

# CSE331 - Computer Organization

## Homework 2 - Report

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### Procedures and Their Brief Explanation:

#### findLIS

- findLIS is the main procedure which calculates the subsequences of the given array. It also calls **isIncrementing**, **copyArray**, **printArray** procedures. (nested)
- **Arguments** : a1 (array address), a2 (array size), a3 (sequence array address)
- **Return Value** : v0 (size of the longest increasing subsequence array)

#### printLongest

- It just calls **printArray** but this procedure also adds "Longest Increasing Subsequence Array : " string in front of the array before calling printArray.
- **Arguments** : a0 (lisArr address), a1 (size of the lisArr)
- **Return Value** : No return value
- lisArr : Longest Increasing Subsequence Array

#### copyArray

- This procedure gets two array addresses and copies first array elements to the second array. It is used in **findLIS** to copy seqArr elements to lisArr.
- **Arguments** : a0 (first array address), a1 (second array address), a2 (size of the first array)
- **Return Value** : No return value
- seqArr : Subsequence Array

#### isIncrementing

- If numbers inside of the given array address is in increasing order returns 1, else returns 0
- **Arguments** : a0 (array address), a1 (size of the array)
- **Return Value** : v0 (equals 1 if incrementing, else 0)

#### getSize

- Looks current number in the array, if it is between  $0 < x < 100$ , then it increments current size(0 at first) and address of the array until it finds a number outside of this boundry.
- In .data part when you put **arr2** right after **arr** in memory there will be no value outside boundry between **arr** - **arr2** because of that size will be calculated wrong. To prevent it, I put empty values between them.
- **Arguments** : a0 (array address)
- **Return Value** : v0 (size of the array)

### printArray

- It calls a basic while loop. Until t2 (which is 'i' in c code) is not equal to size of the given array, then print value at the current location and increment index
- **Arguments** : a1 (size of the array), a2 (address of the array)
- **Return Value** : No return value

## Procedures' Pseudo Codes and their Time&Space Complexity

### findLIS:

arr : address of the array  
size : size of that array  
index : 0 (index for that array)  
sNum : 0 (number of subsequences so far)  
maxS :  $2^{\text{size}}$  (max possible num of subseq)  
seqArr: address of the subSequence array  
sSize : 0 (size of the seqArr)  
ISize : size of the longest increasing subsequence array

```
while(sNum < maxS)
    for(i = 0; i < size; i++)
        if(sNum AND (1 << i))
            seqArr[sSize] = arr[i]
```

```
flag = isIncrementing(seqArr)
if(flag == 1)
    printArray(seqArr, sSize)
if(ISize < sSize)
    ISize = sSize
    copyArray(seqArr, lisArr, sSize)
sNum++
sSize = 0
```

**Time Complexity** :  $O(2^n * n)$  ! Because outer while loop is iterated  $2^n$  times and inner for loop is iterated n times. **printArray** is  $O(n)$  and **isIncrementing** is  $O(n^2)$  but

$O(2^n * n) + O(n^2) + O(n) = O(2^n * n)$  so they don't change complexity of the procedure.

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space. They always use what is given from the start

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### printLongest:

arr : address of the array  
size : size of the array  
print("Longest Increasing Subsequence Array : ")  
printArray(arr, size)

```
print("\n")
```

**Time Complexity** :  $O(n)$  ! Because print operation is  $O(1)$  and **printArray** procedure is  $O(n)$ . If we sum,  $O(1) + O(n) = O(n)$

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space.

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### copyArray:

```
arr    : address of the array
cArr   : address of the copy array
size   : size of the base array
i      : 0 (index for loop)
```

```
while (i != size)
    cArr[i] = arr[i]
    i++
```

**Time Complexity** :  $O(n)$  ! Because print operation is  $O(1)$  and while loop iterates  $n$  times.

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space.

