argono			
.bashrc	1	structures/hull_online.h	45
.vimrc	2	structures/max_queue.h	46
template.cpp	3	structures/pairing_heap.h	47
<pre>geometry/convex_hull.h</pre>	4	structures/rmq.h	48
<pre>geometry/convex_hull_dist.h</pre>	5	structures/segment_tree.h	49
<pre>geometry/convex_hull_sum.h</pre>	6	structures/segment_tree_beats.	h 50
geometry/line2.h	7	structures/segment_tree_point.	h 51
geometry/rmst.h	8	structures/treap.h	52
geometry/segment2.h	9	structures/ext/hash_table.h	53
geometry/vec2.h	10	structures/ext/rope.h	54
graphs/2sat.h	11	structures/ext/tree.h	55
graphs/bellman_inequalities.h	12	structures/ext/trie.h	56
graphs/bridges_online.h	13	text/aho_corasick.h	57
graphs/dense_dfs.h	14	text/kmp.h	58
graphs/edmonds_karp.h	15	text/kmr.h	59
graphs/push_relabel.h	16	text/lcp.h	60
graphs/scc.h	17	text/lyndon_factorization.h	61
graphs/turbo_matching.h	18	text/manacher.h	62
math/bit_gauss.h	19	text/min_rotation.h	63
math/bit_matrix.h	20	text/palindromic_tree.h	64
math/crt.h	21	text/suffix_array_linear.h	65
math/discrete_logarithm.h	22	text/z_function.h	66
math/fft_complex.h	23	trees/centroid_decomp.h	67
math/fft_mod.h	24	trees/heavylight_decomp.h	68
math/gauss.h	25	trees/lca.h	69
math/miller_rabin.h	26	trees/link_cut_tree.h	70
math/modinv_precompute.h	27	util/bit_hacks.h	71
math/modular.h	28	util/bump_alloc.h	72
math/modular64.h	29	util/compress_vec.h	73
math/montgomery.h	30	util/inversion_vector.h	74
math/nimber.h	31	util/longest_increasing_sub.h	75
math/phi_large.h	32	util/max_rects.h	76
math/phi_precompute.h	33	util/mo.h	77
math/pi_large_precomp.h	34	util/parallel_binsearch.h	78
math/pollard_rho.h	35	util/radix_sort.h	79
math/polynomial_interp.h	36		
math/sieve.h	37		
math/sieve_factors.h	38		
math/sieve_segmented.h	39		
structures/bitset_plus.h	40		
structures/fenwick_tree.h	41		
structures/fenwick tree 2d.h	42		
structures/find_union.h	43		
structures/hull_offline.h	44		
011110.11	11		
		I	- 1

```
.bashrc
b()(
 q++ $0 -o $1.e -DLOC -O2 -q -std=c++11
      -Wall -Wextra -Wfatal-errors -Wshadow \
      -Wlogical-op -Wconversion -Wfloat-equal
d()(b $0 -fsanitize=address.undefined \
          -D GLIBCXX DEBUG )
cmp()(
  set -e; $1 $2; $1 $3; $1 $4
  for ((::)) {
    ./$4.e > gen.in;
                             echo -n 0
    ./$2.e < gen.in > p1.out; echo -n 1
    ./$3.e < gen.in > p2.out; echo -n 2
   diff pl.out p2.out;
                             echo -n Y
# Other flags:
# -Wformat=2 -Wshift-overflow=2 -Wcast-qual
# -Wcast-align -Wduplicated-cond
# -D GLIBCXX DEBUG PEDANTIC -D FORTIFY SOURCE=2
# -fno-sanitize-recover -fstack-protector
.vimrc
se ai aw cin cul ic is nocp nohls nu rnu sc scs
se bg=dark so=7 sw=4 ttm=9 ts=4
sy on
colo delek
template.cpp
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
using Vi = vector<int>;
using Pii = pair<int,int>;
#define pb push_back
#define x first
#define y second
#define rep(i,b,e) for(int i=(b); i<(e); i++)
#define each (a,x) for (auto& a : (x))
#define all(x)
                   (x).begin(),(x).end()
#define sz(x)
                   int((x).size())
int main() {
  cin.sync_with_stdio(0); cin.tie(0);
  cout << fixed << setprecision(18);
  return 0;
// > Debug printer
#define tem template<class t,class u,class...w>
#define pri(x,y)tem auto operator<<(t& o,u a) \
  ->decltype(x,o) { o << v; return o; }
pri(a.print(), "{"; a.print(); o << "}")</pre>
pri(a.v, "(" << a.x << ", " << a.v << ")")
pri(all(a), "["; auto d=""; for (auto i : a)
```

```
(o << d << i, d = ", "); o << "]")
void DD(...) {}
tem void DD(t s, u a, w... k) {
 int b = 44;
 while (*s && *s != b) {
   b += (*s == 40 ? 50 : *s == 41 ? -50 : 0);
   cerr << *s++;
 cerr << ": " << a << *s++; DD(s, k...);
tem vector<t> span(const t* a, u n) {
 return {a, a+n};
#ifdef LOC
#define deb(...) (DD("#, "# VA ARGS , \
  __LINE__, __VA_ARGS__), cerr << endl)
#else
#define deb(...)
#endif
#define DBP(...) void print() { \
 DD (#__VA_ARGS__, __VA_ARGS__); }
// > Utils
// #pragma GCC optimize("Ofast, unroll-loops,
                         no-stack-protector")
// #pragma GCC target("avx")
// while (clock() < time*CLOCKS_PER_SEC)</pre>
// using namespace rel ops;
// Return smallest k such that 2^k > n
// Undefined for n = 0!
int uplg(int n) { return 32-__builtin_clz(n); }
int uplq(ll n) { return 64- builtin clzll(n); }
// Compare with certain epsilon (branchless)
// Returns -1 if a < b; 1 if a > b; 0 if equal
// a and b are assumed equal if |a-b| <= eps
int cmp(double a, double b, double eps=1e-10) {
 return (a > b+eps) - (a+eps < b);
geometry/convex hull.h
#include "vec2.h"
// Translate points such that lower-left point
// is (0, 0). Returns old point location; O(n)
vec2 normPos(vector<vec2>& points) {
 auto g = points[0].vxPair();
 each(p, points) q = min(q, p.yxPair());
 vec2 ret{q.y, q.x};
 each (p, points) p = p-ret;
 return ret;
// Find convex hull of points; time: O(n la n)
// Points are returned counter-clockwise.
vector<vec2> convexHull(vector<vec2> points) {
 vec2 pivot = normPos(points);
 sort (all (points));
 vector<vec2> hull:
```

```
// Returns start position.
 each(p, points) {
   while (sz(hull) >= 2) {
     vec2 = hull.back() - hull[sz(hull)-2];
     vec2 b = p - hull.back();
     if (a.cross(b) > 0) break;
     hull.pop back();
   hull.pb(p);
  // Translate back, optional
 each(p, hull) p = p+pivot;
 return hull:
geometry/convex hull dist.h
#include "vec2.h"
// Check if p is inside convex polygon. Hull
// must be given in counter-clockwise order.
// Returns 2 if inside, 1 if on border,
// 0 if outside; time: O(n)
int insideHull(vector<vec2>& hull, vec2 p) {
 int ret = 1:
 rep(i, 0, sz(hull)) {
   auto v = hull[(i+1)%sz(hull)] - hull[i];
   auto t = v.cross(p-hull[i]);
   ret = min(ret, cmp(t, 0)); // For doubles
    //ret = min(ret, (t>0) - (t<0)); // Ints
 return int(max(ret+1, 0));
#include "segment2.h"
// Get distance from point to hull; time: O(n)
double hullDist(vector<vec2>& hull, vec2 p) {
 if (insideHull(hull, p)) return 0;
 double ret = 1e30;
 rep(i, 0, sz(hull)) {
   seq2 seq{hull[(i+1)%sz(hull)], hull[i]};
   ret = min(ret, seg.distTo(p));
 return ret;
// Compare distance from point to hull
// with sqrt(d2); time: O(n)
// -1 if smaller, 0 if equal, 1 if greater
int cmpHullDist(vector<vec2>& hull,
               vec2 p, 11 d2) {
 if (insideHull(hull,p)) return (d2<0)-(d2>0);
 int ret = 1;
 rep(i, 0, sz(hull)) {
   seg2 seg{hull[(i+1)%sz(hull)], hull[i]};
   ret = min(ret, seg.cmpDistTo(p, d2));
 return ret;
                                                 };
geometry/convex hull sum.h
#include "vec2.h"
// Get edge sequence for given polygon
// starting from lower-left vertex; time: O(n)
```

```
vec2 edgeSeg(vector<vec2> points,
             vector<vec2>& edges) {
  int i = 0, n = sz(points);
  rep(i, 0, n) {
    if (points[i].yxPair()>points[j].yxPair())
      i = j;
  rep(j, 0, n) edges.pb(points[(i+j+1)%n] -
                        points[(i+j)%n]);
 return points[i];
// Minkowski sum of given convex polygons.
// Vertices are required to be in
// counter-clockwise order; time: O(n+m)
vector<vec2> hullSum(vector<vec2> A.
                     vector<vec2> B) {
  vector\langle vec2 \rangle sum, e1, e2, es(sz(A) + sz(B));
  vec2 pivot = edgeSeq(A, e1) + edgeSeq(B, e2);
  merge(all(e1), all(e2), es.begin());
  sum.pb(pivot);
  each(e, es) sum.pb(sum.back() + e);
  sum.pop_back();
  return sum;
geometry/line2.h
#include "vec2.h"
// 2D line structure: PARTIALLY TESTED
// Base class of versions for ints and doubles
template < class T, class P, class S>
struct bline2 { // norm*point == off
 P norm; // Normal vector [A; B]
 T off; // Offset (C parameter of equation)
  // Line through 2 points
  static S through (P a, P b) {
    return { (b-a).perp(), b.cross(a) };
  // Parallel line through point
  static S parallel(P a, S b) {
    return { b.norm, b.norm.dot(a) };
  // Perpendicular line through point
  static S perp(P a, S b) {
    return { b.norm.perp(), b.norm.cross(a) };
  // Distance from point to line
  double distTo(P a) {
    return fabs(norm.dot(a)-off) / norm.len();
// Version for integer coordinates (long long)
struct line2i : bline2<11, vec2i, line2i> {
 line2i() : bline2{{}, 0} {}
 line2i(vec2i n, 11 c) : bline2{n, c} {}
 int side(vec2i a) {
```

};

11 d = norm.dot(a);

