



CSE312

Operating Systems

Homework 2 / 2021 - 2022

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Global Variables

```

154 #define TABLE_SIZE 5
155 // Page list to test with the algorithm
156 char pages[] = "dgj hj";
157 int size = sizeof(pages) - 1;
158 // Disk & Page Table Representation
159 char page_table[TABLE_SIZE];
160 // Other variables
161 int first_empty_page = 0; // Shows index of the first empty place in page table
162 // Result variables
163 int hit = 0;
164 int page_fault = 0;
165 int pages_loaded = 0;
166 int disk_written = 0;

```

I will start my report by explaining global variables and define values.

- **TABLE_SIZE** = Table size shows the how many pages can fit into page table. An upper limit let's say.
- **pages[]** = Pages are actually char values need to be sorted. Name 'pages' a bit misleading but, these values are stored in disk. This is our disk.
- **size** = sizeof(pages) value takes into account '\0' character which indicates end of the array. But since we are concerned about the size of the disk, we will calculate it by removing 1 from sizeof(pages).
- **page_table** = Page table is our table which we store pages. This page_table stored in our RAM. When we use page algorithms, we actually move pages from pages[] array to page_table array(from disk to array) or vice versa.
- **first_empty_page** = First empty page, as it name implies, shows first index that is empty on the page_table. So we can use it while there are empty spaces on page_table. Once every place is occupied, then this value is set to -1 and so we can understand that, our page_table is full so we need to get rid of one of the pages.
- **hit** = Stores number of hits performed during the sorting algorithms. Hit value incremented when a page is exist in the page_table.
- **page_fault** = Stores number of page faults performed during the sorting algorithms. Page fault occurs when a page doesn't exist in the page_table
- **pages_loaded** = Number of pages loaded. It is incremented every time we call a page replacement algorithm. So it is actually the sum of both hit + page_fault

- **disk_written** = Number of pages written back to the disk. This value incremented every time we get rid of a page to replace it with the new page.

Sorting Algorithms

Bubble Sort

```

171 void bubbleSort() {
172     printArray("Before Bubble Sort : ");
173     int temp;
174     for(int i=0; i<size; i++) {
175         for(int j=0; j<size-(i+1); j++) {
176             lru        (pages[j]);        pages_loaded++;
177             lru        (pages[j+1]);      pages_loaded++;
178             //fifo      (pages[j]);        pages_loaded++;
179             //fifo      (pages[j+1]);      pages_loaded++;
180             //secondChance (pages[j]);      pages_loaded++;
181             //secondChance (pages[j+1]);    pages_loaded++;
182             if(pages[j] > pages[j+1]) {
183                 temp      = pages[j];
184                 pages[j]   = pages[j+1];
185                 pages[j+1] = temp;
186             }
187         }
188     }
189     printArray("After Bubble Sort  : ");
190     printResult();
191 }

```

Bubble sort algorithm is easy to implement. All you have to do is control elements adjacent to each other and if element with the big index is smaller than the low index, than swap them. Do this until every element with big index is bigger than the low index. Bubble sort function doesn't get any parameter because I assigned these functions to keyboard. And when you assign them to keyboard, you can't use any parameters on that function.

After explaining the bubble sort, lets talk about lines between 176-181. These lines shows which paging algorithm will be called. pages array indicates the disk like a said in

the Global Variables section. We are trying to access disk value at `pages[j]` and `pages[j+1]`. So we have to put these two values into our page table. Everytime one of these functions are called I also incremented the `pages_loaded` value by 1. After placing `pages[j]` and `pages[j+1]` to `page_table` (to memory) we are ready to swap them. We will not write these swapped values back to disk until new page takes the place of one of these pages.

