Jagiellonian University - Jagiellonian 1

Jagienoman emversity Jagienoman i					
.bashrc	2	math/linear_rec.h	12	text/alcs.h	20
.vimrc	2	math/linear_rec_fast.h	12	text/kmp.h	20
template.cpp	2	math/matrix.h	12	text/kmr.h	20
template.java	2	math/miller_rabin.h	12	text/lcp.h	20
various.h	2	math/modinv_precompute.h	12	text/lyndon_factorization.h	20
various.py	2	math/modular.h	12	text/main_lorentz.h	20
geometry/convex_hull.h	2	math/modular64.h	12	text/manacher.h	21
<pre>geometry/convex_hull_dist.h</pre>	3	math/modular_generator.h	13	text/min_rotation.h	21
geometry/convex_hull_sum.h	3	math/modular_log.h	13	text/monge.h	21
geometry/halfplanes.h	3	math/modular_sqrt.h	13	text/palindromic_tree.h	21
geometry/line2.h	3	math/montgomery.h	13	text/suffix_array_linear.h	22
geometry/rmst.h	3	math/nimber.h	13	text/suffix_automaton.h	22
geometry/segment2.h	4	math/phi_large.h	13	text/suffix_tree.h	22
geometry/vec2.h	4	math/phi_precompute.h	13	text/z_function.h	23
graphs/2sat.h	4	math/phi_prefix_sum.h	13	trees/centroid_decomp.h	23
graphs/bellman_ineq.h	4	math/pi_large.h		trees/centroid_offline.h	23
graphs/biconnected.h	4	math/pi_large_precomp.h	14	trees/heavylight_decomp.h	23
graphs/bip_edge_coloring.h	5	math/pollard_rho.h		trees/lca.h	24
graphs/boski_matching.h	5	math/polynomial.h		trees/lca_linear.h	24
graphs/bridges_online.h	5	math/polynomial_interp.h		trees/link_cut_tree.h	24
graphs/dense_dfs.h	5	math/sieve.h	15	util/arc_interval_cover.h	24
graphs/directed_mst.h	5	math/sieve_factors.h	15	_	25
graphs/dominators.h	5	math/sieve_segmented.h	15	_	25
graphs/edge_coloring.h	6	math/simplex.h	15	_	25
graphs/edmonds_karp.h	6	structures/bitset_plus.h	16	util/inversion_vector.h	25
graphs/flow_with_demands.h	6	structures/fenwick_tree.h	16		25
graphs/general_matching.h	6	structures/fenwick_tree_2d.h	16		25
graphs/general_matching_w.h	6	structures/find_union.h	16		25
graphs/global_min_cut.h	7	structures/find_union_undo.h	16	_	25
graphs/gomory_hu.h	7	structures/hull_offline.h		util/radix_sort.h	25
graphs/kth_shortest.h	8	structures/hull_online.h	16		26
graphs/matroids.h	8	structures/li_chao_tree.h	17	xyz/uj.h	27
graphs/min_cost_max_flow.h	9	structures/max_queue.h	17		
graphs/push_relabel.h	9	structures/pairing_heap.h	17		
graphs/scc.h	9	structures/rmq.h	17		
graphs/turbo_matching.h	9	structures/segtree_config.h	17		
graphs/weighted_matching.h	10	structures/segtree_general.h	18		
math/berlekamp_massey.h	10	structures/segtree_persist.h	18		
math/bit_gauss.h	10	structures/segtree_point.h	18		
math/bit_matrix.h	10	structures/treap.h	19		
math/crt.h	10	structures/wavelet_tree.h	19		
math/fft_complex.h	10	structures/ext/hash_table.h	19		
math/fft_mod.h	11	structures/ext/rope.h	19		
math/fwht.h	11	structures/ext/tree.h	19		
math/gauss.h	11	structures/ext/trie.h	19		
math/gauss_ortho.h	11	text/aho_corasick.h	19		

Jagiellonian University - Jagiellonian 1

```
.bashrc
build()(
 g++ $0 -o $1.e -DLOC -std=c++14
      -Wall -W -Wfatal-errors -Wshadow \
      -Wlogical-op -Wconversion -Wfloat-equal
b() ( build $@ -02 )
d() ( build $@ -fsanitize=address, undefined \
              -D_GLIBCXX_DEBUG -q )
run()($1 $2 && echo start >&2 && time ./$2.e)
  set -e; $1 $2; $1 $3
  for ((;;)) {
    ./$3.e > gen.in
   time ./$2.e < gen.in > gen.out
cmp()(
  set -e; $1 $2; $1 $3; $1 $4
  for ((;;)) {
    ./$4.e > gen.in;
                             echo -n 0
    ./$2.e < gen.in > p1.out; echo -n 1
    ./$3.e < gen.in > p2.out; echo -n 2
   diff pl.out p2.out;
                              echo -n Y
# Other flags:
# -Wformat=2 -Wshift-overflow=2 -Wcast-qual
# -Wcast-align -Wduplicated-cond
# -D GLIBCXX DEBUG PEDANTIC -D FORTIFY SOURCE=2
# -fno-sanitize-recover -fstack-protector
# Print optimization info: -fopt-info-all
# Stacktrace on STL assert:
   q++ -D GLIBCXX DEBUG -H test.cpp 2>&1
     >/dev/null | grep "debug/macros.h"
   #ifdef ___SANITIZE_ADDRESS__
    extern "C"
      void __sanitizer_print_stack_trace();
.vimrc
se ai aw cin cul ic is nocp nohls nu sc scs
se bg=dark sw=4 ts=4 so=7 ttm=9
vn _ :w !cpp -dD -P -fpreprocessed \|
  sed -z sg\\sggg \| md5sum \| cut -c-4 <cr>
template.cpp
#include <bits/stdc++.h>
using namespace std:
using 11 = long long;
using Vi = vector<int>;
using Pii = pair<int,int>;
#define mp make_pair
#define pb push back
#define x first
#define y second
#define rep(i,b,e) for(int i=(b); i<(e); i++)
#define each (a, x) for (auto \& a : (x))
#define all(x)
                   (x).begin(),(x).end()
#define sz(x)
                   int((x).size())
int main() {
 cin.sync_with_stdio(0); cin.tie(0);
```

```
cout << fixed << setprecision(12);</pre>
  // Don't call destructors:
  cout << flush; _Exit(0);</pre>
} // dd2c
// > Debug printer
#define tem \
 template < class t, class u, class...w > auto
#define pri(x, v, z) \
 tem operator << (t&o, u a) -> decltype(z,o) \
 { o << *#x; y; z; return o << #x+1; }
pri({},, a.print())
pri((),, o << a.x << ", " << a.y)
pri([], auto d=""; for (auto i : a)
  (o \ll d \ll i, d = ", "), all(a))
void DD(...) {}
tem DD(t s, u a, w... k) {
  for (int b=1; cerr << *s++, *s && *s - b*44;)
   b += 2 / (*s*2 - 81);
  cerr << ": " << a; DD(s, k...);
} // dbb7
#ifdef LOC
#define deb(x...) \
 DD("[,\b :] "#x, __LINE__, x), cerr << endl
#define deb(...)
#endif
#define DBP(x...) void print() { DD(\#x, x); }
// > Utils
// Return smallest k such that 2^k > n
// Undefined for n = 0!
int upla(int n) { return 32- builtin clz(n); }
int uplg(ll n) { return 64-__builtin_clzll(n); }
// Compare with certain epsilon (branchless)
// Returns -1 if a < b; 1 if a > b; 0 if equal
// a and b are assumed equal if |a-b| <= eps</pre>
int cmp(double a, double b, double eps=1e-9) {
 return (a > b+eps) - (a+eps < b);
} // 81c1
template.java
                                          4cea
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.io.InputStreamReader;
import java.util.StringTokenizer;
import java.io.BufferedReader;
import java.io.InputStream;
public class Main {
 public static void main(String[] args) {
    try (PrintWriter out =
         new PrintWriter(System.out)) {
      new Task().solve(
        new InputReader(System.in), out);
   } // f642
  } // cbc2
  static class InputReader {
    BufferedReader r;
    StringTokenizer t = null;
    public InputReader(InputStream s) {
      r = new BufferedReader(
```

```
new InputStreamReader(s), 32768);
    } // Oa0f
    public String next() {
      while (t == null || !t.hasMoreTokens())
          t = new StringTokenizer(
            r.readLine());
        } catch (IOException e) {
          throw new RuntimeException (e);
        } // fdda
      return t.nextToken();
    } // d64b
    public int nextInt() {
      return Integer.parseInt(next());
    } // 7737
  1 // 94b9
  static class Task {
    public void solve (InputReader in,
                      PrintWriter out) {
      int n = in.nextInt();
      out.printf("hello %d", n);
    } // 477d
 } // 30cf
} // 75cd
various.h
// If math constants like M_PI are not found
// add this at the beginning of file
#define _USE_MATH_DEFINES
// Pragmas
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("popent, avx, tune=native")
// Clock
while (clock() < duration*CLOCKS_PER_SEC)</pre>
// Automatically implement operators:
// 1. != if == is defined
// 2. >, <= and >= if < is defined
using namespace rel_ops;
// Mersenne twister for randomization.
mt19937 64 rnd(chrono::steady clock::now()
  .time_since_epoch().count());
// To shuffle randomly use:
shuffle(all(vec), rnd);
// To pick random integer from [A;B] use:
uniform_int_distribution <> dist(A, B);
int value = dist(rnd);
// To pick random real number from [A;B] use:
uniform_real_distribution 	⇔ dist(A, B);
double value = dist(rnd);
// Floats can represent integers up to 19*10^6
// Doubles can represent integers up to 9*10^15
various.py
                                          a 9e 9
input().split(' ') # Read and split into words
print('abc', end='') # Print without newline
>>> from fractions import *
>>> Fraction (16, -10)
Fraction (-8, 5)
>>> Fraction (123)
Fraction (123, 1)
>>> Fraction()
Fraction(0, 1)
```

```
>>> Fraction('3/7')
        Fraction (3, 7)
        >>> Fraction(' -3/7 ')
       Fraction(-3, 7)
        >>> Fraction('1.414213 \t\n')
        >>> Fraction('7e-6')
       Fraction (7, 1000000)
        >>> Fraction (2.25)
       Fraction (9, 4)
        >>> Fraction (1.1)
        Fraction (2476979795053773, 2251799813685248)
        >>> Fraction('1/2') * Fraction('4/3')
        Fraction (2, 3)
        >>> Fraction (16, 5).numerator
       16
       >>> Fraction (16, 5).denominator
        >>> from decimal import *
        >>> getcontext().prec = 28
        >>> Decimal (10)
       Decimal('10')
        >>> Decimal('3.14')
       Decimal('3.14')
       >>> Decimal (3.14)
       Decimal('3.140000000000001243449787580175...')
        >>> Decimal((0, (3, 1, 4), -2))
1a4d | Decimal('3.14')
        >>> Decimal(str(2.0 ** 0.5))
       Decimal('1.4142135623730951')
       >>> Decimal(2) ** Decimal('0.5')
       Decimal('1.414213562373095048801688724')
        >>> Decimal('NaN')
       Decimal('NaN')
        >>> Decimal('-Infinity')
       Decimal ('-Infinity')
        geometry/convex_hull.h
                                                  511a
        #include "vec2.h"
        // Translate points such that lower-left point
        // is (0, 0). Returns old point location; O(n)
        vec2 normPos(vector<vec2>& points) {
         auto q = points[0].yx();
          each(p, points) q = min(q, p.yx());
         vec2 ret{q.v, q.x};
         each (p, points) p = p-ret;
         return ret:
        } // ee96
        // Find convex hull of points; time: O(n lq n)
        // Points are returned counter-clockwise,
        // first point is the lowest-left.
        vector<vec2> convexHull(vector<vec2> points) {
         vec2 pivot = normPos(points);
         sort(all(points));
         vector<vec2> hull:
         each (p, points) {
            while (sz(hull) >= 2) {
              vec2 a = hull.back() - hull[sz(hull)-2];
              vec2 b = p - hull.back();
              if (a.cross(b) > 0) break;
             hull.pop_back();
            } // ad91
           hull.pb(p);
         } // 5908
         // Translate back, optional
         each (p, hull) p = p+pivot;
         return hull;
```

```
} // 62ed
// Find point p that minimizes dot product p*q.
// Returns point index in hull; time: O(lg n)
// If multiple points have same dot product
// one with smallest index is returned.
// Points are expected to be in the same order
// as output from convexHull function.
int minDot(const vector<vec2>& hull, vec2 q) {
  auto C = [](vec2 a, vec2 b) {
    return mp(a.dot(b), a.cross(b));
  }; // 27c4
  auto search = [&](int b, int e, vec2 p) {
    int f = b, q = e;
    while (b+1 < e) {
     int m = (b+e) / 2;
      (C(p, hull[m-1]) > C(p, hull[m])
        ? b : e) = m;
    } // 618b
    rep(i, 0, min(g-f, 2)) {
      if (C(p, hull[f+i]) < C(p, hull[b]))
      if (C(p, hull[g-i-1]) < C(p, hull[b]))
       b = a-i-1;
    } // 9413
    return b;
  }; // e993
  int m = search(0, sz(hull), \{0, -1\});
  int i = search(0, m, q);
  int j = search(m, sz(hull), q);
  return C(q, hull[i]) > C(q, hull[j])
    ? j : i;
} // 2fba
geometry/convex hull dist.h 2859
#include "vec2.h"
// Check if p is inside convex polygon. Hull
// must be given in counter-clockwise order.
// Returns 2 if inside, 1 if on border,
// 0 if outside; time: O(n)
int insideHull(vector<vec2>& hull, vec2 p) {
  int ret = 1;
  rep(i, 0, sz(hull)) {
    auto v = hull[(i+1)%sz(hull)] - hull[i];
    auto t = v.cross(p-hull[i]);
    ret = min(ret, cmp(t, 0)); // For doubles
    //ret = min(ret, (t>0) - (t<0)); // Ints
  return int(max(ret+1, 0));
} // 1f39
#include "segment2.h"
// Get distance from point to hull; time: O(n)
double hullDist(vector<vec2>& hull, vec2 p) {
  if (insideHull(hull, p)) return 0;
  double ret = 1e30;
  rep(i, 0, sz(hull)) {
    seg2 seg{hull[(i+1)%sz(hull)], hull[i]};
    ret = min(ret, seg.distTo(p));
  } // f3be
  return ret;
} // a00c
// Compare distance from point to hull
// with sqrt(d2); time: O(n)
// -1 if smaller, 0 if equal, 1 if greater
int cmpHullDist(vector<vec2>& hull,
                vec2 p, 11 d2) {
 if (insideHull(hull,p)) return (d2<0)-(d2>0);
```

```
int ret = 1;
 rep(i, 0, sz(hull)) {
   seg2 seg{hull[(i+1)%sz(hull)], hull[i]};
   ret = min(ret, seq.cmpDistTo(p, d2));
 } // 28cb
 return ret:
} // 30f3
geometry/convex hull sum.h
#include "vec2.h"
// Get edge sequence for given polygon
// starting from lower-left vertex; time: O(n)
// Returns start position.
vec2 edgeSeg(vector<vec2> points.
             vector<vec2>& edges) {
 int i = 0, n = sz(points);
  rep(j, 0, n)
   if (points[i].yx() > points[j].yx()) i = j;
  rep(j, 0, n) edges.pb(points[(i+j+1)%n] -
                        points[(i+j)%n]);
 return points[i];
1 // 3aa7
// Minkowski sum of given convex polygons.
// Vertices are required to be in
// counter-clockwise order; time: O(n+m)
vector<vec2> hullSum(vector<vec2> A.
                    vector<vec2> B) {
 vector\langle vec2 \rangle sum, e1, e2, es(sz(A) + sz(B));
 vec2 pivot = edgeSeg(A, e1) + edgeSeg(B, e2);
 merge(all(e1), all(e2), es.begin());
  sum.pb(pivot);
 each (e, es) sum.pb(sum.back() + e);
  sum.pop_back();
 return sum;
} // f183
geometry/halfplanes.h
                                          356a
#include "vec2.h"
#include "line2.h"
// Intersect halfplanes given by `lines`
// and output hull vertices to 'out'
// in counter-clockwise order. Returns true
// if intersection is non-empty and bounded.
// Unbounded cases are not supported,
// add bounding-box if necessary. Works only
// with floating point vec2/line2; O(n lg n)
// PARTIALLY TESTED
bool intersectHalfplanes(vector<line2> in.
                         vector<vec2>€ out) {
  sort(all(in), [](line2 a, line2 b) {
   return (a.v.angleCmp(b.v) ?:
            a.c*b.v.len() - b.c*a.v.len()) < 0;
  1): // 82fb
  int a = 0, b = 0, n = sz(in);
 vector<line2> dq(n+5);
 out.resize(n+5);
 dq[0] = in[0];
  rep(i, 1, n+1) {
   if (i == n) in.pb(dq[a]);
    if (!in[i].v.angleCmp(in[i-1].v)) continue;
    while (a < b && in[i].side(out[b-1]) > 0)
     b--;
    while (i!=n && a < b && in[i].side(out[a])>0)
   if (in[i].intersect(dq[b], out[b]))
```

```
dq[++b] = in[i];
         } // b9ba
         out.resize(b);
         out.erase(out.begin(), out.begin()+a);
         return b-a > 2;
       } // f334
7f53 | geometry/line2.h
                                                 9207
       #include "vec2.h"
       // 2D line/halfplane structure
       // PARTIALLY TESTED
       // Base class of versions for ints and doubles
       template < class T, class P, class S>
       struct bline2 {
        // For lines: v * point == c
         // For halfplanes: v * point <= c
         // (i.e. normal vector points outside)
         P v; // Normal vector [A; B]
         T c; // Offset (C parameter of equation)
         DBP (v, c);
         // Line through 2 points; normal vector
         // points to the right of ab vector
         static S through(P a, P b) {
           return { (a-b).perp(), a.cross(b) };
         } // 9ac7
         // Parallel line through point
         static S parallel (P a, S b) {
           return { b.v, b.v.dot(a) };
         } // 8e1c
         // Perpendicular line through point
         static S perp (P a, S b) {
           return { b.v.perp(), b.v.cross(a) };
         } // 7b75
         // Distance from point to line
         double distTo(P a) {
           return fabs(v.dot(a)-c) / v.len();
         } // 79e6
       }; // ee4f
       // Version for integer coordinates (long long)
       struct line2i : bline2<11, vec2i, line2i> {
         line2i() : bline2{{}, 0} {}
         line2i(vec2i a, ll b) : bline2{a, b} {}
         // Returns 0 if point a lies on the line,
         // 1 if on side where normal vector points,
         // -1 if on the other side.
         int side(vec2i a) {
           11 d = v.dot(a):
           return (d > c) - (d < c);
         } // 18a7
       }; // fc9c
        // Version for double coordinates
       // Requires cmp() from template
       struct line2d : bline2<double, vec2d, line2d> {
         line2d() : bline2{{}, 0} {}
         line2d(vec2d a, double b) : bline2{a, b} {}
         // Returns 0 if point a lies on the line,
         // 1 if on side where normal vector points.
         // -1 if on the other side.
         int side(vec2d a) { return cmp(v.dot(a),c); }
         // Intersect this line with line a, returns
         // true on success (false if parallel).
         // Intersection point is saved to `out`.
```

```
3
  bool intersect (line2d a, vec2d& out) {
    double d = v.cross(a.v);
    if (!cmp(d, 0)) return 0;
    out = (v*a.c - a.v*c).perp() / d;
    return 1;
 } // 2e68
}; // ab54
using line2 = line2d;
geometry/rmst.h
                                          476a
#include "../structures/find_union.h"
// Rectilinear Minimum Spanning Tree
// (MST in Manhattan metric); time: O(n lg n)
// Returns MST weight. Outputs spanning tree
// to G, vertex indices match point indices.
// Edge in G is pair (target, weight).
ll rmst (vector<Pii>€ points,
        vector<vector<Pii>> G) {
  int n = sz(points);
  vector<pair<int, Pii>> edges;
  vector<Pii> close;
  Vi ord(n), merged(n);
  iota(all(ord), 0);
  function<void(int,int)> octant =
      [&] (int begin, int end) {
    if (begin+1 >= end) return;
    int mid = (begin+end) / 2;
    octant (begin, mid);
    octant (mid, end);
    int j = mid;
    Pii best = \{INT_MAX, -1\};
    merged.clear();
    rep(i, begin, mid) {
      int v = ord[i];
      Pii p = points[v];
      while (j < end) {
        int e = ord[j];
        Pii q = points[e];
        if (q.x-q.y > p.x-p.y) break;
        best = min(best, make_pair(q.x+q.y, e));
        merged.pb(e);
        j++;
      } // 8576
      if (best.v != -1) {
        int alt = best.x-p.x-p.y;
        if (alt < close[v].x)</pre>
          close[v] = {alt, best.y};
      } // 4208
      merged.pb(v);
    } // f3ff
    while (j < end) merged.pb(ord[j++]);</pre>
    copy(all(merged), ord.begin()+begin);
  }; // a4e1
  rep(i, 0, 4) {
    rep(j, 0, 2) {
      sort(all(ord), [&](int 1, int r) {
        return points[1] < points[r];</pre>
      }); // fe33
      close.assign(n, {INT_MAX, -1});
      octant(0, n);
      rep(k, 0, n) {
       Pii p = close[k];
        if (p.y != -1) edges.pb({p.x,{k,p.y}});
```

```
points[k].x \star = -1;
     } // 1c1d
   } // 9b38
    each (p, points) p = \{p.y, -p.x\};
  } // d06f
  11 sum = 0;
  FAU fau(n);
  sort (all (edges));
  G.assign(n, {});
  each(e, edges) if (fau.join(e.y.x, e.y.y)) {
   sum += e.x;
   G[e.y.x].pb({e.y.y, e.x});
   G[e.y.y].pb({e.y.x, e.x});
  } // b04a
  return sum;
} // f586
                                          6504
geometry/segment2.h
#include "vec2.h"
// 2D segment structure; PARTIALLY TESTED
// Base class of versions for ints and doubles
template < class P, class S> struct bseq2 {
  P a, b; // Endpoints
  // Distance from segment to point
  double distTo(P p) const {
    if ((p-a).dot(b-a) < 0) return (p-a).len();</pre>
    if ((p-b).dot(a-b) < 0) return (p-b).len();</pre>
    return double(abs((p-a).cross(b-a)))
                  / (b-a).len();
  } // 62a2
}; // 85bc
// Version for integer coordinates (long long)
struct seq2i : bseq2<vec2i, seq2i> {
  seq2i() {}
  seq2i(vec2i c, vec2i d) : bseq2{c, d} {}
  // Check if segment contains point p
  bool contains(vec2i p) {
    return (a-p).dot(b-p) <= 0 &&
           (a-p).cross(b-p) == 0;
  // Compare distance to p with sqrt(d2)
  // -1 if smaller, 0 if equal, 1 if greater
  int cmpDistTo(vec2i p, 11 d2) const {
    if ((p-a).dot(b-a) < 0) {
     11 1 = (p-a).len2();
      return (1 > d2) - (1 < d2);
    } // dla6
    if ((p-b).dot(a-b) < 0) {
     11 1 = (p-b).len2();
     return (1 > d2) - (1 < d2);
    1 // 9e65
    11 c = abs((p-a).cross(b-a));
    d2 = (b-a).len2();
    return (c*c > d2) - (c*c < d2);
  } // 726d
}; // 4df2
// Version for double coordinates
// Requires cmp() from template
struct seg2d : bseg2<vec2d, seg2d> {
  seg2d() {}
  seg2d(vec2d c, vec2d d) : bseg2{c, d} {}
  bool contains (vec2d p) {
    return cmp((a-p).dot(b-p), 0) <= 0 &&
```

```
cmp((a-p).cross(b-p), 0) == 0;
 } // b507
}; // 2036
using seg2 = seg2d;
geometry/vec2.h
                                          6e47
// 2D point/vector structure; PARTIALLY TESTED
// Base class of versions for ints and doubles
template < class T, class S> struct bvec2 {
 Тх, у;
  S operator+(S r) const {return{x+r.x,y+r.y};}
  S operator-(S r) const {return{x-r.x,y-r.y};}
  S operator*(T r) const { return {x*r, y*r}; }
  S operator/(T r) const { return {x/r, y/r}; }
 T dot(S r) const { return x*r.x + y*r.y; }
 T cross(S r) const { return x*r.y - y*r.x; }
 T len2()
               const { return x*x + y*y; }
  double len() const { return hypot(x, y); }
  S perp()
               const { return {-y,x}; } // CCW
  pair<T, T> yx() const { return {y, x}; }
  double angle() const { //[0;2*PI] CCW from OX
    double a = atan2(y, x);
    return (a < 0 ? a+2*M_PI : a);</pre>
  } // 7095
}; // 17ed
// Version for integer coordinates (long long)
struct vec2i : bvec2<11, vec2i> {
 vec2i() : bvec2{0, 0} {}
  vec2i(11 a, 11 b) : bvec2{a, b} {}
  bool upper() const { return (y ?: x) >= 0; }
  int angleCmp(vec2i r) const {
    11 c = cross(r);
    return r.upper() -upper() ?: (c<0) - (c>0);
  } // b35f
  // Compare by angle, length if angles equal
  bool operator<(vec2i r) const {</pre>
    return (angleCmp(r) ?:
            len2() - r.len2()) < 0;
  bool operator == (vec2i r) const {
   return x == r.x && y == r.y;
  1 // 136e
}; // d3f4
// Version for double coordinates
// Requires cmp() from template
struct vec2d : bvec2<double, vec2d> {
 vec2d() : bvec2\{0, 0\} \{\}
  vec2d(double a, double b) : bvec2{a, b} {}
  bool upper() const {
   return (cmp(y, 0) ?: cmp(x, 0)) >= 0;
  int angleCmp(vec2d r) const {
    return r.upper() - upper() ?:
           cmp(0, cross(r));
  } // 12f3
  // Compare by angle, length if angles equal
  bool operator<(vec2d r) const {</pre>
    return (angleCmp(r) ?:
            cmp(len2(), r.len2())) < 0;
  bool operator==(vec2d r) const {
```

```
return !cmp(x, r.x) && !cmp(y, r.y);
  } // 81cd
  vec2d unit() const { return *this / len(); }
  vec2d rotate(double a) const { // CCW
   return {x*cos(a) - y*sin(a),
            x*sin(a) + y*cos(a); // 1890
 } // 97e3
}; // 08e9
using vec2 = vec2d;
graphs/2sat.h
                                         2443
// 2-SAT solver; time: O(n+m), space: O(n+m)
// Variables are indexed from 1 and
// negative indices represent negations!
// Usage: SAT2 sat(variable_count);
// (add constraints...)
// bool solution_found = sat.solve();
// sat[i] = value of i-th variable, 0 or 1
            (also indexed from 1!)
// (internally: positive = i*2-1, neg. = i*2-2)
struct SAT2 : Vi {
 vector<Vi> G:
 Vi order, flags;
  // Init n variables, you can add more later
  SAT2 (int n = 0) : G(n*2) {}
  // Add new var and return its index
  int addVar() {
   G.resize(sz(G)+2); return sz(G)/2;
  } // 98f3
  // Add (i => j) constraint
  void imply(int i, int j) {
   i = i*2^ i >> 31;
    j = j*2 ^ j >> 31;
   G[--i].pb(--j); G[j^1].pb(i^1);
  1 // 8e25
  // Add (i v j) constraint
  void either(int i, int j) { imply(-i, j); }
  // Constraint at most one true variable
  void atMostOne(Vi& vars) {
    int y, x = addVar();
    each(i, vars) {
      imply(x, y = addVar());
      imply(i, -x); imply(i, x = y);
   } // 24aa
  } // 3ed7
  // Solve and save assignments in `values`
  bool solve() { // O(n+m), Kosaraju is used
    assign(sz(G)/2+1, -1);
    flags.assign(sz(G), 0);
    rep(i, 0, sz(G)) dfs(i);
    while (!order.empty()) {
      if (!propag(order.back()^1, 1)) return 0;
      order.pop_back();
    1 // 5594
    return 1;
 } // 1e58
  void dfs(int i) {
   if (flags[i]) return;
    flags[i] = 1;
    each(e, G[i]) dfs(e);
    order.pb(i);
  } // d076
  bool propag(int i, bool first) {
```

```
if (!flags[i]) return 1;
   flags[i] = 0;
   if (at(i/2+1) >= 0) return first;
   at (i/2+1) = i&1;
   each(e, G[i]) if (!propag(e, 0)) return 0;
   return 1:
 } // 4c1b
}; // d74c
graphs/bellman ineq.h
                                         cd51
struct Ineq {
 11 a, b, c; // a - b >= c
}; // 663a
// Solve system of inequalities of form a-b>=c
// using Bellman-Ford; time: O(n*m)
bool solveIneg(vector<Ineg>& edges,
              vector<ll>& vars) {
 rep(i, 0, sz(vars)) each(e, edges)
   vars[e.b] = min(vars[e.b], vars[e.a]-e.c);
 each (e, edges)
   if (vars[e.a]-e.c < vars[e.b]) return 0;</pre>
 return 1:
} // 241e
graphs/biconnected.h
                                         2b9e
// Biconnected components; time: O(n+m)
// Usage: Biconnected bi(graph);
// bi[v] = indices of components containing v
// bi.verts[i] = vertices of i-th component
// bi.edges[i] = edges of i-th component
// Bridges <=> components with 2 vertices
// Articulation points <=> vertices that belong
                           to > 1 component
// Isolated vertex <=> empty component list
struct Biconnected : vector<Vi> {
 vector<Vi> verts;
 vector<vector<Pii>>> edges;
 vector<Pii> S;
 Biconnected() {}
 Biconnected(vector<Vi>& G) : S(sz(G)) {
   resize(sz(G));
   rep(i, 0, sz(G)) S[i].x ?: dfs(G, i, -1);
   rep(c, 0, sz(verts)) each(v, verts[c])
     at (v) .pb(c);
 } // cfce
 int dfs(vector<Vi>& G, int v, int p) {
   int low = S[v].x = sz(S)-1;
   S.pb(\{v, -1\});
   each(e, G[v]) if (e != p) {
     if (S[e].x < S[v].x) S.pb({v, e});</pre>
      low = min(low, S[e].x ?: dfs(G, e, v));
   } // 446d
   if (p+1 && low >= S[p].x) {
     verts.pb({p}); edges.pb({});
     rep(i, S[v].x, sz(S)) {
       if (S[i].y == -1)
         verts.back().pb(S[i].x);
         edges.back().pb(S[i]);
     1 // 4fab
     S.resize(S[v].x);
   } // 6d66
   return low:
 } // 7fcc
```