return (d > off) - (d < off);</pre>

11

```
// Version for double coordinates
// Requires cmp() from template
struct line2d : bline2<double, vec2d, line2d> {
  line2d() : bline2{{}, 0} {}
  line2d(vec2d n, double c) : bline2{n, c} {}
  int side(vec2d a) {
   return cmp(norm.dot(a), off);
  bool intersect(line2d a, vec2d& out) {
    double d = norm.cross(a.norm);
   if (cmp(d, 0) == 0) return false;
   out = (norm*a.off-a.norm*off).perp() / d;
   return true:
};
using line2 = line2d;
geometry/rmst.h
#include "../structures/find union.h"
// Rectilinear Minimum Spanning Tree
// (MST in Manhattan metric); time: O(n lg n)
// Returns MST weight. Outputs spanning tree
// to G, vertex indices match point indices.
// Edge in G is pair (target, weight).
11 rmst(vector<Pii>% points,
        vector<vector<Pii>>& G) {
  int n = sz(points);
  vector<pair<int, Pii>> edges;
  vector<Pii> close;
  Vi ord(n), merged(n);
  iota(all(ord), 0);
  function<void(int,int)> octant =
     [&] (int begin, int end) {
    if (begin+1 >= end) return;
    int mid = (begin+end) / 2;
   octant (begin, mid);
   octant (mid, end);
    int i = mid;
   Pii best = {INT\_MAX, -1};
    merged.clear();
    rep(i, begin, mid) {
     int v = ord[i];
     Pii p = points[v];
      while (j < end) {
       int e = ord[i];
       Pii q = points[e];
       if (q.x-q.y > p.x-p.y) break;
       best = min(best, make_pair(q.x+q.y, e));
        merged.pb(e);
        j++;
```

```
if (best.v != -1) {
        int alt = best.x-p.x-p.y;
        if (alt < close[v].x)</pre>
          close[v] = {alt, best.v};
     merged.pb(v);
    while (j < end) merged.pb(ord[j++]);</pre>
   copy(all(merged), ord.begin()+begin);
 rep(i, 0, 4) {
   rep(j, 0, 2) {
      sort(all(ord), [&](int 1, int r) {
        return points[1] < points[r];</pre>
      close.assign(n, {INT_MAX, -1});
      octant(0, n);
      rep(k, 0, n) {
        Pii p = close[k];
        if (p.y != -1) edges.pb({p.x,{k,p.y}});
        points[k].x \star = -1;
    each (p, points) p = \{p.y, -p.x\};
  11 sum = 0;
 FAU fau(n);
  sort (all (edges));
 G.assign(n, {});
  each(e, edges) if (fau.join(e.y.x, e.y.y)) {
   sum += e.x;
   G[e.y.x].pb({e.y.y, e.x});
   G[e.y.y].pb({e.y.x, e.x});
  return sum;
geometry/segment2.h
```

```
#include "vec2.h"
// 2D segment structure; NOT HEAVILY TESTED
// Base class of versions for ints and doubles
template < class P, class S> struct bseq2 {
 P a, b; // Endpoints
  // Distance from segment to point
 double distTo(P p) const {
   if ((p-a).dot(b-a) < 0) return (p-a).len();</pre>
   if ((p-b).dot(a-b) < 0) return (p-b).len();</pre>
   return double (abs ((p-a).cross(b-a)))
                  / (b-a).len();
// Version for integer coordinates (long long)
struct seg2i : bseg2<vec2i, seg2i> {
 seq2i() {}
  seg2i(vec2i c, vec2i d) : bseg2{c, d} {}
  // Check if segment contains point p
 bool contains (vec2i p) {
```

```
return (a-p).dot(b-p) <= 0 &&
           (a-p).cross(b-p) == 0;
 // Compare distance to p with sgrt(d2)
  // -1 if smaller, 0 if equal, 1 if greater
 int cmpDistTo(vec2i p, 11 d2) const {
   if ((p-a).dot(b-a) < 0) {
     11 \ 1 = (p-a).len2();
     return (1 > d2) - (1 < d2);
   if ((p-b).dot(a-b) < 0) {
     11 1 = (p-b).len2();
     return (1 > d2) - (1 < d2);
   11 c = abs((p-a).cross(b-a));
   d2 *= (b-a).len2();
   return (c*c > d2) - (c*c < d2);
// Version for double coordinates
// Requires cmp() from template
struct seq2d : bseq2<vec2d, seq2d> {
 seq2d() {}
 seg2d(vec2d c, vec2d d) : bseg2{c, d} {}
 bool contains (vec2d p) {
   return cmp((a-p).dot(b-p), 0) <= 0 &&
          cmp((a-p).cross(b-p), 0) == 0;
};
using seg2 = seg2d;
                                            10
geometry/vec2.h
```

```
// 2D point/vector structure; PARTIALLY TESTED
```

```
using vec2 = vec2d;
// Base class of versions for ints and doubles
template<class T, class S> struct bvec2 {
 T x, y;
 S operator+(S r) const {return{x+r.x,y+r.y};}
 S operator-(S r) const {return{x-r.x,y-r.y};}
 S operator*(T r) const { return {x*r, y*r}; }
 S operator/(T r) const { return {x/r, y/r}; }
 T dot(S r) const { return x*r.x+y*r.y; }
 T cross(S r) const { return x*r.y-y*r.x; }
 T len2()
              const { return x*x + y*y; }
 double len() const { return sqrt(len2()); }
 S perp()
              const { return {-y,x}; } //90deg
 pair<T, T> yxPair() const { return {y,x}; }
 double angle() const { //[0;2*PI] CCW from OX
   double a = atan2(v, x);
   return (a < 0 ? a+2*M_PI : a);
// Version for integer coordinates (long long)
struct vec2i : bvec2<11, vec2i> {
 vec2i() : bvec2{0, 0} {}
```

vec2i(11 a, 11 b) : bvec2{a, b} {}

```
bool operator==(vec2i r) const {
    return x == r.x && v == r.v;
  // Sort by angle, length if angles equal
 bool operator<(vec2i r) const {
    if (upper() != r.upper()) return upper();
    auto t = cross(r);
    return t > 0 || (!t && len2() < r.len2());
 bool upper() const {
    return y > 0 || (y == 0 && x >= 0);
};
// Version for double coordinates
// Requires cmp() from template
struct vec2d : bvec2<double, vec2d> {
 vec2d() : bvec2{0, 0} {}
 vec2d(double a, double b) : bvec2{a, b} {}
  vec2d unit() const { return *this/len(); }
 vec2d rotate(double a) const { // CCW
    return {x*cos(a) - y*sin(a),
            x*sin(a) + y*cos(a);
 bool operator==(vec2d r) const {
    return !cmp(x, r.x) && !cmp(y, r.y);
  // Sort by angle, length if angles equal
 bool operator<(vec2d r) const {</pre>
    int t = cmp(angle(), r.angle());
    return t < 0 || (!t && len2() < r.len2());</pre>
```

graphs/2sat.h

```
// 2-SAT solver; time: O(n+m), space: O(n+m)
// Variables are indexed from 1 and
// negative indices represent negations!
// Usage: SAT2 sat(variable count);
// (add constraints...)
// bool solution found = sat.solve();
// sat[i] = value of i-th variable, 0 or 1
            (also indexed from 1!)
// (internally: positive = i*2-1, neg. = i*2-2)
struct SAT2 : Vi {
  vector<Vi> G:
  Vi order, flags;
  // Init n variables, you can add more later
  SAT2(int n = 0) : G(n*2) {}
  // Add new var and return its index
  int addVar() {
    G.resize(sz(G)+2); return sz(G)/2;
  // Add (i => j) constraint
  void imply(int i, int j) {
```

 $i = \max(i * 2 - 1, -i * 2 - 2);$

16

```
j = \max(j*2-1, -j*2-2);
  G[i].pb(j); G[j<sup>1</sup>].pb(i<sup>1</sup>);
// Add (i v i) constraint
void either(int i, int j) { imply(-i, j); }
// Constraint at most one true variable
void atMostOne(Vi& vars) {
  int x = addVar();
  each(i, vars) {
    int y = addVar();
    imply(x, y); imply(i, -x); imply(i, y);
    x = y;
// Solve and save assignments in 'values'
bool solve() { // O(n+m), Kosaraju is used
  assign(sz(G)/2+1, -1);
  flags.assign(sz(G), 0);
  rep(i, 0, sz(G)) dfs(i);
  while (!order.emptv()) {
    if (!propag(order.back()^1, 1)) return 0;
    order.pop_back();
  return 1;
void dfs(int i) {
  if (flags[i]) return;
  flags[i] = 1;
  each(e, G[i]) dfs(e);
  order.pb(i);
bool propag(int i, bool first) {
  if (!flags[i]) return 1;
  flags[i] = 0;
  if (at(i/2+1) >= 0) return first;
  at (i/2+1) = i\&1;
  each(e, G[i]) if (!propag(e, 0)) return 0;
  return 1;
11 a, b, c; // a - b >= c
             vector<11>& vars) {
rep(i, 0, sz(vars)) each(e, edges)
  vars[e.b] = min(vars[e.b], vars[e.a]-e.c);
each (e, edges)
  if (vars[e.a]-e.c < vars[e.b]) return 0;</pre>
return 1:
```

graphs/bellman_inequalities.h 12

```
struct Ineq {
// Solve system of inequalities of form a-b>=c
// using Bellman-Ford; time: O(n*m)
bool solveIneq(vector<Ineq>& edges,
```

graphs/bridges online.h

```
// Dynamic 2-edge connectivity gueries
// Usage: Bridges bridges(vertex count);
// - bridges.addEdge(u, v); - add edge (u, v)
```

```
// - bridges.cc[v] = connected component ID
// - bridges.bi(v) = 2-edge connected comp ID
struct Bridges {
 vector<Vi> G; // Spanning forest
 Vi cc, size, par, bp, seen;
 int cnt{0};
  // Initialize structure for n vertices; O(n)
  Bridges (int n = 0) : G(n), cc(n), size(n, 1),
                       par(n, -1), bp(n, -1),
                       seen(n) {
    iota(all(cc), 0);
  // Add edge (u, v); time: amortized O(lg n)
 void addEdge(int u, int v) {
   if (cc[u] == cc[v]) {
      int r = lca(u, v);
      while ((v = root(v)) != r)
        v = bp[bi(v)] = par[v];
      while ((u = root(u)) != r)
        u = bp[bi(u)] = par[u];
      G[u].pb(v); G[v].pb(u);
     if (size[cc[u]] > size[cc[v]]) swap(u,v);
     size[cc[v]] += size[cc[u]];
     dfs(u, v);
  // Get 2-edge connected component ID
 int bi(int v) { // amortized time: < O(lq n)</pre>
   return bp[v] == -1 ? v : bp[v] = bi(bp[v]);
  int root(int v) {
    return par[v] == -1 || bi(par[v]) != bi(v)
     ? v : par[v] = root(par[v]);
  void dfs(int v, int p) {
   par[v] = p; cc[v] = cc[p];
   each(e, G[v]) if (e != p) dfs(e, v);
  int lca(int u, int v) { // Don't use this!
    for (cnt++;; swap(u, v)) if (u != -1) {
     if (seen[u = root(u)] == cnt) return u;
      seen[u] = cnt; u = par[u];
```

graphs/dense dfs.h

13

```
#include "../math/bit matrix.h"
// DFS over adjacency matrix; time: O(n^2/64)
// G = graph, V = not visited vertices masks
// UNTESTED
struct DenseDFS {
 BitMatrix G. V: // space: O(n^2/64)
 DenseDFS (int n = 0) : G(n, n), V(1, n) {
   reset();
```

14

```
void reset() { each(x, V.M) x = -1; }
 void setVisited(int i) { V.set(0, i, 0); }
 bool isVisited(int i) { return !V(0, i); }
 // DFS step: func is called on each unvisited
 // neighbour of i. You need to manually call
 // setVisited(child) to mark it visited.
 template < class T > // Single step: O(n/64)
 void step(int i, T func) {
   ull* E = G.row(i):
   for (int w = 0; w < G.stride;) {</pre>
     ull x = E[w] & V.row(0)[w];
     if (x) func((w<<6) | __builtin_ctzll(x));</pre>
     else w++;
};
```

graphs/edmonds karp.h constexpr int INF = 1e9+10;

```
// Edmonds-Karp algorithm for finding
// maximum flow in graph; time: O(V*E^2)
// NOT HEAVILY TESTED
struct MaxFlow {
 using T = int;
  struct Edge {
   int dst, inv;
   T flow, cap;
  vector<vector<Edge>> G;
  vector<T> add:
 Vi prev;
  // Initialize for n vertices
  MaxFlow(int n = 0) : G(n) {}
  // Add new vertex
  int addVert() {
    G.emplace_back(); return sz(G)-1;
  // Add edge between u and v with capacity cap
  // and reverse capacity rcap
  void addEdge(int u, int v, T cap, T rcap=0) {
    G[u].pb({ v, sz(G[v]), 0, cap });
    G[v].pb({u, sz(G[u])-1, 0, rcap});
  // Compute maximum flow from src to dst.
  // Flow values can be found in edges.
  // vertices with 'add' >= 0 belong to
  // cut component containing 's'.
  T maxFlow(int src, int dst) {
   T f = 0;
      queue<int> Q;
      O.push(src);
      prev.assign(sz(G), -1);
      add.assign(sz(G), -1);
      add[src] = INF;
```

while (!O.emptv()) {

```
int i = 0.front();
       T m = add[i];
       Q.pop();
       if (i == dst) {
          while (i != src) {
            auto& e = G[i][prev[i]];
            e.flow -= m;
            G[e.dst][e.inv].flow += m;
            i = e.dst;
         f += m;
         break:
        each(e, G[i]) if (add[e.dst] < 0) {</pre>
         if (e.flow < e.cap) {</pre>
            O.push(e.dst);
            prev[e.dst] = e.inv;
            add[e.dst] = min(m, e.cap-e.flow);
   } while (prev[dst] != -1);
   return f;
};
```

graphs/push relabel.h

15

```
constexpr int64_t INF = 1e18;
// Push-relabel algorithm with global relabel
// heuristic for finding maximum flow; O(V^3),
// but very fast in practice.
// Preflow is not converted to flow!
struct MaxFlow {
 using T = int64 t;
 struct Vert {
   int head{0}, cur{0}, label;
   T excess;
 struct Edge {
   int dst, nxt;
   T avail, cap;
 vector<Vert> V:
 vector<Edge> E;
 queue<int> que, bfs;
 // Initialize for n vertices
 MaxFlow(int n = 0) {
   V.assign(n, {});
   E.resize(2);
```