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### isIncrementing:

```
arr    : address of the array
size   : size of the array
i      : 0 (index for outer loop)
j      : 0 (index for inner loop)
flag   : true
```

```
while (i != size)
    j = i + 1
    while (j != size)
        if(arr[j] < arr[i])
            flag = false
        j++
    i++
```

```
return flag
```

**Time Complexity** :  $O(n^2)$  ! Because outer loop iterates  $n$  times and inner loop iterates  $n/2$  times which is still  $O(n)$ . So it becomes  $O(n*n) = O(n^2)$

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space.

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## getSize:

MAX\_VAL : 100 (Max value that can be assigned to an array)  
arr : address of the array  
cSize : 0 (current size of the array)  
cVal : 1 (current value)

```
while (cVal != 0) and (cVal < 100):  
    cSize++  
    cVal = arr[cSize]
```

return cSize

**Time Complexity** :  $O(n)$  ! Because while loop iterates n times which equals size of the array

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space.

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## printArray:

arr : address of the array  
size : 0 (size of the array)  
cSize : 0 (current size of the array)  
cVal : 0 (current value)

```
while cSize != size:  
    cVal = arr[cSize]  
    print(" " + cVal)  
    cSize++
```

print("Size : " + size + "\n")

**Time Complexity** :  $O(n)$  ! Because while loop iterates n times which equals size of the array and print and assign operation takes  $O(1)$  time

**Space Complexity** :  $O(1)$  ! Because procedure doesn't allocate any extra space.

## Result of Test Cases and Explanation

```
main:  
  
## ----- 1. Array ----- ##  
# 1) Find size of the array  
la    $a0,    arr           # a0 = address of the array  
jal    getSize             # v0 = size of the array  
# 2) Find subsequences of arr and print them  
la    $a1,    arr           # a1 = address of the array  
move   $a2,    $v0          # a2 = v0 (came from getSize)  
la    $a3,    seqArr        # a3 = address of the seqArr  
jal    findLIS             # Finds Longest Increasing Subsequence (prints rest of them)  
# 3) Print Longest Increasing Subsequence Array  
la    $a0,    lisArr        # a0 = lisArr address  
move   $a1,    $v0          # v0 = size of the lisArr (returned from findLIS)  
jal    printLongest        # Prints lisArr and it's size to the console
```

In main, first I got size of the input array using **getSize** procedure. After that, by using **v0** (which return value from **getSize**) I called **findLIS** procedure to find longest increasing subsequence array(**lisArr**) and it's size. Finally, I printed the result. I repeated this 6 times for different arr inputs. These are the arrays I tried:

#### # Input Arrays

```
arr:           .word    3, 10, 7, 9, 4, 11
empty:         .word    0
arr2:          .word    2, 5, 3, 95, 24, 1, 22, 10, 88
empty2:        .word    0
arr3:          .word    4, 11, 59, 12, 53, 51 , 22 , 44 , 5
empty3:        .word    0
arr4:          .word    11, 4, 9, 10, 3, 7, 51, 14, 22
empty4:        .word    0
arr5:          .word    90, 80, 70, 60, 50, 40, 30, 20, 10, 1
empty5:        .word    0
arr6:          .word    1, 2, 3, 4, 5
empty6:        .word    0
```

I explained why I put empty values between them in the code. Because I didn't implement reading and writing file part, I directly put values here. Other than that, my code prints all increasing subsequences and their sizes. Here is output for first array (3, 10, 7, 9, 4, 11):

```
, Size : 0
3 , Size : 1
10 , Size : 1
3 10 , Size : 2
7 , Size : 1
3 7 , Size : 2
9 , Size : 1
3 9 , Size : 2
7 9 , Size : 2
3 7 9 , Size : 3
4 , Size : 1
3 4 , Size : 2
11 , Size : 1
3 11 , Size : 2
10 11 , Size : 2
3 10 11 , Size : 3
7 11 , Size : 2
3 7 11 , Size : 3
9 11 , Size : 2
3 9 11 , Size : 3
7 9 11 , Size : 3
3 7 9 11 , Size : 4
4 11 , Size : 2
3 4 11 , Size : 3
Longest Increasing Subsequence Array : 3 7 9 11 , Size : 4
```