack(b);
} // 9000
void blossom(int b) {
  each(e, flo[b]) setSt(e, e);
  int xr = floFrom[b][q[b][pa[b]].u];
  int pr = getPr(b, xr);
  for (int i = 0; i < pr; i += 2) {</pre>
    int xs = flo[b][i], xns = flo[b][i+1];
    pa[xs] = q[xns][xs].u;
    S[xs] = 1; S[xns] = slack[xs] = 0;
    setSlack(xns); push(xns);
  } // f26f
  S[xr] = 1; pa[xr] = pa[b];
  rep(i, pr+1, sz(flo[b])) {
    int xs = flo[b][i];
    S[xs] = -1; setSlack(xs);
  } // a12a
  st[b] = 0;
} // f750
bool found (const edge& e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u; S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = S[nu] = 0;
    push (nu);
  } else if (!S[v]) {
    int lca = getLca(u, v);
    if (!lca) return augment(u, v),
      augment (v, u), 1;
    else blossom(u, lca, v);
  } // ddbb
  return 0;
} // 1c00
bool matching() {
  fill(S.begin(), S.begin()+nx+1, -1);
  fill(slack.begin(), slack.begin()+nx+1, 0);
  q = {};
  rep(x, 1, nx+1)
    if (st[x] == x &  imatch[x])
      pa[x] = S[x] = 0, push(x);
  if (q.empty()) return 0;
  while (1) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      rep (v, 1, n+1)
        if (g[u][v].w > 0 && st[u] != st[v]){
          if (!delta(g[u][v])) {
            if (found(g[u][v])) return 1;
          } else updateSlack(u, st[v]);
        } // b782
    } // 4d33
    int d = INT_MAX;
    rep (b, n+1, nx+1)
      if (st[b] == b && S[b] == 1)
```

```
d = \min(d, lab[b]/2);
                                                    } // d2a7
      rep(x, 1, nx+1)
                                                   bool dfs(int v, int q) {
        if (st[x] == x && slack[x]) {
                                                     if (vis[v] < k) vis[v] = k, seen.pb(v);
          if (S[x] == -1)
                                                      while (low[v] < q) {</pre>
            d = min(d, delta(g[slack[x]][x]));
                                                       int e = adj[v][pos[v]];
          else if (!S[x])
                                                        if (at(e) != v && low[v] == rank[e]) {
            d = min(d, delta(g[slack[x]][x])/2);
                                                          rank[e]++;
       } // 2a0e
                                                          if (at (e) == -1 || dfs(at (e), rank[e]))
      rep(u, 1, n+1) {
                                                            return at (e) = v, 1;
        if (!S[st[u]]) {
                                                        } else if (++pos[v] == sz(adj[v])) {
          if (lab[u] <= d) return 0;</pre>
                                                          pos[v] = 0; low[v]++;
          lab[u] -= d;
                                                       } // e532
       } else if (S[st[u]] == 1) lab[u] += d;
                                                     } // 3d88
      } // 4601
                                                     return 0;
      rep(b, n+1, nx+1) if (st[b] == b) {
                                                   } // 8561
        if (!S[st[b]]) lab[b] += d*2;
                                                 }; // 4560
        else if (S[st[b]] == 1) lab[b] -= d*2;
                                                 graphs/matching_turbo.h
     } // e09b
     q = \{\};
                                                  // Find maximum bipartite matching; time: ?
      rep(x, 1, nx+1)
                                                 // G must be bipartite graph!
        if (st[x] == x && slack[x] &&
                                                 // Returns matching size (edge count).
          st[slack[x]] != x &&
                                                  // match[v] = vert matched to v or -1
          !delta(g[slack[x]][x]) &&
                                                  int matching(vector<vi>& G, vi& match) {
          found(q[slack[x]][x])) return 1;
                                                   vector<bool> seen;
      rep(b, n+1, nx+1)
                                                   int n = 0, k = 1;
        if (st[b] == b && S[b] == 1 && !lab[b])
                                                   match.assign(sz(G), -1);
          blossom(b);
    } // a122
                                                    auto dfs = [&](auto f, int i)->int {
    return 0;
                                                     if (seen[i]) return 0;
 } // e966
                                                      seen[i] = 1;
1: // 35de
                                                      each(e, G[i]) {
                                                       if (match[e] < 0 || f(f, match[e])) {</pre>
graphs/matching boski.h
                                          c8ac
                                                          match[i] = e; match[e] = i;
                                                          return 1;
// Bosek's algorithm for partially online
                                                       } // 893d
// bipartite maximum matching - white vertices
                                                      ) // 1c44
// are fixed, black vertices are added
                                                     return 0;
// one by one; time: O(E*sqrt(V))
                                                    1: // c8bd
// Usage: Matching match(num white);
// match[v] = index of black vertex matched to
                                                    while (k) {
              white vertex v or -1 if unmatched
                                                      seen.assign(sz(G), 0);
// match.add(indices of white neighbours);
// Black vertices are indexed in order they
                                                     rep(i, 0, sz(G)) if (match[i] < 0)
// were added, the first black vertex is 0.
                                                       k += dfs(dfs, i);
struct Matching : vi {
                                                     n += k;
  vector<vi> adj;
                                                   } // 62a7
  vi rank, low, pos, vis, seen;
                                                   return n;
  int k = 0:
                                                  } // 616f
  // Initialize structure for n white vertices
                                                  // Convert maximum matching to vertex cover
  Matching (int n = 0): vi(n, -1), rank(n) {}
                                                  // time: O(n+m)
                                                  vi vertexCover(vector<vi>& G, vi& match) {
  // Add new black vertex with its neighbours
                                                   vi ret, col(sz(G)), seen(sz(G));
  // given by 'vec'. Returns true if maximum
  // matching is increased by 1.
                                                    auto dfs = [&](auto f, int i, int c)->void {
  bool add(vi vec) {
                                                     if (col[i]) return;
    adi.pb (move (vec));
                                                     col[i] = c+1:
    low.pb(0); pos.pb(0); vis.pb(0);
                                                     each(e, G[i]) f(f, e, !c);
    if (!adj.back().emptv()) {
                                                   }; // b718
     int i = k;
                                                    auto aug = [&](auto f, int i)->void {
    nxt:
                                                     if (seen[i] || col[i] != 1) return;
      seen.clear();
                                                      seen[i] = 1;
      if (dfs(sz(adj)-1, ++k-i)) return 1;
                                                     each(e, G[i]) seen[e] = 1, f(f, match[e]);
      each (v, seen) each (e, adj[v])
                                                   1: // 3452
       if (rank[e] < 1e9 && vis[at(e)] < k)</pre>
                                                    rep(i, 0, sz(G)) dfs(dfs, i, 0);
          goto nxt:
                                                   rep(i, 0, sz(G)) if (match[i]<0) aug(aug, i);
      each (v, seen) each (w, adj[v])
        rank[w] = low[v] = 1e9;
                                                   rep(i, 0, sz(G))
    } // 6aec
                                                     if (seen[i] == col[i]-1) ret.pb(i);
    return 0;
                                                    return ret;
```

```
} // a4c1
       graphs/matching weighted.h
       // Minimum cost bipartite matching; O(n^2*m)
       // Input is n x m cost matrix, where n <= m.
       // Returns matching weight.
       // L[i] = right vertex matched to i-th left
       // R[i] = left vertex matched to i-th right
       ll hungarian(const vector<vector<ll>>& cost.
                    vi& L, vi& R) {
         if (cost.empty())
           return L.clear(), R.clear(), 0;
         int b, c = 0, n = sz(cost), m = sz(cost[0]);
         assert(n <= m);
         vector<11> x(n), y(m+1);
         L.assign(n, -1);
         R.assign(m+1, -1);
6439
         rep(i, 0, n) {
           vector<ll> sla(m, INT64_MAX);
           vi vis(m+1), prv(m, -1);
           for (R[b = m] = i; R[b]+1; b = c) {
             int a = R[b];
             11 d = INT64 MAX;
             vis[b] = 1;
             rep(j, 0, m) if (!vis[j]) {
               ll cur = cost[a][j] - x[a] - y[j];
               if (cur < sla[j])</pre>
                 sla[j] = cur, prv[j] = b;
               if (sla[j] < d) d = sla[j], c = j;</pre>
             1 // 6717
             rep(j, 0, m+1) {
               if (vis[j]) x[R[j]] \leftarrow d, y[j] \leftarrow d;
               else sla[j] -= d;
             } // 8bb3
           } // 01c6
           while (b-m) c = b, R[c] = R[b = prv[b]];
         } // 50bb
         rep(j, 0, m) if (R[j]+1) L[R[j]] = j;
         R.resize(m);
         return -y[m];
       1 // 0430
       graphs/matroids.h
       // Find largest subset S of [n] such that
       // S is independent in both matroid A and B.
       // A and B are given by their oracles,
       // see example implementations below.
       // Returns vector V such that V[i] = 1 iff
       // i-th element is included in found set;
       // time: O(r^2*init + r^2*n*add),
       // where r is max independent set,
       // 'init' is max time of oracles init
       // and 'add' is max time of oracles canAdd.
       vector<bool> intersectMatroids(
             auto& A, auto& B, int n) {
         vector<bool> ans(n);
         bool ok = 1;
         // NOTE: for weighted matroid intersection
         // find shortest augmenting paths
         // first by weight change, then by length
         // using Bellman-Ford, and skip this speedup:
         A.init(ans):
         B.init(ans);
         rep(i, 0, n) if (A.canAdd(i) && B.canAdd(i))
           ans[i] = 1, A.init(ans), B.init(ans);
         while (ok) {
```

```
vector<vi> G(n);
           vector<bool> good(n);
ed77
           queue<int> que;
           vi prev(n, -1);
           A.init(ans);
           B.init(ans);
           ok = 0;
           rep(i, 0, n) if (!ans[i]) {
             if (A.canAdd(i)) que.push(i), prev[i]=-2;
             good[i] = B.canAdd(i);
           } // 9581
           rep(i, 0, n) if (ans[i]) {
             ans[i] = 0;
             A.init(ans);
             B.init(ans);
             rep(j, 0, n) if (i != j && !ans[j]) {
               if (A.canAdd(j)) G[i].pb(j);
               if (B.canAdd(j)) G[j].pb(i);
             } // bd2a
             ans[i] = 1;
           } // bf3e
           while (!que.emptv()) {
             int i = que.front();
             que.pop();
             if (good[i]) {
               ans[i] = 1;
               while (prev[i] >= 0) {
                 ans[i = prev[i]] = 0;
                 ans[i = prev[i]] = 1;
               ) // 51c8
               ok = 1;
               break:
             } // 384b
             each(j, G[i]) if (prev[j] == -1)
               que.push(j), prev[j] = i;
           } // 6eb6
         } // 3c97
         return ans;
       } // 774e
ca31
       // Matroid where each element has color
        // and set is independent iff for each color c
       // #{elements of color c} <= maxAllowed[c].</pre>
       struct LimOracle {
         vi color; // color[i] = color of i-th element
         vi maxAllowed; // Limits for colors
         // Init oracle for independent set S; O(n)
         void init(vector<bool>& S) {
           tmp = maxAllowed:
           rep(i, 0, sz(S)) tmp[color[i]] -= S[i];
         } // 4dfb
         // Check if S+{k} is independent; time: O(1)
         bool canAdd(int k) {
           return tmp[color[k]] > 0;
         } // e312
       }; // c7d0
        // Graphic matroid - each element is edge,
       // set is independent iff subgraph is acyclic.
       struct GraphOracle {
         vector<pii> elems; // Ground set: graph edges
         int n; // Number of vertices, indexed [0;n-1]
         vi par;
```

```
int find(int i) {
    return par[i] == -1 ? i
     : par[i] = find(par[i]);
  // Init oracle for independent set S; ~O(n)
  void init(vector<bool>& S) {
   par.assign(n, -1);
   rep(i, 0, sz(S)) if (S[i])
     par[find(elems[i].x)] = find(elems[i].y);
  // Check if S+{k} is independent; time: ~O(1)
 bool canAdd(int k) {
   return
      find(elems[k].x) != find(elems[k].y);
1: // 19d3
// Co-graphic matroid - each element is edge,
// set is independent iff after removing edges
// from graph number of connected components
// doesn't change.
struct CographOracle {
  vector<pii> elems; // Ground set: graph edges
  int n; // Number of vertices, indexed [0;n-1]
  vi pre, low;
  int cnt;
  int dfs(int v, int p) {
   pre[v] = low[v] = ++cnt;
   each(e, G[v]) if (e != p)
     low[v] = min(low[v], pre[e] ?: dfs(e,v));
   return low[v];
  // Init oracle for independent set S; O(n)
  void init(vector<bool>& S) {
   G.assign(n, {});
   pre.assign(n, 0);
   low.resize(n);
    cnt = 0;
    rep(i, 0, sz(S)) if (!S[i]) {
     pii e = elems[i];
     G[e.x].pb(e.y);
     G[e.y].pb(e.x);
   } // f4e8
   rep(v, 0, n) if (!pre[v]) dfs(v, -1);
  } // dfe1
  // Check if S+{k} is independent; time: O(1)
 bool canAdd(int k) {
   pii e = elems[k];
   return max(pre[e.x], pre[e.y])
     != max(low[e.x], low[e.y]);
 } // f6c5
}; // 4149
// Matroid equivalent to linear space with XOR
struct XorOracle {
  vector<1l> elems; // Ground set: numbers
  vector<11> base:
  // Init for independent set S; O(n+r^2)
  void init(vector<bool>& S) {
   base.assign(63, 0);
   rep(i, 0, sz(S)) if (S[i]) {
     11 e = elems[i];
     rep(j, 0, sz(base)) if ((e >> j) & 1) {
       if (!base[j]) {
         base[j] = e;
```

```
break:
        } // 1df5
        e ^= base[j];
      } // 8495
    } // 655e
  } // b68c
  // Check if S+{k} is independent; time: O(r)
  bool canAdd(int k) {
   11 e = elems[k];
    rep(i, 0, sz(base)) if ((e >> i) & 1) {
      if (!base[i]) return 1;
      e ^= base[i];
    } // 49d1
    return 0;
 } // 66ff
}; // 4af3
graphs/max_clique.h
// Quickly finds a maximum clique of a graph
// (given as symmetric bitset matrix;
// self-edges not allowed).
// time: ~1s for n=155 and worst case random
// graphs (p=.90). Faster for sparse graphs.
typedef vector<bitset<200>> vb;
struct MaxClique {
  double limit=0.025, pk=0;
  struct Vertex { int i, d=0; };
  typedef vector<Vertex> vv;
  vb e;
  vv V;
  vector<vi> C;
  vi qmax, q, S, old;
  void init(vv& r) {
    for (auto& v : r) v.d = 0;
    for (auto& v : r) for (auto j : r)
     v.d += e[v.i][j.i];
    sort (all (r),
      [](auto a, auto b) {return a.d > b.d;});
    int mxD = r[0].d;
    rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
  void expand(vv& R, int lev = 1) {
    S[lev] += S[lev - 1] - old[lev];
    old[lev] = S[lev - 1];
    while (sz(R)) {
      if (sz(q) + R.back().d \le sz(qmax))
      g.push back (R.back ().i);
      vv T;
      for(auto v:R) if (e[R.back().il[v.il)
        T.push back({v.i});
      if (sz(T)) {
        if (S[lev]++ / ++pk < limit) init(T);</pre>
        int j = 0, mxk = 1;
        int mnk = max(sz(qmax) - sz(q) + 1, 1);
        C[1].clear(), C[2].clear();
        for (auto v : T) {
          int k = 1:
          auto f=[&](int i){return e[v.i][i];};
          while (any_of(all(C[k]), f)) k++;
          if (k > mxk) mxk=k, C[mxk+1].clear();
          if (k < mnk) T[j++].i = v.i;</pre>
          C[k].push_back(v.i);
        } // e825
        if (j > 0) T[j - 1].d = 0;
        rep(k, mnk, mxk + 1) for (int i : C[k])
          T[j].i = i, T[j++].d = k;
        expand(T, lev + 1);
                                                  using Edge = pair<int, 11>;
```

```
} else if (sz(q) > sz(qmax)) qmax = q;
     q.pop_back(), R.pop_back();
   } // dea6
 } // f0ce
 vi solve() {
   init(V), expand(V); return qmax; } // 2243
 MaxClique (vb conn)
     : e(conn), C(sz(e)+1), S(sz(C)), old(S) {
   rep(i, 0, sz(e)) V.push_back({i});
 } // cd99
}; // b944
graphs/max clique chinese.h beca
constexpr int N = 405:
// Max clique heuristic that seems to work well
// with geometric packing problems. Vertices
// should be ordered by (X,Y), not shuffled.
struct MaxClique {
 bool q[N][N];
 int n, dp[N], st[N][N], ans, res[N], stk[N];
 void init(int n_) {
   n = n;
   memset(g, 0, sizeof(g));
 1 // 5413
 void addEdge(int u, int v, int w) {
   g[u][v] = w;
 } // 6fb6
 bool dfs(int sz, int num) {
   if (sz == 0) {
     if (num > ans) {
       ans = num;
        copy(stk+1, stk+1+num, res+1);
       return 1;
     } // 9ad5
     return 0;
   } // c6b6
   for (int i = 0; i < sz; i++) {
     if (sz-i+num <= ans) return 0;</pre>
     int u = st[num][i];
     if (dp[u]+num <= ans) return 0;</pre>
     int cnt = 0;
     rep(j, i+1, sz)
       if (q[u][st[num][j]])
         st[num+1][cnt++] = st[num][j];
     stk[num+1] = u;
     if (dfs(cnt, num + 1)) return 1;
   } // fddd
   return 0;
 } // aae9
 int solve() {
   ans = 0;
   memset(dp, 0, sizeof(dp));
   for (int i = n; i >= 1; i--) {
     int cnt = 0;
     rep(j, i+1, n+1)
      if (q[i][j]) st[1][cnt++] = j;
     stk[1] = i;
     dfs(cnt, 1);
     dp[i] = ans;
   } // 2361
   return ans;
 } // dcc6
}; // 7599
graphs/spfa.h
                                         2179
```

```
// SPFA with subtree erasure heuristic;
// time: pessimistic O(nm), on random O(m)
// Returns array of distances or empty array
// if negative cycle is reachable from source.
// par[v] = parent in shortest path tree
vector<ll> spfa(vector<vector<Edge>>& G,
                vi& par, int src) {
 int n = sz(G);
 vi que, prv(n+1);
 iota(all(prv), 0);
 vi nxt = prv;
 vector<ll> dist(n, INT64_MAX);
 par.assign(n, -1);
  auto add = [&] (int v, int p, 11 d) {
   par[v] = p;
    dist[v] = d;
    prv[n] = nxt[prv[v] = prv[nxt[v] = n]] = v;
 }; // aeb1
  auto del = [&](int v) {
   nxt[prv[nxt[v]] = prv[v]] = nxt[v];
   prv[v] = nxt[v] = v;
 }; // df30
  for (add(src, -2, 0); nxt[n] != n;) {
    int v = nxt[n];
    del(v);
    each(e, G[v]) {
     ll alt = dist[v] + e.v;
      if (alt < dist[e.x]) {</pre>
        que = \{e.x\};
        rep(i, 0, sz(que)) {
          int w = que[i];
         par[w] = -1:
          del(w);
          each(f, G[w])
           if (par[f.x] == w) que.pb(f.x);
        } // c58d
        if (par[v] == -1) return {};
       add(e.x, v, alt);
     } // fd17
   } // 0e38
 } // b029
 return dist;
} // Of19
graphs/strongly connected.h 72ba
// Tarjan's SCC algorithm; time: O(n+m)
// Usage: SCC scc(graph);
// scc[v] = index of SCC for vertex v
// scc.comps[i] = vertices of i-th SCC
// Components are in reversed topological order
struct SCC : vi {
 vector<vi> comps:
 vi S;
 SCC() {}
 SCC (vector\langle vi \rangle \& G) : vi(sz(G), -1), S(sz(G)) {
   rep(i, 0, sz(G)) if (!S[i]) dfs(G, i);
 } // f0fa
 int dfs(vector<vi>& G, int v) {
    int low = S[v] = sz(S);
    S.pb(v);
    each(e, G[v]) if (at(e) < 0)
     low = min(low, S[e] ?: dfs(G, e));
    if (low == S[v]) {
```

```
comps.pb({});
     rep(i, S[v], sz(S)) {
       at (S[i]) = sz(comps)-1;
       comps.back().pb(S[i]);
     } // 8ed0
     S.resize(S[v]);
   } // ecc7
   return low;
 } // f3c6
}; // a7ca
math/berlekamp massev.h
                                         7d12
constexpr int MOD = 998244353;
ll modInv(ll a, ll m) { // a^{(-1)} \mod m
 if (a == 1) return 1;
 return ((a - modInv(m%a, a))*m + 1) / a;
} // c437
// Find shortest linear recurrence that matches
// given starting terms of recurrence; O(n^2)
// Returns vector C such that for each i >= |C|
// A[i] = sum A[i-j-1] *C[j] for j = 0..|C|-1
vector<ll> massey (vector<ll>& A) {
  if (A.emptv()) return {};
  int n = sz(A), len = 0, k = 0;
 11 s = 1:
  vector<ll> B(n), C(n), tmp;
  B[0] = C[0] = 1;
  rep(i, 0, n) {
   11 d = 0;
    rep(j, 0, len+1)
     d = (d + C[j] * A[i-j]) % MOD;
     11 q = d * modInv(s, MOD) % MOD;
     tmp = C:
      rep(j, k, n)
       C[j] = (C[j] - q * B[j-k]) % MOD;
     if (len*2 <= i) {</pre>
       B.swap(tmp);
       len = i-len+1;
       s = d + (d < 0) * MOD;
       k = 0;
     } // c350
   } // 79c7
  } // f70c
  C.resize(len+1);
  C.erase(C.begin());
  each(x, C) x = (MOD - x) % MOD;
  return C;
} // 20ce
math/bit gauss.h
                                         4b1a
constexpr int MAX COLS = 2048;
// Solve system of linear equations over Z 2
// time: O(n^2*m/W), where W is word size
// - A - extended matrix, rows are equations,
```

columns are variables.

// - ans - output for variables values

int bitGauss (vector < bitset < MAX COLS >> & A,

// 2 if more than 1 solution exist.

// - m - variable count

m-th column is equation result

// Returns 0 if no solutions found, 1 if one,

(A[i][j] - i-th row and j-th column)

```
vi col;
 ans.assign(m, 0);
 rep(i, 0, sz(A)) {
   int c = int(A[i]._Find_first());
   if (c >= m) {
     if (c == m) return 0;
     continue;
   } // a6bb
   rep(k, i+1, sz(A)) if (A[k][c]) A[k]^=A[i];
   swap(A[i], A[sz(col)]);
   col.pb(c);
 } // a953
 for (int i = sz(col); i--;) if (A[i][m]) {
   ans[col[i]] = 1;
   rep(k,0,i) if(A[k][col[i]]) A[k][m].flip();
 return sz(col) < m ? 2 : 1;</pre>
} // 986e
math/bit matrix.h
using ull = uint64 t;
// Matrix over Z 2 (bits and xor)
struct BitMatrix {
 vector<ull> M:
 int rows, cols, stride;
 // Create matrix with n rows and m columns
 BitMatrix(int n = 0, int m = 0) {
   rows = n; cols = m;
   stride = (m+63)/64;
   M.resize(n*stride);
 } // 7ef0
 // Get pointer to bit-packed data of i-th row
 ull* row(int i) { return &M[i*stride]; }
  // Get value in i-th row and j-th column
 bool operator()(int i, int j) {
   return (row(i)[j/64] >> (j%64)) & 1;
 } // 28bd
 // Set value in i-th row and j-th column
 void set(int i, int i, bool val) {
   ull &w = row(i)[j/64], m = 1ull << (j%64);
   if (val) w |= m;
   else w &= ~m;
 } // 98a8
}; // 4df7
                                         883b
math/continued fractions.h
// for N \sim 1e7; long double for N \sim 1e9
using dbl = double;
// Given N and a real number x \ge 0, finds the
// closest rational approximation p/g with
// p, q < N. It will obey |p/q - x| < 1/qN.
// For consecutive convergents.
// p \{k+1\}q k - q \{k+1\}p k = (-1)^k.
// (p k/q k alternates between >x and <x.)
// If x is rational, y eventually becomes inf;
// if x is the root of a degree 2 polynomial
// the a's eventually become cyclic; O(lq n)
pair<ll, ll> approximate(dbl x, ll N) {
 ll LP=0, LQ=1, P=1, Q=0, inf = LLONG_MAX;
 for (dbl y = x;;) {
   11 lim = min(P ? (N-LP) / P : inf,
      Q ? (N-LQ) / Q : inf),
```

vector<bool>& ans, int m) {