Insertion Sort

```

193 void insertionSort() {
194     printArray("Before Insertion Sort : ");
195     int page, j;
196     for(int i=1; i<size; i++) {
197         lru          (pages[i]);          pages_loaded++;
198         //fifo        (pages[i]);          pages_loaded++;
199         //secondChance (pages[i]);          pages_loaded++;
200         page = pages[i];
201         j = i - 1;
202         while(page < pages[j] && j >= 0) {
203             lru          (pages[j]);          pages_loaded++;
204             lru          (pages[j+1]);        pages_loaded++;
205             //fifo        (pages[j]);          pages_loaded++;
206             //fifo        (pages[j+1]);        pages_loaded++;
207             //secondChance (pages[j]);          pages_loaded++;
208             //secondChance (pages[j+1]);        pages_loaded++;
209             pages[j+1] = pages[j];
210             j = j-1;
211         }
212         lru          (pages[i]);          pages_loaded++;
213         //fifo        (pages[i]);          pages_loaded++;
214         //secondChance (pages[i]);          pages_loaded++;
215         pages[j+1] = page;
216     }
217     printArray("After Insertion Sort : ");
218     printResult();
219 }

```

Insertion sort says that, start from index 0, take an element and move it to the left until there is no element that is lower than this element and repeat this until everything is sorted.

Different from bubble sort, we have much more call for a page replacement algorithm. Rule is simple: if any line contains any pages[] value, before using it, call a page replacement algorithm for that page. This is what I did here and in other sorting algorithms. printArray function prints the current status of the disk and printResult function prints the hit/miss/ values, etc.

Quick Sort

```

193 > void insertionSort() { ...
220
221 void quickSort() {
222     printArray("Before Quick Sort : ");
223     int low = 0;
224     int high = size - 1;
225     quickSortStart(low, high);
226     printArray("After Quick Sort : ");
227     printResult();
228 }
229
230 void quickSortStart(int low, int high) {
231     if(low < high) {
232         int pivot = partition(low, high);
233
234         quickSortStart(low, pivot - 1);
235         quickSortStart(pivot + 1, high);
236     }
237 }
238
239 > int partition(int low, int high) { ...
278

```

Partition function is very long so I didn't put the whole code of it here but all it does partitioning the pages according to pivot value. Different from other sorting algorithms, this sorting algorithm calls a page replacement algorithm A LOT. Bubble sort calls 2 times, Insertion Sort calls 4 times and Quick sort calls 6 times. 3 times more than Bubble sort.

Page Replacement Algorithms

FIFO (First In First Out)

```

323 int nextIndex = 0;
324 void fifo(char page) {
325     printTable(page);
326     int index = 0;
327     // First loop all the page table to find given page
328     for(int i=0; i<TABLE_SIZE; i++) {
329         // Increment hit and reference count by one if the page is present and return
330         if(page_table[i] == page) {
331             hit++;
332             reference_count[i]++;
333             printTable(page);
334             return;
335         }
336     }
337     // If the page is not present in page_table, look if there is an empty place
338     if(first_empty_page != -1) {
339         page_table[first_empty_page] = page;    // Insert in the first empty space
340         page_fault++;                          // Increment the page fault by 1
341         if(first_empty_page == (TABLE_SIZE-1)) // If this is the last empty space at the page table
342             first_empty_page = -1;             // Assign first empty page to -1
343         else
344             first_empty_page++;                // Else, increment the first_empty_page by 1
345     }
346     else {
347         index = nextIndex % TABLE_SIZE;
348         page_table[index] = page;
349         nextIndex++;
350         disk_written++;
351         page_fault++;                          // Incrementing the page fault by 1
352     }
353     printTable(page);
354 }
355

```

There are 3 situations in FIFO page replacement algorithm. These are:

1. Finding the page in the page_table before retrieving it from disk : Lines : 328-336
2. If page is not in the page_table, then looking for empty space in page_table to put new page from the disk to this place(Lines: 338-345). There could be 2 options now:
 - a. **There are more than one empty space in the page table (Lines: 341-342)**
 - b. **There is only one space left in the page table (Lines: 343-344)**
3. If page table is full, then according to “nextIndex” value defined in line 323 removing the page with the longest staying time. (Lines: 346-352)