}; // 4fa4

647c

graphs/bip edge coloring.h

```
// Bipartite edge coloring; time: O(nm)
// `edges` is list of (left vert, right vert),
// where vertices on both sides are indexed
// from 0 to n-1. Returns number of used colors
// (which is equal to max degree).
// col[i] = color of i-th edge [0..max_deg-1]
int colorEdges(vector<Pii>€ edges,
               int n, Vi& col) {
  int m = sz(edges), c[2] = {}, ans = 0;
 Vi deg[2];
  vector<vector<Pii>>> has[2];
  col.assign(m, 0);
  rep(i, 0, 2) {
   deg[i].resize(n+1);
   has[i].resize(n+1, vector<Pii>(n+1));
  function<void(int,int)> dfs =
    [&] (int x, int p) {
     Pii i = has[p][x][c[!p]];
     if (has[!p][i.x][c[p]].y) dfs(i.x, !p);
     else has[!p][i.x][c[!p]] = {};
     has[p][x][c[p]] = i;
     has[!p][i.x][c[p]] = \{x, i.y\};
     if (i.y) col[i.y-1] = c[p]-1;
   }; // 19ad
  rep(i, 0, m) {
    int x[2] = \{edges[i].x+1, edges[i].y+1\};
    rep(d, 0, 2) {
     deg[d][x[d]]++;
     ans = max(ans, deg[d][x[d]]);
     for (c[d] = 1; has[d][x[d]][c[d]].v;)
       c[d]++;
   } // 9454
    if (c[0]-c[1]) dfs(x[1], 1);
    rep(d, 0, 2)
     has[d][x[d]][c[0]] = \{x[!d], i+1\};
   col[i] = c[0]-1;
  } // 5678
 return ans;
} // 1bfc
graphs/boski matching.h
                                         dbd2
// Bosek's algorithm for partially online
// bipartite maximum matching - white vertices
// are fixed, black vertices are added
// one by one; time: O(E*sqrt(V))
// Usage: Matching match (num white);
// match[v] = index of black vertex matched to
             white vertex v or -1 if unmatched
// match.add(indices_of_white_neighbours);
// Black vertices are indexed in order they
// were added, the first black vertex is 0.
struct Matching : Vi {
  vector<Vi> adj;
 Vi rank, low, pos, vis, seen;
  int k{0}:
  // Initialize structure for n white vertices
  Matching (int n = 0) : Vi(n, -1), rank(n) {}
  // Add new black vertex with its neighbours
  // given by `vec`. Returns true if maximum
  // matching is increased by 1.
  bool add (Vi vec) {
   adi.pb (move (vec));
    low.pb(0); pos.pb(0); vis.pb(0);
   if (!adj.back().empty()) {
```

```
int i = k;
    nxt:
      seen.clear();
      if (dfs(sz(adj)-1, ++k-i)) return 1;
      each (v, seen) each (e, adj[v])
        if (rank[e] < 1e9 && vis[at(e)] < k)
          goto nxt;
      each (v, seen) each (w, adj[v])
        rank[w] = low[v] = 1e9;
    } // 6aec
    return 0;
  } // d2a7
  bool dfs(int v, int q) {
    if (vis[v] < k) vis[v] = k, seen.pb(v);
    while (low[v] < q) {</pre>
      int e = adj[v][pos[v]];
      if (at(e) != v && low[v] == rank[e]) {
        rank[e]++:
        if (at (e) == -1 || dfs(at (e), rank[e]))
          return at (e) = v, 1;
      } else if (++pos[v] == sz(adj[v])) {
        pos[v] = 0; low[v]++;
      } // e532
    } // 3d88
   return 0;
 } // 8561
}; // aac1
graphs/bridges_online.h
                                          4124
// Dynamic 2-edge connectivity queries
// Usage: Bridges bridges(vertex_count);
// - bridges.addEdge(u, v); - add edge (u, v)
// - bridges.cc[v] = connected component ID
// - bridges.bi(v) = 2-edge connected comp ID
struct Bridges {
 vector<Vi> G; // Spanning forest
  Vi cc, size, par, bp, seen;
  int cnt{0};
  // Initialize structure for n vertices; O(n)
  Bridges (int n = 0) : G(n), cc(n), size(n, 1),
                       par(n, -1), bp(n, -1),
                       seen(n) {
    iota(all(cc), 0);
  } // ed70
  // Add edge (u, v); time: amortized O(lg n)
  void addEdge(int u, int v) {
   if (cc[u] == cc[v]) {
      int r = lca(u, v);
      for (int x : \{u, v\})
        while ((x = root(x)) != r)
          x = bp[bi(x)] = par[x];
    } else {
      G[u].pb(v); G[v].pb(u);
      if (size[cc[u]] > size[cc[v]]) swap(u,v);
      size[cc[v]] += size[cc[u]];
      dfs(u, v);
   } // abc7
 } // a6fd
  // Get 2-edge connected component ID
  int bi(int v) { // amortized time: < O(lq n)</pre>
   return bp[v] + 1 ? bp[v] = bi(bp[v]) : v;
 } // 3206
  int root(int v) {
    return par[v] == -1 || bi(par[v]) != bi(v)
      ? v : par[v] = root(par[v]);
  } // 2d27
```

```
void dfs(int v, int p) {
   cc[v] = cc[par[v] = p];
   each(e, G[v]) if (e != p) dfs(e, v);
 } // 85f5
 int lca(int u, int v) { // Don't use this!
   for (cnt++;; swap(u, v)) if (u != -1) {
     if (seen[u = root(u)] == cnt) return u;
     seen[u] = cnt; u = par[u];
   } // afed
 } // 7f56
}; // bd70
                                         4fbd
graphs/dense dfs.h
#include "../math/bit matrix.h"
// DFS over bit-packed adjacency matrix
// G = NxN adjacency matrix of graph
// G(i,j) <=> (i,j) is edge
// V = 1xN matrix containing unvisited vertices
// V(0,i) <=> i-th vertex is not visited
// Total DFS time: O(n^2/64)
struct DenseDFS {
 BitMatrix G, V; // space: O(n^2/64)
 // Initialize structure for n vertices
 DenseDFS (int n = 0) : G(n, n), V(1, n) {
   reset():
 } // 79e4
  // Mark all vertices as unvisited
 void reset() { each(x, V.M) x = -1; }
 // Get/set visited flag for i-th vertex
 void setVisited(int i) { V.set(0, i, 0); }
 bool isVisited(int i) { return !V(0, i); }
 // DFS step: func is called on each unvisited
 // neighbour of i. You need to manually call
 // setVisited(child) to mark it visited
 // or this function will call the callback
 // with the same vertex again.
 template<class T>
 void step(int i, T func) {
   ull* E = G.row(i);
   for (int w = 0; w < G.stride;) {</pre>
     ull x = E[w] & V.row(0)[w];
     if (x) func((w<<6) | __builtin_ctzll(x));</pre>
     else w++;
   } // 4c0a
 } // f045
}; // 8edb
                                         f935
graphs/directed mst.h
#include "../structures/find union undo.h"
struct Edge {
 int a, b;
 11 w:
}; // 88dd
struct Node {
 Edge key;
 Node *1, *r;
 11 delta;
 void prop() {
   kev.w += delta:
   if (1) 1->delta += delta;
   if (r) r->delta += delta;
   delta = 0;
 } // 3f0f
```

```
Edge top() { prop(); return key; }
}; // 019f
Node* merge (Node* a, Node* b) {
 if (!a || !b) return a ?: b;
  a->prop(); b->prop();
  if (a->key.w > b->key.w) swap(a, b);
  swap(a->1, (a->r = merge(b, a->r)));
  return a;
} // 34a3
void pop(Node*& a) {
 a\rightarrow prop(); a = merge(a\rightarrow 1, a\rightarrow r);
} // 666b
// Find directed minimum spanning tree
// rooted at vertex `root`; O(m log n)
// Returns weight of found spanning tree.
// out[i] = parent of i-th vertex in the tree,
// out[root] = -1
ll dmst (vector<Edge>& edges,
        int n, int root, Vi& out) {
  RollbackFAU uf (n):
  vector<Node*> heap(n);
  each (e, edges)
    heap[e.b] = merge(heap[e.b],
      new Node{e, 0, 0, 0});
  11 \text{ res} = 0;
  Vi seen(n, -1), path(n);
  seen[root] = root;
  struct Cvcle { int u, t; vector<Edge> e; };
  vector\langle Edge \rangle Q(n), in(n, \{-1, -1, 0\}), comp;
  vector<Cycle> cycs;
  rep(s, 0, n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      if (!heap[u]) return out.clear(), -1;
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w; pop(heap[u]);
      O[qi] = e; path[qi++] = u; seen[u] = s;
      res += e.w; u = uf.find(e.a);
      if (seen[u] == s) {
        Node* cvc = 0;
        int end = qi, time = uf.time();
        do cyc = merge(cyc, heap[w=path[--qi]]);
        while (uf.join(u, w));
        heap[u = uf.find(u)] = cyc;
        seen[u] = -1;
        cycs.pb({u, time, {&Q[qi], &Q[end]}});
      } // 4ff6
    } // 05d8
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
 } // fc95
  reverse (all (cycs));
  each(c, cycs) {
    uf.rollback(c.t);
    Edge tmp = in[c.u];
    each (e, c.e) in [uf.find(e.b)] = e;
    in[uf.find(tmp.b)] = tmp;
  } // 957d
  out.resize(n);
  rep(i, 0, n) out[i] = in[i].a;
 return res;
} // b30d
graphs/dominators.h
                                           aec6
// Tarjan's algorithm for finding dominators
```

```
// in directed graph; time: O(m log n)
                                                       swap(adj[at][cd], adj[end=at][cd^c^d]);
// Returns array of immediate dominators idom.
                                                     while (adj[fan[i]][d] + 1) {
                                                       int x = fan[i], y = fan[++i], f = cc[i];
// idom[root] = root
// idom[v] = -1 if v is unreachable from root
                                                       adj[u][f] = x; adj[x][f] = u;
Vi dominators (const vector < Vi>& G, int root) {
                                                       adi[v][f] = -1; fre[v] = f;
  int n = sz(G):
  vector<Vi> in(n), bucket(n);
                                                     adj[u][d] = fan[i];
  Vi pre(n, -1), anc(n, -1), par(n), best(n);
                                                     adj[fan[i]][d] = u;
  Vi ord, idom(n, -1), sdom(n, n), rdom(n);
                                                     for (int y : {fan[0], u, end})
                                                       for (int& z = fre[y] = 0; adj[y][z]+1;)
  function<void(int,int)> dfs =
                                                         z++;
   [&] (int v, int p) {
                                                   } // 4240
     if (pre[v] == -1) {
                                                   rep(i, 0, sz(edges))
       par[v] = p;
                                                     for (tie(u, v) = edges[i];
       pre[v] = sz(ord);
                                                          adj[u][ret[i]] != v;) ++ret[i];
       ord.pb(v);
                                                   return ret;
       each(e, G[v]) in[e].pb(v), dfs(e, v);
                                                 } // c259
     } // 1182
   }; // ffd2
                                                 graphs/edmonds karp.h
  function<Pii(int)> find = [&](int v) {
                                                 using flow t = int;
   if (anc[v] == -1) return mp(best[v], v);
                                                 constexpr flow t INF = 1e9+10;
    int b; tie(b, anc[v]) = find(anc[v]);
                                                 // Edmonds-Karp algorithm for finding
   if (sdom[b] < sdom[best[v]]) best[v] = b;</pre>
                                                 // maximum flow in graph; time: O(V*E^2)
    return mp(best[v], anc[v]);
                                                 struct MaxFlow {
  1: // c07b
                                                   struct Edge {
  rdom[root] = idom[root] = root;
                                                     int dst, inv;
  iota(all(best), 0);
                                                     flow t flow, cap;
  dfs(root, -1);
                                                   }; // a53c
  rep(i, 0, sz(ord)) {
                                                   vector<vector<Edge>> G;
   int v = ord[sz(ord)-i-1], b = pre[v];
                                                   vector<flow t> add;
    each(e, in[v])
                                                   Vi prev;
     b = min(b, pre[e] < pre[v] ? pre[e] :</pre>
                                                   // Initialize for n vertices
                 sdom[find(e).x]);
                                                   MaxFlow(int n = 0) : G(n) {}
    each(u, bucket[v]) rdom[u] = find(u).x;
   sdom[v] = b:
                                                    // Add new vertex
    anc[v] = par[v];
                                                   int addVert() { G.pb({}); return sz(G)-1; }
   bucket[ord[sdom[v]]].pb(v);
                                                   // Add edge from u to v with capacity cap
  ) // 54f4
                                                   // and reverse capacity rcap.
  each (v, ord) idom [v] = (rdom[v] == v ?
                                                   // Returns edge index in adjacency list of u.
   ord[sdom[v]] : idom[rdom[v]]);
                                                   int addEdge(int u, int v,
 return idom:
                                                               flow_t cap, flow_t rcap = 0) {
} // 0656
                                                     G[u].pb({ v, sz(G[v]), 0, cap });
                                                     G[v].pb({u, sz(G[u])-1, 0, rcap});
graphs/edge_coloring.h
                                          439b
                                                     return sz(G[u])-1;
// General graph edge coloring; time: O(nm)
                                                   } // c96a
// Finds (D+1)-edge-coloring of given graph,
                                                    // Compute maximum flow from src to dst.
// where D is max vertex degree.
                                                   flow_t maxFlow(int src, int dst) {
// Returns vector of edge colors `col`.
                                                     flow t i, m, f = 0;
// col[i] = color of i-th edge [0..D]
                                                     each (v, G) each (e, v) e.flow = 0;
Vi vizing(vector<Pii>€ edges, int n) {
 Vi cc(n+1), ret(sz(edges)),
     fan(n), fre(n), loc:
                                                     queue<int> 0:
  each(e, edges) cc[e.x]++, cc[e.y]++;
                                                     O.push (src);
  int u, v, cnt = *max element(all(cc)) + 1;
                                                     prev.assign(sz(G), -1);
  vector<Vi> adj(n, Vi(cnt, -1));
                                                     add.assign(sz(G), -1);
  each (e, edges) {
                                                     add[src] = INF;
   tie(u, v) = e;
                                                     while (!O.empty()) {
    fan[0] = v;
                                                       m = add[i = Q.front()];
   loc.assign(cnt, 0);
                                                       Q.pop();
   int at = u, end = u, d, c = fre[u],
                                                       if (i == dst) {
       ind = 0, i = 0;
                                                         while (i != src) {
    while (d = fre[v].
           !loc[d] && (v = adj[u][d]) != -1)
                                                           auto& e = G[i][prev[i]];
     loc[d] = ++ind, cc[ind] = d, fan[ind] = v;
                                                           e.flow -= m;
    cc[loc[d]] = c:
                                                           G[i = e.dst][e.inv].flow += m;
    for (int cd = d; at+1; cd ^= c ^ d,
                                                         } // 1f86
        at = adj[at][cd])
                                                         f += m;
```

```
goto nxt;
     } // 43a2
     each(e, G[i])
       if (add[e.dst] < 0 && e.flow < e.cap) {</pre>
         Q.push(e.dst);
         prev[e.dst] = e.inv;
         add[e.dst] = min(m, e.cap-e.flow);
       } // 4cdb
   } // 887e
   return f;
 } // cec0
 // Get flow through e-th edge of vertex v
 flow_t getFlow(int v, int e) {
   return G[v][e].flow;
 } // Ofaf
 // Get if v belongs to cut component with src
 bool cutSide(int v) { return add[v] >= 0; }
1; // c5ef
graphs/flow with demands.h
                                         0153
#include "edmonds karp.h"
//#include "push_relabel.h" // if you need
// Flow with demands; time: O(maxflow)
struct FlowDemands {
 MaxFlow net;
 vector<vector<flow t>> demands;
 flow_t total{0};
  // Initialize for k vertices
 FlowDemands(int k = 0): net(2) {
   while (k--) addVert();
 } // 7bdf
  // Add new vertex
 int addVert() {
   int v = net.addVert();
   demands.pb({});
   net.addEdge(0, v, 0);
   net.addEdge(v, 1, 0);
   return v-2;
 1 // 48b6
 // Add edge from u to v with demand dem
 // and capacity cap (dem <= flow <= cap).
  // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
              flow t dem, flow t cap) {
   demands[u].pb(dem);
   demands[v].pb(0);
   total += dem;
   net.G[0][v].cap += dem;
   net.G[u+2][1].cap += dem;
   return net.addEdge(u+2, v+2, cap-dem) - 2;
 1 // a403
 // Check if there exists a flow with value f
 // for source src and destination dst.
 // For circulation, you can set args to 0.
 bool canFlow(int src, int dst, flow_t f) {
   net.addEdge(dst += 2, src += 2, f);
   f = net.maxFlow(0, 1);
   net.G[src].pop_back();
   net.G[dst].pop_back();
   return f == total;
 } // 6285
 // Get flow through e-th edge of vertex v
 flow_t getFlow(int v, int e) {
```

8326

```
return net.getFlow(v+2,e+2)+demands[v][e];
 } // 6cf6
}; // db37
graphs/general_matching.h
                                         f42c
// Edmond's Blossom algorithm for maximum
// matching in general graphs; time: O(nm)
// Returns matching size (edge count).
// match[v] = vert matched to v or -1
int blossom(vector<Vi>& G, Vi& match) {
 int n = sz(G), cnt = -1, ans = 0;
 match.assign(n, -1);
 Vi lab(n), par(n), orig(n), aux(n, -1), q;
  auto blos = [&](int v, int w, int a) {
    while (orig[v] != a) {
      par[v] = w; w = match[v];
      if (lab[w] == 1) lab[w] = 0, q.pb(w);
      orig[v] = orig[w] = a; v = par[w];
    1 // 319e
 }; // ab9e
 rep(i, 0, n) if (match[i] == -1)
    each(e, G[i]) if (match[e] == -1) {
      match[match[e] = i] = e; ans++; break;
    } // a22a
  rep(root, 0, n) if (match[root] == -1) {
    fill(all(lab), -1);
   iota(all(orig), 0);
    lab[root] = 0;
   q = \{root\};
    rep(i, 0, sz(q)) {
      int v = q[i];
      each(x, G[v]) if (lab[x] == -1) {
       lab[x] = 1; par[x] = v;
       if (match[x] == -1) {
         for (int y = x; y+1;) {
           int p = par[y], w = match[p];
           match[match[p] = y] = p; y = w;
         } // 30c1
         ans++;
         goto nxt;
       } // 6fd0
       lab[match[x]] = 0; q.pb(match[x]);
      } else if (lab[x] == 0 &&
                orig[v] != orig[x]) {
       int a = orig[v], b = orig[x];
       for (cnt++;; swap(a, b)) if (a+1) {
         if (aux[a] == cnt) break;
         aux[a] = cnt;
         a = (match[a]+1 ?
           orig[par[match[a]]] : -1);
       } // 2776
       blos(x, v, a); blos(v, x, a);
      } // 45a1
   } // d488
   nxt:;
 } // 8d8a
 return ans:
} // d872
graphs/general matching w.h 8fd8
// Edmond's Blossom algorithm for weighted
// maximum matching in general graphs; O(n^3)?
// Weights must be positive (I believe).
struct WeightedBlossom {
 struct edge { int u, v, w; };
 int n, s, nx;
```

```
vector<vector<edge>> g;
Vi lab, match, slack, st, pa, S, vis;
vector<Vi> flo, floFrom;
queue<int> q;
// Initialize for k vertices
WeightedBlossom(int k)
   : n(k), s(n*2+1),
     q(s, vector<edge>(s)),
     lab(s), match(s), slack(s), st(s),
     pa(s), S(s), vis(s), flo(s),
     floFrom(s, Vi(n+1)) {
  rep(u, 1, n+1) rep(v, 1, n+1)
   g[u][v] = \{u, v, 0\};
} // 5e51
// Add edge between u and v with weight w
void addEdge(int u, int v, int w) {
 q[u][v].w = q[v][u].w = max(q[u][v].w, w);
// Compute max weight matching.
// `count` is set to matching size,
// `weight` is set to matching weight.
// Returns vector `match` such that:
// match[v] = vert matched to v or -1
Vi solve(int& count, ll& weight) {
 fill(all(match), 0);
 nx = n;
 weight = count = 0:
  rep(u, 0, n+1) flo[st[u] = u].clear();
  int tmp = 0:
  rep(u, 1, n+1) rep(v, 1, n+1) {
   floFrom[u][v] = (u-v ? 0 : v);
   tmp = max(tmp, q[u][v].w);
  rep(u, 1, n+1) lab[u] = tmp;
  while (matching()) count++;
  rep(u, 1, n+1)
   if (match[u] && match[u] < u)
     weight += g[u][match[u]].w;
 Vi ans(n);
  rep(i, 0, n) ans[i] = match[i+1]-1;
  return ans;
} // d611
int delta(edge& e) {
 return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
void updateSlack(int u, int x) {
 if (!slack[x] || delta(q[u][x]) <</pre>
   delta(g[slack[x]][x])) slack[x] = u;
} // 1f7f
void setSlack(int x) {
  slack[x] = 0;
  rep(u, 1, n+1) if (q[u][x].w > 0 &&
   st[u] != x && !S[st[u]])
     updateSlack(u, x);
} // ee9c
void push(int x) {
 if (x \le n) q.push(x);
 else rep(i, 0, sz(flo[x])) push(flo[x][i]);
} // 594d
void setSt(int x, int b) {
 st[x] = b;
 if (x > n) rep(i, 0, sz(flo[x]))
   setSt(flo[x][i],b);
} // c5c8
int getPr(int b, int xr) {
```

```
int pr = int(find(all(flo[b]), xr) -
    flo[b].begin());
  if (pr % 2) {
    reverse(flo[b].begin()+1, flo[b].end());
    return sz(flo[b]) - pr;
  } else return pr;
} // 399f
void setMatch(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = floFrom[u][e.u], pr = getPr(u,xr);
  rep(i, 0, pr)
    setMatch(flo[u][i], flo[u][i^1]);
  setMatch(xr, v);
  rotate(flo[u].begin(), flo[u].begin()+pr,
    flo[u].end());
void augment(int u, int v) {
  while (1) {
    int xnv = st[match[u]];
    setMatch(u, v);
    if (!xnv) return;
    setMatch(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 } // bd23
} // e61a
int getLca(int u, int v) {
  static int t = 0;
  for (++t; u||v; swap(u, v)) {
    if (!u) continue:
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  1 // aa78
  return 0;
1 // 9f28
void blossom(int u, int lca, int v) {
  int b = n+1;
  while (b <= nx && st[b]) ++b;
  if (b > nx) ++nx:
  lab[b] = S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].pb(lca);
  for (int x=u, y; x != lca; x = st[pa[y]]) {
    flo[b].pb(x);
    flo[b].pb(y = st[match[x]]);
    push (y);
  } // 63e0
  reverse(flo[b].begin()+1, flo[b].end());
  for (int x=v, y; x != lca; x = st[pa[y]]) {
    flo[b].pb(x);
    flo[b].pb(y = st[match[x]]);
   push (y);
  } // 63e0
  setSt(b, b);
  rep(x, 1, nx+1) g[b][x].w = g[x][b].w = 0;
  rep(x, 1, n+1) floFrom[b][x] = 0;
  rep(i, 0, sz(flo[b])) {
    int xs = flo[b][i];
    rep(x, 1, nx+1) if (!g[b][x].w ||
      delta(g[xs][x]) < delta(g[b][x]))
        g[b][x]=g[xs][x], g[x][b]=g[x][xs];
    rep(x, 1, n+1) if (floFrom[xs][x])
      floFrom[b][x] = xs;
  } // 5833
```

```
setSlack(b);
} // 9000
void blossom(int b) {
  each(e, flo[b]) setSt(e, e);
  int xr = floFrom[b][g[b][pa[b]].u];
  int pr = getPr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i+1];
    pa[xs] = q[xns][xs].u;
    S[xs] = 1; S[xns] = slack[xs] = 0;
    setSlack(xns); push(xns);
  } // f26f
  S[xr] = 1; pa[xr] = pa[b];
  rep(i, pr+1, sz(flo[b])) {
    int xs = flo[b][i];
    S[xs] = -1; setSlack(xs);
  } // a12a
  st[b] = 0;
} // f750
bool found (const edge& e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u; S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = S[nu] = 0;
    push (nu);
  } else if (!S[v]) {
    int lca = getLca(u, v);
    if (!lca) return augment(u, v),
      augment (v, u), 1;
    else blossom(u, lca, v):
  } // ddbb
  return 0;
} // 1c00
bool matching() {
  fill(S.begin(), S.begin()+nx+1, -1);
  fill(slack.begin(), slack.begin()+nx+1, 0);
  q = \{\};
  rep(x, 1, nx+1)
    if (st[x] == x && !match[x])
      pa[x] = S[x] = 0, push(x);
  if (q.empty()) return 0;
  while (1) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      rep(v, 1, n+1)
        if (g[u][v].w > 0 && st[u] != st[v]){
          if (!delta(g[u][v])) {
            if (found(g[u][v])) return 1;
          } else updateSlack(u, st[v]);
        1 // b782
    } // 4d33
    int d = INT MAX;
    rep (b, n+1, nx+1)
      if (st[b] == b && S[b] == 1)
        d = min(d, lab[b]/2);
    rep(x, 1, nx+1)
      if (st[x] == x && slack[x]) {
        if (S[x] == -1)
          d = min(d, delta(g[slack[x]][x]));
        else if (!S[x])
          d = min(d, delta(q[slack[x]][x])/2);
      } // 2a0e
    rep(u, 1, n+1) {
      if (!S[st[u]]) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
```

```
} else if (S[st[u]] == 1) lab[u] += d;
     } // 4601
      rep(b, n+1, nx+1) if (st[b] == b) {
       if (!S[st[b]]) lab[b] += d*2;
       else if (S[st[b]] == 1) lab[b] \rightarrow d*2;
     } // e09b
     q = {};
     rep(x, 1, nx+1)
       if (st[x] == x && slack[x] &&
         st[slack[x]] != x &&
         !delta(g[slack[x]][x]) &&
         found(q[slack[x]][x])) return 1;
      rep (b, n+1, nx+1)
       if (st[b] == b && S[b] == 1 && !lab[b])
         blossom(b);
   } // a122
   return 0:
 } // e966
}; // 3ee2
graphs/global min cut.h
                                         0eac
// Find a minimum cut in an undirected graph
// with non-negative edge weights
// given its adjacency matrix M; time: O(n^3)
// 'out' contains vertices on one side.
ll minCut(vector<vector<ll>>> M, Vi& out) {
 int n = sz(M):
 ll ans = INT64 MAX;
 vector<Vi> co(n);
 rep(i, 0, n) co[i].pb(i);
 out.clear();
 rep(ph, 1, n) {
   auto w = M[0];
   size t s = 0, t = 0;
   // O(V^2) -> O(E log V) with priority queue
   rep(it, 0, n-ph) {
     w[t] = INT64_MIN; s = t;
     t = max_element(all(w)) - w.begin();
     rep(i, 0, n) w[i] += M[t][i];
   } // 0831
   ll alt = w[t] - M[t][t];
   if (alt < ans) ans = alt, out = co[t];</pre>
   co[s].insert(co[s].end(), all(co[t]));
   rep(i, 0, n) M[s][i] += M[t][i];
   rep(i, 0, n) M[i][s] = M[s][i];
   M[0][t] = INT64_MIN;
 } // df69
 return ans;
} // 2a85
graphs/gomory hu.h
                                         a520
#include "edmonds karp.h"
//#include "push_relabel.h" // if you need
struct Edge {
 int a, b; // vertices
 flow_t w; // weight
}; // c331
// Build Gomory-Hu tree; time: O(n*maxflow)
// Gomory-Hu tree encodes minimum cuts between
// all pairs of vertices: mincut for u and v
// is equal to minimum on path from u and v
// in Gomory-Hu tree. n is vertex count.
// Returns vector of Gomory-Hu tree edges.
vector<Edge> gomoryHu (vector<Edge>& edges,
                      int n) {
 MaxFlow flow(n);
```