```
a = (l1) floor(y), b = min(a, lim),
                                                          cmpl operator*=(cmpl& a,cmpl b) {return a=a*b;}
              NP = b*P + LP, NQ = b*Q + LQ;
                                                          // Compute DFT over complex numbers; O(n lq n)
           if (a > b) {
                                                          // Input size must be power of 2!
             // If b > a/2, we have a semi-convergent
                                                         void fft(vector<cmpl>& a) {
             // that gives us a better approximation;
                                                           static vector<cmpl> w(2, 1);
             // if b = a/2, we *may* have one.
                                                           int n = sz(a);
             // Return {P, Q} here for a more
             // canonical approximation.
                                                           for (int k = sz(w); k < n; k *= 2) {
             return (abs(x - (dbl)NP / (dbl)NQ)
                                                             w.resize(n);
               < abs(x - (dbl)P / (dbl)0)) ?
                                                             rep(i,0,k) w[k+i] = \exp(\text{cmpl}(0, M_PI*i/k));
               make_pair(NP, NQ) : make_pair(P, Q);
                                                           } // 92a9
                                                           vi rev(n);
           if (abs(y = 1/(y - (dbl)a)) > 3*N)
                                                           rep(i,0,n) rev[i] = (rev[i/2] | i%2*n) / 2;
             return {NP, NQ};
                                                           rep(i,0,n) if(i<rev[i]) swap(a[i],a[rev[i]]);
           LP = P; P = NP;
                                                           for (int k = 1; k < n; k *= 2) {
           LQ = Q; Q = NQ;
        } // 2fd0
                                                             for (int i=0; i < n; i += k*2) rep(j,0,k) {
       ) // f41e
                                                               auto d = a[i+j+k] * w[j+k];
                                                               a[i+j+k] = a[i+j] - d;
       math/crt.h
                                                 4e5f
                                                               a[i+j] += d;
       using pll = pair<11, 11>;
                                                             } // b389
                                                           } // 84bf
2e3f | 11 egcd(11 a, 11 b, 11& x, 11& v) {
                                                         1 // 9dc8
         if (!a) return x=0, y=1, b;
         11 d = egcd(b%a, a, y, x);
                                                         // Convolve complex-valued a and b,
                                                         // store result in a; time: O(n lq n), 3x FFT
         x = b/a*y;
        return d;
                                                         void convolve(vector<cmpl>& a, vector<cmpl> b) {
                                                           int len = sz(a) + sz(b) - 1;
       1 // 23c8
                                                           if (len <= 0) return a.clear();</pre>
       // Chinese Remainder Theorem; time: O(lg lcm)
                                                           int n = 2 << lq(len);</pre>
       // Solves x = a.x \pmod{a.v}, x = b.x \pmod{b.v}
                                                           a.resize(n); b.resize(n);
       // Returns pair (x mod lcm, lcm(a.v, b.v))
                                                           fft(a); fft(b);
       // or (-1, -1) if there's no solution.
                                                           rep(i, 0, n) a[i] *= b[i] / dbl(n);
       // WARNING: a.x and b.x are assumed to be
                                                           reverse(a.begin()+1, a.end());
       // in [0;a.v) and [0;b.v) respectively.
                                                           fft(a):
       // Works properly if lcm(a.y, b.y) < 2^63.
                                                           a.resize(len);
       pll crt(pll a, pll b) {
                                                         } // 1796
         if (a.y < b.y) swap(a, b);</pre>
                                                          // Convolve real-valued a and b, returns result
         ll x, y, q = \operatorname{egcd}(a.y, b.y, x, y);
         11 c = b.x-a.x, d = b.y/g, p = a.y*d;
                                                         // time: O(n lg n), 2x FFT
         if (c % g) return {-1, -1};
                                                         // Rounding to integers is safe as long as
                                                         // (max coeff<sup>2</sup>) *n*log 2(n) < 9*10<sup>14</sup>
         11 s = (a.x + c/q*x % d * a.y) % p;
                                                          // (in practice 10^16 or higher).
         return {s < 0 ? s+p : s, p};
       } // 35a8
                                                         vector<dbl> convolve (vector<dbl> € a,
                                                                              vector<dbl>& b) {
       math/fast mod.h
                                                 d65b
                                                           int len = \max(sz(a) + sz(b) - 1, 0);
       using ull = uint64 t;
                                                           int n = 2 << lq(len);</pre>
                                                           vector<cmpl> in(n), out(n);
       // Compute a % b faster, where b is constant,
                                                           rep(i, 0, sz(a)) in[i].real(a[i]);
       // but not known at compile time.
       // Returns value in range [0,2b).
                                                           rep(i, 0, sz(b)) in[i].imag(b[i]);
       struct FastMod {
                                                           fft(in);
         ull b. m:
                                                           each(x, in) x \star = x;
         FastMod(ull a) : b(a), m(-1ULL / a) {}
                                                           rep(i,0,n) out[i] = in[-i&(n-1)]-conj(in[i]);
         ull operator()(ull a) { // a % b + (0 or b)
                                                           fft (out);
           return a - ull((__uint128_t (m) *a) >> 64) * b;
                                                           vector<dbl> ret(len);
        } // f27d
                                                           rep(i, 0, len) ret[i] = imag(out[i]) / (n*4);
       }; // 09d4
                                                           return ret;
       math/fft complex.h
                                                 0d46
                                                         } // 41bb
       using dbl = double;
                                                         constexpr 11 MOD = 1e9+7;
       using cmpl = complex<dbl>;
                                                          // High precision convolution of integer-valued
        // Default std::complex multiplication is slow.
                                                         // a and b mod MOD; time: O(n lg n), 4x FFT
       // You can use this to achieve small speedup.
                                                         // Input is expected to be in range [0; MOD)!
       cmpl operator*(cmpl a, cmpl b) {
                                                          // Rounding is safe if MOD*n*log_2(n) < 9*10^14
         dbl ax = real(a), ay = imag(a);
                                                          // (in practice 10^16 or higher).
         dbl bx = real(b), by = imag(b);
                                                         vector<ll> convMod(vector<ll>& a,
         return {ax*bx-ay*by, ax*by+ay*bx};
                                                                            vector<ll>& b) {
       1 // 3b78
                                                           vector<11> ret(sz(a) + sz(b) - 1);
```

```
int n = 2 << __lq(sz(ret));</pre>
  11 cut = 11(sqrt(MOD))+1;
  vector<cmpl> c(n), d(n), g(n), f(n);
  rep(i, 0, sz(a))
   c[i] = {dbl(a[i]/cut), dbl(a[i]%cut)};
  rep(i, 0, sz(b))
   d[i] = {dbl(b[i]/cut), dbl(b[i]%cut)};
  fft(c); fft(d);
  rep(i, 0, n) {
   int j = -i & (n-1);
    f[j] = (c[i] + conj(c[j])) * d[i] / (n*2.0);
      (c[i]-conj(c[j])) * d[i] / cmpl(0, n*2);
  } // e877
  fft(f); fft(q);
  rep(i, 0, sz(ret)) {
   11 t = llround(real(f[i])) % MOD * cut;
   t += llround(imag(f[i]));
   t = (t + llround(real(g[i]))) % MOD * cut;
   t = (t + 1 | f(x))  MOD;
   ret[i] = (t < 0 ? t+MOD : t);
  } // e75d
  return ret;
} // df22
math/fft mod.h
                                          7f8c
// Number Theoretic Tranform (NTT)
// For functions below you can choose 2 params:
// 1. M - prime modulus that MUST BE of form
         a*2^k+1, computation is done in Z_M
// 2. R - generator of Z M
// Modulus often seen on Codeforces:
// M = (119<<23)+1, R = 62; M is 998244353
// Parameters for 11 computation with CRT:
// M = (479<<21)+1, R = 62; M is > 10^9
// M = (483<<21)+1, R = 62; M is > 10^9
11 modPow(11 a, 11 e, 11 m) {
 11 t = 1 \% m;
  while (e) {
    if (e % 2) t = t*a % m;
   e /= 2; a = a*a % m;
  } // 66ca
  return t;
} // 1973
// Compute DFT over Z M with generator R.
// Input size must be power of 2; O(n lq n)
// Input is expected to be in range [0; MOD)!
// dit == true <=> inverse transform * 2^n
                   (without normalization)
template<11 M, 11 R, bool dit>
void ntt(vector<ll>& a) {
  static vector<11> w(2, 1);
  int n = sz(a):
  for (int k = sz(w); k < n; k *= 2) {
   w.resize(n, 1);
   11 c = modPow(R, M/2/k, M);
   if (dit) c = modPow(c, M-2, M);
   rep(i, k+1, k*2) w[i] = w[i-1]*c % M;
  } // 0d98
  for (int t = 1; t < n; t *= 2) {
    int k = (dit ? t : n/t/2);
    for (int i=0; i < n; i += k*2) rep(j,0,k) {</pre>
```

```
ll &c = a[i+j], &d = a[i+j+k];
      ll e = w[j+k], f = d;
      d = (dit ? c - (f=f*e%M) : (c-f)*e % M);
      if (d < 0) d += M;
      if ((c += f) >= M) c -= M;
   } // e4a6
 } // 8d38
} // 01f5
// Convolve a and b mod M (R is generator),
// store result in a; time: O(n lg n), 3x NTT
// Input is expected to be in range [0; MOD)!
template<11 M = (119<<23)+1, 11 R = 62>
void convolve(vector<11>& a, vector<11> b) {
 int len = sz(a) + sz(b) - 1;
 if (len <= 0) return a.clear();</pre>
 int n = 2 << __lg(len);</pre>
 11 t = modPow(n, M-2, M);
 a.resize(n); b.resize(n);
 ntt < M, R, 0 > (a); ntt < M, R, 0 > (b);
 rep(i, 0, n) a[i] = a[i] *b[i] % M * t % M;
 ntt<M,R,1>(a);
 a.resize(len);
} // 24fe
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 11 d = egcd(b%a, a, y, x);
 x = b/a*y;
 return d;
} // 23c8
// Convolve a and b with 64-bit output,
// store result in a; time: O(n lg n), 6x NTT
// Input is expected to be non-negative!
void convLong(vector<11>& a, vector<11> b) {
 const 11 M1 = (479 << 21) +1, M2 = (483 << 21) +1;
 const 11 MX = M1*M2, R = 62;
  auto c = a, d = b;
 each(k, a) k %= M1;
  each(k, b) k %= M1;
 each(k, c) k %= M2;
  each(k, d) k %= M2;
  convolve<M1, R>(a, b);
  convolve<M2, R>(c, d);
 11 x, y; egcd(M1, M2, x, y);
  rep(i, 0, sz(a)) {
   a[i] += (c[i]-a[i]) *x % M2 * M1;
   if ((a[i] %= MX) < 0) a[i] += MX;</pre>
 1 // 2279
} // c493
// Big-integer multiplication note:
// - use convLong with base 10<sup>6</sup> for n < 10<sup>6</sup>
// - use convLong with base 10^5 for n < 10^8
math/fft online.h
#include "modular.h"
#include "fft mod.h"
// Online convolution helper. Ensures that:
// \text{ out } [m] = sum \{ f(i) * q(m-i) : 1 <= i <= m-1 \}
// See usage example below.
void onlineConv(vector<Zp>& out, int m,
                auto f, auto q) {
 int len = m & \sim (m-1), b = m-len;
 int e = min(m+len, sz(out));
  auto apply = [&](auto r, auto s) {
   vector<1l> P(m-b+1), Q(min(e-b, m));
```

```
rep(i, max(b, 1), m) P[i-b] = r(i).x;
   rep(i, 1, sz(Q)) Q[i] = s(i).x;
   convolve (P, Q);
   rep(i, m, e) out[i] += P[i-b];
 }; // d14b
 apply(f, g);
 if (b) apply(g, f);
} // b6d6
// h[m] = 1 + sum h(i)*i * h(m-i)/(m-i)
void example(int n) {
 vector<Zp> h(n);
 for (int m = 1; m < n; m++) {
   onlineConv(h, m,
     [&] (int i) { return h[i] * i; },
     [&] (int i) { return h[i] / i; });
   h[m] += 1;
 } // 11ee
1 // 369c
math/fwht.h
                                         a4d3
// Fast Walsh-Hadamard Transform; O(n lg n)
// Input must be power of 2!
// Uncommented version is for XOR.
// OR version is equivalent to sum-over-subsets
// (Zeta transform, inverse is Moebius).
// AND version is same as sum-over-supersets.
template<bool inv>
void fwht(auto& b) {
 for (int s = 1; s < sz(b); s *= 2) {
   for (int i = 0; i < sz(b); i += s*2) {
     rep(j, i, i+s) {
       auto &x = b[j], &y = b[j+s];
       tie(x, y) = make_pair(x+y, x-y); // XOR
                                        // AND
       x += inv ? -y : y;
       y += inv ? -x : x;
                                         // OR
     } // ceb0
   } // f260
 } // b094
 if (inv) each (e, b) e /= sz(b); // ONLY XOR
} // 45b6
// Compute convolution of a and b such that
// ans[i#i] += a[i]*b[i], where # is OR, AND
// or XOR, depending on FWHT version.
// Stores result in a; time: O(n lg n)
// Both arrays must be of same size = 2^n!
void bitConv(auto& a, auto b) {
 fwht<0>(a);
 fwht<0>(b);
 rep(i, 0, sz(a)) a[i] *= b[i];
 fwht<1>(a);
} // 7b82
math/gauss.h
                                         8469
constexpr double eps = 1e-9;
// Solve system of linear equations; O(n^2*m)
// - A - extended matrix, rows are equations,
        columns are variables.
        m-th column is equation result
         (A[i][i] - i-th row and i-th column)
// - ans - output for variables values
// - m - variable count
// Returns 0 if no solutions found, 1 if one,
// 2 if more than 1 solution exist.
int gauss (vector < vector < double >> & A,
         vector<double>& ans, int m) {
 vi col:
```

ans.assign(m, 0);

```
rep(i, 0, sz(A)) {
    int c = 0;
    while (c <= m && fabs(A[i][c]) < eps) c++;</pre>
    // For Zp:
    //while (c <= m && !A[i][c].x) c++;
    if (c >= m) {
      if (c == m) return 0;
      continue;
    } // a6bb
    rep(k, i+1, sz(A)) {
      auto mult = A[k][c] / A[i][c];
      rep(j, 0, m+1) A[k][j] -= A[i][j]*mult;
    } // 8dd5
    swap(A[i], A[sz(col)]);
    col.pb(c);
  } // 470a
  for (int i = sz(col); i--;) {
    ans [col[i]] = A[i][m] / A[i][col[i]];
    rep(k, 0, i)
      A[k][m] = ans[col[i]] * A[k][col[i]];
  return sz(col) < m ? 2 : 1;</pre>
} // fcf5
math/gauss ortho.h
                                          754f
using Row = vector<double>;
using Matrix = vector<Row>;
constexpr double eps = 1e-9;
// Given a system of n linear equations A
// over m variables, find dimensionality D
// of solution subspace, matrix M and vector t
// such that:
// - matrix M is orthogonal (i.e. M*M^T = I)
// - x is a solution \langle = \rangle (Mx+t)[D...] = 0
// - x[D..] = 0 \iff M^T(x-t) is a solution
// (in particular -M^T*t is a solution)
// Returns number of dimensions D. or -1 if
// there is no solution; time: O(n^2*m + n*m^2)
// Warning: numerical stability is kinda sus
int orthoGauss (Matrix& A, Matrix& M,
               Row& t, int m) {
  int d = m;
  t.assign(m, 0);
  M.assign(m, Row(m));
  rep(i, 0, m) M[i][i] = 1;
  rep(i, 0, sz(A)) {
    auto& w = A[i];
    double s = 0;
    rep(j, 0, d) s += w[j]*w[j];
    if (fabs(s) < eps) {</pre>
      if (fabs(w[m]) > eps) return -1;
      continue;
    } // e92f
    double r = sqrt(s);
    if (w[d-1] < 0) r = -r;
    s = sqrt((s + w[d-1]*r)*2);
    w[d-1] += r;
    rep(j, 0, d) w[j] /= s;
    r = w[m] / (w[d-1] * r * 2);
    rep(j, i+1, sz(A)) {
      s = 0;
      rep(k, 0, d) s += A[j][k] * w[k];
      s *= 2;
```

} // 5b8d

```
rep(k, 0, d) A[j][k] -= s * w[k];
                                                 math/matrix.h
                                                                                          9hf7
                                                                                                       swap(A[i], A[rank]);
                                                                                                                                                      11 b = modPow(a%p, d, p);
                                                                                                       sign *=-1;
                                                                                                                                                      if (b == 1 || b == p-1) continue;
     A[j][m] = s*r;
                                                 #include "modular.h"
   } // 69fe
                                                                                                       break:
                                                                                                                                                      rep(i, 1, t)
                                                                                                     } // f98a
                                                 using Row = vector<Zp>;
                                                                                                                                                        if ((b = modMul(b, b, p)) == p-1) break;
    rep(j, 0, m) {
                                                 using Matrix = vector < Row>;
                                                                                                      if (A[rank][c].x) {
                                                                                                                                                      if (b != p-1) return 0;
     s = 0;
                                                                                                                                                    1 // 9342
                                                                                                       rep(i, rank+1, sz(A)) {
     rep(k, 0, d) s += M[k][j] * w[k];
                                                 // Create n x n identity matrix
                                                                                                          auto mult = A[i][c] / A[rank][c];
                                                 Matrix ident(int n) {
                                                                                                                                                    return 1:
     s *= 2;
                                                                                                          rep(j, 0, sz(A[0]))
     rep(k, 0, d) M[k][j] -= s * w[k];
                                                  Matrix ret(n, Row(n));
                                                                                                                                                  } // bec2
                                                                                                           A[i][j] -= A[rank][j]*mult;
                                                  rep(i, 0, n) ret[i][i] = 1;
   } // 692b
                                                                                                                                                  math/modiny precompute.h
                                                                                                                                                                                            2427
                                                                                                       } // f519
                                                  return ret;
                                                                                                       rank++:
    s = -r;
                                                                                                                                                  constexpr 11 MOD = 234567899;
                                                 } // adld
                                                                                                     } // 4cd8
    rep(k, 0, d) s += t[k] * w[k];
                                                 // Add matrices
    s *= 2;
                                                                                                   } // 36e9
                                                                                                                                                   // Precompute modular inverses; time: O(n)
    rep(k, 0, d) t[k] -= s * w[k];
                                                 Matrix& operator+= (Matrix& 1, const Matrix& r) {
                                                                                                   return rank;
                                                                                                                                                  auto modInv = [] {
   d--:
                                                  rep(i, 0, sz(l)) rep(k, 0, sz(l[0]))
                                                                                                  1 // 6882
                                                                                                                                                    vector<ll> v (MOD, 1); // You can lower size
 } // 789e
                                                    l[i][k] += r[i][k];
                                                                                                                                                    rep(i, 2, sz(v))
                                                                                                  // Compute matrix rank; time: O(n^3)
                                                  return 1:
                                                                                                                                                      v[i] = (MOD - (MOD/i) * v[MOD i]) % MOD;
  return d;
                                                                                                  #define rank rank
                                                                                                                                                    return v;
                                                 } // b6bf
                                                                                                 int rank (Matrix A) {
1 // a6c9
                                                 Matrix operator+(Matrix 1, const Matrix& r) {
                                                                                                                                                  ); // 7806
                                                                                                  int s; return echelon(A, s);
                                                 return 1 += r:
                                         60be
math/linear rec.h
                                                                                                 } // c599
                                                                                                                                                  math/modular.h
                                                                                                                                                                                            72a7
                                                 } // d9b3
constexpr 11 MOD = 998244353;
                                                                                                  // Compute square matrix determinant: O(n^3)
                                                                                                                                                   // Modulus often seen on Codeforces:
                                                 // Subtract matrices
                                                                                                 Zp det (Matrix A) {
                                                                                                                                                  constexpr int MOD = 998244353;
using Poly = vector<11>;
                                                 Matrix& operator = (Matrix& 1, const Matrix& r) {
                                                                                                   int s; echelon(A, s);
                                                                                                                                                   // Some big prime: 15*(1<<27)+1 ~ 2*10^9
                                                  rep(i, 0, sz(l)) rep(k, 0, sz(l[0]))
// Compute k-th term of an n-order linear
                                                                                                   Zp ret = s;
                                                    l[i][k] = r[i][k];
                                                                                                                                                  ll modInv(ll a, ll m) { // a^{(-1)} \mod m
// recurrence C[i] = sum C[i-j-1]*D[j],
                                                                                                   rep(i, 0, sz(A)) ret \star= A[i][i];
                                                  return 1;
                                                                                                                                                    if (a == 1) return 1;
// given C[0..n-1] and D[0..n-1]; O(n^2 \log k)
                                                                                                   return ret;
                                                 } // 90a1
                                                                                                                                                    return ((a - modInv(m%a, a)) *m + 1) / a;
ll linearRec (const Poly € C,
                                                 Matrix operator-(Matrix 1, const Matrix& r) {
                                                                                                                                                  } // c437
            const Poly& D, 11 k) {
                                                  return 1 -= r:
                                                                                                  // Invert square matrix if possible: O(n^3)
  int n = sz(D);
                                                                                                                                                  11 modPow(11 a, 11 e, 11 m) { // a^e mod m
                                                 } // dc4f
                                                                                                  // Returns true if matrix is invertible.
                                                                                                                                                    ll t = 1 % m;
  auto mul = [&] (Poly a, Poly b) {
                                                                                                 bool invert (Matrix& A) {
                                                 // Multiply matrices
                                                                                                                                                    while (e) {
   Poly ret(n*2+1);
                                                                                                   int s, n = sz(A):
                                                 Matrix operator*(const Matrix& 1,
                                                                                                                                                      if (e % 2) t = t*a % m;
   rep(i, 0, n+1) rep(j, 0, n+1)
                                                                                                   rep(i, 0, n) A[i].resize(n*2), A[i][n+i] = 1;
                                                                 const Matrix& r) {
                                                                                                                                                      e /= 2; a = a*a % m;
     ret[i+j] = (ret[i+j] + a[i]*b[j]) % MOD;
                                                                                                   echelon(A, s):
                                                  Matrix ret(sz(1), Row(sz(r[0])));
                                                                                                                                                    } // 66ca
    for (int i = n*2; i > n; i--) rep(j, 0, n)
                                                                                                   for (int i = n; i--;) {
                                                  rep(i, 0, sz(l)) rep(j, 0, sz(r[0]))
                                                                                                                                                    return t;
     ret[i-j-1] =
                                                                                                     if (!A[i][i].x) return 0;
                                                    rep(k, 0, sz(r))
                                                                                                                                                  1 // 1973
        (ret[i-j-1] + ret[i]*D[j]) % MOD;
                                                                                                     auto mult = A[i][i].inv();
                                                      ret[i][j] += l[i][k] * r[k][j];
    ret.resize(n+1);
                                                                                                     each(k, A[i]) k *= mult;
                                                                                                                                                   // Wrapper for modular arithmetic
                                                  return ret:
   return ret;
                                                                                                     rep(k, 0, i) rep(j, 0, n)
                                                                                                                                                  struct Zp {
  }; // e722
                                                                                                       A[k][n+j] -= A[i][n+j] *A[k][i];
                                                                                                                                                    11 x; // Contained value, in range [0; MOD-1]
                                                 Matrix& operator *= (Matrix& 1, const Matrix& r) {
                                                                                                   } // 1e97
                                                                                                                                                    Zp() : x(0) {}
  Poly pol(n+1), e(n+1);
                                                  return 1 = 1*r;
                                                                                                   each(r, A) r.erase(r.begin(), r.begin()+n);
 pol[0] = e[1] = 1;
                                                                                                                                                    Zp(11 a) : x(a\%MOD) { if (x < 0) x += MOD; }
                                                 1 // da8a
                                                                                                   return 1;
                                                                                                                                                     for (k++; k; k /= 2) {
                                                 // Square matrix power; time: O(n^3 * 1g e)
                                                                                                 } // 65b9
                                                                                                                                                        x = x d; return *this; } \
   if (k % 2) pol = mul(pol, e);
                                                 Matrix pow (Matrix a, 11 e) {
                                                                                                 math/miller rabin.h
                                                                                                                                           2d52
                                                                                                                                                       Zp operator c(Zp r) const { \
   e = mul(e, e);
                                                  Matrix t = ident(sz(a));
                                                                                                                                                        Zp t = *this; return t c##= r; } // e986
  } // 13af
                                                  while (e) {
                                                                                                  #include "modular64.h"
                                                    if (e % 2) t *= a;
                                                                                                                                                    OP(+, +r.x - MOD*(x+r.x >= MOD));
 ll ret = 0;
                                                                                                  // Miller-Rabin primality test
                                                    e /= 2; a *= a;
                                                                                                                                                    OP(-, -r.x + MOD*(0 > x-r.x));
  rep(i,0,n) ret = (ret + pol[i+1] \starC[i]) % MOD;
                                                                                                  // time O(k*lg^2 n), where k = number of bases
                                                  1 // 4400
  return ret;
                                                                                                                                                    OP(*, *r.x % MOD);
                                                  return t:
                                                                                                 // Deterministic for p <= 1'050'535'501
                                                                                                                                                    OP(/, *r.inv().x % MOD);
} // 3fd1
                                                 } // 65ea
                                                                                                  // constexpr 11 BASES[] = {
                                                                                                                                                    Zp operator-() const { return Zp()-*this; }
math/linear rec fast.h
                                         58a8
                                                                                                  // 336'781'006'125, 9'639'812'373'923'155
                                                 // Transpose matrix
                                                                                                                                                     // For composite modulus use modInv, not pow
                                                                                                  // }; // d41d
                                                 Matrix transpose (const Matrix & m) {
#include "polvnomial.h"
                                                                                                                                                    Zp inv() const { return pow(MOD-2); }
                                                  Matrix ret(sz(m[0]), Row(sz(m)));
                                                                                                  // Deterministic for p <= 2^64
                                                                                                                                                    Zp pow(ll e) const{ return modPow(x,e,MOD); }
// Compute k-th term of an n-order linear
                                                  rep(i, 0, sz(m)) rep(j, 0, sz(m[0]))
                                                                                                                                                    void print() { cerr << x; } // For deb()</pre>
                                                                                                  constexpr 11 BASES[] = {
// recurrence C[i] = sum C[i-j-1]*D[i],
                                                    ret[j][i] = m[i][j];
                                                                                                                                                   1: // f730
                                                                                                   2, 325, 9'375, 28'178,
// given C[0..n-1] and D[0..n-1];
                                                  return ret;
                                                                                                   450'775. 9'780'504. 1'795'265'022
                                                                                                                                                   // Extended Euclidean Algorithm
// time: O(n log n log k)
                                                 ) // 5650
                                                                                                 }; // 0eld
Zp linearRec(const Poly& C,
                                                                                                                                                  ll egcd(ll a, ll b, ll& x, ll& y) {
                                                 // Transform matrix to echelon form
                                                                                                                                                    if (!a) return x=0, y=1, b;
            const Poly& D, 11 k) {
                                                                                                  bool isPrime(ll p) {
                                                 // and compute its determinant sign and rank.
                                                                                                                                                    11 d = eqcd(b%a, a, y, x);
 Poly f(sz(D)+1, 1);
                                                                                                   if (p <= 2) return p == 2;</pre>
                                                 int echelon(Matrix& A, int& sign) { // O(n^3)
                                                                                                   if (p%2 == 0) return 0;
                                                                                                                                                    x = b/a*v;
  rep(i, 0, sz(D)) f[i] = -D[sz(D)-i-1];
                                                  int rank = 0;
 f = pow({0, 1}, k, f);
                                                                                                                                                    return d;
                                                                                                   11 d = p-1, t = 0;
                                                  sign = 1;
                                                                                                                                                  1 // 23c8
  Zp ret = 0;
                                                                                                   while (d%2 == 0) d /= 2, t++;
                                                  rep(c, 0, sz(A[0])) {
  rep(i, 0, sz(f)) ret += f[i]*C[i];
                                                                                                                                                  math/modular64.h
                                                                                                                                                                                            4b73
                                                    if (rank >= sz(A)) break;
  return ret;
                                                                                                   each(a, BASES) if (a%p) {
```

// 11 a = rand() % (p-1) + 1;

// Modular arithmetic for modulus < 2^62

rep(i, rank+1, sz(A)) if (A[i][c].x) {