// Add edge between u and v with capacity cap

// Add new vertex

V.emplace back();

return sz(V)-1;

int addVert() {

for (int i = sz(col); i--;) if (A[i][m]) {

rep(k,0,i) if(A[k][col[i]]) A[k][m].flip();

ans[col[i]] = 1;

```
// and reverse capacity rcap
void addEdge(int u, int v, T cap, T rcap=0) {
 E.pb({ v, V[u].head, 0, cap });
 E.pb({ u, V[v].head, 0, rcap });
 V[u].head = sz(E)-2;
 V[v].head = sz(E)-1;
void push(int v, int e) {
 T f = min(V[v].excess, E[e].avail);
 E[e].avail -= f;
 E[e^1].avail += f;
 V[v].excess -= f;
 if ((V[E[e].dst].excess += f) == f)
   que.push(E[e].dst);
// Compute maximum flow from src to dst
T maxFlow(int src, int dst) {
 each(v, V) v.excess = v.label = v.cur = 0;
 each(e, E) e.avail = max(e.cap, T(0));
 int cnt, n = cnt = V[src].label = sz(V);
 V[src].excess = INF;
 for (int e = V[src].head; e; e = E[e].nxt)
   push(src, e);
 for (; !que.empty(); que.pop()) {
   if (cnt >= n/2) {
     each(v, V) v.label = n;
     V[dst].label = 0;
     bfs.push(dst);
     cnt = 0;
     for (; !bfs.empty(); bfs.pop()) {
       auto& v = V[bfs.front()];
       for (int e=v.head; e; e = E[e].nxt) {
         int x = E[e].dst;
         if (E[e^1].avail &&
             V[x].label > v.label+1) {
           V[x].label = v.label+1;
           bfs.push(x);
   int v = que.front(), &l = V[v].label;
   if (v == dst) continue;
   while (V[v].excess && 1 < n) {
     if (!V[v].cur) {
       1 = n;
       for (int e=V[v].head; e; e=E[e].nxt){
         if (E[e].avail)
           l = min(l, V[E[e].dst].label+1);
       V[v].cur = V[v].head;
       cnt++;
     int e = V[v].cur;
     V[v].cur = E[e].nxt;
     if (E[e].avail &&
       l == V[E[e].dst].label+1) push(v, e);
```

```
return V[dst].excess;
  // Get if v belongs to cut component with src
  bool cutSide(int v) {
   return V[v].label >= sz(V);
};
graphs/scc.h
                                             17
// Tarjan's SCC algorithm; time: O(n+m)
// Usage: SCC scc(graph);
// scc[v] = index of SCC for vertex v
// scc.comps[i] = vertices of i-th SCC
struct SCC : Vi {
 vector<Vi> comps:
 Vi S:
 int cnt{0};
 SCC() {}
  SCC(vector < Vi > \& G) : Vi(sz(G), -1), S(sz(G)) 
   rep(i, 0, sz(G)) if (!S[i]) dfs(G, i);
  int dfs(vector<Vi>& G, int v) {
   int low = S[v] = ++cnt, t = -1;
    S.pb(v);
    each(e, G[v]) if (at(e) < 0)
     low = min(low, S[e] ?: dfs(G, e));
    if (low == S[v]) {
     comps.emplace_back();
      for (; t != v; S.pop back()) {
        at (t = S.back()) = sz(comps) - 1;
        comps.back().pb(t);
   return low;
graphs/turbo matching.h
                                            18
// Find maximum bipartite matching; time: ?
// G must be bipartite graph!
// Returns matching size (edge count).
// match[v] = vert matched to v or -1
int matching(vector<Vi>& G, Vi& match) {
 vector<bool> seen;
 int n = 0, k = 1;
  match.assign(sz(G), -1);
  function<int(int)> dfs = [&](int i) {
   if (seen[i]) return 0;
    seen[i] = 1;
   each(e, G[i]) {
     if (match[e] < 0 || dfs(match[e])) {</pre>
        match[i] = e; match[e] = i;
        return 1:
    return 0:
```

```
while (k) {
   seen.assign(sz(G), 0);
   rep(i, 0, sz(G)) if (match[i] < 0)</pre>
     k += dfs(i);
   n += k;
 return n;
// Convert maximum matching to vertex cover
// time: O(n+m)
Vi vertexCover(vector<Vi>& G, Vi& match) {
 Vi ret, col(sz(G)), seen(sz(G));
 function<void(int. int)> dfs =
     [&](int i, int c) {
   if (col[i]) return;
   col[i] = c+1;
   each(e, G[i]) dfs(e, !c);
  function < void (int) > aug = [&] (int i) {
   if (seen[i] || col[i] != 1) return;
   seen[i] = 1;
   each(e, G[i]) seen[e] = 1, aug(match[e]);
 rep(i, 0, sz(G)) dfs(i, 0);
 rep(i, 0, sz(G)) if (match[i] < 0) aug(i);</pre>
 rep(i, 0, sz(G))
   if (seen[i] == col[i]-1) ret.pb(i);
 return ret;
math/bit gauss.h
constexpr int MAX COLS = 2048;
// Solve system of linear equations over Z_2
// time: O(n^2*m/W), where W is word size
// - A - extended matrix, rows are equations,
        columns are variables.
//
        m-th column is equation result
        (A[i][i] - i-th row and i-th column)
//
// - ans - output for variables values
// - m - variable count
// Returns 0 if no solutions found, 1 if one,
// 2 if more than 1 solution exist.
int bitGauss(vector<bitset<MAX COLS>>& A,
            vector<bool>& ans, int m) {
 Vi col:
 ans.assign(m, 0);
 rep(i, 0, sz(A)) {
   int c = int(A[i]. Find first());
   if (c >= m) {
     if (c == m) return 0;
     continue;
   rep(k, i+1, sz(A)) if (A[k][c]) A[k]^=A[i];
   swap(A[i], A[sz(col)]);
   col.pb(c);
```

```
return sz(col) < m ? 2 : 1;
                                             20
math/bit matrix.h
using ull = uint64 t;
// Matrix over Z 2 (bits and xor)
// UNTESTED and UNFINISHED
struct BitMatrix {
 vector<ull> M;
 int rows, cols, stride;
  BitMatrix(int n = 0, int m = 0) {
    rows = n; cols = m;
    stride = (m+63)/64;
    M.resize(n*stride);
  ull* row(int i) { return &M[i*stride]; }
  bool operator()(int i, int i) {
    return (row(i)[j/64] >> (j%64)) & 1;
  void set(int i, int i, bool val) {
    ull &w = row(i)[\frac{1}{64}], m = 1 << (\frac{1}{864});
    if (val) w |= m;
    else w &= ~m:
};
math/crt.h
                                             21
using Pll = pair<11, 11>;
11 egcd(11 a, 11 b, 11& x, 11& v) {
 if (!a) return x=0, y=1, b;
 11 d = eqcd(b%a, a, y, x);
 x -= b/a*v;
 return d;
// Chinese Remainder Theoerem; time: O(lq lcm)
// Solves x = a.x \pmod{a.v}, x = b.x \pmod{b.v}
// Returns pair (x mod lcm, lcm(a.v, b.v))
// or (-1, -1) if there's no solution.
// WARNING: a.x and b.x are assumed to be
// in [0;a.v) and [0;b.v) respectively.
// Works properly if lcm(a.y, b.y) < 2^63.
Pll crt(Pll a, Pll b) {
 if (a.v < b.v) swap(a, b);
 11 x, y, q = eqcd(a.y, b.y, x, y);
 11 c = b.x-a.x, d = b.y/g, p = a.y*d;
 if (c % q) return {-1, -1};
 11 s = (a.x + c/q*x % d * a.y) % p;
 return {s < 0 ? s+p : s, p};
math/discrete logarithm.h
#include "../modular.h"
```

```
// Baby-step giant-step algorithm; O(sgrt(p))
// Finds smallest x such that a^x = b \pmod{p}
// or returns -1 if there's no solution.
11 dlog(11 a, 11 b, 11 p) {
  int m = int(min(llround(sqrt(p))+1, p-1));
  unordered map<11, int> small;
  11 t = 1:
  rep(i, 0, m) {
   int& k = small[t];
   if (!k) k = i+1;
   t = t*a % p;
  t = modInv(t, p);
  rep(i, 0, m) {
   int i = small[b];
   if (j) return i*m + j - 1;
   b = b*t % p;
 return -1;
```

```
math/fft complex.h
using cmpl = complex<double>;
using Vfft = vector<cmpl>;
// Compute DFT over complex numbers; O(n lq n)
// DFT is in bit-reversed order!
// Input size must be power of 2!
template<int dir> // 1 for DFT, -1 for inverse
void fft(Vfft& buf) {
  assert(__builtin_popcount(sz(buf)) == 1);
  int n = sz(buf), bits = 31-__builtin_clz(n);
  int i = (dir > 0 ? 0 : bits-1);
  Vfft bases = {1};
  auto c = exp(cmpl(0, 2*M_PI/n));
  rep(k, 0, n) bases.pb(bases.back()*c);
  for (; i >= 0 && i < bits; i += dir) {
   int shift = 1 << (bits-i-1);</pre>
    rep(j, 0, 1 << i) rep(k, 0, shift) {</pre>
     int a = (j << (bits-i)) | k, b = a|shift;</pre>
      auto v1 = buf[a], v2 = buf[b];
     auto base = bases[(dir*(k<<i)) & (n-1)];</pre>
     if (dir > 0) {
       buf[b] = (v1 - v2) * base;
      } else {
       v2 *= base;
       buf[b] = v1 - v2;
      buf[a] = v1 + v2;
  if (dir < 0) each (x, buf) x /= n;
// Convolve a and b, store result in a;
// time: 0(n lq n)
void convolve(Vfft& a, Vfft b) {
 int len = sz(a) + sz(b) - 1;
```

```
int n = 1 \ll (32 - builtin clz(len));
  a.resize(n); b.resize(n);
  fft<1>(a); fft<1>(b);
  rep(i, 0, n) a[i] *= b[i];
  fft<-1>(a);
  a.resize(len);
math/fft mod.h
constexpr 11 MOD = 15*(1<<27)+1; // ^{2}e9
constexpr 11 ROOT = 440564289; // order = 1<<27
11 modInv(11 a, 11 m) {
 if (a == 1) return 1;
  return ((a - modInv(m%a, a))*m + 1) / a;
// Compute DFT over integers mod MOD; O(n lg n)
// DFT is in bit-reversed order!
// Input size must be power of 2!
template<int dir> // 1 for DFT, -1 for inverse
void fft(vector<11>& buf) {
  assert(__builtin_popcount(sz(buf)) == 1);
  int n = sz(buf), bits = 31- builtin clz(n);
  int i = (dir > 0 ? 0 : bits-1);
  vector<11> bases = {1};
  11 c = ROOT:
  rep(k, bits, 27) c = c*c % MOD;
  rep(k, 0, n) bases.pb(bases.back()*c % MOD);
  for (; i >= 0 && i < bits; i += dir) {
    int shift = 1 << (bits-i-1);</pre>
    rep(j, 0, 1 << i) rep(k, 0, shift) {</pre>
      int a = (j << (bits-i)) | k, b = a|shift;</pre>
      auto v1 = buf[a], v2 = buf[b];
      auto base = bases[(dir*(k<<i)) & (n-1)];
      if (dir > 0) {
        buf[b] = (v1-v2)*base % MOD;
      } else {
       v2 *= base;
        buf[b] = (v1-v2) % MOD;
      buf[a] = (v1+v2) % MOD;
  if (dir < 0) {
   11 v = modInv(n, MOD);
    each (x, buf) x = x*y % MOD;
  each(x, buf) if (x < 0) x += MOD;
// Convolve a and b, store result in a;
// time: 0(n lq n)
void convolve(vector<11>& a, vector<11> b) {
  int len = sz(a) + sz(b) - 1;
  int n = 1 \ll (32 - builtin clz(len));
  a.resize(n); b.resize(n);
  fft<1>(a); fft<1>(b);
  rep(i, 0, n) a[i] = a[i] *b[i] % MOD;
  fft<-1>(a);
  a.resize(len);
```

```
math/gauss.h
                                             25
// Solve system of linear equations; O(n^2*m)
// - A - extended matrix, rows are equations,
         columns are variables.
         m-th column is equation result
         (A[i][i] - i-th row and i-th column)
// - ans - output for variables values
// - m - variable count
// Returns 0 if no solutions found, 1 if one,
 // 2 if more than 1 solution exist.
int gauss (vector<vector<double>>% A.
          vector<double>& ans, int m) {
  Vi col:
  ans.assign(m, 0);
  rep(i, 0, sz(A)) {
    int c = 0;
    while (c <= m && !cmp(A[i][c], 0)) c++;</pre>
    // For Zp:
    //while (c <= m && !A[i][c].x) c++;
    if (c >= m) {
      if (c == m) return 0;
      continue:
    rep(k, i+1, sz(A)) {
      auto mult = A[k][c] / A[i][c];
      rep(j, 0, m+1) A[k][j] -= A[i][j] *mult;
    swap(A[i], A[sz(col)]);
    col.pb(c);
  for (int i = sz(col); i--;) {
    ans[col[i]] = A[i][m] / A[i][col[i]];
    rep(k, 0, i)
      A[k][m] = ans[col[i]] * A[k][col[i]];
  return sz(col) < m ? 2 : 1;
math/miller rabin.h
#include "modular64.h"
// Miller-Rabin primality test
// time O(k*lg^2 n), where k = number of bases
// Deterministic for p <= 10^9
// constexpr 11 BASES[] = {
// 336781006125, 9639812373923155
// 1;
// Deterministic for p <= 2^64
constexpr 11 BASES[] = {
 2,325,9375,28178,450775,9780504,1795265022
bool isPrime(11 p) {
  if (p == 2) return true;
  if (p <= 1 || p%2 == 0) return false;
  11 d = p-1, times = 0;
```