Second Chance

```

356 int chance_table[TABLE_SIZE] = {0};
357 void secondChance(char page) {
358     // First try to find page in the page table
359     for(int i=0; i<TABLE_SIZE; i++) {
360         // Increment hit and reference count by one if the page is present and return
361         if(page_table[i] == page) {
362             hit++;
363             chance_table[i] = 1;          // Set reference value to 1 to give it a second chance
364             printTable(page);
365             return;
366         }
367     }
368     // If page not found in the page table, then check if the table is empty or not
369     if(first_empty_page != -1) {
370         page_table[first_empty_page] = page;    // Insert in the first empty space
371         page_fault++;                          // Increment the page fault by 1
372         chance_table[first_empty_page] = 0;    // Set page with 0 reference point (yet)
373         if(first_empty_page == (TABLE_SIZE-1)) // If this is the last empty space at the page table
374             first_empty_page = -1;           // Assign first empty page to -1
375         else
376             first_empty_page++;              // Else, increment the first_empty_page by 1
377     }
378     // Else, if page table is not empty and page is not in the page table, then traverse the page table
379     // To find a page with no reference point
380     else {
381         for(int i=0; i<TABLE_SIZE*2; i++) {
382             if(chance_table[i] == 0) {
383                 page_table[i] = page;
384                 printTable(page);
385                 return;
386             }
387             i = i % TABLE_SIZE;
388             chance_table[i] = 0;
389         }
390         page_fault++;
391         disk_written++;
392     }
393     printTable(page);
394 }
395

```

Like FIFO, there are 3 situations in FIFO page replacement algorithm. These are:

1. Finding the page in the page_table before retrieving it from disk : Lines : 359-367
2. If page is not in the page_table, then looking for empty space in page_table to put new page from the disk to this place(Lines: 369-377). There could be 2 options now:
 - a. **There are more than one empty space in the page table** (Lines: 373-374)
 - b. **There is only one space left in the page table** (Lines: 375-376)
3. If page is not in the page_table and page_table is full, then according to chance_table, find a page with the reference value 0. Each time you came across with a page with reference value 1, then make it value back to 0 and continue

searching. Start over again if you can't find any page with reference 0 (this time there will be because we set their reference values back to 0)