```
11 modAdd(11 x, 11 y, 11 m) {
 x += y;
  return x < m ? x : x-m;
} // b653
11 modSub(11 x, 11 y, 11 m) {
  return x >= 0 ? x : x+m;
} // b073
// About 4x slower than normal modulo
11 modMul(11 a, 11 b, 11 m) {
  11 c = 11((long double) a * b / m);
 ll r = (a*b - c*m) % m;
  return r < 0 ? r+m : r;
} // 1815
11 modPow(11 x, 11 e, 11 m) {
  11 t = 1:
  while (e) {
   if (e \& 1) t = modMul(t, x, m);
   e >>= 1;
   x = modMul(x, x, m);
  } // bd61
  return t;
} // c8ba
                                          f203
math/modular generator.h
#include "modular.h" // modPow
// Get unique prime factors of n; O(sqrt n)
vector<ll> factorize(ll n) {
  vector<11> fac;
  for (11 i = 2; i*i <= n; i++) {
   if (n\%i == 0) {
      while (n%i == 0) n /= i;
      fac.pb(i);
   } // 6069
  } // a0cc
  if (n > 1) fac.pb(n);
  return fac;
} // 4a2a
// Find smallest primitive root mod n;
// time: O(sqrt(n) + g*log^2 n)
// Returns -1 if generator doesn't exist.
// For n \le 10^7 smallest generator is \le 115.
// You can use faster factorization algorithm
// to get rid of sgrt(n).
11 generator(11 n) {
  if (n \le 1 \mid | (n > 4 \&\& n \% 4 == 0)) return -1;
  vector<ll> fac = factorize(n);
  if (sz(fac) > (fac[0] == 2)+1) return -1;
  11 phi = n;
  each (p, fac) phi = phi / p \star (p-1);
  fac = factorize(phi);
  for (ll q = 1;; q++) if (qcd(q, n) == 1) {
   each (f, fac) if (modPow(q, phi/f, n) == 1)
     goto nxt;
    return q;
   nxt:;
  } // db24
} // 55e6
                                          ac62
math/modular log.h
#include "modular.h" // modInv
// Baby-step giant-step algorithm; O(sqrt(p))
```

// Finds smallest x such that $a^x = b \pmod{p}$

```
// or returns -1 if there's no solution.
11 dlog(ll a, ll b, ll p) {
 int m = int(min(ll(sqrt(p))+2, p-1));
 unordered_map<11, int> small;
 11 t = 1;
 rep(i, 0, m) {
   int& k = small[t];
   if (!k) k = i+1;
   t = t*a % p;
 } // f1d0
 t = modInv(t, p);
 rep(i, 0, m) {
   int j = small[b];
   if (j) return i*ll(m) + j - 1;
   b = b*t % p;
 } // c7ed
 return -1;
} // 5c26
math/modular sgrt.h
#include "modular.h" // modPow
// Tonelli-Shanks algorithm for modular sqrt
// modulo prime; O(lg^2 p), O(lg p) for most p
// Returns -1 if root doesn't exists or else
// returns square root x (the other one is -x).
11 modSqrt(11 a, 11 p) {
 a %= p;
 if (a < 0) a += p;
 if (a <= 1) return a;
 if (modPow(a, p/2, p) != 1) return -1;
 if (p%4 == 3) return modPow(a, p/4+1, p);
 ll s = p-1, n = 2;
 int r = 0, j;
  while (s%2 == 0) s /= 2, r++;
  while (modPow(n, p/2, p) != p-1) n++;
 ll x = modPow(a, (s+1)/2, p);
 ll b = modPow(a, s, p), g = modPow(n, s, p);
  for (;; r = j) {
   11 t = b;
   for (j = 0; j < r && t != 1; j++)
     t = t*t % p;
    if (!j) return x;
   ll gs = modPow(q, 1LL \ll (r-j-1), p);
   q = qs * qs % p;
   x = x*qs % p;
   b = b * q % p;
 } // f83f
} // 7a97
math/nimber.h
                                          d22e
// Arithmetic over 64-bit nimber field.
// Operations on nimbers are defined as:
// a+b = mex(\{a'+b : a' < a\} u \{a+b' : b' < b\})
// ab = mex({a'b+ab'+a'b' : a' < a, b' < b})
// Nimbers smaller than 2^2^k
// form a field of characteristic 2.
// Addition is equivalent to bitwise xor.
using ull = uint64 t;
uint16_t npw[1<<16], nlg[1<<16];
// Multiply 64-bit nimbers a and b.
template<int half = 32, bool prec = 0>
ull nimMul(ull a, ull b) {
 if (a < 2 || b < 2) return a * b;</pre>
```

```
if (!prec && half <= 8)
   return npw[(nlg[a] + nlg[b]) % 0xFFFF];
 constexpr ull tot = 1ull << half;</pre>
 ull c = a % tot, d = a >> half;
 ull e = b % tot, f = b >> half;
 ull p = nimMul<half/2, prec>(c, e);
 ull r = nimMul<half/2, prec>(d, f);
 ull s = nimMul<half/2, prec>(c^d, e^f);
 ull t = nimMul<half/2, prec>(r, tot/2);
 return p ^ t ^ (p ^ s) << half;</pre>
} // df59
int dummy = ([]() {
 rep(i, npw[0] = 1, 0xFFFF) {
   ull v = nimMul<16, 1>(npw[i-1], -1);
   nlq[npw[i] = uint16 t(v)] = uint16 t(i);
 1 // 43d9
}(), 0);
// Compute a^e under nim arithmetic;
// O(lq M) nimber multiplications
ull nimPow(ull a, ull e) {
 ull t = 1;
 while (e) {
   if (e % 2) t = nimMul(t, a);
   e \neq 2; a = nimMul(a, a);
 } // da53
 return t;
} // c06c
// Compute inverse of a in 2^64 nim-field;
// O(lq M) nimber multiplications
ull nimInv(ull a) {
 return nimPow(a, -2);
} // 4d01
math/phi large.h
                                         8703
#include "pollard rho.h"
// Compute Euler's totient of large numbers
// time: O(n^{(1/4)}) < - factorization
ll phi(ll n) {
 each (p, factorize (n)) n = n / p.x * (p.x-1);
 return n;
} // 798e
math/phi precompute.h
                                          544a
constexpr int MAX_PHI = 1e7;
// Precompute Euler's totients; time: O(n lg n)
vi phi = [] {
 vi p(MAX PHI+1);
 iota(all(p), 0);
 rep(i, 2, sz(p)) if (p[i] == i)
   for (int j = i; j < sz(p); j \leftarrow i)
     p[j] = p[j] / i * (i-1);
 return p;
}(); // d94b
math/phi_prefix_sum.h
                                         89f6
#include "phi precompute.h"
constexpr int MOD = 998244353;
// Precompute Euler's totient prefix sums
// for small values; time: O(n lg n)
// phiSum[k] = sum from 0 to k-1
auto phiSum = [] {
 vector<ll> s(sz(phi)+1);
 rep(i, 0, sz(phi))
   s[i+1] = (s[i] + phi[i]) % MOD;
```

```
return s:
}(); // b078
// Get prefix sum of phi(0) + ... + phi(n-1).
// For MOD > 4*10^9, answer will overflow.
ll getPhiSum(ll n) { // time: O(n^{(2/3)})
  static unordered map<11, 11> big;
  if (n < sz(phiSum)) return phiSum[n];</pre>
  if (big.count(--n)) return big[n];
  ll ret = (n\%2 ? n\%MOD * ((n+1)/2 % MOD)
                : n/2\%MOD * (n\%MOD+1)) \% MOD;
  for (ll s, i = 2; i <= n; i = s+1) {
    s = n / (n/i);
    ret -= (s-i+1) %MOD*getPhiSum(n/i+1) % MOD;
  return big[n] = ret = (ret%MOD + MOD) % MOD;
} // 1d5f
math/pi large.h
                                           c04b
// Precompute prime counting function
// for small values; time: O(n lq lq n)
vector<ll> prl, pis = [] {
  constexpr int MAX_P = 1e7;
  vector<ll> p (MAX_P+1, 1);
  p[0] = p[1] = 0;
  for (int i = 2; i*i <= MAX_P; i++) if (p[i])</pre>
    for (int j = i*i; j <= MAX_P; j += i)</pre>
      p[j] = 0;
  rep(i, 1, sz(p)) {
    if (p[i]) prl.pb(i);
    p[i] += p[i-1];
 } // d28e
 return p;
}(); // f6a5
ll partial(ll x, ll a) {
  static vector<unordered_map<11, 11>> big;
  big.resize(sz(prl));
  if (!a) return (x+1) / 2;
  if (big[a].count(x)) return big[a][x];
  ll ret = partial(x, a-1)
    - partial(x / prl[a], a-1);
  return big[a][x] = ret;
} // 774f
// Count number of primes <= x;
// \text{ time: } O(n^{(2/3)} * log(n)^{(1/3)})
// Set MAX P to be > sart(x) before using!
ll pi(ll x) {
  static unordered_map<11, 11> big;
  if (x < sz(pis)) return pis[x];</pre>
  if (big.count(x)) return big[x];
  while (prl[a]*prl[a]*prl[a] < x) a++;</pre>
  11 ret = 0, b = --a;
  while (++b < sz(prl) && prl[b]*prl[b] < x) {
    11 w = x / prl[b];
    ret -= pi(w);
    for (ll j = b; prl[j]*prl[j] <= w; j++)</pre>
      ret -= pi(w / prl[j]) - j;
  } // a584
  ret += partial(x, a) + (b+a-1)*(b-a)/2;
  return big[x] = ret;
} // eald
```

```
math/pi large precomp.h
                                         e93e
#include "sieve.h"
// Count primes in given interval
// using precomputed table.
// Set MAX P to sart (MAX N)!
// Precomputed table will contain N_BUCKETS
// elements - check source size limit.
// If you need to pack more values,
// you can use 'utils/packing.h'.
// Precomputed table size:
// MAX_N=1e11, N_BUCKETS=1e4 -> 43.96 KB
// MAX N=1e11, N BUCKETS=2e4 -> 85.55 KB
constexpr ll MAX N = 1e11:
constexpr 11 N_BUCKETS = 2e4;
constexpr 11 BUCKET SIZE = (MAX N/N BUCKETS)+1;
constexpr ll precomputed[] = {/* ... */};
// Unpack precomputed data.
// Warning: comment out during precomputing.
vector<ll> buckets = [] {
  vector<11> ret(N BUCKETS+1);
  11 d = 0:
  rep(i, 0, N_BUCKETS)
    ret[i+1] = ret[i] + (d += precomputed[i]);
  return ret:
}(); // db13
// Count primes in range [b;e] naively.
ll sieveRange(ll b, ll e) {
  bitset < BUCKET SIZE > elems:
  b = max(b, 2LL);
  e = max(b, e);
  each (p, primesList) {
   ll c = max((b+p-1) / p, 2LL);
    for (11 i = c*p; i < e; i += p)
      elems.set(i-b);
  1 // 9f7f
 return e-b-elems.count();
} // f028
// Run once on local computer to precompute
// table. Takes about 10 minutes for n = 1e11.
// First and last values for default params:
// 348513, -32447, -9941, -6221, -4585,
// ..., -162, -162, 563, -286, -949
void localPrecompute() {
  11 last = 0:
  for (11 i = 0; i <= MAX N; i += BUCKET SIZE) {</pre>
    11 to = min(i+BUCKET SIZE, MAX N+1);
   11 cur = sieveRange(i, to);
    cout << cur-last << ',' << flush;
   last = cur;
  1 // 93f9
  cout << endl;
} // e009
// Count number of primes <= x;
// time: O(BUCKET_SIZE*lg lg n + sqrt(n)/lg(n))
ll pi(ll x) {
 11 b = x/BUCKET_SIZE, j = b*BUCKET_SIZE;
  return buckets[b] + sieveRange(j, x+1);
} // 8582
                                         1d22
math/pollard rho.h
#include "modular64.h"
#include "miller rabin.h"
using Factor = pair<11, int>;
```

```
void rho(vector<11>& out, 11 n) {
 if (n <= 1) return;</pre>
 if (isPrime(n)) out.pb(n);
  else if (n^2 == 0) rho(out,2), rho(out,n/2);
  else for (11 a = 2;; a++) {
   11 x = 2, y = 2, d = 1;
    while (d == 1) {
     x = modAdd(modMul(x, x, n), a, n);
     y = modAdd(modMul(y, y, n), a, n);
     y = modAdd(modMul(y, y, n), a, n);
     d = gcd(abs(x-y), n);
   } // 20e5
    if (d != n) return rho(out,d),rho(out,n/d);
 } // 423f
} // 0c30
// Pollard's rho factorization algorithm
// Las Vegas version; time: n^(1/4)
// Returns pairs (prime, power), sorted
vector<Factor> factorize(ll n) {
 vector<Factor> ret;
 vector<11> raw:
 rho(raw, n);
  sort(all(raw));
  each(f, raw) {
   if (ret.emptv() | ret.back().x != f)
      ret.pb({ f, 1 });
      ret.back().y++;
 } // 2ab1
 return ret;
} // 471c
                                          6e75
math/polynomial.h
#include "modular.h"
#include "fft mod.h"
using Poly = vector<Zp>;
// Cut off trailing zeroes; time: O(n)
void norm(Poly& P) {
 while (!P.empty() && !P.back().x)
   P.pop back();
} // 8a8a
// Evaluate polynomial at x; time: O(n)
Zp eval(const Poly& P, Zp x) {
 Zp n = 0, v = 1;
  each (a, P) n += a*y, y *= x;
 return n;
} // d865
// Add polynomial; time: O(n)
Polv& operator+= (Polv& 1, const Polv& r) {
 1.resize(max(sz(1), sz(r)));
 rep(i, 0, sz(r)) l[i] += r[i];
 norm(1);
 return 1;
} // 656e
Poly operator+(Poly 1, const Poly& r) {
 return | += r:
} // d9b3
// Subtract polynomial; time: O(n)
Poly& operator == (Poly& 1, const Poly& r) {
 1.resize(max(sz(1), sz(r)));
 rep(i, 0, sz(r)) l[i] = r[i];
 norm(1);
 return 1:
} // c68b
Poly operator-(Poly 1, const Poly& r) {
```

```
return 1 -= r:
} // dc4f
// Multiply by polynomial; time: O(n lg n)
Poly& operator*=(Poly& 1, const Poly& r) {
 if (min(sz(1), sz(r)) < 50) {
    // Naive multiplication
    Poly p(sz(1)+sz(r));
    rep(i, 0, sz(l)) rep(i, 0, sz(r))
     p[i+j] += l[i]*r[j];
    l.swap(p);
 } else {
    // FFT multiplication
    // Choose appropriate convolution method.
    // see fft_mod.h and fft_complex.h
    using v = vector<11>;
   convolve<MOD, 62>(*(v*)&l, *(const v*)&r);
 } // 30c9
 norm(1);
 return 1;
} // e8b3
Poly operator*(Poly 1, const Poly& r) {
 return 1 *= r;
1 // 2de3
// Compute inverse series mod x^n; O(n lg n)
// Requires P(0) != 0.
Poly invert (const Poly € P, int n) {
 assert(!P.empty() && P[0].x);
 Poly tmp{P[0]}, ret = {P[0].inv()};
  for (int i = 1; i < n; i *= 2) {
    rep(j, i, min(i\star2, sz(P))) tmp.pb(P[j]);
    (ret \star= Polv{2} - tmp\starret).resize(i\star2);
 } // 904e
  ret.resize(n);
 return ret;
1 // 9293
// Floor division by polynomial; O(n lg n)
Poly& operator/=(Poly& 1, Poly r) {
 norm(1); norm(r);
 int d = sz(1) - sz(r) + 1;
  if (d <= 0) return l.clear(), l;</pre>
  reverse (all(1));
  reverse (all(r));
 l.resize(d):
 l \star = invert(r, d);
 l.resize(d);
 reverse (all(1));
 return 1:
1 // cf5e
Poly operator/(Poly 1, const Poly& r) {
 return 1 /= r;
} // 152d
// Remainder modulo a polynomial; O(n lg n)
Poly operator% (const Poly € 1, const Poly € r) {
return 1 - r*(1/r);
} // 4fc8
Poly& operator%=(Poly& 1, const Poly& r) {
 return 1 -= r*(1/r);
} // 80bb
// Compute a^e mod x^n, where a is polynomial;
// time: O(n log n log e)
Poly pow(Poly a, ll e, int n) {
 Poly t = \{1\};
 while (e) {
    if (e % 2) (t *= a).resize(n);
    e /= 2; (a \star= a).resize(n);
 } // d0c6
```

```
norm(t);
 return t;
} // ada1
// Compute a^e mod m, where a and m are
// polynomials; time: O(|m| log |m| log e)
Poly pow (Poly a, 11 e, const Poly € m) {
 Poly t = \{1\};
 while (e) {
   if (e % 2) t = t*a % m;
   e /= 2; a = a*a % m;
 } // 66ca
 return t;
} // 6f9c
// Derivate polynomial; time: O(n)
Poly derivate (Poly P) {
 if (!P.empty()) {
   rep(i, 1, sz(P)) P[i-1] = P[i]*i;
    P.pop back();
 1 // bd78
 return P;
} // c6c5
// Integrate polynomial; time: O(n)
Poly integrate (Poly P) {
 if (!P.emptv()) {
    P.pb(0);
    for (int i = sz(P); --i;) P[i] = P[i-1]/i;
   P[0] = 0;
 } // eec1
 return P;
} // e2f3
// Compute ln(P) mod x^n; time: O(n log n)
Poly log(const Poly& P, int n) {
 Poly a = integrate (derivate (P) *invert (P, n));
 a.resize(n);
 return a;
1 // 5d6b
// Compute exp(P) mod x^n; time: O(n lg n)
// Requires P(0) = 0.
Poly exp(Poly P, int n) {
 Poly tmp{P[0]+1}, ret = {1};
 for (int i = 1; i < n; i *= 2) {
    rep(j, i, min(i\star2, sz(P))) tmp.pb(P[j]);
    (ret \star= (tmp - log(ret, i\star2))).resize(i\star2);
 } // c28a
 ret.resize(n);
 return ret:
1 // bd42
// Compute sgrt(P) mod x^n; time: O(n log n)
#include "modular sgrt.h"
bool sqrt(Poly& P, int n) {
 norm(P);
 if (P.empty()) return P.resize(n), 1;
 int tail = 0;
 while (!P[tail].x) tail++;
 if (tail % 2) return 0:
  11 sq = modSqrt(P[tail].x, MOD);
  if (sq == -1) return 0;
  Poly tmp{P[tail]}, ret = {sq};
 for (int i = 1; i < n - tail/2; i *= 2) {
    rep(j, i, min(i\star2, sz(P)-tail))
      tmp.pb(P[tail+j]);
    (ret += tmp * invert(ret, i*2)).resize(i*2);
    each (e, ret) e /= 2;
```

```
1 // 2d41
  P.resize(tail/2);
  P.insert (P.end(), all(ret));
  P.resize(n):
  return 1:
} // b9b3
// Compute polynomial P(x+c); time: O(n lg n)
Poly shift (Poly P, Zp c) {
  int n = sz(P);
  Poly Q(n, 1);
  Zp fac = 1;
  rep(i, 1, n) {
   P[i] \star= (fac \star= i);
   Q[n-i-1] = Q[n-i] * c / i;
  } // 1c20
  P *= 0;
  if (sz(P) < n) return {};</pre>
  P.erase (P.begin(), P.begin()+n-1);
  rep(i, 1, n) P[i] /= (fac *= i);
  return P:
} // b11f
// Compute values P(x^0), ..., P(x^{n-1});
// time: 0(n lg n)
Poly chirpz (Poly P, Zp x, int n) {
  if (P.emptv()) return Polv(n);
  if (!x.x) {
   Polv O(n, P[01);
    rep(i, 1, sz(P)) Q[0] += P[i];
   return Q;
  } // ab77
  int k = sz(P);
  Polv O(n+k):
  rep(i, 0, n+k) Q[i] = x.pow(i*ll(i-1)/2);
  rep(i, 0, k) P[i] /= O[i];
  reverse (all (P));
  P *= 0;
  P.resize(n+k):
  rep(i, 0, n) P[i] = P[k+i-1] / Q[i];
  P.resize(n);
 return P:
} // 5c3c
// Evaluate polynomial P in given points;
// time: O(n 1g^2 n)
Poly eval(const Poly& P, Poly points) {
  int len = 1:
  while (len < sz(points)) len ★= 2;
  vector<Poly> tree(len*2, {1});
  rep(i, 0, sz(points))
   tree[len+i] = \{-points[i], 1\};
  for (int i = len; --i;)
   tree[i] = tree[i\star2] \star tree[i\star2+1];
  tree[0] = P;
  rep(i, 1, len*2)
   tree[i] = tree[i/2] % tree[i];
  rep(i, 0, sz(points)) {
   auto& vec = tree[len+i];
   points[i] = vec.empty() ? 0 : vec[0];
  } // c1c2
  return points;
} // 69b0
// Given n points (x, f(x)) compute n-1-degree
// polynomial f that passes through them;
// time: O(n lg^2 n)
                                                    primes.set();
```