while (d%2 == 0) d /= 2, times++;

```
each(a, BASES) if (a%p) {
    // 11 a = rand() % (p-1) + 1;
    11 b = modPow(a%p, d, p);
    if (b == 1 || b == p-1) continue;
    rep(i, 1, times) {
     b = modMul(b, b, p);
      if (b == p-1) break;
    if (b != p-1) return false;
 return true:
math/modinv precompute.h
constexpr 11 MOD = 234567899;
vector<11> modInv(MOD); // You can lower size
// Precompute modular inverses; time: O(MOD)
void initModInv() {
 modInv[1] = 1:
 rep(i, 2, sz(modInv)) modInv[i] =
    (MOD - (MOD/i) * modInv[MOD%i]) % MOD;
math/modular.h
// Big prime number, about 2*10^9
constexpr int MOD = 15*(1<<27)+1;
11 modInv(11 a, 11 m) { // a^(-1) mod m
 if (a == 1) return 1;
 return ((a - modInv(m%a, a))*m + 1) / a;
11 modPow(11 a, 11 e, 11 m) { // a^e mod m
 11 t = 1 % m;
 while (e) {
    if (e % 2) t = t*a % m;
    e /= 2; a = a*a % m;
 return t;
// Wrapper for modular arithmetic
struct Zp {
 11 x; // Contained value, in range [0; MOD-1]
 Zp() {}
 Zp(11 \ a) : x(a\%MOD) { if (x < 0) x += MOD; }
  #define OP(c,d) Zp& operator c##=(Zp r) { \
      x = x d; return *this; } \
    Zp operator c(Zp r) const { \
      Zp t = *this; return t c##= r; }
  OP(+, +r.x - MOD*(x+r.x >= MOD));
 OP(-, -r.x + MOD*(x-r.x < 0));
 OP(*, *r.x % MOD);
 OP(/, *r.inv().x % MOD);
  // For composite modulus use modInv, not pow
  Zp inv() const { return pow(MOD-2); }
  Zp pow(11 e) const{ return modPow(x,e,MOD); }
 void print() { cerr << x; } // For deb()</pre>
```

```
// Extended Euclidean Algorithm
11 egcd(11 a, 11 b, 11& x, 11& y) {
  if (!a) return x=0, v=1, b;
 11 d = eqcd(b%a, a, y, x);
  x -= b/a*v;
 return d;
math/modular64.h
// Modular arithmetic for modulus < 2^62
11 modAdd(11 x, 11 y, 11 m) {
  return x < m ? x : x-m;
11 modSub(11 x, 11 y, 11 m) {
 return x >= 0 ? x : x+m;
11 modMul(11 a, 11 b, 11 m) {
 11 c = 11((long double) a * b / m);
 11 r = (a*b - c*m) % m;
  return r < 0 ? r+m : r;
11 modPow(11 x, 11 e, 11 m) {
  11 t = 1;
  while (e)
   if (e & 1) t = modMul(t, x, m);
   e >>= 1:
   x = modMul(x, x, m);
  return t:
math/montgomery.h
#include "modular.h"
// Montgomery modular multiplication
// MOD < MG MULT, gcd(MG MULT, MOD) must be 1
// Don't use if modulo is constexpr; UNTESTED
constexpr 11 MG SHIFT = 32;
constexpr 11 MG MULT = 1LL << MG SHIFT;</pre>
constexpr 11 MG MASK = MG MULT - 1;
const 11 MG INV = MG MULT-modInv(MOD, MG MULT);
// Convert to Montgomery form
11 MG(11 x) { return (x*MG MULT) % MOD; }
// Montgomery reduction
// redc(mg * mg) = Montgomery-form product
11 redc(11 x) {
 11 q = (x * MG INV) & MG MASK;
  x = (x + q*MOD) >> MG_SHIFT;
  return (x \ge MOD ? x-MOD : x);
                                            31
math/nimber.h
// Nimbers are defined as sizes of Nim heaps.
// Operations on nimbers are defined as:
```

```
// a+b = mex({a'+b : a' < a} u {a+b' : b' < b})
// ab = mex(\{a'b+ab'+a'b' : a' < a, b' < b\})
// Nimbers smaller than M = 2^2 k form a field.
// Addition is equivalent to xor, meanwhile
// multiplication can be evaluated
// in O(lg^2 M) after precomputing.
using ull = uint64 t;
ull nbuf[64][64]; // Nim-products for 2^i * 2^j
// Multiply nimbers; time: O(1g^2 M)
// WARNING: Call initNimMul() before using.
ull nimMul(ull a, ull b) {
  ull ret = 0;
  for (ull s = a; s; s &= (s-1))
    for (ull t = b; t; t &= (t-1))
      ret ^= nbuf[ builtin ctzll(s)]
                 [ builtin ctzll(t)];
  return ret:
// Initialize nim-products lookup table
void initNimMul() {
  rep(i, 0, 64)
    nbuf[i][0] = nbuf[0][i] = 1ull << i;
  rep(b, 1, 64) rep(a, 1, b+1) {
    int i = 1 << (63 - __builtin_clzll(a));</pre>
    int j = 1 << (63 - __builtin_clzll(b));</pre>
    ull t = nbuf[a-i][b-i];
    if (i < i)
     t = nimMul(t, 1ull << i) << j;
    else
      t = nimMul(t, 1ull << (i-1)) ^ (t << i);
    nbuf[a][b] = nbuf[b][a] = t;
// Compute a^e under nim arithmetic; O(lg^3 M)
// WARNING: Call initNimMul() before using.
ull nimPow(ull a, ull e) {
 ull t = 1;
  while (e) {
   if (e % 2) t = nimMul(t, a);
   e \neq 2; a = nimMul(a, a);
  return t:
// Compute inverse of a in 2^64 nim-field;
// time: O(1g^3 M)
// WARNING: Call initNimMul() before using.
ull nimInv(ull a) {
  return nimPow(a, ull(-2));
// If you need to multiply many nimbers by
// the same value you can use this to speedup.
// WARNING: Call initNimMul() before using.
struct NimMult {
  ull M[64] = \{0\};
  // Initialize lookup; time: O(lg^2 M)
  NimMult(ull a) {
    for (ull t=a; t; t &= (t-1)) rep(i, 0, 64)
      M[i] ^= nbuf[ builtin ctzll(t)][i];
```

```
ull operator()(ull b) {
   ull ret = 0;
   for (ull t = b; t; t \&= (t-1))
     ret ^= M[__builtin_ctzll(t)];
   return ret;
math/phi large.h
#include "pollard rho.h"
// Compute Euler's totient of large numbers
// time: O(n^{(1/4)}) \leftarrow factorization
11 phi(11 n) {
 each(p, factorize(n)) n = n / p.x * (p.x-1);
math/phi precompute.h
Vi phi(10e6+1);
// Precompute Euler's totients; time O(n lq n)
void calcPhi() {
 iota(all(phi), 0);
 rep(i, 2, sz(phi)) if (phi[i] == i)
   for (int j = i; j < sz(phi); j += i)</pre>
     phi[j] = phi[j] / i * (i-1);
math/pi large precomp.h
#include "sieve.h"
// Count primes in given interval
// using precomputed table.
// Set MAX P to sgrt (MAX N) and run sieve()!
// Precomputed table will contain N BUCKETS
// elements - check source size limit.
constexpr 11 MAX_N = 1e11+1;
constexpr 11 N BUCKETS = 10000;
constexpr 11 BUCKET_SIZE = (MAX_N/N_BUCKETS)+1;
constexpr 11 precomputed[] = {/* ... */};
11 sieveRange(11 from, 11 to) {
 bitset<BUCKET_SIZE> elems;
 from = max(from, 2LL);
 to = max(from, to);
 each(p, primesList) {
   11 c = max((from+p-1) / p, 2LL);
   for (11 i = c*p; i < to; i += p)
     elems.set(i-from);
 return to-from-elems.count();
// Run once on local computer to precompute
// table. Takes about 10 minutes for n = 1e11.
// Sanity check (for default params):
// 664579, 606028, 587253, 575795, ...
void localPrecompute() {
 for (11 i = 0; i < MAX N; i += BUCKET SIZE) {</pre>
   11 to = min(i+BUCKET SIZE, MAX N);
    cout << sieveRange(i, to) << ',' << flush;
```

// Multiply by b; time: O(lq M)

```
cout << endl;
// Count primes in [from:to] using table.
// O(N BUCKETS + BUCKET SIZE*la la n + sart(n))
11 countPrimes(11 from, 11 to) {
 11 bFrom = from/BUCKET SIZE+1,
     bTo = to/BUCKET SIZE;
 if (bFrom > bTo) return sieveRange(from, to);
 11 ret = accumulate(precomputed+bFrom,
                      precomputed+bTo, 0);
 ret += sieveRange(from, bFrom*BUCKET SIZE);
 ret += sieveRange(bTo*BUCKET SIZE, to);
 return ret:
                                             35
math/pollard rho.h
#include "modular64.h"
#include "miller rabin.h"
using Factor = pair<11, int>;
void rho(vector<11>& out, 11 n) {
 if (n <= 1) return;</pre>
 if (isPrime(n)) out.pb(n);
 else if (n\%2 == 0) rho(out,2), rho(out,n/2);
 else for (11 a = 2;; a++) {
   11 x = 2, y = 2, d = 1;
   while (d == 1) {
     x = modAdd(modMul(x, x, n), a, n);
     y = modAdd(modMul(y, y, n), a, n);
     y = modAdd(modMul(y, y, n), a, n);
     d = \underline{gcd(abs(x-y), n)};
   if (d != n) {
     rho(out, d);
      rho(out, n/d);
      return:
// Pollard's rho factorization algorithm
// Las Vegas version; time: n^(1/4)
// Returns pairs (prime, power), sorted
vector<Factor> factorize(11 n) {
 vector<Factor> ret;
 vector<11> raw:
 rho(raw, n);
 sort(all(raw));
 each(f, raw) {
   if (ret.empty() || ret.back().x != f)
     ret.pb({ f, 1 });
     ret.back().v++;
 return ret:
math/polynomial interp.h
                                             36
// Interpolates set of points (i, vec[i])
// and returns it evaluated at x; time: O(n^2)
// TODO: Improve to linear time
template<typename T>
```

T polvExtend(vector<T>& vec, T x) {

```
T ret = 0;
  rep(i, 0, sz(vec)) {
   T a = vec[i], b = 1;
    rep(j, 0, sz(vec)) if (i != j) {
     a *= x-i; b *= i-i;
    ret += a/b;
  return ret;
math/sieve.h
                                              37
constexpr int MAX P = 1e6;
bitset<MAX_P+1> primes;
Vi primesList;
// Erathostenes sieve; time: O(n lg lg n)
void sieve() {
  primes.set();
  primes.reset(0);
  primes.reset(1);
  for (int i = 2; i*i <= MAX P; i++)</pre>
   if (primes[i])
      for (int j = i*i; j <= MAX_P; j += i)</pre>
       primes.reset(j);
  rep(i, 0, MAX P+1) if (primes[i])
    primesList.pb(i);
math/sieve factors.h
constexpr int MAX_P = 1e6;
Vi factor (MAX P+1);
// Erathostenes sieve with saving smallest
// factor for each number; time: O(n lq lq n)
void sieve() {
  for (int i = 2; i*i <= MAX P; i++)</pre>
    if (!factor[i])
      for (int i = i*i; i <= MAX P; i += i)</pre>
        if (!factor[i])
          factor[j] = i;
  rep(i, 0, MAX P+1) if (!factor[i]) factor[i]=i;
// Factorize n <= MAX P; time: O(lq n)
// Returns pairs (prime, power), sorted
vector<Pii> factorize(11 n) {
  vector<Pii>> ret;
  while (n > 1) {
    int f = factor[n]:
    if (ret.empty() || ret.back().x != f)
     ret.pb({ f, 1 });
    else
     ret.back().y++;
    n /= f;
  return ret;
```

math/sieve segmented.h

bitset < MAX P/2+1> primes: // Only odd numbers

constexpr int MAX P = 1e9;

39

```
// Cache-friendly Erathostenes sieve
// ~1.5s on Intel Core i5 for MAX P = 10^9
// Memory usage: MAX P/16 bytes
void sieve() {
  constexpr int SEG SIZE = 1<<18;</pre>
  int pSqrt = int(sqrt(MAX_P)+0.5);
  vector<Pii>> dels;
  primes.set();
 primes.reset(0);
  for (int i = 3; i <= pSqrt; i += 2) {</pre>
    if (primes[i/2]) {
      int i:
      for (j = i*i; j \le pSqrt; j += i*2)
        primes.reset(j/2);
      dels.pb({ i, j/2 });
  for (int seg = pSqrt/2;
       seg <= sz(primes); seg += SEG_SIZE) {</pre>
    int lim = min(seg+SEG_SIZE, sz(primes));
    each(d, dels) for (;d.y < lim; d.y += d.x)</pre>
      primes.reset(d.v);
bool isPrime(int x) {
  return x == 2 || (x%2 && primes[x/2]);
structures/bitset plus.h
                                             40
// Undocumented std::bitset features:
// - Find first() - returns first bit = 1 or N
// - _Find_next(i) - returns first bit = 1
11
                     after i-th bit
11
                     or N if not found
// Bitwise operations for vector<bool>
// UNTESTED
#define OP(x) vector<bool>& operator x##=(
   vector<bool>& 1, const vector<bool>& r) { \
  assert(sz(1) == sz(r));
  auto a = 1.begin(); auto b = r.begin();
  while (a<1.end()) *a._M_p++ x##= *b._M_p++; \</pre>
  return 1; }
OP (&) OP (|) OP (^)
structures/fenwick tree.h
                                             41
// Fenwick tree (BIT tree); space: O(n)
// Default version: prefix sums
struct Fenwick {
  using T = int;
  static const T ID = 0;
  T f(T a, T b) { return a+b; }
  vector<T> s;
  Fenwick(int n = 0) : s(n, ID) {}
  // A[i] = f(A[i], v); time: O(lg n)
  void modify(int i, T v) {
    for (; i < sz(s); i |= i+1) s[i]=f(s[i],v);</pre>
```

```
T query(int i) {
    T v = ID;
    for (; i > 0; i \&= i-1) v = f(v, s[i-1]);
    return v:
  // Find smallest i such that
  // f(A[0],...,A[i-1]) >= val; time: O(lq n)
  // Prefixes must have non-descreasing values.
  int lowerBound(T val) {
    if (val <= ID) return 0;</pre>
    int i = -1, mask = 1;
    while (mask \leq sz(s)) mask *= 2:
    T \circ ff = TD:
    while (mask /= 2) {
      int k = mask+i:
      if (k < sz(s)) {
        T x = f(off, s[k]);
        if (val > x) i=k, off=x;
    return i+2;
};
structures/fenwick tree 2d.h
// Fenwick tree 2D (BIT tree 2D); space: O(n*m)
// Default version: prefix sums 2D
// Change s to hashmap for O(q lg^2 n) memory
struct Fenwick2D {
  using T = int;
  static constexpr T ID = 0;
 T f(T a, T b) { return a+b; }
  vector<T> s:
  int w, h;
  Fenwick2D(int n = 0, int m = 0)
   : s(n*m, ID), w(n), h(m) {}
  // A[i,j] = f(A[i,j], v); time: O(lg^2 n)
  void modify(int i, int j, T v) {
    for (; i < w; i |= i+1)
      for (int k = j; k < h; k | = k+1)
        s[i*h+k] = f(s[i*h+k], v);
  // Query prefix; time: O(lg^2 n)
 T query(int i, int j) {
    T v = ID;
    for (; i>0; i&=i-1)
      for (int k = j; k > 0; k \&= k-1)
       v = f(v, s[i*h+k-h-1]);
    return v:
structures/find union.h
// Disjoint set data structure; space: O(n)
// Operations work in amortized O(alfa(n))
struct FAU {
```