each(e, edges) flow.addEdge(e.a,e.b,e.w,e.w);

```
vector<Edge> ret(n-1);
  rep(i, 1, n) ret[i-1] = {i, 0, 0};
  rep(i, 1, n) {
   ret[i-1].w = flow.maxFlow(i, ret[i-1].b);
    rep(j, i+1, n)
     if (ret[j-1].b == ret[i-1].b &&
          flow.cutSide(j)) ret[j-1].b = i;
  } // 5ae4
  return ret;
} // afdb
graphs/kth shortest.h
                                          1f40
constexpr ll INF = 1e18;
// Eppstein's k-th shortest path algorithm;
// time and space: O((m+k) log (m+k))
struct Eppstein {
  using T = 11; // Type for edge weights
  using Edge = pair<int, T>;
  struct Node {
   int E[2] = \{\}, s\{0\};
   Edge x;
  }; // 013b
  T shortest; // Shortest path length
  priority_queue<pair<T, int>> Q;
  vector<Node> P{1};
  Vi h;
  // Initialize shortest path structure for
  // weighted graph G, source s and target t;
  // time: O(m log m)
  Eppstein (vector < vector < Edge >> & G,
          int s, int t) {
    int n = sz(G);
    vector<vector<Edge>> H(n);
    rep(i,0,n) each(e,G[i]) H[e.x].pb(\{i,e.y\});
    Vi ord, par(n, -1);
    vector<T> d(n, -INF);
    Q.push(\{d[t] = 0, t\});
    while (!O.emptv()) {
     auto v = Q.top();
      Q.pop();
      if (d[v.v] == v.x) {
       ord.pb(v.y);
        each (e, H[v.y]) if (v.x-e.y > d[e.x]) {
          Q.push(\{d[e.x] = v.x-e.y, e.x\});
          par[e.x] = v.y;
       } // 5895
     } // 1b62
    } // 1a6d
    if ((shortest = -d[s]) >= INF) return;
   h.resize(n);
    each (v, ord) {
     int p = par[v];
      if (p+1) h[v] = h[p];
      each(e, G[v]) if (d[e.x] > -INF) {
       T k = e.y - d[e.x] + d[v];
       if (k || e.x != p)
         h[v] = push(h[v], \{e.x, k\});
        else
         p = -1;
     } // 5e05
    } // 31b9
    P[0].x.x = s;
    Q.push({0, 0});
```

```
} // e00e
  int push(int t, Edge x) {
   P.pb(P[t]);
   if (!P[t = sz(P)-1].s || P[t].x.y >= x.y)
      swap(x, P[t].x);
    if (P[t].s) {
      int i = P[t].E[0], j = P[t].E[1];
      int d = P[i].s > P[j].s;
     int k = push(d ? j : i, x);
     P[t].E[d] = k; // Don't inline k!
    } // 10e1
   P[t].s++;
   return t:
 } // a2dc
  // Get next shortest path length,
 // the first call returns shortest path.
 // Returns -1 if there's no more paths;
  // time: O(\log k), where k is total count
  // of nextPath calls.
 11 nextPath() {
   if (Q.empty()) return -1;
   auto v = Q.top();
   0.pop();
    for (int i : P[v.y].E) if (i)
     Q.push({ v.x-P[i].x.y+P[v.y].x.y, i });
    int t = h[P[v.v].x.x];
    if (t) Q.push({ v.x - P[t].x.y, t });
    return shortest - v.x;
 } // 08af
}; // 5326
graphs/matroids.h
                                          55ef
// Find largest subset S of [n] such that
// S is independent in both matroid A and B.
// A and B are given by their oracles,
// see example implementations below.
// Returns vector V such that V[i] = 1 iff
// i-th element is included in found set;
// time: O(r^2*init + r^2*n*add),
// where r is max independent set,
// `init` is max time of oracles init
// and `add` is max time of oracles canAdd.
template < class T, class U>
vector<bool> intersectMatroids (T& A, U& B,
                               int n) {
  vector<bool> ans(n);
 bool ok = 1;
  // NOTE: for weighted matroid intersection
  // find shortest augmenting paths
  // first by weight change, then by length
  // using Bellman-Ford, and skip this speedup:
 A.init(ans):
 B.init(ans):
 rep(i, 0, n) if (A.canAdd(i) && B.canAdd(i))
   ans[i] = 1, A.init(ans), B.init(ans);
  while (ok) {
   vector<Vi> G(n);
    vector<bool> good(n);
    queue<int> que;
   Vi prev(n, -1);
   A.init(ans);
   B.init(ans);
   ok = 0;
    rep(i, 0, n) if (!ans[i]) {
     if (A.canAdd(i)) que.push(i), prev[i]=-2;
```

```
good[i] = B.canAdd(i);
   } // 9581
   rep(i, 0, n) if (ans[i]) {
     ans[i] = 0;
     A.init(ans);
     B.init(ans);
     rep(j, 0, n) if (i != j && !ans[j]) {
       if (A.canAdd(j)) G[i].pb(j);
       if (B.canAdd(j)) G[j].pb(i);
     } // bd2a
     ans[i] = 1;
   } // bf3e
   while (!que.empty()) {
     int i = que.front();
     que.pop();
     if (good[i]) {
        ans[i] = 1;
        while (prev[i] >= 0) {
         ans[i = prev[i]] = 0;
         ans[i = prev[i]] = 1;
        } // 51c8
        ok = 1:
       break;
     } // 384b
     each(j, G[i]) if (prev[j] == -1)
       que.push(j), prev[j] = i;
   } // 6eb6
 } // e092
 return ans;
} // ae8e
// Matroid where each element has color
// and set is independent iff for each color c
// #{elements of color c} <= maxAllowed[c].
struct LimOracle {
 Vi color; // color[i] = color of i-th element
 Vi maxAllowed; // Limits for colors
 // Init oracle for independent set S; O(n)
 void init(vector<bool>& S) {
   tmp = maxAllowed;
   rep(i, 0, sz(S)) tmp[color[i]] -= S[i];
 } // 4dfb
  // Check if S+{k} is independent; time: O(1)
 bool canAdd(int k) {
   return tmp[color[k]] > 0;
 } // e312
}; // 7b5d
// Graphic matroid - each element is edge,
// set is independent iff subgraph is acyclic.
struct GraphOracle {
 vector<Pii> elems; // Ground set: graph edges
 int n; // Number of vertices, indexed [0;n-1]
 Vi par;
 int find(int i) {
   return par[i] == -1 ? i
     : par[i] = find(par[i]);
 } // b8b7
 // Init oracle for independent set S; ~O(n)
 void init(vector<bool>& S) {
   par.assign(n, -1);
   rep(i, 0, sz(S)) if (S[i])
     par[find(elems[i].x)] = find(elems[i].y);
 } // 1827
```

```
// Check if S+{k} is independent; time: ~O(1)
 bool canAdd(int k) {
   return
     find(elems[k].x) != find(elems[k].y);
 } // 8ca4
}; // c506
// Co-graphic matroid - each element is edge,
// set is independent iff after removing edges
// from graph number of connected components
// doesn't change.
struct CographOracle {
 vector<Pii> elems; // Ground set: graph edges
 int n; // Number of vertices, indexed [0;n-1]
 vector<Vi> G;
 Vi pre, low;
 int cnt;
 int dfs(int v, int p) {
   pre[v] = low[v] = ++cnt;
   each(e, G[v]) if (e != p)
     low[v] = min(low[v], pre[e] ?: dfs(e,v));
   return low[v];
 } // 9d30
 // Init oracle for independent set S; O(n)
 void init(vector<bool>& S) {
   G.assign(n, {});
   pre.assign(n, 0);
   low.resize(n):
   cnt = 0;
   rep(i, 0, sz(S)) if (!S[i]) {
     Pii e = elems[i]:
     G[e.x].pb(e.v);
     G[e.y].pb(e.x);
   } // 79a1
   rep(v, 0, n) if (!pre[v]) dfs(v, -1);
 } // 1200
 // Check if S+{k} is independent; time: O(1)
 bool canAdd(int k) {
   Pii e = elems[k]:
   return max(pre[e.x], pre[e.y])
     != max(low[e.x], low[e.y]);
 1 // 2550
}; // a5cc
// Matroid equivalent to linear space with XOR
struct XorOracle {
 vector<ll> elems; // Ground set: numbers
 vector<11> base;
 // Init for independent set S; O(n+r^2)
 void init(vector<bool>& S) {
   base.assign(63, 0);
   rep(i, 0, sz(S)) if (S[i]) {
     ll e = elems[i];
     rep(j, 0, sz(base)) if ((e >> j) & 1) {
       if (!base[j]) {
         base[j] = e;
         break;
       } // 1df5
       e ^= base[j];
     } // 8495
   } // 655e
 } // b68c
  // Check if S+{k} is independent; time: O(r)
 bool canAdd(int k) {
   11 e = elems[k];
   rep(i, 0, sz(base)) if ((e >> i) & 1) {
```

```
if (!base[i]) return 1;
     e ^= base[i];
   } // 49d1
   return 0;
 } // 66ff
1: // 4af3
graphs/min cost max flow.h
using flow_t = 11;
constexpr flow t INF = 1e18;
// Min cost max flow using cheapest paths;
// time: O(nm + |f| * (m log n))
// or O(|f|*(m log n)) if costs are nonnegative
struct MCMF {
  struct Edge {
   int dst, inv;
   flow_t flow, cap, cost;
 }; // 20f7
  vector<vector<Edge>> G;
  vector<flow t> add;
  // Initialize for n vertices
  MCMF(int n = 0) : G(n) {}
  // Add new vertex
  int addVert() { G.pb({}); return sz(G)-1; }
  // Add edge from u to v.
  // Returns edge index in adjacency list of u.
  int addEdge(int u, int v,
              flow_t cap, flow_t cost) {
   G[u].pb({v, sz(G[v]), 0, cap, cost });
   G[v].pb({u, sz(G[u])-1, 0, 0, -cost });
   return sz(G[u])-1;
  1 // 1095
  // Compute minimum cost maximum flow
  // from src to dst. `f` is set to flow value,
  // `c` is set to total cost value.
  // Returns false iff negative cycle
  // is reachable from from source.
  bool maxFlow(int src, int dst,
               flow t& f, flow t& c) {
    flow t i, m, d:
   f = c = 0;
   each (v, G) each (e, v) e.flow = 0;
    // [If costs are nonnegative]
    // vector<flow_t> pot(sz(G));
    // [/end]
    // [If costs can be negative] O(n*m)
   vector<flow t> pot(sz(G), INF);
   pot[src1 = 0:
    int it = sz(G), ch = 1;
    while (ch-- && it--)
     rep(s, 0, sz(G)) if (pot[s] != INF)
       each(e, G[s]) if (e.cap)
          if ((d = pot[s]+e.cost) < pot[e.dst])
           pot[e.dst] = d, ch = 1;
   if (it < 0) return 0;</pre>
   // [/end]
    Vi prev(sz(G), -1);
   vector<flow_t> dist(sz(G), INF);
   priority_queue<pair<flow_t, int>> Q;
   add.assign(sz(G), -1);
   Q.push({0, src});
    add[src] = INF;
   dist[src] = 0;
```

```
while (!Q.empty()) {
     tie(d, i) = Q.top();
     Q.pop();
     if (d != -dist[i]) continue;
     m = add[i];
     if (i == dst) {
       f += m;
       c += m * (dist[i]-pot[src]+pot[i]);
       while (i != src) {
         auto& e = G[i][prev[i]];
         e.flow -= m;
         G[i = e.dst][e.inv].flow += m;
       } // 1f86
       rep(j, 0, sz(G))
         pot[j] = min(pot[j]+dist[j], INF);
       goto nxt;
     1 // 36d4
     each(e, G[i]) if (e.flow < e.cap) {
       d = dist[i]+e.cost+pot[i]-pot[e.dst];
       if (d < dist[e.dst]) {</pre>
         Q.push({-d, e.dst});
         prev[e.dst] = e.inv;
         add[e.dst] = min(m, e.cap-e.flow);
         dist[e.dst] = d;
       } // 5ee6
     } // b6b2
   } // d47c
   return 1;
 } // Obc3
 // Get flow through e-th edge of vertex v
 flow t getFlow(int v, int e) {
   return G[v][e].flow;
 } // Ofaf
 // Get if v belongs to cut component with src
 bool cutSide(int v) { return add[v] >= 0; }
graphs/push relabel.h
                                         07d1
using flow t = int;
// Push-relabel algorithm for maximum flow;
// O(V^2*sqrt(E)), but very fast in practice.
struct MaxFlow {
 struct Edge {
   int to, inv;
   flow t rem, cap;
 }; // bc77
 vector < basic string < Edge >> G;
 vector<flow t> extra:
 Vi hei, arc, prv, nxt, act, bot;
 queue<int> 0:
 int n, high, cut, work;
 // Initialize for k vertices
 MaxFlow(int k = 0) : G(k) {}
 // Add new vertex
 int addVert() { G.pb({}); return sz(G)-1; }
 // Add edge from u to v with capacity cap
 // and reverse capacity rcap.
 // Returns edge index in adjacency list of u.
 int addEdge(int u, int v,
             flow_t cap, flow_t rcap = 0) {
   G[u].pb({ v, sz(G[v]), 0, cap });
   G[v].pb({u, sz(G[u])-1, 0, rcap });
   return sz(G[u])-1;
```

```
1 // c96a
void raise(int v, int h) {
  prv[nxt[prv[v]] = nxt[v]] = prv[v];
  hei[v] = h;
  if (extra[v] > 0) {
    bot[v] = act[h]; act[h] = v;
    high = max(high, h);
  } // d7ee
  if (h < n) cut = max(cut, h+1);
  nxt[v] = nxt[prv[v] = h += n];
  prv[nxt[nxt[h] = v]] = v;
} // 5274
void global(int s, int t) {
 hei.assign(n, n*2);
  act.assign(n*2, -1);
  iota(all(prv), 0);
  iota(all(nxt), 0);
  hei[t] = high = cut = work = 0;
  hei[s] = n;
  for (int x : {t, s})
    for (Q.push(x); !Q.empty(); Q.pop()) {
      int v = Q.front();
      each(e, G[v])
        if (hei[e.to] == n*2 &&
            G[e.to][e.inv].rem)
          Q.push(e.to), raise(e.to,hei[v]+1);
    } // 1901
} // 3181
void push(int v, Edge& e, bool z) {
  auto f = min(extra[v], e.rem);
  if (f > 0) {
    if (z && !extra[e.to]) {
      bot[e.to] = act[hei[e.to]];
      act[hei[e.to]] = e.to;
    } // 9d90
    e.rem -= f; G[e.to][e.inv].rem += f;
    extra[v] -= f; extra[e.to] += f;
  } // Offb
} // da44
void discharge(int v) {
  int h = n*2, k = hei[v];
  rep(j, 0, sz(G[v])) {
    auto& e = G[v][arc[v]];
    if (e.rem) {
      if (k == hei[e.to]+1) {
        push (v, e, 1);
        if (extra[v] <= 0) return;</pre>
      } else h = min(h, hei[e.to]+1);
    if (++arc[v] >= sz(G[v])) arc[v] = 0;
  1 // 9741
  if (k < n \& \& nxt[k+n] == prv[k+n]) {
    rep(j, k, cut) while (nxt[j+n] < n)
      raise(nxt[j+n], n);
    cut = k;
  } else raise(v, h), work++;
1 // b64f
// Compute maximum flow from src to dst
flow_t maxFlow(int src, int dst) {
  extra.assign(n = sz(G), 0);
  arc.assign(n, 0);
  prv.resize(n*3);
  nxt.resize(n*3);
  bot.resize(n):
  each (v, G) each (e, v) e.rem = e.cap;
```

```
each(e, G[src])
      extra[src] = e.cap, push(src, e, 0);
    global(src, dst);
    for (; high; high--)
      while (act[high] != -1) {
        int v = act[high];
        act[high] = bot[v];
        if (v != src && hei[v] == high) {
          discharge(v);
          if (work > 4*n) global(src, dst);
        } // 7dcc
      ) // 26d4
    return extra[dst]:
 } // aa5e
  // Get flow through e-th edge of vertex v
  flow t getFlow(int v, int e) {
    return G[v][e].cap - G[v][e].rem;
 } // 812c
  // Get if v belongs to cut component with src
 bool cutSide(int v) { return hei[v] >= n; }
}; // 2d6b
graphs/scc.h
                                          1c43
// Tarjan's SCC algorithm; time: O(n+m)
// Usage: SCC scc(graph);
// scc[v] = index of SCC for vertex v
// scc.comps[i] = vertices of i-th SCC
// Components are in reversed topological order
struct SCC : Vi {
 vector<Vi> comps:
  Vi S;
  SCC() {}
  SCC(vector\langle \text{Vi} \rangle \& \text{G}) : Vi(sz(G),-1), S(sz(G)) {
    rep(i, 0, sz(G)) if (!S[i]) dfs(G, i);
  } // f0fa
  int dfs(vector<Vi>& G, int v) {
    int low = S[v] = sz(S);
    S.pb(v);
    each (e, G[v]) if (at (e) < 0)
      low = min(low, S[e] ?: dfs(G, e));
    if (low == S[v]) {
      comps.pb({});
      rep(i, S[v], sz(S)) {
        at(S[i]) = sz(comps)-1;
        comps.back().pb(S[i]);
      } // 8ed0
      S.resize(S[v]);
    } // ecc7
    return low:
 } // f3c6
}; // 215e
graphs/turbo matching.h
                                          d400
// Find maximum bipartite matching; time: ?
// G must be bipartite graph!
// Returns matching size (edge count).
// match[v] = vert matched to v or -1
int matching(vector<Vi>& G, Vi& match) {
 vector<bool> seen:
  int n = 0, k = 1;
  match.assign(sz(G), -1);
```

function<int(int)> dfs = [&](int i) {

```
if (seen[i]) return 0;
    seen[i] = 1;
    each(e, G[i]) {
     if (match[e] < 0 || dfs(match[e])) {</pre>
       match[i] = e; match[e] = i;
        return 1:
     } // 893d
   } // 9532
   return 0;
  }; // d332
  while (k) {
    seen.assign(sz(G), 0);
    rep(i, 0, sz(G)) if (match[i] < 0)
     k += dfs(i);
   n += k;
  } // 1128
  return n:
} // 0d38
// Convert maximum matching to vertex cover
// time: O(n+m)
Vi vertexCover(vector<Vi>& G, Vi& match) {
  Vi ret, col(sz(G)), seen(sz(G));
  function<void(int, int)> dfs =
     [&] (int i, int c) {
    if (col[i]) return;
    col[i] = c+1;
    each(e, G[i]) dfs(e, !c);
  }; // 1f1b
  function<void(int)> aug = [&](int i) {
   if (seen[i] || col[i] != 1) return;
    seen[i] = 1;
   each(e, G[i]) seen[e] = 1, aug(match[e]);
  ): // 2465
  rep(i, 0, sz(G)) dfs(i, 0);
  rep(i, 0, sz(G)) if (match[i] < 0) aug(i);
  rep(i, 0, sz(G))
   if (seen[i] == col[i]-1) ret.pb(i);
  return ret;
} // 6f72
graphs/weighted matching.h
                                          8264
// Minimum cost bipartite matching; O(n^2*m)
// Input is n x m cost matrix, where n <= m.
// Returns matching weight.
// L[i] = right vertex matched to i-th left
// R[i] = left vertex matched to i-th right
ll hungarian (const vector < vector < ll >> cost.
             Vi& L, Vi& R) {
  if (cost.empty())
    return L.clear(), R.clear(), 0;
  int b, c = 0, n = sz(cost), m = sz(cost[0]);
  assert (n <= m);
  vector<ll> x(n), y(m+1);
  L.assign(n, -1);
  R.assign(m+1, -1);
  rep(i, 0, n) {
   vector<11> sla(m, INT64_MAX);
    Vi vis(m+1), prv(m, -1);
    for (R[b = m] = i; R[b]+1; b = c) {
     int a = R[b];
     11 d = INT64_MAX;
      vis[b] = 1:
      rep(j, 0, m) if (!vis[j]) {
       ll cur = cost[a][j] - x[a] - y[j];
```

```
if (cur < sla[j])</pre>
          sla[j] = cur, prv[j] = b;
        if (sla[j] < d) d = sla[j], c = j;</pre>
      } // 6717
      rep(i, 0, m+1) {
        if (vis[j]) x[R[j]] \leftarrow d, y[j] \leftarrow d;
        else sla[j] -= d;
                                                    Vi col;
     } // 8bb3
    } // 01c6
    while (b-m) c = b, R[c] = R[b = prv[b]];
 } // 71f5
  rep(j, 0, m) if (R[j]+1) L[R[j]] = j;
 R.resize(m);
  return -y[m];
} // 349d
math/berlekamp massey.h
                                          7d12
constexpr int MOD = 998244353;
                                                    } // a953
ll modInv(ll a, ll m) { // a^{(-1)} \mod m
 if (a == 1) return 1:
 return ((a - modInv(m%a, a)) *m + 1) / a;
} // c437
                                                    } // 4ca1
// Find shortest linear recurrence that matches
// given starting terms of recurrence; O(n^2)
                                                  } // 996e
// Returns vector C such that for each i >= |C|
// A[i] = sum A[i-j-1]*C[j] for j = 0..|C|-1
vector<ll> massey (vector<ll> € A) {
 if (A.empty()) return {};
  int n = sz(A), len = 0, k = 0;
 11 s = 1:
  vector<ll> B(n), C(n), tmp;
 B[0] = C[0] = 1;
  rep(i, 0, n) {
   11 d = 0;
    k++;
    rep(j, 0, len+1)
      d = (d + C[j] * A[i-j]) % MOD;
                                                    } // 7ef0
      11 q = d * modInv(s, MOD) % MOD;
      tmp = C:
      rep(j, k, n)
       C[j] = (C[j] - q * B[j-k]) % MOD;
      if (len*2 <= i) {</pre>
        B.swap(tmp);
                                                    } // 28bd
        len = i-len+1;
        s = d + (d < 0) * MOD;
       k = 0;
      } // c350
   } // 79c7
  } // f70c
                                                   } // 98a8
  C.resize(len+1);
                                                  }; // 4df7
  C.erase(C.begin());
  each (x, C) x = (MOD - x) % MOD;
  return C:
} // 20ce
math/bit_gauss.h
                                          13eb
constexpr int MAX COLS = 2048:
// Solve system of linear equations over Z 2
                                                   return d;
// time: O(n^2*m/W), where W is word size
                                                  } // 23c8
// - A - extended matrix, rows are equations,
         columns are variables.
         m-th column is equation result
         (A[i][j] - i-th row and j-th column)
```

```
// - ans - output for variables values
// - m - variable count
// Returns 0 if no solutions found, 1 if one,
// 2 if more than 1 solution exist.
int bitGauss (vector < bitset < MAX COLS >> & A,
            vector<bool>& ans, int m) {
 ans.assign(m, 0);
 rep(i, 0, sz(A)) {
   int c = int(A[i]._Find_first());
   if (c >= m) {
     if (c == m) return 0;
     continue:
   } // a6bb
    rep(k, i+1, sz(A)) if (A[k][c]) A[k]^=A[i];
   swap(A[i], A[sz(col)]);
   col.pb(c);
 for (int i = sz(col); i--;) if (A[i][m]) {
   ans [col[i]] = 1;
   rep(k,0,i) if(A[k][col[i]]) A[k][m].flip();
 return sz(col) < m ? 2 : 1;</pre>
                                         2e3f
math/bit matrix.h
using ull = uint64 t;
// Matrix over Z_2 (bits and xor)
// TODO: arithmetic operations
struct BitMatrix {
 vector<ull> M;
 int rows, cols, stride;
  // Create matrix with n rows and m columns
 BitMatrix(int n = 0, int m = 0) {
   rows = n; cols = m;
   stride = (m+63)/64;
   M.resize(n*stride);
 // Get pointer to bit-packed data of i-th row
 ull* row(int i) { return &M[i*stride]; }
  // Get value in i-th row and j-th column
 bool operator()(int i, int j) {
   return (row(i)[j/64] >> (j%64)) & 1;
 // Set value in i-th row and j-th column
 void set(int i, int j, bool val) {
   ull &w = row(i)[j/64], m = 1ull << (j^64);
   if (val) w |= m;
   else w &= ~m;
math/crt.h
                                         8a85
using Pll = pair<11, 11>;
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 ll d = egcd(b%a, a, y, x);
 x = b/a*y;
// Chinese Remainder Theoerem; time: O(lg lcm)
// Solves x = a.x \pmod{a.y}, x = b.x \pmod{b.y}
// Returns pair (x mod lcm, lcm(a.y, b.y))
```

```
// or (-1, -1) if there's no solution.
// WARNING: a.x and b.x are assumed to be
// in [0;a.y) and [0;b.y) respectively.
// Works properly if lcm(a.y, b.y) < 2^63.
Pll crt (Pll a, Pll b) {
 if (a.y < b.y) swap(a, b);
  ll x, y, g = egcd(a.y, b.y, x, y);
  11 c = b.x-a.x, d = b.y/g, p = a.y*d;
  if (c % q) return {-1, -1};
 11 s = (a.x + c/g*x % d * a.y) % p;
  return {s < 0 ? s+p : s, p};
} // 35a8
math/fft complex.h
                                          945b
using dbl = double;
using cmpl = complex<dbl>;
// Default std::complex multiplication is slow.
// You can use this to achieve small speedup.
cmpl operator*(cmpl a, cmpl b) {
 dbl ax = real(a), av = imag(a);
  dbl bx = real(b), by = imag(b);
  return {ax*bx-ay*by, ax*by+ay*bx};
} // 3b78
cmpl operator*=(cmpl& a,cmpl b) {return a=a*b;}
// Compute DFT over complex numbers; O(n lg n)
// Input size must be power of 2!
void fft(vector<cmpl>& a) {
  static vector<cmpl> w(2, 1);
  int n = sz(a);
  for (int k = sz(w); k < n; k *= 2) {
    w.resize(n);
    rep(i,0,k) w[k+i] = \exp(\text{cmpl}(0, M_PI*i/k));
  } // 92a9
  rep(i,0,n) rev[i] = (rev[i/2] | i\%2*n) / 2;
  rep(i,0,n) if(i<rev[i]) swap(a[i],a[rev[i]]);
  for (int k = 1; k < n; k *= 2) {
    for (int i=0; i < n; i += k*2) rep(j,0,k) {
      auto d = a[i+j+k] * w[j+k];
      a[i+j+k] = a[i+j] - d;
      a[i+j] += d;
    } // b389
 } // 84bf
} // adf8
// Convolve complex-valued a and b,
// store result in a; time: O(n lg n), 3x FFT
void convolve(vector<cmpl>& a, vector<cmpl> b) {
 int len = sz(a) + sz(b) - 1;
  if (len <= 0) return a.clear();</pre>
  int n = 1 << (32 - __builtin_clz(len));</pre>
  a.resize(n); b.resize(n);
  fft(a); fft(b);
  rep(i, 0, n) a[i] *= b[i] / dbl(n);
  reverse (a.begin()+1, a.end());
  fft(a);
 a.resize(len);
} // 8fe9
// Convolve real-valued a and b, returns result
// time: O(n lg n), 2x FFT
// Rounding to integers is safe as long as
// (max\_coeff^2)*n*log\_2(n) < 9*10^14
// (in practice 10^16 or higher).
vector<dbl> convolve (vector<dbl> € a,
                     vector<dbl>& b) {
```

```
int len = \max(sz(a) + sz(b) - 1, 0);
  int n = 1 << (32 - __builtin_clz(len));</pre>
  vector<cmpl> in(n), out(n);
  rep(i, 0, sz(a)) in[i].real(a[i]);
  rep(i, 0, sz(b)) in[i].imag(b[i]);
  fft(in);
  each(x, in) x \star = x;
  rep(i,0,n) out[i] = in[-i&(n-1)]-conj(in[i]);
  fft(out);
  vector<dbl> ret(len);
  rep(i, 0, len) ret[i] = imag(out[i]) / (n*4);
  return ret:
} // 19ed
constexpr 11 MOD = 1e9+7;
// High precision convolution of integer-valued
// a and b mod MOD; time: O(n lg n), 4x FFT
// Input is expected to be in range [0; MOD)!
// Rounding is safe if MOD*n*log_2(n) < 9*10^14
// (in practice 10^16 or higher).
vector<ll> convMod(vector<ll> € a,
                  vector<ll>& b) {
  vector<ll> ret(sz(a) + sz(b) - 1);
  int n = 1 \ll (32 - builtin clz(sz(ret)));
  11 cut = 11(sqrt(MOD))+1;
  vector<cmpl> c(n), d(n), q(n), f(n);
  rep(i, 0, sz(a))
   c[i] = {dbl(a[i]/cut), dbl(a[i]%cut)};
  rep(i, 0, sz(b))
   d[i] = {dbl(b[i]/cut), dbl(b[i]%cut)};
  fft(c); fft(d);
  rep(i, 0, n) {
    int j = -i & (n-1);
    f[j] = (c[i] + conj(c[j])) * d[i] / (n*2.0);
   g[j] =
      (c[i]-conj(c[j])) * d[i] / cmpl(0, n*2);
  } // e877
  fft(f); fft(q);
  rep(i, 0, sz(ret)) {
   11 t = llround(real(f[i])) % MOD * cut;
   t += llround(imag(f[i]));
   t = (t + llround(real(g[i]))) % MOD * cut;
   t = (t + llround(imag(g[i]))) % MOD;
   ret[i] = (t < 0 ? t+MOD : t);
 } // e75d
  return ret;
} // 513f
math/fft mod.h
                                          0a8a
// Number Theoretic Tranform (NTT)
// For functions below you can choose 2 params:
// 1. M - prime modulus that MUST BE of form
         a*2^k+1, computation is done in Z_M
// 2. R - generator of Z M
// Modulus often seen on Codeforces:
// M = (119<<23)+1, R = 62; M is 998244353
// Parameters for 11 computation with CRT:
// M = (479 << 21) + 1, R = 62; M is > 10^9
// M = (483<<21)+1, R = 62; M is > 10^9
ll modPow(ll a, ll e, ll m) {
 ll t = 1 % m;
 while (e) {
```

```
if (e % 2) t = t*a % m;
   e /= 2; a = a*a % m;
 } // 66ca
  return t;
} // 1973
// Compute DFT over Z_M with generator R.
// Input size must be power of 2; O(n lq n)
// Input is expected to be in range [0; MOD)!
// dit == true <=> inverse transform * 2^n
                   (without normalization)
template<11 M, 11 R, bool dit>
void ntt(vector<ll>& a) {
  static vector<1l> w(2, 1);
 int n = sz(a);
  for (int k = sz(w); k < n; k *= 2) {
   w.resize(n, 1);
    11 c = modPow(R, M/2/k, M);
    if (dit) c = modPow(c, M-2, M);
    rep(i, k+1, k*2) w[i] = w[i-1]*c % M;
  } // 0d98
  for (int t = 1; t < n; t *= 2) {
    int k = (dit ? t : n/t/2);
    for (int i=0; i < n; i += k*2) rep(j,0,k) {
      ll &c = a[i+j], &d = a[i+j+k];
      11 e = w[j+k], f = d;
      d = (dit ? c - (f=f*e%M) : (c-f)*e % M);
      if (d < 0) d += M;
      if ((c += f) \geq= M) c -= M;
   } // e4a6
 } // 8d38
} // 01f5
// Convolve a and b mod M (R is generator),
// store result in a; time: O(n lg n), 3x NTT
// Input is expected to be in range [0:MOD]!
template<11 M = (119<<23)+1, 11 R = 62>
void convolve(vector<11>& a, vector<11> b) {
  int len = sz(a) + sz(b) - 1;
  if (len <= 0) return a.clear();</pre>
  int n = 1 << (32 - __builtin_clz(len));</pre>
  ll t = modPow(n, M-2, M);
  a.resize(n); b.resize(n);
  ntt < M, R, 0 > (a); ntt < M, R, 0 > (b);
  rep(i, 0, n) a[i] = a[i] *b[i] % M * t % M;
  ntt < M, R, 1 > (a);
 a.resize(len);
} // b413
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 11 d = egcd(b%a, a, y, x);
 x = b/a*v;
return d:
} // 23c8
// Convolve a and b with 64-bit output,
// store result in a; time: O(n lg n), 6x NTT
// Input is expected to be non-negative!
void convLong(vector<ll>& a, vector<ll> b) {
 const 11 M1 = (479 << 21) +1, M2 = (483 << 21) +1;
  const 11 MX = M1*M2, R = 62;
  auto c = a, d = b;
  each(k, a) k %= M1;
  each(k, b) k %= M1;
  each(k, c) k %= M2;
  each (k, d) k %= M2;
  convolve<M1, R>(a, b);
```

```
convolve<M2, R>(c, d);
  ll x, y; egcd(M1, M2, x, y);
  rep(i, 0, sz(a)) {
   a[i] += (c[i]-a[i]) *x % M2 * M1;
    if ((a[i] %= MX) < 0) a[i] += MX;</pre>
 } // 2279
} // c493
math/fwht.h
                                          3e6f
// Fast Walsh-Hadamard Transform; O(n lg n)
// Input must be power of 2!
// Uncommented version is for XOR.
// OR version is equivalent to sum-over-subsets
// (Zeta transform, inverse is Moebius).
// AND version is same as sum-over-supersets.
template<bool inv, class T>
void fwht (vector<T>& b) {
 for (int s = 1; s < sz(b); s *= 2) {
    for (int i = 0; i < sz(b); i += s*2) {
      rep(j, i, i+s) {
        auto &x = b[j], &y = b[j+s];
        tie(x, y) =
          mp (x+y, x-y);
                                          //XOR
        // inv ? mp(x-y, y) : mp(x+y, y); //AND
        // inv ? mp(x, y-x) : mp(x, x+y); //OR
      } // eea9
   } // a3d5
 } // 95ed
  // ONLY FOR XOR:
 if (inv) each(e, b) e /= sz(b);
} // 0779
// Compute convolution of a and b such that
// ans[i#j] += a[i]*b[j], where # is OR, AND
// or XOR, depending on FWHT version.
// Stores result in a; time: O(n lg n)
// Both arrays must be of same size = 2^n!
template<class T>
void bitConv(vector<T>& a, vector<T> b) {
 fwht<0>(a);
  fwht<0>(b);
  rep(i, 0, sz(a)) a[i] *= b[i];
 fwht<1>(a);
} // 7b82
                                          7f0b
math/gauss.h
// Solve system of linear equations; O(n^2*m)
// - A - extended matrix, rows are equations,
         columns are variables.
        m-th column is equation result
         (A[i][j] - i-th row and j-th column)
// - ans - output for variables values
// - m - variable count
// Returns 0 if no solutions found, 1 if one,
// 2 if more than 1 solution exist.
int gauss (vector < vector < double >> & A,
         vector<double>& ans, int m) {
  Vi col;
  ans.assign(m, 0);
  rep(i, 0, sz(A)) {
   int c = 0:
    while (c <= m && !cmp(A[i][c], 0)) c++;</pre>
    // For Zp:
    //while (c <= m && !A[i][c].x) c++;
    if (c >= m) {
      if (c == m) return 0;
```

```
continue:
    } // a6bb
    rep(k, i+1, sz(A)) {
      auto mult = A[k][c] / A[i][c];
      rep(j, 0, m+1) A[k][j] -= A[i][j]*mult;
    } // 8dd5
    swap(A[i], A[sz(col)]);
    col.pb(c);
 } // ea2c
  for (int i = sz(col); i--;) {
    ans[col[i]] = A[i][m] / A[i][col[i]];
    rep(k, 0, i)
     A[k][m] = ans[col[i]] * A[k][col[i]];
 } // 31b9
 return sz(col) < m ? 2 : 1;
} // Ob76
math/gauss ortho.h
                                         c6ae
using Row = vector<double>;
using Matrix = vector<Row>;
// Given a system of n linear equations A
// over m variables, find dimensionality D
// of solution subspace, matrix M and vector t
// such that:
// - matrix M is orthogonal (i.e. M*M^T = I)
// - x is a solution \iff (Mx+t)[D..] = 0
// - x[D..] = 0 \iff M^T(x-t) is a solution
// (in particular -M^T*t is a solution)
// Returns number of dimensions D, or -1 if
// there is no solution; time: O(n^2*m + n*m^2)
// Warning: numerical stability is kinda sus
int orthoGauss (Matrix A. Matrix M.
              Row& t, int m) {
 int d = m;
 t.assign(m, 0);
 M.assign(m, Row(m));
 rep(i, 0, m) M[i][i] = 1;
  rep(i, 0, sz(A)) {
    auto& w = A[i];
    double s = 0:
    rep(j, 0, d) s += w[j]*w[j];
    if (!cmp(s, 0)) {
     if (cmp(w[m], 0)) return -1;
      continue:
    } // 7462
    double r = sqrt(s);
    if (w[d-1] < 0) r = -r:
   s = sqrt((s + w[d-1]*r)*2);
    w[d-1] += r;
    rep(j, 0, d) w[j] /= s;
    r = w[m] / (w[d-1] * r * 2);
    rep(j, i+1, sz(A)) {
      s = 0;
      rep(k, 0, d) s += A[j][k] * w[k];
      s *= 2;
      rep(k, 0, d) A[j][k] -= s * w[k];
      A[j][m] = s*r;
    } // 69fe
    rep(j, 0, m) {
      s = 0;
      rep(k, 0, d) s += M[k][j] * w[k];
      s *= 2;
      rep(k, 0, d) M[k][j] -= s * w[k];
    } // 692b
```