```
// For O(n^2) version see polynomial_interp.h.
Poly interpolate (const vector <pair < Zp, Zp >> & P) {
 int len = 1;
 while (len < sz(P)) len \star= 2;
 vector<Poly> mult(len*2, {1}), tree(len*2);
 rep(i, 0, sz(P))
   mult[len+i] = {-P[i].x, 1};
 for (int i = len; --i;)
   mult[i] = mult[i*2] * mult[i*2+1];
 tree[0] = derivate(mult[1]);
 rep(i, 1, len*2)
   tree[i] = tree[i/2] % mult[i];
 rep(i, 0, sz(P))
   tree[len+i][0] = P[i].y / tree[len+i][0];
 for (int i = len; --i;)
   tree[i] = tree[i\star2]\starmult[i\star2+1]
           + mult[i*2]*tree[i*2+1];
 return tree[1];
1 // b706
math/polynomial interp.h
                                          a4cc
// Interpolate set of points (i, vec[i])
// and return it evaluated at x; O(n lq MOD)
template<class T>
T polyExtend(vector<T>& vec, T x) {
 int n = sz(vec);
 vector<T> fac(n, 1), suf(n, 1);
 rep(i, 1, n) fac[i] = fac[i-1] * i;
 for (int i=n; --i;) suf[i-1] = suf[i]*(x-i);
 T pref = 1, ret = 0;
 rep(i, 0, n) {
   T d = fac[i] * fac[n-i-1] * ((n-i) %2*2-1);
   ret += vec[i] * suf[i] * pref / d;
   pref *= x-i;
 } // 681d
 return ret;
} // dd92
// Given n points (x, f(x)) compute n-1-degree
// polynomial f that passes through them;
// time: O(n^2 lg MOD)
// For O(n lg^2 n) version see polynomial.h
template<class T>
vector<T> polyInterp(vector<pair<T, T>> P) {
 int n = sz(P);
 vector<T> ret(n), tmp(n);
 T last = 0;
 tmp[0] = 1;
 rep(k, 0, n-1) rep(i, k+1, n)
   P[i].y = (P[i].y-P[k].y) / (P[i].x-P[k].x);
 rep(k, 0, n) rep(i, 0, n) {
   ret[i] += P[k].y * tmp[i];
   swap(last, tmp[i]);
   tmp[i] = last * P[k].x;
 } // af1c
 return ret:
} // 7c2c
math/sieve.h
                                          a3cc
constexpr int MAX_P = 1e6;
bitset<MAX_P+1> primes;
// Erathostenes sieve; time: O(n lg lg n)
vi primesList = [] {
```

```
primes.reset(0):
 primes.reset(1);
  for (int i = 2; i*i <= MAX_P; i++)</pre>
   if (primes[i])
      for (int j = i*i; j <= MAX_P; j += i)</pre>
        primes.reset(j);
  vi ret:
  rep(i, 0, MAX_P+1) if (primes[i]) ret.pb(i);
 return ret;
}(); // c997
math/sieve factors.h
                                          3cff
constexpr int MAX P = 1e6:
// Erathostenes sieve that saves smallest
// factor for each number; time: O(n lq lq n)
vi factor = [] {
 vi f(MAX P+1);
 iota(all(f), 0);
 for (int i = 2; i*i <= MAX P; i++)
   if (f[i] == i)
      for (int j = i*i; j <= MAX_P; j += i)</pre>
        f[j] = min(f[j], i);
 return f;
}(); // ac6f
// Factorize n <= MAX_P; time: O(lg n)</pre>
// Returns pairs (prime, power), sorted
vector<pii> factorize(ll n) {
 vector<pii> ret;
 while (n > 1) {
    int f = factor[n];
    if (ret.empty() || ret.back().x != f)
      ret.pb({ f, 1 });
    else
      ret.back().y++;
   n /= f;
 } // 664c
 return ret:
} // 56cb
                                          655c
math/sieve segmented.h
constexpr int MAX P = 1e9:
// Cache-friendly Erathostenes sieve
// \sim 1.5s on Intel Core i5 for MAX_P = 10^9
// Memory usage: MAX P/16 bytes
// The bitset stores only odd numbers.
auto primes = [] {
 constexpr int SEG = 1<<18;</pre>
  int i, sq = int(sqrt(MAX P))+1;
  vector<pii> dels;
 bitset<MAX_P/2+1> ret;
  ret.set():
  ret.reset(0);
  for (int i = 3; i <= sq; i += 2) {
   if (ret[i/21) {
      for (j = i*i; j \le sq; j += i*2)
        ret.reset(j/2);
      dels.pb(\{i, j/2\});
   } // d26d
 } // d26d
  for (int i = sq/2; i \le sz(ret); i += SEG) {
   j = min(i+SEG, sz(ret));
    each(d, dels) for (; d.y < j; d.y += d.x)
     ret.reset(d.y);
  } // 6676
  return ret;
```

```
}(); // 7490
bool isPrime(int n) {
 return n == 2 || (n%2 && primes[n/2]);
} // eb6c
math/simplex.h
                                          ab7a
using dbl = double;
using Row = vector<dbl>;
using Matrix = vector<Row>;
#define mp make_pair
#define ltj(X) if (s == -1 // \
 mp(X[j], N[j]) < mp(X[s], N[s])) s = j
// Simplex algorithm; time: O(nm * pivots)
// Given m x n matrix A, vector b of length m,
// vector c of length n solves the following:
// maximize c^T x, Ax \le b, x \ge 0
// Output vector 'x' contains optimal solution
// or some feasible solution in unbounded case.
// Returns objective value if bounded,
// +inf if unbounded, and -inf if no solution.
// You can test if double is inf using `isinf`.
// PARTIALLY TESTED
dbl simplex (const Matrix& A,
            const Row& b, const Row& c,
            Row& x, dbl eps = 1e-8) {
 int m = sz(b), n = sz(c);
 x.assign(n, 0);
 if (!n) return
   *min_element(all(b)) < -eps ? -1/.0 : 0;
 vi N(n+1), B(m);
 Matrix D(m+2, Row(n+2));
  auto pivot = [&](int r, int s) {
   dbl inv = 1 / D[r][s];
    rep(i, 0, m+2)
     if (i != r && abs(D[i][s]) > eps) {
        dbl tmp = D[i][s] * inv;
        rep(j,0,n+2) D[i][j] -= D[r][j] * tmp;
       D[i][s] = D[r][s] * tmp;
      } // 5281
    each(k, D[r]) k \star= inv;
    each(k, D) k[s] \star = -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
 1: // f56b
  auto solve = [&](int phase) {
    for (int y = m+phase-1;;) {
      int s = -1, r = -1;
      rep(j, 0, n+1)
       if (N[j] != -phase) ltj(D[y]);
      if (D[y][s] >= -eps) return 1;
      rep(i, 0, m)
        if (D[i][s] > eps && (r == -1 ||
          mp(D[i][n+1] / D[i][s], B[i]) <
          mp(D[r][n+1] / D[r][s], B[r]))) r=i;
      if (r == -1) return 0;
     pivot(r, s);
   } // 3bef
 }; // 614a
  rep(i, 0, m) {
    copy(all(A[i]), D[i].begin());
    B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i];
 } // b705
 rep(j, 0, n) D[m][N[j] = j] = -c[j];
 N[n] = -1; D[m+1][n] = 1;
```

```
int r = 0:
  rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
  if (D[r][n+1] < -eps) {</pre>
    pivot(r, n);
    if (!solve(2) || D[m+1][n+1] < -eps)</pre>
      return -1/.0;
    rep(i, 0, m) if (B[i] == -1) {
     int s = 0;
      rep(j, 1, n+1) ltj(D[i]);
     pivot(i, s);
   } // 78fd
  } // b52b
  bool ok = solve(1);
  rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
  return ok ? D[m][n+1] : 1/.0;
math/subset sum.h
                                          aa1b
#include "polvnomial.h"
// Count number of possible subsets that sum
// to t for each t = 1, ..., n; O(n log n)
// Input elements are given by frequency array,
// i.e. counts[x] = how many times elements x
// is contained in the multiset.
// Requires counts[0] == 0.
Poly subsetSum (Poly counts, int n) {
  assert (counts [0].x == 0);
  Poly mul(n);
  rep(i, 0, n)
   \text{mul}[i] = \text{Zp}(i).inv() * (i*2 ? 1 : -1);
  counts.resize(n);
  for (int i = n-2; i > 0; i--)
    for (int j = 2; i * j < n; j++)
      counts[i*j] += mul[j] * counts[i];
  return exp(counts, n);
} // c6ac
math/subset sum mod.h
                                          e358
// Shift-tree with splitmix64 hashing.
struct ShiftTree {
  vector<uint64 t> H;
  int len, delta:
  // Init tree of size n = 2^d.
  ShiftTree(int n) : H(n*2), len(n), delta(0) {
   assert (n && ! (n & (n-1)));
  1 // 5236
  // Set a[i] := 1; time: O(log n)
  void set(int i) {
    H[i = (i+len-delta) % len + len] = 1;
    for (int d = delta; i > 1; d /= 2)
      update(i = parent(i, d%2), d%2);
  } // d5e1
  // Cyclically shift by k to the right;
  // time: O(n / 2^i), where i max s.t. 2^i k
  void shift(int k) {
    if (k %= len) {
      delta = (delta+len+k) % len;
      int div = k \in \sim (k-1), d = delta / div;
      for (int t = len/div/2; t >= 1; t /= 2) {
        rep(i, t, t*2) update(i, d%2);
       d /= 2;
     } // 45ce
    } // b582
  } // 1a6d
  // Find mismatches between T[a:b) and O[a:b);
 // time: O((|D|+1) log n)
```

```
void diff(vi& out, const ShiftTree& T,
            int vb, int ve, int lvl = -1,
            int b = 0, int e = -1,
            int i = 1, int j = 1) {
    if (e < 0) lvl = __lg(e=len)-1;</pre>
    if (b >= ve || vb >= e || H[i] == T.H[j])
      return;
    if (e-b == 1) return out.push back(b);
    int m = (b+e) / 2;
    int s1 = (delta >> lv1) & 1;
    int s2 = (T.delta >> lvl) & 1;
    diff(out, T, vb, ve, lvl-1, b, m,
      left(i, s1), left(j, s2));
    diff(out, T, vb, ve, lvl-1, m, e,
      right(i, s1), right(j, s2));
  1 // b60c
  void update(int i, int s) {
    auto x = H[left(i, s)] +
     H[right(i, s)] * 0x9E37'79B9'7F4A'7C15;
    x = (x ^ (x>>30)) * 0xBF58'476D'1CE4'E5B9;
    x = (x ^ (x>>27)) * 0x94D0'49BB'1331'11EB;
    H[i] = x ^ (x >> 31);
  } // 3447
  int parent(int i, int s) {
    int k = i + s;
    return k&i ? k/2 : k/4;
  } // 314f
  int left(int i, int s) {
    int k = i * 2, j = k - s;
    return k&j ? j : k|j;
  } // b4eb
  int right(int i, int s) {
   return i*2 + !s:
 } // e440
}; // 6f54
int bitrev(int n, int bits) {
  int ret = 0;
  rep(i, 0, bits)
   ret |= ((n >> i) & 1) << (bits-i-1);
  return ret;
} // 23d1
// Find all attainable subset sums modulo m;
// time: O(m log m)
// Input elements are given by frequency array,
// i.e. counts[x] = how many times element x
// is contained in the input multiset.
// Size of 'counts' is the modulus m.
// The returned array encodes solutions.
// which can be recovered using 'recover'.
// ans[x] != -1 <=> subset with sum x exists
vi subsetSumMod(const vi& counts) {
  int mod = sz(counts), len = 1, k = 0;
  while (len < mod*2) len *= 2, k++;
  vi tmp, ans (mod, -1);
  ShiftTree T(len), Q(len);
  ans[0] = 0;
  T.set(0);
  Q.set(0);
  Q.set (-mod);
  rep(i, 1, len) {
    int x = bitrev(i, k);
    if (x >= mod || !counts[x]) continue;
```

```
O.shift(x - O.delta);
   rep(j, 0, counts[x]) {
     tmp.clear();
     T.diff(tmp, Q, 0, mod);
     if (tmp.empty()) break;
     each(d, tmp) if (ans[d] == -1) {
       ans[d] = x;
       T.set(d);
       0.set(d+x);
       Q.set (d+x-mod);
     } // ce75
   } // c2d1
 } // 9204
 return ans;
} // Oaa2
vi recoverSubset (const vi& dp, int s) {
 assert (dp[s] != -1);
 vi ret:
 while (s) {
   ret.pb(dp[s]);
   s = (s - dp[s] + sz(dp)) % sz(dp);
 } // ea17
 return ret;
} // 7103
segtree/general config.h
// Segment tree configurations to be used
// in general_fixed and general_persistent.
// See comments in TREE PLUS version
// to understand how to create custom ones.
// Capabilities notation: (update; query)
#if TREE_PLUS // (+; sum, max, max count)
 // time: O(lq n)
 using T = int; // Data type for update
                // operations (lazv tag)
  static constexpr T ID = 0; // Neutral value
                // for updates and lazv tags
  // This structure keeps aggregated data
 struct Agg {
   // Aggregated data: sum, max, max count
    // Default values should be neutral
    // values, i.e. "aggregate over empty set"
   T sum = 0, vMax = INT_MIN, nMax = 0;
   int cnt = 0: // And node count.
    // Initialize as leaf (single value)
   void leaf() { sum=vMax=0; nMax=cnt=1; }
   // Combine data with aggregated data
    // from node to the right
   void merge(const Agg& r) {
     if (vMax < r.vMax) nMax = r.nMax;</pre>
     else if (vMax == r.vMax) nMax += r.nMax;
     vMax = max(vMax, r.vMax);
     sum += r.sum;
     cnt += r.cnt;
   } // 262d
   // Apply update provided in `x`:
   // - update aggregated data and 'lazy' tag
   // - return 0 if update should branch
   // (can be used in "segment tree beats")
   // - if you change value of 'x' it will be
       passed to next node to the right
        during updates
   bool apply (T& lazy, T& x) {
     lazy += x;
```

```
sum += x*cnt;
     vMax += x;
     return 1;
   } // 4a4e
}; // f11d
#elif TREE_MAX // (max; max, max count)
 // time: 0(lg n)
 using T = int;
 static constexpr T ID = INT_MIN;
 struct Agg {
   // Aggregated data: max value, max count
   T vMax = INT_MIN, nMax = 0, cnt = 0;
   void leaf() { vMax = 0; nMax = cnt = 1; }
   void merge(const Agg& r) {
     if (vMax < r.vMax) nMax = r.nMax;</pre>
     else if (vMax == r.vMax) nMax += r.nMax;
     vMax = max(vMax, r.vMax);
     cnt += r.cnt;
   1 // 8561
   bool apply (T& lazy, T& x) {
     if (vMax <= x) nMax = cnt;</pre>
     lazy = max(lazy, x);
     vMax = max(vMax, x);
     return 1;
   } // 118c
 }; // 9643
#elif TREE SET // (=; sum, max, max count)
 // time: 0(lg n)
 // Set ID to some unused value.
 using T = int;
 static constexpr T ID = INT MIN;
 struct Agg {
   // Aggregated data: sum, max, max count
   T sum = 0, vMax = INT MIN, nMax = 0, cnt=0;
   void leaf() { sum=vMax=0; nMax=cnt=1; }
   void merge(const Agg& r) {
     if (vMax < r.vMax) nMax = r.nMax;</pre>
     else if (vMax == r.vMax) nMax += r.nMax;
     vMax = max(vMax, r.vMax);
     sum += r.sum;
     cnt += r.cnt;
   } // 262d
   bool apply (T& lazy, T& x) {
     if (x != ID) {
       lazv = x;
       sum = x*cnt;
       vMax = x:
       nMax = cnt:
     1 // Of7e
     return 1;
   } // f684
 }; // 895c
#elif TREE BEATS // (+, min; sum, max)
// time: amortized O(lg n) if not using +
          amortized O(lg^2 n) if using +
 // Lazy tag is pair (add, min).
 // To add x: run update with {x, INT_MAX},
 // to min x: run update with {0, x}.
 // If both parts are provided, addition
 // is applied first, then minimum.
 using T = pii;
 static constexpr T ID = {0, INT_MAX};
 struct Agg {
   // Aggregated data: max value, max count,
                       second max value, sum
```

```
int vMax = INT_MIN, nMax = 0;
   int max2 = INT_MIN, sum = 0, cnt = 0;
   void leaf() { sum=vMax=0; nMax=cnt=1; }
   void merge(const Agg& r) {
     if (r.vMax > vMax) {
       max2 = vMax;
       vMax = r.vMax;
       nMax = r.nMax;
     } else if (r.vMax == vMax) {
       nMax += r.nMax;
     } else if (r.vMax > max2) {
       max2 = r.vMax;
     ) // b074
     max2 = max(max2, r.max2);
     sum += r.sum;
     cnt += r.cnt;
   } // 1a7f
   bool apply (T& lazy, T& x) {
     if (max2 != INT_MIN && max2+x.x >= x.y)
     lazv.x += x.x;
     sum += x.x*cnt;
     vMax += x.x;
     if (max2 != INT_MIN) max2 += x.x;
     if (x.v < vMax) {
       sum -= (vMax-x.v) * nMax;
       vMax = x.y;
     } // 7025
     lazy.y = vMax;
     return 1:
   } // 46b3
 }; // 507e
#endif
                                         c33c
```

segtree/general fixed.h

```
// Highly configurable statically allocated
// interval-interval segment tree; space: O(n)
struct SegTree {
  // Choose/write configuration
  #include "general_config.h"
  // Root node is 1, left is i*2, right i*2+1
  vector<Agg> agg; // Aggregated data for nodes
  vector<T> lazy; // Lazy tags for nodes
  int len = 1:
                 // Number of leaves
  // Initialize tree for n elements; time: O(n)
  SegTree(int n = 0) {
   while (len < n) len ★= 2;
   agg.resize(len*2);
   lazv.resize(len*2, ID);
    rep(i, 0, n) agg[len+i].leaf();
    for (int i = len; --i;) pull(i);
  1 // 0769
  void pull(int i) {
    (agg[i] = agg[i*2]).merge(agg[i*2+1]);
 } // ebdf
  void push(int i) {
   rep(c, 0, 2)
     agg[i*2+c].apply(lazy[i*2+c], lazy[i]);
   lazy[i] = ID;
  } // e5c9
  template<bool U>
  void go (int vb, int ve, int i, int b, int e,
          auto fn) {
    if (vb < e && b < ve)
     if (b < vb || ve < e || !fn(i)) {</pre>
```

```
int m = (b+e) / 2;
        push(i);
        qo<U>(vb, ve, i*2, b, m, fn);
        qo<U>(vb, ve, i*2+1, m, e, fn);
        if (U) pull(i);
     } // 5ff2
 } // 399a
  // Modify interval [b;e) with val; O(lq n)
 void update(int b, int e, T val) {
   go<1>(b, e, 1, 0, len, [&](int i) {
     return agg[i].apply(lazy[i], val);
   }); // 2828
 } // e3e4
  // Query interval [b;e); time: O(lq n)
 Agg query(int b, int e) {
   Agg t; go<0>(b, e, 1, 0, len, [&] (int i) {
     return t.merge(agg[i]), 1;
   }); // c9dd
   return t:
 } // 1c6e
  // Find smallest 'j' such that
  // g(aggregate of [0, j)) is true; O(lg n)
 // The predicate 'g' must be monotonic.
  // Returns -1 if no such prefix exists.
 int lowerBound(auto q) {
   if (!q(aqq[1])) return -1;
   Agg x, s;
   int i = 1:
   for (; i < len; g(s) \mid | (x = s, i++))
     push(i), (s = x).merge(agg[i \star= 2]);
   return i - len + !q(x);
 } // f732
}; // c6c9
```

segtree/general persistent.h f85c

```
// Highly configurable interval-interval
// persistent segment tree; space: O(q lg n)
// First tree version number is 0.
struct SegTree {
 // Choose/write configuration
  #include "general config.h"
 vector<Agg> agg{{}}; // Aggregated data
 vector<T> lazy{ID}; // Lazy tags
 vector<bool> cow{0}; // Copy children on push
                      // Children links
 vi L{0}, R{0};
 int len{1};
                       // Number of leaves
 // Initialize tree for n elements; O(lq n)
 SegTree(int n = 0) {
   int k = 3:
   while (len < n) len \star= 2, k += 3;
   rep(i, 1, k) fork(0);
   iota(all(R)-3, 3);
   L = R;
   if (n--) {
     agg[k -= 3].leaf();
     agg[k+1].leaf();
     for (int i = k-3; i \ge 0; i = 3, n \ne 2)
       (n%2 ? L[i] : ++R[i])++;
     while (k--) pull(k);
   } // 13a7
 } // 4a93
 // New version from version 'i'; time: O(1)
 int fork(int i) {
   L.pb(L[i]); R.pb(R[i]); cow.pb(cow[i] = 1);
   agg.pb(agg[i]); lazy.pb(lazy[i]);
```

```
} // a21b
 void pull(int i) {
   (agg[i] = agg[L[i]]).merge(agg[R[i]]);
 } // 359c
 void push(int i, bool w) {
   if (w || lazy[i] != ID) {
     if (cow[i]) {
       int x = fork(L[i]), y = fork(R[i]);
       L[i] = x; R[i] = y; cow[i] = 0;
     agg[L[i]].apply(lazy[L[i]], lazy[i]);
     agg[R[i]].apply(lazy[R[i]], lazy[i]);
     lazy[i] = ID;
   } // 9f41
 } // 678e
 template<bool U>
 void go(int vb, int ve, int i, int b, int e,
         auto fn) {
   if (vb < e && b < ve)
     if (b < vb || ve < e || !fn(i)) {</pre>
       int m = (b+e) / 2;
        push(i, U);
        go<U>(vb, ve, L[i], b, m, fn);
       go<U>(vb, ve, R[i], m, e, fn);
       if (U) pull(i);
     } // 3fd0
 1 // 3a95
 // Modify interval [b:e) with val
 // in tree version `j`; time: O(lq n)
 void update(int j, int b, int e, T val) {
   go<1>(b, e, j, 0, len, [&](int i) {
     return agg[i].apply(lazy[i], val);
   1): // 2828
 } // 9f22
 // Query interval [b;e) in tree version `i`;
 Agg query(int j, int b, int e) { // O(lg n)
   Agg t; go<0>(b, e, j, 0, len, [&](int i) {
     return t.merge(agg[i]), 1;
   1): // c9dd
   return t;
 } // 2c98
 // Find smallest 'j' such that
 // g(aggregate of [0, j)) is true
 // in tree version 'i'; time: O(lg n)
 // The predicate 'g' must be monotonic.
 // Returns -1 if no such prefix exists.
 int lowerBound(int i, auto g) {
   if (!g(agg[i])) return -1;
   Agg x, s;
   int p = 0, k = len;
   while (L[i]) {
     push(i, 0);
     (s = x).merge(agg[L[i]]);
     k /= 2;
     i = g(s)? L[i]: (x = s, p += k, R[i]);
   } // 7d74
   return p + !q(x);
 } // f7d7
}; // 21f9
segtree/point fixed.h
                                         14b6
// Point-interval segment tree
```

// - T - stored data type

// - ID - neutral element for query operation

return sz(L)-1;