 $FAU(int n = 0) : G(n, -1) {}$

// Get f(A[0], ..., A[i-1]); time: O(lg n)

```
// Get size of set containing i
  int size(int i) { return -G[find(i)]; }
  // Find representative of set containing i
 int find(int i) {
    return G[i] < 0 ? i : G[i] = find(G[i]);</pre>
  // Union sets containing i and i
 bool join(int i, int j) {
    i = find(i); j = find(j);
    if (i == j) return 0;
   if (G[i] > G[j]) swap(i, j);
    G[i] += G[j]; G[j] = i;
    return 1;
};
structures/hull offline.h
                                             44
constexpr 11 INF = 2e18;
// constexpr double INF = 1e30;
// constexpr double EPS = 1e-9;
// MAX of linear functions; space: O(n)
// Use if you add lines in increasing 'a' order
// Default uncommented version is for int64
struct Hull {
 using T = 11; // Or change to double
 struct Line {
   T a, b, end;
   T intersect (const Line& r) const {
      // Version for double:
      //if (r.a-a < EPS) return b>r.b?INF:-INF;
      //return (b-r.b) / (r.a-a);
      if (a==r.a) return b > r.b ? INF : -INF;
      11 u = b-r.b, d = r.a-a;
      return u/d + ((u^d) >= 0 || !(u%d));
 };
  vector<Line> S;
 Hull() { S.pb({ 0, -INF, INF }); }
  // Insert f(x) = ax+b; time: amortized O(1)
 void push(T a, T b) {
    Line 1{a, b, INF};
    while (true) {
      T e = S.back().end=S.back().intersect(1);
      if (sz(S) < 2 \mid | S[sz(S)-2].end < e)
       break:
      S.pop_back();
    S.pb(1);
  // Query max(f(x) for each f): time: O(lg n)
 T query(T x) {
    auto t = *upper_bound(all(S), x,
      [](int 1, const Line& r) {
        return 1 < r.end;</pre>
    return t.a*x + t.b;
};
```

algolib

```
structures/hull online.h
                                             45
constexpr 11 INF = 2e18;
// MAX of linear functions online; space: O(n)
struct Hull {
  static bool mode0; // Toggles operator< mode
  struct Line {
   mutable 11 a, b, end;
    11 intersect (const Line& r) const {
     if (a==r.a) return b > r.b ? INF : -INF;
     11 u = b-r.b, d = r.a-a;
     return u/d + ((u^d) >= 0 || !(u%d));
   bool operator<(const Line& r) const {
     return modeQ ? end < r.end : a < r.a;</pre>
  };
  multiset<Line> S:
  Hull() { S.insert({ 0, -INF, INF }); }
  // Updates segment end
  bool update(multiset<Line>::iterator it) {
    auto cur = it++; cur->end = INF;
    if (it == S.end()) return false;
   cur->end = cur->intersect(*it);
   return cur->end >= it->end;
  // Insert f(x) = ax+b; time: O(lq n)
  void insert(11 a, 11 b) {
   auto it = S.insert({ a, b, INF });
    while (update(it)) it = --S.erase(++it);
    rep(i, 0, 2)
     while (it != S.begin() && update(--it))
       update(it = --S.erase(++it));
  // Query max(f(x) for each f): time: O(lq n)
  11 query(11 x) {
   modeQ = 1;
    auto 1 = *S.upper_bound({ 0, 0, x });
   mode0 = 0;
   return l.a*x + l.b;
};
bool Hull::modeO = false;
                                             46
```

```
structures/max queue.h
// Queue with max query on contained elements
struct MaxOueue {
  using T = int;
 deque<T> O, M;
  // Add v to the back; time: amortized O(1)
  void push(T v) {
   while (!M.emptv() && M.back() < v)</pre>
     M.pop back();
   M.pb(v); Q.pb(v);
 // Pop from the front; time: O(1)
```

```
void pop() {
   if (M.front() == O.front()) M.pop front();
   Q.pop_front();
 // Get max element value; time: O(1)
 T max() const { return M.front(); }
structures/pairing heap.h
                                             47
// Pairing heap implementation; space O(n)
// Elements are stored in vector for faster
// allocation. It's MINIMUM queue.
// Allows to merge heaps in O(1)
template<class T, class Cmp = less<T>>
struct PHeap {
 struct Node {
   T val;
   int child{-1}, next{-1}, prev{-1};
   Node (T x = T()) : val(x) \{ \}
 };
 using Vnode = vector<Node>;
 Vnode& M:
 int root{-1};
 int unlink(int& i) {
   if (i >= 0) M[i].prev = -1;
   int x = i; i = -1;
   return x;
 void link(int host, int& i, int val) {
   if (i >= 0) M[i].prev = -1;
   i = val;
   if (i >= 0) M[i].prev = host;
 int merge(int 1, int r) {
   if (1 < 0) return r;
   if (r < 0) return 1;
   if (Cmp()(M[1].val, M[r].val)) swap(l, r);
   link(1, M[1].next, unlink(M[r].child));
   link(r, M[r].child, 1);
   return r;
 int mergePairs(int v) {
   if (v < 0 || M[v].next < 0) return v;</pre>
   int v2 = unlink(M[v].next);
   int v3 = unlink(M[v2].next);
   return merge(merge(v, v2), mergePairs(v3));
 // ---
 // Initialize heap with given node storage
  // Just declare 1 Vnode and pass it to heaps
 PHeap(Vnode& mem) : M(mem) {}
  // Add given key to heap, returns index; O(1)
 int push(const T& x) {
   int index = sz(M);
                                                 // propagation. Configure by modifying:
   M.emplace back(x);
                                                 // - T - data type for updates (stored type)
```

```
root = merge(root, index);
   return index;
 // Change key of i to smaller value; O(1)
 void decrease(int i, T val) {
   assert(!Cmp()(M[i].val, val));
   M[i].val = val;
   int prev = M[i].prev;
   if (prev < 0) return;</pre>
   auto& p = M[prev];
   link(prev, (p.child == i ? p.child
        : p.next), unlink(M[i].next));
   root = merge(root, i);
 bool empty() { return root < 0; }</pre>
 const T& top() { return M[root].val; }
 // Merge with other heap. Must use same vec.
 void merge(PHeap& r) { // time: O(1)
   assert(&M == &r.M);
   root = merge(root, r.root); r.root = -1;
 // Remove min element; time: O(lq n)
 } () gog biov
   root = mergePairs(unlink(M[root].child));
};
                                             48
structures/rmg.h
// Range Minimum Query; space: O(n lg n)
struct RMQ {
 using T = int;
 static constexpr T ID = INT_MAX;
 T f(T a, T b) { return min(a, b); }
 vector<vector<T>> s;
 // Initialize RMQ structure; time: O(n lg n)
 RMO(const vector<T>& vec = {}) {
   s = \{vec\};
   for (int h = 1; h \le sz(vec); h *= 2) {
     s.emplace back();
     auto\& prev = s[sz(s)-2];
     rep(i, 0, sz(vec) -h \star 2+1)
       s.back().pb(f(prev[i], prev[i+h]));
 // Query f(s[b], ..., s[e-1]); time: O(1)
 T query(int b, int e) {
   if (b >= e) return ID;
   int k = 31 - builtin clz(e-b);
   return f(s[k][b], s[k][e - (1<<k)]);
};
structures/segment_tree.h
// Optionally dynamic segment tree with lazy
```

```
// - ID - neutral element for extra
// - Node - details in comments
struct SeamentTree {
 using T = int;
 static constexpr T ID = 0;
 // static constexpr T ID = INT_MIN; // max/=
 struct Node {
   T extra{ID}; // Lazy propagated value
   // Aggregates: sum, max, count of max
   T sum{0}, great{INT_MIN}, nGreat{0};
   // Initialize node with default value x
   void init(T x, int size) {
     sum = x*size; great = x; nGreat = size;
   // Merge with node R on the right
   void merge(const Node& R) {
     else if(great==R.great) nGreat+=R.nGreat;
     sum += R.sum;
     great = max(great, R.great);
   // + version
   // Apply modification to node, return
    // value to be applied to node on right
   T apply(T x, int size) {
     extra += x;
     sum += x*size;
     great += x;
     return x;
   // T apply(T x, int size) {
   // if (great <= x) nGreat = size;</pre>
   // extra = max(extra, x);
   // great = max(great, x);
   // // sum doesn't work here
   // return x;
   1/ }
   // T apply(T x, int size) {
    // extra = x;
   // sum = x*size;
   // great = x;
   // nGreat = size;
   // return x;
   1/ }
 };
 vector<Node> V;
 int len:
 // vector<array<int, 3>> links; // [DYNAMIC]
 // T defVal:
                                 // [DYNAMIC]
 SegmentTree(int n=0, T def=ID) {init(n,def);}
 void init(int n, T def) {
   for (len = 1; len < n; len *= 2);</pre>
   // [STATIC] version
```

algolib

```
9
```

```
V.assign(len*2, {});
  rep(i, len, len+n) V[i].init(def, 1);
  for (int i = len-1; i > 0; i--) update(i);
  // [DYNAMIC] version
  // defVal = def;
  // links.assign(2, {-1, -1, len});
  // V.assign(2, {});
  // V[1].init(def, len);
// [STATIC] version
int getChild(int i, int j) { return i*2+j; }
// [DYNAMIC] version
// int getChild(int i, int j) {
// if (links[i][i] < 0) {
    int size = links[i][2] / 2;
    links[i][i] = sz(V);
    links.push_back({ -1, -1, size });
     V.emplace back();
     V.back().init(defVal, size);
// }
// return links[i][j];
1/ }
int L(int i) { return getChild(i, 0); }
int R(int i) { return getChild(i, 1); }
void update(int i) {
  int \bar{a} = L(i), b = R(i);
 V[i] = {};
 V[i].merge(V[a]);
 V[i].merge(V[b]);
void push(int i, int size) {
 T e = V[i].extra;
 if (e != ID) {
   e = V[L(i)].apply(e, size/2);
   V[R(i)].apply(e, size/2);
   V[i].extra = ID;
// Modify [vBegin; end) with x; time: O(lq n)
T modify(int vBegin, int vEnd, T x,
         int i = 1,
         int begin = 0, int end = -1) {
  if (end < 0) end = len;
 if (vEnd <= begin || end <= vBegin)</pre>
    return x;
  if (vBegin <= begin && end <= vEnd) {</pre>
   return V[i].apply(x, end-begin);
  int mid = (begin + end) / 2;
 push (i, end-begin);
  x = modify(vBegin, vEnd, x, L(i), begin, mid);
  x = modify(vBegin, vEnd, x, R(i), mid, end);
 update(i);
  return x:
// Ouerv [vBegin; vEnd); time: O(lg n)
```

```
// Returns base nodes merged together
 Node guery (int vBegin, int vEnd, int i = 1,
             int begin = 0, int end = -1) {
    if (end < 0) end = len;</pre>
    if (vEnd <= begin || end <= vBegin)
      return {};
    if (vBegin <= begin && end <= vEnd)</pre>
     return V[i];
    int mid = (begin + end) / 2;
    push(i, end-begin);
    Node x = query(vBegin, vEnd, L(i), begin, mid);
   x.merge(query(vBegin, vEnd, R(i), mid, end));
   return x;
  // TODO: generalize?
  // Find longest suffix of given interval
  // such that max value is smaller than val.
  // Returns suffix begin index; time: O(lg n)
 T search (int vBegin, int vEnd, int val,
           int i=1, int begin=0, int end=-1) {
    if (end < 0) end = len;</pre>
    if (vEnd <= begin || end <= vBegin)</pre>
      return begin:
    if (vBegin <= begin && end <= vEnd) {</pre>
     if (V[i].great < val) return begin;</pre>
     if (begin+1 == end) return end;
    int mid = (begin+end) / 2;
   push(i, end-begin);
    int ind = search(vBegin, vEnd, val,
                     R(i), mid, end);
   if (ind > mid) return ind;
   return search (vBegin, vEnd, val,
                  L(i), begin, mid);
structures/segment tree beats.h 50
constexpr 11 INF = 1e18;
// Segment tree with min/+ update and
// sum/max query; time: amortized O(n lg^2 n)
// or O(n lg n) if not using + operation
struct SegmentTree {
 using T = 11:
  vector<T> sum, plus, max1, max2, cnt1;
 int len;
  SegmentTree(int n = 0) {
   for (len = 1; len < n; len *= 2);
    sum.resize(len*2);
    plus.resize(len*2);
    max1.resize(len*2);
    max2.assign(len*2, -INF);
    cnt1.assign(len*2, 1);
   for (int i = len-1; i > 0; i--) update(i);
```

void apply(int i, T m, T p, int size) {

plus[i] += p; sum[i] += p*size;