rep(k, 0, d) s += t[k] * w[k];

s = -r;

s *= 2;

```
rep(k, 0, d) t[k] -= s * w[k];
   d--;
  } // a688
  return d;
} // c093
math/linear rec.h
                                         60be
constexpr 11 MOD = 998244353;
using Poly = vector<11>;
// Compute k-th term of an n-order linear
// recurrence C[i] = sum C[i-j-1]*D[j],
// given C[0..n-1] and D[0..n-1]; O(n^2 \log k)
ll linearRec(const Poly& C,
            const Poly& D, 11 k) {
  int n = sz(D);
  auto mul = [&] (Poly a, Poly b) {
   Poly ret (n*2+1);
    rep(i, 0, n+1) rep(j, 0, n+1)
     ret[i+j] = (ret[i+j] + a[i]*b[j]) % MOD;
    for (int i = n*2; i > n; i--) rep(j, 0, n)
     ret[i-j-1] =
        (ret[i-j-1] + ret[i]*D[j]) % MOD;
    ret.resize(n+1);
   return ret;
  1: // e722
  Poly pol(n+1), e(n+1);
  pol[0] = e[1] = 1;
  for (k++; k; k /= 2) {
   if (k % 2) pol = mul(pol, e);
   e = mul(e, e);
  } // 13af
  11 ret = 0;
  rep(i,0,n) ret = (ret + pol[i+1] \starC[i]) % MOD;
 return ret:
} // 3fd1
                                         58a8
math/linear rec fast.h
#include "polynomial.h"
// Compute k-th term of an n-order linear
// recurrence C[i] = sum C[i-j-1]*D[i],
// given C[0..n-1] and D[0..n-1];
// time: O(n log n log k)
Zp linearRec(const Polv& C.
            const Poly& D, 11 k) {
  Polv f(sz(D)+1, 1);
  rep(i, 0, sz(D)) f[i] = -D[sz(D)-i-1];
  f = pow({0, 1}, k, f);
 Zp ret = 0:
  rep(i, 0, sz(f)) ret += f[i]*C[i];
 return ret:
} // 5b8d
                                         9bf7
math/matrix.h
#include "modular.h"
using Row = vector<Zp>;
using Matrix = vector<Row>;
// Create n x n identity matrix
Matrix ident(int n) {
 Matrix ret(n, Row(n));
```

```
rep(i, 0, n) ret[i][i] = 1;
  return ret;
} // ad1d
// Add matrices
Matrix& operator+= (Matrix& 1, const Matrix& r) {
 rep(i, 0, sz(l)) rep(k, 0, sz(l[0]))
   l[i][k] += r[i][k];
 return 1;
} // b6bf
Matrix operator+ (Matrix 1, const Matrix& r) {
return 1 += r;
} // d9b3
// Subtract matrices
Matrix& operator == (Matrix& 1, const Matrix& r) {
 rep(i, 0, sz(l)) rep(k, 0, sz(l[0]))
   l[i][k] = r[i][k];
  return 1;
} // 90a1
Matrix operator-(Matrix 1, const Matrix& r) {
 return 1 -= r;
} // dc4f
// Multiply matrices
Matrix operator*(const Matrix& 1,
                 const Matrix& r) {
 Matrix ret(sz(1), Row(sz(r[0])));
 rep(i, 0, sz(l)) rep(j, 0, sz(r[0]))
   rep(k, 0, sz(r))
      ret[i][j] += l[i][k] * r[k][j];
  return ret;
} // 52ca
Matrix& operator *= (Matrix& 1, const Matrix& r) {
  return 1 = 1*r;
} // da8a
// Square matrix power; time: O(n^3 * 1g e)
Matrix pow (Matrix a, 11 e) {
 Matrix t = ident(sz(a));
  while (e) {
   if (e % 2) t *= a;
   e /= 2; a *= a;
  1 // 4400
  return t:
} // 65ea
// Transpose matrix
Matrix transpose (const Matrix € m) {
 Matrix ret(sz(m[0]), Row(sz(m)));
 rep(i, 0, sz(m)) rep(j, 0, sz(m[0]))
   ret[j][i] = m[i][j];
 return ret;
1 // 5650
// Transform matrix to echelon form
// and compute its determinant sign and rank.
int echelon(Matrix& A, int& sign) { // O(n^3)
  int rank = 0;
  sign = 1:
  rep(c, 0, sz(A[0])) {
    if (rank >= sz(A)) break;
    rep(i, rank+1, sz(A)) if (A[i][c].x) {
      swap(A[i], A[rank]);
      sign *=-1;
     break;
    } // f98a
    if (A[rank][c].x) {
      rep(i, rank+1, sz(A)) {
        auto mult = A[i][c] / A[rank][c];
        rep(j, 0, sz(A[0]))
```

```
A[i][j] -= A[rank][j] *mult;
      } // f519
      rank++;
   } // 4cd8
 } // 36e9
 return rank;
} // 6882
// Compute matrix rank; time: O(n^3)
#define rank rank
int rank(Matrix A) {
int s; return echelon(A, s);
} // c599
// Compute square matrix determinant; O(n^3)
Zp det (Matrix A) {
 int s; echelon(A, s);
 Zp ret = s:
 rep(i, 0, sz(A)) ret \star= A[i][i];
 return ret;
1 // b252
// Invert square matrix if possible; O(n^3)
// Returns true if matrix is invertible.
bool invert (Matrix& A) {
 int s, n = sz(A);
  rep(i, 0, n) A[i].resize(n*2), A[i][n+i] = 1;
  echelon(A, s);
  for (int i = n; i--;) {
    if (!A[i][i].x) return 0;
    auto mult = A[i][i].inv();
    each(k, A[i]) k *= mult;
    rep(k, 0, i) rep(j, 0, n)
      A[k][n+j] -= A[i][n+j]*A[k][i];
  } // 1e97
 each(r, A) r.erase(r.begin(), r.begin()+n);
 return 1;
} // 65b9
math/miller rabin.h
                                         7005
#include "modular64.h"
// Miller-Rabin primality test
// time O(k*lg^2 n), where k = number of bases
// Deterministic for p <= 10^9
// constexpr 11 BASES[] = {
// 336781006125, 9639812373923155
// }; // d41d
// Deterministic for p <= 2^64
constexpr 11 BASES[] = {
 2.325.9375.28178.450775.9780504.1795265022
}; // b8e0
bool isPrime(ll p) {
 if (p <= 2) return p == 2;
  if (p%2 == 0) return 0:
  11 d = p-1, t = 0;
  while (d%2 == 0) d /= 2, t++;
  each(a, BASES) if (a%p) {
   // 11 a = rand() % (p-1) + 1;
    11 b = modPow(a%p, d, p);
    if (b == 1 || b == p-1) continue;
    rep(i, 1, t)
     if ((b = modMul(b, b, p)) == p-1) break;
   if (b != p-1) return 0;
  } // 9342
  return 1:
} // bec2
```

```
math/modiny precompute.h
                                         2882
constexpr 11 MOD = 234567899;
vector<ll> modInv(MOD); // You can lower size
// Precompute modular inverses; time: O(MOD)
void initModInv() {
 modInv[1] = 1;
 rep(i, 2, sz(modInv)) modInv[i] =
    (MOD - (MOD/i) ★ modInv[MOD%i]) % MOD;
} // 22c1
math/modular.h
                                         72a7
// Modulus often seen on Codeforces:
constexpr int MOD = 998244353:
// Some big prime: 15*(1<<27)+1 ~ 2*10^9
11 modInv(11 a, 11 m) { // a^(-1) mod m
 if (a == 1) return 1:
 return ((a - modInv(m%a, a)) *m + 1) / a;
} // c437
11 modPow(11 a, 11 e, 11 m) { // a^e mod m
 11 t = 1 % m;
 while (e) {
   if (e % 2) t = t*a % m;
   e /= 2; a = a*a % m;
 } // 66ca
 return t:
1 // 1973
// Wrapper for modular arithmetic
struct Zp {
 11 x; // Contained value, in range [0; MOD-1]
 Zp() : x(0) {}
 Zp(11 a) : x(a\%MOD) { if (x < 0) x += MOD; }
  #define OP(c,d) Zp\& operator c##=(Zp r) \{ \
     x = x d; return *this; } \
    Zp operator c(Zp r) const { \
      Zp t = *this; return t c##= r; } // e986
 OP(+, +r.x - MOD*(x+r.x >= MOD));
 OP(-, -r.x + MOD*(0 > x-r.x));
 OP(*, *r.x % MOD);
 OP(/, *r.inv().x % MOD);
 Zp operator-() const { return Zp()-*this; }
  // For composite modulus use modInv, not pow
 Zp inv() const { return pow(MOD-2); }
 Zp pow(ll e) const{ return modPow(x,e,MOD); }
 void print() { cerr << x; } // For deb()</pre>
1: // f730
// Extended Euclidean Algorithm
ll egcd(ll a, ll b, ll& x, ll& y) {
 if (!a) return x=0, y=1, b;
 ll d = egcd(b%a, a, y, x);
 x = b/a*v;
 return d;
1 // 23c8
math/modular64.h
                                         4b73
// Modular arithmetic for modulus < 2^62
11 modAdd(11 x, 11 y, 11 m) {
 x += y;
 return x < m ? x : x-m;
} // b653
ll modSub(ll x, ll y, ll m) {
 x -= y;
 return x >= 0 ? x : x+m;
} // b073
```

```
// About 4x slower than normal modulo
11 modMul(11 a, 11 b, 11 m) {
  11 c = 11((long double) a * b / m);
  11 r = (a*b - c*m) % m;
  return r < 0 ? r+m : r:
} // 1815
11 modPow(11 x, 11 e, 11 m) {
 11 t = 1;
  while (e) {
   if (e \& 1) t = modMul(t, x, m);
   e >>= 1;
   x = modMul(x, x, m);
  } // bd61
  return t;
} // c8ba
                                         845b
math/modular generator.h
#include "modular.h" // modPow
// Get unique prime factors of n; O(sqrt n)
vector<ll> factorize(ll n) {
  vector<ll> fac;
  for (11 i = 2; i*i <= n; i++) {
   if (n\%i == 0) {
      while (n%i == 0) n /= i;
     fac.pb(i);
   } // 6069
  } // a0cc
  if (n > 1) fac.pb(n);
  return fac;
} // 4a2a
// Find smallest primitive root mod n;
// time: O(sqrt(n) + q*log^2 n)
// Returns -1 if generator doesn't exist.
// For n \le 10^7 smallest generator is \le 115.
// You can use faster factorization algorithm
// to get rid of sqrt(n).
11 generator(11 n) {
 if (n \le 1 \mid | (n > 4 \&\& n%4 == 0)) return -1;
  vector<ll> fac = factorize(n);
  if (sz(fac) > (fac[0] == 2)+1) return -1;
  11 phi = n:
  each (p, fac) phi = phi / p * (p-1);
  fac = factorize(phi);
  for (ll q = 1;; q++) if (__qcd(q, n) == 1) {
   each (f, fac) if (modPow(q, phi/f, n) == 1)
     goto nxt;
    return q;
   nxt::
  } // db24
1 // 7641
math/modular log.h
                                         ac62
#include "modular.h" // modInv
// Baby-step giant-step algorithm; O(sgrt(p))
// Finds smallest x such that a^x = b \pmod{p}
// or returns -1 if there's no solution.
11 dlog(11 a, 11 b, 11 p) {
  int m = int(min(ll(sqrt(p))+2, p-1));
  unordered map<11, int> small;
  11 t = 1:
```

rep(i, 0, m) {

int& k = small[t];

if (!k) k = i+1;

```
t = t*a % p;
 } // f1d0
 t = modInv(t, p);
 rep(i, 0, m) {
   int j = small[b];
   if (j) return i*ll(m) + j - 1;
   b = b \star t % p;
 } // c7ed
 return -1:
} // 5c26
math/modular sgrt.h
                                          db16
#include "modular.h" // modPow
// Tonelli-Shanks algorithm for modular sqrt
// modulo prime; O(lq^2 p), O(lq p) for most p
// Returns -1 if root doesn't exists or else
// returns square root x (the other one is -x).
11 modSgrt(11 a, 11 p) {
 a %= p;
 if (a < 0) a += p;
 if (a <= 1) return a;</pre>
 if (modPow(a, p/2, p) != 1) return -1;
 if (p%4 == 3) return modPow(a, p/4+1, p);
 ll s = p-1, n = 2;
 int r = 0, i:
 while (s%2 == 0) s /= 2, r++;
 while (modPow(n, p/2, p) != p-1) n++;
 ll x = modPow(a, (s+1)/2, p);
 11 b = modPow(a, s, p), q = modPow(n, s, p);
 for (;; r = j) {
   11 t = b;
   for (j = 0; j < r && t != 1; j++)
    t = t*t % p;
   if (!i) return x;
   ll qs = modPow(q, 1LL \ll (r-j-1), p);
   q = qs*qs % p;
   x = x*qs % p;
   b = b \star g % p;
 } // f83f
} // 7a97
math/montgomery.h
#include "modular.h" // modInv
// Montgomery modular multiplication
// MOD < MG_MULT, gcd (MG_MULT, MOD) must be 1
// Don't use if modulo is constexpr; UNTESTED
constexpr 11 MG SHIFT = 32;
constexpr ll MG MULT = 1LL << MG SHIFT;</pre>
constexpr ll MG_MASK = MG_MULT - 1;
const 11 MG_INV = MG_MULT-modInv(MOD, MG_MULT);
// Convert to Montgomery form
11 MG(11 x) { return (x*MG MULT) % MOD; }
// Montgomery reduction
// redc(mg * mg) = Montgomery-form product
11 redc(11 x) {
 11 g = (x * MG INV) & MG MASK;
 x = (x + q*MOD) >> MG SHIFT;
 return (x >= MOD ? x-MOD : x);
} // d0f5
                                          474f
math/nimber.h
// Nimbers are defined as sizes of Nim heaps.
// Operations on nimbers are defined as:
```

```
// a+b = mex(\{a'+b : a' < a\} u \{a+b' : b' < b\})
       // ab = mex({a'b+ab'+a'b' : a' < a, b' < b})
       // Nimbers smaller than M = 2^2k form a field.
       // Addition is equivalent to xor, meanwhile
       // multiplication can be evaluated
       // in O(lg^2 M) after precomputing.
       using ull = uint64_t;
       ull nbuf[64][64]; // Nim-products for 2^i * 2^j
       // Multiply nimbers; time: O(lg^2 M)
       // WARNING: Call initNimMul() before using.
       ull nimMul(ull a, ull b) {
        ull ret = 0;
         for (ull s = a; s; s &= (s-1))
           for (ull t = b; t; t &= (t-1))
             ret ^= nbuf[__builtin_ctzll(s)]
                         [__builtin_ctzll(t)];
         return ret;
        // Initialize nim-products lookup table
       void initNimMul() {
         rep(i, 0, 64)
           nbuf[i][0] = nbuf[0][i] = 1ull << i;
         rep(b, 1, 64) rep(a, 1, b+1) {
           int i = 1 << (63 - __builtin_clzll(a));</pre>
           int j = 1 << (63 - __builtin_clzll(b));</pre>
           ull t = nbuf[a-i][b-j];
           if (i < j)
             t = nimMul(t, 1ull << i) << i;
             t = nimMul(t, 1ull << (i-1)) ^ (t << i);
           nbuf[a][b] = nbuf[b][a] = t;
        } // ca24
       } // 1811
       // Compute a^e under nim arithmetic; O(1g^3 M)
       // WARNING: Call initNimMul() before using.
       ull nimPow(ull a, ull e) {
         ull t = 1:
         while (e) {
           if (e % 2) t = nimMul(t, a);
           e /= 2; a = nimMul(a, a);
         1 // da53
        return t:
a4ba | } // c06c
       // Compute inverse of a in 2^64 nim-field:
       // time: 0(1g^3 M)
       // WARNING: Call initNimMul() before using.
       ull nimInv(ull a) {
        return nimPow(a, ull(-2));
       } // c6d9
       // If you need to multiply many nimbers by
        // the same value you can use this to speedup.
        // WARNING: Call initNimMul() before using.
       struct NimMult {
         ull M[64] = \{0\};
         // Initialize lookup; time: O(lg^2 M)
         NimMult(ull a) {
           for (ull t=a; t; t &= (t-1)) rep(i, 0, 64)
             M[i] ^= nbuf[__builtin_ctzll(t)][i];
         } // ea88
         // Multiply by b; time: O(lg M)
         ull operator() (ull b) {
           ull ret = 0:
           for (ull t = b; t; t &= (t-1))
             ret ^= M[__builtin_ctzll(t)];
```

```
return ret:
 } // e480
}; // 1b80
                                          8703
math/phi large.h
#include "pollard rho.h"
// Compute Euler's totient of large numbers
// time: O(n^{(1/4)}) \leftarrow factorization
ll phi(ll n) {
 each (p, factorize (n)) n = n / p.x * (p.x-1);
 return n:
} // 798e
math/phi precompute.h
                                          728b
Vi phi(1e7+1):
// Precompute Euler's totients; time: O(n la n)
void calcPhi() {
 iota(all(phi), 0);
  rep(i, 2, sz(phi)) if (phi[i] == i)
    for (int j = i; j < sz(phi); j += i)</pre>
      phi[j] = phi[j] / i * (i-1);
} // 3c65
math/phi prefix sum.h
                                          39a5
#include "phi_precompute.h"
constexpr int MOD = 998244353;
vector<11> phiSum; // [k] = sum \ from \ 0 \ to \ k-1
// Precompute Euler's totient prefix sums
// for small values; time: O(n lg n)
void calcPhiSum() {
  calcPhi():
  phiSum.resize(sz(phi)+1);
  rep(i, 0, sz(phi))
    phiSum[i+1] = (phiSum[i] + phi[i]) % MOD;
} // bcf5
// Get prefix sum of phi(0) + ... + phi(n-1).
// WARNING: Call calcPhiSum first!
// For MOD > 4*10^9, answer will overflow.
ll getPhiSum(ll n) { // time: O(n^{2/3})
  static unordered_map<11, 11> big;
  if (n < sz(phiSum)) return phiSum[n];</pre>
  if (big.count(--n)) return big[n];
  11 ret = (n\%2 ? n\%MOD * ((n+1)/2 % MOD)
                : n/2%MOD * (n%MOD+1)) % MOD;
  for (ll s, i = 2; i <= n; i = s+1) {
    s = n / (n/i);
    ret -= (s-i+1) %MOD*getPhiSum(n/i+1) % MOD;
  return big[n] = ret = (ret%MOD + MOD) % MOD;
} // 1d5f
math/pi large.h
                                          fcbd
constexpr int MAX P = 1e7;
vector<ll> pis, prl;
// Precompute prime counting function
// for small values; time: O(n lg lg n)
void initPi() {
 pis.assign(MAX_P+1, 1);
  pis[0] = pis[1] = 0;
  for (int i = 2; i*i <= MAX_P; i++)</pre>
   if (pis[i])
      for (int j = i*i; j <= MAX_P; j += i)</pre>
        pis[j] = 0;
```

```
rep(i, 1, sz(pis)) {
    if (pis[i]) prl.pb(i);
    pis[i] += pis[i-1];
  } // 0672
1 // 6d92
ll partial(ll x, ll a) {
  static vector<unordered map<11, 11>> big;
  big.resize(sz(prl));
  if (!a) return (x+1) / 2;
  if (big[a].count(x)) return big[a][x];
  ll ret = partial(x, a-1)
    - partial(x / prl[a], a-1);
  return big[a][x] = ret;
} // 774f
// Count number of primes <= x;</pre>
// \text{ time: } O(n^{(2/3)} * log(n)^{(1/3)})
// Set MAX P to be > sqrt(x) and call initPi
// before using!
ll pi(ll x) {
  static unordered_map<11, 11> big;
  if (x < sz(pis)) return pis[x];</pre>
  if (big.count(x)) return big[x];
  while (prl[a]*prl[a]*prl[a]*prl[a] < x) a++;</pre>
  11 ret = 0, b = --a;
  while (++b < sz(prl) && prl[b]*prl[b] < x) {</pre>
    11 w = x / prl[b];
    ret -= pi(w);
    for (ll j = b; prl[j]*prl[j] <= w; j++)</pre>
      ret -= pi(w / prl[j]) - j;
  } // a584
  ret += partial(x, a) + (b+a-1)*(b-a)/2;
  return big[x] = ret;
} // eald
                                           7fc0
math/pi large precomp.h
#include "sieve.h"
// Count primes in given interval
// using precomputed table.
// Set MAX P to sgrt (MAX N) and run sieve()!
// Precomputed table will contain N BUCKETS
// elements - check source size limit.
constexpr ll MAX N = 1e11+1;
constexpr 11 N BUCKETS = 10000;
constexpr 11 BUCKET SIZE = (MAX N/N BUCKETS)+1;
constexpr ll precomputed[] = {/* ... */};
ll sieveRange(ll from, ll to) {
  bitset < BUCKET SIZE > elems:
  from = max(from, 2LL);
  to = max(from, to);
  each (p. primesList) {
   ll c = max((from+p-1) / p, 2LL);
    for (11 i = c*p; i < to; i += p)
      elems.set(i-from);
  } // a29f
  return to-from-elems.count():
1 // c646
// Run once on local computer to precompute
// table. Takes about 10 minutes for n = 1e11.
// Sanity check (for default params):
// 664579, 606028, 587253, 575795, ...
void localPrecompute() {
```

```
for (11 i = 0; i < MAX_N; i += BUCKET_SIZE) {</pre>
   11 to = min(i+BUCKET_SIZE, MAX_N);
    cout << sieveRange(i, to) << ',' << flush;</pre>
 } // f6a7
  cout << endl;
} // 2b1e
// Count primes in [from; to) using table.
// O(N BUCKETS + BUCKET SIZE*lq lq n + sqrt(n))
ll countPrimes(ll from, ll to) {
 11 bFrom = from/BUCKET_SIZE+1,
    bTo = to/BUCKET SIZE;
  if (bFrom > bTo) return sieveRange(from, to);
  11 ret = accumulate (precomputed+bFrom.
                      precomputed+bTo, 0);
  ret += sieveRange (from, bFrom*BUCKET_SIZE);
  ret += sieveRange (bTo*BUCKET SIZE, to);
  return ret;
} // cced
math/pollard rho.h
                                          ef01
#include "modular64.h"
#include "miller rabin.h"
using Factor = pair<11, int>;
void rho(vector<11>& out, 11 n) {
 if (n <= 1) return;</pre>
  if (isPrime(n)) out.pb(n);
  else if (n%2 == 0) rho(out,2), rho(out,n/2);
  else for (11 a = 2;; a++) {
    11 x = 2, y = 2, d = 1;
    while (d == 1) {
      x = modAdd(modMul(x, x, n), a, n);
      y = modAdd (modMul (y, y, n), a, n);
     y = modAdd(modMul(y, y, n), a, n);
      d = \underline{gcd(abs(x-y), n)};
    1 // 3378
    if (d != n) return rho(out,d), rho(out, n/d);
 } // 047e
} // ba89
// Pollard's rho factorization algorithm
// Las Vegas version; time: n^(1/4)
// Returns pairs (prime, power), sorted
vector<Factor> factorize(ll n) {
 vector<Factor> ret;
  vector<11> raw:
  rho(raw, n);
  sort(all(raw));
  each(f, raw) {
    if (ret.empty() | ret.back().x != f)
      ret.pb({ f, 1 });
    else
      ret.back().y++;
  } // 2ab1
  return ret:
} // 471c
math/polynomial.h
                                          a449
#include "modular.h"
#include "fft mod.h"
using Poly = vector<Zp>;
// Cut off trailing zeroes; time: O(n)
void norm(Poly& P) {
  while (!P.empty() && !P.back().x)
   P.pop_back();
} // 8a8a
// Evaluate polynomial at x; time: O(n)
```

```
Zp eval(const Poly& P, Zp x) {
  Zp n = 0, y = 1;
  each(a, P) n += a*y, y *= x;
  return n;
} // d865
// Add polynomial; time: O(n)
Poly& operator+=(Poly& 1, const Poly& r) {
1.resize(max(sz(1), sz(r)));
  rep(i, 0, sz(r)) l[i] += r[i];
  norm(1);
  return 1;
} // 656e
Poly operator+(Poly 1, const Poly& r) {
 return 1 += r;
} // d9b3
// Subtract polynomial; time: O(n)
Poly& operator -= (Poly& 1, const Poly& r) {
 1.resize(max(sz(1), sz(r)));
  rep(i, 0, sz(r)) l[i] -= r[i];
  norm(1);
  return 1:
1 // c68b
Poly operator-(Poly 1, const Poly& r) {
 return 1 -= r:
// Multiply by polynomial; time: O(n lg n)
Poly& operator *= (Poly& 1, const Poly& r) {
 if (\min(sz(1), sz(r)) < 50) {
    // Naive multiplication
    Poly p(sz(1)+sz(r));
    rep(i, 0, sz(l)) rep(j, 0, sz(r))
      p[i+j] += l[i]*r[j];
    1.swap(p);
  } else {
    // FFT multiplication
    // Choose appropriate convolution method,
    // see fft_mod.h and fft_complex.h
    using v = vector<11>;
    convolve<MOD, 62>(*(v*)&l, *(const v*)&r);
  } // 30c9
 norm(1);
 return 1;
1 // e8b3
Poly operator*(Poly 1, const Poly& r) {
 return | *= r:
} // 2de3
// Compute inverse series mod x^n; O(n lq n)
// Requires P(0) != 0.
Poly invert (const Poly € P, int n) {
 assert(!P.empty() && P[0].x);
  Poly tmp{P[0]}, ret = {P[0].inv()};
  for (int i = 1; i < n; i *= 2) {
    rep(j, i, min(i\star2, sz(P))) tmp.pb(P[j]);
    (ret \star= Poly{2} - tmp\starret).resize(i\star2);
  } // 904e
  ret.resize(n);
  return ret;
1 // 9293
// Floor division by polynomial; O(n lg n)
Poly& operator/=(Poly& 1, Poly r) {
 norm(1); norm(r);
  int d = sz(1)-sz(r)+1;
  if (d <= 0) return l.clear(), l;</pre>
  reverse (all(1));
  reverse (all(r));
  l.resize(d);
```

```
1 *= invert(r, d);
 l.resize(d);
 reverse (all(1));
 return 1;
} // cf5e
Poly operator/(Poly 1, const Poly& r) {
 return 1 /= r;
} // 152d
// Remainder modulo a polynomial; O(n lg n)
Poly operator% (const Poly& 1, const Poly& r) {
 return 1 - r*(1/r);
} // 4fc8
Poly& operator%=(Poly& 1, const Poly& r) {
 return 1 -= r*(1/r);
} // 80bb
// Compute a^e mod x^n, where a is polynomial;
// time: O(n log n log e)
Poly pow(Poly a, ll e, int n) {
 Poly t = \{1\};
 while (e) {
   if (e % 2) (t *= a).resize(n);
   e /= 2; (a *= a).resize(n);
 } // d0c6
 norm(t);
 return t;
} // ada1
// Compute a^e mod m, where a and m are
// polynomials: time: O(|m| log |m| log e)
Poly pow (Poly a, 11 e, const Poly € m) {
 Poly t = \{1\};
 while (e) {
    if (e % 2) t = t*a % m;
    e /= 2; a = a*a % m;
 } // 66ca
 return t;
} // 6f9c
// Derivate polynomial; time: O(n)
Poly derivate (Poly P) {
 if (!P.empty()) {
    rep(i, 1, sz(P)) P[i-1] = P[i]*i;
    P.pop_back();
 } // bd78
 return P;
} // c6c5
// Integrate polynomial; time: O(n)
Poly integrate (Poly P) {
 if (!P.empty()) {
    P.pb(0):
    for (int i = sz(P); --i;) P[i] = P[i-1]/i;
   P[0] = 0;
 } // eec1
 return P;
} // e2f3
// Compute ln(P) mod x^n; time: O(n log n)
Poly log(const Poly& P, int n) {
 Poly a = integrate (derivate (P) *invert (P, n));
 a.resize(n);
 return a:
} // 5d6b
// Compute exp(P) mod x^n; time: O(n lg n)
// Requires P(0) = 0.
Poly exp(Poly P, int n) {
 Poly tmp{P[0]+1}, ret = {1};
 for (int i = 1; i < n; i *= 2) {
```

```
rep(j, i, min(i\star2, sz(P))) tmp.pb(P[j]);
    (ret \star= (tmp - log(ret, i\star2))).resize(i\star2);
  } // c28a
  ret.resize(n);
  return ret;
} // bd42
// Compute sqrt(P) mod x^n; time: O(n log n)
#include "modular sgrt.h"
bool sgrt (Poly& P, int n) {
  norm(P);
  if (P.empty()) return P.resize(n), 1;
  int tail = 0;
  while (!P[tail].x) tail++;
  if (tail % 2) return 0:
  11 sq = modSqrt(P[tail].x, MOD);
  if (sq == -1) return 0;
  Poly tmp{P[tail]}, ret = {sq};
  for (int i = 1; i < n - tail/2; i *= 2) {
   rep(j, i, min(i\star2, sz(P)-tail))
     tmp.pb(P[tail+j]);
    (ret += tmp * invert(ret, i*2)).resize(i*2);
   each (e, ret) e /= 2;
  } // 2d41
  P.resize(tail/2);
  P.insert(P.end(), all(ret));
  P.resize(n);
  return 1:
} // b9b3
// Compute polynomial P(x+c); time: O(n lq n)
Poly shift (Poly P, Zp c) {
  int n = sz(P);
  Polv O(n, 1);
  Zp fac = 1:
  rep(i, 1, n) {
   P[i] *= (fac *= i);
   Q[n-i-1] = Q[n-i] * c / i;
  } // 1c20
  P *= 0;
  if (sz(P) < n) return {}:
  P.erase(P.begin(), P.begin()+n-1);
  fac = 1:
  rep(i, 1, n) P[i] /= (fac *= i);
  return P.
} // b11f
// Compute values P(x^0), ..., P(x^{n-1});
// time: O(n la n)
Poly chirpz (Poly P, Zp x, int n) {
  int k = sz(P);
  Polv O(n+k);
  rep(i, 0, n+k) Q[i] = x.pow(i*(i-1)/2);
  rep(i, 0, k) P[i] /= O[i];
  reverse(all(P));
  P *= Q;
  rep(i, 0, n) P[i] = P[k+i-1] / Q[i];
  P.resize(n);
 return P;
} // 4e8b
// Evaluate polynomial P in given points;
// time: O(n 1g^2 n)
Poly eval(const Poly € P, Poly points) {
  int len = 1;
  while (len < sz(points)) len ★= 2;
  vector<Poly> tree(len*2, {1});
  rep(i, 0, sz(points))
```

```
tree[len+i] = {-points[i], 1};
  for (int i = len: --i:)
   tree[i] = tree[i\star2] \star tree[i\star2+1];
  tree[0] = P;
  rep(i, 1, len*2)
   tree[i] = tree[i/2] % tree[i];
  rep(i, 0, sz(points)) {
    auto& vec = tree[len+i];
   points[i] = vec.empty() ? 0 : vec[0];
  } // c1c2
  return points;
1 // 69b0
// Given n points (x, f(x)) compute n-1-degree
// polynomial f that passes through them;
// time: O(n la^2 n)
// For O(n^2) version see polynomial_interp.h.
Poly interpolate (const vector < pair < Zp, Zp>>& P) {
 int len = 1;
  while (len < sz(P)) len \star= 2;
  vector<Poly> mult(len*2, {1}), tree(len*2);
  rep(i, 0, sz(P))
   mult[len+i] = {-P[i].x, 1};
  for (int i = len; --i;)
    mult[i] = mult[i*2] * mult[i*2+1];
  tree[0] = derivate(mult[1]);
  rep(i, 1, len*2)
   tree[i] = tree[i/2] % mult[i];
  rep(i, 0, sz(P))
   tree[len+i][0] = P[i].y / tree[len+i][0];
  for (int i = len: --i;)
   tree[i] = tree[i\star2]\starmult[i\star2+1]
           + mult[i*2]*tree[i*2+1];
  return tree[1]:
} // b706
math/polynomial interp.h
                                          a4cc
// Interpolate set of points (i, vec[i])
// and return it evaluated at x; time: O(n)
template<class T>
T polyExtend(vector<T>& vec, T x) {
  int n = sz(vec);
  vector<T> fac(n, 1), suf(n, 1);
  rep(i, 1, n) fac[i] = fac[i-1] \star i;
  for (int i=n; --i;) suf[i-1] = suf[i]*(x-i);
 T pref = 1, ret = 0;
  rep(i, 0, n) {
   T d = fac[i] * fac[n-i-1] * ((n-i) *2*2-1);
    ret += vec[i] * suf[i] * pref / d;
   pref \star = x-i;
  1 // 681d
  return ret;
} // dd92
// Given n points (x, f(x)) compute n-1-degree
// polynomial f that passes through them;
// time: 0(n^2)
// For O(n lg^2 n) version see polynomial.h
template<class T>
vector<T> polyInterp(vector<pair<T, T>> P) {
 int n = sz(P):
  vector<T> ret(n), tmp(n);
 T last = 0;
  tmp[0] = 1;
```

```
rep(k, 0, n-1) rep(i, k+1, n)
   P[i].y = (P[i].y-P[k].y) / (P[i].x-P[k].x);
 rep(k, 0, n) rep(i, 0, n) {
   ret[i] += P[k].y * tmp[i];
   swap(last, tmp[i]);
   tmp[i] = last * P[k].x;
 } // af1c
 return ret;
} // 7c2c
math/sieve.h
                                          3f3d
constexpr int MAX P = 1e6;
bitset<MAX_P+1> primes;
Vi primesList;
// Erathostenes sieve; time: O(n lg lg n)
void sieve() {
 primes.set();
 primes.reset(0);
 primes.reset(1);
 for (int i = 2; i*i <= MAX_P; i++)</pre>
   if (primes[i])
     for (int j = i*i; j <= MAX_P; j += i)</pre>
        primes.reset(j);
 rep(i, 0, MAX_P+1) if (primes[i])
   primesList.pb(i);
} // d5ca
math/sieve factors.h
                                          312d
constexpr int MAX P = 1e6;
Vi factor (MAX_P+1);
// Erathostenes sieve with saving smallest
// factor for each number; time: O(n lg lg n)
void sieve() {
 for (int i = 2; i*i <= MAX_P; i++)</pre>
   if (!factor[i])
     for (int j = i*i; j <= MAX P; j += i)</pre>
        if (!factor[j])
         factor[j] = i;
 rep(i,0,MAX P+1) if (!factor[i]) factor[i]=i;
} // 82b6
// Factorize n <= MAX_P; time: O(lq n)</pre>
// Returns pairs (prime, power), sorted
vector<Pii> factorize(ll n) {
 vector<Pii> ret:
 while (n > 1) {
   int f = factor[n];
   if (ret.empty() | ret.back().x != f)
     ret.pb({ f, 1 });
     ret.back().y++;
   n /= f;
 1 // 664c
 return ret;
1 // bb65
math/sieve segmented.h
                                          849b
constexpr int MAX P = 1e9;
bitset<MAX_P/2+1> primes; // Only odd numbers
// Cache-friendly Erathostenes sieve
// ~1.5s on Intel Core i5 for MAX_P = 10^9
// Memory usage: MAX_P/16 bytes
void sieve() {
 constexpr int SEG SIZE = 1<<18;</pre>
```