```
// - f(a, b) - associative aggregate function
struct SegTree {
 using T = int;
 static constexpr T ID = INT_MIN;
 T f(T a, T b) { return max(a, b); }
 vector<T> V;
 int len = 1;
  // Initialize tree for n elements; time: O(n)
 SegTree (int n = 0, T def = 0) {
   while (len < n) len *= 2;
   V.resize(len+n, def);
   V.resize(len*2, ID);
   for (int i = len; --i;)
     V[i] = f(V[i*2], V[i*2+1]);
 } // ac47
 // Set element 'i' to 'val'; time: O(lg n)
 void set(int i, T val) {
   V[i += len] = val;
   while (i /= 2) V[i] = f(V[i*2], V[i*2+1]);
 // Query interval [b;e); time: O(lq n)
 T query(int b, int e) {
   T \times = ID, y = ID;
   for (e += len; b < e; b /= 2, e /= 2) {
     if (b % 2) x = f(x, V[b++]);
     if (e % 2) y = f(V[--e], y);
   } // 4ed0
   return f(x, v);
 } // 7816
 // Find smallest 'j' such that
 // g(aggregate of [0,i)) is true; O(lg n)
 // The predicate 'g' must be monotonic.
 // Returns -1 if no such prefix exists.
 int lowerBound(auto g) {
   if (!q(V[1])) return -1;
   T s. x = ID:
   int j = 1;
   while (j < len)
     if (!g(s = f(x, V[j \star= 2]))) x = s, j++;
   return j - len + !q(x);
 } // 6cc5
}; // a0c7
segtree/point persistent.h
                                         4113
// Point-interval persistent segment tree
// - T - stored data type
// - ID - neutral element for query operation
// - f(a, b) - associative aggregate function
// First tree version number is 0.
struct SegTree {
 using T = int;
 static constexpr T ID = INT MIN;
 T f(T a, T b) { return max(a, b); }
 vector<T> agg{ID}; // Aggregated data
 vector<bool> cow{1}; // Copy children on push
 vi L{0}, R{0};
                      // Children links
 int len{1};
                      // Number of leaves
 // Initialize tree for n elements; O(lq n)
 SegTree (int n = 0, T def = 0) {
   int k = 3:
   while (len < n) len \star= 2, k += 3;
   rep(i, 1, k) fork(0);
```

iota(all(R)-3, 3);

L = R;

```
if (n--) {
     k -= 3;
     agg[k] = agg[k+1] = def;
      for (int i = k-3; i \ge 0; i = 3, n \ne 2)
        (n%2 ? L[i] : ++R[i])++;
     while (k--)
       agg[k] = f(agg[L[k]], agg[R[k]]);
   } // 6fde
  } // 3cfb
  // New version from version 'i'; time: O(1)
  int fork(int i) {
   L.pb(L[i]); R.pb(R[i]);
    agg.pb(agg[i]); cow.pb(cow[i] = 1);
   return sz(L)-1;
  } // bb75
  // Set element 'pos' to 'val' in version 'i';
  // time: O(lg n)
  void set (int i, int pos, T val,
           int b = 0, int e = 0) {
    if (L[i]) {
     if (!e) e = len;
     if (cow[i]) {
       int x = fork(L[i]), y = fork(R[i]);
       L[i] = x; R[i] = y; cow[i] = 0;
     } // 82ec
     int m = (b+e) / 2;
     if (pos < m) set(L[i], pos, val, b, m);</pre>
     else set(R[i], pos, val, m, e);
     agg[i] = f(agg[L[i]], agg[R[i]]);
   } else {
     agg[i] = val;
   } // 23c8
  } // 7a55
  // Query interval [b;e) in tree version 'i';
  // time: 0(lq n)
  T query(int i, int vb, int ve,
          int b = 0, int e = 0) {
   if (!e) e = len;
   if (vb >= e || b >= ve) return ID;
   if (b >= vb && e <= ve) return agg[i];</pre>
   int m = (b+e) / 2;
    return f (query (L[i], vb, ve, b, m),
             query(R[i], vb, ve, m, e));
  } // 2664
  // Find smallest 'j' such that
  // g(aggregate of [0,j)) is true
  // in tree version 'i'; time: O(lq n)
  // The predicate 'g' must be monotonic.
  // Returns -1 if no such prefix exists.
  int lowerBound(int i, auto g) {
   if (!g(agg[i])) return -1;
   T x = ID;
   int p = 0, k = len;
   while (L[i]) {
     T s = f(x, agg[L[i]]);
     k /= 2;
     i = g(s) ? L[i] : (x = s, p += k, R[i]);
   } // Ofba
   return p + !q(x);
 } // 1a9a
}; // e4ec
structures/bitset plus.h
                                          6737
// Undocumented std::bitset features:
```

// - Find first() - returns first bit = 1 or N

// - _Find_next(i) - returns first bit = 1

```
after i-th bit
                    or N if not found
// Bitwise operations for vector<bool>
// UNTESTED
#define OP(x) vector<bool>& operator x##=(
   vector<bool>& 1, const vector<bool>& r) { \
 assert(sz(1) == sz(r));
 auto a = 1.begin(); auto b = r.begin();
 while (a<1.end()) *a._M_p++ x##= *b._M_p++; \
 return 1; } // f164
OP (&) OP (|) OP (^)
structures/fenwick tree.h
                                         ec21
// Fenwick tree (BIT tree); space: O(n)
// Default version: prefix sums
struct Fenwick {
 using T = 11:
 static constexpr T ID = 0;
 T f(T a, T b) { return a+b; }
 vector<T> s:
 Fenwick(int n = 0) : s(n, ID) {}
 // A[i] = f(A[i], v); time: O(lg n)
 void modify(int i, T v) {
   for (; i < sz(s); i |= i+1) s[i]=f(s[i],v);
 } // a047
 // Get f(A[0], ..., A[i-1]); time: O(lg n)
 T query(int i) {
  T v = ID;
   for (; i > 0; i \&= i-1) v = f(v, s[i-1]);
   return v;
 } // 9810
 // Find smallest i such that
 // f(A[0],...,A[i-1]) >= val; time: O(lg n)
 // Prefixes must have non-descreasing values.
 int lowerBound(T val) {
   if (val <= ID) return 0;</pre>
   int i = -1, mask = 1;
   while (mask \leq sz(s)) mask \star= 2;
   T off = ID;
   while (mask /= 2) {
     int k = mask+i;
     if (k < sz(s)) {
       T \times = f(off, s[k]);
       if (val > x) i=k, off=x;
     } // de7f
   } // 929c
   return i+2;
 } // 4be9
}; // eb2e
structures/fenwick tree 2d.h 9f31
// Fenwick tree 2D (BIT tree 2D); space: O(n*m)
// Default version: prefix sums 2D
// Change s to hashmap for O(q lg^2 n) memory
struct Fenwick2D {
 using T = int:
 static constexpr T ID = 0;
 T f(T a, T b) { return a+b; }
 vector<T> s:
 int w, h;
 Fenwick2D(int n = 0, int m = 0)
   : s(n*m, ID), w(n), h(m) {}
 // A[i,j] = f(A[i,j], v); time: O(1g^2 n)
```

```
void modify(int i, int j, T v) {
   for (; i < w; i |= i+1)
     for (int k = j; k < h; k = k+1)
        s[i*h+k] = f(s[i*h+k], v);
 } // d46b
 // Query prefix; time: O(lg^2 n)
 T query(int i, int j) {
   T v = ID;
   for (; i>0; i&=i-1)
     for (int k = j; k > 0; k &= k-1)
       v = f(v, s[i*h+k-h-1]);
   return v;
} // 08cf
}; // e570
structures/find union.h
                                         f9a4
// Disjoint set data structure; space: O(n)
// Operations work in amortized O(alfa(n))
struct FAU {
 vi G:
 FAU(int n = 0) : G(n, -1) {}
 // Get size of set containing i
 int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
  int find(int i) {
   return G[i] < 0 ? i : G[i] = find(G[i]);</pre>
 } // 5bc1
  // Union sets containing i and j
 bool join(int i, int j) {
   i = find(i); j = find(j);
   if (i == j) return 0;
   if (G[i] > G[j]) swap(i, j);
   G[i] += G[j]; G[j] = i;
   return 1;
 } // c721
}; // 3839
structures/find union undo.h 399f
// Disjoint set data structure
// with rollback; space: O(n)
// Operations work in O(log(n)) time.
struct RollbackFAU {
 vi G;
 vector<pii> his;
 RollbackFAU(int n = 0) : G(n, -1) {}
  // Get size of set containing i
 int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
 int find(int i) {
   return G[i] < 0 ? i : find(G[i]);</pre>
 } // e478
  // Current time (for rollbacks)
 int time() { return sz(his); }
 // Rollback all operations after time 't'
 void rollback(int t) {
   for (int i = time(); t < i--;)</pre>
     G[his[i].x] = his[i].y;
   his.resize(t);
 } // 3ef3
 // Union sets containing i and j
 bool join(int i, int j) {
   i = find(i); j = find(j);
   if (i == j) return 0;
```

```
if (G[i] > G[j]) swap(i, j);
    his.pb({i, G[i]});
    his.pb({j, G[j]});
    G[i] += G[j]; G[j] = i;
    return 1;
 } // 1491
}; // 18ef
structures/hull offline.h
                                         ed05
constexpr 11 INF = 2e18;
// constexpr double INF = 1e30;
// constexpr double EPS = 1e-9;
// MAX of linear functions; space: O(n)
// Use if you add lines in increasing 'a' order
// Default uncommented version is for int64
struct Hull {
  using T = 11: // Or change to double
  struct Line {
   T a, b, end;
    T intersect (const Line& r) const {
      // Version for double:
      //if (r.a-a < EPS) return b>r.b?INF:-INF;
      //return (b-r.b) / (r.a-a);
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b, d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u^d));
    } // f27f
  }; // 10dc
  vector<Line> S;
  Hull() { S.pb({ 0, -INF, INF }); }
  // Insert f(x) = ax+b; time: amortized O(1)
  void push(T a, T b) {
    Line 1{a, b, INF};
    while (1) {
      T e = S.back().end=S.back().intersect(1);
      if (sz(S) < 2 | | S[sz(S)-2].end < e)
        break:
      S.pop_back();
    } // 044f
    S.pb(1);
  } // 3022
  // Query max(f(x) for each f): time: O(lg n)
 T query (T x) {
    auto t = *upper_bound(all(S), x,
      [](int 1, const Line& r) {
        return 1 < r.end;</pre>
      }); // de77
    return t.a*x + t.b;
 } // b8de
}; // fa73
structures/hull online.h
                                          6884
constexpr 11 INF = 2e18;
// MAX of linear functions online; space: O(n)
struct Hull {
  static bool modeQ; // Toggles operator< mode</pre>
  struct Line {
    mutable 11 a, b, end;
    ll intersect (const Line& r) const {
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b, d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u%d));
    bool operator<(const Line& r) const {</pre>
```

Hull() { S.insert({ 0, -INF, INF }); }

auto cur = it++; cur->end = INF;

if (it == S.end()) return false;

cur->end = cur->intersect(*it);

bool update(multiset<Line>::iterator it) {

} // cfab

multiset < Line > S:

// Updates segment end

}; // 6046

return mode0 ? end < r.end : a < r.a;

```
return cur->end >= it->end;
  } // 63b8
  // Insert f(x) = ax+b; time: O(\lg n)
  void insert(ll a, ll b) {
    auto it = S.insert({ a, b, INF });
    while (update(it)) it = --S.erase(++it);
    rep(i, 0, 2)
      while (it != S.begin() && update(--it))
       update(it = --S.erase(++it));
  // Query max(f(x)) for each f): time: O(\lg n)
  11 query(11 x) {
   mode0 = 1:
    auto 1 = \starS.upper bound({ 0, 0, x });
   mode0 = 0;
   return l.a*x + l.b;
 } // 7533
}; // 037e
bool Hull::mode0 = 0;
structures/intset.h
// Bitset with fast predecessor and successor
// gueries. Can handle 50-200mln operations
// per second. Assumes X86 shift overflows.
template<int N>
struct IntSet {
  uint64 t V[N/64+1] = \{\};
  IntSet<(N < 65 ? 0 : N/64+1)> up;
  // Is 'i' contained in the set?
  bool has(int i) const {
    return (V[i/64] >> i) & 1;
  } // abab
  // Add 'i' to the set.
  void add(int i) {
   if (!V[i/64]) up.add(i/64);
   V[i/64] |= 1ull << i;
  } // 342e
  // Delete 'i' from the set.
  void del(int i) {
   if (!(V[i/64] &= ~(1ull<<i))) up.del(i/64);</pre>
  // Find first element > i, or return -1.
  // 'i' must be in range [0;N).
  int next(int i) {
   auto x = V[i/64] \gg i;
    if (x &= ~1) return i+__builtin_ctzll(x);
    return (i = up.next(i/64)) < 0 ? i :</pre>
     i*64+ builtin ctzll(V[i]);
  } // 8160
  // Find last element < i, or return -1.
  // 'i' must be in range [0:N).
  int prev(int i) {
    auto x = V[i/64] << (63-i);
```

```
if (x &= INT64 MAX)
      return i-__builtin_clzll(x);
    return (i = up.prev(i/64)) < 0 ? i :</pre>
      i*64+63-__builtin_clzll(V[i]);
 } // 4b0d
}; // 6ba3
template<>
struct IntSet<0> {
 void add(int) {}
  void del(int) {}
 int next(int) { return -1; }
 int prev(int) { return -1; }
1: // ace7
structures/li chao tree.h
                                          5559
// Extended Li Chao tree; space: O(n)
// Let F be a family of functions,
// closed under function addition, such that
// for every f != g from the family F
// there exists x such that:
// f(z) \ll g(z) for z \ll x, else f(z) \gg g(z)
// or
// g(z) \ll f(z) for z \ll x, else g(z) \gg f(z).
// Typically F is family of linear functions.
// DS maintains a sequence c[0], \ldots, c[n-1]
// under operations max, add, query
// (see comments below for explanations).
// Configure by modifying:
// - T - type of sequence elements
// - Func - represents function from family F
// - ID ADD - neutral element for function add
// - ID_MAX - neutral element for function max
// TESTED ON RANDS
struct LiChao {
  struct Func {
   ll a, b; // a*x + b
    // Evaluate function in point x
    11 operator()(11 x) const { return a*x+b; }
    // Sum of two functions
    Func operator+ (Func r) const {
     return {a+r.a, b+r.b};
   } // f911
  }; // 633c
  static constexpr Func ID ADD{0, 0};
  static constexpr Func ID_MAX{0, ll(-1e9)};
  vector<Func> val, lazy;
  int len;
  // Initialize tree for n elements; time: O(n)
  LiChao (int n = 0) {
    for (len = 1: len < n: len *= 2):
    val.resize(len*2, ID MAX);
    lazv.resize(len*2, ID ADD);
 } // c0ba
  void push(int i) {
    if (i < len) rep(j, 0, 2) {</pre>
      lazy[i*2+j] = lazy[i*2+j] + lazy[i];
      val[i*2+j] = val[i*2+j] + lazy[i];
    } // 54fc
    lazy[i] = ID_ADD;
  } // 1777
  // For each x in [vb; ve)
  // set c[x] = max(c[x], f(x));
  // time: O(log^2 n) in general case,
          O(\log n) if [vb; ve) = [0; len)
```

```
void max(int vb, int ve, Func f,
           int i = 1, int b = 0, int e = -1) {
    if (e < 0) e = len;</pre>
    if (vb >= e || b >= ve || i >= len*2)
      return;
    int m = (b+e) / 2;
    push(i);
    if (b >= vb && e <= ve) {
      auto& g = val[i];
      if (q(m) < f(m)) swap(q, f);
      if (q(b) < f(b))
       max(vb, ve, f, i*2, b, m);
        max(vb, ve, f, i*2+1, m, e);
      max(vb, ve, f, i*2, b, m);
      \max(vb, ve, f, i*2+1, m, e);
   } // f2c0
 } // 03ed
 // For each x in [vb; ve)
  // set c[x] = c[x] + f(x);
  // time: O(log^2 n) in general case,
          O(1) if [vb; ve) = [0; len)
  void add(int vb, int ve, Func f,
           int i = 1, int b = 0, int e = -1) {
    if (e < 0) e = len;
    if (vb >= e || b >= ve) return;
    if (b >= vb && e <= ve) {
      lazv[i] = lazv[i] + f;
      val[i] = val[i] + f;
    } else {
      int m = (b+e) / 2;
      push(i):
      max(b, m, val[i], i*2, b, m);
      max(m, e, val[i], i*2+1, m, e);
      val[i] = ID MAX;
      add (vb, ve, f, i*2, b, m);
      add(vb, ve, f, i*2+1, m, e);
   } // bbe5
 } // 259f
  // Get value of c[x]; time: O(log n)
  auto query(int x) {
   int i = x+len:
    auto ret = val[i](x);
    while (i \neq 2)
     ret = ::max(ret+lazy[i](x), val[i](x));
    return ret;
 } // dfe4
}; // 0104
structures/max queue.h
                                          3e9e
// Oueue with max query on contained elements
struct MaxOueue {
 using T = int;
 deque<T> 0, M;
  // Add v to the back; time: amortized O(1)
  void push(T v) {
   while (!M.empty() && M.back() < v)</pre>
     M.pop_back();
   M.pb(v); Q.pb(v);
 } // 57a2
  // Pop from the front; time: O(1)
  void pop() {
   if (M.front() == Q.front()) M.pop_front();
   Q.pop_front();
```

```
} // 101c
 // Get max element value; time: O(1)
 T max() const { return M.front(); }
}; // b6c4
structures/rmq.h
                                         b828
// Range Minimum Query; space: O(n lg n)
struct RMO {
 using T = int;
 static constexpr T ID = INT MAX;
 T f(T a, T b) { return min(a, b); }
 vector<vector<T>>> s;
 // Initialize RMO structure; time: O(n lq n)
 RMO(const vector<T>& vec = {}) {
   s = \{vec\};
   for (int h = 1; h \le sz(vec); h *= 2) {
     s.pb({});
     auto& prev = s[sz(s)-2];
     rep(i, 0, sz(vec)-h*2+1)
       s.back().pb(f(prev[i], prev[i+h]));
   1 // 7c37
 } // 14ed
 // Query f(s[b], ..., s[e-1]); time: O(1)
 T query(int b, int e) {
   if (b >= e) return ID;
   int k = __lg(e-b);
   return f(s[k][b], s[k][e - (1<<k)]);</pre>
 ) // bb12
}; // c8f0
structures/treap.h
                                         6156
// "Set" of implicit keyed treaps; space: O(n)
// Nodes are keyed by their indices in array
// of all nodes. Treap key is key of its root.
// "Node x" means "node with key x".
// "Treap x" means "treap with kev x".
// Kev -1 is "null".
// Put any additional data in Node struct.
struct Treap {
 struct Node {
   // E[0] = left child, E[1] = right child
   // weight = node random weight (for treap)
    // size = subtree size, par = parent node
   int E[2] = \{-1, -1\}, weight = rand();
   int size = 1, par = -1;
   bool flip = 0; // Is interval reversed?
 }; // 3036
 vector<Node> G: // Array of all nodes
 // Initialize structure for n nodes
 // with keys 0, ..., n-1; time: O(n)
 // Each node is separate treap.
 // use ioin() to make sequence.
 Treap(int n = 0) : G(n) {}
 // Create new treap (a single node),
 // returns its key; time: O(1)
 int make() { G.pb({}); return sz(G)-1; }
 // Get size of node x subtree. x can be -1.
 int size(int x) { // time: O(1)
   return (x \ge 0 ? G[x].size : 0);
 } // 81cf
 // Propagate down data (flip flag etc).
 // x can be -1; time: O(1)
 void push(int x) {
   if (x >= 0 && G[x].flip) {
```

```
G[x].flip = 0;
    swap(G[x].E[0], G[x].E[1]);
    each (e, G[x].E) if (e>=0) G[e].flip ^= 1;
  } // + any other lazy operations
} // ed19
// Update aggregates of node x.
// x can be -1; time: O(1)
void update(int x) {
  if (x >= 0) {
   int & s = G[x].size = 1;
   G[x].par = -1;
    each (e, G[x].E) if (e >= 0) {
     s += G[e].size;
     G[e].par = x;
   } // f7a7
  } // + any other aggregates
} // 46a3
// Split treap x into treaps 1 and r
// such that 1 contains first i elements
// and r the remaining ones.
// x, 1, r can be -1; time: \sim O(\log n)
void split(int x, int& l, int& r, int i) {
 push (x); 1 = r = -1;
  if (x < 0) return;
  int key = size(G[x].E[0]);
  if (i <= kev) {
   split(G[x].E[0], 1, G[x].E[0], i);
   r = x;
   split(G[x].E[1], G[x].E[1], r, i-key-1);
   1 = x;
 } // fe19
  update(x);
} // 8211
// Join treaps 1 and r into one treap
// such that elements of 1 are before
// elements of r. Returns new treap.
// l, r and returned value can be -1.
int join(int 1, int r) { // time: ~O(lg n)
  push(1); push(r);
 if (1 < 0 || r < 0) return max(1, r);</pre>
  if (G[l].weight < G[r].weight) {</pre>
   G[1].E[1] = join(G[1].E[1], r);
   update(1);
   return 1:
  } // 18c7
  G[r].E[0] = join(l, G[r].E[0]);
 update(r):
 return r;
} // b559
// Find i-th node in treap x.
// Returns its key or -1 if not found.
// x can be -1; time: ~O(lq n)
int find(int x, int i) {
  while (x \ge 0) {
    push(x);
    int key = size(G[x].E[0]);
   if (key == i) return x;
    x = G[x].E[key < i];
   if (key < i) i -= key+1;</pre>
  } // 054c
  return -1;
} // 0b9b
// Get key of treap containing node x
// (key of treap root). x can be -1.
```

```
int root(int x) { // time: ~O(lg n)
   while (G[x].par \ge 0) x = G[x].par;
   return x;
 } // be8b
  // Get position of node x in its treap.
  // x is assumed to NOT be -1; time: \sim O(\log n)
 int index(int x) {
   int p, i = size(G[x].E[G[x].flip]);
    while ((p = G[x].par) >= 0) {
     if (G[p].E[1] == x) i+=size(G[p].E[0])+1;
     if (G[p].flip) i = G[p].size-i-1;
     x = p;
   } // 3f81
   return i;
 } // ddad
  // Reverse interval [1;r) in treap x.
 // Returns new key of treap; time: ~O(lg n)
 int reverse(int x, int 1, int r) {
   int a, b, c;
    split(x, b, c, r);
    split(b, a, b, 1);
   if (b >= 0) G[b].flip ^= 1;
   return join(join(a, b), c);
 } // e418
}; // 17cc
structures/wavelet tree.h
                                          80d3
// Wavelet tree ("merge-sort tree over values")
// Each node represent interval of values.
// seg[1]
             = original sequence
// seq[i]
              = subsequence with values
                represented by i-th node
// left[i][j] = how many values in seq[0:j)
                go to left subtree
struct WaveletTree {
 vector<vi> seq, left;
 int len;
  WaveletTree() {}
  // Build wavelet tree for sequence 'elems';
  // time and space: O((n+maxVal) log maxVal)
  // Values are expected to be in [0; maxVal).
 WaveletTree(const vi& elems, int maxVal) {
    for (len = 1; len < maxVal; len *= 2);</pre>
    seq.resize(len*2);
   left.resize(len*2);
    seq[1] = elems;
   build(1, 0, len);
 1 // a5e9
  void build(int i, int b, int e) {
   if (i >= len) return;
    int m = (b+e) / 2;
    left[i].pb(0);
    each(x, seq[i]) {
     left[i].pb(left[i].back() + (x < m));
     seq[i*2 + (x >= m)].pb(x);
    } // ac25
   build(i*2, b, m);
   build(i*2+1, m, e);
 1 // 8153
 // Find k-th smallest element in [begin; end)
  // [begin;end); time: O(log maxVal)
  int kth(int begin, int end, int k, int i=1) {
   if (i >= len) return seq[i][0];
    int x = left[i][begin], y = left[i][end];
    if (k < y-x) return kth(x, y, k, i*2);
```

```
return kth (begin-x, end-y, k-y+x, i*2+1);
 } // 7861
  // Count number of elements >= vb and < ve
  // in [begin;end); time: O(log maxVal)
  int count (int begin, int end, int vb, int ve,
           int i = 1, int b = 0, int e = -1) {
    if (e < 0) e = len;
    if (b >= ve || vb >= e) return 0;
    if (b >= vb && e <= ve) return end-begin;
    int m = (b+e) / 2;
    int x = left[i][begin], y = left[i][end];
    return count (x, y, vb, ve, i*2, b, m) +
      count (begin-x, end-y, vb, ve, i*2+1, m, e);
 } // 71cf
}; // 49a9
structures/ext/hash table.h 2d30
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
// gp_hash_table<K, V> = faster unordered_set
// Anti-anti-hash
const size_t HXOR = mt19937_64(time(0))();
template < class T> struct SafeHash {
 size_t operator()(const T& x) const {
    return hash<T>()(x ^ T(HXOR));
 } // 3a78
}; // 7d0e
structures/ext/heap.h
                                         d41d
#include <ext/pb_ds/priority_queue.hpp>
// Pairing heap: push O(1), pop O(1g n)
// __gnu_pbds::priority_queue<T, Cmp>
// Standard priority_queue methods and:
// 1. Iterable
// 2. t.erase(iterator)
                                       0(la n)
// 3. t.modifv(iterator, value)
                                       0(la n)
// 4. t1.join(t2) - merge t2 into t1
                                         051f
structures/ext/rope.h
#include <ext/rope>
using namespace gnu cxx;
// rope<T> = persistent implicit cartesian tree
structures/ext/tree.h
                                         a3bc
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template < class T, class Cmp = less < T>>
using ordered_set = tree<</pre>
 T, null_type, Cmp, rb_tree_tag,
 tree_order_statistics_node_update
// Standard set functions and:
// t.order_of_key(key) - index of first >= key
// t.find by order(i) - find i-th element
// t1.join(t2) - assuming t1<>t2 merge t2 to t1
structures/ext/trie.h
                                         5cc2
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/trie_policy.hpp>
using namespace __gnu_pbds;
using pref_trie = trie<</pre>
 string, null_type,
  trie_string_access_traits , pat_trie_tag,
```

trie_prefix_search_node_update