 $\max 1[i] += p; \max 2[i] += p;$

```
if (m < max1[i]) {
    sum[i] -= (max1[i]-m)*cnt1[i];
    \max 1[i] = m;
void update(int i) {
  int a = i * 2, b = i * 2 + 1;
  sum[i] = sum[a] + sum[b];
  0 = 11 = 0:
 \max 1[i] = \max 1[a];
  \max 2[i] = \max 2[a];
  cnt1[i] = cnt1[a];
  if (max1[b] > max1[i]) {
    \max 2[i] = \max 1[i];
    \max[i] = \max[b];
    cnt1[i] = cnt1[b];
  } else if (max1[b] == max1[i]) {
    cnt1[i] += cnt1[b];
  } else if (max1[b] > max2[i]) {
    max2[i] = max1[b];
  max2[i] = max(max2[i], max2[b]);
void push(int i, int s) {
  rep(j, 0, 2)
    apply(i*2+j, max1[i], plus[i], s/2);
  plus[i] = 0;
// Apply min with x on [vBegin; vEnd)
// time: amortized O(lg n) or O(lg^2 n)
void setMin(int vBegin, int vEnd, T x,
            int i = 1,
            int begin = 0, int end = -1) {
  if (end < 0) end = len;
  if (vEnd <= begin || end <= vBegin ||</pre>
      max1[i] < x) return;</pre>
  if (begin >= vBegin && end <= vEnd &&
      \max 2[i] < x)
    return apply(i, x, 0, end-begin);
  int mid = (begin+end) / 2;
  push(i, end-begin);
  setMin(vBegin, vEnd, x, i*2, begin, mid);
  setMin(vBegin, vEnd, x, i*2+1, mid, end);
  update(i);
// Add x on [vBegin; vEnd); time: O(lg n)
void add(int vBegin, int vEnd, T x,
         int i = 1,
         int begin = 0, int end = -1) {
  if (end < 0) end = len;</pre>
  if (vEnd <= begin || end <= vBegin) return;</pre>
  if (begin >= vBegin && end <= vEnd)</pre>
    return apply(i, INF, x, end-begin);
  int mid = (begin+end) / 2;
  push(i, end-begin);
  add(vBegin, vEnd, x, i*2, begin, mid);
  add (vBegin, vEnd, x, i*2+1, mid, end);
  update(i);
```

```
// Ouerv sum of [vBegin; vEnd]; time: O(lg n)
 T getSum(int vBegin, int vEnd, int i = 1,
           int begin = 0, int end = -1) {
    if (end < 0) end = len;
    if (vEnd<=begin || end<=vBegin) return 0;</pre>
    if (vBegin <= begin && end <= vEnd)</pre>
      return sum[i];
    int mid = (begin+end) / 2;
    push(i, end-begin);
    return getSum(vBegin, vEnd, i*2, begin, mid) +
           getSum(vBegin, vEnd, i*2+1, mid, end);
  // Query max of [vBegin; vEnd); time: O(lq n)
 T getMax(int vBegin, int vEnd, int i = 1,
           int begin = 0, int end = -1) {
    if (end < 0) end = len;
    if (vEnd <= begin || end <= vBegin)</pre>
      return -INF;
    if (vBegin <= begin && end <= vEnd)</pre>
      return max1[i];
    int mid = (begin+end) / 2;
    push(i, end-begin);
    return max (
      getMax(vBegin, vEnd, i*2, begin, mid),
      getMax(vBegin, vEnd, i*2+1, mid, end)
};
structures/segment tree point.h 51
// Segment tree (point, interval)
// Configure by modifying:
// - T - stored data type
// - ID - neutral element for query operation
// - merge(a, b) - combine results
struct SegmentTree {
 using T = int;
 static constexpr T ID = INT_MIN;
```

```
static T merge(T a, T b) { return max(a,b); }
vector<T> V;
int len;
SegmentTree(int n = 0, T def = ID) {
  for (len = 1; len < n; len *= 2);</pre>
  V.resize(len*2, ID);
  rep(i, 0, n) V[len+i] = def;
  for (int i = len-1; i > 0; i--)
    V[i] = merge(V[i*2], V[i*2+1]);
void set(int i, T val) {
  V[i+=len] = val;
  while ((i/=2) > 0)
    V[i] = merge(V[i*2], V[i*2+1]);
T query(int begin, int end) {
  begin += len; end += len-1;
  if (begin > end) return ID;
  if (begin == end) return V[begin];
  T \times = merge(V[begin], V[end]);
```

update(x);

```
while (begin/2 < end/2) {
     if (~begin&1) x = merge(x, V[begin^1]);
     if (end\&1) x = merge(x, V[end^1]);
                                                   // Join two treaps in given order; O(lg n)
     begin /= 2; end /= 2;
                                                   int join(int 1, int r) {
                                                     push(1); push(r);
                                                     if (1 < 0 || r < 0) return max(1, r);</pre>
   return x;
                                                     if (G[1].weight < G[r].weight) {</pre>
};
                                                       G[1].E[1] = join(G[1].E[1], r);
constexpr SegmentTree::T SegmentTree::ID;
                                                       update(1):
                                                       return 1:
structures/treap.h
// "Set" of implicit keyed treaps; space: O(n)
                                                     G[r].E[0] = join(l, G[r].E[0]);
// Treaps are distinguished by roots indices
                                                     update(r);
// Put any additional data in Node struct.
                                                     return r:
struct Treap {
 struct Node {
   int E[2] = \{-1, -1\}, \text{ weight}\{\text{rand}()\};
                                                   // Find node with index i in treap x; O(lg n)
                                                   int find(int x, int i) {
   int size{1}, par{-1};
   bool flip{false}; // Is interval reversed?
                                                     while (x >= 0)
                                                       push(x);
                                                       int key = size(G[x].E[0]);
                                                       if (key == i) return x;
 vector<Node> G:
                                                       x = G[x].E[key < i];
 // Initialize structure for n nodes; O(n)
                                                       if (key < i) i -= key+1;
 // Each node is separate treap,
  // use join() to construct sequence.
                                                     return x;
 Treap(int n = 0) : G(n) {}
  int size(int x) { // Returns subtree size
                                                   // Reverse interval [1;r) in treap x; O(lg n)
                                                   int reverse(int x, int 1, int r) {
   return (x \ge 0 ? G[x].size : 0);
                                                     int a, b, c;
                                                     split(x, b, c, r);
  void push(int x) { // Propagates down stuff
                                                     split(b, a, b, 1);
   if (x >= 0 && G[x].flip) {
                                                     if (b >= 0) G[b].flip ^= 1;
     G[x].flip = 0;
                                                     return join(join(a, b), c);
     swap(G[x].E[0], G[x].E[1]);
     each(e, G[x].E) if (e>=0) G[e].flip ^= 1;
   } // + any other lazy operations
                                                   // Find root of treap containing x; O(lg n)
                                                   int root(int x) {
                                                     while (G[x].par >= 0) x = G[x].par;
  void update(int x) { // Updates aggregates
                                                     return x;
   if (x >= 0) {
     int & s = G[x].size = 1;
     G[x].par = -1;
                                                 structures/ext/hash table.h
     each(e, G[x].E) if (e >= 0) {
       s += G[e].size;
                                                 #include <ext/pb_ds/assoc_container.hpp>
       G[e].par = x;
                                                 using namespace __gnu_pbds;
                                                 // gp hash table<K, V> = faster unordered set
    } // + any other aggregates
                                                 // Anti-anti-hash
                                                 const size t HXOR = mt19937 64(time(0))();
  // Split treap x by index i into 1 and r
                                                 template < class T > struct SafeHash {
  // average time: O(lg n)
                                                   size t operator()(const T& x) const {
  void split(int x, int& l, int& r, int i) {
                                                     return hash<T>()(x ^ T(HXOR));
   push(x); 1 = r = -1;
   if (x < 0) return;</pre>
                                                 };
   int key = size(G[x].E[0]);
   if (i <= key) {
                                                 structures/ext/rope.h
     split(G[x].E[0], 1, G[x].E[0], i);
     r = x:
                                                 #include <ext/rope>
                                                 using namespace __gnu_cxx;
     split(G[x].E[1], G[x].E[1], r, i-key-1);
                                                 // rope<T> = implicit cartesian tree
     1 = x;
                                                 structures/ext/tree.h
```

```
#include <ext/pb ds/assoc container.hpp>
    #include <ext/pb_ds/tree_policy.hpp>
    using namespace __gnu_pbds;
    template<class T, class Cmp = less<T>>
    using ordered set = tree<
     T, null type, Cmp, rb tree tag,
     tree_order_statistics_node_update
    // Standard set functions and:
    // t.order of key(key) - index of first >= key
    // t.find_by_order(i) - find i-th element
    // t1.join(t2) - assuming t1<>t2 merge t2 to t1
    structures/ext/trie.h
    #include <ext/pb ds/assoc container.hpp>
    #include <ext/pb_ds/trie_policy.hpp>
    using namespace __gnu_pbds;
    using pref trie = trie<
      string, null_type,
      trie_string_access_traits<>, pat_trie_tag,
     trie_prefix_search_node_update
    text/aho corasick.h
    constexpr char AMIN = 'a'; // Smallest letter
    constexpr int ALPHA = 26; // Alphabet size
    // Aho-Corasick algorithm for linear-time
    // multiple pattern matching.
    // Add patterns using add(), then call build().
    struct Aho {
      vector<array<int, ALPHA>> nxt{1};
      Vi suf = \{-1\}, accLink = \{-1\};
      vector<Vi> accept{1};
      // Add string with given ID to structure
      // Returns index of accepting node
      int add(const string& str, int id) {
        int i = 0;
        each(c, str) {
          if (!nxt[i][c-AMIN]) {
            nxt[i][c-AMIN] = sz(nxt);
            nxt.pb({}); suf.pb(-1);
            accLink.pb(1); accept.pb({});
          i = nxt[i][c-AMIN];
        accept[i].pb(id);
        return i;
      // Build automata; time: O(V*ALPHA)
      void build() {
        queue<int> que;
        each(e, nxt[0]) if (e) {
          suf[e] = 0; que.push(e);
54
        while (!que.emptv()) {
          int i = que.front(), s = suf[i], j = 0;
          que.pop();
          each(e, nxt[i]) {
            if (e) que.push(e);
55
            (e ? suf[e] : e) = nxt[s][i++];
```

```
accLink[i] = (accept[s].empty() ?
          accLink[s] : s);
  // Append 'c' to state 'i'
  int next(int i, char c) {
    return nxt[i][c-AMIN];
  // Call 'f' for each pattern accepted
  // when in state 'i' with its ID as argument.
  // Return true from 'f' to terminate early.
  // Calls are in descreasing length order.
  template<class F> void accepted(int i, F f) {
    while (i !=-1) {
      each(a, accept[i]) if (f(a)) return;
      i = accLink[i];
text/kmp.h
// Computes prefsuf array; time: O(n)
// ps[i] = max prefsuf of [0;i); ps[0] := -1
template < class T > Vi kmp (const T& str) {
  Vi ps; ps.pb(-1);
  each(x, str) {
    int k = ps.back();
    while (k \ge 0 \&\& str[k] != x) k = ps[k];
    ps.pb(k+1);
  return ps;
// Finds occurences of pat in vec; time: O(n)
// Returns starting indices of matches.
template<class T>
Vi match(const T& str, T pat) {
 int n = sz(pat);
  pat.pb(-1); // SET TO SOME UNUSED CHARACTER
  pat.insert(pat.end(), all(str));
  Vi ret, ps = kmp(pat);
  rep(i, 0, sz(ps)) {
    if (ps[i] == n) ret.pb(i-2*n-1);
  return ret;
text/kmr.h
                                             59
// KMR algorithm for O(1) lexicographical
// comparison of substrings.
struct KMR {
 vector<Vi> ids:
  KMR() {}
  // Initialize structure; time: O(n lg^2 n)
  // You can change str type to Vi freely.
  explicit KMR(const string& str) {
    ids.clear();
    ids.pb(Vi(all(str)));
    for (int h = 1; h \le sz(str); h *= 2) {
      vector<pair<Pii, int>> tmp;
```