int pSqrt = int(sqrt(MAX_P)+0.5);

```
vector<Pii> dels:
  primes.set();
  primes.reset(0);
  for (int i = 3; i <= pSqrt; i += 2) {</pre>
    if (primes[i/2]) {
      int ;
      for (j = i*i; j <= pSqrt; j += i*2)</pre>
        primes.reset(j/2);
      dels.pb(\{ i, j/2 \});
    } // 9e62
 } // ff49
  for (int seg = pSgrt/2;
       seq <= sz(primes); seq += SEG_SIZE) {</pre>
    int lim = min(seg+SEG_SIZE, sz(primes));
    each(d, dels) for (;d.y < lim; d.y += d.x)
      primes.reset(d.v);
 } // 97ae
} // 6456
bool isPrime(int x) {
 return x == 2 \mid \mid (x \cdot 2 \cdot \xi \cdot primes[x/2]);
} // 422c
math/simplex.h
                                          c4cf
using dbl = double;
using Row = vector<dbl>;
using Matrix = vector<Row>;
#define lti(X) if (s == -1)
 mp(X[j], N[j]) < mp(X[s], N[s])) s = j
// Simplex algorithm; time: O(nm * pivots)
// Given m x n matrix A, vector b of length m,
// vector c of length n solves the following:
// maximize c^T x, Ax <= b, x >= 0
// Output vector `x` contains optimal solution
// or some feasible solution in unbounded case.
// Returns objective value if bounded.
// +inf if unbounded, and -inf if no solution.
// You can test if double is inf using `isinf`.
// PARTIALLY TESTED
dbl simplex (const Matrix& A,
            const Row& b, const Row& c,
            Row& x, dbl eps = 1e-8) {
  int m = sz(b), n = sz(c);
  x.assign(n, 0);
  if (!n) return
    *min_element(all(b)) < -eps ? -1/.0 : 0;
  Vi N(n+1), B(m);
  Matrix D(m+2, Row(n+2)):
  auto pivot = [&](int r, int s) {
    dbl inv = 1 / D[r][s];
    rep(i, 0, m+2)
      if (i != r && abs(D[i][s]) > eps) {
        dbl tmp = D[i][s] * inv;
        rep(j,0,n+2) D[i][j] -= D[r][j] * tmp;
        D[i][s] = D[r][s] * tmp;
      } // 5281
    each(k, D[r]) k \star= inv;
    each(k, D) k[s] \star= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  }; // f56b
  auto solve = [&](int phase) {
    for (int y = m+phase-1;;) {
      int s = -1, r = -1;
      rep(j, 0, n+1)
```

```
if (N[j] != -phase) ltj(D[y]);
      if (D[y][s] >= -eps) return 1;
      rep(i, 0, m)
       if (D[i][s] > eps && (r == -1 ||
          mp(D[i][n+1] / D[i][s], B[i]) <
          mp(D[r][n+1] / D[r][s], B[r]))) r=i;
     if (r == -1) return 0;
     pivot(r, s);
   } // 3bef
  }; // 614a
  rep(i, 0, m) {
   copy(all(A[i]), D[i].begin());
   B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i];
  rep(j, 0, n) D[m][N[j] = j] = -c[j];
  N[n] = -1; D[m+1][n] = 1;
  int r = 0;
  rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
  if (D[r][n+1] < -eps) {</pre>
   pivot(r, n);
   if (!solve(2) || D[m+1][n+1] < -eps)
     return -1/.0;
    rep(i, 0, m) if (B[i] == -1) {
     int s = 0;
     rep(j, 1, n+1) ltj(D[i]);
     pivot(i, s);
   } // 78fd
  } // b52b
 bool ok = solve(1);
  rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
  return ok ? D[m][n+1] : 1/.0;
} // fe6f
structures/bitset plus.h
                                         6737
// Undocumented std::bitset features:
// - Find first() - returns first bit = 1 or N
// - _Find_next(i) - returns first bit = 1
                    after i-th bit
                    or N if not found
// Bitwise operations for vector<bool>
// UNTESTED
#define OP(x) vector<bool>& operator x##=(
   vector<bool>& 1, const vector<bool>& r) {
  assert(sz(1) == sz(r));
  auto a = 1.begin(); auto b = r.begin();
  while (a<1.end()) *a._M_p++ x##= *b._M_p++; \
  return 1; } // f164
OP (&) OP (|) OP (^)
structures/fenwick tree.h
                                         b467
// Fenwick tree (BIT tree); space: O(n)
// Default version: prefix sums
struct Fenwick {
  using T = int;
 T ID = 0;
 T f(T a, T b) { return a+b; }
  vector<T> s;
 Fenwick(int n = 0) : s(n, ID) {}
  // A[i] = f(A[i], v); time: O(lg n)
  void modify(int i, T v) {
   for (; i < sz(s); i |= i+1) s[i]=f(s[i],v);
 } // a047
 // Get f(A[0], ..., A[i-1]); time: O(lg n)
 T query(int i) {
   T v = ID;
                                                  bool join(int i, int j) {
```

```
for (; i > 0; i \in [i-1]) v = f(v, s[i-1]);
   return v;
 } // 9810
  // Find smallest i such that
 // f(A[0],...,A[i-1]) >= val; time: O(lq n)
  // Prefixes must have non-descreasing values.
 int lowerBound(T val) {
   if (val <= ID) return 0;</pre>
   int i = -1, mask = 1;
    while (mask \leq sz(s)) mask \star= 2;
   T \circ ff = TD:
    while (mask /= 2) {
     int k = mask+i;
     if (k < sz(s)) {
       T x = f(off, s[k]);
        if (val > x) i=k, off=x;
     } // de7f
   } // 929c
   return i+2;
 } // 4be9
): // 1470
structures/fenwick tree 2d.h 4945
// Fenwick tree 2D (BIT tree 2D); space: O(n*m)
// Default version: prefix sums 2D
// Change s to hashmap for O(q lg^2 n) memory
struct Fenwick2D {
 using T = int;
 T ID = 0:
 T f(T a, T b) { return a+b; }
 vector<T> s;
 int w, h;
 Fenwick2D(int n = 0, int m = 0)
   : s(n*m, ID), w(n), h(m) {}
  // A[i,j] = f(A[i,j], v); time: O(1g^2 n)
 void modify(int i, int j, T v) {
   for (; i < w; i |= i+1)
      for (int k = j; k < h; k = k+1)
        s[i*h+k] = f(s[i*h+k], v);
 1 // d46b
  // Query prefix; time: O(1g^2 n)
 T query(int i, int j) {
   T v = ID;
   for (; i>0; i&=i-1)
     for (int k = j; k > 0; k \le k-1)
       v = f(v, s[i*h+k-h-1]);
   return v;
 } // 08cf
}; // 36b4
structures/find union.h
// Disjoint set data structure; space: O(n)
// Operations work in amortized O(alfa(n))
struct FAU {
 Vi G;
 FAU(int n = 0) : G(n, -1) {}
  // Get size of set containing i
 int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
 int find(int i) {
   return G[i] < 0 ? i : G[i] = find(G[i]);</pre>
 } // 5bc1
  // Union sets containing i and i
```

```
i = find(i); j = find(j);
    if (i == j) return 0;
    if (G[i] > G[j]) swap(i, j);
   G[i] += G[j]; G[j] = i;
    return 1;
 } // c721
}; // 62a4
structures/find union undo.h ce37
// Disjoint set data structure
                                                     S.pb(1);
// with rollback; space: O(n)
// Operations work in O(log(n)) time.
                                                  } // 978e
struct RollbackFAU {
 Vi G:
 vector<Pii> his;
  RollbackFAU(int n = 0) : G(n, -1) {}
  // Get size of set containing i
  int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
                                                  } // b8de
  int find(int i) {
                                                 }; // 1d64
    return G[i] < 0 ? i : find(G[i]);</pre>
  } // e478
  // Current time (for rollbacks)
  int time() { return sz(his); }
  // Rollback all operations after time `t`
  void rollback(int t) {
    for (int i = time(); t < i--;)</pre>
      G[his[i].x] = his[i].y;
   his.resize(t);
 } // 3ef3
  // Union sets containing i and i
 bool join(int i, int j) {
   i = find(i); j = find(j);
    if (i == j) return 0;
    if (G[i] > G[j]) swap(i, j);
    his.pb({i, G[i]});
    his.pb({j, G[j]});
   G[i] += G[j]; G[j] = i;
   return 1;
 } // 1491
1: // da8c
structures/hull offline.h
                                         3030
constexpr 11 INF = 2e18;
// constexpr double INF = 1e30;
// constexpr double EPS = 1e-9;
// MAX of linear functions; space: O(n)
                                                  } // 63b8
// Use if you add lines in increasing `a` order
// Default uncommented version is for int64
struct Hull {
  using T = 11; // Or change to double
  struct Line {
   Ta, b, end:
   T intersect (const Line& r) const {
      // Version for double:
                                                  } // 4f69
      //if (r.a-a < EPS) return b>r.b?INF:-INF;
      //return (b-r.b) / (r.a-a);
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b. d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u%d));
   } // f27f
  }; // 10dc
                                                  } // 7533
  vector<Line> S;
                                                 }; // 037e
  Hull() { S.pb({ 0, -INF, INF }); }
```

```
// Insert f(x) = ax+b; time: amortized O(1)
 void push(T a, T b) {
    Line 1{a, b, INF};
    while (true) {
     T e = S.back().end=S.back().intersect(1);
     if (sz(S) < 2 | | S[sz(S)-2].end < e)
       break;
      S.pop_back();
    } // 044f
  // Query max(f(x)) for each f): time: O(\lg n)
 T query (T x) {
    auto t = *upper_bound(all(S), x,
      [](int 1, const Line& r) {
        return 1 < r.end;</pre>
      1): // de77
    return t.a*x + t.b;
structures/hull online.h
                                          2a7b
constexpr 11 INF = 2e18;
// MAX of linear functions online; space: O(n)
struct Hull {
  static bool modeQ; // Toggles operator< mode</pre>
  struct Line {
    mutable 11 a, b, end;
    ll intersect (const Line& r) const {
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b, d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u^d));
    } // f27f
    bool operator < (const Line& r) const {
     return mode0 ? end < r.end : a < r.a;
    } // cfab
 }; // 6046
  multiset<Line> S;
 Hull() { S.insert({ 0, -INF, INF }); }
  // Updates segment end
 bool update(multiset<Line>::iterator it) {
    auto cur = it++; cur->end = INF;
    if (it == S.end()) return false;
    cur->end = cur->intersect(*it);
    return cur->end >= it->end;
  // Insert f(x) = ax+b; time: O(\lg n)
 void insert(ll a, ll b) {
    auto it = S.insert({ a, b, INF });
    while (update(it)) it = --S.erase(++it);
    rep(i, 0, 2)
      while (it != S.begin() && update(--it))
        update(it = --S.erase(++it));
  // Query max(f(x) for each f): time: O(lq n)
 ll query(ll x) {
   mode0 = 1:
   auto 1 = \starS.upper_bound({ 0, 0, x });
    mode0 = 0;
    return 1.a*x + 1.b;
```

```
bool Hull::mode0 = false;
                                           9b83
structures/li chao tree.h
// Extended Li Chao tree; space: O(n)
// Let F be a family of functions,
// closed under function addition, such that
// for every f != g from the family F
// there exists x such that:
// f(z) \ll g(z) for z \ll x, else f(z) \gg g(z)
// q(z) \le f(z) for z \le x, else q(z) >= f(z).
// Typically F is family of linear functions.
// DS maintains a sequence c[0], \ldots, c[n-1]
// under operations max, add, query
// (see comments below for explanations).
// Configure by modifying:
// - T - type of sequence elements
// - Func - represents function from family F
// - ID_ADD - function f that is neutral
             element for function addition
// - ID MAX - function f that is neutral
             element for function max
// TESTED ON RANDS
struct LiChao {
  using T = 11;
  struct Func {
   T a, b; // a*x + b
    // Evaluate function in point x
   T operator()(int x) const { return a*x+b; }
    // Sum of two functions
   Func operator+ (Func r) const {
     return {a+r.a, b+r.b};
    ) // f911
  }; // Obed
  Func ID_ADD{0, 0}, ID_MAX{0, T(-1e9)};
  vector<Func> val. lazv:
  int len;
  // Initialize tree for n elements; time: O(n)
  LiChao(int n = 0) {
    for (len = 1; len < n; len *= 2);
    val.resize(len*2, ID MAX);
   lazy.resize(len*2, ID_ADD);
  } // c0ba
  void push(int i) {
   if (i < len) rep(j, 0, 2) {</pre>
     lazy[i*2+j] = lazy[i*2+j] + lazy[i];
     val[i*2+j] = val[i*2+j] + lazv[i];
    } // 54fc
   lazy[i] = ID_ADD;
  } // 1777
  // For each x in [vb;ve)
  // set c[x] = max(c[x], f(x));
  // time: O(log^2 n) in general case,
           O(\log n) if \lceil vb; ve \rangle = \lceil 0; len \rangle
  void max(int vb, int ve, Func f,
           int i = 1, int b = 0, int e = -1) {
    if (e < 0) e = len;
    if (vb >= e || b >= ve || i >= len\star2)
     return:
    int m = (b+e) / 2;
   push(i);
    if (b >= vb && e <= ve) {
     auto& g = val[i];
     if (g(m) < f(m)) swap(g, f);</pre>
```

```
if ((q(b) < f(b)) != (q(m) < f(m)))
        max(vb, ve, f, i*2, b, m);
      else
        max(vb, ve, f, i*2+1, m, e);
    } else {
      max(vb, ve, f, i*2, b, m);
      max(vb, ve, f, i*2+1, m, e);
   } // f2c0
 } // dec2
  // For each x in [vb; ve)
  // set c[x] = c[x] + f(x);
 // time: O(log^2 n) in general case,
       O(1) if [vb; ve) = [0; len)
  void add(int vb, int ve, Func f,
          int i = 1, int b = 0, int e = -1) {
    if (e < 0) e = len;
    if (vb >= e || b >= ve) return;
    if (b >= vb && e <= ve) {
      lazy[i] = lazy[i] + f;
      val[i] = val[i] + f;
    } else {
      int m = (b+e) / 2;
      push(i);
      max(b, m, val[i], i*2, b, m);
      max(m, e, val[i], i*2+1, m, e);
      val[i] = ID MAX;
      add(vb, ve, f, i*2, b, m);
      add(vb, ve, f, i*2+1, m, e);
   } // bbe5
  1 // 259f
  // Get value of c[x]; time: O(\log n)
 T query(int x) {
    int i = x+len;
   T ret = val[i](x);
    while (i \neq 2)
      ret = ::max(ret+lazy[i](x), val[i](x));
    return ret;
 } // 0c22
}; // c95e
                                         3e9e
structures/max queue.h
// Queue with max query on contained elements
struct MaxQueue {
 using T = int;
  deque<T> Q, M;
  // Add v to the back; time: amortized O(1)
  void push(T v) {
    while (!M.empty() && M.back() < v)</pre>
      M.pop_back();
   M.pb(v); Q.pb(v);
  } // 57a2
  // Pop from the front; time: O(1)
  void pop() {
   if (M.front() == Q.front()) M.pop_front();
   Q.pop_front();
 } // 101c
  // Get max element value; time: O(1)
 T max() const { return M.front(); }
1: // b6c4
structures/pairing_heap.h
                                         b2a7
// Pairing heap implementation; space O(n)
// Elements are stored in vector for faster
// allocation. It's MINIMUM queue.
```

// Allows to merge heaps in O(1)

```
template < class T, class Cmp = less < T>>
struct PHeap {
 struct Node {
   T val:
   int child{-1}, next{-1}, prev{-1};
   Node(T x = T()) : val(x) \{\}
 }; // 11ee
 using Vnode = vector<Node>;
 Vnode& M;
 int root{-1};
 int unlink(int& i) {
   if (i >= 0) M[i].prev = -1;
   int x = i; i = -1;
   return x;
 } // d9f6
 void link(int host, int& i, int val) {
   if (i >= 0) M[i].prev = -1;
   i = val;
   if (i >= 0) M[i].prev = host;
 } // 47d5
 int merge(int 1, int r) {
   if (1 < 0) return r;</pre>
   if (r < 0) return 1;
   if (Cmp()(M[1].val, M[r].val)) swap(1, r);
   link(1, M[1].next, unlink(M[r].child));
   link(r, M[r].child, 1);
   return r;
 } // fc42
 int mergePairs(int v) {
   if (v < 0 || M[v].next < 0) return v;</pre>
   int v2 = unlink(M[v].next);
   int v3 = unlink(M[v2].next);
   return merge(merge(v, v2), mergePairs(v3));
 } // 2eea
 // Initialize heap with given node storage
  // Just declare 1 Vnode and pass it to heaps
 PHeap(Vnode& mem) : M(mem) {}
  // Add given key to heap, returns index; O(1)
  int push (const T& x) {
   int index = sz(M);
   M.emplace back(x):
   root = merge(root, index);
   return index:
 } // e744
  // Change key of i to smaller value; O(1)
 void decrease(int i, T val) {
   assert(!Cmp()(M[i].val, val));
   M[i].val = val;
   int prev = M[i].prev;
   if (prev < 0) return;</pre>
   auto& p = M[prev];
   link (prev, (p.child == i ? p.child
         : p.next), unlink(M[i].next));
   root = merge(root, i);
 } // 1a67
 bool empty() { return root < 0; }</pre>
 const T& top() { return M[root].val; }
  // Merge with other heap. Must use same vec.
 void merge(PHeap& r) { // time: O(1)
```

```
assert (&M == &r.M);
    root = merge(root, r.root); r.root = -1;
 } // 9623
  // Remove min element; time: O(lg n)
 void pop() {
   root = mergePairs(unlink(M[root].child));
 } // 5b13
}; // 09f3
                                         f749
structures/rmq.h
// Range Minimum Query; space: O(n lg n)
struct RMQ {
 using T = int;
 T ID = INT MAX;
 T f(T a, T b) { return min(a, b); }
 vector<vector<T>>> s:
  // Initialize RMQ structure; time: O(n lq n)
 RMO(const vector<T>& vec = {}) {
    s = \{vec\};
    for (int h = 1; h \le sz(vec); h *= 2) {
      s.pb({});
      auto& prev = s[sz(s)-2];
      rep(i, 0, sz(vec)-h*2+1)
        s.back().pb(f(prev[i], prev[i+h]));
   1 // 7c37
 } // 14ed
  // Query f(s[b], ..., s[e-1]); time: O(1)
 T query(int b, int e) {
   if (b >= e) return ID;
   int k = 31 - __builtin_clz(e-b);
    return f(s[k][b], s[k][e - (1<<k)]);</pre>
 } // caaa
}; // elce
structures/segtree config.h 7ef8
// Segment tree configurations to be used
// in segtree_general and segtree_persistent.
// See comments in TREE PLUS version
// to understand how to create custom ones.
// Capabilities notation: (update; query)
#if TREE PLUS // (+; sum, max, max count)
  // time: 0(lg n)
 using T = int; // Data type for update
                 // operations (lazy tag)
 T ID = 0;
                 // Neutral value for
                 // updates and lazy tags
  // This structure keeps aggregated data
  struct Agg {
    // Aggregated data: sum, max, max count
    // Default values should be neutral
    // values, i.e. "aggregate over empty set"
    T sum{0}, vMax{INT_MIN}, nMax{0};
    // Initialize as leaf (single value)
    void leaf() { sum = vMax = 0; nMax = 1; }
    // Combine data with aggregated data
    // from node to the right
    void merge(const Agg& r) {
      if (vMax < r.vMax) nMax = r.nMax;</pre>
      else if (vMax == r.vMax) nMax += r.nMax;
      vMax = max(vMax, r.vMax);
      sum += r.sum;
    } // 8850
    // Apply update provided in `x`:
    // - update aggregated data
```

```
// - update lazy tag `lazy`
    // - `size` is amount of elements
    // - return 0 if update should branch
    // (to be used in "segement tree beats")
    // - if you change value of `x` changed
    // value will be passed to next node
    // to the right during updates
   bool apply (T& lazy, T& x, int size) {
     lazy += x;
     sum += x*size;
     vMax += x;
     return 1;
   } // 4858
  }; // 9bf5
#elif TREE_MAX // (max; max, max count)
  // time: 0(lg n)
  using T = int:
 T ID = INT MIN;
  struct Agg {
    // Aggregated data: max value, max count
   T vMax{INT_MIN}, nMax{0};
   void leaf() { vMax = 0; nMax = 1; }
    void merge (const Agg& r) {
     if (vMax < r.vMax) nMax = r.nMax;</pre>
     else if (vMax == r.vMax) nMax += r.nMax;
     vMax = max(vMax, r.vMax);
   bool apply (T& lazy, T& x, int size) {
     if (vMax <= x) nMax = size;</pre>
     lazy = max(lazy, x);
     vMax = max(vMax, x);
     return 1;
   } // 8bd5
  }; // 15b6
#elif TREE SET // (=; sum, max, max count)
 // time: 0(la n)
  // Set ID to some unused value.
  using T = int:
 T ID = INT MIN:
  struct Agg {
   // Aggregated data: sum, max, max count
   T sum{0}, vMax{INT_MIN}, nMax{0};
   void leaf() { sum = vMax = 0; nMax = 1; }
    void merge(const Agg& r) {
     if (vMax < r.vMax) nMax = r.nMax;</pre>
     else if (vMax == r.vMax) nMax += r.nMax;
     vMax = max(vMax, r.vMax);
     sum += r.sum;
    } // 8850
   bool apply (T& lazy, T& x, int size) {
     lazy = x;
     sum = x*size;
     vMax = x;
     nMax = size;
     return 1;
   } // 845b
 }; // 7488
#elif TREE_BEATS // (+, min; sum, max)
  // time: amortized O(lg n) if not using +
         amortized O(lg^2 n) if using +
  // Lazy tag is pair (add, min).
  // To add x: run update with {x, INT MAX},
  // to min x: run update with \{0, x\}.
  // When both parts are provided addition
  // is applied first, then minimum.
```

```
int sum{0};
   void leaf() { sum = vMax = 0; nMax = 1; }
   void merge(const Agg& r) {
     if (r.vMax > vMax) {
       max2 = vMax;
       vMax = r.vMax;
       nMax = r.nMax;
     } else if (r.vMax == vMax) {
       nMax += r.nMax;
     } else if (r.vMax > max2) {
       max2 = r.vMax;
     } // b074
     max2 = max(max2, r.max2);
     sum += r.sum;
   } // 3124
   bool apply (T& lazy, T& x, int size) {
     if (max2 != INT_MIN && max2+x.x >= x.y)
     lazv.x += x.x;
     sum += x.x*size;
     vMax += x.x;
     if (max2 != INT_MIN) max2 += x.x;
     if (x.y < vMax) {
       sum -= (vMax-x.y) * nMax;
       vMax = x.v;
     1 // 7025
     lazy.y = vMax;
     return 1;
   } // fe0c
 }; // 2924
#endif
structures/segtree general.h 725a
// Highly configurable statically allocated
// (interval; interval) segment tree;
// space: 0(n)
struct SegTree {
 // Choose/write configuration
 #include "segtree config.h"
 // Root node is 1, left is i*2, right i*2+1
 vector<Agg> agg: // Aggregated data for nodes
 vector<T> lazy; // Lazy tags for nodes
 int len{1};
                  // Number of leaves
  // Initialize tree for n elements; time: O(n)
 SegTree(int n = 0) {
   while (len < n) len *= 2;
   agg.resize(len*2);
   lazy.resize(len*2, ID);
   rep(i, 0, n) agg[len+i].leaf();
   for (int i = len; --i;)
     (agg[i] = agg[i*2]).merge(agg[i*2+1]);
 1 // 4417
 void push(int i, int s) {
   if (lazy[i] != ID) {
     agg[i*2].apply(lazy[i*2], lazy[i], s/2);
     agg[i*2+1].apply(lazy[i*2+1],
                      lazy[i], s/2);
     lazy[i] = ID;
```

// Aggregated data: max value, max count,

int vMax{INT MIN}, nMax{0}, max2{INT MIN};

second max value, sum

using T = Pii:

struct Agg {

 $T ID = \{0, INT_MAX\};$

```
} // 3ba9
 } // 5d19
 // Modify interval [vb;ve) with val; O(lg n)
 T update(int vb, int ve, T val, int i = 1,
           int b = 0, int e = -1) {
   if (e < 0) e = len;
   if (vb >= e || b >= ve) return val;
   if (b >= vb && e <= ve &&
        agg[i].apply(lazy[i], val, e-b))
     return val:
   int m = (b+e) / 2;
   push(i, e-b);
   val = update(vb, ve, val, i*2, b, m);
   val = update(vb, ve, val, i*2+1, m, e);
   (agg[i] = agg[i*2]).merge(agg[i*2+1]);
   return val;
 } // aa8e
  // Query interval [vb;ve); time: O(lg n)
 Agg guery (int vb, int ve, int i = 1,
           int b = 0, int e = -1) {
   if (e < 0) e = len;</pre>
   if (vb >= e || b >= ve) return {};
   if (b >= vb && e <= ve) return agg[i];</pre>
   int m = (b+e) / 2;
   push(i, e-b);
   Agg t = query(vb, ve, i*2, b, m);
   t.merge(query(vb, ve, i*2+1, m, e));
   return t;
 } // 1a1e
1: // db5c
structures/segtree_persist.h dcfc
// Highly configurable (interval; interval)
// persistent segment tree;
// space: O(queries lq n)
// First tree version number is 0.
struct SegTree {
 // Choose/write configuration
  #include "segtree config.h"
  vector<Agg> agg; // Aggregated data for nodes
 vector<T> lazy; // Lazy tags for nodes
 vector<bool> cow; // Copy children on push?
 Vi L, R;
                   // Children links
 int len{1};
                  // Number of leaves
  // Initialize tree for n elements; O(lq n)
 SegTree(int n = 0) {
   int k = 1;
   while (len < n) len \star= 2, k++;
   agg.resize(k):
   lazy.resize(k, ID);
   cow.resize(k, 1);
   L.resize(k);
   R.resize(k);
   agg[--k].leaf();
   while (k--) {
     (agg[k] = agg[k+1]).merge(agg[k+1]);
     L[k] = R[k] = k+1;
   } // 211f
 } // 83cf
 // New version from version `i`; time: O(1)
 // First version number is 0.
 int fork(int i) {
   L.pb(L[i]); R.pb(R[i]); cow.pb(cow[i] = 1);
```

```
agg.pb(agg[i]); lazy.pb(lazy[i]);
   return sz(L)-1;
 } // a21b
 void push(int i, int s, bool w) {
   bool has = (lazy[i] != ID);
   if ((has || w) && cow[i]) {
     int a = fork(L[i]), b = fork(R[i]);
     L[i] = a; R[i] = b; cow[i] = 0;
   } // 1a3e
   if (has) {
     agg[L[i]].apply(lazy[L[i]],lazy[i],s/2);
     agg[R[i]].apply(lazy[R[i]],lazy[i],s/2);
     lazv[i] = ID;
   } // eca6
 } // 9f84
 // Modify interval [vb;ve) with val
 // in tree version `i`; time: O(lg n)
 T update(int i, int vb, int ve, T val,
           int b = 0, int e = -1) {
   if (e < 0) e = len;
   if (vb >= e || b >= ve) return val;
   if (b >= vb && e <= ve &&
       agg[i].apply(lazy[i], val, e-b))
      return val:
   int m = (b+e) / 2;
   push(i, e-b, 1);
   val = update(L[i], vb, ve, val, b, m);
   val = update(R[i], vb, ve, val, m, e);
    (agg[i] = agg[L[i]]).merge(agg[R[i]]);
   return val;
 } // 776e
 // Ouerv interval [vb:ve)
 // in tree version `i`; time: O(lq n)
 Agg query(int i, int vb, int ve,
           int b = 0, int e = -1) {
   if (e < 0) e = len;
   if (vb >= e || b >= ve) return {};
   if (b >= vb && e <= ve) return agg[i];</pre>
   int m = (b+e) / 2;
   push(i, e-b, 0);
   Agg t = query(L[i], vb, ve, b, m);
   t.merge(query(R[i], vb, ve, m, e));
   return t;
 } // abf4
}; // 8a44
structures/segtree point.h
                                         2b0f
// Segment tree (point, interval)
// Configure by modifying:
// - T - stored data type
// - ID - neutral element for query operation
// - f(a, b) - combine results
struct SegTree {
 using T = int;
 T ID = INT MIN;
 T f(T a, T b) { return max(a,b); }
 vector<T> V;
 int len:
 // Initialize tree for n elements; time: O(n)
 SegTree(int n = 0, T def = 0) {
   for (len = 1; len < n; len *= 2);</pre>
   V.resize(len*2, ID);
   rep(i, 0, n) V[len+i] = def;
   for (int i = len; --i;)
```