```
text/aho corasick.h
                                         fc9b
constexpr char AMIN = 'a'; // Smallest letter
constexpr int ALPHA = 26; // Alphabet size
// Aho-Corasick algorithm for linear-time
// multiple pattern matching.
// Add patterns using add(), then call build().
struct Aho {
 vector<arrav<int, ALPHA>> nxt{1};
 vi suf = \{-1\}, accLink = \{-1\};
 vector<vi> accept{1};
 // Add string with given ID to structure
 // Returns index of accepting node
 int add(const string& str, int id) {
   int i = 0:
   each(c, str) {
     if (!nxt[i][c-AMIN]) {
       nxt[i][c-AMIN] = sz(nxt);
       nxt.pb({}); suf.pb(-1);
       accLink.pb(1); accept.pb({});
     } // 5ead
     i = nxt[i][c-AMIN];
   1 // ace9
   accept[i].pb(id);
   return i;
 1 // 27c8
 // Build automata; time: O(V*ALPHA)
 void build() {
   queue<int> que:
   each(e, nxt[0]) if (e) {
     suf[e] = 0; que.push(e);
   } // c34d
   while (!que.empty()) {
     int i = que.front(), s = suf[i], j = 0;
     que.pop();
     each(e, nxt[i]) {
       if (e) que.push(e);
        (e ? suf[e] : e) = nxt[s][j++];
      accLink[i] = (accept[s].empty() ?
         accLink[s] : s);
   } // 1e8a
 } // 2561
  // Append 'c' to state 'i'
 int next(int i, char c) {
   return nxt[i][c-AMIN];
 } // 6bb7
 // Call 'f' for each pattern accepted
 // when in state 'i' with its ID as argument.
 // Return true from 'f' to terminate early.
 // Calls are in descreasing length order.
 void accepted(int i, auto f) {
   while (i !=-1) {
     each(a, accept[i]) if (f(a)) return;
     i = accLink[i];
   } // c175
 } // 1f0d
}; // 5c1c
text/alcs.h
                                         a97c
// All-substrings common sequences algorithm.
// Given strings A and B, algorithm computes:
// C(i, j, k) = |LCS(A[:i), B[j:k))|
// in compressed form; time and space: O(n^2)
```

// To describe the compression, note that:

```
// 1. C(i, j, k-1) \le C(i, j, k) \le C(i, j, k-1)+1
// 2. If j < k and C(i, j, k) = C(i, j, k-1)+1,
     then C(i, j+1, k) = C(i, j+1, k-1) + 1
// 3. If j >= k, then C(i, j, k) = 0
// This allows us to store just the following:
// ih (i, k) = min j s.t. C(i, j, k-1) < C(i, j, k)
struct ALCS {
 string A, B;
  vector<vi> ih;
  ALCS() {}
  // Precompute compressed matrix; time: O(nm)
  ALCS(string s, string t) : A(s), B(t) {
   int n = sz(A), m = sz(B);
   ih.resize(n+1, vi(m+1));
   iota(all(ih[0]), 0);
    rep(1, 1, n+1) {
     int iv = 0;
     rep(j, 1, m+1) {
       if (A[1-1] != B[j-1]) {
         ih[l][j] = max(ih[l-1][j], iv);
          iv = min(ih[l-1][j], iv);
       } else {
          ih[1][j] = iv;
          iv = ih[1-1][j];
       } // 7af8
     } // d115
   } // baff
  } // b761
  // Compute | LCS(A[:i), B[i:k))|; time: O(k-i)
  // Note: You can precompute data structure
  // to answer these queries in O(log n)
  // or compute all answers for fixed 'i'.
  int operator()(int i, int j, int k) {
   int ret = 0;
   rep(q, j, k) ret += (ih[i][q+1] <= j);
   return ret:
  } // dabf
  // Compute subsequence LCS(A[:i), B[j:k));
  // time: O(k-i)
  string recover(int i, int j, int k) {
   string ret;
   while (i > 0 && j < k) {
     if (ih[i][k--] <= j) {</pre>
       ret.pb(B[k]);
       while (A[--i] != B[k]);
     } // 9d77
   } // ledf
   reverse (all (ret));
   return ret;
  } // 738c
  // Compute LCS'es of given prefix of A,
  // and all prefixes of given suffix of B.
  // Returns vector L of length |B|+1 s.t.
  // L[k] = |LCS(A[:i), B[j:k))|; time: O(|B|)
  vi row(int i, int j) {
   vi ret(sz(B)+1);
   rep(k, j+1, sz(ret))
     ret[k] = ret[k-1] + (ih[i][k] \leftarrow j);
   return ret;
  } // 9167
  // Compute LCS'es of given prefix of A,
  // and all substrings of B; time: O(n^2)
  // Return matrix M such that:
  //M[i][k] = |LCS(A[:i), B[i:i+k))|
  vector<vi> matrix(int i) {
```

```
vector<vi> ret:
    rep(j, 0, sz(B)+1) ret.pb(row(i, j));
    return ret;
 } // 15f7
}; // fd6b
                                          e912
text/hashing.h
using ull = uint64_t;
// Arithmetic mod 2^64-1.
// Around 2x slower than mod 2^64.
struct Hash {
  constexpr Hash(ull v = 0) : x(v) {}
  Hash operator+ (Hash r) {
   return x + r.x + (x + r.x < x);
  } // b42e
  Hash operator-(Hash r) {
   return *this + ~r.x;
  } // e855
  Hash operator* (Hash r) {
    auto m = uint128 t(x) * r.x;
   return Hash(ull(m)) + ull(m>>64);
  1 // 1241
  auto get() const { return x + !~x; }
  bool operator==(Hash r) const {
   return get() == r.get();
  } // d4d5
  bool operator<(Hash r) const {</pre>
    return get() < r.get();</pre>
 } // 34a9
  void print() { cerr << x; }</pre>
}; // 6064
// Base for hashing (prime, big order).
constexpr Hash C = 11(1e11-981);
Hash powC(int n) { // C^n
  static vector<Hash> vec = {1};
  while (sz(vec) <= n) vec.pb(vec.back() * C);</pre>
 return vec[n];
} // 3ff8
// Precompute prefix hashes for a string.
struct HashInterval : vector<Hash> {
 HashInterval (auto& s) {
    pb(0);
    rep(i, 0, sz(s)) pb(at(i)*C + s[i]);
  } // 2b42
  // Get hash of interval [b;e)
  Hash operator()(int b, int e) {
    return at (e) - at (b) * powC(e-b);
 } // 4adc
}; // 6c35
text/kmp.h
                                          a014
// Computes prefsuf array; time: O(n)
// ps[i] = max prefsuf of [0;i); ps[0] := -1
vi kmp(auto& str) {
 vi ps; ps.pb(-1);
  each(x, str) {
    int k = ps.back();
    while (k \ge 0 \&\& str[k] != x) k = ps[k];
   ps.pb(k+1);
  1 // 05aa
 return ps;
} // fa90
// Finds occurences of pat in vec; time: O(n)
// Returns starting indices of matches.
vi match (auto& str, T pat) {
```

```
int n = sz(pat);
 pat.pb(-1); // SET TO SOME UNUSED CHARACTER
 pat.insert(pat.end(), all(str));
 vi ret, ps = kmp(pat);
 rep(i, 0, sz(ps)) {
   if (ps[i] == n) ret.pb(i-2*n-1);
 } // ale9
 return ret;
} // c6e8
text/kmr.h
                                          7b40
// KMR algorithm for O(1) lexicographical
// comparison of substrings.
struct KMR {
 vector<vi> ids;
 KMR() {}
 // Initialize structure; time: O(n 1g^2 n)
  // You can change str type to vi freely.
 KMR (const string& str) {
   ids.clear();
   ids.pb(vi(all(str)));
    for (int h = 1; h \le sz(str); h *= 2) {
     vector<pair<pii, int>> tmp;
     rep(j, 0, sz(str)) {
       int a = ids.back()[i], b = -1;
        if (j+h < sz(str)) b = ids.back()[j+h];
       tmp.pb({ {a, b}, j });
     } // a210
      sort(all(tmp));
     ids.emplace_back(sz(tmp));
     int n = 0;
     rep(j, 0, sz(tmp)) {
       if (j > 0 \&\& tmp[j-1].x != tmp[j].x)
        ids.back()[tmp[j].y] = n;
     } // bd2e
   } // cf37
 } // d7a7
  // Get representative of [begin; end); 0(1)
 pii get (int begin, int end) {
   if (begin >= end) return {0, 0};
   int k = lg(end-begin);
   return {ids[k][begin], ids[k][end-(1<<k)]};</pre>
 1 // 6ele
 // Compare [b1;e1) with [b2;e2); O(1)
  // Returns -1 if <, 0 if ==, 1 if >
 int cmp(int b1, int e1, int b2, int e2) {
   int 11 = e1-b1, 12 = e2-b2;
   int 1 = min(11, 12);
   pii x = get(b1, b1+1), y = get(b2, b2+1);
   if (x == y) return (11 > 12) - (11 < 12);</pre>
   return (x > y) - (x < y);
 1 // 5d4e
 // Compute suffix array of string; O(n)
 vi sufArrav() {
   vi sufs(sz(ids.back()));
   rep(i, 0, sz(ids.back()))
     sufs[ids.back()[i]] = i;
   return sufs:
 } // 455e
}; // 2fb3
```

```
text/lcp.h
                                         e309
// Compute Longest Common Prefix array for
// given string and it's suffix array; O(n)
// In order to compute suffix array use kmr.h
// or suffix_array_linear.h
vi lcpArray(auto& str, vi& sufs) {
 int n = sz(str), k = 0;
 vi pos(n), lcp(n-1);
 rep(i, 0, n) pos[sufs[i]] = i;
 rep(i, 0, n) {
   if (pos[i] < n-1) {</pre>
     int j = sufs[pos[i]+1];
     while (i+k < n && j+k < n &&
         str[i+k] == str[j+k]) k++;
     lcp[pos[i]] = k;
   } // 2cba
   if (k > 0) k--;
 1 // 8b22
 return lcp;
} // 4202
text/lyndon factorization.h 688c
// Compute Lyndon factorization for s; O(n)
// Word is simple iff it's stricly smaller
// than any of it's nontrivial suffixes.
// Lyndon factorization is division of string
// into non-increasing simple words.
// It is unique.
vector<string> duval(const string& s) {
 int n = sz(s), i = 0;
 vector<string> ret;
 while (i < n) {
   int j = i+1, k = i;
   while (j < n && s[k] <= s[j])
     k = (s[k] < s[j] ? i : k+1), j++;
   while (i <= k)
     ret.pb(s.substr(i, j-k)), i += j-k;
 } // 3f17
 return ret:
} // 0e48
text/main lorentz.h
                                         401c
#include "z_function.h"
struct Sqr {
 int begin, end, len;
}; // f012
// Main-Lorentz algorithm for finding
// all squares in given word; time: O(n lq n)
// Results are in compressed form:
// (b, e, 1) means that for each b <= i < e
// there is square at position i of size 21.
// Each square is present in only one interval.
vector<Sqr> lorentz(const string& s) {
 vector<Sgr> ans;
 vi pos(sz(s)/2+2, -1);
 rep(mid, 1, sz(s)) {
   int part = mid & ~(mid-1), off = mid-part;
   int end = min(mid+part, sz(s));
   auto a = s.substr(off, part);
   auto b = s.substr(mid, end-mid);
   string ra(a.rbegin(), a.rend());
   string rb(b.rbegin(), b.rend());
   rep(j, 0, 2) {
     // Set # to some unused character!
     vi z1 = prefPref(ra);
```

```
vi z2 = prefPref(b+"#"+a);
      z1.pb(0); z2.pb(0);
      rep(c, 0, sz(a)) {
        int 1 = sz(a)-c;
        int x = c - \min(1-1, z1[1]);
        int y = c - max(1-z2[sz(b)+c+1], j);
        if (x > y) continue;
        int sb = (j ? end-y-1*2 : off+x);
        int se = (i ? end-x-1*2+1 : off+v+1);
        int& p = pos[1];
        if (p != -1 && ans[p].end == sb)
          ans[p].end = se;
          p = sz(ans), ans.pb({sb, se, 1});
      } // af4b
     a.swap(rb);
     b.swap(ra);
   } // 193e
  } // 4fa7
  return ans;
} // 5b80
text/manacher.h
                                         4dbe
// Manacher algorithm; time: O(n)
// Finds largest radiuses for palindromes:
// p[0][i] for center between i-1 and i
// p[1][i] for center at i (single letter = 0)
array<vi, 2> manacher(auto& s) {
  int n = sz(s), l = 0, r = 0;
  array<vi, 2> p = {vi(n+1), vi(n)};
  rep(i, 0, n) rep(z, 0, 2) {
   int t = r-i+!z, &x = p[z][i];
   if (i < r) x = min(t, p[z][l+t]);
   int b = i-x-1, e = i+x+z;
    while (b >= 0 && e < n && s[b] == s[e])
     x++, b--, e++;
   if (r < e) l = b+1, r = e-1;</pre>
  } // 6fd0
 return p:
} // e211
text/min rotation.h
                                         e4d6
// Find lexicographically smallest
// rotation of s: time: O(n)
// Returns index where shifted word starts.
// You can use std::rotate to get the word:
// rotate(s.begin(), s.begin()+minRotation(s),
         s.end());
int minRotation(string s) {
  int a = 0, n = sz(s); s += s;
  rep(b, 0, n) rep(i, 0, n) {
   if (a+i == b || s[a+i] < s[b+i]) {
     b += max(0, i-1); break;
    } // 865b
   if (s[a+i] > s[b+i]) {
     a = b; break;
   } // 7628
  } // 40be
  return a;
} // 9ed8
text/monge.h
                                         e6a5
// NxN matrix A is simple (sub-)unit-Monge
// iff there exists a (sub-)permutation
// (N-1) \times (N-1) matrix P such that:
```

// A[x,y] = sum i>=x, j<y: P[i,j]

```
// The first column and last row are always 0.
// We represent these matrices implicitly
// using permutations p s.t. P[i,p(i)] = 1.
// (min, +) product of simple unit-Monge
// matrices represented by permutations P, Q,
// is also a simple unit-Monge matrix.
// The permutation that describes the product
// can be obtained by the following procedure:
// 1. Decompose P, O into minimal sequences of
// elementary transpositions.
// 2. Concatenate the transposition sequences.
// 3. Scan from left to right and remove
   transpositions that decrease
     inversion count (i.e. second crossings).
// 4. The reduced sequence represents result.
// Invert sub-permutation with values [0:n).
// Missing values should have value `def`.
vi invert (const vi& P, int n, int def) {
 vi ret(n, def);
 rep(i, 0, sz(P)) if (P[i] != def)
   ret[P[i]] = i;
 return ret;
} // 035e
// Split permutation P into half 'lo'
// with values less than 'k', and half 'hi'
// with remaining values, shifted by 'k'.
// Missing rows from 'lo' and 'hi' are removed,
// original indices are in 'loInd' and 'hiInd'.
void split (const vi& P, int k, vi& lo, vi& hi,
          vi& loInd, vi& hiInd) {
 int i = 0:
 each (e, P) {
   if (e < k) lo.pb(e), loInd.pb(i++);
   else hi.pb(e-k), hiInd.pb(i++);
 } // c3a6
} // 7bb7
// Map sub-permutation into sub-permutation
// of length 'n' on given indices sets.
vi expand (const vi& P, vi& indl, vi& ind2,
         int n, int def) {
 vi ret(n, def);
 rep(k, 0, sz(P)) if (P[k] != def)
   ret[ind1[k]] = ind2[P[k]];
 return ret;
} // 0da7
// Compute (min, +) product of square
// simple unit-Monge matrices given their
// permutation representations; time: O(n lg n)
// Permutation of second matrix is inverted!
vi comb (const vi& P, const vi& invO) {
 int n = sz(P):
 if (n < 100) {
   // 5s -> 1s speedup for ALIS for n = 10^5
   vi ret = invert(P, n, -1);
   rep(i, 0, sz(invQ)) {
     int from = invQ[i];
     rep(j, 0, i) from += invQ[j] > invQ[i];
     for (int j = from; j > i; j--)
       if (ret[j-1] < ret[j])</pre>
         swap (ret[j-1], ret[j]);
   } // 7cd1
   return invert(ret, n, -1);
 } // 679e
 vi p1, p2, q1, q2, i1, i2, j1, j2;
 split (P, n/2, p1, p2, i1, i2);
```

```
split(invQ, n/2, q1, q2, j1, j2);
 p1 = expand(comb(p1, q1), i1, j1, n, -1);
 p2 = expand(comb(p2, q2), i2, j2, n, n);
 q1 = invert(p1, n, -1);
 q2 = invert(p2, n, n);
 vi ans(n, -1);
 int delta = 0, j = n;
  rep(i, 0, n) {
   ans[i] = (p1[i] < 0 ? p2[i] : p1[i]);
   while (j > 0 && delta >= 0)
     delta -= (q2[--j] < i || q1[j] >= i);
   if (p2[i] < j || p1[i] >= j)
     if (delta++ < 0)
        if (q2[j] < i || q1[j] >= i)
         ans[i] = j;
 } // c396
 return ans;
} // c059
// Helper function for `mongeMul`.
void padPerm(const vi& P, vi& has, vi& pad,
           vi& ind, int n) {
 vector<bool> seen(n);
 rep(i, 0, sz(P)) if (P[i] != -1) {
   ind.pb(i);
   has.pb(P[i]);
   seen[P[i]] = 1;
 } // 157e
 rep(i, 0, n) if (!seen[i]) pad.pb(i);
} // 103b
// Compute (min, +) product of
// simple sub-unit-Monge matrices given their
// permutation representations; time: O(n lg n)
// Left matrix has size sz(P) x sz(O).
// Right matrix has size sz(0) x n.
// Output matrix has size sz(P) x n.
// NON-SOUARE MATRICES ARE NOT TESTED!
vi mongeMul(const vi& P, const vi& Q, int n) {
 vi h1, p1, i1, h2, p2, i2;
 padPerm(P, h1, p1, i1, sz(Q));
 padPerm(invert(Q, n, -1), h2, p2, i2, sz(Q));
 h1.insert(h1.begin(), all(p1));
 h2.insert(h2.end(), all(p2));
 vi ans(sz(P), -1), tmp = comb(h1, h2);
 rep(i, 0, sz(i1)) {
   int j = tmp[i+sz(p1)];
   if (j < sz(i2)) {
     ans[i1[i]] = i2[j];
   } // 4d16
 } // c8a0
 return ans;
} // 3326
// Range Longest Increasing Subsequence Query;
// preprocessing: O(n 1g^2 n), query: O(lg n)
#include "../structures/wavelet tree.h"
struct ALIS {
 WaveletTree tree:
 ALIS() {}
  // Precompute data structure; O(n 1g^2 n)
 ALIS(const vi& seq) {
   vi P = build(seq);
   each(k, P) if (k == -1) k = sz(seq);
   tree = \{P, sz(seq)+1\};
 } // f00f
```

```
// Query LIS of s[b;e); time: O(lg n)
 int operator()(int b, int e) {
   return e - b -
     tree.count(b, sz(tree.seg[1]), 0, e);
 } // fb4a
 vi build(const vi& seq) {
   int n = sz(seq);
   if (!n) return {};
   int lo = *min_element(all(seq));
   int hi = *max element(all(seq));
   if (10 == hi) {
     vi tmp(n);
     iota(all(tmp), 1);
     tmp.back() = -1;
     return tmp:
   } // 989d
   int mid = (lo+hi+1) / 2;
   vi p1, p2, i1, i2;
   split(seq, mid, p1, p2, i1, i2);
   p1 = expand(build(p1), i1, i1, n, -1);
   p2 = expand(build(p2), i2, i2, n, -1);
   each(j, i1) p2[j] = j;
   each(i, i2) p1[i] = i;
   return mongeMul(p1, p2, n);
 } // 6517
}; // 27ea
text/palindromic tree.h
                                         f86e
constexpr int ALPHA = 26; // Set alphabet size
// Tree of all palindromes in string,
// constructed online by appending letters.
// space: O(n*ALPHA); time: O(n)
struct PalTree {
 vi txt: // Text for which tree is built
 // Node 0 = empty palindrome (root of even)
 // Node 1 = "-1" palindrome (root of odd)
 vi len{0, -1}; // Lengths of palindromes
 vi link{1, 0}; // Suffix palindrome links
 // Edges to next palindromes
 vector<array<int, ALPHA>> to{ {}, {} };
 int last{0}; // Current node (max suffix pal)
#if MIN PALINDROME PARTITION
 // An extension that computes minimal
 // palindromic partition in O(n log n).
 vi diff{0, 0}; // len[i]-len[link[i]]
 vi slink{0, 0}; // Serial links
 vi series{0, 0}; // Series DP answer
                  // DP answer for prefix
 vi ans{0};
#endif
 int ext(int i) {
   while (len[i]+2 > sz(txt) ||
          txt[sz(txt)-len[i]-2] != txt.back())
      i = link[i];
   return i;
 } // d442
 // Append letter from [0;ALPHA]; time: O(1)
 // (or O(lg n) for MIN PALINDROME PARTITION)
 void add(int x) {
   txt.pb(x);
   last = ext(last);
   if (!to[last][x]) {
```

len.pb(len[last]+2);

```
link.pb(to[ext(link[last])][x]);
     to[last][x] = sz(to);
     to.pb({});
    #if MIN PALINDROME PARTITION
     diff.pb(len.back() - len[link.back()]);
     slink.pb(diff.back() == diff[link.back()]
       ? slink[link.back()] : link.back());
     series.pb(0);
    #endif
    } // e432
    last = to[last][x];
  #if MIN PALINDROME PARTITION
    ans.pb(INT_MAX);
    for (int i=last; len[i] > 0; i=slink[i]) {
     series[i] = ans[sz(ans) - len[slink[i]]
                     - diff[i] - 1];
     if (diff[i] == diff[link[i]])
       series[i] = min(series[i],
                       series[link[i]]);
     // If you want only even palindromes
     // set ans only for sz(txt) %2 == 0
     ans.back() = min(ans.back(), series[i]+1);
   } // ab3b
  #endif
 } // 14a4
}; // f8d9
```

text/suffix array linear.h

```
#include "../util/radix sort.h"
// KS algorithm for suffix array; time: O(n)
// Input values are assumed to be in [1;k]
vi sufArray(vi str, int k) {
 int n = sz(str);
  vi suf(n):
  str.resize(n+15);
  if (n < 15) {
   iota(all(suf), 0);
    rep(j, 0, n) countSort(suf,
      [&] (int i) { return str[i+n-j-1]; }, k);
    return suf:
  } // 5fcf
  // Compute triples codes
  vi tmp, code (n+2);
  rep(i, 0, n) if (i % 3) tmp.pb(i);
  rep(j, 0, 3) countSort(tmp,
   [&] (int i) { return str[i-j+2]; }, k);
  int mc = 0, j = -1;
  each(i, tmp) {
   code[i] = mc += (i == -1)
       str[i] != str[i] ||
       str[i+1] != str[i+1] ||
       str[i+2] != str[j+2]);
   j = i;
  } // bfdc
  // Compute suffix array of 2/3
  tmp.clear();
  for (int i=1; i < n; i += 3) tmp.pb(code[i]);</pre>
  tmp.pb(0):
  for (int i=2; i < n; i += 3) tmp.pb(code[i]);</pre>
  tmp = sufArray(move(tmp), mc);
  // Compute partial suffix arrays
  vi third:
 int th = (n+4) / 3;
```