rep(j, 0, sz(str)) {

int a = ids.back()[i], b = -1;

if (j+h < **sz**(str)) b = ids.back()[j+h];

```
tmp.pb({ {a, b}, j });
      sort(all(tmp));
      ids.emplace back(sz(tmp));
      int n = 0;
      rep(j, 0, sz(tmp)) {
       if (j > 0 && tmp[j-1].x != tmp[j].x)
        ids.back()[tmp[j].y] = n;
  // Get representative of [begin; end); O(1)
  Pii get (int begin, int end) {
    if (begin >= end) return {0, 0};
    int k = 31 - builtin clz(end-begin);
    return {ids[k][begin], ids[k][end-(1<<k)]};</pre>
  // Compare [b1;e1) with [b2;e2); O(1)
  // Returns -1 if <, 0 if ==, 1 if >
  int cmp(int b1, int e1, int b2, int e2) {
   int 11 = e1-b1, 12 = e2-b2;
    int 1 = \min(11, 12);
   Pii x = get(b1, b1+1), y = get(b2, b2+1);
    if (x == y) return (11 > 12) - (11 < 12);
    return (x > y) - (x < y);
  // Compute suffix array of string; O(n)
  Vi sufArray() {
    Vi sufs(sz(ids.back()));
    rep(i, 0, sz(ids.back()))
     sufs[ids.back()[i]] = i;
    return sufs:
};
text/lcp.h
// Compute Longest Common Prefix array for
// given string and it's suffix array; O(n)
// In order to compute suffix array use kmr.h
// or suffix_array_linear.h
template < class T>
Vi lcpArray(const T& str, const Vi& sufs) {
  int n = sz(str), k = 0;
  Vi pos(n), lcp(n-1);
  rep(i, 0, n) pos[sufs[i]] = i;
  rep(i, 0, n) {
    if (pos[i] < n-1) {
      int j = sufs[pos[i]+1];
      while (i+k < n \&\& j+k < n \&\&
          str[i+k] == str[i+k]) k++;
     lcp[pos[i]] = k;
    if (k > 0) k--;
  return lcp;
```

return a;

text/palindromic tree.h

constexpr int ALPHA = 26; // Set alphabet size

```
// Tree of all palindromes in string,
                                                 // constructed online by appending letters.
text/lvndon factorization.h
                                                // space: O(n*ALPHA); time: O(n)
                                                 // Code marked with [EXT] is extension for
// Compute Lyndon factorization for s: O(n)
                                                 // calculating minimal palindrome partition
// Word is simple iff it's stricly smaller
                                                 // in O(n lg n). Can also be modified for
// than any of it's nontrivial suffixes.
                                                 // similar dynamic programmings.
// Lyndon factorization is division of string
                                                 struct PalTree {
// into simple words. It is unique.
                                                  Vi txt; // Text for which tree is built
vector<string> duval(const string& s) {
 int n = sz(s), i = 0;
                                                   // Node 0 = empty palindrome (root of even)
 vector<string> ret:
                                                   // Node 1 = "-1" palindrome (root of odd)
  while (i < n) {</pre>
                                                   Vi len{0, -1}; // Lengths of palindromes
   int i = i+1, k = i:
                                                   Vi link{1, 0}; // Suffix palindrome links
    while (j < n \&\& s[k] <= s[j])
                                                   // Edges to next palindromes
     k = (s[k] < s[j] ? i : k+1), j++;
                                                   vector<array<int, ALPHA>> to{ {}, {} };
    while (i <= k)
                                                   int last{0}; // Current node (max suffix pal)
      ret.pb(s.substr(i, j-k)), i += j-k;
                                                   Vi diff{0, 0};  // len[i]-len[link[i]] [EXT]
 return ret;
                                                  Vi slink{0, 0}; // Serial links
                                                  Vi series{0, 0}; // Series DP answer
                                                  Vi ans{0};
                                                                   // DP answer for prefix[EXT]
                                            62
text/manacher.h
// Manacher algorithm; time: O(n)
                                                   int ext(int i) {
// Finds largest radiuses for palindromes:
                                                     while (len[i]+2 > sz(txt) ||
// r[2*i] = for center at i (single letter = 1)
                                                            txt[sz(txt)-len[i]-2] != txt.back())
// r[2*i+1] = for center between i and i+1
                                                      i = link[i];
template < class T > Vi manacher (const T& str) {
                                                     return i;
 int n = sz(str) *2, c = 0, e = 1;
 Vi r(n, 1);
  auto get = [&](int i) { return i%2 ? 0 :
                                                   // Append letter from [0; ALPHA); time: O(1)
   (i \ge 0 \&\& i < n ? str[i/2] : i); };
                                                   // (or O(lq n) if [EXT] is enabled)
                                                   void add(int x) {
  rep(i, 0, n) {
                                                     txt.pb(x);
   if (i < e) r[i] = min(r[c*2-i], e-i);
                                                     last = ext(last);
    while (get(i-r[i]) == get(i+r[i])) r[i]++;
   if (i+r[i] > e) c = i, e = i+r[i]-1;
                                                     if (!to[last][x]) {
                                                      len.pb(len[last]+2);
                                                      link.pb(to[ext(link[last])][x]);
 rep(i, 0, n) r[i] /= 2;
                                                      to[last][x] = sz(to);
 return r;
                                                      to.emplace_back();
text/min rotation.h
                                            63
                                                      diff.pb(len.back() - len[link.back()]);
                                                      slink.pb(diff.back() == diff[link.back()]
// Find lexicographically smallest
                                                        ? slink[link.back()] : link.back());
// rotation of s; time: O(n)
                                                      series.pb(0);
// Returns index where shifted word starts.
                                                      // [/EXT]
// You can use std::rotate to get the word:
// rotate(s.begin(), v.begin()+minRotation(v),
                                                     last = to[last][x];
        v.end());
int minRotation(string s) {
 int a = 0, n = sz(s); s += s;
                                                     ans.pb(INT MAX);
  rep(b, 0, n) rep(i, 0, n) {
                                                     for (int i=last; len[i] > 0; i=slink[i]) {
   if (a+i == b || s[a+i] < s[b+i]) {
                                                      series[i] = ans[sz(ans) - len[slink[i]]
     b += max(0, i-1); break;
                                                                      - diff[i] - 1];
                                                      if (diff[i] == diff[link[i]])
   if (s[a+i] > s[b+i]) {
                                                         series[i] = min(series[i],
     a = b; break;
                                                                         series[link[i]]);
```

```
text/suffix array linear.h
                                            65
#include "../util/radix sort.h"
// KS algorithm for suffix array; time: O(n)
// Input values are assumed to be in [1;k]
Vi sufArrav(Vi str, int k) {
 int n = sz(str);
 Vi suf(n):
 str.resize(n+15);
 if (n < 15) {
   iota(all(suf), 0);
   rep(j, 0, n) countSort(suf,
     [&](int i) { return str[i+n-j-1]; }, k);
   return suf:
 // Compute triples codes
 Vi tmp, code(n+2);
 rep(i, 0, n) if (i % 3) tmp.pb(i);
 rep(j, 0, 3) countSort(tmp,
   [&] (int i) { return str[i-j+2]; }, k);
 int mc = 0, j = -1;
 each(i, tmp) {
   code[i] = mc += (j == -1)
       str[i] != str[j] ||
       str[i+1] != str[i+1] ||
       str[i+2] != str[j+2]);
  // Compute suffix array of 2/3
 tmp.clear();
 for (int i=1; i < n; i += 3) tmp.pb(code[i]);</pre>
 tmp.pb(0);
 for (int i=2; i < n; i += 3) tmp.pb(code[i]);</pre>
 tmp = sufArray(move(tmp), mc);
  // Compute partial suffix arrays
 Vi third;
 int th = (n+4) / 3;
 if (n%3 == 1) third.pb(n-1);
 rep(i, 1, sz(tmp)) {
   int e = tmp[i];
   tmp[i-1] = (e 
   code[tmp[i-1]] = i;
   if (e < th) third.pb(e*3);
 tmp.pop back();
 countSort (third,
   [&] (int i) { return str[i]; }, k);
 // Merge suffix arravs
 merge(all(third), all(tmp), suf.begin(),
   [&] (int 1, int r) {
     while (1\%3 == 0 | | r\%3 == 0) {
       if (str[1] != str[r])
         return str[l] < str[r];</pre>
       1++; r++;
      return code[1] < code[r];</pre>
```

[EXT]

[EXT]

ans.back() = min(ans.back(), series[i]+1);

// [/EXT]

};

```
});
  return suf;
// KS algorithm for suffix array; time: O(n)
Vi sufArray(const string& str) {
  return sufArray(Vi(all(str)), 255);
                                             66
text/z function.h
// Computes Z function array; time: O(n)
// zf[i] = max common prefix of str and str[i:]
template < class T > Vi prefPref(const T& str) {
  int n = sz(str), b = 0, e = 1;
  Vi zf(n):
 rep(i, 1, n) {
   if (i < e) zf[i] = min(zf[i-b], e-i);</pre>
    while (i+zf[i] < n \&\&
     str[zf[i]] == str[i+zf[i]]) zf[i]++;
   if (i+zf[i] > e) b = i, e = i+zf[i];
 zf[0] = n;
  return zf;
```

trees/centroid decomp.h

int root; // Root centroid

CentroidTree() {}

```
// Centroid decomposition; space: O(n lg n)
// UNTESTED
struct CentroidTree {
 // child[v] = children of v in centroid tree
 // ind[v][i] = index of vertex v in
                i-th centroid from root
 // subtree[v] = vertices in centroid subtree
  // dists[v] = distances from v to vertices
               in centroid subtree
 // par[v] = parent of v in centroid tree
  // depth[v] = depth of v in centroid tree
  // size[v] = size of centroid subtree of v
  vector<Vi> child, ind, dists, subtree;
 Vi par, depth, size;
```

```
CentroidTree (vector < Vi>& G)
   : child(sz(G)), ind(sz(G)), dists(sz(G)),
      subtree (\mathbf{sz}(G)), par (\mathbf{sz}(G), -2),
      depth(sz(G)), size(sz(G)) {
  root = decomp(G, 0, 0);
int dfs(vector<Vi>& G, int i, int p) {
  size[i] = 1;
  each(e, G[i]) if (e != p && par[e] == -2)
    size[i] += dfs(G, e, i);
  return size[i];
```

```
void layer (vector < Vi>& G, int i,
          int p, int c, int d) {
 ind[i].pb(sz(subtree[c]));
 subtree[c].pb(i);
 dists[c].pb(d);
 each(e, G[i]) if (e != p && par[e] == -2)
   laver(G, e, i, c, d+1);
```

```
int decomp(vector<Vi>& G, int v, int d) {
    int p = -1, s = dfs(G, v, -1);
   bool ok = 1;
    while (ok) {
     ok = 0;
     each(e, G[v]) {
       if (e != p && par[e] == -2 &&
            size[e] > s/2)  {
          p = v; v = e; ok = 1;
          break;
    par[v] = -1;
    size[v] = s;
    depth[v] = d;
    layer(G, v, -1, v, 0);
    each(e, G[v]) if (par[e] == -2) {
     int j = decomp(G, e, d+1);
     child[v].pb(j);
     par[j] = v;
   return v;
};
```

trees/heavylight decomp.h

```
#include "../structures/segment tree point.h"
// Heavy-Light Decomposition of tree
// with subtree query support; space: O(n)
struct HLD {
 // Subtree of v = [pos[v]; pos[v]+size[v])
 // Chain with v = [chBegin[v]; chEnd[v])
 Vi par;
             // Vertex parent
             // Vertex subtree size
 Vi size;
 Vi depth; // Vertex distance to root
             // Vertex position in "HLD" order
 Vi pos;
 Vi chBegin; // Begin of chain with vertex
 Vi chEnd; // End of chain with vertex
 Vi order; // "HLD" preorder of vertices
 SegmentTree tree; // Verts are in HLD order
 HLD() {}
 // Initialize structure for tree G
 // and root r; time: O(n lq n)
 // MODIFIES ORDER OF EDGES IN G!
 HLD(vector<Vi>& G, int r)
     : par(sz(G)), size(sz(G)),
       depth(sz(G)), pos(sz(G)),
       chBegin(sz(G)), chEnd(sz(G)) {
   dfs(G, r, -1);
   decomp(G, r, -1, 0);
   tree = {sz(order)};
 void dfs(vector<Vi>& G, int i, int p) {
   par[i] = p;
   size[i] = 1;
   depth[i] = p < 0 ? 0 : depth[p]+1;
```

```
int& fs = G[i][0];
 if (fs == p) swap(fs, G[i].back());
 each(e, G[i]) if (e != p) {
   dfs(G, e, i);
   size[i] += size[e];
   if (size[e] > size[fs]) swap(e, fs);
void decomp(vector<Vi>& G,
            int i, int p, int chb) {
 pos[i] = sz(order);
 chBegin[i] = chb;
 chEnd[i] = pos[i]+1;
 order.pb(i);
 each(e, G[i]) if (e != p) {
   if (e == G[i][0]) {
      decomp(G, e, i, chb);
      chEnd[i] = chEnd[e];
   } else {
      decomp(G, e, i, sz(order));
// Get root of chain containing v
int chRoot(int v) {return order[chBeqin[v]];}
// Level Ancestor Query; time: O(lg n)
int lag(int i, int level) {
 while (true) {
   int k = pos[i] - depth[i] + level;
   if (k >= chBegin[i]) return order[k];
   i = par[chRoot(i)];
// Lowest Common Ancestor; time: O(lg n)
int lca(int a, int b) {
 while (chBegin[a] != chBegin[b]) {
   int ha = chRoot(a), hb = chRoot(b);
   if (depth[ha] > depth[hb]) a = par[ha];
   else b = par[hb];
 return depth[a] < depth[b] ? a : b;</pre>
// Call func(chBegin, chEnd) on each path
// segment; time: O(lg n * time of func)
template<class T>
void iterPath(int a, int b, T func) {
 while (chBegin[a] != chBegin[b]) {
   int ha = chRoot(a), hb = chRoot(b);
   if (depth[ha] > depth[hb]) {
      func(chBegin[a], pos[a]+1);
      a = par[ha];
    } else {
      func(chBegin[b], pos[b]+1);
      b = par[hb];
 if (pos[a] > pos[b]) swap(a, b);
  // Remove +1 from pos[a]+1 for vertices
```

```
// queries (with +1 -> edges).
   func(pos[a]+1, pos[b]+1);
 // Ouerv path between a and b; O(1g^2 n)
 SegmentTree::T guervPath(int a, int b) {
   auto ret = SegmentTree::ID;
   iterPath(a, b, [&](int i, int i) {
     ret = SegmentTree::merge(ret,
         tree.query(i, j));
   return ret;
 // Query subtree of v; time: O(lg n)
 SegmentTree::T querySubtree(int v) {
   return tree.query(pos[v], pos[v]+size[v]);
};
```