// Set element `i` to `val`; time: O(lg n)

V[i] = f(V[i*2], V[i*2+1]);

} // 459e

```
void set(int i, T val) {
   V[i += len] = val;
    while (i \neq 2)
     V[i] = f(V[i*2], V[i*2+1]);
  } // 4bcd
  // Ouery interval [b;e); time: O(lq n)
  T query(int b, int e) {
   b += len; e += len-1;
    if (b > e) return ID;
    if (b == e) return V[b];
   T \times = f(V[b], V[e]);
    while (b/2 < e/2) {
     if (^{\sim}b\&1) x = f(x, V[b^{\sim}1]);
     if (e&1) x = f(x, V[e^1]);
     b /= 2; e /= 2;
    } // 444a
    return x:
 } // de36
}; // d571
structures/treap.h
                                          0da3
// "Set" of implicit keyed treaps; space: O(n)
// Nodes are keyed by their indices in array
// of all nodes. Treap key is key of its root.
// "Node x" means "node with key x".
// "Treap x" means "treap with key x".
// Key -1 is "null".
// Put any additional data in Node struct.
struct Treap {
  struct Node {
    // E[0] = left child, <math>E[1] = right child
    // weight = node random weight (for treap)
    // size = subtree size, par = parent node
    int E[2] = \{-1, -1\}, weight {rand()};
   int size{1}, par{-1};
   bool flip{0}; // Is interval reversed?
  }; // c082
  vector<Node> G; // Array of all nodes
  // Initialize structure for n nodes
  // with keys 0, ..., n-1; time: O(n)
  // Each node is separate treap,
  // use join() to make sequence.
  Treap(int n = 0) : G(n) {}
  // Create new treap (a single node),
  // returns its key; time: O(1)
  int make() { G.pb({}); return sz(G)-1; }
  // Get size of node x subtree. x can be -1.
  int size(int x) { // time: O(1)
   return (x \ge 0 ? G[x].size : 0);
  // Propagate down data (flip flag etc).
  // x can be -1; time: O(1)
  void push(int x) {
   if (x >= 0 && G[x].flip) {
     G[x].flip = 0;
      swap(G[x].E[0], G[x].E[1]);
      each(e, G[x].E) if (e>=0) G[e].flip ^= 1;
   } // + any other lazy operations
  } // ed19
  // Update aggregates of node x.
 // x can be -1; time: O(1)
```

```
void update(int x) {
  if (x >= 0) {
    int& s = G[x].size = 1;
    G[x].par = -1;
    each (e, G[x].E) if (e >= 0) {
      s += G[e].size;
      G[e].par = x;
    } // f7a7
  } // + any other aggregates
} // 46a3
// Split treap x into treaps 1 and r
// such that 1 contains first i elements
// and r the remaining ones.
// x, 1, r can be -1; time: ~O(lg n)
void split(int x, int& l, int& r, int i) {
  push (x); 1 = r = -1;
  if (x < 0) return;
  int key = size(G[x].E[0]);
  if (i <= key) {
    split(G[x].E[0], 1, G[x].E[0], i);
  } else {
    split(G[x].E[1], G[x].E[1], r, i-key-1);
    1 = x;
  } // fe19
  update(x);
} // 8211
// Join treaps 1 and r into one treap
// such that elements of 1 are before
// elements of r. Returns new treap.
// 1, r and returned value can be -1.
int join(int 1, int r) { // time: ~O(lg n)
  push(1); push(r);
  if (1 < 0 || r < 0) return max(1, r);</pre>
  if (G[1].weight < G[r].weight) {</pre>
    G[1].E[1] = join(G[1].E[1], r);
    update(1);
    return 1:
  } // 18c7
  G[r].E[0] = join(1, G[r].E[0]);
  update(r);
  return r:
} // b559
// Find i-th node in treap x.
// Returns its key or -1 if not found.
// x can be -1; time: ~O(lq n)
int find(int x, int i) {
  while (x \ge 0) {
    push(x);
    int key = size(G[x].E[0]);
    if (key == i) return x;
    x = G[x].E[kev < i];
    if (key < i) i -= key+1;</pre>
  } // 054c
  return -1;
} // 0b9b
// Get key of treap containing node x
// (key of treap root). x can be -1.
int root(int x) { // time: ~O(lg n)
  while (G[x].par >= 0) x = G[x].par;
  return x;
1 // be8b
// Get position of node x in its treap.
// x is assumed to NOT be -1; time: ~O(lq n)
int index(int x) {
```

```
if (G[p].E[1] == x) i+=size(G[p].E[0])+1;
     if (G[p].flip) i = G[p].size-i-1;
     x = p;
   } // 3f81
   return i;
 } // ddad
 // Reverse interval [1;r) in treap x.
 // Returns new key of treap; time: ~O(lq n)
 int reverse(int x, int 1, int r) {
   int a, b, c;
   split(x, b, c, r);
   split (b, a, b, 1);
   if (b >= 0) G[b].flip ^= 1;
   return join(join(a, b), c);
 } // e418
}; // 73f2
structures/wavelet tree.h
// Wavelet tree ("merge-sort tree over values")
// Each node represent interval of values.
// seg[1]
             = original sequence
// seg[i]
             = subsequence with values
                represented by i-th node
// left[i][j] = how many values in seq[0:j)
                go to left subtree
struct WaveletTree {
 vector<Vi> seq, left;
 int len;
 WaveletTree() {}
 // Build wavelet tree for sequence `elems`;
 // time and space: O((n+maxVal) log maxVal)
  // Values are expected to be in [0:maxVal).
 WaveletTree(const Vi& elems, int maxVal) {
   for (len = 1: len < maxVal: len *= 2);</pre>
   seq.resize(len*2);
   left.resize(len*2);
   seq[1] = elems;
   build(1, 0, len);
  } // a5e9
 void build(int i, int b, int e) {
   if (i >= len) return;
   int m = (b+e) / 2;
   left[i].pb(0);
   each(x, seq[i]) {
     left[i].pb(left[i].back() + (x < m));
     seq[i*2 + (x >= m)1.pb(x);
   } // ac25
   build(i \star 2, b, m);
   build(i*2+1, m, e);
 } // 8153
 // Find k-th smallest element in [begin; end)
  // [begin;end); time: O(log maxVal)
 int kth(int begin, int end, int k, int i=1) {
   if (i >= len) return seq[i][0];
   int x = left[i][begin], y = left[i][end];
   if (k < y-x) return kth(x, y, k, i*2);
   return kth (begin-x, end-y, k-y+x, i*2+1);
 // Count number of elements >= vb and < ve
 // in [begin;end); time: O(log maxVal)
 int count (int begin, int end, int vb, int ve,
            int i = 1, int b = 0, int e = -1) {
   if (e < 0) e = len;
```

int p, i = size(G[x].E[G[x].flip]);

while ((p = G[x].par) >= 0) {

```
if (b >= ve || vb >= e) return 0;
           if (b >= vb && e <= ve) return end-begin;
           int m = (b+e) / 2;
           int x = left[i][begin], y = left[i][end];
           return count (x, y, vb, ve, i*2, b, m) +
             count (begin-x, end-y, vb, ve, i*2+1,m,e);
        } // 71cf
       }; // 3f09
       structures/ext/hash table.h 2d30
       #include <ext/pb ds/assoc container.hpp>
       using namespace __gnu_pbds;
       // gp_hash_table<K, V> = faster unordered_set
       // Anti-anti-hash
       const size_t HXOR = mt19937_64(time(0))();
       template<class T> struct SafeHash {
         size_t operator()(const T& x) const {
           return hash<T>() (x ^ T(HXOR));
         } // 3a78
1133 | }; // 7d0e
       structures/ext/rope.h
                                                051f
       #include <ext/rope>
       using namespace __gnu_cxx;
       // rope<T> = implicit cartesian tree
       structures/ext/tree.h
                                                a3bc
       #include <ext/pb_ds/assoc_container.hpp>
       #include <ext/pb_ds/tree_policy.hpp>
       using namespace __qnu_pbds;
       template<class T, class Cmp = less<T>>
       using ordered set = tree<</pre>
        T, null_type, Cmp, rb_tree_tag,
        tree order statistics node update
       // Standard set functions and:
       // t.order_of_key(key) - index of first >= key
       // t.find_by_order(i) - find i-th element
       // t1.join(t2) - assuming t1<>t2 merge t2 to t1
       structures/ext/trie.h
                                                5cc2
       #include <ext/pb_ds/assoc_container.hpp>
       #include <ext/pb_ds/trie_policy.hpp>
       using namespace __gnu_pbds;
       using pref trie = trie<</pre>
         string, null_type,
         trie_string_access_traits >> , pat_trie_tag ,
         trie_prefix_search_node_update
       text/aho corasick.h
                                                697d
       constexpr char AMIN = 'a'; // Smallest letter
       constexpr int ALPHA = 26: // Alphabet size
       // Aho-Corasick algorithm for linear-time
       // multiple pattern matching.
       // Add patterns using add(), then call build().
       struct Aho {
         vector<array<int, ALPHA>> nxt{1};
         Vi suf = \{-1\}, accLink = \{-1\};
         vector<Vi> accept{1};
         // Add string with given ID to structure
         // Returns index of accepting node
         int add(const string& str. int id) {
           int i = 0:
           each(c, str) {
```

nxt[i][c-AMIN] = sz(nxt);

accLink.pb(1); accept.pb({});

nxt.pb({}); suf.pb(-1);

if (!nxt[i][c-AMIN]) {

```
} // 5ead
     i = nxt[i][c-AMIN];
   } // ace9
    accept[i].pb(id);
   return i;
  } // 27c8
  // Build automata; time: O(V*ALPHA)
  void build() {
   queue<int> que:
   each(e, nxt[0]) if (e) {
     suf[e] = 0; que.push(e);
    } // c34d
    while (!que.emptv()) {
     int i = que.front(), s = suf[i], j = 0;
     each(e, nxt[i]) {
       if (e) que.push(e);
        (e ? suf[e] : e) = nxt[s][j++];
     accLink[i] = (accept[s].empty() ?
         accLink[s] : s);
   } // 1e8a
  } // 2561
  // Append `c` to state `i`
  int next(int i, char c) {
   return nxt[i][c-AMIN];
  } // 6bb7
  // Call `f` for each pattern accepted
  // when in state `i` with its ID as argument.
  // Return true from `f` to terminate early.
  // Calls are in descreasing length order.
  template<class F> void accepted(int i, F f) {
    while (i != -1) {
     each(a, accept[i]) if (f(a)) return;
     i = accLink[i];
   } // c175
 } // 1f0d
}; // 2768
                                         16fd
text/alcs.h
// All-substrings common sequences algorithm.
// Given strings A and B, algorithm computes:
// C(i,i,k) = |LCS(A[:i), B[i:k))|
// in compressed form; time and space: O(n^2)
// To describe the compression, note that:
// 1. C(i, j, k-1) \le C(i, j, k) \le C(i, j, k-1)+1
// 2. If j < k and C(i, j, k) = C(i, j, k-1)+1,
// then C(i, j+1, k) = C(i, j+1, k-1) + 1
// 3. If i >= k, then C(i, i, k) = 0
// This allows us to store just the following:
// ih(i,k) = min j s.t. C(i,j,k-1) < C(i,j,k)
struct ALCS {
 string A. B:
 vector<Vi>ih;
  ALCS() {}
  // Precompute compressed matrix; time: O(nm)
  ALCS(string s, string t) : A(s), B(t) {
   int n = sz(A), m = sz(B);
   ih.resize(n+1, Vi(m+1));
   iota(all(ih[0]), 0);
    rep(1, 1, n+1) {
     int iv = 0;
```

```
iv = min(ih[l-1][j], iv);
       } else {
         ih[l][j] = iv;
         iv = ih[l-1][j];
       } // 7af8
     } // d115
   } // baff
 } // 3e1d
 // Compute | LCS(A[:i), B[j:k))|; time: O(k-j)
 // Note: You can precompute data structure
 // to answer these queries in O(log n)
 // or compute all answers for fixed `i`.
 int operator()(int i, int j, int k) {
   int ret = 0;
   rep(q, j, k) ret += (ih[i][q+1] <= j);
   return ret;
 } // dabf
 // Compute subsequence LCS(A[:i), B[j:k));
 // time: O(k-j)
 string recover(int i, int j, int k) {
   string ret;
   while (i > 0 && j < k) {
     if (ih[i][k--] <= j) {</pre>
       ret.pb(B[k]);
       while (A[--i] != B[k]);
     } // 9d77
   } // ledf
   reverse (all (ret));
   return ret;
 } // 738c
 // Compute LCS'es of given prefix of A,
 // and all prefixes of given suffix of B.
 // Returns vector L of length |B|+1 s.t.
 // L[k] = |LCS(A[:i), B[j:k))|; time: O(|B|)
 Vi row(int i, int i) {
   Vi ret(sz(B)+1):
   rep(k, j+1, sz(ret))
     ret[k] = ret[k-1] + (ih[i][k] <= j);
   return ret;
 } // e5bd
 // Compute LCS'es of given prefix of A,
 // and all substrings of B; time: O(n^2)
 // Return matrix M such that:
 // M[j][k] = |LCS(A[:i), B[j:j+k))|
 vector<Vi> matrix(int i) {
   vector<Vi> ret;
   rep(j, 0, sz(B)+1) ret.pb(row(i, j));
   return ret;
 } // b4b4
}; // 30e0
                                          5729
text/kmp.h
// Computes prefsuf array; time: O(n)
// ps[i] = max prefsuf of [0;i); ps[0] := -1
template < class T > Vi kmp (const T& str) {
 Vi ps; ps.pb(-1);
 each(x, str) {
   int k = ps.back();
   while (k \ge 0 \&\& str[k] != x) k = ps[k];
   ps.pb(k+1);
 } // 05aa
 return ps;
} // 8a6c
```

rep(j, 1, m+1) {

if (A[1-1] != B[j-1]) {

ih[l][j] = max(ih[l-1][j], iv);

```
// Finds occurences of pat in vec; time: O(n)
// Returns starting indices of matches.
template<class T>
Vi match (const T& str, T pat) {
 int n = sz(pat);
  pat.pb(-1); // SET TO SOME UNUSED CHARACTER
 pat.insert(pat.end(), all(str));
 Vi ret, ps = kmp(pat);
  rep(i, 0, sz(ps)) {
   if (ps[i] == n) ret.pb(i-2*n-1);
 } // ale9
 return ret;
} // f986
text/kmr.h
                                          ee8c
// KMR algorithm for O(1) lexicographical
// comparison of substrings.
struct KMR {
 vector<Vi> ids;
  KMR() {}
  // Initialize structure; time: O(n lg^2 n)
  // You can change str type to Vi freely.
  KMR (const string € str) {
    ids.clear();
    ids.pb(Vi(all(str)));
    for (int h = 1; h \le sz(str); h *= 2) {
      vector<pair<Pii, int>> tmp;
      rep(j, 0, sz(str)) {
        int a = ids.back()[j], b = -1;
        if (j+h < sz(str)) b = ids.back()[j+h];
        tmp.pb({ {a, b}, j });
      } // a210
      sort(all(tmp));
      ids.emplace_back(sz(tmp));
      int n = 0;
      rep(j, 0, sz(tmp)) {
        if (j > 0 \&\& tmp[j-1].x != tmp[j].x)
        ids.back()[tmp[j].y] = n;
      } // bd2e
    } // 969a
 } // a40e
  // Get representative of [begin;end); O(1)
  Pii get (int begin, int end) {
    if (begin >= end) return {0, 0};
    int k = 31 - __builtin_clz(end-begin);
    return {ids[k][begin], ids[k][end-(1<<k)]};</pre>
  } // 85f3
  // Compare [b1;e1) with [b2;e2); O(1)
  // Returns -1 if <, 0 if ==, 1 if >
  int cmp(int b1, int e1, int b2, int e2) {
    int 11 = e1-b1, 12 = e2-b2;
    int 1 = min(11, 12);
   Pii x = \text{get}(b1, b1+1), y = \text{get}(b2, b2+1);
    if (x == y) return (11 > 12) - (11 < 12);
    return (x > y) - (x < y);
  // Compute suffix array of string; O(n)
  Vi sufArrav() {
   Vi sufs(sz(ids.back()));
    rep(i, 0, sz(ids.back()))
      sufs[ids.back()[i]] = i;
    return sufs;
```

```
} // d98d
}; // 457e
text/lcp.h
                                         0c65
// Compute Longest Common Prefix array for
// given string and it's suffix array; O(n)
// In order to compute suffix array use kmr.h
// or suffix_array_linear.h
template < class T>
Vi lcpArray (const T& str, const Vi& sufs) {
 int n = sz(str), k = 0;
 Vi pos(n), lcp(n-1);
  rep(i, 0, n) pos[sufs[i]] = i;
 rep(i, 0, n) {
   if (pos[i] < n-1) {</pre>
      int j = sufs[pos[i]+1];
      while (i+k < n && j+k < n &&
          str[i+k] == str[j+k]) k++;
      lcp[pos[i]] = k;
    } // 2cba
    if (k > 0) k--;
 1 // 8b22
 return lcp;
1 // d438
text/lyndon factorization.h 688c
// Compute Lyndon factorization for s; O(n)
// Word is simple iff it's stricly smaller
// than any of it's nontrivial suffixes.
// Lyndon factorization is division of string
// into non-increasing simple words.
// It is unique.
vector<string> duval(const string& s) {
 int n = sz(s), i = 0;
 vector<string> ret;
 while (i < n) {</pre>
    int j = i+1, k = i;
    while (j < n \&\& s[k] <= s[j])
     k = (s[k] < s[j] ? i : k+1), j++;
    while (i \le k)
      ret.pb(s.substr(i, j-k)), i += j-k;
 } // 3f17
 return ret;
} // 0e48
text/main lorentz.h
                                         70b7
#include "z function.h"
struct Sgr {
 int begin, end, len;
}; // f012
// Main-Lorentz algorithm for finding
// all squares in given word; time: O(n lq n)
// Results are in compressed form:
// (b, e, 1) means that for each b <= i < e
// there is square at position i of size 21.
// Each square is present in only one interval.
vector<Sgr> lorentz(const string& s) {
 vector<Sqr> ans:
 Vi pos(sz(s)/2+2, -1);
  rep(mid, 1, sz(s)) {
    int part = mid & ~ (mid-1), off = mid-part;
    int end = min(mid+part, sz(s));
    auto a = s.substr(off, part);
    auto b = s.substr(mid, end-mid);
    string ra(a.rbegin(), a.rend());
    string rb(b.rbegin(), b.rend());
```

```
rep(j, 0, 2) {
     // Set # to some unused character!
     Vi z1 = prefPref(ra);
     Vi z2 = prefPref(b+"#"+a);
      z1.pb(0); z2.pb(0);
      rep(c, 0, sz(a)) {
       int l = sz(a)-c;
        int x = c - min(1-1, z1[1]);
        int y = c - max(1-z2[sz(b)+c+1], j);
        if (x > y) continue;
        int sb = (j ? end-y-1*2 : off+x);
        int se = (j ? end-x-1*2+1 : off+y+1);
        int& p = pos[1];
        if (p != -1 \&\& ans[p].end == sb)
         ans[p].end = se;
          p = sz(ans), ans.pb({sb, se, 1});
     } // af4b
     a.swap(rb);
     b.swap(ra);
   } // d147
  } // a9bb
  return ans;
} // 0562
text/manacher.h
                                         8680
```

```
// Manacher algorithm; time: O(n)
// Finds largest radiuses for palindromes:
// r[2*i] = for center at i (single letter = 1)
// r[2*i+1] = for center between i and i+1
template < class T > Vi manacher (const T& str) {
  int n = sz(str) \star 2, c = 0, e = 1;
  Vi r(n, 1);
  auto get = [&](int i) { return i%2 ? 0 :
    (i \ge 0 \&\& i < n ? str[i/2] : i); }; // 3d98
  rep(i, 0, n) {
   if (i < e) r[i] = min(r[c*2-i], e-i);
    while (get(i-r[i]) == get(i+r[i])) r[i]++;
    if (i+r[i] > e) c = i, e = i+r[i]-1;
  } // Of87
  rep(i, 0, n) r[i] /= 2;
  return r;
} // a300
```

text/min rotation.h

e4d6

```
// Find lexicographically smallest
// rotation of s; time: O(n)
// Returns index where shifted word starts.
// You can use std::rotate to get the word:
// rotate(s.begin(), s.begin()+minRotation(s),
         s.end());
int minRotation(string s) {
  int a = 0, n = sz(s); s += s;
  rep(b, 0, n) rep(i, 0, n) {
   if (a+i == b || s[a+i] < s[b+i]) {</pre>
     b += max(0, i-1); break;
    } // 865b
   if (s[a+i] > s[b+i]) {
     a = b; break;
   } // 7628
  } // 40be
  return a;
} // 9ed8
```

```
text/monge.h
                                          b5e3
// NxN matrix A is simple (sub-)unit-Monge
// iff there exists a (sub-)permutation
// (N-1) \times (N-1) matrix P such that:
// A[x,y] = sum i>=x, j<y: P[i,j]
// The first column and last row are always 0.
// We represent these matrices implicitly
// using permutations p s.t. P[i,p(i)] = 1.
// (min, +) product of simple unit-Monge
// matrices represented by permutations P, Q,
// is also a simple unit-Monge matrix.
// The permutation that describes the product
// can be obtained by the following procedure:
// 1. Decompose P. O into minimal sequences of
// elementary transpositions.
// 2. Concatenate the transposition sequences.
// 3. Scan from left to right and remove
     transpositions that decrease
     inversion count (i.e. second crossings).
// 4. The reduced sequence represents result.
// Invert sub-permutation with values [0;n).
// Missing values should have value `def`.
Vi invert (const Vi& P, int n, int def) {
 Vi ret(n, def);
 rep(i, 0, sz(P)) if (P[i] != def)
   ret[P[i]] = i;
return ret;
1 // b5a6
// Split permutation P into half 'lo'
// with values less than 'k', and half 'hi'
// with remaining values, shifted by `k`.
// Missing rows from 'lo' and 'hi' are removed,
// original indices are in `loInd` and `hiInd`.
void split (const Vi& P, int k, Vi& lo, Vi& hi,
          Vi& loInd, Vi& hiInd) {
 int i = 0:
 each (e, P) {
   if (e < k) lo.pb(e), loInd.pb(i++);</pre>
   else hi.pb(e-k), hiInd.pb(i++);
 } // c3a6
} // 7bb7
// Map sub-permutation into sub-permutation
// of length `n` on given indices sets.
Vi expand (const Vi& P, Vi& ind1, Vi& ind2,
          int n, int def) {
 Vi ret(n, def);
 rep(k, 0, sz(P)) if (P[k] != def)
   ret[ind1[k]] = ind2[P[k]];
 return ret;
} // 7f23
// Compute (min, +) product of square
// simple unit-Monge matrices given their
// permutation representations; time: O(n lq n)
// Permutation of second matrix is inverted!
Vi comb (const Vi& P, const Vi& invQ) {
 int n = sz(P);
 if (n < 100) {
    // 5s -> 1s speedup for ALIS for n = 10^5
   Vi ret = invert(P, n, -1);
   rep(i, 0, sz(inv0)) {
     int from = invQ[i];
      rep(j, 0, i) from += invQ[j] > invQ[i];
     for (int j = from; j > i; j--)
        if (ret[j-1] < ret[j])</pre>
```

swap(ret[j-1], ret[j]);

```
1 // 7cd1
    return invert(ret, n, -1);
  } // a886
  Vi p1, p2, q1, q2, i1, i2, j1, j2;
  split (P, n/2, p1, p2, i1, i2);
  split(invQ, n/2, q1, q2, j1, j2);
  p1 = expand (comb (p1, q1), i1, j1, n, -1);
  p2 = expand(comb(p2, q2), i2, j2, n, n);
  q1 = invert(p1, n, -1);
  q2 = invert(p2, n, n);
  Vi ans(n, -1);
  int delta = 0, j = n;
  rep(i, 0, n) {
   ans[i] = (p1[i] < 0 ? p2[i] : p1[i]);
    while (j > 0 && delta >= 0)
     delta -= (q2[--j] < i || q1[j] >= i);
    if (p2[i] < j || p1[i] >= j)
      if (delta++ < 0)
        if (q2[j] < i || q1[j] >= i)
          ans[i] = j;
  } // c396
 return ans;
} // cbff
// Helper function for `mongeMul`.
void padPerm(const Vi& P, Vi& has, Vi& pad,
            Vi& ind, int n) {
  vector<bool> seen(n);
  rep(i, 0, sz(P)) if (P[i] != -1) {
    ind.pb(i):
   has.pb(P[i]);
   seen[P[i]] = 1;
 } // 157e
 rep(i, 0, n) if (!seen[i]) pad.pb(i);
} // 103b
// Compute (min, +) product of
// simple sub-unit-Monge matrices given their
// permutation representations; time: O(n lg n)
// Left matrix has size sz(P) x sz(O).
// Right matrix has size sz(0) x n.
// Output matrix has size sz(P) x n.
// NON-SQUARE MATRICES ARE NOT TESTED!
Vi mongeMul(const Vi& P, const Vi& Q, int n) {
 Vi h1, p1, i1, h2, p2, i2;
  padPerm(P, h1, p1, i1, sz(Q));
  padPerm(invert(Q, n, -1), h2, p2, i2, sz(Q));
  h1.insert(h1.begin(), all(p1));
  h2.insert(h2.end(), all(p2));
  Vi ans(sz(P), -1), tmp = comb(h1, h2);
  rep(i, 0, sz(i1)) {
   int j = tmp[i+sz(p1)];
    if (j < sz(i2)) {</pre>
      ans[i1[i]] = i2[j];
   } // 4d16
 } // c8a0
 return ans;
} // fda1
// Range Longest Increasing Subsequence Query;
// preprocessing: O(n lg^2 n), query: O(lg n)
#include "../structures/wavelet_tree.h"
struct ALIS {
 WaveletTree tree;
 ALIS() {}
```

```
// Precompute data structure; O(n 1g^2 n)
 ALIS(const Vi& seq) {
   Vi P = build(seq);
   each(k, P) if (k == -1) k = sz(seq);
   tree = \{P, sz(seq)+1\};
 } // aaed
 // Query LIS of s[b;e); time: O(lq n)
 int operator()(int b, int e) {
   return e - b -
     tree.count(b, sz(tree.seq[1]), 0, e);
 } // fb4a
 Vi build (const Vi& seg) {
   int n = sz(seq);
   if (!n) return {};
   int lo = *min_element(all(seq));
   int hi = *max_element(all(seq));
   if (lo == hi) {
     Vi tmp(n);
     iota(all(tmp), 1);
     tmp.back() = -1;
     return tmp;
   } // 16d1
   int mid = (lo+hi+1) / 2;
   Vi p1, p2, i1, i2;
   split(seq, mid, p1, p2, i1, i2);
   p1 = expand(build(p1), i1, i1, n, -1);
   p2 = expand(build(p2), i2, i2, n, -1);
   each(i, i1) p2[i] = i;
   each(j, i2) p1[j] = j;
   return mongeMul(p1, p2, n);
 } // fe71
): // 5740
text/palindromic tree.h
                                         1207
// Tree of all palindromes in string,
```

constexpr int ALPHA = 26; // Set alphabet size // constructed online by appending letters. // space: O(n*ALPHA); time: O(n) // Code marked with [EXT] is extension for // calculating minimal palindrome partition // in O(n lg n). Can also be modified for // similar dynamic programmings. struct PalTree { Vi txt; // Text for which tree is built // Node 0 = empty palindrome (root of even) // Node 1 = "-1" palindrome (root of odd) Vi len{0, -1}; // Lengths of palindromes Vi link{1, 0}; // Suffix palindrome links // Edges to next palindromes vector<arrav<int, ALPHA>> to{ {}, {} }; int last{0}; // Current node (max suffix pal) Vi diff{0, 0}; // len[i]-len[link[i]] [EXT] Vi slink{0, 0}; // Serial links Vi series{0, 0}; // Series DP answer [EXT] Vi ans{0}; // DP answer for prefix[EXT] int ext(int i) { while (len[i]+2 > sz(txt) || txt[sz(txt)-len[i]-2] != txt.back())i = link[i];return i; } // d442 // Append letter from [0; ALPHA); time: O(1)