```
if (n%3 == 1) third.pb(n-1);
 rep(i, 1, sz(tmp)) {
   int e = tmp[i];
   tmp[i-1] = (e 
   code[tmp[i-1]] = i;
   if (e < th) third.pb(e*3);
 } // f9f1
 tmp.pop_back();
 countSort (third,
   [&] (int i) { return str[i]; }, k);
 // Merge suffix arrays
 merge(all(third), all(tmp), suf.begin(),
   [&](int 1, int r) {
     while (1%3 == 0 | | r%3 == 0) {
       if (str[1] != str[r])
         return str[1] < str[r];</pre>
       1++: r++:
     } // 2f8a
     return code[l] < code[r];</pre>
   }); // 4cb3
return suf:
} // 671f
// KS algorithm for suffix array; time: O(n)
vi sufArray(const string& str) {
return sufArray(vi(all(str)), 255);
} // 2f32
text/suffix automaton.h
                                        e45b
constexpr char AMIN = 'a'; // Smallest letter
constexpr int ALPHA = 26; // Set alphabet size
// Suffix automaton - minimal DFA that
```

```
// recognizes all suffixes of given string
// (and encodes all substrings);
// space: O(n*ALPHA); time: O(n)
// Paths from root are equivalent to substrings
struct SufDFA {
 // State v represents endpos-equivalence
 // class that contains words of all lengths
 // between link[len[v]]+1 and len[v].
 // len[v] = longest word of equivalence class
 // link[v] = link to state of longest suffix
              in other equivalence class
 // to[v][c] = automaton edge c from v
 vi len{0}, link{-1};
 vector<array<int, ALPHA>> to{ {} };
 int last{0}; // Current node (whole word)
#if COUNT SUBSTR OCCURENCES
 vector<vi> inSufs: // Suffix-link tree
 vi cnt{0};
                    // Occurence count
#endif
#if COUNT OUTGOING PATHS
 vector<ll> paths; // Out-path count
 SufDFA() {}
 // Build suffix automaton for given string
 // and compute extended stuff; time: O(n)
 SufDFA(const string& s) {
   each(c, s) add(c);
   finish();
 } // ec2e
 // Append letter to the back
 void add(char c) {
```

int v = last, x = c-AMIN;

```
last = sz(len);
   len.pb(len[v]+1);
   link.pb(0);
   to.pb({});
   cnt.pb(1); // COUNT SUBSTR OCCURENCES
   while (v != -1 && !to[v][x]) {
     to[v][x] = last;
     v = link[v];
   } // 4cfc
   if ( \lor != -1 )  {
     int q = to[v][x];
     if (len[v]+1 == len[q]) {
       link[last] = q;
     } else {
       len.pb(len[v]+1);
       link.pb(link[q]);
       to.pb(to[q]);
       cnt.pb(0); // COUNT SUBSTR OCCURENCES
       link[last] = link[q] = sz(len)-1;
       while (v != -1 \&\& to[v][x] == g) {
         to[v][x] = link[q];
         v = link[v];
       } // 784f
     } // 90aa
   } // af69
 } // 345a
 // Go using edge 'c' from state 'i'.
 // Returns 0 if edge doesn't exist.
 int next(int i, char c) {
   return to[i][c-AMIN];
 1 // c363
 // Compute extended stuff (offline)
 void finish() {
 #if COUNT SUBSTR OCCURENCES
   inSufs.resize(sz(len));
   rep(i, 1, sz(link)) inSufs[link[i]].pb(i);
   dfsSufs(0):
 #if COUNT OUTGOING PATHS
   paths.assign(sz(len), 0);
   dfs(0);
 #endif
 } // d3dc
#if COUNT SUBSTR OCCURENCES
 void dfsSufs(int v) {
   each(e, inSufs[v]) {
     dfsSufs(e);
     cnt[v] += cnt[e];
   } // 2469
} // 0c60
#endif
#if COUNT OUTGOING PATHS
 void dfs(int v) {
   if (paths[v1) return;
   paths[v] = 1;
   each(e, to[v]) if (e) {
     dfs(e);
     paths[v] += paths[e];
   } // 22b3
 } // d004
 // Get lexicographically k-th substring
 // of represented string; time: O(|substr|)
 // Empty string has index 0.
 string lex(ll k) {
   string s;
```

```
int v = 0:
    while (k--) rep(i, 0, ALPHA) {
      int e = to[v][i];
      if (e) {
       if (k < paths[e]) {</pre>
          s.pb(char(AMIN+i));
          v = e;
         break;
       } // f307
       k -= paths[e];
     } // 29be
    } // 4600
   return s;
 } // e4af
#endif
}; // ef50
text/suffix tree.h
                                         8a6e
constexpr int ALPHA = 26;
// Ukkonen's algorithm for online suffix tree
// construction; space: O(n*ALPHA); time: O(n)
// Real tree nodes are called dedicated nodes.
// "Nodes" lying on compressed edges are called
// implicit nodes and are represented
// as pairs (lower node, label index).
// Labels are represented as intervals [L;R)
// which refer to substrings [L;R) of txt.
// Leaves have labels of form [L; infinity),
// use getR to get current right endpoint.
// Suffix links are valid only for internal
// nodes (non-leaves).
struct SufTree {
 vi txt; // Text for which tree is built
 // to[v][c] = edge with label starting with c
               from node v
 vector<array<int, ALPHA>> to{ {} };
 vi L{0}, R{0}; // Parent edge label endpoints
              // Parent link
 vi par{0};
 vi link{0}; // Suffix link
 pii cur{0, 0}; // Current state
  // Get current right end of node label
 int getR(int i) { return min(R[i],sz(txt)); }
  // Follow edge 'e' of implicit node 's'.
  // Returns (-1, -1) if there is no edge.
 pii next(pii s, int e) {
   if (s.v < getR(s.x))
      return txt[s.y] == e ? pii(s.x, s.y+1)
                           : pii(-1, -1);
    e = to[s.x][e];
    return e ? pii(e, L[e]+1) : pii(-1, -1);
 // Create dedicated node for implicit node
 // and all its suffixes
 int split(pii s) {
   if (s.v == R[s.x]) return s.x;
    int t = sz(to); to.pb({});
    to[t][txt[s.y]] = s.x;
    L.pb(L[s.x]);
    R.pb(L[s.x] = s.y);
    par.pb(par[s.xl);
    par[s.x] = to[par[t]][txt[L[t]]] = t;
    link.pb(-1);
    int v = link[par[t]], l = L[t] + !par[t];
```

while (1 < R[t]) {

v = to[v][txt[1]];

```
Jagiellonian University - Jagiellonian 1
     1 += getR(v) - L[v];
   } // 0393
   v = split(\{v, getR(v)-l+R[t]\});
   link[t] = v;
   return t;
  } // 10bb
  // Append letter from [0; ALPHA] to the back
  void add(int x) { // amortized time: 0(1)
   pii t; txt.pb(x);
    while ((t = next(cur, x)).x == -1) {
     int m = split(cur);
     to[m][x] = sz(to);
     to.pb({});
     par.pb(m):
     L.pb(sz(txt)-1);
     R.pb(INT_MAX);
     link.pb(-1);
     cur = {link[m], getR(link[m])};
     if (!m) return:
   } // 60c2
   cur = t;
 } // 1e43
}; // 8926
text/z function.h
                                         770b
// Computes Z function array; time: O(n)
// zf[i] = max common prefix of str and str[i:]
vi prefPref(auto& str) {
 int n = sz(str), b = 0, e = 1;
  vi zf(n);
  rep(i, 1, n) {
   if (i < e) zf[i] = min(zf[i-b], e-i);</pre>
   while (i+zf[i] < n &&
     str[zf[i]] == str[i+zf[i]]) zf[i]++;
   if (i+zf[i] > e) b = i, e = i+zf[i];
```

trees/centroid decomp.h

vi par, depth, size;

int root; // Root centroid

} // e906

zf[0] = n;

return zf:

} // b7a7

// Centroid decomposition; space: O(n lg n) struct CentroidTree { // child[v] = children of v in centroid tree // par[v] = parent of v in centroid tree (-1 for root) // depth[v] = depth of v in centroid tree (0 for root) // ind[v][i] = index of vertex v in i-th centroid subtree from root // size[v] = size of centroid subtree of v // subtree[v] = list of vertices in centroid subtree of v // dists[v] = distances from v to vertices in its centroid subtree (in the order of subtree[v]) // neigh[v] = neighbours of v in its centroid subtree // dir[v][i] = index of centroid neighbour that is first vertex on path from centroid v to i-th vertex of centroid subtree (-1 for centroid) vector<vi> child, ind, dists, subtree, neigh, dir;

607a

```
CentroidTree() {}
 CentroidTree (vector<vi>& G)
     : child(sz(G)), ind(sz(G)), dists(sz(G)),
       subtree(sz(G)), neigh(sz(G)),
       dir(sz(G)), par(sz(G), -2),
       depth(sz(G)), size(sz(G)) {
   root = decomp(G, 0, 0);
 } // 026c
 void dfs(vector<vi>& G, int v, int p) {
   size[v] = 1;
   each (e, G[v]) if (e != p && par[e] == -2)
     dfs(G, e, v), size[v] += size[e];
 } // bbed
 void layer(vector<vi>& G, int v,
            int p, int c, int d) {
   ind[v].pb(sz(subtree[c]));
   subtree[c].pb(v);
   dists[c].pb(d);
   dir[c].pb(sz(neigh[c])-1);
   each(e, G[v]) if (e != p && par[e] == -2) {
     if (v == c) neigh[c].pb(e);
     layer(G, e, v, c, d+1);
   } // dc82
 } // 37ee
 int decomp(vector<vi>& G, int v, int d) {
   dfs(G, v, -1);
   int p = -1, s = size[v];
   each (e, G[v]) {
     if (e != p && par[e] == -2 &&
         size[e] > s/2) {
       p = v; v = e; goto loop;
     } // e0a5
   1 // 3533
   par[v] = -1;
   size[v] = s;
   depth[v] = d;
   layer(G, v, -1, v, 0);
   each(e, G[v]) if (par[e] == -2) {
     int j = decomp(G, e, d+1);
     child[v].pb(j);
     par[j] = v;
   } // 70b5
   return v:
 } // 217c
}; // 1253
```

trees/centroid offline.h

void computeSize(int v, int p) {

```
dd93
// Helper for offline centroid decomposition
// Usage: CentroidDecomp(G);
// Constructor calls method `process`
// for each centroid subtree.
struct CentroidDecomp {
 vector<vi>← G; // Reference to target graph
 vector<bool> on: // Is vertex enabled?
 vi size; // Used internally
 // Run centroid decomposition for graph g
 CentroidDecomp(vector<vi>& a)
     : G(q), on (sz(q), 1), size (sz(q)) {
   decomp(0);
 1 // 8677
 // Compute subtree sizes for subtree rooted
 // at v, ignoring p and disabled vertices
```

```
size[v] = 1:
   each(e, G[v]) if (e != p && on[e])
     computeSize(e, v), size[v] += size[e];
 } // 1c0d
 void decomp(int v) {
   computeSize(v, -1);
   int p = -1, s = size[v];
   each(e, G[v]) {
     if (e != p && on[e] && size[e] > s/2) {
       p = v; v = e; goto loop;
     } // e0a5
   ) // f31d
   process(v);
   on[v] = 0;
   each (e, G[v]) if (on[e]) decomp (e);
 } // f170
 // Process current centroid subtree:
 // - v is centroid
 // - boundary vertices have on[x] = 0
 // Formally: Let H be subgraph induced
 // on vertices such that on [v] = 1.
 // Then current centroid subtree is
 // connected component of H that contains v
 // and v is its centroid.
 void process(int v) {
   // Do your stuff here...
 } // d41d
}; // f598
trees/compress tree.h
                                         12da
#include "lca.h" // or lca linear.h
using vpi = vector<pair<int, int>>;
// Given a rooted tree and a subset S of nodes,
// compute the minimal subtree that contains
// all the nodes by adding all pairwise LCA's
// and compressing edges; time: O(|S| log |S|)
// Returns a list of (par, orig\_index)
// representing a tree rooted at 0.
// The root points to itself.
vpi compressTree(LCA& lca, const vi& subset) {
 static vi rev; rev.resize(sz(lca.pre));
 vi li = subset, &T = lca.pre;
 auto cmp = [&](int a, int b) {
   return T[a] < T[b];</pre>
 }; // df37
 sort(all(li), cmp);
 int m = sz(li)-1;
 rep(i, 0, m) {
   int a = li[i], b = li[i+1];
   li.push_back(lca(a, b));
 } // 8757
 sort(all(li), cmp);
 li.erase(unique(all(li)), li.end());
 rep(i, 0, sz(li)) rev[li[i]] = i;
 vpi ret = {pii(0, li[0])};
 rep(i, 0, sz(li)-1) {
   int a = li[i], b = li[i+1];
   ret.emplace_back(rev[lca(a, b)], b);
 } // 5101
 return ret:
} // ef6b
trees/heavylight_decomp.h
                                         5562
#include "../seqtree/point fixed.h"
```

// Heavy-Light Decomposition of tree

```
// with subtree query support; space: O(n)
struct HLD {
 // G[v] = children of v (no parents!)
 // G[v][0] = heavy child of v
 // par[v] = parent of vertex v
 // size[v] = size of subtree rooted at v
 // depth[v] = distance from root to v
 // pos[v] = index of v in "HLD preorder"
 // head[v] = first vertex of chain with v
 // len[v] = length of chain starting at v
              (0 if v is not head of chain)
 // order[i] = i-th vertex in "HLD preorder"
 vector<vi> G;
 vi par, size, depth, pos, head, len, order;
 SegTree tree: // Vertices are in HLD order
 // Initialize structure for tree G
 // and given root v; time: O(n lg n)
 HLD (vector<vi> H, int v)
     : G(move(H)), par(sz(G), -1),
       size(sz(G), 1), depth(sz(G)),
       pos(sz(G)), head(sz(G)), len(sz(G)) {
   dfs(v);
   go(v, v);
   tree = {sz(order)};
 } // Oadc
 void dfs(int v) {
   erase(G[v], par[v]);
   each(e, G[v]) {
     depth[e] = depth[par[e] = v] + 1;
     dfs(e):
     size[v] += size[e];
     if (size[e] > size[G[v][0]])
       swap(G[v][0], e);
   } // 2ffe
 } // 5f1c
 void go(int v, int h) {
   pos[v] = sz(order);
   len[head[v] = h1++;
   order.pb(v):
   each(e, G[v]) go(e, G[v][0] == e ? h : e);
 } // 89bf
 // Level Ancestor Query; time: O(lq n)
 int lag(int v, int level) {
   for (;; v = par[v]) {
     int k = level - depth[v = head[v]];
     if (k >= 0) return order[pos[v]+k];
   } // 45a8
 1 // c3a8
 // Lowest Common Ancestor; time: O(lg n)
 int lca(int a, int b) {
   for (;;) {
     int ha = head[a], hb = head[b];
     if (ha == hb)
       return depth[a] < depth[b] ? a : b;</pre>
      if (depth[ha] > depth[hb]) a = par[ha];
     else b = par[hb];
   } // 1341
 } // 493e
 // Call func (begin, end, isAscending)
 // for each path segment in order
 // from a to b; time: O(lg n * time of func)
 // func can be called on empty intervals!
 void iterPath(int a, int b, auto func) {
   for (static vector<pii> tmp;;) {
```

```
int ha = head[a], hb = head[b];
      if (ha == hb) {
       bool f = (pos[a] > pos[b]);
       if (f) swap(a, b);
       // Remove +1 from pos[a]+1 for vertex
        // queries (with +1 -> edges).
        func (pos[a]+1, pos[b]+1, !f);
        reverse(all(tmp));
        each (e, tmp) func (e.x, e.y, 0);
        return tmp.clear();
     } // 5b4d
     if (depth[ha] > depth[hb]) {
        func (pos[ha], pos[a]+1, 1);
       a = par[ha];
     } else {
       tmp.pb({pos[hb], pos[b]+1});
       b = par[hb];
     } // 1a37
   } // af03
  ) // 771c
  // Query path between a and b; O(1g^2 n)
  SegTree::T gueryPath(int a, int b) {
   auto ret = tree.ID:
   iterPath(a, b, [&](int i, int j, bool) {
     ret = tree.f(ret, tree.guerv(i, j));
   }); // 1113
   return ret;
  } // 4221
  // Query subtree of v; time: O(lq n)
  SegTree::T guervSubtree(int v) {
   return tree.query(pos[v], pos[v]+size[v]);
}; // 7f6f
                                         048e
```

```
trees/lca.h
// LAQ and LCA using jump pointers
// space: O(n lq n)
struct LCA {
  vector<vi> jumps;
  vi level, pre, post;
  int cnt = 0, depth;
  LCA() {}
  // Initialize structure for tree G
  // and root r; time: O(n lg n)
  LCA(vector<vi>& G, int root)
      : jumps(sz(G)), level(sz(G)),
       pre(sz(G)), post(sz(G)) {
    dfs(G, root, root);
    depth = int(log2(sz(G))) + 2;
    rep(j, 0, depth) each(v, jumps)
     v.pb(jumps[v[j]][j]);
  void dfs(vector<vi>& G, int v, int p) {
   level[v] = p == v ? 0 : level[p]+1;
    jumps[v].pb(p);
   pre[v] = ++cnt;
    each(e, G[v]) if (e != p) dfs(G, e, v);
    post[v] = ++cnt;
  } // e286
  // Check if a is ancestor of b; time: O(1)
  bool isAncestor(int a, int b) {
    return pre[a] <= pre[b] &&
           post[b] <= post[a];</pre>
  } // 5514
```

```
// Lowest Common Ancestor; time: O(lg n)
  int operator()(int a, int b) {
   for (int j = depth; j--;)
     if (!isAncestor(jumps[a][j], b))
        a = jumps[a][j];
   return isAncestor(a, b) ? a : jumps[a][0];
 } // 27d8
  // Level Ancestor Ouery; time: O(lq n)
 int lag(int a, int lvl) {
   for (int j = depth; j--;)
     if (lvl <= level[jumps[a][j]])</pre>
        a = jumps[a][j];
   return a:
 } // 75b3
  // Get distance from a to b; time: O(lg n)
 int distance(int a, int b) {
   return level[a] + level[b] -
          level[operator()(a, b)]*2;
  // Get k-th vertex on path from a to b,
 // a is 0, b is last; time: O(lq n)
  // Returns -1 if k > distance(a, b)
 int kthVertex(int a, int b, int k) {
    int c = operator()(a, b);
   if (level[a]-k >= level[c])
     return lag(a, level[a]-k);
   k += level[c] *2 - level[a];
   return (k > level[b] ? -1 : lag(b, k));
 } // 46c9
}; // 2861
                                          22f7
```

trees/lca linear.h

} // 0a9a

void rot(int p, int i) {

```
// LAQ and LCA using jump pointers
// with linear memory; space: O(n)
struct LCA {
 vi par, jmp, depth, pre, post;
 int cnt = 0;
 LCA() {}
 // Initialize structure for tree G
  // and root v; time: O(n lg n)
 LCA(vector<vi>& G, int v)
     : par(sz(G), -1), jmp(sz(G), v),
       depth(sz(G)), pre(sz(G)), post(sz(G)) {
   dfs(G, v):
 } // 94cf
 void dfs(vector<vi>& G, int v) {
   int j = jmp[v], k = jmp[j], x =
     depth[v]+depth[k] == depth[j]*2 ? k : v;
   pre[v] = ++cnt;
   each(e, G[v]) if (!pre[e]) {
     par[e] = v; jmp[e] = x;
     depth[e] = depth[v]+1;
     dfs(G, e);
   } // b123
   post[v] = ++cnt;
 } // 3280
  // Level Ancestor Query; time: O(lg n)
 int lag(int v, int d) {
   while (depth[v] > d)
     v = depth[jmp[v]] < d ? par[v] : jmp[v];
   return v;
 } // f509
 // Lowest Common Ancestor; time: O(lq n)
 int operator()(int a, int b) {
```

```
if (depth[a] > depth[b]) swap(a, b);
   b = laq(b, depth[a]);
   while (a != b) {
     if (jmp[a] == jmp[b])
       a = par[a], b = par[b];
        a = jmp[a], b = jmp[b];
   } // fe08
   return a:
 } // 25ff
 // Check if a is ancestor of b; time: O(1)
 bool isAncestor(int a, int b) {
   return pre[a] <= pre[b] &&</pre>
           post[b] <= post[a];</pre>
 // Get distance from a to b; time: O(lg n)
 int distance(int a, int b) {
   return depth[a] + depth[b] -
           depth[operator()(a, b)]*2;
 // Get k-th vertex on path from a to b,
 // a is 0, b is last; time: O(lq n)
 // Returns -1 if k > distance(a, b)
 int kthVertex(int a, int b, int k) {
   int c = operator()(a, b);
   if (depth[a]-k >= depth[c])
     return lag(a, depth[a]-k);
   k += depth[c]*2 - depth[a];
   return (k > depth[b] ? -1 : laq(b, k));
 } // 34ed
}; // c19e
trees/link cut_tree.h
                                          23eb
constexpr int INF = 1e9;
// Link/cut tree; space: O(n)
// Represents forest of (un)rooted trees.
struct LinkCutTree {
 vector<array<int, 2>> child;
 vi par, prev, flip, size;
 // Initialize structure for n vertices; O(n)
 // At first there's no edges.
 LinkCutTree(int n = 0)
      : child(n, \{-1, -1\}), par(n, -1),
        prev(n, -1), flip(n, -1), size(n, 1) {}
 void push(int x) {
   if (x >= 0 && flip[x]) {
     flip[x] = 0;
     swap(child[x][0], child[x][1]);
     each(e, child[x]) if (e>=0) flip[e] ^= 1;
   } // + any other lazy path operations
 } // bae2
 void update(int x) {
   if (x >= 0) {
     size[x] = 1;
     each(e, child[x]) if (e \geq= 0)
        size[x] += size[e];
   } // + any other path aggregates
 } // 8ec0
 void auxLink(int p, int i, int ch) {
   child[p][i] = ch;
   if (ch \ge 0) par[ch] = p;
   update(p);
```

```
int x = child[p][i], g = par[x] = par[p];
  if (g \ge 0) child[g][child[g][1] == p] = x;
  auxLink(p, i, child[x][!i]);
  auxLink(x, !i, p);
  swap(prev[x], prev[p]);
  update(q);
} // 4c76
void splay(int x) {
  while (par[x] \geq= 0) {
    int p = par[x], g = par[p];
    push(g); push(p); push(x);
    bool f = (child[p][1] == x);
    if (a >= 0) {
      if (child[g][f] == p) { // zig-zig}
        rot(q, f); rot(p, f);
      } else { // zig-zag
        rot(p, f); rot(q, !f);
      } // 2ebb
    } else { // zig
      rot (p, f);
    } // f8a2
  } // 446b
  push(x);
} // 55a7
// After this operation x becomes the end
// of preferred path starting in root;
void access(int x) { // amortized O(lg n)
  while (1) {
    splay(x);
    int p = prev[x];
    if (p < 0) break;
    prev[x] = -1;
    splay(p);
    int r = child[p][1];
    if (r \ge 0) swap(par[r], prev[r]);
    auxLink(p, 1, x);
  1 // 2b87
} // d224
// Make x root of its tree; amortized O(lq n)
void makeRoot(int x) {
  access(x);
  int& 1 = child[x][0];
  if (1 >= 0) {
    swap(par[1], prev[1]);
    flip[1] ^= 1;
    update(1);
    1 = -1;
    update(x):
 } // 0064
} // b246
// Find root of tree containing x
int find(int x) { // time: amortized O(lq n)
  access(x);
  while (child[x][0] >= 0)
    push(x = child[x][0]);
  splay(x);
 return x;
1 // d78d
// Add edge x-y; time: amortized O(lg n)
// Root of tree containing y becomes
// root of new tree.
void link(int x, int y) {
  makeRoot(x); prev[x] = y;
} // fb4f
```

```
// Remove edge x-y; time: amortized O(lq n)
  // x and y become roots of new trees!
  void cut(int x, int y) {
   makeRoot(x); access(y);
   par[x] = child[v][0] = -1;
   update(y);
  } // 1908
  // Get distance between x and v,
  // returns INF if x and v there's no path.
  // This operation makes x root of the tree!
  int dist(int x, int y) { // amortized O(lg n)
   makeRoot(x):
   if (find(v) != x) return INF;
   access(y);
   int t = child[y][0];
   return t >= 0 ? size[t] : 0;
 } // ae69
}; // 0197
util/arc interval cover.h
                                         5209
using dbl = double;
// Find size of smallest set of points
// such that each arc contains at least one
// of them; time: O(n la n)
int arcCover(vector<pair<dbl, dbl>>& inters,
            dbl wrap) {
  int n = sz(inters);
  rep(i, 0, n) {
   auto& e = inters[i];
   e.x = fmod(e.x, wrap);
   e.y = fmod(e.y, wrap);
   if (e.x < 0) e.x += wrap, e.y += wrap;</pre>
   if (e.x > e.y) e.y += wrap;
   inters.pb({e.x+wrap, e.y+wrap});
  1 // a73b
  vi nxt(n);
  deque<dbl> que;
  dbl r = wrap*4;
  sort (all (inters));
  for (int i = n*2-1; i--;) {
   r = min(r, inters[i].y);
   que.push front(inters[i].x);
   while (!que.empty() && que.back() > r)
     que.pop_back();
   if (i < n) nxt[i] = i+sz(que);</pre>
  } // 5e6c
  int a = 0, b = 0;
  do {
   a = nxt[a] % n;
   b = nxt[nxt[b]%n] % n;
  } while (a != b);
  int ans = 0;
  while (b < a+n) {
   b += nxt[b%n] - b%n;
   ans++;
  } // 7350
  return ans;
} // e6b2
util/bit hacks.h
                                         2a84
// builtin popcount - count number of 1 bits
// __builtin_clz - count most significant 0s
// __builtin_ctz - count least significant 0s
// __builtin_ffs - like ctz, but indexed from 1
                  returns 0 for 0
```