LRU (Least Recently Used)

```

283 int reference_count[TABLE_SIZE] = {0}; // How many times a page is referenced
284 void lru(char page) {
285     // First loop all the page table to find given page
286     for(int i=0; i<TABLE_SIZE; i++) {
287         // Increment hit and reference count by one if the page is present and return
288         if(page_table[i] == page) {
289             hit++;
290             reference_count[i]++;
291             printTable(page);
292             return;
293         }
294     }
295     // If the page is not present in page_table, look if there is an empty place
296     if(first_empty_page != -1) {
297         page_table[first_empty_page] = page; // Insert in the first empty space
298         page_fault++; // Increment the page fault by 1
299         if(first_empty_page == (TABLE_SIZE-1)) // If this is the last empty space at the page table
300             first_empty_page = -1; // Assign first empty page to -1
301         else
302             first_empty_page++; // Else, increment the first_empty_page by 1
303     }
304     // If there is no empty space in the page table, then find the page with the least reference
305     else {
306         // Finding index of the least referenced table
307         int maxReference = 10000;
308         int index = 0;
309         for(int i=0; i<TABLE_SIZE; i++) {
310             if(reference_count[i] < maxReference) {
311                 maxReference = reference_count[i];
312                 index = i;
313             }
314         }
315         // After finding the index of the least referenced page, we replace it with the new page
316         page_table[index] = page;
317         page_fault++; // Incrementing the page fault by 1
318         disk_written++;
319     }
320     printTable(page);
321 }

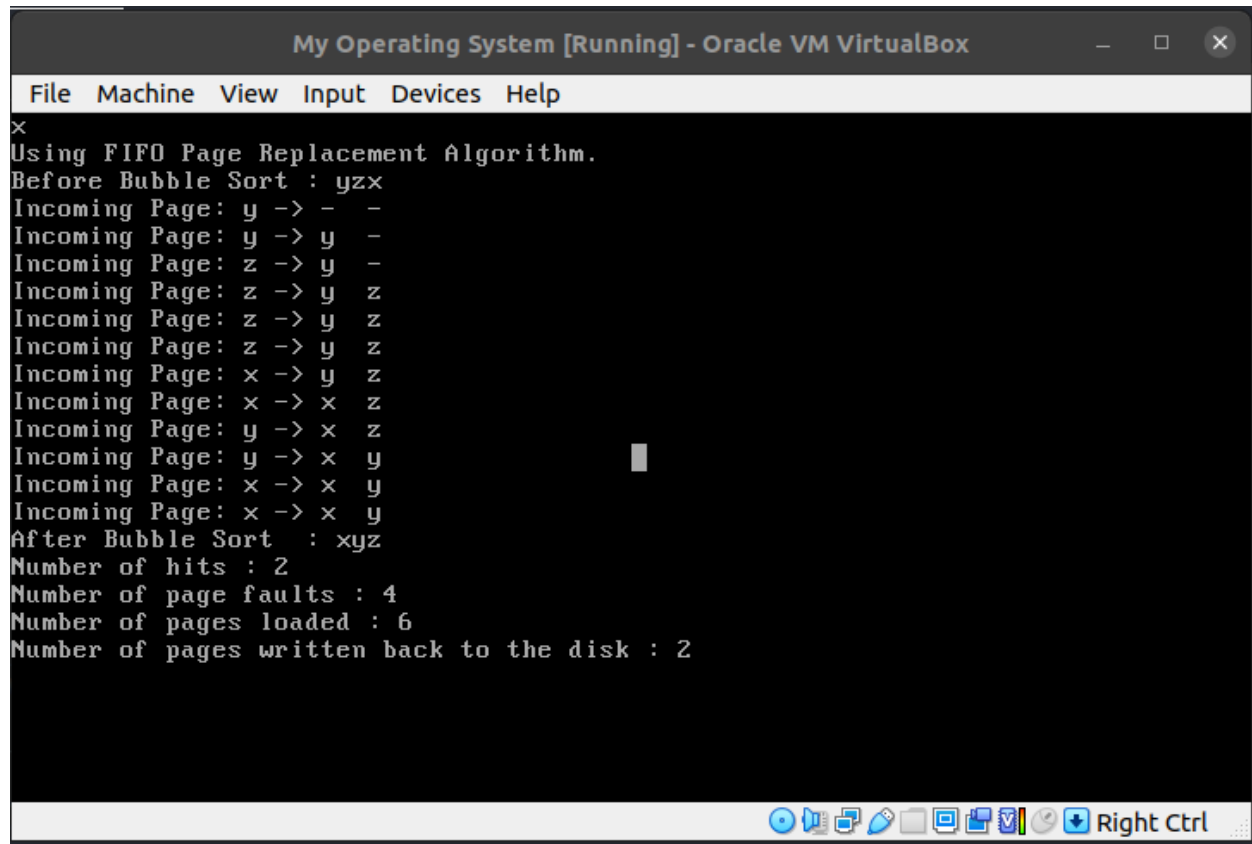
```

First 2 steps are same with the other paging algorithms. But last step is different in each of them. For LRU (Least Recently Used) algorithm I can explain it like this:

1. First create a value with max reference possible (there is no INF so i used 1000)
2. Then starting from beginning of the page table, looking every page and changing maxReference value if there is any page with reference value less than **maxReference**.
3. After finding page with the least reference value, we can insert our new page into this page's place and increment page_fault and disk_written by 1.