69

```
trees/lca.h
// LAQ and LCA using jump pointers
// space: 0(n lg n)
struct LCA {
 vector<Vi> jumps;
 Vi level, pre, post;
 int cnt{0}, depth;
 LCA() {}
  // Initialize structure for tree G
  // and root r; time: O(n lg n)
 LCA (vector<Vi>& G, int r)
      : jumps(sz(G)), level(sz(G)),
        pre(sz(G)), post(sz(G)) {
    dfs(G, r, -1);
    depth = int(log2(sz(G))) + 2;
    rep(j, 0, depth) rep(i, 0, sz(G))
      jumps[i].pb(jumps[jumps[i][j]][j]);
  void dfs(vector<Vi>& G, int i, int p) {
    level[i] = p < 0 ? 0 : level[p]+1;
    jumps[i].pb(p < 0 ? i : p);
    pre[i] = ++cnt;
    each(e, G[i]) if (e != p) dfs(G, e, i);
    post[i] = ++cnt;
  // Check if a is ancestor of b; time: O(1)
 bool isAncestor(int a, int b) {
    return pre[a] <= pre[b] &&</pre>
           post[b] <= post[a];</pre>
  // Lowest Common Ancestor; time: O(lq n)
 int operator()(int a, int b) {
    for (int j = depth; j >= 0; j--) {
      if (!isAncestor(jumps[a][j], b))
        a = iumps[a][i];
    return isAncestor(a, b) ? a : jumps[a][0];
```

// Level Ancestor Ouerv; time: O(lg n)

```
int lag(int a, int lvl) {
    for (int j = depth; j >= 0; j--) {
      if (lvl < level[jumps[a][j]])</pre>
        a = jumps[a][j];
    return level[a] <= lvl ? a : jumps[a][0];</pre>
  // Get distance from a to b; time: O(lq n)
  int distance(int a. int b) {
    return level[a] + level[b] -
           level[operator()(a, b)]*2;
};
trees/link cut tree.h
                                             70
// Link/cut tree; space: O(n)
// Represents forest of (un)rooted trees.
```

struct LinkCutTree { vector<arrav<int, 2>> child; Vi par, prev, flip; // Initialize structure for n vertices; O(n) // At first there's no edges. LinkCutTree(int n = 0) : $child(n, \{-1, -1\}), par(n, -1),$

prev(n, -1), flip(n, -1) {} void auxLink(int p, int i, int ch) { child[p][i] = ch;**if** (ch >= 0) par[ch] = p;

void push(int x) { if (x >= 0 && flip[x]) { flip[x] = 0;swap(child[x][0], child[x][1]); each(e, child[x]) if (e>=0) flip[e] ^= 1;

void rot(int p, int i) { int x = child[p][i], g = par[x] = par[p]; if (g >= 0) child[g][child[g][1] == p] = x; auxLink(p, i, child[x][!i]); auxLink(x, !i, p); swap(prev[x], prev[p]);

void splay(int x) { while (par[x] >= 0) { int p = par[x], q = par[p]; push(q); push(p); push(x);

bool f = (child[p][1] == x);**if** (g >= 0) { **if** (child[g][f] == p) { // zig-zig rot(q, f); rot(p, f); } **else** { // zig-zag rot(p, f); rot(q, !f);

} **else** { // ziq rot(p, f);

push(x);

using ull = uint64 t; #define T64(s,up) for (ull i=0; i<64; i+=s*2) for (ull i = i; i < i+s; i++) { **ull** a = (M[j] >> s) & up; **ull** b = (M[j+s] & up) << s;M[i] = (M[i] & up) | b;

void access(int x) { while (true) { splav(x): int p = prev[x]; if (p < 0) break; prev[x] = -1;splay(p); int r = child[p][1]; if (r >= 0) swap(par[r], prev[r]); auxLink(p, 1, x);void makeRoot(int x) { access(x): int& 1 = child[x][0]; **if** (1 >= 0) { swap(par[1], prev[1]); flip[1] ^= 1; 1 = -1: // Find representative of tree containing x int find(int x) { // time: amortized O(lq n) access(x); while (child[x][0] >= 0)push(x = child[x][0]);splay(x);return x; // Add edge x-y; time: amortized O(lg n) void link(int x, int y) { makeRoot(x); prev[x] = y;// Remove edge x-y; time: amortized O(lg n) void cut(int x, int y) { makeRoot(x); access(y); par[x] = child[y][0] = -1;};

util/bit hacks.h

```
// __builtin_popcount - count number of 1 bits
// __builtin_clz - count most significant 0s
// __builtin_ctz - count least significant 0s
// __builtin_ffs - like ctz, but indexed from 1
                  returns 0 for 0
// For 11 version add 11 to name
     M[j+s] = (M[j+s] & (up << s)) | a;
```

```
// Transpose 64x64 bit matrix
void transpose64(arrav<ull, 64>& M) {
 T64(1, 0x555555555555555);
 T64(2, 0x333333333333333333);
 T64(4, 0xF0F0F0F0F0F0F0F);
 T64(8, 0xFF00FF00FF00FF);
 T64(16, 0xFFFF0000FFFF);
 T64(32, 0xFFFFFFFFLL);
// Lexicographically next mask with same
// amount of ones.
int nextSubset(int v) {
 int t = v | (v - 1):
 return (t + 1) | (((~t & -~t) - 1) >>
      (\underline{\phantom{a}}builtin_ctz(v) + 1));
```

util/bump_alloc.h

// Allocator, which doesn't free memory. char mem[400<<20]; // Set memory limit</pre> size t nMem; void* operator new(size t n) { nMem += n; return &mem[nMem-n];

util/compress vec.h

void operator delete(void*) {}

```
// Compress integers to range [0:n) while
// preserving their order; time: O(n lg n)
// Returns mapping: compressed -> original
Vi compressVec(vector<int*>& vec) {
 sort (all (vec),
    [](int* 1, int* r) { return *1 < *r; });</pre>
  Vi old:
  each(e, vec) {
   if (old.empty() || old.back() != *e)
      old.pb(*e);
    \star e = sz(old)-1;
  return old;
```

util/inversion vector.h

```
// Get inversion vector for sequence of
// numbers in [0;n); ret[i] = count of numbers
// smaller than perm[i] to the left; O(n lg n)
Vi encodeInversions(Vi perm) {
 Vi odd, ret(sz(perm));
 int cont = 1;
 while (cont) {
   odd.assign(sz(perm)+1, 0);
   cont = 0;
   rep(i, 0, sz(perm)) {
     if (perm[i] % 2) odd[perm[i]]++;
     else ret[i] += odd[perm[i]+1];
     cont += perm[i] /= 2;
 return ret:
```

```
// Count inversions in sequence of numbers
// in [0;n); time: O(n lq n)
11 countInversions(Vi perm) {
 11 ret = 0, cont = 1;
 Vi odd;
  while (cont) {
    odd.assign(sz(perm)+1, 0);
    cont = 0;
    rep(i, 0, sz(perm)) {
      if (perm[i] % 2) odd[perm[i]]++;
      else ret += odd[perm[i]+1];
      cont += perm[i] /= 2;
 return ret:
```

util/longest increasing sub.h 75

```
// Longest Increasing Subsequence; O(n lg n)
int lis(const Vi& seq) {
 Vi dp(sz(seq), INT_MAX);
 each(c, seq) *lower_bound(all(dp), c) = c;
 return int(lower_bound(all(dp), INT_MAX)
            - dp.begin());
```

73 util/max rects.h

return ans:

```
76
struct MaxRect {
 // begin = first column of rectangle
 // end = first column after rectangle
 // hei = height of rectangle
 // touch = columns of height hei inside
 int begin, end, hei;
 Vi touch; // sorted increasing
};
// Given consecutive column heights find
// all inclusion-wise maximal rectangles
// contained in "drawing" of columns; time O(n)
vector<MaxRect> getMaxRects(Vi hei) {
 hei.insert(hei.begin(), -1);
 hei.pb(-1);
 Vi reach (sz(hei), sz(hei)-1);
 vector<MaxRect> ans;
 for (int i = sz(hei)-1; --i;) {
   int j = i+1, k = i;
   while (hei[j] > hei[i]) j = reach[j];
   reach[i] = j;
   while (hei[k] > hei[i-1]) {
     ans.pb({ i-1, 0, hei[k], {} });
     auto& rect = ans.back();
      while (hei[k] == rect.hei) {
       rect.touch.pb(k-1);
       k = reach[k];
     rect.end = k-1;
```

util/mo.h

```
// Modified MO's queries sorting algorithm,
// slightly better results than standard.
// Allows to process q queries in O(n*sqrt(q))
struct Ouerv {
 int begin, end;
// Get point index on Hilbert curve
11 hilbert (int x, int y, int s, 11 c = 0) {
  if (s <= 1) return c;</pre>
  s /= 2; c *= 4;
  if (y < s)
   return hilbert (x&(s-1), y, s, c+(x>=s)+1);
  if (x < s)
   return hilbert (2*s-y-1, s-x-1, s, c);
  return hilbert (y-s, x-s, s, c+3);
// Get good order of queries; time: O(n lg n)
Vi moOrder(vector<Query>& queries, int maxN) {
  int s = 1:
  while (s < maxN) s \star= 2;
  vector<11> ord;
  each (q, queries)
   ord.pb(hilbert(q.begin, q.end, s));
  Vi ret(sz(ord));
  iota(all(ret), 0);
  sort(all(ret), [&](int 1, int r) {
   return ord[l] < ord[r];</pre>
  return ret;
util/parallel binsearch.h
                                             78
// Run 'count' binary searches on [begin; end),
// 'cmp' arguments:
// 1) vector<Pii>& - pairs (value, index)
// which are queries if value of index is
     greater or equal to value,
     sorted by value
// 2) vector<bool>& - true at index i means
     value of i-th query is >= queried value
// Returns vector of found values.
// Time: O((n+c) lg n), where c is cmp time.
template<class T>
Vi multiBS(int begin, int end, int count, T cmp) {
  vector<Pii> ranges(count, {begin, end});
  vector<Pii>> queries(count);
  vector<bool> answers(count);
  rep(i, 0, count) queries[i] = { (begin+end) /2, i};
  for (int k = uplg(end-begin); k > 0; k--) {
    int last = 0, j = 0;
    cmp(queries, answers);
    rep(i, 0, sz(queries)) {
     Pii &q = queries[i], &r = ranges[q.y];
     if (q.x != last) last = q.x, j = i;
```

```
(answers[i] ? r.x : r.y) = q.x;
     q.x = (r.x+r.y) / 2;
      if (!answers[i])
        swap(queries[i], queries[j++]);
 Vi ret;
  each(p, ranges) ret.pb(p.x);
 return ret;
                                             79
util/radix sort.h
Vi buf, cnt;
// Stable countingsort; time: O(k+sz(vec))
// See example usage in radixSort for pairs.
template < class F>
void countSort (Vi& vec, F key, int k) {
 vec.swap(buf);
 vec.resize(sz(buf));
  cnt.assign(k+1, 0);
 each (e, buf) cnt[key(e)]++;
 rep(i, 1, k+1) cnt[i] += cnt[i-1];
  for (int i = sz(vec)-1; i >= 0; i--)
   vec[--cnt[key(buf[i])]] = buf[i];
// Compute order of elems, k is max key; O(n)
Vi radixSort(const vector<Pii>& elems, int k) {
 Vi order (sz(elems));
 iota(all(order), 0);
  countSort (order,
    [&](int i) { return elems[i].y; }, k);
  countSort (order,
   [&](int i) { return elems[i].x; }, k);
 return order:
```