// (or O(lg n) if [EXT] is enabled)

```
void add(int x) {
   txt.pb(x);
    last = ext(last);
    if (!to[last][x]) {
     len.pb(len[last]+2);
     link.pb(to[ext(link[last])][x]);
     to[last][x] = sz(to);
     to.pb({});
     diff.pb(len.back() - len[link.back()]);
     slink.pb(diff.back() == diff[link.back()]
       ? slink[link.back()] : link.back());
     series.pb(0);
     // [/EXT]
    } // 8c1b
    last = to[last][x];
    // [EXT]
    ans.pb(INT_MAX);
    for (int i=last; len[i] > 0; i=slink[i]) {
     series[i] = ans[sz(ans) - len[slink[i]]
                    - diff[i] - 1];
     if (diff[i] == diff[link[i]])
       series[i] = min(series[i],
                       series[link[i]]);
     // If you want only even palindromes
     // set ans only for sz(txt) %2 == 0
     ans.back() = min(ans.back(), series[i]+1);
   } // ab3b
   // [/EXT]
 } // 66d3
1: // 595d
```

text/suffix array linear.h 1341

```
#include "../util/radix sort.h"
// KS algorithm for suffix array; time: O(n)
// Input values are assumed to be in [1;k]
Vi sufArray(Vi str, int k) {
  int n = sz(str);
  Vi suf(n);
  str.resize(n+15);
  if (n < 15) {
   iota(all(suf), 0);
    rep(j, 0, n) countSort(suf,
     [&] (int i) { return str[i+n-j-1]; }, k);
    return suf;
  } // 5fcf
  // Compute triples codes
  Vi tmp, code(n+2);
  rep(i, 0, n) if (i % 3) tmp.pb(i);
  rep(j, 0, 3) countSort(tmp,
   [&] (int i) { return str[i-j+2]; }, k);
  int mc = 0, j = -1;
  each(i, tmp) {
   code[i] = mc += (j == -1)
       str[i] != str[j] ||
       str[i+1] != str[j+1] ||
       str[i+2] != str[j+2]);
   j = i;
  } // bfdc
  // Compute suffix array of 2/3
  tmp.clear();
  for (int i=1; i < n; i += 3) tmp.pb(code[i]);</pre>
 tmp.pb(0);
```

```
for (int i=2; i < n; i += 3) tmp.pb(code[i]);</pre>
 tmp = sufArray(move(tmp), mc);
 // Compute partial suffix arrays
 Vi third:
 int th = (n+4) / 3;
 if (n%3 == 1) third.pb(n-1);
 rep(i, 1, sz(tmp)) {
   int e = tmp[i];
   tmp[i-1] = (e 
   code[tmp[i-1]] = i;
   if (e < th) third.pb(e*3);
 } // f9f1
 tmp.pop_back();
 countSort (third,
   [&] (int i) { return str[i]; }, k);
 // Merge suffix arrays
 merge(all(third), all(tmp), suf.begin(),
   [&] (int 1, int r) {
     while (1\%3 == 0 \mid | r\%3 == 0) {
       if (str[1] != str[r])
         return str[l] < str[r];</pre>
       1++; r++;
     } // 2f8a
     return code[1] < code[r];</pre>
   }); // 4cb3
 return suf;
) // 9165
// KS algorithm for suffix array; time: O(n)
Vi sufArray(const string& str) {
 return sufArray(Vi(all(str)), 255);
```

text/suffix automaton.h

d00d

```
constexpr char AMIN = 'a'; // Smallest letter
constexpr int ALPHA = 26; // Set alphabet size
// Suffix automaton - minimal DFA that
// recognizes all suffixes of given string
// (and encodes all substrings);
// space: O(n*ALPHA); time: O(n)
// Paths from root are equivalent to substrings
// Extensions:
// - [OCC] - count occurences of substrings
// - [PATHS] - count paths from node
struct SufDFA {
 // State v represents endpos-equivalence
 // class that contains words of all lengths
 // between link[len[v]]+1 and len[v].
 // len[v] = longest word of equivalence class
 // link[v] = link to state of longest suffix
              in other equivalence class
 // to [v][c] = automaton edge c from v
 Vi len{0}, link{-1};
 vector<array<int, ALPHA>> to{ {} };
 int last{0}; // Current node (whole word)
 vector<Vi> inSufs; // [OCC] Suffix-link tree
                    // [OCC] Occurence count
 Vi cnt{0};
 vector<ll> paths; // [PATHS] Out-path count
 SufDFA() {}
 // Build suffix automaton for given string
 // and compute extended stuff; time: O(n)
 SufDFA(const string& s) {
   each(c, s) add(c);
   finish();
```

```
} // ec2e
// Append letter to the back
void add(char c) {
  int v = last, x = c-AMIN;
  last = sz(len);
  len.pb(len[v]+1);
  link.pb(0);
  to.pb({});
  cnt.pb(1); // [OCC]
  while (v != -1 \&\& !to[v][x]) {
   to[v][x] = last;
    v = link[v];
  } // 4cfc
  if ( \lor != -1 )  {
    int q = to[v][x];
    if (len[v]+1 == len[q]) {
      link[last] = q;
      len.pb(len[v]+1);
      link.pb(link[q]);
      to.pb(to[q]);
      cnt.pb(0); // [OCC]
      link[last] = link[q] = sz(len)-1;
      while (v != -1 && to[v][x] == q) {
       to[v][x] = link[q];
       v = link[v];
      } // 784f
    } // 90aa
 } // af69
} // 345a
// Compute some additional stuff (offline)
void finish() {
  inSufs.resize(sz(len));
  rep(i, 1, sz(link)) inSufs[link[i]].pb(i);
  dfsSufs(0);
  // [PATHS]
  paths.assign(sz(len), 0);
  dfs(0);
  // [/PATHS]
} // 3f75
// Only for [OCC]
void dfsSufs(int v) {
  each(e, inSufs[v]) {
    dfsSufs(e);
    cnt[v] += cnt[e];
  } // 2469
} // Oc60
// Only for [PATHS]
void dfs(int v) {
  if (paths[v]) return;
  paths[v] = 1;
  each(e, to[v]) if (e) {
    dfs(e);
    paths[v] += paths[e];
  } // 22b3
} // d004
// Go using edge `c` from state `i`.
// Returns 0 if edge doesn't exist.
int next(int i, char c) {
  return to[i][c-AMIN];
} // c363
// Get lexicographically k-th substring
```

```
// of represented string; time: O(|substr|)
 // Empty string has index 0.
 // Requires [PATHS] extension.
 string lex(ll k) {
   string s;
   int v = 0:
   while (k--) rep(i, 0, ALPHA) {
     int e = to[v][i];
     if (e) {
       if (k < paths[e]) {</pre>
         s.pb(char(AMIN+i));
         v = e;
         break;
       } // f307
       k -= paths[e];
     } // 29be
   } // 4600
   return s;
 } // e4af
}; // 7135
text/suffix tree.h
                                          40a4
constexpr int ALPHA = 26;
// Ukkonen's algorithm for online suffix tree
// construction; space: O(n*ALPHA); time: O(n)
// Real tree nodes are called dedicated nodes.
// "Nodes" lying on compressed edges are called
// implicit nodes and are represented
// as pairs (lower node, label index).
// Labels are represented as intervals [L;R)
// which refer to substrings [L;R) of txt.
// Leaves have labels of form [L; infinity],
// use getR to get current right endpoint.
// Suffix links are valid only for internal
// nodes (non-leaves).
struct SufTree {
 Vi txt; // Text for which tree is built
 // to[v][c] = edge with label starting with c
               from node v
 vector<array<int, ALPHA>> to{ {} };
 Vi L{0}, R{0}; // Parent edge label endpoints
              // Parent link
// Suffix link
 Vi par{0};
 Vi link{0};
 Pii cur{0, 0}; // Current state
 // Get current right end of node label
 int getR(int i) { return min(R[i],sz(txt)); }
 // Follow edge `e` of implicit node `s`.
 // Returns (-1, -1) if there is no edge.
 Pii next (Pii s, int e) {
   if (s.v < qetR(s.x))
     return txt[s.y] == e ? mp(s.x, s.y+1)
                           : mp(-1, -1);
   e = to[s.x][e];
   return e ? mp(e, L[e]+1) : mp(-1, -1);
 // Create dedicated node for implicit node
 // and all its suffixes
 int split(Pii s) {
   if (s.y == R[s.x]) return s.x;
   int t = sz(to); to.pb({});
   to[t][txt[s.y]] = s.x;
   L.pb(L[s.x]);
   R.pb(L[s.x] = s.y);
   par.pb(par[s.x]);
```

par[s.x] = to[par[t]][txt[L[t]]] = t;

link.pb(-1);

```
int v = link[par[t]], l = L[t] + !par[t];
    while (1 < R[t]) {</pre>
     v = to[v][txt[l]];
     1 += getR(v) - L[v];
   } // 0393
   v = split(\{v, getR(v)-l+R[t]\});
   link[t] = v;
   return t;
  } // 10bb
  // Append letter from [0; ALPHA) to the back
  void add(int x) { // amortized time: 0(1)
   Pii t; txt.pb(x);
    while ((t = next(cur, x)).x == -1) {
     int m = split(cur);
     to[m][x] = sz(to);
     to.pb({});
     par.pb(m);
     L.pb(sz(txt)-1);
     R.pb(INT MAX);
     link.pb(-1);
     cur = {link[m], getR(link[m])};
     if (!m) return;
   } // 60c2
   cur = t;
 } // 6f4e
}; // dbfb
text/z function.h
                                         0466
```

```
// Computes Z function array; time: O(n)
// zf[i] = max common prefix of str and str[i:]
template < class T > Vi prefPref (const T& str) {
  int n = sz(str), b = 0, e = 1;
  Vi zf(n);
  rep(i, 1, n) {
    if (i < e) zf[i] = min(zf[i-b], e-i);</pre>
    while (i+zf[i] < n &&
     str[zf[i]] == str[i+zf[i]]) zf[i]++;
   if (i+zf[i] > e) b = i, e = i+zf[i];
  } // e906
  zf[0] = n;
  return zf;
} // b88d
```

trees/centroid decomp.h

5247

```
// Centroid decomposition; space: O(n lg n)
struct CentroidTree {
 // child[v] = children of v in centroid tree
  // par[v] = parent of v in centroid tree
              (-1 for root)
  // depth[v] = depth of v in centroid tree
                (0 for root)
  // ind[v][i] = index of vertex v in i-th
                centroid subtree from root
  // size[v] = size of centroid subtree of v
 // subtree[v] = list of vertices
                 in centroid subtree of v
  // dists[v] = distances from v to vertices
               in its centroid subtree
                (in the order of subtree[v])
 // neigh[v] = neighbours of v
                in its centroid subtree
  // dir[v][i] = index of centroid neighbour
                that is first vertex on path
                from centroid v to i-th vertex
                of centroid subtree
                (-1 for centroid)
```

```
vector<Vi> child, ind, dists, subtree,
            neigh, dir;
 Vi par, depth, size;
 int root; // Root centroid
 CentroidTree() {}
 CentroidTree (vector<Vi>& G)
     : child(sz(G)), ind(sz(G)), dists(sz(G)),
       subtree(sz(G)), neigh(sz(G)),
       dir(sz(G)), par(sz(G), -2),
       depth(sz(G)), size(sz(G)) {
   root = decomp(G, 0, 0);
 } // 026c
 void dfs(vector<Vi>& G, int v, int p) {
   size[v] = 1;
   each (e, G[v]) if (e != p && par[e] == -2)
     dfs(G, e, v), size[v] += size[e];
 } // bbed
 void layer (vector < Vi>& G, int v,
            int p, int c, int d) {
   ind[v].pb(sz(subtree[c]));
   subtree[c].pb(v);
   dists[c].pb(d);
   dir[c].pb(sz(neigh[c])-1);
   each(e, G[v]) if (e != p && par[e] == -2) {
     if (v == c) neigh[c].pb(e);
     layer(G, e, v, c, d+1);
   } // dc82
 } // 37ee
 int decomp(vector<Vi>& G, int v, int d) {
   dfs(G, v, -1);
   int p = -1, s = size[v];
 loop:
   each(e, G[v]) {
     if (e != p && par[e] == -2 &&
         size[e] > s/2) {
       p = v; v = e; goto loop;
     } // e0a5
   } // 3533
   par[v] = -1;
   size[v] = s;
   depth[v] = d;
   layer(G, v, -1, v, 0);
   each (e, G[v]) if (par[e] == -2) {
     int j = decomp(G, e, d+1);
     child[v].pb(j);
     par[j] = v;
   } // 70b5
   return v;
 } // 217c
1: // 71d6
trees/centroid offline.h
// Helper for offline centroid decomposition
// Usage: CentroidDecomp(G);
// Constructor calls method `process`
```

// for each centroid subtree.

Vi size; // Used internally

CentroidDecomp (vector < Vi> € q)

vector<bool> on: // Is vertex enabled?

: G(g), on(sz(g), 1), size(sz(g)) {

struct CentroidDecomp {

decomp(0);

```
ac92
vector<Vi>& G; // Reference to target graph
// Run centroid decomposition for graph g
```

```
1 // 8677
 // Compute subtree sizes for subtree rooted
 // at v, ignoring p and disabled vertices
 void computeSize(int v, int p) {
   size[v] = 1;
   each(e, G[v]) if (e != p && on[e])
     computeSize(e, v), size[v] += size[e];
 } // 1c0d
 void decomp(int v) {
   computeSize(v, -1);
   int p = -1, s = size[v];
 loop:
   each(e, G[v]) {
     if (e != p && on[e] && size[e] > s/2) {
       p = v; v = e; goto loop;
     } // e0a5
   } // f31d
   process(v);
   on[v] = 0;
   each(e, G[v]) if (on[e]) decomp(e);
 } // f170
 // Process current centroid subtree:
 // - v is centroid
 // - boundary vertices have on[x] = 0
 // Formally: Let H be subgraph induced
 // on vertices such that on [v] = 1.
 // Then current centroid subtree is
 // connected component of H that contains v
 // and v is its centroid.
 void process(int v) {
   // Do your stuff here...
 } // d41d
1: // a923
trees/heavylight decomp.h
                                         ebb8
#include "../structures/segtree_point.h"
// Heavy-Light Decomposition of tree
// with subtree query support; space: O(n)
struct HLD {
 // Subtree of v = [pos[v]; pos[v]+size[v])
 // Chain with v = [chBegin[v]; chEnd[v])
 Vi par;
             // Vertex parent
 Vi size;
             // Vertex subtree size
 Vi depth; // Vertex distance to root
             // Vertex position in "HLD" order
 Vi pos:
 Vi chBegin; // Begin of chain with vertex
 Vi chEnd; // End of chain with vertex
 Vi order; // "HLD" preorder of vertices
 SegTree tree; // Verts are in HLD order
 HLD() {}
 // Initialize structure for tree G
 // and given root; time: O(n lg n)
  // MODIFIES ORDER OF EDGES IN G!
 HLD (vector<Vi>& G, int root)
     : par(sz(G)), size(sz(G)),
        depth(sz(G)), pos(sz(G)),
        chBegin(sz(G)), chEnd(sz(G)) {
   dfs(G, root, -1);
   decomp(G, root, -1, 0);
   tree = {sz(order)};
 } // 8263
 void dfs(vector<Vi>& G, int v, int p) {
   par[v] = p;
   size[v] = 1;
   depth[v] = p < 0 ? 0 : depth[p]+1;
```

```
if (G[v].empty()) return;
  int& fs = G[v][0];
  if (fs == p) swap(fs, G[v].back());
  each(e, G[v]) if (e != p) {
    dfs(G, e, v);
    size[v] += size[e];
    if (size[e] > size[fs]) swap(e, fs);
  } // 9872
} // e25f
void decomp(vector<Vi>& G,
            int v, int p, int chb) {
  pos[v] = sz(order);
  chBegin[v] = chb;
  chEnd[v] = pos[v]+1;
  order.pb(v);
  each(e, G[v]) if (e != p) {
    if (e == G[v][0]) {
      decomp (G, e, v, chb);
      chEnd[v] = chEnd[e];
      decomp(G, e, v, sz(order));
    } // c84a
  } // f707
} // eb89
// Get root of chain containing v
int chRoot(int v) {return order[chBegin[v]];}
// Level Ancestor Query; time: O(lq n)
int lag(int v, int level) {
  while (true) {
    int k = pos[v] - depth[v] + level;
    if (k >= chBegin[v]) return order[k];
    v = par[chRoot(v)];
  } // 8c18
1 // 675e
// Lowest Common Ancestor; time: O(lq n)
int lca(int a, int b) {
  while (chBegin[a] != chBegin[b]) {
    int ha = chRoot(a), hb = chRoot(b);
    if (depth[ha] > depth[hb]) a = par[ha];
    else b = par[hb];
  } // 5620
  return depth[a] < depth[b] ? a : b;</pre>
} // c168
// Call func (chBegin, chEnd) on each path
// segment; time: O(lg n * time of func)
template<class T>
void iterPath(int a, int b, T func) {
  while (chBegin[a] != chBegin[b]) {
    int ha = chRoot(a), hb = chRoot(b);
    if (depth[ha] > depth[hb]) {
      func(chBegin[a], pos[a]+1);
      a = par[ha];
    } else {
      func(chBegin[b], pos[b]+1);
      b = par[hb];
    } // f9a5
  } // 563c
  if (pos[a] > pos[b]) swap(a, b);
  // Remove +1 from pos[a]+1 for vertices
  // queries (with +1 -> edges).
  func(pos[a]+1, pos[b]+1);
} // 17e5
// Query path between a and b; O(lg^2 n)
```

iterPath(a, b, [&](int i, int j) {

auto ret = tree.ID;

```
ret = tree.f(ret, tree.query(i, j));
    }); // 1113
    return ret:
  } // 1bc9
  // Ouerv subtree of v; time: O(lq n)
  SegTree::T querySubtree(int v) {
   return tree.query(pos[v], pos[v]+size[v]);
 } // 23db
}; // 59bf
trees/lca.h
                                          294f
// LAQ and LCA using jump pointers
// space: O(n lg n)
struct LCA {
  vector<Vi> jumps;
  Vi level, pre, post;
  int cnt{0}, depth;
  LCA() {}
  // Initialize structure for tree G
  // and root r; time: O(n lg n)
  LCA(vector<Vi>& G, int root)
      : jumps(sz(G)), level(sz(G)),
        pre(sz(G)), post(sz(G)) {
    dfs(G, root, root);
    depth = int(log2(sz(G))) + 2;
    rep(j, 0, depth) each(v, jumps)
      v.pb(jumps[v[j]][j]);
  } // d6ce
  void dfs(vector<Vi>& G, int v, int p) {
   level[v] = p == v ? 0 : level[p]+1;
    jumps[v].pb(p);
    pre[v] = ++cnt;
   each(e, G[v]) if (e != p) dfs(G, e, v);
   post[v] = ++cnt;
  } // e286
  // Check if a is ancestor of b; time: O(1)
  bool isAncestor(int a, int b) {
    return pre[a] <= pre[b] &&
           post[b] <= post[a];</pre>
  } // 5514
  // Lowest Common Ancestor; time: O(lq n)
  int operator()(int a, int b) {
    for (int j = depth; j--;)
      if (!isAncestor(jumps[a][j], b))
       a = iumps[a][i]:
    return isAncestor(a, b) ? a : jumps[a][0];
  1 // 27d8
  // Level Ancestor Query; time: O(lg n)
  int lag(int a, int lvl) {
    for (int j = depth; j--;)
      if (lvl <= level[jumps[a][j]])</pre>
       a = jumps[a][j];
    return a;
  // Get distance from a to b; time: O(lg n)
  int distance(int a, int b) {
    return level[a] + level[b] -
           level[operator()(a, b)]*2;
 // Get k-th vertex on path from a to b,
```

```
// a is 0, b is last; time: O(lq n)
  // Returns -1 if k > distance(a, b)
  int kthVertex(int a, int b, int k) {
    int c = operator()(a, b);
    if (level[a]-k >= level[c])
      return lag(a, level[a]-k);
    k \leftarrow level[c] \times 2 - level[a];
    return (k > level[b] ? -1 : laq(b, k));
 } // 46c9
}; // 2254
trees/lca linear.h
                                          7aa5
// LAQ and LCA using jump pointers
// with linear memory; space: O(n)
struct LCA {
  Vi par, jmp, depth, pre, post;
 int cnt{0};
  LCA() {}
  // Initialize structure for tree G
  // and root v; time: O(n lg n)
  LCA(vector<Vi>& G, int v)
      : par(sz(G), -1), jmp(sz(G), v),
        depth(sz(G)), pre(sz(G)), post(sz(G)) {
    dfs(G, v);
  } // 94cf
  void dfs(vector<Vi>& G, int v) {
    int j = jmp[v], k = jmp[j], x =
      depth[v]+depth[k] == depth[j]*2 ? k : v;
    pre[v] = ++cnt;
    each(e, G[v]) if (!pre[e]) {
      par[e] = v; jmp[e] = x;
      depth[e] = depth[v]+1;
     dfs(G, e);
    } // b123
    post[v] = ++cnt;
  1 // 3280
  // Level Ancestor Query; time: O(lg n)
  int laq(int v, int d) {
    while (depth[v] > d)
      v = depth[jmp[v]] < d ? par[v] : jmp[v];
    return v;
  } // f509
  // Lowest Common Ancestor; time: O(lq n)
  int operator()(int a, int b) {
    if (depth[a] > depth[b]) swap(a, b);
    b = laq(b, depth[a]);
    while (a != b) {
      if (jmp[a] == jmp[b])
        a = par[a], b = par[b];
        a = jmp[a], b = jmp[b];
    } // fe08
    return a;
  } // 25ff
  // Check if a is ancestor of b; time: O(1)
  bool isAncestor(int a, int b) {
    return pre[a] <= pre[b] &&</pre>
           post[b] <= post[a];</pre>
  // Get distance from a to b; time: O(lg n)
```