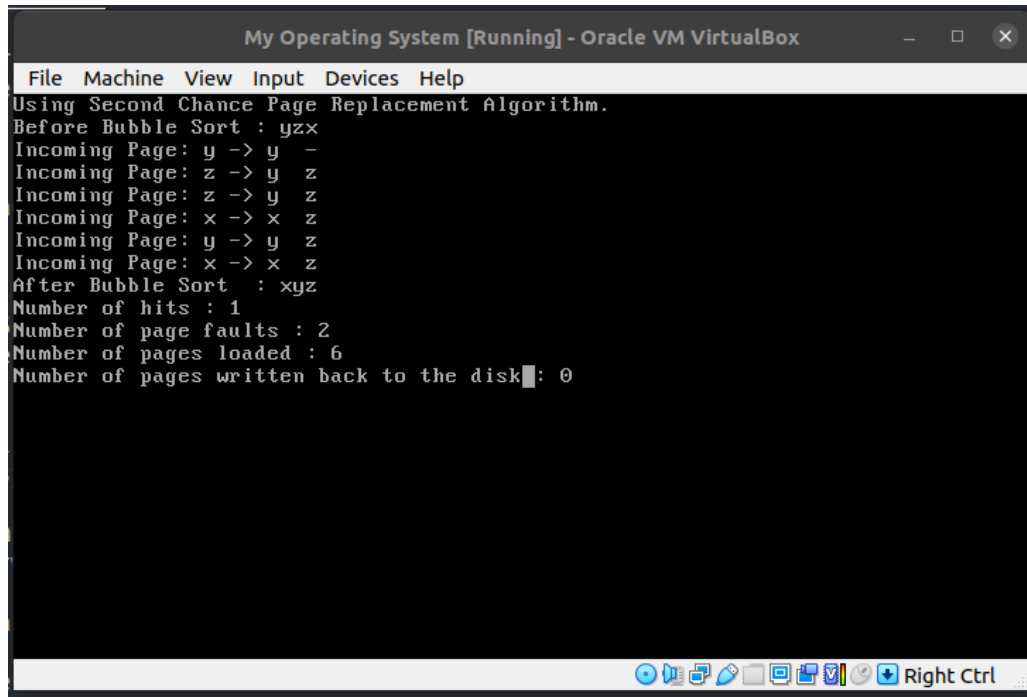
Example Outputs

Bubble Sort - FIFO



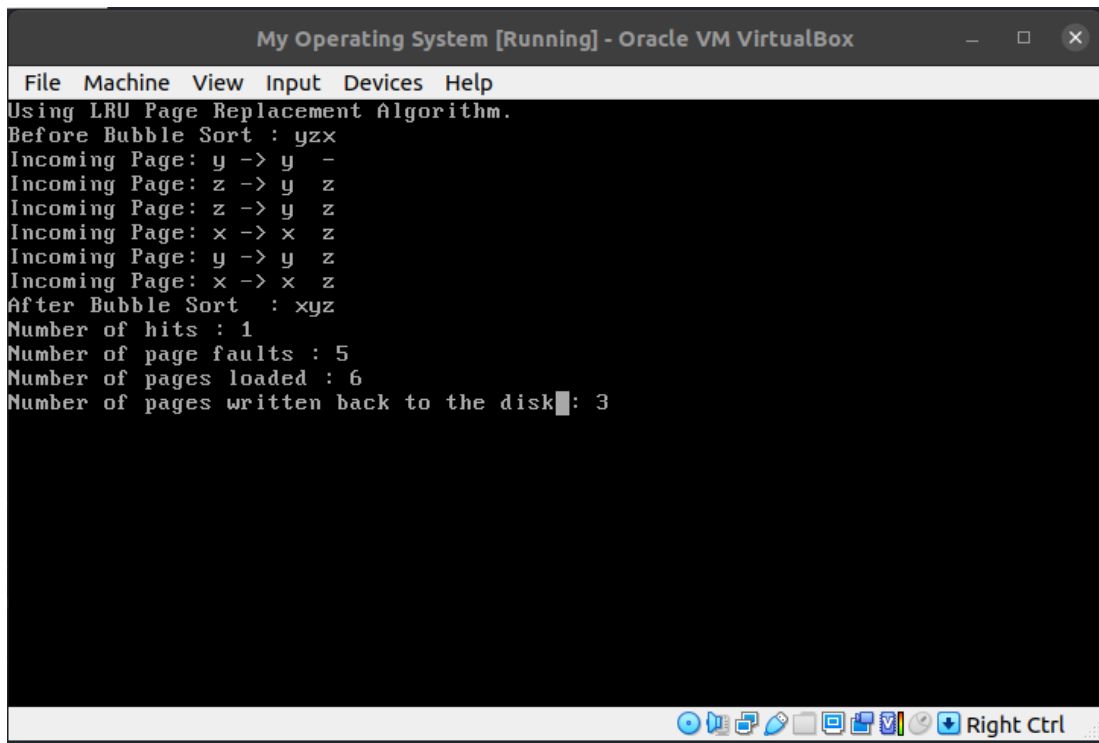
```
My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
x
Using FIFO Page Replacement Algorithm.
Before Bubble Sort : yzx
Incoming Page: y -> - -
Incoming Page: y -> y -
Incoming Page: z -> y -
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: x -> y z
Incoming Page: x -> x z
Incoming Page: y -> x z
Incoming Page: y -> x y
Incoming Page: x -> x y
Incoming Page: x -> x y
After Bubble Sort : xyz
Number of hits : 2
Number of page faults : 4
Number of pages loaded : 6
Number of pages written back to the disk : 2
```