```
// For 11 version add 11 to name
using ull = uint64 t:
// Transpose 64x64 bit matrix
void transpose64 (array<ull, 64>& M) {
 #define T(s,up)
   for (ull i=0; i<64; i+=s*2)
      for (ull j = i; j < i+s; j++) {
        ull \ a = (M[j] >> s) \& up;
        ull \ b = (M[j+s] \& up) << s;
        M[j] = (M[j] \& up) / b;
        M[j+s] = (M[j+s] & (up << s)) | a; 
      } // a290
 T(1, 0x5555'5555'5555'5555);
 T(2, 0x3333'3333'3333'3333);
 T(4, 0xF0F'0F0F'0F0F'0F0F);
 T(8,
         0xFF'00FF'00FF'00FF);
             0xFFFF'0000'FFFF);
 T(16,
 T (32,
                  0xFFFF'FFFFLL);
 #undef T
} // cba2
// Lexicographically next mask with same
// amount of ones.
int nextSubset(int v) {
 int t = v | (v - 1);
 return (t + 1) | (((~t & -~t) - 1) >>
      ( builtin ctz(v) + 1));
} // 4c0c
// Permutation -> integer conversion.
int permToInt(vi& v) { // Not order preserving!
 int use = 0, i = 0, r = 0;
 each (x, v) {
   r = r * + + i +
     __builtin_popcount(use & -(1<<x));
   use l = 1 \ll x:
 } // d764
 return r:
) // 6574
util/bump alloc.h
                                          09f9
// Allocator, which doesn't free memory.
char mem[400<<201: // Set memory limit</pre>
size t nMem;
void* operator new(size_t n) {
 nMem += n; return &mem[nMem-n];
} // fba6
void operator delete(void*) {}
util/compress vec.h
                                          33ee
// Compress integers to range [0;n) while
// preserving their order; time: O(n lg n)
// Returns mapping: compressed -> original
vi compressVec(vector<int*>& vec) {
 sort (all (vec).
   [](int* l, int* r) { return *l < *r; });
 vi old:
 each (e, vec) {
   if (old.empty() | old.back() != *e)
     old.pb(*e);
   \star e = sz(old)-1;
 } // 7eb0
 return old;
} // d53e
util/deque undo.h
// Deque-like undoing on data structures with
```

```
// amortized O(log n) overhead for operations.
// Maintains a deque of objects alongside
// a data structure that contains all of them.
// The data structure only needs to support
// insertions and undoing of last insertion
// using the following interface:
// - insert (...) - insert an object to DS
// - time() - returns current version number
// - rollback(t) - undo all operations after t
// Assumes time() == 0 for empty DS.
struct DequeUndo {
 DataStructure ds; // Configure DS type here.
 vector<tuple<int, int>> elems[2];
 vector<pii> his{{0,0}};
 // Push object to front or back of deque,
 // depending on side parameter.
 void push(auto val, bool side) {
   elems[side].pb(val);
   doPush(0, side);
 } // df9f
 // Pop object from front or back of deque,
 // depending on side parameter.
 void pop(int side) {
   auto &A = elems[side], &B = elems[!side];
   int cnt[2] = {};
   if (A.empty()) {
     assert(!B.emptv());
     auto it = B.begin() + sz(B)/2 + 1;
     A.assign(B.begin(), it);
     B.erase (B.begin(), it);
     reverse (all(A));
     his.resize(1);
     cnt[0] = sz(A);
     cnt[1] = sz(B);
   } else {
     do {
        cnt[his.back().v ^ side]++;
        his.pop back():
     } while (cnt[0] *2 < cnt[1] &&
               cnt[0] < sz(A));
   } // b4ef
   cnt[0]--;
   A.pop back();
   ds.rollback(his.back().x);
   for (int i : {1, 0})
     while (cnt[i]) doPush(--cnt[i], i^side);
 } // 6eba
 void doPush(int i, bool s) {
   apply([&](auto... x) { ds.insert(x...); },
     elems[s].rbegin()[i]);
   his.pb({ds.time(), s});
 } // 4fed
}; // 189b
util/int128 io.h
                                          a481
istream& operator>>(istream& i, int128& x) {
 char s[\overline{50}], *p = s;
 for (i >> s, \bar{x} = 0, p += *p < 48; *p;)
   x = x*10 + *p++ - 48;
 if (*s == 45) x = -x;
 return i:
1 // 6015
// Note: Doesn't work for INT128 MIN!
ostream& operator<<(ostream& o, __int128 x) {
if (x < 0) \circ << '-', x = -x;
```

```
char s[50] = \{\}, *p = s+49;
 for (; x > 9; x /= 10) *--p = char(x%10+48);
 return o << 11(x) << p;
1 // b9ed
util/inversion vector.h
                                        dcdb
// Get inversion vector for sequence of
// numbers in [0;n); ret[i] = count of numbers
// greater than perm[i] to the left; O(n lq n)
vi encodeInversions(vi perm) {
 vi odd, ret(sz(perm));
 int cont = 1;
 while (cont) {
   odd.assign(sz(perm)+1, 0);
   cont = 0;
   rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
      else ret[i] += odd[perm[i]+1];
     cont += perm[i] /= 2;
   } // 4ed0
 ) // a4f0
 return ret;
} // 86e4
// Count inversions in sequence of numbers
// in [0;n); time: O(n lq n)
11 countInversions(vi perm) {
 11 ret = 0, cont = 1;
 vi odd:
 while (cont) {
   odd.assign(sz(perm)+1, 0);
   cont = 0;
   rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret += odd[perm[i]+1];
     cont += perm[i] /= 2;
   } // 916f
 } // c9b5
 return ret;
} // 4bd8
util/longest inc subseq.h
                                         98cb
// Longest Increasing Subsequence; O(n lg n)
vi lis(const vi& seq) {
 vi dp(sz(seq)+1, INT_MAX);
 vi ind(sz(dp), -1), prv(sz(dp));
 rep(i, 0, sz(seq)) {
   int j = int(lower_bound(1+all(dp), seq[i])
               - dp.begin());
   prv[i] = ind[i-1];
   dp[j] = seq[ind[j] = i];
 } // b229
 int i = *--find(1+all(ind), -1);
 while (i !=-1) ret.pb(i), i = prv[i];
 reverse (all (ret));
 return ret:
} // c0fc
util/max rects.h
                                         4b65
struct MaxRect {
 // begin = first column of rectangle
 // end = first column after rectangle
 // hei = height of rectangle
 // touch = columns of height hei inside
 int begin, end, hei;
```

```
vi touch; // sorted increasing
// Given consecutive column heights find
// all inclusion-wise maximal rectangles
// contained in "drawing" of columns; time O(n)
vector < MaxRect > getMaxRects (vi hei) {
  hei.insert(hei.begin(), -1);
  hei.pb(-1);
  vi reach (sz(hei), sz(hei)-1);
  vector<MaxRect> ans;
  for (int i = sz(hei)-1; --i;) {
   int j = i+1, k = i;
    while (hei[j] > hei[i]) j = reach[j];
    reach[i] = i:
    while (hei[k] > hei[i-1]) {
     ans.pb(\{ i-1, 0, hei[k], \{\} \});
     auto& rect = ans.back();
      while (hei[k] == rect.hei) {
       rect.touch.pb(k-1);
       k = reach[k];
     } // 6e7e
     rect.end = k-1;
   } // e03f
  } // 2796
  return ans;
} // f8f9
util/mo.h
// Modified MO's queries sorting algorithm,
// slightly better results than standard.
// Allows to process q queries in O(n*sqrt(q))
struct Query {
 int begin, end;
1: // b76d
// Get point index on Hilbert curve
11 hilbert(int x, int y, int s, 11 c = 0) {
  if (s <= 1) return c;</pre>
  s /= 2; c \star = 4;
  if (v < s)
   return hilbert (x \in (s-1), y, s, c+(x>=s)+1);
  if (x < s)
   return hilbert (2*s-y-1, s-x-1, s, c);
  return hilbert (y-s, x-s, s, c+3);
// Get good order of queries; time: O(n lq n)
vi moOrder(vector<Ouery>€ queries, int maxN) {
  int s = 1:
  while (s < maxN) s \star= 2;
  vector<ll> ord;
  each (q, queries)
   ord.pb(hilbert(q.begin, q.end, s));
  vi ret(sz(ord));
  iota(all(ret), 0);
  sort(all(ret), [&](int 1, int r) {
   return ord[l] < ord[r];</pre>
  }); // 9aea
  return ret:
} // 29f4
util/multinomial.h
                                          a0a3
// Computes n! / (k1! * .. * kn!)
ll multinomial(vi& v) {
 ll c = 1, m = v.empty() ? 1 : v[0];
 rep(i, 1, sz(v)) rep(j, 0, v[i])
```

```
c = c * ++m / (j+1);
 return c;
} // d07d
util/packing.h
// Utilities for packing precomputed tables.
// Encodes 13 bits using two characters.
// Example usage:
// Writer out:
    out.ints(-123, 8);
    out.done();
    cout << out.buf;
struct Writer {
 string buf:
 int cur = 0, has = 0;
  void done() {
   buf.pb(char(cur%91 + 35));
   buf.pb(char(cur/91 + 35));
   cur = has = 0;
  // Write unsigned b-bit integer.
 void intu(uint64 t v, int b) {
   assert (b == 64 \mid \mid v < (1ul1 << b));
   while (b--) {
     cur |= (v & 1) << has;
     if (++has == 13) done();
     v >>= 1;
   } // f132
 } // Of64
  // Write signed b-bit integer (sign included)
 void ints(ll v, int b) {
   intu(v < 0 ? -v*2+1 : v*2, b);
 } // 08d0
}; // 7d0d
// Example usage:
// Reader in("packed_data");
// int firstValue = in.ints(8);
struct Reader {
 const char *buf;
 11 \text{ cur} = 0;
 Reader (const char *s) : buf (s) {}
 // Read unsigned b-bit integer.
 uint64_t intu(int b) {
   uint64_t n = 0;
    rep(i, 0, b) {
      if (cur < 2) {
        cur = *buf++ + 4972;
        cur += *buf++ * 91;
     } // a930
     n = (cur & 1) << i;
     cur >>= 1;
   } // f12f
   return n;
  // Read signed b-bit integer (sign included).
 11 ints(int b) {
   auto v = intu(b);
   return (v%2 ? -1 : 1) * 11(v/2);
} // 1fc9
1: // 2217
util/parallel binsearch.h
                                          02bb
// Run 'n' binary searches on [b;e) parallely.
```

// 'cmp' should be lambda with arguments:

```
// 1) vector<pii>& - pairs (v, i)
     which are queries if value for index i
     is greater or equal to v;
     pairs are sorted by v
// 2) vector<bool>& - output vector,
     set true at index i if value
     for i-th query is >= queried value
// Returns vector of found values;
// time: O((n+c) lg range), where c is cmp time
vi multiBS(int b, int e, int n, auto cmp) {
 if (b >= e) return vi(n, b);
 vector<pii> que(n), rng(n, {b, e});
 vector<bool> ans(n);
 rep(i, 0, n) que[i] = \{(b+e)/2, i\};
 for (int k = __lg(e-b); k >= 0; k--) {
   int last = 0, j = 0;
   cmp (que, ans);
   rep(i, 0, sz(que)) {
     pii &q = que[i], &r = rng[q.y];
     if (q.x != last) last = q.x, j = i;
     (ans[i] ? r.x : r.y) = q.x;
     q.x = (r.x+r.y) / 2;
     if (!ans[i]) swap(que[i], que[j++]);
   } // 4765
 } // 8bc8
 vi ret;
 each (p, rng) ret.pb(p.x);
 return ret:
} // 638f
```

util/radix sort.h

0573

```
// Stable countingsort: time: O(k+sz(vec))
// See example usage in radixSort for pairs.
void countSort(vi& vec, auto key, int k) {
 static vi buf, cnt;
 vec.swap(buf);
 vec.resize(sz(buf));
  cnt.assign(k+1, 0);
  each (e, buf) cnt [key (e)]++;
  rep(i, 1, k+1) cnt[i] += cnt[i-1];
  for (int i = sz(vec)-1; i >= 0; i--)
   vec[--cnt[key(buf[i])]] = buf[i];
// Compute order of elems, k is max key; O(n)
vi radixSort(const vector<pii>& elems, int k) {
 vi order(sz(elems));
 iota(all(order), 0);
 countSort (order,
   [&] (int i) { return elems[i].y; }, k);
  countSort (order,
   [&] (int i) { return elems[i].x; }, k);
 return order;
} // f272
```

Pick's theorem

For a simple polygon with integer vertices, area A, i grid points in the interior, and b grid points on the boundary: A = i + b/2 - 1.

Tutte matrix (perfect matching test)

$$M_{ij} = \begin{cases} x_{ij} & \text{if } ij \in E, i < j \\ -x_{ji} & \text{if } ij \in E, i > j \\ 0 & \text{otherwise} \end{cases}$$

 $det(M) = 0 \iff$ no perfect matching w.h.p.

Kirchhoff's theorem (# of spanning trees)

$$M_{ij} = \begin{cases} \deg_{\text{in}}(i) & \text{if } i = j \\ -\#(ij \text{ edges}) & \text{if } i \neq j \end{cases}$$

M' = M with *i*-th row and column removed det(M') = # of oriented spanning trees rooted at *i*

Cayley's formula (# of labelled trees)

For degree sequence $d_1, ..., d_n$:

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

 $\begin{array}{l} n_1n_2...n_kn^{k-2}=\text{for }k\text{ existing trees of size }n_i\\ kn^{n-k-1}=\text{forests on }n\text{ vertices with }k\text{ components}\\ \text{such that }1,...,k\text{ belong to different components}\\ x_1\ldots x_n(x_1+\ldots+x_n)^{n-2}=\sum_T x_1^{d_1(T)}\ldots x_n^{d_n(T)} \end{array}$

of partitions into positive integers

$$p(0) = 1$$

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

Bernoulli numbers

EGF of Bernoulli numbers is $B(t) = \frac{t}{e^t - 1}$ (FFT-able). $B[0, \ldots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \ldots]$ Sums of powers:

$$\sum_{i=1}^{n} n^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} \cdot (n+1)^{m+1-k}$$

Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k)$$

$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$
$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

Bell numbers

Total number of partitions of n distinct elements. $B(n)=1,1,2,5,15,52,203,877,4140,21147,\ldots$ For p prime,

$$B(p^{m} + n) \equiv mB(n) + B(n+1) \pmod{p}$$
$$B(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot B(k)$$

Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \ C_{n+1} = \sum_{i=1}^{n} C_i C_{n-i}$$

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, \dots$$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with with n+1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines
- permutations of [n] with no 3-term increasing subseq.

Catalan convolution: find the count of balanced parentheses sequences consisting of n+k pairs of parentheses where the first k symbols are open brackets.

$$C^k = \frac{k+1}{n+k+1} {2n+k \choose n}$$

Burnside's lemma

G= group that acts on a set X $X^g=$ set of elements fixed by $g\in G$ X/G= set of orbits, i.e. equivalence classes by G

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$

to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

Sums

$$c^{a} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1} \quad \text{if } c \neq 1$$

$$1 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

Series

$$\begin{split} e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \, (|x| < \infty) \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \, (x \le 1) \\ \sqrt{1+x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \, (|x| \le 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \, (|x| < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \, (|x| < \infty) \end{split}$$

Trigonometry

$$\sin(v \pm w) = \sin v \cos w \pm \cos v \sin w$$

$$\cos(v \pm w) = \cos v \cos w \mp \sin v \sin w$$

$$\sin v + \sin w = 2 \sin \frac{v + w}{2} \cos \frac{v - w}{2}$$

$$\cos v + \cos w = 2 \cos \frac{v + w}{2} \cos \frac{v - w}{2}$$

$$|\sin \frac{x}{2}| = \sqrt{\frac{1 - \cos x}{2}} \quad |\cos \frac{x}{2}| = \sqrt{\frac{1 + \cos x}{2}}$$

$$\tan(v \pm w) = \frac{\tan v \pm \tan w}{1 \mp \tan v \tan w} \quad |\tan \frac{x}{2}| = \sqrt{\frac{1 - \cos x}{1 + \cos x}}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \operatorname{atan2}(b, a)$.

Spherical coordinates



$$\begin{split} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \operatorname{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \operatorname{atan2}(y, x) \end{split}$$

Integrals

$$\int \sqrt{a^2 + x^2} dx = \frac{x}{2} \sqrt{a^2 + x^2} + \frac{a^2}{2} \ln(x + \sqrt{a^2 + x^2}) dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{|a|} dx$$

$$\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{|a|} = -\arccos \frac{x}{|a|} dx$$

$$\int \frac{dx}{\sqrt{a^2 + x^2}} = \ln(x + \sqrt{a^2 + x^2}) dx$$
Sub $s = \tan(x/2)$ to get: $dx = \frac{2 ds}{1 + s^2}$,
$$\sin x = \frac{2s}{1 + s^2}, \cos x = \frac{1 - s^2}{1 + s^2} dx$$

$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$
(Integration by parts)
$$\int \tan ax = -\frac{\ln|\cos ax|}{a} dx$$

$$\int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2} dx$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x), \quad \int xe^{ax} dx = \frac{e^{ax}}{a^2}(ax - 1) dx$$

$$\frac{d}{dx} \tan x = 1 + \tan^2 x, \quad \frac{d}{dx} \arctan x = \frac{1}{1 + x^2} dx$$
Curve length: $\int_a^b \sqrt{(X'(t))^2 + (Y'(t))^2} dx$
When $X(t), Y(t) : \int_a^b \sqrt{(X'(t))^2 + (Y'(t))^2} dx$
Solid of revolution vol: $\pi \int_a^b (f(x))^2 dx$
Surface area: $2\pi \int_a^b |f(x)| \sqrt{1 + (f'(x))^2} dx$

Markov chains

A <u>Markov chain</u> is a discrete random process with the property that the next state depends only on the current state. Let X_1, X_2, \ldots be a sequence of random variables generated by the Markov process. Then there is a transition matrix $\mathbf{P} = (p_{ij})$, with $p_{ij} = \Pr(X_n = i | X_{n-1} = j)$, and $\mathbf{p}^{(n)} = \mathbf{P}^n \mathbf{p}^{(0)}$ is the probability distribution for X_n (i.e., $p_i^{(n)} = \Pr(X_n = i)$), where $\mathbf{p}^{(0)}$ is the initial distribution. π is a stationary distribution if $\pi = \pi \mathbf{P}$. If the Markov

chain is irreducible (it is possible to get to any state from any state), then $\pi_i = \frac{1}{\mathbb{E}(T_i)}$ where $\mathbb{E}(T_i)$ is the expected time between two visits in state i. π_j/π_i is the expected number of visits in state j between two visits in state i.

For a connected, undirected and non-bipartite graph, where the transition probability is uniform among all neighbors, π_i is proportional to node i's degree.

A Markov chain is <u>ergodic</u> if the asymptotic distribution is independent of the initial distribution. A finite Markov chain is ergodic iff it is irreducible and <u>aperiodic</u> (i.e., the gcd of cycle lengths is 1). $\lim_{k \to \infty} \mathbf{P}^k = \mathbf{1}\pi$.

A Markov chain is an A-chain if the states can be partitioned into two sets \mathbf{A} and \mathbf{G} , such that all states in \mathbf{A} are absorbing $(p_{ii}=1)$, and all states in \mathbf{G} leads to an absorbing state in \mathbf{A} . The probability for absorption in state $i \in \mathbf{A}$, when the initial state is j, is $a_{ij}=p_{ij}+\sum_{k\in\mathbf{G}}a_{ik}p_{kj}$. The expected time until absorption, when the initial state is i, is $t_i=1+\sum_{k\in\mathbf{G}}p_{ki}t_k$.

Pythagorean Triples

Uniquely generated by

$$a = k \cdot (m^2 - n^2), \ b = k \cdot (2mn), \ c = k \cdot (m^2 + n^2),$$

m > n > 0, k > 0, $m \perp n$, and either m or n even.

Estimates

The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

Mobius Function

$$\mu(n) = \begin{cases} 0 & n \text{ is not square free} \\ 1 & n \text{ has even number of prime factors} \\ -1 & n \text{ has odd number of prime factors} \end{cases}$$

$$g(n) = \sum_{d \mid n} f(d) \Leftrightarrow f(n) = \sum_{d \mid n} \mu(d)g(n/d)$$

$$\sum_{d \mid n} \mu(d) = [n = 1] \text{ (very useful)}$$

$$g(n) = \sum_{n \mid d} f(d) \Leftrightarrow f(n) = \sum_{n \mid d} \mu(d/n)g(d)$$

$$g(n) = \sum_{1 < m < n} f(\lfloor \frac{n}{m} \rfloor) \Leftrightarrow f(n) = \sum \mu(m)g(\lfloor \frac{n}{m} \rfloor)$$

Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write $n = n_k p^k + ... + n_1 p + n_0$ and $m = m_k p^k + ... + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$.

CPUID Vendors: 1: intel; 2: amd Types: 1: bonnell, atom; 2: core2; 3: corei7 4: amdfam10h; 5: amdfam15h, shanghai; 6: silvermont, slm, istanbul 7: knl, bdver1; 8: bdver2; 9: btver2 10: amdfam17h; 11: knm; 12: goldmont 13: goldmont-plus; 14: tremont 15: amdfam19h; 18: grandridge 19: clearwaterforest 1: nehalem; 2: westmere; 3: sandybridge; 4: barcelona; 7: btver1; 9: bdver3; 10: bdver4; 11: znver1; 12: ivvbridge; 13: haswell; 14: broadwell; 15: skylake; 16: skylake-avx512; 17: cannonlake, sierraforest; 18: icelake-client; 19: icelake-server; 20: znver2; 21: cascadelake; 22: tigerlake; 23: cooperlake; 24: sapphirerapids, emeraldrapids; 25: alderlake, raptorlake, meteorlake, ...; 26: znver3; 27: rocketlake; 28: lujiazui; 29: znver4; 30: graniterapids; 31: graniterapids-d; 32: arrowlake; 33: arrowlake-s, lunarlake; 34: pantherlake; 35: yongfeng; 36: znver5; Check CPU features using 'man g++'. Verify: __builtin_cpu_is __builtin_cpu_supports

CPUID submit

```
#include "cpuid.h"
extern "C" struct {
 int vendor, type, subtype, features;
} __cpu_model;
int main() {
  char brand[50] = {};
  auto b = (unsigned*) brand;
  rep(i, 2, 5) {
    __get_cpuid(INT_MIN+i, b, b+1, b+2, b+3);
   b += 4;
  auto m = __cpu_model;
  cout << brand << endl << m.vendor << ' ';</pre>
  cout << m.type << ' ' << m.subtype << endl;</pre>
  // Extract CPU subtype using 4 submissions.
  int submitID = 0; // Set to 0, 1, 2, 3.
  int t = m.subtype;
  while (submitID--) t /= 3;
  if (t%3 == 2) for (volatile int c=0;;) c=c;
  return t%3;
```

CPUID recovery

```
int main() {
    // 0 = ANS, 1 = RTE, 2 = TLE
    int id = 0, status[4] = {0, 2, 1, 0};
    rep(i, 0, 4) id = id*3 + status[3-i];
    cout << id << endl;
}</pre>
```

Checklist

- .vimrc
- .bashrc
- template.cpp
- Hash verification
- Java
- Python
- Printing
- Clarifications
- Documentation
- Submit script
- Whitespace/case insensitivity
- Source code limit
- CPU on local machine
- CPU on checker
- Test Dijkstra speed
- clock()
- Judge errors

List binaries

```
echo $PATH | tr ':' ' | xargs ls | grep -v / | sort | uniq
```