int distance(int a, int b) {

} // a340

return depth[a] + depth[b] -

depth[operator()(a, b)]*2;

```
// Get k-th vertex on path from a to b.
  // a is 0, b is last; time: O(lg n)
  // Returns -1 if k > distance(a, b)
  int kthVertex(int a, int b, int k) {
    int c = operator()(a, b);
    if (depth[a]-k >= depth[c])
      return laq(a, depth[a]-k);
    k \leftarrow depth[c] \times 2 - depth[a];
    return (k > depth[b] ? -1 : lag(b, k));
 } // 34ed
}; // a221
trees/link cut tree.h
                                          6bd6
constexpr int INF = 1e9:
// Link/cut tree; space: O(n)
// Represents forest of (un)rooted trees.
struct LinkCutTree {
 vector<array<int, 2>> child;
 Vi par, prev, flip, size;
  // Initialize structure for n vertices; O(n)
  // At first there's no edges.
  LinkCutTree(int n = 0)
      : child(n, \{-1, -1\}), par(n, -1),
        prev(n, -1), flip(n, -1), size(n, 1) {}
  void push(int x) {
   if (x >= 0 && flip[x]) {
      flip[x] = 0;
      swap(child[x][0], child[x][1]);
      each(e, child[x]) if (e>=0) flip[e] ^= 1;
   } // + any other lazy path operations
 } // bae2
  void update(int x) {
   if (x >= 0) {
      size[x] = 1;
      each(e, child[x]) if (e \geq= 0)
        size[x] += size[e];
   } // + any other path aggregates
  } // 8ec0
  void auxLink(int p, int i, int ch) {
    child[p][i] = ch;
    if (ch >= 0) par[ch] = p;
    update(p);
  } // Oa9a
  void rot(int p, int i) {
    int x = child[p][i], g = par[x] = par[p];
    if (q \ge 0) child [q] [child [q] [1] == p] = x;
    auxLink(p, i, child[x][!i]);
    auxLink(x, !i, p);
    swap(prev[x], prev[p]);
    update(g);
  1 // 4c76
  void splav(int x) {
    while (par[x] >= 0) {
      int p = par[x], g = par[p];
      push(q); push(p); push(x);
      bool f = (child[p][1] == x);
      if (q >= 0) {
        if (child[g][f] == p) { // zig-zig
          rot(g, f); rot(p, f);
        } else { // zig-zag
         rot(p, f); rot(g, !f);
        } // 2ebb
      } else { // zig
        rot (p, f);
      } // f8a2
```

```
1 // 446b
   push(x);
 } // 55a7
 // After this operation x becomes the end
 // of preferred path starting in root;
 void access(int x) { // amortized O(lg n)
   while (true) {
     splav(x);
     int p = prev[x];
     if (p < 0) break;
     prev[x] = -1;
     splay(p);
     int r = child[p][1];
     if (r \ge 0) swap(par[r], prev[r]);
     auxLink(p, 1, x);
   } // 2b87
 } // 30be
 // Make x root of its tree; amortized O(lq n)
 void makeRoot(int x) {
   access(x):
   int& 1 = child[x][0];
   if (1 >= 0) {
      swap(par[1], prev[1]);
     flip[1] ^= 1;
     update(1);
     1 = -1;
     update(x);
   } // 0064
 1 // b246
 // Find root of tree containing x
 int find(int x) { // time: amortized O(lg n)
   access(x):
   while (child[x][0] \geq= 0)
     push(x = child[x][0]);
   splay(x);
   return x;
 1 // d78d
 // Add edge x-v; time: amortized O(lg n)
 // Root of tree containing v becomes
 // root of new tree.
 void link(int x, int y) {
   makeRoot(x); prev[x] = y;
 } // fb4f
 // Remove edge x-v; time: amortized O(lg n)
 // x and v become roots of new trees!
 void cut(int x, int y) {
   makeRoot(x); access(y);
   par[x] = child[y][0] = -1;
   update(v);
 1 // 1908
 // Get distance between x and y,
 // returns INF if x and y there's no path.
 // This operation makes x root of the tree!
 int dist(int x, int y) { // amortized O(lg n)
   makeRoot(x);
   if (find(y) != x) return INF;
   access(y);
   int t = child[y][0];
   return t >= 0 ? size[t] : 0;
 } // ae69
}; // 4480
util/arc interval cover.h
                                         7507
using dbl = double;
```

```
// Find size of smallest set of points
// such that each arc contains at least one
// of them; time: O(n lq n)
int arcCover(vector<pair<dbl, dbl>>& inters,
             dbl wrap) {
  int n = sz(inters);
  rep(i, 0, n) {
   auto& e = inters[i];
   e.x = fmod(e.x, wrap);
   e.y = fmod(e.y, wrap);
   if (e.x < 0) e.x += wrap, e.y += wrap;</pre>
   if (e.x > e.y) e.x += wrap;
   inters.pb({e.x+wrap, e.y+wrap});
  } // b87d
 Vi nxt(n);
  deque<dbl> que:
  dbl r = wrap*4;
  sort(all(inters));
  for (int i = n*2-1; i--;) {
   r = min(r, inters[i].v);
   que.push front(inters[i].x);
   while (!que.empty() && que.back() > r)
     que.pop_back();
    if (i < n) nxt[i] = i+sz(que);
  } // 5e6c
  int a = 0, b = 0;
  do {
   a = nxt[a] % n;
   b = nxt[nxt[b]%n] % n;
  } while (a != b);
  int ans = 0;
  while (b < a+n) {
   b += nxt[b%n] - b%n;
   ans++;
  ) // 7350
 return ans;
) // 7871
util/bit hacks.h
                                         599a
// __builtin_popcount - count number of 1 bits
// builtin clz - count most significant Os
// __builtin_ctz - count least significant 0s
// __builtin_ffs - like ctz, but indexed from 1
                  returns 0 for 0
// For 11 version add 11 to name
using ull = uint64_t;
#define T64(s.up)
 for (ull i=0; i<64; i+=s*2)
    for (ull j = i; j < i+s; j++) {
     ull \ a = (M[j] >> s) \& up;
      ull\ b = (M[j+s] \& up) << s;
     M[j] = (M[j] \& up) | b;
     M[j+s] = (M[j+s] & (up << s)) | a; 
    } // a290
// Transpose 64x64 bit matrix
void transpose64(arrav<ull, 64>& M) {
 T64(1, 0x55555555555555);
 T64(2, 0x3333333333333333);
  T64(4, 0xF0F0F0F0F0F0F0F);
  T64(8, OxFF00FF00FF00FF);
  T64(16, 0xFFFF0000FFFF);
 T64 (32, OxFFFFFFFLL);
// Lexicographically next mask with same
```

```
// amount of ones.
int nextSubset(int v) {
 int t = v | (v - 1);
 return (t + 1) | ((("t & -"t) - 1) >>
      ( builtin ctz(v) + 1));
} // 4c0c
util/bump alloc.h
                                         09f9
// Allocator, which doesn't free memory.
char mem[400<<20]; // Set memory limit</pre>
size t nMem;
void* operator new(size_t n) {
 nMem += n; return &mem[nMem-n];
} // fba6
void operator delete(void*) {}
util/compress vec.h
                                         bc5d
// Compress integers to range [0;n) while
// preserving their order; time: O(n lg n)
// Returns mapping: compressed -> original
Vi compressVec(vector<int*>& vec) {
 sort (all (vec),
   [](int* 1, int* r) { return *1 < *r; });
 Vi old;
 each(e, vec) {
   if (old.empty() || old.back() != *e)
     old.pb(*e);
   \star e = sz(old)-1;
 } // 7eb0
 return old;
} // 2b60
util/inversion vector.h
                                         0.1 \pm 9
// Get inversion vector for sequence of
// numbers in [0;n); ret[i] = count of numbers
// greater than perm[i] to the left; O(n lg n)
Vi encodeInversions (Vi perm) {
 Vi odd, ret(sz(perm));
 int cont = 1;
  while (cont) {
   odd.assign(sz(perm)+1, 0);
    cont = 0;
   rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret[i] += odd[perm[i]+1];
     cont += perm[i] /= 2;
   } // 4ed0
 } // a4f0
 return ret;
} // c2e1
// Count inversions in sequence of numbers
// in [0;n); time: O(n lq n)
ll countInversions (Vi perm) {
 ll ret = 0, cont = 1;
 Vi odd;
  while (cont) {
   odd.assign(sz(perm)+1, 0);
    cont = 0;
   rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret += odd[perm[i]+1];
     cont += perm[i] /= 2;
   } // 916f
 } // c9b5
 return ret;
```

```
} // laaf
util/longest inc subseq.h
                                         0bd1
// Longest Increasing Subsequence; O(n lg n)
int lis(const Vi& seg) {
 Vi dp(sz(seq), INT_MAX);
 each(c, seg) *lower bound(all(dp), c) = c;
 return int(lower_bound(all(dp), INT_MAX)
             - dp.begin());
1 // d0e9
util/max rects.h
                                         2a16
struct MaxRect {
 // begin = first column of rectangle
 // end = first column after rectangle
 // hei = height of rectangle
 // touch = columns of height hei inside
 int begin, end, hei;
 Vi touch; // sorted increasing
}; // 41fe
// Given consecutive column heights find
// all inclusion-wise maximal rectangles
// contained in "drawing" of columns; time O(n)
vector < MaxRect > getMaxRects (Vi hei) {
 hei.insert(hei.begin(), -1);
 hei.pb(-1);
 Vi reach(sz(hei), sz(hei)-1);
 vector < MaxRect > ans;
  for (int i = sz(hei)-1; --i;) {
   int j = i+1, k = i;
   while (hei[j] > hei[i]) j = reach[j];
   reach[i] = j;
   while (hei[k] > hei[i-1]) {
     ans.pb(\{ i-1, 0, hei[k], \{\} \});
     auto& rect = ans.back();
     while (hei[k] == rect.hei) {
       rect.touch.pb(k-1);
       k = reach[k];
     } // 6e7e
     rect.end = k-1;
   } // e03f
 } // 2796
 return ans;
} // 0e49
util/mo.h
                                         caeb
// Modified MO's queries sorting algorithm,
// slightly better results than standard.
// Allows to process q queries in O(n*sqrt(q))
struct Query {
int begin, end;
}; // b76d
// Get point index on Hilbert curve
11 hilbert(int x, int v, int s, 11 c = 0) {
 if (s <= 1) return c;</pre>
 s /= 2; c *= 4;
 if (v < s)
   return hilbert (x \in (s-1), y, s, c+(x>=s)+1);
 if (x < s)
   return hilbert (2*s-y-1, s-x-1, s, c);
 return hilbert(y-s, x-s, s, c+3);
// Get good order of gueries; time: O(n lg n)
Vi moOrder(vector<Ouery>€ queries, int maxN) {
 int s = 1;
```

```
while (s < maxN) s \star= 2;
  vector<ll> ord:
 each (q, queries)
   ord.pb(hilbert(q.begin, q.end, s));
  Vi ret(sz(ord)):
 iota(all(ret), 0);
 sort(all(ret), [&](int l, int r) {
   return ord[1] < ord[r];</pre>
 }); // 9aea
 return ret;
} // ecec
util/parallel binsearch.h
                                          fdd4
// Run `n` binary searches on [b;e) parallely.
// `cmp` should be lambda with arguments:
// 1) vector<Pii>& - pairs (v, i)
     which are queries if value for index i
     is greater or equal to v:
     pairs are sorted by v
// 2) vector<bool>& - output vector,
    set true at index i if value
      for i-th query is >= queried value
// Returns vector of found values:
// time: O((n+c) lg range), where c is cmp time
template<class T>
Vi multiBS(int b, int e, int n, T cmp) {
 if (b >= e) return Vi(n, b);
 vector<Pii> que(n), rng(n, {b, e});
 vector<bool> ans(n):
  rep(i, 0, n) que[i] = \{(b+e)/2, i\};
  for (int k = 32-__builtin_clz(e-b); k--;) {
   int last = 0, j = 0;
    cmp(que, ans);
    rep(i, 0, sz(que)) {
     Pii &q = que[i], &r = rng[q.y];
      if (q.x != last) last = q.x, j = i;
      (ans[i] ? r.x : r.y) = q.x;
      q.x = (r.x+r.y) / 2;
     if (!ans[i]) swap(que[i], que[j++]);
   } // 6c4e
 } // 622c
 Vi ret;
 each (p, rng) ret.pb(p.x);
 return ret;
} // 9e29
util/radix sort.h
                                          6fb4
// Stable countingsort; time: O(k+sz(vec))
// See example usage in radixSort for pairs.
template<class F>
void countSort(Vi& vec, F kev, int k) {
 static Vi buf, cnt;
 vec.swap(buf);
 vec.resize(sz(buf));
 cnt.assign(k+1, 0);
 each(e, buf) cnt[key(e)]++;
  rep(i, 1, k+1) cnt[i] += cnt[i-1];
 for (int i = sz(vec)-1; i >= 0; i--)
    vec[--cnt[key(buf[i])]] = buf[i];
// Compute order of elems, k is max key; O(n)
Vi radixSort (const vector < Pii> € elems, int k) {
 Vi order(sz(elems)):
 iota(all(order), 0);
 countSort (order,
```

```
[&](int i) { return elems[i].y; }, k);
  countSort (order,
    [&](int i) { return elems[i].x; }, k);
  return order;
} // e8f6
xyz/kactl.h
                                          b5ec
// --- POINT 3D
template < class T> struct Point 3D {
  typedef Point3D P;
  typedef const P& R;
  T x, y, z;
  explicit Point3D(T a=0, T b=0, T c=0)
    : x(a), y(b), z(c) {}
  bool operator<(R p) const {</pre>
    return tie(x, y, z) < tie(p.x, p.y, p.z); }</pre>
  bool operator==(R p) const {
    return tie(x, y, z) == tie(p.x, p.y, p.z);}
  P operator+(R p) const {
    return P(x+p.x, y+p.y, z+p.z); } // 6cad
  P operator-(R p) const {
    return P(x-p.x, y-p.y, z-p.z); } // 01e2
  P operator*(T d) const {
    return P(x*d, y*d, z*d); } // 071d
  P operator/(T d) const {
    return P(x/d, y/d, z/d); } // 40df
  T dot (R p) const {
    return x*p.x + y*p.y + z*p.z; } // 466c
  P cross(R p) const {
    return P (y*p.z - z*p.y, z*p.x - x*p.z,
             x \star p.y - y \star p.x);
  T dist2() const { return x*x + y*y + z*z; }
  double dist() const {
    return sqrt((double)dist2()); } // de26
  //Azimuthal angle (longitude) to x-axis
  double phi() const { // in interval [-pi, pi]
    return atan2(v, x); }
  //Zenith angle (latitude) to the z-axis
  double theta() const { // in interval [0, pi]
    return atan2(sqrt(x*x+y*y),z); } // ed29
  P unit() const {
    return *this/(T) dist(); } //makes dist()=1
  //returns unit vector normal to *this and p
  P normal (P p) const {return cross(p).unit();}
  //returns point rotated 'angle' radians
  // ccw around axis
  P rotate (double angle, P axis) const {
    double s = sin(angle), c = cos(angle);
    P u = axis.unit():
    return u*dot (u) * (1-c) +
      (*this) *c - cross(u) *s;
  } // bc6c
}; // ce68
// --- HULL 3D (requires POINT 3D) O(n^2)
typedef Point3D<double> P3:
struct PR {
  void ins(int x) { (a == -1 ? a : b) = x; }
  void rem(int x) { (a == x ? a : b) = -1; }
  int cnt() { return (a != -1) + (b != -1); }
  int a, b;
}; // 8ad1
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vector<P3>& A) {
  assert (sz(A) >= 4);
  vector<vector<PR>>> E(sz(A), vector<PR>(sz(A),
```

```
\{-1, -1\}));
#define E(x,y) E[f.x][f.y]
 vector<F> FS;
  auto mf = [&] (int i, int j, int k, int l) {
   P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
    if (q.dot(A[1]) > q.dot(A[i]))
     q = q \star -1;
    F f{q, i, j, k};
    E(a,b).ins(k); E(a,c).ins(j);
    E(b,c).ins(i);
    FS.push_back(f);
  }; // 51be
  rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)
   mf(i, j, k, 6 - i - j - k);
  rep(i,4,sz(A)) {
    rep(j, 0, sz(FS)) {
      F f = FS[i];
      if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
        E(a,b).rem(f.c);
        E(a,c).rem(f.b);
        E(b,c).rem(f.a);
        swap(FS[j--], FS.back());
        FS.pop_back();
     } // 5d69
    } // eeb6
    int nw = sz(FS);
    rep(j,0,nw) {
     F f = FS[i];
#define C(a, b, c) if (E(a,b),cnt() != 2)
        mf(f.a, f.b, i, f.c);
      C(a, b, c); C(a, c, b); C(b, c, a);
   ) // 9578
  } // c66b
  for (F& it : FS) if ((A[it.b]-A[it.a]).cross(
   A[it.c] - A[it.a]).dot(it.q) <= 0)
      swap(it.c, it.b);
 return FS:
}; // 2116
// --- DELAUNAY (requires POINT 3D and HULL 3D)
// Calls trifun for every triangle
template<class P, class F> // O(n^2)
void delaunay(vector<P>& ps, F trifun) {
  if (sz(ps) == 3) { int d =
    ((ps[1]-ps[0]).cross(ps[2]-ps[0]) < 0);
   trifun(0,1+d,2-d); } // 1266
  vector<P3> p3;
  for (P p : ps)
   p3.emplace_back(p.x, p.y, p.len2());
  if (sz(ps) > 3) for (auto t:hull3d(p3))
    if ((p3[t.b]-p3[t.a]).
      cross(p3[t.c]-p3[t.a]).dot(P3(0,0,1)) < 0)
    trifun(t.a, t.c, t.b);
} // 95dc
// --- FAST DELAUNAY (requires our vec2)
// O(n log n)
#include "../geometry/vec2.h"
typedef vec2i P;
typedef struct Quad* Q;
// (can be 11 if coords are < 2e4)
typedef __int128_t 111;
// not equal to any other point
P arb(LLONG_MAX, LLONG_MAX);
struct Ouad {
 Q rot, o; P p = arb; bool mark;
  Quad(Q q) {rot=q;}
  P& F() { return r()->p; }
```

```
Q& r() { return rot->rot; }
 Q prev() { return rot->o->rot; }
 Q next() { return r()->prev(); }
ll cross(P x, P a, P b) {
return (a-x).cross(b-x);
1 // 0647
bool circ(P p, P a, P b, P c) {
 111 p2 = p.len2(), A = a.len2()-p2,
      B = b.len2()-p2, C = c.len2()-p2;
  return cross(p,a,b) *C + cross(p,b,c) *A +
    cross (p,c,a)*B > 0;
Q makeEdge (P orig, P dest) {
 Q r = H ? H : new Quad{new Quad{new Quad{
    new Quad{0}}}}; // 08c9
  H = r \rightarrow 0; r \rightarrow r() \rightarrow r() = r;
  rep(i,0,4) r = r->rot, r->p = arb,
    r \rightarrow 0 = i & 1 ? r : r \rightarrow r();
  r\rightarrow p = orig; r\rightarrow F() = dest;
 return r:
} // 8da8
void splice(Q a, Q b) {
 swap(a->o->rot->o, b->o->rot->o);
 swap (a->0, b->0);
} // 6867
0 connect(0 a, 0 b) {
 Q q = makeEdge(a->F(), b->p);
 splice(g, a->next());
  splice(q->r(), b);
 return q;
} // af0b
pair<0,0> rec(const vector<P>& s) {
 if (sz(s) \le 3) {
    Q = makeEdge(s[0], s[1]), b =
      makeEdge(s[1], s.back());
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = cross(s[0], s[1], s[2]);
    0 c = side ? connect(b, a) : 0;
    return {side < 0 ? c->r() : a, side < 0 ?
      c : b->r() }; // d148
 } // d2fb
#define H(e) e^{->F()}, e^{->p}
\#define\ valid(e)\ (cross(e->F(),\ H(base))\ >\ 0)
 O A, B, ra, rb;
 int half = sz(s) / 2;
  tie(ra, A) = rec({all(s) - half});
  tie(B, rb) = rec(\{sz(s) - half + all(s)\});
  while ((cross(B->p, H(A)) < 0 &&
    (A = A->next()))
         (cross(A->p, H(B)) > 0 \& \&
            (B = B->r()->o)));
  Q base = connect(B->r(), A);
  if (A\rightarrow p == ra\rightarrow p) ra = base->r();
  if (B\rightarrow p == rb\rightarrow p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; \
 if (valid(e)) \
    while (circ(e->dir->F(),H(base),e->F())) {\
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e->o = H; H = e; e = t; 
    } // a3f9
  for (;;) {
    DEL(LC, base->r(), o); DEL(RC,base,prev());
```

```
if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) &&
        circ(H(RC), H(LC))))
      base = connect(RC, base->r());
      base = connect(base->r(), LC->r());
 } // 53db
  return { ra, rb };
} // 4baa
vector<P> triangulate(vector<P> pts) {
  sort(all(pts)); assert(unique(all(pts))
    == pts.end());
  if (sz(pts) < 2) return {};</pre>
  Q e = rec(pts).first;
  vector < Q > q = {e};
  int qi = 0;
  while (cross(e->o->F(), e->F(), e->p)
    < 0) e = e->0;
#define ADD { Q c = e; do { c->mark = 1; \
    pts.push_back(c->p); \
    q.push\_back(c->r()); c = c->next(); 
    while (c != e); } // 889e
  ADD; pts.clear();
  while (qi < sz(q))</pre>
   if (!(e = q[qi++])->mark) ADD;
 return pts;
} // ef04
// --- CIRCUMCIRCLE (requires our vec2)
#include "../geometry/vec2.h"
typedef vec2d P;
double ccRadius (const P& A, const P& B,
    const P& C) {
  return (B-A).len() * (C-B).len() *
    (A-C).len()/abs((B-A).cross(C-A))/2;
} // 95ee
P ccCenter (const P& A, const P& B, const P& C) {
 P b = C-A, c = B-A;
  return A + (b*c.len2()-c*b.len2()).perp()
    / b.cross(c)/2;
1 // 2880
// --- MINIMUM ENCLOSING CIRCLE
// requires CIRCUMCIRCLE; expected O(n)
pair<P, double> mec(vector<P> ps) {
  shuffle(all(ps), mt19937(uint32_t(time(0))));
  P \circ = ps[0];
  double r = 0, EPS = 1 + 1e-8;
  rep(i,0,sz(ps)) if ((o - ps[i]).len() >
     r * EPS) {
    o = ps[i], r = 0;
    rep(j,0,i) if ((o - ps[j]).len() >
       r * EPS) {
      o = (ps[i] + ps[j]) / 2;
      r = (o - ps[i]).len();
      rep(k,0,j) if ((o - ps[k]).len() >
          r * EPS) {
        o = ccCenter(ps[i], ps[j], ps[k]);
       r = (o - ps[i]).len();
      } // b992
    } // 942d
  } // Oa5c
  return {o, r};
} // faa0
// --- POLYGON CENTER OF MASS
typedef vec2d P; // (requires our vec2d)
P polygonCenter(const vector<P>& v) {
 P res(0, 0); double A = 0;
```

```
for (int i = 0, j = sz(v) - 1;
     i < sz(v); j = i++) {
    res = res + (v[i]+v[j]) * v[j].cross(v[i]);
   A += v[j].cross(v[i]);
  } // eaf6
  return res / A / 3;
} // 405a
xyz/uj.h
                                          33d1
// P = punkt 2D lub 3D, K = double
// --- GEOMETRIA OKREGOW W 2D ~Bartosz Walczak
struct circle { // okrag w 2D
  P c; K r; // srodek, promien
  circle(const P &ci, K ri=0) : c(ci),r(ri) {}
  circle() {}
  K length() const { return 2*M_PI*r; }
  K area() const { return M_PI*r*r; }
}; // 2a8d
// Czy punkt lezy na okregu?
bool on_circle (const P &a, const circle &c)
  { return fabs((a-c.c).norm()-c.r*c.r) < EPS;}
// Czy kolo/punkt lezy wewnatrz lub na brzegu?
bool operator < (const circle &a, const circle &b)
  { return b.r+EPS > a.r && (a.c-b.c).norm() <
  (b.r-a.r)*(b.r-a.r)+EPS; } // 1f58
// Srodek okragu opisanego na trojkacie
circle circumcircle (P a, P b, P c) {
  if ((a-b).norm() > (c-b).norm()) swap(a,c);
  if ((b-c).norm() > (a-c).norm()) swap(a,b);
  if (fabs(det(b-a, c-b)) < EPS)</pre>
    throw "zdegenerowany";
  P v=intersection(median(a, b), median(b,c));
  return circle(v, sqrt((a-v).norm()));
} // 1327
// Przeciecie okregu i prostej.
// Zwraca liczbe punktow
int intersection(const circle &c,
      const line &p, P I[]/*OUT*/) {
  K d = p.n.norm(), a = (p.n*c.c-p.c)/d;
  P u = c.c-p.n*a; a *= a; K r = c.r*c.r/d;
  if (a >= r+EPS) return 0;
  if (a > r-EPS) { I[0]=u; return 1; }
  K h = sgrt(r-a):
  I[0] = u+cross(p.n)*h; I[1] = u-cross(p.n)*h;
  return 2;
} // 51f4
// Przeciecie dwoch okregow.
// Zwraca liczbe punktow. Zalozenie: c1.c!=c2.c
int intersection (const circle &c1,
    const circle &c2, P I[]/*OUT*/) {
  K d = (c2.c-c1.c).norm(), r1 =
    c1.r*c1.r/d, r2 = c2.r*c2.r/d;
  P u = c1.c*((r2-r1+1)*0.5) +
    c2.c*((r1-r2+1)*0.5);
  if (r1 > r2) swap(r1, r2);
  K a = (r1-r2+1)*0.5; a *= a;
  if (a >= r1+EPS) return 0;
  if (a > r1-EPS) { I[0] = u; return 1; }
  P v = cross(c2.c - c1.c); K h = sqrt(r1-a);
  I[0] = u+v*h; I[1] = u-v*h; return 2;
1 // 45ea
// --- GEOMETRIA 3D ~Bartosz Walczak
// Kat pomiedzy dwoma wektorami. Zawsze >=0.
// Zalozenie: a.b!=0
K angle3 (const xyz &a, const xyz &b)
  { return atan2(sqrt(cross(b,a).norm()),a*b);}
struct plane { // plaszczyzna {v: n*v=c}
 xyz n; K c; // (n - wektor normalny)
```

```
plane(const xyz &ni, K ci) : n(ni), c(ci) {}
  plane() {}
}; // 93a3
// Czy punkt lezy na plaszczyznie?
bool on plane (const xyz &a, const plane &p)
 { return fabs(p.n*a - p.c) < EPS; }
// Plaszczyzna rozpieta przez 3 punkty ccw.
// Zalozenie: a,b,c niezalezne
plane span3 (const xyz &a, const xyz &b,
    const xyz &c)
  { xyz n = cross(c-a, b-a);
    return plane(n, n*a); } // 2d8e
// Plaszczyzna symetralna odcinka.
// Zalozenie: a!=b
plane median3(const xyz &a, const xyz &b)
{ return plane(b-a, (b-a)*(b+a)*0.5); }
// Plaszczyzna rownolegla przez punkt
plane parallel3 (const xyz &a, const plane &p)
 { return plane(p.n, p.n*a); }
// Odleglosc punktu od plaszczyzny
K dist3(const xyz &a, const plane &p)
  { return fabs(p.n*a-p.c)/sqrt(p.n.norm()); }
struct line3 { // prosta {v: cross(v, u) = w}
    xyz u, w; // (u - wektor kierunku)
    // UWAGA! konstruktor dwuargumentowy
    // nie tworzy prostej przechodzacej
    // przez 2 punkty,
    // w tym celu nalezy uzyc span3!
    line3(const xyz &ui, const xyz &wi)
      : u(ui), w(wi) {}
    line3() {}
}; // 8ba4
// Czy punkt lezy na prostej?
bool on line3 (const xyz &a, const line3 &p)
 { return (cross(a,p.u)-p.w).norm() < EPS; }
// Prosta rozpieta przez 2 punkty. Za. l.: a!=b
line3 span3 (const xyz &a, const xyz &b)
 { return line3(b-a, cross(a, b-a)); }
// Plaszczyzna rozpieta przez prosta i
// punkt ccw. Zalozenie: cross(a,p.u)!=p.w
plane span3 (const line3 &p, const xyz &a)
 { return plane(cross(a,p.u)-p.w, p.w*a); }
// Prosta przeciecia dwoch plaszczyzn
line3 intersection3 (const plane &p,
    const plane &q) {
  xyz u=cross(q.n, p.n);
  if (u.norm() < EPS) throw "rownolegle";</pre>
  return line3(u, q.n*p.c-p.n*q.c);
1 // 08c7
// Punkt przeciecia plaszczyzny i prostej
xyz intersection3 (const plane &p,
   const line3 &q) {
  K d = q.u*p.n;
  if (fabs(d) < EPS) throw "rownolegle";</pre>
  return (q.u*p.c + cross(p.n, q.w))/d;
} // 7f03
// Prosta prostopadla do plaszczyzny
// przechodzaca przez punkt
line3 perp3 (const xyz &a, const plane &p)
 { return line3(p.n, cross(a, p.n)); }
// Plaszczyzna prostopadla do prostej
// przechodzaca przez punkt
plane perp3 (const xyz &a, const line3 &p)
  { return plane(p.u, p.u*a); }
// Odleglosc punktu od prostej
K dist3(const xyz &a, const line3 &p)
  { return sqrt((cross(a,p.u)-p.w).norm()) /
    sgrt (p.u.norm()); } // a713
```

```
// Odleglosc 2 prostych od siebie.
// Zalozenie: cross(q.u,p.u)!=0
// (niestabilne przy bliskim 0)
K dist3(const line3 &p, const line3 &q)
 { return fabs(p.u*q.w + q.u*p.w) /
   sqrt(cross(q.u, p.u).norm()); } // 88d8
// --- GEOMETRIA SFER W 3D ~Bartosz Walczak
struct sphere {
 xyz c; K r; // srodek, promien
 sphere(const xyz &ci, K ri=0)
   : c(ci), r(ri) {}
 sphere() {}
 // pole powierzchni
 K area() const { return 4*M_PI*r*r; }
 // objetosc kuli
 K volume() const { return 4*M_PI*r*r*r/3; }
}; // 029e
// Czy punkt lezy na sferze?
bool on_sphere(const xyz &a, const sphere &s)
{ return fabs((a-s.c).norm()-s.r*s.r) < EPS;}
// Czy sfera/punkt lezy wewnatrz lub na brzegu?
bool in sphere (const sphere &a, const sphere &b)
 { return b.r+EPS > a.r && (a.c-b.c).norm() <
    (b.r-a.r)*(b.r-a.r)*EPS; } // 1f58
// Przeciecie sfery i prostej.
// Zwraca liczbe punktow przeciecia
int intersection3 (const sphere &s,
 const line3 &p, xvz I[]/*OUT*/) {
 K d = p.u.norm(), a = (cross(s.c.p.u)-p.w).
   norm()/(d*d), r = s.r*s.r/d;
 if (a >= r+EPS) return 0;
 xyz u = (p.u*(p.u*s.c)+cross(p.u,p.w))/d;
 if (a > r-EPS) { I[0] = u; return 1; }
 K h = sgrt(r-a);
 I[0] = u+p.u*h; I[1] = u-p.u*h; return 2;
} // 14d7
// Przeciecie sfery i plaszczyzny.
// Zwraca true, jesli sie przecinaja. Wtedy u,r
// sa odp. srodkiem i promieniem okregu
// przeciecia. Zalozenie: s1.c!=s2.c
bool intersection3(const sphere &s,
   const plane &p, xyz &u, K &r) {
 K d = p.n.norm(), a = (p.n*s.c-p.c)/d;
 u = s.c-p.n*a; a *= a; K r1 = s.r*s.r/d;
 if (a >= r1+EPS) return false;
 r = a > r1-EPS ? 0 : sqrt(r1-a)*sqrt(d);
 return true;
1 // d90d
// Przeciecie dwoch sfer.
// Zwraca true, jesli sie przecinaja.
// Wtedy u,r sa odp. srodkiem i promieniem
// okregu przeciecia. Zalozenie: s1.c!=s2.c
bool intersection3 (const sphere &s1,
   const sphere &s2, xyz &u, K &r) {
 K d = (s2.c-s1.c).norm(), r1 = s1.r*s1.r/d,
   r2 = s2.r \star s2.r/d;
 u = s1.c*((r2-r1+1)*0.5) +
   s2.c*((r1-r2+1)*0.5);
 if (r1 > r2) swap(r1, r2);
 K = (r1-r2+1)*0.5; a *= a;
 if (a >= r1+EPS) return false;
 r = a > r1-EPS ? 0 : sqrt(r1-a)*sqrt(d);
 return true;
} // db4b
// --- GEOMETRIA NA SFERZE ~Bartosz Walczak
// Odleglosc dwoch punktow na sferze
K distS(const xyz &a, const xyz &b)
```

```
{ return atan2(sqrt(cross(b,a).norm()),a*b);}
struct circleS { // okrag na sferze
 xyz c; K r; // srodek, promien katowy
 circleS(const xyz &ci, K ri) : c(ci), r(ri) {}
 circleS() {}
 K area() const { return 2*M_PI*(1-cos(r)); }
}; // faf2
// Okrag rozpiety przez 3 punkty.
// Zalozenie: punkty sa parami rozne
circleS spanS(xyz a, xyz b, xyz c) {
    int tmp = 1:
    if ((a-b).norm() > (c-b).norm())
      { swap(a, c); tmp = -tmp; }
    if ((b-c).norm() > (a-c).norm())
     { swap(a, b); tmp = -tmp; }
    xyz v = cross(c-b, b-a);
    v = v*(tmp/sqrt(v.norm()));
    return circleS(v, distS(a,v));
// Przeciecie 2 okregow na sferze.
// Zalozenie: cross(c2.c,c1.c)!=0
int intersectionS(const circleS &cl,
    const circleS &c2, xyz I[]/*OUT*/) {
 xyz n = cross(c2.c, c1.c),
    w = c2.c*cos(c1.r)-c1.c*cos(c2.r);
 K d = n.norm(), a = w.norm()/d;
 if (a >= 1+EPS) return 0;
 xyz u = cross(n, w)/d;
 if (a > 1-EPS) { I[0] = u; return 1; }
 K h = sqrt(1-a)/sqrt(d);
 I[0] = u+n*h; I[1] = u-n*h; return 2;
} // cbcb
```

Tutte Matrix. For a simple undirected graph G, Let M be a matrix with entries $A_{i,j} = 0$ if $(i,j) \notin E$ and $A_{i,j} = -A_{j,i} = X$ if $(i,j) \in E$. X could be any random value. If the determinants are non-zero, then a perfect matching exists, while other direction might not hold for very small probability.

Kirchhoff's Theorem. For a multigraph G with no loops, define Laplacian matrix as L = D - A. D is a diagonal matrix with $D_{i,i} = deg(i)$, and A is an adjacency matrix. If you remove any row and column of L, the determinant gives a number of spanning trees.

Burnside's lemma / Pólya enumeration theorem, let G and H be groups of permutations of finite sets X and Y. Let $c_m(g)$ denote the number of cycles of length m in $g \in G$ when permuting X. The number of colorings of X into |Y| = n colors with exactly r_i occurrences of the i-th color is the coefficient of $w_1^{r_1} \dots w_n^{r_n}$ in the following polynomial:

$$P(w_1, \dots, w_n) = \frac{1}{|H|} \sum_{h \in H} \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (\sum_{h^m(b) = b} (w_b^m))^{c_m(g)}$$

When
$$H = \{I\}$$
 (No color permutation):
$$P(w_1, \dots, w_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m>1} (w_1^m + \dots + w_n^m)^{c_m(g)}$$

Total number of partitions of n distinct elements. B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$ For p prime,

Without the occurrence restriction:

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

$$P(1,...,1) = \frac{1}{|G|} \sum_{g \in G} n^{c(g)}$$

where c(g) could also be interpreted as the number of elements in X that are fixed up to g.

Pick's Theorem. $A = i + \frac{b}{2} - 1$, where: P is a simple polygon whose vertices are grid points, A is area of P, i is # of grid points in the interior of P, and b is # of grid points on the boundary of P. If h is # of holes of P (h + 1 simple closed curves in total), $A = i + \frac{b}{2} + h - 1$.

Xudyh Sieve.
$$F(n) = \sum_{d|n} f(d)$$

$$S(n) = \sum_{i \le n} f(i) = \sum_{i \le n} F(i) - \sum_{d=2}^{n} S\left(\left\lfloor \frac{n}{d} \right\rfloor\right)$$

Preprocess S(1) to S(M) (Set $M = n^{\frac{2}{3}}$ for complexity)

$$S(n) = \sum f(i) = \sum_{i \le n} \left[F(i) - \sum_{j|i,j \ne i} f(j) \right] = \sum F(i) - \sum_{i/j = d = 2}^{n} \sum_{dj \le n} f(j)$$

$$S(n) = \sum_{i \leq n} i \left[F(i) - \sum_{j|i,j \neq i} f(j) \right] = \sum_{i \leq n} i F(i) - \sum_{i/j=d=2}^{n} \sum_{dj \leq n} dj f(j)$$

$$\sum_{d|n} \varphi(d) = n \qquad \sum_{d|n} \mu(d) = \text{if } (n > 1) \text{ then } 0 \text{ else } 1 \qquad \sum_{d|n} (\mu(\frac{n}{d}) \sum_{e|d} f(e)) = f(n)$$

Labeled unrooted trees

on n vertices: n^{n-2}

on k existing trees of size n_i : $n_1 n_2 \cdots n_k n^{k-2}$

with degrees d_i : $(n-2)!/((d_1-1)!\cdots(d_n-1)!)$

$$x_1 x_2 x_3 \dots x_n (x_1 + x_2 + \dots x_n)^{n-2} = \sum_{T} x_1^{d_{T(1)}} x_2^{d_{T(2)}} \dots x_n^{d_{T(n)}}$$

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

where the sum is over all spanning trees T in K_n and $d_{T(i)}$ is the degree of i in T

Taylor's theorem^{[4][5][6]} — Let $k \ge 1$ be an integer and let the function $f: \mathbb{R} \to \mathbb{R}$ be k times differentiable at the point $a \in \mathbb{R}$.

Then there exists a function $h_k: \mathbf{R} \to \mathbf{R}$ such that

$$f(x) = f(a) + f'(a)(x-a) + rac{f''(a)}{2!}(x-a)^2 + \dots + rac{f^{(k)}(a)}{k!}(x-a)^k + h_k(x)(x-a)^k,$$

Jesli
$$f(n) = \sum_{d|n} g(d)$$
, to $g(n) = \sum_{d|n} \mu(d) f\left(\frac{n}{d}\right)$

gdzie
$$\mu(1) = 1$$
, $\mu(p^2 \cdot a) = 0$, $\mu(p_1 \cdot p_2 \cdot \ldots \cdot p_k) = (-1)^k$ dla p , p_i pierwszych.

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^{2}}{8} + \frac{2x^{3}}{32} - \frac{5x^{4}}{128} + \dots, (-1 \le x \le 1)$$

$$\log \frac{1}{1-x} = \sum_{n \ge 0} x^{n} \qquad \sin x = \sum_{n \ge 0} (-1)^{n} \frac{x^{2n+1}}{(2n+1)!} \qquad \tan^{-1} x = \sum_{n \ge 0} (-1)^{n} \frac{x^{2n+1}}{2n+1} \qquad \frac{1}{\sqrt{1-4x}} \left(\frac{1-\sqrt{1-4x}}{2x}\right)^{k} = \sum_{n} \left(\frac{2n+k}{n}\right) x^{n}$$

$$\log \frac{1}{1-x} = \sum_{n \ge 1} \frac{x^{n}}{n} \qquad \cos x = \sum_{n \ge 1} (-1)^{n} \frac{x^{2n+1}}{(2n+1)!} \qquad \tan^{-1} x = \sum_{n \ge 0} (-1)^{n} \frac{x^{2n+1}}{2n+1} \qquad \frac{1}{\sqrt{1-4x}} \left(\frac{1-\sqrt{1-4x}}{2x}\right)^{k} = \sum_{n} \left(\frac{2n+k}{n}\right) x^{n}$$

$$\log \frac{1}{1-x} = \sum_{n \ge 1} \frac{x^{n}}{n} \qquad \cos x = \sum_{n \ge 1} (-1)^{n} \frac{x^{2n+1}}{(2n+1)!} \qquad \frac{1}{2x} (1-\sqrt{1-4x}) = \sum_{n} \frac{1}{n+1} \left(\frac{2n}{n}\right) x^{n} \qquad \frac{1}{\sqrt{1-4x}} = \sum_{n \ge 1} \left(\frac{2k}{n}\right) x^{n}$$

Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1 \qquad 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$
$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6} \qquad 1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$
$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{p}$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan\frac{\alpha+\beta}{2}}{\tan\frac{\alpha-\beta}}$ Law of sines: $\frac{\sin\alpha}{a} = \frac{\sin\beta}{b} = \frac{\sin\gamma}{c} = \frac{1}{2R}$ Law of cosines: $a^2 = b^2 + c^2 - 2bc\cos\alpha$

Length of median

(divides triangle into two equal-area triangles):

 $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(rac{a}{b+c}
ight)^2
ight]}$$

Symbol Newtona:

$$\sum_{k=0}^{n} \binom{n}{k}^2 = \binom{2n}{n} \qquad \sum_{k=1}^{n} k \binom{n}{k}^2 = n2^{n-1} \qquad \sum_{k=0}^{n} \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{m+n}$$

Liczby Stirlinga I rodzaju (liczba permutacji n elementow o k cyklach):

$$\begin{bmatrix} {n+1} \\ k \end{bmatrix} = n \begin{bmatrix} n \\ k \end{bmatrix} + \begin{bmatrix} n \\ k-1 \end{bmatrix} \qquad \qquad \sum_{p=k}^n \begin{bmatrix} n \\ p \end{bmatrix} \binom{p}{k} = \begin{bmatrix} {n+1} \\ k+1 \end{bmatrix}$$

Liczby Stirlinga II rodzaju (liczba podzialow zbioru n-elementowego na k klas)

Jesli kazde 2 elementy zbioru musza byc odlegle o co najmniej d:

$$S^{d}(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$$