Bubble Sort - Second Chance



```
My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using Second Chance Page Replacement Algorithm.
Before Bubble Sort : yzx
Incoming Page: y -> y -
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: x -> x z
Incoming Page: y -> y z
Incoming Page: x -> x z
After Bubble Sort : xyz
Number of hits : 1
Number of page faults : 2
Number of pages loaded : 6
Number of pages written back to the disk : 0
```

Bubble Sort - LRU



```
My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using LRU Page Replacement Algorithm.
Before Bubble Sort : yzx
Incoming Page: y -> y -
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: x -> x z
Incoming Page: y -> y z
Incoming Page: x -> x z
After Bubble Sort : xyz
Number of hits : 1
Number of page faults : 5
Number of pages loaded : 6
Number of pages written back to the disk : 3
```

Insertion Sort - FIFO

```

My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
X
Using FIFO Page Replacement Algorithm.
Before Insertion Sort : yzx
Incoming Page: z -> - -
Incoming Page: z -> z -
Incoming Page: z -> z -
Incoming Page: z -> z -
Incoming Page: x -> z -
Incoming Page: x -> z x
Incoming Page: z -> z x
Incoming Page: z -> z x
Incoming Page: x -> z x
Incoming Page: x -> z x
Incoming Page: y -> z x
Incoming Page: y -> y x
Incoming Page: z -> y x
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: z -> y z
After Insertion Sort : xyz
Number of hits : 4
Number of page faults : 4
Number of pages loaded : 8
Number of pages written back to the disk : 2

```

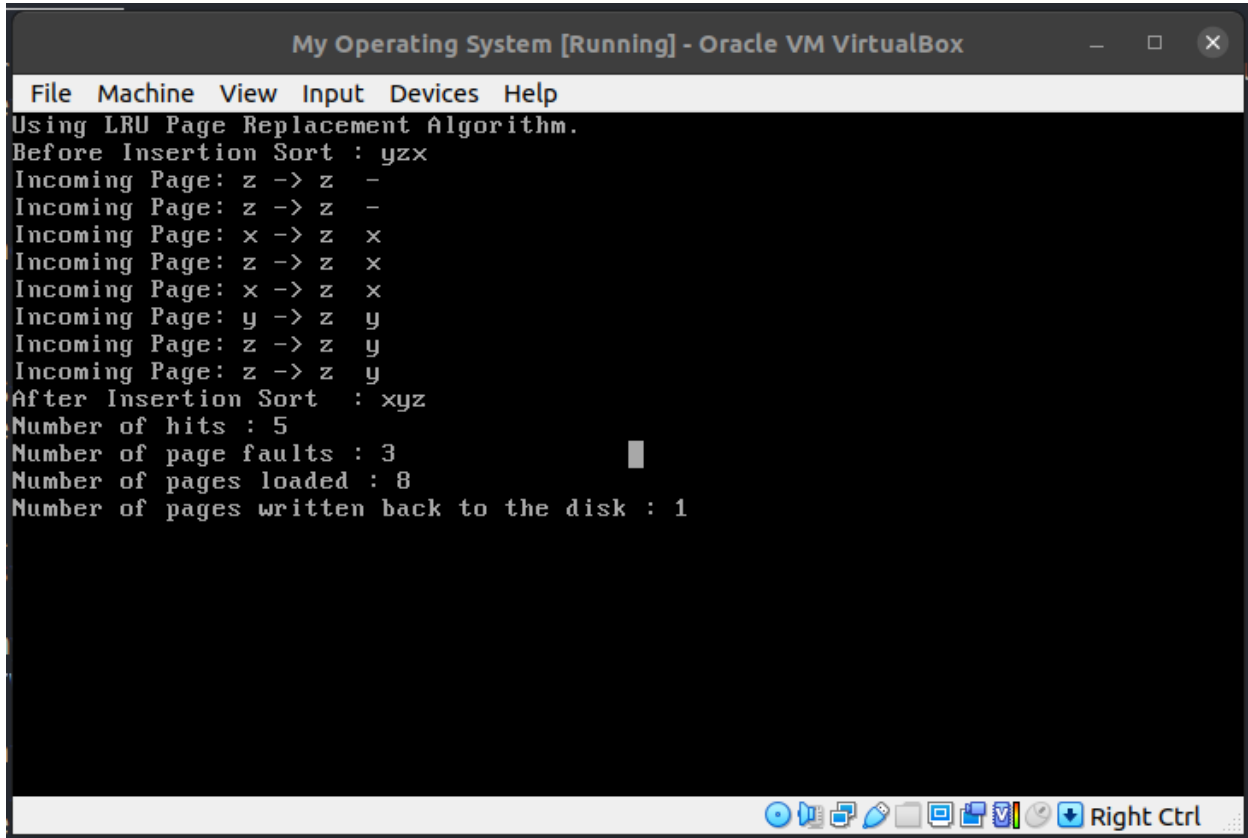
Insertion Sort - Second Chance

```

My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using Second Chance Page Replacement Algorithm.
Before Insertion Sort : yzx
Incoming Page: z -> z -
Incoming Page: z -> z -
Incoming Page: x -> z x
Incoming Page: z -> z x
Incoming Page: x -> z x
Incoming Page: y -> z y
Incoming Page: z -> z y
Incoming Page: z -> z y
After Insertion Sort : xyz
Number of hits : 5
Number of page faults : 2
Number of pages loaded : 8
Number of pages written back to the disk : 0

```

Insertion Sort - LRU



```
My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using LRU Page Replacement Algorithm.
Before Insertion Sort : yzx
Incoming Page: z -> z -
Incoming Page: z -> z -
Incoming Page: x -> z x
Incoming Page: z -> z x
Incoming Page: x -> z x
Incoming Page: y -> z y
Incoming Page: z -> z y
Incoming Page: z -> z y
After Insertion Sort : xyz
Number of hits : 5
Number of page faults : 3
Number of pages loaded : 8
Number of pages written back to the disk : 1
```

Quick Sort - FIFO

```

My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using FIFO Page Replacement Algorithm.
Before Quick Sort : yzx
Incoming Page: x -> - -
Incoming Page: x -> x -
Incoming Page: y -> x -
Incoming Page: y -> x y
Incoming Page: z -> x y
Incoming Page: z -> z y
Incoming Page: y -> z y
Incoming Page: y -> z y
Incoming Page: x -> z y
Incoming Page: x -> z x
Incoming Page: y -> z x
Incoming Page: y -> y x
Incoming Page: z -> y x
Incoming Page: z -> y z
Incoming Page: z -> y z
Incoming Page: y -> y z
Incoming Page: y -> y z
After Quick Sort : xyz
Number of hits : 3
Number of page faults : 6
Number of pages loaded : 9
Number of pages written back to the disk : 4

```

Quick Sort - Second Chance

```

My Operating System [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Using Second Chance Page Replacement Algorithm.
Before Quick Sort : yzx
Incoming Page: x -> x -
Incoming Page: y -> x y
Incoming Page: z -> z y
Incoming Page: y -> z y
Incoming Page: x -> x y
Incoming Page: y -> x y
Incoming Page: z -> z y
Incoming Page: z -> z y
Incoming Page: y -> z y
After Quick Sort : xyz
Number of hits : 4
Number of page faults : 2
Number of pages loaded : 9
Number of pages written back to the disk : 0

```

